Sure! Here's a simplified version of the video transcript using easy words:

[MUSIC]

Welcome to my course **Numerical Methods for Engineers**.

In this course, we'll use math from other courses—like matrix algebra, differential equations, and vector calculus—and apply it in **scientific computing** using **MATLAB**.

♦ Week 1:

We'll learn the basics:

- How computers store numbers
- How to write simple programs in MATLAB

♦ Week 2:

We'll learn **root finding**—this means solving equations like f(x) = 0 when we can't solve them by hand

We'll learn the common methods, how fast they are, and how to use MATLAB to find roots.

♦ Week 3:

We'll study matrix algebra:

- How to solve linear systems
- How to find eigenvalues and eigenvectors
- How to solve nonlinear systems
 This is important for solving PDEs later.

♦ Week 4:

We'll learn:

- Quadrature (numerical integration—finding the area under a curve)
- Interpolation (finding values between known data points)
 We'll use MATLAB for both, including cubic spline interpolation (which works very well).

♦ Week 5:

We'll work with ordinary differential equations (ODEs).

These are equations used in many science problems.

We'll learn easy methods and also Runge-Kutta methods, and use MATLAB's ode45.m tool.

We'll also briefly talk about boundary value problems, which use both ODE solving and root finding.

♦ Week 6:

We'll learn about partial differential equations (PDEs).

This is the hardest part.

We'll need matrix algebra to solve big systems of equations.

We'll also learn iterative methods and some theory like stability.

- Z Each week includes:
 - Short video lessons
 - Practice problems
 - A quiz and a small coding project

particular in Thanks for joining me in learning Numerical Methods for Engineers. Let's get started!

Let me know if you want a PDF or notes version!

Here's a simplified version of your Week 1 introduction in easy words:

Welcome to Week 1 of Numerical Methods for Engineers!

This week, we'll start learning the basics of scientific computing.

- First, we'll learn:
 - How **computers store numbers** using **binary** (just 0s and 1s).
 - What **double precision numbers** are these are the main type of numbers we'll use.
- These numbers have limits:
 - Some fractions (like 1/3 or 1/5) can't be stored exactly.
 - Numbers can't be too big or too small, or they may overflow or underflow.
 - This can cause small **errors** in our results.
- Then, we'll learn MATLAB basics:
 - How to use it like a calculator
 - How to make vectors and matrices
 - How to draw plots (line graphs)
 - How to write scripts and functions
 - How to use logic, if-else statements, and loops

If you already know how to program, this will be easier.

If not, you can learn the basics of MATLAB this week.

Your first project is to create a bifurcation diagram for something called the logistic map. It's a simple equation, but it creates very complex patterns.

Join me in this exciting first week of learning numerical methods!

Here's a **simple version** of the lecture transcript using easy words:

[MUSIC]

Before we start using computers to do math, we need to learn **how computers store numbers**. We'll cover that in this and the next lecture.

Decimal Numbers (Base-10):

- We normally use **decimal numbers** (also called **base-10**).
- They use digits from 0 to 9.
- We likely use base-10 because we have 10 fingers (digits = fingers!).

Example:

524.503

Each digit's **position** matters:

- $5 \times 10^2 = 500$
- $2 \times 10^1 = 20$
- $4 \times 10^{\circ} = 4$
- $5 \times 10^{-1} = 0.5$
- $0 \times 10^{-2} = 0$
- $3 \times 10^{-3} = 0.003$

So, 524.503 = 500 + 20 + 4 + 0.5 + 0 + 0.003

This is called a **positional number system**.

Decimal numbers made math like adding and multiplying **easy**, compared to old systems like **Roman numerals**.

Binary Numbers (Base-2):

- Computers use binary, not decimals.
- Binary uses only **0s and 1s**. These are called **bits**.
- Binary is a base-2 positional system.

Example:

101.011 (binary)

- $\bullet 1 \times 2^2 = 4$
- $0 \times 2^1 = 0$
- $1 \times 2^{\circ} = 1 \rightarrow \text{Total before point} = 5$

- Binary point = like the decimal point
- After the point:
 - $0 \times 2^{-1} = 0$
 - $0.1 \times 2^{-2} = 0.25$
 - $0.1 \times 2^{-3} = 0.125$
 - → Total after point = 0.375
 - **So,** 101.011 (binary) = **5.375** in decimal

Some numbers can't be stored exactly:

In base-10:

- 1/2 = 0.5
- 1/3 = 0.333... **(repeats forever, inexact)**
- 1/5 = 0.2
- 1/7 = 0.142857142857... **(repeats forever)**

In binary:

- 1/2 = 0.1
- 1/3 = 0.010101... **X** (repeats)
- 1/4 = 0.01
- 1/7 = 0.001001... **(repeats)**

These repeating numbers can't be stored exactly on a computer.

We call them inexact numbers.

↑ What does this mean for you?

- When numbers can't be stored exactly, **small errors** can happen.
- These are called round-off errors.
- Most of the time, they don't matter much.
- But it's important to **understand** them when writing programs.

Summary:

- Decimal: base-10, digits 0-9
- Binary: base-2, digits 0 and 1 (bits)
- Some fractions can't be stored exactly in decimal or binary → inexact numbers

- This can lead to **rounding errors** when computers do math
- I'm Jeff Chesnoff. Thanks for watching, and I'll see you in the next video!

Here's a clear **summary** of the video lecture you just watched on **IEEE Floating Point Arithmetic** (**Double-Precision**) by **Jeff Chasnov**:

(64-bit) Key Concepts of IEEE 754 Double-Precision Floating Point Representation (64-bit)

1. Bit Structure

A double-precision number uses **64 bits**:

- 1 bit for Sign (S)
 - $0 \rightarrow Positive$
 - 1 → Negative
- 11 bits for Exponent (E)
 - Stored as biased exponent: actual exponent = E 1023
- o 52 bits for Fraction (F) (also called the significand or mantissa)
 - Value is stored as 1.F in binary (with implicit leading 1)

Formula:

(-1)S×1.F×2(E-1023)(-1)^S \times 1.F \times 2^{(E - 1023)}

2. Special/Reserved Values

- o Infinity (±∞)
 - E = all 1s (2047) and F = all 0s
 - Example: Division by 0 (e.g., 1 / 0)
- NaN (Not a Number)
 - E = all 1s and F ≠ all 0s
 - Example: 0 / 0, ∞ ∞
- Denormal/Subnormal Numbers
 - E = all 0s
 - Value is computed as:

(-1)S×0.F×2-1022(-1)^S \times 0.F \times 2^{-1022}

Used to represent numbers very close to zero (graceful underflow)

- o Zero
 - All bits = 0 (positive zero)
 - Sign bit = 1 (negative zero)

3. Important Limits

o realmax (Maximum Representable Number)

≈1.797×10308\approx 1.797 \times 10^{308}

o realmin (Smallest Normal Positive Number)

≈2.2251×10-308\approx 2.2251 \times 10^{-308}

4. Machine Epsilon (ε)

Smallest difference such that:

1+ε≠11 + \varepsilon \neq 1

o For double precision:

 $\varepsilon \approx 2.22 \times 10 - 16 \text{ varepsilon } 2.22 \text{ times } 10^{-16}$

Used to detect roundoff errors

5. Sources of Roundoff Error

- o Adding a very small number to a large one
 - If $yx<\epsilon 2\frac{y}{x} < \frac{x}{x}$, then $x+y\approx xx + y \approx x$
- Subtracting nearly equal numbers
 - Can lose significant digits ("catastrophic cancellation")

6. Why IEEE Standard is Powerful

- o Represents:
 - Very large numbers (up to 1030810^{308})
 - Very small numbers (down to 10–30810^{-308})
 - Precise values with up to 16 decimal digits of accuracy
 - Special values: ∞, NaN, subnormals, etc.

Let me know if you want a diagram, practice questions, or help with binary-to-decimal examples using IEEE 754 format.

Here's a summary of the key points from the MATLAB video tutorial by Jeff Chasnov on writing mathematical expressions in MATLAB:

Basic Operators in MATLAB

Symbol	Operation	Example
+	Addition	3 + 2
-	Subtraction	5 - 1
*	Multiplication	4 * 3
/	Division	10/2
^	Exponentiation	2^3 = 8
•	Transpose (for matrices)	A'

Common Functions and Constants

Function/Constan	t Meaning	Example
sin(x), $cos(x)$, $tan(x)$ Trigonometric functions $sin(pi/2) = 1$		
log(x)	Natural logarithm (In)	log(exp(1)) = 1
log10(x)	Base-10 logarithm	log10(1000) = 3
sqrt(x)	Square root	sqrt(16) = 4
abs(x)	Absolute value or modulus of complex number abs(-5) = 5	
pi	π = 3.1416	
iorj	Imaginary unit	1 + 2i
inf	Infinity	1/0 = Inf
NaN	Not a Number (e.g. 0/0, Inf - Inf)	

Examples Demonstrated

- **Euler's Number:**
- exp(1) % ≈ 2.7183

- log(exp(1)) % = 1
- Golden Ratio (Ф):
- Phi = (sqrt(5) + 1) / 2; % ≈ 1.618
- phi = (sqrt(5) 1) / 2; % ≈ 0.618
- Phi * phi % ≈ 1 (with a small roundoff error)
- **Roundoff Errors:**
 - o Multiplying irrational numbers may give slight errors (e.g., not exactly 1).
 - Use format hex to view floating-point representation.

↑ Special Numbers

Operation Result 1/0 Inf -1/0 -Inf 0 / 0 or Inf - Inf NaN

Formatting and Output

Ηερεэσ α στρυχτυρεδ συμμαρψ οφ τηε τηιρδ ΜΑΤΛΑΒ σιδεο βψ θεφφ Χηασνοσ, φοχυσεδ ον σεχ τορσ ανδ πεχτορ οπερατιονσ ιν ΜΑΤΛΑΒ:

T .

🗆 Βασιχ Δατα Τψπε: Τηε Ματριξ

Ιν ΜΑΤΛΑΒ, επερψτηινή ισ α ματριξ \Box ινχλυδινή πεχτορσ, ωηιχή αρε 1Δ ματρίχεσ.

Τψπε	Δεφινιτιον	Εζαμπλε
Ροω ςεχτορ	1 ροω, μυλτιπλε χολυμνσ	$\xi = [1 \ 2 \ 3] \text{ op } \xi = 1 \ 2 \ 3$

Χολυμν ςεχτορ Μυλτιπλε ροωσ, 1 χολυμν $\xi = [1; 2; 3]$

Ψου χαν αχχεσσ ελεμεντσ τηε σαμε ωαψ φορ βοτη:

ξ(3) % ρετυρνσ τηε 3ρδ ελεμεντ

□ Τρανσποσε ςεχτορσ

Υσε ϶ (αποστροπηε) το τρανσποσε:

 $\xi = [1; 2; 3];$ % χολυμν πεχτορ

% βεχομεσ ροω πεχτορ [1 2 3]

🗆 ςεχτορ Γενερατιον Φυνχτιονσ

Φυνχτιον	Πυρποσε	Εξαμπλε
ζεροσ(μ,ν)	Ματριξ (ορ σεχτορ) οφ αλλ ζεροσ	$\xi = \zeta \epsilon \rho o \sigma(1,3) \square [0 \ 0 \ 0]$
ονεσ(μ,ν)	Ματριξ (ορ σεχτορ) οφ αλλ ονεσ	$\xi = \text{ones}(1,3) \square [1 \ 1 \ 1]$
λινσπαχε(α,β, ν)	ν εθυαλλψ σπαχεδ παλυεσ βετωεεν α α νδ β	λινσπαχε $(0,1,6)$ \square $[0\ 0.2\ 0.4\ 0.6\ 0.8\ 1]$
Χολον νοτατι ον	Δεφινε σταρτ:στεπ:ενδ	ξ = 0:0.2:1 \Box σαμε ασ αβοψε
• λινσπανε(0.1)	6) Thurse of market	

- $\lambda inspace(0,1,6) \square$ uumber of points
- 0:0.2:1 □ στεπ σιζε

□ Ελεμεντ-Ωισε Οπερατιονσ

Φορ χομπονεντ-ωισε οπερατιονσ ον πεχτορσ, υσε α δοτ (.) βεφορε τηε οπερατορ:

Οπερατιον	Σψνταξ	Εξαμπλε	Ρεσυλτ
Μυλτιπλιχατιον	.*	ξ.* ξ	[1 4 9]
Ποωερ	.1	$\xi \perp 2$	[1 4 9]
Δισισιον	./	ξ./ξ	[1 1 1]
Εξαμπλε:			
$\xi = [1 \ 2 \ 3];$			

🗆 Φυνχτιονσ Αππλιεδ το ςεχτορσ

ξ.* ξ % [149]

 $\xi \perp 2 \qquad \% [1 \ 4 \ 9]$

ξ./ξ %[111]

Μοστ ΜΑΤΛΑΒ φυνχτιονσ αχχεπτ πεχτορσ:

 $\xi = 0:\pi i/2:2*\pi i; \% [0 \pi i/2 \pi i 3\pi i/2 2\pi i]$

χοσ(ξ) % ρετυρνσ [1 0 –1 0 1] (αππροξ.)

Λ Τηέσε αρέ νοτ **εξαχτ** δυε το **ρουνδ-οφφ ερρορσ** (ε.γ., χοσ(πι/2) Η 6.1232ε-17 ινστεαδ οφ 0).

□ ςεχτορσ ιν Πλοττινγ

• Ιν νυμεριχαλ χομπυτινγ, φυνχτιονσ αρε **επαλυατεδ ασ πεχτορσ**.

Φορ πλοττινγ:

Ψου τψπιχαλλψ πλοτ ψ σαλυεσ (φυνχτιον σαλυεσ) αγαινστ ξ σαλυεσ (δομαιν ποιντσ).

□ Χονχλυσιον

 $\forall M$ ανψ νυμεριχαλ σολυτιονσ τακε τηε φορμ οφ ϖ εχτορσ \square Ω ε \ni λλ σεε ηοω το π λοτ φυνχτιονσ νεξτ

🗆 θεφφ Χηασνοω

Ωουλδ ψου λικε α πραχτιχε σηεετ ον σεχτορσ ορ σομε μινι-εξερχισεσ το αππλψ ωηατ ψουθεε λε αρνεδ?

Here's a structured **summary** of Jeff Chasnov's MATLAB video on **plotting graphs (specifically a circle)** using plot and customizing it:

✓ Plotting Basics in MATLAB

- Most solutions are visualized using line plots.
- Commonly, we plot y vs x, i.e., a graph of y = f(x).

Example: Plotting a Circle

1. Set up polar coordinates

theta = linspace(0, 2*pi); % 100 points from 0 to 2π (default) x = cos(theta);

y = sin(theta);

2. Plot the circle

plot(x, y)

▲ But this appears oval, because MATLAB scales the x and y axes differently by default.

K Fixing Plot Appearance

3. Make axes equally scaled

axis equal

Now it looks like a circle, but spacing might still be off.

4. Set axis limits manually

axis([-1.1 1.1 -1.1 1.1]) % [xmin xmax ymin ymax]

Now both axes show the same range.

5. Customize Tick Marks

Use gca (Get Current Axis) and dot notation to access axis properties:

```
% get current axis handle
ax = gca;
ax.XTick = [-1 -0.5 0 0.5 1]; % set x-axis tick marks
ax.YTick = [-1 -0.5 0 0.5 1]; % set y-axis tick marks
```

Add Labels and Title (Using LaTeX Style)

Use dollar signs to format as LaTeX math in labels:

```
xlabel('$x$', 'interpreter', 'latex', 'fontsize', 14)
ylabel('$y$', 'interpreter', 'latex', 'fontsize', 14)
title('Circle Plot', 'fontsize', 16)
```

You can skip LaTeX formatting if unfamiliar, or just use simple strings.

MATLAB Help and Documentation

- Use help or doc to learn about functions:
- help plot
- doc axis
- doc gca
- Most plotting customization requires checking documentation or Google.

Final Code (Clean Circle Plot)

```
theta = linspace(0, 2*pi);
x = cos(theta);
y = sin(theta);
plot(x, y)
axis equal
axis([-1.1 1.1 -1.1 1.1])
```

```
ax = gca;
ax.XTick = [-1 -0.5 0 0.5 1];
ax.YTick = [-1 -0.5 0 0.5 1];
xlabel('$x$', 'interpreter', 'latex', 'fontsize', 14)
ylabel('$y$', 'interpreter', 'latex', 'fontsize', 14)
title('Circle Plot', 'fontsize', 16)
```

Would you like a **hands-on task** to practice plotting functions like sin(x), e^x, or even a spiral? Here's a **structured summary** of Jeff Chasnov's video on **matrix operations in MATLAB**:

♦ What is MATLAB?

- MATLAB stands for Matrix Laboratory.
- Matrices are the core data structure in MATLAB.
- MATLAB provides many built-in **functions** for creating and manipulating matrices.

Creating Matrices

Function	Description	Example
zeros(m,n)	m×n matrix of all zeros	zeros(3,2) \rightarrow 3×2 matrix
ones(m,n)	m×n matrix of all ones	ones(2,3) \rightarrow 2×3 matrix
eye(n)	n×n identity matrix	eye(2) \rightarrow 2×2 identity
diag(v)	Diagonal matrix with elements of vector v	diag([1 2]) \rightarrow 2×2 with 1, 2 on diagonal

Accessing Matrix Elements

Example Matrix A:

Element Access:

Action	Syntax	Result
Access single element	A(2,3)	6
Submatrix	A(1:2, 2:3)	[2 3; 5 6]
First row	A(1,:)	[1 2 3]
First column	A(:, 1)	[1; 4; 7]
Single index (column-wise)	A(5)	5 (fifth element down column-wise)
Modify element	A(1,1) = 10	A is updated

Colon Operator:

- Used to select entire rows/columns or ranges.
- A(1,:) % all elements in row 1
- A(:,3) % all elements in column 3
- A(1, 2:end) % elements from 2nd to last in row 1

Matrix Multiplication

A = [1 2; 3 4];

B = [-3 -4; 4 5];

C = A * B;

- ✓ Valid because inner dimensions match: $(2\times2)*(2\times2) \rightarrow 2\times2$
- \times Invalid: $x = [1 \ 2 \ 3]; x * x; \rightarrow Error! (1×3)*(1×3) not valid.$
- ✓ Valid scalar product: $x * x' \rightarrow (1\times3)*(3\times1) = \text{scalar } 14$

Special Matrix: Banded Matrix

Create banded matrix (main diagonal and off-diagonals):

n = 4;

A = diag(2*ones(1,n)) + diag(-1*ones(1,n-1), 1) + diag(-1*ones(1,n-1), -1);

- ✓ Produces a 4×4 matrix with:
- 2 on main diagonal
- -1 above and below the diagonal

• Applications

- Matrix operations are **fundamental** in:
 - Solving nonlinear systems
 - Solving PDEs
 - o Performing numerical methods

Syntax

x * x'

a Summary Table

Dot/scalar product

Operation

- •	
Create zero matrix	zeros(m,n)
Create ones matrix	ones(m,n)
Identity matrix	eye(n)
Diagonal matrix	diag([a b c])
Access row/column	A(1,:), A(:,2)
Multiply matrices	A * B
Transpose	x'

Would you like a practice task or code quiz based on matrix manipulation in MATLAB?

Here's a structured summary of Jeff Chasnov's video on **Logical Variables and Operators in MATLAB**:

✓ Logical Variables in MATLAB

- Logical variables can take only **two values**:
 - \circ 1 \rightarrow true
 - \circ 0 \rightarrow false

Logical Operators

Operator	Meaning	Example	Result
>	Greater than	3 > 2	1
<	Less than	2 < 3	1
>=	Greater than or equal to	3 >= 3	1
<=	Less than or equal to	2 <= 1	0
==	Equal to	4 == 4	1
~=	Not equal to	4 ~= 5	1
~	NOT	~(1 == 1)	0
&	AND	1 & 0	0
`	•	OR	`1

Mote: == is for comparison, = is for assignment.

Examples in MATLAB

p = 0 < 1; % p = 1 (true)

q = 0 > 1; % q = 0 (false)

• These are **logical variables** and can be used in computations:

p * 5 % Result: 5, because 1 * 5 = 5

q * 5 % Result: 0, because 0 * 5 = 0

Q Logical Arrays (Element-wise Comparison)

 $x = [0 \ 1 \ 2 \ 0];$

 $y = [0 \ 2 \ 4 \ 0];$

z = (x == y); % Compare element-wise

Result:

z = [1 0 0 1]

- 0 == 0 → 1
- 1 == 2 → 0
- 2 == 4 → 0
- $0 == 0 \rightarrow 1$

Operator Precedence: Logical vs Arithmetic

Example:

result = 10 / 2 == 5; % TRUE (1)

- 10 / 2 is done **first**, then compared to 5.
- To make it clear, use parentheses:

result = (10 / 2) == 5;

Recommended: Use extra parentheses for clarity and safety.

Why Logical Variables Matter

- Logical expressions are used in **conditional statements**:
 - o if, else, while, etc.
 - o These control the **flow of your programs**.

Would you like a mini MATLAB quiz or practice exercise on logical operations?

Here's a clear summary of Jeffrey Chasnov's video on Conditional Statements in MATLAB:

What Are Conditional Statements?

- Conditional statements allow MATLAB to **decide which code to execute** based on whether a **logical expression is true or false**.
- These are **core elements** of programming logic.

Types of Conditional Statements in MATLAB

1. if ... end

if condition

% Code runs if condition is true

end

- If the condition is **true** → execute the code.
- If the condition is false → do nothing.

2. if ... else ... end

if condition

% Code runs if condition is true

else

% Code runs if condition is false

end

• Either one **or** the other block of code will run, depending on the condition.

Example: sinc Function Fix

We want to define:

MATLAB Code:

function $y = my_sinc(x)$

```
if x ~= 0
  y = sin(x) / x;
else
  y = 1;
end
```

end

- If $x \neq 0 \rightarrow$ normal formula
- If $x = 0 \rightarrow \text{return 1}$ (the correct mathematical limit)

Test Cases

```
my_sinc(pi/4) \% \approx 0.9036 (same as sin(pi/4)/pi/4)
my_sinc(0) \% = 1
```

Simplified Version (Optional)

```
function y = my_sinc(x)

if x

y = sin(x) / x;
```

else

```
y = 1;
```

end

end

- if x is **true** when $x \neq 0$
- Still behaves the same way
- More concise, but slightly less readable

Why Conditional Statements Matter

- Used to control the program's flow
- Essential in functions, loops, and decision-making logic

Would you like a short **practice problem or quiz** using conditional statements in MATLAB? Here is a structured summary of **Jeffrey Chasnov's video** on **Iteration in MATLAB**:

☑ Iteration in MATLAB

Iteration is essential in mathematical algorithms for repeatedly performing computations.

MATLAB provides two main loop structures:

1. for Loop (Fixed number of iterations)

Syntax:

for n = 1:N

% Code to execute N times

end

- n takes values from 1 to N.
- Commonly used for **predefined number of steps**.

Example: Logistic Map with for Loop

```
n_times = 10000;
x = 0.5;
r = 2.5;
```

```
for n = 1:n_times
    x = r * x * (1 - x);
end
```

disp(x)



- Change initial condition: $x = 0.1 \rightarrow \text{still converges to } 0.6$
- Increase n_times: still 0.6
- Change r to 2.6 → different result (shows sensitivity to parameters)

2. while Loop (Condition-based iteration)

Syntax:

while condition

% Code to execute while condition is true

end

- Runs until a condition becomes false.
- Often used for **convergence checks** or **unknown number of iterations**.

Example: Convergence with while Loop

```
x = 0.1;
r = 2.5;
tol = 1e-8;
error = 2 * tol;

while error > tol
    x_old = x;
    x = r * x * (1 - x);
    error = abs(x - x_old);
end
```

disp(x)

- Iterates until the change between x values is smaller than tol.
- Also converges to 0.6 for r = 2.5

break and continue

- break: exits the loop immediately
- continue: skips to the next iteration

Rarely used, but useful in certain control flow conditions.

Wey Insights

- Most computational code runs **inside loops**, often for thousands of iterations.
- Loop code must be efficient and correct.
- Prefer for loop for fixed iterations, and while loop for convergence-based iteration.

Would you like a practice exercise on loops in MATLAB to solidify your understanding?

Here's a **clear summary** of Professor Jeff Chasnov's lecture on the **bifurcation diagram for the logistic map** and the **theory behind it**:

Project Overview: Bifurcation Diagram for the Logistic Map

The Logistic Map Equation

 $xn+1=\mu xn(1-xn)x_{n+1} = \mu x_n (1-x_n)$

- A nonlinear recurrence relation (or iterative map)
- μ\mu is a parameter, typically varied from 0 to 4
- Used to model population dynamics

Iteration Concept

- Start with an initial value x0x_0
- Apply the function repeatedly to get:

$$x1=f(x0), x2=f(x1), x3=f(x2), ...x_1 = f(x_0), \quad x_2 = f(x_1), \quad x_3 = f(x_2), \quad x_4 = f(x_1), \quad x_4 = f(x_2), \quad x_4 = f(x_2), \quad x_4 = f(x_2), \quad x_4 = f(x_3), \quad x_4 = f(x_4), \quad x_4 = f(x_4$$

A **fixed point** x*x^* satisfies:

$$x*=f(x*)=\mu x*(1-x*)x^* = f(x^*) = \mu x^* (1-x^*)$$

Solving gives:

- $x*=0x^*=0$
- $x*=1-1\mu x^* = 1 \frac{1}{\mu u} (when \mu \neq 0 \mu \le 0)$

Stability of Fixed Points (Linear Stability Analysis)

Let:

$$xn=x*+\epsilon nx_n = x^* + \epsilon n_n$$

Apply Taylor expansion and linearize:

 $\epsilon n+1\approx f'(x*)\cdot \epsilon n \Rightarrow |\epsilon n+1||\epsilon n|=|f'(x*)|\cdot \epsilon n+1 \Rightarrow |\epsilon n+1||\epsilon n|=|f'(x*)|\cdot \epsilon n+1|$ \approx f'(x^*) \cdot \epsilon_n \Rightarrow \frac{|\epsilon_{n+1}|}{|\epsilon_n|} = |f'(x^*)|

- Stable: if $|f'(x^*)| < 1 |f'(x^{*})| < 1$
- Unstable: if |f'(x*)|>1|f'(x^*)| > 1

Derivatives and Stability

 $f(x)=\mu x(1-x) \Rightarrow f'(x)=\mu (1-2x)f(x) = \mu (1-x) Rightarrow f'(x) = \mu (1-2x)$

1. For $x*=0x^*=0$:

 $f'(0)=\mu \Rightarrow Stable if \mu < 1f'(0) = \mu \times \{Stable if \} \setminus 1$

2. For $x*=1-1\mu x^* = 1 - \frac{1}{\mu u}$:

 $f'(x*)=\mu(1-2(1-1\mu))=2-\mu \Rightarrow Stable if 1<\mu<3f'(x^*) = \mu \left(1-2(1-1\mu)\right)=2-\mu \Rightarrow Stable if 1<\mu<3f'(x^*) = \mu \left(1-2(1-1\mu)\right)=2-\mu \Rightarrow Stable if 1<\mu<3f'(x^*) = \mu(1-2(1-1\mu))=2-\mu \Rightarrow Stable if 1<\mu<3f'(x^*) = \mu(1-2\mu)$

Bifurcation and Chaos

When μ >3\mu > 3:

- The second fixed point becomes unstable.
- The system **bifurcates**: exhibits a **period-2 cycle**: alternates between two values.

Further increases in μ \mu:

- Period-4 cycle
- Period-8 cycle
- Eventually, the system becomes **chaotic**: unpredictable and sensitive to initial conditions.

This phenomenon is called a **period-doubling route to chaos**.

Bifurcation Diagram

- Horizontal axis: μ\mu (0 to 4)
- Vertical axis: Long-term values of xnx_n
- Shows the transition from:
 - Stability (fixed point)
 - to periodicity (cycles)
 - to chaos

A Your Task

Compute and **plot** the bifurcation diagram of the logistic map:

- 1. Vary μ \mu from 0 to 4 (e.g. in small steps like 0.001)
- 2. For each μ \mu, iterate the logistic map (e.g., 1000 iterations)
- 3. Discard the first few hundred iterations (transients)
- 4. Plot the remaining xnx_n values vs μ\mu

Would you like a sample MATLAB code to start with for plotting the bifurcation diagram?

Thanks for sharing the project instructions. Here's a **clear summary** of what you're being asked to do and **how to structure your code** for computing the **bifurcation diagram of the logistic map**:

Objective:

Compute and visualize the **bifurcation diagram** of the logistic map:

$$xn+1=\mu xn(1-xn)x_{n+1} = \mu x_n (1 - x_n)$$

For values of μ \mu from **2.4 to 4**.

Three-Loop Structure:

V Outer Loop (Loop 1): Loop over μ\mu

- Range: μ∈[2.4,4.0]\mu \in [2.4, 4.0]
- Use linspace(2.4, 4.0, num_mu) to define the resolution in μ \mu

• Set x = 0.5 (initial condition) at the start of each μ \mu

Middle Loop (Loop 2): Transient iterations

- Iterate the logistic map for num_transient steps
- Purpose: discard early iterations to let system settle

Inner Loop (Loop 3): Data collection

- Iterate num_data more steps
- Store x values for plotting (or binning if using density plot)

Parameters to Set:

- 1. num_mu how many μ\mu values you simulate (e.g. 1000)
- 2. num_transient number of transient iterations to discard (e.g. 100)
- 3. num_data number of iterations to record (e.g. 100)
- 4. x0 = 0.5 initial value of xx
- 5. (Optional for pixel plotting) num_bins_x number of bins along the x-axis

Plotting (Two Options):

1. Simple Scatter Plot (Good for learning)

- Just plot x values for each μ\mu
- In MATLAB:

plot(mu_val * ones(1, num_data), x_data, '.k', 'MarkerSize', 1)

2. Density Plot (Better visualization)

- Bin the values of xx at each μ\mu
- Count how many values fall in each bin
- Shade pixels according to density
- Use provided code to generate the density image

✓ Minimal MATLAB Code Skeleton:

```
mu_vals = linspace(2.4, 4.0, 1000);
x0 = 0.5;
num_transient = 100;
num_data = 100;
```

```
all_mu = [];
all_x = [];
for mu = mu_vals
  x = x0;
  % Loop 2: Transient iterations
  for i = 1:num_transient
    x = mu * x * (1 - x);
  end
  % Loop 3: Data collection
  for i = 1:num_data
    x = mu * x * (1 - x);
    all_mu(end+1) = mu; %#ok<SAGROW>
    all_x(end+1) = x; %#ok<SAGROW>
  end
end
plot(all_mu, all_x, '.k', 'MarkerSize', 1);
xlabel('\mu');
ylabel('x');
title('Bifurcation diagram of the logistic map');
```

Summary of What You Learned:

- The logistic map has two fixed points: x=0x = 0 and $x=1-1\mu x = 1 \frac{1}{\mu x}$
- Stability of fixed points is determined by $|f'(x^*)| < 1|f'(x^{n^*})| < 1$
- As μ\mu increases:
 - Stable fixed point \rightarrow Period 2 \rightarrow Period 4 \rightarrow Chaos
- Bifurcation diagram shows how the long-term behavior of xnx_n changes with μ \mu

If you'd like, I can help you:

- Translate this to Python code instead of MATLAB
- Add the density plot version
- Add a color-coded image version for clarity

Let me know how you'd like to proceed!

Thanks for sharing the Week 2 overview of **Numerical Methods for Engineers** — this week focuses on **root-finding**, a fundamental part of computational mathematics and engineering. Here's a structured summary with key concepts, algorithms, and what to expect in your second project:

Root-Finding Overview

Root-finding is about solving equations of the form:

$$f(x)=0f(x)=0$$

Many numerical problems (like optimization, nonlinear equations, or physical simulations) can be **reformulated** into root-finding problems.

Three Root-Finding Algorithms:

1. Bisection Method

- Guaranteed to converge
- X Slow (linear convergence)
- Requires that f(a)f(a) and f(b)f(b) have **opposite signs** (i.e., the root is bracketed)
- Procedure:
 - o Compute midpoint: c=a+b2c = \frac{a + b}{2}
 - Replace aa or bb based on the sign of f(c)f(c)

2. Newton's Method

- **\$\rightarrow\$** Fast (quadratic convergence)
- Requires:
 - A good initial guess x0x 0
 - The derivative f'(x)f'(x)
- Iteration formula:

```
xn+1=xn-f(xn)f'(xn)x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}
```

• Can fail if derivative is 0 or initial guess is poor

3. Secant Method

- Section almost as fast as Newton
- Doesn't require f'(x)f'(x)
- Uses two initial guesses: x0x_0, x1x_1
- Iteration formula:

```
xn+1=xn-f(xn)\cdot xn-xn-1f(xn)-f(xn-1)x_{n+1} = x_n - f(x_n)\cdot f(x_{n-1})
```

MATLAB Root-Finding Tools:

roots(poly_coeffs)

- Computes all roots of a polynomial
- Input: coefficient vector of a polynomial
- roots([1 -3 2]) % Solves $x^2 3x + 2 = 0$

fzero(function_handle, initial_guess)

- Finds a single root of a **nonlinear** function
- $f = @(x) x^3 x 2;$
- root = fzero(f, 1.5);

Week 2 Project: Compute the Feigenbaum Delta

- You'll use **Newton's method** to compute the **Feigenbaum delta**, a universal constant in chaos theory.
- It's related to the geometric convergence of bifurcation points in the logistic map.
- You'll likely:
 - o Write Newton's method code in MATLAB
 - o Apply it to a recurrence or fixed-point equation derived from the bifurcation points

Feigenbaum Delta (Quick Insight):

• Describes the rate at which period-doubling bifurcations accumulate in maps like:

$$xn+1=\mu xn(1-xn)x_{n+1} = \mu x_n (1 - x_n)$$

Delta value:

✓ Summary of Skills You'll Gain in Week 2:

- Implement bisection, Newton, and secant methods
- Understand convergence rates and tradeoffs
- Use MATLAB built-in functions for root-finding
- Apply Newton's method to chaos theory and discover Feigenbaum constants

Would you like help implementing any of the methods (e.g., Newton's method) in MATLAB for your project?

Great! Here's a clear and structured summary of **the Bisection Method** as explained by Jeff Chasnov in the video:

Bisection Method - Root Finding

6 Goal

Solve the equation:

$$f(x)=0f(x)=0$$

Find the root x=rx=r, where the function crosses the x-axis (i.e., where y=f(x)=0y=f(x)=0).

📌 Key Idea

- You need two initial guesses: x0x_0 and x1x_1
- These two values must bracket the root, i.e.:

 $f(x0)\cdot f(x1) < 0f(x_0) \cdot cdot f(x_1) < 0$

(i.e., the signs of $f(x0)f(x_0)$ and $f(x1)f(x_1)$ must be opposite)

Iteration Process

1. Compute the midpoint:

 $xmid=x0+x12x_{\text{mid}} = \frac{x_0 + x_1}{2}$

- Evaluate f(xmid)f(x_{\text{mid}})
- 3. Replace either x0x_0 or x1x_1 with xmidx_{\text{mid}} so that the root remains bracketed
- 4. Repeat until:

 $|x_1-x_0| < (desired tolerance) | x_1 - x_0| < (desired tolerance) |$

Example: Find 2\sqrt{2}

Let $f(x)=x^2-2f(x)=x^2-2$, so the root is $2\approx 1.41421 \setminus \{2\} \setminus \{2\}$

- Initial guesses: x0=1x 0 = 1, x1=2x 1 = 2
- Midpoint $x2=(1+2)/2=1.5x_2=(1+2)/2=1.5$, $f(1.5)=0.25f(1.5)=0.25 \rightarrow root$ is between 1 and 1.5
- $x3=(1+1.5)/2=1.25x_3=(1+1.5)/2=1.25$, $f(1.25)=-0.4375f(1.25)=-0.4375 \rightarrow$ root is between 1.25 and 1.5
- $x4=(1.25+1.5)/2=1.375x_4=(1.25+1.5)/2=1.375$, f(1.375)=-0.109f(1.375)=-0.109
- $x5=(1.375+1.5)/2=1.4375x_5=(1.375+1.5)/2=1.4375$, etc.

As you keep iterating, you converge slowly toward 2\sqrt{2}.

Pros and Cons

✓ Pros **X** Cons

Guaranteed to converge Slow convergence (linear rate)

Simple to implement Not ideal for high-speed applications

Doesn't require derivative

Summary

- Bisection method uses interval halving to converge to a root
- Always preserves bracketing
- Converges linearly: interval size halves each iteration
- Very robust, though slow

Would you like a simple MATLAB or Python code example for the Bisection method?

Here's a clear and complete **summary of Newton's Method** for root finding, as explained by Jeff Chasnov:

♦ Newton's Method – Root Finding Using Calculus



Solve:

$$f(x)=0f(x)=0$$

Find a root x=rx = r, where the function intersects the x-axis.



***** Key Idea

Use the tangent line at a current point xnx_n to approximate the function and find the next guess $xn+1x_{n+1}$.

This uses **derivatives** (calculus) to get rapid convergence.

Derivation

From the tangent line at xnx n:

$$y-f(xn)=f'(xn)(x-xn)y - f(x n) = f'(x n)(x-x n)$$

Set y=0y=0 (since we want the root), and solve for xx:

$$0-f(xn)=f'(xn)(xn+1-xn)0-f(x_n)=f'(x_n)(x_{n+1}-x_n) \ xn+1=xn-f(xn)f'(xn)x_{n+1}=x_n-f($$

Newton's Iteration Formula

 $xn+1=xn-f(xn)f'(xn)x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$

Requirements

- A function f(x)f(x)
- Its derivative f'(x)f'(x)
- A good initial guess x0x_0 close to the root (If too far, Newton's method can diverge.)



Example: Find 2\sqrt{2}

Let:

$$f(x)=x2-2, f'(x)=2xf(x)=x^2-2, \quad f'(x)=2x$$

Iteration Formula:

$$xn+1=xn-xn2-22xn=xn2+22xnx_{n+1} = x_n - \frac{x_n^2 - 2}{2x_n} = \frac{x_n^2 + 2}{2x_n}$$

Start with x0=1x_0 = 1:

• $x1=12+22\cdot1=32=1.5x 1 = \frac{1^2 + 2}{2 \cdot 1} = \frac{3}{2} = 1.5$

- x2=(3/2)2+22·3/2=1712≈1.416x_2 = \frac{(3/2)^2 + 2}{2 \cdot 3/2} = \frac{17}{12} \approx 1.416
- x3=(17/12)2+22·17/12=577408≈1.41426x_3 = \frac{(17/12)^2 + 2}{2 \cdot 17/12} = \frac{577}{408} \approx 1.41426
- ✓ After only **3 iterations**, the approximation is very close to:

2×1.41421\sqrt{2}\approx 1.41421

Pros and Cons

✓ Pros X Cons

Very fast convergence (quadratic) Requires derivative f'(x)f'(x)

Simple and elegant iteration Can diverge with bad initial guess

Ideal for problems needing speed Doesn't always converge globally

Summary

- Newton's method uses the slope of the tangent line to find the next approximation
- Formula:

 $xn+1=xn-f(xn)f'(xn)x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$

- Fastest converging method if the initial guess is good
- Widely used in numerical applications requiring high precision quickly

Would you like a MATLAB or Python implementation for Newton's method as well?

Here's a **clear summary of the Secant Method** for root finding, as explained by Jeff Chasnov:

Secant Method – Root Finding Without Derivatives

6 Goal

Solve:

f(x) = Of(x) = 0

Find a root (x-intercept), like with Newton's method, but without needing the derivative f'(x)f'(x).



Use the **Secant Method** when:

- You want fast convergence.
- You cannot compute the derivative f'(x)f'(x) analytically.
- You can evaluate f(x)f(x) at any point.

Idea

Instead of the tangent line (Newton), use a secant line — a line passing through two previous points on the curve.

Approximate the derivative:

 $f'(xn) \approx f(xn) - f(xn-1)xn - xn-1f'(x_n) \cdot f(x_n) - f(x_{n-1}) \times x_n - x_{n-1}$

Plug into Newton's formula:

 $xn+1=xn-f(xn)f(xn)-f(xn-1)xn-xn-1x_{n+1} = x_n - \frac{f(x_n)}{\frac{f(x_n)}{\frac{f(x_n)-f(xn-1)}}}x_n - \frac{f(x_n-1)}{\frac{f(x_n-1)}}$ x_{n-1}}}

Simplifies to:

 $xn+1=xn-f(xn)\cdot xn-xn-1f(xn)-f(xn-1)x_{n+1} = x_n - f(x_n) \cdot f(x_n) - x_{n-1} = x_n - f(x_n) \cdot f(x_n) - x_n - x_n$ $f(x_{n-1})$



Example: Approximating 2\sqrt{2}

Let:

- $f(x)=x^2-2f(x)=x^2-2$
- The root is 2≈1.41421\sqrt{2}\approx 1.41421

Start with:

- x0=1x 0=1
- x1=2x 1=2

Iterations:

1. X₂:

 $x2=x1\cdot x0+2x1+x0=2\cdot 1+22+1=43\approx 1.333...x_2 = \frac{x_1 \cdot x_0+2}{x_1+x_0} = \frac{2 \cdot x_1}{x_1+x_0} = \frac{2 \cdot x_0+2x_1+x_0}{x_1+x_0} = \frac{2 \cdot x_0+x_0}{x_1+x_0} = \frac{2 \cdot x_0+x_$ 1 + 2{2 + 1} = \frac{4}{3} \approx 1.333...

2. x₃:

 $x3=(4/3)\cdot 2+24/3+2=8/3+210/3=14/310/3=1410=1.4x_3 = \frac{(4/3) \cdot 2+24/3+2=8/3+210/3=14/310/3=1410=1.4x_3 = \frac{(4/3) \cdot 2+24/3+2=8/3+210/3=14/310/3=1410=1.4x_3 = \frac{(4/3) \cdot 2+24/3+2=8/3+210/3=14/310/3=1410=1.4x_3}{(4/3) \cdot 2+24/3+2=8/3+210/3=14/310/3=1410=1.4x_3}$ $\frac{8/3 + 2}{10/3} = \frac{14/3}{10/3} = \frac{14}{10} = 1.4$

3. **x**₄:

 $x4=(7/5)\cdot(4/3)+27/5+4/3=5841\approx1.41463x_4=\frac{(7/5)\cdot(4/3)+2}{7/5+4/3}=$ \frac{58}{41} \approx 1.41463

Very close to 2≈1.41421\sqrt{2} \approx 1.41421 in just 3 steps!

Comparison of Methods

Method	Requires Derivative?	Initial Guesses	Convergence Speed	Always Converges?
Bisection	X No	2 (bracketing)	Slow (linear)	✓ Yes
Newton's	✓ Yes	1	Very Fast (quadratic)	X No
Secant	X No	2 (not bracketed)	Fast (super- linear)	X No

Summary

- Secant Method is a **derivative-free alternative** to Newton's method.
- Uses two points to approximate the slope (secant line).
- Formula:

 $xn+1=xn-f(xn)\cdot xn-xn-1f(xn)-f(xn-1)x_{n+1} = x_n - f(x_n) \cdot f(x_n - x_{n-1}) f(x_n) - f(x_n - x_n - x_n$ $f(x_{n-1})$

- Converges faster than bisection, but not as fast as Newton.
- Needs 2 initial guesses, but can be very effective in practice.

Would you like a side-by-side code implementation for Newton's vs Secant in Python or MATLAB?

Here's a **summary of the concept of Order of Convergence**, as explained by Jeff Chasnov:



Order of Convergence – How Fast Do Root-Finding Methods Converge?



To quantify how fast a root-finding method converges to the actual root rr of a function f(x)=0f(x)=0.



Error Definition

Let:

- xnx_n be the nth approximation
- rr be the actual root

Then the **error** at step nn is:

 $\varepsilon n = |xn-r| \vee r = |x_n - r|$

Order of Convergence (p)

We say a method has order of convergence pp if:

εn+1≈C·εnp\varepsilon_{n+1} \approx C \cdot \varepsilon_n^p

Where:

- CC is a constant,
- εn→0\varepsilon_n \to 0 as n→∞n \to \infty

Interpretation:

- If p=1p = 1: Linear convergence (slow)
- If p=2p = 2: Quadratic convergence (fast)
- If p>1p > 1: Superlinear convergence (very good)

Orders of Common Methods

Method	Convergence Order pp	Description

Bisection p=1p=1 **Linear** (slowest)

Secant p≈1.618p \approx 1.618 Superlinear (faster than linear)

Newton's p=2p=2 **Quadratic** (fastest)

Example: Bisection Method

Each iteration **halves** the interval, so:

 $\varepsilon n+1=12\varepsilon n \Rightarrow p=1 \vee n-1 = \frac{1}{2} \vee n-1 = 1$

Key Takeaway

- Higher p → faster convergence.
- **Bisection** always converges, but slowly.

- **Newton's method** is fastest, but needs derivative and a good starting guess.
- **Secant** is a great compromise: faster than bisection, doesn't need derivative.

Would you like to see this visualized in a graph or demonstrated with a small code simulation comparing convergence rates?

Here's a step-by-step summary of how the order of convergence of Newton's Method is derived, as explained by Jeff Chasnov:

Newton's Method: Order of Convergence

We aim to prove that Newton's method has quadratic convergence — i.e., order p=2p = 2 under suitable conditions.

Step 1: Newton's Iteration

The iteration formula is:

$$xn+1=xn-f(xn)f'(xn)x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

Let rr be the **root**, such that f(r)=0 f(r) = 0, and define the **error** at step nn:

$$\varepsilon n = xn - r \cdot varepsilon_n = x_n - r$$

So we want to express $\epsilon n+1$ \varepsilon_{n+1} in terms of ϵn \varepsilon_n.

Step 2: Error Form of the Iteration

Start from the Newton update:

```
xn+1=xn-f(xn)f'(xn)x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}
```

Then:

Now express $f(xn)f(x_n)$ and $f'(xn)f'(x_n)$ as **Taylor expansions around rr**, because $\varepsilon n = x - r \cdot varepsilon_n = x_n - r$ is small.

Step 3: Taylor Expansions

Since f(r)=0f(r)=0, expand $f(xn)=f(r+\epsilon n)f(x_n)=f(r+ \sqrt{x^2-1})$:

Numerator:

```
f(xn)=f(r+\epsilon n)=f'(r)\epsilon n+12f''(r)\epsilon n2+\cdots f(x_n)=f(r+varepsilon_n)=f'(r)varepsilon_n+12f''(r)\epsilon n2+\cdots f(x_n)=f(r+\epsilon n)=f'(r)\epsilon n2+\cdots f(x_n)=f'(r)\epsilon n2+\cdots f'(r)\epsilon 
\frac{1}{2}f''(r)\varepsilon_n^2 + \cdot cdots
```

Denominator:

 $f'(xn)=f'(r+\epsilon n)=f'(r)+f''(r)\epsilon n+\cdots f'(x_n)=f'(r+varepsilon_n)=f'(r)+f''(r)varepsilon_n+varepsilon_n$

Step 4: Plug into Error Formula

Divide numerator and denominator by f'(r)f'(r) to simplify:

Now use the Taylor expansion of $11+\delta \approx 1-\delta+\delta 2-\cdots \text{ }1$ \delta\\2 - \cdots

So:

Step 5: Multiply and Simplify

After careful expansion and dropping higher-order terms:

 $\epsilon n+1\approx 12f''(r)f'(r)\epsilon n2\varepsilon_{n+1} \approx \frac{1}{2}\frac{f''(r)}{f'(r)}\varepsilon_{n^2}$

✓ Conclusion: Quadratic Convergence

|\en+1|≈K|\en|2|\varepsilon_{n+1}| \approx K |\varepsilon_n|^2

Where $K=|12 \cdot f''(r)f'(r)|K = \left| \frac{1}{2} \cdot \frac{f''(r)}{f'(r)}\right|$

Thus, Newton's method has order of convergence p=2p = 2 when:

- The root rr is **simple** (i.e., $f'(r) \neq 0f'(r) \setminus neq 0$)
- The function is sufficiently smooth (differentiable)

⚠ Important Notes

- If f'(r)=0f'(r)=0, convergence is not quadratic it may be slower or fail.
- Newton's method is very sensitive to the initial guess.

Would you like a numerical example or Python code to visualize the convergence rate?

Here's a concise summary of the second video explaining how to generate a fractal using Newton's Method for complex cube roots of unity:

Mathematical Setup

• The equation to solve is:

$$f(z)=z3-1=0$$
 $f(z)=z^3-1=0$

- The solutions (cube roots of unity) are:
 - o r1=1r_1 = 1
 - \circ r2=-12+i32r_2 = -\frac{1}{2} + i \frac{\sqrt{3}}{2}
 - o $r3=-12-i32r_3 = -\frac{1}{2} i \frac{3}{2}$

These lie on the **unit circle** in the complex plane, evenly spaced at 120°.

Newton's Method

To find roots of f(z)f(z), use the iterative formula:

 $zn+1=zn-f(zn)f'(zn)=zn-zn3-13zn2z_{n+1}=z_n- \\ f'(z_n)_{f'(z_n)}=z_n- \\ f'(z_n)_{f'(z_n)}=z_n$

Generating the Fractal

- 1. **Grid the complex plane** Treat each complex number z0z_0 as a pixel.
- 2. For each pixel:
 - Apply Newton's method starting from z0z_0.
 - See which root it converges to (i.e., whether final value is close to r1r_1, r2r_2, or r3r_3).
 - o Color the pixel:
 - Red if it converges to r1r_1
 - Green for r2r_2
 - Blue for r3r_3
- 3. **Borderline pixels**, where it's not clear which root they'll converge to, result in extremely intricate patterns **this is where the fractal appears**.

- In regions near the **boundaries between basins of attraction**, Newton's method becomes sensitive to small changes in the initial value.
- This sensitivity produces self-similar structures as you zoom in a defining property of fractals.

What's Next

In the next video (mentioned by Jeff Chasnoff), he'll show how to implement this fractal visualization using **MATLAB**.

If you'd like, I can help you implement this in **Python (e.g., with NumPy and Matplotlib)** or explain the MATLAB version.

Sure! Here's the English translation of the video transcript that explains how to code the Newton Fractal step-by-step:

[MUSIC]

In this video, I'll show you how to code the Newton fractal.

We want to apply Newton's method to find the cube roots of unity.

So, we need to find the zeros of the function:

$$f(z) = z^3 - 1.$$

We scan over the complex plane and iterate each grid point using Newton's method:

$$z_{n+1} = z_n - f(z_n)/f'(z_n)$$

Here, f'(z) is just the derivative of f(z), which is $3z^2$.

At **0:42**:

Since Newton's method will be the core of our code, we'll start with it.

0:52:

Here is Newton's iteration. **Z** will be a matrix. The initial values of Z are points on the complex plane where we want to compute the fractal.

At the end of the iteration, each element of Z should have converged to one of the three cube roots of unity.

1:15:

We will iterate Newton's method 40 times — it's fast and that should be enough for all starting points to converge.

1:37:

Now let's begin coding from the start. We want to set all initial values in Z and define the necessary functions.

1:53:

Our function f is: $f(Z) = Z^3 - 1$. Since Z is a matrix, we need to use the dot operator to cube element-wise.

Its derivative, \mathbf{f}' , is defined as 3 * Z^2 . So that defines the two functions we'll use in Newton's method.

2:31:

We might as well define the cube roots at the beginning of the code.

2.36

The cube roots of unity are:

- root1 = 1,
- $root2 = -1/2 + i\sqrt{3}/2$,
- $root3 = -1/2 i\sqrt{3}/2$.

When I write the imaginary unit i, I prefer to use 1i because I sometimes use i as a variable or loop index in code.

So using 1i avoids bugs.

3:17:

Next, I want to set up the grid.

3:23:

We start with 1000 points in the **x-direction** (real axis), and 1000 in the **y-direction** (imaginary axis).

3:35:

We're visualizing a fractal from -2 to 2 on both the real and imaginary axes — this gives a nice picture.

3:58:

Use linspace to divide xmin to xmax into nx points (and similarly for y).

4:15:

Create the grid using meshgrid(X, Y) — it returns two matrices:

- X holds the x-values of the grid
- Y holds the y-values of the grid

4:47:

What we need is the complex grid $\mathbf{Z} = \mathbf{X} + \mathbf{i}\mathbf{Y}$. This gives the initial values for Newton's method.

5:29

We now run Newton's method — it's very fast, as it processes all grid values simultaneously.

5:43:

Next, we determine which root each grid point converged to.

5:49:

We define a small threshold, **epsilon = 0.001**, and then define four logical matrices: **Z1, Z2, Z3, Z4**.

- **Z1** will be true (1) where abs(Z root1) < epsilon
- So if Z converged to root1, this condition is true, else false (0)

6:24:

So **Z1** is a matrix of ones where the starting points converged to **root1**, and zeros elsewhere.

7:19:

Similarly:

- **Z2** = 1 where values converged to **root2**
- **Z3** = 1 where values converged to **root3**

7:57

Z4 = not (Z1 + Z2 + Z3) — so it indicates non-convergence.

8:00+:

Now, we are ready to plot the result.

We open a figure and define a color map:

- First row: how to color value 1
- Second: how to color **2**, and so on...
- [1, 0, 0] = red
- [0, 1, 0] = green
- [0, 0, 1] = blue
- [0, 0, 0] = black (for non-convergence)

Then set the colormap of the figure.

We now create the final image data:

Z = Z1 + 2Z2 + 3Z3 + 4*Z4

This results in a matrix of 1s (red), 2s (green), 3s (blue), and 4s (black).

Finally, plot the image using the **image** command, setting axes and labels.

Note:

- Images are by default shown top-down (like in image processing)
- Numerical plots usually have Y increasing upward

So, to match expectations, we use set(gca, 'YDir', 'normal')

12:56:

When writing your code, use MATLAB's debugging tools.

Set a smaller grid (like 10x10) for easier debugging.

Place breakpoints after assigning Z, after Newton's method, and before image creation.

16:16:

Now Z should contain only 1, 2, 3, or 4 (depending on which root it converged to or if it didn't).

17:03:

MATLAB also has a **profiler** to measure how long each line of code takes.

If we use a larger grid (e.g., 8000x8000), we can see that most of the runtime is spent in Newton's iteration — as expected.

18:13

For example, 33 out of 35 seconds were spent on the line:

$$Z = Z - f(Z)./fp(Z);$$

This confirms that Newton's method is the "engine" of the fractal code.

Let me know if you'd like the full MATLAB code version in English or explained line-by-line!

Here's the **full English transcript** of the video, cleaned up and organized for clarity:

Introduction

0:03

In this video, let's learn about the **root-finding functions** provided by **MATLAB**.

Most of the time, you don't need to write your own root-finding algorithm, like we did when drawing the Newton fractal.

Instead, you can use MATLAB's built-in functions.

There are two main root-finding functions in MATLAB:

- 1. roots finds all real and complex roots of a polynomial.
- 2. fzero finds one real root of a non-linear function.

Using the roots Function

To use roots, you just give it the **coefficients of the polynomial**.

Suppose we have a polynomial of the form:

$$p_1x^n + p_2x^{n-1} + ... + p_{n+1} = 0$$

This polynomial has n+1 coefficients.

You must write the coefficients in descending order of powers, and include zeros for any missing powers.

Syntax:

r = roots(p)

Here, p is a vector of coefficients, and r will be a column vector of **complex roots**.

Example

Let's solve the polynomial:

$$x^3 - 3x^2 + 4x - 2 = 0$$

We define the coefficients in MATLAB:

```
p = [1 -3 4 -2];
```

r = roots(p);

Run this in the Live Editor.

The output is:

r =

1.0000 + 0.0000i

0.0000 + 1.0000i

0.0000 - 1.0000i

So the roots are 1, i, and -i.

Note: MATLAB may show very small imaginary parts due to round-off errors.

We can verify that **1** is a root:

 $1^3 = 1$

 $-3 \times 1 = -3$

 $+4 \times 1 = 4$

-2 = -2

So: 1 - 3 + 4 - 2 = 0



Using the fzero Function

fzero finds a **single real root** of a non-linear function:

$$f(x) = 0$$

Function Definition

You need to define f(x) as either:

- An anonymous function: f = @(x) ...
- Or a separate subfunction

If the function depends on a parameter (like a), you can define:

```
f = @(x,a) x - exp(-a*x);

To solve x = e^{(-a*x)}, rewrite it as:

f(x) = x - exp(-a \cdot x) = 0

Let's say a = 0.5, then:

a = 0.5;

f = @(x) x - exp(-a*x);
```

You need to give fzero:

- A function of one variable
- An initial guess or a bracketing interval

If your function has parameters, you can wrap it like this:

```
r = fzero(@(x) f(x, a), x0);
```

Example

```
a = 0.5;

f = @(x, a) \times - \exp(-a^*x);

x0 = 0; % initial guess

r = fzero(@(x) f(x, a), x0);

Running this, we get:

r = 0.7035
```

This is the solution to $x = e^{(-0.5x)}$.

Summary

8:41

We introduced two root-finding functions in MATLAB:

- roots.m finds all roots of a polynomial
- fzero.m finds one root of a non-linear function

"I've used fzero many times in my research. It's powerful and simple to use. You'll get a chance to try it out in course exercises."

– Jeff [last name not heard clearly]

Let me know if you'd like the code from these examples or want a practice problem based on this!

Great! Here's a concise **summary of the first lecture** on the theoretical background for computing the **Feigenbaum delta**:

Topic: Computing the Feigenbaum Delta

Context: Based on the logistic map:

$$xi+1=\mu xi(1-xi)x_{i+1} = \mu x_{i}(1-x_{i})$$

This map exhibits **bifurcations** leading to **chaos** as the parameter μ \mu increases.

Bifurcation Behavior

- At small μ \mu: the map converges to a **fixed point** (period-1).
- As μ\mu increases:
 - Period-2 cycle emerges at μ0\mu_0
 - Period-4 at μ1\mu_1
 - Period-8 at μ2\mu_2, and so on...
- Eventually, the system exhibits **chaotic behavior** after infinite period-doubling.

Feigenbaum Delta Definition

This constant captures how the bifurcation points cluster closer as the period doubles:

 $\delta=\lim_n\to\infty\mu n-\mu n-1\mu n+1-\mu n \leq \lim_n \to \min_n - \mu n-1 \} \{ \mu_n = \lim_n \to \mu_n-1 \}$

Problem with Direct Computation

- Near bifurcation points, the **transient** (settling time) is very long.
- Requires many iterations to accurately compute bifurcation values μη\mu n.

✓ Better Approach: Superstable Cycles

Instead of bifurcation points, we compute at superstable cycles, where:

- The orbit includes x=12x = \frac{1}{2}
- These values of μ\mu are denoted mnm_n
- Superstable cycles **converge fast**, so they're more computationally efficient.

Examples & Definitions

- $m0=2m_0 = 2$: because x=0.5x = 0.5 is a **fixed point** when $\mu=2\mu=2$
- m1m_1: for period-2 superstable cycle
 - Derived from solving a cubic:

```
\mu 3-4\mu 2+8=0\mu^3 - 4\mu^2 + 8 = 0
```

o Solution:

 $m1=1+5m 1=1+\sqrt{5}$

★ Final Delta Definition (Practical Version)

This converges to the same Feigenbaum delta, but is faster to compute.

% What's Next

- In future lectures:
 - 1. How to compute mnm_n using Newton's method
 - 2. How to implement this in MATLAB

Let me know if you'd like a visual diagram or notes formatted for printing or LaTeX!

Here's a concise summary and explanation of the **algorithm** described in the video for computing the **Feigenbaum delta** (δ):

Goal

Compute the Feigenbaum delta:

 $\delta=\lim_n\to\infty mn-mn-1-mn-2\ delta=\lim_n \to \min_n-1-mn-2\ delta=\lim_n \to m_{n-1}=m_{n-2}$

Where mnm_n is the value of μ \mu such that starting from x0=12x_0 = \frac{1}{2}, the logistic map returns to 12\frac{1}{2} after 2n2^n iterations. That means x2n=12x_{2^n} = \frac{1}{2}.

Logistic Map

 $xi+1=\mu xi(1-xi)x_{i+1} = \mu x_{i}(1-x_{i})$

Step-by-Step Algorithm

1. Define Function for Root Finding

We want:

$$g(\mu)=xN(\mu)-12g(\mu) = x_N(\mu) - \frac{1}{2}$$

Where $N=2nN = 2^n$, and $x0=12x 0 = \frac{1}{2}$.

We seek roots of $g(\mu)g(\mu)$, i.e., values of μ such that $xN=12x_N = \frac{1}{2}$.

2. Use Newton's Method

Newton's iteration:

$$\mu j + 1 = \mu j - g(\mu j) g'(\mu j) = \mu j - \lambda frac\{g(\mu j)\}\{g'(\mu j)\} = \mu j - \lambda frac\{x_N - \lambda frac\{1\}\{2\}\{x_N'\}\}$$

So we need both xNx_N and xN'x_N' — the derivative of xNx_N w.r.t. μ \mu.

3. Compute Derivative Recursively

Let:

- $x0=12x_0 = \frac{1}{2}$
- x0'=0x 0'=0

Then iterate:

- $xi+1=\mu xi(1-xi)x_{i+1} = \mu x_i(1-x_i)$
- $xi+1'=xi(1-xi)+\mu xi'(1-2xi)x_{i+1}'=x_i(1-x_i)+\mu x_i'(1-2x_i)$

These two coupled equations must be computed together (simultaneous iteration) up to $N=2nN = 2^n$.

4. Initial Guess for Newton's Method

We don't know mnm_n yet, so we **guess** it based on the **previous values** and an estimate of delta:

Initial guess for mn=mn-1+mn-1-mn-2 δ n-1\text{Initial guess for } m_n = m_{n-1} + \frac{m_{n-2}}{\left(\frac{n-1}{n-2}\right)}

Where:

$$\delta n-1=mn-1-mn-2mn-2-mn-3 \cdot \{n-1\} = \frac{n-1}{m_{n-2}} \cdot m_{n-2} \cdot m_{n-3}$$

5. Bootstrapping

To start computing $m2m_2$, you only have $m0=2m_0 = 2$ and $m1=1+5m_1 = 1 + \sqrt{5}$. You don't yet have a delta value.

So:

- Pick an initial delta estimate (say, 5.0).
- Use it to guess m2m 2
- Compute m2m_2 using Newton's method
- Then compute updated δ1\delta_1
- Use $\delta1\$ delta_1 to guess m3m_3, and so on.

Limitations

Eventually, round-off error in floating point arithmetic will prevent accurate computation of further mnm_n's — especially as N=2nN = 2^n becomes large.

✓ Final Summary

To compute Feigenbaum delta:

- Define a root-finding problem for $g(\mu)=xN-12g(\mu)=x$ N f(1){2}
- Use Newton's method:

 $\mu j+1=\mu j-xN-0.5xN' \mu_{j+1} = \mu_{j} - \frac{x_N - 0.5}{x_N'}$

- Simultaneously compute both xix_i and xi'x_i'
- Use previous mm's and delta estimates to generate initial guesses
- Iterate to get new mnm n's and better approximations to delta.

If you'd like, I can help you implement this algorithm in Python step-by-step.

Thanks for sharing the transcript. This final explanation from Prof. Jeff Chasnoff summarizes the implementation strategy for computing the Feigenbaum delta numerically using Newton's method on the logistic map. Here's a breakdown of the computational plan based on the lecture:

What You're Computing

- The **Feigenbaum delta** δ≈4.66920160910299...\delta \approx 4.66920160910299\ldots a universal constant characterizing the geometric rate of period-doubling bifurcations in certain dynamical systems.
- Specifically, you're computing a sequence of bifurcation points mnm_n, and using them to estimate:

Key Concepts

• Logistic Map:

 $xn+1=\mu xn(1-xn)x_{n+1} = \mu x_n (1-x_n)$

Bifurcation Point mnm n:

Value of μ \mu at which the system undergoes a period-doubling to 2n2^n-period.

• Newton's Method:

Solve $g(\mu)=x2n(\mu)-x2n-1(\mu)=0g(\mu)=x_{2^n}(\mu)-x_{2^n}(\mu)=0$ Newton step:

 $\mu new = \mu old - g(\mu)g'(\mu) \setminus \{ \text{text}\{new\} \} = \mu \{ \text{text}\{old\} \} - \{ \text{frac}\{g(\mu)\} \} \}$

Three Loops

1. Outer Loop (Loop 1): Iterate Over n

- For n=2n = 2 to max_n (e.g., 14–15)
- o Compute mnm_n
- Use prior mnm_n values and δ\delta to guess next mm

2. Middle Loop (Loop 2): Newton's Method Iteration

- o Until convergence (e.g., relative change < 10−1210^{-12})
- Each iteration:
 - Call logistic map function to compute x2nx {2^n}
 - Also compute dgdµ\frac{dg}{d\mu} (requires chain rule through logistic iterations)

3. Inner Loop (Loop 3): Logistic Map Iteration

- Run the map 2n2ⁿ times for given μ\mu
- Compute both xx and derivative x'x' via:

 $xi+1=\mu xi(1-xi)x_{i+1} = \mu x_i (1 - x_i) xi+1'=xi(1-xi)+\mu(1-2xi)xi'x'_{i+1} = x_i (1 - x_i) + \mu(1 - 2x_i)x$

Scode Structure Sketch (Pseudocode/Algorithm)

% Initialization

m = zeros(1, max_n); % store bifurcation points

delta = zeros(1, max_n); % store Feigenbaum delta estimates

```
m(1) = known_m1;
m(2) = known_m2;
for n = 2:max_n-1
  % Initial guess for m_n based on previous deltas
  if n == 2
    mu = m(n) + 1e-4; % crude guess
  else
    mu = m(n) + (m(n) - m(n-1)) / delta(n-1);
  end
  % Newton's method to refine mu
  for newton_iter = 1:max_iter
    [g, dg] = compute_g_and_derivative(mu, n); \% g = f^2^n(x) - x
    mu_new = mu - g / dg;
    if abs(mu_new - mu) < 1e-12
      break;
    end
    mu = mu_new;
  end
  m(n+1) = mu;
  % Compute delta
  delta(n) = (m(n) - m(n-1)) / (m(n+1) - m(n));
end
```

What You Should Implement

- A logistic_map(mu, n) function that returns x2nx_{2^n}
- A derivative computation: $dx2nd\mu \frac{x_{2^n}}{d \mu}$

• Newton's method step using both

Final Advice

- Save all computed mnm_n values and delta estimates.
- Plot or print the convergence of delta it should stabilize around 4.669.
- Double precision only gives you around **15–16 digits**, so don't expect more accuracy than that.

Would you like help implementing any specific part of this in MATLAB?

Here's a concise summary of what you'll learn in **Week 3** of *Numerical Methods for Engineers*, based on the transcript:

Week 3 Overview: Numerical Linear Algebra

***** Topics Covered:

- 1. Systems of Linear Equations
 - o Gaussian Elimination with Partial Pivoting
 - o LU Decomposition
 - Counting the number of floating-point operations (flops)
- 2. Eigenvalues and Eigenvectors
 - Power Method for dominant eigenvalue/eigenvector computation
- 3. MATLAB Tools
 - o \ (backslash operator): Solves Ax=bAx = b
 - o lu(A): Computes LU decomposition
 - o eig(A): Computes eigenvalues and eigenvectors
- 4. Systems of Nonlinear Equations
 - Use of Newton's Method with matrix algebra



- Apply Newton's Method to compute fractals from solutions of the Lorenz equations
- Lorenz system is a classic example of a nonlinear system that exhibits chaotic behavior

Would you like help getting started with any specific part (e.g. implementing LU decomposition in MATLAB, or applying the power method)?

Here's a summary and key takeaways from this lecture on Gaussian Elimination and Round-Off Error:



Lecture Summary: Gaussian Elimination Without Pivoting

What is Gaussian Elimination?

- A method for solving systems of linear equations.
- Involves transforming the system's matrix into an upper triangular form using row operations.
- Solving then proceeds via back substitution.



↑ The Issue with Large Systems

Gaussian elimination works well on small systems (e.g., 3×3), but breaks down for large systems (e.g., 1000×1000) due to round-off error.

P Understanding Round-Off Error

- **Machine epsilon** (ϵ) is the smallest number that distinguishes two floating-point numbers.
 - o In double precision:

 $\varepsilon=2-52\approx2.2\times10-16$ \varepsilon = 2^{-52} \approx 2.2 \times 10^{-16}

- Example:
 - $2+\varepsilon=22 + \text{varepsilon} = 2$,
 - \circ 4- ε =44 \varepsilon = 4, due to limits in floating-point representation.



Contrived 2×2 Example

System of Equations:

 $\{ \epsilon x_1 + 2x_2 = 4x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_2 = 1 \setminus \{ \epsilon x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_2 = 1 \setminus \{ \epsilon x_1 - x_2 = 1 \setminus \{$

Expected (approximate) solution:

- x2≈2x_2 \approx 2
- x1≈3x_1 \approx 3

Gaussian elimination result (due to round-off):

- x2=2x_2 = 2 (correct)
- x1=0x_1 = 0 (totally wrong)

> What went wrong?

Elimination step produced an expression like:

 $\epsilon x1=0$ \varepsilon $x_1=0$

Should have been:

 $\epsilon x_1= very small number \Rightarrow x_1 \approx 3 varepsilon x_1 = \text{very small number} \$ Rightarrow x_1 \approx 3

But due to round-off, it became 0, giving a catastrophically incorrect result.

Rey Insight:

- If the pivot element (diagonal entry) is very small but non-zero, and we don't swap rows, we lose significant digits.
- In large matrices, small round-off errors can accumulate → completely wrong answers.

Solution: Partial Pivoting

- Swap rows to bring the largest available element into the pivot position.
- This avoids dividing by small numbers and reduces round-off errors.

Coming Up:

The next video will explain **partial pivoting** in detail.

Let me know if you'd like help implementing Gaussian elimination with pivoting in MATLAB or Python.

Here's a summary and key points from the lecture on Gaussian Elimination with Partial **Pivoting:**



Lecture Summary: Gaussian Elimination with Partial Pivoting

Recap:

In the last video, we saw that applying Gaussian elimination without pivoting can lead to catastrophic errors due to round-off error—especially in large systems.

% What is Partial Pivoting?

Partial Pivoting = Row Swapping Only

- At each step of elimination:
 - Look down the current column.
 - Identify the entry with the largest magnitude.
 - Swap that row to the top (pivot position).
- Full Pivoting = Row and column interchanges
 - o Rarely used in practice; partial pivoting is sufficient for stability in most cases.



Example: Solving the System with Pivoting

Given system:

 $\{ \epsilon x_1 + 2x_2 = 4x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_2 = 4 \setminus \{ \epsilon x_1 - x_2 = 1 \setminus \{ \epsilon x_1 + 2x_2 = 4 \setminus x_2 = 4 \setminus \{ \epsilon x_1 - x_2 = 1 \setminus \{ \epsilon x_1 - x_$

Augmented matrix:

[ϵ 2|41–1|1]\begin{bmatrix} \varepsilon & 2 & | & 4 \\ 1 & -1 & | & 1 \end{bmatrix}

Step 1: Pivoting

- Compare first column: ɛ\varepsilon vs. 11
- Swap rows: put **row with 1** at the top

[1-1|1\(\epsilon\) | 2 & | & 4 \end{bmatrix} 1 & -1 & | & 1 \\ \varepsilon & 2 & | & 4 \end{bmatrix}

Step 2: Elimination

- Eliminate ε \varepsilon below the pivot (first column):
 - ο Multiply first row by -ε-\varepsilon, add to second row:

Second row becomes:

- $\varepsilon \varepsilon = 0 \setminus \text{varepsilon} = 0$
- 2+ε2 + \varepsilon
- 4-ε4 \varepsilon

 \Rightarrow [1-1|102+ ϵ |4- ϵ]\Rightarrow\begin{bmatrix} 1 & -1 & | & 1 \\ 0 & 2 + \varepsilon & | & 4 -\varepsilon \end{bmatrix}

Due to round-off:

- 2+ε≈22 + \varepsilon \approx 2
- 4-ε≈44 \varepsilon \approx 4



Back Substitution

1. Solve second equation:

 $2x2=4 \Rightarrow x2=22x_2=4 \land Rightarrow x_2=2$

2. Plug into first equation:

$$x1-x2=1 \Rightarrow x1=3x_1 - x_2 = 1 \setminus Rightarrow x_1 = 3$$

Correct result achieved, despite round-off approximations.

6 Key Insights

- Partial pivoting avoids dividing by very small numbers, which reduces round-off errors.
- Without it, even simple 2×2 problems can fail numerically.
- With it, large systems (1000+ equations) can be solved accurately and stably.
- This was a breakthrough in numerical methods, making large-scale computations practical.

Summary Formula:

At step kk of elimination:

1. Find i≥ki \ge k such that:

 $|A[i,k]|=\max_{i\in J}\geq k|A[j,k]||A[i,k]|=\max_{i\in J} |A[i,k]|$

- 2. Swap row ii and row kk
- 3. Proceed with Gaussian elimination as normal

Let me know if you want a code implementation of Gaussian elimination with partial pivoting (in Python or MATLAB), or help with LU decomposition, which builds on this method.

Here's a clear and structured summary of the lecture on LU Decomposition with Partial **Pivoting**, as explained by Prof. Jeff Chasnov:



Lecture Summary: LU Decomposition with Partial Pivoting

What is LU Decomposition?

• LU Decomposition factors a matrix A into:

A=LUA = LU

where:

L = Lower triangular matrix

- **U** = Upper triangular matrix
- **Purpose**: Makes solving Ax=bAx = b faster, especially when:
 - o The same matrix AA is used with multiple different right-hand sides bb
 - o Common in solving **PDEs** numerically

↑ Why Use Partial Pivoting?

- To avoid division by small numbers (which causes round-off errors)
- We perform **row interchanges** during Gaussian elimination

Step-by-Step: LU with Partial Pivoting

Step 1: Apply Partial Pivoting

- Start with a matrix AA, scan the first column.
- Find the entry with largest absolute value (e.g., 6)
- Swap rows if necessary record this as a **permutation matrix** P12P_{12}

✓ Step 2: Gaussian Elimination

• Eliminate below the pivot using **elementary row operations**, represented as matrices:

M1,M2,...M_1, M_2, \dots

- For example:
 - Use row 1 to eliminate entries in rows 2 and 3
 - o New entries calculated by subtracting a multiple of row 1 from other rows

✓ Step 3: Second Pivoting Step

- In the second column, apply partial pivoting again:
 - o Find largest magnitude entry in remaining submatrix
 - Swap rows (e.g., P23P_{23}) to avoid zero pivot

Step 4: Record Final Matrix

• After full Gaussian elimination, the resulting matrix is:

U=Upper Triangular MatrixU = \text{Upper Triangular Matrix}

Step 5: Constructing the L Matrix

- Combine the inverse of elementary matrices and account for row swaps
- Trick: since PP matrices are **their own inverses** (P-1=PP^{-1} = P):

 $A=P-1LUA=P^{-1}LU$

• Rearranged:

PA=LU(standard form)PA = LU \quad \text{(standard form)}

• Example LL matrix (with multipliers used during elimination):

 $L=[1001210-1301]L = \begin{array}{c} L=[1001210-1301]L = \begin{array}{c} L=[1001210-1301]L$

• Example UU matrix from elimination:

 $U=[6-670-5120043]U = \left[6-670-5120043]U = \left[6-670-5120043\right]U = \left[6-670-5120044\right]U = \left[6-670$

In MATLAB

- MATLAB uses a slightly different convention:
 - o It writes:

A=PLUA = P L U

where PP is a permutation matrix

- The **L matrix is not exactly lower triangular** due to permutations
 - MATLAB calls it "psychologically lower triangular"
 - Meaning: it behaves like a lower triangular matrix, even if rows are permuted

Summary Formula

PA=LUPA = LU

- P: product of all permutation matrices (row swaps)
- L: lower triangular matrix (from elimination multipliers)
- U: final upper triangular matrix

Why It Matters

- LU decomposition is **much faster** than Gaussian elimination if you're solving the same system with different right-hand sides.
- LU with pivoting ensures numerical stability.

For large systems, let MATLAB do the work!

Would you like to see an actual example of LU decomposition with pivoting in Python or MATLAB? I can walk you through the code and outputs.

Here's a clear and concise summary of the video on operation count and algorithm scaling in matrix algebra by Prof. Jeff Chasnov:



Lecture Summary: Why LU is Faster – Operation Count and Scaling

Main Claim Recap

- LU decomposition helps **speed up** solving systems like Ax=bAx = b, especially when:
 - AA is fixed
 - bb changes frequently (e.g. in time-stepping PDEs)
- Why? Because once A=LUA = LU is computed, solving becomes cheaper via forward and backward substitution rather than full Gaussian elimination each time.



We use operation counting:

- Count multiplications and additions an algorithm performs
- Called the operation count or computational complexity

P Big O Notation:

- Denotes how the number of operations scales with problem size
- We care about the **leading-order term** only (e.g. n3n^3, not constants or lower powers)

Example: Matrix Multiplication

Let's say AA and BB are both n×nn \times n matrices.

- **Operation Count:**
- For one element of result matrix:
 - o nn multiplications
 - o n-1n 1 additions
- **Total number of elements** in output matrix = n2n^2

✓ Total Operations:

 $Multiplications=n2 \cdot n=n3 \\ Additions=n3 \ (since n2(n-1)\approx n3) \\ text{Multiplications} = n^2 \\ cdot n=n^3 \\ text{Additions} \\ approx n^3 \\ text{(since } n^2(n-1) \\ approx n^3) \\$

So, matrix multiplication scales as:

O(n3)\boxed{ $O(n^3)$ }

Real-World Application (MATLAB)

- Multiply two random 4096×40964096 \times 4096 matrices:
 - o Takes about 1.1 seconds
- Now multiply two 8192×81928192 \times 8192 matrices:
 - Size is doubled \rightarrow 2n2n
 - o Expected time:

 $T2n\approx 8 \cdot Tn=8.8 \text{ secondsT}_{2n} \operatorname{approx } 8 \cdot T_n = 8.8 \cdot \text{text}$

- o Actual result: 7.9 seconds
- Close enough → scaling holds

Why This Matters

- Helps estimate compute time in high-performance computing
- If a simulation takes 30 mins, you want to know:
 - o Can I double resolution and finish it **overnight**?
 - o If complexity is O(n3)O(n^3), doubling nn makes time ≈8×\approx 8 \times

What's Next?

- Next videos:
 - 1. Operation count for Gaussian elimination
 - 2. Operation count for forward/backward substitution
 - 3. Final justification: **LU + substitution is faster** than repeated elimination

Review Summary

Task	Scaling	Operations
Matrix Multiplication	O(n3)O(n^3)	~n3n^3
Gaussian Elimination	Coming soon	TBD
LU + Substitution	Coming soon	TBD

Would you like a visual chart or Python demo showing how run-time changes with nn in matrix multiplication?

Here's a clear and structured summary of the video on Operation Count for Gaussian **Elimination** by Prof. Jeff Chasnov:



★ Lecture Summary: Counting Operations in Gaussian Elimination



To compute the number of operations (additions + multiplications) required to perform Gaussian elimination, and understand how it scales with matrix size nn.

Key Identities Used

To simplify the count, we use standard summation formulas:

- 1. $\sum k=1n1=n\sum_{k=1}^n 1=n$
- 2. $\sum k=1nk=n(n+1)2\sum_{k=1}^n k = \frac{n(n+1)}{2}$
- 3. $\sum k=1nk2=n(n+1)(2n+1)6\sum_{k=1}^n k^2 = \frac{n(n+1)(2n+1)}{6}$

Step-by-Step Operation Count



- Consider an n×nn \times n matrix AA
- We perform elimination **pivot-by-pivot** from row 1 to n-1n-1

At Pivot Step ii:

- Pivot is at position (i,i)(i, i)
- To eliminate below this pivot:
 - o Operate on n-in i rows below
 - Each row has n-in i elements to the right

Each step needs:

(n-i)2 multiplications and additions $(n-i)^2 \text{ text} \{ \text{ multiplications and additions} \}$

Total Operations:

 $\sum_{i=1}^{n-1}(n-i)^2 = k=1n-1k^2=(n-1)n(2n-1)6\sum_{i=1}^{n-1}(n-i)^2 = \sum_{i=1}^{n-1}k^2=1$ \frac{(n-1)n(2n-1)}{6}

This expression grows like 13n3\frac{1}{3}n^3 for large nn

Scaling Result

Gaussian Elimination is O(n3) boxed{\text{Gaussian Elimination is } $O(n^3)$ }

- Doubling matrix size \rightarrow 8× computation time
- This is **why LU decomposition** is expensive upfront

Conclusion

Method **Operation Count** Scaling

Gaussian Elimination \sim 13n3\sim \frac{1}{3}n^3 $O(n3)O(n^3)$

- This confirms Gaussian elimination is costly, especially for large matrices
- Next: Show that forward and backward substitution are only O(n2)O(n^2), making LU + substitution much faster when solving with multiple right-hand sides

Would you like a visual animation or table comparing Gaussian Elimination vs LU + Substitution over different values of nn?

Here's a concise summary of this video on operation count for forward and backward substitution, which follows up the earlier discussion on Gaussian elimination:



📌 Lecture Summary: Why LU Decomposition is Efficient for Multiple Right-Hand Sides

Ø Objective

To justify why solving systems using LU decomposition (i.e., solving $A\vec{x} = \vec{b} A \cdot \vec{x} = \vec{x} = \vec{x}$ $LUx^2 = b^2 LU \cdot vec\{x\} = vec\{b\}$ is far more efficient than performing Gaussian elimination each time for multiple right-hand sides b \vec{b}.

Operation Count for Back Substitution

Scenario:

- You already have Ux =y U\vec{x} = \vec{y}, where UU is upper triangular
- Solving this via back substitution

At row ii:

• You compute:

 $xi=1aii(bi-\sum_{j=i+1}naijxj)x_i = \frac{1}{a_{ii}}\left(b_i - \sum_{j=i+1}^n a_{ij}x_j\right)$

- Requires:
 - o n-in i multiplications (for the sum)
 - o 1 subtraction
 - 1 division (can count as a multiplication)
- Total: (n-i+1)(n i + 1) operations per row

Total operations:

 $\sum_{i=1}^{n}(n-i+1)=\sum_{k=1}^{n}(n+1)2\sum_{i=1}^{n}(n-i+1)=\sum_{k=1}^{n}k=n(n+1)2\sum_{i=1}^{n}(n-i+1)=\sum_{k=1}^{n}($

 \Diamond Back substitution = $12n2+12n\sim O(n2) \frac{1}{2}n^2 + \frac{1}{2}n \sim O(n^2)$

Forward Substitution (for solving Ly = b L\vec{y} = \vec{b})

- Same structure and cost as back substitution
- So total operations for solving $LUx = bLU \le x = vec\{b\}$:

O(n2) (forward)+O(n2) (backward)=O(n2)\boxed{ $O(n^2)$ \text{ (forward)} + $O(n^2)$ \text{ (backward)} = $O(n^2)$ }

Comparison with Gaussian Elimination

Step	Operation Count	Scaling
LU Decomposition	13n3\frac{1}{3}n^3	O(n3)O(n^3)
Forward + Back Substitution	n2n^2	O(n2)O(n^2)

Why This Matters

- If solving many systems $Ax^2=b^21,b^22,...A \le x \le b \le 1$, $x \le b \le 1$.
- For example:
 - o n=1000n = 1000: LU is **300**× slower than one substitution
 - o n=10,000n = 10,000: LU is **3000**× slower

- So decompose once, solve fast many times.
- Conclusion
- Gaussian elimination = costly: $O(n3)O(n^3)$
- LU + forward/backward substitution = efficient: O(n2)O(n^2)
- For multiple right-hand sides, LU decomposition is dramatically faster

Would you like a comparison chart or code example in Python or MATLAB to see this efficiency in practice?

Here's a summary of the Eigenvalue Power Method explained in the lecture by Jeff Chasnov:



Q Topic: The Power Method for Finding the Dominant Eigenvalue



To compute the largest eigenvalue (in absolute value) and its associated eigenvector of a real matrix AA, using the Power Method.

* Assumptions

- $A \in Rn \times nA \in \mathbb{R}^{n} \times nA$
- All eigenvalues λ1,λ2,...,λn\lambda 1, \lambda 2, \ldots, \lambda n are **real and distinct**
- $|\lambda 1| > |\lambda 2| > ... > |\lambda n| |\lambda 2| > |\lambda 2| > \lambda 3|$
- Corresponding eigenvectors e1,e2,...,ene_1, e_2, \ldots, e_n are linearly independent

Key Idea

Start with an arbitrary initial vector x0x 0 (not orthogonal to e1e 1), and repeatedly multiply by AA:

$$xp=Apx0x_p = A^p x_0$$

Since any vector can be written as:

$$x0=c1e1+c2e2+...+cnenx_0 = c_1 e_1 + c_2 e_2 + \label{eq:c1e1} dots + c_n e_n$$

Then:

$$xp=Apx0=\sum_{i=1}^{i=1}nci\lambda_{i}p=ix_{p}=A^{p}x_{0}=\sum_{i=1}^{n}c_{i}\lambda_{i}p=i$$

As $p \rightarrow \infty p \to \inf_{k \in \mathbb{Z}_i} \left(\frac{\lambda i \lambda 1}{p\left(\frac{\lambda i \lambda 1}{p \cdot \frac{1}{p \cdot \frac{1$

```
xp≈c1λ1pe1x_p \approx c_1 \lambda_1^p e_1
```

Thus, xpx_p tends toward the direction of e1e_1, and its **magnitude** grows/shrinks by λ 1p\lambda_1^p.

III Estimating λ1\lambda_1

You can estimate $\lambda 1 \times 1$ using the **Rayleigh quotient**:

 $\lambda 1 \approx xp Txp + 1xp Txp \lambda_1 \approx xp Tx$

This approximates the dominant eigenvalue once xpx_p is close to e1e_1.

Numerical Stability: Overflow / Underflow

- If |λ1|>1 \lambda_1| > 1, xpx_p becomes too large (overflow)
- If |λ1|<1|\lambda_1| < 1, xpx_p becomes too small (underflow)
- Fix: Normalize at each iteration:

```
xk+1=Axk\|Axk\|x_{k+1} = \frac{A x_k}{A x_k}
```

This keeps the vector at a manageable size.

Algorithm Summary: Power Method

```
x = random_vector()
```

for k in range(max_iters):

```
x_new = A @ x
```

x_new = x_new / norm(x_new) # Normalize to avoid overflow

lambda approx = (x.T @ A @ x) / (x.T @ x)

x = x new

Result

- xx converges to the eigenvector e1e_1 of the dominant eigenvalue
- λ1\lambda_1 can be approximated by:

 $\lambda 1 \approx xkTAxkxkTxk \cdot 1 \cdot x_k \cdot x_k$

Conclusion

- Use MATLAB's eig() for general problems.
- But the **Power Method** teaches:
 - o Insight into matrix behavior
 - A fast way to find the largest eigenvalue/eigenvector for large, sparse, or structured matrices.

Would you like a simple worked-out numerical example in Python or MATLAB next?

Here is a clear and concise **summary of the example** Professor Jeff Chasnov worked through using the **Power Method** to find the **dominant eigenvalue and eigenvector**:

Example: Power Method with a 2×2 Matrix

Matrix A

 $A=[6545]A = \left[6545 \right]A = \left[$

This matrix has:

Step-by-step Iteration

Initial vector:

 $x0=[10]x_0 = \left[10 \right] x_0 = \left[10$

Multiply repeatedly by AA:

1.

x1=Ax0=[64]x 1 = A x 0 = \begin{bmatrix} 6 \\ 4 \end{bmatrix}

2.

 $x2=Ax1=A[64]=[6\times6+5\times44\times6+5\times4]=[5644]x_2=Ax_1=A \begin{bmatrix} 6 \ 4 \end{bmatrix} = \begin{bmatrix} 56 \ 44 \end{bmatrix}$

3.

 $x3=Ax2=A[5644]=[6\times56+5\times444\times56+5\times44]=[556444]x_3=Ax_2=A \begin{bmatrix} 56 \ 44 \end{bmatrix} = \begin{bmatrix} 6\times56 + 5\times44 \ 4\times56 + 5\times44 \end{bmatrix} = \begin{bmatrix} 556 \ 444 \end{bmatrix}$

4.

 $x4=Ax3=A[5564444]=[55564444]x_4=Ax_3=A \begin{bmatrix} 556 \ 444 \end{bmatrix} = \begin{bmatrix} 5556 \ 4444 \end{bmatrix}$

5.

 $x5=Ax4=A[555644444]=[5555644444]x_5 = A x_4 = A \end{bmatrix} 5556 \ 44444 \end{bmatrix} = \end{bmatrix}$

Pattern: The digits clearly show the dominant eigenvalue effect.

Estimating the Dominant Eigenvalue (λ1\lambda_1)

Using the Rayleigh Quotient:

 $\lambda 1 \approx x4Tx5x4Tx4 = 506178271506180272 \approx 9.99991 \ | 1 \approx x4Tx5x4Tx4 = \frac{x_4^T x_5}{x_4^T x_4} = \frac{506178271}{506180272} \ | 1 \approx x4Tx5x4Tx4 = \frac{x_4^T x_5}{x_4^T x_4} = \frac{1}{x_4^T x_4^T x_4} = \frac{1}{x_4^T x_4^T x_4^T x_4} = \frac{1}{x_4^T x_4^T x_4^T x_4^T x_4} = \frac{1}{x_4^T x_4^T x$

Approximation is very close to 10

Approximating the Dominant Eigenvector

From:

 $x5 = [5555644444] \Rightarrow e1 \approx [55556444441] = [1.250021] \\ x_5 = \left[5555644444] \right] = \left[1.250021] \\ x_5 = \left[5555644444] \right] = \left[1.25002 \right] \\ x_5 = \left[5555644444] \right] = \left[1.25002 \right] \\ x_6 = \left[1.25002 \right] \\ x_7 = \left[1.25002$

This is close to:

[541]\begin{bmatrix} \frac{5}{4} \\ 1 \end{bmatrix}

Summary

- Start with x0=[10]x_0 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}
- After just 5 iterations, the approximation of the dominant eigenvalue converged to ~10
- The eigenvector converged to approximately [54,1]T\left[\frac{5}{4}, 1 \right]^T
- Simple method, useful for educational purposes and understanding matrix behavior

Would you like a Python or MATLAB implementation of this example?

Here's a clear **summary of the key MATLAB concepts** demonstrated in the video by Professor Jeff Chasnov, organized into four major matrix algebra topics:

- Matrix Algebra in MATLAB: Key Operations
- 1 Solving a System of Linear Equations

Goal: Solve Ax=bA x = b, where AA is a square invertible matrix.

MATLAB Code:

```
A = [-3 2 -1; 6 -6 7; 3 -4 4];
```

x = A\b; % Uses Gaussian elimination internally

Output:

x =

2

2

-1

✓ The \ operator automatically chooses the best method to solve the system.

2 Solving a Least-Squares Problem

Goal: Solve Ax=bA x = b when AA is overdetermined (more rows than columns), no exact solution.

MATLAB Code:

```
A = [11; 12; 13];
```

$$b = [1; 3; 2];$$

x = A\b; % Finds the least-squares solution

Output:

x =

1.0000

0.5000

 \checkmark Equivalent to solving the **normal equations**: ATAx=ATbA^T A x = A^T b

3 LU Decomposition

Goal: Factor A=LUA = L U, useful for solving many systems with the same matrix AA but different right-hand sides bb

MATLAB Code:

$$[L, U] = Iu(A);$$

 $x = U \setminus (L \setminus b)$; % Forward and backward substitution

Notes:

- L is a **permuted lower triangular matrix**, not strictly lower triangular due to pivoting.
- This approach is **faster** for large matrices and repeated solves.

4 Eigenvalues and Eigenvectors

```
Goal: Solve Ax = \lambda x A x = \lambda x A
```

MATLAB Code:

```
A = [0 1; 1 0]; % Simple matrix
```

lambda = eig(A); % Just eigenvalues

[V, D] = eig(A); % V: eigenvectors, D: diagonal matrix of eigenvalues

Output:

lambda =

-1

1

V = % Columns are eigenvectors

-0.7071 0.7071

0.7071 0.7071

D = % Diagonal eigenvalues

- -1 0
- 0 1

Check:

A * V(:,1) % Should equal -1 * V(:,1)

✓ Confirms AV=VDA V = V D, which is matrix diagonalization.

Conclusion

- MATLAB makes matrix computations easy and efficient.
- Description of the backslash operator \ is a powerful tool for both exact and approximate solutions.
- Functions like lu and eig provide deeper insights into matrix structure.

Would you like this in a PDF handout or an interactive MATLAB Live Script?

This video introduces **Newton's method for solving a system of nonlinear equations**, generalizing from the scalar (single-equation) case. Here's a concise summary of the concepts covered:

Goal:

Solve a system of two nonlinear equations:

$$f(x,y)=0g(x,y)=0f(x, y) = 0 \setminus g(x, y) = 0$$

This generalizes to *N* equations in *N* variables.

Newton's Iteration for Systems:

We seek a sequence $(xn,yn)(x_n, y_n)$ that converges to a root of the system.

Derivation via Taylor Expansion:

Use first-order Taylor expansion of $f(xn+1,yn+1)f(x_{n+1}, y_{n+1})$ and $g(xn+1,yn+1)g(x_{n+1}, y_{n+1})$ around $(xn,yn)(x_n, y_n)$:

 $f(xn+1,yn+1)\approx f(xn,yn)+\partial f\partial x\Delta xn+\partial f\partial y\Delta yng(xn+1,yn+1)\approx g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow f(xn,yn)+\partial f\partial x\Delta xn+\partial f\partial y\Delta yng(xn+1,yn+1)\approx g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial y\Delta ynf(x_{n+1}, y_{n+1}) \Rightarrow g(xn,yn)+\partial g\partial x\Delta xn+\partial g\partial x\Delta xn+$

where:

 $\Delta x_n = x_n + 1 - x_n, \Delta y_n = y_n + 1 - y_n \Delta x_n = x_{n+1} - x_n, \quad \Delta x_n = x_n + 1 - x_n + 1 - x_n, \quad \Delta x_n = x_n + 1 - x_n +$

We want f(xn+1,yn+1)=0 $f(x_{n+1}, y_{n+1}) = 0$ and g(xn+1,yn+1)=0 $g(x_{n+1}, y_{n+1}) = 0$, so we get:

 $\label{thm:linear_continuous_c$

\) Jacobian Matrix:

The matrix of partial derivatives is called the **Jacobian**:

 $J(xn,yn)=[fxfygxgy]J(x_n,y_n) = \left\{begin\{bmatrix\} f_x & f_y \right\} g_x & g_y \left\}$

Iteration Update Rule:

1. Solve the linear system:

 $J(xn,yn)[\Delta xn\Delta yn] = -[f(xn,yn)g(xn,yn)]J(x_n,y_n) \cdot begin\{bmatrix\} \cdot Delta x_n \cdot \nabla Delta y_n \cdot B(xn,yn) \cdot B$ $\end{bmatrix} = -\begin{bmatrix} f(x_n, y_n) \ g(x_n, y_n) \end{bmatrix}$

2. Update:

 $xn+1=xn+\Delta xnyn+1=yn+\Delta ynx_{n+1}=x_n+\Delta x_n \setminus y_{n+1}=y_n+\Delta ynx_{n+1}=x_n+\Delta x_n \setminus y_{n+1}=y_n+\Delta ynx_{n+1}=x_n+\Delta x_n \setminus y_{n+1}=y_n+\Delta ynx_{n+1}=x_n+\Delta x_n \setminus y_{n+1}=x_n+\Delta x_n+\Delta x_n \setminus y_{n+1}=x_n+\Delta x_n+\Delta x_n+$

3. Repeat until convergence.

Convergence Tip:

- Newton's method converges fast (quadratically), if the initial guess is close to a solution.
- You must recompute the Jacobian and the function values at each iteration.

Application:

You'll likely be asked to:

- Implement this method.
- Compute Jacobians manually or symbolically.
- Use matrix algebra techniques like Gaussian elimination to solve the update equations.

If you'd like, I can provide a MATLAB or Python implementation of this method or help you solve a specific system using it.

Here's a clear breakdown of the **Newton's method application to the Lorenz system** from the video:

S Lorenz Equations (for chaotic dynamics)

 $x = \sigma(y-x)y = x(r-z) - yz = xy - \beta z \setminus \{x\} \& = \sigma(y-x)y = x(r-z) - y \setminus \{x\} \& = x(r-z) - y \setminus \{x\} \& = x(r-z) - y \in x(r-z) - x(r-z) - y \in x(r-z) - x(r-z)$ \dot{z} &= xy - \beta z \end{aligned}

We're looking for **fixed points**, so set the derivatives to **zero**:

 $f(x,y,z) = \sigma(y-x) = 0g(x,y,z) = x(r-z) - y = 0h(x,y,z) = xy - \beta z = 0$ legin{aligned} f(x, y, z) &= sigma(y - x) = 0

Goal: Use Newton's method to find (x,y,z)(x, y, z) such that:

 $f(x,y,z)=0,g(x,y,z)=0,h(x,y,z)=0f(x,y,z)=0,\quad g(x,y,z)=0,\quad h(x,y,z)=0$

Take partial derivatives of f,g,hf, g, h with respect to x,y,zx, y, z:

All values are evaluated at the current guess $(xn,yn,zn)(x_n, y_n, z_n)$.

Mewton's Iteration Step:

We solve for:

Where:

- $f(xn,yn,zn)=\sigma(yn-xn)f(x_n,y_n,z_n) = sigma(y_n x_n)$
- g(xn,yn,zn)=xn(r-zn)-yng(x_n, y_n, z_n) = x_n(r z_n) y_n
- $h(xn,yn,zn)=xnyn-\beta znh(x_n,y_n,z_n)=x_ny_n \beta znh(x_n,y_n,z_n)$

Update:

 $xn+1=xn+\Delta xnyn+1=yn+\Delta ynzn+1=zn+\Delta zn\begin{aligned} x_{n+1} &= x_n + \Delta z_n \\ &= y_n + \Delta z_n + \Delta z_n \\ &= z$

Repeat this until the change $(\Delta x n, \Delta y n, \Delta z n)$ (\Delta x_n, \Delta y_n, \Delta z_n) is sufficiently small.

Suggested Homework Implementation (in MATLAB or Python):

- 1. Choose initial guess $(x_0,y_0,z_0)(x_0,y_0,z_0)$ for example: (1,1,1)(1,1,1)
- 2. Set parameters σ ,r, β \sigma, r, \beta (e.g., σ =10,r=28, β =83\sigma = 10, r = 28, \beta = \frac{8}{3})
- 3. Iterate Newton's method:
 - o Evaluate the Jacobian and the function values
 - Solve the matrix equation for deltas
 - Update x,y,zx, y, z

Reminder:

- You can analytically solve for the fixed points (try it as an exercise!).
- Newton's method converges fast, but only if your initial guess is close to a root.

Let me know if you want help implementing this in MATLAB or Python, or if you'd like help solving for the **exact fixed points** analytically.

Here's a complete summary and action plan for your Project 3: Fractals from Newton's Method on Lorenz Equations.



Project Overview: Newton Fractals in 3D (Lorenz Fixed Points)

You're building a 3D Newton fractal using the Lorenz system's fixed points and visualizing convergence behavior in the **xz-plane** for a **fixed y**.

The Lorenz Fixed Point Equations

Set time derivatives to 0 (fixed points):

 $f(x,y,z) = \sigma(y-x) = O(x,y,z) = x(r-z) - y = O(x,y,z) = xy - \beta z = O \cdot \{aligned\} \ f(x,y,z) \ \& = sigma(y-x) = O(x,y,z) + (x,y,z) + (x,$

You already have the Newton iteration set up:

 $\end{bmatrix} = -\begin{bmatrix} f(x, y, z) \ g(x, y, z) \ h(x, y, z) \end{bmatrix}$

Update:

z &\leftarrow z + \Delta z \end{aligned}

The 3 Fixed Points

1. Origin:

(0,0,0)(0,0,0)

2. Positive branch:

$$x=y=\beta(r-1), z=r-1x = y = \sqrt{\frac{1}{2}}, \quad z = r-1$$

3. Negative branch:

$$x=y=-\beta(r-1), z=r-1x = y = -\sqrt{\frac{r-1}{r-1}}, \quad z = r-1$$

- σ=10\sigma = 10
- r=28r = 28
- $\beta=83$ \beta = \frac{8}{3}

Then the square root term becomes:

 $\beta(r-1)=83(27)=72=62 \operatorname{total} = \operatorname{sqrt} \{frac\{8\}\{3\}(27)\} = \operatorname{sqrt}\{72\} = 6 \operatorname{total} = 6 \operatorname{total$

So the two non-zero fixed points are:

- $x=y=\pm 62\approx \pm 8.485x = y = \pm 6\sqrt{2} \approx \pm 8.485$
- z=27z=27

Fractal Mapping Plan

- You'll map a grid in the **xz-plane** (e.g., $-20 \le x \le 20-20 \le x \le 20, -20 \le x \le 40-20 \le x \le 40-20$ 40)
- For each grid point (x0,z0)(x_0, z_0), fix:

 $y0=32\approx4.24y_0 = 3\sqrt{2} \alpha 4.24$

- Use **Newton's method** to determine to which of the 3 fixed points (x,y,z)(x, y, z) converges.
- Assign colors:
 - o **Red** → origin
 - o **Green** \rightarrow +62,+62,27+6\sqrt{2}, +6\sqrt{2}, 27
 - o **Blue** \rightarrow -62,-62,27-6\sqrt{2}, -6\sqrt{2}, 27

Speed Optimization Strategy

Don't use a fixed number of iterations. Use a while loop.

At each grid point:

Iterate until:

 $\|\Delta x\| + \|\Delta y\| + \|\Delta z\| < \epsilon \setminus \Delta x + \|\Delta y\| + \|\Delta z\| < \epsilon$

- Use a max iteration cap (e.g., 50) to avoid infinite loops.
- Convergence is **slower near fractal boundaries**, so most time is spent there.

Steps to Implement in MATLAB

- 1. Define constants: sigma = 10; beta = 8/3; r = 28;
- 2. Define grid for x0x_0 and z0z_0

- 3. Fix $y0=3*sqrt(2)y_0 = 3*sqrt(2)$
- 4. For each grid point:
 - Use while loop for Newton's method
 - o Compute Jacobian and function values
 - Solve linear system for Δx,Δy,Δz\Delta x, \Delta y, \Delta z
 - Update values
 - o Check convergence
- 5. Once converged, check final (x,y,z)(x, y, z) against fixed points
 - Use norm(x x_fixed) to match to one of the 3 roots
 - o Assign color index: 1 = red, 2 = green, 3 = blue
- 6. Plot fractal with imagesc, imshow, or scatter

Tips

- Use norm([dx, dy, dz]) < tol for convergence check.
- Use a grid size of 200×200 or smaller if performance is an issue.
- Store output in a 2D matrix where each cell stores the color (1, 2, 3).
- Preallocate the matrix for speed.
- For identifying convergence, use a tolerance (e.g., 1e-6).

Example: MATLAB Color Mapping Snippet

```
colors = [1 0 0; 0 1 0; 0 0 1]; % RGB: Red, Green, Blue
image_data = zeros(Nx, Nz, 3);

for i = 1:Nx
    for j = 1:Nz
        % Newton iteration...
    if converged_to_root1
        image_data(i, j, :) = colors(1, :);
    elseif converged_to_root2
        image_data(i, j, :) = colors(2, :);
    else
```

```
image_data(i, j, :) = colors(3, :);
end
end
end
imshow(image_data);
```

Would you like a full **MATLAB script** to start with? I can provide a complete template including initialization, iteration, convergence test, and plotting.

format compact — Makes output cleaner with less spacing.

- CLC Clears the **Command Window**.
- ; (semicolon) Suppresses output when assigning values:
- x = 0; % No output shown
- y = 0; % No output shown
- 2*x % Output: 0

What's Next

- This video focused on interacting with MATLAB via the command line.
- In real-world projects, you will write **scripts and functions** for complex tasks.

Let me know if you want practice problems or a cheatsheet for MATLAB operators and functions.

Here's a structured **summary** of the second MATLAB video by **Jeff Chasnov**, focused on **scripts**, **functions**, **and programming in MATLAB**:

Scripts vs. Functions in MATLAB

Feature	Script	Function
Use Case	Series of commands, like a calculator	Reusable code with input and output
Variables	Shared with the workspace	Local (isolated from the workspace)
Inputs/Outputs None required		Can have both inputs and outputs

Feature	Script	Function
Typical Usage	Simple tasks, plotting, testing ideas	Repetitive or modular code (like defining equations)

Creating a Script

1. How to Create:

- Click New Script (toolbar or File > New > Script)
- Save with a .m extension

2. Start with Clean Slate (Recommended Initial Line):

- 3. clear all; close all; clc;
- 4. Example Golden Ratio:
- 5. clear all; close all; clc;
- 6.
- 7. Phi = (sqrt(5) + 1) / 2;
- 8. phi = (sqrt(5) 1) / 2;

Run the script → MATLAB will ask to name the file (e.g., golden_ratio.m)

9. Command Window Check:

- 10. Phi * phi % ≈ 1 with small roundoff error
- 11. 1 ans % Shows the tiny error: -2e-16

Creating a Function

1. How to Create:

- Click **New > Function**
- o Template appears automatically

2. Function Example – Euler's Number e:

MATLAB doesn't have a built-in e constant like pi. So:

- 3. function e = e()
- 4. e = exp(1);
- 5. end

Save as: e.m

6. Using the Function:

- Call in the Command Window:
- % returns 2.7183...
- log(e) % returns 1

MATLAB will check for a matching file (e.m) in the current directory.

Objection State of Contract Contract State of Contract C

- A **local function** is a helper function written **inside a script**.
- Common in solving **differential equations** or for organizing code.

Best Practices

- Use **scripts** for building programs step by step.
- Use **functions** for modular, reusable code.
- Always start fresh with:
- clear all; close all; clc;

Conclusion

"You'll write many scripts in this course. By the end I hope you'll be a pretty good MATLAB programmer."

Jeff Chasnov

Let me know if you'd like help writing your first function, practicing exercises, or creating a template for scripts or functions in MATLAB.

Here's a clear summary and explanation of what Professor Jeff Chasnov derived in this video: the Midpoint Rule for numerical integration using Taylor series.

Goal: Approximate the integral

We want to compute:

Ih= $\{0hf(x) dx | h = \}$ oh $\{f(x) \}$, dx

This is a small "building block" integral which can later be used in larger composite rules to integrate over intervals like [a,b][a, b].

Step-by-step Derivation

1. Expand f(x)f(x) as a Taylor Series around x=h/2x = h/2 (the midpoint):

 $f(x) = f(h2) + (x-h2)f'(h2) + (x-h2)22f''(h2) + \cdots \\ f(x) = f(h\{2\} right) + (x-h\{2\} right$

He writes more terms than needed (to analyze error):

 $f(x) = f(h2) + (x-h2)f'(h2) + (x-h2)22!f''(h2) + (x-h2)33!f(3)(h2) + (x-h2)44!f(4)(h2) + \cdots \cdot f(x) = f(f(x) + f(x))^2 + f(x) +$

2. Substitute this into the integral

Now plug the Taylor expansion into the integral:

 $Ih = \int 0hf(x) dx = \int 0h(Taylor series) dxI_h = \int 0hf(x) dx = \int 0hf(x)$

3. Change of Variables

Let:

 $y=x-h2 \Rightarrow dx=dyy = x - \frac{h}{2} \operatorname{Rightarrow} dx = dy$

Limits change from:

- $x=0 \Rightarrow y=-h/2x = 0 \setminus Rightarrow y = -h/2$
- $x=h \Rightarrow y=+h/2x = h \setminus Rightarrow y = +h/2$

Now the integral becomes symmetric:

 $Ih=\int -h/2h/2(Taylor series in y) dyl_h = \int -h/2^{h/2} \text{(Taylor series in } y) \, dy$

4. Use symmetry to eliminate odd terms

Since we're integrating over [-h/2,h/2][-h/2,h/2], **odd powers of yy** integrate to 0:

Only even terms contribute:

 $Ih= \int -h/2h/2[f(h2)+y22f''(h2)+y424f(4)(h2)+\cdots] dyl_h = \int_{-h/2}^{h/2} \left[f\left(\frac{h}{2}\right) + \frac{y^2}{2} f''\left(\frac{h}{2}\right) + \frac{y^4}{24} f^{(4)}\left(\frac{h}{2}\right) + \frac{y^4}{2$

5. Evaluate Integrals

You need:

```
•
\int_{-h/2}^{h/2} dy = h
]
```

 $\int_{-h/2}^{h/2} y^2 , dy = \frac{h^3}{12}$ $\int_{-h/2}^{h/2} y^4$, dy = $\int_{-h/2}^{80}$

Plug back into the formula:

 $Ih=hf(h2)+h324f''(h2)+h51920f(4)(h2)+\cdots \setminus I_h &= h f\left(\frac{h}{2}\right)+h51920f(4)(h2)+\cdots \setminus I_h &= h f\left(\frac{h}{2}\right)+h51920f(4)(h2)$ $\frac{h^3}{24} f'' \left(\frac{h}{2}\right) + \frac{h^5}{1920} f^{(4)}\left(\frac{h}{2}\right) + \cdot cdots$ \end{aligned}

Final Result: The Midpoint Rule

 $0^h(x) dx \approx h \cdot f(h2) \cdot f(h2$

Error term: O(h3)\mathcal{O}(h^3)

 $Error = h324f''(h2) + \cdots + \{Error\} = \frac{h^3}{24} f'' \left(\frac{h}{2}\right) + \cdots$

Interpretation

- This is a simple and powerful **numerical integration rule**.
- You're using **only one function evaluation** at the midpoint of the interval.
- Accuracy is better than Left or Right Riemann sums (which are O(h)\mathcal{O}(h)).

Let me know if you'd like the code implementation of this in MATLAB or Python—or help deriving **Trapezoidal Rule** or **Simpson's Rule**, which are coming next in the course.

Here's a clear summary and explanation of how Professor Jeffrey Chasnov derives the Trapezoidal Rule for numerical integration from a Taylor series expansion—building on the Midpoint Rule.



We want to compute:

 $Ih=\int Ohf(x) dxI_h = \int O^h f(x), dx$

Step-by-Step Derivation

1. Recall the Taylor Series about the midpoint x=h/2x = h/2

 $f(x) = f(h2) + (x-h2)f'(h2) + (x-h2)22!f''(h2) + (x-h2)33!f(3)(h2) + (x-h2)44!f(4)(h2) + \cdots \cdot f(x) + \frac{h}{2} \cdot f(x) + \frac{h}{$

2. Evaluate the Taylor series at x=0x=0

 $f(0) = f(h2) - h2f'(h2) + h28f''(h2) - h348f(3)(h2) + h4384f(4)(h2) + \cdots \cdot h2f'(h2) + f(0) &= f(h^2f'(h^2) - \frac{h}{2} f'\left(\frac{h}{2}\right) + \frac{h^2f'(h^2) + \frac{h^2f'(h^2)}{eft(\frac{h}{2}\right)} + \frac{h^2f'(h^2)}{eft(\frac{h}{2}\right)} + \frac{h^2f'(h^2)}{eft(\frac{h^2f'(h^2)}{eft(\frac{h}{2}\right)} + \frac{h^2f'(h^2)}{eft(\frac{h^2f'(h^2)}{eft(\frac{h}{2}\right)} + \frac{h^2f'(h^2)}{eft(\frac{h^2}{eft(\frac{h^2f'(h^2)}{e$

3. Evaluate the Taylor series at x=hx = h

Same structure, but the signs alternate **positively**:

 $f(h) = f(h2) + h2f''(h2) + h28f''(h2) + h348f(3)(h2) + h4384f(4)(h2) + \cdots \cdot h348f''(h2) + h348f''($

4. Add f(0)+f(h)f(0) + f(h)

 $f(0) + f(h) = 2f(h2) + 2 \cdot h28f''(h2) + 2 \cdot h4384f(4)(h2) + \cdots = 2f(h2) + h24f''(h2) + h4192f(4)(h2) + \cdots \cdot h2f''(h2) + h2192f(4)(h2) + \cdots \cdot h2f''(h2) + h2192f(4)(h2) + \cdots \cdot h2f''(h2) + h2192f(4)(h2) + \dots \cdot h2192f(4)(h2) + \dots$

Now multiply by $h2\frac{h}{2}$:

5. Use Midpoint Rule for Substitution

From the Midpoint Rule, recall:

 $Ih=hf(h2)+h324f''(h2)+h51920f(4)(h2)+\cdots I_h = h f\left(\frac{h}{2}\right) + \frac{h^3}{24} f''\left(\frac{h}{2}\right) + \frac{h^5}{1920} f^{(4)}\left(\frac{h}{2}\right) + \cdot cdots$

Now, rearrange:

Plug into the expression for $h2(f(0)+f(h))\frac{h}{2}(f(0)+f(h))$:

 $h2(f(0)+f(h))=Ih+(h38-h324)f''(h2)+(h5384-h51920)f(4)(h2)+\cdots frac{h}{2}(f(0)+f(h))=I_h+(h38-h324)f''(h2)+(h5384-h51920)f(4)(h2)+\cdots frac{h}{2}(f(0)+f(h))=I_h+(h38-h324)f''(h2)+(h5384-h51920)f(4)(h2)+\cdots frac{h}{2}(f(0)+f(h))=I_h+(h38-h324)f''(h2)+(h5384-h51920)f(4)(h2)+\cdots frac{h}{2}(f(0)+f(h))=I_h+(h38-h324)f''(h2)+(h5384-h51920)f(4)(h2)+\cdots frac{h}{2}(f(0)+f(h))=I_h+(h38-h324)f''(h2)+(h5384-h51920)f(4)(h2)+\cdots frac{h}{2}(f(0)+f(h))=I_h+(h38-h324)f''(h2)+(h5384-h51920)f(4)(h2)+\cdots frac{h}{2}(f(0)+f(h))=I_h+(h38-h324)f''(h2)+(h38-h324)f'''$ $\left(\frac{h^3}{8} - \frac{h^3}{24}\right) f''\left(\frac{h^2}{2}\right) + \left(\frac{h^5}{384} - \frac{h^5}{384}\right) - \frac{h^5}{384} - \frac{h^5}$ $\frac{h^5}{1920}\right) f^{(4)}\left(\frac{h}{2}\right) + \cdots$

Simplify:

- $h38-h324=h312\frac\{h^3\}\{8\} \frac\{h^3\}\{24\} = \frac\{h^3\}\{12\}$
- $h5384-h51920=h5480\frac\{h^5\}{384} frac\{h^5\}{1920} = frac\{h^5\}{480}$

So:

 $h2(f(0)+f(h))=Ih+h312f''(h2)+h5480f(4)(h2)\frac{h}{2}(f(0)+f(h))=I_h+\frac{h^3}{12}$ f"\left(\frac{h}{2}\right) + \frac{h^5}{480} f^{(4)}\left(\frac{h}{2}\right)

Solve for IhI_h:

 $lh=h2(f(0)+f(h))-h312f''(h2)-h5480f(4)(h2)+\cdots$ boxed{ I h = \frac{h}{2}(f(0) + f(h)) - \frac{h^3}{12} f"\left(\frac{h}{2}\right) - \frac{h^5}{480} f^{(4)}\left(\frac{h}{2}\right) + \cdots }

Final Result: The Trapezoidal Rule

Error term: O(h3)\mathcal{O}(h^3)

 $Error = -h312f''(h2) + \cdots + \{Error\} = -\{frac\{h^3\}\{12\} f'' \setminus \{frac\{h\}\{2\} \setminus frac\{h\}\{2\} \} \}$

Intuition

- Approximates area under the curve using a **trapezoid** instead of a rectangle.
- Uses two evaluations: one at each endpoint.
- Better than Left/Right Riemann sum, slightly worse than Midpoint Rule in error constant.

Let me know if you want to see the Simpson's Rule next (which uses three points and has error $O(h5)\mathbb{O}(h^5)$.

Here's a clear summary and explanation of how Simpson's Rule is derived by combining the Midpoint Rule and the Trapezoidal Rule, as explained by Professor Jeffrey Chasnov.



We want a more **accurate numerical integration** rule for:

 $\int Ohf(x) dx \int O^h f(x) dx$

Step-by-Step Derivation

1. Recall the Midpoint and Trapezoidal Rules

• Midpoint Rule:

 $Ih=h f(h2)-h324f''(h2)+h51920f(4)(h2)+\cdots I_h = h\,f\left(\frac{h}{2}\right) - \frac{h^3}{24} f''\left(\frac{h}{2}\right) + \frac{h^5}{1920} f^{(4)}\left(\frac{h}{2}\right) + \cdot f'(4)$

• Trapezoidal Rule:

 $Ih=h2(f(0)+f(h))-h312f''(h2)+h5480f(4)(h2)+\cdots I_h = \frac{h}{2}\left(f(0)+f(h)\right)-\frac{h^3}{12} f''\left(\frac{h}{2}\right) + \frac{h^3}{480} f^{(4)}\left(\frac{h}{2}\right) + \cdot f''(h)^{1}\left(\frac{h^2}{480}\right) f^{(4)}\left(\frac{h^2}{480}\right) f^{(4)}\left(\frac{h^2$

2. Cancel the h3h^3 Error Term by Combining

We add twice the midpoint rule and once the trapezoidal rule:

Let: $2 \cdot (Midpoint) + 1 \cdot (Trapezoidal) = 3 I_h + (higher-order error) \cdot \{Let: } 2 \cdot (\text{Midpoint}) + 1 \cdot (\text{Lext}\{Trapezoidal}) = 3 I_h + \text{(higher-order error)}$

Left-hand side becomes:

 $h2[f(0)+f(h)]+2hf(h2)\frac{h}{2} \left[f(0)+f(h)\right]+2h f\left(\frac{h}{2}\right]$

Right-hand side:

 $=31h+O(h5)=31_h + \mathcal{O}(h^5)$

Now solve for IhI_h by dividing by 3:

3. Simpson's Rule Formula

 $Ih=h6[f(0)+4f(h2)+f(h)]+O(h5)I_h = \frac{h}{6} \left[f(0) + 4f\left(\frac{h}{2}\right) + f(h)\right] + \frac{h}{6} \left[f(0) + f(h)\right] + \frac{h}{6} \left[f(0$

This is Simpson's Rule over the interval [0,h][0, h].

4. Standard Form Over [0,2h][0, 2h]

To make the formula look more standard, consider:

 $\int 02hf(x) dx \int 0^{2h} f(x) dx$

Now:

- Interval length = 2h2h
- Midpoint = hh

Apply Simpson's Rule:

• **Error**: O(h5)\mathcal{O}(h^5)

Summary of All Three Elementary Rules

Rule	Formula	Error Term	Points Used
Midpoint	$\label{lem:lem:hamma} $$h f(h2)h\f(\left(\frac{h}{2}\right))$$$	O(h3)\mathcal{O}(h^3	Midpoint only
Trapezoida	l h2[f(0)+f(h)]\frac{h}{2}[f(0) + f(h)]	O(h3)\mathcal{O}(h^3	Endpoints only
Simpson's	$ \begin{tabular}{ll} $h6[f(0)+4f(h2)+f(h)]\\ f(ac_{h}_{2}) + f(h)] (or h3····\frac_{h}_{3} \cdot \\ cdots) \end{tabular} $	O(h5)\mathcal{O}(h^5	Endpoints + Midpoint

P Insight

- Simpson's Rule cleverly eliminates the h3h^3 error by combining the Midpoint and Trapezoidal Rules.
- It gives much better accuracy using only three points.
- Often used in **composite form** to integrate over longer intervals.

Let me know if you'd like to see **Composite Simpson's Rule** next!

Here's a complete summary and explanation of the video on composite quadrature rules, including how to construct and apply the composite trapezoidal and composite Simpson's rule for numerically integrating a function over an interval [a,b][a, b].



6 Goal

We want to compute the integral:

$$\int abf(x) dx \int a^b f(x), dx$$

by breaking the interval [a,b][a, b] into smaller subintervals and applying our elementary integration formulas (Trapezoidal and Simpson's rule) over each subinterval.

Step 1: Divide the Interval

• Divide [a,b][a, b] into nn subintervals:

$$a=x0$$

Each subinterval is $[xi,xi+1][x i, x \{i+1\}]$ with width:

$$hi=xi+1-xih_i = x_{i+1} - x_i$$

Function values:

 $fi=f(xi)f_i=f(x_i)$

Then:

Apply elementary rules to each subinterval.



Name of the Composite Trapezoidal Rule

For arbitrary (unequal) subintervals:

 $\int x_i dx = 12[f(x_i) + f(x_i) - \{x_i\}^{x_i} f(x) dx \approx 12[f(x_i) + f(x_i) - \{x_i\}^{x_i} f(x) dx = 12[f(x_i) + f(x_i) +$ $f(x_{i+1})\right)$

So the full integral:

 $\left(f(x_i) + f(x_{i+1}) \right)$

Works well for experimental data with uneven spacing.

For equally spaced points:

Let $h=b-anh = \frac{b - a}{n}$, then:

 $[abf(x) dx \approx h2[f0+2f1+2f2+\cdots+2fn-1+fn]$ a^b $f(x) dx \approx h2[f0+2f1+2f2+\cdots+2fn-1+fn]$ $\dots + 2f_{n-1} + f_n \right]$

- Endpoints f0f_0 and fnf_n appear once
- All interior points f1f_1 to fn-1f_{n-1} appear **twice**

Error term:

O(h2) mathcal $\{O\}(h^2)$

(because error from each subinterval is O(h3)) and we have $n=(b-a)hn = \frac{(b-a)hn}{n}$ a)}{h} terms, so total error is O(h2)\mathcal{O}(h^2))

Composite Simpson's Rule

Apply the elementary Simpson's rule over pairs of subintervals, so:

- Use even number of subintervals: nn must be even
- Interval width: h=b-anh = \frac{b a}{n}

Formula:

 $[abf(x) dx \approx h3[f0+4f1+2f2+4f3+2f4+\cdots+4fn-1+fn]$ a^b $f(x)\, dx \approx \frac{h}{3} \left[0 + 4f 1 + 2f2+4f3+2f4+\cdots+4fn-1+fn \right]$ $+ 2f_2 + 4f_3 + 2f_4 + \det + 4f_{n-1} + f_n \right]$

Pattern:

- First and last: f0f_0 and fnf_n appear once
- Odd indices f1,f3,...f_1, f_3, \dots: coefficient 4
- Even indices f2,f4,...f_2, f_4, \dots: coefficient 2

Error term:

O(h4) \mathcal{O}(h^4)

Very accurate compared to trapezoidal.

✓ Summary Table

Rule	Formula	Error	Requirements
Composite Trapezoidal	$ \begin{array}{l} \mbox{$h2[f0+2f1+2f2+\cdots+2fn-1+fn]\\frac\{h\}\{2\}\ [f_0+2f_1+2f_2+\dots+2f_\{n-1\}+f_n]$} \end{array} $	$O(h2)\mathbb{C}_{0}$	Any nn, even or odd
Composite Simpson's	$ \begin{array}{l} \mbox{h3[f0+4f1+2f2+4f3+\cdots+4fn-1+fn]\\frac{h}{3}} \\ \mbox{[f_0+4f_1+2f_2+4f_3+\dots+4f_{n-1}+f_n]} \end{array} $	O(h4)\mathcal{O}(h^4)	nn must be

Practical Use

- Composite trapezoidal: Great for experimental data where xix_i values are uneven.
- Composite Simpson's: Ideal when you can choose function points, and want more accuracy.

Let me know if you'd like to see a Python or MATLAB example applying these rules!

Here's a clear **summary** of the second part of the lecture video on **Gaussian Quadrature**, presented by Jeff Chesnoff:

IIII Gaussian Quadrature: Overview

Gaussian quadrature is a **numerical integration** technique that differs from the trapezoidal and Simpson's rules. It's **more accurate**, especially for polynomials, and widely used in computational tools like MATLAB (integral.m uses it internally).

✓ General Form of Gaussian Quadrature

It approximates the integral of a weighted function:

 $\label{linear_sum_simple} \int abW(x)f(x)\ dx \approx \sum i=1nwif(xi) \int a^b W(x)\ f(x)\ dx \ approx \sum i=1nwif(xi) \int a^b W(x)\ f(x)\ dx \ dx \ dx = 1$

• W(x)W(x): Weight function (varies based on the type of Gaussian quadrature)

- wiw_i: Quadrature weights
- xix_i: Quadrature nodes (points at which the function is evaluated)

Types of Gaussian Quadrature

Each type uses a specific **weight function** and **interval**:

Type Weight Function W(x)W(x) Interval

Legendre-Gauss 11 [-1,1][-1,1]

Chebyshev-Gauss $11-x2\frac{1}{\sqrt{1 - x^2}} [-1,1][-1, 1]$

Laguerre-Gauss e-xe^{-x} $[0,\infty)[0, \inf y)$

 $(-\infty,\infty)(-\infty, \infty)$ **Hermite-Gauss** $e-x2e^{-x^2}$

Legendre-Gauss is the most general-purpose and often used in numerical integration.

Example: 2-Point Legendre-Gauss Quadrature

To integrate $\int -11f(x) dx \int -1^{1} f(x) \ dx$, day, Gaussian quadrature finds weights and nodes such that the formula is exact for polynomials up to degree 3.

 $\int -11f(x)dx \approx w1f(x1) + w2f(x2) \int -11f(x)dx = w1f(x1) + w1f(x1) + w2f(x2) \int -11f(x)dx = w1f(x1) + w1f(x1) + w2f(x2) = w1f(x1) + w1f(x1$

Solving the system for $f(x)=1,x,x2,x3f(x)=1,x,x^2,x^3$ gives:

- Weights: w1=w2=1w_1 = w_2 = 1
- **Nodes**: $x1=-13, x2=13x_1 = -\frac{1}{\sqrt{3}}$, $\qquad x_2 = \frac{1}{\sqrt{3}}$

Final rule:

 $\int -11f(x)dx \approx f(-13) + f(13) \int_{-1}^{1} f(x) dx \approx f\left(-\frac{1}{\sqrt{3}}\right) + f(13) \int_{-1}^{1} f(x) dx \approx f\left(-\frac{1}{\sqrt{3$ f\left(\frac{1}{\sqrt{3}}\right)

Practical Notes

- Higher-order Gaussian quadrature uses more points and has precomputed weights and nodes (in tables or software).
- MATLAB's integral.m function uses adaptive Gaussian quadrature internally.
- You don't typically calculate nodes/weights manually—just understand the concept.



Gaussian quadrature:

- Is highly accurate for smooth functions.
- Uses specially chosen nodes and weights for efficient integration.
- Is widely implemented in software—users benefit from its accuracy without needing to manually compute anything.

Let me know if you want a comparison chart with error rates of Trapezoidal, Simpson's, and Gaussian methods, or examples in Python or MATLAB.

Here's a **clear summary** of the lecture on **Adaptive Quadrature** by Jeff Chesnoff:

Adaptive Quadrature: Concept and Motivation

- **Goal**: Automatically determine where to place function evaluation points (x_i) to achieve a desired **accuracy** without the user needing to specify grid spacing.
- This is what MATLAB's integral.m function does behind the scenes.
- It adapts the grid based on how the function behaves, refining where needed.

To explain adaptive quadrature, the **trapezoidal rule** is used as an example.

♦ Step 1: Start with Level 1 (coarse)

• Integrate from aa to bb using just the endpoints:

 $S1=h2[f(a)+f(b)],h=b-aS_1 = \frac{h}{2}[f(a) + f(b)], \quad h=b-aS_1 = \frac{h}{2}[f(a) + f(b)],$

• Error estimate:

♦ Step 2: Move to Level 2 (refined)

• Add a midpoint c=a+b2c = \frac{a+b}{2} and compute:

 $S2=h4[f(a)+2f(c)+f(b)]S_2 = \frac{h}{4}[f(a) + 2f(c) + f(b)]$

- Now you have two trapezoidal areas: from aa to cc, and cc to bb.
- Assume the second derivatives across both subintervals are roughly equal, giving:

E2~14E1E_2 \approx \frac{1}{4} E_1

Estimating the Error

Since both integrals should approximate the same true value:

 $S1+E1=S2+E2S_1+E_1=S_2+E_2$

Using the assumption E1=4E2E_1 = 4E_2, you get:

 $S1+4E2=S2+E2 \Rightarrow 3E2=S2-S1S_1 + 4E_2 = S_2 + E_2 \setminus Rightarrow 3E_2 = S_2 - S_1$



Error Estimate:

 $|E2|=|S2-S1|3|E_2| = \frac{|S_2-S_1|}{3}$

Recursive Refinement (Adaptive Loop)

If the estimated error E2E_2 exceeds your tolerance:

- **Subdivide** again: from aa to cc and from cc to bb.
- Apply the same logic **independently to each subinterval**.
- Keep going until the local error on each subinterval is below a decreasing fraction of the total tolerance (e.g., half the tolerance for two subintervals).

* Key Points About Adaptive Quadrature

- It automatically refines where the function is difficult (wiggly).
- It uses **fewer points** where the function is smooth (almost constant).
- **Efficient and precise**, especially for irregular functions.
- MATLAB's integral.m function:
 - Uses adaptive quadrature.
 - Has a **default tolerance**, or you can specify your own.
 - You don't need to choose a step size or grid spacing—the algorithm handles it.

Summary

Concept	Meaning
S ₁ , S ₂	Two estimates using different numbers of points
E ₁ , E ₂	Corresponding error estimates
Adaptive process	s Keep subdividing intervals where error > tolerance
Efficiency	Saves time and resources by focusing on problem areas
Used in	MATLAB integral.m, SciPy quad(), and many other numerical libraries

Let me know if you'd like a **visual flowchart**, Python code for adaptive quadrature, or MATLAB example!

Here's a **concise summary** of the video tutorial on how to compute **definite integrals (quadrature)** using MATLAB:



Evaluate the definite integral:

 $\int 0 - ax2(\log x)^2 dx = e^{-ax^2} (\log x)^2 \ dx$

with the parameter $a=\pi a = \pi a = \pi$, using MATLAB.

Key Concepts

- MATLAB function: integral
- Syntax:
- integral(@(x) f(x), lower_limit, upper_limit)
- The @ symbol defines an anonymous function.

Step-by-Step MATLAB Implementation

- 1. **Define the integrand** as a function of x and parameter a:
- 2. $f = @(x, a) \exp(-a * x.^2) .* (log(x)).^2;$
 - \circ .^2 \rightarrow element-wise square for vectors
 - \circ .* \rightarrow element-wise multiplication
- 3. Set parameter a to π :
- 4. a = pi;
- 5. Define the integral using integral and fix a:
- 6. I = integral(@(x) f(x, a), 0, Inf);
- 7. Run it, and you'll get the result:
- 8. I = 1.8245

Key Takeaways

Concept Description

@ (anonymous function) Needed to pass functions to integral

Concept	Description
---------	-------------

integral() MATLAB's adaptive quadrature function

Handles Inf limits No need to truncate manually

Supports parameters Use extra arguments like a via anonymous function

Automatically adapts grid Uses adaptive quadrature internally



Use Case

Useful for integrals involving parameters, infinite limits, or complex functions like Bessel functions (as mentioned for later in the course).

Let me know if you'd like to see a visual plot, a parameter sweep, or a Python equivalent of this

Here's a **summary** of the video on **Piecewise Linear Interpolation**:



***** What is Interpolation?

Interpolation is a numerical method where:

- You are given n+1n+1 points: $(x0,y0),(x1,y1),...,(xn,yn)(x_0,y_0),(x_1,y_1), dots,(x_n,y_n)$
- These yiy_i values are from an unknown function y=f(x)y = f(x), evaluated at known xix_i points.
- The **goal** is to estimate yy at some new value xx **not in the list** $x0,x1,...,xnx_0, x_1,...,x_n$.

Types of Interpolation

- 1. Polynomial Interpolation (Global)
 - Fit a **single polynomial** of degree nn through all n+1n+1 points.
 - Accurate for **2–3 points**, but:
 - For many points, causes unwanted oscillations ("wiggles").
 - Not preferred for large datasets.
- 2. Piecewise Linear Interpolation (Local)



- Use a **straight line between each pair** of consecutive points.
- Simple, fast, and avoids wiggles.

- Used by MATLAB's plot() function.
- 3. Cubic Spline Interpolation (explained in next videos)
 - o Uses piecewise cubic polynomials.
 - o Results in a **smoother** curve (continuous derivatives).

Piecewise Linear Interpolation – How It Works

Define a piecewise linear function g(x)g(x), such that:

g(x)=gi(x) for $x \in [xi,xi+1]g(x) = g_i(x) \quad \text{(in } [x_i,x_{i+1}]$

Each segment gi(x)g_i(x) is a line:

$$gi(x)=ai(x-xi)+big_i(x) = a_i(x - x_i) + b_i$$

Now to Find aia_i and bib_i

You have 2 constraints for each line:

- 1. $gi(xi)=yig_i(x_i)=y_i$
- 2. $gi(xi+1)=yi+1g_i(x_{i+1}) = y_{i+1}$

From these:

- bi=yib_i = y_i
- $ai=yi+1-yixi+1-xia_i = \frac{y_{i+1}-y_i}{x_{i+1}-x_i} \leftarrow this is the slope$

Graphical Understanding

- Points: (x0,y0),(x1,y1),(x2,y2),...(x_0, y_0), (x_1, y_1), (x_2, y_2), \dots
- Connect each pair with a straight line \rightarrow forms a **piecewise linear** curve g(x)g(x).
- Function is continuous, but the slope (derivative) is not.

Summary

Concept	Explanation
Given	n+1n+1 data points (xi,yi)(x_i, y_i)
Goal	Estimate yy for xx not in the data
Method	Piecewise linear functions gi(x)g_i(x)
Continuity	g(x)g(x) is continuous

Concept **Explanation**

Derivatives Discontinuous at data points

Form of Each Segment $gi(x)=ai(x-xi)+yig_i(x)=a_i(x-x_i)+y_i$

Used When Simple, quick interpolation with few assumptions

Let me know if you want a MATLAB implementation, code to visualize it, or a comparison with cubic spline interpolation.

Here's a summary of the main points from the video on Cubic Spline Interpolation (Part A) by Prof. Jeff Chasnov:



Q Objective:

To interpolate a function using **piecewise cubic polynomials** — a smoother alternative to piecewise linear interpolation.

Problem Setup:

Given n+1n+1 data points:

 $(x0,y0),(x1,y1),...,(xn,yn)(x_0,y_0),(x_1,y_1), dots,(x_n,y_n)$

 Goal: Find a smooth function g(x)g(x) such that g(xi)=yig(x_i) = y_i, and also get good estimates of g(x)g(x) at other x-values.

Cubic Spline Interpolation:

- Break the x-intervals into **n subintervals**.
- On each interval [xi,xi+1][x_i, x_{i+1}], define a **cubic polynomial**:

 $gi(x)=ai(x-xi)3+bi(x-xi)2+ci(x-xi)+dig_i(x) = a_i(x-x_i)^3+b_i(x-x_i)^2+c_i(x-x_i)+d_i$

Total unknowns: 4n4n (since each of the nn intervals has 4 coefficients).

Nonstraints (to determine the unknowns):

- 1. Interpolation Constraints (2n):
 - o Pass through both endpoints of each interval:

 $gi(xi)=yi,gi(xi+1)=yi+1g_i(x_i) = y_i,\quad g_i(x_{i+1}) = y_{i+1}$

2. First Derivative Continuity (n-1):

o Ensure slope continuity at internal nodes:

$$gi'(xi+1)=gi+1'(xi+1)g'_i(x_{i+1}) = g'_{i+1}(x_{i+1})$$

- 3. Second Derivative Continuity (n-1):
 - o Ensure curvature continuity at internal nodes:

$$gi''(xi+1)=gi+1''(xi+1)g''_i(x_{i+1})=g''_{i+1}(x_{i+1})$$

Counting Constraints:

- Interpolation: 2n2n
- First derivative continuity: n-1n-1
- Second derivative continuity: n-1n-1

Total so far: 2n+2(n-1)=4n-22n+2(n-1)=4n-2

→ Missing 2 constraints to fully determine the spline.

Two Extra Constraints (needed):

 Will be discussed in the next video (e.g., "not-a-knot" condition or specifying second derivatives at the boundaries).

Defining Helpful Quantities:

- Interval width: hi=xi+1-xih_i = x_{i+1} x_i
- Output difference: ηi=yi+1-yi\eta_i = y_{i+1} y_i

Summary of Coefficient Determination:

- 1. di=yid_i = y_i
- 2. $aihi3+bihi2+cihi+di=yi+1a_i h_i^3 + b_i h_i^2 + c_i h_i + d_i = y_{i+1}$
- 3. First derivative continuity at xi+1x_{i+1}:

$$3aihi2+2bihi+ci=ci+13a_i h_i^2 + 2b_i h_i + c_i = c_{i+1}$$

4. Second derivative continuity at xi+1x_{i+1}:

$$6aihi+2bi=2bi+16a_i h_i + 2b_i = 2b_{i+1}$$

Insight:

You don't usually compute the spline by hand — MATLAB or Python can do it using built-in functions — but understanding the derivation helps you understand how the interpolation is smooth and continuous.

Let me know if you'd like a simplified derivation or code implementation in Python or MATLAB.

Here's a **summary and explanation** of the full cubic spline interpolation analysis from both parts of Jeff Chasnov's lecture:



To construct a **smooth interpolating function** that passes through n+1n+1 known points $(x0,y0),...,(xn,yn)(x_0,y_0), \dots, (x_n,y_n)$ using **piecewise cubic polynomials** (called cubic splines).

Structure of the Interpolant

We build nn cubic polynomials:

gi(x)=ai(x-xi)3+bi(x-xi)2+ci(x-xi)+di, for i=0 to $n-1g_i(x)=a_i(x-x_i)^3+b_i(x-x_i)^2+c_i(x-x_i)+d_i$, \quad \text{for } i=0 \text{ to } n-1

Each has 4 coefficients, leading to 4n unknowns.

- **✓** Constraints (Total: 4n 2 before final fix)
- 1. Function passes through data points (2n constraints):
 - gi(xi)=yi⇒di=yig_i(x_i) = y_i \Rightarrow d_i = y_i
 - gi(xi+1)=yi+1⇒gives equation in ai,bi,ci,dig_i(x_{i+1}) = y_{i+1} \Rightarrow \text{gives equation in } a_i, b_i, c_i, d_i
- 2. First derivative continuity at internal knots (n 1 constraints):

$$gi'(xi+1)=gi+1'(xi+1)g_i'(x_{i+1})=g_{i+1}'(x_{i+1})$$

3. Second derivative continuity at internal knots (n – 1 constraints):

$$gi''(xi+1)=gi+1''(xi+1)g_i''(x_{i+1}) = g_{i+1}''(x_{i+1})$$

▼ Total so far: 2n+(n-1)+(n-1)=4n-22n + (n - 1) + (n - 1) = 4n - 2

Still Need 2 More Constraints

To fully solve for the 4n coefficients, 2 more equations are needed.

☑ Not-a-Knot Condition (Default in MATLAB)

This assumes the spline is "extra smooth" at the first and last interior points:

Require the third derivatives to be equal at the second and second-to-last point:

```
g0'''(x1)=g1'''(x1) \Rightarrow a0=a1g_0'''(x_1) = g_1'''(x_1) \setminus Rightarrow a_0 = a_1
gn-2'''(xn-1)=gn-1'''(xn-1) \Rightarrow an-2=an-1g \{n-2\}'''(x \{n-1\}) = g \{n-1\}'''(x \{n-1\}) \setminus Rightarrow a \{n-2\} = gn-1 = g
a_{n-1}
```



✓ Now total constraints = 4n

Final Form: Matrix Equation

After substituting and eliminating aia_i, cic_i, and did_i, we get a tridiagonal matrix equation for the bib i coefficients:

```
13hibi+23(hi+hi+1)bi+1+13hi+1bi+2=\eta i+1hi+1-\eta ihi frac{1}{3} h_i b_i + \frac{2}{3}(h_i + h_{i+1})
b_{i+1} + \frac{1}{3} h_{i+1} b_{i+2} = \frac{(i+1)}{h_{i+1}} - \frac{1}{h_i}
```

Where:

- $hi=xi+1-xih_i = x_{i+1}-x_i$
- ηi=yi+1-yi\eta_i = y_{i+1} y_i

This leads to a **linear system** of size $(n+1)\times(n+1)(n+1)$ \times (n+1).

Solving It

- This tridiagonal system is solved using matrix operations (e.g., LU decomposition).
- Once bib i's are found, plug them back to find aia i, cic i, and did i.

👨 💻 In Practice

- You don't solve this by hand.
- Use built-in software routines (e.g., spline() in MATLAB or scipy.interpolate.CubicSpline() in Python).
- MATLAB uses not-a-knot by default, but can be changed to clamped, natural, or other types of boundary conditions.

Key Advantages of Cubic Splines

- Smoothness: Continuous first and second derivatives.
- Stability: Less oscillatory than high-degree single polynomials.
- **Accuracy**: Great for smooth interpolation of real-world data.

Would you like me to generate the matrix form visually or in code (MATLAB/Python)?

Here's a **summary and breakdown** of Jeff Chasnov's MATLAB demonstration on **cubic spline interpolation using interp1**:



Use MATLAB's interp1 function to interpolate known data points and evaluate how well **cubic spline interpolation** reconstructs a smooth function—in this case, the sine function.

Step-by-Step in MATLAB

1. Define Sample Data Points

Create x values from 0 to $2\pi 2\pi$, spaced by $\pi 4\frac{\pi}{4}$, and compute their sine values.

x = 0:pi/4:2*pi;

 $y = \sin(x);$

2. Plot Original Points

plot(x, y, 'o'); % circle markers for original data

hold on;

You'll see widely spaced points lying on the sine curve.

3. Create a Fine Grid for Interpolation

xx = linspace(0, 2*pi, 1000); % finer x-grid

This gives a smoother appearance when interpolating.

4. Perform Interpolation Using Cubic Spline

yy = interp1(x, y, xx, 'spline');

- 'spline' uses **cubic spline interpolation** with **not-a-knot** condition.
- You can also use 'linear' if you want piecewise linear interpolation.

5. Plot the Interpolated Curve

plot(xx, yy, 'r'); % red interpolated curve

4 6. Compare with Actual Sine Function

plot(xx, sin(xx), 'b'); % actual sine in blue

When you run this, you'll notice:

- The **red** spline-interpolated curve follows the **blue** actual sine curve very closely.
- With zooming in, small differences can be observed—especially at the edges.

Conclusion

- interp1 with 'spline' gives a **very accurate and smooth interpolation**, especially useful for approximating smooth functions like sin(x).
- MATLAB's default spline method is not-a-knot, making the interpolated curve extra smooth at boundaries.

% Useful Command Summary

```
x = 0:pi/4:2*pi;
y = sin(x);
xx = linspace(0, 2*pi, 1000);
yy = interp1(x, y, xx, 'spline');
plot(x, y, 'o'); hold on;
plot(xx, yy, 'r');
plot(xx, sin(xx), 'b');
```

Would you like a version of this in Python with scipy.interpolate as well?

Here's a clear **summary and guidance** for your project on **finding zeros of Bessel functions** based on Jeff Chasnov's video:

OProject Objective

Find the first five zeros of the first six Bessel functions $Jn(x)J_n(x)$ for n=0,1,...,5n=0,1, \ldots, 5.

What You Need to Do

1. Understand the Bessel function definition:

The $nthn^{\text{text}}$ Bessel function of the first kind $Jn(x)J_n(x)$ is defined by the integral:

 $Jn(x)=1\pi \int 0\pi \cos(x\sin(\theta-n\theta)) d\theta. \\ J_n(x)=\frac{1}{\pi} \int 0\pi \cos(x\sin(\theta-n\theta)) d\theta.$

2. Plot the Bessel functions to understand their shape:

- For each nn, plot Jn(x)J_n(x) for a range of xx.
- The zeros of $Jn(x)J_n(x)$ are the points where it crosses zero on the xx-axis.
- 3. Find zeros (roots) of Jn(x)J_n(x):
- You want to find the first five zeros for each JnJ_n.
- This is a **root-finding problem** where you find xx such that $Jn(x)=0J_n(x)=0$.
- 4. Tools to use in MATLAB:
- integral.m to numerically evaluate the integral defining Jn(x)J_n(x).
- fzero.m to find roots of Jn(x)J_n(x).
- 5. Code Structure:
- Use **nested loops**: outer loop over n=0,...,5n = 0, \ldots, 5 (which Bessel function), inner loop over zeros j=1,...,5j=1, \ldots, 5.
- Use initial guesses for the zeros to feed into fzero. The zeros can be estimated by plotting or taken from tables (provided initial guesses in your template).
- Refine each root with fzero to get accurate numerical values.

MATLAB Implementation Sketch

```
% Function to compute J_n(x) via integral
function val = Jn(n, x)
  integrand = @(theta) cos(x .* sin(theta) - n .* theta);
  val = (1/pi) * integral(integrand, 0, pi);
end

% Number of zeros and functions
numZeros = 5;
maxN = 5;

% Initial guesses for zeros (can be from a table or approximate)
initialGuesses = [...]; % 6x5 matrix or cell array, each row for n, each col for zero index

for n = 0:maxN
  fprintf('Zeros for J_%d(x):\n', n);
```

```
for k = 1:numZeros
  root = fzero(@(x) Jn(n, x), initialGuesses(n+1, k));
  fprintf(' Zero %d: %.6f\n', k, root);
  end
end
```

P Tips

- Plot Jn(x)J_n(x) first for n=0,...,5n=0,...,5 and xx in a reasonable range to visually identify where zeros lie.
- Use linspace and a vectorized form of your integral function for plotting.
- The function fzero requires an initial guess near the root.
- The **not-a-knot** condition or interpolation is not used here; this project focuses on integral evaluation and root finding.

If you want, I can help you build a **complete MATLAB script** for this project or even translate it into **Python** with scipy.integrate.quad and scipy.optimize.root_scalar. Would that help?

Here's a clear summary and some guidance for your **Week 5** content in Numerical Methods for Engineers, focusing on **ODEs**:

Week 5 Overview: Solving Ordinary Differential Equations (ODEs)

Why ODEs matter:

- Many engineering problems are modeled by ODEs.
- Example: Newton's Second Law F=maF = ma leads to a second-order ODE describing motion.

Topics you will cover:

- 1. **Initial Value Problems (IVPs)** solving ODEs given initial conditions.
- 2. **Euler's Method** the simplest numerical integration method.
- 3. **Modified Euler Method** improves accuracy by correcting the slope estimate.
- 4. **Runge-Kutta Methods** more accurate, widely used methods (including adaptive versions).
- 5. **Adaptive Runge-Kutta Methods** algorithms that automatically adjust the step size for efficiency and accuracy.
- 6. MATLAB's ode45 an implementation of an adaptive Runge-Kutta method you'll use.
- 7. **Higher-order ODEs and systems** transforming them into first-order systems for numerical solving.

8. **Shooting Method** — to solve two-point boundary value problems by converting them to root-finding problems.

X Your Week 5 Project: The Two-Body Problem

- Model the motion of two masses interacting gravitationally via Newton's Laws.
- You will solve the system of ODEs describing their positions and velocities.
- This will involve numerical integration of the equations of motion.
- Likely you will use methods like Runge-Kutta (or MATLAB's ode45) to simulate their trajectories.

Tips for Success:

- Review how to convert second-order ODEs into first-order systems (by defining velocity and position variables).
- Practice implementing Euler and Runge-Kutta methods on simpler problems.
- Familiarize yourself with MATLAB's ode45 syntax and options.
- Understand the shooting method basics if you encounter boundary value problems.

If you want, I can help you by:

- Providing example MATLAB code for Euler, modified Euler, or Runge-Kutta methods.
- Guiding you step-by-step on how to implement the two-body problem solver.
- Explaining the shooting method or adaptive time stepping.

Would you like to start with any of those?

Here's a concise breakdown of the **Euler Method** explanation from the video transcript, plus some key points:

Euler Method for Solving ODEs: Overview

Problem Setup:

Solve an ODE of the form:

 $dxdt=f(t,x)\left(f(t,x)\right)$

Initial condition:

$$x(t0)=x0x(t_0) = x_0$$

Goal: Find x(t)x(t) for t>t0t > t_0.

How Euler's Method Works:

- 1. Start at initial point $(t0,x0)(t_0,x_0)$.
- 2. Compute slope (derivative) at this point:

 $f(t0,x0)=dxdt|t=t0,x=x0f(t_0,x_0) = \frac{dx}{dt}\Big[= t_0,x=x_0\Big]$

3. Approximate solution at the next time step $t1=t0+\Delta tt_1=t_0+\Delta tt_1=t_0$ + \Delta t by moving along the tangent line:

 $x1=x0+\Delta t \cdot f(t0,x0)x_1 = x_0 + \Delta t \cdot f(t_0,x_0)$

4. Repeat for subsequent steps:

 $xn+1=xn+\Delta t \cdot f(tn,xn)x_{n+1} = x_n + \Delta t \cdot f(t_n,x_n)$

Intuition:

- Euler's method marches forward by using the slope at the current point to predict the next value.
- It assumes the slope stays constant over the small interval $\Delta t \triangle t$.
- This is like stepping along the tangent line at each point.

Limitations:

- Since slope can change over the interval, this method accumulates error.
- Euler's method is **first order accurate** error per step scales with $\Delta t \geq \Delta t^2$, but overall error scales with $\Delta t \geq \Delta t$.
- Not very accurate for stiff or complex problems.

Next Steps:

- More accurate methods improve by sampling slopes at multiple points in the interval.
- This leads to **Runge-Kutta methods**, which you will study next.
- Runge-Kutta methods reduce error and improve stability.

If you want, I can help you:

- Write MATLAB code for Euler's method on a simple ODE.
- Explain the error behavior with examples.
- Introduce the modified Euler or basic Runge-Kutta methods next.

Would you like to start coding Euler's method or learn more theory?

Here's a clear summary of the **Modified Euler Method** (also known as the second-order Runge-Kutta method, RK2) explained in the video:

Modified Euler Method (RK2)

Starting Point:

We want to solve:

 $dxdt=f(t,x)\frac{dx}{dt} = f(t,x)$

with initial condition xnx_n at time tnt_n.

Problem with Euler Method:

- Euler uses only the slope at the start of the interval tnt_n.
- But slope generally changes over the interval $[tn,tn+\Delta t][t_n,t_n+\Delta t]$.
- So Euler's method is not very accurate.

Idea of Modified Euler:

- Use two slope evaluations:
 - o f(tn,xn)f(t_n, x_n), the slope at the start.
 - o $f(tn+\Delta t,xn+1)f(t_n + \Delta t,x_{n+1})$, the slope at the end.
- Take the average slope to get better accuracy:

 $xn+1=xn+\Delta t2[f(tn,xn)+f(tn+\Delta t,xn+1)]x_{n+1}=x_n+ \frac{1}{2} \left[f(t_n,x_n)+f(t_n+\Delta t,x_{n+1})\right] x_{n+1}=x_n+\frac{1}{2} \left[f(t_n,x_n)+f(t_n+\Delta t,x_{n+1})\right] x_{n+1}$

Issue:

- xn+1x_{n+1} appears on both sides of the equation implicit formula.
- To solve, use a **predictor-corrector approach**:
 - 1. Predictor (Euler step):

 $xn+1(p)=xn+\Delta t \cdot f(tn,xn)x_{n+1}^{(p)} = x_n + \Delta t \cdot f(t_n,x_n)$

2. Corrector (use predicted value):

 $xn+1=xn+\Delta t2[f(tn,xn)+f(tn+\Delta t,xn+1(p))]x_{n+1}=x_n+\frac{1}{2}\left[f(t_n,x_n)+f(t_n+\Delta t,x_{n+1}^{(p)})\right] x_{n+1}=x_n+\frac{1}{2}\left[f(t_n,x_n)+f(t_n+\Delta t,x_n+1(p))\right] x_{n+1}=x_n+\frac{1}{2}\left[f(t_n,x_n+1(t_n+\Delta t,x_n+1(p))\right] x_{n+1}=x_n+\frac{1}{2}\left[f(t_n,x_n+1(t_n+\Delta t,x_n+1(p))\right] x_{n+1}=x_n+\frac{1}{2}\left[f(t_n,x_n+1(t_n+\Delta t,x_n+1(t_n+\Delta t,x_n$

This improves accuracy by averaging the slopes at the start and predicted end of the interval.

Runge-Kutta Notation:

Define stages K1K_1 and K2K_2:

 $K1=\Delta t \cdot f(tn,xn)K_1 = \Delta t \cdot f(tn,x_n)K_2 = \Delta t \cdot f(tn+\Delta t,x_n+K_1)K_2 = \Delta t \cdot f(tn+\Delta t,x_n+K_1)K_2 = \Delta t \cdot f(tn,x_n)K_1 = \Delta t \cdot f(tn$

• Then:

 $xn+1=xn+12(K1+K2)x_{n+1} = x_n + \frac{1}{2}(K_1 + K_2)$

Summary:

- Modified Euler method is a simple 2nd order Runge-Kutta method.
- More accurate than Euler because it uses the slope at both the beginning and the end of the interval.
- It introduces the idea of "stages" KiK_i, foundational for more advanced Runge-Kutta methods.
- Next step: generalize to other RK methods that sample slope at other points and have higher order accuracy.

If you'd like, I can help you:

- Write MATLAB or Python code for this Modified Euler method.
- Show how the error compares with Euler's method on an example ODE.
- Explain how to extend this to 4th order Runge-Kutta (RK4), the most common method.

Would you like to do one of these?

Here's a detailed but clear explanation of the **full family of second-order Runge-Kutta (RK2) methods**, based on the video:

General Second-Order Runge-Kutta Methods (RK2)

We want to solve the ODE:

 $dxdt=f(t,x)\frac{dx}{dt} = f(t,x)$

from time tnt_n to $tn+1=tn+\Delta tt_{n+1} = t_n + \Delta tt_n$, starting with value xnx_n .

The General RK2 Scheme

1. Stage 1:

 $K1=\Delta t \cdot f(tn,xn)K_1 = \Delta t \cdot f(t_n,x_n)$

2. Stage 2:

 $K2=\Delta t \cdot f(tn+\alpha \Delta t, xn+\beta K1)K_2 = \Delta t \cdot f(t_n + \alpha t, \alpha x_n + \beta K1)K_2 = \Delta t \cdot f(t_n + \alpha t, \alpha x_n + \beta K1)K_2 = \Delta t \cdot f(t_n + \alpha t, \alpha t, \alpha t, \alpha t)$ where $\alpha \cdot f(t_n + \alpha t, \alpha t)$ where $\alpha \cdot f(t_n + \alpha t, \alpha t)$ where $\alpha \cdot f(t_n + \alpha t, \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ and $\beta \cdot f(t_n + \alpha t)$ where $\alpha \cdot f(t_n + \alpha t)$ and $\beta \cdot f($

3. Update:

$$xn+1=xn+aK1+bK2x_{n+1} = x_n + aK_1 + bK_2$$

where aa and bb are also parameters to be determined.

Goal: Determine α,β,a,b \alpha, \beta, a, b such that

The method is **second-order accurate**, i.e., error is proportional to $(\Delta t)3(\Delta t)^3$.

How to find constraints on the parameters?

Step 1: Taylor series expansion of the true solution

 $x(tn+\Delta t)=xn+\Delta t\cdot f(tn,xn)+\Delta t22\cdot x''(tn)+O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(t_n,x_n) + \Delta t^2 \cdot x''(t_n) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(t_n) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t^2 \cdot x''(tn) + O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t\cdot f(tn,xn) + \Delta t\cdot f(tn,xn)$

where $x''(tn)x''(t_n)$ is the second derivative with respect to time.

Using the chain rule:

 $x''(tn) = ddtf(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot dxdt = f_t + f_x \cdot f_x''(t_n) = \frac{d}{dt} f(t,x) = f_t + f_x \cdot f_x''(t_n) = f_t + f_x \cdot f_x''(t_n) = f_t + f_x \cdot f_x''(t_n) = f_t$

with partial derivatives $ft=\partial f \partial t f_t = \frac{f}{\pi x} = \frac{f}{\pi x}.$

Thus,

 $x(tn+\Delta t)=xn+\Delta tf+\Delta t 22(ft+fxf)+O(\Delta t3)x(t_n + \Delta t) = x_n + \Delta t f + frac(\Delta t^2)\{2\} (f_t + f_x + G(\Delta t^3))$

Step 2: Taylor series expansion of the RK2 method

Expand the second term $f(tn+\alpha\Delta t,xn+\beta K1)f(t_n + \alpha t,x_n + \beta K1)f(t_n,x_n)(t_n,x_n)$:

 $f(tn+\alpha\Delta t,xn+\beta K1)=f+\alpha\Delta tft+\beta K1fx+O(\Delta t2)f(t_n + \alpha t,x_n + \beta K1)=f+\alpha\Delta tft+\beta K1fx+O(\Delta t2)f(t_n + \beta t,x_n + \beta K1)=f+\alpha\Delta tft+\beta K1fx+O(\Delta t2)f(t_n + \beta t,x_n + \beta K1)=f+\alpha\Delta tft+\beta K1fx+O(\Delta t2)f(t_n + \beta t,x_n + \beta K1)=f+\alpha\Delta tft+\beta K1fx+O(\Delta t2)f(t_n + \beta t,x_n + \beta K1)=f+\alpha\Delta tft+\beta K1fx+O(\Delta t2)f(t_n + \beta t,x_n + \beta K1)=f+\alpha\Delta tft+\beta K1fx+O(\Delta t2)f(t_n + \beta t,x_n + \beta K1)=f+\alpha\Delta tft+\beta K1fx+O(\Delta t2)f(t_n + \beta t,x_n + \beta t,x$

Since $K1=\Delta tfK$ 1 = \Delta t f, we have:

 $f(tn+\alpha\Delta t,xn+\beta K1)=f+\alpha\Delta tft+\beta\Delta tffx+O(\Delta t2)f(t_n+\alpha t,x_n+\beta K1)=f+\alpha\Delta tft+\beta\Delta tft+\Delta tft+\Delta$

Substitute back:

 $xn+1=xn+a\Delta tf+b\Delta t(f+\alpha\Delta tft+\beta\Delta tfxf)+O(\Delta t3)x_{n+1}=x_n+a \Delta tf+b\Delta tfxf)+O(\Delta t5)x_{n+1}=x_n+a \Delta tf+b\Delta tf+$

Collect terms:

 $xn+1=xn+(a+b)\Delta tf+b\alpha\Delta t2ft+b\beta\Delta t2fxf+O(\Delta t3)x_{n+1}=x_n+(a+b) \Delta tf+b\alpha\Delta t2ft+b\alpha\Delta t2$

Step 3: Equate coefficients with the true solution expansion

Match terms with the exact solution's Taylor series:

Term From exact solution From RK2 method

 $\Delta tf \setminus Delta t f$ ff (a+b)f(a+b) f

 $\Delta t22ft\frac{1}{2} f_t = 12ft\frac{1}{2} f_t = b\alpha ftb \alpha f_t$

 $\Delta t22fxf\frac{1}{2} f_x f 12fxf\frac{1}{2} f_x f b\beta fxfb \beta f_x f$

This gives the system of equations:

Interpretation

- We have **four parameters** $(a,b,\alpha,\beta)(a,b,\lambda)$ beta) but only **three constraints**.
- So there is one degree of freedom, meaning infinitely many RK2 methods.
- Different choices lead to famous RK2 variants.

Examples of RK2 methods satisfying these conditions:

Method	α\alpha	β\beta	aa	bb
Modified Euler (Heun's)	1	1	0.5	0.5
Midpoint method	0.5	0.5	0	1
Ralston's method	2/32/3	3/43/4	1/41/4	3/43/4

Summary

- The **general RK2 method** is a weighted average of two slopes.
- The slopes are evaluated at:

- The beginning of the interval (tn,xn)(t_n, x_n)
- O A point inside the interval $(tn+\alpha\Delta t,xn+\beta K1)(t_n + \alpha Delta t,x_n + \beta K_1)$
- The weights a,ba, b and parameters α,β \alpha, \beta satisfy constraints ensuring **second-order accuracy**.
- This creates a whole **family of methods** with the same order but slightly different properties.

If you want, I can help you:

- See a step-by-step algebra derivation from the Taylor expansions.
- Implement one or more of these RK2 methods in code.
- Compare the numerical performance of different RK2 variants on a test problem.

Would you like to explore any of those?

Great! Here's a concise summary of what you just described about deriving specific second-order Runge-Kutta methods from the general family:

Recap: General second-order Runge-Kutta method

- We solve $\dot{x} = f(t,x) \cdot dot\{x\} = f(t,x)$ from tnt n to tn+1=tn+ Δ tt $\{n+1\} = t$ n + Δ tt $\{n+1\} = t$ n
- Two stages:

 $k1=\Delta t \cdot f(tn,xn)k_1 = \Delta t \cdot f(tn+\alpha\Delta t,xn+\beta k1)k_2 = \Delta t \cdot f(tn+$

• The update step:

 $xn+1=xn+ak1+bk2x_{n+1} = x_n + ak_1 + bk_2$

• The parameters α,β,a,b heta, a, b satisfy these constraints for second-order accuracy:

 $a+b=1,\alpha b=12,\beta b=12a+b=1, \quad b=\frac{1}{2}, \quad b=1,\alpha b=1,\alpha b=12a+b=1,\alpha b=12a+b=12a+b=1,\alpha b=12a+b=12a+b=1,\alpha b=12a+b=12a+b=1,\alpha b=12a+b=12a+b=12a$

1. Modified Euler method

- Known as the improved Euler or Heun's method.
- Parameters:

 $a=12,b=12,\alpha=1,\beta=1a = \frac{1}{2}, \quad b=1,\beta=1a = 1, \quad b=1,\beta=1a = 1$

• So:

 $k1=\Delta t \cdot f(tn,xn)k_1 = \Delta t \cdot f(tn+\Delta t,xn+k1)k_2 = \Delta t \cdot f(tn+\Delta t,xn+k1)k_2 = \Delta t \cdot f(tn+\Delta t,xn+k1)k_1 = \Delta t \cdot f(tn+\Delta t,xn+k1)k_2 = \Delta t \cdot f(tn+\Delta t,xn+k1)k_$

• Update:

2. Midpoint method

- Evaluates the slope at the midpoint of the interval.
- Choose:

 $\alpha=12$ \alpha = \frac{1}{2}

• From constraints:

 $b=1,\beta=12,a=0b=1, \quad b=1,\beta=12,a=0b=1, \quad a=0$

So:

 $k1=\Delta t \cdot f(tn,xn)k_1 = \Delta t \cdot f(tn,x_n) k2=\Delta t \cdot f(tn+\Delta t^2,x_n+12k_1)k_2 = \Delta t \cdot f(tn+\Delta t^2,x_n+12k_1)k_2 = \Delta t \cdot f(tn,x_n)k_1 = \Delta t$

• Update:

 $xn+1=xn+k2x_{n+1} = x_n + k_2$

Summary

- Both methods are second-order accurate, but they differ in how they weigh the intermediate slopes.
- Modified Euler averages slopes at the start and end of the interval.
- Midpoint method uses the slope at the midpoint for the update.
- You can generate other second-order Runge-Kutta methods by choosing one parameter and solving the constraints accordingly.

If you want, I can help you write out code implementations for these methods or explain how they compare in practice! Would you like that?

Here's a clear summary of the key points from your detailed explanation about higher-order Runge-Kutta methods, focusing especially on the third and fourth order methods:

Generalizing Runge-Kutta Methods Beyond Second Order

1. Setup for Third-Order Runge-Kutta (RK3)

- We solve $x'=f(t,x) \cdot dot\{x\} = f(t,x)$ from tnt_n to tn+1=tn+ $\Delta tt_{n+1} = t_n + \Delta tt_{n+1}$.
- Three stages (three evaluations of ff):

 $k1=\Delta t \cdot f(tn,xn)k_1 = \Delta t \cdot f(t_n,x_n) \ k2=\Delta t \cdot f(tn+\alpha \Delta t,xn+\beta k1)k_2 = \Delta t \cdot f(t_n+\alpha t,xn+\beta k1)k_2 = \Delta t \cdot f(t_n+\alpha t,xn+\beta t,xn+\beta t)k_2 = \Delta t \cdot f(t_n+\alpha t,xn+\beta t)k_2 =$

• Update formula combines the three slopes:

 $xn+1=xn+ak1+bk2+ck3x_{n+1} = x_n + ak_1 + bk_2 + ck_3$

- Parameters $\alpha, \beta, \gamma, \delta, \epsilon, a, b, c$ alpha, \beta, \gamma, \delta, \epsilon, a, b, c must satisfy constraints derived via Taylor series expansions for third-order accuracy.
- Third-order methods require 3 stages (3 evaluations of ff).

2. Number of Stages vs. Order

• Second order: 2 stages

• Third order: 3 stages

• Fourth order: 4 stages

• **Fifth order**: 6 stages (not 5 — it requires more)

• Sixth order: 7 stages

• Seventh order: 9 stages

• Eighth order: 11 stages

Higher-order methods quickly require more function evaluations, increasing computational cost.

3. The Popular Fourth-Order Runge-Kutta Method (RK4)

- Four stages, evaluating ff at:
 - 1. The start of the interval: tnt_n
 - 2. The midpoint (twice): $tn+\Delta t2t_n + \frac{\Delta t}{2}$
 - 3. The end of the interval: $tn+\Delta tt_n + \Delta tt_n$
- Stages:

 $k1=\Delta t \cdot f(tn,xn)k_1 = \Delta t \cdot f(t_n,x_n) \ k2=\Delta t \cdot f(tn+\Delta t_2,xn+k_{12})k_2 = \Delta t \cdot f(t_n+\Delta t_2,xn+k_{12})k_2 = \Delta t \cdot f(t_n+\Delta t_2,xn+k_{12})k_3 = \Delta t \cdot f(tn+\Delta t_2,xn+k_{12})k_3 = \Delta t \cdot f(t_n+\Delta t_2,xn+k_{12})k_3 = \Delta t \cdot f(t_n+\Delta t_2,xn+k_{12})k_3 = \Delta t \cdot f(t_n+\Delta t_2,xn+k_{12})k_4 = \Delta t \cdot f(t_n+\Delta t_2,xn+k_{12})$

• Update:

 $xn+1=xn+16(k1+2k2+2k3+k4)x_{n+1} = x_n + \frac{1}{6}(k_1 + 2k_2 + 2k_3 + k_4)$

- This method is fourth-order accurate (local error $\sim \Delta t \le \pi \cdot \Delta t \le \pi \cdot$
- Became very popular due to good balance of accuracy and computational cost.

• Still widely taught and used, though modern software like MATLAB often uses even higherorder adaptive methods.

4. Summary and Practical Notes

- RK4 is a "sweet spot" in accuracy and complexity.
- Deriving coefficients for third and higher-order RK methods requires careful Taylor series expansion and solving constraints.
- Higher-order methods need more stages, increasing computational effort.
- Software packages often handle the choice of method internally.

If you want, I can help you:

- Derive or explain the RK3 coefficients and formulas.
- Write sample code implementing RK3 or RK4.
- Compare RK4 with other integration methods or adaptive solvers.

Would you like to dive deeper into any of those?

Here's a clear summary and explanation of the video content about integrating **second order ODEs** and **systems of ODEs** using Runge-Kutta methods:

Integrating Second Order ODEs by Converting to First Order Systems

The Problem:

• You often encounter **second order ODEs**, e.g., Newton's law:

 $d2xdt2=f(t,x,x')\frac\{d^2 x\}\{dt^2\} = f(t,x,\dot\{x\})$

where the second derivative depends on time tt, position xx, and velocity $x' \cdot dot\{x\}$.

The Key Trick: Define a New Variable

• Introduce a new variable for the first derivative of xx:

 $u=dxdt=x\dot{u} = \frac{dx}{dt} = \det\{x\}$

• Now rewrite the original second order ODE as a system of two first order ODEs:

 $\{dxdt=ududt=f(t,x,u)\geq f(t,x,u)\} = u \setminus f(t,x,u) = f(t,x,u) = f(t,x,u)$

• This converts one second order ODE into two coupled first order ODEs.

Generalizing to Systems of Coupled First Order ODEs

Consider a system of two coupled first order ODEs:

 $\{dxdt = f(t,x,y)dydt = g(t,x,y) \land g(t,x,y$

• You want to integrate **both simultaneously**.

Applying a Runge-Kutta Method (Modified Euler as example) to Systems

Parallel integration approach:

- For the variable xx, define stages with kk terms.
- For the variable yy, define stages with II terms.

First stage:

 $k1=\Delta t \cdot f(tn,xn,yn)k_1 = \Delta t \cdot g(tn,xn,yn)l_1 = \Delta t \cdot g(tn,xn,yn)$

Second stage:

```
k2=\Delta t \cdot f(tn+\Delta t, xn+k1, yn+l1)k_2 = \Delta t \cdot g(tn+\Delta t, xn+k1, yn+k1, yn
```

Update step:

```
 xn+1=xn+12(k1+k2)x_{n+1} = x_n + \frac{1}{2}(k_1 + k_2) yn+1=yn+12(l1+l2)y_{n+1} = y_n + \frac{1}{2}(l_1 + l_2)
```

Summary:

- **Second order ODEs** can be transformed into **systems of first order ODEs** by introducing auxiliary variables.
- When integrating systems of coupled first order ODEs, Runge-Kutta methods are applied in parallel to each variable.
- Each variable has its own set of stages (kk's for xx, II's for yy), but the updates depend on all variables because the derivatives are coupled.
- This approach extends naturally to larger systems or higher order equations converted to first order systems.

If you'd like, I can help you:

- Write example Python code for RK2 or RK4 integration of such systems.
- Show how to handle larger systems or include higher order methods.
- Explain how this applies to physics problems like Newton's equations or coupled oscillators.

Would you like to see that?

Here's a detailed explanation and summary of the video about **MATLAB's adaptive Runge-Kutta method for solving ODEs**, specifically focusing on the Dormand-Prince method used by ode45:

MATLAB's Adaptive Runge-Kutta ODE Solver (ode45)

The Problem with Fixed Time Step (Δt)

- Classical Runge-Kutta methods require you to specify a fixed time step Δt\Delta t.
- Choosing Δt\Delta t is tricky:
 - Too large → large errors (solution inaccurate)
 - Too small → very slow, expensive computation
- For complex or large systems, manual tuning of $\Delta t \setminus Delta t$ is inefficient.

Adaptive Time Step and Error Control

- MATLAB uses an adaptive time step algorithm that automatically adjusts Δt\Delta t to control error.
- Instead of specifying Δt\Delta t, you specify the error tolerance ε\epsilon, which is the maximum allowed error per step.
- Smaller tolerance \rightarrow smaller $\Delta t \setminus Delta t \rightarrow more$ accurate but more expensive.
- Larger tolerance \rightarrow larger $\Delta t \setminus Delta t \rightarrow faster but less accurate.$

Dormand-Prince Method (Used in ode45)

- It is a six-stage Runge-Kutta method that provides:
 - A 5th order accurate solution (more accurate)
 - A 4th order accurate solution (less accurate)
- Both solutions are computed **simultaneously** from the same six derivative evaluations.
- This allows estimating the **local error** without extra derivative calculations.

How Error is Estimated

- Let:
 - o $xn+1x_{n+1} = 5$ th order solution (more accurate, error $\sim \Delta t6$ sim \Delta t^6)
 - $Xn+1X_{n+1} = 4th \text{ order solution (less accurate, error } \sim \Delta t5 \sin \Delta t^5$

• The difference between these two solutions approximates the actual error ee of the step:

 $e \approx |xn+1-Xn+1|e \cdot |x_{n+1} - X_{n+1}|$

• Since the 4th order solution has larger error, the difference is dominated by Δt5\Delta t^5, giving a reliable error estimate.

Adjusting the Time Step

- Suppose you did a step with time step Δt\Delta t, got error estimate ee, but want error less than tolerance ϵ\epsilon.
- The error scales like Δt5\Delta t^5, so:

 $e \in -(\Delta t \Delta \tau) \int \frac{e}{\position} = \left(\frac{\Delta t \Delta \tau}{\Delta \tau}\right)^5$

where $\Delta \tau \backslash Delta \$ tau is the *ideal* time step that would yield error $\epsilon \backslash epsilon$.

Solving for Δτ\Delta \tau:

 $\Delta \tau = \Delta t(\epsilon e) 1/5 \Delta t = \Delta t (\epsilon e) 1/5 \Delta t = \Delta t(\epsilon e) 1/5 \Delta t = \Delta t($

- Interpretation:
 - o If e>εe > \epsilon, error too large → reduce time step → reject current solution and redo with smaller Δt = Δt \Delta t = \Delta \tau.
 - If e<∈e < \epsilon, error acceptable → accept solution and maybe increase time step for next step.

Safety Factor

- To avoid too frequent rejections and unstable step sizes, multiply by a **safety factor ss < 1**, typically s=0.9s = 0.9.
- New time step formula becomes:

 $\Delta \tau = s \Delta t \times (\epsilon e) 1/5 \Delta t = s \times (\epsilon e) 1/5 \Delta t = s \Delta t \times (\epsilon e) 1/5 \Delta t = s$

This makes the algorithm more robust.

Summary of MATLAB ode45 Workflow

- 1. User specifies error tolerance $\epsilon \neq \Delta t \Delta t$.
- 2. MATLAB tries a step with some $\Delta t \setminus Delta t$.
- 3. Computes both 4th and 5th order solutions and estimates error ee.
- 4. If $e > \epsilon e > \$ reduces $\Delta t \setminus Delta t$, retries.
- 5. If $e \le e$ \le \epsilon, accepts step, adjusts Δt \Delta t for next step.
- 6. Process repeats until integration completes.

Why This is Important

- Automatically balances accuracy and efficiency.
- User doesn't have to guess the time step.
- Ensures error stays below the specified tolerance.
- Crucial for solving stiff or complex differential equations.

If you want, I can also help you:

- Show example MATLAB code using ode45 and setting tolerances.
- Explain how to interpret solver output and tweak tolerances.
- Compare with fixed time step methods in code.

Would you like that?

Got it! You shared a detailed walkthrough of how to solve ODEs in MATLAB using ode45, covering:

1. Single first-order ODE (Logistic equation):

 $dxdt=rx(1-x)\left(frac\left(dx\right) \right) =rx\left(1-x\right)$

- o Define the right-hand side as an anonymous function.
- Set parameter rr, initial condition x0x_0, and time span.
- Call ode45(@(t,x) r*x*(1-x), tspan, x0).
- \circ Plot the solution x(t)x(t).
- 2. **System of two first-order ODEs** (Lotka-Volterra predator-prey model):

- Define a sub-function that returns a two-component column vector for the derivatives.
- o Set parameters and initial conditions for both xx and yy.
- Call ode45(@(t,xy) lv(xy, r), tspan, [x0; y0]).
- Plot both x(t)x(t) and y(t)y(t), and also yy vs xx phase plot.

3. Notes on accuracy and smoothness:

- o The default solver steps and interpolation might make plots appear less smooth.
- Over longer integration times, numerical errors accumulate, affecting periodicity and plot overlap.
- You can improve accuracy with solver options (not covered yet).

If you want, I can help you write a clean MATLAB script based on this, explain how to add solver options to improve accuracy, or help with any other part of this process. Would you like me to?

Here's a concise summary of how to **improve accuracy and smoothness** when solving ODEs with ode45 in MATLAB, based on the video:

Problem 1: Plots Are Not Smooth

- Cause: ode45 returns solution at a relatively small number of time points (e.g., 57 points), leading to coarse plots.
- **Fix:** Instead of specifying tspan = [t0 tf], specify tspan as a vector with many points using linspace.
- tspan = linspace(0, 6*pi, 500); % 500 points instead of just 2
- [t, xy] = ode45(@(t,xy) lv(xy,r), tspan, [x0; y0]);
- This makes MATLAB interpolate the solution at the requested finer time points, giving a smoother plot.

Problem 2: Solution Is Not Exactly Periodic (Accuracy Issues)

- Cause: Default relative tolerance of ode45 is 10–310^{-3}, which might not be precise enough for subtle periodicity in the solution.
- **Fix:** Use odeset to create options with tighter relative tolerance, e.g., 10–410^{-4} or smaller:
- options = odeset('RelTol', 1e-4);
- [t, xy] = ode45(@(t,xy) lv(xy,r), tspan, [x0; y0], options);
- This improves the numerical accuracy, reducing cumulative errors and ensuring periodicity appears correctly.

Summary MATLAB example incorporating both fixes:

```
r = 1;

x0 = 1.1; y0 = 1;

tspan = linspace(0, 6*pi, 500); % many points for smooth plotting

options = odeset('RelTol',1e-4); % tighter tolerance for accuracy

[t, xy] = ode45(@(t,xy) lv(xy, r), tspan, [x0; y0], options);
```

```
% Plot solutions
figure;
plot(t, xy(:,1), t, xy(:,2));
xlabel('Time t');
ylabel('Population');
legend('Prey (x)', 'Predator (y)');
title('Lotka-Volterra Solution');
figure;
plot(xy(:,1), xy(:,2));
xlabel('Prey (x)');
ylabel('Predator (y)');
title('Phase plot of Lotka-Volterra');
```

Additional Notes:

- Iv function defines the system of ODEs as before.
- Increasing the number of output points smooths plots by forcing interpolation at those points.
- Reducing the error tolerances makes the numerical solver more accurate but might increase computation time.

Would you like me to help you implement these improvements into your full MATLAB code or explain how to tune other ode45 options?

Here's a clear explanation of the **shooting method** for solving two-point boundary value problems (BVPs) as described in the video:

What is a Boundary Value Problem (BVP)?

Given a second-order ODE like

 $d2ydx2=f(x,y,dydx)\left\{ d^2y\right\} = f(x,y,\left\{ dx\right\})$

• Boundary conditions specify values of yy at two points:

```
y(x0)=y0,y(xf)=yfy(x_0) = y_0, \quad y(x_f) = y_f
```

• Goal: find a function y(x)y(x) that satisfies the ODE and these two boundary conditions.

Why is BVP Different from Initial Value Problem (IVP)?

• To solve a second order ODE numerically with methods like ode45, we rewrite as a system of first order ODEs:

 $dydx=z,dzdx=f(x,y,z)\frac\{dy\}\{dx\}=z,\quad\frac\{dz\}\{dx\}=f(x,y,z)$

- For IVP, we know both $y(x0)y(x_0)$ and $z(x0)=dydx(x0)z(x_0)=\frac{dy}{dx}(x_0)$ (initial slope).
- For BVP, $y(x0)=y0y(x_0)=y_0$ is known, but $z(x0)z(x_0)$ is unknown. We call this unknown initial slope parameter α alpha (or *kosai* in the video).

Shooting Method Idea

- The problem becomes: Find α alpha such that when we solve the ODE from $x0x_0$ to xfx_f with initial slope α alpha, the solution satisfies $y(xf) = y_f$.
- Define a function

 $F(\alpha)=y\alpha(xf)-yfF(\alpha)=y_{\alpha}(x_f)-y_f$

where $y\alpha(xf)y_{\alpha}(x_f)$ is the solution at xfx_f after integrating with initial slope α

We want to solve for

 $F(\alpha)=0F(\alpha)=0$

This is a **root-finding problem** for α alpha.

How the Shooting Method Works

- 1. **Guess** α **\alpha:** Pick an initial guess for the unknown initial slope $z(x_0) = \alpha z(x_0) = -\alpha z$
- 2. **Integrate ODE:** Solve the ODE IVP with $y(x0)=y0y(x_0)=y_0$, $z(x0)=\alpha z(x_0)=\lambda alpha using ode45.$
- 3. Check boundary condition at xfx_f: Calculate $F(\alpha) = y_{\alpha}(xf) y_{f}(\alpha) = y_{\alpha}(xf) = y_{\alpha}(xf) y_{f}(\alpha)$
- 4. **Update** α **\alpha:** Use a root-finding method (like fzero in MATLAB) to find a better α \alpha based on the sign and value of $F(\alpha)F(\alpha)$.
- 5. **Repeat:** Continue guessing and integrating until $F(\alpha)F(\alpha)$ is close enough to zero, i.e., the boundary condition at xfx f is satisfied.

Why Is It Called "Shooting"?

Imagine the unknown slope α\alpha as the angle you "aim" a gun at x0x 0.

- Solving the ODE with this slope is like "firing" the bullet along a trajectory (the solution curve).
- Your goal is to "hit the target" yfy_f at xfx_f.
- If you miss, adjust the angle α\alpha and shoot again.
- Repeat until you hit the target.

Summary in MATLAB Context

- Use ode45 inside a function F(α)F(\alpha) that:
 - o Integrates from x0x_0 to xfx_f with initial slope α \alpha,
 - Returns the difference y(xf)-yfy(x_f) y_f.
- Use MATLAB's fzero or other root-finders to find α \alpha such that $F(\alpha)=0F(\alpha)=0$.

Would you like me to help you with a MATLAB code example implementing the shooting method for a sample BVP?

Here's a clear summary of the key points from the video on the **two-body problem** and its reduction to a one-body problem for numerical solution:

Two-Body Problem Setup

- You have two masses, m1m_1 and m2m_2, at positions r1\mathbf{r}_1 and r2\mathbf{r}_2.
- Newton's second law for each mass:

 $m1r"1=F, m2r"2=-Fm_1 \dot{\mathbf{r}}_1 = \mathbb{F}_1 \dot{\mathbf{r}}_2 = -\mathbb{F}_1 \dot{\mathbf{r}}_2 = -\mathbb{F}_1$

• The gravitational force F\mathbf{F} follows Newton's universal law of gravitation:

 $F=-Gm1m2|r1-r2|3(r1-r2)\mathbb{F} = -G \frac{m_1 m_2}{|\mathbb{F}_1 - \mathbb{F}_2|^3} (\mathbf{F}_1 - \mathbb{F}_2)$

Force acts along the vector between the two masses with an inverse-square law.

Coordinate Transformation: From Two-Body to One-Body

Define the relative position vector:

 $r=r1-r2\mathbb{r} - r^2\mathbb{r} -$

• Define the center of mass (COM) position vector:

 $R=m1r1+m2r2m1+m2 \setminus mathbf{R} = \frac{m_1 \mathbb{F}_1 + m_2 \mathbb{F}_1}{m_1 + m_2}$

• Using r\mathbf{r} and R\mathbf{R}, the system can be rewritten:

The acceleration of the center of mass is zero:

 $R''=0\dot{\mathbb{R}} = 0$

Meaning the COM moves uniformly or can be fixed by choice of reference frame.

The relative vector r\mathbf{r} satisfies:

 $\mu r = F \setminus dot{mathbf{r}} = \mathbb{F}$

where μ \mu is the **reduced mass**:

 $\mu=m1m2m1+m2\mu = \frac{m_1 m_2}{m_1 + m_2}$

• This reduces the two-body problem to a single particle with mass μ \mu moving under the central force F\mathbf{F}.

Simplified Equation for Numerical Solution

Writing the equation explicitly:

 $r''=-kr|r|3\dot{\mathbb{r}} = -k \frac{r}{r}$

where

 $k=G(m1+m2)k = G(m_1 + m_2)$

• By non-dimensionalizing, you can scale kk to 1 for easier computation.

Recovering Original Positions

• After solving for r(t)\mathbf{r}(t), you find the original positions by:

Summary

- The two-body problem reduces to solving one second-order ODE for the relative vector r\mathbf{r}.
- The center of mass moves with zero acceleration and can be treated separately.
- After solving r(t)\mathbf{r}(t), original trajectories r1(t)\mathbf{r}_1(t) and r2(t)\mathbf{r}_2(t) follow.
- The numerical solution to the reduced problem lets you plot the orbits of the two masses under gravity.

If you want, I can help you set up the numerical solution in code next, or clarify any of these steps!

Here's a structured summary and explanation of the continuation for numerically solving the two-body problem, based on the video transcript you shared:

Setting Up the Numerical Problem in the Center of Mass Frame

1. Fixing the Center of Mass (COM) Reference Frame

- Since the acceleration of the center of mass R\mathbf{R} is zero (R"=0\ddot{\mathbf{R}} = 0), the center of mass moves with constant velocity.
- Choose a coordinate system moving with this velocity so that:

 $R'=0,R=0\dot{\mathbb{R}}=0,\quad\mathbf{R}=\mathbf{0}$

• This means the center of mass is fixed at the origin for the numerical solution.

2. Reducing to 2D Motion

- If the initial positions r1,r2\mathbf{r}_1, \mathbf{r}_2 and velocity vectors lie in the xy-plane, the motion stays in the xy-plane due to symmetry.
- So, we can **set z=0z = 0** and solve the problem in two dimensions only.

3. Non-Dimensionalization and Choosing Units

- Use the gravitational constant and masses combined into a constant kk.
- Choose units so that:
 - o k=1k = 1
 - o The minimum distance between the two masses is 1 unit.
- This simplification (non-dimensionalization) simplifies the numerical solution by reducing constants.

4. Coordinate System Orientation and Initial Conditions

- Define the coordinate system so that:
 - At t=0t=0, the relative position vector is at its minimum distance along the x-axis:

$$x(0)=-1,y(0)=0x(0)=-1, \quad y(0)=0$$

 The velocity at t=0t=0 is perpendicular to the position vector and directed along the y-axis:

 $x'(0)=0,y'(0)=1+e \cdot \{x\}(0)=0, \quad dot\{y\}(0)= \cdot \{1+e\}$

- Here, **ee** is the eccentricity of the orbit:
 - \circ e=0e = 0 → circular orbit,

5. Final Equations to Solve Numerically

The second-order coupled ODE system for the relative position $r=(x,y) \cdot r=(x,y)$:

```
x"=-x(x2+y2)3/2, y"=-y(x2+y2)3/2 \cdot \{x^2 + y^2\}^{3/2}, \quad \dot\{y\} = -x(x^2+y^2)^{3/2} \cdot \{x^2 + y^2\}^{3/2}
```

This can be rewritten as four first-order ODEs:

6. Initial Conditions for the ODE System

```
x(0)=-1,y(0)=0,vx(0)=0,vy(0)=1+ex(0)=-1, \quad y(0)=0, \quad v_x(0)=0, \quad v_y(0)=0, \quad
```

7. Orbit Period

• The period TT of the orbit depends on ee by:

 $T=2\pi(1-e)3/2T = \frac{2\pi(1-e)^{3/2}}$

You can use this TT to set how long to run the numerical solver to complete one full orbit.

8. Recovering Original Mass Positions

After solving for x(t), y(t), y(t), you can find the positions of the two masses:

```
 r1=m2m1+m2r=m2m1+m2(x,y)\mathbb{f}_1 = \frac{m_2}{m_1+m_2} \mathbb{f}_r = \frac{m_2}{m_1+m_2}(x,y)\mathbb{f}_r = \frac{m_2}{m_1+m_2}\mathbb{f}_r = \frac{m_2}{m_1+m_2}\mathbb{f}_r = \frac{m_1}{m_1+m_2}\mathbb{f}_r = \frac{m_1}{m_1+m_2}(x,y)\mathbb{f}_r = \frac{m_1}{m_1+m_2}(x,y)\mathbb{f}_r = \frac{m_1}{m_1+m_2}(x,y)\mathbb{f}_r = \frac{m_1}{m_1+m_2}(x,y)\mathbb{f}_r = \frac{m_1}{m_1+m_2}\mathbb{f}_r = \frac{m_1}{m_1+m_2}\mathbb
```

Summary

- Use the COM frame to fix R=0\mathbf{R}=0.
- Solve the coupled ODEs for (x,y)(x,y) with given initial conditions related to orbit eccentricity
- Translate the solution back to positions of the two bodies.
- Plot trajectories for visualization.

If you want, I can help you write MATLAB code to implement this numerical solver with these initial conditions and equations! Would you like that?

Nice overview! Here's a quick summary of what you'll cover in Week 6 of Numerical Methods for Engineers, focusing on PDEs and how to solve them computationally:

Week 6: Introduction to Solving Partial Differential Equations (PDEs)

1. Importance of PDEs

- PDEs describe many real-world phenomena (fluid mechanics, electromagnetism, chemistry, heat flow, etc.).
- Week 6 focuses on **two key PDEs** and their numerical solution.

2. Two-Dimensional Laplace Equation

- This is a boundary value problem (BVP):
 - o The solution inside a domain depends only on boundary values.
- Use the **finite difference method (FDM)** to discretize the equation:
 - o This converts the PDE into a system of algebraic equations.
- Leads to a matrix equation:
 - Natural ordering of points forms a linear system Au=bA \mathbf{u} = \mathbf{b}.
 - Can solve this system directly in MATLAB using the backslash operator (\).
- Alternatively, solve it iteratively via:
 - Jacobi method
 - o Gauss-Seidel method
 - Successive Overrelaxation (SOR) method
- Iterative methods are convenient to code with while loops.

3. One-Dimensional Diffusion Equation

- This is an initial value problem (IVP):
 - You start with initial conditions and evolve the solution over time.
- Numerical methods:
 - Explicit methods (forward in time):
 - Simple but can be unstable (depends on timestep and grid size).
 - o Implicit methods (backward in time):

- More stable, often require solving linear systems at each timestep.
- O Von Neumann stability analysis:
 - Used to analyze the stability of explicit methods.
- The Crank-Nicolson method:
 - o A popular implicit method.
 - o Second-order accurate in time.
 - Combines accuracy and stability.
- Implement solutions in MATLAB.

4. Final Project

- Solve the two-dimensional diffusion equation using:
 - o The Crank-Nicolson method
 - Matrix algebra techniques
- This project combines methods learned from both Laplace and diffusion equation solving.

If you'd like, I can help you understand the finite difference discretization, write MATLAB code for Jacobi/Gauss-Seidel or Crank-Nicolson methods, or explain Von Neumann stability analysis in detail. Just say the word!

Great explanation! Here's a clear summary to help you organize the main ideas from the video about **types of PDE problems** and their numerical solutions:

Understanding PDEs: Boundary Value Problems vs. Initial Value Problems

Why PDEs are harder than ODEs:

- One PDE corresponds to infinitely many ODEs.
- So, solving PDEs numerically requires more care and computational effort.

1. Boundary Value Problem (BVP)

- Example: Laplace equation.
- You have a fixed spatial domain (e.g., a 2D grid in x and y).
- The **boundary values** of the solution (function values at edges of the domain) are known and given.
- The goal is to find the values of the solution at interior points.

- This means solving for the entire spatial field all at once.
- Boundary conditions specify the function value (or derivatives, but in this course, function values) at the edges.
- Numerical methods find the solution by solving a system over all interior points simultaneously.

2. Initial Value Problem (IVP)

- Example: Diffusion equation.
- The domain involves space (e.g., x) and time (t).
- You specify initial conditions: the value of the solution at **t = 0** across space.
- You also have boundary conditions in space (e.g., at smallest and largest x).
- The solution evolves forward in time.
- You compute the solution step-by-step: first at the next time point, then the next, etc.
- Unlike BVPs, you find solutions one time slice at a time rather than all points at once.

Key Differences:

Aspect	Boundary Value Problem	Initial Value Problem
Domain	Spatial only (fixed grid)	Space + Time
Known conditions	Boundary values at domain edges	Initial values at time zero + boundaries in space
Solution approach	Solve for all interior points simultaneously	Step forward in time, solving sequentially
Example	Laplace equation	Diffusion equation

If you'd like, I can help you write MATLAB code for either problem type, explain how to implement boundary or initial conditions, or work through specific examples of discretization methods. Just let me know!

This is a great explanation of how finite difference approximations for derivatives are derived using Taylor series expansions! Here's a concise summary to help organize the key points:

Finite Difference Method: Central Difference Approximations

Setup:

• Consider three points along the x-axis: x-hx - h, xx, and x+hx + h.

- The function values at these points are y(x-h)y(x-h), y(x)y(x), and y(x+h)y(x+h).
- Goal: Approximate derivatives of yy at xx using these function values.

1. Taylor Series Expansions

• Around xx, expand y(x+h)y(x+h) and y(x-h)y(x-h):

```
 y(x+h) = y(x) + hy'(x) + h22y''(x) + h36y(3)(x) + h424y(4)(x) + \cdots + y(x) + hy'(x) + frac\{h^2\}\{2\} y''(x) + frac\{h^3\}\{6\} y^{(3)}(x) + frac\{h^4\}\{24\} y^{(4)}(x) + \cdot y(x-h) = y(x) - hy'(x) + h22y''(x) - h36y(3)(x) + h424y(4)(x) + \cdots + y(x-h) = y(x) - hy'(x) + frac\{h^2\}\{2\} y''(x) - frac\{h^3\}\{6\} y^{(3)}(x) + frac\{h^4\}\{24\} y^{(4)}(x) + \cdot y(x-h) = y(x) - hy'(x) + frac\{h^4\}\{24\} y^{(4)}(x) + \cdot y(x-h) = y(x) - hy'(x) + frac\{h^4\}\{24\} y^{(4)}(x) + \cdot y(x-h) = y(x) - hy'(x) + frac\{h^4\}\{24\} y^{(4)}(x) + \cdot y(x-h) = y(x) - hy'(x) + frac\{h^4\}\{24\} y^{(4)}(x) + y(x) +
```

2. First Derivative Approximation (Central Difference)

• Subtract y(x-h)y(x-h) from y(x+h)y(x+h):

 $y(x+h)-y(x-h)=2hy'(x)+2h36y(3)(x)+\cdots y(x+h)-y(x-h)=2hy'(x)+\frac{1}{2}h^3{6}y^{(3)}(x)+\frac{1}{2}h^3$

Solve for y'(x)y'(x):

 $y'(x)\approx y(x+h)-y(x-h)2hy'(x) \approx \frac{y(x+h)-y(x-h)}{2h}$

• **Error term:** O(h2)O(h^2), meaning the approximation error scales as h2h^2.

3. Second Derivative Approximation (Central Difference)

Add y(x+h)y(x+h) and y(x-h)y(x-h):

```
y(x+h)+y(x-h)=2y(x)+h2y''(x)+h412y(4)(x)+\cdots y(x+h) + y(x-h) = 2 y(x) + h^2 y''(x) + \frac{h^4}{12} y^{(4)}(x) + \cdot cdots
```

• Solve for y"(x)y"(x):

 $y''(x) \approx y(x+h) - 2y(x) + y(x-h)h2y''(x) \approx \frac{y(x+h) - 2 y(x) + y(x-h)}{h^2}$

• **Error term:** O(h2)O(h^2) again, indicating a good accuracy.

Why This Matters

- The **central difference formulas** give second-order accurate approximations to derivatives, better than forward or backward difference formulas which are typically first-order.
- This means you can use coarser grids (larger hh) while keeping reasonable accuracy.
- These approximations are the key step in turning partial differential equations (PDEs) into discrete algebraic equations that can be solved numerically.

If you want, I can help you implement these approximations in MATLAB or Python for some example PDEs! Just say the word.

Here's a clear summary of the key steps in discretizing the 2D Laplace equation on a rectangular grid using finite differences, based on your video transcript:

Solving the 2D Laplace Equation Using Finite Differences

1. Problem Setup:

- Solve for $\Phi(x,y)$ inside a 2D domain $[0,Lx] \times [0,Ly][0,L_x]$ \times $[0,L_y]$.
- Boundary values of Φ\Phi are known (Dirichlet conditions).
- Goal: Find Φ\Phi at the interior grid points.

2. Grid Construction:

• Define grid points along x:

```
xi=(i-1)\Delta x, i=1,2,...,nxx_i = (i-1)\Delta x, \quad i=1,2,...,n_x
```

• Define grid points along y:

 $yj=(j-1)\Delta y, j=1,2,...,nyy_j = (j-1)\Delta y, \quad j=1,2,...,n_y$

- Typically take $\Delta x = \Delta y \setminus Delta x = \setminus Delta y$ for simplicity (square grid).
- Φ \Phi values on the grid are denoted as Φ i,j= Φ (xi,yj)\Phi {i,j} = \Phi(x i,y j).

3. Discretizing the Laplace Equation:

The Laplace equation:

 $\partial 2\Phi dx^2 + \partial 2\Phi dy^2 = 0 \frac{\pi x^2} + \frac{x^2} +$

Use central difference approximations for second derivatives:

4. Combine and Simplify:

Assuming $\Delta x = \Delta y \triangle x = \Delta y$ Delta x = \Delta y, multiply through by $\Delta x \ge \Delta x$

$$\Phi i + 1, j + \Phi i - 1, j + \Phi i, j + 1 + \Phi i, j - 1 - 4\Phi i, j = 0 \\ Phi_{i+1,j} + Phi_{i-1,j} + Phi_{i,j+1} + Phi_{i,j-1} - 4Phi_{i,j} = 0 \\ O$$

Or equivalently:

```
4\Phi_{i,j}=\Phi_{i+1,j}+\Phi_{i-1,j}+\Phi_{i,j}+1+\Phi_{i,j}-14\Phi_{i,j}=\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,j}+\Phi_{i+1,
```

5. Interpretation:

- At each interior grid point, the value Φi,j\Phi_{i,j} is the average of its four nearest neighbors.
- This discrete equation forms a system of linear equations for all interior points.
- The boundary values Φ\Phi are known and fixed.

6. Next Steps:

- Assemble all these equations for interior points into a matrix equation.
- Solve using direct methods (e.g., matrix backslash in MATLAB) or iterative methods (Jacobi, Gauss-Seidel, SOR).
- This will give Ф\Phi values inside the domain consistent with boundary conditions.

If you'd like, I can help you with MATLAB or Python code for setting up and solving this system! Would you want me to do that?

Great! Here's a detailed summary of what you just described about turning the discretized Laplace equation into a matrix system suitable for direct linear algebra solution (like in MATLAB):

From the Discrete Laplace Equation to a Matrix Equation

1. Discrete Laplace Equation Recap:

For interior points, the equation is:

 $4\Phi i, j - \Phi i + 1, j - \Phi i - 1, j - \Phi i, j + 1 - \Phi i, j - 1 = 04 \ Phi_{i,j} - Phi_{i+1,j} - Phi_{i-1,j} - Phi_{i,j+1} - Phi_{i,j-1} = 0$

2. Need for Vectorization:

- Φ\Phi is naturally arranged as a 2D array over indices (i,j)(i,j).
- To solve with matrix methods, we reshape Φ\Phi into a **1D column vector**.
- The matrix equation form:

 $A\Phi=bA \mathbb{Phi} = \mathbb{b}$

where:

- AA is the coefficient matrix,
- o $\Phi\mathbb{P}_{\Phi}$ is the unknown vector of interior $\Phi\mathbb{P}_{\Phi}$
- b\mathbf{b} includes known boundary terms.

3. Natural Ordering of Grid Points:

- The ordering goes **row-wise** starting from the bottom row (j=1j=1), moving left to right across i=1,2,...,nxi = 1, 2, ..., n_x.
- Then move to the next row j=2j=2, again left to right.
- Continue up to the top row j=nyj=n_y.

Example ordering for Φ\Phi values:

 Φ 1,1, Φ 2,1,..., Φ nx,1, Φ 1,2, Φ 2,2,..., Φ nx,2,..., Φ nx,ny\Phi_{1,1}, \Phi_{2,1}, \Idots, \Phi_{n_x,1}, \Phi_{1,2}, \Phi_{2,2}, \Idots, \Phi_{n_x,2}, \Idots, \Phi_{n_x,n_y}

4. Mapping 2D Indices (i,j)(i,j) to 1D Index kk:

 $k=i+(j-1)\times nxk = i + (j-1) \times n_x$

- kk runs from 1 to nx×nyn_x \times n_y.
- So the 2D grid point (i,j)(i,j) corresponds to vector element Φk\Phi_k.

5. Rewrite Discrete Equation Using kk:

Substitute i,j→ki,j \to k:

 $4\Phi k - \Phi k + 1 - \Phi k - 1 - \Phi k - nx = 04 \ Phi_k - Phi_{k+1} - Phi_{k-1} - Phi_{k+n_x} - Phi_{k - n_x} = 0$

- $\Phi k+1\Phi_{k+1}$ and $\Phi k-1\Phi_{k-1}$ correspond to neighbors in the same row (ii direction).
- Φk+nx\Phi_{k+n_x} and Φk-nx\Phi_{k n_x} correspond to neighbors in the rows above and below (jj direction).

6. Matrix AA Structure:

- The matrix AA will be **sparse** with a block structure.
- The diagonal entries are 4.
- The entries –1-1 appear immediately to the left and right of the diagonal (for neighbors in the same row).
- The entries -1-1 appear nxn_x positions away on the diagonals above and below the main diagonal (for neighbors in adjacent rows).
- AA is symmetric and banded.

7. Incorporating Boundary Conditions:

- Boundary values Φ\Phi at edges (where i=1i=1, i=nxi=n_x, j=1j=1, j=nyj=n_y) are known.
- These known values contribute to the right-hand side vector b\mathbf{b}.
- For example, if a neighbor is on the boundary, move its Φ\Phi value to b\mathbf{b}.
- Unknowns correspond only to interior points.

8. Summary of the Process:

- 1. Discretize PDE to difference equations on 2D grid.
- 2. Convert 2D indices (i,j)(i,j) to 1D index kk (natural ordering).
- 3. Write the difference equations in terms of vector Φk\Phi_k.
- 4. Construct matrix AA reflecting the Laplace stencil and ordering.
- 5. Form right side b\mathbf{b} from boundary conditions.
- 6. Solve $A\Phi = bA \ Phi = b$ with a linear solver.

If you want, I can help you **write MATLAB or Python code** that builds matrix AA and vector b\mathbf{b} based on these steps and solves the Laplace equation numerically. Would you like me to do that?

Sure! Here's a clear explanation and step-by-step outline of the approach Jeff Chasnov explained for solving the 2D Laplace equation on a small grid, which you can generalize to larger grids:

Solving the 2D Laplace Equation via Matrix Methods on a Grid

Problem Setup:

- Solve $\nabla 2\Phi = 0 \cdot \Delta^2 \cdot \Phi$ on a 2D domain discretized into a grid.
- Grid size: nx×nyn_x \times n_y points.
- Use **finite difference** approximation for derivatives.
- Use natural ordering to convert Φ(i,j)\Phi(i,j) (2D array) into a vector Φk\Phi_k (1D).

Natural Ordering:

• Index grid points row-wise from bottom to top:

 $\Phi 1 = \Phi 1, 1, \Phi 2 = \Phi 2, 1, ..., \Phi nx, 1, \Phi nx + 1 = \Phi 1, 2, ..., \Phi nx, ny \land Phi_1 = \Phi_{1,1}, \Phi_2 = \Phi_{2,1}, \A Phi_{n_x,1}, \Phi_{n_x+1} = \Phi_{1,2}, \A Phi_{n_x,n_y} = \Phi_{n_x,n_y}$

Total number of points: N=nx×nyN = n_x \times n_y.

Discrete Laplace Equation in Vector Form:

For interior points (ignoring boundaries for now), the equation is:

 $4\Phi k - \Phi k - 1 - \Phi k + 1 - \Phi k - nx - \Phi k + nx = 04 \ Phi_k - Phi_{k-1} - Phi_{k+1} - Phi_{k-n_x} - Phi_{k+n_x} = 0$

- Here k-1,k+1k-1, k+1 correspond to neighbors in the same row.
- k-nx,k+nxk-n_x, k+n_x correspond to neighbors one row below and one row above, respectively.
- Note: Boundary points do **not** use this equation; their values are given by boundary conditions.

Example: Small 3×33 \times 3 Grid

- Number of points: N=9N = 9.
- Points are ordered as:

1,2,34,5,67,8,91, 2, 3 \\ 4, 5, 6 \\ 7, 8, 9

- Interior point: only k=5k=5 (center point).
- Boundary points: k=1,2,3,4,6,7,8,9k=1,2,3,4,6,7,8,9.

Matrix System Setup:

We want:

 $A\Phi = bA \ Phi = b$

- AA: 9×99 \times 9 matrix.
- Φ\Phi: 9×19 \times 1 vector of unknowns.
- bb: 9×19 \times 1 right-hand side vector.

Boundary Conditions Enforcement:

- For boundary points k≠5k \neq 5, Φk=bk\Phi_k = b_k (known values).
- Implement this by making corresponding rows in AA **identity rows**, i.e.:

Akk=1, and zeros elsewhere in row kA_{kk} = 1, \quad \text{and zeros elsewhere in row } k

• The right side bkb_k is set to the known boundary value.

Interior Point Equation (Row 5):

The discrete Laplace equation at k=5k=5:

 $4\Phi 5 - \Phi 4 - \Phi 6 - \Phi 2 - \Phi 8 = 04 \ Phi_5 - Phi_4 - Phi_6 - Phi_2 - Phi_8 = 0$

• Translate into matrix row:

Column index 1 2 3 4 5 6 7 8 9

Matrix row 5 0 -1 0 -1 4 -1 0 -1 0

• Right side b5=0b_5 = 0 since interior Laplace equals zero with no source.

Complete Matrix AA (Rows 1-9):

- Rows for boundary points: identity rows
- Row 5: Laplace stencil as above

Right-hand Side Vector bb:

 $b=[b1b2b3b40b6b7b8b9]b = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \\ 0 \\ b_6 \\ b_7 \\ b_8 \\ b_9 \\ end{bmatrix}$

where bkb k are known boundary values.

Generalization for Large Grids:

- Matrix size: N×NN \times N with N=nx×nyN = n_x \times n_y.
- Matrix AA is sparse with five diagonals:
 - o Main diagonal with 4's.
 - o Two diagonals just above and below with −1-1 (for horizontal neighbors).
 - Two diagonals nxn_x above and below main diagonal with −1-1 (for vertical neighbors).
- Boundary rows replaced with identity rows for boundary conditions.

MATLAB Implementation Note:

- Use spdiags to construct sparse matrix with five diagonals.
- Then replace boundary rows by rows of the identity matrix.
- Solve $A\Phi = bA \ Phi = b$ with sparse solvers.

If you want, I can help write code snippets for MATLAB or Python to build this matrix and solve the system! Would you like that?

Here's a concise summary of the video tutorial you shared on solving the 2D Laplace equation in MATLAB using the direct method:

2D Laplace Equation in MATLAB — Direct Method Overview

Problem setup:

- Solve Laplace's equation on a unit square domain.
- Boundary conditions:
 - \circ Φ = 0 on bottom and side edges.
 - $\Phi = 4x(1 x)$ on the top edge (creates a maximum $\Phi = 1$ at x = 0.5).

Step-by-step outline:

1. Discretization:

- Use central difference approximation for the Laplace equation.
- Define a grid with Nx x Ny intervals (e.g., 100x100).
- o Grid points are (Nx+1) x (Ny+1).

2. Natural ordering:

- o Flatten 2D grid into a vector using column-wise ordering (MATLAB style).
- o Identify boundary indices for bottom, top, left, and right edges.

3. Construct Laplacian matrix A:

- Banded sparse matrix with 5 diagonals:
 - Main diagonal = 4
 - Two diagonals immediately above and below = -1
 - Two diagonals offset by Nx above and below = -1
- Use spdiags for memory-efficient sparse matrix.

4. Apply boundary conditions:

- Replace rows in A corresponding to boundary points with identity matrix rows (speye).
- This enforces fixed Φ values on boundaries.

5. Right-hand side vector b:

- o Initialize zero matrix b (size Nx x Ny).
- Assign boundary values to b accordingly.
- o Reshape b into a vector consistent with natural ordering.

6. Solve the linear system:

- Use MATLAB's backslash operator: phi = A \ b;
- o This step uses Gaussian elimination optimized for sparse matrices.

7. Reshape and plot:

- o Reshape phi vector back to matrix form.
- Use meshgrid to create X and Y matrices for plotting.
- Plot contours with contour(X, Y, phi') (transpose needed due to MATLAB matrix indexing).
- Customize contour levels, axis scaling (axis equal), ticks, and fonts.

Notes:

- MATLAB goes column-wise, so physics coordinates (x horizontal, y vertical) and matrix indices (row first, column second) have a mismatch, handled by transposing φ before plotting.
- The direct solve (backslash operator) is computationally heavy but straightforward.
- Plotting code often takes more time to get visually right than the numerical solution code.
- MATLAB environment features (breakpoints, workspace variables) are helpful for debugging.

If you want, I can help you write the MATLAB code based on this or explain any step in more detail! Would you like that?

Sure! Here's a clear summary and explanation of the lecture on iterative methods for solving the discrete Laplace equation, based on the transcript you provided:

Overview:

The lecture introduces **iterative methods** for solving the discrete Laplace equation, contrasting them with the direct method like Gaussian elimination.

The Problem Setup:

- We want to solve the Laplace equation $\nabla 2\Phi = 0 \cdot \Delta^2 \cdot \Phi = 0$.
- Using finite differences, the discrete Laplace equation at grid point (i,j)(i,j) is:

 $\Phi_{i+1,j+\Phi_{i-1,j+\Phi_{i,j+1}+\Phi_{i,j-1}-4\Phi_{i,j=0}}} + \Phi_{i,j+1,j} + \Phi_{i,j+1,j+1,j} + \Phi_{i,j+1,j+1,j} + \Phi_{i,j+1,j+1,j} + \Phi_{i,j+1,j+1,j+1,j+1} + \Phi_{i,j+1,j+1,j+1} + \Phi_{i,j+1,j+1} +$

Iterative Methods:

Instead of solving the linear system $A\Phi=bA\Phi=b$ directly (e.g., by Gaussian elimination), **iterative methods** update guesses for $\Phi\Phi$ repeatedly until convergence.

Jacobi Method:

Rearrange the discrete equation to solve for Φi,j\Phi_{i,j}:

 $\Phi_{i,j=14}(\Phi_{i+1,j}+\Phi_{i,j+1}+\Phi_{i,j-1}).\$ = \frac{1}{4} \left(\Phi_{i+1,j} + \Phi_{i-1,j} + \Phi_{i,j+1} + \Phi_{i,j-1} \right).

• Define iteration index nn, then update:

- Steps:
 - 1. Initialize $\Phi(0)$ \Phi^{(0)}, e.g., all zeros.
 - 2. Keep boundary values fixed.
 - 3. Update interior Φi,j\Phi_{i,j} using above formula.
 - 4. Repeat until $\Phi(n)$ Phi^{(n)} converges.

Gauss-Seidel Method:

- Similar to Jacobi but updates points in place as you sweep through the grid.
- When updating Φi,j\Phi_{i,j}, use the most recent values:
 - Values from the current iteration for indices less than i,ji,j.
 - Values from the previous iteration for indices greater than i, ji, j.
- This often converges faster because it uses fresher information during the sweep.

Successive Over-Relaxation (SOR) Method:

- Modify the update to accelerate convergence.
- Write the Jacobi update as:

 $\Phi_{i,j}(n+1) = \Phi_{i,j}(n) + 14(\sum_{n=1}^{n} -4\Phi_{i,j}(n)). \\ Ph_{i,j}^{(n+1)} = Ph_{i,j}^{(n)} + \frac{1}{4} \left(\sum_{n=1}^{n} -4 \right) - 4 Ph_{i,j}^{(n)} \right) .$

Introduce relaxation parameter λ\lambda:

 $\Phi_{i,j}(n+1) = \Phi_{i,j}(n) + \lambda \cdot 14(\sum_{n \in \mathbb{N}^{(n)} - 4\Phi_{i,j}(n)}. \Phi_{i,j}^{(n+1)} = \Phi_{i,j}^{(n)} + \lambda \cdot 14(\sum_{n \in \mathbb{N}^{(n)} - 4\Phi_{i,j}(n)}. \Phi_{i,j}^{(n)} - 4\Phi_{i,j}^{(n)} - 4\Phi_{i,j}^{(n)} + \lambda \cdot 14(\sum_{n \in \mathbb{N}^{(n)} - 4\Phi_{i,j}}^{(n)} - 4\Phi_{i,j}^{(n)} - 4\Phi$

- When $\lambda>1$ \lambda > 1: **over-relaxation** (speeds up convergence, but $\lambda<2$ \lambda < 2 for stability).
- When λ <1\lambda < 1: **under-relaxation** (used for stability if iterations are unstable).

Summary:

- **Jacobi**: simple, uses only values from the previous iteration.
- **Gauss-Seidel**: updates values in place, using the latest available.
- **SOR**: modifies the update step with a relaxation factor λ lambda to speed up or stabilize convergence.

If you'd like, I can help you write out the code for these methods or explain any step in more detail! Would you like to see how these iterative methods are implemented in MATLAB or Python?

Got it! Here's a clear summary of the **Red-Black Ordering** iterative method for solving the discrete Laplace equation, based on the video transcript you shared:

Red-Black Ordering Method (also called Red-Black Gauss-Seidel)

Concept:

- Imagine your 2D grid as a checkerboard with alternating **red** and **black** points (here purple is used to represent red, and blue to represent black).
- Each grid point's new value depends on the values of its four neighbors (up, down, left, right).

How It Works:

- Red points update their values using the values of the black points from the previous iteration.
- Black points update their values using the values of the red points from the previous iteration.

Thus, updates happen in two stages per iteration:

- 1. **Update all red points** simultaneously, using the black points' values from the previous iteration.
- 2. **Update all black points** simultaneously, using the newly updated red points.

Then repeat: red \rightarrow black \rightarrow red \rightarrow black ...

Why Use Red-Black Ordering?

- This splitting allows for a form of parallelism and can speed up convergence slightly compared to updating all points at once.
- Historically, this was an important optimization for performance on slower computers.
- Modern computers are fast enough that for simple problems like Laplace's equation, the speedup may not be noticeable.
- However, it remains a useful conceptual and practical method, especially in parallel implementations.

Summary:

Step Update grid points using values from previous step at other color

- 1 Update **Red** points using **Black** points at iteration nn
- 2 Update **Black** points using **Red** points at iteration n+1n+1

This approach usually converges a bit faster than the plain Jacobi or Gauss-Seidel methods updating all points simultaneously.

If you want, I can help you implement the red-black ordering iteration in code or explain its connection to parallel computing. Would you like that?

Great! Here's a clear, step-by-step summary and explanation of the **MATLAB implementation of the Jacobi method** for solving the 2D Laplace equation based on your video transcript:

Problem Setup

Equation:

 $\nabla 2\Phi = 0 \cdot abla^2 \cdot Phi = 0$

(Laplace equation in 2D)

- **Domain:** Rectangular grid, discretized.
- Boundary conditions:
 - \circ Φ =0\Phi = 0 on bottom, left, and right boundaries.
 - \circ $\Phi=4x(1-x)$ Phi = 4x(1-x) on the top boundary (a function of xx).

Jacobi Iteration Scheme

• Discretize the Laplace equation using finite differences.

• Rearrange the equation so that the new value Φi,jn+1\Phi_{i,j}^{n+1} depends on the average of neighboring values at iteration nn:

• The iteration updates **all interior points simultaneously** based on the previous iteration values.

MATLAB Code Algorithm Summary

- 1. **Define the grid and domain parameters** (same as in the direct method).
- 2. Set initial conditions and boundary conditions:
 - Initialize Φ\Phi to zero everywhere.
 - \circ Fix boundary values: zero on bottom, left, and right; 4x(1-x)4x(1-x) on top.
- 3. **Define convergence criteria:**
 - \circ ϵ =10-6\epsilon = 10^{-6} (error tolerance).
- 4. Indexing interior points:
 - o index_x = 2 : nx-1 (exclude boundaries in x)
 - o index_y = 2 : ny-1 (exclude boundaries in y)
- 5. **Iteration loop:**
 - Initialize error to something larger than ∈\epsilon, e.g., 2*epsilon.
 - Initialize iteration counter ncount = 0.
 - While error > epsilon:
 - Increment ncount.
 - Save current Φ\Phi as Φold\Phi_{\text{old}}.
 - Update Φ\Phi at interior points using the Jacobi formula vectorized over the grid indices:
 - Phi(index_x, index_y) = 0.25 * (Phi_old(index_x+1, index_y) + Phi_old(index_x-1, index_y) + Phi_old(index_x, index_y+1) + Phi_old(index_x, index_y-1));
 - Compute error as max absolute difference between Φ\Phi and Φold\Phi {\text{old}} over all points:
 - error = max(abs(Phi(:) Phi_old(:)));
 - Optional: print error at each iteration to monitor convergence.
- 6. Stop when error $< \epsilon \setminus epsilon$.

7. **Print number of iterations and plot the solution** (reuse plotting code from direct method).

Additional Notes

- The code includes a check for divergence (NaN or Inf in Φ \Phi), which helps catch bugs or unstable iteration schemes.
- Printing the error after each iteration is useful during debugging but slows down the code and can be commented out once stable.
- The Jacobi method converged after about **10,311 iterations** for this problem with this tolerance.
- The final Φ\Phi matrix is available for inspection or further analysis.

If you want, I can help you write the full MATLAB code based on this summary or help with other iterative methods like Gauss-Seidel or Red-Black ordering. Would you like that?

Here's a concise summary and explanation of the key points from the video transcript about solving the diffusion equation as an initial value problem using the Forward Time Centered Space (FTCS) scheme:

Topic: Initial Value Problem for PDEs — The Diffusion Equation

1. The Diffusion Equation:

- We consider u(x,t)u(x,t), a function of position xx (one-dimensional) and time tt.
- The diffusion equation is:

where DD is the diffusion constant.

2. Boundary Conditions (BC):

- Values of uu are fixed at the boundaries of the spatial domain for all time.
- Example:

u(-Lx,t)=0, u(Lx,t)=0 for all $tu(-L_x,t)=0$, $quad u(L_x,t)=0$ t=0

3. Initial Condition (IC):

• Specifies u(x,0)u(x,0) at time zero as a known function:

 $u(x,0)=u0(x)u(x,0) = u_0(x)$

4. Discretizing the Domain:

- Space is discretized into grid points xjx_j with spacing Δx\Delta x.
- Time is discretized into time steps tlt 1 with time step size $\Delta t \Delta t$.

Notation:

 $ujl=u(xj,tl)u_j^l=u(x_j,t^l)$

where jj indexes space and ll indexes time.

5. The Forward Time Centered Space (FTCS) Scheme:

• Approximate the time derivative (forward difference, Euler method):

 $\partial u \partial t \approx u | l+1-u | l \Delta t \leq u | l+1 - u | l+1 - u | l+1 \leq u |$

• Approximate the second spatial derivative (centered difference):

 $d2udx2≈uj-1l-2ujl+uj+1l(Δx)2\frac{v}{\partial}^2 u}{\partial}^1-2ujl+u_{j-1}^1-2u_{j-1}^1+u_{j+1}^1}{(\Delta x)^2}$

6. Combining these into the discretized diffusion equation:

 $ujl+1-ujl\Delta t=Duj-1l-2ujl+uj+1l(\Delta x)2\frac{u_j^{l+1}-u_j^{l}}{Delta\ t}=D\frac{u_{j-1}^{l}-2u_j^{l+1}-u_j^{l}}{Delta\ x}^2}$

7. Rearranged as an explicit iteration scheme:

 $u_j|+1=u_j|+D\Delta t(\Delta x)^2(u_j-1|-2u_j|+u_j+1|)u_j^{l+1} = u_j^l + \frac{D \Delta t}{(\Delta x)^2} \left(u_{j-1}^l - 2u_j^l + u_{j+1}^l \right)$

• This allows calculation of the solution at the next time step l+1l+1 explicitly from the known values at time II.

8. Explicit Method:

- "Explicit" means ujl+1u_j^{l+1} depends only on known values at the previous time step II.
- This makes the method easy to implement and understand.

9. Stability Consideration:

- The next video will discuss the Von Neumann stability analysis.
- If the time step $\Delta t \triangle t$ is too large relative to $\Delta x \triangle t$, this explicit method can become unstable (solution can blow up).

Summary:

- You define spatial and temporal grids.
- You apply initial and boundary conditions.
- Use FTCS to advance the solution explicitly in time.
- Careful choice of Δt\Delta t relative to Δx\Delta x and DD is necessary to ensure stability.

If you want, I can help you write MATLAB code implementing this FTCS scheme for the diffusion equation or explain the stability criterion next! Would you like that?

Here's a clear summary and explanation of the video's key points on stability analysis of the Forward Time Centered Space (FTCS) scheme for the 1D diffusion equation:

1. The Setup: FTCS Scheme for Diffusion Equation

- The diffusion equation describes how something like heat or a chemical concentration diffuses over time.
- The Forward Time Centered Space (FTCS) scheme is an explicit numerical method.
- Explicit means: the solution at the next time step depends only on known values at the current time step.
- The update formula for the solution uu at grid point jj and time step l+1l+1 depends on values at time step ll.

2. Stability Issue with Explicit Schemes

- Explicit schemes can become **unstable** if the time step $\Delta t \triangle t$ is too large.
- Instability means the numerical solution "blows up," grows without bound, which is non-physical for diffusion (where solutions should smooth out and decay over time).
- We want the scheme stable so that numerical solutions behave physically (diffusion causes concentration to decrease smoothly).

3. Von Neumann Stability Analysis

- To analyze stability, we assume solutions are made of **Fourier modes**: waves with different wave numbers kk.
- The trial solution (ansatz) is:

ujl= ζ leikj Δ xu_j^l = ζ eta^l e^{i k j Delta x}

where ζ \zeta is a growth factor per time step, i=-1i = \sqrt{-1}, kk is wave number, and Δx \Delta x is spatial grid size.

• If $|\zeta|>1$ \zeta | > 1 for **any** kk, the scheme is unstable because that mode grows exponentially with time step II.

4. Deriving the Growth Factor ζ\zeta

- Substitute the ansatz into the FTCS scheme.
- After algebraic simplification, using Euler's formula:

 $eik\Delta x + e - ik\Delta x = 2cos(k\Delta x)e^{i k \Delta x} + e^{-i k \Delta x} = 2 cos(k\Delta x)$

• The growth factor comes out as:

```
\zeta=1+2r(\cos(k\Delta x)-1)\cdot = 1 + 2r(\cos(k\Delta x)-1) where r=D\Delta t(\Delta x)^2 = \frac{D \det t}{(\Delta x)^2} and DD is the diffusion coefficient.
```

5. Interpreting Stability Condition

- Since $cos (k\Delta x) cos(k \Delta x) ranges between -1-1 and +1+1:$
 - O Maximum ζ\zeta is at $cos (k\Delta x)=1 cos(k \Delta x)=1$, yielding ζ=1\zeta = 1 (neutral stability).
 - Minimum $\zeta = 1 \cos(k \Delta x) = 1 \cos(k \Delta x) = 1$:

```
\zeta \min_{0 \le 1+2r(-1-1)=1-4r} = 1 + 2r(-1-1) = 1 - 4r
```

For stability, require:

```
-1 \le \zeta \min_{0 \le 1-1 \le 1} \le 1
```

SO

$$-1 \le 1-4r \implies 4r \le 2 \implies r \le 12-1 \le 1-4r \le 4r \le 2 \le r \le 12-1 \le 1-4r \le 4r \le 2 \le r \le 12-1 \le 1-4r \le 12-1 \le 12-1$$

• Therefore, the stability condition is:

 $D\Delta t(\Delta x)2\leq 12 \Rightarrow \Delta t\leq (\Delta x)22D \text{ $$\Gamma t}(\Delta x)^2 \leq 12 \Rightarrow \Delta t\leq (\Delta x)^2$

• If $\Delta t \setminus Delta t$ is larger than this, the scheme becomes unstable.

6. Summary

- The FTCS explicit scheme for diffusion is **conditionally stable**.
- The time step must satisfy:

 $\Delta t \leq (\Delta x) = 1 + (\Delta x)^2$

- Larger time steps cause the scheme to blow up.
- In the next part (not in this video), implicit schemes will be shown to be unconditionally stable, meaning you can pick any time step size without instability.

If you want, I can also help write out the detailed algebra or explain how to implement this scheme or do a similar stability analysis for other PDEs. Would you like that?

Great summary! Let me break down the key points and explain the stability difference between the explicit Forward Time Centered Space (FTCS) scheme and the implicit scheme for the diffusion equation:

Explicit FTCS Scheme Recap:

Scheme:

 $ujl+1=ujl+\Delta tD\Delta x2(uj+1l-2ujl+uj-1l)u_j^{l+1} = u_j^{l} + \frac{Delta\ t\ D}{Delta\ x^2} \left(u_{j+1}^{l} - 2 u_j^{l} + u_{j-1}^{l} \right)$

- Here, the spatial second derivative is computed at the **current time step II**.
- Stability:

The von Neumann stability analysis shows that for stability,

 $\Delta t \leq \Delta x \geq 2D$ Delta t \leq \frac{\Delta x^2}{2D}

Otherwise, the scheme becomes unstable and solutions can blow up.

 This means the time step size is severely limited by the spatial grid size and diffusion coefficient.

Implicit Scheme:

Scheme:

- Now, the second derivative is calculated at the future time step l+1l+1.
- Because the right-hand side contains unknowns at l+1l+1, you must solve a linear system (matrix equation) for all grid points at each time step.
- Stability:

Von Neumann analysis yields:

 $\cos\theta=11+2r(1-\cos(k\Delta x))\cos \theta=11+2r(1-\cos(k\Delta x))$

where

 $r=\Delta t D \Delta x 2r = \frac{\Delta t D}{\Delta x^2}$

- Since the denominator is always ≥1\geq 1, the magnitude of cos
 ⊕\cos \theta is always ≤1\leq 1.
- **Result:** The implicit method is **unconditionally stable** no matter how large Δt\Delta t is, the scheme will not blow up.

Why is this important?

 Explicit methods are simple but require very small time steps to maintain stability, especially for small spatial grid sizes.

- **Implicit methods** are more complex computationally (solve a system of equations each step) but allow larger time steps without instability.
- This makes implicit methods preferable for stiff problems or when long time simulations are needed.

Next Step: Crank-Nicolson Method

- Combines explicit and implicit ideas.
- It is second-order accurate in time and unconditionally stable.
- We'll discuss that next!

If you'd like, I can help you write out the matrix form for the implicit scheme or guide you through the Crank-Nicolson method next. How do you want to proceed?

Great detailed explanation! Let me summarize and clarify the Crank-Nicolson method and its matrix formulation for solving the 1D diffusion equation:

Problem Setup:

We want to solve the diffusion equation

 $\label{lem:dudt=Dd2udx2\frac{\hat u}{\hat x^2} = D \frac{1}{partial x^2}$

on a spatial grid indexed by jj and time steps indexed by II.

Recap of Methods:

• Explicit FTCS:

Uses spatial second derivative at time II only, conditionally stable.

Fully Implicit:

Uses spatial second derivative at time l+1l+1, unconditionally stable but requires solving a matrix system each step.

Crank-Nicolson Method:

- **Idea:** Average the spatial second derivative between time levels II and I+1I+1. This makes the method:
 - Second-order accurate in both space and time.
 - Unconditionally stable like the fully implicit method.
- Finite difference equation:

Rearranged form:

Define

 $\alpha = \Delta t D \Delta x 2$ alpha = $\frac{D}{Delta x^2}$

Multiply both sides by 2 to clear denominators and rearrange terms to isolate ujl+1u_j^{l+1}:

 $-\alpha uj+1l+1+2(1+\alpha)ujl+1-\alpha uj-1l+1=\alpha uj+1l+2(1-\alpha)ujl+\alpha uj-1l-\alpha\ u_{j+1}^{l+1}+2(1+\alpha\ u_{j+1}^{l+1}-\alpha\ u_{j-1}^{l+1}=\alpha uj+1l+2(1-\alpha)ujl+\alpha uj-1l-\alpha\ u_{j+1}^{l+1}+2(1+\alpha)ujl+1l-\alpha uj-1l+1=\alpha uj+1l+2(1-\alpha)ujl+\alpha uj-1l-\alpha\ u_{j+1}^{l+1}+2(1+\alpha)ujl+1-\alpha uj-1l+1=\alpha uj+1l+2(1-\alpha)ujl+\alpha uj-1l-\alpha\ u_{j+1}^{l+1}+2(1+\alpha)ujl+1-\alpha uj-1l+1=\alpha uj+1l+2(1-\alpha)ujl+\alpha uj-1l-\alpha\ u_{j+1}^{l+1}+2(1+\alpha)ujl+1-\alpha uj-1l+1=\alpha uj+1l+2(1-\alpha)ujl+\alpha uj-1l-\alpha\ u_{j+1}^{l+1}+2(1+\alpha)ujl+1-\alpha uj-1l-\alpha\ u_{j+1}^{l+1}+2(1+\alpha)uj-1l-\alpha\ u_{j+1}^$

Matrix Form:

For a grid with 4 points (including boundaries), assume boundary conditions $u1|=u4|=0u_1^l = u_4^l = 0$ (Dirichlet zero BCs):

- The first and last rows implement the boundary conditions.
- The interior rows form a **tridiagonal system** that you solve each timestep.

Important notes:

- The matrix is **constant in time**, so you can factorize it once (LU decomposition or Thomas algorithm) for efficient solving.
- At each time step, the right-hand side vector updates based on the previous solution ulu^l.
- The tridiagonal structure means solving the system is very fast (order O(N)O(N)).

Summary:

- Crank-Nicolson method is an average of explicit and implicit schemes.
- It is unconditionally stable.
- It is **second-order accurate** in both space and time.
- Leads to a **tridiagonal matrix system** to solve each timestep.
- Boundary conditions are directly incorporated into the matrix system.

If you'd like, I can help you write MATLAB or Python code to implement this, or explain the Thomas algorithm for efficiently solving the tridiagonal system. Would you like to see that next?

Great! Here's a concise summary of the MATLAB implementation of the Crank-Nicolson method for the 1D diffusion equation based on the video and transcript you shared:

MATLAB Code Outline for Crank-Nicolson Diffusion Solver

1. Define domain and parameters:

- Domain length Lx=1L_x = 1 (so $x \in [-1,1]x \in [-1,1]$)
- Number of intervals $Nx=500N_x = 500$, grid points $nx=Nx+1=501n_x = N_x + 1 = 501$
- Grid spacing Δx=2LxNx\Delta x = \frac{2 L_x}{N_x}
- Diffusion constant D=1D = 1 (unit diffusion)
- Time step $\Delta t = \Delta x = \Delta x = \Delta t = \Delta t$
- Number of time steps (e.g., 10,000)
- Output interval (e.g., plot every 500 steps)

2. Create spatial grid:

x = linspace(-Lx, Lx, nx); % row vector of grid points

3. Define $\alpha = D\Delta t \Delta x 2 = \frac{D \Delta t \Delta x^2}{\rho x^2}$

4. Construct the time-independent tridiagonal matrix AA:

- Main diagonal: $2(1+\alpha)2(1+\alpha)$
- Off-diagonals (upper and lower): $-\alpha$ -\alpha
- Use MATLAB's spdiags for efficient sparse matrix creation
- Impose Dirichlet boundary conditions by replacing first and last rows with identity rows

```
main_diag = 2*(1 + alpha) * ones(nx,1);

off_diag = -alpha * ones(nx-1,1);

A = spdiags([off_diag main_diag off_diag], [-1 0 1], nx, nx);
```

% Boundary conditions (u=0 at boundaries)

$$A(1,:) = 0; A(1,1) = 1;$$

 $A(end,:) = 0; A(end,end) = 1;$

5. Initialize uu as a Gaussian centered in the domain:

```
sigma = Lx / 16;
```

```
u = (1/(sigma*sqrt(2*pi))) * exp(-0.5 * (x/sigma).^2);
u = u'; % convert to column vector
6. Plot initial condition:
plot(x, u, 'r');
xlabel('x'); ylabel('u');
title('Initial Condition');
drawnow;
```

7. Time-stepping loop:

For each time step:

- Construct the right-hand side vector BB using values of uu at previous step, according to Crank-Nicolson scheme
- Impose boundary conditions on BB (zero at boundaries)
- Solve the linear system Aul+1=BA u^{l+1} = B using MATLAB's backslash operator
- Plot solution every nout steps

```
for m = 1:ntime

B = zeros(nx,1);

% Boundary conditions for B

B(1) = 0; B(end) = 0;

% Fill B for interior points (2 to nx-1) using previous u and alpha

B(2:end-1) = alpha*u(1:end-2) + 2*(1 - alpha)*u(2:end-1) + alpha*u(3:end);

% Solve linear system

u = A \ B;

% Plot every nout steps

if mod(m, nout) == 0

hold on;
plot(x, u);
drawnow;
end
```

Key Points:

- The matrix AA is sparse, tridiagonal, and time-independent → efficient setup.
- Boundary conditions (Dirichlet zero) are enforced by fixing first and last rows of AA and BB.
- The Crank-Nicolson method is implicit but stable and second-order accurate.
- MATLAB's \ operator makes solving the system simple and fast.
- The solution uu starts as a sharp Gaussian and diffuses over time, spreading out and flattening.

If you'd like, I can help you write the **full MATLAB script** ready to run, or help with any specific part of the code or explanation! Would you like that?

Great! This is an excellent summary and guidance on tackling the **2D diffusion equation using the Crank-Nicolson method** in MATLAB. Let me break down the key steps and the main ideas from your project description, so you have a clear roadmap for coding:

Overview: 2D Diffusion Equation

- u=u(x,y,t)u=u(x,y,t)
- Domain: square [-L,L]×[-L,L][-L, L] \times [-L, L]
- Boundary conditions: u=0u=0 on all boundaries
- Initial condition: large concentration at the center, decaying to zero at boundaries

Step-by-Step Solution Outline

1. Define the grid

- Create a uniform grid for xx and yy with spacing $h=\Delta x=\Delta yh=\Delta yh=\Delta x=\Delta yh=\Delta x=\Delta yh$
- Total number of points Nx=Ny=NN_x = N_y = N (fixed for grading).
- Use **natural ordering** (e.g., row-major or column-major) to flatten the 2D grid ui,ju_{i,j} into a 1D vector uku_k where k=1,2,...,N2k = 1, 2, \ldots, N^2.

2. Discretize the PDE with Crank-Nicolson

- Time discretization: implicit, average of the spatial discretization at time nn and n+1n+1.
- Spatial discretization: standard 5-point stencil for Laplacian

3. Construct the matrix AA

• The Crank-Nicolson scheme can be expressed as

 $Aun+1=Bun+bA u^{n+1} = B u^n + b$

where AA and BB are large sparse matrices of size N2×N2N^2 \times N^2.

- AA includes diagonal elements and off-diagonals that connect points in the grid according to the 5-point stencil.
- Apply boundary conditions by setting rows corresponding to boundary nodes to enforce u=0u=0.

4. LU Decomposition

- To accelerate solving Aun+1=RHSA u^{n+1} = \text{RHS} at every time step, compute LU factorization once outside the time loop.
- Inside the loop, solve two triangular systems instead of inverting or directly solving AA.

5. Initial condition

- Initialize uu as a large peak at the center of the domain (e.g., Gaussian).
- Flatten it into a vector matching the ordering scheme.

6. Time stepping loop

- For each time step:
 - Compute the right-hand side vector b=Bunb = B u^n plus boundary conditions.
 - Solve for un+1u^{n+1} using the LU factors.
 - Optionally save or plot un+1u^{n+1}.

7. Graphics

- Plot uu as a 3D surface or contour plot.
- Animate or update plot every few steps to show diffusion.

Summary of main MATLAB implementation ideas:

- Use sparse matrices for AA and BB.
- Use natural ordering: e.g., flatten 2D indices (i,j)(i,j) into 1D index k=i+(j-1)*Nk = i + (j-1) * N.
- Enforce boundary conditions by replacing rows in AA with identity rows.
- Use lu factorization in MATLAB once outside time loop.
- Use \ operator with LU factors for fast solve inside time loop.

• Plot results using surf or mesh for 3D visualization.

If you'd like, I can help you by writing a **basic MATLAB template** for this 2D diffusion problem based on the above outline. Would you like me to do that?

This is a great course summary highlighting the main topics you covered in your Numerical Methods for Engineers class, including all six projects! Here's a quick recap and the big picture for each week/project:

Week 1: Basics & Roundoff Errors

- Learned MATLAB programming fundamentals.
- Explored roundoff errors.
- **Project:** Compute the bifurcation diagram for the logistic map.

Week 2: Root Finding Methods

- Learned Newton's method and fzero.m for root finding.
- **Project:** Use Newton's method to compute Feigenbaum Delta.

Week 3: Numerical Linear Algebra

- Focused on Gaussian elimination with partial pivoting (MATLAB's \ operator).
- Applied to solving nonlinear systems and PDEs.
- **Project:** Compute a fractal from the Lorenz equations.

Week 4: Quadrature & Interpolation

- Learned numerical integration with integral.m.
- Learned interpolation methods such as cubic splines with interp1.m.
- **Project:** Compute zeros of Bessel functions.

Week 5: Ordinary Differential Equations

- Learned the Runge-Kutta method, especially MATLAB's ode45.
- Learned shooting methods for boundary value problems.
- **Project:** Compute orbits of two gravitationally interacting masses using Newton's laws.

Week 6: Finite Difference Methods for PDEs

- Learned finite difference schemes for PDEs including diffusion equations.
- Solved 1D and 2D diffusion equations.
- **Project:** Solve the two-dimensional diffusion equation using Crank-Nicolson.

Final Thoughts

- These numerical methods and MATLAB skills are essential for engineering and scientific computing.
- They're highly valued by professors and research groups who need capable programmers for solving complex numerical problems.

If you want, I can help you review or practice any particular topic or code from these weeks. Which part would you like to dive into?