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DISTRIBUTING FIRST-CLASS REACTORS IN A CLUSTER ENVIRONMENT

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June 2024

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Supervisor: Bjarno Oeyen

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Abstract

Reactive programming has gained significant traction for its efficacy in handling asynchronous and event-driven scenarios, particularly in distributed systems. This paradigm emphasizes responsiveness, scalability, and resilience, making it an ideal choice for real-time applications. However, distributing reactive programming introduces challenges such as consistency, coordination, fault tolerance, and communication overhead. This thesis investigates these challenges and proposes a system designed to address them. By leveraging a shared-nothing architecture and deploying an Erlang cluster, we develop a runtime for the purely functional reactive language Haai. This system aims to eliminate known issues in distributed reactive programming, such as the Reactive Thread Hijacking Problem and Reactive/Imperative Impedance Mismatch. We detail the implementation of the Haai virtual machine (Hvm), the deployment of reactors, and the management of distributed environments to enhance scalability, fault tolerance, and performance in distributed reactive systems.

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Chapter 1

Introduction

Reactive programming has been around for a while, gaining traction in various domains due to its ability to handle asynchronous and event-driven scenarios effectively. One of its notable advantages lies in its suitability for distributed systems. By its nature, reactive programming facilitates the development of systems that can easily scale across distributed environments, allowing components to communicate asynchronously and react to events in real-time. However, distributing reactive programming poses several challenges (Bainomugisha et al., 2013) such as consistency, coordination, fault tolerance and many more. These challenges highlight the need for further research and study in understanding how reactive programming principles can be effectively applied and scaled in distributed environments. By investigating these challenges, we can gain insights into best practices and architectural approaches for building robust and scalable distributed reactive systems.

1.1 Reactive programming

Reactive Programming (RP) is a broader programming paradigm that encompasses the idea of reacting to changes and events in a system. It involves modeling the flow of data as streams of events and using declarative and composable abstractions to handle asynchronous and event-driven scenarios. The canonical reactive program model (ref?) is based on time-varying values and propagation of change.

Spreadsheets offer a good example to illustrate this. When a cell containing a formula that includes other cells it is dependent on the cells that are used in that formula. As an example lets take a cell A3 that contains the formula " $=A1*A2$ " whenever the value changes in the cells A1 or A2, A3 will automatically update. We call each spreadsheet cell a time-varying value or signal because their concrete value changes over time due to events. In this example, an event could be when a user changes the value A1 or A2 using the gui of the spreadsheet program.

The two fundamental abstractions defined by RP to represent time-varying values are behaviors and events. Both can be seen as time-varying values but behaviors are said to be continuous over time whilst events are discrete over time. Behaviors have a value at any moment (continuous). Events have a value at specific moments in time (discrete). Listing 1.1 shows the spreadsheet example represented in the RP language Haai (Oeyen et al., 2024). In this example A3 is defined as a behavior that continuously contains the result of $a1 * a2$. For each time that an event occurs updating a1 or a2, A3 will be automatically updated. One could say that the tempo or rhythm of the propagation of change for the behavior A3 depends on the speed that events have a new discrete value.

Listing 1.1: Haai code

```
(defr (A3 a1 a2)
  (out (* a1 a2)))
```

1.2 Problem Statement

In a distributed system where single components or nodes communicate in an asynchronous way RP is an interesting programming paradigm since it deals well with asynchronous data. The distribution of reactive programming (RP) introduces a myriad of complexities and challenges. In this work we would like to focus on the following two challenges. Developing a declarative specification to deploy the distributed reactive program and simulating a virtual general clock between the individual nodes by temporizing the occurrence of events or discrete time-varying values.

1.3 Overview

The Structure of this thesis starts with a situation of (distributed) reactive programming. After explaining the strategy of the thesis we explain purely reactive programming in section 4. More detail on the runtime is found in chapter 5. Chapter 6 explains the distribution of the runtime and the tools used to build it. Before ending with the conclusion we have chapter 7 with the evaluation of the distribution and the use case of the distributed reactive melody generator.

Chapter 2

Reactive programming

Reactive programming (RP) and reactive streams (RS) are terms that are often used interchangeably. Both adhere to the same canonical model, explained in section 1.1, showing there similarities. Two fundamental differences instead are the way dependencies between signals are established and how updates to signals are orchestrated.

2.0.1 Signal Dependencies

The way signal dependencies are established in the source code differs in the following way. In RP languages these dependencies are automatically taken care of by the language. Programmers implement there desired functionality in RP similar as they would in a non-reactive language. In RS the dependencies are explicitly programmed by the programmer by making use of operators provided by the RS library. The programmer creates a linear application of those operators often called a pipe as shown in listing 2.1 where an observable exiting of the numbers 1 to 10 will be the source of the pipe. The second or last operator in the pipe is Sum, who will produce one value to be printed on screen.

2.0.2 Signal Orchestration

The operators used in RS define independently for each signal how these will be updated. In contrary to RP where it is again the system that updates signals in a global way. As can be seen in listing 1.1, no special code defines how or in what order the values a1 and a2 need to be updated. Both approaches can fail and result in a glitch, if for example the RP language updates signals in the wrong order, dependent signals end up having the wrong value, this is called a glitch. On the other hand in RS, if the programmer does not exactly know what an operator does and uses it in a pipe or collection of operators, inconsistencies might arise and also produce a glitch.

The responsibility to handle signal dependencies and orchestration correct is at the system side for RP and in the developers hand for RS.

Listing 2.1: Python, RxPy library

```
rx.of(1,2,3,4,5,6,7,8,9,10).pipe(  
    operators.filter(lambda i: i %2 == 0),  
    operators.sum()  
) .subscribe(lambda x: print("Value is {0}".format(x)))
```

2.1 State of the art

In this section and the rest of this thesis we will focus on reactive programming. Reactive programming has many applications, for example Fran (Elliott and Hudak, 1997) the first reactive programming language, originated to compose richly interactive multimedia animations based on behaviors and events. ReactiFi (Sterz et al., 2021) a high-level reactive programming language to program Wi-Fi chips on mobile consumer devices without expert knowledge of Wi-Fi chips. Or the domain-specific language Yampa (Courtney et al., 2003) embedded in Haskell for programming hybrid (mixed discrete-time and continuous-time) systems. Most of the reactive programming languages are built on top of a sequential host language. Fran and Yampa are built on Haskell, ReactiFi is built on Scala and C. Host languages provide functions that can be called from the reactive programming language, this is called lifting the function from the host language into the RP language.

All of the mentioned frameworks and libraries have the downside of being composed of two components. The reactive system is built on top of an imperative language. This makes them prone to the *Reactive Thread Hijacking Problem* and the *Reactive/Imperative Impedance Mismatch* problems (Vonder et al., 2020).

2.2 Distributed reactive programming

When distributing a reactive program, issues related to distributed systems come apparent in the distributed reactive program (drp). Resolving those issues becomes part of the distributed reactive program and the distributed system architecture. In this thesis we focus on the following challenges.

Consistency and Coordination: In distributed environments, maintaining consistency across different nodes can be challenging. Reactive programming relies on the propagation of changes and events, and ensuring that all nodes in a distributed system are in sync requires careful coordination and management of data flows. Different nodes may receive events in different orders, leading to potential inconsistencies that need to be addressed through synchronization mechanisms and event ordering techniques.

Fault Tolerance and Resilience: Distributed systems are prone to failures, whether due to network partitions, node crashes, or other issues. Reactive programming emphasizes responsiveness and fault tolerance, but distributing these principles across a network introduces new challenges in handling failures gracefully and ensuring system resilience.

Network Communication Overhead: Asynchronous communication between distributed components in a reactive system can introduce overhead in terms of network latency and message passing. Efficiently managing this communication overhead while maintaining responsiveness is crucial for the performance of distributed reactive systems.

In the literature inconsistencies in the data are called glitches. A glitch is a momentary inconsistency in a time-varying value. For example in a financial application we do not want to read inconsistent values. We want to read the new value after the update on the values this new value depends on has been finished. Being free of glitches is called glitch-freedom. According to Margara and Salvaneschi (Margara and Salvaneschi, 2018) there are two types of glitch-freedom. A 'single-source' and 'complete' way for the program to be free of glitches. In the single-source case we assume one node, all pending updates to that node are performed before the updated values are available. Complete glitch freedom is considered when we have more than one node. To

have complete glitch freedom we must make sure that updates occurred in the same order for all nodes considered by this update.

Chapter 3

Strategy and Methodology

Can we use the tempo of events arriving at the source to steer the global update orchestrator (introduce guo)?

Our approach in this work is to remove as many of the issues known to rp and drp by developing a system that by design solves the challenges known to drp. The distributed system is build as a shared nothing architecture (sna) and hosts an Erlang cluster. We have chosen this architecture for basic simplicity as it allows to easily scale and monitor fault tolerance since each node is independent and self-sufficient.

To deploy the system we make use of code mobility, allowing us to declaratively define which rp needs to run on what node or group of nodes. We have build a runtime for the purely functional reactive language Haai on top of the sna. Each node represents a runtime towards which a (master) node can send the code to be deployed.

All the different nodes or runtimes in the cluster have a rhythmical organization (eg: the loop time of a reactor is a fixed value in milliseconds). The 'tempo' of each runtime will allow us to monitor the cluster for failures of nodes in the sna.

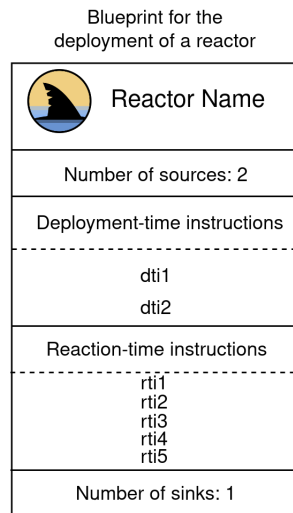
Chapter 4

Purely reactive programming

4.1 Reactors

In Haai, a purely reactive computer language, the only construct available for expressing computations is a reactor. A reactor is a sort of blue print that will be used to deploy any number of deployments based on that reactor. We can write a single reactor or combine multiple reactors in one 'main' reactor or program. The main program is build out of reactors, these can be native or user defined. Every computation in a reactor is done by a reactor, some are native to the runtime others are user defined. For example the runtime could provide native reactors for basic arithmetic. 4.1 show the concept of a reactor.

Figure 4.1: Graphical representation of a reactor



The deployment of a reactor reacts on a data stream, that data altered by the reaction is found in the sink of the reactor. The deployment of a reactor is explained in the next chapter.

4.2 Haai

We are making use of the novel purely reactive language Haai. Haai is reactive all-the-way-through. As such the only construct available for expressing computations is a reactor. A reactor has sources or inputs and sinks or outputs. For each time the sources are updated the reactor will react and produce a new sink value. The number of sources and sinks are variable per reactor definition. In Haai there are no functions, only reactors. This allows us to build reactive programs without functions (Oeyen et al., 2024). Different from other rp is that Haai exists out of only one layer. By design Haai solves the Reactive thread hijacking problem and the reactive/imperative impedance mismatch.

Levels of reactivity

There is a trade off where more strict levels of reactivity allow for a less wide pallet of reactive programs to be expressed. For Haai the level of reactivity will be partly defined by the computational time of the native reactors. In this implementation of the Haai virtual machine we only have native reactors that have a computational complexity of $O(1)$. This guarantees a strong level of reactivity (Oeyen et al., 2024).

Code example

This Haai code in listing 4.1 represents three user defined reactors: consonance, duration and main. The main reactor is build out of the two user defined reactors consonance and duration and makes no use of native reactors. Consonance and duration make use of multiple native reactors to perform arithmetic. The arguments to the reactor are the sources and the keyword 'out' defines the sink of the reactor.

Listing 4.1: Haai code

```
(defr (consonance f)
  (out (* ci f)))

(defr (duration bpm)
  (def q (/ 60000 bpm))
  (out (* lm q)))

(defr (main f bpm)
  (out (consonance f))
  (out (note_length bpm)))
```

4.3 Remus instruction set

The Remus instruction set is the set of low-level operations that the virtual machine can execute. A Haai program or Reactor will after compilation exist out of a combination of instructions. Instructions form the Remus instruction set as listed here under. The Haai virtual machine is an implementation of this instruction set in Elixir. The following example program is build up out of a combination of the Remus instruction set. It is the compilation of the main reactor shown in the code snippet above.

Instruction	Arguments	Description
I-ALLOCMONO	[reactor]	Allocate memory for the given reactor.
I-LOOKUP	[signal]	Lookup the value for signal
I-SUPPLY	[from, destination, index]	Move a value.
I-REACT	[memory_location]	Apply the reaction found in memory.
I-CONSUME	[memory_location, index]	Position the result of react in run-time-memory.
I-SINK	[rti_location, sink_index]	Position a value from run-time-memory into sink.

Table 4.1: Remus instruction Set

4.4 Example bytecode program

After parsing the bytecode that was compiled from the original Haai code in listing 4.1 with the Haai compiler into an Elixir readable format of nested lists a program could look like listing 4.2. In this example we see two user defined reactors named `consonance` and `note.length` followed by the main reactor that represents the actual program that will be deployed.

Listing 4.2: Parsed remus bytecode code

```
[
  [:consonance, 1, 1,
    [
      ["I-ALLOCMONO", :multiply]
    ],
    [
      ["I-LOOKUP", :ci],
      ["I-SUPPLY", [%RREF, 1], [%DREF, 1], 1],
      ["I-SUPPLY", [%SRC, 1], [%DREF, 1], 2],
      ["I-REACT", [%DREF, 1]],
      ["I-CONSUME", [%DREF, 1], 1],
      ["I-SINK", [%RREF, 5], 1]
    ],
  ],
  [:duration, 1, 1,
    [
      ["I-ALLOCMONO", :divide],
      ["I-ALLOCMONO", :multiply]
    ],
    [
      ["I-SUPPLY", 60000, [%DREF, 1], 1],
      ["I-SUPPLY", [%SRC, 1], [%DREF, 1], 2],
      ["I-REACT", [%DREF, 1]],
      ["I-LOOKUP", :lm],
      ["I-SUPPLY", [%RREF, 4], [%DREF, 2], 1],
      ["I-CONSUME", [%DREF, 1], 1],
      ["I-SUPPLY", [%RREF, 6], [%DREF, 2], 2],
    ],
  ],
]
```

```

        ["I-REACT",["%DREF",2]],
        ["I-CONSUME",["%DREF",2],1],
        ["I-SINK",["%RREF",9],1]]
    ],
    [:main, 2, 2,
    [
        ["I-ALLOCMONO", :consonance],
        ["I-ALLOCMONO", :note_length]
    ],
    [
        ["I-SUPPLY",["%SRC",1],["%DREF",1],1],
        ["I-REACT",["%DREF",1]],
        ["I-SUPPLY",["%SRC",2],["%DREF",2],1],
        ["I-REACT",["%DREF",2]],
        ["I-CONSUME",["%DREF",1],1],
        ["I-SINK",["%RREF",5],1],
        ["I-CONSUME",["%DREF",2],1],
        ["I-SINK",["%RREF",7],2]
    ]
    ]
]

```

Chapter 5

Haai virtual machine

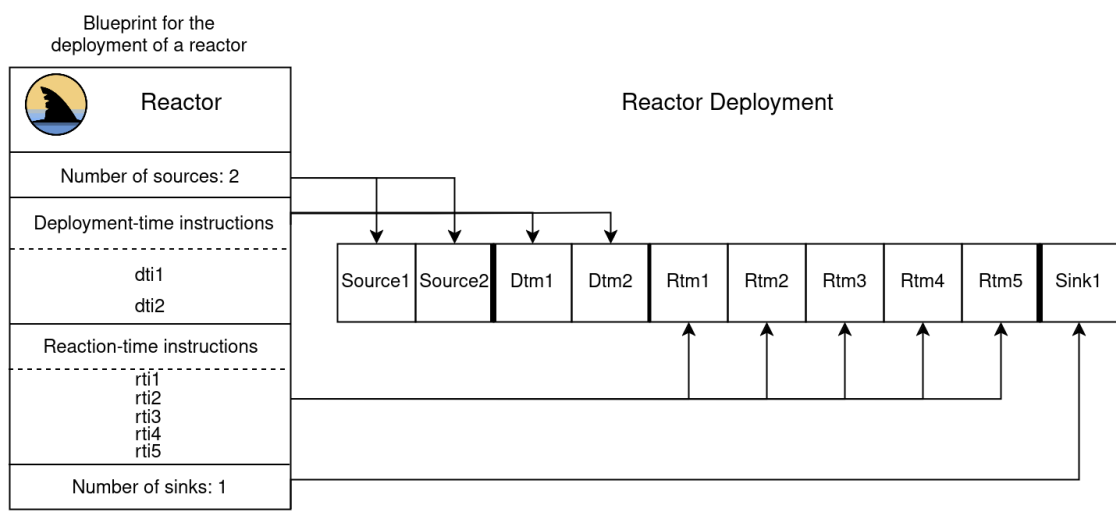
To make it possible for Haai code to work on multiple platforms a virtual machine for each required platform is necessary. After the example of Remus, "a virtual machine that has been carefully designed to be usable for running reactive programs in low-powered computing environments." (Oeyen et al., 2022) We implement in this thesis a virtual machine for Haai. The purpose of the Haai virtual machine (Hvm) is to enable the possibility of distributing Haai programs in a cluster of machines. There are multiple reasons to distribute a reactive program like: scalability, low latency and many others. Firstly the Hvm is a prototype to investigate distribution in general and to see if at all the language works well in a distributed setting. The Hvm is developed in the Elixir computer language. The chapter on distribution arguments our choice for Elixir. The Haai bytecode as compiled for Remus will be the representation of Haai reactors in the Hvm.

5.1 Deploying a reactor

A reactor can be seen as a blue print for its deployment, similar how an object created from a class in object-oriented programming is an instantiation of the blueprint that the class represents. The deployment of the reactor will request a certain size of memory. That total block of memory exists out of four parts: Sources, deployment time memory (dtm), reaction time memory (rtm) and the sinks. When we go from a reactor to a deployment we can find the following information in the reactor or blue print for the deployment. The number of sources and sinks, the deployment time instructions (dti) and the reaction time instructions (rti). For each source and each sink a position is reserved in the memory block we are constructing. From the deployment time instructions in the reactor we can prepare the deployment time memory (dtm). The Dtm part of the memoryblock we are constructing will reserve space for the deployment of the reactors found in the dti.

Sofar the space for the sources, the sinks and dtm has been reserved, the last step is to reserve space for the reaction time memory (rtm). The space required for rtm is equal to the number of instructions found in the reaction time instructions. For each instruction in rti there is one block reserved in rtm. A deployed reactor uses the four memory parts: sources, dtm, rtm and sinks to react according to the reaction time instructions found in rti. Figure 5.1 expresses the deployment of a reactor.

Figure 5.1: Reactor deployment, from blueprint to memory block



5.2 Running the reactor

Running the reactor encompasses starting a Haai virtual machine with all requested arguments for the Hvm to function. The Haai virtual machine needs the following arguments to start:

Reactor bytecode Bytecode obtained from the Haai compiler (eg: listing 4.2).

Connect source A connector function that the Hvm can use to connect some input to the sources of the reactor. At each start of the next iteration this function will be called to update the sources of the deployed reactor. For example connecting to a WebSocket server to receive the data it broadcasts.

handle sink A sink function, that after each iteration is called to handle the new sinks produced by the deployed reactor. This function is tailored for the specific reactor or program that has been send as bytecode.

5.3 Managing the reactors state

Managing the state of a reactor is done by one actor per program or reactor deployment. The actor is build around the GenServer module in Elixir which is a behavior module that provides a generic server implementation. In this implementation the name of the GenServer is Memory. Memory is a single actor that encapsulates multiple functions to respond to various messages or calls. Together these define the behaviour of the Memory GenServer.

Memory will mainly accept calls from the runtime instructions found in the reactor deployment. Available calls are mapped on the Haai instruction set. In addition there are calls to initialize or to get or set sources and sinks into or from the Memory GenServer. One instance of Memory represents the memory block of a deployed reactor.

The GenServer module in Elixir works with calls or casts, we use calls since they are synchronous requests to the GenServer. By calling Memory in a synchronous way we make sure

that the values in Memory are only accessed by one operation at a time, assuring correct values.

5.4 Native reactors

The Haai virtual machine makes use of native reactors to perform simple arithmetic operations like addition, division, multiplication and subtraction. Native reactors are recognized by the Hvm and applied when required.

Haai is a one layer reactive programming language, contrary to two layer system where the top reactive layer can lift functions from the bottom imperative layer, Haai only uses reactors for all computation. The Hvm has a collection of frequent reactors. Thanks to the native reactors we do not need to implement the most basic reactors when composing a program. By design Haai prevents the Reactive/imperative impedance mismatch as explained in the chapter on Haai. Having the native reactors smartly covers the lack of lifting possibilities.

Chapter 6

Distribution

6.1 Distributing the Hvm

The Elixir cluster consists of a number of Elixir nodes. An idle Elixir node is not resource-intensive, making it feasible to run tenfold of nodes on a single machine to simulate a cluster. Once multiple nodes are operational, one node will be selected as the controller or master node. The master node will be made aware of all nodes within the cluster. With all nodes connected to the master node, but not directly to each other, we can commence the deployment of reactors.

Deploying a reactor on the cluster involves initiating the Hvm on the selected node with the bytecode of the requested reactor. This allows each node to run either the same reactor or different reactors, depending on the bytecode it receives. The ability to send the actual reactor to a node enables us to dynamically replace a node that has gone down or to balance the load by adjusting the number of nodes running a particular reactor. Code mobility operates from the master node to any of the nodes in the cluster.

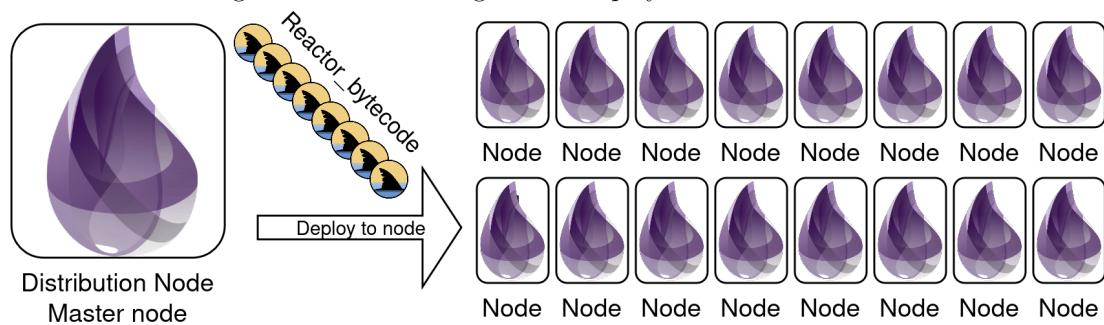
Once the cluster is established, with one master connected to all other nodes, the deployment process can be summarized as follows, chose a node in the cluster, a reactor bytecode, functions to connect the sources and handle the sinks. The reactor or our reactive program is started and will keep on reacting for as long as the node exist in this cluster. This setup ensures that the reactive program can respond to changes and maintain operations effectively across the distributed system. Figure 6.1 shows the concept of deploying the reactors from the master node into a chosen number of nodes. One reactor can be deployed to an 'unlimited' number of nodes. Each deployment will be an individual instance of the same reactor send towards the runtime on the node.

Once the deployments are active each deployment will start sending messages towards the monitor of the cluster. The monitor is aware of the duration each single deployment uses to operate. Comparing the actual timing with the timings in the monitors directory we can asses about the individual deployments. For example if the actual timing is slower then the recorded timings in the monitor we might want to liberate more resources. The differences in timing are the starting poit of an investigation into what might cause this difference.

6.2 Distributed environment

We explain here in more detail the details of the system on which the Hvm is running.....?

Figure 6.1: Distributing reactor deployment on Elixir cluster



6.2.1 Shared-nothing architecture

To organize the distribution of the Haai virtual machine we have opted for a shared-nothing architecture. (Stonebraker, 1986) This allows us to build the cluster with the following characteristics, suitable for our research in distributing the reactive language Haai.

Independent Nodes Each node can execute its own specific code and perform computations independently since no memory, storage or processing power is shared.

No shared state Each node manages its state independently. Cluster communication between nodes is done with messages. we use Erlang distributed mechanisms in the implementation.

Horizontal scalability Shared-nothing architectures are highly scalable horizontally, this makes it straightforward to add or remove nodes from the cluster. It allows us to run clusters of all sizes. Practically the cluster size will be bound to the available hardware.

Fault isolation Faults or failures in one node do not affect the operation of other nodes in the system. We can monitor node failure with Erlang distributed mechanisms and act accordingly.

Flexibility and Modularity Changes to one node do not impact the operation of other nodes, allowing for maintenance, upgrades, and system evolution of the cluster.'

6.2.2 Erlang virtual machine

We implemented the shared-nothing architecture in Elixir a functional, concurrent, and fault-tolerant programming language built on the Erlang Virtual Machine (BEAM). The Erlang Virtual Machine provides a robust platform due to several key features and design principles:

Concurrency model Erlang has a lightweight concurrency model which is based on actor-like processes. Since each process runs independently and can communicate with other processes through message passing the Erlang virtual machine is inherently decentralized by design.

Fault tolerance The BEAM VM provides mechanisms for process supervision, fault isolation, and hot code swapping, allowing systems to recover from failures without affecting the overall operation of the cluster.

Distribution transparency Erlang’s distribution features enable processes to communicate effortlessly across network boundaries. Within an Erlang cluster, nodes can seamlessly interact through remote procedure calls (RPCs) and message passing. Different types of distribution models can be build on the Erlang platform.

Scalability Horizontal scaling or adding nodes to a cluster is straightforward. This makes it easy to scale clusters to need.

Tooling Common tasks are simplified thanks to a comprehensive tool set. tools like the Mix build tool and the OTP framework are core tools used in the implementation of the distributed Haai cluster.

The example program is explained in more detail in the use case chapter. The same program or reactor can be run many times. The distributed setup we make exists out of a number of elixir nodes that all send there messages to one sound server. The bigger the distributed network is the more messages that are send to the sound server. The sound server it self has virtually no limit to receive messages. The main parts of the distribution are taken care of by the underlying Beam virtual machine. One main reason to develop the Hvm on this platform is to be able to distribute in a solid and proven way.

6.2.3 Actor model

The actor system in Erlang implements the processes actor model. (De Koster et al., 2016) These actors are modeled as processes and can define there receive primitive to specify what messages it responds to. Thanks to the isolated turn principle the actor model guarantees to be free of data races and deadlocks by design. In Erlang actors are typically used when you need to encapsulate state and behavior within a single unit of computation., actors are spawned using constructs like GenServer, Task, or Agent within a single node. All actors communicate with each other using message passing within the same node. One can use these actors for modeling concurrent and stateful components in an application.

6.2.4 Distributed Erlang mechanisms

Distributed Erlang mechanisms come into play when you need to communicate between actors residing on different nodes within a cluster. Nodes communicate with each other using distributed Erlang protocols, allowing actors on different nodes to exchange messages. Distributed Erlang mechanisms provide fault tolerance, distribution transparency, and location transparency across the cluster.

Distributed Erlang provides location transparency, allowing actors to communicate regardless of their physical location within the cluster. This transparency hides the underlying network details, promoting loose coupling (Carreton et al., 2010) between nodes and actors. Actors on different nodes can interact seamlessly through message passing, without needing to be aware of the specific nodes they are communicating with. Loose coupling in distributed Erlang enables the system to scale horizontally by adding or removing nodes without significantly impacting the overall architecture.

Chapter 7

Evaluation

7.1 Evaluating the distribution model

The master-slave model used in our distribution works well. We have a cluster of actors, with basically two different types of actors. The master actor and the slave actors. This model allows us to handle issues with the slaves in a timely fashion by reacting as soon as inconsistencies between the recorded and live timings exist in the cluster monitor. The challenges mentioned in section 2.2 are controllable within this model. We can guarantee glitch freedom on the individual basis of a node. Each individual node guarantees glitch freedom and in our master-slave model where we have no values exchanged between the nodes we do not expect other sources of inconsistency.

The individual nature of the nodes in the cluster guarantees that when one node goes down no other nodes are impacted by this. The monitor of the cluster can based on timing measurements decide on a tactic to solve the possible failure of a node. For example if a node actual time is not updating the node must have trouble. we can spin up a node to replicate or solve the actual problem the initial node is having. the organization of the applicable tactics the monitor can perform are not part of this work but provide ... for future work.

Working with the cluster monitor and the individual nodes requires only a small amount of network traffic. The cluster monitor will receive one message on a per duration time tempo for the deployed reactor. In the case of an issue the second process of the cluster monitor needs only to communicate with one single (new) node. The shared nothing architecture makes it simple to add or remove individual nodes.

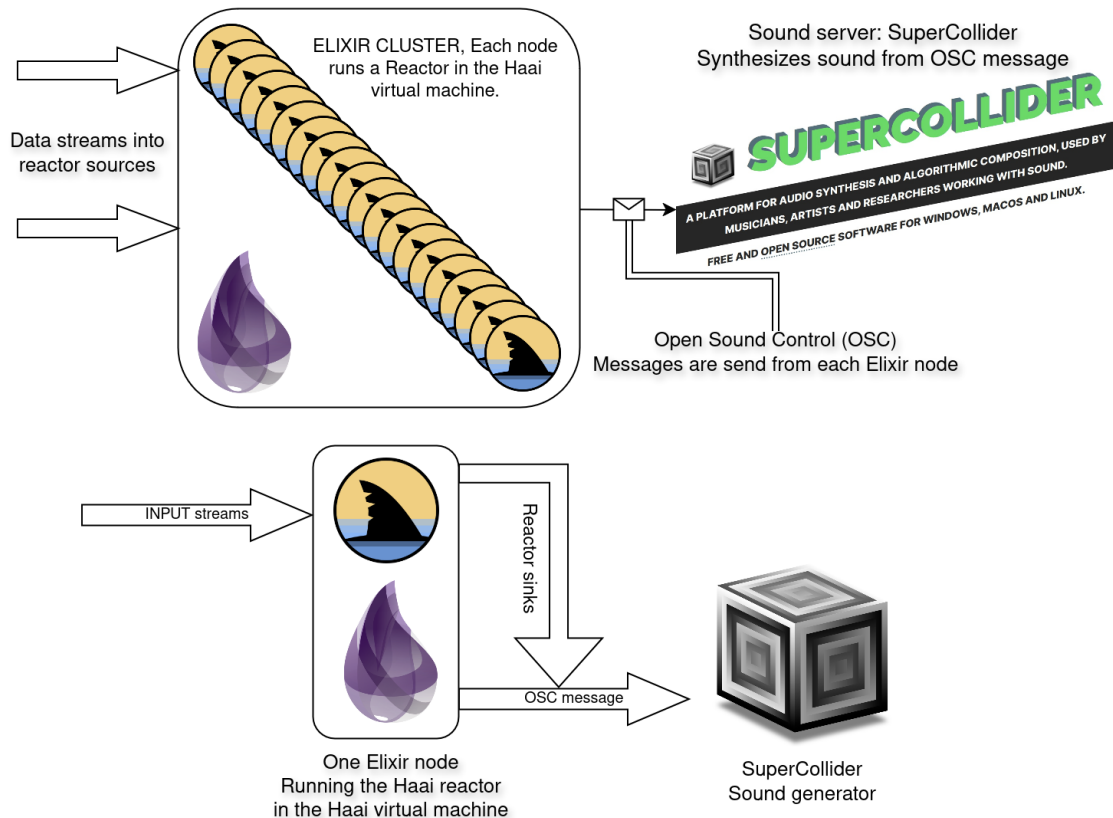
The number of nodes that simulate the cluster can easily be a tenfold. The reactors itself have a very light computational cost existing out of some basic arithmetic operations. One Elixir node would consumes about 90MB of memory and virtually no processor time when idle and about 3% when performing the calculation. This allowed to simulate a cluster on one machine with up to a hundred nodes.

The next section explains the Distributed reactive melody generator (drmg) figuring here as a stand in for the cluster monitor. The drmg receives for each deployment a message on a predefined tempo per deployment. What we do here is having a synthesizer play a note equal to the duration of the deployment tempo for each node. All these single notes come together to form a melody.

7.2 Distributed reactive melody generator

The distributed Haai cluster, existing out of many distributed reactors provides the sink values form the reactors as values inserted into OSC messages (Schmeder et al., 2010) that are send to a sound server to trigger the synthesis of some sound. It is the reactors that, by reacting on incoming streams of data, produce the values that result in a rhythmical flow of OSC messages to the sound server.

Figure 7.1: Distriuted Reactive Melody Generator



The flow of rhythmical messages is related to the duration of each sound. That same duration that makes the sound disappear after it was started determines the iteration speed for the reactor or program that runs on each individual node. Each reactor iterates over and over but each iteration has a defined duration, this duration is virtually equal to the duration for a sound calculated by the reactor. A reactor will calculate the new values, sleep for the calculated duration and then provide the new values in the sink. At this point the Hvm calls a function that it was given to handle the sink values. For this use case that function constructs the message and sends it to the sound server that propagates it to the requested synthesizer.

7.2.1 Sound server

The synthesizer that produces the sound is a Synthdef (the definition of a synthesizer) inside the sound server named Supercollider. Supercollider is an open-source platform for audio synthesis and algorithmic composition. Developed by James McCartney in the late 1990s (Wilson et al., 2011), it has since become a powerful tool in the fields of music technology, computer music, and sound art. Supercollider provides a flexible and expressive environment for creating and manipulating sound in real-time, making it an invaluable resource for both research and artistic exploration. Some main features of Supercollider include:

Audio Synthesis Supercollider offers a wide range of synthesis techniques, including additive, subtractive, granular, and physical modeling synthesis. Users can create complex sounds by combining these techniques and modulating parameters in real-time.

Real-time Processing One of Supercollider’s key strengths is its ability to process audio in real-time. This makes it suitable for live performances, interactive installations, and other time-sensitive applications.

Algorithmic Composition Supercollider provides tools for generating music algorithmically, allowing users to create compositions based on mathematical algorithms, rulesets, or generative processes.

Integration with External Hardware Supercollider can interface with external MIDI controllers, audio interfaces, and other hardware devices, enabling users to incorporate physical instruments and sensors into their sound projects.

Community and Documentation Supercollider has a vibrant online community of users who share code, tutorials, and resources. Additionally, comprehensive documentation is available, including tutorials, reference guides, and examples to help users learn and master the software.

SuperCollider has been and still is actively used in academia for many purposes, for example:

Research in Music Technology Digital signal processing, human-computer interaction, and machine learning for music. Researchers leverage its flexibility and programmability to prototype new algorithms, experiment with novel synthesis techniques, and investigate the perceptual and cognitive aspects of sound.

Composition and Sound Design Supercollider is used to create innovative works that push the boundaries of traditional music and sound art. Its ability to generate complex and evolving textures, as well as its support for algorithmic composition, makes it a valuable tool for exploring new sonic territories.

Teaching and Learning Supercollider is increasingly incorporated into music technology and computer music curricula at universities and colleges worldwide. It provides students with hands-on experience in sound synthesis, programming, and digital audio processing.

7.2.2 Sound server communication

The communication between the nodes in the Elixir cluster and the sound server is done over the network using the UDP network protocol. The sound server can handle a virtual unlimited amount of messages, but is practically bounded to the udp buffer size on the actual machine that runs the sound server. As such the Elixir cluster of reactors can be of any feasible size and the sound server will respond on any number of messages by synthesizing the requested sound.

7.2.3 Musical values

The reactor or program that runs on all nodes in the cluster does for each note the following. It calculates two values given two input streams of numbers. The first source is a stream of base frequencies, represented as a float number, the second source is again a stream of numbers representing a tempo in beats per minute (bpm)

Consonant notes

Consonant notes sound harmonious when played together. consonance can be expressed as a ratio between two notes or frequencies. Ratios like the perfect fifth (frequency ratio of 3:2) and the major third (frequency ratio of 5:4) are very commonly used ratios in a musical setting. One of the two reactors in the main program used for this use case calculates such a note for a given base note, both are expressed in there actual frequency represented as a float number.

Tempo

We can express the speed of a musical note progression in beats per minute (bpm). From that number we can then calculate the duration of individual notes in the musical performance.

We calculate the time in milliseconds. Since one minute is 60000 milliseconds one can calculate the length of a quarter note with $q = \frac{60000}{bpm}$ with the assumption that we use a 4/4 time signature. From the length of the quarter note it is easy to find other durations. For example, the duration of a halve note is double that of a quarter note or the duration of a eight note is half the time of a quarter note.

One of the two reactors in the main program used for this use case calculates exactly that, the length for a quarter note and improvises a deviser or multiplier to produce some duration.

Chapter 8

Conclusion

In conclusion, this thesis presents a comprehensive exploration of distributing reactive programming within the framework of the purely functional reactive language Haai. By addressing critical challenges such as consistency, fault tolerance, and communication overhead, we developed a system capable of efficiently scaling across distributed environments. The adoption of a shared-nothing architecture, coupled with the implementation of an Erlang cluster, facilitated the creation of a robust and resilient reactive system. The Haai virtual machine (Hvm) demonstrated effective management of distributed reactive programs, leveraging native reactors and ensuring glitch-free operation. This research contributes to the broader understanding of distributed reactive programming and provides a solid foundation for future developments in building scalable, fault-tolerant reactive systems.

8.1 Further work

The cluster monitor would be an addition to this work, different distribution models could be explored towards models that do communicate

Further work can be organized to explore different distribution models on the same Elixir set of nodes. Where all nodes did react independently of each other, in future work communication between nodes can be explored and more different programs or reactors can be setup and linked together... Study glitch avoidance

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