

# Basic Blocks and Control Flow Graphs

## Basic Blocks

### Introduction:

- A **basic block** is a sequence of consecutive instructions which are always executed in sequence without halt or possibility of branching.
- The basic block does not have any jump statements among them.

# Basic Blocks and Control Flow Graphs

## Basic Blocks

- When the first instruction is executed , all the instructions in the same basic block will be executed in their sequence of appearance without losing the flow control from the program.

Examples



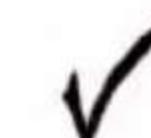
$$a = b + c + d$$

Three address code-

$$t_1 = b + c$$

$$t_2 = t_1 + d$$

$$a = t_2$$



If  $A < B$  then 1 else 0

(1)  $q_E(A < B)$  goto (4)

(2)  $T1 = 0$

(3) goto (5)

(4)  $T1 = 1$

(5)



# Basic Blocks Construction Algorithm

- *Input:* A sequence of three-address statements.
- *Output:* A list of basic blocks with each three-address statement in exactly one block.

# Basic Blocks Construction Algorithm

- 1. Determine the set of *leaders*,**
  - a. The first statement is the leader.
  - b. Any statement that is the target of a conditional or goto is a leader.
  - c. Any statement that immediately follows conditional or goto is a leader.
- 2. For each leader, its basic block consist of the leader and all statements up to but not including the next leader or the end of the program**

# Basic Blocks Construction Example

Consider the following three address code statements .Compute the basic blocks.

- 1) PROD = 0
- 2) I = 1
- 3) T2=addr(A)-4
- 4) T4 =addr(B)-4
- 5) T1 =4\*I
- 6) T3= T2[T1]
- 7) T5 = T4[T1]
- 8) T6 =T3\*T5
- 9) PROD =PROD + T6
- 10) I =I+1
- 11) IF I<= 20 GOTO (5)

# Basic Blocks Construction Example

## Solution:

- Because first statement is a leader ,so -  
**PROD =0 is a leader.**
- Because the target statement of conditional or unconditional goto statement is a leader, so-  
**T1 =4\*I is also a leader**

# Basic Blocks Construction Example

So the given code can be portioned in to 2 blocks as :

B1:

```
PROD = 0  
I = 1  
T2=addr(A)-4  
T4 =addr(B)-4
```

B2:

```
T1 =4*I  
T3= T2[T1]  
T5 = T4[T1]  
T6 =T3*T5  
PROD =PROD + T6  
I =I+1  
IF I<= 20 GOTO B2
```

# Basic Blocks and Control Flow Graphs

## Flow Graphs

### Introduction:

- A *flow graph* is a graphical representation of a sequence of instructions with control flow edges.
- A flow graph can be defined at the intermediate code level or target code level.
- The nodes of flow graphs are the basic blocks and flow-of-control to immediately follow node connected by directed arrow.

# Flow Graphs

## **Points to remember:**

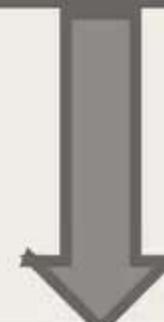
- The basic blocks are the nodes to the flow graph .
- The block whose leader is the first statement is called initial block.
- There is a directed edge from block B1 to B2 if B2 immediately follows B1 in the given sequence
- Then we say that B1 is the predecessor of B2.

# Flow Graphs

The flow graph for the above three address code is given below:

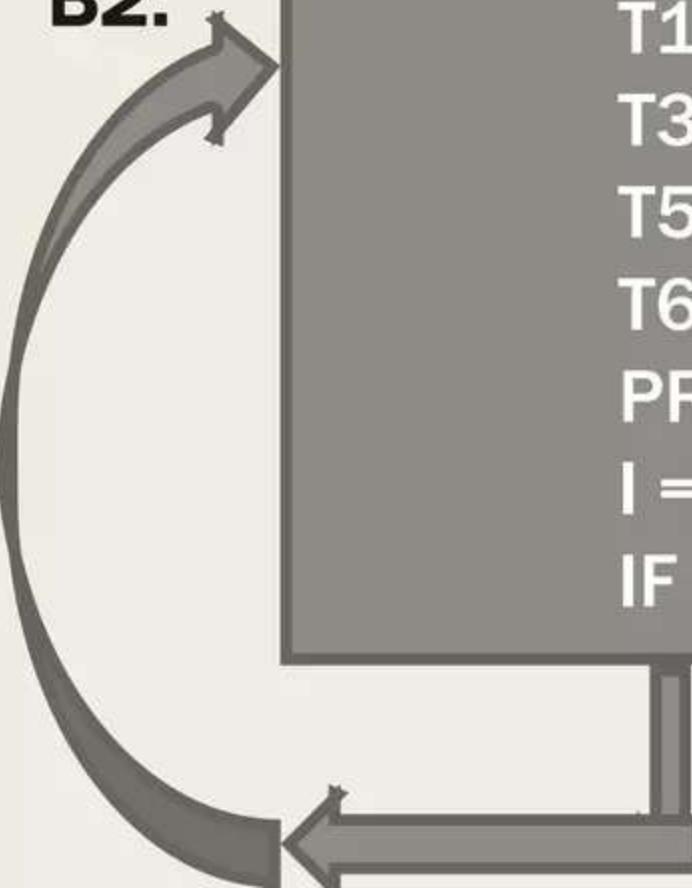
**B1:**

```
PROD = 0  
I = 1  
T2=addr(A)-4  
T4 =addr(B)-4
```



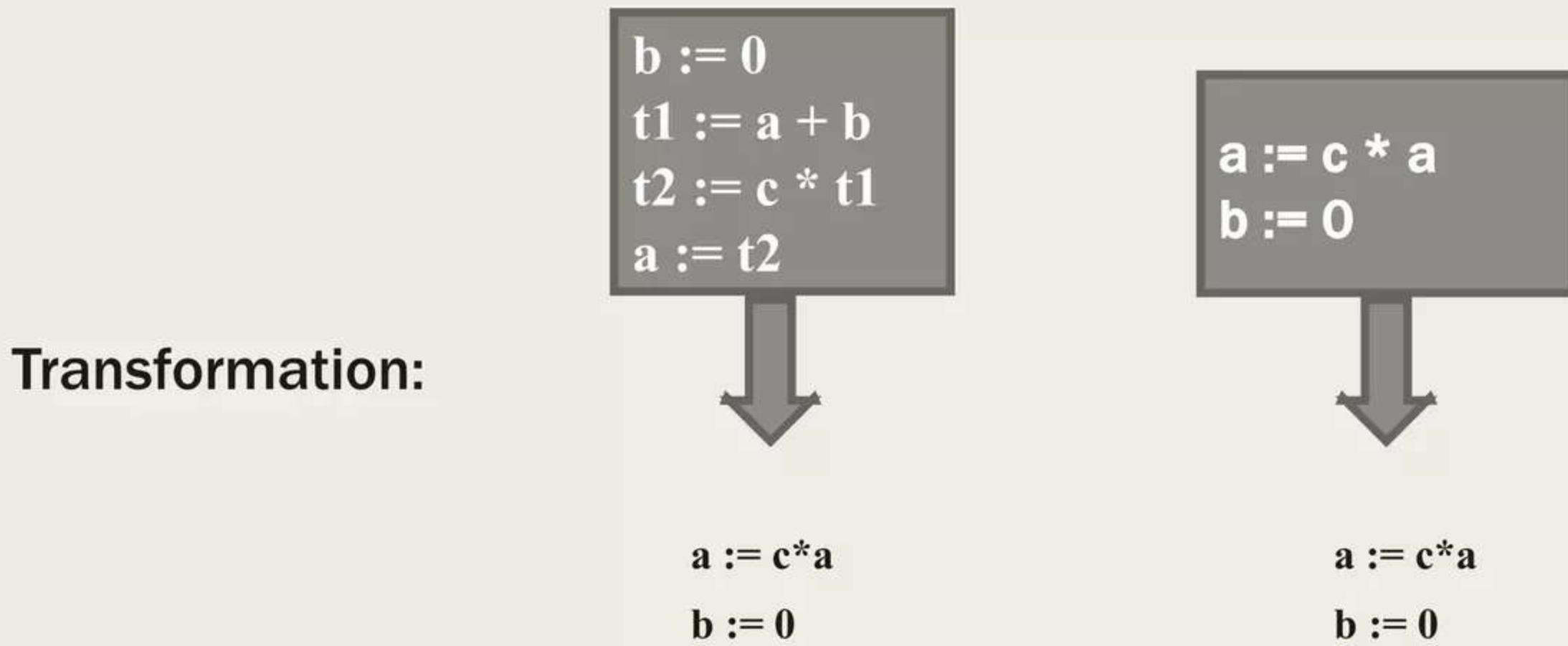
**B2:**

```
T1 =4*I  
T3= T2[T1]  
T5 = T4[T1]  
T6 =T3*T5  
PROD =PROD + T6  
I =I+1  
IF I<= 20 GOTO B2
```



# Equivalence of Basic Blocks

- Two basic blocks are (semantically) *equivalent* if they compute the same set of expressions



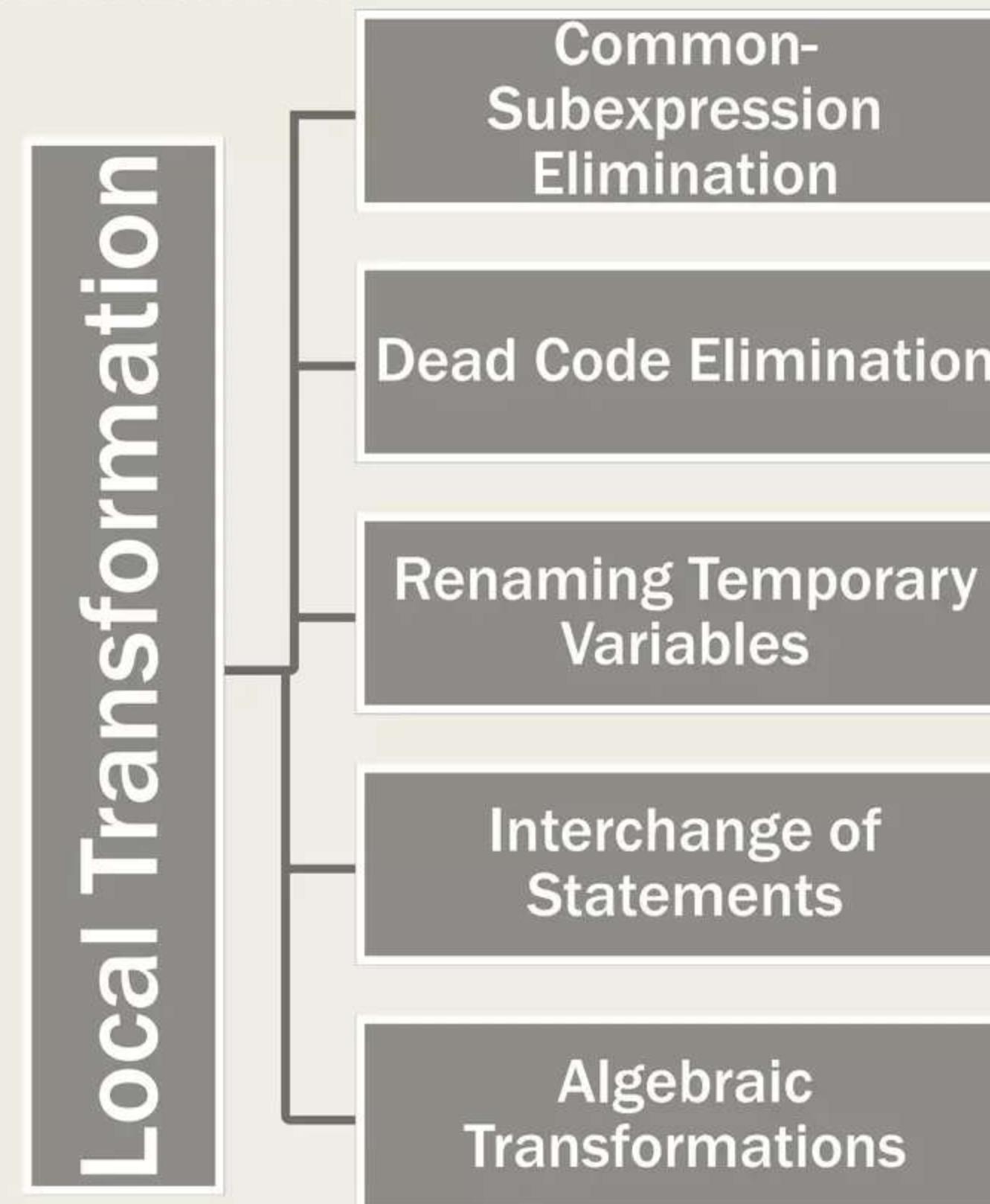
- Blocks are equivalent, assuming **t1** and **t2** are *dead*: no longer used (no longer *live*)

# Transformations on Basic Blocks

- A *code-improving transformation* is a **code optimization** to improve speed or reduce code size.
  - **Global transformations** are performed across basic blocks.
  - **Local transformations** are only performed on single basic blocks.
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- Transformations must be safe and preserve the meaning of the code .
  - A local transformation is safe if the transformed basic block is guaranteed to be equivalent to its original form

# Transformations on Basic Blocks

Some local transformation are:



# Common-Subexpression Elimination

- Remove redundant computations
- 2nd and 4th:compute same expression in fig:1(a)
- Look at 1st and 3<sup>rd</sup> in fig:1(b) :b is redefine in 2<sup>nd</sup> therefore different in 3<sup>rd</sup>, not the same expression

```
a := b + c  
b := a - d  
c := b + c  
d := a - d
```

fig:1(a)

```
a := b + c  
b := a - d  
c := b + c  
d := b
```

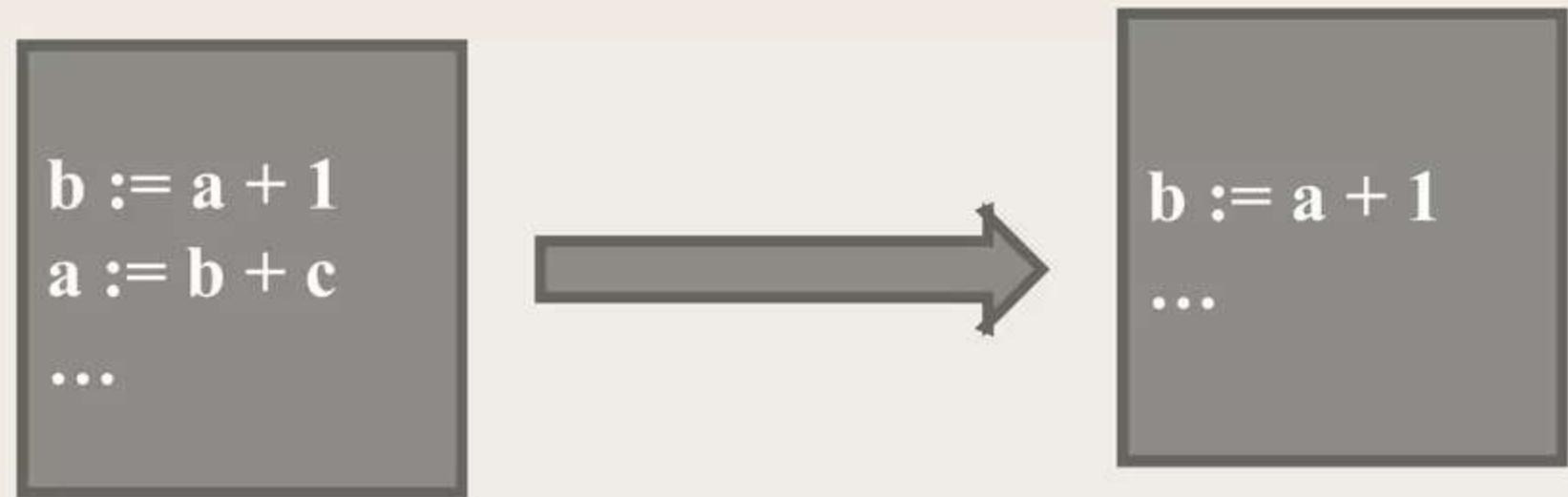
fig:1(b)

```
t1 := b * c  
t2 := a - t1  
t3 := b * c  
t4 := t2 + t3
```

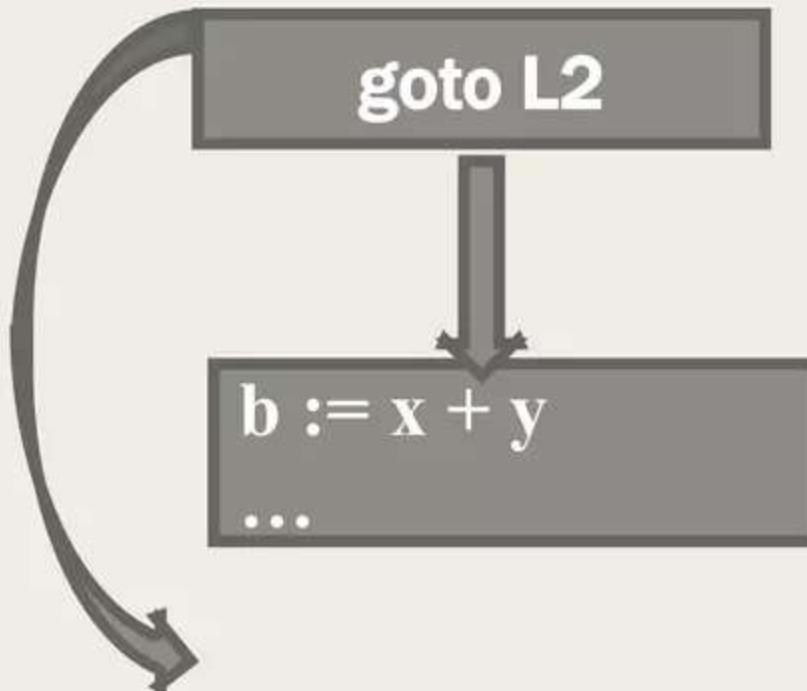
```
t1 := b * c  
t2 := a - t1  
t4 := t2 + t1
```

# Dead Code Elimination

- Remove unused statements



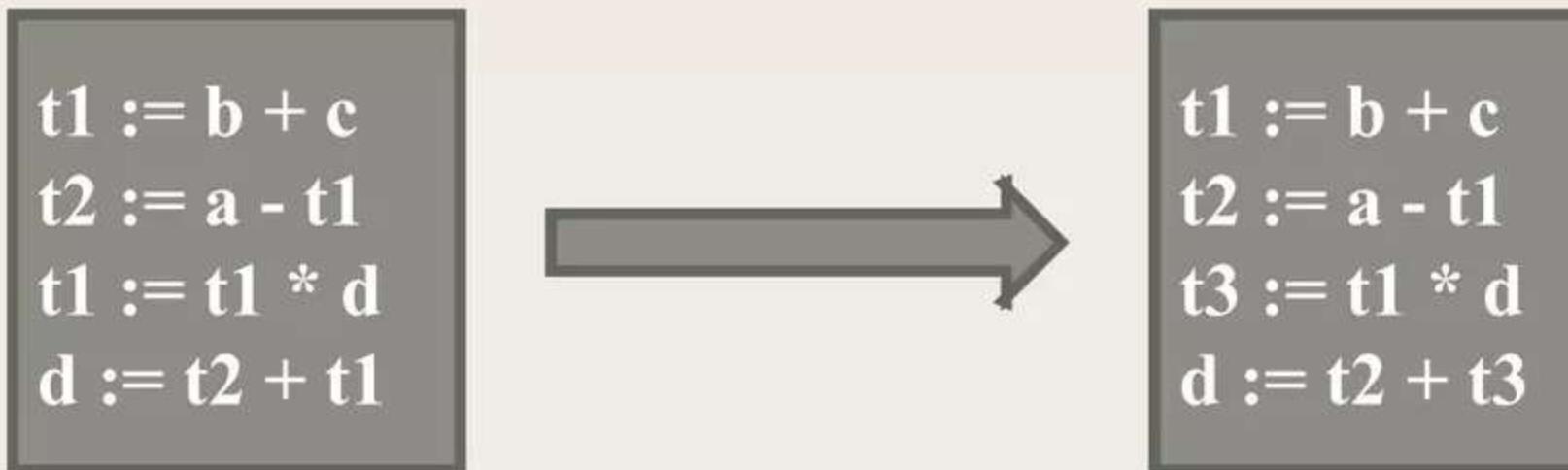
Assuming  $a$  is *dead* (not used)



Remove unreachable code

# Renaming Temporary Variables

- Temporary variables that are dead at the end of a block can be safely renamed.
- The basic block transforms into an equivalent block in which each statement that defines a temporary defines a new temporary.
- Such a basic block is called *normal-form block* or *simple block*.



Normal-form block

- Note that normal-form blocks permit all statement interchanges that are possible.

# Interchange of Statements

- Independent statements can be reordered without effecting the value of block to make its optimal use.

```
t1 := b + c  
t2 := a - t1  
t3 := t1 * d  
d := t2 + t3
```

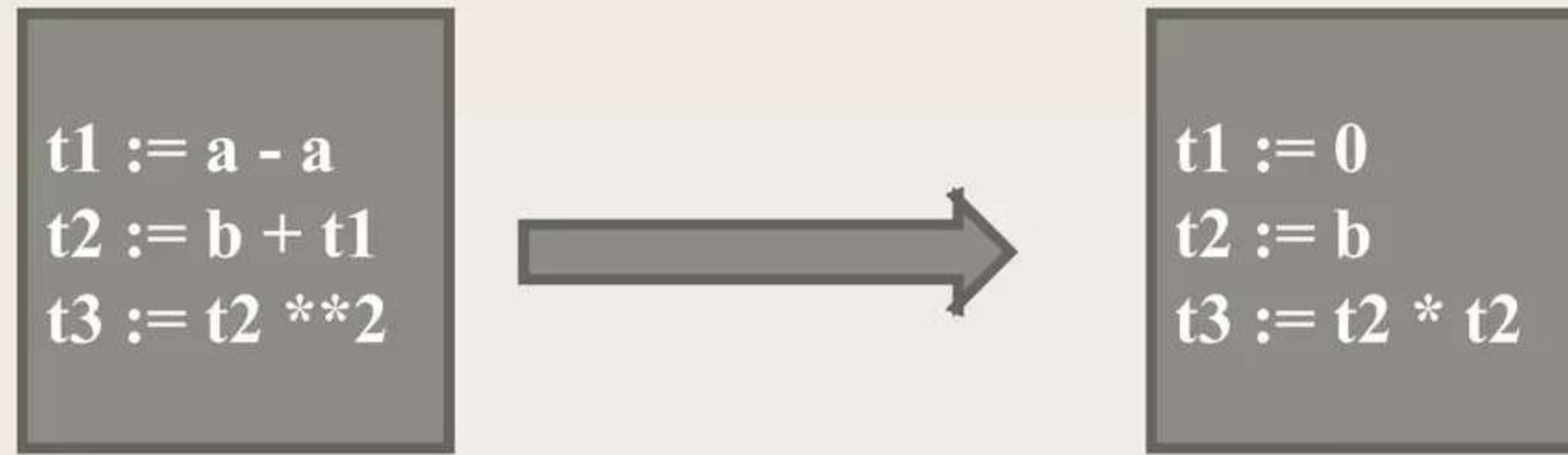


```
t1 := b + c  
t3 := t1 * d  
t2 := a - t1  
d := t2 + t3
```

- Note that normal-form blocks permit all statement interchanges that are possible

# Algebraic Transformations

- Change arithmetic operations to transform blocks to algebraic equivalent forms.
- Simplify expression or replace expensive expressions by cheaper ones.



- In statement 3, usually require a function call
- Transforms to simple and equivalent statement

**!! Queries ??**

**Q & A**

😊Thank You 😊