

Bottom-Up Parsing

Lecture 11

LR Parsing

One Parsing Algorithm

Several Ways to Build the Tables

SLR (or “Simple LR”)

- May fail to build a table for some LR grammars
- SLR Grammars \subset LR Grammars
- Easiest to understand

LR (or “ Canonical LR”)

- The general algorithm
- Will work for any LR Grammar

LALR (or “ Lookahead LR”)

- Will build smaller tables
- May fail for some LR Grammars

SLR Parsing

- An LR(0) state is a set of LR(0) items
- An LR(0) item is a production with a • (dot) in the right-hand side
- Build the LR(0) DFA by
 - *Closure operation* to construct LR(0) items
 - *Goto operation* to determine transitions
- Construct the SLR parsing table from the DFA
- LR parser program uses the SLR parsing table to determine shift/reduce operations

Constructing SLR Parsing Tables

4

Augment the grammar with $S' \rightarrow S$

1. Construct the set $C = \{I_0, I_1, \dots, I_n\}$ of $LR(0)$ items
2. State i is constructed from I_i . Parsing state determined as follows
 - a. If $[A \rightarrow \alpha \bullet a \beta] \in I_i$ and $goto(I_i, a) = I_j$ then set $action[i, a] = \text{shift } j$
 - b. If $[A \rightarrow \alpha \bullet] \in I_i$ then set $action[i, a] = \text{reduce } A \rightarrow \alpha$ for all $a \in FOLLOW(A)$ (apply only if $A \neq S'$)
 - c. If $[S' \rightarrow S \bullet]$ is in I_i then set $action[i, \$] = \text{accept}$
3. If $goto(I_i, A) = I_j$ then set $goto[i, A] = j$
4. Repeat 3-6 until no more entries added
5. The initial state i is the I_i holding item $[S' \rightarrow \bullet S]$

LR(0) Item set

Grammar:

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \text{id}$

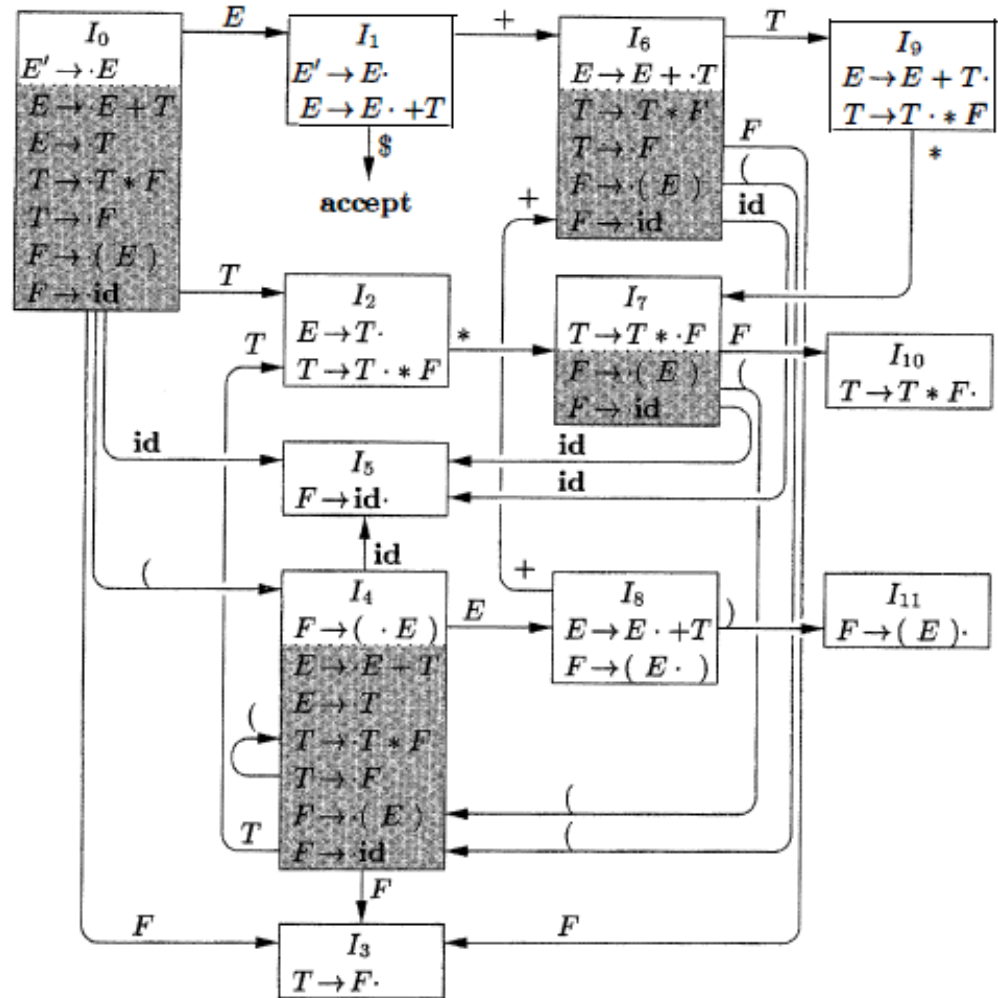


Figure 4.31: LR(0) automaton for the expression grammar (4.1)

$$\begin{array}{l} \textcolor{red}{I}_0 \\ [E' \rightarrow \bullet E] \\ [E \rightarrow \bullet E + T] \\ [E \rightarrow \bullet T] \\ [T \rightarrow \bullet T * F] \\ [T \rightarrow \bullet F] \\ [F \rightarrow \bullet (E)] \\ [F \rightarrow \bullet \text{id}] \end{array}$$

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \mathbf{id}$

Action [0,(]=?

[illegible]

$$\begin{array}{l} \textcolor{red}{I}_0 \\ [E' \rightarrow \bullet E] \\ [E \rightarrow \bullet E + T] \\ [E \rightarrow \bullet T] \\ [T \rightarrow \bullet T * F] \\ [T \rightarrow \bullet F] \\ [F \rightarrow \bullet (E)] \\ [F \rightarrow \bullet \text{id}] \end{array}$$
$$[F \rightarrow \bullet \text{id}] \}$$

Action [0,()]=s4

SLR Parse Table

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \mathbf{id}$

[illegible]

$$I_0$$

SLR Parse Table

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \mathbf{id}$

Action [0,(]=s4
Action [0,id]=s5

[illegible]

$$\begin{array}{l} \textcolor{red}{I}_0 \\ [E' \rightarrow \bullet E] \\ [E \rightarrow \bullet E + T] \\ [E \rightarrow \bullet T] \\ [T \rightarrow \bullet T * F] \\ [T \rightarrow \bullet F] \\ [F \rightarrow \bullet (E)] \\ [F \rightarrow \bullet \text{id}] \end{array}$$

Action [0,()]=s4
Action [0,id]=s5

GoTo(0,E)=1

SLR Parse Table

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \text{id}$

[illegible]

$$\begin{array}{l} \textcolor{red}{I}_0 \\ [E' \rightarrow \bullet E] \\ [E \rightarrow \bullet E + T] \\ [E \rightarrow \bullet T] \\ [T \rightarrow \bullet T * F] \\ [T \rightarrow \bullet F] \\ [F \rightarrow \bullet (E)] \\ [F \rightarrow \bullet \text{id}] \end{array}$$

Action [0,(]=s4
Action [0,id]=s5

GoTo(0,E)=1
GoTo(0,T)=2
GoTo(0,F)=3

SLR Parse Table

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \text{id}$

[illegible]

$$\begin{array}{l} \textcolor{red}{I_1} \\ [E' \rightarrow E \bullet] \\ [E \rightarrow E \bullet + T] \end{array}$$

SLR Parse Table

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \text{id}$

Action [1,\$]=acc

[illegible]

$$\begin{array}{l} \textcolor{red}{I_1} \\ [E' \rightarrow E \bullet] \\ [E \rightarrow E \bullet + T] \end{array}$$

SLR Parse Table

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \text{id}$

Action [1,\$]=acc
Action [1,+]=s6

[illegible]

$$\begin{array}{l} \text{I}_2 \\ [E \rightarrow T \bullet] \\ [T \rightarrow T \bullet * F] \end{array}$$

SLR Parse Table

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \text{id}$

Follow(E)={+,),}\$}

Action [2,\$]=

Action [2,+]=

Action [2,)=r2

[illegible]

$$I_2$$

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow \text{id}$

Follow(E)={+,),\$}

Action [2,\$]=

Action [2,+]=

Action [2,)=r2

Action [2,*]=s7

[illegible]

SLR Parse Table

$$I_4$$
$$E' \rightarrow (\bullet \ E)$$
$$E \rightarrow \bullet E + T$$
$$E \rightarrow \bullet T$$
$$T \rightarrow \bullet \ T \ * \ F$$
$$T \rightarrow \bullet F$$
$$F \rightarrow \bullet (E)$$
$$F \rightarrow \bullet \text{ id}$$

(1) $E \rightarrow E + T$

(2) $E \rightarrow T$

(3) $T \rightarrow T * F$

(4) $T \rightarrow F$

(5) $F \rightarrow (E)$

(6) $F \rightarrow \text{id}$

Action [4,id]=s4

[illegible]

SLR Parse Table

$$I_4$$
$$E' \rightarrow (\bullet \ E)$$
$$E \rightarrow \bullet E + T$$
$$E \rightarrow \bullet T$$
$$T \rightarrow \bullet T * F$$
$$T \rightarrow \bullet F$$
$$F \rightarrow \bullet (E)$$
$$F \rightarrow \bullet \text{ id}$$

(1) $E \rightarrow E + T$

(2) $E \rightarrow T$

(3) $T \rightarrow T * F$

(4) $T \rightarrow F$

(5) $F \rightarrow (E)$

(6) $F \rightarrow \mathbf{id}$

Action [4,id]=s4

Action [4,)=s5

[illegible]

I_{11}

$F \rightarrow (E) \cdot$

SLR Parse Table

- (1) $E \rightarrow E + T$
- (2) $E \rightarrow T$
- (3) $T \rightarrow T * F$
- (4) $T \rightarrow F$
- (5) $F \rightarrow (E)$
- (6) $F \rightarrow id$

$Follow(F) =$
 $Follow(T) = \{*, +,), \$\}$

$Action[8, \$] =$
 $Action[8, *] =$
 $Action[8, +] =$
 $Action[8,)] = r5$

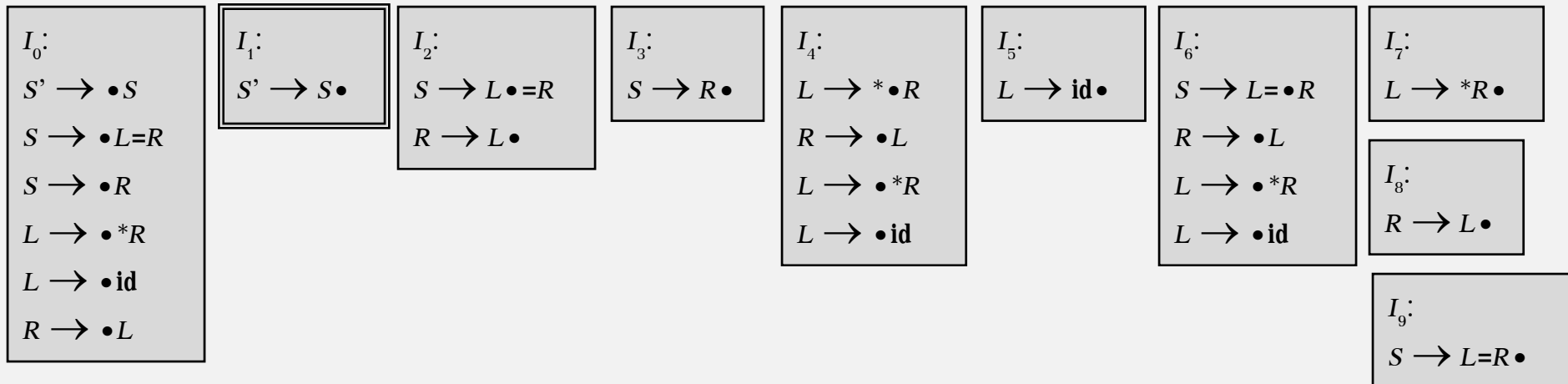
State	Action						GOTO		
	id	+	*	()	\$	E	T	F
0	S5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s4				s5		8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

SLR and Ambiguity

- Every SLR grammar is unambiguous, but **not** every unambiguous grammar is SLR
- Consider for example the unambiguous grammar

$$\begin{aligned} S &\rightarrow L = R \mid R \\ L &\rightarrow * R \mid \text{id} \\ R &\rightarrow L \end{aligned}$$

(1) $S \rightarrow L = R$
 (2) $S \rightarrow R$
 (3) $L \rightarrow * R$
 (4) $L \rightarrow \text{id}$
 (5) $R \rightarrow L$

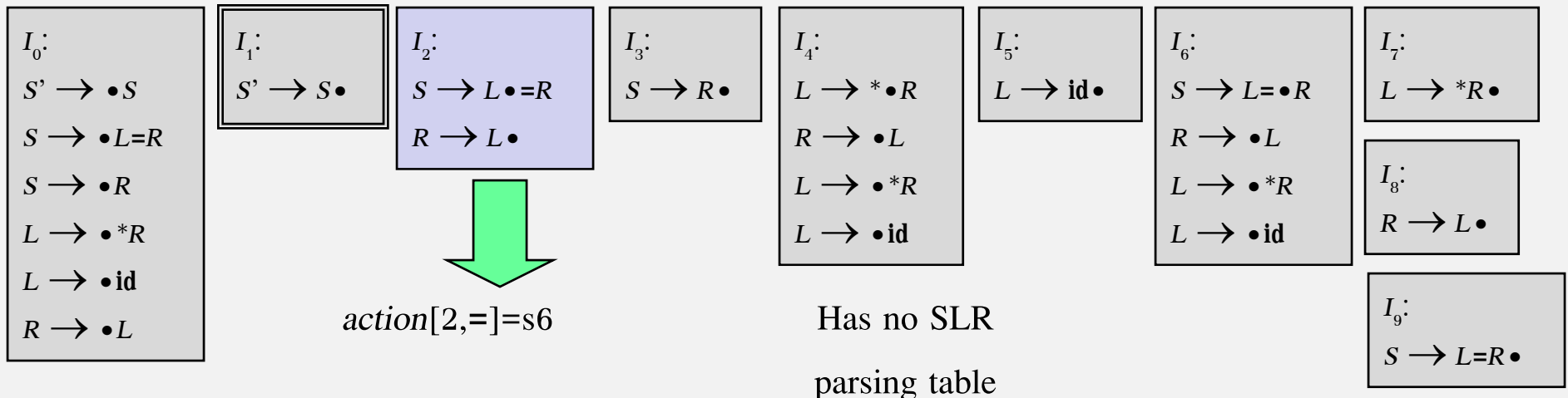


SLR and Ambiguity

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$$\begin{aligned} S &\rightarrow L = R \mid R \\ L &\rightarrow * R \mid \text{id} \\ R &\rightarrow L \end{aligned}$$

(1) $S \rightarrow L = R$
 (2) $S \rightarrow R$
 (3) $L \rightarrow * R$
 (4) $L \rightarrow \text{id}$
 (5) $R \rightarrow L$



SLR and Ambiguity

- Every SLR grammar is unambiguous, but **not** every unambiguous grammar is SLR
- Consider for example the unambiguous grammar

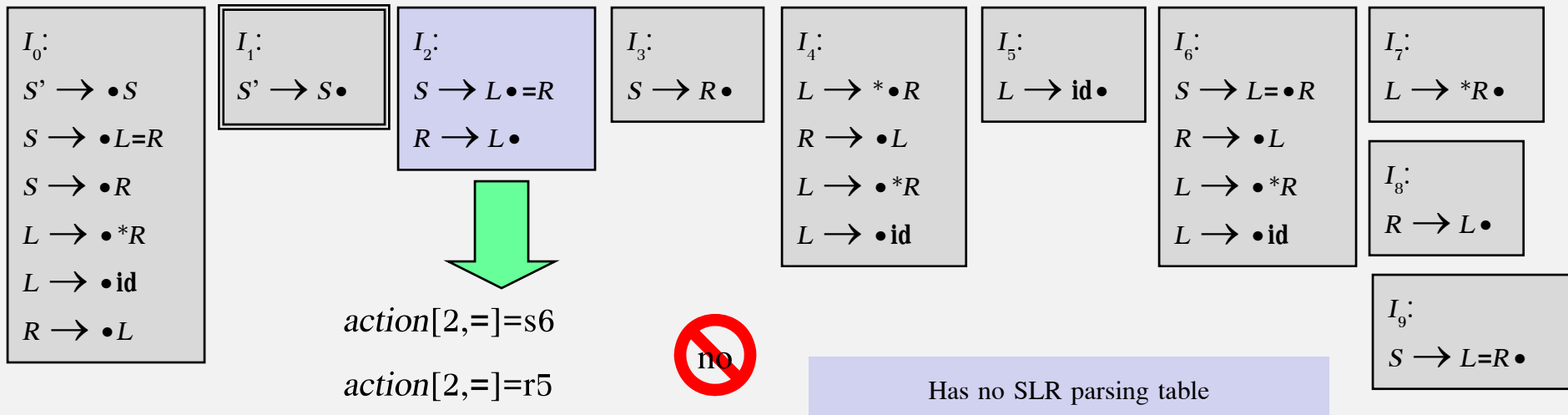
$$S \rightarrow L = R \mid R$$

$$L \rightarrow * R \mid \text{id}$$

$$R \rightarrow L$$

Follow(L) = Follow (R)
 =Follow(S)
 $S \rightarrow L=R$
 $S \rightarrow *R=R$

(1) $S \rightarrow L = R$
 (2) $S \rightarrow R$
 (3) $L \rightarrow * R$
 (4) $L \rightarrow \text{id}$
 (5) $R \rightarrow L$



Thank You