

Introduction to CMOS VLSI Design

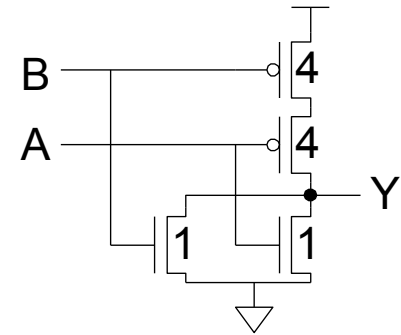
Lecture 9: Circuit Families

Outline

- ☐ Pseudo-nMOS Logic
- ☐ Dynamic Logic
- ☐ Pass Transistor Logic

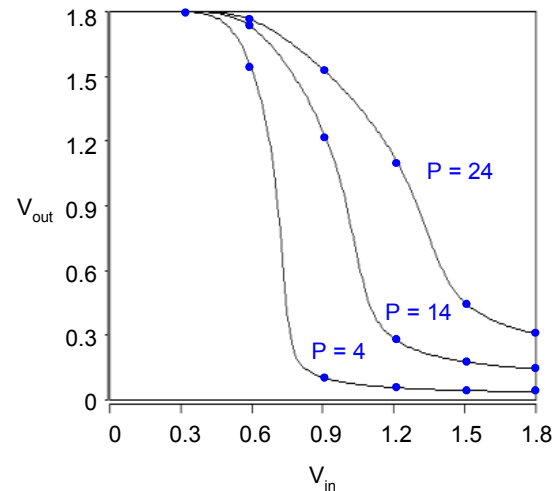
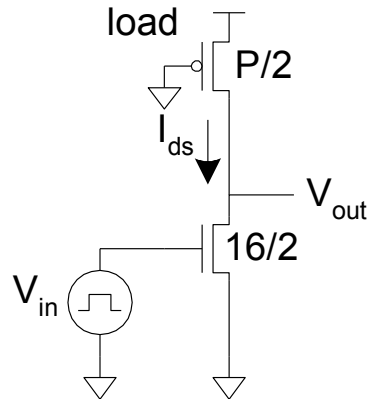
Introduction

- ❑ What makes a circuit fast?
 - $I = C \, dV/dt \rightarrow t_{pd} \propto (C/I) \Delta V$
 - low capacitance
 - high current
 - small swing
- ❑ Logical effort is proportional to C/I
- ❑ pMOS are the enemy!
 - High capacitance for a given current
- ❑ Can we take the pMOS capacitance off the input?
- ❑ Various circuit families try to do this...



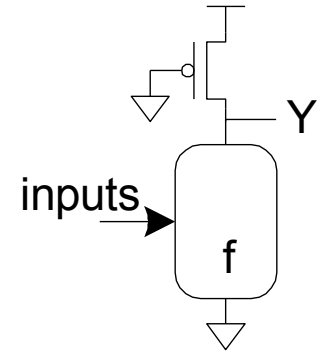
Pseudo-nMOS

- ❑ In the old days, nMOS processes had no pMOS
 - Instead, use pull-up transistor that is always ON
- ❑ In CMOS, use a pMOS that is always ON
 - *Ratio* issue
 - Make pMOS about $\frac{1}{4}$ effective strength of pulldown network

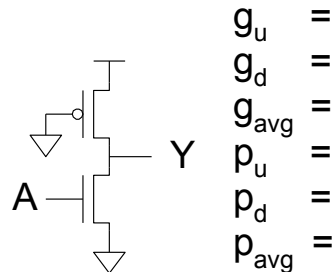


Pseudo-nMOS Gates

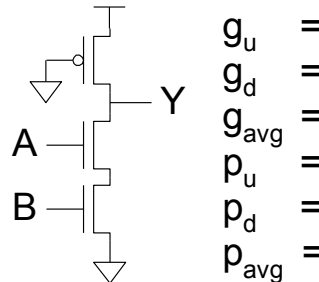
- ❑ Design for unit current on output to compare with unit inverter.
- ❑ pMOS fights nMOS



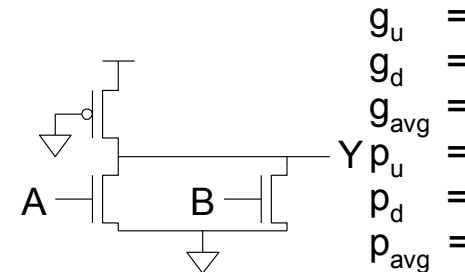
Inverter



NAND2

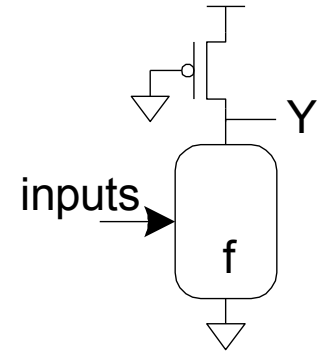


NOR2

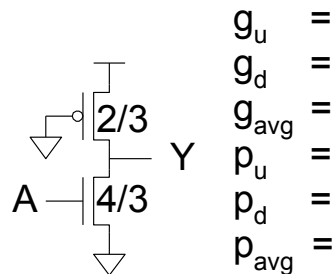


Pseudo-nMOS Gates

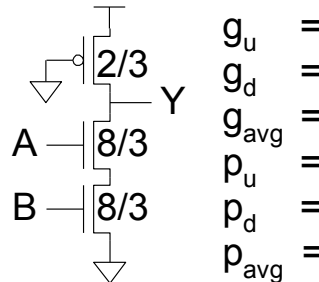
- ❑ Design for unit current on output to compare with unit inverter.
- ❑ pMOS fights nMOS



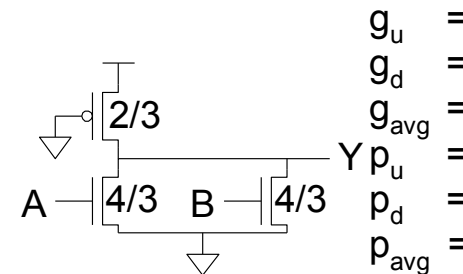
Inverter



NAND2

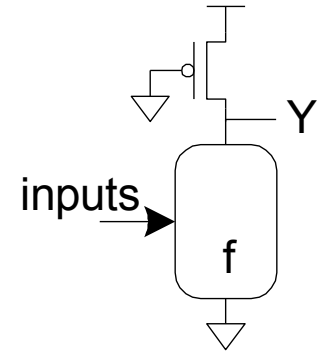


NOR2

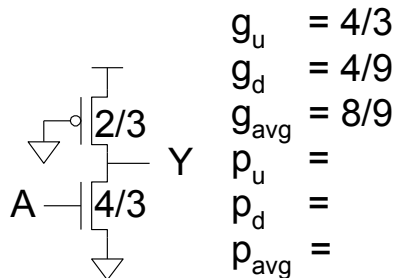


Pseudo-nMOS Gates

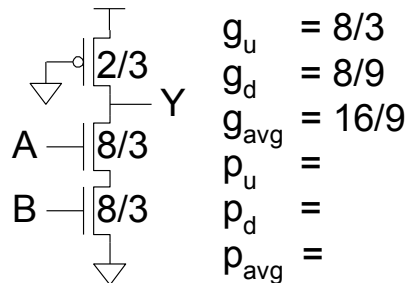
- ❑ Design for unit current on output to compare with unit inverter.
- ❑ pMOS fights nMOS



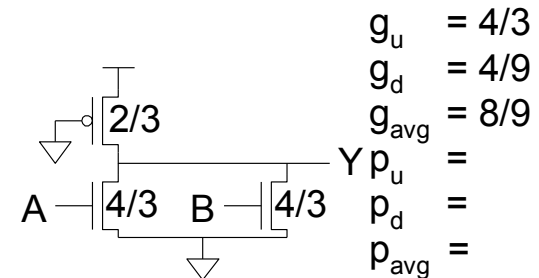
Inverter



NAND2

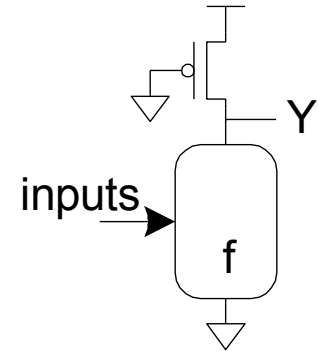


NOR2

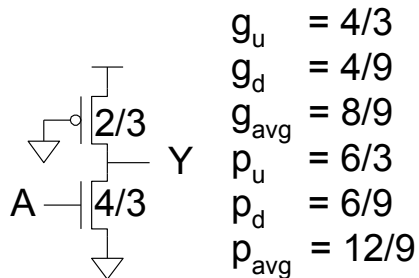


Pseudo-nMOS Gates

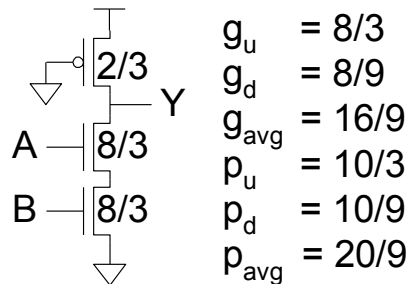
- ❑ Design for unit current on output to compare with unit inverter.
- ❑ pMOS fights nMOS



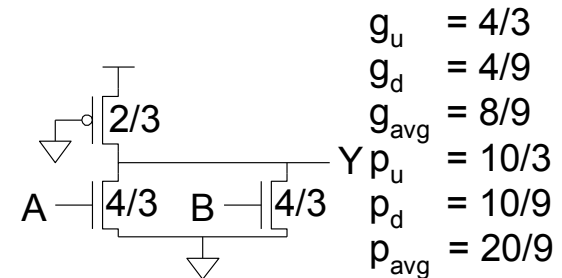
Inverter



NAND2



NOR2



Pseudo-nMOS Design

- ❑ Ex: Design a k-input AND gate using pseudo-nMOS. Estimate the delay driving a fanout of H

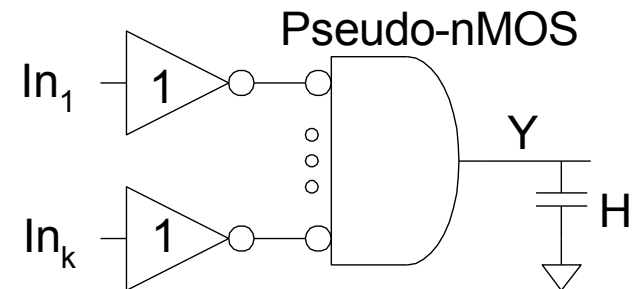
❑ $G =$

❑ $F =$

❑ $P =$

❑ $N =$

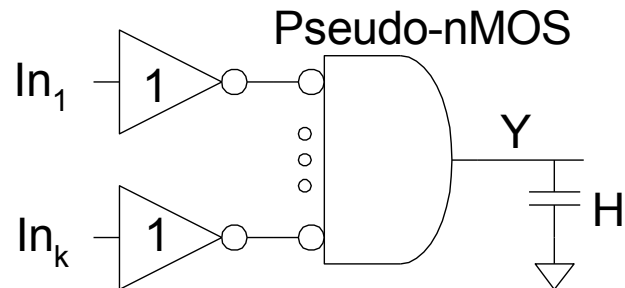
❑ $D =$



Pseudo-nMOS Design

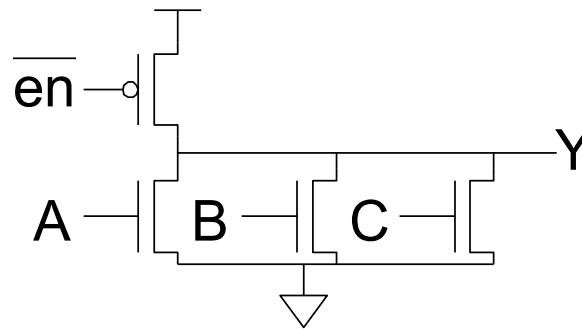
- ❑ Ex: Design a k-input AND gate using pseudo-nMOS. Estimate the delay driving a fanout of H

- ❑ $G = 1 * 8/9 = 8/9$
- ❑ $F = GBH = 8H/9$
- ❑ $P = 1 + (4+8k)/9 = (8k+13)/9$
- ❑ $N = 2$
- ❑ $D = NF^{1/N} + P = \frac{4\sqrt{2H}}{3} + \frac{8k+13}{9}$



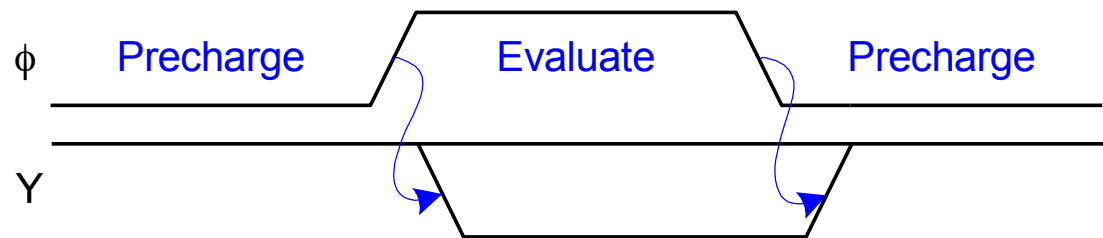
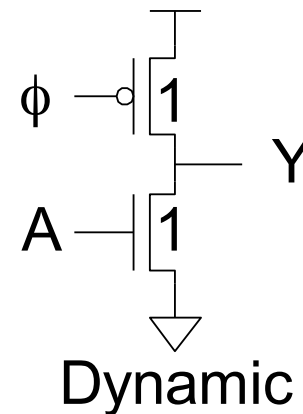
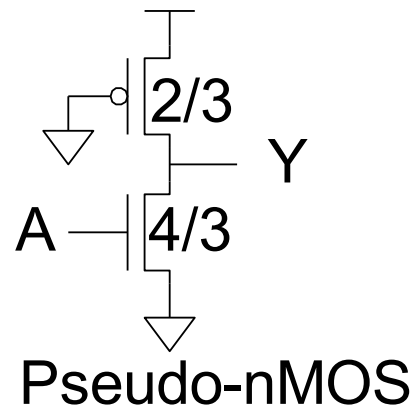
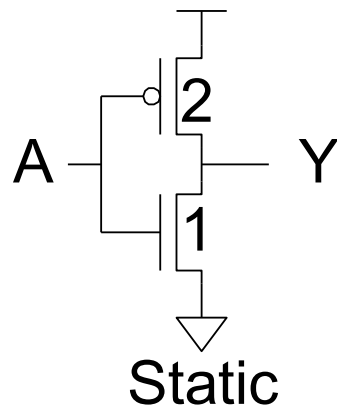
Pseudo-nMOS Power

- ❑ Pseudo-nMOS draws power whenever $Y = 0$
 - Called static power $P = I \cdot V_{DD}$
 - A few mA / gate * 1M gates would be a problem
 - This is why nMOS went extinct!
- ❑ Use pseudo-nMOS sparingly for wide NORs
- ❑ Turn off pMOS when not in use



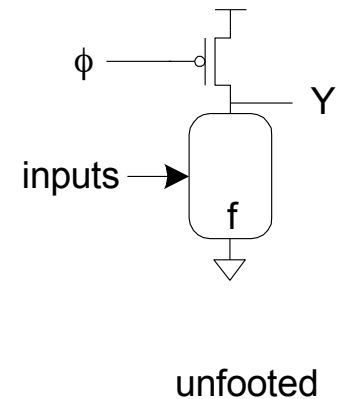
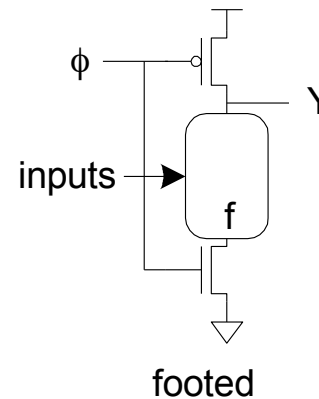
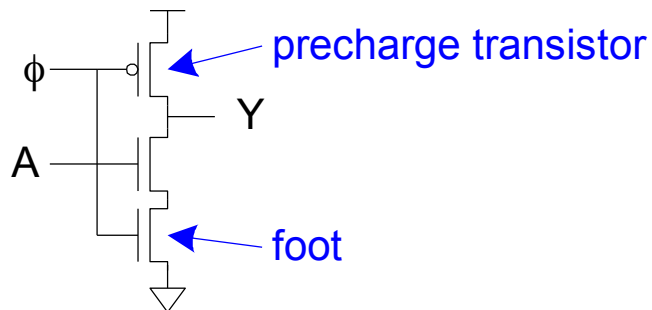
Dynamic Logic

- ❑ *Dynamic* gates uses a clocked pMOS pullup
- ❑ Two modes: *precharge* and *evaluate*



The Foot

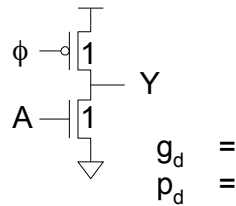
- ❑ What if pulldown network is ON during precharge?
- ❑ Use series evaluation transistor to prevent fight.



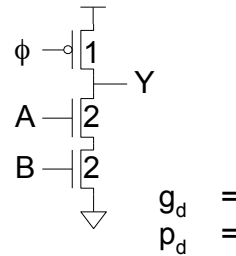
Logical Effort

unfooted

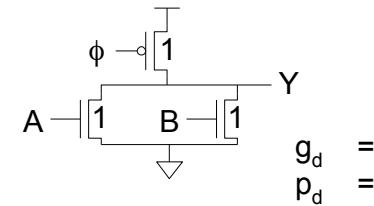
Inverter



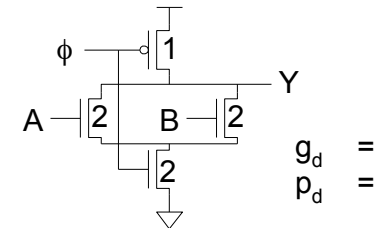
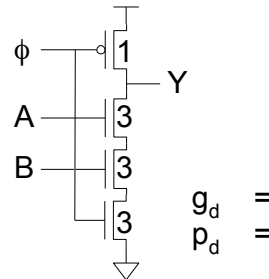
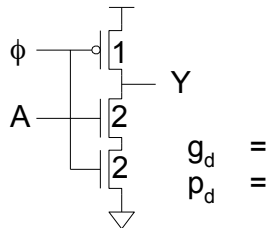
NAND2



NOR2



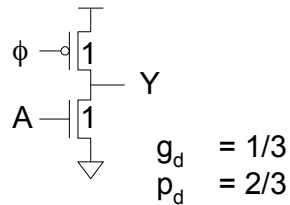
footed



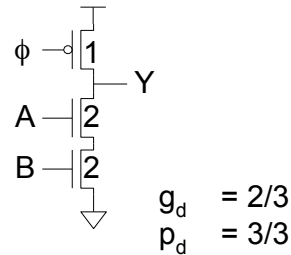
Logical Effort

unfooted

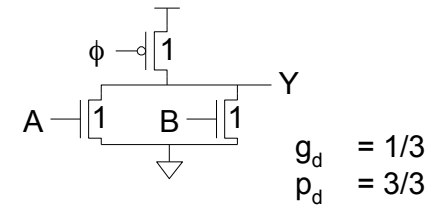
Inverter



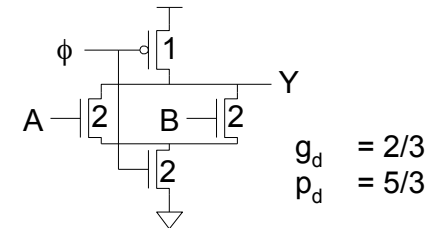
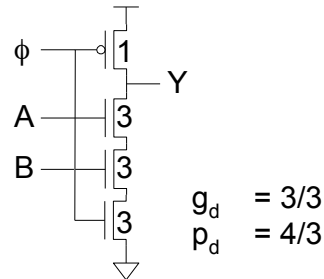
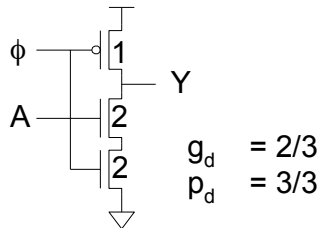
NAND2



NOR2

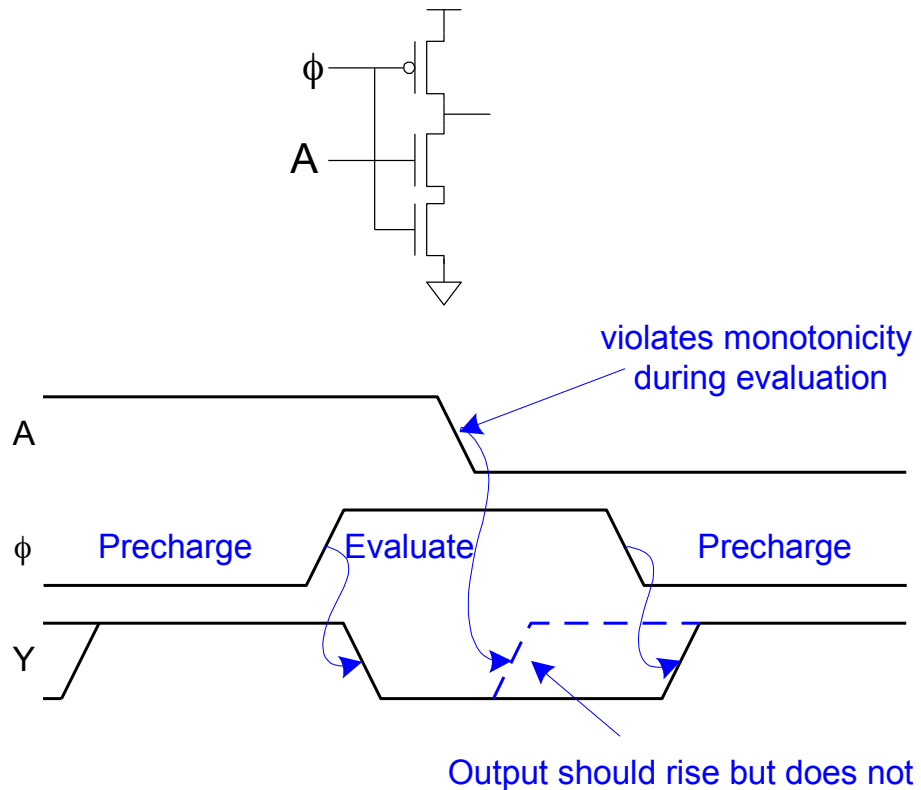


footed



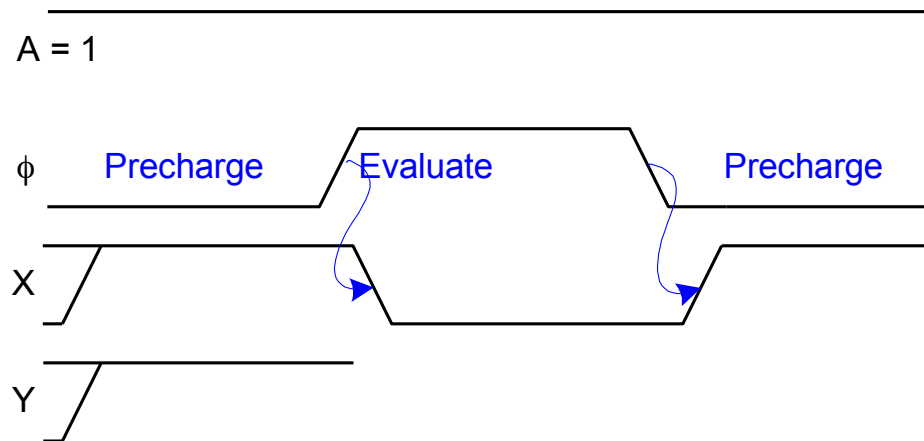
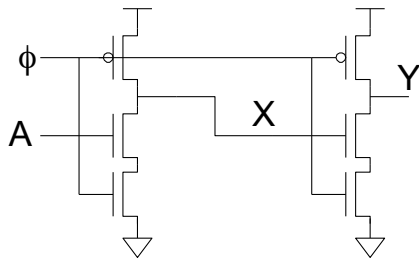
Monotonicity

- ❑ Dynamic gates require *monotonically rising* inputs during evaluation
 - 0 \rightarrow 0
 - 0 \rightarrow 1
 - 1 \rightarrow 1
 - But not 1 \rightarrow 0



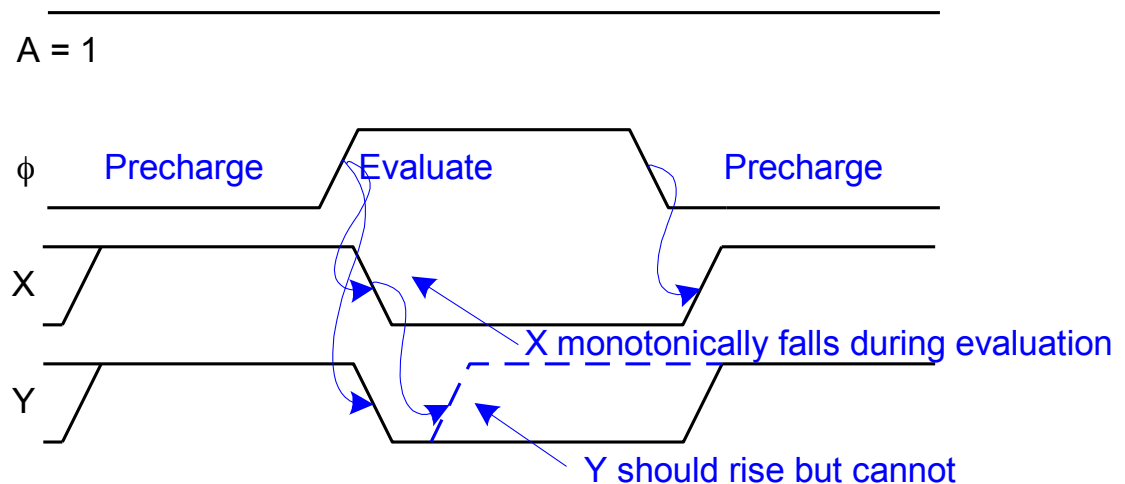
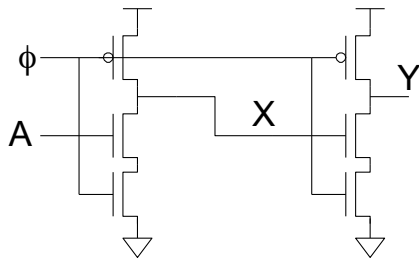
Monotonicity Woes

- ❑ But dynamic gates produce monotonically falling outputs during evaluation
- ❑ Illegal for one dynamic gate to drive another!



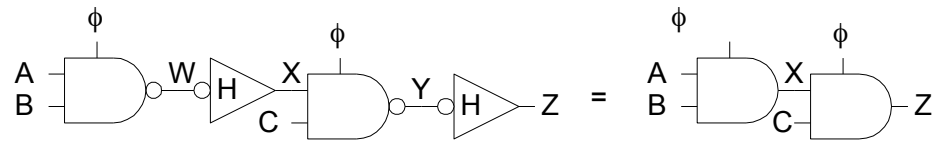
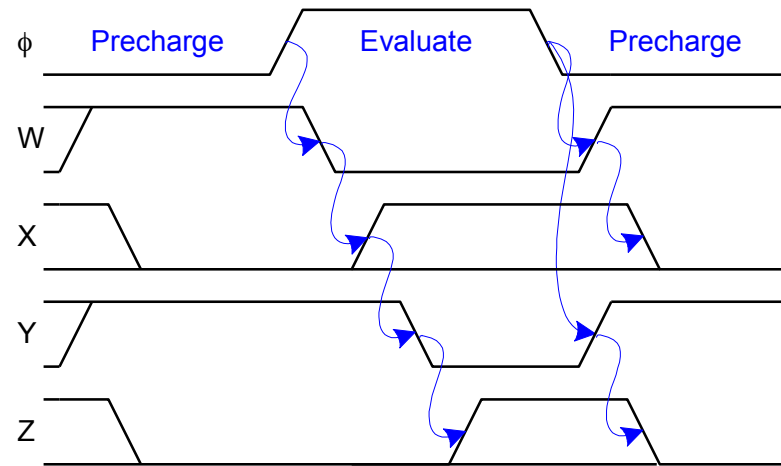
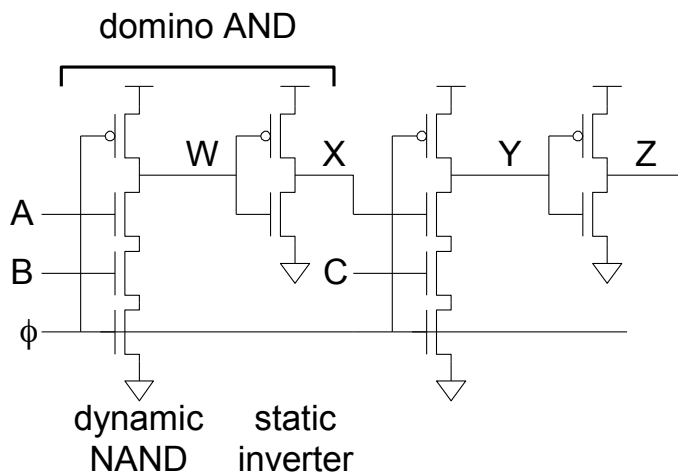
Monotonicity Woes

- ❑ But dynamic gates produce monotonically falling outputs during evaluation
- ❑ Illegal for one dynamic gate to drive another!



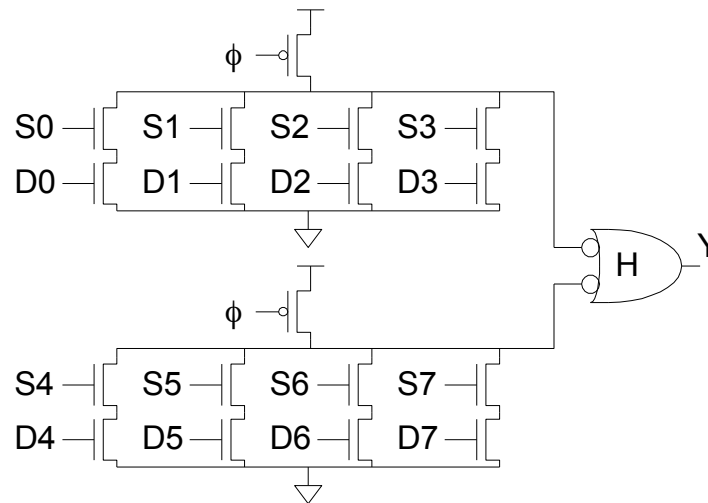
Domino Gates

- ❑ Follow dynamic stage with inverting static gate
 - Dynamic / static pair is called domino gate
 - Produces monotonic outputs



Domino Optimizations

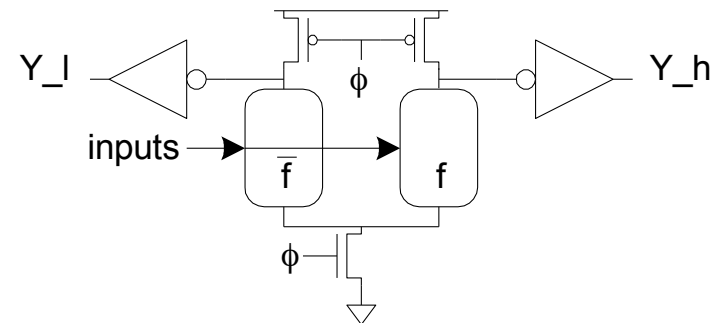
- ❑ Each domino gate triggers next one, like a string of dominos toppling over
- ❑ Gates evaluate sequentially but precharge in parallel
- ❑ Thus evaluation is more critical than precharge
- ❑ HI-skewed static stages can perform logic



Dual-Rail Domino

- ❑ Domino only performs noninverting functions:
 - AND, OR but not NAND, NOR, or XOR
- ❑ Dual-rail domino solves this problem
 - Takes true and complementary inputs
 - Produces true and complementary outputs

sig_h	sig_l	Meaning
0	0	Precharged
0	1	'0'
1	0	'1'
1	1	invalid

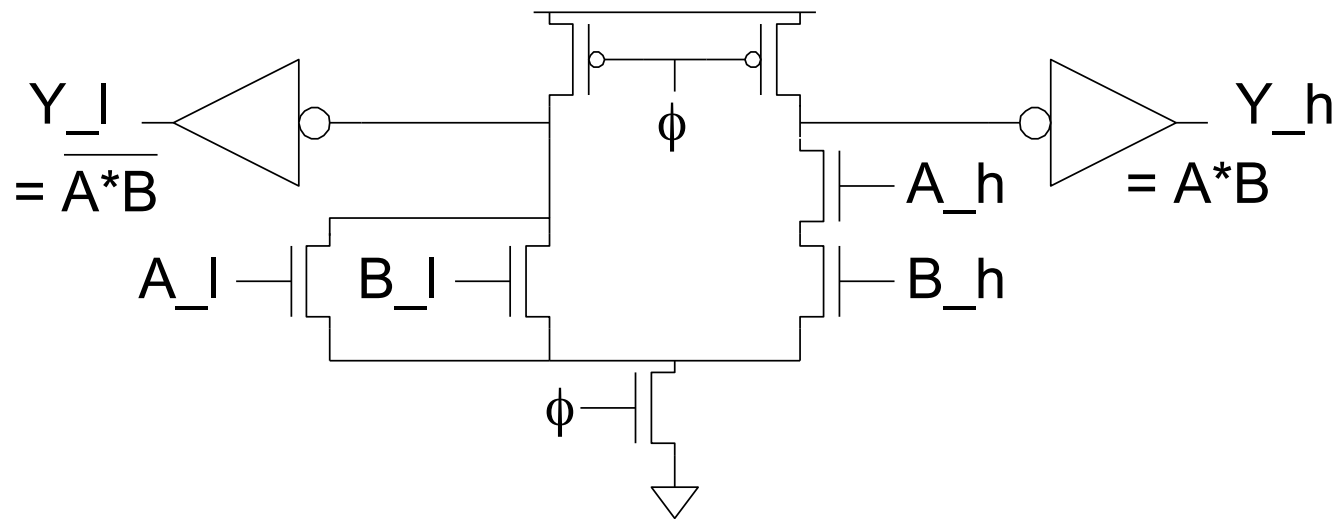


Example: AND/NAND

- ❑ Given A_h, A_l, B_h, B_l
- ❑ Compute $Y_h = A * B, Y_l = \sim(A * B)$

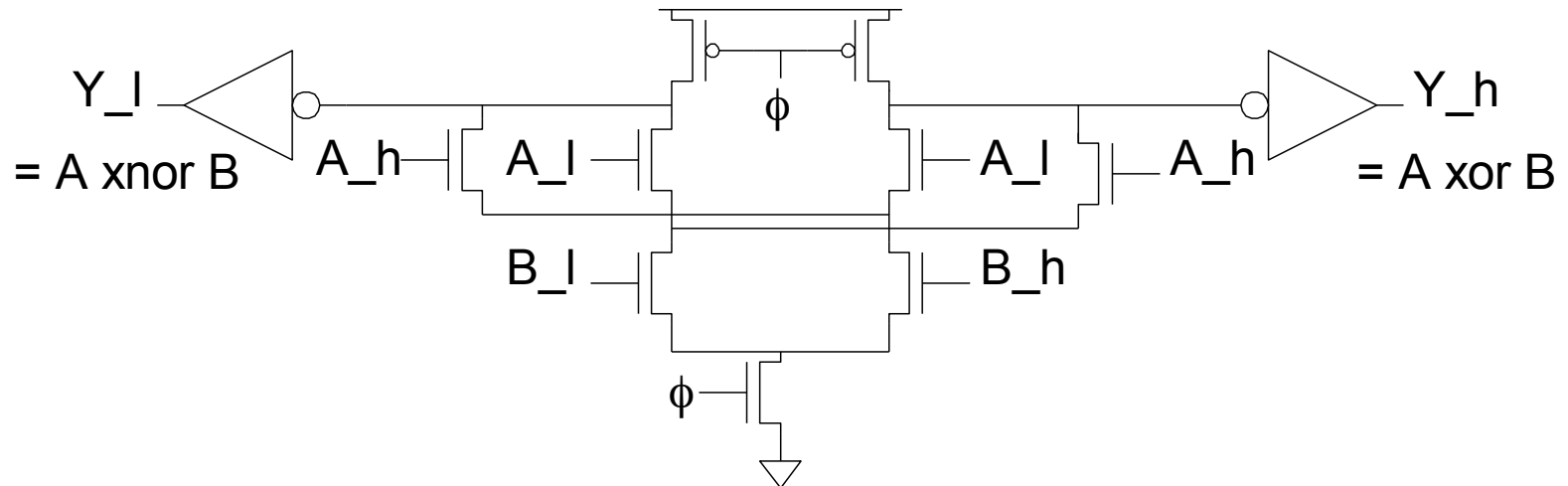
Example: AND/NAND

- ❑ Given A_h, A_l, B_h, B_l
- ❑ Compute $Y_h = A * B, Y_l = \sim(A * B)$
- ❑ Pulldown networks are conduction complements



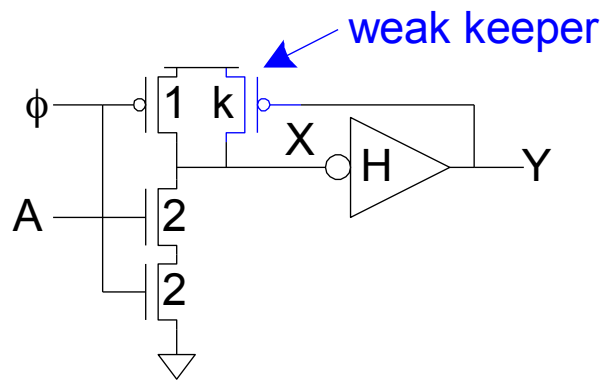
Example: XOR/XNOR

- Sometimes possible to share transistors



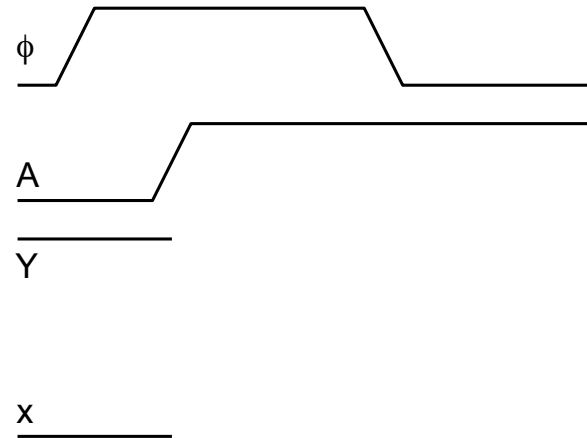
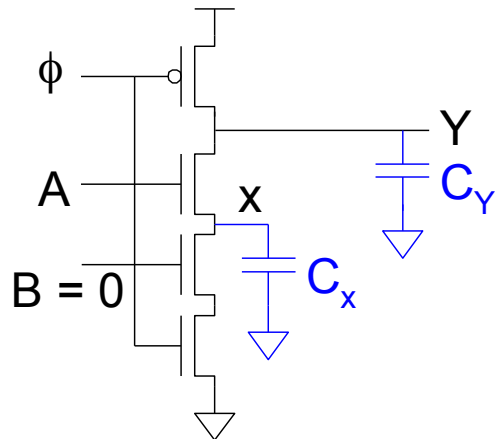
Leakage

- ❑ Dynamic node floats high during evaluation
 - Transistors are leaky ($I_{OFF} \neq 0$)
 - Dynamic value will leak away over time
 - Formerly milliseconds, now nanoseconds!
- ❑ Use keeper to hold dynamic node
 - Must be weak enough not to fight evaluation



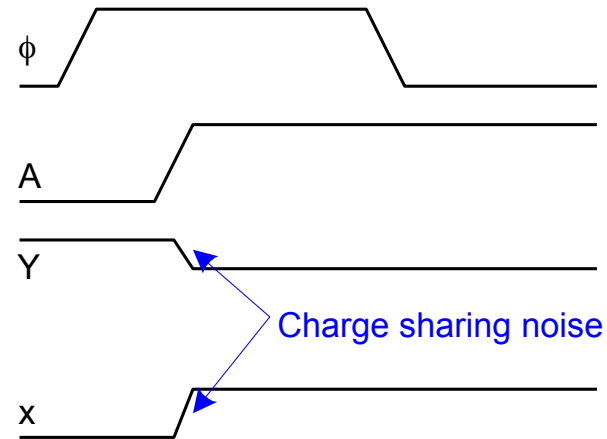
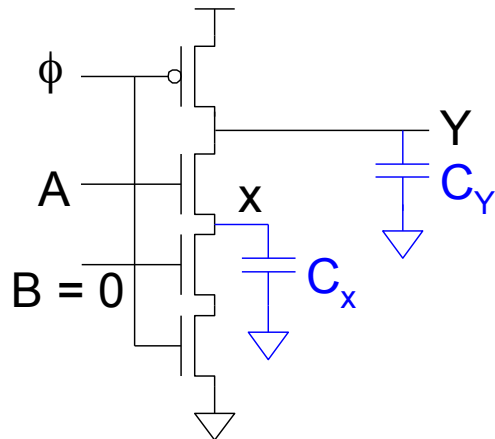
Charge Sharing

- ❑ Dynamic gates suffer from charge sharing



Charge Sharing

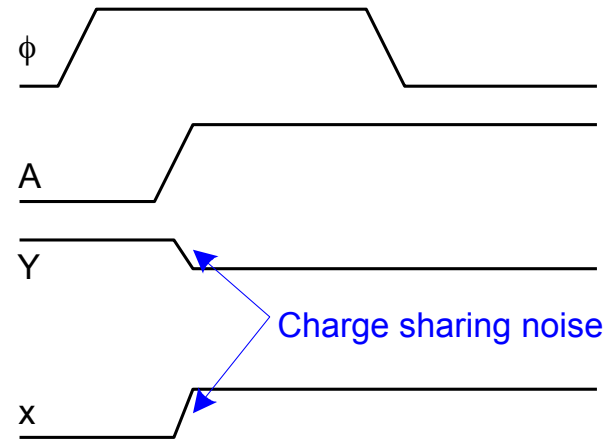
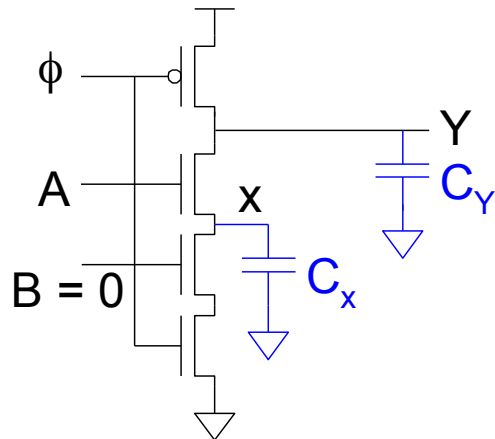
- ❑ Dynamic gates suffer from charge sharing



$$V_x = V_Y =$$

Charge Sharing

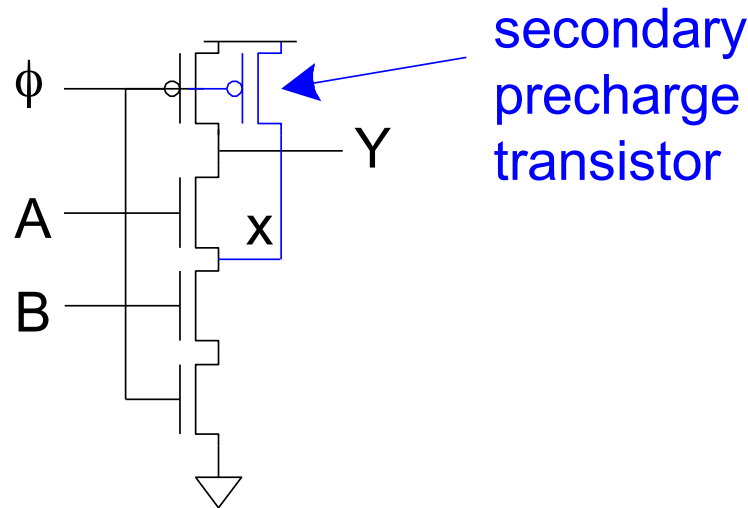
- ❑ Dynamic gates suffer from charge sharing



$$V_x = V_Y = \frac{C_Y}{C_x + C_Y} V_{DD}$$

Secondary Precharge

- ❑ Solution: add secondary precharge transistors
 - Typically need to precharge every other node
- ❑ Big load capacitance C_Y helps as well



Noise Sensitivity

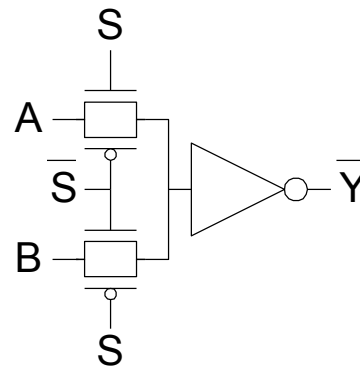
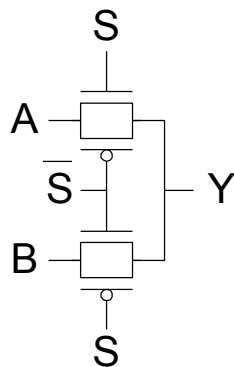
- ❑ Dynamic gates are very sensitive to noise
 - Inputs: $V_{IH} \approx V_{tn}$
 - Outputs: floating output susceptible noise
- ❑ Noise sources
 - Capacitive crosstalk
 - Charge sharing
 - Power supply noise
 - Feedthrough noise
 - And more!

Domino Summary

- ❑ Domino logic is attractive for high-speed circuits
 - 1.5 – 2x faster than static CMOS
 - But many challenges:
 - Monotonicity
 - Leakage
 - Charge sharing
 - Noise
- ❑ Widely used in high-performance microprocessors

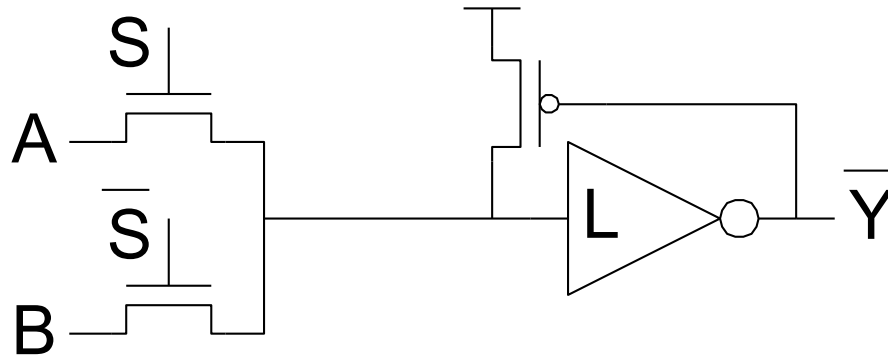
Pass Transistor Circuits

- ❑ Use pass transistors like switches to do logic
- ❑ Inputs drive diffusion terminals as well as gates
- ❑ CMOS + Transmission Gates:
 - 2-input multiplexer
 - Gates should be restoring



LEAP

- ❑ **LEA**n integration with **P**ass transistors
- ❑ Get rid of pMOS transistors
 - Use weak pMOS feedback to pull fully high
 - Ratio constraint



CPL

- ❑ Complementary Pass-transistor Logic
 - Dual-rail form of pass transistor logic
 - Avoids need for ratioed feedback
 - Optional cross-coupling for rail-to-rail swing

