

② Suppose  $n$  activities apply for using a common resource. Activity  $a_i$  ( $1 \leq i \leq n$ ) has a starting time  $S[i]$  and a finish time  $F[i]$  such that  $0 < S[i] < F[i]$ . Two activities  $a_i$  and  $a_j$  ( $1 \leq i, j \leq n$ ) are compatible if intervals  $[S[i], F[i])$  and  $[S[j], F[j])$  do not overlap. We assume the activities have been sorted such that  $S[1] \leq S[2] \leq \dots \leq S[n]$ .

A Design an  $\mathcal{O}(n^2)$  dynamic programming algorithm to find a set of compatible activities such that the total amount of time the resource is used by these compatible activities is maximized. You need to define the sub-problems, establish the inductive formula and show the initial conditions. Pseudocode is not required.

B Apply your algorithm to the following set of activities:

i	1	2	3	4	5	6	7	8	9	10	11
S[i]	2	3	5	6	7	9	10	12	13	14	16
F[i]	6	5	7	10	8	13	16	14	14	18	20

#### Part A.

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**Algorithm 1** A dynamic programming algorithm usable to solve the activity problem above in  $\mathcal{O}(n^2)$  time. In this algorithm...

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1: function MAXACTIVITIES( $S[1..n]$ ,  $F[1..n]$ )
2:    $M \leftarrow P \leftarrow \emptyset$ 
3:   INITIALIZE( $M, P$ )
4:
5:   for  $i$  from 1 to  $n$  do                                      $\triangleright$  Locate max for  $m_i$ 
6:      $max \leftarrow M[i]$ 
7:     for  $j$  from 1 to  $i$  do
8:       if  $max < M[j]$  and  $F[j] < S[i]$  then
9:          $max \leftarrow M[j]$ 
10:         $P[i] \leftarrow j$ 
11:      end if
12:    end for
13:  end for
14:
15:   $max \leftarrow 1$                                               $\triangleright$  Find global maximum
16:  for  $i$  from 2 to  $n$  do
17:    if  $M[max] < M[i]$  then
18:       $max \leftarrow i$ 
19:    end if
20:  end for
21:
22:  return ( $max, P$ )
23: end function

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With this algorithm in mind, the answer is as follows:

**Inductive Formula:**  $\{m(i) = \max_{1 \leq j \leq i} (m_j) + 1 \mid F[j] < S[i]\}$

**Initial Conditions:**  $m(1) = 1$

The **subproblem** can be thought of as follows. Each activity,  $i$ , will be added to a chain of activities resulting in the maximum set that includes activity  $i$ . Activity  $i$  is always included in its maximum. Its maximum,  $m_i$ , must also be based off of previous activities which are compatible.

**Part B.**