

Integer programming approaches for solving variants of Sudoku puzzle

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 - Some variants of Sudoku
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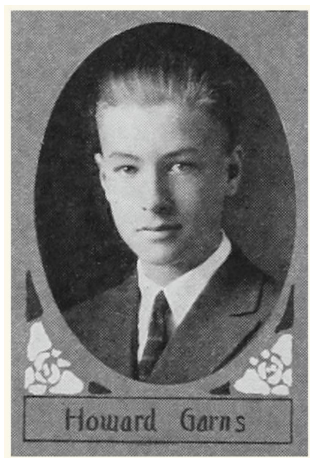
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Introduction



Howard Garns (1905 - 1989): father of modern Sudoku

Classical Sudoku rule for size 9×9

- (S1) Each cell of grid must be filled by exactly one digit from 1 to 9.
- (S2) Each digit from 1 to 9 appears exactly once in each row of the grid.
- (S3) Each digit from 1 to 9 appears exactly once in each column of the grid.
- (S4) Each digit from 1 to 9 appears exactly once in each of the 3×3 blocks of the grid.

			5				6	
8		9					1	
1	6			8	7			
3				2	6			
		7		1		6		
			8	5				3
			4	7			2	1
	4					9		8
	8				3			

4	7	2	5	3	1	8	6	9
8	5	9	6	4	2	3	1	7
1	6	3	9	8	7	2	5	4
3	1	8	7	2	6	4	9	5
5	9	7	3	1	4	6	8	2
6	2	4	8	5	9	1	7	3
9	3	6	4	7	8	5	2	1
7	4	1	2	6	5	9	3	8
2	8	5	1	9	3	7	4	6

Variants of Sudoku

- X-Sudoku
- Windoku
- Killer Sudoku
- Jigsaw Sudoku

X-Sudoku

(X1) Each digit from 1 to 9 appears exactly once in each of the two main diagonals of the grid.

5						3	4
	7				5		
4				8			1
			4		6		9 2
6	4	2		3			1
9			1	2			
			6				
						6	
						3	

5	8	9	7	6	1	2	3	4
1	7	3	2	4	5	9	8	6
4	2	6	9	8	3	7	5	1
3	1	8	4	7	6	5	9	2
6	4	2	5	3	9	8	1	7
9	5	7	1	2	8	4	6	3
7	3	5	6	9	2	1	4	8
8	9	4	3	1	7	6	2	5
2	6	1	8	5	4	3	7	9

Windoku

(W1) Each digit from 1 to 9 appears exactly once in each of the four additional blocks of the grid.

9			4		3			8
	1			9				
3		4				6		2
				7			4	
			5		8			
	2			3				
2		5				3		6
				4			8	
1			3		6			7

9	6	2	4	5	3	1	7	8
8	1	7	6	9	2	5	3	4
3	5	4	8	1	7	6	9	2
6	9	3	2	7	1	8	4	5
4	7	1	5	6	8	9	2	3
5	2	8	9	3	4	7	6	1
2	4	5	7	8	9	3	1	6
7	3	6	1	4	5	2	8	9
1	8	9	3	2	6	4	5	7

Jigsaw Sudoku

(J1) Each digit from 1 to 9 appears exactly once in each of the nine nonominos.

	5				1			
	8					9		
3					6			7
		5						
	1	2				3		
				1		5		
2								5
					9		4	6
			1				9	

9	5	7	4	2	1	6	8	3
6	8	3	2	4	7	9	5	1
3	9	1	5	8	6	4	2	7
8	6	5	7	3	9	2	1	4
5	1	2	9	6	4	3	7	8
7	3	4	8	1	2	5	6	9
2	4	9	6	7	8	1	3	5
1	2	8	3	9	5	7	4	6
4	7	6	1	5	3	8	9	2

Killer Sudoku

- (K1) No digit appears more than once in a cage.
- (K2) The sum of all digits filled in a cage must equal the associated number of the cage.

7	19			25		12	25	
7		9						10
17	12		9					
	10			10	20		12	
			12			17		
14	15				12		12	
	7		11			19		10
18			3				11	
12			28					

7	19			25		12	25	
7		9						10
17	12		9					
	10			10	20		12	
			12			17		
14	15				12		12	
	7		11			19		10
18			3				11	
12			28					

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Framework

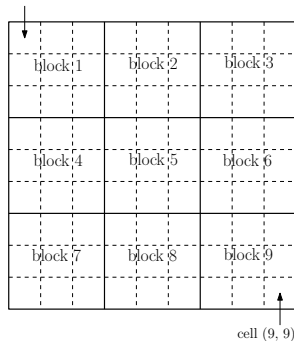
Formulate Sudoku problem as an IP

$$\begin{array}{ll}\min & c^T x \\ \text{s.t.} & Ax = b \\ & x \in \mathbb{Z}^n\end{array}$$

then solve numerically by some IP solver

Classical Sudoku

cell (1, 1)



$$S = \{1, \dots, 9\}$$

$$B := \{(1, 1), (1, 4), (1, 7), (4, 1), (4, 4), (4, 7), (7, 1), (7, 4), (7, 7)\}$$

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A binary integer linear program

$$G = \{(i, j, k) \mid k \text{ is clue in cell } (i, j)\}, \quad x_{ijk} = \begin{cases} 1 & \text{if } k \text{ is filled in cell } (i, j) \\ 0 & \text{otherwise} \end{cases}$$

$$(SUDOKU - BLP) \quad \min \sum_{i, j, k \in S} 0 \cdot x_{ijk} \quad (1)$$

$$\text{s.t.} \quad x_{ijk} = 1 \quad \forall (i, j, k) \in G \quad (2)$$

$$\sum_{k \in S} x_{ijk} = 1 \quad \forall i \in S, j \in S \quad (3)$$

$$\sum_{i \in S} x_{ijk} = 1 \quad \forall j \in S, k \in S \quad (4)$$

$$\sum_{j \in S} x_{ijk} = 1 \quad \forall i \in S, k \in S \quad (5)$$

$$\sum_{i=p}^{p+2} \sum_{j=q}^{q+2} x_{ijk} = 1 \quad \forall k \in S, (p, q) \in B \quad (6)$$

$$x_{ijk} \in \{0, 1\} \quad \forall i \in S, j \in S, k \in S \quad (7)$$

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A non-linear integer programming formulation

F is the set of clue cells
 $g(i, j)$ is clue inside cell (i, j)

$$(SUDOKU - NLIP) \quad \min \sum_{i,j \in S} 0 \cdot x_{ij} \quad (8)$$

$$\text{s.t.} \quad x_{ij} = g(i, j) \quad \forall (i, j) \in F \quad (9)$$

$$|x_{ij} - x_{ik}| \geq 1 \quad \forall i, j, k \in S, j < k \quad (10)$$

$$|x_{ij} - x_{kj}| \geq 1 \quad \forall i, j, k \in S, i < k \quad (11)$$

$$|x_{p+i, q+j} - x_{p+k, q+l}| \geq 1 \quad \forall (p, q) \in B, \\ i, j, k, l \in \{0, 1, 2\} : 3i + j < 3k + l \quad (12)$$

$$x_{ij} \in S \quad \forall i, j \in S \quad (13)$$

Linearize non-linear constraints

$$y^+ = \begin{cases} y & \text{if } y > 0, \\ 0 & \text{if } y \leq 0, \end{cases} \quad y^- = \begin{cases} 0 & \text{if } y \geq 0, \\ -y & \text{if } y < 0, \end{cases}$$
$$a^+ = \begin{cases} 1 & \text{if } y > 0, \\ 0 & \text{if } y \leq 0, \end{cases} \quad a^- = \begin{cases} 1 & \text{if } y < 0, \\ 0 & \text{if } y \geq 0. \end{cases}$$

Since $|y| = y^+ + y^-$, non-linear constraint

$$1 \leq |y| \leq 8, y \in \mathbb{Z}$$

can be linearized by

$$y^+ - y^- = y \tag{14}$$

$$y^+ + y^- \geq 1 \tag{15}$$

$$a^+ \leq y^+ \leq 8a^+ \tag{16}$$

$$a^- \leq y^- \leq 8a^- \tag{17}$$

$$a^+ + a^- = 1 \tag{18}$$

$$y^+, y^- \in \mathbb{Z} \tag{19}$$

$$a^+, a^- \in \{0, 1\} \tag{20}$$

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X-Sudoku

5							3	4
	7				5			
4				8				1
			4		6		9	2
6	4	2		3			1	
9			1	2				
			6					
						6		
						3		

(X1) Each digit from 1 to 9 appears exactly once in each of the two main diagonals of the grid.

$$\sum_{i \in S} x_{iik} = 1 \quad \forall k \in S \quad (21)$$

$$\sum_{i \in S} x_{i,10-i,k} = 1 \quad \forall k \in S \quad (22)$$

Windoku

9			4		3			8
	1			9				
3		4				6		2
				7			4	
			5		8			
	2			3				
2		5				3		6
				4			8	
1			3		6			7

(W1) Each digit from 1 to 9 appears exactly once in each of the four additional blocks of the grid.

$$B_W := \{(2, 2), (2, 6), (6, 2), (6, 6)\}$$

$$\sum_{i=p}^{p+2} \sum_{j=q}^{q+2} x_{ijk} = 1 \quad \forall k \in S, (p, q) \in B_W \quad (23)$$

Jigsaw Sudoku

	5			1	
8				9	
3				6	7
	5				
1	2			3	
			1	5	
2					5
			9	4	6
		1		9	

(J1) Each digit from 1 to 9 appears exactly once in each of the nine jigsaw blocks.

Let \mathcal{H} be the set of jigsaw blocks.

$$\sum_{(i,j) \in H} x_{ijk} = 1 \quad \forall k \in S, H \in \mathcal{H} \quad (24)$$

Killer Sudoku

(K1) No digit appears more than once in a cage.

7	19			25		12	25	
7		9						10
17	12		9					
	10			10	20		12	
			12			17		
14	15				12		12	
	7		11			19		10
18			3				11	
12			28					

$$\sum_{(i,j) \in C} x_{ijk} \leq 1 \quad \forall C \in \mathcal{C}, k \in S \quad (25)$$

(K2) The sum of all digits filled in a cage must equal the associated number of the cage.

$$\sum_{(i,j) \in C} \sum_{k \in S} kx_{ijk} = n(C) \quad \forall C \in \mathcal{C} \quad (26)$$

where \mathcal{C} is the set of all cages,
 $n(C)$ is the number associated with cage C

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X-Sudoku

5							3	4
	7				5			
4				8				1
			4		6		9	2
6	4	2		3			1	
9			1	2				
			6					
						6		
						3		

(X1) 'Each digit from 1 to 9 appears exactly once in each of the two main diagonals of the grid'

$$|x_{ii} - x_{jj}| \geq 1 \quad \forall i, j \in S : i < j \quad (27)$$

$$|x_{i,10-i} - x_{j,10-j}| \geq 1 \quad \forall i, j \in S : i < j \quad (28)$$

Windoku

9			4		3			8
	1			9				
3		4				6		2
				7			4	
			5		8			
	2			3				
2		5				3		6
				4			8	
1			3		6			7

(W1) 'Each digit from 1 to 9 appears exactly once in each of the four additional blocks of the grid'

$$|x_{p+i,q+j} - x_{p+k,q+l}| \geq 1$$

$$\forall (p, q) \in B \cup B_W,$$

$$i, j, k, l \in \{0, 1, 2\} : 3i + j < 3k + l \quad (29)$$

Jigsaw Sudoku

	5				1			
	8					9		
3					6			7
		5						
	1	2				3		
				1		5		
2								5
				9			4	6
			1				9	

(J1) 'Each digit from 1 to 9 appears exactly once in each of the nine jigsaw blocks'

$$|x_{ij} - x_{kl}| \geq 1$$

$$\forall (i,j), (k,l) \in H : 9i + j < 9k + l \quad (30)$$

Killer Sudoku

7	19			25		12	25	
7		9						10
17	12		9					
	10			10	20		12	
			12			17		
14	15				12		12	
	7		11			19		10
18			3				11	
12			28					

(K1) No digit appears more than once in a cage.

$$|x_{ij} - x_{kl}| \geq 1$$

$$\forall (i, j), (k, l) \in C : 9i + j < 9k + l, \quad (31)$$

(K2) The sum of all digits filled in a cage must equal the associated number of the cage.

$$\sum_{(i,j) \in C} \sum_{k \in S} kx_{ijk} = n(C) \quad \forall C \in \mathcal{C}. \quad (32)$$

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Numerical experiments

- Model by ZIMPL; IP solver: SCIP
- Intel Core i3-4005U CPU, 64 GHz Processor, 4 GB RAM
- Performance (average running time in seconds) of IP formulations:

Variant	(<i>BLP</i>)	(<i>NLIP</i>)
Classical Sudoku	-	31.7
X-Sudoku	-	26.7
Windoku	-	298.8
Jigsaw Sudoku	-	220.5
Killer Sudoku	7.6	475.5

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Summary

- Classical Sudoku and some variants
 - X-Sudoku, Windoku, Jigsaw Sudoku, Killer Sudoku
- IP formulations
 - BLP: three - index binary variables
 - NLIP: two-index integer variables
- Numerical experiments
 - BLP outperforms NLIP
 - Killer Sudoku is hardest
 - Classical Sudoku & X-Sudoku is easiest

Thank you for your attention!