

A Collection of Quick Tours

Hassium

1 Derived Categories

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In this note, we will sketch some proofs but not in detailed.

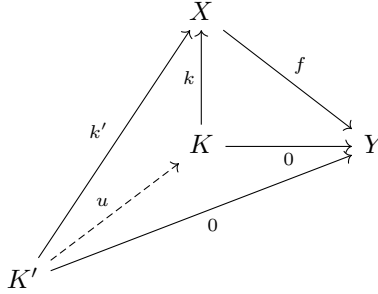
Definition 1.1. Let \mathcal{C} be a category. A category \mathcal{D} is a *subcategory* of \mathcal{C} if the $\text{Ob}(\mathcal{D})$ and the $\text{Hom}_{\mathcal{D}}(X, Y)$ are subcollections of $\text{Ob}(\mathcal{C})$ and $\text{Hom}_{\mathcal{C}}(X, Y)$, respectively, for all objects X and Y in \mathcal{C} . The subcategory \mathcal{D} is said to be *full* if for all objects X and Y in \mathcal{D} , $\text{Hom}_{\mathcal{D}}(X, Y)$ is exactly $\text{Hom}_{\mathcal{C}}(X, Y)$.

Definition 1.2. Let \mathcal{C} be a category. A morphism $f : A \rightarrow B$ is a *monomorphism* if for all $g, h : C \rightarrow A$, $f \circ g = f \circ h$ implies $g = h$. A morphism $f : A \rightarrow B$ is called an *epimorphism* if for all $i, j : B \rightarrow D$, $i \circ f = j \circ f$ implies $i = j$.

Definition 1.3. Let \mathcal{C} be a category. A morphism $f : X \rightarrow Y$ in \mathcal{C} is a *constant morphism* if for all object Z and morphisms $g, h : Z \rightarrow X$, $fg = fh$. A morphism is a *zero morphism* if it is both a constant morphism and a coconstant morphism.

Definition 1.4. Let \mathcal{C} be a category. A *kernel* of a morphism $f : X \rightarrow Y$ is a pair (K, k) , where K is an object and $k : K \rightarrow X$ is a morphism, such that

1. $fk = 0_{KY}$.
2. For all (K', k') such that $k'f = 0_{K'Y}$, there exists a unique $u : K' \rightarrow K$.



Definition 1.5. Let F and G be functors between \mathcal{C} and \mathcal{D} . A *natural transformation* η from F to G is a family of morphisms such that

1. For all object X in \mathcal{C} , there exists a morphism $\eta_X : F(X) \rightarrow G(X)$, called the *component* at X .
2. For all $f : X \rightarrow Y$, $\eta_Y \circ F(f) = G(f) \circ \eta_X$.

A functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is a *natural functor* if for all $G : \mathcal{C} \rightarrow \mathcal{D}$, there exists a natural transformation η .

Definition 1.6. A category \mathcal{C} is an *additive category* if every hom-set of \mathcal{C} is an abelian group, composition of morphisms is bilinear, and \mathcal{C} admits finite coproduct.

Definition 1.7. A functor $f : \mathcal{A} \rightarrow \mathcal{B}$ between additive categories is an *additive functor* if it preserves the finite coproduct.

Definition 1.8. A category \mathcal{C} is a *pre-abelian category* if \mathcal{C} is an additive category and every morphism in \mathcal{C} has a kernel and a cokernel. A pre-abelian category \mathcal{C} is an *abelian category* if every monomorphism is a kernel and every epimorphism is a cokernel.

Definition 1.9. Let \mathcal{A} be an additive category. A *cochain complex* is a collection X^* of objects X^n in \mathcal{A} with maps $d^n : X^n \rightarrow X^{n+1}$ such that $d^{n+1}d^n = 0$. The morphism between complexes X^* and Y^* is a collection of maps $f^n : X^n \rightarrow Y^n$ such that $f^{n+1}d_X^n = d_Y^n f^n$. This defines the category of cochain complexes, denoted $\mathcal{C}(\mathcal{A})$.

Definition 1.10. Let X^* and Y^* be cochain complexes. The maps $f, g : X^* \rightarrow Y^*$ are *homotopic* if there exists a collection of maps $k^n : X^n \rightarrow Y^{n-1}$ such that $f^n - g^n = d_Y^{n-1}k^n + k^{n+1}d_X^n$.

It is trivial that homotopy is an equivalence relation.

Definition 1.11. Let \mathbf{A} be an additive category. The objects of the *homotopy category*, denoted $K(\mathbf{A})$, are cochain complexes and the morphisms are homotopic class of morphisms of cochain complexes.

Let \mathbf{A} be an additive category and $K(\mathbf{A})$ be the homotopy category. Define the finite coproduct to be the direct sum $(X^* \oplus Y^*)^n = X^n \oplus Y^n$. Define the composition to be pairwise addition. Since composition in \mathbf{A} is bilinear, it is bilinear in $K(\mathbf{A})$. It is not hard to check that $K(\mathbf{A})$ is additive.

Definition 1.12. A *triangulated category* is an additive category \mathbf{C} with:

1. an additive automorphism $T : \mathbf{C} \rightarrow \mathbf{C}$ called the *translation functor*,
2. a collection of *triangles* (X, Y, Z, u, v, w) , where X, Y , and Z are objects of \mathbf{C} and $u : X \rightarrow Y$, $v : Y \rightarrow Z$, and $w : Z \rightarrow T(X)$ are morphisms,
3. morphisms $(f, g, h) : (X, Y, Z, u, v, w) \rightarrow (X', Y', Z', u', v', w')$ such that the following diagram commutes.

$$\begin{array}{ccccccc}
 X & \xrightarrow{u} & Y & \xrightarrow{v} & Z & \xrightarrow{w} & T(X) \\
 \downarrow f & & \downarrow g & & \downarrow h & & \downarrow T(f) \\
 X' & \xrightarrow{u'} & Y' & \xrightarrow{v'} & Z' & \xrightarrow{w'} & T(X')
 \end{array}$$

The data subject to the following rules:

1. The sextuple $(X, X, 0, \text{id}_X, 0, 0)$ is a triangle and for all $f : A \rightarrow B$, there exists a triangle (A, B, C, f, g, h) .
2. A sextuple (A, B, C, f, g, h) is a triangle if and only if $(B, C, T(A), g, h, -T(f))$ is a triangle.
3. Let (X, Y, Z, u, v, w) and (X', Y', Z', u', v', w') be triangles. Let $f : X \rightarrow X'$ and $g : Y \rightarrow Y'$ be morphisms in \mathbf{C} that commutes with u and u' , then there exists $h : Z \rightarrow Z'$ such that (f, g, h) is a morphism between triangles.

$$\begin{array}{ccccccc}
 X & \xrightarrow{u} & Y & \longrightarrow & Z & \longrightarrow & T(X) \\
 \downarrow f & & \downarrow g & & \downarrow h & & \downarrow T(f) \\
 X' & \xrightarrow{u'} & Y' & \longrightarrow & Z' & \longrightarrow & T(X')
 \end{array}$$

4. Let $(X, Y, Z', u, j,)$, $(Y, Z, X', v, , i)$, and $(X, Z, Y', vu, ,)$ be triangles, then there exists morphisms $f : Z' \rightarrow Y'$ and $g : Y' \rightarrow X'$ such that $(Z', Y', X', f, g, T(j)i)$ is a triangle and the following diagram commutes.

$$\begin{array}{ccc}
 Z' & \xrightarrow{\quad f \quad} & Y' \\
 \downarrow & \swarrow & \downarrow \\
 X & & X'
 \end{array}
 \qquad
 \begin{array}{ccc}
 Y' & \xleftarrow{\quad g \quad} & Z \\
 \downarrow & \swarrow & \downarrow \\
 X' & & X
 \end{array}$$

This is called the *octohedral axiom*.

$$\begin{array}{ccccc}
 & & Y' & & \\
 & \swarrow f & & \searrow g & \\
 Z' & & & & X' \\
 \downarrow & & \downarrow T(j)i & & \downarrow \\
 X & & & & Z \\
 \swarrow u & & \swarrow v & & \swarrow \\
 & & Y & &
 \end{array}$$

Definition 1.13. Let \mathcal{C} and \mathcal{D} be additive categories. An additive functor $f : \mathcal{C} \rightarrow \mathcal{D}$ is a *covariant ∂ -functor* if it commutes with the translation functor and it preserves triangles. A *contravariant ∂ -functor* takes triangles into triangles with the arrows reversed, and sends the translation functor into its inverse.

Homotopy categories can be triangulated.

Proposition. Let $T(X^*)^p = X^{p+1}$ and $d_{T(X)} = -d_X$. Given a morphism $u : X^* \rightarrow Y^*$, we define Z^* to be the *mapping cone* $T(X^*) \oplus Y^*$. The differential maps in Z^* are matrices $\begin{pmatrix} T(d_X) & T(u) \\ 0 & d_Y \end{pmatrix}$. We pick $v : Y^* \rightarrow Z^*$ and $w : Z^* \rightarrow X^*$ to be the natural inclusion and projection. A triangle is defined to be a sextuple that induces by a morphism $u : X^* \rightarrow Y^*$. This defines a triangulated category.

Proof. (TR1) Take $\text{id}_X : X^* \rightarrow X^*$, then $Z^* = T(X^*) \oplus X^*$ and the differentials are $d_Z^n = \begin{pmatrix} -d_X^n & \text{id}_{X^{n+1}} \\ 0 & d_X^n \end{pmatrix}$. Define $k^n : Z^n \rightarrow Z^{n-1}$ by $k^n((a, b)) = (0, a)$. Then $k^{n+1}d_Z^n((a, b)) = k^{n+1}((-d_X^n(a) + b, d_X^n(b))) = (0, -d_X^n(a) + b)$ and $d_Z^{n-1}k^n((a, b)) = d_Z^{n-1}((0, a)) = (a, d_X^n(a))$. We have $(k^{n+1}d_Z^n + d_Z^{n-1}k^n)((a, b)) = (a, b)$, hence $\text{id}_Z \sim 0$. The sextuple $(X, X, 0, \text{id}_X, 0, 0)$ is the triangle induced by $\text{id}_X : X^* \rightarrow X^*$. \square (TR2) Let (X^*, Y^*, Z^*, u, v, w) be a triangle. Denote $T(Y^*) \oplus Z^*$ by A^* . Consider the sextuple (Y^*, Z^*, A^*, v, s, t) , \square

Definition 1.14. Let \mathcal{C} be a category and let S be a collection of morphisms in \mathcal{C} , then the *localization* of \mathcal{C} with respect to S is a category \mathcal{C}_S together with a functor $Q : \mathcal{C} \rightarrow \mathcal{C}_S$ such that

1. $Q(s)$ is an isomorphism for every $s \in S$.
2. Any functor $F : \mathcal{C} \rightarrow \mathcal{D}$ such that $F(s)$ is an isomorphism for all $s \in S$ factors uniquely through Q .

If $Q' : \mathcal{C} \rightarrow \mathcal{C}_S$ is another functor satisfies these properties, then the two localizations are considered to be equivalent.

Proposition. The localization \mathcal{C}_S exists.

Proof. Let $\text{Ob}(\mathcal{C}_S) = \text{Ob}(\mathcal{C})$ and let Q be the identity on objects. For all $s \in S$, introduce some variable x_s , then construct an oriented graph Γ as follows: let the vertices be objects in \mathcal{C} , let the edges be morphisms in \mathcal{C} union $\{x_s\}$, the edge has the direction $X \rightarrow Y$ if there exists a morphism $X \rightarrow Y$ in \mathcal{C} , and the edge x_s has the same vertices as the edges but the opposite orientation. Define a morphism to be an equivalence class of paths that share the same initial and terminal in Γ . Two paths are equivalent if they can be joined by a chain of elementary equivalences that two consecutive arrows in a path can be replaced by their composition and arrows $X \xrightarrow{s} Y \xrightarrow{s_x} X$ can be replaced by $X \xrightarrow{\text{id}_X} X$. The composition of morphisms is induced by the conjunction of paths. The functor Q takes a morphism to a corresponding path of length 1. The inverse of a morphism $s \in S$ is x_s . Suppose we have another functor $\mathcal{C} \rightarrow \mathcal{C}'$ satisfies the universal property. We construct $G : \mathcal{C}_S \rightarrow \mathcal{C}'$ by $G(X) = F(X)$, $G(f) = F(f)$, and $G((x_s)) = F(s)^{-1}$ for all $X \in \text{Ob}(\mathcal{C})$, $f \in \text{Mor}_{\mathcal{A}}$, and $s \in S$. This defines a localization but we will not check the axioms here. \square

Definition 1.15. Let \mathcal{A} be an abelian category. A *quasi-isomorphism* is a morphism $f : X^* \rightarrow Y^*$ in $\mathcal{K}(\mathcal{A})$ which induces an isomorphism on cohomology. The collection of all quasi-isomorphisms is denoted by Qis .

Definition 1.16. The *derived category* $\mathcal{D}(\mathcal{A})$ of an abelian category \mathcal{A} is the localization $\mathcal{K}(\mathcal{A})_{\text{Qis}}$.

Let $s \in S$ and let $f : X' \rightarrow Y$ be a morphism in \mathcal{A} . A *roof* is a diagram (s, f) of the form

$$\begin{array}{ccc} & X' & \\ s \swarrow & & \searrow f \\ X & & Y \end{array}$$

The morphisms in \mathcal{C}_S are equivalence classes of roofs. One could check this is exactly the same as we described in the previous proposition.

Proposition. Derived categories are additive.

Proof. Let $\varphi, \varphi' : X \rightarrow Y$ be morphisms in $D(A)$.

$$\begin{array}{ccc} & U & \\ t \swarrow & & \searrow g \\ X & & Y \end{array} \qquad \begin{array}{ccc} & U' & \\ t' \swarrow & & \searrow g' \\ X & & Y \end{array}$$

Since t and t' are quasi-isomorphisms, we can construct V by homotopy pullback such that the following diagram commutes up to homotopy. We define the addition $\varphi + \varphi'$ to be the equivalence class of roofs of the form below.

$$\begin{array}{ccc} V & \xrightarrow{q'} & U' \\ q \downarrow & & \downarrow t' \\ U & \xrightarrow{t} & X \end{array} \qquad \begin{array}{ccc} & V & \\ t \circ q \swarrow & & \searrow g \circ q + g' \circ q' \\ X & & Y \end{array}$$

It is left to check that addition does define an additive structure, but we will skip the details. \square

Definition 1.17. Let \mathcal{C} be a triangulated category. A collection S of morphisms in \mathcal{C} is said to be *compatible with triangulation* if

1. $s \in S$ if and only if $T(s) \in S$.
2. consider the diagram (TR3), if $f, g \in S$, then the complementing morphism $h \in S$.

Derived categories can be triangulated. Let \mathcal{C} be a triangulated category and consider some localization \mathcal{C}_S that is compatible with triangulation. Define T to be the localizing functor and define $\text{Ob}(\mathcal{C}_S) = \text{Ob}(\mathcal{C})$. A triangle in \mathcal{C}_S is a image of a triangle in \mathcal{C} under the localization $\mathcal{C} \rightarrow \mathcal{C}_S$. This defines a triangulated category. See [Gel03] IV.2 for a detailed proof. It is not hard to check that Qis is compatible with $K(A)$, hence $D(A)$ is triangulated.

Definition 1.18. Let A be an abelian category and let $K^*(A)$ be a triangulated subcategory of $K(A)$. If the natural functor $K^*(A)_{K^*(A) \cap \text{Qis}} \rightarrow D(A)$ is fully-faithful, then $D^*(A)$ is called a *localizing subcategory* and we denote it by $D^*(A)$.

We show the definition of a right derived covariant functor and the other definitions can be obtained similarly.

Definition 1.19. Let A and B be abelian categories. Let $F : K^*(A) \rightarrow K(B)$ be a ∂ -functor. The *right derived covariant functor* of F is a ∂ -functor $R^*F : D^*(A) \rightarrow D(B)$ together with a natural transformation $\eta : Q \circ F \rightarrow R^*F \circ Q$ from $K^*(A) \rightarrow D(B)$ such that if $G : D^*(A) \rightarrow D(B)$ is a ∂ -functor and $\xi : Q \circ F \rightarrow G \circ Q$ is a natural transformation, then there exists unique natural transformation $\eta' : R^*F \rightarrow G$ such that $\xi = (\eta' \circ Q) \circ \eta$.

References

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