

Assignment 3

Formal Methods & Software Engineering



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Question 1 – Cinema Ticket Booking System

Basic Types and Report

[USER, SHOW, SEAT, BOOKING]

Report ::= ok | userExists | notRegistered | notLoggedIn
 | pastShow | seatUnavailable | bookingUnknown

System State

TicketSystem

users	: P USER
loggedIn	: P USER
shows	: P SHOW
pastShows	: P SHOW
showSeats	: SHOW \mapsto P SEAT
booked	: (SHOW \times SEAT) \mapsto USER
bookings	: BOOKING \mapsto (USER \times SHOW \times P SEAT)

dom showSeats \subseteq shows
pastShows \subseteq shows
dom booked \subseteq shows \times SEAT

Operation: DoRegisterUser

DoRegisterUser

Δ TicketSystem

u?	: USER
rep!	: Report

(
 u? \notin users \wedge
 users' = users \cup {u?} \wedge
 loggedIn' = loggedIn \wedge
 shows' = shows \wedge
 pastShows' = pastShows \wedge
 showSeats' = showSeats \wedge
 booked' = booked \wedge
 bookings' = bookings \wedge
 rep! = ok
)
 \vee
(
 u? \in users \wedge
 users' = users \wedge
 loggedIn' = loggedIn \wedge
 shows' = shows \wedge
 pastShows' = pastShows \wedge
 showSeats' = showSeats \wedge
 booked' = booked \wedge

```

    bookings' = bookings  $\wedge$ 
    rep!       = userExists
)

```

Operation: DoLogin

DoLogin

```

 $\Delta$  TicketSystem
u?      : USER
rep!     : Report

```

```

(
    u?  $\in$  users  $\wedge$ 

    users'      = users  $\wedge$ 
    loggedIn'    = loggedIn  $\cup$  {u?}  $\wedge$ 
    shows'       = shows  $\wedge$ 
    pastShows'   = pastShows  $\wedge$ 
    showSeats'   = showSeats  $\wedge$ 
    booked'      = booked  $\wedge$ 
    bookings'    = bookings  $\wedge$ 
    rep!         = ok
)
 $\vee$ 
(
    u?  $\notin$  users  $\wedge$ 

    users'      = users  $\wedge$ 
    loggedIn'    = loggedIn  $\wedge$ 
    shows'       = shows  $\wedge$ 
    pastShows'   = pastShows  $\wedge$ 
    showSeats'   = showSeats  $\wedge$ 
    booked'      = booked  $\wedge$ 
    bookings'    = bookings  $\wedge$ 
    rep!         = notRegistered
)

```

Operation: DoBook

DoBook

```

 $\Delta$  TicketSystem
user?      : USER
show?      : SHOW
seats?     : P SEAT
bid!       : BOOKING
rep!       : Report

```

```

(
    user?  $\in$  users  $\wedge$ 
    user?  $\in$  loggedIn  $\wedge$ 
    show?  $\in$  shows  $\wedge$ 
    show?  $\notin$  pastShows  $\wedge$ 
    seats?  $\subseteq$  showSeats show?  $\wedge$ 
    seats?  $\neq \emptyset$   $\wedge$ 

```

```

(∀ s : SEAT • s ∈ seats? ⇒ (show?, s) ∉ dom booked) ∧
bid! ∉ dom bookings ∧

users'      = users ∧
loggedIn'   = loggedIn ∧
shows'      = shows ∧
pastShows'  = pastShows ∧
showSeats'  = showSeats ∧

booked' =
  booked ∪ { (show?, s) ↦ user? | s ∈ seats? } ∧

bookings' =
  bookings ∪ { bid! ↦ (user?, show?, seats?) } ∧

rep! = ok
)
∨
(
  user? ∉ users ∧
  users'      = users ∧
  loggedIn'   = loggedIn ∧
  shows'      = shows ∧
  pastShows'  = pastShows ∧
  showSeats'  = showSeats ∧
  booked'     = booked ∧
  bookings'   = bookings ∧
  rep!        = notRegistered
)
∨
(
  user? ∈ users ∧ user? ∉ loggedIn ∧
  users'      = users ∧
  loggedIn'   = loggedIn ∧
  shows'      = shows ∧
  pastShows'  = pastShows ∧
  showSeats'  = showSeats ∧
  booked'     = booked ∧
  bookings'   = bookings ∧
  rep!        = notLoggedIn
)
∨
(
  show? ∈ pastShows ∧
  users'      = users ∧
  loggedIn'   = loggedIn ∧
  shows'      = shows ∧
  pastShows'  = pastShows ∧
  showSeats'  = showSeats ∧
  booked'     = booked ∧
  bookings'   = bookings ∧
  rep!        = pastShow
)
∨
(

```

```

 $\exists s : \text{SEAT} \bullet s \in \text{seats?} \wedge (\text{show?}, s) \in \text{dom booked} \wedge$ 
users'      = users  $\wedge$ 
loggedIn'   = loggedIn  $\wedge$ 
shows'      = shows  $\wedge$ 
pastShows'  = pastShows  $\wedge$ 
showSeats'  = showSeats  $\wedge$ 
booked'     = booked  $\wedge$ 
bookings'   = bookings  $\wedge$ 
rep!        = seatUnavailable
)

```

Operation: DoCancel

DoCancel

Δ TicketSystem

user? : USER
bid? : BOOKING
rep! : Report

```

(
  user?  $\in$  users  $\wedge$ 
  user?  $\in$  loggedIn  $\wedge$ 
  bid?  $\in$  dom bookings  $\wedge$ 

  LET (u, sh, Ss) == bookings bid? IN
    u = user?  $\wedge$ 
    sh  $\notin$  pastShows  $\wedge$ 

    users'      = users  $\wedge$ 
    loggedIn'   = loggedIn  $\wedge$ 
    shows'      = shows  $\wedge$ 
    pastShows'  = pastShows  $\wedge$ 
    showSeats'  = showSeats  $\wedge$ 

    booked' =
      { (sh2,s2)  $\mapsto$  u2 | (sh2,s2)  $\mapsto$  u2  $\in$  booked  $\wedge$   $\neg$ (sh2 = sh  $\wedge$  s2  $\in$  Ss) }  $\wedge$ 

    bookings' = bookings  $\wedge$ 

    rep! = ok
)
 $\vee$ 
(
  (user?  $\notin$  users  $\vee$  user?  $\notin$  loggedIn  $\vee$  bid?  $\notin$  dom bookings)  $\wedge$ 
  users'      = users  $\wedge$ 
  loggedIn'   = loggedIn  $\wedge$ 
  shows'      = shows  $\wedge$ 
  pastShows'  = pastShows  $\wedge$ 
  showSeats'  = showSeats  $\wedge$ 
  booked'     = booked  $\wedge$ 
  bookings'   = bookings  $\wedge$ 
  rep!        = bookingUnknown
)

```

Operation: DoResetShow

DoResetShow

Δ TicketSystem

show? : SHOW

rep! : Report

(
 show? \in shows \wedge
 booked' = { (sh,s) \mapsto u | (sh,s) \mapsto u \in booked \wedge sh \neq show? } \wedge
 bookings' = bookings \wedge
 users' = users \wedge
 loggedIn' = loggedIn \wedge
 shows' = shows \wedge
 pastShows' = pastShows \wedge
 showSeats' = showSeats \wedge
 rep! = ok
)
 \vee
(
 show? \notin shows \wedge
 users' = users \wedge
 loggedIn' = loggedIn \wedge
 shows' = shows \wedge
 pastShows' = pastShows \wedge
 showSeats' = showSeats \wedge
 booked' = booked \wedge
 bookings' = bookings \wedge
 rep! = pastShow
)

Operation: ViewBookingHistory (admin view)

ViewBookingHistory

\exists TicketSystem

bh! : BOOKING \mapsto (USER \times SHOW \times P SEAT)

bh! = bookings

Question 2 – University Course Registration System

Basic Types

[STUDENT, COURSE]

System State

CourseReg

students : P STUDENT

$\text{courses} : P \text{ COURSE}$
 $\text{completed} : \text{STUDENT} \mapsto P \text{ COURSE}$
 $\text{prereq} : \text{COURSE} \mapsto P \text{ COURSE}$
 $\text{capacity} : \text{COURSE} \mapsto \mathbb{N}$
 $\text{reg} : \text{STUDENT} \mapsto P \text{ COURSE}$

$\text{dom completed} \subseteq \text{students}$
 $\text{dom reg} \subseteq \text{students}$
 $\text{dom prereq} \subseteq \text{courses}$
 $\text{dom capacity} \subseteq \text{courses}$

Operation: AddStudent

AddStudent

$\Delta \text{ CourseReg}$

$s? : \text{STUDENT}$

$($
 $s? \notin \text{students} \wedge$

 $\text{students}' = \text{students} \cup \{s?\} \wedge$
 $\text{courses}' = \text{courses} \wedge$
 $\text{completed}' = \text{completed} \wedge$
 $\text{prereq}' = \text{prereq} \wedge$
 $\text{capacity}' = \text{capacity} \wedge$
 $\text{reg}' = \text{reg}$
 $)$
 \vee
 $($
 $s? \in \text{students} \wedge$

 $\text{students}' = \text{students} \wedge$
 $\text{courses}' = \text{courses} \wedge$
 $\text{completed}' = \text{completed} \wedge$
 $\text{prereq}' = \text{prereq} \wedge$
 $\text{capacity}' = \text{capacity} \wedge$
 $\text{reg}' = \text{reg}$
 $)$

Operation: AddCourse

AddCourse

$\Delta \text{ CourseReg}$

$c? : \text{COURSE}$
 $\text{prereq}? : P \text{ COURSE}$
 $\text{cap}? : \mathbb{N}$

$($
 $c? \notin \text{courses} \wedge$
 $\text{prereq}? \subseteq \text{courses} \wedge$

 $\text{students}' = \text{students} \wedge$
 $\text{courses}' = \text{courses} \cup \{c?\} \wedge$

```

    completed' = completed  $\wedge$ 
    reg'       = reg  $\wedge$ 

    prereq'    = prereq  $\cup$  {  $c? \mapsto \text{prereq?}$  }  $\wedge$ 
    capacity'  = capacity  $\cup$  {  $c? \mapsto \text{cap?}$  }
  )
 $\vee$ 
  (
     $c? \in \text{courses} \wedge$ 

    students'  = students  $\wedge$ 
    courses'   = courses  $\wedge$ 
    completed' = completed  $\wedge$ 
    prereq'    = prereq  $\wedge$ 
    capacity'  = capacity  $\wedge$ 
    reg'       = reg
  )

```

Operation: RegisterCourse

RegisterCourse

```

 $\Delta$  CourseReg
s?      : STUDENT
c?      : COURSE

```

```

  (
     $s? \in \text{students} \wedge$ 
     $c? \in \text{courses} \wedge$ 
     $s? \in \text{dom reg} \wedge$ 
     $c? \notin \text{reg } s? \wedge$ 

     $c? \in \text{dom prereq} \wedge$ 
     $c? \in \text{dom capacity} \wedge$ 
     $s? \in \text{dom completed} \wedge$ 
     $\text{prereq } c? \subseteq \text{completed } s? \wedge$ 

    # { st : STUDENT | st  $\in$  dom reg  $\wedge$   $c? \in \text{reg st}$  } < capacity  $c? \wedge$ 

    students' = students  $\wedge$ 
    courses'   = courses  $\wedge$ 
    completed' = completed  $\wedge$ 
    prereq'    = prereq  $\wedge$ 
    capacity'  = capacity  $\wedge$ 

    reg'       = reg  $\oplus$  {  $s? \mapsto (\text{reg } s? \cup \{c?\})$  }
  )
 $\vee$ 
  (
     $\neg$ (
       $s? \in \text{students} \wedge$ 
       $c? \in \text{courses} \wedge$ 
       $s? \in \text{dom reg} \wedge$ 
       $c? \notin \text{reg } s? \wedge$ 
       $c? \in \text{dom prereq} \wedge$ 
       $c? \in \text{dom capacity} \wedge$ 

```



```

    s? ∈ dom completed ∧
    prereq c? ⊆ completed s? ∧
    # { st : STUDENT | st ∈ dom reg ∧ c? ∈ reg st } < capacity c?
  ) ∧

  students' = students ∧
  courses' = courses ∧
  completed' = completed ∧
  prereq' = prereq ∧
  capacity' = capacity ∧
  reg' = reg
)

```

Operation: DropCourse

DropCourse

```

Δ CourseReg
s?      : STUDENT
c?      : COURSE

```

```

(
  s? ∈ dom reg ∧
  c? ∈ reg s? ∧

  let Cs == reg s? \ {c?} in
    students' = students ∧
    courses' = courses ∧
    completed' = completed ∧
    prereq' = prereq ∧
    capacity' = capacity ∧

    reg' =
      (if Cs = ∅
       then { st ↦ R | st ↦ R ∈ reg ∧ st ≠ s? }
       else reg ⊕ { s? ↦ Cs }
      )
)
∨
(
  ¬(s? ∈ dom reg ∧ c? ∈ reg s?) ∧

  students' = students ∧
  courses' = courses ∧
  completed' = completed ∧
  prereq' = prereq ∧
  capacity' = capacity ∧
  reg' = reg
)

```

Question 3 – Reasoning About Loops

Part 1 – Sum = 1 + 2 + ... + n

Program:

```

sum := 0;
i   := 1;

while i ≤ n do
    sum := sum + i;
    i   := i + 1
od

```

Loop invariant:

I: $\text{sum} = \sum_{k=1}^{i-1} k$ and $1 \leq i \leq n + 1$.

Sketch: initially $\text{sum} = 0$ and $i = 1$, so the invariant holds.

Each iteration adds i to sum and then increases i , preserving $\text{sum} = \sum_{k=1}^{i-1} k$.

When the loop exits, $i = n + 1$ and therefore $\text{sum} = \sum_{k=1}^n k$.

Part 2 – Relationship Between i and j

Example program:

```

i := 0;
j := 0;

while i < n do
    i := i + 1;
    j := j + 2*i - 1
od

```

Tracing a few steps gives:

$i = 1, 2, 3, 4 \rightarrow j = 1, 4, 9, 16 = 1^2, 2^2, 3^2, 4^2$.

So when the loop terminates, $i = n$ and $j = n^2$.

Loop invariant:

I: $j = i^2$ and $0 \leq i \leq n$.

Sketch: initially $i = 0$ and $j = 0$, so $j = i^2$. Each iteration increases i to $i+1$ and adds $2*i - 1$ (with the new i) to j , which turns i^2 into $(i+1)^2$. When the loop exits, $i = n$ and $j = n^2$.