CSC429 – Computer Security

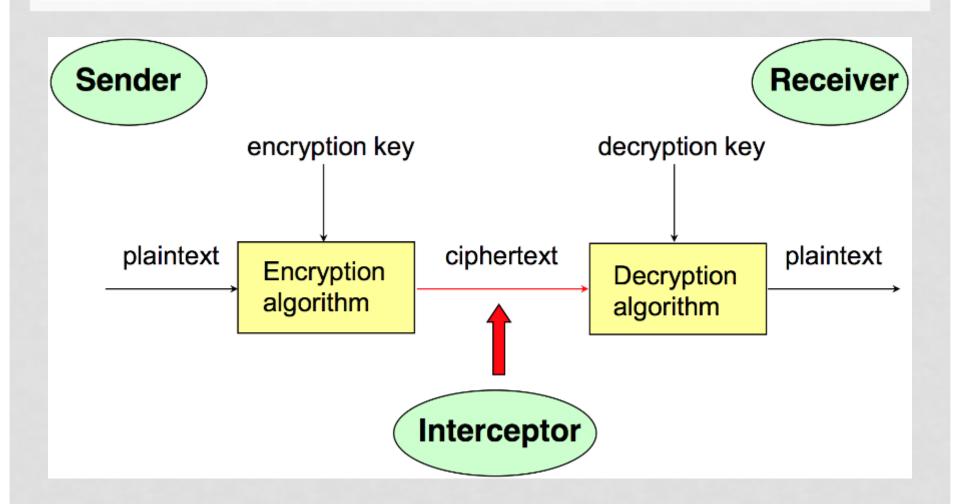
LECTURE 2 <u>INTRODUCTION TO CR</u>YPTOGRAPHY

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Notes

Make sure you have access to Piazza.

A Cryptosystem



Three Important Questions

- Can cryptography prevent a communication from being intercepted?
- Which of the following must be kept secret?
 - 1. Encryption algorithm
 - 2. Decryption algorithm
 - 3. Encryption key
 - 4. Decryption key
- Does using a good encryption algorithm guarantee the confidentiality of a message?

Models for Evaluating Security

- Unconditional (information-theoretic) security:
 - Adversary has unlimited resources.
 - Scheme that achieve such level is perfectly secret.
 - Analysis is done using probability theory.

Computational Security:

- Measures the amount of computational effort to defeat the system.
- Usually based on difficult mathematical problems (e.g. discrete logarithm, factoring, etc).

Classical Ciphers

Shift Cipher

- Each letter is shifted by K positions.
 - Can be modeled as a addition modulo 26.
- Encryption:
 - Shift to the right by K.
- Decryption:
 - Shift to the left by K.
- History:
 - Caesar cipher [K = 3].

Shift Cipher - 2

- A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
- 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25
- Example:
 - P = CRYPTOGRAPHYISFUN
 - K = 11
 - C = NCJAVZRCLASJTDQFY
- What is the key space?
- How can you brake it?

Mono-alphabetical Substitution Cipher

- The key space: all permutations of $\Sigma = \{A, B, C, ..., Z\}$
- Encryption given a key π :
 - each letter X in the plaintext P is replaced with $\pi(X)$
- Decryption given a key π :
 - each letter Y in the ciphertext P is replaced with $\pi^{-1}(Y)$

Substitution Cipher - 2

Example:

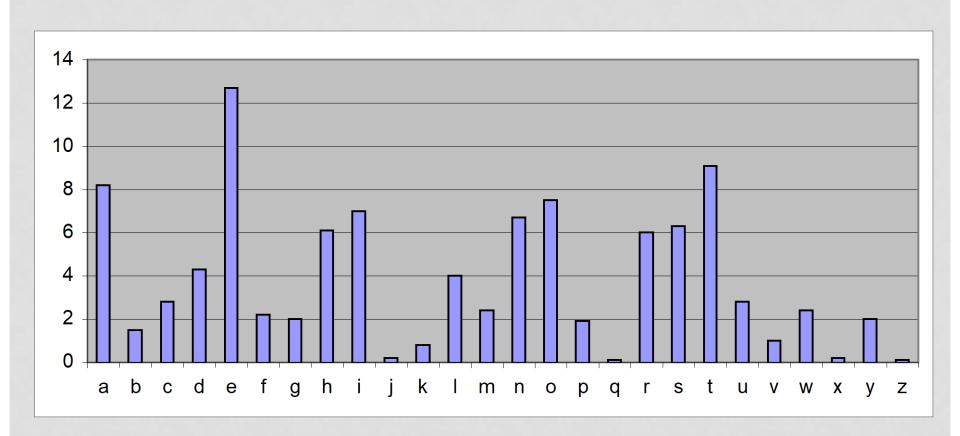
A B C D E F G H I J K L M N O P Q R S T U V W X Y Z π = B A D C Z H W Y G O Q X S V T R N M S K J I P F E U

- BECAUSE → AZDBJSZ
- What is the key space?
- Can it be broken?

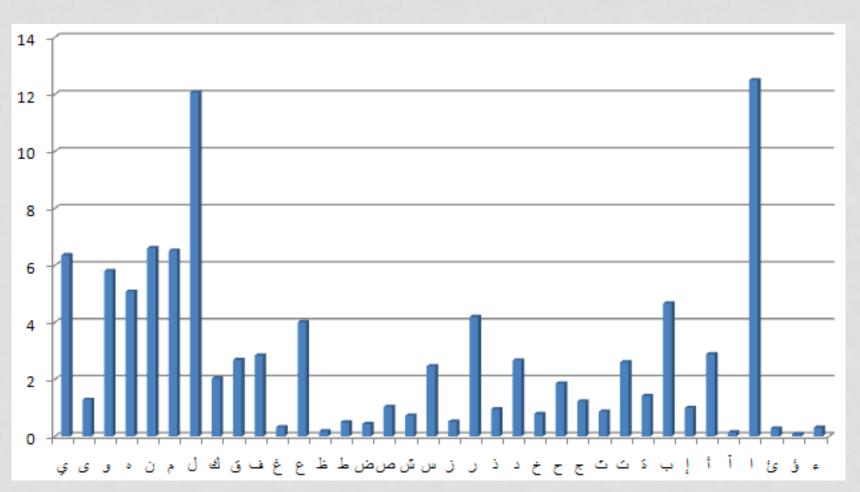
Breaking Substitution Cipher

- Each language has certain features frequency of letters.
- Substitution ciphers preserve the language features.
- Substitution ciphers are vulnerable to frequency analysis attacks.

Frequency of Letters in English



Frequency of Letters in Arabic



Polyalphabetic Substitution Ciphers

- Main weaknesses of mono-alphabetic substitution ciphers:
 - Each letter in the ciphertext corresponds to only one letter in the plaintext letter
- Lesson:
 - A large key space alone doesn't guarantee security.
- Lead to the development Vigenère cipher.

The Vigenère Cipher

- Given m, a positive integer, $P = C = (Z_{26})^n$, and $K = (k_1, k_2, ..., k_m)$ a key, we define:
 - Encryption:
 - $E_k(p_1, p_2...p_m) = (p_1+k_1, p_2+k_2...p_m+k_m) \pmod{26}$
 - Decryption:
 - $D_k(c_1, c_2... c_m) = (c_1-k_1, c_2-k_2... c_m-k_m) \pmod{26}$
- Example:
 - Plaintext: CRYPTOGRAPHY
 - Key: LUCKLUCKLUCK
 - Ciphertext: NLAZEIIBLJJI

Security of Vigenère Cipher

 Vigenère masks the frequency with which a character appears in a language: one letter in the ciphertext corresponds to multiple letters in the plaintext. Makes the direct use of frequency analysis more difficult.

• Is it secure?

Cryptanalysis of Vigenère Cipher

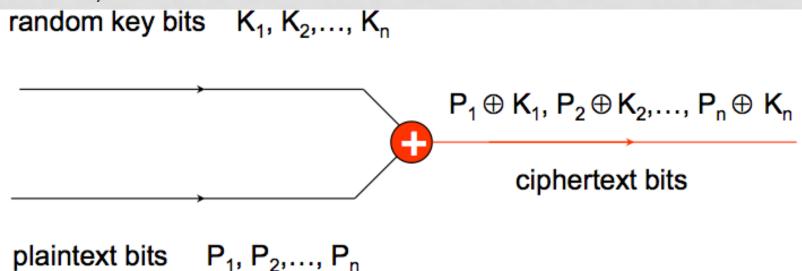
- Find the length of the key: (e.g. using Kasisky test).
- Divide the message into that many shift cipher encryptions.
- Use frequency analysis to solve the resulting shift ciphers.

One-Time Pad (OTP)

- Key is a random string that is at least as long as the plaintext.
- Encryption is similar to shift cipher.
- Let $Z_m = \{0,1,...,m-1\}$ be the alphabet.
 - Plaintext space = Ciphtertext space = Key space = (Z_m)ⁿ
 - The key is chosen uniformly randomly
 - Plaintext $X = (x_1 x_2 ... x_n)$
 - Key $K = (k_1 k_2 ... k_n)$
 - Ciphertext $Y = (y_1 y_2 ... y_n)$
 - $E_k(X) = (x_1+k_1 \ x_2+k_2 \dots x_n+k_n) \mod m$
 - $D_k(Y) = (y_1-k_1 \ y_2-k_2 ... \ y_n-k_n) \mod m$

One-Time Pad (OTP) - 2

- Key must be:
 - As long as the plaintext.
 - · Random.
 - Not be re-used.
- Binary Version:



Next Lecture

- We will start discussing modern cryptography.
- Readings for next lecture:
 - Anderson's book (5.3.2) and (5.3.3).