

National University of Computer and Emerging Sciences (Lahore Campus)

Quiz 6: Link Layer (Chapter 6)

Name: _____

Roll No: _____

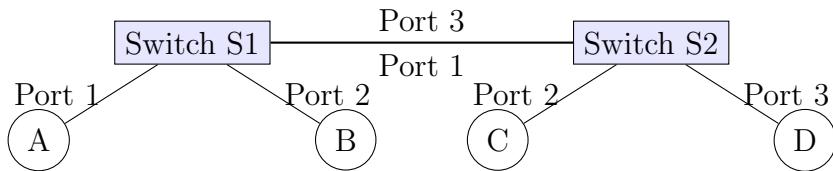
Section: BSE-6B1 (Spring 2026)

1. (5 points) Switch Tables

Consider the switched LAN topology shown below. The switch tables are initially empty. The following sequence of frame transmissions occurs:

1. Node A sends a frame to Node D.
2. Node D replies with a frame to Node A.
3. Node C sends a frame to Node D.

Show the state of the Switching Table for **Switch S2** after these three events. Format: (MAC Address, Interface). Also, show the same for **Switch S1**.



Solution:

- **Switch S1 Table:**

MAC Address	Interface
A	Port 1
D	Port 3

Note: S1 learns A from step 1. S1 learns D from step 2. S1 never sees the frame from C (step 3) because S2 knows D is on Port 3 and does not flood the frame back to S1.

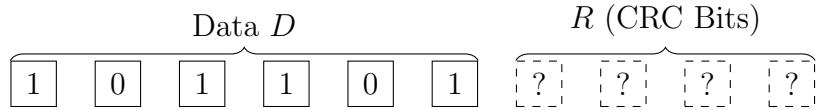
- **Switch S2 Table:**

MAC Address	Interface
A	Port 1
D	Port 3
C	Port 2

Note: S2 learns A from step 1 (incoming on Port 1). S2 learns D from step 2 (incoming on Port 3). S2 learns C from step 3 (incoming on Port 2).

2. (10 points) Cyclic Redundancy Check (CRC) and Burst Errors

Consider a data transmission scenario using a CRC generator polynomial $G(x) = x^4 + x + 1$. The data string to be transmitted is $D = 101101$.



- Analytically calculate the 4-bit CRC remainder R . Show the long division in binary modulo-2 arithmetic. What is the actual bit-string transmitted by sender?
- A "burst error" of length k is a contiguous sequence of bits in which the first and last bits are errors, and the intermediate bits may or may not be errors. Can this specific generator $G(x)$ detect **all** burst errors of length $k = 3$. Give reason?
- Suppose a burst error occurs during transmission such that the received bit string has the 3rd and 4th bits (counting from the left, 1-based index) inverted. Does the receiver accept or reject this frame?

Solution:

1. CRC Calculation:

- Generator $G(x) = x^4 + x + 1 \Rightarrow 10011$ (5 bits).
- Data $D = 101101$. Append 4 zeros: 1011010000 .
- Perform Binary Division (XOR):

$$\begin{array}{r}
 101101 \\
 \hline
 10011) 1011010000 \\
 \quad\quad\quad 10011 \\
 \hline
 \quad\quad\quad 01011 \\
 \quad\quad\quad 00000 \\
 \hline
 \quad\quad\quad 10110 \\
 \quad\quad\quad 10011 \\
 \hline
 \quad\quad\quad 01010 \\
 \quad\quad\quad 00000 \\
 \hline
 \quad\quad\quad 10100 \\
 \quad\quad\quad 10011 \\
 \hline
 \quad\quad\quad 01110 \\
 \quad\quad\quad 00000 \\
 \hline
 \quad\quad\quad 1110 \quad \leftarrow \text{Remainder (R)}
 \end{array}$$

The remainder R is **1110**.

The transmitted string is **1011011110**.

2. Burst Error Detection:

Yes. A generator polynomial of degree r (here $r = 4$) detects **all** burst errors of length $k \leq r$. Since the burst length $k = 3$ and $3 \leq 4$, this error is guaranteed to be detected.

3. Receiver Check:

The 3rd and 4th bits are inverted.

Original Transmitted: 10**1**1011110

Received Frame: 1000011110

Performing the division of 1000011110 by 10011:

$1000011110 \div 10011$ yields a remainder of **1011**.

Since the remainder is non-zero ($\neq 0000$), the receiver **rejects** the frame.