approach

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Monday

Median of two sorted Arrays

arr1[] = [1 3 4 7 10 12]

arr2[] = [2 3 6 15]

1 3 4 7 | 10 12
2 | 3 6 15
(5) (5)

First here we taken 4 elements from array 1 and 1 element from array 2 and checking conditions

Here max is 7 and in rgt side low/min is 3 so it does not form

Nao take 3 elements from arr1 and 2 elements from arr2

1 3 4 | 7 10 12
2 3 | 6 15

Now we check conditions
max in left side is 4 and min in right side is 6 so it is in sorted order

Now take 2 elements from arr1 and 3 elements from arr2

1 3 | 4 7 10 12
2 3 6 | 15

Now we check conditions
max in left is 6 and min in right side is 4 so it is not in sorted order.

Now to Calculate median

l_1 | r_1
1 3 4 | 7 10 12
2 3 l_2 | r_2 6 15

we take $\frac{\max(l_1, l_2) + \min(r_1, r_2)}{2}$

if $n_1 + n_2 \% 2 == 0$ even

otherwise
only $\max(l_1, l_2)$

NOW (4)

1 3 4 7 l_1 | r_1 10 12
2 l_2 | r_2 3 6 15

$l_1 > r_2$

high = mid - 1

1 3 l_1 | r_1 4 7 10 12
2 3 6 l_2 | r_2 15

$l_2 > r_1$

low = mid + 1

if ($l_1 <= r_2$ && $l_2 <= r_1$)
Valid

1 3 4 | 7 10
2 3 | 6 15

we do in binary search

I will take smaller array bcuz for max I can put smaller elements at one side so when compared to bigger array time complexity is small

arr1[] = [7 12 14 15]

arr2[] = [1 2 3 4 9 11]

low 0, mid 2, high 4

So. left over are 5-2 = 3 mid2 = 3

7 12 l_1 | 14 r_1 15

1 2 3 l_2 | 4 r_2 9 11

$l_1 > r_2$ condition fails I mean 12 < 4 false

So ~~low~~ high = mid - 1

low	mid	high
0	0	1

left over are 5

0 are these

l_1 | 7 r_1 12 14 15

1 2 3 4 9 l_2 | 11 r_2

So we assume Integer. min_value

So here $l_2 > r_1$

So we move right side

low = mid + 1 mid1 = 1

and leftover are 5-1 mid2 = 4

low	mid	high
1	1	

mid1

r_1

7 l_1 | 12 14 15

r_2

1 2 3 4 l_2 | 9 11

mid2

7

1 2 14 15

Here $l_1 < r_2$ and $l_2 < r_1$

So valid

So now based on total length we find median

if length is even we find

$$\frac{\max(l_1, l_2) + \min(r_1, r_2)}{2}$$

else
we take max(l1, l2)

```
public int solve (int[] a, int [] b) int n;int
{   int n1=a.length;
    int n2=b.length
  if(n1>n2)
     return solve(b,a)

   int low=0
   int high= a.length;        -> int left= (n1+n2+1)/2;
   while(low<=high)           -> int n= n1+n2;
   {
      int mid1 = low+(high-low)/2;
      int mid2 = left-mid1;
      int l1 = Integer. MIN_VALUE;
      int l2 =  ''      ''   ;
      int r1 = Integer. MAX_VALUE;
      int r2 =  ''      ''   ''
      if(mid1<n1)  r1 = a[mid1];
      if(mid2<n2) r2 = b[mid2];
       if(mid1-1>=0)  l1 = a[mid1-1];
       if(mid2-1>=0)  l2=a b[mid2-1];

      if (l1<=r2 && l2<=r1){
         if(n%2==1)  return
                      Math.max(l1,l2);
        return (double)(Math.max(l1,l2)+
                        Math.min(r1,r2))/2)
      }
```

```
else if(l1>s2)
 {
    high=mid1-1;
 }
 else
 {
   low=mid1+1;
 }
}
return 0;
}
```