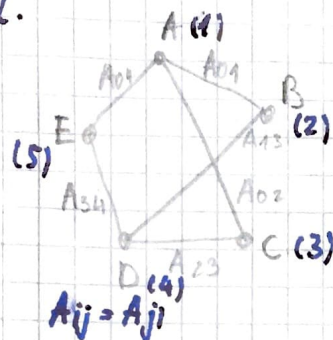


Problem 1-1 Adjacency Matrix

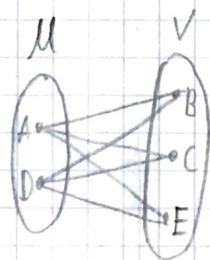
$$A_G = \begin{pmatrix} 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 \end{pmatrix}$$

1. G is undirected, as A_G is symmetric (i.e. $(A_G)_{ij} = (A_G)_{ji}$)
 $\forall i, j = 0, \dots, 5$

2.



3. Yes, we can divide the nodes N into two disjoint sets $U = \{A, D\}$ and $V = \{B, C, E\}$ so that every link connects only a node in U to a node in V (but not to a node within the same set). This means that U and V are independent sets. An illustration can be seen here:



4. adjacency list:

node	linked to
A	B, C, E
B	A, D
C	A, D
D	B, C, E
E	A, D

edge list

pair of edges
(A, B)
(A, C)
(A, E)
(B, D)
(C, D)
(D, E)

Problem 1-2 Average Degree of a Growing Network

1. total number of nodes N :

$$N(t) = t \rightarrow N(t=T) = \underline{T} \quad (\text{with } N(1)=1, N(2)=2, \dots)$$

2. total number of links L :

$$L(t) = t-1 \rightarrow L(t=T) = \underline{T-1} \quad (\text{with } N(1)=0, N(2)=1, \dots)$$

3. average degree $\langle k \rangle$:

$$\langle k \rangle(t) = \frac{2L}{N} = \frac{2(t-1)}{t} \rightarrow \langle k \rangle(T) = \frac{2(T-1)}{T} \quad (\text{with } \langle k \rangle(1)=0, \langle k \rangle(2)=1, \dots)$$

4. average limit degree in the limit $T \rightarrow \infty$

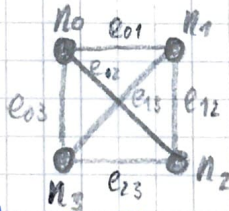
$$\lim_{T \rightarrow \infty} \langle k \rangle(T) = \lim_{T \rightarrow \infty} \frac{2(T-1)}{T} = \lim_{T \rightarrow \infty} 2\left(1 - \frac{1}{T}\right) = \underline{\underline{2}}$$

Problem 1-3 Difficulty of an Exhaustive Search

Number of walks that need to be checked for an Eulerian trail for:

1. complete graph with $N=4$ nodes:

sketch:



$$\begin{aligned} \text{average degree } \langle k \rangle &= N-1 \\ \text{total number of links } L &= 6 = \frac{\langle k \rangle \cdot N}{2} \end{aligned} \quad (k = \langle k \rangle \text{ for nodes } n_0, \dots, n_3)$$

Procedure for checking for Eulerian trails as described in the problem description (exhaustive search):

(a) start at any of the nodes n_i (4 possibilities)

(b) traverse an edge $e_{i' i} = (n_{i'}, n_i)$ ($i \neq i'$)

($k = \langle k \rangle = 3$ possibilities)

(c) repeat (b) $L=L$ times so that a walk of length $l=L$ is performed ($L=6$ times)

$$\Rightarrow \# \text{ walks checked} = N \cdot \langle k \rangle^L = 4 \cdot 3^6 = \underline{\underline{2916}}$$
$$= N \cdot k^L$$