23k-2001 BCS-3J Date____20_

COAL: Assignment 01

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	Answer#1:
	a. High level languages Assembly language
	· Closer resemblance to human · Closer resemblance to
	natural language nabive machine language. • A code-line is called as • A code-line is called as
	'statement'. 'instruction'.
	· Multiply operations can be · Only one operation can be
	achieved in one statement. performed by an instruction.
	· Single statement corresponds · Each instruction corresponds
	to multiply instructions (one-to- to a single machine-language
	many) instruction. (one-to-one)
4.5	· High-level languages have data · Assembly language has no data
	types such as int, char, float. type, rather data sizes such as
	byte, word, dword.
	and often system independent. for a specific architecture, hence
	they are not portable.
	Relationship and similarities:
经	-> Most high-level languages are compiled to a lower-level
	(assembly language) before being translated to machine code. -> Assembly language is used in parallel with high-level languages in applications to optimize or target specific hardware.
	-> Assembly language is used in parallel with high-level languages
	in applications to optimize or target specific hardware.
C-1	-> All programming concepts such as conditions, selection,
幕	iteration & recursion are facilitated by both type of languages.
	-> High-level languages are useful for faster development by
	abstracting hardware details, whereas assembly language optimizes
	performance by direct control over hardware.

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b. Assemblers:	Compilers:
· Assembler translates assembly	· Compilers translate high-
language directly into machine	level languages into a lower
code, ready for execution.	level languages into a lower to level (often assembly) which is
	passed on to be converted to
	machine language.
· Translates instructions in	· Abstract details at the high
correspondence to hardware-detail	1 . A . A . A . A . V//
as assembly language is closely	task that require further
bied to the cpu's architecture.	interpretation to operate on hardware.
· Error - handling is limited, and code has to be proof read	· Compilers are highly
before - hand.	and handling e.g. syntax errors
before many.	corto ricorigirig.
c. Portability of Alls as	compared to Assembly:
High-level languages are regard	led as more portable as they
are designed to be hardw	are - independent. The compilers
for Hlls abstract away the	hardware details and make
code ready for execution on a	inu machina unharens Assembly
	A HILL THE CONTROL
language is directly assembled	I into machine code that is
specific to the architecture	it was written for.
specific to the architecture → E.g: A high-level language	program such as Java
specific to the architecture > E.g: A high-level language can be executed on a	program such as Java ny system with JVM
specific to the architecture > E.g: A high-level language can be executed on a irrespective of the hardwar	program such as Java ny system with JVM
specific to the architecture > E.g: A high-level language can be executed on a irrespective of the hardwar > E.g: A program written in	program such as Java ny system with JVM e. C/C++ does not need to be
specific to the architecture > E.g: A high-level language can be executed on a irrespective of the hardwar > E.g: A program written in rewritten for each platform	program such as Java ny system with JVM c. C/C++ does not need to be m. It can be compiled and run
specific to the architecture > E.g: A high-level language can be executed on a irrespective of the hardwar > E.g: A program written in	program such as Java ny system with JVM c. C/C++ does not need to be m. It can be compiled and run
specific to the architecture > E.g: A high-level language can be executed on a irrespective of the hardwar > E.g: A program written in rewritten for each platfor by any C++ compiler avail	program such as Java ny system with JVM c. C/C++ does not need to be m. It can be compiled and run able.
specific to the architecture > E.g: A high-level language can be executed on a irrespective of the hardwar > E.g: A program written in rewritten for each platfor by any C++ compiler avail	program such as Java ny system with JVM c. C/C++ does not need to be m. It can be compiled and run able.
specific to the architecture > E.g: A high-level language can be executed on a irrespective of the hardwar > E.g: A program written in rewritten for each platform	program such as Java ny system with JVM c. C/C++ does not need to be m. It can be compiled and run able. Len for a specific hordware a architecture without

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	Δ
	Answer#2:
	Virtual Machine: A software program that can emulate the functionality of another physical or virtual computer is known as a Virtual Machine. According to the given example:
per.	A software program that can emulate the functionality
	of another physical or virtual computer is known as
	a Virtual Machine.
-	According to the given example:
	level-3: High-level (anguage) compilation VM#1
	According to the given example: level-3: High-level language compilation Level-2: Assembly language translation translation
_	translation
	level - 1: Instruction-set computer architecture VM#0
1	level-0: Digital logic D'execute instructions
	· Ill is compiled to Assembly large at 1/M#1
1	· Then translated to TSA based on VM#O susten and
10	 Hll is compiled to Assembly languag at VM#1. Then, translated to TSA based on VM#O system and executed at level-0 which is the hardware level
	on VM#O.
	A 4 3
	Answer#3:
	a. MOV EAX, F2F0F0F1h
	ADD EXX, 10000100h
	Sol:
	F2F0F0F1 1111 0010 1111 0000 1111 0000 1111 0001
	+10000100 +0001 0000 0000 0000 0001 0000 0000
	0001 0000 0010 1111 0000 1111 0001
	⇒ O2FOF1F1h
	OF: 0
	ZF: 0 CF: 0
	31.
	AF. O
	PAPER PRODUCTS

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b. Keal Address Modes:	
8086 processors operating in real-address made	
can address 20-bit addresses following the range	Marry
can address 20-bit addresses following the range O to FFFFF (hex), which is why it can only access 1MB of RAM at a time.	
access dag 1 RAM at a bine	-
access I will of will an a ame.	
c. Physical Address: 1 2AB: 025F	******
C. Physical Huciress	
1 2AB: 025F	
	-
Unear physical address = segment x 10h + offset	
	-
⇒ 12 AB	4
x 10 in the line of the second	
12 ABOh Physical address = 12DOFh	
+ 025F	_
12 DOF Ans.	
12001	
d. 8086 is a 16-bit architecture that utilizes	-
16-bit registers for efficient usage. This is due	
16-bit registers for efficient orage. Mis is and	-
to the fact that 16-bit data transfers were	
sufficient for most general-purpose tasks and 20-bit address bus allow it to access IMB at a time	
address bus allow it to access INB at a time	
providing a balance between memory readability and	
transfer efficiency. The CPU is capable to store	
transfer efficiency. The CPU is capable to store the 20-bit linear address based on the formula	
above.	



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	Λ 11 / 1
	Answer#4
	TITLE Question 4
	INCLUDE Invine32.inc
	· data
	sunday = 0
	monday = 1
	hesday = 2
-	wednesday = 3
	thursday = 4
	friday = 5 salurday = 6
	salurday = 6
C	week BYTE sunday, monday, tuesday, wednesday, thursday, friday, saturday
	· code
A	elder main PROC
	call DumRegs
	exit
5	main ENDP
	END main 1 + 5
	Answer#5
44	. data
-	varA DWORD 5 DUP (?)
	yarB BYTE 2 DUP(?)
	varC BYTE 15 BUP("&") varD BYTE 7 DUP("%")
	0/5-11)
	varE BYTE 1 DUP('M')
- Channel	



Answer#6:

111300	
i) mov al, 88h	
1) mov al, 88h add al, 90h	
add oi, Ton	
Sol:	
88h 1,000	1000
+ 90h => + 1001	0000
1 18h 0001 0001	1000
→ 18h	
Ans: al = 18h	
OF = 1	(result of the arithmetic
SF = 0	too large to fit in
CF = 1	destination)
11) mov d,5	
add al, 123	
Sol:	
5 '0000 '0101	
$\pm 123 \Rightarrow \pm 0111 \ 1011$	
128 1000 0000	
<u>⇒</u> 80 k	
Ans. al = 80	(MSB=1 implies a sign
OF = 1	change which also incurs
SF = 1	an overflow condition as 128
CF = 0	is now being interpreted as a
-1	re number while the actual
a	(MSB=1 implies a sign change which also incurs an overflow condition as 128 is now being interpreted as a ve number while the actual number should have been +128)

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	A17
	Answer#7:
E	A) EAX = 20001000 h dwlist is of type DWORD, hence first 4 bytes will be stored
F-	B) EBX = 00200010h offset +1(byk), moves forward storing och from 3000h of the previous 3 bytes in order
	C) ECX = 30002000 h offset +2 (bytes) stores 4 bytes ahead of the offset (1000h)
E	D) EDX = 11 300020h Afset+3(bytes), moves forward storing 11h from lower half of 111h and the previous 3 bytes in order.
	V
1	Memory Allocation for data segment:
	100020003000111122223333
	(Little Endian representation)
	000 0002 0003 111 2222 3333
1	

