

GRAPH THEORY

Chapter 10

INTRODUCTION

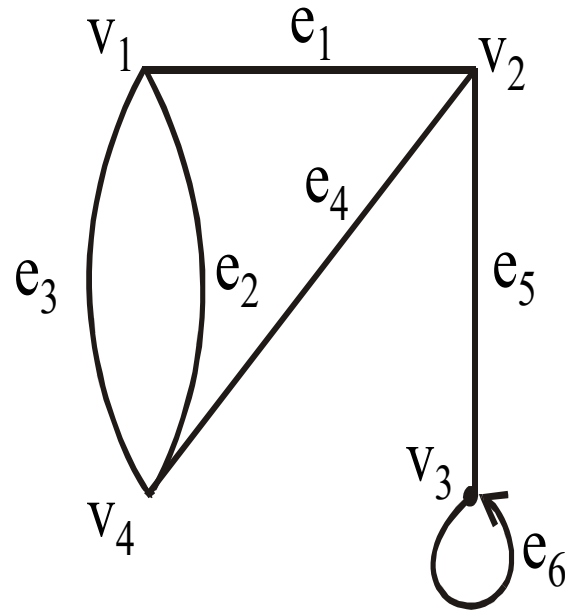
- ▶ **Graph theory** plays an important role in several areas of **computer science** such as:
 - ▶ **Switching Theory and Logical Design**
 - ▶ **Artificial Intelligence**
 - ▶ **Formal Languages**
 - ▶ **Computer Graphics**
 - ▶ **Operating Systems**
 - ▶ **Compiler Writing**
 - ▶ **Information Organization and Retrieval.**

GRAPH

- ▶ A **graph** is a non-empty set of points called **vertices** and A set of line segments joining pairs of **vertices** called **edges**.
- ▶ Formally, a **graph** $G = (V, E)$ consists of two finite sets:
 - ▶ A set $V = V(G)$ of **vertices** (or points or nodes)
 - ▶ A set $E = E(G)$ of **edges**;

where each **edge** corresponds to a pair of **vertices**.

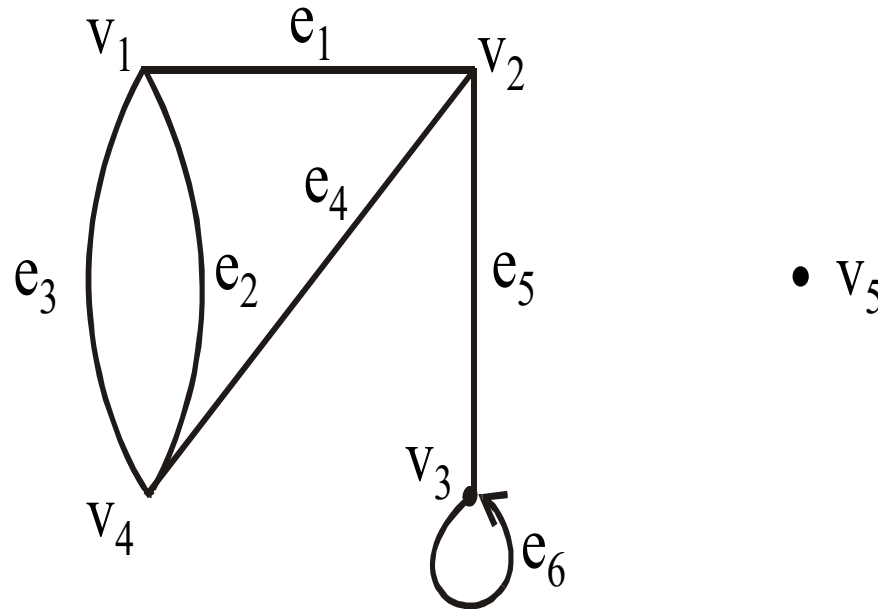
EXAMPLE



- ▶ We have **five vertices** labeled by v_1, v_2, v_3, v_4, v_5 .
- ▶ We have **edges** $e_1, e_2, e_3, e_4, e_5, e_6$.

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- ▶ e_1 edge is for vertices v_1 and v_2 .
 - ▶ e_2 and e_3 end points are v_1 and v_4 .
 - ▶ e_4 has end points v_2 and v_4 .
 - ▶ e_5 has end points v_2 and v_3 .
 - ▶ e_6 is a loop .

SOME TERMINOLOGY



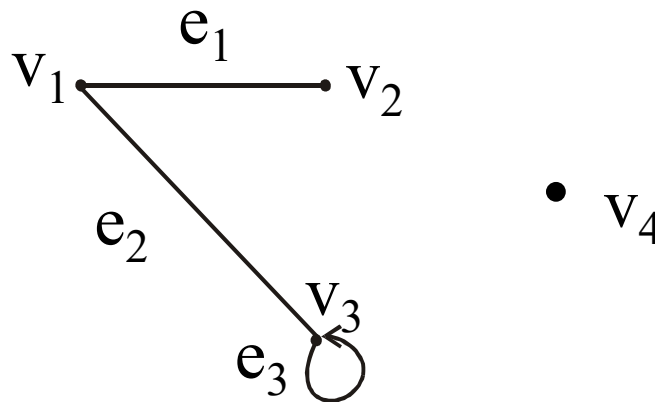
- ▶ An **edge** connects either one or two **vertices** called its **endpoints** (edge e_1 connects vertices v_1 and v_2 described as $\{v_1, v_2\}$ i.e v_1 and v_2 are the endpoints of an edge e_1).

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- ▶ An **edge** with just **one endpoint** is called a **loop**. Thus a **loop** is an **edge** that connects a **vertex** to itself (e.g., edge e_6).
 - ▶ Two **vertices** that are connected by an **edge** are called **adjacent**; and a **vertex** that is an **endpoint** of a **loop** is said to be **adjacent** to itself.
 - ▶ An **edge** is said to be **incident** on each of its **endpoints** (i.e. e_1 is incident on v_1 and v_2).

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- ▶ A **vertex** on which no **edges** are incident is called **isolated** (e.g., v_5)
 - ▶ Two **distinct edges** with the same set of **end points** are said to be **parallel** (i.e. e_2 & e_3).

EXAMPLE

- Define the following **graph** formally by specifying its **vertex** set, its edge set, and a table giving the **edge endpoint function**.



SOLUTION

Vertex Set = $\{v_1, v_2, v_3, v_4\}$

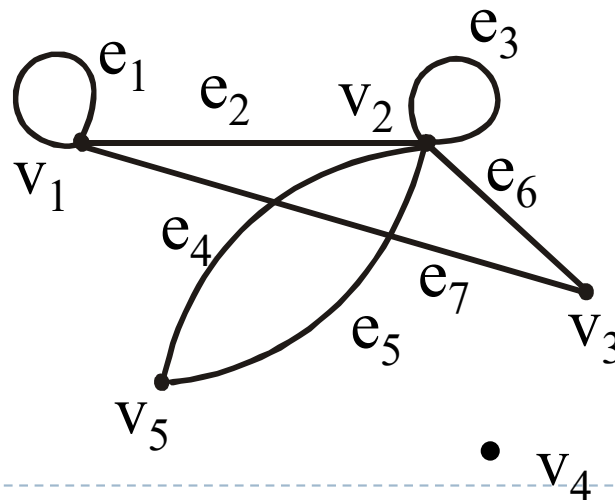
Edge Set = $\{e_1, e_2, e_3\}$

Edge - endpoint function is:

Edge	Endpoint
e_1	$\{v_1, v_2\}$
e_2	$\{v_1, v_3\}$
e_3	$\{v_3\}$

EXAMPLE

- ▶ For the **graph** shown below:
 - ▶ Find all **edges** that are **incident** on v_1 ;
 - ▶ Find all **vertices** that are **adjacent** to v_3 ;
 - ▶ Find all **loops**;
 - ▶ Find all **parallel** edges;
 - ▶ Find all **isolated** vertices;



SOLUTION

- ▶ Find all **edges** that are **incident** on v_1 ?
 - ▶ v_1 is **incident** with edges e_1 , e_2 and e_7
- ▶ Find all **vertices** that are **adjacent** to v_3 ?
 - ▶ vertices **adjacent** to v_3 are v_1 and v_2
- ▶ Find all **loops**?
 - ▶ **Loops** are e_1 and e_3

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- ▶ Find all **parallel** edges?
 - ▶ Only edges e_4 and e_5 are **parallel**
 - ▶ Find all **isolated** vertices?
 - ▶ The only **isolated** vertex is v_4 in this Graph.

EXERCISE

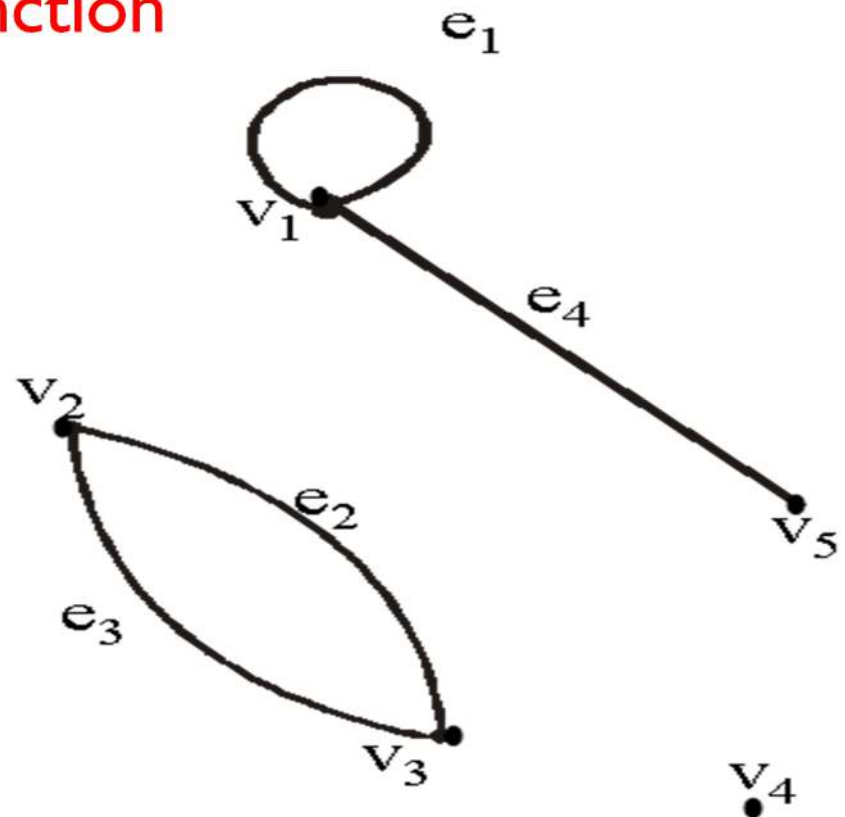
- Draw picture of **Graph H** having **vertex** set $\{v_1, v_2, v_3, v_4, v_5\}$ and **edge** set $\{e_1, e_2, e_3, e_4\}$ with **edge endpoint function**:

Edge	Endpoint
e_1	$\{v_1\}$
e_2	$\{v_2, v_3\}$
e_3	$\{v_2, v_3\}$
e_4	$\{v_1, v_5\}$

EXERCISE

- Given $V(H) = \{v_1, v_2, v_3, v_4, v_5\}$ and
 $E(H) = \{e_1, e_2, e_3, e_4\}$
with edge **endpoint function**

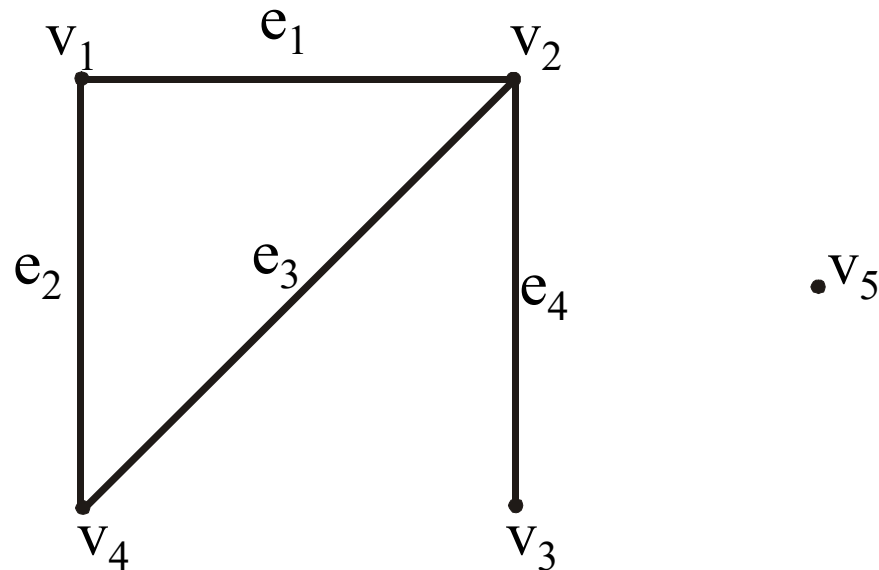
Edge	Endpoint
e_1	$\{v_1\}$
e_2	$\{v_2, v_3\}$
e_3	$\{v_2, v_3\}$
e_4	$\{v_1, v_5\}$



SIMPLE GRAPH

- ▶ A simple graph is a graph that does not have any loop or parallel edges.

- ▶ Example:



- ▶ It is a simple graph H

$$V(H) = \{v_1, v_2, v_3, v_4, v_5\} \quad \& \quad E(H) = \{e_1, e_2, e_3, e_4\}$$

EXERCISE

- ▶ Draw all **simple graphs** with the four vertices $\{u, v, w, x\}$ and two edges, one of which is $\{u, v\}$.

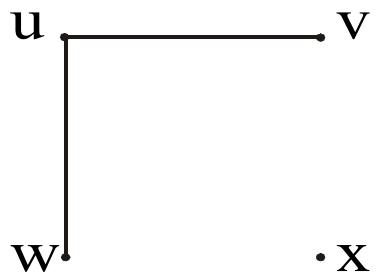
- ▶ **Solution:**

We are given four vertices $\{u, v, w, x\}$ and one edge is $\{u, v\}$.

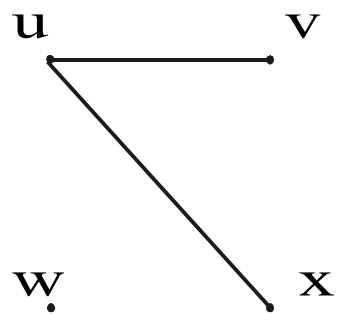
Since we are interested in **simple graph** so we cannot take $\{u, v\}$ again.

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- ▶ There are $C(4,2) = 6$ ways of choosing two vertices from 4 vertices.
 - ▶ These edges may be listed as:
 $\{u, v\}, \{u, w\}, \{u, x\}, \{v, w\}, \{v, x\}, \{w, x\}$
 - ▶ One edge of the graph is specified to be $\{u, v\}$, so any of the remaining **five** from this list may be chosen to be the second edge.
 - ▶ This required **graphs** are:

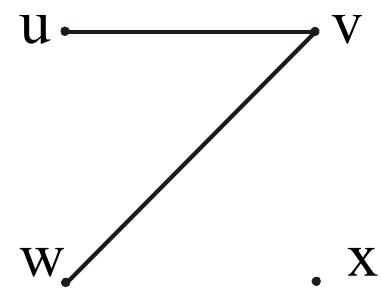
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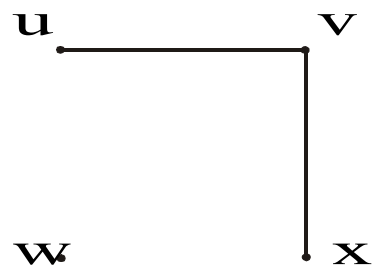
2.



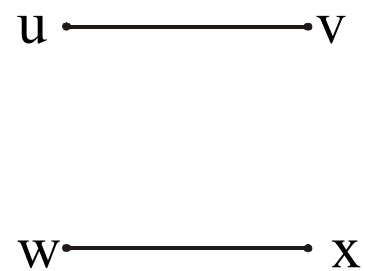
3.



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5.

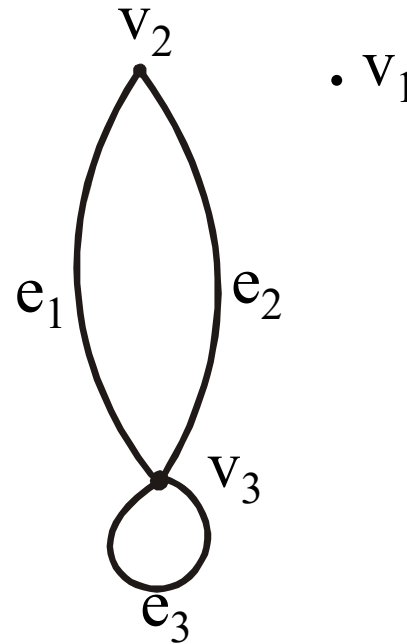


DEGREE OF A VERTEX

- ▶ Let G be a graph and v a vertex of G . The degree of v , denoted $\deg(v)$, equals the number of edges that are incident on v , with an edge that is a loop counted twice.
- ▶ The total degree of G is the sum of the degrees of all the vertices of G .
- ▶ The degree of a loop is counted twice.

EXAMPLE

- ▶ For the graph shown



- ▶ $\deg(v_1) = 0$, since v_1 is isolated vertex.
- ▶ $\deg(v_2) = 2$, since v_2 is incident on e_1 and e_2 .
- ▶ $\deg(v_3) = 4$, since v_3 is incident on e_1, e_2 and the loop e_3 .
- ▶ Total degree of $G = \deg(v_1) + \deg(v_2) + \deg(v_3)$
 $= 0 + 2 + 4 = 6$

NOTE

- ▶ The **total degree** of **G**, which is **6**, equals **twice** the **number** of **edges** of **G**, which is **3**.
- ▶ This is always the case the **total degree** of **graph** is always **twice** the **number** of **edges** in that **graph**. This is actually the theorem called Handshaking Theorem.

THE HANDSHAKING THEOREM

- ▶ If G is any graph, then the sum of the degrees of all the vertices of G equals twice the number of edges of G .
- ▶ Specifically, if the vertices of G are v_1, v_2, \dots, v_n , where n is a positive integer, then
- ▶ The total degree of
$$G = \deg(v_1) + \deg(v_2) + \dots + \deg(v_n)$$
$$= 2 \cdot (\text{the number of edges of } G)$$

COROLLARY

- ▶ The total degree of G is an even number

EXERCISE

- ▶ Draw a **graph** with the specified properties or explain why no such **graph** exists.

- (i) Graph with **four vertices** of **degrees 1, 2, 3 and 3**
- (ii) Graph with **four vertices** of **degrees 1, 2, 3 and 4**
- (iii) **Simple graph** with **four vertices** of **degrees 1, 2, 3 and 4**.

(i) Graph with four vertices of degrees 1, 2, 3 and 3

$$\begin{aligned}\text{Total degree of graph} &= 1 + 2 + 3 + 3 \\ &= 9 \text{ an odd integer}\end{aligned}$$

Hence by Hand-Shaking Theorem, first graph is not possible.

(ii) Graph with four vertices of degrees 1, 2, 3 and 4

$$\begin{aligned}\text{Total degree of graph} &= 1 + 2 + 3 + 4 \\ &= 10 \text{ an even integer}\end{aligned}$$

► The vertices a, b, c, d have degrees 1, 2, 3, and 4 respectively (i.e. graph exists).

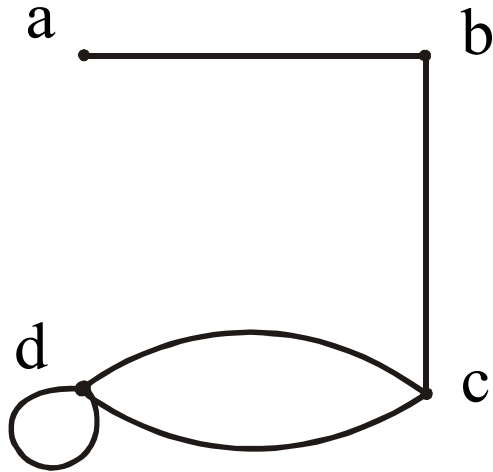
$$\deg(a) = 1$$

$$\deg(b) = 2$$

$$\deg(c) = 3$$

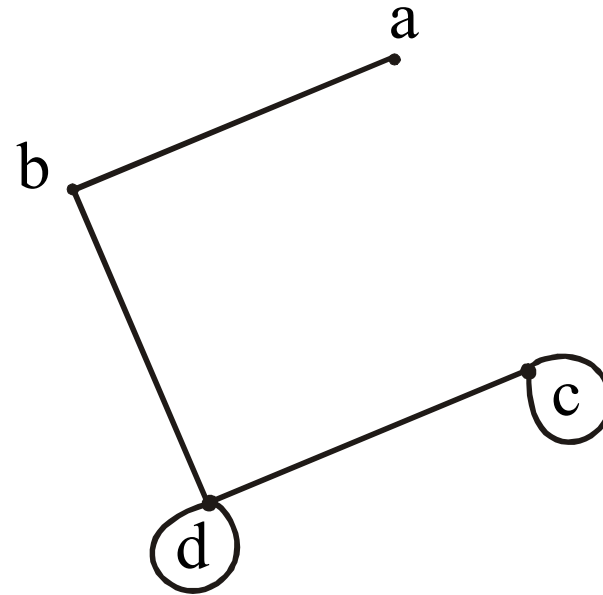
$$\deg(d) = 4$$

1.



$$\begin{aligned} \deg(a) &= 1 & \deg(b) &= 2 \\ \deg(c) &= 3 & \deg(d) &= 4 \end{aligned}$$

2.



$$\begin{aligned} \deg(a) &= 1 & \deg(b) &= 2 \\ \deg(c) &= 3 & \deg(d) &= 4 \end{aligned}$$

(iii) Simple graph with four vertices of degrees 1, 2, 3 and 4.

- ▶ Suppose there was a simple graph with four vertices of degrees 1, 2, 3, and 4. Then the vertex of degree 4 would have to be connected by edges to four distinct vertices other than itself because of the assumption that the graph is simple (and hence has no loop or parallel edges.) This contradicts the assumption that the graph has four vertices in total.
- ▶ Hence there is no simple graph with four vertices of degrees 1, 2, 3, and 4, so simple graph is not possible in this case.

EXERCISE

- ▶ Suppose a graph has vertices of degrees 1, 1, 4, 4 and 6. How many edges does the graph have?

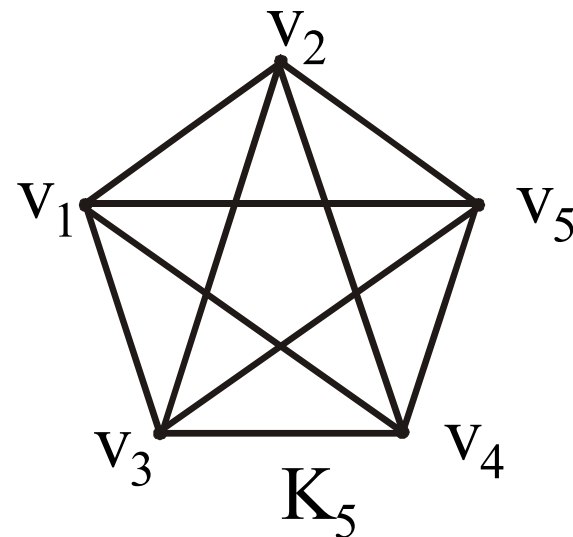
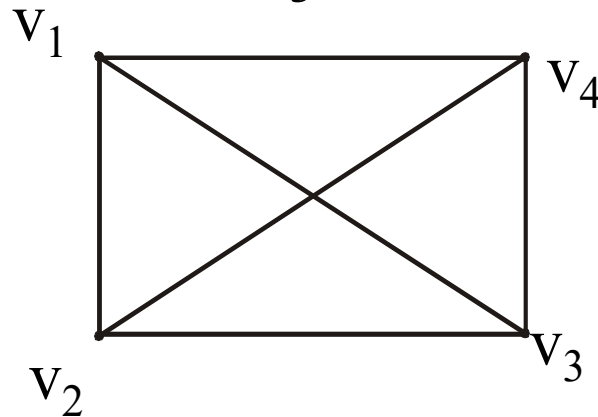
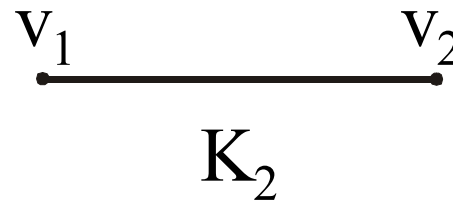
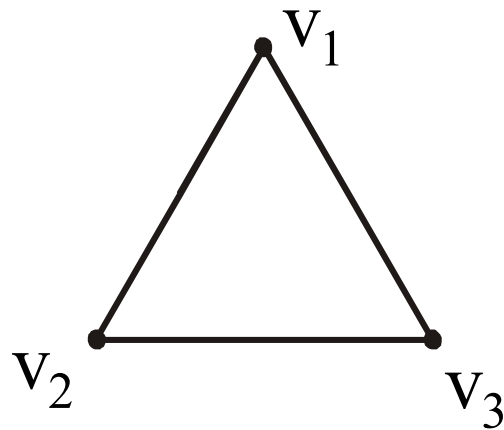
- ▶ SOLUTION:

$$\begin{aligned}\text{The total degree of graph} &= 1 + 1 + 4 + 4 + 6 \\ &= 16\end{aligned}$$

$$\Rightarrow \text{Number of edges of graph} = \frac{16}{2} = 8$$

COMPLETE GRAPH

- ▶ A **complete graph** on “ n ” vertices is a **simple graph** in which each **vertex** is **connected** to every other **vertex** and is denoted by K_n (K_n means that there are n vertices).



K_4

K_5

REGULAR GRAPH

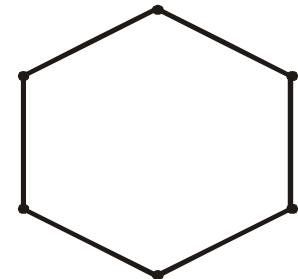
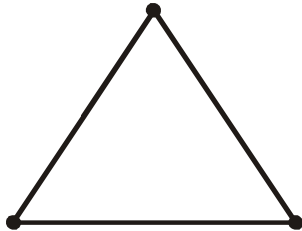
- ▶ A graph G is regular graph of degree k or k -regular if every vertex of G has degree k .
- ▶ In other words, a graph is regular if every vertex has the same degree.
- ▶ Following are some regular graphs.



(i) 0-regular



(ii) 1-regular

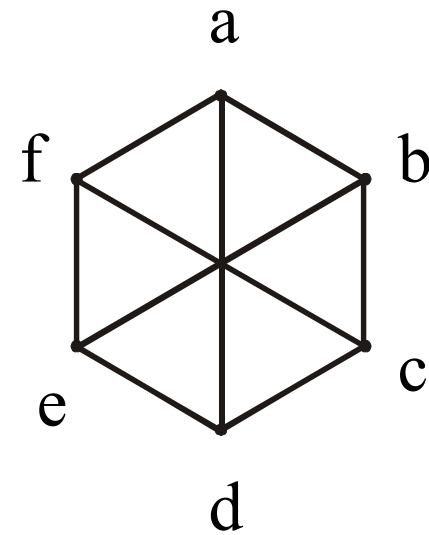
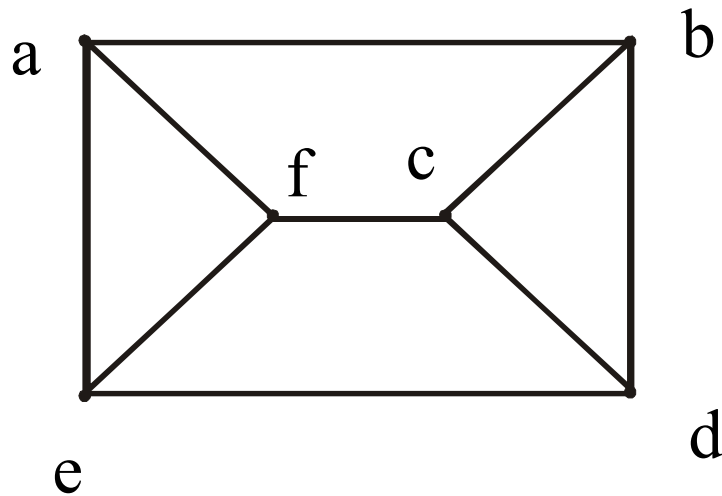


(iii) 2-regular

EXERCISE

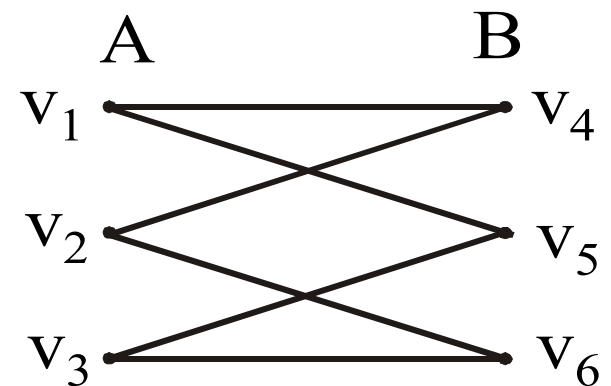
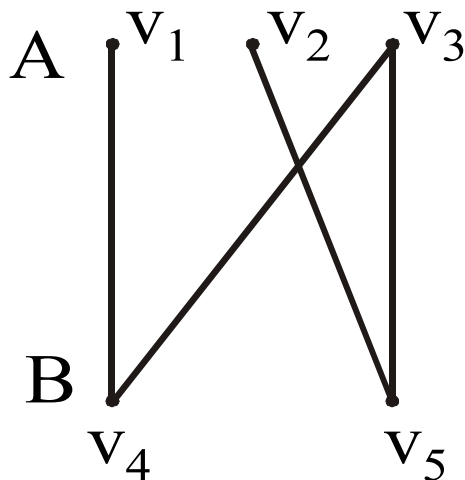
- ▶ Draw two 3-regular graphs with six vertices.

- ▶ SOLUTION:



BIPARTITE GRAPH

- ▶ A bipartite graph G is a simple graph whose vertex set can be partitioned into two mutually disjoint non empty subsets A and B such that the vertices in A may be connected to vertices in B , but no vertices in A are connected to vertices in A and no vertices in B are connected to vertices in B .



MATRIX REPRESENTATIONS OF GRAPHS

- ▶ To store graph in computer with pictorial representation is not possible rather you will store the graph with matrix representation.
- ▶ It is difficult to analyze a big complex graph with hundreds of vertices and thousands of edges, but in matrix form you can analyze big graph better.

MATRIX

- ▶ An $m \times n$ matrix A over a set S is a rectangular array of elements of S arranged into m rows and n columns:

$$A = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1j} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2j} & \cdots & a_{2n} \\ \vdots & \vdots & & \vdots & & \vdots \\ a_{i1} & a_{i2} & \cdots & a_{ij} & \cdots & a_{in} \\ \vdots & \vdots & & \vdots & & \vdots \\ a_{m1} & a_{m2} & \cdots & a_{mj} & \cdots & a_{mn} \end{bmatrix}$$

j^{th} column of A

Briefly, it is written as: $A = [a_{ij}]_{m \times n}$

EXAMPLE

$$A = \begin{bmatrix} 4 & -2 & 0 & 6 \\ 2 & -3 & 1 & 9 \\ 0 & 7 & 5 & -1 \end{bmatrix}$$

- ▶ **A** is a **matrix** having **3 rows** and **4 columns**. We call it a **3×4 matrix**, or matrix of **size 3×4** (or we say that a **matrix** having an **order 3×4**).

- ▶ Note:

$a_{11} = 4$ (11 means 1st row and 1st column),

$a_{12} = -2$ (12 means 1st row and 2nd column), $a_{13} = 0$,

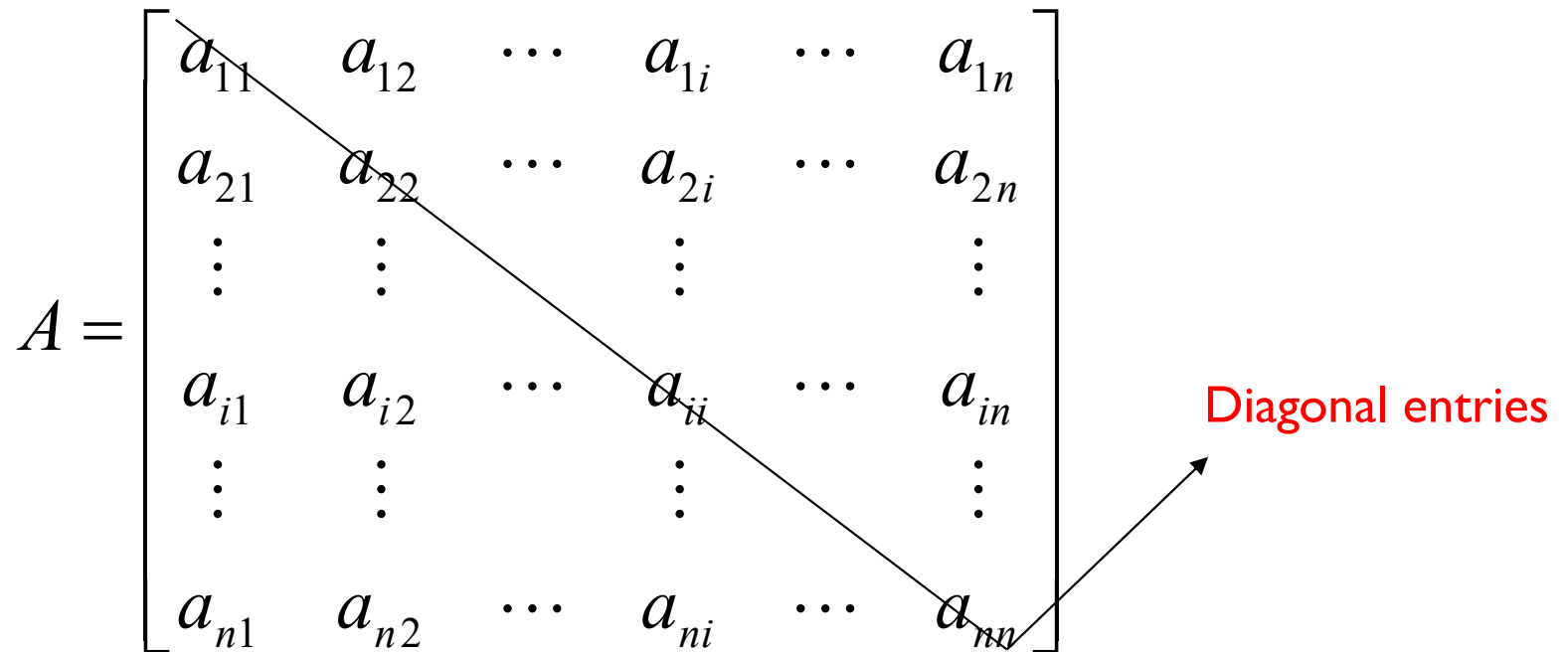
$a_{14} = 6$ $a_{21} = 2$, $a_{22} = -3$, $a_{23} = 1$, $a_{24} = 9$ etc.

SQUARE MATRIX

- ▶ A **matrix** for which the number of rows and columns are equal is called a **square matrix**.

$$A = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1i} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2i} & \cdots & a_{2n} \\ \vdots & \vdots & & \vdots & & \vdots \\ a_{i1} & a_{i2} & \cdots & a_{ii} & \cdots & a_{in} \\ \vdots & \vdots & & \vdots & & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{ni} & \cdots & a_{nn} \end{bmatrix}$$

Diagonal entries



- ▶ The main **diagonal** of **A** consists of all the entries

$$a_{11}, a_{22}, a_{33}, \dots, a_{ii}, \dots, a_{nn}$$

TRANSPOSE OF A MATRIX

- ▶ The **transpose** of a **matrix A** of **size $m \times n$** , is the **matrix** denoted by **A^t** of **size $n \times m$** , obtained by writing the **rows of A**, in order, as columns.

$$A = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{m1} & a_{m2} & \cdots & a_{mn} \end{bmatrix}, \quad \text{then } A^t = \begin{bmatrix} a_{11} & a_{21} & \cdots & a_{m1} \\ a_{12} & a_{22} & \cdots & a_{m2} \\ \vdots & \vdots & & \vdots \\ a_{1n} & a_{2n} & \cdots & a_{mn} \end{bmatrix}$$

EXAMPLE

$$A = \begin{bmatrix} 4 & -2 & 0 & 6 \\ 2 & -3 & 1 & 9 \\ 0 & 7 & 5 & -1 \end{bmatrix}$$

► Then

$$A^t = \begin{bmatrix} 4 & 2 & 0 \\ -2 & -3 & 7 \\ 0 & 1 & 5 \\ 6 & 9 & -1 \end{bmatrix}$$

SYMMETRIC MATRIX

- ▶ A square matrix $A = [a_{ij}]$ of size $n \times n$ is called **symmetric** if, and only if, $A^t = A$ i.e., for all $i, j = 1, 2, \dots, n$, $a_{ij} = a_{ji}$

EXAMPLE

Let $A = \begin{bmatrix} 1 & 3 & 7 \\ 5 & 2 & 9 \end{bmatrix}$, and $B = \begin{bmatrix} 4 & 2 & 0 \\ 2 & -3 & 1 \\ 0 & 1 & 5 \end{bmatrix}$

Then $A^t = \begin{bmatrix} 1 & 5 \\ 3 & 2 \\ 7 & 9 \end{bmatrix}$, and $B^t = \begin{bmatrix} 4 & 2 & 0 \\ 2 & -3 & 1 \\ 0 & 1 & 5 \end{bmatrix}$

- Note that $B^t = B$, so that B is a **symmetric matrix**.

MATRIX MULTIPLICATION

$$\begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1p} \\ \vdots & \vdots & & \vdots \\ a_{i1} & a_{i2} & \cdots & a_{ip} \\ \vdots & \vdots & & \vdots \\ a_{m1} & a_{m2} & \cdots & a_{mp} \end{bmatrix} \begin{bmatrix} b_{11} & \cdots & b_{1j} & \cdots & b_{1n} \\ b_{21} & \cdots & b_{2j} & \cdots & b_{2n} \\ \vdots & & & & \\ b_{p1} & \cdots & b_{pj} & \cdots & b_{pn} \end{bmatrix} = \begin{bmatrix} c_{11} & \cdots & c_{1j} & \cdots & c_{1n} \\ \vdots & & & & \\ c_{i1} & \cdots & c_{ij} & \cdots & c_{in} \\ \vdots & & & & \\ c_{m1} & \cdots & c_{mj} & \cdots & c_{mn} \end{bmatrix}$$

► Note:

If the **number of columns** of **A** is **not equal** to the **number of rows** of **B**, then the **product AB** is **not defined**.

EXAMPLE

- Find the **product AB** and **BA** of the **matrices**

$$A = \begin{bmatrix} 1 & 3 \\ 2 & -1 \end{bmatrix}$$

and

$$B = \begin{bmatrix} 2 & 0 & -4 \\ 3 & -2 & 6 \end{bmatrix}$$

► **SOLUTION:**

Size of **A** is **2×2** and of **B** is **2×3** , the **product AB** is defined as a **2×3 matrix**.

$$AB = \begin{bmatrix} 1 & 3 \\ 2 & -1 \end{bmatrix} \begin{bmatrix} 2 & 0 & -4 \\ 3 & -2 & 6 \end{bmatrix}$$

$$= \begin{bmatrix} (1)(2) + (3)(3) & (1)(0) + (3)(6) & (1)(-4) + (3)(6) \\ (2)(2) + (-1)(3) & (2)(0) + (-1)(-2) & (2)(-4) + (-1)(6) \end{bmatrix}$$

$$= \begin{bmatrix} 11 & -6 & 14 \\ 1 & 2 & -14 \end{bmatrix}$$

ADJACENCY MATRIX OF A GRAPH

- ▶ Let G be a graph with ordered vertices v_1, v_2, \dots, v_n . The adjacency matrix of G is the matrix $A = [a_{ij}]$ over the set of non-negative integers such that

a_{ij} = the number of edges connecting v_i and v_j
for all $i, j = 1, 2, \dots, n$.

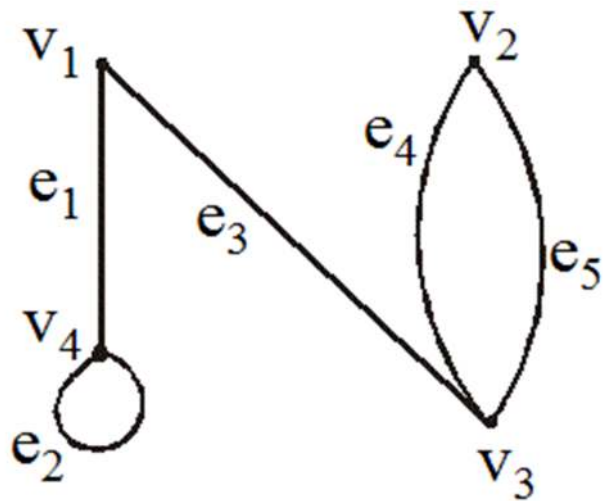
OR

- ▶ The adjacency matrix say $A = [a_{ij}]$ is also defined as

$$a_{ij} = \begin{cases} 1 & \text{if } \{v_i, v_j\} \text{ is an edge of } G \\ 0 & \text{otherwise} \end{cases}$$

EXAMPLE

- ▶ A graph with its adjacency matrix is shown.



$$A = \begin{matrix} & \begin{matrix} v_1 & v_2 & v_3 & v_4 \end{matrix} \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix} & \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 2 & 0 \\ 1 & 2 & 0 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$

- ▶ Clearly graph has four vertices. It means that the corresponding square matrix will be order 4×4 .

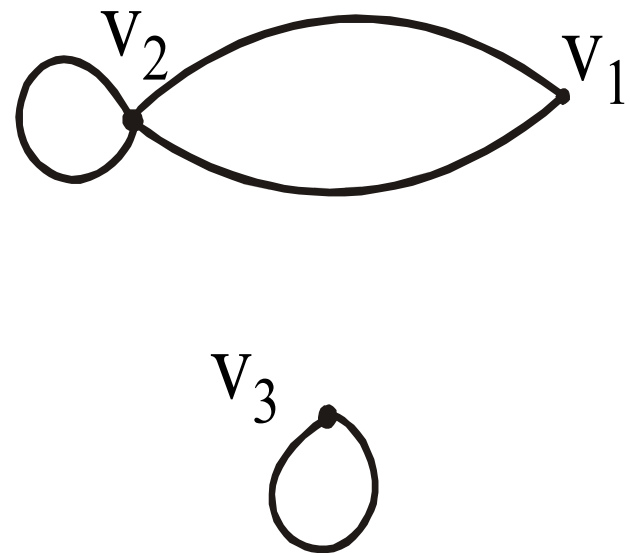
EXERCISE

- ▶ Find a **graph** that have the following **adjacency matrix**.

$$\begin{bmatrix} 0 & 2 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

- ▶ Its order is **3 x 3**, it means its corresponding graph has **three vertices**.
- ▶ Let the **three vertices** of the **graph** be named **v_1 , v_2 and v_3**

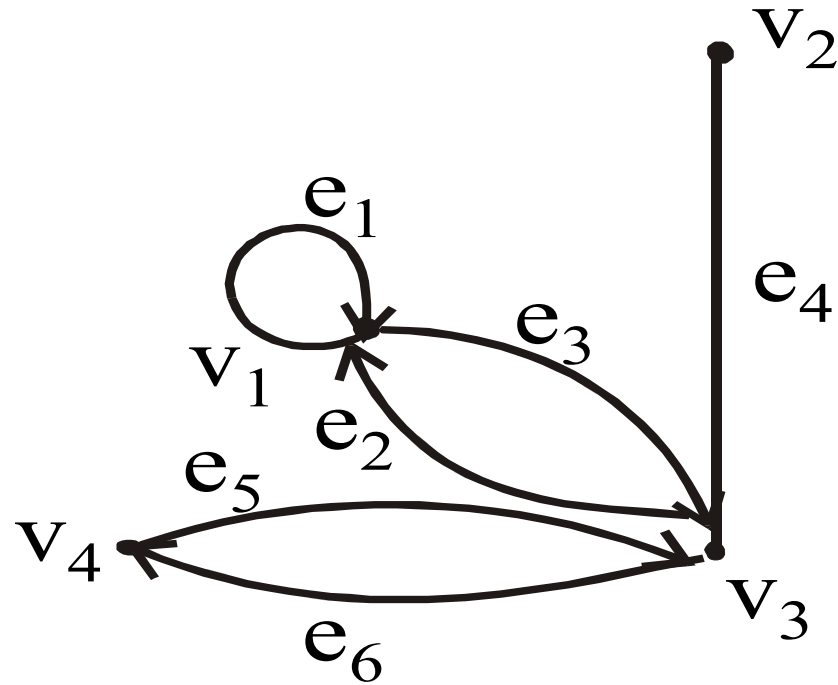
$$\begin{array}{c} v_1 \\ v_2 \\ v_3 \end{array} \begin{bmatrix} v_1 & v_2 & v_3 \\ 0 & 2 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$



DIRECTED GRAPH

- ▶ A **directed graph** or **digraph**, consists of two finite **sets**: a set $V(G)$ of **vertices** and a set $D(G)$ of **directed edges**,
- ▶ where each **edge** is associated with an **ordered pair** of **vertices** called its **end points**.
- ▶ If **edge** e is associated with the **pair** (v, w) of **vertices**, then e is said to be the **directed edge** from v to w and is represented by drawing an **arrow** from v to w .

EXAMPLE OF DIGRAPH



ADJACENCY MATRIX OF A DIRECTED GRAPH

- ▶ Let G be a graph with ordered vertices

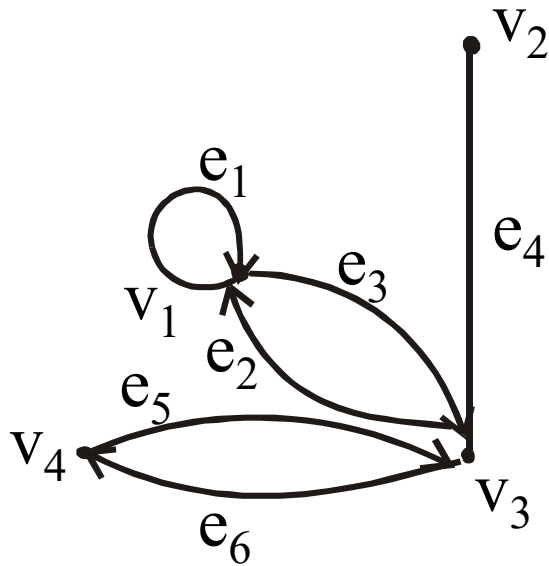
$$v_1, v_2, \dots, v_n.$$

- ▶ The adjacency matrix of G is the matrix $A = [a_{ij}]$ over the set of non-negative integers such that

a_{ij} = the number of arrows from v_i to v_j
for all $i, j = 1, 2, \dots, n$.

EXAMPLE

- ▶ A directed graph with its adjacency matrix is shown



$$A = \begin{matrix} & v_1 & v_2 & v_3 & v_4 \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix} & \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

Adjacency matrix

EXERCISE

- ▶ Find **directed graph** that has the **adjacency matrix**

$$\begin{bmatrix} 1 & 0 & 1 & 2 \\ 0 & 0 & 1 & 0 \\ 0 & 2 & 1 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$$

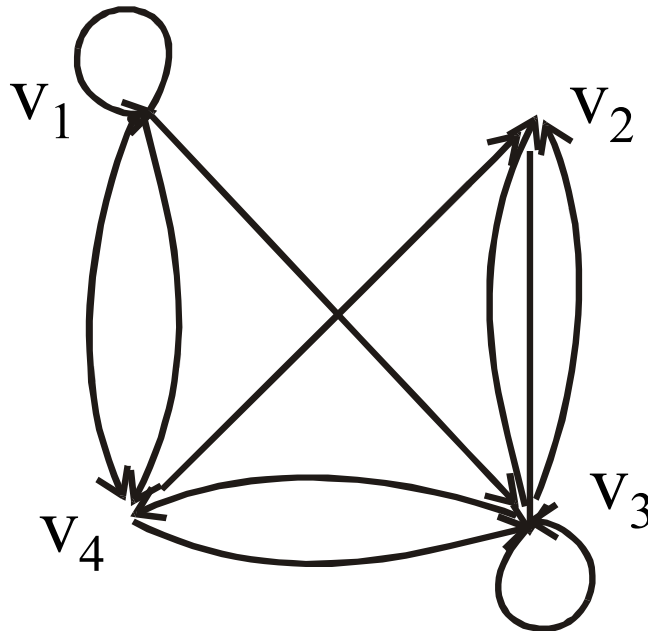
- ▶ The order of **matrix** is **4 x 4**, it means it has **4 vertices**.

SOLUTION

- ▶ The 4×4 adjacency matrix shows that the graph has 4 vertices say v_1, v_2, v_3 and v_4 labeled across the top and down the left side of the matrix.

$$A = \begin{array}{c} v_1 \\ v_2 \\ v_3 \\ v_4 \end{array} \begin{array}{cc} & \begin{array}{cccc} v_1 & v_2 & v_3 & v_4 \end{array} \\ \left[\begin{array}{cccc} 1 & 0 & 1 & 2 \\ 0 & 0 & 1 & 0 \\ 0 & 2 & 1 & 1 \\ 0 & 1 & 1 & 0 \end{array} \right] \end{array}$$

-
- ▶ A corresponding **directed graph** is



- ▶ It means that a **loop** exists from **v_1** and **v_3** , **two** arrows go from **v_1** to **v_4** and **two** from **v_3** and **v_2** and **one arrow** go from **v_1** to **v_3** , **v_2** to **v_3** , **v_3** to **v_4** , **v_4** to **v_2** and **v_3** .

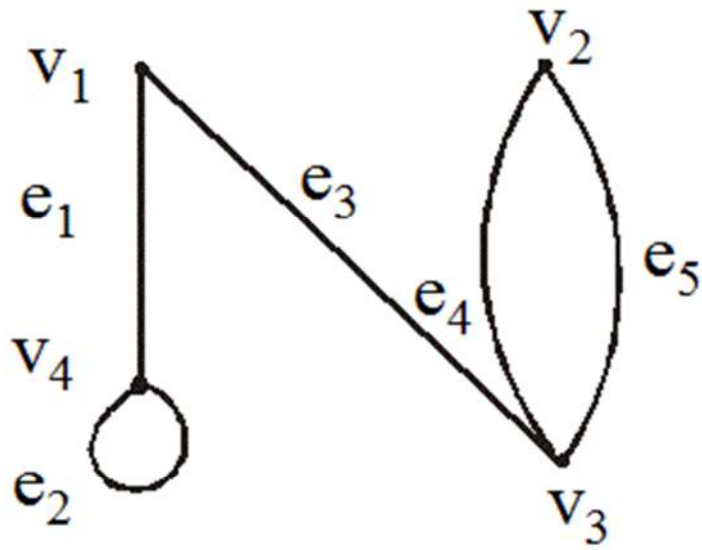
INCIDENCE MATRIX OF A SIMPLE GRAPH

- ▶ Let G be a graph with vertices v_1, v_2, \dots, v_n and edges e_1, e_2, \dots, e_m .
- ▶ The incidence matrix of G is the matrix $M = [m_{ij}]$ of size $n \times m$ defined by:

$$m_{ij} = \begin{cases} 1 & \text{if the vertex } v_i \text{ is incident on the edge } e_j \\ 0 & \text{otherwise} \end{cases}$$

EXERCISE

- A graph with its incidence matrix is shown.



$$M = \begin{matrix} & e_1 & e_2 & e_3 & e_4 & e_5 \\ \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{matrix} & \begin{bmatrix} 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

REMARK

- ▶ In the **incidence matrix**:
 - ▶ **Parallel edges** are represented by columns with **identical entries** (in this matrix e_4 & e_5 are **parallel edges**).
 - ▶ **Loops** are represented using a **column** with **exactly one** entry equal to **1**, corresponding to the vertex that is incident with this loop and **other zeros** (here e_2 is only a loop).