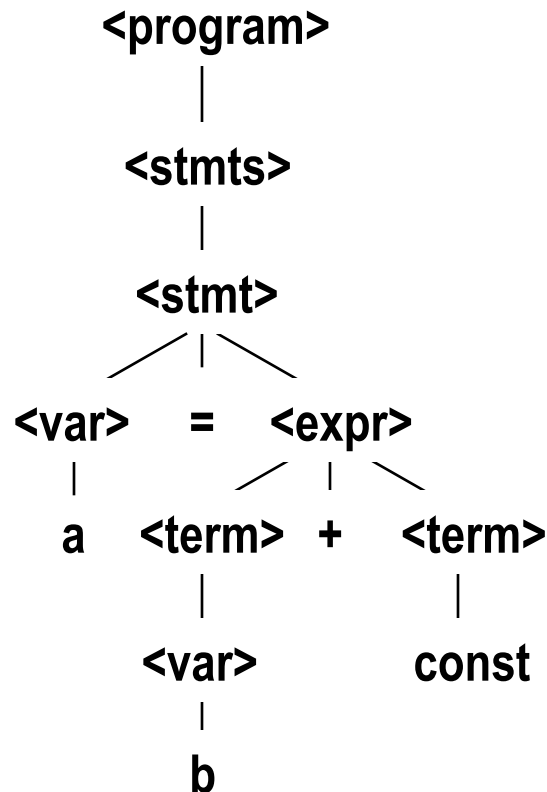


Tutorial-3: part-I

Ambiguity

Parse Tree

- A hierarchical representation of a derivation



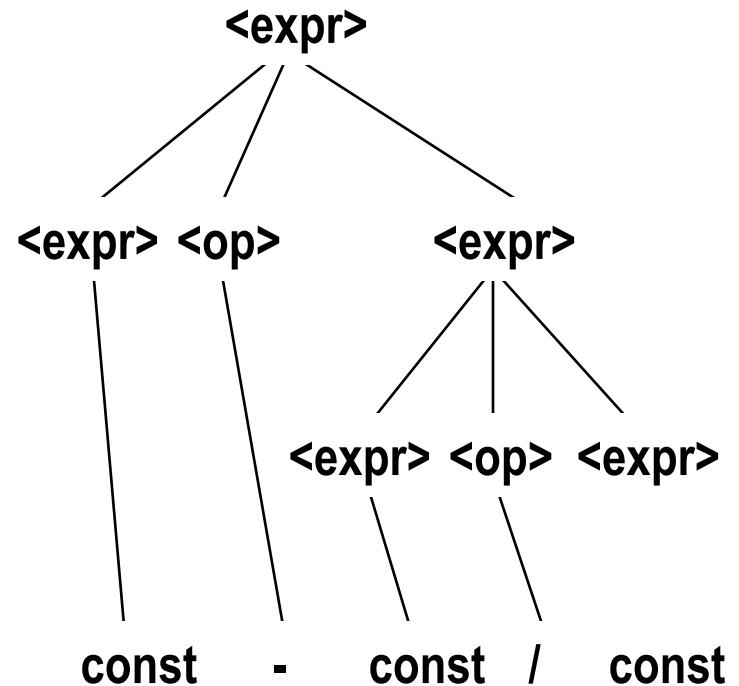
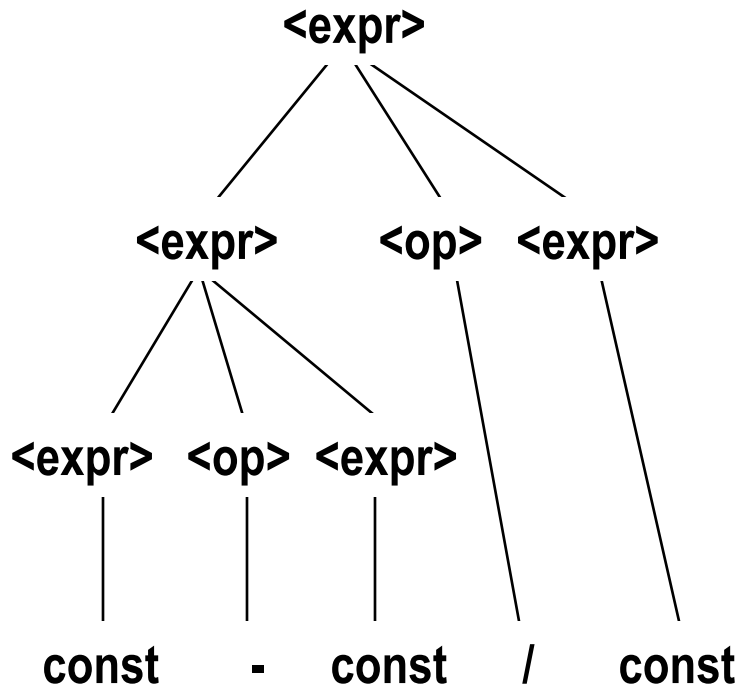
Ambiguity in Grammars

- A grammar is *ambiguous* if and only if it generates a sentential form that has two or more distinct parse trees

An Ambiguous Expression Grammar

$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle \langle \text{op} \rangle \langle \text{expr} \rangle \quad | \quad \text{const}$

$\langle \text{op} \rangle \rightarrow / \quad | \quad -$

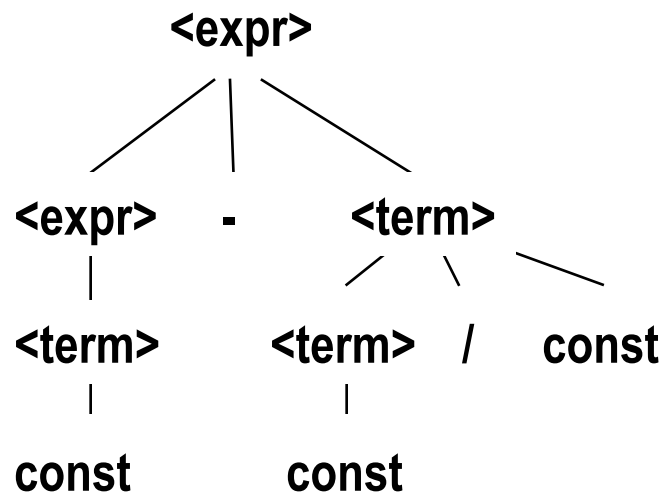


An Unambiguous Expression Grammar

- If we use the parse tree to indicate precedence levels of the operators, we cannot have ambiguity (i.e., it must be eliminated)

$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle - \langle \text{term} \rangle \mid \langle \text{term} \rangle$

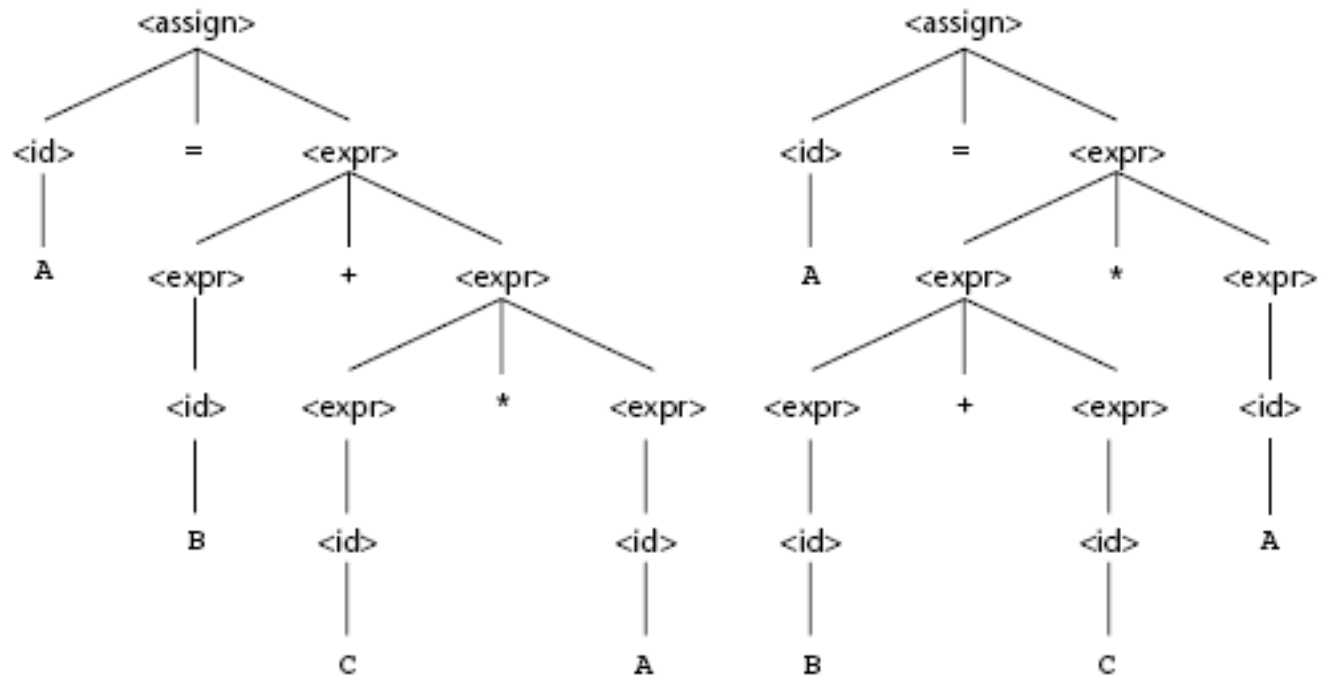
$\langle \text{term} \rangle \rightarrow \langle \text{term} \rangle / \text{const} \mid \text{const}$



EXAMPLE 3.3**An Ambiguous Grammar for Simple Assignment Statements**
$$\begin{aligned}\langle \text{assign} \rangle &\rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle \\ \langle \text{id} \rangle &\rightarrow \mathbf{A} \mid \mathbf{B} \mid \mathbf{C} \\ \langle \text{expr} \rangle &\rightarrow \langle \text{expr} \rangle + \langle \text{expr} \rangle \\ &\quad \mid \langle \text{expr} \rangle * \langle \text{expr} \rangle \\ &\quad \mid (\langle \text{expr} \rangle) \\ &\quad \mid \langle \text{id} \rangle\end{aligned}$$

Figure 3.2

Two distinct parse
trees for the same
sentence,
 $A = B + C * A$



- Notice that the previous grammar allow the parse tree of an expression to grow on both the left and right sides.
- Ambiguity is a problem for compilers because they often base the semantics of these structures (e.g. the precedence of operators) on their syntactic structure.

An un-ambiguous grammar that describes the same language

EXAMPLE 3.2

A Grammar for Simple Assignment Statements

$\langle \text{assign} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$

$\langle \text{id} \rangle \rightarrow A \mid B \mid C$

$\langle \text{expr} \rangle \rightarrow \langle \text{id} \rangle + \langle \text{expr} \rangle$

$\mid \langle \text{id} \rangle * \langle \text{expr} \rangle$

$\mid (\langle \text{expr} \rangle)$

$\mid \langle \text{id} \rangle$

Precedence

- Notice that the above grammar will always put the rightmost operation on the lowest level.
- So in the sentence $A=A*B+C$
- $+$ will be in the lowest level and will have higher precedence than $*$

Associativity

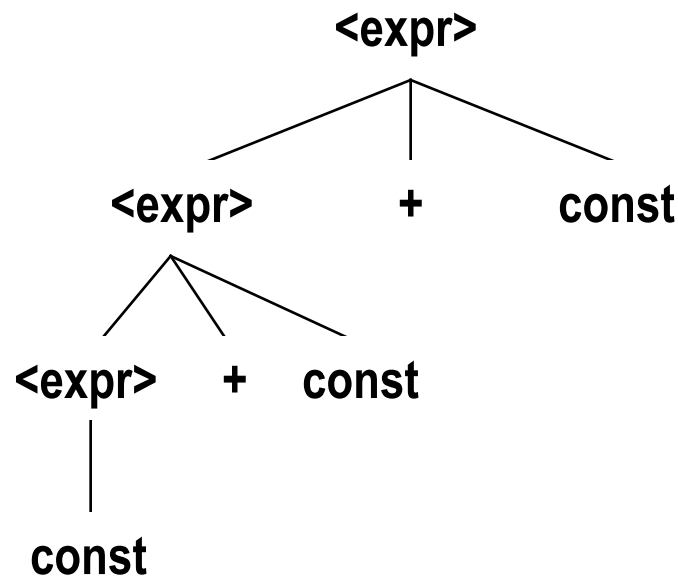
- Also in the sentence
- $A=B+C+D$
- The rightmost + will be at the lowest level
➔ incorrect associativity

Associativity of Operators

- Operator associativity can also be indicated by a grammar

`<expr> -> <const>+<expr>|const` (right associativity)

`<expr> -> <expr> + const | const` (left associativity)



Unambiguous and Correct Associativity (Left to right)

EXAMPLE 3.4

An Unambiguous Grammar for Expressions

```
<assign> → <id> = <expr>
<id> → A | B | C
<expr> → <expr> + <term>
        | <term>
<term> → <term> * <factor>
        | <factor>
<factor> → ( <expr> )
          | <id>
```

EXAMPLE 3.3

An Ambiguous Grammar for Simple Assignment Statements

```
<assign> → <id> = <expr>
<id> → A | B | C
<expr> → <expr> + <expr>
        | <expr> * <expr>
        | ( <expr> )
        | <id>
```