

# An Optimum Closed Form ALOHA Frame Length for Multiple Collision Recovery Coefficients

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**Abstract**—Minimizing the reading time of large tag populations is a critical issue in RFID systems. The usual approach to reduce the reading time is to select the frame size attaining the highest throughput per frame. Previous studies have focused on the frame length calculations using the conventional Framed Slotted ALOHA (FSA) algorithms. In such systems, only the answer of a single tag is considered as a successful slot, and if multiple tags respond simultaneously, a collision occurs. Then all tag replies are discarded. However, modern systems have the capability of recovering these collisions and converting collided slots into successful slots. Recent studies focused on calculating the optimal frame length taking into consideration the collision recovery probability. However, these studies have assumed a constant collision recovery probability coefficient, i.e. the probability to recover one tag from  $i$  collided tags per slot is constant, regardless the number of collided tags  $i$ . In this work we propose a novel closed form solution for the optimal FSA frame length which considers the differences in the collision recovery probabilities. The values of the collision recovery coefficients are extracted from the physical layer parameters. Timing comparisons are presented in simulation results to show the mean reduction in reading time using the proposed frame length compared to the other proposals.

## I. INTRODUCTION

Radio Frequency Identification (RFID) is an identification technology that wirelessly transmits the identity of a tag that is attached to an object or a person. Recently the number of applications that use RFID technology has increased, and the reading time is one of the most critical issues in applications with many tags. In RFID systems, the tags typically share a common communication channel. Thus, there is a certain probability of tag collisions, i.e. multiple tags answer simultaneously. This collision probability naturally increases in dense networks. According to EPCglobal class 1 gen 2 standards [1], the reader is responsible of coordinating the network and has to avoid tags collisions using specific anti-collision algorithms. The conventional anti-collision algorithm is the Framed Slotted ALOHA (FSA) algorithm [2]. Using this algorithm, only the single tag reply (successful slot) can be decoded and identified. Therefore, the conventional definition of the expected reading efficiency  $\eta_{conv}$  is equivalent to the probability of success  $P(1)$ , which is expressed as:

$$\eta_{conv} = P(1) = \frac{n}{L} \left(1 - \frac{1}{L}\right)^{n-1}, \quad (1)$$

where  $n$  presents the number of tags in the reading area, and  $L$  is the frame length.

The main goal is to find the optimal frame length  $L$ , which maximizes the reading efficiency  $\eta_{conv}$ . Based on (1), the reading efficiency  $\eta_{conv}$  is maximized to  $\eta_{conv(max)} = 36\%$  when  $L = n$  as shown in [2].

Some research groups (e.g. [3], [4]) have concentrated on the effect of the collision resolving probability and the optimum value of the FSA frame length. They have proposed the reading efficiency equation:

$$\eta_{[3]} = P(1) + \alpha \cdot \sum_{i=2}^n P(i), \quad (2)$$

where  $\sum_{i=2}^n P(i)$  is the probability of collision,  $\alpha$  is the average collision resolving probability coefficient. In this efficiency equation, the RFID reader can convert part of the collided slots into successful slots. Based on this efficiency, the authors have calculated a closed form solution for the optimum frame length:

$$L_{[3]} = \alpha + (1 - \alpha) \cdot n. \quad (3)$$

The authors here have assumed unlimited and equal collision resolving probabilities coefficients. For example, the probability to resolve two collided tags is identical to the probability to resolve ten collided tags. However, this is not an accurate assumption, as it is clear that the collision recovery probabilities decrease when the number of collided tags increase. Moreover, there is a practical saturation limit for the collision recovery capability.

Another research group [5] considered the limited RFID reader capability of collision resolving. They have proposed a limited reading efficiency expressed as:

$$\eta_{[5]} = \sum_{i=1}^M P(i), \quad (4)$$

where  $P(i) = \binom{n}{i} \left(\frac{1}{L}\right)^i \left(1 - \frac{1}{L}\right)^{n-i}$ , and  $M$  represents the number of collided tags that the reader is capable to recover. The authors assumed that the probability to recover one tag from  $i$  collided tags equals to 100%, independent of  $i$ , which is not a practical assumption. In reality, the probability to recover a single tag from  $i$  collided tags varies based on the receiver type, the number of collided tags  $i$ , and the Signal to Noise Ratio (SNR). According to the efficiency equation in (4), the authors proposed fixed values for the optimum FSA frame length based on the collision resolving capability of the RFID reader. They calculated these values numerically

by searching for the frame length which maximizes their proposed reading efficiency.

In this paper we propose a new reading efficiency metric called Multiple Collision Recovery Coefficients Reading Efficiency  $\eta_{MCRC}$ . This efficiency includes a unique collision resolving coefficient for each number of collided tags, which is closer to the RFID practical environment. Then, we calculate these coefficients for a strongest tag reply RFID reader (as an example) to show how the proposed system could be applied on real-life applications. Hence, we propose a novel closed form solution for the optimum FSA frame length which maximize the proposed efficiency metric. The proposed solution gives a direct relation between the frame length and the number of tags  $n$  in the reading area in addition to the collision recovery coefficients. This paper is organized as follows: Section II presents the system model under multiple collision recovery coefficients. In section III, we analytically derive a new closed form equation for the optimal frame length based on the multiple collision recovery coefficients. Afterwards, we show how to calculate the collision recovery coefficients under a simple RFID receiver in section IV. Finally, section V gives numerical results on the improvements of the new optimization criterion, before we conclude in section VI.

## II. SYSTEM MODEL UNDER MULTIPLE COLLISION RECOVERY COEFFICIENTS

In this section, we present a new FSA efficiency metric called Multiple Collision Recovery Coefficients Reading Efficiency  $\eta_{MCRC}$ . The main contribution in this efficiency is: It contains a unique collision recovery coefficient  $\alpha_i$  for each probability of collision  $P_{col.}(i)$ . These new coefficients indicates the ability of the reader to recover one tag from  $i$  collided tags. The proposed reading efficiency  $\eta_{MCRC}$  is expressed as:

$$\eta_{MCRC} = P(1) + \sum_{i=2}^n \alpha_i P_{col.}(i). \quad (5)$$

Figure 1 presents the distribution of collision probability for a collided slot in FSA. This simulation is under condition of uniformly distributed  $L$  with  $\frac{n}{2} \leq L \leq 2n$ , which is the practical range of the frame length in RFID systems [6]. According to figure 1, the probability that the collided slot comes from two, three, or four collided tags is equal to  $P_{col.}(2) + P_{col.}(3) + P_{col.}(4) \simeq 96\%$ , and the remaining tag collisions  $\sum_{i=5}^n P_{col.}(i) \simeq 4\%$ . Moreover, the values of the collision recovery coefficients  $\alpha_i$ , where  $i > 4$  are practically almost neglected.

Therefore, our proposed  $\eta_{MCRC}$  for the practical RFID environment can be expressed as shown:

$$\eta_{MCRC} = P(1) + \alpha_2 P_{col.}(2) + \alpha_3 P_{col.}(3) + \alpha_4 P_{col.}(4), \quad (6)$$

where  $\alpha_2$ ,  $\alpha_3$ , and  $\alpha_4$  are respectively the second, third, and fourth collision recovery coefficients.

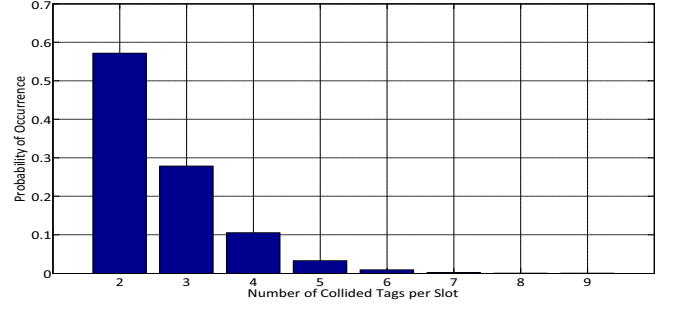


Figure 1: Distribution of collision probability for a collided slot in FSA, under condition of  $\frac{n}{2} \leq L \leq 2n$

## III. PROPOSED CLOSED FORM FRAME LENGTH

The next step is to derive a closed form solution for the new optimum frame length  $L_{MCRC}$  under multiple collision recovery coefficients environment.  $L_{MCRC}$  can be optimized by finding the value of  $L$  which maximizes  $\eta_{MCRC}$ . According to [7], if  $L \gg 1$ , and  $n \gg i$ ,  $P(i)$  is considered Poisson distribution. So  $P(i)$  can be represented as:

$$P(i) \simeq \frac{1}{i!} \cdot \beta^{-i} \cdot e^{-\frac{1}{\beta}}, \quad (7)$$

where  $\beta = \frac{L}{n}$ . After substituting by (7) in (6) we get:

$$\eta_{MCRC} = e^{-\frac{1}{\beta}} \cdot \left( \beta^{-1} + \frac{\alpha_2}{2} \beta^{-2} + \frac{\alpha_3}{6} \beta^{-3} + \frac{\alpha_4}{24} \beta^{-4} \right). \quad (8)$$

Now we have to find the value of  $\beta$  which maximizes  $\eta_{MCRC}$ . This is achieved by differentiating the reading efficiency in (8) with respect to the  $\beta$  and equate the result to zero. After differentiating, the equation can be simplified as:

$$-e^{-\frac{1}{\beta}} \cdot \left( \beta^{-2} + \beta^{-3}(\alpha_2 - 1) + \frac{\beta^{-4}}{2}(\alpha_3 - \alpha_2) + \frac{\beta^{-5}}{6}(\alpha_4 - \alpha_3) - \frac{\beta^{-6} \cdot \alpha_4}{24} \right) = 0. \quad (9)$$

After multiplying the equation by  $-e^{\frac{1}{\beta}} \cdot \beta^6$ , the equation finally results in:

$$\underbrace{1}_a \beta^4 + \underbrace{(\alpha_2 - 1)}_b \beta^3 + \underbrace{\frac{(\alpha_3 - \alpha_2)}{2}}_c \beta^2 + \underbrace{\frac{(\alpha_4 - \alpha_3)}{6}}_d \beta - \underbrace{\frac{\alpha_4}{24}}_e = 0. \quad (10)$$

Equation (10) has four roots [8]:

$$\begin{aligned} \beta_{1,2} &= -\frac{b}{4a} - S \pm 0.5 \sqrt{\underbrace{-4S^2 - 2P + \frac{q}{S}}_X} \\ \beta_{3,4} &= -\frac{b}{4a} + S \pm 0.5 \sqrt{\underbrace{-4S^2 - 2P - \frac{q}{S}}_Y}, \end{aligned} \quad (11)$$

where  $P = \frac{8ac - 3b^2}{8a^2}$ ,  $q = \frac{b^3 - 4abc + 8a^2d}{8a^3}$

and,  $S = 0.5 \sqrt{-\frac{2}{3}P + \frac{1}{3a} \left( Q + \frac{\Delta_0}{Q} \right)}$ ,  $Q = \sqrt[3]{\frac{\Delta_1 + \sqrt{\Delta_1^2 - 4\Delta_0^3}}{2}}$ .

with,  $\Delta_0 = c^2 - 3bd + 12ae$ ,  $\Delta_1 = 2c^3 - 9bcd + 27ad^2 - 72ace$

According to practical ranges for the collision recovery coefficients,  $\alpha_i$ ,  $0 \leq \alpha_i \leq 1$ , and  $\alpha_2 \geq \alpha_3 \geq \alpha_4$ . Therefore, we can proof that the signs of the polynomial coefficients

are constant and not changing in all ranges of  $\alpha_i$  and can be expressed as follows:  $a = +$ ,  $b = -$ ,  $c = -$ ,  $d = -$ , and  $e = -$ .

Using Descartes' rules of sign [9], we can count the number of real positive solutions that a polynomial has. Assume that the polynomial in (10) is  $P(\beta)$ , and let  $\nu$  be the number of variations in the sign of the coefficients  $a, b, c, d, e$ , so  $\nu = 1$ . Let  $n_p$  be the number of real positive solutions. According to Descartes' rules of sign [9]:

- $n_p \leq \nu$  which means that  $n_p = 0$  or  $1$ .
- $\nu - n_p$  must be an even integer. Therefore,  $n_p = 1$ .

According to the above two Descartes' rules of sign, there is only one valid real positive solution for equation (10). Now, we will identify which solution is the valid one. we have one real solution, so there are two possibilities for the solutions:

1. One positive real solution and the remaining three solutions are negative. In this case, all solutions are real and we need just to identify what is the root which has the largest values from the four solutions. According to (11), the value of the square roots  $\sqrt{X}$  and  $\sqrt{Y}$  are positive reals, because we do not have complex solutions. This means,  $\beta_1 > \beta_2$  and also  $\beta_3 > \beta_4$ . Thus, the solution will be either  $\beta_1$  or  $\beta_3$ . Moreover, the value of  $S$  should be also positive real, and  $q$  has always negative real value. so  $\beta_3 > \beta_1$  which means in this case that our solution is  $\beta_3$ .

2. Two complex solutions, one real positive solution, and one negative solution. In this case, we have either  $\beta_{1,2}$  or  $\beta_{3,4}$  real solutions.  $S$  should be positive real number, and the complex value comes only from the square roots  $\sqrt{X}$  and  $\sqrt{Y}$ . Moreover,  $q$  has always negative real value. Therefore, in (11) the value of  $X < Y$ . So  $\beta_{1,2}$  must be the complex roots, and as mentioned before that  $\beta_3 > \beta_4$ , so  $\beta_4$  is the negative root and  $\beta_3$  is the positive real root.

Based on the above discussion, the proposed closed form optimum frame length  $L_{MCRC}$  under the multiple collision recovery coefficients leads to:

$$L_{MCRC} = \left( -\frac{b}{4a} + S + 0.5\sqrt{-4S^2 - 2P - \frac{q}{S}} \right) \cdot n \quad (12)$$

According to (12), the proposed equation gives a linear relation wrt. the number of tags  $n$ , and includes the effect of different collision recovery coefficients. In case that, the RFID reader has no collision resolving capability (i.e.  $\alpha_2 = \alpha_3 = \alpha_4 = 0$ ), the proposed formula gives  $L_{MCRC} = n$ , which is equal to the frame length in the conventional case. When the RFID reader has a full and equal collision resolving capability for the two, three, and four collided tags per slot i.e.  $\alpha_2 = \alpha_3 = \alpha_4 = 1$ , the proposed formula gives  $L_{MCRC} = 0.452 \cdot n$  which match the results in [5].

#### IV. COLLISION RECOVERY COEFFICIENTS CALCULATIONS

Now the calculation of the the collision recovery coefficients  $\alpha_2$ ,  $\alpha_3$  and  $\alpha_4$  will be clarified, which are the main optimization variables in our proposal. The values of these coefficients are strongly depend on the receiver type. Therefore, calculations of the collision recovery coefficients are done

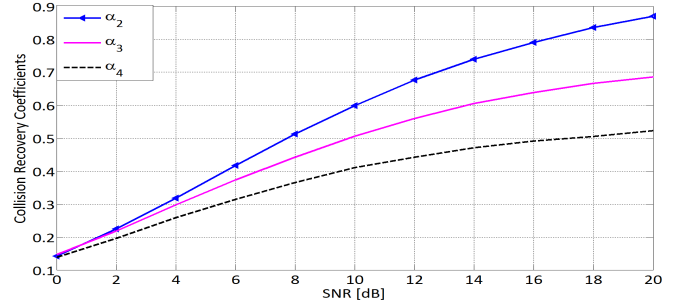


Figure 2: Collision recovery coefficients versus SNR using the strongest tag reply receiver.

based on the strongest tag reply reader. The reader resolves the collision based on the strongest tag reply. Thus, the collision will be resolved if the strongest tag reply is stronger than the summation of the power for the other collided tags within the same slot. The main advantage of this reader is that it does not need any channel state information (CSI) to recover the strongest tag. According to EPCglobal C1G2 standard [1], collisions in RFID systems are occurred only at 16 bits packet called *RN16*. If any single bit error occurred in this packet, the total packed is considered lost. Therefore, the meaning of the the collision recovery probability coefficients  $\alpha_i$  is the probability that the RFID reader can identify a complete *RN16* packet from  $i$  collided tags. Therefore, the collision recovery coefficient can be expressed as:

$$\alpha_i = (1 - PER_i), \quad (13)$$

where  $PER_i$  is the Packet Error Rate for  $i$  collided tags. In this work, we measure the SNR for each slot, then we calculate the average SNR per frame as  $E \left\{ \frac{|h|^2 \cdot x^2}{\sigma^2} \right\}$ , where  $\sigma$  is the standard deviation of the Additive White Gaussian Noise (AWGN) per slot, and we used normalized signal power i.e.  $E \{x^2\} = 1$ . In addition, the number of samples per symbols is normalized. Based on [10], we assumed that the equivalent channel coefficients  $h$  follow Rayleigh fading. The single Rayleigh channel coefficients are independent zero mean circularly symmetric complex Gaussian random variables with normalized energy  $E \{|h|^2\} = 1$ , and all tags are statistically identical, which means all of them experience the same path loss. Therefore, the average SNR per frame is  $E \left\{ \frac{1}{\sigma^2} \right\}$ . Figure 2 shows the values of the collision recovery coefficients versus the average SNR per frame using the strongest tag reply receiver.

#### V. SIMULATION RESULTS

For every type of RFID reader, each value of the SNR leads to corresponding values of the collision recovery coefficients  $\alpha_2$ ,  $\alpha_3$  and  $\alpha_4$ . Figure 3 shows a comparison frame length comparison between the proposed formula in (12) and the numerical solution, which maximizes the reading efficiency in (8) versus the SNR. Both simulations used the same receiver type which is the strongest tag reply receiver. According to figure 3, the proposed formula gives identical results

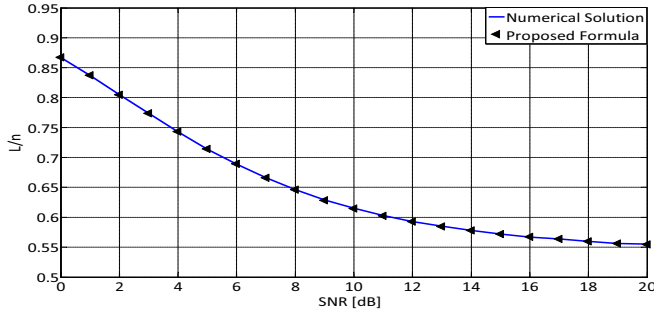


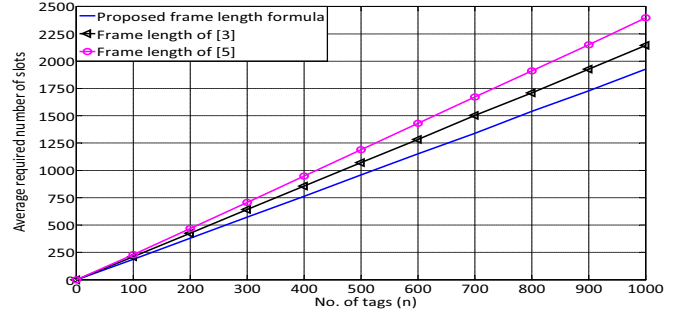
Figure 3: Frame length comparison between the proposed formula and the numerical solution vs. the SNR assuming a strongest tag reply receiver.

compared to the numerical solution in the complete SNR range.

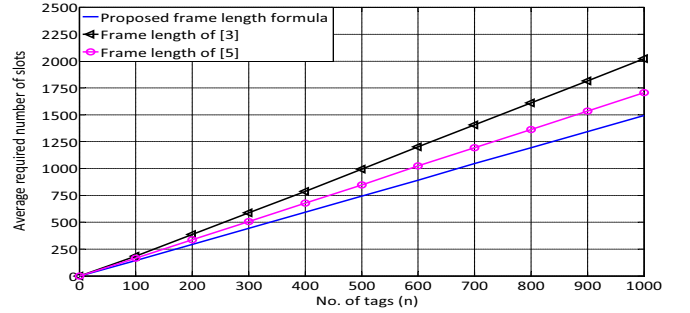
According to practical measurements, the SNR typical range for successful slots is between 4 dB and 12 dB. Therefore, we calculated the total average number of slots needed to identify a complete bunch of tags using the strongest tag reply receiver. Then, we compared the results using the proposed formula and using the frame lengths proposed by [3], and [5]. These calculations are done for the boundary of the practical range of the SNR (i.e. 4 dB and 12 dB). The authors of [3] used the reading efficiency equation in (2) assuming constant collision recovery coefficient  $\alpha = \alpha_2$  regardless the number of collided slots. In [5], the authors assumed limited collision recovery capability. However, they have assumed a unity collision recovery coefficients  $\alpha_2 = \alpha_3 = \alpha_4 = 1$  as shown in (4). Based on figure 4a and 4b, the proposed formula gives better performance than the frame lengths of [3] and [5]. As shown in figure 4a, the performance of [3] is closer than [5] to our proposed performance, because the values of the collision recovery coefficients are small. So the effect of the remaining terms in [3] is smaller than the effect of assuming  $\alpha_2 = \alpha_3 = \alpha_4 = 1$ . However, in case of SNR 12 dB, the values of the collision recovery coefficients are bigger to make the performance of [5] is closer than [3] to our proposed performance as shown in figure 4b.

## VI. CONCLUSION

This paper proposes a new performance metric called Multiple Collision Recovery Coefficients Reading Efficiency  $\eta_{MCRC}$ . The proposed reading efficiency consider the effect of the practical limitations and the differences in collision recovery probability coefficients of the RFID readers. Then, a novel closed form solution for the optimum value of FSA frame length in RFID systems. The theoretical derivations lead to a new optimization criterion that can be easily implemented in RFID readers. The proposed frame length equation gives the most accurate results in the complete range of the collision recovery coefficients  $\alpha_i$  compared to the other solutions. Timing comparisons are presented to shows the saving in reading time using the proposed frame length compared to the other proposals. Finally the proposed equation can be applied for all applications using Framed Slotted ALOHA.



(a) SNR = 4 dB



(b) SNR = 12 dB

Figure 4: Average number of slots needed to identify a complete bunch of tags using the proposed formula, frame length proposed by [3], and [5].

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