

## Homework 7

1. A)

$Dec(k, m)$
$r \leftarrow \{0, 1\}^\lambda$
$m := F(k, c \  r)$
return $m$

b) To show that the scheme has CPA security, we will start with the cpa-L library.

$\mathcal{L}_{\text{cpa-L}}$
$k \leftarrow \text{KeyGen}$
$EAVESDROP(m_L, m_R)$
$c := \text{Enc}(k, m_L)$
return $c$

First, we will factor out the call to Enc() into a subroutine, changing nothing about the operation of the library.

$EAVESDROP(m_L, m_R)$
$c := \text{CTXT}(m_L)$
return $c$

 $\diamond$ 

$\mathcal{L}_{\text{cpa-real}}$
$k \leftarrow \text{KeyGen}$
$\text{CTXT}(m)$
$c := \text{Enc}(k, m)$
return $c$

Second, we will replace cpa-real with cpa-rand. This change should be indistinguishable.

$EAVESDROP(m_L, m_R)$
$c := \text{CTXT}(m_L)$
return $c$

 $\diamond$ 

$\mathcal{L}_{\text{cpa-rand}}$
$\text{CTXT}(m)$
$c := C$
return $c$

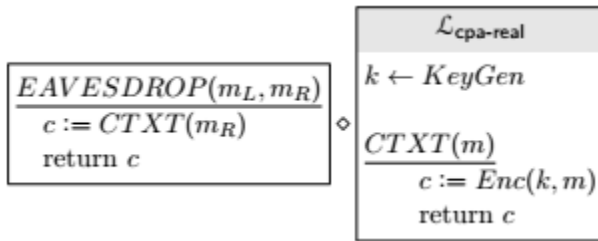
Next, we will change the argument that is passed into CTXT. This won't affect the operation of the library.

$EAVESDROP(m_L, m_R)$
$c := \text{CTXT}(m_R)$
return $c$

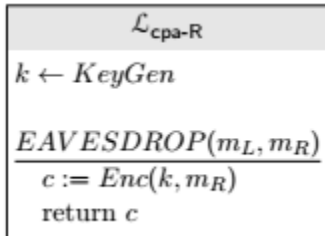
 $\diamond$ 

$\mathcal{L}_{\text{cpa-rand}}$
$\text{CTXT}(m)$
$c := C$
return $c$

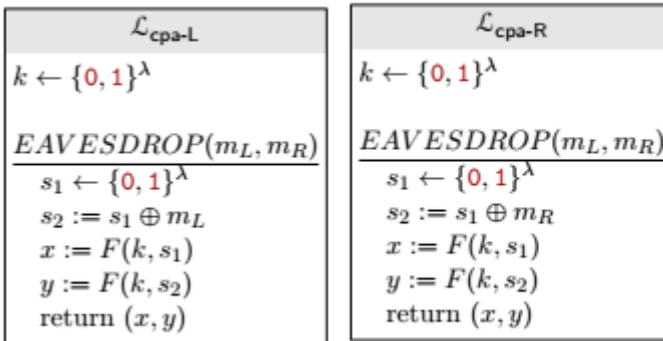
Now we can change cpa-rand back to cpa-real.



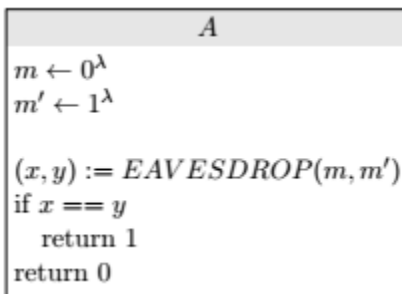
And finally, we can inline the cpa library.



## 2. Libraries:



Calling program:



$$\Pr[A \diamond \text{cpa-L} = 1] = 1$$

$$\Pr[A \diamond \text{cpa-R} = 1] = 0$$

$$\text{Advantage: } |\Pr[A \diamond \text{cpa-L} = 1] - \Pr[A \diamond \text{cpa-R} = 1]| = 1 - 0 = 1$$