



# **Theory of Computation**

CSC 339 – Spring 2021

## **Chapter-1: part3**

Regular Languages

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**Department of Computer Science**  
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# Outline

- **Recap**
- **Introduction**
- **Regular Expressions**

# Recap

## ➤ Non-deterministic Finite Automata (NFA)

- Given a word w, the automaton can end up in different state each time.
- NFA provides us with an easier way to design finite automata.. But, we can always convert it to equivalent DFA.

# Regular Expressions: Introduction

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➤ In a similar fashion, we can use regular operations to build up expressions describing languages.

➤  $(0 \cup 1)0^*$

# Regular Expressions: Introduction

- **Regular expressions are useful in many text-based applications**
  - **Compilers, search engines, log processing, information retrieval, etc**
- **In arithmetic, we know that  $\times$  has precedence over  $+$**
- **In regular expressions, order of operations:**
  - **Star**
  - **Concatenation**
  - **Union**

# Regular Expressions: Definition 1.52

- Say that  $R$  is a regular expression if  $R$  is
  - $a$  for some  $a$  in the alphabet  $\Sigma$ ,
  - $\epsilon$ ,
  - $\emptyset$ ,
  - $(R1 \cup R2)$ , where  $R1$  and  $R2$  are regular expressions,
  - $(R1 \circ R2)$ , where  $R1$  and  $R2$  are regular expressions, or
  - $(R1^*)$ , where  $R1$  is a regular expression.

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Represent the languages  
 $\{a\}$  and  $\{\epsilon\}$



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Represent the  
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# Regular Expressions: Examples

➤  **$R^*$**

➤ All strings that are zero or more concatenations of strings from  $R$

➤  $R^+ \cup \epsilon = R^*$

➤  **$R^+$**

➤ All strings that are one or more concatenations of strings from  $R$

➤ Shorthand for  $RR^*$

# Regular Expressions: Examples

- $0^*10^* = \{w \mid w \text{ contains a single } 1\}$
- $\Sigma^*1\Sigma^* = \{w \mid w \text{ has at least one } 1\}$
- $\Sigma^*001\Sigma^* = \{w \mid w \text{ contains the string } 001 \text{ as a substring}\}$
- $(\Sigma \Sigma)^* = \{w \mid w \text{ is a string of even length}\}$
- $(\Sigma \Sigma \Sigma)^* = \{w \mid \text{the length of } w \text{ is a multiple of } 3\}$
- $(0 \cup \varepsilon)(1 \cup \varepsilon) = \{\varepsilon, 0, 1, 01\}$

# Regular Expressions: Identities

‣  $R \cup \emptyset = R$

‣  $R \circ \varepsilon = R$

‣ What if we flip the  $\emptyset$  and  $\varepsilon$ ?

# Regular Expressions: Use-cases

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Example for defining the format of numerical constants with a fractional part and/or a sign

$$(+ \cup - \cup \varepsilon) (D^+ \cup D^+ . D^* \cup D^* . D^+)$$

# Regular Expressions: Use-cases

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Example for defining the format of numerical constants with a fractional part and/or a sign

$$(+ \cup - \cup \varepsilon) (D^+ \cup D^+ . D^* \cup D^* . D^+)$$

Examples of accepted expressions: **72**, **3.14159**, **+7.**, **-0.1**



# Regular Expressions: Use-cases

- **Regular expressions have numerous text-based applications in computer science.**
  - **Compilers for programming languages**
  - **Log analysis and debugging large systems**
  - **Natural language processing (NLP)**

# Regular Expressions: equivalence w/ finite automata

- **Regular expressions are equivalent with finite automata.**
- **Any regular expression can be converted into a finite automaton that recognizes the language it describes.**

## Regular Expressions: theorem 1.54

*“A language is regular if and only if some regular expression describes it.”*

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‣ **Two-direction theorem (iff)**

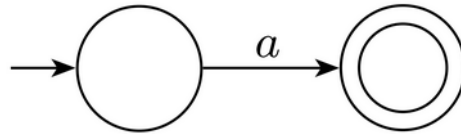
‣ **First direction:**

‣ **If we can convert a regular expression  $R$  into NFA, then we can prove that the language  $R$  describes is regular.**

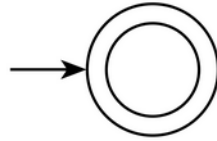
‣ **Back to definition of a regular expression.**

## Regular Expressions: theorem 1.54

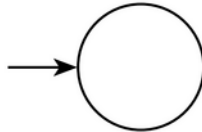
›  $R = a$  for some  $a \in \Sigma$ . Then  $L(R) = \{a\}$  and the following NFA recognizes  $L(R)$ .



›  $R = \varepsilon$ . Then  $L(R) = \{\varepsilon\}$



›  $R = \Phi$ . Then  $L(R) = \{\Phi\}$

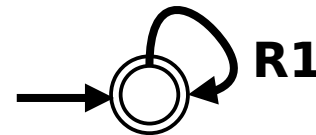
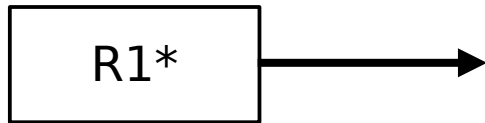
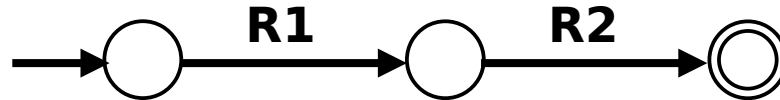
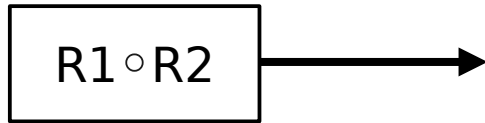
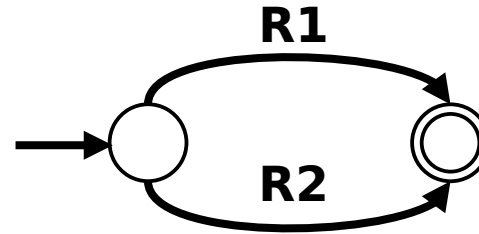
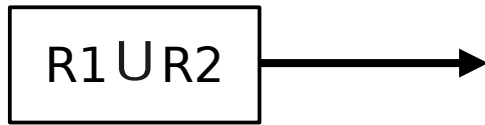


## Regular Expressions: theorem 1.54

- $(R1 \cup R2)$ , where  $R1$  and  $R2$  are regular expressions,
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  - $(R1^*)$ , where  $R1$  is a regular expression.
- For these three cases, we can use constructions given in previous proofs that the class of regular languages is closed under the regular operations.

# Regular Expressions: Regex to NFA

## Three main operations



# Regular Expressions: Regex to NFA - Examples

## › Three main operations

**1.  $ab^*a$**

**2.  $(a \cup b)a$**

**3.  $a(bb)^*$**

**4.  $(ba)^*$**

**5.  $(ba)^+$**

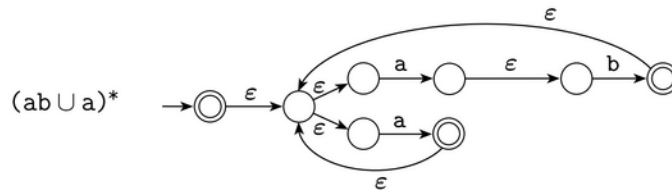
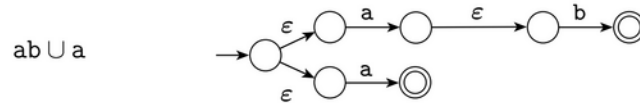
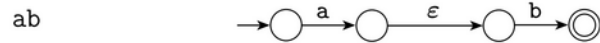


## Regular Expressions: Example 1.56

› Let's convert the regular expression  $(ab \cup a)^*$  into an NFA.

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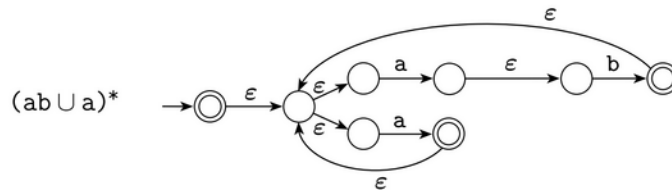
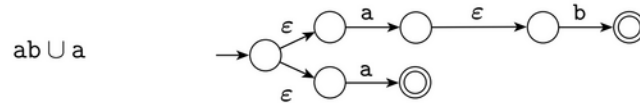
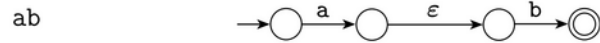


# Regular Expressions: Example 1.56

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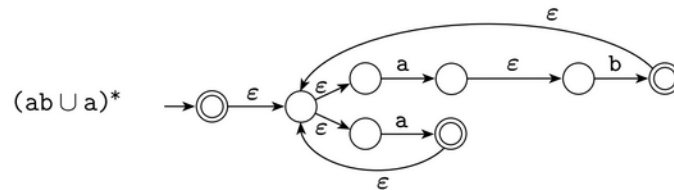
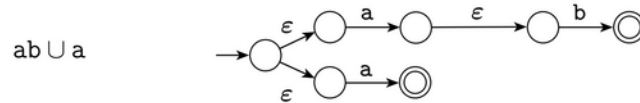
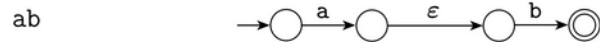


Can we simplify  
this NFA?



# Regular Expressions: Example 1.56

► Let's convert the regular expression  $(ab \cup a)^*$  into an NFA.



**Also see example 1.58**

# Regular Expressions: GNFA

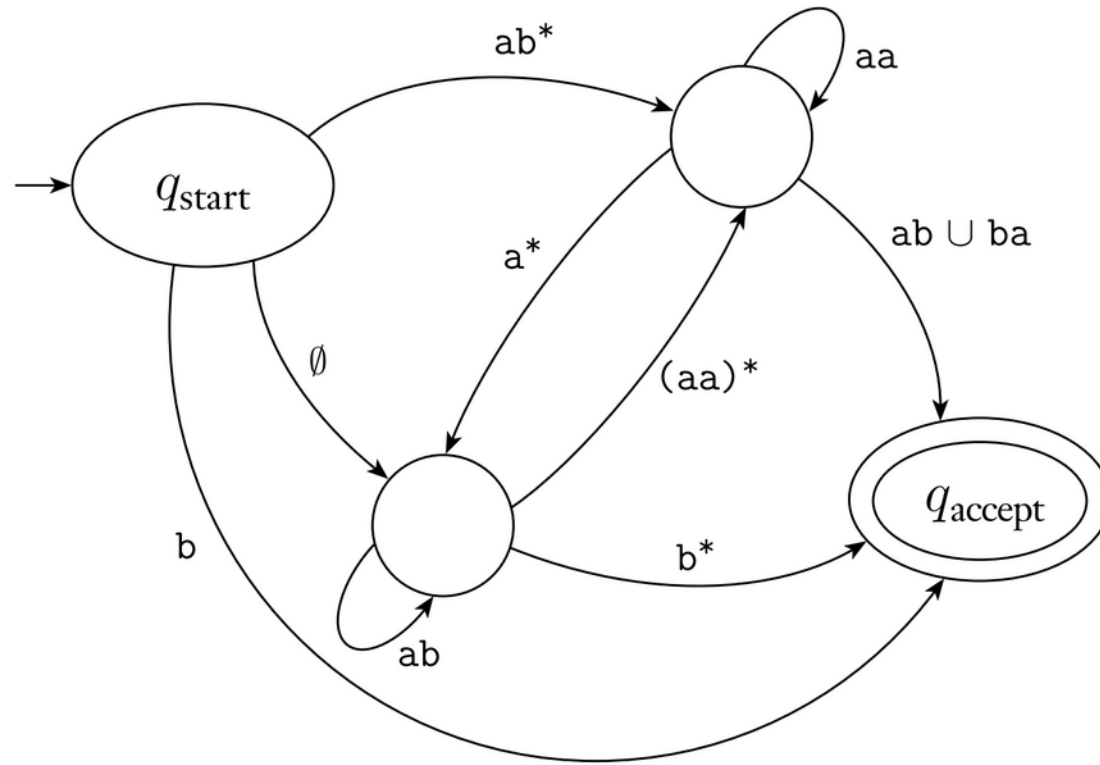
- **The other direction of theorem 1.54**
  - *“If a language is regular, then it is described by a regular expression”*
- **Converting an NFA into a regular expression**
- **Generalized nondeterministic finite automata (GNFA)**
  - **Transition arrows may have regular expressions as labels**
  - **GNFA reads blocks of symbols from the input** – as opposed to reading symbol by symbol

# Regular Expressions: GNFA

## ➤ Special conditions

- **Start state has outgoing transition arrows to every other state (but, no incoming arrows).**
- **Only single accept state that has arrows from every other state.**
- **Except for start and accept states, one arrow goes from every state to every other state, and from each state to itself.**

# Regular Expressions: GNFA



# Regular Expressions: GNFA

## ➤ Converting DFA into GNFA

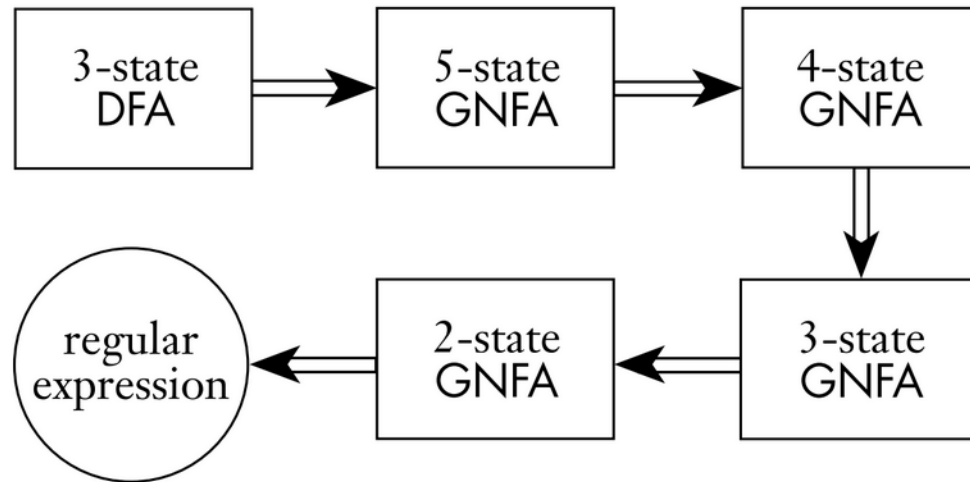
- Add new start state with an  $\epsilon$  arrow to the old start state.
- Add new accept state with  $\epsilon$  arrows from old accept states.



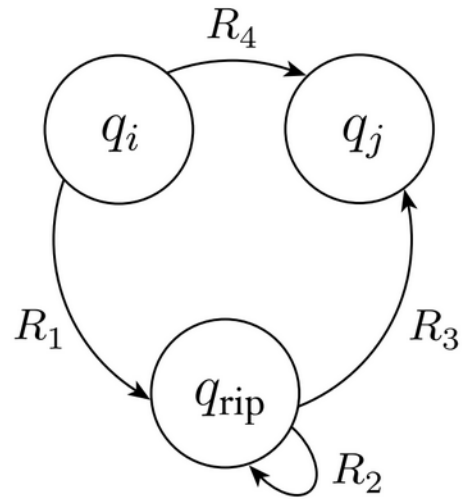
# Regular Expressions: DFA to Regex

- **We will follow a state elimination technique**
- **Convert the DFA into a GNFA, and then, start eliminating states by converting them into corresponding regular expressions.**

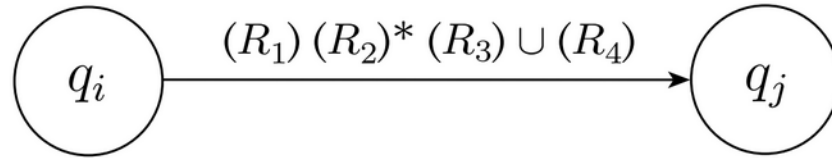
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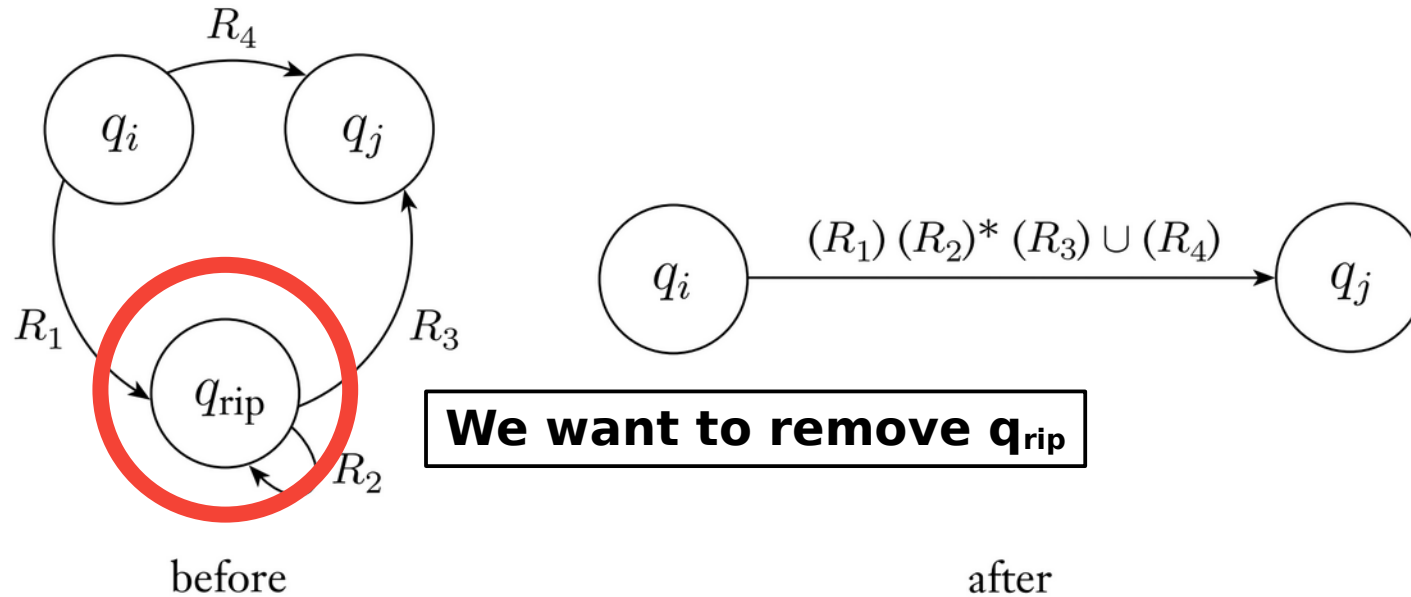


before

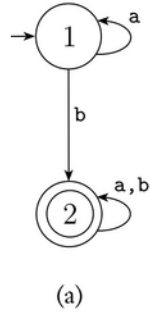


after

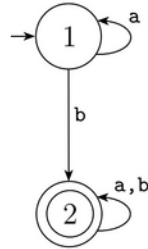
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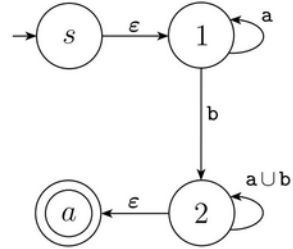
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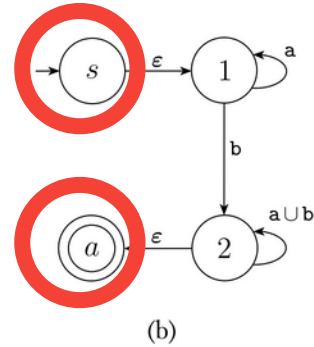
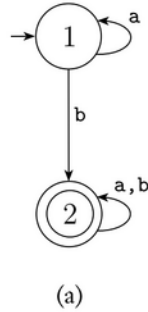
(a)



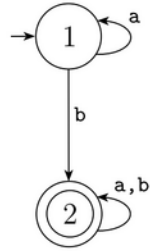
(b)

# Regular Expressions: GNFA

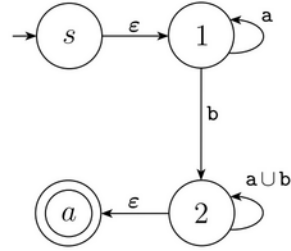
**Add new start and accept states**



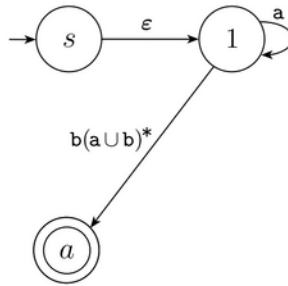
# Regular Expressions: GNFA



(a)



(b)

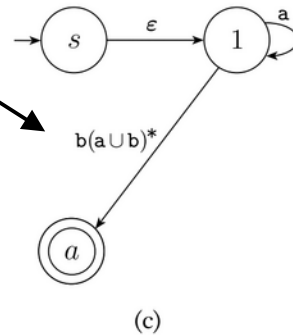
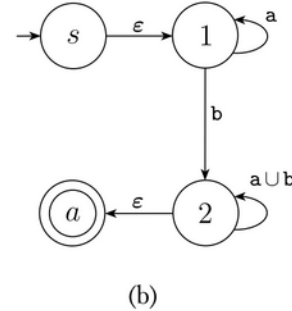
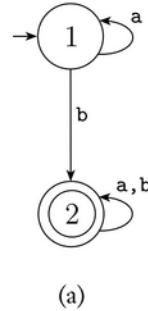


(c)

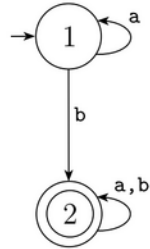


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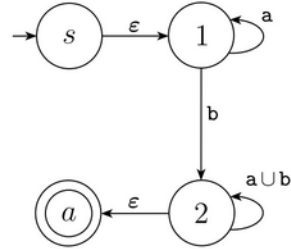
**We've eliminated state(2)  
and replaced it with  
a regular expression that  
describes its transitions**



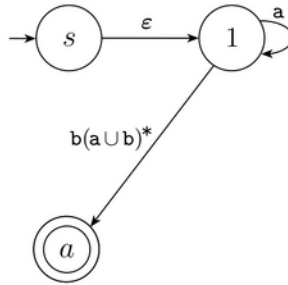
# Regular Expressions: GNFA



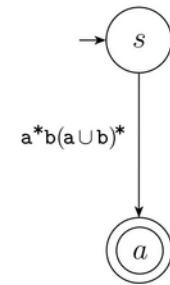
(a)



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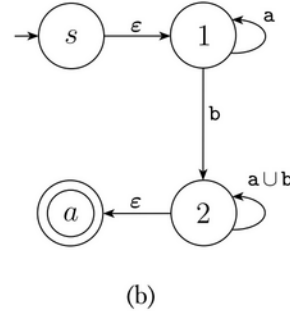
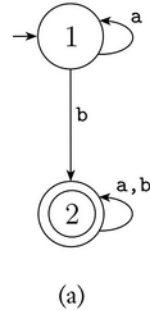


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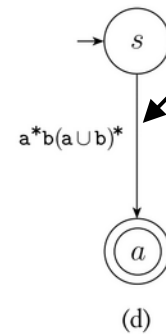
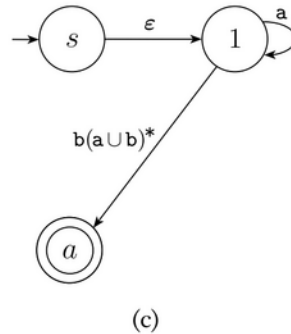


(d)

# Regular Expressions: GNFA

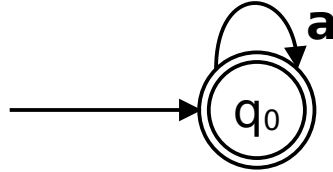


**Now, we eliminate state (1) and replace it with a regex that describes its transition**



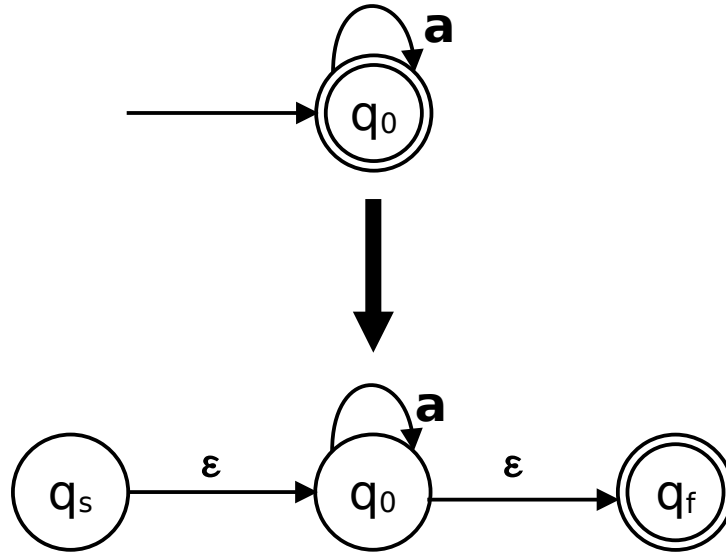
# Regular Expressions: DFA to Regex

- › Convert the following DFA into regular expressions
- › We want to reserve the same language.



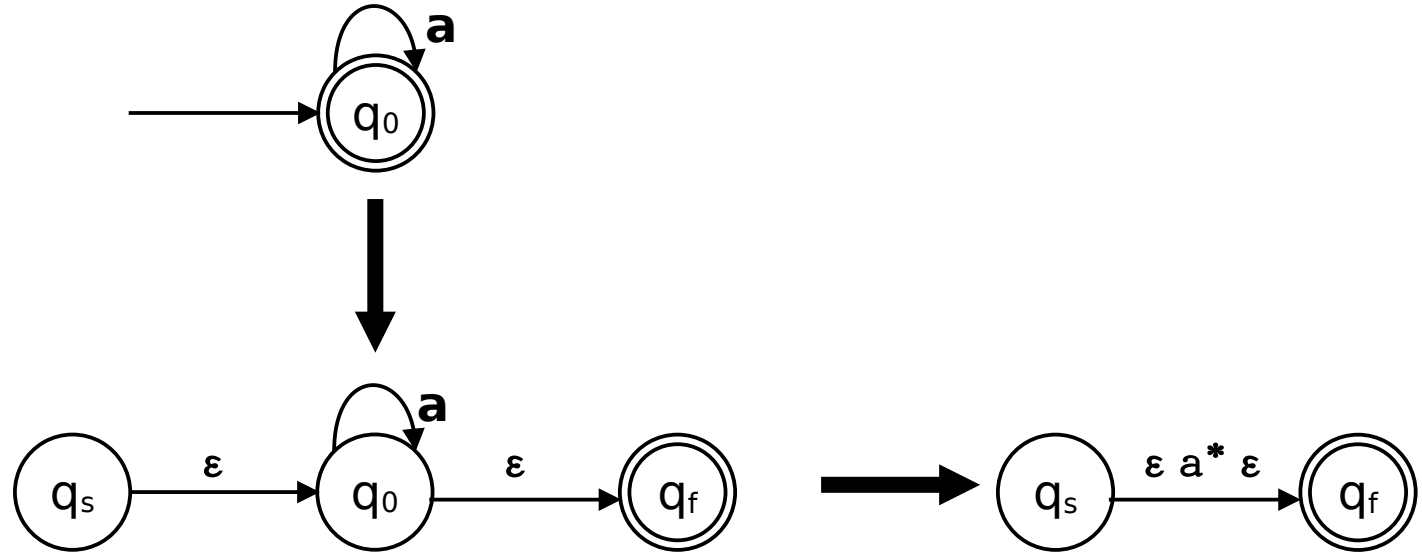
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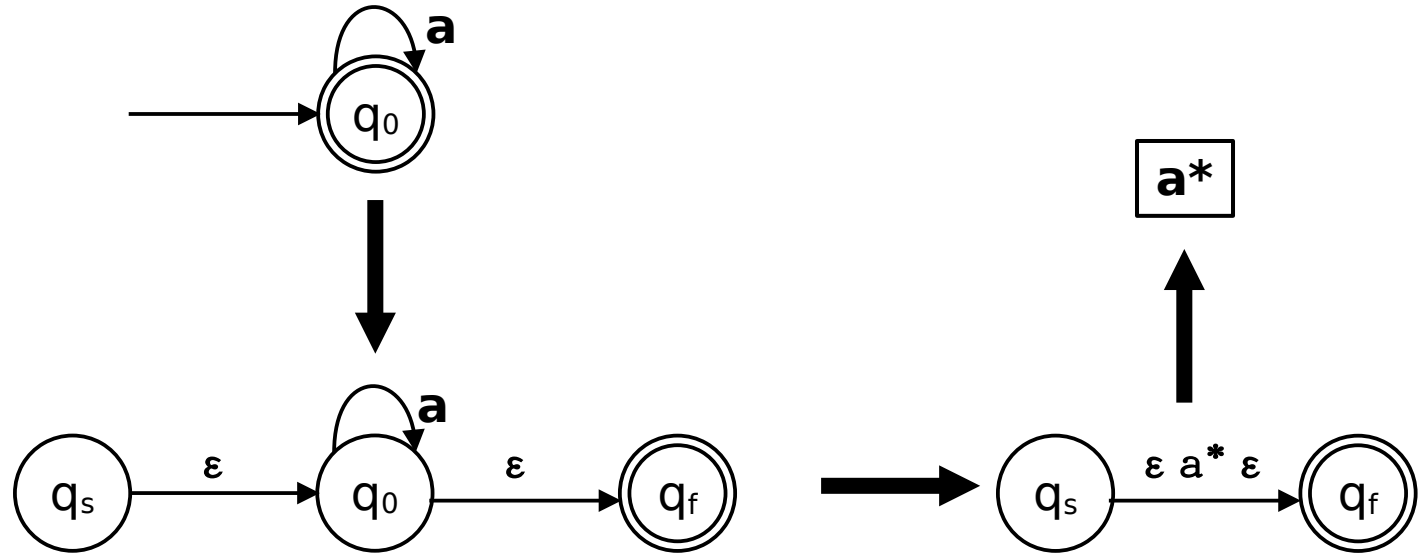
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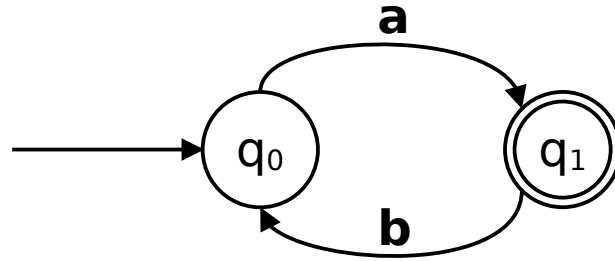
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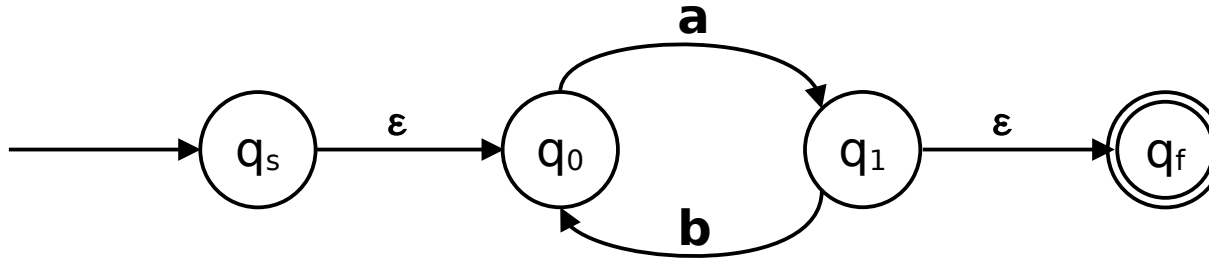
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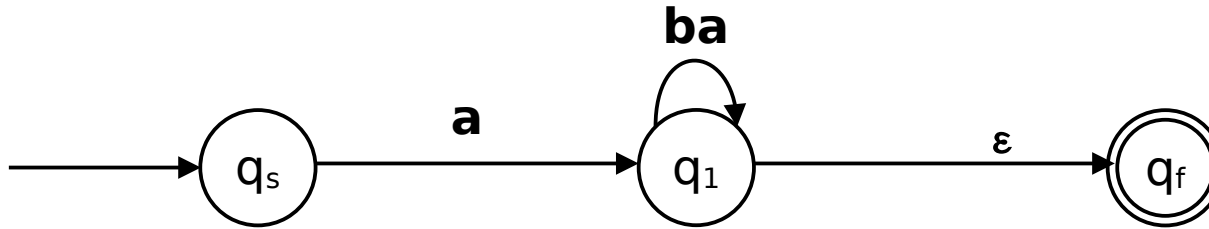
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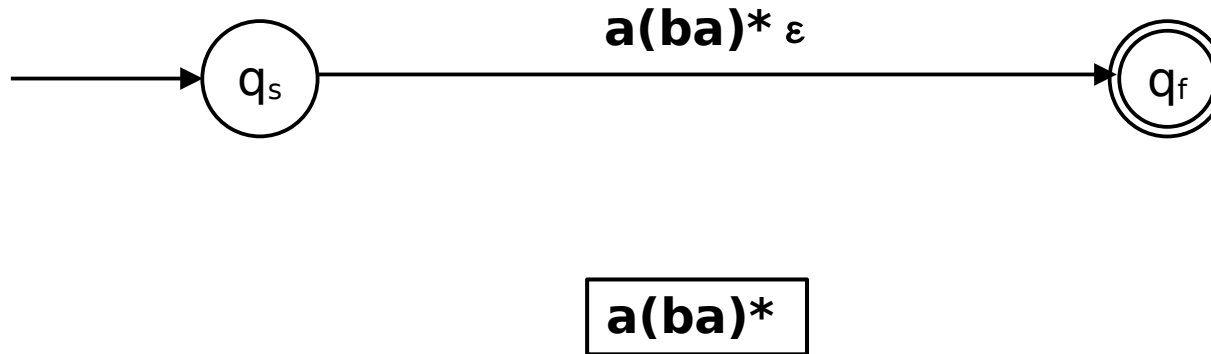
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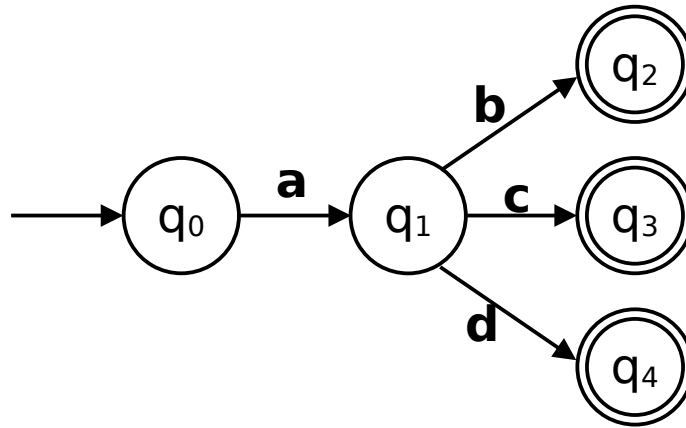
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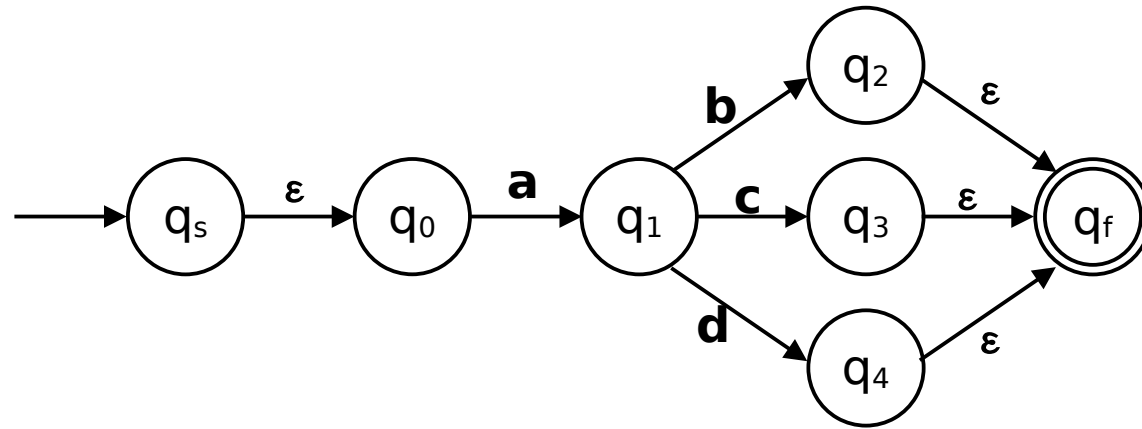
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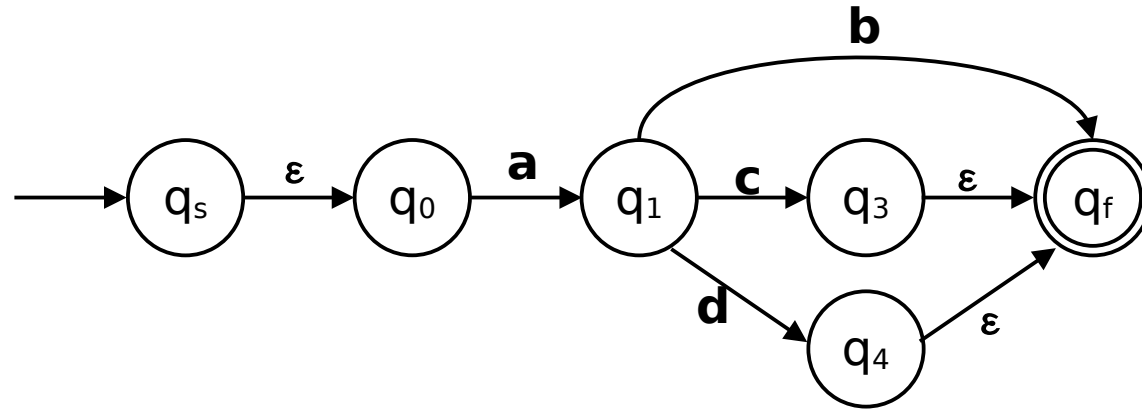
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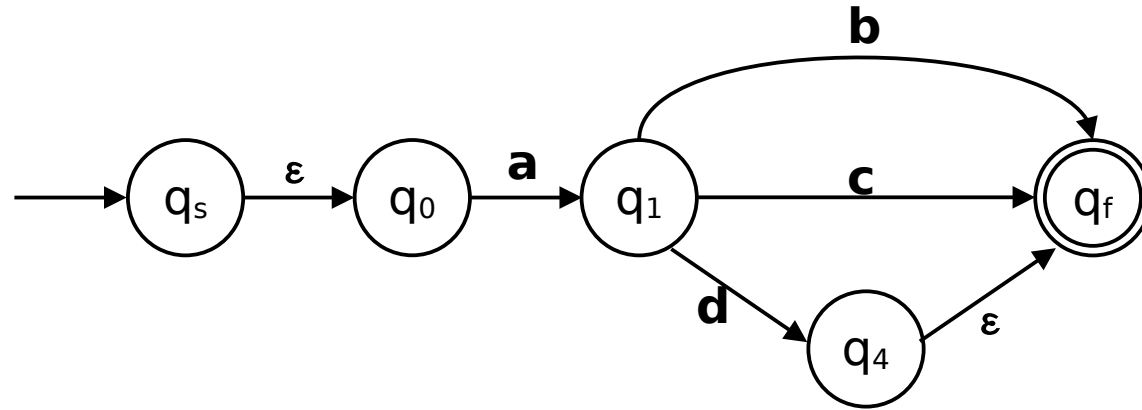
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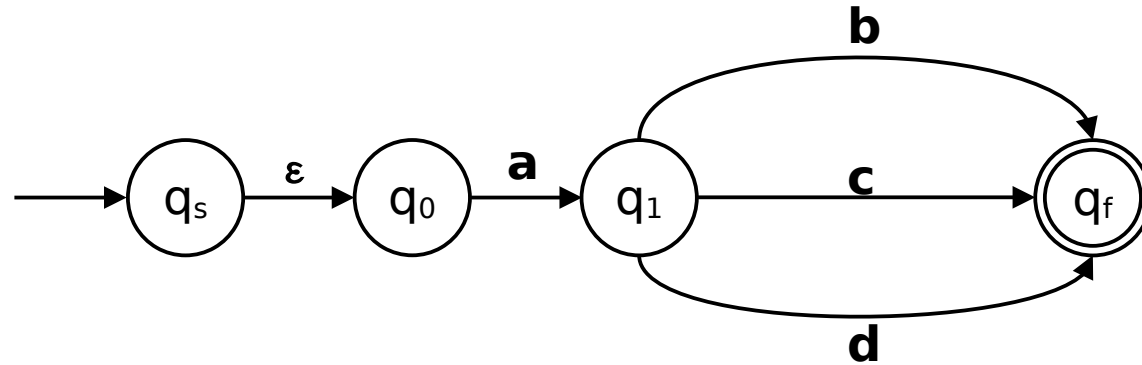
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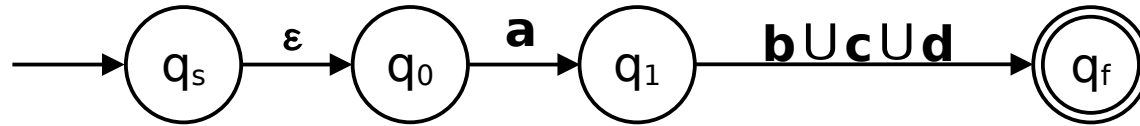
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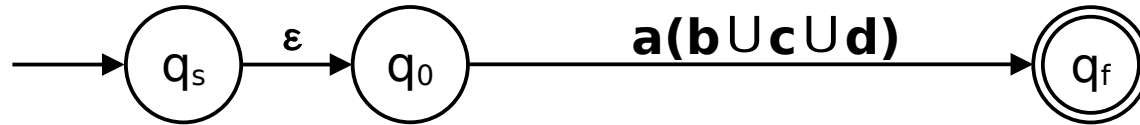
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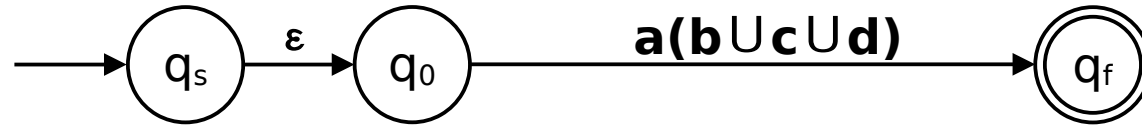
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# Regular Expressions: DFA to Regex

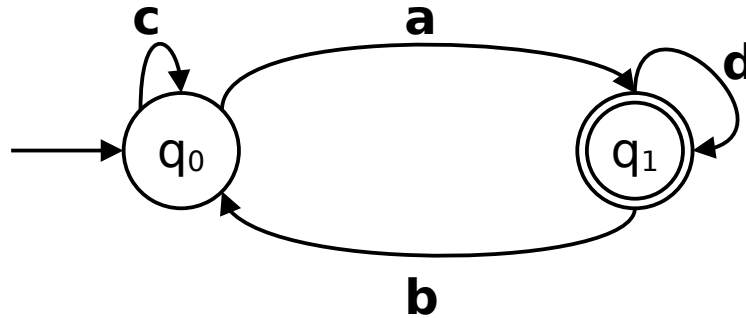
- Convert the following DFA into regular expressions
- We want to reserve the same language.



**$a(b \cup c \cup d)$**

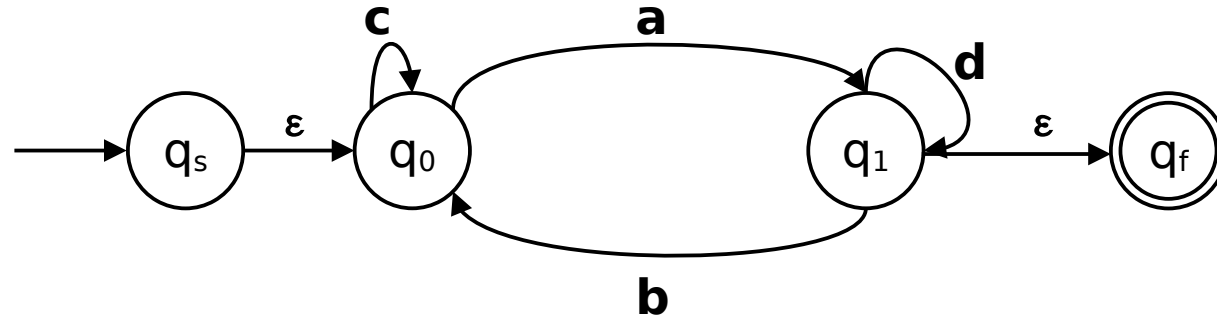
# Regular Expressions: DFA to Regex

- › Convert the following DFA into regular expressions
- › We want to reserve the same language.



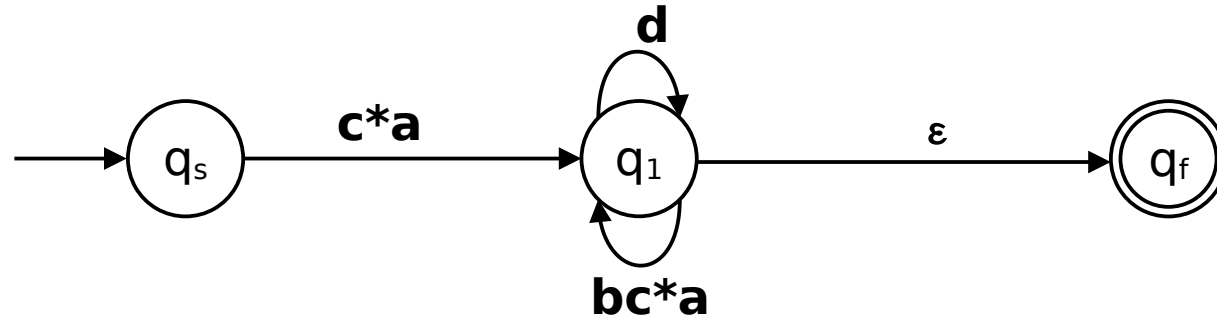
# Regular Expressions: DFA to Regex

- Convert the following DFA into regular expressions
- We want to reserve the same language.



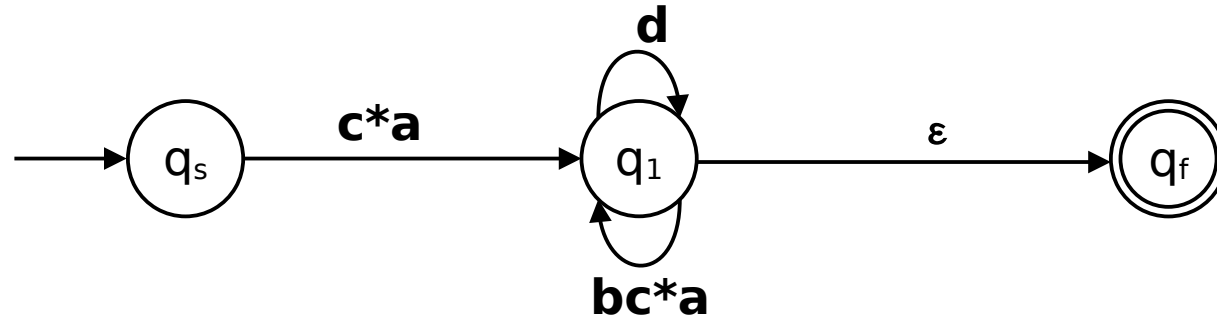
# Regular Expressions: DFA to Regex

- Convert the following DFA into regular expressions
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# Regular Expressions: DFA to Regex

- Convert the following DFA into regular expressions
- We want to reserve the same language.



$c^*a(d \cup bc^*a)^*$

# Regular Expressions: GNFA

Read example **1.68**



# Homework

## ➤ Reading

### ➤ 1.4