

## 8. LINK-BASED RANKING

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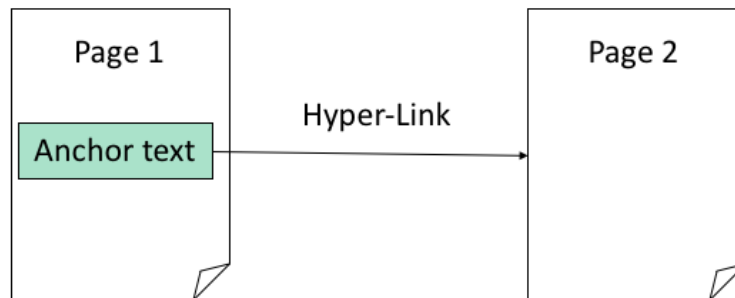
Link Analysis - 1

todo: what is this for

# Web is a Hypertext

Web documents are connected through hyperlinks

1. **Anchor text** describes content of referred document
2. **Hyperlink** is a quality signal



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Link Analysis - 2

Beyond the textual contents, Web documents consist also of hyperlinks. A hyperlink can be exploited for information retrieval in two ways:

1. The link is surrounded by some **textual information that presumably refers to the content of the document** the link is pointing to. Thus this text can complement the content of the referred document.
2. The link can also be **considered as an endorsement** of the referenced document by the author of the referring document. Thus the link can be used as a signal for quality and importance of the referred document.

# INDEXING ANCHOR TEXT

## Indexing Anchor Text

Anchor text is loosely defined as the text surrounding a hyperlink

*Example:* “You can find cheap cars [here](#).”

Anchor text: “You can find cheap cars here”

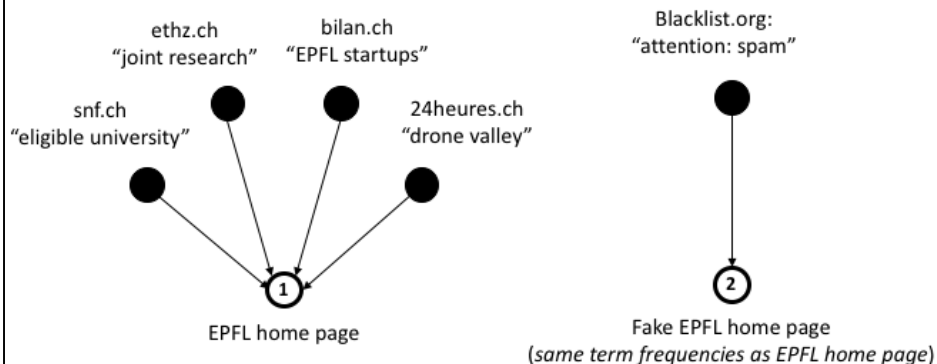
Anchor text may contain a lot of additional content on the referred page

- It might be a better description of the page than the page content itself

Anchor text can be defined in different ways. In general, it is considered as the text that is surrounding the link, and not only the text contained as part of the link tag (in the example, this text would simply be “here”.) This text can contain valuable information on the referred page and thus be very helpful in retrieval.

## Example

When indexing a document  $D$ , include (with some weight) anchor text from links pointing to  $D$



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Link Analysis - 5

This example illustrates the use of anchor text in retrieval. Typically a home page is very graphical and contains often very little relevant content. If we consider a home page, such as the EPFL home page, we would probably find many pages that very well characterize EPFL, such as pages mentioning topics related to research and technology transfer. Typically these links would point to the EPFL home page. Thus it would also let the home page stand out from other EPFL pages (such as pages on the laws and ordinances of EPFL).

Assume that a malicious Internet user would create a fake EPFL home page. Then chances that such a page is referred by reputed organizations, such as SNF, is very low. On the other hand pages listing spam pages might point to such a page and indicate its true character.

Therefore it makes a lot of sense to **include such anchor text in the representation of the document**. This is usually done by adding it with a given weight.

## Scoring of Anchor Text

Score anchor text with a weight depending on the authority of the anchor page's website

- E.g., if we were to assume that content from cnn.com or yahoo.com is authoritative, then trust (more) the anchor text from them

Score anchor text from other sites (domains) higher than text from the same site

- non-nepotistic scoring

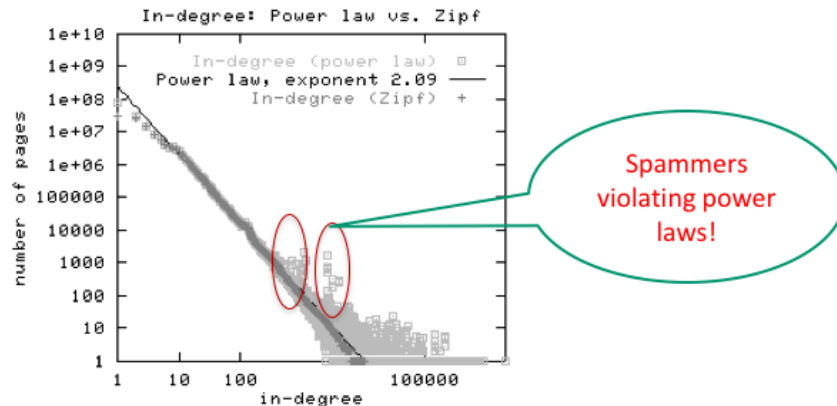
to prevent own domain from scoring itself

To fight spam, one can adapt the weighting function to the reputation of the page that is referring. This could be done by having a directory of highly reputed sites and to give high weights to its members. A more fundamental of producing such reputation scores we will see in the following part on link-based ranking.

In order to avoid self-promotion, another method to fight link spamming is to give lower weights to links within the same site (nepotistic scoring = promoting your own family members).

# Indexing Anchor Text

Can sometimes lead to unexpected effects, e.g., easily **spammable**



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Link Analysis - 7

One of the risks of including anchor text is that it makes pages spammable. Malicious users could create spam pages that point to web pages and try to relate it to contents that serve their interests (e.g., higher the quality of preferred pages by adding links, lower the quality of the undesired page by attaching negative anchor text). That this is happening can be seen from analyzing the in-degree distribution of Web pages. The figure shows a standard log-log representation of the in-degree vs. the frequency of pages. Normally this relationship should follow a power-law, which shows in a log-log representation as a linear dependency. In real Web data, we see that this power law is violated, and that certain levels of in-degrees are over-represented. This can be attributed to link spamming, which does create moderate numbers of additional links on Web pages.

todo

# LINK-BASED RANKING: PAGERANK



## Citation Analysis

Bibliometry: analysis of citations in scientific publications

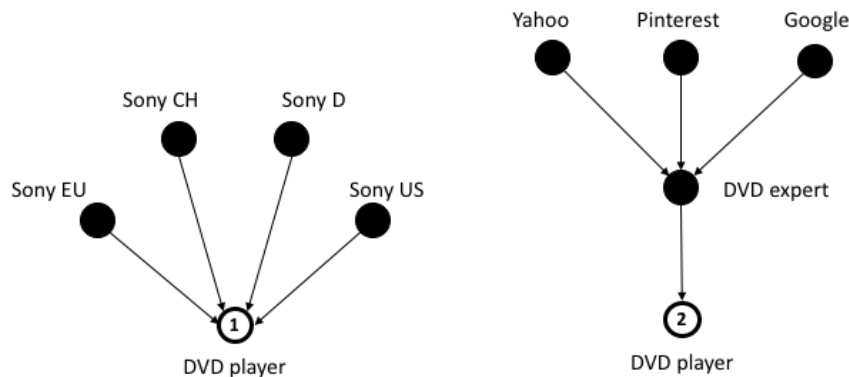
- Citation frequency: how important is a paper of author?
- Co-citation analysis: articles that co-cite the same articles are related
- Citation indexing: who is this author cited by?
- Impact factor: Authority of sources, such as journals

The use of links in order to evaluate the quality of information sources has a long tradition, in particular in science. The scientific discipline of bibliometry is fully devoted to the problem of evaluating the quality of research through citation analysis. Different ideas can be exploited to that end:

- The frequency of citations to a paper, indicating how popular / visible it is
- Co-citation analysis in order to identify research working in related disciplines
- Citation indexing in order to explore the profile of researchers referring to an author (as indicator of both the discipline and the quality)
- Analysis of the authority of sources of scientific publications, e.g. journals, publishers, conferences. This measure can then in turn be used to weight the relevance of publications.

All these ideas could also be exploited for any other document collections that have references, in particular for Web document collections with hyperlinks.

## Citations on the Web



Full text retrieval result with equal ranking; which page is more relevant ?

- relevance related to number of referrals (incoming links)
- relevance related to number of referrals with high relevance

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Link Analysis - 10

When retrieving documents from the Web, the link structure bears important information on the relevance of documents. Generally speaking, a document that is referred more often by other documents through Web links, is likely to be of higher interest and therefore relevance. So a possibility to rank documents is considering the number of incoming links. Doing this allows to distinguish documents that otherwise would be ranked equally or similarly when relying on text retrieval alone.

However, when doing this, also the importance of documents having a link to the document to be ranked may be different. Therefore not only counting then number of incoming links, but also weighing the links by the relevance of documents that contain these links appears to be appropriate. The same reasoning of course again applies then for evaluating the relevance of documents pointing to the document and so forth.

## The Web isn't Scholarly Citation

Millions of participants, each with self interests

- Spamming is widespread

Once search engines began to use links for ranking (roughly 1998), link spam grew

- You can join a link farm – a group of websites that heavily link to one another

Simple link counting might not be appropriate!

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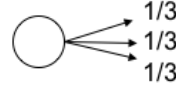
Link Analysis - 11

In the web, analysis of **links servers** not only (or necessarily) the purpose of finding authoritative sources, such as in science. It also and in particular can help to fight the aforementioned problem of link spamming. When using such methods, such as in any arms race, of course counter-measures have been devised by the adversaries. These resulted in the creation of link farms, that try to boost the relevance of Web pages by **creating many (artificial) links** to them. So **just counting incoming links is probably not such a good idea in order to evaluate the quality of a Web page.**

## Link-based Ranking: Idea

Imagine a user doing a **random walk** on Web pages:

- Start at a random page
- At each step, leave the current page along one of the links on that page, with same probability



“In the long run” each page has a long-term visit rate - use this as the page’s score

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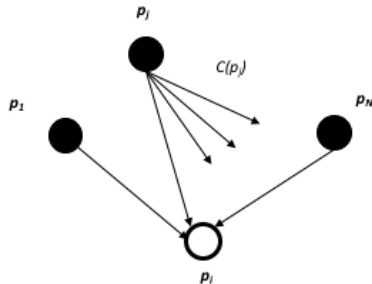
Link Analysis - 12

The fundamental idea of link-based scoring while attempting to fight spam is based on the concept of a random walker on the Web. Assume a random walker would visit a Web page. Then it would follow each of the outgoing links with equal probability. If walking for a very long time in the Web, it would have a certain fraction of visits it passes by every page.

One of the consequences of this model would be that pages that have few in-links, would be relatively infrequently visited. Since link farms usually have not many links pointing to them, in this way their influence in terms of link spamming would be moderated.

Inlinks = Links on other websites that send traffic to your site.  
Random walker works because link farms usually do not have many links pointing to them.  
(Inlinks)

# Random Walker Model



relevance of referreals  
and no. of referrals

$$P(p_i) = \sum_{p_j | p_j \rightarrow p_i} \frac{P(p_j)}{C(p_j)}$$

$p_j$  = probability of  
visiting page  $p_j$   
divided by no of.  
outgoing links of  $p_j$

$N$  is the number of Web pages

$C(p)$  is the number of outgoing links of  
page  $p$

$P(p_i)$  probability to visit page  $p_i$ , where  
page  $p_i$  is pointed to by pages  $p_1$  to  $p_N$  =  
**relevance**

## Result

- If a random walker visits a page more often it is more relevant
- takes into account the number of referrals AND the relevance of referrals

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Link Analysis - 13

Here we describe the random walker model more formally. The long term probability of visiting a page, depends on the long-term probability of having visited a page that is referring to the page. Thus the formulation of the random walker model becomes recursive.

## Transition Matrix for Random Walker

The definition of  $P(p_i)$  can be reformulated as matrix equation

$r_{ij}$  = to i from j  
= 1 / outgoing links from j

$$R_{ij} = \begin{cases} \frac{1}{C(p_j)}, & \text{if } p_j \rightarrow p_i \\ 0, & \text{otherwise} \end{cases}$$

R = number of referrals  
p = relevance

$$\vec{p} = (P(p_1), \dots, P(p_n))$$

$$\vec{p} = R \cdot \vec{p}, \quad \|\vec{p}\|_1 = \sum_{i=1}^n p_i = 1$$

The vector of page relevance values is an Eigenvector of the matrix R

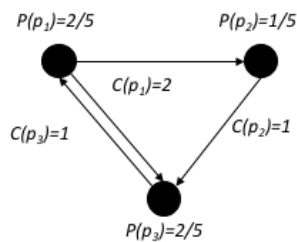
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Link Analysis - 14

In order to determine the solution to the recursive formulation of the probabilities of a random walker to visit a page, we introduce a transition probability matrix R, which captures the probability of transitioning from one page to another. We also require that the long-term probabilities of visiting a page add up to 1. With this representation the **long-term visiting probabilities** become the Eigenvector of matrix R. More precisely, they are the Eigenvector with the largest Eigenvalue.

p = long term probability vector, computed after many iterations of  $p = R \cdot p$ . also can be found via the eigenvectors

## Example



Links from  $p_1$

Links to  $p_1$

$$L = \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \end{pmatrix}$$

Link Matrix

divide by sum of column

$$R = \begin{pmatrix} 0 & 0 & 1 \\ \frac{1}{2} & 0 & 0 \\ \frac{1}{2} & 1 & 0 \end{pmatrix}, \quad \vec{p} = \begin{pmatrix} \frac{2}{5} \\ \frac{1}{5} \\ \frac{2}{5} \end{pmatrix}$$

why largest eigenvector ?

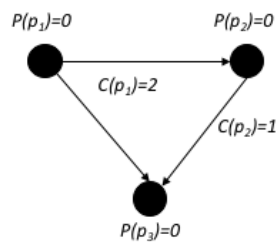
Ranking  $p_1, p_3 > p_2$

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Link Analysis - 15

This example illustrates the computation of the probabilities for visiting a specific Web page. The values  $C(p_i)$  correspond to the transition probabilities. They can be derived from the link matrix, which is the matrix with entries 1 at  $(i,j)$  if there exists a link from  $p_i$  to  $p_j$ , by dividing the values in the columns by the sum of the values found in the column. The probability of a random walker being in a node is then obtained from the Eigenvector of this matrix.

## Modified Example



Links from  $p_1$

$$L = \begin{pmatrix} 0 & 0 & 0 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \end{pmatrix}$$

Links to  $p_1$

Link Matrix

$$R = \begin{pmatrix} 0 & 0 & 0 \\ \frac{1}{2} & 0 & 0 \\ \frac{1}{2} & 1 & 0 \end{pmatrix}, \quad \vec{p} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$$

why is  $p_3$  not 1  
- just multiply  $Rp$ . still ends up  
as  $(0,0,0)$

No Ranking

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Link Analysis - 16

This example illustrates the problem of dead ends. We see that there exists a node  $p_3$  that is a "sink of rank". Any random walk ends up in this sink, and therefore the other nodes do not receive any ranking weight. Consequently also the rank of sink does not. Therefore the only solution to the equation  $p=Rp$  is the zero vector.

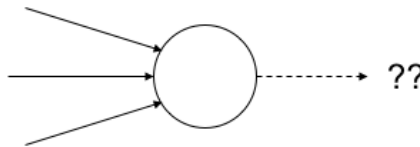
why is the eigenvector of this matrix  
0 ? What about the second largest ?



## Pure Random Walker Does Not Work

The web is full of dead-ends

- Random walk can get stuck in dead-ends
- Makes no sense to talk about long-term visit rates



### Teleporting

- At a dead end, jump to a random web page
- At any non-dead end, jump to a random web page with some probability (e.g. 15%)
- Result: Now cannot get stuck locally, there is a long-term rate at which any page is visited

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Link Analysis - 17

A practical problem with the random walker is the fact that there exist Web pages that have no outgoing links. Thus the random walker would get stuck. To fix this, the concept of teleporting is introduced, where the random walker jumps to any randomly selected Web page. In case of arriving at a dead end, it jumps in any case, otherwise with a given probability instead of following an outgoing link.

Another problem are pages that have no incoming links: they would never be reached by the random walker, and the weight that they could provide to other pages would not be considered. Also this problem is solved by teleporting.

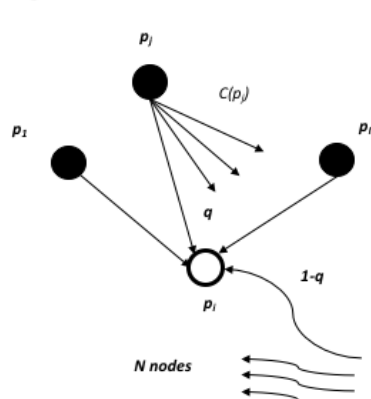
teleport fix pages with  
no incoming links

## Source of Rank: Teleporting

### Assumption

- random walker jumps with probability  $1-q$  to an arbitrary node
- thus it can leave dead ends and nodes without incoming links are reached

### PageRank method



$$P(p_i) = c \left( \frac{(1-q)}{N} + q \sum_{p_j | p_j \rightarrow p_i} \frac{P(p_j)}{C(p_j)} \right), c \leq 1$$

$$\vec{p} = c((qR + (1-q)E)) \cdot \vec{p}, \quad E = \left[ \frac{1}{N} \right]_{N \times N}$$

$$\vec{p} = c(qR \cdot \vec{p} + \frac{(1-q)}{N} \vec{e}), \quad \vec{e} = (1, \dots, 1)$$

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Link Analysis - 18

what is c for ?

$E\vec{p} = \vec{e}/N$   
where  $\vec{e}$  is the 1 vector

$E\vec{p} = \vec{E}$

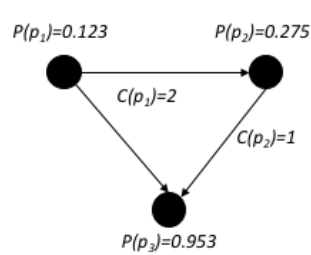
To avoid the previously described problem, we add a "source of rank". The idea is that a random walker in each step can, rather than following a link, jump to any page with probability  $1-q$ . Therefore the random walker can leave pages that have no outgoing links and also pages without incoming links can be reached by the random walk and give weight to other pages. In the mathematical formulation of the random walk this is resulting in an additional term for the source of rank. Since at each step the random walker makes a jump with probability  $1-q$  and any of the  $N$  pages is reached with the same probability the additional term is  $(1-q)/N$ . Reformulating this again as matrix equation requires adding a  $N \times N$  Matrix  $E$  with all entries being  $1/N$ . This is equivalent to saying that with probability  $1/N$  transitions among any pairs of nodes (including transition from a node to itself) are performed. Since the vector  $\vec{p}$  has norm 1, i.e., the sum of the components is exactly 1,  $E \cdot \vec{p} = \vec{e}$ , where  $\vec{e}$  is the unit vector, the matrix equation can be reformulated in the second form shown below. The method described is called PageRank and is used by Google. By modifying the values of the matrix  $E$  also a priori knowledge about the relative importance of pages can be added to the algorithm.

Original formulation:  $\vec{p} = R\vec{p}$

With teleport:  $\vec{p} = q \cdot R\vec{p} + (1-q) \cdot E\vec{p}$

$E$  can be used to incorporate priori knowledge

## Modified Example



$$R = \begin{pmatrix} 0 & 0 & 0 \\ \frac{1}{2} & 0 & 0 \\ \frac{1}{2} & 1 & 0 \end{pmatrix}, E = \begin{pmatrix} \frac{1}{3} & \frac{1}{3} & \frac{1}{3} \\ \frac{1}{3} & \frac{1}{3} & \frac{1}{3} \\ \frac{1}{3} & \frac{1}{3} & \frac{1}{3} \end{pmatrix}, q = 0.9$$

$$qR + (1-q)E = \begin{pmatrix} \frac{1}{30} & \frac{1}{30} & \frac{1}{30} \\ \frac{29}{60} & \frac{1}{30} & \frac{1}{30} \\ \frac{29}{60} & \frac{14}{15} & \frac{1}{30} \end{pmatrix} \rightarrow \vec{p} = \begin{pmatrix} 0.123 \\ 0.275 \\ 0.953 \end{pmatrix}$$

Ranking  $p_3 > p_2 > p_1$

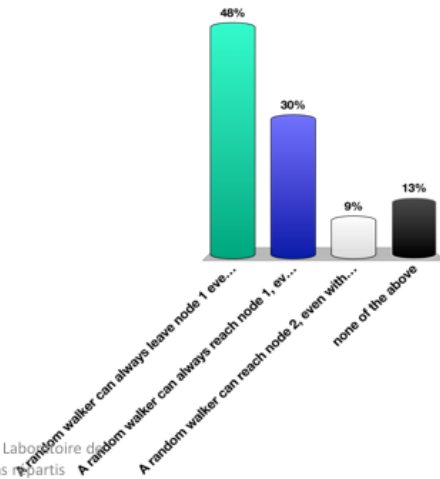
With the modification of rank computation using a source of rank, we obtain for our example again a non-trivial ranking which appears to be appropriate.

Consider the following matrix for assigning random jump probabilities

$$\begin{pmatrix} \frac{1}{3} & 0 & 0 \\ \frac{1}{3} & 0 & 0 \\ \frac{1}{3} & 0 & 0 \end{pmatrix}$$

It means

1. A random walker can always leave node 1 even without outgoing edges
2. A random walker can always reach node 1, even without incoming edges
3. A random walker can always leave node 2, even without outgoing edges
4. none of the above



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## Practical Computation of PageRank

Iterative computation     $\vec{p}_0 \leftarrow \vec{s}$   
while  $\delta > \epsilon$   
     $\vec{p}_{i+1} \leftarrow qR \bullet \vec{p}_i$   $q \times R_p$ , where  $q$  is a scalar  
     $\vec{p}_{i+1} \leftarrow \vec{p}_{i+1} + \frac{(1-q)}{N} \vec{e}$   
     $\delta \leftarrow \|\vec{p}_{i+1} - \vec{p}_i\|_1$

$\epsilon$  termination criterion

$s$  arbitrary start vector, e.g.  $s = e$

For the practical computation of the PageRank ranking an iterative approach can be used. It is derived from the second form of the formulation of the visiting probabilities of the random walker that we have given. The vector  $e$  used to add a source of rank has not necessarily to assign uniform weights to all pages, but might reflect itself a ranking of Web pages.

## Example: ETHZ Page Rank

| Doc. ID | Rank Value | URL   |
|---------|------------|---|
| 1       | 0.002536   | <a href="http://www.ethz.ch/">http://www.ethz.ch/</a>   |
| 146     | 0.002292   | <a href="http://www.ethz.ch/ir_amb/">http://www.ethz.ch/ir_amb/</a>   |
| 10      | 0.000654   | <a href="http://www.ethz.ch/default_de.asp">http://www.ethz.ch/default_de.asp</a>   |
| 35      | 0.000511   | <a href="http://www.rereth.ethz.ch/">http://www.rereth.ethz.ch/</a>   |
| 376124  | 0.000503   | <a href="http://computing.ee.ethz.ch/sepp/matlab-5.2-to/helpdesk.html">http://computing.ee.ethz.ch/sepp/matlab-5.2-to/helpdesk.html</a>   |
| 67378   | 0.000497   | <a href="http://computing.ee.ethz.ch/sepp/">http://computing.ee.ethz.ch/sepp/</a>   |
| 59887   | 0.000485   | <a href="http://www.computing.ee.ethz.ch/sepp/">http://www.computing.ee.ethz.ch/sepp/</a>   |
| 89307   | 0.000485   | <a href="http://www.isg.inf.ethz.ch/docu/documents/java/jdk1.2.2ref/docs/api/overview-summary.html">http://www.isg.inf.ethz.ch/docu/documents/java/jdk1.2.2ref/docs/api/overview-summary.html</a> |
| 216716  | 0.000485   | <a href="http://www.isg.inf.ethz.ch/docu/documents/java/jdk1.2/api/overview-summary.html">http://www.isg.inf.ethz.ch/docu/documents/java/jdk1.2/api/overview-summary.html</a>                     |
| 147932  | 0.000484   | <a href="http://isg.inf.ethz.ch/docu/documents/java/jdk1.2ref/docs/api/overview-summary.html">http://isg.inf.ethz.ch/docu/documents/java/jdk1.2ref/docs/api/overview-summary.html</a>             |
| 175544  | 0.000484   | <a href="http://www.isg.inf.ethz.ch/docu/documents/java/jdk1.2ref/docs/api/overview-summary.html">http://www.isg.inf.ethz.ch/docu/documents/java/jdk1.2ref/docs/api/overview-summary.html</a>     |
| 186766  | 0.000478   | <a href="http://isg.inf.ethz.ch/docu/documents/java/jdk1.2/api/overview-summary.html">http://isg.inf.ethz.ch/docu/documents/java/jdk1.2/api/overview-summary.html</a>                             |
| 228634  | 0.000477   | <a href="http://isg.inf.ethz.ch/docu/documents/java/jdk1.2.1ref/docs/api/overview-summary.html">http://isg.inf.ethz.ch/docu/documents/java/jdk1.2.1ref/docs/api/overview-summary.html</a>         |
| 228421  | 0.000464   | <a href="http://isg.inf.ethz.ch/docu/documents/java/jdk1.2.2ref/docs/api/overview-summary.html">http://isg.inf.ethz.ch/docu/documents/java/jdk1.2.2ref/docs/api/overview-summary.html</a>         |
| 3161    | 0.00045    | <a href="http://www.ethz.ch/ir_amb/reto_ambuehler.html">http://www.ethz.ch/ir_amb/reto_ambuehler.html</a>   |
| 215673  | 0.000447   | <a href="http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/files.html">http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/files.html</a>                   |
| 259672  | 0.000447   | <a href="http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/globals.html">http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/globals.html</a>               |
| 259671  | 0.000447   | <a href="http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/functions.html">http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/functions.html</a>           |
| 259670  | 0.000447   | <a href="http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/annotated.html">http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/annotated.html</a>           |
| 259669  | 0.000447   | <a href="http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/classes.html">http://www.vision.ee.ethz.ch/computing/statlinks/sepp.sun/vxl-1.2b-mo/classes.html</a>               |

Figure 1: Top 20 of ETH Zurich Web Documents

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Link Analysis - 22

These are the top documents from the PageRank ranking of all Web pages at ETHZ (Data from 2001). It is interesting to see that documents related to Java documentation receive high ranking values. This is related to the fact that these documents have many internal cross-references.

# Web Search

PageRank is part of the ranking method used by Google

- Compute the global PageRank for all Web pages
- Given a keyword-based query retrieve a ranked set of documents using standard text retrieval methods
- Merge the ranking with the result of PageRank to both achieve high precision (text retrieval) and high quality (PageRank)
- Google uses also many other methods to improve ranking

Technical challenges

- Crawling the Web
- Efficient computation of Page Rank for large link databases
- Combination with other ranking methods (text)

PageRank is used as one metrics to rank result documents in Google. Essentially Google uses text retrieval methods to retrieve relevant documents and then applies PageRank to create a more appropriate ranking. Google uses also many other methods to improve ranking, e.g., by giving different weights to different parts of Web documents and user relevance feedbacks. For example, title elements are given higher weight. The details of the ranking methods are trade secrets of the Web search engine providers.

Other issues that Web search engines have to deal with, are **crawling the Web**, which requires techniques that can explore the Web without revisiting pages too frequently. Also the enormous size of the document and link database poses implementation challenges in order to keep the ranking computations scalable. One of the outcomes of solving these challenges are the recent “cloud computing” infrastructures, which are large-scale computing clusters constructed from commodity hardware.

# LINK-BASED RANKING: HITS



## Hyperlink-Induced Topic Search (HITS)

Key Idea: in response to a query, instead of an ordered list of pages, find two sets of inter-related pages:

- **Hub pages** are good lists of links on a subject
  - e.g., “World top universities”
- **Authoritative pages** are referred recurrently on good hubs on the subject
  - e.g., “EPFL”

todo: what does this mean

Best suited for “broad topic” queries rather than for page-finding queries

- Understand common perception of quality

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Link Analysis - 25

Hub-Authority ranking identifies not only pages that have a high authority, as measured by the number of incoming links, but also pages that have a substantial "referential" value, having many outgoing links (to pages of high importance). Different to the PageRank algorithm this technique has been originally proposed to post-process query results (rather than to rank pages from the complete Web graph). It can be used as a add-on to existing search engines, but as well as an alternative to the PageRank method.

high authority = high number of incoming links  
high hub = high number of outgoing links

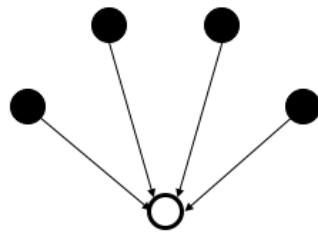
- A good hub points to pages with high authority
- A good authority is pointed to by pages that have high hub

- PageRank was more of a high authority based approach
- HITS uses both authority and hub. High hub ensures substantial "referential" value, having many outgoing links (to pages of high importance)

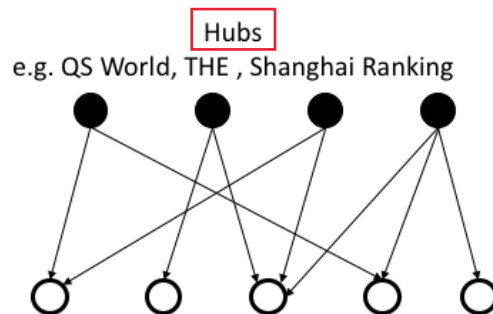
# Hub-Authority Ranking

## Approach

- **Hubs** are pages that point to many/relevant authorities
- **Authorities** are pages that are pointed to by many/relevant hubs



page with large in-degree  
e.g. EPFL



Authorities  
e.g. EPFL, MIT, Stanford, ETHZ

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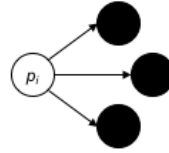
Link Analysis - 26

In order to realize the idea of distinguish authorities from hubs a simple approach is taken. Hub pages are consider as those that are referred a lot by authority pages and vice versa. The example illustrates of how this would ideally separate authority pages, in the case of universities well known universities, from hub pages, such as university ranking and portal sites.

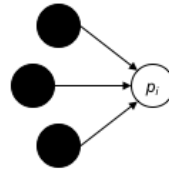
## Computing Hubs and Authorities

Repeat the following updates, for all  $p$

$$H(p_i) = \sum_{p_j \in N | p_i \rightarrow p_j} A(p_j)$$



$$A(p_i) = \sum_{p_j \in N | p_j \rightarrow p_i} H(p_j)$$



Normalize values (scaling)

$$\sum_{p_j \in N} A(p_j)^2 = 1$$

$$\sum_{p_j \in N} H(p_j)^2 = 1$$

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Link Analysis - 27

More precisely, the scores are recomputed by simply adding the scores of all incoming edges. For computing authority scores this is the hub scores and for hub scores the authority scores. In order to avoid that the scores grow continuously, they are rescaled to in each step, by normalizing the score vectors to one.

- A good hub points to pages with high authority
- A good authority is pointed to by pages that have high hub

## HITS algorithm

$n := |N|; (a_0, h_0) := \frac{1}{n^2}((1, \dots, 1), (1, \dots, 1))$

normalized this way  
because L2 norm must  
= 1

while  $l < k$

$l := l + 1$

$a_l := (\sum_{p_i \rightarrow p_1} h_{l-1,i}, \dots, \sum_{p_i \rightarrow p_n} h_{l-1,i})$

$h_l := (\sum_{p_1 \rightarrow p_i} a_{l,i}, \dots, \sum_{p_n \rightarrow p_i} a_{l,i})$

$(a_l, h_l) := \left( \frac{a_l}{|a_l|}, \frac{h_l}{|h_l|} \right)$

In practice, 5 iterations  
sufficient to converge!

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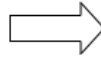
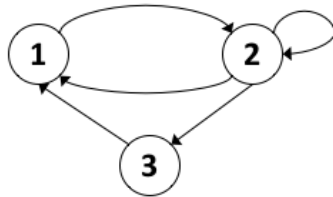
Link Analysis - 28

As for PageRank the Hub/Authority values can be iteratively computed.  $x$  corresponds to the authority weight and  $y$  to the hub weight. At each iteration the values are renormalized to length 1.

$x$  - authority  
 $y$  - hub

## Convergence of HITS

$n \times n$  link matrix  $L$



|   | 1 | 2 | 3 |
|---|---|---|---|
| 1 | 0 | 1 | 0 |
| 2 | 1 | 1 | 1 |
| 3 | 1 | 0 | 0 |

from 1  
to 1,2,3

to 1  
from 1,2,3

Up to normalization

$$y = Lx, x = L^t y, \text{ thus}$$

$y^*$  is an Eigenvector of  $LL^t$

$x^*$  is an Eigenvector of  $L^t L$

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Link Analysis - 29

We can interpret the iterative algorithm for computing HITS in terms of computing Eigenvectors for the link matrix. If  $L$  is the link matrix then the computation of  $x$  from  $y$  is defined as  $x = Ly$  and the computation of  $y$  from  $x$  is defined as  $y = L^t x$ . Therefore  $x^*$ , the authority vector obtained after convergence, is the principal Eigenvector of Matrix  $LL^t$  and  $y^*$ , the hub vector obtained after convergence, is the principal Eigenvector of matrix  $L^t L$ .

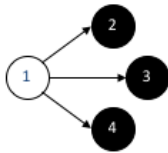
Remains the question whether this algorithm always converges. The algorithm is a particular algorithm for computing eigenvectors: the **power iteration** method. It is proven to converge against the principal Eigenvalue. Important is the normalization of the vectors obtained in each step.

$x$  - authority  
 $y$  - hub

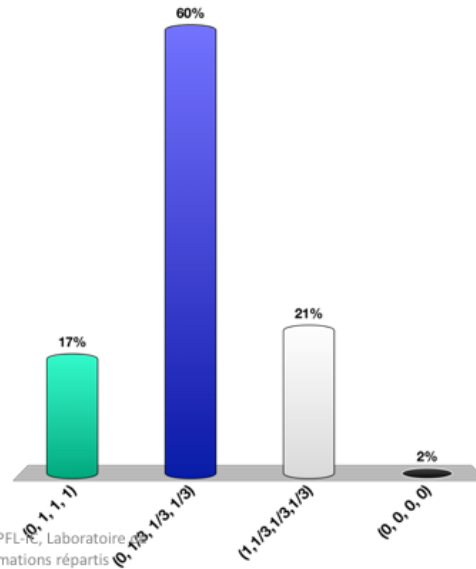
sub  $y$  into  $x$   
or  $x$  into  $y$

what does  $LL^t$  mean

The authority values of this graph are



1.  $(0, 1, 1, 1)$
2.  $(0, 1/3, 1/3, 1/3)$
3.  $(1, 1/3, 1/3, 1/3)$
4.  $(0, 0, 0, 0)$



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## Obtaining the Base Set

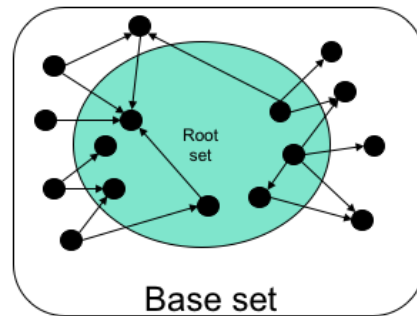
Given a query (e.g. “EPFL”), obtain all pages mentioning the query

- call this the **root set** of pages.

Add page that either

- points to a page in the root set, or
- is pointed to by a page in the root set.

Use this set as **base set**



## Sample Results Obtained

### (java) Authorities

.328 <http://www.gamelan.com/>

.251 <http://java.sun.com/>

.190 <http://www.digitalfocus.com/digitalfocus/faq/howdoi.html>

.190 <http://lightyear.ncsu.unc.edu/~srp/java/javabooks.html>

.183 <http://sunsite.unc.edu/javafaq/javafaq.html>

*Gamelan*

*JavaSoft Home Page*

*The Java Developer: How Do I...*

*The Java Book Pages*

*comp.lang.java FAQ*

### (censorship) Authorities

.378 <http://www.eff.org/>

.344 <http://www.eff.org/blueribbon.html>

.238 <http://www.cdt.org/>

.235 <http://www.vfw.org/>

.218 <http://www.aclu.org/>

*EFFweb - The Electronic Frontier Foundation*

*The Blue Ribbon Campaign for Online Free Speech*

*The Center for Democracy and Technology*

*Veters Telecommunications Watch*

*ACLU: American Civil Liberties Union*

These are two example results that have been obtained by applying this method. (It is interesting to compare those to the ones you would obtain by using a search engine alone). In particular Gamelan will not show up in the result, whereas the Java Sun page is usually among the top ranked.



# HITS Conclusions

## Potential issues

topic drift due to off topic pages

- Topic Drift: off-topic pages can cause off-topic “authorities” to be returned
- Mutually Reinforcing Affiliates: clusters of affiliated pages/sites can boost each others’ scores

## Social Network Analysis

- PageRank and HITS are examples of SN Analysis algorithms
- SNs contain a lot of other interesting structure (see later)

## Implementation

- How to efficiently obtain the base set?

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Link Analysis - 33

Similarly as for PageRank, the HITS algorithm is also vulnerable structural anomalies of the link structure, be it caused by conscious manipulation or by chance. Topic drift would imply that the set of base pages refers to a topic that is different from the original intended one, expressed by a search query. For example, if we would search for top European universities, it could well happen that the topic would drift to top US universities, when ranking hubs come into play that strongly refer to those. We could have also clusters of related pages, that in the iterative computation mutually boost each others scores, and thus give too high weight to pages that would not merit it. This problem is similar to the one of link spamming with link farms.

In terms of implementation, link-based ranking algorithms require the capability to efficiently retrieve the link graph for the Web. This problem is addressed by so-called connectivity servers that we will introduce next.

# LINK INDEXING

## Connectivity Server

Support for fast queries on the web graph

- Which URLs point to a given URL?
- Which URLs does a given URL point to?

Stores mappings in memory from

- URL to outlinks, URL to inlinks

Applications

- Link analysis (PageRank, HITS)
- Web crawl control
- Web graph analysis
  - Connectivity, crawl optimization

In order to efficiently analyze the Web graph the Web links need to be stored in a database, respectively index. A connectivity server is such an index. In essence, it stores for each URL all the URLs that the Web page at the URL is pointing to, and the URLs of other Web pages that point to this URL. Apart from link-based ranking algorithms as described before the applications are manifold.

Connectivity server  
- to efficiently get the link graph from the Web

## Adjacency Lists

The set of URLs a node is pointing to (or pointed to) sorted in *lexicographical order*

*Example:* outgoing links from [www.epfl.ch](http://www.epfl.ch)

[actu.epfl.ch/feeds/rss/mediacom/en/](http://actu.epfl.ch/feeds/rss/mediacom/en/)  
[bachelor.epfl.ch/studies](http://bachelor.epfl.ch/studies)  
[futuretudiant.epfl.ch/en](http://futuretudiant.epfl.ch/en)  
[futuretudiant.epfl.ch/mobility](http://futuretudiant.epfl.ch/mobility)  
[master.epfl.ch/page-94489-en.html](http://master.epfl.ch/page-94489-en.html)  
[phd.epfl.ch/home](http://phd.epfl.ch/home)  
[www.epfl.ch/navigate.en.shtml](http://www.epfl.ch/navigate.en.shtml)

A connectivity server has to store all outgoing (and incoming) links to a web page. For example, the home page of EPFL contains a large set of outgoing links, some of which are shown here. As a first step, the lists of links are sorted in **lexicographical order**. As a result we obtain the adjacency list, which is similar to the posting list of a document.

## Representation of Adjacency Lists

Assume each URL represented by an integer

- For a 4 billion page web, we need 32 bits per node  $(2^{32} - 1)$
- Naively, this demands 64 bits to represent each hyperlink (source and destination node)
- For the current Web: 320 GB
- Can we do better (for main memory storage)?

x2 because of source and destination

| Node | Outdegree | Successors                                     |
|------|-----------|--|
| ...  | ...       | ...  |
| 15   | 11        | 13, 15, 16, 17, 18, 19, 23, 24, 203, 315, 1034 |
| 16   | 10        | 15, 16, 17, 22, 23, 24, 315, 316, 317, 3041    |
| 17   | 0         |  |
| 18   | 5         | 13, 15, 16, 17, 50                             |
| ...  | ...       | ...  |

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Link Analysis - 37

As a first means to optimize the representation of adjacency lists, we represent each URL by one integer, instead of storing its textual form. From this we can estimate that for the current Web, we need 320 GB to represent all links:

- The current (crawled) Web has about 4 billion pages (<http://www.worldwidewebsize.com/>)
- It is estimated that a page contains on average 10 links
- We need 32 bits for each URL, which demands 64 bits for the storage of a single link.

Even with large memories this still a significant index size. In the following, we will show how to get this down to approximately ~3 bits/link which makes the index much more manageable. This will be achieved by systematically compressing the adjacency lists.

## Properties of Adjacency Lists

### Locality (within lists)

- Most links contained in a page have a navigational nature, thus their indices are close in lexicographical order
- E.g. for [www.epfl.ch](http://www.epfl.ch)
  - [futuretudiant.epfl.ch/en](http://futuretudiant.epfl.ch/en)
  - [futuretudiant.epfl.ch/mobility](http://futuretudiant.epfl.ch/mobility)

### Similarity (between lists)

- Either two lists have almost nothing in common, or they share large segments of their successor lists
- Pages that occur close to each other in lexicographic order tend to have similar lists
  - E.g. [futuretudiant.epfl.ch/en](http://futuretudiant.epfl.ch/en) and [futuretudiant.epfl.ch/mobility](http://futuretudiant.epfl.ch/mobility)

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Link Analysis - 38

For compressing adjacency lists we base ourselves on the following observations:

**Locality.** Most links contained in a page have a **navigational nature**: they lead the user to some other pages within the same host (as we can see clearly from the example of the EPFL home page). If we compare the **source and target URLs** of these links, we observe that they **share often a long common prefix**. Thus, if URLs are sorted lexicographically, the **index of source and target are close** to each other. **Locality is a property of a list, thus addresses intra-list similarity.**

**Similarity.** Pages that occur close to each other (in lexicographic order) tend to have many common successors. This is because many navigational links are the same within the same local cluster of pages, and even non-navigational links are often copied from one page to another within the same host. Similarity is a property concerning different lists, thus addresses inter-list similarity.

## Exploiting Locality

### Use Gap Encoding (as in inverted files)

- $S(x) = (s_1, \dots, s_k)$  will be represented as  
 $(s_1 - x, s_2 - s_1 - 1, \dots, s_k - s_{k-1} - 1)$
- Use of varying length encoding

what does x stand for ?

todo: how to get 3 ? first row

| Node | Outdegree | Successors                                     |
|------|-----------|--|
| ...  | ...       | ...  |
| 15   | 11        | 13, 15, 16, 17, 18, 19, 23, 24, 203, 315, 1034 |
| 16   | 10        | 15, 16, 17, 22, 23, 24, 315, 316, 317, 3041    |
| 17   | 0         |  |
| 18   | 5         | 13, 15, 16, 17, 50                             |
| ...  | ...       | ...  |



| Node | Outdegree | Successors                            |
|------|-----------|---------------------------------------|
| ...  | ...       | ...                                   |
| 15   | 11        | 3, 1, 0, 0, 0, 0, 3, 0, 178, 111, 718 |
| 16   | 10        | 1, 0, 0, 4, 0, 0, 290, 0, 0, 2723     |
| 17   | 0         |                                       |
| 18   | 5         | 9, 1, 0, 0, 32                        |
| ...  | ...       | ...                                   |

Locality can be exploited in a way analogous of how compression of posting lists for text indexing has been performed. Instead of storing the absolute values of the URL identifiers, differences are stored. In other words, we perform **gap encoding**. The resulting differences are then encoded using a varying length compression scheme, such as gamma coding.

## Exploiting Similarity

### Copy data from similar lists

- Reference list: reference to another list
  - Searched in a neighboring window of nodes
- Copy list: indicate which node is copied from reference list
- Extra nodes: additional nodes not in reference list

| Node | Outdegree | Successors                            | Node | Outd. | Ref. | Copy list   | Extra nodes                                    |
|------|-----------|---------------------------------------|------|-------|------|-------------|--|
| ...  | ...       | ...                                   | ...  | ...   | ...  | ...         | ...  |
| 15   | 11        | 3, 1, 0, 0, 0, 0, 3, 0, 178, 111, 718 | 15   | 11    | 0    |             | 13, 15, 16, 17, 18, 19, 23, 24, 203, 315, 1034 |
| 16   | 10        | 11, 0, 0, 4, 0, 0, 290, 0, 0, 2723    | 16   | 10    | 1    | 01110011010 | 22, 316, 317, 3041                             |
| 17   | 0         |                                       | 17   | 0     |      |             |  |
| 18   | 5         | 9, 1, 0, 0, 32                        | 18   | 5     | 3    | 11110000000 | 50   |
| ...  | ...       | ...                                   | ...  | ...   | ...  | ...         | ...  |

i cant use successors if I reference from another list ?

011100...

0 - because i dont copy from 1st element of that list

111 - because i copy from 2nd to 4th element from that list

Result: about 3 bytes / link (with some further compression)

- copy from list before me
- look at previous page for original list

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Link Analysis - 40

For exploiting similarity we apply a similar idea as gap encoding, but now applied to the differences among different adjacency lists. If we consider a reference list (in the example the list of Node 15), subsequent adjacency lists can store a bitlist that indicates which nodes of the reference lists are retained in the subsequent adjacency list, and which are not. 0 indicates that the node in the reference lists is not retained, and 1 indicates that it is retained. Since the subsequent adjacency list can also contain other nodes, that are not stored in the reference list, those extra nodes are stored explicitly.

Candidates for potential reference lists are searched among neighboring lists using a window of predefined size. The choice of the window size is critical, as larger windows increase chances of finding good candidates, but also increase the cost of compression

Together with some further compression applied to the copy lists and the extra nodes, this index compression achieves about 3 Bytes/link cost in the representation of the Web graph.



1. Exploiting locality with gap encoding may increase the size of an adjacency list
2. Exploiting similarity with reference lists may increase the size of an adjacency list
3. Both of the above is true
4. None of the above is true

very small number of links or the difference between is very little. Then the encoding can increase size

if all the list do not have anything or extremely little in common

# Resources

## Course material based on

- Christopher D. Manning, Prabhakar Raghavan and Hinrich Schütze, Introduction to Information Retrieval, Cambridge University Press. 2008 (<http://www-nlp.stanford.edu/IR-book/>)
- Course Information Retrieval by TU Munich (<http://www.cis.lmu.de/~hs/teach/14s/ir/>)

## Relevant articles

- Sergey Brin , Lawrence Page, The anatomy of a large-scale hypertextual Web search engine, Computer Networks and ISDN Systems, v.30 n.1-7, p.107-117, April 1, 1998.
- Jon M. Kleinberg: Authoritative Sources in a Hyperlinked Environment. JACM 46(5): 604-632 (1999)
- Boldi, Paolo, and Sebastiano Vigna. "The webgraph framework I: compression techniques." Proceedings of the 13th international conference on World Wide Web. ACM, 2004.