

Assignment 3

Formal Methods & Software Engineering



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Question 1 – Cinema Ticket Booking System

Basic Types and Report

[USER, SHOW, SEAT, BOOKING]

```
Report ::= ok | userExists | notRegistered | notLoggedIn  
          | pastShow | seatUnavailable | bookingUnknown
```

System State

TicketSystem

```
users      : P USER  
loggedIn   : P USER  
  
shows      : P SHOW  
pastShows  : P SHOW  
  
showSeats  : SHOW ↠ P SEAT  
  
booked     : (SHOW × SEAT) ↠ USER  
bookings   : BOOKING ↠ (USER × SHOW × P SEAT)
```

```
dom showSeats ⊆ shows  
pastShows ⊆ shows  
dom booked ⊆ shows × SEAT
```

Operation: DoRegisterUser

DoRegisterUser

```
Δ TicketSystem  
u?       : USER  
rep!    : Report
```

```
(  
  u? ∈ users ∧  
  users'   = users ∪ {u?} ∧  
  loggedIn' = loggedIn ∧  
  shows'   = shows ∧  
  pastShows' = pastShows ∧  
  showSeats' = showSeats ∧  
  booked'   = booked ∧  
  bookings' = bookings ∧  
  rep!      = ok  
)  
V  
(  
  u? ∈ users ∧  
  users'   = users ∧  
  loggedIn' = loggedIn ∧  
  shows'   = shows ∧  
  pastShows' = pastShows ∧  
  showSeats' = showSeats ∧  
  booked'   = booked ∧
```

```

bookings' = bookings ∧
rep!      = userExists
)

```

Operation: DoLogin

DoLogin

```

Δ TicketSystem
u?          : USER
rep!        : Report

```

```

(
  u? ∈ users ∧

    users'     = users ∧
    loggedIn'  = loggedIn ∪ {u?} ∧
    shows'     = shows ∧
    pastShows' = pastShows ∧
    showSeats' = showSeats ∧
    booked'    = booked ∧
    bookings'  = bookings ∧
    rep!       = ok
)
∨
(
  u? ∉ users ∧

    users'     = users ∧
    loggedIn'  = loggedIn ∧
    shows'     = shows ∧
    pastShows' = pastShows ∧
    showSeats' = showSeats ∧
    booked'    = booked ∧
    bookings'  = bookings ∧
    rep!       = notRegistered
)

```

Operation: DoBook

DoBook

```

Δ TicketSystem
user?       : USER
show?      : SHOW
seats?     : P SEAT
bid!       : BOOKING
rep!       : Report

```

```

(
  user? ∈ users ∧
  user? ∈ loggedIn ∧
  show? ∈ shows ∧
  show? ∉ pastShows ∧
  seats? ⊆ showSeats show? ∧
  seats? ≠ ∅ ∧

```

```

(∀ s : SEAT • s ∈ seats? ⇒ (show?, s) ∉ dom booked) ∧
bid! ∉ dom bookings ∧

users'      = users ∧
loggedIn'   = loggedIn ∧
shows'      = shows ∧
pastShows'  = pastShows ∧
showSeats'   = showSeats ∧

booked'     =
    booked ∪ { (show?, s) ↦ user? | s ∈ seats? } ∧

bookings'   =
    bookings ∪ { bid! ↦ (user?, show?, seats?) } ∧

rep! = ok
)
∨
(
    user? ∉ users ∧
    users'      = users ∧
    loggedIn'   = loggedIn ∧
    shows'      = shows ∧
    pastShows'  = pastShows ∧
    showSeats'   = showSeats ∧
    booked'     = booked ∧
    bookings'   = bookings ∧
    rep!        = notRegistered
)
∨
(
    user? ∈ users ∧ user? ∉ loggedIn ∧
    users'      = users ∧
    loggedIn'   = loggedIn ∧
    shows'      = shows ∧
    pastShows'  = pastShows ∧
    showSeats'   = showSeats ∧
    booked'     = booked ∧
    bookings'   = bookings ∧
    rep!        = notLoggedIn
)
∨
(
    show? ∈ pastShows ∧
    users'      = users ∧
    loggedIn'   = loggedIn ∧
    shows'      = shows ∧
    pastShows'  = pastShows ∧
    showSeats'   = showSeats ∧
    booked'     = booked ∧
    bookings'   = bookings ∧
    rep!        = pastShow
)
∨
(

```

```

 $\exists s : SEAT \bullet s \in seats? \wedge (show?, s) \in \text{dom booked} \wedge$ 
users' = users  $\wedge$ 
loggedIn' = loggedIn  $\wedge$ 
shows' = shows  $\wedge$ 
pastShows' = pastShows  $\wedge$ 
showSeats' = showSeats  $\wedge$ 
booked' = booked  $\wedge$ 
bookings' = bookings  $\wedge$ 
rep! = seatUnavailable
)

```

Operation: DoCancel

DoCancel

Δ TicketSystem

| | | |
|-------|---|---------|
| user? | : | USER |
| bid? | : | BOOKING |
| rep! | : | Report |

```

(
  user?  $\in$  users  $\wedge$ 
  user?  $\in$  loggedIn  $\wedge$ 
  bid?  $\in$  dom bookings  $\wedge$ 

  LET (u, sh, Ss) == bookings bid? IN
    u = user?  $\wedge$ 
    sh  $\notin$  pastShows  $\wedge$ 

    users' = users  $\wedge$ 
    loggedIn' = loggedIn  $\wedge$ 
    shows' = shows  $\wedge$ 
    pastShows' = pastShows  $\wedge$ 
    showSeats' = showSeats  $\wedge$ 

    booked' =
      { (sh2,s2)  $\mapsto$  u2 | (sh2,s2)  $\mapsto$  u2  $\in$  booked  $\wedge$   $\neg$ (sh2 = sh  $\wedge$  s2  $\in$  Ss) }  $\wedge$ 

    bookings' = bookings  $\wedge$ 

    rep! = ok
)
V
(
  (user?  $\notin$  users  $\vee$  user?  $\notin$  loggedIn  $\vee$  bid?  $\notin$  dom bookings)  $\wedge$ 
  users' = users  $\wedge$ 
  loggedIn' = loggedIn  $\wedge$ 
  shows' = shows  $\wedge$ 
  pastShows' = pastShows  $\wedge$ 
  showSeats' = showSeats  $\wedge$ 
  booked' = booked  $\wedge$ 
  bookings' = bookings  $\wedge$ 
  rep! = bookingUnknown
)

```

Operation: DoResetShow

DoResetShow

Δ TicketSystem

show? : SHOW
rep! : Report

```

(
  show? ∈ shows ∧
  booked' = { (sh,s) ↦ u | (sh,s) ↢ u ∈ booked ∧ sh ≠ show? } ∧
  bookings' = bookings ∧
  users' = users ∧
  loggedIn' = loggedIn ∧
  shows' = shows ∧
  pastShows' = pastShows ∧
  showSeats' = showSeats ∧
  rep! = ok
)
∨
(
  show? ∉ shows ∧
  users' = users ∧
  loggedIn' = loggedIn ∧
  shows' = shows ∧
  pastShows' = pastShows ∧
  showSeats' = showSeats ∧
  booked' = booked ∧
  bookings' = bookings ∧
  rep! = pastShow
)

```

Operation: ViewBookingHistory (admin view)

ViewBookingHistory

Ξ TicketSystem

bh! : BOOKING ↨ (USER × SHOW × P SEAT)

bh! = bookings

Question 2 – University Course Registration System

Basic Types

[STUDENT, COURSE]

System State

CourseReg

students : P STUDENT

```

courses      : P COURSE
completed    : STUDENT ↪ P COURSE
prereq       : COURSE ↪ P COURSE
capacity     : COURSE ↪ N
reg          : STUDENT ↪ P COURSE

```

```

dom completed ⊆ students
dom reg ⊆ students
dom prereq ⊆ courses
dom capacity ⊆ courses

```

Operation: AddStudent

AddStudent

```

Δ CourseReg
s?           : STUDENT

```

```

(
  s? ∈ students ∧

  students' = students ∪ {s?} ∧
  courses' = courses ∧
  completed' = completed ∧
  prereq' = prereq ∧
  capacity' = capacity ∧
  reg' = reg
)
∨
(
  s? ∈ students ∧

  students' = students ∧
  courses' = courses ∧
  completed' = completed ∧
  prereq' = prereq ∧
  capacity' = capacity ∧
  reg' = reg
)

```

Operation: AddCourse

AddCourse

```

Δ CourseReg
c?           : COURSE
prereq?     : P COURSE
cap?         : N

```

```

(
  c? ∈ courses ∧
  prereq? ⊆ courses ∧

  students' = students ∧
  courses' = courses ∪ {c?} ∧

```

```

completed' = completed ∧
reg'      = reg ∧

prereq'    = prereq ∪ { c? ↦ prereq? } ∧
capacity'  = capacity ∪ { c? ↦ cap? } ∧
)
∨
(
c? ∈ courses ∧

students'  = students ∧
courses'   = courses ∧
completed' = completed ∧
prereq'    = prereq ∧
capacity'  = capacity ∧
reg'       = reg
)

```

Operation: RegisterCourse

RegisterCourse

Δ CourseReg

s? : STUDENT
c? : COURSE

```

(
s? ∈ students ∧
c? ∈ courses ∧
s? ∈ dom reg ∧
c? ∉ reg s? ∧

c? ∈ dom prereq ∧
c? ∈ dom capacity ∧
s? ∈ dom completed ∧
prereq c? ⊆ completed s? ∧

# { st : STUDENT | st ∈ dom reg ∧ c? ∈ reg st } < capacity c? ∧

students'  = students ∧
courses'   = courses ∧
completed' = completed ∧
prereq'    = prereq ∧
capacity'  = capacity ∧

reg'       = reg ⊕ { s? ↦ (reg s? ∪ {c?}) }
)
∨
(
¬(
s? ∈ students ∧
c? ∈ courses ∧
s? ∈ dom reg ∧
c? ∉ reg s? ∧
c? ∈ dom prereq ∧
c? ∈ dom capacity ∧

```

```

s? ∈ dom completed ∧
prereq c? ⊑ completed s? ∧
# { st : STUDENT | st ∈ dom reg ∧ c? ∈ reg st } < capacity c?
) ∧

students' = students ∧
courses' = courses ∧
completed' = completed ∧
prereq' = prereq ∧
capacity' = capacity ∧
reg' = reg
)

```

Operation: DropCourse

DropCourse

Δ CourseReg

s? : STUDENT
c? : COURSE

```

(
s? ∈ dom reg ∧
c? ∈ reg s? ∧

let Cs == reg s? \ {c?} in
  students' = students ∧
  courses' = courses ∧
  completed' = completed ∧
  prereq' = prereq ∧
  capacity' = capacity ∧

  reg' =
    (if Cs = ∅
      then { st ↦ R | st ↦ R ∈ reg ∧ st ≠ s? }
      else reg ⊕{ s? ↦ Cs }
    )
)
∨
(
¬(s? ∈ dom reg ∧ c? ∈ reg s?) ∧

students' = students ∧
courses' = courses ∧
completed' = completed ∧
prereq' = prereq ∧
capacity' = capacity ∧
reg' = reg
)
```

Question 3 – Reasoning About Loops

Part 1 – Sum = 1 + 2 + ... + n

Program:

```

sum := 0;
i    := 1;

while i ≤ n do
    sum := sum + i;
    i    := i + 1
od
Loop invariant:
```

$$I: \text{sum} = \sum_{k=1}^{i-1} k \text{ and } 1 \leq i \leq n + 1.$$

Sketch: initially sum = 0 and i = 1, so the invariant holds.

Each iteration adds i to sum and then increases i, preserving $\text{sum} = \sum_{k=1}^{i-1} k$.

When the loop exits, i = n + 1 and therefore $\text{sum} = \sum_{k=1}^n k$.

Part 2 – Relationship Between i and j

Example program:

```

i := 0;
j := 0;

while i < n do
    i := i + 1;
    j := j + 2*i - 1
od
```

Tracing a few steps gives:

$$i = 1, 2, 3, 4 \rightarrow j = 1, 4, 9, 16 = 1^2, 2^2, 3^2, 4^2.$$

So when the loop terminates, i = n and j = n².

Loop invariant:

$$I: j = i^2 \text{ and } 0 \leq i \leq n.$$

Sketch: initially i = 0 and j = 0, so j = i². Each iteration increases i to i+1 and adds 2*i - 1 (with the new i) to j, which turns i² into (i+1)². When the loop exits, i = n and j = n².