

Theory of Computation

CSC 339 – Spring 2021

Chapter-7: part1

Time Complexity

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Introduction

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- **What's the difference between computability and complexity theories?**

Complexity Theory

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- **In complexity theory, we would like to take a closer look at solvable problems and see how difficult the solution would be.**
- **In this part, we will study time complexity.**

Complexity Theory: Time Complexity

‣ **How much time does a given TM take to process some input?**

Complexity Theory: Time Complexity

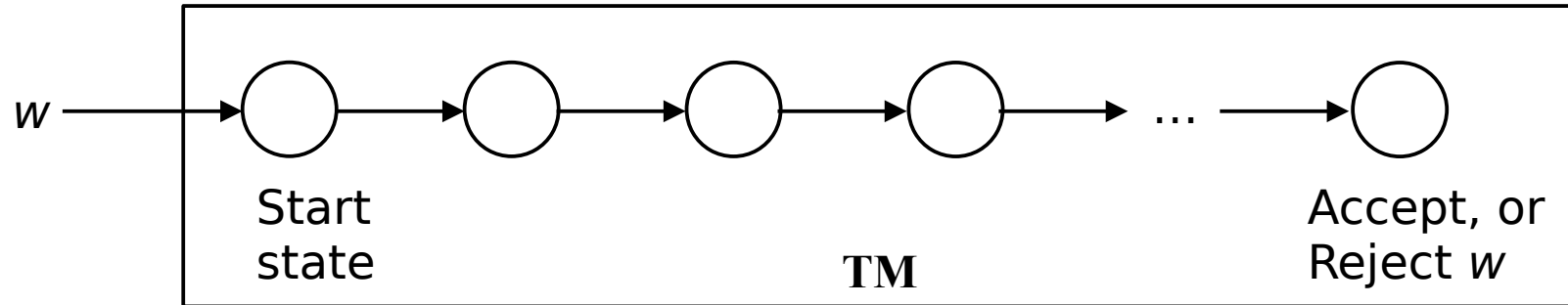
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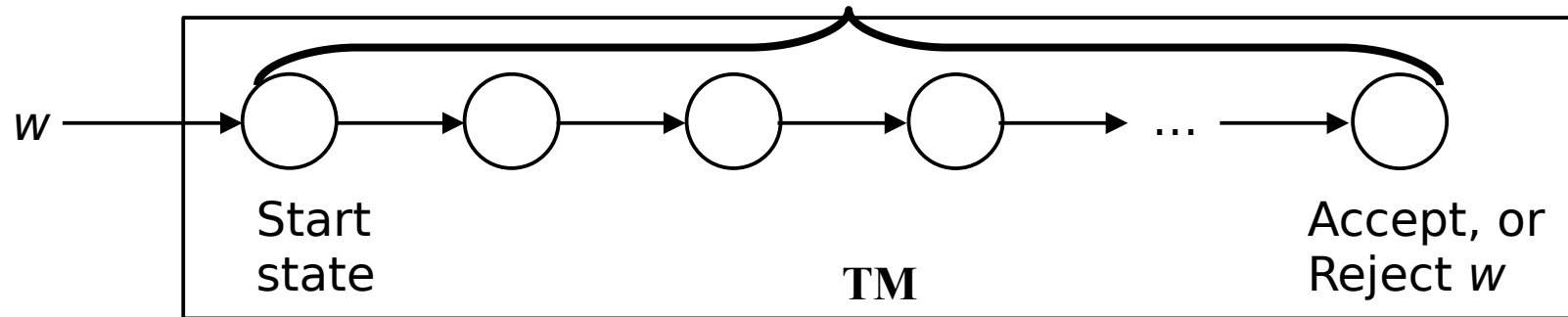
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Running Time = number of transitions

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 - aaabbb → may take n number of steps
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- We compute the running time of an algorithm (or a TM) as a function of the length of the string representing the input.

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- **Worst-case analysis**

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- **Average-case analysis**

- **We consider the average running time of all inputs of a particular length.**

Complexity Theory: Time Complexity

Definition 7.1

Let M be a deterministic Turing machine that halts on all inputs. The *running time* or *time complexity* of M is the function $f: N \rightarrow N$, where $f(n)$ is the maximum number of steps that M uses on any input of length n . If $f(n)$ is the running time of M , we say that M runs in time $f(n)$ and that M is an $f(n)$ time Turing machine. Customarily, we use n to represent the length of the input.

Complexity Theory: Asymptotic Analysis

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- Asymptotic notation is also known as big-O notation.

Complexity Theory: Asymptotic Analysis

Definition 7.2

Let f and g be functions $f, g : N \rightarrow R^+$. Say that $f(n) = O(g(n))$ if positive integers c and n_0 exist such that for every integer $n \geq n_0$,

$$f(n) \leq cg(n).$$

When $f(n) = O(g(n))$, we say that $g(n)$ is an upper bound for $f(n)$, or more precisely, that $g(n)$ is an asymptotic upper bound for $f(n)$, to emphasize that we are suppressing constant factors.

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Intuitively, $f(n) = O(g(n))$ means that f is less than or equal to g if we disregard differences up to a constant factor.

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Is $f(n) = O(n^4)$?

Complexity Theory: Asymptotic Analysis

› Examples

› $\sqrt{n} = O(n)$.

› $n = O(n \log \log n)$.

› $n \log \log n = O(n \log n)$.

› $n \log n = O(n^2)$.

› $n^2 = O(n^3)$.