# Combinatorial Auction Mechanism(VCG) for Computation Distribution

for

Advanced Topics in Intelligent Systems - CSE6369

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## **Problem Statement:**

Given a range of available resources on a computational cluster (divided into 3 types: processing, memory, disk space, each with a specific amount), implement an efficient (VCG) auction mechanism that can distribute these to N bidding agents that each want to schedule a randomly constructed job (i.e. with random resource requirements) with a true valuation that is additive in terms of the values of the required resources.

# **Vickrey Clarke Groves Auction:**

In auction theory, a Vickrey–Clarke–Groves (VCG) auction is a type of sealed-bid auction of multiple items. Bidders submit bids that report their valuations for the items, without knowing the bids of the other people in the auction. The auction system assigns the items in a *socially optimal manner*.

# **VCG Auction:**

For any set of auctioned items  $M=\{t_1,\ldots,t_m\}$  and any set of bidders  $N=\{b_1,\ldots,b_n\}$ , let  $V_N^M$  be the social value of the VCG auction for a given bid-combination. That is, how much each person values the items they've just won, added up across everyone. The value of the item is zero if they do not win. For a bidder  $b_i$  and item  $t_j$ , let the bidder's bid for the item be  $v_i(t_j)$ . The notation  $A\setminus B$  means the set of elements of A which are not elements of B.

### **Assignment**

A bidder  $b_i$  whose bid for an item  $t_j$  is an "overbid", namely  $v_i(t_j)$ , wins the item, but pays  $V_{N\setminus\{b_i\}}^M - V_{N\setminus\{b_i\}}^{M\setminus\{t_j\}}$ , which is the social cost of their winning that is incurred by the rest of the agents.

### **Explanation**

Indeed, the set of bidders other than  $b_i$  is  $N\setminus\{b_i\}$ . When item  $t_j$  is available, they could attain welfare  $V_{N\setminus\{b_i\}}^M$ . The winning of the item by  $b_i$  reduces the set of available items to  $M\setminus\{t_j\}$ , however, so that the attainable welfare is now  $V_{N\setminus\{b_i\}}^{M\setminus\{t_j\}}$ . The difference between the two levels of welfare is therefore the loss in attainable welfare suffered by the rest of the bidders, as predicted, **given** the winner  $b_i$  got the item  $t_j$ . This quantity depends on the offers of the rest of the agents and is unknown to agent  $b_i$ .

### Winner's utility:

The winning bidder whose bid is the true value A for the item  $t_j$ ,  $v_j\{t_j\}$ =A derives maximum  $A - \left(V_{N\setminus\{b_i\}}^M - V_{N\setminus\{b_i\}}^{M\setminus\{t_j\}}\right)$ 

### Project: vcg-auction.ipnyb

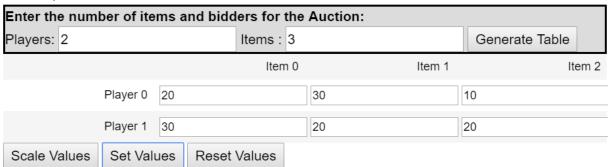
```
pip3 install jupyter
jupyter notebook
```

### **Code Structure:**

Enter the number of *players* and number of *items* and click on *generate table*. Here players are the agents and items are the resources.

After giving the inputs click on Set Values.

### For Example,



- First compute all different one-to-one allocations of items to bidders.
- valuations are: [array([20, 30, 10]), array([30, 20, 20])]
- Identify all possible allocations of items to bidders
- The value of auction matrix is :

```
[[ 0 0 0 60 0 60]
[ 0 0 1 50 20 70]
[ 0 1 0 30 20 50]
[ 0 1 1 20 40 60]
[ 1 0 0 40 30 70]
[ 1 0 1 30 50 80]
[ 1 1 0 10 50 60]
[ 1 1 1 0 70 70]]
```

In the following auction matrix i0,i1 and i2 are the items and the v0, v1 are the valuations and v\_tot is the total valuation. There are total of  $8(2^3)$  combinations for which the valuations of the players if they get the resources are displayed.

```
The auction matrix is
    i0
       i1 i2 v0
                      v1
                          v tot
0
    ()
         \cap
             0
                 60
                      0
                             60
    0
         0
             1
                 50
                             70
1
                     20
2
    0
        1
             0
                30
                     30
                             60
3
    0
         1
             1
                 20
                     50
                             70
4
    1
        0
             0
                40
                     30
                             70
5
    1
        0
            1
                30
                     50
                             80
6
    1
        1
             0
                10
                     60
                             70
    1
         1
                             80
             1
                 0
                     80
```

Here we are using the Clarke-Pivot Rule

$$h_i(v_{-i}) = -\max_{x \in X} \sum_{j 
eq i} v_j(x)$$

(value of others if agent/player were absent) - (value of others when agent/player were present).

• By using the Clarke Pivot Rule, we only take the cases where we get the max evaluation

• Items: ['i0', 'i1', 'i2']

Valuation of Player to their allocation: ['v0', 'v1'] Payment of Player for their allocation: ['p0', 'p1']

	i0	i1	i2	v0	v1	v_tot	p0	p1
0	1	0	1	30	50	80	20	30

### VCG Payments:

The VCG Payments are given by

Player 1: v0-p0 Player 2: v1-p1

> VCG Cost for Player 0: 10 30 - 20 VCG Cost for Player 1: 20 50 - 30

### **References:**

https://pathtogeek.com/vickrey-clarke-groves-vcg-auction

https://en.wikipedia.org/wiki/Vickrey%E2%80%93Clarke%E2%80%93Groves\_auction

https://www.youtube.com/watch?v=8-t BmeNWgs