代码库

上海交通大学

September 1, 2015

Contents

1	数论		4
	1.1	快速求逆元	 . 4
	1.2	扩展欧几里德算法	
	1.3	中国剩余定理	
	1.4	Miller Rabin 素数测试	
	1.5	Pollard Rho 大数分解	 . 5
	1.6		
	1.7	原根	
	1.8	离散对数	
	1.9	离散平方根	
	1.10	佩尔方程求解	
	1.11	牛顿迭代法	
	1.12	直线下整点个数	 . 7
_			_
2	数值		7
	2.1	高斯消元	
	2.2	快速傅立叶变换	
	2.3	单纯形法求解线性规划	 . 8
	2.4	自适应辛普森	 . 9
	2.5	多项式方程求解	 . 10
	2.6	最小二乘法	
3	数据	结构	10
	3.1	平衡的二叉查找树	 . 10
		3.1.1 Treap	
		3.1.2 Splay	
	3.2	Sina Spagiiiiiiiiiiiiiiiiiiiiiiiiiiiiiiiiiiii	
	J.L	3.2.1 坚固的线段树	
	2.2	3.2.4 坚固的左偏树	
	3.3	树上的魔术师	
		3.3.1 轻重树链剖分	
		3.3.2 Link Cut Tree	
		3.3.3 AAA Tree	
	3.4	k-d 树	 . 14
4	图论		14
	4.1	强连通分量	
	4.2	双连通分量	
		4.2.1 点双连通分量	 . 14
		4.2.2 边双连通分量	
	4.3	2-SAT 问题	- 4
		二分图最大匹配	
		441 Hungary 算法	

			Hopcroft																
	4.5		最大权匹																
		4.5.1	KM 算法																
		4.5.2	扩展 KM																
	4.6	最大流																	
	4.7		用最大流																
		4.7.1	稀疏图。																
	4.0	4.7.2	稠密图 .																
	4.8		最大匹配					 	 			 	 						 21
	4.9		全局最小																
	4.10	- 17 - 2 - 1 - 1						 	 			 	 						 24
	4.11																		
			」生成树 .																
	4.13	弦图相	大					 	 			 	 						 25
			弦图的判																
	111		弦图的团		 # == ŕ														
	4.14	归名')	(顿回路 (ORE 1	土灰店)	 	 			 				 •	 •		 25
5	字符	串																	26
J	5.1	· 単 模式四	两己																
	J.1	5.1.1	KMP 算法																
		5.1.2	扩展 KM																
		5.1.3	AC 自动																
	5.2	后缀三																	
	٥	5.2.1	后缀数组																
		5.2.2	后缀自动																
	5.3	回文三																	
	5.5	5.3.1	Manache																
		5.3.2	回文树 .																
	5.4	循环串	最小表示																
6		几何																	29
	6.1	二维基																	
		6.1.1										 							 29
			点类					 	 				 						29
		6.1.2	凸包					 	 			 	 			 			
		6.1.3	凸包 半平面交					 	 			 	 			 			 30
	6.2	6.1.3 三维基	凸包 半平面交 础	· · · · · · · · · · · · · · · · · · ·				 	 			 	 	 	· ·	 	 		 30 30
	6.2	6.1.3 三维基 6.2.1	凸包 半平面交 础 点类				· · · · · · · · · · · · · · · · · · ·	 	 	 	· · · · · · · · · · · · · · · · · · ·	 	 	· ·	 	 	 	· ·	 30 30 30
	6.2	6.1.3 三维基 6.2.1 6.2.2	凸包 半平面交 础 点类 凸包					 	 		· · · · · · · · · · · · · · · · · · ·	 	 		· · · · · · · · · · · · · · · · · · ·		 		 30 30 30 30
		6.1.3 三维基 6.2.1 6.2.2 6.2.3	凸包 半平面交 础 点类 点类 凸包 绕轴旋转					 	 	· · · · · · · · · · · · · · · · · · ·					· · · · · · · · · · · · · · · · · · ·		 		 30 30 30 30 30
	6.2	6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形	凸包 半平 点之 点类 点色 点色 。 。 。 。 。 。 。 。 。 。 。 。 。 。 。 。 。					 	 · · · · · · · · · · · · · · · · · · ·	· · · · · · · · · · · · · · · · · · ·		 		 30 30 30 30 30
		6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1	凸包 半平 础 点 点 点 色 轴 。 。 。 。 。 。 。 。 。 。 。 。 。 。 。 。 。 。	·····································	·····································	· · · · · · · · · · · ·			 					· · · · · · · · · · · · · · · · · · ·		 	· · · · · · · · · · · · · · · · · · ·	 30 30 30 30 30 30
		6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2	凸半础点凸绕, 判旋包平, 类包轴, 断转, 症点卡, 症, 点卡,	·····································	·····································	· · · · · · · · · · · ·			 								 		30 30 30 30 30 30 30
		6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3	凸半础点凸绕,判旋动包平,类包轴,断转态加速,点十分的,点十分的,有一个,一个一个一个一个一个一个一个一个一个一个一个一个一个一个一个一个一个一	·····································	·····································	· · · · · · · · · · · · · · · · · · ·			 								· · · · · · · · · · · · · · · · · · ·		30 30 30 30 30 30 31 31
		6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4	凸半础点凸绕,判旋动点包平,类包轴,断转态到流,上上上上上上上上上上上上上上上上上上上上上上上上上上上上上上上上上上上上	· · · · · · · · · · · · · · · · · · ·	····· ····· ···· ···· ···· ···· ····	· · · · · · · · · · · · · · · · · · ·			 										30 30 30 30 30 30 31 31
		6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4 6.3.5	凸半础点凸绕,判旋动点直包平,类包轴,断转态到线,面; 旋,点卡凸凸与一交。	·····································		· · · · · · · · · · · · · · · · · · ·													30 30 30 30 30 30 31 31 31
		6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.6	凸半础点凸绕,判旋动点直凸包平,类包轴,断转态到线多面, 放,点卡凸凸与边方边,交 、 、 、 、 、 、 、 、 、 、 、 、 、 、 、 、	· · · · · · · · · · · · · · · · · · ·		· · · · · · · · · · · · · · · · · · ·													30 30 30 30 30 30 31 31 31
	6.3	6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.6 6.3.7	凸半础点凸绕,判旋动点直包平,类包轴,断转态到线,面; 旋,点卡凸凸与一交。	· · · · · · · · · · · · · · · · · · ·	・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・														30 30 30 30 30 31 31 31 31
		6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.6 6.3.7 圆	凸半础点凸绕,判旋动点直凸凸:包平,类包轴,断转态到线多多;面, 旋,点卡凸凸与边边,交。转	· · · · · · · · · · · · · · · · · · ·	・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・														30 30 30 30 30 30 31 31 31 31 31
	6.3	6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.6 6.3.7 圆	凸半础点凸绕,判旋动点直凸凸、圆包平、类包轴、断转态到线多多、类面、	· · · · · · · · · · · · · · · · · · ·															30 30 30 30 30 30 31 31 31 31 31
	6.3	6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.6 6.3.7 圆 6.4.1 6.4.2	凸半础点凸绕,判旋动点直凸凸,圆圆包平,类包轴,断转态到线多多,类的面, 一、 一、 一、 一、 一、 一、 一、 一、 一、 、 、 、 、 、 、	· · · · · · · · · · · · · · · · · · ·															30 30 30 30 30 31 31 31 31 31 31
	6.3	6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.7 圆 6.4.1 6.4.2 6.4.3	凸半础点凸绕,判旋动点直凸凸:圆圆最包平:类包轴:断转态到线多多:类的小面; 一旋,点卡凸凸与边边; 一交覆,交;; 转,在壳包包凸形形;;集盖	· · · · · · · · · · · · · · · · · · ·															30 30 30 30 30 30 31 31 31 31 31 31 31 31 31
	6.3	6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.6 6.3.7 圆 6.4.1 6.4.2 6.4.3 6.4.4	凸半础点凸绕,判旋动点直凸凸,圆圆最最包平,类包轴,断转态到线多多,类的小小面,一旋,点卡凸凸与边边,一交覆覆,交,,转,在壳包包凸形形,,集盖盖	· · · · · · · · · · · · · · · · · · ·															30 30 30 30 30 30 31 31 31 31 31 31 31 31
	6.3	6.1.3 三维基 6.2.1 6.2.2 6.2.3 多3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.6 6.3.7 回 6.4.1 6.4.2 6.4.3 6.4.4 6.4.5	凸半础点凸绕,判旋动点直凸凸,圆圆最最判包平,类包轴,断转态到线多多,类的小小断面。	·	・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・														30 30 30 30 30 31 31 31 31 31 31 31 31
	6.3	6.1.3 三维基 6.2.1 6.2.2 6.2.3 多边形 6.3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.6 6.3.7 圆 6.4.1 6.4.2 6.4.3 6.4.4 6.4.5 6.4.6	凸半础点凸绕,判旋动点直凸凸,圆圆最最判圆包平,类包轴,断转态到线多多,类的小小断与一面; 一旋,点卡凸凸与边边; 一交覆覆圆多,变;;转;在壳包包凸形形;;集盖盖存达	·	・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・														30 30 30 30 30 31 31 31 31 31 31 31 31 31 31
	6.3	6.1.3 三维基 6.2.1 6.2.2 6.2.3 多3.1 6.3.2 6.3.3 6.3.4 6.3.5 6.3.6 6.3.7 回 6.4.1 6.4.2 6.4.3 6.4.4 6.4.5	凸半础点凸绕,判旋动点直凸凸,圆圆最最判圆包平,类包轴,断转态到线多多,类的小小断与一面; 一旋,点卡凸凸与边边; 一交覆覆圆多,变;;转;在壳包包凸形形;;集盖盖存达		・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・・														30 30 30 30 30 31 31 31 31 31 31 31 31 31 31

	6.6	6.5.3 三角形的垂心3黑暗科技36.6.1 平面图形的转动惯量36.6.2 平面区域处理36.6.3 Vonoroi 图3
7	其他 7.1 7.2 7.3	3 某年某月某日是星期几 3 动态规划 32 搜索 32 7.3.1 Dancing Links X 32
8	Java 8.1 8.2 8.3	基础模板 32 BigInteger 33 BigDecimal 33
9	数学 9.1 9.2 9.3	第用积分表 33 常用积分表 33 9.2.1 求和公式 33 9.2.2 斐波那契数列 33 9.2.3 错排公式 33 平面几何公式 33 9.3.1 三角形 33 9.3.1 三角形 33 9.3.2 四边形 34 9.3.3 正 n 边形 34 9.3.4 園 34 9.3.5 棱柱 34 9.3.6 棱锥 34 9.3.7 棱台 34 9.3.8 圆柱 34 9.3.9 圆锥 34 9.3.10 圆台 34 常用数表 34
		9.4.1 梅森数 32

1 数论

```
1.1 快速求逆元
```

```
返回结果:
                                       x^{-1} (mod)
   使用条件: x \in [0, mod) 并且 x = mod 互质。
long long inverse(const long long &x, const long long &mod) {
    if (x == 1) {
        return 1;
    } else {
        return (mod - mod / x) * inverse(mod % x, mod) % mod;
}
1.2 扩展欧几里德算法
  返回结果:
                                   ax + by = gcd(a, b)
  时间复杂度: O(nlogn)
void solve(const long long &a, const long long &b, long long &x, long long &y) {
    if (b == 0) {
        x = 1;
        y = 0;
    } else {
        solve(b, a \% b, x, y);
        x = a / b * y;
        std::swap(x, y);
}
1.3 中国剩余定理
  返回结果:
                               x \equiv r_i \pmod{p_i} \ (0 \le i < n)
   使用条件:p_i 无需两两互质
   时间复杂度: O(nlogn)
bool solve(int n, std::pair<long long, long long> input[], std::pair<long long, long long> &output)
    output = std::make_pair(1, 1);
    for (int i = 0; i < n; ++i) {
        long long number, useless;
        euclid(output.second, input[i].second, number, useless);
        long long divisor = std::__gcd(output.second, input[i].second);
        if ((input[i].first - output.first) % divisor) {
            return false;
        number *= (input[i].first - output.first) / divisor;
        fix(number, input[i].second);
        output.first += output.second * number;
        output.second *= input[i].second / divisor;
        fix(output.first, output.second);
    return true;
}
```

```
1.4 Miller Rabin 素数测试
```

```
const int BASE[12] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
bool check(const long long &prime, const long long &base) {
    long long number = prime - 1;
   for (; ~number & 1; number >>= 1);
    long long result = power_mod(base, number, prime);
   for (; number != prime - 1 && result != 1 && result != prime - 1; number <<= 1) {
        result = multiply_mod(result, result, prime);
   return result == prime - 1 || (number & 1) == 1;
}
bool miller_rabin(const long long &number) {
    if (number < 2) {
        return false;
   if (number < 4) {
        return true;
   if (~number & 1) {
        return false;
    for (int i = 0; i < 12 && BASE[i] < number; ++i) {</pre>
        if (!check(number, BASE[i])) {
            return false;
   return true;
}
1.5 Pollard Rho 大数分解
  时间复杂度: O(n^{1/4})
long long pollard_rho(const long long &number, const long long &seed) {
    long long x = rand() \% (number - 1) + 1, y = x;
    for (int head = 1, tail = 2; ; ) {
        x = multiply_mod(x, x, number);
        x = add_mod(x, seed, number);
        if (x == y) {
            return number;
        long long answer = std::__gcd(abs(x - y), number);
        if (answer > 1 && answer < number) {</pre>
            return answer;
        }
        if (++head == tail) {
            y = x;
            tail <<= 1;
        }
   }
}
void factorize(const long long &number, std::vector<long long> &divisor) {
    if (number > 1) {
        if (miller rabin(number)) {
            divisor.push_back(number);
```

```
} else {
            long long factor = number;
            for (; factor >= number; factor = pollard_rho(number, rand() % (number - 1) + 1));
            factorize(number / factor, divisor);
            factorize(factor, divisor);
        }
    }
}
1.6 快速数论变换
  返回结果:
                            c_i = \sum_{0 \le j \le i} a_j \cdot b_{i-j}(mod) \ (0 \le i < n)
   使用说明:magic 是 mod 的原根
   时间复杂度:O(nlogn)
void solve(long long number[], int length, int type) {
        for (int i = 1, j = 0; i < length - 1; ++i) {
                for (int k = length; j ^= k >>= 1, ~j & k; );
                if (i < j) {
                         std::swap(number[i], number[j]);
        }
        long long unit_p0;
        for (int turn = 0; (1 << turn) < length; ++turn) {</pre>
                int step = 1 << turn, step2 = step << 1;</pre>
                if (type == 1) {
                         unit_p0 = power_mod(MAGIC, (MOD - 1) / step2, MOD);
                } else {
                         unit_p0 = power_mod(MAGIC, MOD - 1 - (MOD - 1) / step2, MOD);
                for (int i = 0; i < length; i += step2) {
                         long long unit = 1;
                         for (int j = 0; j < step; ++j) {
                                 long long &number1 = number[i + j + step];
                                 long long &number2 = number[i + j];
                                 long long delta = unit * number1 % MOD;
                                 number1 = (number2 - delta + MOD) % MOD;
                                 number2 = (number2 + delta) % MOD;
                                 unit = unit * unit_p0 % MOD;
                         }
                }
        }
}
void multiply() {
        for (; lowbit(length) != length; ++length);
        solve(number1, length, 1);
        solve(number2, length, 1);
        for (int i = 0; i < length; ++i) {</pre>
                number[i] = number1[i] * number2[i] % MOD;
        solve(number, length, -1);
        for (int i = 0; i < length; ++i) {</pre>
                answer[i] = number[i] * power_mod(length, MOD - 2, MOD) % MOD;
        }
}
```

- 1.7 原根
- 1.8 离散对数
- 1.9 离散平方根
- 1.10 佩尔方程求解
- 1.11 牛顿迭代法
- 1.12 直线下整点个数

返回结果:

$$\sum_{0 \le i < n} \lfloor \frac{a + b \cdot i}{m} \rfloor$$

使用条件:n, m > 0, $a, b \ge 0$ 时间复杂度:O(nlogn)

```
long long solve(const long long &n, const long long &a, const long long &b, const long long &m) {
   if (b == 0) {
     return n * (a / m);
}
   if (a >= m) {
     return n * (a / m) + solve(n, a % m, b, m);
}
   if (b >= m) {
     return (n - 1) * n / 2 * (b / m) + solve(n, a, b % m, m);
}
   return solve((a + b * n) / m, (a + b * n) % m, m, b);
}
```

- 2 数值
- 2.1 高斯消元
- 2.2 快速傅立叶变换

返回结果:

$$c_i = \sum_{0 \le j \le i} a_j \cdot b_{i-j} \ (0 \le i < n)$$

时间复杂度: O(nlogn)

```
void solve(Complex number[], int length, int type) {
    for (int i = 1, j = 0; i < length - 1; ++i) {
        for (int k = length; j ^= k >>= 1, ~j & k; );
        if (i < j) {
            std::swap(number[i], number[j]);
        }
    }
    Complex unit_p0;
    for (int turn = 0; (1 << turn) < length; ++turn) {</pre>
        int step = 1 << turn, step2 = step << 1;</pre>
        double p0 = PI / step * type;
        sincos(p0, &unit_p0.imag(), &unit_p0.real());
        for (int i = 0; i < length; i += step2) {</pre>
            Complex unit = 1;
            for (int j = 0; j < step; ++j) {
                Complex &number1 = number[i + j + step];
```

```
Complex &number2 = number[i + j];
                 Complex delta = unit * number1;
                 number1 = number2 - delta;
                 number2 = number2 + delta;
                 unit = unit * unit_p0;
            }
        }
    }
}
void multiply() {
    for (; lowbit(length) != length; ++length);
    solve(number1, length, 1);
    solve(number2, length, 1);
    for (int i = 0; i < length; ++i) {</pre>
        number[i] = number1[i] * number2[i];
    solve(number, length, −1);
    for (int i = 0; i < length; ++i) {</pre>
        answer[i] = (int)(number[i].real() / length + 0.5);
    }
}
2.3 单纯形法求解线性规划
   返回结果:
                     max\{c_{1\times m}\cdot x_{m\times 1}\mid x_{m\times 1}\geq 0_{m\times 1}, a_{n\times m}\cdot x_{m\times 1}\leq b_{n\times 1}\}
std::vector<double> solve(const std::vector<std::vector<double> > &a,
                         const std::vector<double> &b, const std::vector<double> &c) {
    int n = (int)a.size(), m = (int)a[0].size() + 1;
    std::vector<std::vector<double> > value(n + 2, std::vector<double>(m + 1));
    std::vector<int> index(n + m);
    int r = n, s = m - 1;
    for (int i = 0; i < n + m; ++i) {
        index[i] = i;
    for (int i = 0; i < n; ++i) {</pre>
        for (int j = 0; j < m - 1; ++j) {
             value[i][j] = -a[i][j];
        value[i][m - 1] = 1;
        value[i][m] = b[i];
        if (value[r][m] > value[i][m]) {
            r = i;
    for (int j = 0; j < m - 1; ++j) {
        value[n][j] = c[j];
    value[n + 1][m - 1] = -1;
    for (double number; ; ) {
        if (r < n) {
             std::swap(index[s], index[r + m]);
             value[r][s] = 1 / value[r][s];
             for (int j = 0; j \le m; ++j) {
                 if (j != s) {
                     value[r][j] *= -value[r][s];
```

```
}
            for (int i = 0; i <= n + 1; ++i) {
                if (i != r) {
                    for (int j = 0; j \le m; ++j) {
                        if (j != s) {
                            value[i][j] += value[r][j] * value[i][s];
                    value[i][s] *= value[r][s];
                }
            }
        }
        r = s = -1;
        for (int j = 0; j < m; ++j) {
            if (s < 0 || index[s] > index[j]) {
                if (value[n + 1][j] > eps | | value[n + 1][j] > -eps && value[n][j] > eps) {
                    s = j;
            }
        }
        if (s < 0) {
            break;
        for (int i = 0; i < n; ++i) {
            if (value[i][s] < -eps) {</pre>
                if (r < 0)
                    || (number = value[r][m] / value[r][s] - value[i][m] / value[i][s]) < -eps
                    | |  number < eps && index[r + m] > index[i + m]) {
                     r = i;
                }
            }
        }
        if (r < 0) {
            //
                Solution is unbounded.
            return std::vector<double>();
        }
    if (value[n + 1][m] < -eps) {
             No solution.
        return std::vector<double>();
    std::vector<double> answer(m - 1);
   for (int i = m; i < n + m; ++i) {
        if (index[i] < m - 1) {</pre>
            answer[index[i]] = value[i - m][m];
        }
   return answer;
2.4 自适应辛普森
double area(const double &left, const double &right) {
   double mid = (left + right) / 2;
   return (right - left) * (calc(left) + 4 * calc(mid) + calc(right)) / 6;
}
```

}

```
double simpson(const double &left, const double &right, const double &eps, const double &area_sum) {
    double mid = (left + right) / 2;
    double area_left = area(left, mid);
    double area_right = area(mid, right);
    double area_total = area_left + area_right;
    if (std::abs(area_total - area_sum) < 15 * eps) {</pre>
        return area_total + (area_total - area_sum) / 15;
    return simpson(left, mid, eps / 2, area_left) + simpson(mid, right, eps / 2, area_right);
}
double simpson(const double &left, const double &right, const double &eps) {
    return simpson(left, right, eps, area(left, right));
}
2.5
     多项式方程求解
2.6 最小二乘法
3
    数据结构
3.1 平衡的二叉查找树
3.1.1 Treap
class Node {
public:
    Node *child[2];
    int key;
    int size, priority;
    Node(Node *left, Node *right, int key) : key(key), size(1), priority(rand()) {
        child[0] = left;
        child[1] = right;
    }
    void update() {
        size = child[0]->size + 1 + child[1]->size;
};
Node *null;
void rotate(Node *&x, int dir) {
    Node *y = x->child[dir];
    x->child[dir] = y->child[dir ^ 1];
    y \rightarrow child[dir ^ 1] = x;
    x->update();
    y->update();
    x = y;
}
void insert(Node *&x, int key) {
    if (x == null) {
        x = new Node(null, null, key);
    } else {
        insert(x->child[key > x->key], key);
        if (x->child[key > x->key]->priority < x->priority) {
```

```
rotate(x, key > x->key);
        x->update();
    }
}
void remove(Node *&x, int key) {
    if (x->key != key) {
        remove(x->child[key > x->key], key);
    } else if (x->child[0] == null && x->child[1] == null) {
        x = null;
        return;
    } else {
        int dir = x->child[0]->priority > x->child[1]->priority;
        rotate(x, dir);
        remove(x->child[dir ^ 1], key);
    x->update();
}
void build() {
    null = new Node(NULL, NULL, 0);
    null->child[0] = null->child[1] = null;
    null->size = 0;
   null->priority = RAND_MAX;
}
3.1.2 Splay
3.2 坚固的数据结构
3.2.1 坚固的线段树
class Node {
public:
    Node *left, *right;
    int value;
    Node(Node *left, Node *right, int value) : left(left), right(right), value(value) {}
    Node* modify(int 1, int r, int q1, int qr, int value);
    int query(int 1, int r, int qx);
};
Node* null;
Node* Node::modify(int 1, int r, int q1, int qr, int value) {
    if (qr < 1 || r < ql) {
       return this;
    }
    if (ql <= 1 && r <= qr) {
        return new Node(this->left, this->right, this->value + value);
    int mid = 1 + r >> 1;
    return new Node(this->left->modify(1, mid, q1, qr, value),
                    this->right->modify(mid + 1, r, ql, qr, value),
                    this->value);
}
```

```
int Node::query(int 1, int r, int qx) {
    if (qx < 1 | | r < qx) {
        return 0;
    if (qx \le 1 \&\& r \le qx) \{
        return this->value;
    int mid = 1 + r >> 1;
    return this->left->query(1, mid, qx)
         + this->right->query(mid + 1, r, qx)
         + this->value;
}
void build() {
    null = new Node(NULL, NULL, 0);
    null->left = null->right = null;
}
3.2.2 坚固的平衡树
3.2.3 坚固的字符串
3.2.4 坚固的左偏树
3.3 树上的魔术师
3.3.1 轻重树链剖分
int father[N], height[N], size[N], son[N], top[N], pos[N], data[N];
void build(int root) {
    std::vector<int> queue;
    father[root] = -1;
    height[root] = 0;
    queue.push_back(root);
    for (int head = 0; head < (int)queue.size(); ++head) {</pre>
        int x = queue[head];
        for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
            int y = edge[x][i];
            if (y != father[x]) {
                father[y] = x;
                height[y] = height[x] + 1;
                queue.push_back(y);
            }
        }
    }
    for (int index = n - 1; index >= 0; --index) {
        int x = queue[index];
        size[x] = 1;
        son[x] = -1;
        for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
            int y = edge[x][i];
            if (y != father[x]) {
                size[x] += size[y];
                if (son[x] == -1 \mid \mid size[son[x]] < size[y]) {
                    son[x] = y;
                }
            }
```

```
}
    }
    std::fill(top, top + n, 0);
    int counter = 0;
    for (int index = 0; index < n; ++index) {</pre>
        int x = queue[index];
        if (top[x] == 0) {
            for (int y = x; y != -1; y = son[y]) {
                top[y] = x;
                pos[y] = ++counter;
                data[counter] = value[y];
        }
    }
    build(1, 1, n);
void solve(int x, int y) {
    while (true) {
        if (top[x] == top[y]) {
            if (x == y) {
                solve(1, 1, n, pos[x], pos[x]);
            } else {
                if (height[x] < height[y]) {</pre>
                    solve(1, 1, n, pos[x], pos[y]);
                } else {
                    solve(1, 1, n, pos[y], pos[x]);
            }
            break;
        }
        if (height[top[x]] > height[top[y]]) {
            solve(1, 1, n, pos[top[x]], pos[x]);
            x = father[top[x]];
            solve(1, 1, n, pos[top[y]], pos[y]);
            y = father[top[y]];
        }
    }
}
```

```
3.3.2 Link Cut Tree
3.3.3 AAA Tree
3.4 k-d 树
    图论
4
4.1 强连通分量
4.2 双连通分量
4.2.1 点双连通分量
4.2.2 边双连通分量
4.3 2-SAT 问题
4.4 二分图最大匹配
4.4.1 Hungary 算法
  时间复杂度:O(V \cdot E)
int n, m, stamp;
int match[N], visit[N];
bool dfs(int x) {
   for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
        int y = edge[x][i];
        if (visit[y] != stamp) {
           visit[y] = stamp;
           if (match[y] == -1 || dfs(match[y])) {
               match[y] = x;
               return true;
           }
       }
   return false;
}
int solve() {
    std::fill(match, match + m, -1);
    int answer = 0;
   for (int i = 0; i < n; ++i) {
       stamp++;
       answer += dfs(i);
   return answer;
4.4.2 Hopcroft Karp 算法
  时间复杂度:O(\sqrt{V} \cdot E)
int matchx[N], matchy[N], level[N];
bool dfs(int x) {
   for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
       int y = edge[x][i];
```

```
int w = matchy[y];
        if (w == -1 \mid \mid level[x] + 1 == level[w] && dfs(w)) {
            matchx[x] = y;
            matchy[y] = x;
            return true;
        }
    }
    level[x] = -1;
    return false;
}
int solve() {
    std::fill(matchx, matchx + n, -1);
    std::fill(matchy, matchy + m, -1);
    for (int answer = 0; ; ) {
        std::vector<int> queue;
        for (int i = 0; i < n; ++i) {
            if (matchx[i] == -1) {
                level[i] = 0;
                queue.push_back(i);
            } else {
                level[i] = -1;
            }
        for (int head = 0; head < (int)queue.size(); ++head) {</pre>
            int x = queue[head];
            for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
                int y = edge[x][i];
                int w = matchy[y];
                if (w != -1 \&\& level[w] < 0) {
                    level[w] = level[x] + 1;
                    queue.push_back(w);
                }
            }
        }
        int delta = 0;
        for (int i = 0; i < n; ++i) {</pre>
            if (matchx[i] == -1 && dfs(i)) {
                delta++;
            }
        }
        if (delta == 0) {
            return answer;
        } else {
            answer += delta;
        }
    }
}
4.5 二分图最大权匹配
4.5.1 KM 算法
  时间复杂度:O(V^3)
int labelx[N], labely[N], match[N], slack[N];
bool visitx[N], visity[N];
bool dfs(int x) {
```

```
visitx[x] = true;
    for (int y = 0; y < n; ++y) {
        if (visity[y]) {
            continue;
        int delta = labelx[x] + labely[y] - graph[x][y];
        if (delta == 0) {
            visity[y] = true;
            if (match[y] == -1 \mid \mid dfs(match[y])) {
                match[y] = x;
                return true;
            }
        } else {
            slack[y] = std::min(slack[y], delta);
    }
    return false;
}
int solve() {
    for (int i = 0; i < n; ++i) {
        match[i] = -1;
        labelx[i] = INT_MIN;
        labely[i] = 0;
        for (int j = 0; j < n; ++j) {
            labelx[i] = std::max(labelx[i], graph[i][j]);
        }
    }
    for (int i = 0; i < n; ++i) {
        while (true) {
            std::fill(visitx, visitx + n, 0);
            std::fill(visity, visity + n, 0);
            for (int j = 0; j < n; ++j) {
                slack[j] = INT_MAX;
            }
            if (dfs(i)) {
                break;
            }
            int delta = INT_MAX;
            for (int j = 0; j < n; ++j) {
                if (!visity[j]) {
                    delta = std::min(delta, slack[j]);
            }
            for (int j = 0; j < n; ++j) {
                if (visitx[j]) {
                    labelx[j] -= delta;
                if (visity[j]) {
                    labely[j] += delta;
                } else {
                    slack[j] -= delta;
            }
        }
    }
    int answer = 0;
    for (int i = 0; i < n; ++i) {
```

```
answer += graph[match[i]][i];
   return answer;
4.5.2 扩展 KM 算法
4.6 最大流
  时间复杂度: O(V^2 \cdot E)
struct EdgeList {
    int size;
    int last[N];
    int succ[M], other[M], flow[M];
    void clear(int n) {
       size = 0;
       fill(last, last + n, -1);
   void add(int x, int y, int c) {
       succ[size] = last[x];
       last[x] = size;
       other[size] = y;
       flow[size++] = c;
   }
} e;
int n, source, target;
int dist[N], curr[N];
void add(int x, int y, int c) {
    e.add(x, y, c);
    e.add(y, x, 0);
}
bool relabel() {
   std::vector<int> queue;
   for (int i = 0; i < n; ++i) {
        curr[i] = e.last[i];
       dist[i] = -1;
   queue.push_back(target);
   dist[target] = 0;
   for (int head = 0; head < (int)queue.size(); ++head) {</pre>
        int x = queue[head];
        for (int i = e.last[x]; ~i; i = e.succ[i]) {
           int y = e.other[i];
           dist[y] = dist[x] + 1;
               queue.push_back(y);
       }
   return ~dist[source];
int dfs(int x, int answer) {
   if (x == target) {
       return answer;
```

```
}
    int delta = answer;
    for (int &i = curr[x]; ~i; i = e.succ[i]) {
        int y = e.other[i];
        if (e.flow[i] \&\& dist[x] == dist[y] + 1) {
            int number = dfs(y, std::min(e.flow[i], delta));
            e.flow[i] -= number;
            e.flow[i ^ 1] += number;
            delta -= number;
        }
        if (delta == 0) {
            break;
    }
    return answer - delta;
int solve() {
    int answer = 0;
    while (relabel()) {
        answer += dfs(source, INT_MAX));
    return answer;
4.7 最小费用最大流
4.7.1 稀疏图
  时间复杂度:O(V \cdot E^2)
struct EdgeList {
    int size;
    int last[N];
    int succ[M], other[M], flow[M], cost[M];
    void clear(int n) {
        size = 0;
        std::fill(last, last + n, -1);
    void add(int x, int y, int c, int w) {
        succ[size] = last[x];
        last[x] = size;
        other[size] = y;
        flow[size] = c;
        cost[size++] = w;
    }
} e;
int n, source, target;
int prev[N];
void add(int x, int y, int c, int w) {
    e.add(x, y, c, w);
    e.add(y, x, 0, -w);
}
bool augment() {
    static int dist[N], occur[N];
    std::vector<int> queue;
```

```
std::fill(dist, dist + n, INT_MAX);
    std::fill(occur, occur + n, 0);
    dist[source] = 0;
    occur[source] = true;
    queue.push_back(source);
    for (int head = 0; head < (int)queue.size(); ++head) {</pre>
        int x = queue[head];
        for (int i = e.last[x]; ~i; i = e.succ[i]) {
            int y = e.other[i];
            if (e.flow[i] && dist[y] > dist[x] + e.cost[i]) {
                dist[y] = dist[x] + e.cost[i];
                prev[y] = i;
                if (!occur[y]) {
                    occur[y] = true;
                    queue.push_back(y);
            }
        }
        occur[x] = false;
    return dist[target] < INT_MAX;</pre>
}
std::pair<int, int> solve() {
    std::pair<int, int> answer = std::make_pair(0, 0);
    while (augment()) {
        int number = INT MAX;
        for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
            number = std::min(number, e.flow[prev[i]]);
        answer.first += number;
        for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
            e.flow[prev[i]] -= number;
            e.flow[prev[i] ^ 1] += number;
            answer.second += number * e.cost[prev[i]];
        }
    }
    return answer;
4.7.2 稠密图
  时间复杂度: O(V \cdot E^2)
struct EdgeList {
    int size;
    int last[N];
    int succ[M], other[M], flow[M], cost[M];
    void clear(int n) {
        size = 0;
        std::fill(last, last + n, -1);
    void add(int x, int y, int c, int w) {
        succ[size] = last[x];
        last[x] = size;
        other[size] = y;
        flow[size] = c;
        cost[size++] = w;
```

```
}
} e;
int n, source, target, flow, cost;
int slack[N], dist[N];
bool visit[N];
void add(int x, int y, int c, int w) {
    e.add(x, y, c, w);
    e.add(y, x, 0, -w);
}
bool relabel() {
    int delta = INT_MAX;
    for (int i = 0; i < n; ++i) {
        if (!visit[i]) {
            delta = std::min(delta, slack[i]);
        slack[i] = INT_MAX;
    }
    if (delta == INT_MAX) {
        return true;
    for (int i = 0; i < n; ++i) {
        if (visit[i]) {
            dist[i] += delta;
    return false;
}
int dfs(int x, int answer) {
    if (x == target) {
        flow += answer;
        cost += answer * (dist[source] - dist[target]);
        return answer;
    visit[x] = true;
    int delta = answer;
    for (int i = e.last[x]; ~i; i = e.succ[i]) {
        int y = e.other[i];
        if (e.flow[i] > 0 && !visit[y]) {
            if (dist[y] + e.cost[i] == dist[x]) {
                int number = dfs(y, std::min(e.flow[i], delta));
                e.flow[i] -= number;
                e.flow[i ^ 1] += number;
                delta -= number;
                if (delta == 0) {
                    dist[x] = INT_MIN;
                    return answer;
                }
            } else {
                slack[y] = std::min(slack[y], dist[y] + e.cost[i] - dist[x]);
        }
    }
    return answer - delta;
}
```

```
std::pair<int, int> solve() {
    flow = cost = 0;
    std::fill(dist, dist + n, 0);
    do {
        do {
            fill(visit, visit + n, 0);
        } while (dfs(source, INT_MAX));
    } while (!relabel());
    return std::make_pair(flow, cost);
}
4.8 一般图最大匹配
  时间复杂度: O(V^3)
int match[N], belong[N], next[N], mark[N], visit[N];
std::vector<int> queue;
int find(int x) {
    if (belong[x] != x) {
        belong[x] = find(belong[x]);
    return belong[x];
}
void merge(int x, int y) {
    x = find(x);
    y = find(y);
    if (x != y) {
       belong[x] = y;
}
int lca(int x, int y) {
    static int stamp = 0;
    stamp++;
    while (true) {
        if (x != -1) {
            x = find(x);
            if (visit[x] == stamp) {
                return x;
            }
            visit[x] = stamp;
            if (match[x] != -1) {
                x = next[match[x]];
            } else {
                x = -1;
            }
        }
        std::swap(x, y);
    }
}
void group(int a, int p) {
    while (a != p) {
        int b = match[a], c = next[b];
        if (find(c) != p) {
```

```
next[c] = b;
        }
        if (mark[b] == 2) {
            mark[b] = 1;
            queue.push_back(b);
        if (mark[c] == 2) {
            mark[c] = 1;
            queue.push_back(c);
        merge(a, b);
        merge(b, c);
        a = c;
    }
}
void augment(int source) {
    queue.clear();
    for (int i = 0; i < n; ++i) {
        next[i] = visit[i] = -1;
        belong[i] = i;
        mark[i] = 0;
    mark[source] = 1;
    queue.push_back(source);
    for (int head = 0; head < (int)queue.size() && match[source] == -1; ++head) {
        int x = queue[head];
        for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
            int y = edge[x][i];
            if (match[x] == y \mid \mid find(x) == find(y) \mid \mid mark[y] == 2) {
                continue;
            if (mark[y] == 1) {
                int r = lca(x, y);
                if (find(x) != r) {
                    next[x] = y;
                if (find(y) != r) {
                    next[y] = x;
                group(x, r);
                group(y, r);
            } else if (match[y] == -1) {
                next[y] = x;
                for (int u = y; u != -1; ) {
                     int v = next[u];
                     int mv = match[v];
                    match[v] = u;
                    match[u] = v;
                    u = mv;
                break;
            } else {
                next[y] = x;
                mark[y] = 2;
                mark[match[y]] = 1;
                queue.push_back(match[y]);
            }
```

```
}
    }
}
int solve() {
    std::fill(match, match + n, -1);
    for (int i = 0; i < n; ++i) {
        if (match[i] == -1) {
            augment(i);
        }
    }
    int answer = 0;
    for (int i = 0; i < n; ++i) {
        answer += (match[i] != -1);
    return answer;
}
4.9 无向图全局最小割
   时间复杂度:O(V^3)
  注意事项:处理重边时,应该对边权累加
int node[N], dist[N];
bool visit[N];
int solve(int n) {
    int answer = INT_MAX;
    for (int i = 0; i < n; ++i) {
        node[i] = i;
    while (n > 1) {
        int max = 1;
        for (int i = 0; i < n; ++i) {
            dist[node[i]] = graph[node[0]][node[i]];
            if (dist[node[i]] > dist[node[max]]) {
                max = i;
            }
        }
        int prev = 0;
        memset(visit, 0, sizeof(visit));
        visit[node[0]] = true;
        for (int i = 1; i < n; ++i) {
            if (i == n - 1) {
                answer = std::min(answer, dist[node[max]]);
                for (int k = 0; k < n; ++k) {
                    graph[node[k]][node[prev]] = (graph[node[prev]][node[k]] += graph[node[k]][node[
                node[max] = node[--n];
            visit[node[max]] = true;
           prev = max;
           \max = -1;
            for (int j = 1; j < n; ++j) {
                if (!visit[node[j]]) {
                    dist[node[j]] += graph[node[prev]][node[j]];
                    if (max == -1 || dist[node[max]] < dist[node[j]]) {</pre>
                        max = j;
```

```
}
                }
            }
        }
    return answer;
}
4.10 最小树形图
4.11 有根树的同构
  时间复杂度: O(VlogV)
const unsigned long long MAGIC = 4423;
unsigned long long magic[N];
std::pair<unsigned long long, int> hash[N];
void solve(int root) {
    magic[0] = 1;
    for (int i = 1; i <= n; ++i) {
        magic[i] = magic[i - 1] * MAGIC;
    std::vector<int> queue;
    queue.push_back(root);
    for (int head = 0; head < (int)queue.size(); ++head) {</pre>
        int x = queue[head];
        for (int i = 0; i < (int)son[x].size(); ++i) {</pre>
            int y = son[x][i];
            queue.push_back(y);
        }
    }
    for (int index = n - 1; index >= 0; --index) {
        int x = queue[index];
        hash[x] = std::make_pair(0, 0);
        std::vector<std::pair<unsigned long long, int> > value;
        for (int i = 0; i < (int)son[x].size(); ++i) {</pre>
            int y = son[x][i];
            value.push_back(hash[y]);
        }
        std::sort(value.begin(), value.end());
        hash[x].first = hash[x].first * magic[1] + 37;
        hash[x].second++;
        for (int i = 0; i < (int)value.size(); ++i) {</pre>
            hash[x].first = hash[x].first * magic[value[i].second] + value[i].first;
            hash[x].second += value[i].second;
        hash[x].first = hash[x].first * magic[1] + 41;
        hash[x].second++;
    }
}
```

```
4.12 度限制生成树
4.13 弦图相关
4.13.1 弦图的判定
4.13.2 弦图的团数
4.14 哈密尔顿回路(ORE 性质的图)
  ORE 性质:
                        \forall x, y \in V \land (x, y) \notin E \text{ s.t. } deg_x + deg_y \ge n
  返回结果:从顶点1出发的一个哈密尔顿回路
  使用条件: n \ge 3
int left[N], right[N], next[N], last[N];
void cover(int x) {
    left[right[x]] = left[x];
    right[left[x]] = right[x];
int adjacent(int x) {
    for (int i = right[0]; i <= n; i = right[i]) {</pre>
        if (graph[x][i]) {
           return i;
        }
    }
    return 0;
}
std::vector<int> solve() {
    for (int i = 1; i <= n; ++i) {
       left[i] = i - 1;
        right[i] = i + 1;
    int head, tail;
    for (int i = 2; i <= n; ++i) {
        if (graph[1][i]) {
           head = 1;
           tail = i;
            cover(head);
            cover(tail);
           next[head] = tail;
           break;
        }
    }
    while (true) {
        int x;
        while (x = adjacent(head)) {
            next[x] = head;
           head = x;
            cover(head);
        }
        while (x = adjacent(tail)) {
           next[tail] = x;
           tail = x;
            cover(tail);
        }
```

```
if (!graph[head][tail]) {
            for (int i = head, j; i != tail; i = next[i]) {
                if (graph[head][next[i]] && graph[tail][i]) {
                    for (j = head; j != i; j = next[j]) {
                        last[next[j]] = j;
                    }
                    j = next[head];
                    next[head] = next[i];
                    next[tail] = i;
                    tail = j;
                    for (j = i; j != head; j = last[j]) {
                        next[j] = last[j];
                    break;
                }
            }
        }
        next[tail] = head;
        if (right[0] > n) {
            break;
        }
        for (int i = head; i != tail; i = next[i]) {
            if (adjacent(i)) {
                head = next[i];
                tail = i;
                next[tail] = 0;
                break;
            }
        }
    }
    std::vector<int> answer;
    for (int i = head; ; i = next[i]) {
        if (i == 1) {
            answer.push_back(i);
            for (int j = next[i]; j != i; j = next[j]) {
                answer.push_back(j);
            answer.push_back(i);
            break;
        }
        if (i == tail) {
            break;
        }
    }
    return answer;
}
5
    字符串
    模式匹配
5.1.1 KMP 算法
void build(char *pattern) {
    int length = (int)strlen(pattern + 1);
    fail[0] = -1;
    for (int i = 1, j; i <= length; ++i) {</pre>
```

```
for (j = fail[i - 1]; j != -1 && pattern[i] != pattern[j + 1]; j = fail[j]);
        fail[i] = j + 1;
    }
}
void solve(char *text, char *pattern) {
    int length = (int)strlen(text + 1);
    for (int i = 1, j; i <= length; ++i) {
        for (j = match[i - 1]; j != -1 \&\& text[i] != pattern[j + 1]; j = fail[j]);
        match[i] = j + 1;
    }
}
5.1.2 扩展 KMP 算法
5.1.3 AC 自动机,
5.2 后缀三姐妹
5.2.1 后缀数组
5.2.2 后缀自动机
class Node {
public:
    Node *child[256], *parent;
    int length;
    Node(int length = 0) : parent(NULL), length(length) {
        memset(child, NULL, sizeof(child));
    Node* extend(Node *start, int token) {
        Node *p = this;
        Node *np = new Node(length + 1);
        for (; p \&\& !p->child[token]; p = p->parent) {
            p->child[token] = np;
        if (!p) {
            np->parent = start;
        } else {
            Node *q = p->child[token];
            if (p->length + 1 == q->length) {
               np->parent = q;
            } else {
                Node *nq = new Node(p->length + 1);
                memcpy(nq->child, q->child, sizeof(q->child));
                nq->parent = q->parent;
                np->parent = q->parent = nq;
                for (; p \&\& p-> child[token] == q; p = p-> parent) {
                    p->child[token] = nq;
            }
        }
        return np;
    }
};
```

5.3 回文三兄弟

5.3.1 Manacher 算法 void manacher(char *text, int length) { palindrome[0] = 1; for (int i = 1, j = 0; i < length; ++i) { if $(j + palindrome[j] \le i) {$ palindrome[i] = 0; } else { palindrome[i] = std::min(palindrome[(j << 1) - i], j + palindrome[j] - i);</pre> while (i - palindrome[i] >= 0 && i + palindrome[i] < length</pre> && text[i - palindrome[i]] == text[i + palindrome[i]]) { palindrome[i]++; if (i + palindrome[i] > j + palindrome[j]) { } } 5.3.2 回文树 class Node { public: Node *child[256], *fail; int length; Node(int length) : fail(NULL), length(length) { memset(child, NULL, sizeof(child)); }; int size; int text[N]; Node *odd, *even; Node* match(Node *now) { for (; text[size - now->length - 1] != text[size]; now = now->fail); return now; } bool extend(Node *&last, int token) { text[++size] = token; Node *now = last; now = match(now); if (now->child[token]) { last = now->child[token]; return false; last = now->child[token] = new Node(now->length + 2); if (now == odd) { last->fail = even; } else { now = match(now->fail); last->fail = now->child[token]; return true;

```
}
void build() {
    text[size = 0] = -1;
    even = new Node(0), odd = new Node(-1);
    even->fail = odd;
    循环串最小表示
int solve(char *text, int length) {
    int i = 0, j = 1, delta = 0;
    while (i < length && j < length && delta < length) {
        char tokeni = text[(i + delta) % length];
        char tokenj = text[(j + delta) % length];
        if (tokeni == tokenj) {
            delta++;
        } else {
            if (tokeni > tokenj) {
                i += delta + 1;
            } else {
                j += delta + 1;
            }
            if (i == j) {
                j++;
            delta = 0;
        }
    return std::min(i, j);
    计算几何
6.1 二维基础
6.1.1 点类
6.1.2 凸包
std::vector<Point> convex_hull(std::vector<Point> point) {
    if ((int)point.size() < 3) {</pre>
        return point;
    std::sort(point.begin(), point.end());
    std::vector<Point> convex;
    {
        std::vector<Point> stack;
        for (int i = 0; i < (int)point.size(); ++i) {</pre>
            while ((int)stack.size() >= 2 &&
                    sgn(det(stack[(int)stack.size() - 2], stack.back(), point[i])) <= 0) {</pre>
                stack.pop_back());
            stack.push_back(point[i]);
        }
        for (int i = 0; i < (int)stack.size(); ++i) {</pre>
            convex.push_back(stack[i]);
        }
```

```
}
    {
        std::vector<Point> stack;
        for (int i = (int)point.size() - 1; i >= 0; --i) {
            while ((int)stack.size() >= 2 &&
                    sgn(det(stack[(int)stack.size() - 2], stack.back(), point[i])) <= 0) {</pre>
                stack.pop_back());
            }
            stack.push_back(point[i]);
        }
        for (int i = 1; i < (int)stack.size() - 1; ++i) {</pre>
            convex.push_back(stack[i]);
    }
    return convex;
6.1.3 半平面交
6.2 三维基础
6.2.1 点类
6.2.2 凸包
6.2.3 绕轴旋转
6.3 多边形
6.3.1 判断点在多边形内部
bool point_on_line(const Point &p, const Point &a, const Point &b) {
    return sgn(det(p, a, b)) == 0 && sgn(dot(p, a, b)) <= 0;
}
bool point_in_polygon(const Point &p, const std::vector<Point> &polygon) {
    int counter = 0;
    for (int i = 0; i < (int)polygon.size(); ++i) {</pre>
        Point a = polygon[i], b = polygon[(i + 1) % (int)polygon.size()];
        if (point_on_line(p, a, b)) {
                 Point on the boundary are excluded.
            return false;
        }
        int x = sgn(det(a, p, b));
        int y = sgn(a.y - p.y);
        int z = sgn(b.y - p.y);
        counter += (x > 0 \&\& y \le 0 \&\& z > 0);
        counter -= (x < 0 \&\& z <= 0 \&\& y > 0);
    return counter;
}
```

- 6.3.2 旋转卡壳
- 6.3.3 动态凸包
- 6.3.4 点到凸包的切线
- 6.3.5 直线与凸包的交点
- 6.3.6 凸多边形的交集
- 6.3.7 凸多边形内的最大圆
- 6.4 圆
- 6.4.1 圆类
- 6.4.2 圆的交集
- 6.4.3 最小覆盖圆
- 6.4.4 最小覆盖球
- 6.4.5 判断圆存在交集
- 6.4.6 圆与多边形的交集
- 6.5 三角形
- 6.5.1 三角形的内心
- 6.5.2 三角形的外心
- 6.5.3 三角形的垂心
- 6.6 黑暗科技
- 6.6.1 平面图形的转动惯量
- 6.6.2 平面区域处理
- 6.6.3 Vonoroi 图

7 其他

7.1 某年某月某日是星期几

```
int solve(int year, int month, int day) {
   int answer;
   if (month == 1 || month == 2) {
        month += 12;
        year--;
   }
   if ((year < 1752) || (year == 1752 && month < 9) || (year == 1752 && month == 9 && day < 3)) {
        answer = (day + 2 * month + 3 * (month + 1) / 5 + year + year / 4 + 5) % 7;
   } else {
        answer = (day + 2 * month + 3 * (month + 1) / 5 + year + year / 4 - year / 100 + year / 400)
   }
   return answer;
}</pre>
```

```
7.2 动态规划
7.3 搜索
7.3.1 Dancing Links X
8
    Java
8.1 基础模板
import java.io.*;
import java.util.*;
import java.math.*;
public class Main {
    public static void main(String[] args) {
        InputStream inputStream = System.in;
        OutputStream outputStream = System.out;
        InputReader in = new InputReader(inputStream);
        PrintWriter out = new PrintWriter(outputStream);
        Task solver = new Task();
        solver.solve(0, in, out);
        out.close();
    }
}
class Task {
    public void solve(int testNumber, InputReader in, PrintWriter out) {
}
class InputReader {
    public BufferedReader reader;
    public StringTokenizer tokenizer;
    public InputReader(InputStream stream) {
        reader = new BufferedReader(new InputStreamReader(stream), 32768);
        tokenizer = null;
    public String next() {
        while (tokenizer == null || !tokenizer.hasMoreTokens()) {
            try {
                tokenizer = new StringTokenizer(reader.readLine());
            } catch (IOException e) {
                throw new RuntimeException(e);
            }
        }
        return tokenizer.nextToken();
    public int nextInt() {
        return Integer.parseInt(next());
    public long nextLong() {
        return Long.parseLong(next());
```

}

- 8.2 BigInteger
- 8.3 BigDecimal
- 9 数学
- 9.1 常用积分表
- 9.2 常用数学公式
- 9.2.1 求和公式

1.
$$\sum_{k=1}^{n} (2k-1)^2 = \frac{n(4n^2-1)}{3}$$

2.
$$\sum_{k=1}^{n} k^3 = \left[\frac{n(n+1)}{2}\right]^2$$

3.
$$\sum_{k=1}^{n} (2k-1)^3 = n^2(2n^2-1)$$

4.
$$\sum_{k=1}^{n} k^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

5.
$$\sum_{k=1}^{n} k^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12}$$

6.
$$\sum_{k=1}^{n} k(k+1) = \frac{n(n+1)(n+2)}{3}$$

7.
$$\sum_{k=1}^{n} k(k+1)(k+2) = \frac{n(n+1)(n+2)(n+3)}{4}$$

8.
$$\sum_{k=1}^{n} k(k+1)(k+2)(k+3) = \frac{n(n+1)(n+2)(n+3)(n+4)}{5}$$

9.2.2 斐波那契数列

1.
$$fib_0 = 0, fib_1 = 1, fib_n = fib_{n-1} + fib_{n-2}$$

2.
$$fib_{n+2} \cdot fib_n - fib_{n+1}^2 = (-1)^{n+1}$$

3.
$$fib_{-n} = (-1)^{n-1} fib_n$$

4.
$$fib_{n+k} = fib_k \cdot fib_{n+1} + fib_{k-1} \cdot fib_n$$

5.
$$gcd(fib_m, fib_n) = fib_{gcd(m,n)}$$

6.
$$fib_m|fib_n^2 \Leftrightarrow nfib_n|m$$

9.2.3 错排公式

1.
$$D_n = (n-1)(D_{n-2} - D_{n-1})$$

2.
$$D_n = n! \cdot \left(1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + \frac{(-1)^n}{n!}\right)$$

9.3 平面几何公式

9.3.1 三角形

1. 半周长

$$p = \frac{a+b+c}{2}$$

2. 面积

$$S = \frac{a \cdot H_a}{2} = \frac{ab \cdot sinC}{2} = \sqrt{p(p-a)(p-b)(p-c)}$$

3. 中线

$$M_a = \frac{\sqrt{2(b^2 + c^2) - a^2}}{2} = \frac{\sqrt{b^2 + c^2 + 2bc \cdot cosA}}{2}$$

4. 角平分线

$$T_a = \frac{\sqrt{bc \cdot \left[(b+c)^2 - a^2 \right]}}{b+c} = \frac{2bc}{b+c} cos \frac{A}{2}$$

5. 高线

$$H_a = bsinC = csinB = \sqrt{b^2 - (\frac{a^2 + b^2 - c^2}{2a})^2}$$

6. 内切圆半径

$$\begin{split} r &= \frac{S}{p} = \frac{\arcsin\frac{B}{2} \cdot \sin\frac{C}{2}}{\sin\frac{B+C}{2}} = 4R \cdot \sin\frac{A}{2} \sin\frac{B}{2} \sin\frac{C}{2} \\ &= \sqrt{\frac{(p-a)(p-b)(p-c)}{p}} = p \cdot \tan\frac{A}{2} \tan\frac{B}{2} \tan\frac{C}{2} \end{split}$$

7. 外接圆半径

$$R = \frac{abc}{4S} = \frac{a}{2sinA} = \frac{b}{2sinB} = \frac{c}{2sinC}$$

- 9.3.2 四边形
- 9.3.3 正 n 边形
- 9.3.4 圆
- 9.3.5 棱柱
- 9.3.6 棱锥
- 9.3.7 棱台
- 9.3.8 圆柱
- 9.3.9 圆锥
- 9.3.10 圆台
- 9.4 常用数表
- 9.4.1 梅森数