

Design And Analysis Of Algorithms

Bonus Homework Report

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1 Boyer–Moore Implementation

The Boyer–Moore algorithm was implemented using the *bad character heuristic*. The general logic is familiar with the GeeksForGeeks implementation.

1.1 Unicode-Aware Modification

Standard Boyer–Moore implementations typically assume ASCII characters with a fixed-size table (usually 256). However, during testing, cases involving Unicode characters were encountered.

To solve this, the implementation dynamically determines the maximum character value appearing in both the text and the pattern, and constructs the bad character table accordingly.

This ensures algorithm is reliable, but causes overhead as seen in the results table. **Some test cases that are best case for Boyer Moore are not faster with this implementation.**

1.2 Design Summary

- Uses bad character heuristic only
- Dynamically sized bad character table
- Correct handling of empty patterns and edge cases

2 Pre-Analysis Strategy

The selection strategy is based on the following features:

1. **Pattern Length:** Very short patterns use the Naive algorithm.
2. **Repetition Ratio:** Detected via the LPS array; high repetition favors KMP.
3. **Character Diversity:** High diversity favors Boyer–Moore due to better skipping behavior.
4. **Text Entropy:** Shannon entropy is used to estimate randomness. High entropy favors Rabin–Karp, while low entropy favors KMP.

3 Experimental Results

See the output of the ./test.sh command in the Figures 1 and 2.

| Test Case | RabinKarp (μs) | GoCrazy (μs) | BoyerMoore (μs) | KMP (μs) | Naive (μs) | Winner |
|--------------------------------|----------------|--------------|-----------------|----------|------------|--------------|
| Simple Match | 3.473 | N/A | 4.438 | 3.134 | 3.786 | 🏆 KMP |
| No Match | 1.157 | N/A | 1.742 | 1.431 | 0.741 | 🏆 Naive |
| Single Character | 6.863 | N/A | 6.777 | 6.186 | 6.758 | 🏆 KMP |
| Pattern at End | 5.808 | N/A | 3.077 | 1.543 | 1.002 | 🏆 Naive |
| Pattern at Beginning | 2.421 | N/A | 2.140 | 1.423 | 1.074 | 🏆 Naive |
| Overlapping Patterns | 2.802 | N/A | 3.055 | 1.825 | 2.102 | 🏆 KMP |
| Long Text Multiple Matches | 5.738 | N/A | 6.229 | 7.279 | 4.736 | 🏆 Naive |
| Pattern Longer Than Text | 0.160 | N/A | 0.154 | 0.677 | 0.190 | 🏆 BoyerMoore |
| Entire Text Match | 1.270 | N/A | 2.054 | 1.374 | 1.063 | 🏆 Naive |
| Repeating Pattern | 3.483 | N/A | 3.794 | 2.307 | 2.045 | 🏆 Naive |
| Case Sensitive | 2.561 | N/A | 2.878 | 2.400 | 1.768 | 🏆 Naive |
| Numbers and Special Characters | 3.810 | N/A | 3.565 | 3.540 | 2.287 | 🏆 Naive |
| Unicode Characters | 7.004 | N/A | 130.285 | 2.668 | 1.651 | 🏆 Naive |
| Very Long Text | 39.241 | N/A | 30.738 | 41.983 | 28.691 | 🏆 Naive |
| Pattern with Spaces | 4.026 | N/A | 3.796 | 3.570 | 2.625 | 🏆 Naive |
| All Same Character | 14.056 | N/A | 14.122 | 6.874 | 13.418 | 🏆 KMP |
| Alternating Pattern | 21.103 | N/A | 20.233 | 12.492 | 17.430 | 🏆 KMP |
| Long Pattern | 11.901 | N/A | 10.021 | 12.272 | 7.852 | 🏆 Naive |
| Pattern at Boundaries | 2.535 | N/A | 3.058 | 3.374 | 1.934 | 🏆 Naive |
| Near Matches | 1.778 | N/A | 3.664 | 1.978 | 1.652 | 🏆 Naive |
| Empty Pattern | 0.868 | N/A | 0.822 | 1.074 | 1.131 | 🏆 BoyerMoore |
| Empty Text | 0.160 | N/A | 0.163 | 0.403 | 0.179 | 🏆 RabinKarp |
| Both Empty | 0.527 | N/A | 0.316 | 0.542 | 0.326 | 🏆 BoyerMoore |
| Single Character Pattern | 7.551 | N/A | 7.363 | 6.598 | 6.191 | 🏆 Naive |
| Complex Overlap | 3.156 | N/A | 3.855 | 2.357 | 2.713 | 🏆 KMP |
| DNA Sequence | 10.977 | N/A | 9.100 | 8.131 | 9.933 | 🏆 KMP |
| Palindrome Pattern | 6.389 | N/A | 5.141 | 6.409 | 4.171 | 🏆 Naive |
| Worst Case for Naive | 6.859 | N/A | 12.809 | 8.849 | 6.190 | 🏆 Naive |
| Best Case for Boyer-Moore | 3.119 | N/A | 5.957 | 6.594 | 1.477 | 🏆 Naive |
| KMP Advantage Case | 1.535 | N/A | 2.500 | 8.588 | 4.395 | 🏆 RabinKarp |

Figure 1: Execution time comparison of string matching algorithms across test cases

| Test Case | Choice | PreA+Choice (μs) | vs Naive | vs RabinKarp | vs BoyerMoore | vs KMP |
|--------------------------------|------------|------------------|-----------|--------------|---------------|-----------|
| Simple Match | Naive | 1.78 | N/A | +0.72 μs | -1.76 μs | +0.53 μs |
| No Match | Naive | 0.41 | N/A | +0.04 μs | -0.93 μs | -0.28 μs |
| Single Character | Naive | 1.59 | N/A | -0.02 μs | -1.59 μs | +0.15 μs |
| Pattern at End | BoyerMoore | 3.62 | +2.98 μs | +2.86 μs | N/A | +2.00 μs |
| Pattern at Beginning | BoyerMoore | 2.79 | +2.45 μs | +2.32 μs | N/A | +1.73 μs |
| Overlapping Patterns | Naive | 1.03 | N/A | -0.02 μs | -0.67 μs | -0.11 μs |
| Long Text Multiple Matches | Naive | 1.47 | N/A | -0.35 μs | -1.60 μs | -1.15 μs |
| Pattern Longer Than Text | BoyerMoore | 1.40 | +1.28 μs | +1.29 μs | N/A | +0.89 μs |
| Entire Text Match | BoyerMoore | 2.89 | +2.59 μs | +2.45 μs | N/A | +1.89 μs |
| Repeating Pattern | Naive | 1.08 | N/A | -0.12 μs | -0.75 μs | -0.07 μs |
| Case Sensitive | Naive | 0.75 | N/A | -0.25 μs | -0.97 μs | -0.45 μs |
| Numbers and Special Characters | BoyerMoore | 3.53 | +2.79 μs | +2.17 μs | N/A | +2.26 μs |
| Unicode Characters | Naive | 1.10 | N/A | +0.33 μs | -49.09 μs | -0.04 μs |
| Very Long Text | Naive | 9.42 | N/A | -1.27 μs | -2.26 μs | -4.19 μs |
| Pattern with Spaces | BoyerMoore | 3.65 | +2.54 μs | +2.64 μs | N/A | +2.27 μs |
| All Same Character | KMP | 4.77 | -0.22 μs | -0.42 μs | -1.96 μs | N/A |
| Alternating Pattern | Naive | 9.64 | N/A | -1.78 μs | -0.58 μs | +1.97 μs |
| Long Pattern | BoyerMoore | 7.99 | +4.34 μs | +4.49 μs | N/A | +3.25 μs |
| Pattern at Boundaries | Naive | 1.13 | N/A | -0.46 μs | -1.38 μs | -0.65 μs |
| Near Matches | Naive | 1.08 | N/A | +0.25 μs | -0.62 μs | -1.75 μs |
| Empty Pattern | Naive | 1.21 | N/A | +0.50 μs | +0.53 μs | +0.51 μs |
| Empty Text | Naive | 0.44 | N/A | +0.26 μs | +0.27 μs | -0.05 μs |
| Both Empty | Naive | 0.51 | N/A | -0.22 μs | +0.19 μs | -0.18 μs |
| Single Character Pattern | Naive | 4.44 | N/A | -0.10 μs | -1.05 μs | -0.47 μs |
| Complex Overlap | KMP | 2.70 | +0.95 μs | +0.92 μs | -0.06 μs | N/A |
| DNA Sequence | Naive | 4.87 | N/A | -0.38 μs | -1.34 μs | -0.56 μs |
| Palindrome Pattern | RabinKarp | 17.25 | +15.78 μs | N/A | +14.34 μs | +14.17 μs |
| Worst Case for Naive | KMP | 17.64 | +8.22 μs | +15.30 μs | +12.63 μs | N/A |
| Best Case for Boyer-Moore | Naive | 2.45 | N/A | +0.11 μs | -1.01 μs | -0.58 μs |
| KMP Advantage Case | KMP | 19.69 | +13.63 μs | +17.31 μs | +16.25 μs | N/A |

Figure 2: Execution time comparison with pre-analysis strategy (green: pre-analysis was faster, red: pre-analysis was slower).

3.1 Key Observations

- KMP dominates in highly repetitive patterns and structured data (e.g., DNA sequences).
- Boyer–Moore excels when character diversity is high and mismatches occur early.
- Rabin–Karp performs best in high-entropy texts but is sensitive to hashing overhead.

- Pre-analysis algorithm is not perfectly designed since overhead (red labels) is seen in most of the test cases in the Figure 2. It sometimes selects the inappropriate algorithm, too.

References

- GeeksForGeeks, *Boyer Moore Algorithm for Pattern Searching*. <https://www.geeksforgeeks.org/dsa/boyer-moore-algorithm-for-pattern-searching/>