

# CSC429 – Computer Security

LECTURE 3  
MODERN CRYPTOGRAPHY 2

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# Data Encryption Standard (DES)

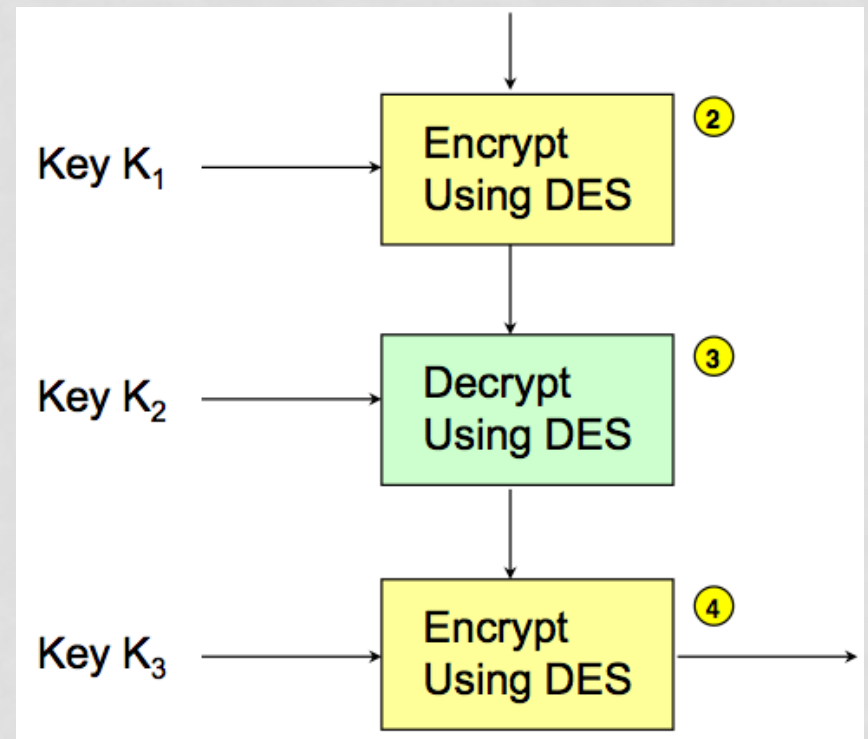
- Designed by IBM, with modifications proposed by the NSA.
- US national standard from 1977 to 2001
- De facto standard
- Block size 64 bits;
- Key size 56 bits
- Designed mostly for hardware implementations
- Considered insecure now
  - vulnerable to brute-force attacks.

# Searching for a DES Key

<b>Year</b>	<b>Source</b>	<b>Implemented?</b>	<b>(Estimated) Cost in US\$</b>	<b>(Estimated) Search time</b>
<b>1977</b>	Diffie Hellman	No	20 million	20 hours
<b>1993</b>	Wiener	No	10.5 million 1.5 million 600 000	21 minutes 3.5 hours 35 hours
<b>1997</b>	Internet	Yes	Unknown	140 days
<b>1998</b>	Deep Crack	Yes	210 000	56 hours
<b>2007</b>	COPACOBANA	Yes	<10,000	<7 days

# Triple DES (3DES)

- Use three different keys
- Key space is  $56 \times 3 = 168$  bits
- No known practical attack against it.



# Advanced Encryption Standard (AES)

- In 1997, NIST made a formal call for algorithms stipulating that the AES would specify an **unclassified, publicly disclosed encryption algorithm, available royalty-free, worldwide.**
- The algorithm must implement symmetric key cryptography as a block cipher and (at a minimum) support **block sizes of 128-bits and key sizes of 128-, 192-, and 256-bits.**
- In 1998, NIST selected 15 AES candidate algorithms.
- On October 2, 2000, NIST selected **Rijndael** (invented by Joan Daemen and Vincent Rijmen) to as the AES.

# Modern Cryptography

Block Ciphers Encryption Modes

# Block Cipher Encryption Modes

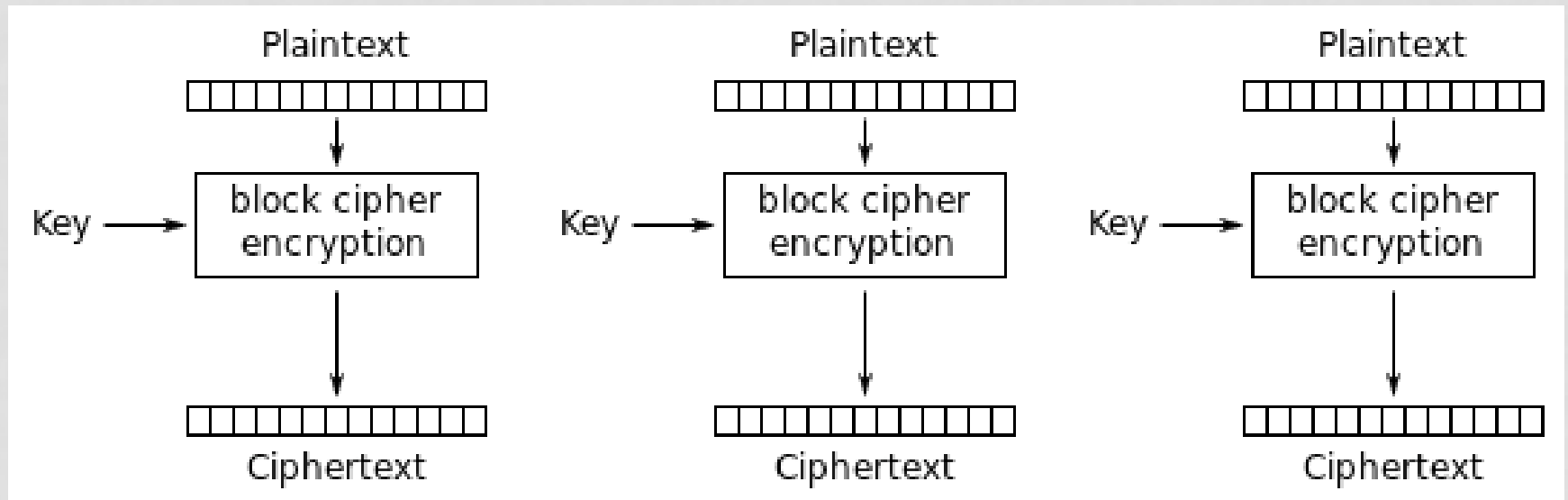
- A block cipher encrypts only one block.
- Needs a way to extend it to encrypt an arbitrarily long message.
- Want to ensure that if the block cipher is secure, then the encryption is secure.
- There are many modes: ECB, CBC, CTR, PCBC, CFB, OFB.
  - We will only discuss the first three.

# Mode 1 – Electronic Code Book (ECB)

- Message is broken into independent blocks of *block\_size* bits;
- Each block encrypted separately.
- **Encryption:**  $c_i = E_k(x_i)$
- **Decryption:**  $x_i = D_k(c_i)$



# ECB Encryption



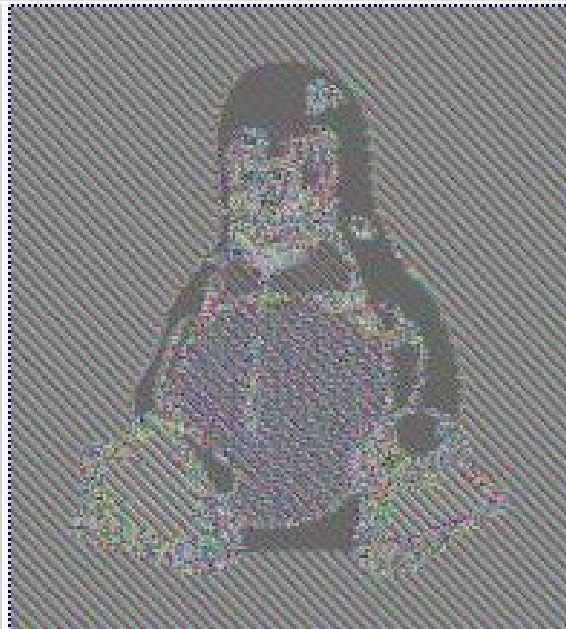
# ECB Properties

- Deterministic: the same data block gets encrypted the same way, **reveals patterns of data when a data block repeats.**
- Malleable: reordering ciphertext results in reordered plaintext.
- Errors in one ciphertext block do not propagate.
- Usage: not recommended to encrypt more than one block of data.

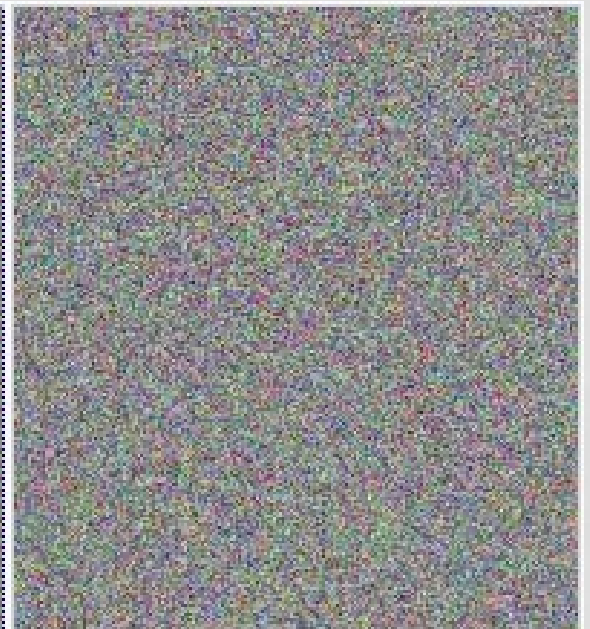
# ECB vs. Other Modes



Original image



Encrypted using ECB mode

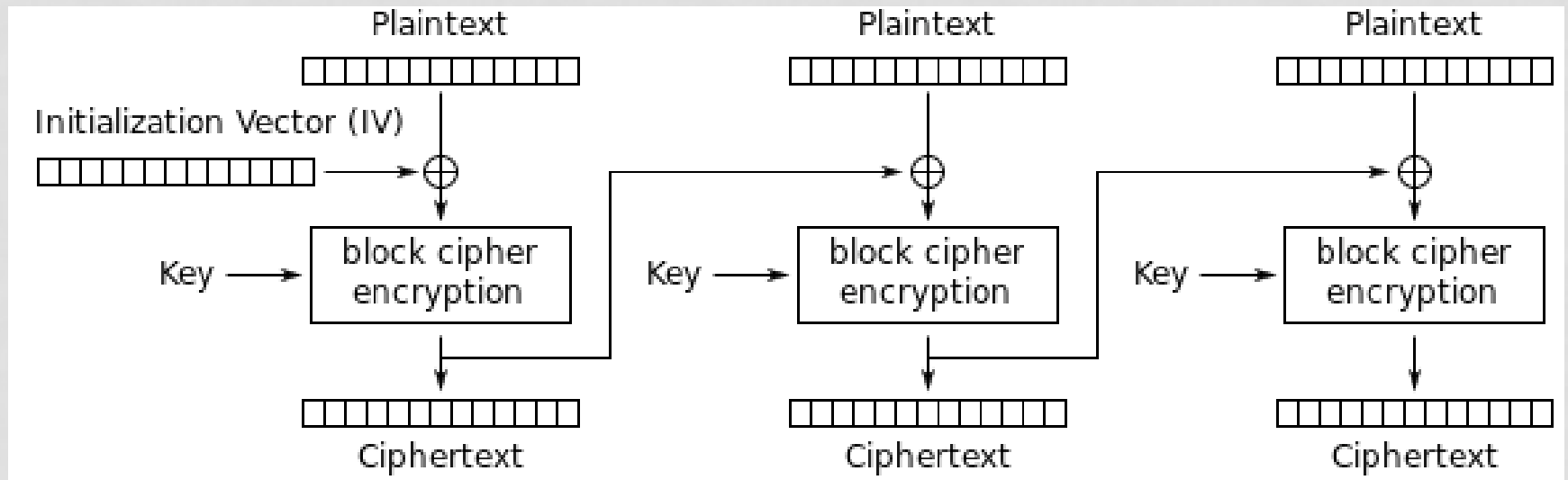


Modes other than ECB result in pseudo-randomness

# Mode 2 – Cipher Block Chaining (CBC)

- Next input depends upon previous output
- **Encryption:**  $C_i = E_k(M_i \oplus C_{i-1})$ , with  $C_0 = IV$
- **Decryption:**  $M_i = C_{i-1} \oplus D_k(C_i)$ , with  $C_0 = IV$

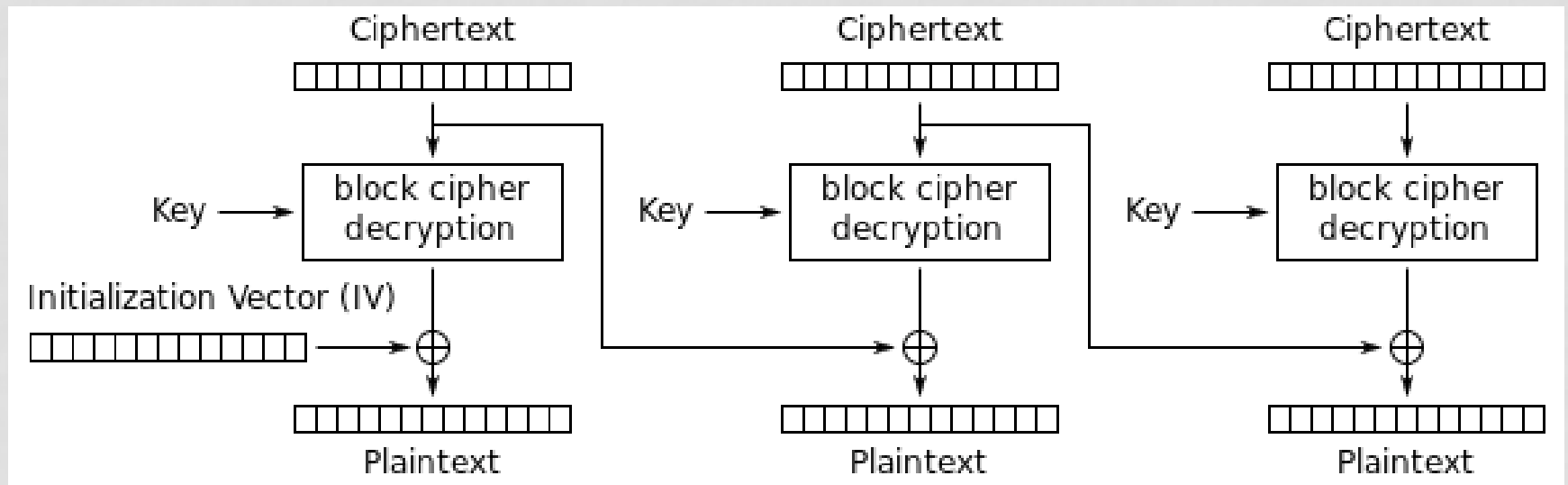
# CBC Encryption



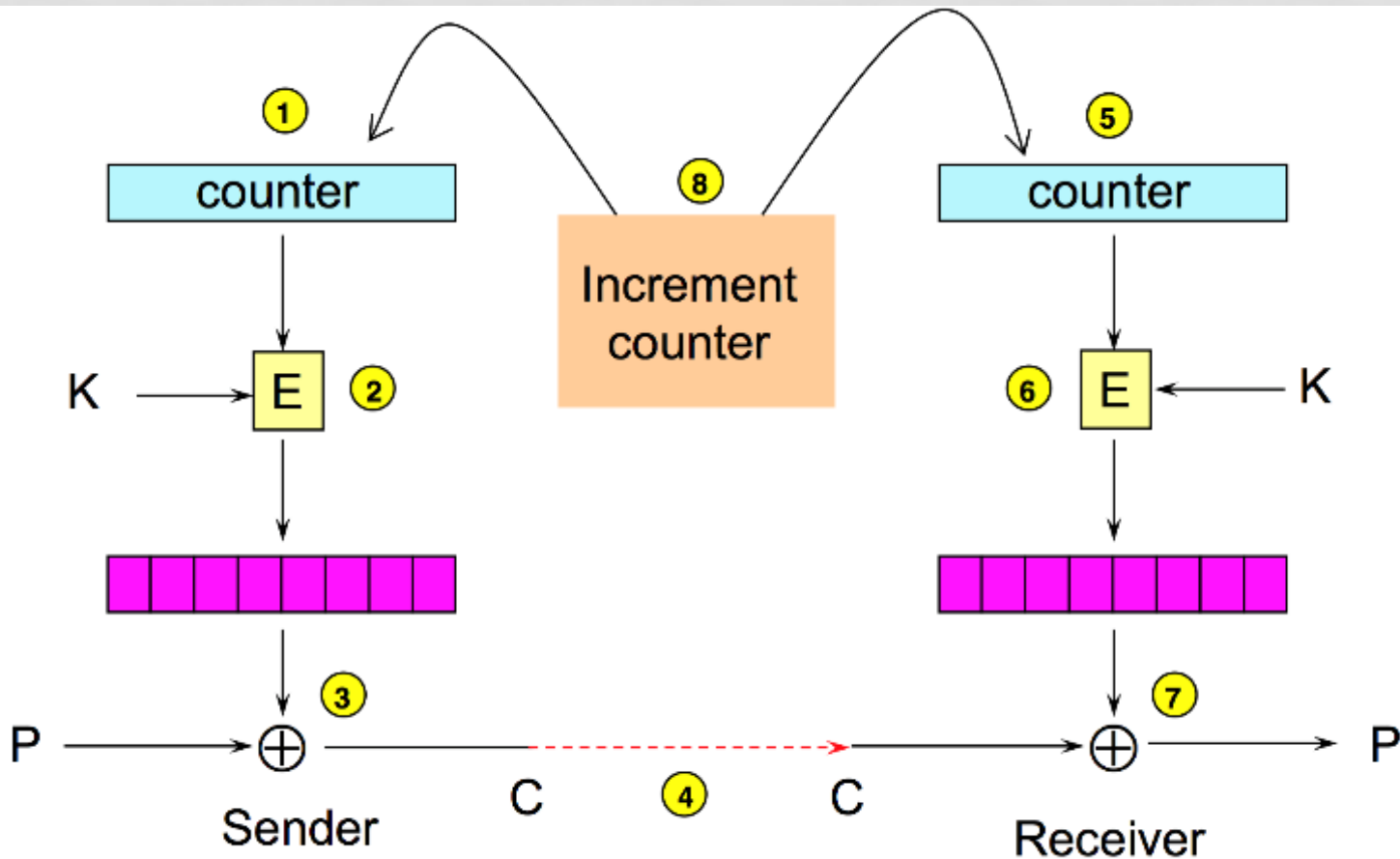
# CBC Properties

- Randomized encryption: repeated text gets mapped to different encrypted data.
- A ciphertext block depends on all preceding plaintext blocks; reorder affects decryption.
- Errors in one block propagate to two blocks
  - one bit error in  $C_j$  affects all bits in  $M_j$  and one bit in  $M_{j+1}$
- Sequential encryption, cannot use parallel hardware.

# CBC Decryption



# Mode 3 – Counter Mode (CTR)





# CTR Mode

- Another way to construct PRNG using a block cipher (e.g. AES):
  - $y_i = E_k[\text{counter}+i]$
- Sender and receiver share: counter (does not need to be secret) and the secret key.

# CTR Properties

- **Software and hardware efficiency**: different blocks can be encrypted in parallel.
- **Preprocessing**: the encryption part can be done offline and when the message is known, just do the XOR.
- **Random access**: decryption of a block can be done in random order, very useful for hard-disk encryption.

# Breaking Cryptosystems

# Attack Scenarios

- Ciphertext-only attack.
- Known-plaintext attack.
- Chosen-plaintext attack.
- Chosen-ciphertext attack.

# Breaking a Cryptosystem

- A cryptosystem is usually broken either by:
  - Finding a way of determining the decryption key
  - Finding a way of determining the plaintext directly
- The term “break” is of course subjective.
- Many cryptosystems are “broken” without “breaking” the encryption algorithm.
- Every algorithm can be broken!