



Introduction to Digital Design

Week 4: Boolean Algebra, Multi-Output Circuits

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Overview

- Boolean Operator
 - $()$, $'$, $*$, $+$.
 - precedence.
- Terminology
 - variable, literal, product term, sum-of-products.
- Properties
 - Commutative, distributive, associative, identity, complement.
 - Null elements, idempotent, involution, DeMorgan.
- Representation
 - Truth table, Boolean equation, circuit.
 - Canonical form, multiple-output circuits.

2

Boolean Algebra

2.5

- By defining logic gates based on Boolean algebra, we can *use algebraic methods to manipulate circuits*
- Notation: Writing a AND b, a OR b, NOT(a) is cumbersome
 - Use symbols: $a * b$ (or just ab), $a + b$, and a'
 - Original: $w = (p \text{ AND NOT}(s) \text{ AND } k) \text{ OR } t$
 - New: $w = ps'k + t$
 - Spoken as "w equals p and s prime k, or t"
 - Or just "w equals p s prime k, or t"
 - s' known as "complement of s"
 - While symbols come from regular algebra, **don't** say "times" or "plus"
 - "product" and "sum" are OK and commonly used

Boolean algebra precedence, highest precedence first.

Symbol	Name	Description
$()$	Parentheses	Evaluate expressions nested in parentheses first
$'$	NOT	Evaluate from left to right
$*$	AND	Evaluate from left to right
$+$	OR	Evaluate from left to right

3

Boolean Algebra Operator Precedence

- Evaluate the following Boolean equations, assuming $a=1$, $b=1$, $c=0$, $d=1$.
 - Q1. $F = a * b + c$.
 - Answer: * has precedence over +, so we evaluate the equation as $F = (1 * 1) + 0 = (1) + 0 = 1 + 0 = 1$.
 - Q2. $F = ab + c$.
 - Answer: the problem is identical to the previous problem, using the shorthand notation for *.
 - Q3. $F = ab'$.
 - Answer: we first evaluate b' because NOT has precedence over AND, resulting in $F = 1 * (1') = 1 * (0) = 1 * 0 = 0$.
 - Q4. $F = (ac)'$.
 - Answer: we first evaluate what is inside the parentheses, then we NOT the result, yielding $(1 * 0)' = (0)' = 0' = 1$.
 - Q5. $F = (a + b') * c + d'$.
 - Answer: Inside left parentheses: $(1 + (1')) = (1 + (0)) = (1 + 0) = 1$. Next, * has precedence over +, yielding $(1 * 0) + 1' = (0) + 1'$. The NOT has precedence over the OR, giving $(0) + (1') = (0) + (0) = 0 + 0 = 0$.

Boolean algebra precedence, highest precedence first.

Symbol	Name	Description
$()$	Parentheses	Evaluate expressions nested in parentheses first
$'$	NOT	Evaluate from left to right
$*$	AND	Evaluate from left to right
$+$	OR	Evaluate from left to right

4

Boolean Algebra Terminology

- Example equation: $F(a,b,c) = a'bc + abc' + ab + c$
- **Variable**
 - Represents a value (0 or 1)
 - Three variables: a, b, and c
- **Literal**
 - Appearance of a variable, in true or complemented form
 - Nine literals: a' , b, c, a, b, c', a, b, and c
- **Product term**
 - Product of literals
 - Four product terms: $a'bc$, abc' , ab , c
- **Sum-of-products**
 - Equation written as OR of product terms only
 - Above equation is in sum-of-products form. "F = (a+b)c + d" is not.

5

Boolean Algebra Properties

- Commutative
 - $a + b = b + a$
 - $a * b = b * a$
- Distributive
 - $a * (b + c) = a * b + a * c$
 - Can write as: $a(b+c) = ab + ac$
 - $a + (b * c) = (a + b) * (a + c)$
 - (This second one is tricky!)
 - Can write as: $a+(bc) = (ab)(ac)$
- Associative
 - $(a + b) + c = a + (b + c)$
 - $(a * b) * c = a * (b * c)$
- Identity
 - $0 + a = a + 0 = a$
 - $1 * a = a * 1 = a$
- Complement
 - $a + a' = 1$
 - $a * a' = 0$
- To prove, just evaluate all possibilities

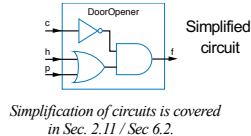
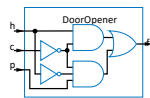
Example uses of the properties

- Show abc' equivalent to $c'b'a$.
 - Use commutative property:
 - $a'b'c' = a'c''b'' = c''a''b'' = c''b''a''$
- Show $abc + abc' = ab$.
 - Use first distributive property
 - $abc + abc' = ab(c+c')$
 - Complement property
 - Replace $c+c'$ by 1: $ab(c+c') = ab(1)$.
 - Identity property
 - $ab(1) = ab * 1 = ab$.
- Show $x + x'z$ equivalent to $x + z$.
 - Second distributive property
 - Replace $x+x'z$ by $(x+x')(x+z)$.
 - Complement property
 - Replace $(x+x')$ by 1.
 - Identity property
 - replace $1*(x+z)$ by $x+z$.

6

Example that Applies Boolean Algebra Properties

- Want automatic door opener circuit (e.g., for grocery store)
 - Output: $f=1$ opens door
 - Inputs:
 - $p=1$: person detected
 - $h=1$: switch forcing hold open
 - $c=1$: key forcing closed
 - Want open door when
 - $h=1$ and $c=0$, or
 - $h=0$ and $p=1$ and $c=0$
 - Equation: $f = hc' + h'pc'$
- Can the circuit be simplified?
 - $f = hc' + h'pc'$
 - $f = c'h + c'h'p$ (by the commutative property)
 - $f = c'(h + h'p)$ (by the first distrib. property)
 - $f = c'((h+h')(h+p))$ (2nd distrib. prop.; tricky one)
 - $f = c'((1)(h+p))$ (by the complement property)
 - $f = c'(h+p)$ (by the identity property)



Simplification of circuits is covered in Sec. 2.11 / Sec 6.2.

7

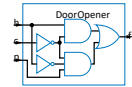
Example that Applies Boolean Algebra Properties



- Found inexpensive chip that computes:
 - $f = c'hp + c'hp' + c'h'p$
 - Can we use it for the door opener?
 - Is it the same as $f = hc' + h'pc'$?
- Apply Boolean algebra:
 - Commutative**
 - $a + b = b + a$
 - $a * b = b * a$
 - Distributive**
 - $a * (b + c) = a * b + a * c$
 - $a + (b * c) = (a + b) * (a + c)$
 - Associative**
 - $(a + b) + c = a + (b + c)$
 - $(a * b) * c = a * (b * c)$
 - Identity**
 - $0 + a = a + 0 = a$
 - $1 * a = a * 1 = a$
 - Complement**
 - $a + a' = 1$
 - $a * a' = 0$

$$\begin{aligned}
 f &= c'hp + c'hp' + c'h'p \\
 &= c'h(p + p') + c'h'p \quad (\text{by the distributive property}) \\
 &= c'h(1) + c'h'p \quad (\text{by the complement property}) \\
 &= c'h + c'h'p \quad (\text{by the identity property}) \\
 &= hc' + h'pc' \quad (\text{by the commutative property})
 \end{aligned}$$

Same! Yes, we can use it.



8

Example that Applies Boolean Algebra Properties

[Video](#)

9

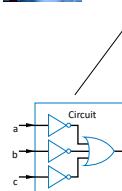
Boolean Algebra: Additional Properties

- Null elements**
 - $a + 1 = 1$
 - $a * 0 = 0$
- Idempotent Law**
 - $a + a = a$
 - $a * a = a$
- Involution Law**
 - $(a')' = a$
- DeMorgan's Law**
 - $(a + b)' = a'b'$
 - $(ab)' = a' + b'$
 - Very useful!
- To prove, just evaluate all possibilities

10

Example Applying DeMorgan's Law

Aircraft lavatory sign example



- Behavior**
 - Three lavatories, each with sensor (a, b, c), equals 1 if door locked
 - Light "Available" sign (S) if any lavatory available
- Equation and circuit**
 - $S = a' + b' + c'$
- Transform**
 - $(abc)' = a' + b' + c'$ (by DeMorgan's Law)
 - $S = (abc)'$
- New circuit**
 - $S' = (a')' * (b')' * (c')'$ (by DeMorgan's Law)
 - $S' = a * b * c$ (by Involution Law)
- Makes intuitive sense**
 - Occupied if all doors are locked

$$\begin{aligned}
 (a + b)' &= a'b' \\
 (ab)' &= a' + b'
 \end{aligned}$$

11

Example Applying Properties

- Commutative**
 - $a + b = b + a$
 - $a * b = b * a$
- Distributive**
 - $a * (b + c) = a * b + a * c$
 - $a + (b * c) = (a + b) * (a + c)$
- Associative**
 - $(a + b) + c = a + (b + c)$
 - $(a * b) * c = a * (b * c)$
- Identity**
 - $0 + a = a + 0 = a$
 - $1 * a = a * 1 = a$
- Complement**
 - $a + a' = 1$
 - $a * a' = 0$
- Null elements**
 - $a + 1 = 1$
 - $a * 0 = 0$
- Idempotent Law**
 - $a + a = a$
 - $a * a = a$
- Involution Law**
 - $(a')' = a$
- DeMorgan's Law**
 - $(a + b)' = a'b'$
 - $(ab)' = a' + b'$

- For door opener $f = c'(h+p)$, prove door stays closed ($f=0$) when $c=1$

$$\begin{aligned}
 f &= c'(h+p) \\
 \text{Let } c &= 1 \quad (\text{door forced closed}) \\
 f &= 1'(h+p) \\
 &= f = 0(h+p) \quad (\text{by the distributive property}) \\
 &= f = 0 + 0 \quad (\text{by the null elements property}) \\
 &= f = 0
 \end{aligned}$$

12

Complement of a Function

- Commonly want to find complement (inverse) of function F
 - 0 when F is 1; 1 when F is 0
- Use DeMorgan's Law repeatedly
 - Note: DeMorgan's Law defined for more than two variables, e.g.:
 - $(a + b + c)' = (abc)'$
 - $(abc)' = (a' + b' + c')$
- Complement of $f = w'xy + wx'y'z'$
 - $f' = (w'xy + wx'y'z')'$
 - $f' = (w'xy)'(wx'y'z')'$ (by DeMorgan's Law)
 - $f' = (w+x'+y')(w'+x+y+z)$ (by DeMorgan's Law)
- Can then expand into sum-of-products form

13

Example that Applies Boolean Algebra Properties

[Video](#)

14

Representations of Boolean Functions

2.6

- English 1: F outputs 1 when a is 0 and b is 0, or when a is 0 and b is 1.
- English 2: F outputs 1 when a is 0, regardless of b 's value
- Equation 1: $F(a,b) = a'b' + a'b$
- Equation 2: $F(a,b) = a'$
- (a)
- (b)
- (c)
- (d) Truth table
- | a | b | F |
|---|---|---|
| 0 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 0 | 0 |
| 1 | 1 | 0 |
- The function F
- A function can be represented in different ways
 - Above shows seven representations of the same functions $F(a,b)$, using four different methods: English, Equation, Circuit, and Truth Table

15

Truth Table Representation of Boolean Functions

- Define value of F for each possible combination of input values
 - 2-input function: 4 rows
 - 3-input function: 8 rows
 - 4-input function: 16 rows
- Q: Use truth table to define function $F(a,b,c)$ that is 1 when abc is 5 or greater in binary

a	b	c	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

(a)

a	b	c	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

(b)

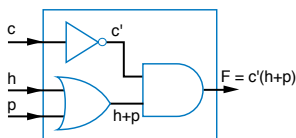
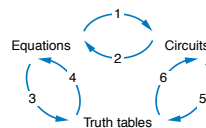
a	b	c	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

(c)

16

Converting among Representations

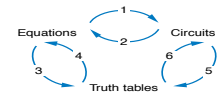
- Can convert from any representation to another
- Common conversions
 - Equation to circuit (we did this earlier)
 - Circuit to equation
 - Start at inputs, write expression of each gate output



17

Converting among Representations

- More common conversions
 - Truth table to equation (which we can then convert to circuit)
 - Easy—just OR each input term that should output 1
 - Equation to truth table
 - Easy—just evaluate equation for each input combination (row)
 - Creating intermediate columns helps



Inputs	Outputs	Term
a	b	F
0	0	1
0	1	1
1	0	0
1	1	0

$F = a'b' + a'b$

Q: Convert to equation

a	b	c	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$F = ab'c + abc' + abc$

Q: Convert to truth table: $F = a'b' + a'b$

Inputs	Outputs
a	b
0	0
0	1
1	0
1	1

18

Example: Converting from Truth Table to Equation

- Parity bit: Extra bit added to data, intended to enable detection of error (a bit changed unintentionally)
 - e.g., errors can occur on wires due to electrical interference
- Even parity: Set parity bit so total number of 1s (data + parity) is even
 - e.g., if data is 001, parity bit is 1 → 0011 has even number of 1s
- Want equation, but easiest to start from truth table for this example

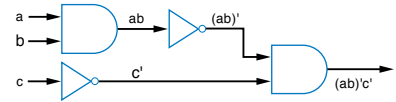
a	b	c	P
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

Convert to eqn.

$$P = a'b'c + a'bc' + ab'c' + abc$$

Example: Converting from Circuit to Truth Table

- First convert to circuit to equation, then equation to table



Inputs						Outputs
a	b	c	ab	(ab)'	c'	F
0	0	0	0	1	1	1
0	0	1	0	1	0	0
0	1	0	0	1	1	1
0	1	1	0	1	0	0
1	0	0	0	1	1	1
1	0	1	0	1	0	0
1	1	0	1	0	1	0
1	1	1	1	0	0	0

20

Example: Converting from Circuit to Truth Table

[Video](#)

21

Standard Representation: Truth Table

- How can we determine if two functions are the same?
 - Recall automatic door example
 - Same as $f = hc' + h'pc'$?
 - Used algebraic methods
 - But if we failed, does that prove *not* equal? No.
- Solution: Convert to truth tables
 - Only ONE truth table representation of a given function
 - Standard representation—for given function, only one version in standard form exists

$$\begin{aligned} f &= c'hp + c'hp' + c'h'p \\ f &= c'h(p + p') + c'h'p \\ f &= c'h(1) + c'h'p \\ f &= c'h + c'h'p \end{aligned}$$

(what if we stopped here?)
 $f = hc' + h'pc'$

Q: Determine if $F=ab+a'$ is same function as $F=a'b'+a'b+ab$, by converting each to truth table first

F = ab + a'			F = a'b' + a'b + ab		
a	b	F	a	b	F
0	0	1	0	0	1
0	1	1	0	1	1
1	0	0	1	0	0
1	1	1	1	1	1

Same

22

Truth Table Canonical Form

- Q: Determine via truth tables whether $ab+a'$ and $(a+b)'$ are equivalent

F = ab + a'			F = (a+b)'		
a	b	F	a	b	F
0	0	1	0	0	1
0	1	1	0	1	0
1	0	0	1	0	0
1	1	1	1	1	0

Not equivalent

23

Canonical Form – Sum of Minterms

- Truth tables too big for numerous inputs
- Use standard form of equation instead
 - Known as **canonical form**
 - Regular algebra: group terms of polynomial by power
 - $ax^2 + bx + c$ ($3x^2 + 4x + 2x^2 + 3 + 1 \rightarrow 5x^2 + 4x + 4$)
 - Boolean algebra: create sum of minterms
 - Minterm**: product term with every function literal appearing exactly once, in true or complemented form
 - Just multiply-out equation until sum of product terms
 - Then expand each term until all terms are minterms

Q: Determine if $F(a,b)=ab+a'$ is equivalent to $F(a,b)=a'b'+a'b+ab$, by converting first equation to canonical form (second already is)

$$\begin{aligned} F &= ab+a' \text{ (already sum of products)} \\ F &= ab + a'(b+b') \text{ (expanding term)} \\ F &= ab + a'b + a'b' \text{ (Equivalent – same three terms as other equation)} \end{aligned}$$

24

Canonical Form – Sum of Minterms

- Q: Determine whether the functions $G(a,b,c,d,e) = abcd + a'bcd$ and $H(a,b,c,d,e) = abcd + abcd' + a'bcd + a'bcd'(a' + c)$ are equivalent.

$$G = abcd + a'bcd$$

$$G = abcd(e+e') + a'bcd$$

$$G = abcd + abcd' + a'bcd$$

$$G = a'bcd + abcd' + abcd \quad (\text{sum of minterms form})$$

Equivalent

$$\begin{aligned} H &= abcd + abcd' + a'bcd + a'bcd'(a' + c) \\ H &= abcd + abcd' + a'bcd + a'bcd'a' + a'bcd'c \\ H &= abcd + abcd' + a'bcd + a'bcd + a'bcd \\ H &= abcd + abcd' + a'bcd \\ H &= a'bcd + abcd' + abcd \end{aligned}$$

25

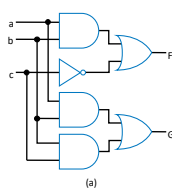
Compact Sum of Minterms Representation

- List each minterm as a number
- Number determined from the binary representation of its variables' values
 - $a'bcd$ corresponds to 01111, or 15
 - $abcd'$ corresponds to 11110, or 30
 - $abcd$ corresponds to 11111, or 31
- Thus, $H = a'bcd + abcd' + abcd$ can be written as:
 - $H = \sum m(15, 30, 31)$
 - "H is the sum of minterms 15, 30, and 31"

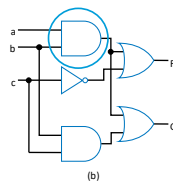
26

Multiple-Output Circuits

- Many circuits have more than one output
- Can give each a separate circuit, or can share gates
- Ex: $F = ab + c'$, $G = ab + bc$



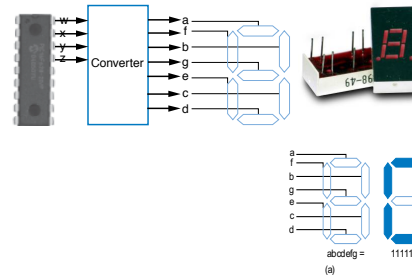
Option 1: Separate circuits



Option 2: Shared gates

27

Multiple-Output Example: BCD to 7-Segment Converter

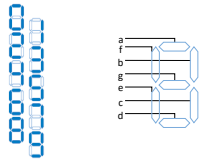


28

Multiple-Output Example: BCD to 7-Segment Converter

TABLE 2-4 4-bit binary number to seven-segment display truth table

w	x	y	z	a	b	c	d	e	f	g
0	0	0	0	1	1	1	1	1	0	0
0	0	0	1	0	1	1	0	0	0	0
0	0	1	0	1	1	0	1	1	0	1
0	0	1	1	1	1	1	1	0	0	1
0	1	0	0	0	1	1	0	0	1	1
0	1	0	1	1	0	1	1	0	1	1
0	1	1	0	1	0	1	1	1	1	1
0	1	1	1	1	1	1	0	0	0	0
1	0	0	0	1	1	1	1	1	1	1
1	0	0	1	1	1	1	1	0	1	1
1	0	1	0	0	0	0	0	0	0	0
1	0	1	1	0	0	0	0	0	0	0
1	1	0	0	0	0	0	0	0	0	0
1	1	0	1	0	0	0	0	0	0	0
1	1	1	0	0	0	0	0	0	0	0
1	1	1	1	0	0	0	0	0	0	0



$$\begin{aligned} a &= w'x'y'z' + w'x'yz' + w'x'y'z + w'xy'z + \\ &w'xyz' + w'xyz + wx'y'z' + wx'y'z \\ b &= w'x'y'z' + w'x'y'z + w'x'yz' + w'x'yz + \\ &w'xy'z' + w'xyz + wx'y'z' + wx'y'z \end{aligned}$$

29

Example: Multiple-Output Circuits

[Video](#)

30

Summary

- Boolean Operator
 - $()$, $'$, $*$, $+$.
 - precedence.
- Terminology
 - variable, literal, product term, sum-of-products.
- Properties
 - Commutative, distributive, associative, identity, complement.
 - Null elements, idempotent, involution, DeMorgan.
- Representation
 - Truth table, Boolean equation, circuit.
 - Canonical form, multiple-output circuits.

31