

# A Collection of Quick Tours

Hassium

- 1 Burnside Group of Odd Exponents: Results by Arkarskaya, A. et al.
- 2 Derived Categories

## 1 Burnside Group of Odd Exponents: Results by Arkarskaya, A. et al.

This note is entirely based on [1] and I will only rephrase the proofs to make them more readable. This note is written for a presentation I will give in Math 490.

**Definition 1.1.** Let  $F = \langle x_1, \dots, x_m \rangle$  be the free group  $F_m$ , where  $m \geq 2$ . The *free Burnside group*  $B(m, n) = F / \langle\langle x_1, \dots, x_m \mid w^n, w \in F \rangle\rangle$ . We say  $B(m, n)$  has *rank*  $m$  and *exponent*  $n$ .

**Definition 1.2.** Let  $F_m$  be a free group of rank  $m$ , then elements of  $\{x_1, \dots, x_m\} \cup \{x_1^{-1}, \dots, x_m^{-1}\}$  are called *words*. A sequence of words is called a *letter*. A word without cancellation is called a *reduced word*.

**Definition 1.3.** We say a word  $A$  is *cyclically contained* in a word  $w$  if  $A$  is a subword of any cyclic shift of  $w$ .

**Example.** Let  $w = abcde$  and let  $A = dea$ . Shift  $w$  cyclically, we obtain  $bcdea$ , hence  $A$  is cyclically contained in  $w$ .

**Definition 1.4.** Let  $w$  be a reduced word. A *prefix* of  $w$  is any initial segment of  $w$ . A *suffix* of  $w$  is any final segment of  $w$ .

**Definition 1.5.** Let  $w$  be a reduced word that is not empty. We say  $w$  is *primitive* if there does not exist  $k \geq 2$  such that  $w = a^k$ .

**Definition 1.6.** Let

**Theorem 1.1.** Induct on  $i$ , then the following properties hold.

1.  $\text{Can}_i \subset \text{Can}_{i-1}$ ;
2. if  $L_1 A^\tau R_1, L_2 A^\tau R_2 \in \text{Can}_i$  for some primitive  $A$  and  $A^n \notin$

If the induction works, we will have the following results immediately.

**Theorem 1.2.**

## 2 Derived Categories

This note is based on my presentation given in Math 501.

**Definition 2.1.** Let  $\mathcal{C}$  be a category. A category  $\mathcal{D}$  is a *subcategory* of  $\mathcal{C}$  if the  $\text{Ob}(\mathcal{D})$  and the  $\text{Hom}_{\mathcal{D}}(X, Y)$  are subcollections of  $\text{Ob}(\mathcal{C})$  and  $\text{Hom}_{\mathcal{C}}(X, Y)$ , respectively, for all objects  $X$  and  $Y$  in  $\mathcal{C}$ . The subcategory  $\mathcal{D}$  is said to be *full* if for all objects  $X$  and  $Y$  in  $\mathcal{D}$ ,  $\text{Hom}_{\mathcal{D}}(X, Y)$  is exactly  $\text{Hom}_{\mathcal{C}}(X, Y)$ .

**Definition 2.2.** Let  $\mathcal{C}$  be a category. A morphism  $f : A \rightarrow B$  is a *monomorphism* if for all  $g, h : C \rightarrow A$ ,  $f \circ g = f \circ h$  implies  $g = h$ . A morphism  $f : A \rightarrow B$  is called an *epimorphism* if for all  $i, j : B \rightarrow D$ ,  $i \circ f = j \circ f$  implies  $i = j$ .

**Definition 2.3.** An *initial object*  $0$  of a category  $\mathcal{C}$  is an object in  $\mathcal{C}$  such that for any object  $A$ , there is a unique morphism  $0 \rightarrow A$ . A *terminal object*  $1$  of a category  $\mathcal{C}$  is an object of  $\mathcal{C}$  such that for any object  $B$ , there is a unique morphism  $A \rightarrow 1$ . We say an object is the *zero object* if it is both an initial object and a terminal object.

**Definition 2.4.** Let  $\mathcal{C}$  be a category with zero object  $0$ . The *zero morphism*  $0_{A,B}$  between objects  $A$  and  $B$  is the unique morphism that factors through  $0$ .

**Definition 2.5.** Let  $\mathcal{C}$  be an object and let  $f : X \rightarrow Y$  be a morphism in  $\mathcal{C}$ . An object  $\ker(f)$  is said to be the *kernel* of  $f$  if for every object  $Z$  and  $h : Z \rightarrow X$  such that  $f \circ h = 0$ , where  $0$  is the zero morphism, there is a unique morphism  $\varphi : Z \rightarrow \ker(f)$  such that  $h = \varphi \circ f$ .

**Definition 2.6.** Let  $F$  and  $G$  be functors between  $\mathcal{C}$  and  $\mathcal{D}$ . A *natural transformation*  $\eta$  from  $F$  to  $G$  is a family of morphisms such that

1. For all object  $X$  in  $\mathcal{C}$ , there exists a morphism  $\eta_X : F(X) \rightarrow G(X)$ , called the *component* at  $X$ .
2. For all  $f : X \rightarrow Y$ ,  $\eta_Y \circ F(f) = G(f) \circ \eta_X$ .

A functor  $F : \mathcal{C} \rightarrow \mathcal{D}$  is a *natural functor* if for all  $G : \mathcal{C} \rightarrow \mathcal{D}$ , there exists a natural transformation  $\eta : F \rightarrow G$ .

**Definition 2.7.** A category  $\mathcal{C}$  is an *additive category* if the following properties hold.

1.  $\mathcal{C}$  has an zero object  $0$ .
2. Every hom-set of  $\mathcal{C}$  is an abelian group.
3. Composition of morphisms is bilinear.
4.  $\mathcal{C}$  admits finite coproduct.

**Definition 2.8.** A functor  $f : \mathcal{A} \rightarrow \mathcal{B}$  between additive categories is an *additive functor* if it preserves the finite coproduct.

**Definition 2.9.** A category  $\mathcal{C}$  is a *pre-abelian category* if  $\mathcal{C}$  is an additive category and every morphism in  $\mathcal{C}$  has a kernel and a cokernel. A pre-abelian category  $\mathcal{C}$  is an *abelian category* if every monomorphism is a kernel and every epimorphism is a cokernel.

**Definition 2.10.** Let  $\mathcal{A}$  be an additive category. A *cochain complex* is a collection  $X^*$  of objects  $X^n$  in  $\mathcal{A}$  with maps  $d^n : X^n \rightarrow X^{n+1}$  such that  $d^{n+1}d^n = 0$ . The morphism between complexes  $X^*$  and  $Y^*$  is a collection of maps  $f^n : X^n \rightarrow Y^n$  such that  $f^{n+1}d_X^n = d_Y^n f^n$ . This defines the category of cochain complexes, denoted  $\mathcal{C}(\mathcal{A})$ .

**Definition 2.11.** Let  $X^*$  and  $Y^*$  be cochain complexes. The maps  $f, g : X^* \rightarrow Y^*$  are *homotopic* if there exists a collection of maps  $k^n : X^n \rightarrow Y^{n-1}$  such that  $f^n - g^n = d_Y^{n-1}k^n + k^{n+1}d_X^n$ .

It is trivial that homotopy is an equivalence relation.

**Definition 2.12.** Let  $\mathcal{A}$  be an additive category. The objects of the *homotopy category*, denoted  $\mathcal{K}(\mathcal{A})$ , are cochain complexes and the morphisms are homotopic class of morphisms of cochain complexes.

Let  $\mathbf{A}$  be an additive category and  $\mathbf{K}(\mathbf{A})$  be the homotopy category. Define the finite coproduct to be the direct sum  $(X^* \oplus Y^*)^n = X^n \oplus Y^n$ . Define the composition to be pairwise addition. Since composition in  $\mathbf{A}$  is bilinear, it is bilinear in  $\mathbf{K}(\mathbf{A})$ . It is not hard to check that  $\mathbf{K}(\mathbf{A})$  is additive.

**Definition 2.13.** A *triangulated category* is an additive category  $\mathbf{C}$  with:

1. an additive automorphism  $T : \mathbf{C} \rightarrow \mathbf{C}$  called the *translation functor*,
2. a collection of *triangles*  $(X, Y, Z, u, v, w)$ , where  $X, Y$ , and  $Z$  are objects of  $\mathbf{C}$  and  $u : X \rightarrow Y$ ,  $v : Y \rightarrow Z$ , and  $w : Z \rightarrow T(X)$  are morphisms,
3. morphisms  $(f, g, h) : (X, Y, Z, u, v, w) \rightarrow (X', Y', Z', u', v', w')$  such that the following diagram commutes.

$$\begin{array}{ccccccc}
 X & \xrightarrow{u} & Y & \xrightarrow{v} & Z & \xrightarrow{w} & T(X) \\
 \downarrow f & & \downarrow g & & \downarrow h & & \downarrow T(f) \\
 X' & \xrightarrow{u'} & Y' & \xrightarrow{v'} & Z' & \xrightarrow{w'} & T(X')
 \end{array}$$

The data subject to the following rules:

1. The sextuple  $(X, X, 0, \text{id}_X, 0, 0)$  is a triangle and for all  $f : A \rightarrow B$ , there exists a triangle  $(A, B, C, f, g, h)$ .
2. A sextuple  $(A, B, C, f, g, h)$  is a triangle if and only if  $(B, C, T(A), g, h, -T(f))$  is a triangle.
3. Let  $(X, Y, Z, u, v, w)$  and  $(X', Y', Z', u', v', w')$  be triangles. Let  $f : X \rightarrow X'$  and  $g : Y \rightarrow Y'$  be morphisms in  $\mathbf{C}$  that commutes with  $u$  and  $u'$ , then there exists  $h : Z \rightarrow Z'$  such that  $(f, g, h)$  is a morphism between triangles.

$$\begin{array}{ccccccc}
 X & \xrightarrow{u} & Y & \longrightarrow & Z & \longrightarrow & T(X) \\
 \downarrow f & & \downarrow g & & \downarrow h & & \downarrow T(f) \\
 X' & \xrightarrow{u'} & Y' & \longrightarrow & Z' & \longrightarrow & T(X')
 \end{array}$$

4. Let  $(X, Y, Z', u, j, -)$ ,  $(Y, Z, X', v, -, i)$ , and  $(X, Z, Y', vu, -, -)$  be triangles, then there exists morphisms  $f : Z' \rightarrow Y'$  and  $g : Y' \rightarrow X'$  such that  $(Z', Y', X', f, g, T(j)i)$  is a triangle and the following diagram commutes.

$$\begin{array}{ccc}
 Z' & \xrightarrow{\quad f \quad} & Y' \\
 \downarrow & \swarrow & \downarrow \\
 X & & X'
 \end{array}
 \qquad
 \begin{array}{ccc}
 Y' & \xleftarrow{\quad g \quad} & Z \\
 \downarrow & \swarrow & \downarrow \\
 X' & & X
 \end{array}$$

This is called the *octohedral axiom*.

$$\begin{array}{ccccc}
 & & Y' & & \\
 & \swarrow f & & \searrow g & \\
 Z' & & & & X' \\
 \downarrow & & \xrightarrow{T(j)i} & & \downarrow \\
 X & & & & Z \\
 \downarrow u & & \downarrow v & & \\
 & & Y & & 
 \end{array}$$

**Definition 2.14.** Let  $\mathbf{C}$  and  $\mathbf{D}$  be additive categories. An additive functor  $f : \mathbf{C} \rightarrow \mathbf{D}$  is a *covariant  $\partial$ -functor* if it commutes with the translation functor and it preserves triangles. A *contravariant  $\partial$ -functor* takes triangles into triangles with the arrows reversed, and sends the translation functor into its inverse.

Homotopy categories can be triangulated.

**Proposition.** Let  $T(X)^p = X^{p+1}$  and  $d_{T(X)} = -d_X$ . Given a morphism  $u : X \rightarrow Y$ , we define  $Z$  to be the *mapping cone*  $T(X) \oplus Y$ . The differential maps in  $Z$  are matrices  $\begin{pmatrix} T(d_X) & T(u) \\ 0 & d_Y \end{pmatrix}$ . We pick  $v : Y \rightarrow Z$  and  $w : Z \rightarrow X$  to be the natural inclusion and projection. A triangle is defined to be a sextuple that induces by a morphism  $u : X \rightarrow Y$ . This defines a triangulated category.

*Proof.* (TR1) Take  $\text{id}_X : X \rightarrow X$ , then  $Z = T(X) \oplus X$  and the differentials are  $d_Z^n = \begin{pmatrix} -d_X^n & \text{id}_{X^{n+1}} \\ 0 & d_X^n \end{pmatrix}$ . Define  $k^n : Z^n \rightarrow Z^{n-1}$  by  $k^n((a, b)) = (0, a)$ . We have  $(k^{n+1}d_Z^n + d_Z^{n-1}k^n)((a, b)) = (a, b)$ , hence  $\text{id}_Z \sim 0$ . The sextuple  $(X, X, 0, \text{id}_X, 0, 0)$  is the triangle induced by  $\text{id}_X : X \rightarrow X$ .  $\square$  (TR2) Let  $(X, Y, Z, u, v, w)$  be a triangle. Denote  $T(Y) \oplus Z$  by  $A$ . Consider the sextuple  $(Y, Z, A, v, s, t)$ . Notice that  $(T(X))^i = X^{i+1}$ , so we pick  $x^{i+1} \in (T(X))^i$ . Define  $\theta^i : (T(X))^i \rightarrow A^i$  by  $x^{i+1} \mapsto (-u^{i+1}(x^{i+1}), x^{i+1}, 0)$ .

$$\begin{array}{ccccccc}
 Y & \xrightarrow{v} & T(X) \oplus Y & \xrightarrow{w} & T(X) & \xrightarrow{T(v)} & T(Y) \\
 \text{id}_Y \downarrow & & \text{id}_{T(X) \oplus Y} \downarrow & & \theta \downarrow & & \text{id}_{T(Y)} \downarrow \\
 Y & \xrightarrow{v} & T(X) \oplus Y & \longrightarrow & T(Y) \oplus T(X) \oplus Y & \longrightarrow & T(Y)
 \end{array}$$

Define  $h^i : Z^i \rightarrow A^{i-1}$  by  $h^i(x^{i+1}, y^i) = (y^i, 0, 0)$ . Then two maps from  $Z$  to  $A$  are homotopic. Let  $\psi : A \rightarrow T(X)$  be the natural projection, then  $\psi \circ \theta = \text{id}_{T(X)}$ . We have  $\theta \circ \psi^i : A^i \rightarrow A^i$  with  $\theta \circ \psi^i : (y^{i+1}, x^{i+1}, y^i) \mapsto (-f^{i+1}(x^{i+1}), x^{i+1}, 0)$ . Pick  $k^i : A^i \rightarrow A^{i-1}$  to be the map  $(y^{i+1}, x^{i+1}, y^i) \mapsto (y^i, 0, 0)$ , then  $\theta \circ \psi \sim \text{id}_A$ .  $\square$  (TR3) We have  $Z = T(X) \oplus Y$  and  $Z' = T(X') \oplus Y'$ , so the complementing morphism is exactly  $h = T(f) \oplus g : Z \rightarrow Z'$ .  $\square$  (TR4) We say an exact sequence of complexes  $0 \rightarrow X \xrightarrow{u} Y \xrightarrow{v} Z \rightarrow 0$  is *semi-split* if for all  $i$ , there exists a morphism  $w^i : Z^i \rightarrow Y^i$  such that  $v^i w^i = \text{id}_{Z^i}$ . Consider a lemma: any semi-split exact sequence in  $\mathbf{C}(\mathbf{A})$  can induce a triangle  $(X, Y, Z, u, v, -)$  in  $\mathbf{K}(\mathbf{A})$ . (lemma)

Now given the upper cap, we want to construct the lower cap by this lemma.

$\square$

**Definition 2.15.** Let  $\mathbf{C}$  be a category and let  $S$  be a collection of morphisms in  $\mathbf{C}$ , then the *localization* of  $\mathbf{C}$  with respect to  $S$  is a category  $\mathbf{C}_S$  together with a functor  $Q : \mathbf{C} \rightarrow \mathbf{C}_S$  such that

1.  $Q(s)$  is an isomorphism for every  $s \in S$ .
2. Any functor  $F : \mathbf{C} \rightarrow \mathbf{D}$  such that  $F(s)$  is an isomorphism for all  $s \in S$  factors uniquely through  $Q$ .

If  $Q' : \mathbf{C} \rightarrow \mathbf{C}_S$  is another functor satisfies these properties, then the two localizations are considered to be equivalent.

**Proposition.** The localization  $\mathbf{C}_S$  exists.

*Proof.* Let  $\text{Ob}(\mathbf{C}_S) = \text{Ob}(\mathbf{C})$  and let  $Q$  be the identity on objects. For all  $s \in S$ , introduce some variable  $x_s$ , then construct an oriented graph  $\Gamma$  as follows: let the vertices be objects in  $\mathbf{C}$ , let the edges be morphisms in  $\mathbf{C}$  union  $\{x_s\}$ , the edge has the direction  $X \rightarrow Y$  if there exists a morphism  $X \rightarrow Y$  in  $\mathbf{C}$ , and the edge  $x_s$  has the same vertices as the edges but the opposite orientation. Define a morphism to be an equivalence class of paths that share the same initial and terminal in  $\Gamma$ . Two paths are equivalent if they can be joined by a chain of elementary equivalences that two consecutive arrows in a path can be replaced by their composition and arrows  $X \xrightarrow{s} Y \xrightarrow{s_x} X$  can be replaced by  $X \xrightarrow{\text{id}_X} X$ . The composition of morphisms is induced by the conjunction of paths. The functor  $Q$  takes a morphism to a corresponding path of length 1. The inverse of a morphism  $s \in S$  is  $x_s$ .  $\square$

**Definition 2.16.** Let  $\mathbf{A}$  be an abelian category. A *quasi-isomorphism* is a morphism  $f : X^* \rightarrow Y^*$  in  $\mathbf{K}(\mathbf{A})$  which induces an isomorphism on cohomology. The collection of all quasi-isomorphisms is denoted by  $\text{Qis}$ .

**Definition 2.17.** The *derived category*  $D(A)$  of an abelian category  $A$  is the localization  $K(A)_{Qis}$ .

Let  $s \in Qis$  and let  $f : X' \rightarrow Y$  be a morphism in  $A$ . A *roof* is a diagram  $(s, f)$  of the form

$$\begin{array}{ccc} & X' & \\ s \swarrow & & \searrow f \\ X & & Y \end{array}$$

The morphisms in  $D(A)$  are equivalence classes of roofs. One could check this is exactly the same as we described in the previous proposition. For a proof, check [GM03] III.2.

**Proposition.** Derived categories are additive.

*Proof.* Let  $\varphi, \varphi' : X \rightarrow Y$  be morphisms in  $D(A)$ .

$$\begin{array}{ccc} & U & \\ t \swarrow & & \searrow g \\ X & & Y \end{array} \quad \begin{array}{ccc} & U' & \\ t' \swarrow & & \searrow g' \\ X & & Y \end{array}$$

We want the roofs to have common top and left side. Since  $t$  and  $t'$  are quasi-isomorphisms, we can construct  $V$  by homotopy pullback such that the following diagram commutes up to homotopy. We define the addition  $\varphi + \varphi'$  to be the equivalence class of roofs of the form below.

$$\begin{array}{ccc} V & \xrightarrow{q'} & U' \\ q \downarrow & & \downarrow t' \\ U & \xrightarrow{t} & X \end{array} \quad \begin{array}{ccc} & V & \\ t \circ q \swarrow & & \searrow g \circ q + g' \circ q' \\ X & & Y \end{array}$$

This defines an additive category. □

**Definition 2.18.** Let  $C$  be a triangulated category. A collection  $S$  of morphisms in  $C$  is said to be *compatible with triangulation* if it satisfies the following property.

1.  $s \in S$  if and only if  $T(s) \in S$ .
2. consider the diagram (TR3), if  $f, g \in S$ , then the complementing morphism  $h \in S$ .

Derived categories can be triangulated.

**Proposition.** Let  $C$  be a triangulated category and consider some localization  $C_S$  that is compatible with triangulation. Define  $T$  to be the localizing functor and define  $Ob(C_S) = Ob(C)$ . A triangle in  $C_S$  is a image of a triangle in  $C$  under the localization  $C \rightarrow C_S$ . This defines a triangulated category.

It is not hard to check that  $Qis$  is compatible with  $K(A)$ , hence  $D(A)$  is triangulated.

**Definition 2.19.** Let  $A$  be an abelian category and let  $K^*(A)$  be a triangulated subcategory of  $K(A)$ . If the natural functor  $K^*(A)_{K^*(A) \cap Qis} \rightarrow D(A)$  is fully-faithful, then  $D^*(A)$  is called a *localizing subcategory* and we denote it by  $D^*(A)$ .

We show the definition of a right derived covariant functor and the other definitions can be obtained similarly.

**Definition 2.20.** Let  $A$  and  $B$  be abelian categories. Let  $F : K^*(A) \rightarrow K(B)$  be a  $\partial$ -functor. The *right derived covariant functor* of  $F$  is a  $\partial$ -functor  $R^*F : D^*(A) \rightarrow D(B)$  together with a natural transformation  $\eta : Q \circ F \rightarrow R^*F \circ Q$  from  $K^*(A) \rightarrow D(B)$  such that if  $G : D^*(A) \rightarrow D(B)$  is a  $\partial$ -functor and  $\xi : Q \circ F \rightarrow G \circ Q$  is a natural transformation, then there exists unique natural transformation  $\eta' : R^*F \rightarrow G$  such that  $\xi = (\eta' \circ Q) \circ \eta$ .

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## References

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