

slido



What game are you looking forward to?

ⓘ Start presenting to display the poll results on this slide.



Recursion



Recursion

Recursion happens in code **when a method calls itself**.

This can happen when a method calls itself directly or when a method calls another method that then calls the original method.

Recursion: Exit Conditions

- Recursion is another kind of loop.
- Recursion however has some overhead which can cause a stack overflow exception if the method calls itself too many times.
- To prevent a stack overflow exception, you need to add an **exit condition** to the method so that it stops calling itself.



Recursion: Exit Conditions

- ALL recursive methods need an exit condition.
- An exit condition happens when a method completes or returns without calling itself again.

Recursion: Exit Conditions

- This recursive method will eventually cause a stack overflow exception because there is **no exit condition**.

```
unsigned long factorial(unsigned int N)
{
    return N * factorial(N - 1);
}
```

Recursion: Exit Conditions

- To fix this recursive method, we need to **add an exit condition** that will stop the method from calling itself.

```
unsigned long factorial(unsigned int N)
{
    if (N <= 1) return 1; //here's the exit condition!
    return N * factorial(N - 1);
}
```



Sorting

Why do we sort?

- Essential for **searching** and **merging**. The list must be sorted before it can be searched or merged efficiently.
- For **UIs** and user experience (ex: sort list by price low-high). Humans like to see lists arranged in ways that help them.
- For **grouping** and **decision making** (ex: Z-order)

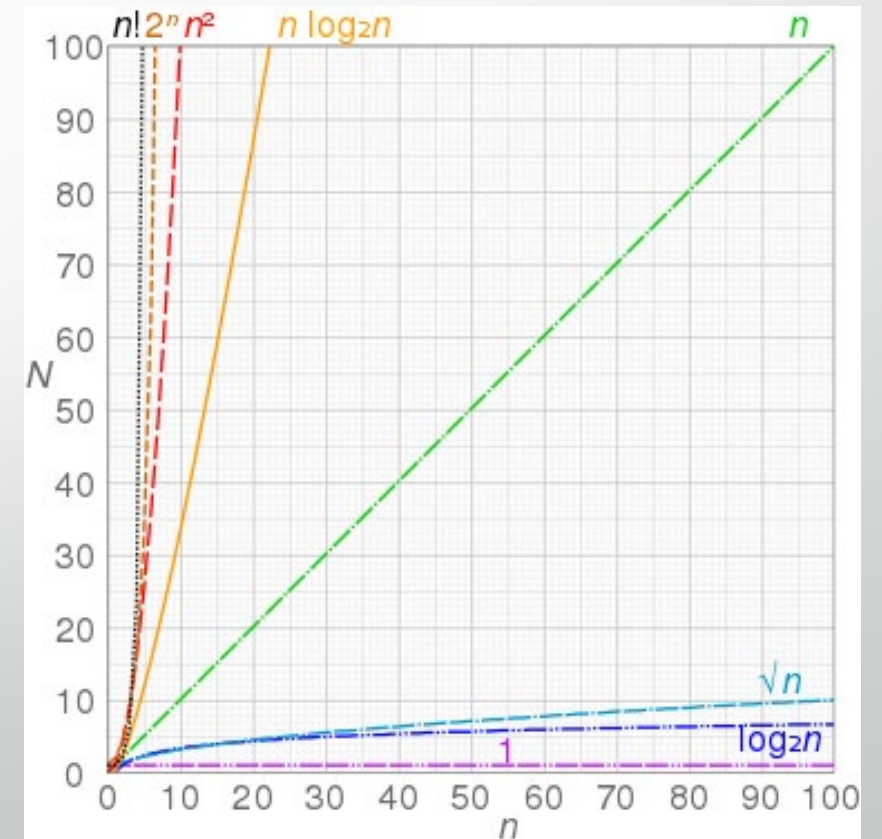
Performance: Big-O notation

Best



Worst

$O(1)$	constant
$O(\log n)$	logarithmic
$O(n)$	linear
$O(n \log n)$	loglinear
$O(n^2)$	quadratic
$O(2^n)$	exponential
$O(n!)$	factorial



Sorting Performance

ALGORITHM	WORST	BEST	AVG
<u>Heap Sort</u>	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
<u>Quick Sort</u>	$O(n^2)$	$O(n \log n)$	$O(n \log n)$
<u>Merge Sort</u>	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
<u>Insertion Sort</u>	$O(n^2)$	$O(n)$	$O(n^2)$
<u>Bubble Sort</u>	$O(n^2)$ comparisons, $O(n^2)$ swaps	$O(n)$ comps, $O(1)$ swaps	$O(n^2)$ comparisons, $O(n^2)$ swaps



Bubble Sort

Bubble Sort Algorithm

- The Algorithm
 1. Compare **each pair** of adjacent elements from the beginning of the array
 - If they are in reversed order, **swap** them
 2. If at least one **swap** was made, repeat step 1

Bubble Sort Pseudocode

procedure bubbleSort(A : list of sortable items)

 n := length(A)

 repeat //this is the start of a loop

 swapped := false

 for i := 1 to n - 1 inclusive do

 if A[i - 1] > A[i] then

 swap(A[i - 1], A[i]) //you'll need swap code here

 swapped = true

 end if

 end for

 n := n - 1

 while we swapped something //this is the end of the loop

end procedure



Swap Logic

Swap logic

- How to swap 2 items (A and B) in a vector
 - Copy A item to a temp variable.
 - Overwrite A item in the vector with the B item.
 - Overwrite the B item in the vector with the temp variable.

EXAMPLE: if vec holds ints...

```
int temp = vec[i - 1];  
vec[i - 1] = vec[i];  
vec[i] = temp;
```


Swap logic

- How to swap 2 items (A and B) in a vector
 - Use the `std::swap` method.
 - Use `#include <string_view>`

```
std::swap(vec[i - 1], vec[i]);
```



Comparing Strings



Comparing Strings

- When comparing strings for the purpose of sorting them, we need **more** information than whether the two strings are equal.
- We need to know “does this string come before the other string?”

Comparing Strings

- There are many ways to compare strings. We'll consider 2 approaches:
 - Using strcmp functions
 - Using std::string::compare
- Both approaches return an int value telling you how the strings compare.
 - So, comparing string s1 to string s2...
 - `< 0` means s1 is **LESS THAN** (comes before) s2
 - `0` means s1 is **EQUAL TO** s2
 - `> 0` means s1 is **GREATER THAN** (comes after) s2

Comparing Strings

- These versions of the strcmp function are case **insensitive**, meaning that it ignores the casing of the strings when comparing. Bob is equal to bob.

[_stricmp, _wcsicmp, _mbstricmp, _stricmp_l, _wcsicmp_l, _mbstricmp_l](#)

[Microsoft Learn](#)

```
std::string s1 = "Batman", s2 = "Aquaman";
```

```
int compResult = _stricmp(s1.c_str(), s2.c_str());
```

**** since _stricmp requires a pointer to a char array, you must call c_str() ****

Comparing Strings

- `std::string::compare`
- Performs a **case-sensitive** comparison

```
std::string s1 = "Batman", s2 = "Aquaman";  
int compareResult = s1.compare(s2);
```

If you want to do a case-insensitive comparison, one way would be to make both strings all uppercase or lowercase. (see next slide on how)

Comparing Strings

- Making a `std::string` uppercase.
- There are quite a few ways to make a string uppercase. Here is one way:
- `for (auto& c : str) c = toupper(c);`
- NOTE: this will modify the string. So if you want to keep the original string as-is, you'll need to make a copy of the string and modify the copy.



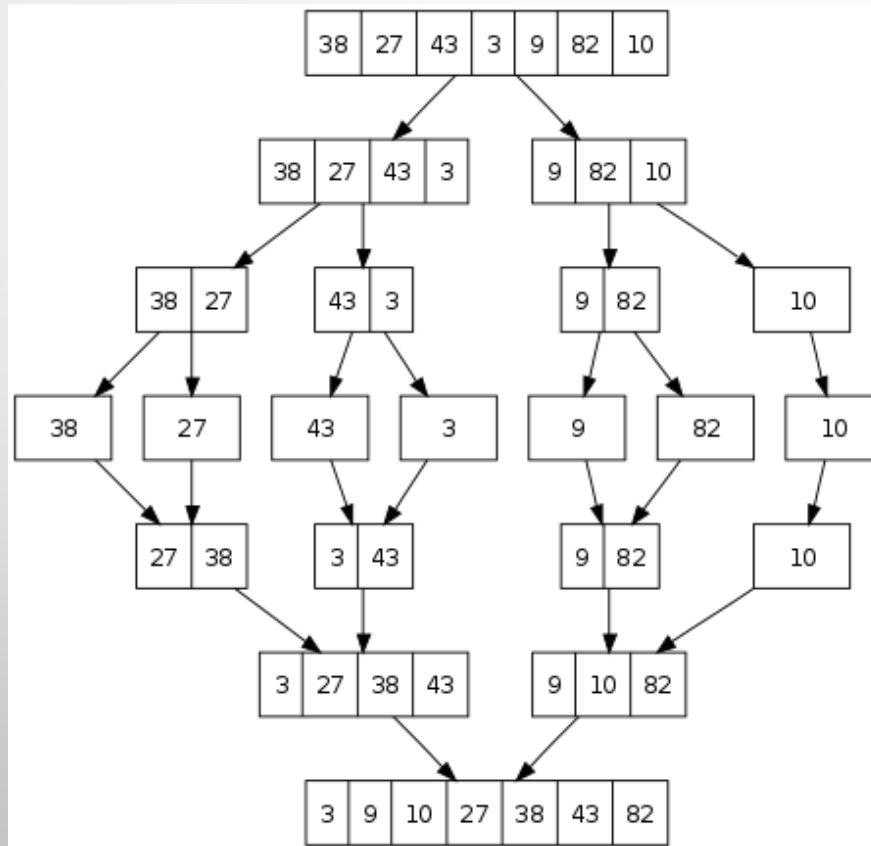
Merge Sort

Merge Sort Algorithm

- The Algorithm
 1. Divide the list into n sublists that are all 1 element long.
 2. Merge the sublists to make new ***sorted*** lists.
 3. Continue merging until only 1 list remains.

Merge Sort Algorithm

- https://en.wikipedia.org/wiki/File:Merge_sort_algorithm_diagram.svg



function merge_sort(list m) is

// exit condition. A list of zero or one elements is sorted, by definition.

if length of m ≤ 1 then

return m

// Recursive case. First, divide the list into sublists

// consisting of the first half and second half of the list. This assumes lists start at index 0.

var left := empty list

var right := empty list

for i = 0 to length(m) do

if i < (length of m)/2 then

add m[i] to left

else

add m[i] to right

// Recursively sort both sublists.

left := merge_sort(left)

right := merge_sort(right)

// Then merge the now-sorted sublists.

return merge(left, right)



```
function merge(left, right) is
```

```
  var result := empty list
```

```
  while left is not empty and right is not empty do
```

```
    if first(left) ≤ first(right) then
```

```
      add first(left) to result
```

```
      remove first from left
```

```
    else
```

```
      add first(right) to result
```

```
      remove first from right
```

```
  // Either left or right may have elements left; consume them.
```

```
  // (Only one of the following loops will actually be entered.)
```

```
  while left is not empty do
```

```
    add first(left) to result
```

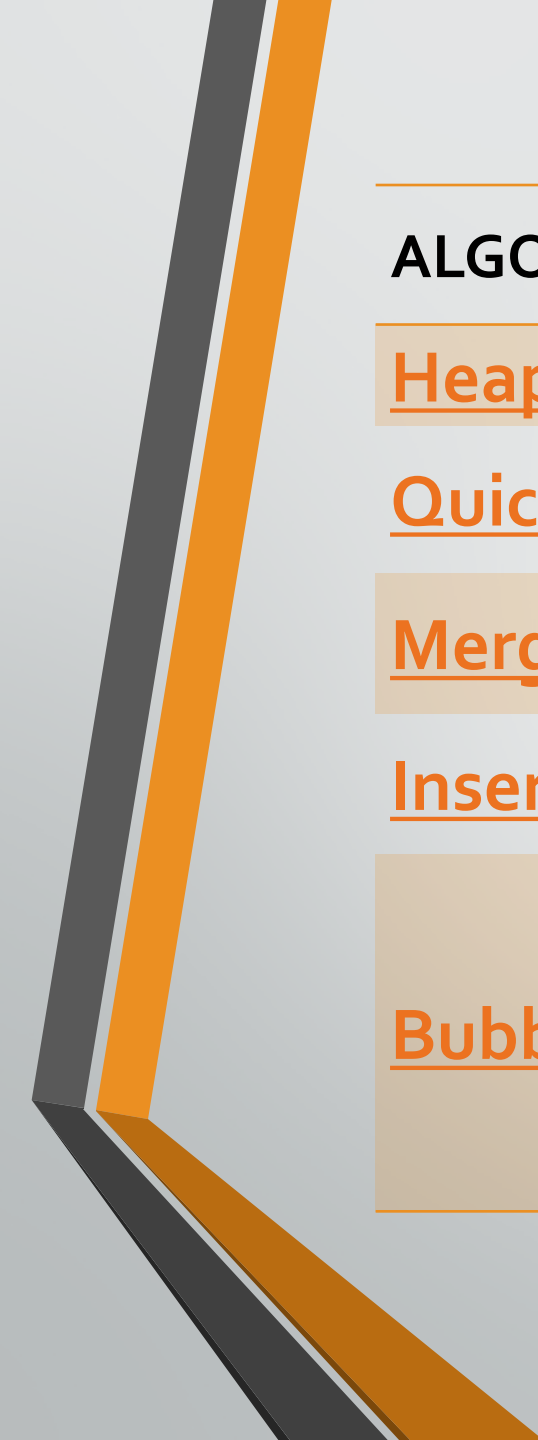
```
    remove first from left
```

```
  while right is not empty do
```

```
    add first(right) to result
```

```
    remove first from right
```

```
  return result
```



ALGORITHM	WORST	BEST	AVG
<u>Heap Sort</u>	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
<u>Quick Sort</u>	$O(n^2)$	$O(n \log n)$	$O(n \log n)$
<u>Merge Sort</u>	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
<u>Insertion Sort</u>	$O(n^2)$	$O(n)$	$O(n^2)$
<u>Bubble Sort</u>	$O(n^2)$ comparisons, $O(n^2)$ swaps	$O(n)$ comps, $O(1)$ swaps	$O(n^2)$ comparisons, $O(n^2)$ swaps

If you have 7 chairs and 10 children, what do you do?



slido



Recursion happens when...

ⓘ Start presenting to display the poll results on this slide.

slido



What do all recursive methods need?

ⓘ Start presenting to display the poll results on this slide.

slido



What happens when a recursive method calls itself too many times?

ⓘ Start presenting to display the poll results on this slide.

slido



What will be returned?

```
std::string s1 = "Batman", s2 = "Batmen";  
int compResult = _stricmp(s1.c_str(), s2.c_str());
```

① Start presenting to display the poll results on this slide.