

# 代码库

上海交通大学

September 5, 2015

## Contents

1	数论	4
1.1	快速求逆元	4
1.2	扩展欧几里德算法	4
1.3	中国剩余定理	4
1.4	Miller Rabin 素数测试	5
1.5	Pollard Rho 大数分解	5
1.6	快速数论变换	6
1.7	原根	7
1.8	离散对数	7
1.9	离散平方根	7
1.10	佩尔方程求解	7
1.11	牛顿迭代法	7
1.12	直线下整点个数	7
2	数值	7
2.1	高斯消元	7
2.2	快速傅立叶变换	7
2.3	单纯形法求解线性规划	8
2.4	自适应辛普森	9
2.5	多项式方程求解	10
2.6	最小二乘法	10
3	数据结构	10
3.1	平衡的二叉查找树	10
3.1.1	Treap	10
3.1.2	Splay	11
3.2	坚固的数据结构	11
3.2.1	坚固的线段树	11
3.2.2	坚固的平衡树	12
3.2.3	坚固的字符串	13
3.2.4	坚固的左偏树	13
3.3	树上的魔术师	13
3.3.1	轻重树链剖分	13
3.3.2	Link Cut Tree	15
3.3.3	AAA Tree	15
3.4	k-d 树	15
4	图论	15
4.1	强连通分量	15
4.2	双连通分量	15
4.2.1	点双连通分量	15
4.2.2	边双连通分量	15
4.3	2-SAT 问题	15
4.4	二分图最大匹配	15
4.4.1	Hungary 算法	15

4.4.2	Hopcroft Karp 算法	15
4.5	二分图最大权匹配	16
4.5.1	KM 算法	16
4.5.2	扩展 KM 算法	18
4.6	最大流	18
4.7	最小费用最大流	19
4.7.1	稀疏图	19
4.7.2	稠密图	20
4.8	一般图最大匹配	22
4.9	无向图全局最小割	24
4.10	最小树形图	25
4.11	有根树的同构	25
4.12	度限制生成树	26
4.13	弦图相关	26
4.13.1	弦图的判定	26
4.13.2	弦图的团数	26
4.14	哈密尔顿回路 (ORE 性质的图)	26
5	字符串	27
5.1	模式匹配	27
5.1.1	KMP 算法	27
5.1.2	扩展 KMP 算法	28
5.1.3	AC 自动机	28
5.2	后缀三姐妹	28
5.2.1	后缀数组	28
5.2.2	后缀自动机	28
5.3	回文三兄弟	29
5.3.1	Manacher 算法	29
5.3.2	回文树	29
5.4	循环串最小表示	30
6	计算几何	30
6.1	二维基础	30
6.1.1	点类	30
6.1.2	凸包	30
6.1.3	半平面交	31
6.2	三维基础	31
6.2.1	点类	31
6.2.2	凸包	31
6.2.3	绕轴旋转	31
6.3	多边形	31
6.3.1	判断点在多边形内部	31
6.3.2	旋转卡壳	32
6.3.3	动态凸包	32
6.3.4	点到凸包的切线	32
6.3.5	直线与凸包的交点	32
6.3.6	凸多边形的交集	32
6.3.7	凸多边形内的最大圆	32
6.4	圆	32
6.4.1	圆类	32
6.4.2	圆的交集	32
6.4.3	最小覆盖圆	32
6.4.4	最小覆盖球	32
6.4.5	判断圆存在交集	32
6.4.6	圆与多边形的交集	32
6.5	三角形	32
6.5.1	三角形的内心	32
6.5.2	三角形的外心	32

6.5.3	三角形的垂心	32
6.6	黑暗科技	32
6.6.1	平面图形的转动惯量	32
6.6.2	平面区域处理	32
6.6.3	Vonoroi 图	32
7	其他	32
7.1	某年某月某日是星期几	32
7.2	动态规划	33
7.3	搜索	33
7.3.1	Dancing Links X	33
8	Java	33
8.1	基础模板	33
8.2	BigInteger	34
8.3	BigDecimal	34
9	数学	34
9.1	常用积分表	34
9.2	常用数学公式	34
9.2.1	求和公式	34
9.2.2	斐波那契数列	34
9.2.3	错排公式	34
9.3	平面几何公式	34
9.3.1	三角形	34
9.3.2	四边形	35
9.3.3	正 $n$ 边形	35
9.3.4	圆	36
9.3.5	棱柱	36
9.3.6	棱锥	36
9.3.7	棱台	36
9.3.8	圆柱	37
9.3.9	圆锥	37
9.3.10	圆台	37
9.3.11	球	37
9.3.12	球台	37
9.3.13	球扇形	38
9.4	常用数表	38
9.4.1	梅森数	38

# 1 数论

## 1.1 快速求逆元

返回结果：

$$x^{-1}(\text{mod})$$

使用条件： $x \in [0, \text{mod})$  并且  $x$  与  $\text{mod}$  互质。

```
long long inverse(const long long &x, const long long &mod) {
    if (x == 1) {
        return 1;
    } else {
        return (mod - mod / x) * inverse(mod % x, mod) % mod;
    }
}
```

## 1.2 扩展欧几里德算法

返回结果：

$$ax + by = \text{gcd}(a, b)$$

时间复杂度： $O(n \log n)$

```
void solve(const long long &a, const long long &b, long long &x, long long &y) {
    if (b == 0) {
        x = 1;
        y = 0;
    } else {
        solve(b, a % b, x, y);
        x -= a / b * y;
        std::swap(x, y);
    }
}
```

## 1.3 中国剩余定理

返回结果：

$$x \equiv r_i(\text{mod } p_i) \quad (0 \leq i < n)$$

使用条件： $p_i$  无需两两互质

时间复杂度： $O(n \log n)$

```
bool solve(int n, std::pair<long long, long long> input[], std::pair<long long, long long> &output) {
    output = std::make_pair(1, 1);
    for (int i = 0; i < n; ++i) {
        long long number, useless;
        euclid(output.second, input[i].second, number, useless);
        long long divisor = std::__gcd(output.second, input[i].second);
        if ((input[i].first - output.first) % divisor) {
            return false;
        }
        number *= (input[i].first - output.first) / divisor;
        fix(number, input[i].second);
        output.first += output.second * number;
        output.second *= input[i].second / divisor;
        fix(output.first, output.second);
    }
    return true;
}
```

## 1.4 Miller Rabin 素数测试

```
const int BASE[12] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};

bool check(const long long &prime, const long long &base) {
    long long number = prime - 1;
    for (; ~number & 1; number >>= 1);
    long long result = power_mod(base, number, prime);
    for (; number != prime - 1 && result != 1 && result != prime - 1; number <<= 1) {
        result = multiply_mod(result, result, prime);
    }
    return result == prime - 1 || (number & 1) == 1;
}

bool miller_rabin(const long long &number) {
    if (number < 2) {
        return false;
    }
    if (number < 4) {
        return true;
    }
    if (~number & 1) {
        return false;
    }
    for (int i = 0; i < 12 && BASE[i] < number; ++i) {
        if (!check(number, BASE[i])) {
            return false;
        }
    }
    return true;
}
```

## 1.5 Pollard Rho 大数分解

时间复杂度： $O(n^{1/4})$

```
long long pollard_rho(const long long &number, const long long &seed) {
    long long x = rand() % (number - 1) + 1, y = x;
    for (int head = 1, tail = 2; ; ) {
        x = multiply_mod(x, x, number);
        x = add_mod(x, seed, number);
        if (x == y) {
            return number;
        }
        long long answer = std::__gcd(abs(x - y), number);
        if (answer > 1 && answer < number) {
            return answer;
        }
        if (++head == tail) {
            y = x;
            tail <<= 1;
        }
    }
}

void factorize(const long long &number, std::vector<long long> &divisor) {
    if (number > 1) {
        if (miller_rabin(number)) {
            divisor.push_back(number);
        }
    }
}
```

```

    } else {
        long long factor = number;
        for (; factor >= number; factor = pollard_rho(number, rand() % (number - 1) + 1));
        factorize(number / factor, divisor);
        factorize(factor, divisor);
    }
}
}

```

## 1.6 快速数论变换

返回结果：

$$c_i = \sum_{0 \leq j \leq i} a_j \cdot b_{i-j}(\text{mod}) \quad (0 \leq i < n)$$

使用说明：*magic* 是 *mod* 的原根

时间复杂度： $O(n \log n)$

```

void solve(long long number[], int length, int type) {
    for (int i = 1, j = 0; i < length - 1; ++i) {
        for (int k = length; j ^= k >>= 1, ~j & k; );
        if (i < j) {
            std::swap(number[i], number[j]);
        }
    }
    long long unit_p0;
    for (int turn = 0; (1 << turn) < length; ++turn) {
        int step = 1 << turn, step2 = step << 1;
        if (type == 1) {
            unit_p0 = power_mod(MAGIC, (MOD - 1) / step2, MOD);
        } else {
            unit_p0 = power_mod(MAGIC, MOD - 1 - (MOD - 1) / step2, MOD);
        }
        for (int i = 0; i < length; i += step2) {
            long long unit = 1;
            for (int j = 0; j < step; ++j) {
                long long &number1 = number[i + j + step];
                long long &number2 = number[i + j];
                long long delta = unit * number1 % MOD;
                number1 = (number2 - delta + MOD) % MOD;
                number2 = (number2 + delta) % MOD;
                unit = unit * unit_p0 % MOD;
            }
        }
    }
}

void multiply() {
    for (; lowbit(length) != length; ++length);
    solve(number1, length, 1);
    solve(number2, length, 1);
    for (int i = 0; i < length; ++i) {
        number[i] = number1[i] * number2[i] % MOD;
    }
    solve(number, length, -1);
    for (int i = 0; i < length; ++i) {
        answer[i] = number[i] * power_mod(length, MOD - 2, MOD) % MOD;
    }
}

```

## 1.7 原根

## 1.8 离散对数

## 1.9 离散平方根

## 1.10 佩尔方程求解

## 1.11 牛顿迭代法

## 1.12 直线下整点个数

返回结果：

$$\sum_{0 \leq i < n} \lfloor \frac{a + b \cdot i}{m} \rfloor$$

使用条件： $n, m > 0, a, b \geq 0$

时间复杂度： $O(n \log n)$

```
long long solve(const long long &n, const long long &a, const long long &b, const long long &m) {
    if (b == 0) {
        return n * (a / m);
    }
    if (a >= m) {
        return n * (a / m) + solve(n, a % m, b, m);
    }
    if (b >= m) {
        return (n - 1) * n / 2 * (b / m) + solve(n, a, b % m, m);
    }
    return solve((a + b * n) / m, (a + b * n) % m, m, b);
}
```

## 2 数值

### 2.1 高斯消元

### 2.2 快速傅立叶变换

返回结果：

$$c_i = \sum_{0 \leq j \leq i} a_j \cdot b_{i-j} \quad (0 \leq i < n)$$

时间复杂度： $O(n \log n)$

```
void solve(Complex number[], int length, int type) {
    for (int i = 1, j = 0; i < length - 1; ++i) {
        for (int k = length; j ^= k >>= 1, ~j & k; );
        if (i < j) {
            std::swap(number[i], number[j]);
        }
    }
    Complex unit_p0;
    for (int turn = 0; (1 << turn) < length; ++turn) {
        int step = 1 << turn, step2 = step << 1;
        double p0 = PI / step * type;
        sincos(p0, &unit_p0.imag(), &unit_p0.real());
        for (int i = 0; i < length; i += step2) {
            Complex unit = 1;
            for (int j = 0; j < step; ++j) {
                Complex &number1 = number[i + j + step];
            }
        }
    }
}
```

```

        Complex &number2 = number[i + j];
        Complex delta = unit * number1;
        number1 = number2 - delta;
        number2 = number2 + delta;
        unit = unit * unit_p0;
    }
}
}

void multiply() {
    for (; lowbit(length) != length; ++length);
    solve(number1, length, 1);
    solve(number2, length, 1);
    for (int i = 0; i < length; ++i) {
        number[i] = number1[i] * number2[i];
    }
    solve(number, length, -1);
    for (int i = 0; i < length; ++i) {
        answer[i] = (int)(number[i].real() / length + 0.5);
    }
}

```

## 2.3 单纯形法求解线性规划

返回结果：

$$\max\{c_{1 \times m} \cdot x_{m \times 1} \mid x_{m \times 1} \geq 0_{m \times 1}, a_{n \times m} \cdot x_{m \times 1} \leq b_{n \times 1}\}$$

```

std::vector<double> solve(const std::vector<std::vector<double>> &a,
                        const std::vector<double> &b, const std::vector<double> &c) {
    int n = (int)a.size(), m = (int)a[0].size() + 1;
    std::vector<std::vector<double>> value(n + 2, std::vector<double>(m + 1));
    std::vector<int> index(n + m);
    int r = n, s = m - 1;
    for (int i = 0; i < n + m; ++i) {
        index[i] = i;
    }
    for (int i = 0; i < n; ++i) {
        for (int j = 0; j < m - 1; ++j) {
            value[i][j] = -a[i][j];
        }
        value[i][m - 1] = 1;
        value[i][m] = b[i];
        if (value[r][m] > value[i][m]) {
            r = i;
        }
    }
    for (int j = 0; j < m - 1; ++j) {
        value[n][j] = c[j];
    }
    value[n + 1][m - 1] = -1;
    for (double number; ; ) {
        if (r < n) {
            std::swap(index[s], index[r + m]);
            value[r][s] = 1 / value[r][s];
            for (int j = 0; j <= m; ++j) {
                if (j != s) {
                    value[r][j] *= -value[r][s];
                }
            }
        }
    }
}

```



```

    }
}
for (int i = 0; i <= n + 1; ++i) {
    if (i != r) {
        for (int j = 0; j <= m; ++j) {
            if (j != s) {
                value[i][j] += value[r][j] * value[i][s];
            }
        }
        value[i][s] *= value[r][s];
    }
}
}
r = s = -1;
for (int j = 0; j < m; ++j) {
    if (s < 0 || index[s] > index[j]) {
        if (value[n + 1][j] > eps || value[n + 1][j] > -eps && value[n][j] > eps) {
            s = j;
        }
    }
}
if (s < 0) {
    break;
}
for (int i = 0; i < n; ++i) {
    if (value[i][s] < -eps) {
        if (r < 0
            || (number = value[r][m] / value[r][s] - value[i][m] / value[i][s]) < -eps
            || number < eps && index[r + m] > index[i + m]) {
            r = i;
        }
    }
}
if (r < 0) {
    // Solution is unbounded.
    return std::vector<double>();
}
}
if (value[n + 1][m] < -eps) {
    // No solution.
    return std::vector<double>();
}
std::vector<double> answer(m - 1);
for (int i = m; i < n + m; ++i) {
    if (index[i] < m - 1) {
        answer[index[i]] = value[i - m][m];
    }
}
return answer;
}
}

```

## 2.4 自适应辛普森

```

double area(const double &left, const double &right) {
    double mid = (left + right) / 2;
    return (right - left) * (calc(left) + 4 * calc(mid) + calc(right)) / 6;
}

```

```

double simpson(const double &left, const double &right, const double &eps, const double &area_sum) {
    double mid = (left + right) / 2;
    double area_left = area(left, mid);
    double area_right = area(mid, right);
    double area_total = area_left + area_right;
    if (std::abs(area_total - area_sum) < 15 * eps) {
        return area_total + (area_total - area_sum) / 15;
    }
    return simpson(left, mid, eps / 2, area_left) + simpson(mid, right, eps / 2, area_right);
}

double simpson(const double &left, const double &right, const double &eps) {
    return simpson(left, right, eps, area(left, right));
}

```

## 2.5 多项式方程求解

## 2.6 最小二乘法

# 3 数据结构

## 3.1 平衡的二叉查找树

### 3.1.1 Treap

```

class Node {
public:
    Node *child[2];
    int key;
    int size, priority;

    Node(Node *left, Node *right, int key) : key(key), size(1), priority(rand()) {
        child[0] = left;
        child[1] = right;
    }

    void update() {
        size = child[0]->size + 1 + child[1]->size;
    }
};

Node *null;

void rotate(Node *&x, int dir) {
    Node *y = x->child[dir];
    x->child[dir] = y->child[dir ^ 1];
    y->child[dir ^ 1] = x;
    x->update();
    y->update();
    x = y;
}

void insert(Node *&x, int key) {
    if (x == null) {
        x = new Node(null, null, key);
    } else {
        insert(x->child[key > x->key], key);
        if (x->child[key > x->key]->priority < x->priority) {

```

```

        rotate(x, key > x->key);
    }
    x->update();
}

void remove(Node *&x, int key) {
    if (x->key != key) {
        remove(x->child[key > x->key], key);
    } else if (x->child[0] == null && x->child[1] == null) {
        x = null;
        return;
    } else {
        int dir = x->child[0]->priority > x->child[1]->priority;
        rotate(x, dir);
        remove(x->child[dir ^ 1], key);
    }
    x->update();
}

void build() {
    null = new Node(NULL, NULL, 0);
    null->child[0] = null->child[1] = null;
    null->size = 0;
    null->priority = RAND_MAX;
}

```

### 3.1.2 Splay

## 3.2 坚固的数据结构

### 3.2.1 坚固的线段树

```

class Node {
public:
    Node *left, *right;
    int value;

    Node(Node *left, Node *right, int value) : left(left), right(right), value(value) {}

    Node* modify(int l, int r, int ql, int qr, int value);
    int query(int l, int r, int qx);
};

Node* null;

Node* Node::modify(int l, int r, int ql, int qr, int value) {
    if (qr < l || r < ql) {
        return this;
    }
    if (ql <= l && r <= qr) {
        return new Node(this->left, this->right, this->value + value);
    }
    int mid = l + r >> 1;
    return new Node(this->left->modify(l, mid, ql, qr, value),
                    this->right->modify(mid + 1, r, ql, qr, value),
                    this->value);
}

```

```

int Node::query(int l, int r, int qx) {
    if (qx < l || r < qx) {
        return 0;
    }
    if (qx <= l && r <= qx) {
        return this->value;
    }
    int mid = l + r >> 1;
    return this->left->query(l, mid, qx)
        + this->right->query(mid + 1, r, qx)
        + this->value;
}

void build() {
    null = new Node(NULL, NULL, 0);
    null->left = null->right = null;
}

```

### 3.2.2 坚固的平衡树

```

class Node {
public:
    Node *left, *right;
    int size;

    Node();
    std::pair<Node*, Node*> split(int size);

    Node* update() {
        size = left->size + 1 + right->size;
        return this;
    }
};

bool random(int a, int b) {
    return rand() % (a + b) < a;
}

Node *null;

Node::Node() : left(null), right(null), size(1) {}

Node* merge(Node *x, Node *y) {
    if (x == null) {
        return y;
    }
    if (y == null) {
        return x;
    }
    if (random(x->size, y->size)) {
        x->right = merge(x->right, y);
        return x->update();
    } else {
        y->left = merge(x, y->left);
        return y->update();
    }
}

```

```

std::pair<Node*, Node*> Node::split(int size) {
    if (this == null) {
        return std::make_pair(null, null);
    }
    if (size <= left->size) {
        std::pair<Node*, Node*> result = left->split(size);
        left = null;
        return std::make_pair(result.first, merge(result.second, this->update()));
    } else {
        std::pair<Node*, Node*> result = right->split(size - left->size - 1);
        right = null;
        return std::make_pair(merge(this->update(), result.first), result.second);
    }
}

void build() {
    null = new Node();
    null->size = 0;
}

```

### 3.2.3 坚固的字符串

### 3.2.4 坚固的左偏树

## 3.3 树上的魔术师

### 3.3.1 轻重树链剖分

```

int father[N], height[N], size[N], son[N], top[N], pos[N], data[N];

void build(int root) {
    std::vector<int> queue;
    father[root] = -1;
    height[root] = 0;
    queue.push_back(root);
    for (int head = 0; head < (int)queue.size(); ++head) {
        int x = queue[head];
        for (int i = 0; i < (int)edge[x].size(); ++i) {
            int y = edge[x][i];
            if (y != father[x]) {
                father[y] = x;
                height[y] = height[x] + 1;
                queue.push_back(y);
            }
        }
    }
    for (int index = n - 1; index >= 0; --index) {
        int x = queue[index];
        size[x] = 1;
        son[x] = -1;
        for (int i = 0; i < (int)edge[x].size(); ++i) {
            int y = edge[x][i];
            if (y != father[x]) {
                size[x] += size[y];
                if (son[x] == -1 || size[son[x]] < size[y]) {
                    son[x] = y;
                }
            }
        }
    }
}

```

```

    }
}
std::fill(top, top + n, 0);
int counter = 0;
for (int index = 0; index < n; ++index) {
    int x = queue[index];
    if (top[x] == 0) {
        for (int y = x; y != -1; y = son[y]) {
            top[y] = x;
            pos[y] = ++counter;
            data[counter] = value[y];
        }
    }
}
build(1, 1, n);
}

void solve(int x, int y) {
    while (true) {
        if (top[x] == top[y]) {
            if (x == y) {
                solve(1, 1, n, pos[x], pos[x]);
            } else {
                if (height[x] < height[y]) {
                    solve(1, 1, n, pos[x], pos[y]);
                } else {
                    solve(1, 1, n, pos[y], pos[x]);
                }
            }
            break;
        }
        if (height[top[x]] > height[top[y]]) {
            solve(1, 1, n, pos[top[x]], pos[x]);
            x = father[top[x]];
        } else {
            solve(1, 1, n, pos[top[y]], pos[y]);
            y = father[top[y]];
        }
    }
}
}

```

### 3.3.2 Link Cut Tree

### 3.3.3 AAA Tree

## 3.4 k-d 树

# 4 图论

## 4.1 强连通分量

## 4.2 双连通分量

### 4.2.1 点双连通分量

### 4.2.2 边双连通分量

## 4.3 2-SAT 问题

## 4.4 二分图最大匹配

### 4.4.1 Hungary 算法

时间复杂度： $O(V \cdot E)$

```
int n, m, stamp;
int match[N], visit[N];

bool dfs(int x) {
    for (int i = 0; i < (int)edge[x].size(); ++i) {
        int y = edge[x][i];
        if (visit[y] != stamp) {
            visit[y] = stamp;
            if (match[y] == -1 || dfs(match[y])) {
                match[y] = x;
                return true;
            }
        }
    }
    return false;
}

int solve() {
    std::fill(match, match + m, -1);
    int answer = 0;
    for (int i = 0; i < n; ++i) {
        stamp++;
        answer += dfs(i);
    }
    return answer;
}
```

### 4.4.2 Hopcroft Karp 算法

时间复杂度： $O(\sqrt{V} \cdot E)$

```
int matchx[N], matchy[N], level[N];

bool dfs(int x) {
    for (int i = 0; i < (int)edge[x].size(); ++i) {
        int y = edge[x][i];
```

```

        int w = matchy[y];
        if (w == -1 || level[x] + 1 == level[w] && dfs(w)) {
            matchx[x] = y;
            matchy[y] = x;
            return true;
        }
    }
    level[x] = -1;
    return false;
}

int solve() {
    std::fill(matchx, matchx + n, -1);
    std::fill(matchy, matchy + m, -1);
    for (int answer = 0; ; ) {
        std::vector<int> queue;
        for (int i = 0; i < n; ++i) {
            if (matchx[i] == -1) {
                level[i] = 0;
                queue.push_back(i);
            } else {
                level[i] = -1;
            }
        }
        for (int head = 0; head < (int)queue.size(); ++head) {
            int x = queue[head];
            for (int i = 0; i < (int)edge[x].size(); ++i) {
                int y = edge[x][i];
                int w = matchy[y];
                if (w != -1 && level[w] < 0) {
                    level[w] = level[x] + 1;
                    queue.push_back(w);
                }
            }
        }
        int delta = 0;
        for (int i = 0; i < n; ++i) {
            if (matchx[i] == -1 && dfs(i)) {
                delta++;
            }
        }
        if (delta == 0) {
            return answer;
        } else {
            answer += delta;
        }
    }
}

```

## 4.5 二分图最大权匹配

### 4.5.1 KM 算法

时间复杂度： $O(V^3)$

```

int labelx[N], labely[N], match[N], slack[N];
bool visitx[N], visity[N];

bool dfs(int x) {

```



```

visitx[x] = true;
for (int y = 0; y < n; ++y) {
    if (visity[y]) {
        continue;
    }
    int delta = labelx[x] + labely[y] - graph[x][y];
    if (delta == 0) {
        visity[y] = true;
        if (match[y] == -1 || dfs(match[y])) {
            match[y] = x;
            return true;
        }
    } else {
        slack[y] = std::min(slack[y], delta);
    }
}
return false;
}

int solve() {
    for (int i = 0; i < n; ++i) {
        match[i] = -1;
        labelx[i] = INT_MIN;
        labely[i] = 0;
        for (int j = 0; j < n; ++j) {
            labelx[i] = std::max(labelx[i], graph[i][j]);
        }
    }
    for (int i = 0; i < n; ++i) {
        while (true) {
            std::fill(visitx, visitx + n, 0);
            std::fill(visity, visity + n, 0);
            for (int j = 0; j < n; ++j) {
                slack[j] = INT_MAX;
            }
            if (dfs(i)) {
                break;
            }
            int delta = INT_MAX;
            for (int j = 0; j < n; ++j) {
                if (!visity[j]) {
                    delta = std::min(delta, slack[j]);
                }
            }
            for (int j = 0; j < n; ++j) {
                if (visitx[j]) {
                    labelx[j] -= delta;
                }
                if (visity[j]) {
                    labely[j] += delta;
                } else {
                    slack[j] -= delta;
                }
            }
        }
    }
}

int answer = 0;
for (int i = 0; i < n; ++i) {

```

```

        answer += graph[match[i]][i];
    }
    return answer;
}

```

#### 4.5.2 扩展 KM 算法

### 4.6 最大流

时间复杂度： $O(V^2 \cdot E)$

```

struct EdgeList {
    int size;
    int last[N];
    int succ[M], other[M], flow[M];
    void clear(int n) {
        size = 0;
        fill(last, last + n, -1);
    }
    void add(int x, int y, int c) {
        succ[size] = last[x];
        last[x] = size;
        other[size] = y;
        flow[size++] = c;
    }
} e;

int n, source, target;
int dist[N], curr[N];

void add(int x, int y, int c) {
    e.add(x, y, c);
    e.add(y, x, 0);
}

bool relabel() {
    std::vector<int> queue;
    for (int i = 0; i < n; ++i) {
        curr[i] = e.last[i];
        dist[i] = -1;
    }
    queue.push_back(target);
    dist[target] = 0;
    for (int head = 0; head < (int)queue.size(); ++head) {
        int x = queue[head];
        for (int i = e.last[x]; ~i; i = e.succ[i]) {
            int y = e.other[i];
            if (e.flow[i ^ 1] && dist[y] == -1) {
                dist[y] = dist[x] + 1;
                queue.push_back(y);
            }
        }
    }
    return ~dist[source];
}

int dfs(int x, int answer) {
    if (x == target) {
        return answer;
    }

```

```

    }
    int delta = answer;
    for (int &i = curr[x]; ~i; i = e.succ[i]) {
        int y = e.other[i];
        if (e.flow[i] && dist[x] == dist[y] + 1) {
            int number = dfs(y, std::min(e.flow[i], delta));
            e.flow[i] -= number;
            e.flow[i ^ 1] += number;
            delta -= number;
        }
        if (delta == 0) {
            break;
        }
    }
    return answer - delta;
}

int solve() {
    int answer = 0;
    while (relabel()) {
        answer += dfs(source, INT_MAX);
    }
    return answer;
}

```

## 4.7 最小费用最大流

### 4.7.1 稀疏图

时间复杂度： $O(V \cdot E^2)$

```

struct EdgeList {
    int size;
    int last[N];
    int succ[M], other[M], flow[M], cost[M];
    void clear(int n) {
        size = 0;
        std::fill(last, last + n, -1);
    }
    void add(int x, int y, int c, int w) {
        succ[size] = last[x];
        last[x] = size;
        other[size] = y;
        flow[size] = c;
        cost[size++] = w;
    }
} e;

int n, source, target;
int prev[N];

void add(int x, int y, int c, int w) {
    e.add(x, y, c, w);
    e.add(y, x, 0, -w);
}

bool augment() {
    static int dist[N], occur[N];
    std::vector<int> queue;

```

```

std::fill(dist, dist + n, INT_MAX);
std::fill(occur, occur + n, 0);
dist[source] = 0;
occur[source] = true;
queue.push_back(source);
for (int head = 0; head < (int)queue.size(); ++head) {
    int x = queue[head];
    for (int i = e.last[x]; ~i; i = e.succ[i]) {
        int y = e.other[i];
        if (e.flow[i] && dist[y] > dist[x] + e.cost[i]) {
            dist[y] = dist[x] + e.cost[i];
            prev[y] = i;
            if (!occur[y]) {
                occur[y] = true;
                queue.push_back(y);
            }
        }
    }
    occur[x] = false;
}
return dist[target] < INT_MAX;
}

std::pair<int, int> solve() {
    std::pair<int, int> answer = std::make_pair(0, 0);
    while (augment()) {
        int number = INT_MAX;
        for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
            number = std::min(number, e.flow[prev[i]]);
        }
        answer.first += number;
        for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
            e.flow[prev[i]] -= number;
            e.flow[prev[i] ^ 1] += number;
            answer.second += number * e.cost[prev[i]];
        }
    }
    return answer;
}

```

#### 4.7.2 稠密图

时间复杂度： $O(V \cdot E^2)$

```

struct EdgeList {
    int size;
    int last[N];
    int succ[M], other[M], flow[M], cost[M];
    void clear(int n) {
        size = 0;
        std::fill(last, last + n, -1);
    }
    void add(int x, int y, int c, int w) {
        succ[size] = last[x];
        last[x] = size;
        other[size] = y;
        flow[size] = c;
        cost[size++] = w;
    }
}

```

```

    }
} e;

int n, source, target, flow, cost;
int slack[N], dist[N];
bool visit[N];

void add(int x, int y, int c, int w) {
    e.add(x, y, c, w);
    e.add(y, x, 0, -w);
}

bool relabel() {
    int delta = INT_MAX;
    for (int i = 0; i < n; ++i) {
        if (!visit[i]) {
            delta = std::min(delta, slack[i]);
        }
        slack[i] = INT_MAX;
    }
    if (delta == INT_MAX) {
        return true;
    }
    for (int i = 0; i < n; ++i) {
        if (visit[i]) {
            dist[i] += delta;
        }
    }
    return false;
}

int dfs(int x, int answer) {
    if (x == target) {
        flow += answer;
        cost += answer * (dist[source] - dist[target]);
        return answer;
    }
    visit[x] = true;
    int delta = answer;
    for (int i = e.last[x]; ~i; i = e.succ[i]) {
        int y = e.other[i];
        if (e.flow[i] > 0 && !visit[y]) {
            if (dist[y] + e.cost[i] == dist[x]) {
                int number = dfs(y, std::min(e.flow[i], delta));
                e.flow[i] -= number;
                e.flow[i ^ 1] += number;
                delta -= number;
                if (delta == 0) {
                    dist[x] = INT_MIN;
                    return answer;
                }
            } else {
                slack[y] = std::min(slack[y], dist[y] + e.cost[i] - dist[x]);
            }
        }
    }
    return answer - delta;
}

```

```

std::pair<int, int> solve() {
    flow = cost = 0;
    std::fill(dist, dist + n, 0);
    do {
        do {
            fill(visit, visit + n, 0);
        } while (dfs(source, INT_MAX));
    } while (!relabel());
    return std::make_pair(flow, cost);
}

```

## 4.8 一般图最大匹配

时间复杂度： $O(V^3)$

```

int match[N], belong[N], next[N], mark[N], visit[N];
std::vector<int> queue;

```

```

int find(int x) {
    if (belong[x] != x) {
        belong[x] = find(belong[x]);
    }
    return belong[x];
}

```

```

void merge(int x, int y) {
    x = find(x);
    y = find(y);
    if (x != y) {
        belong[x] = y;
    }
}

```

```

int lca(int x, int y) {
    static int stamp = 0;
    stamp++;
    while (true) {
        if (x != -1) {
            x = find(x);
            if (visit[x] == stamp) {
                return x;
            }
            visit[x] = stamp;
            if (match[x] != -1) {
                x = next[match[x]];
            } else {
                x = -1;
            }
        }
        std::swap(x, y);
    }
}

```

```

void group(int a, int p) {
    while (a != p) {
        int b = match[a], c = next[b];
        if (find(c) != p) {

```

```

        next[c] = b;
    }
    if (mark[b] == 2) {
        mark[b] = 1;
        queue.push_back(b);
    }
    if (mark[c] == 2) {
        mark[c] = 1;
        queue.push_back(c);
    }
    merge(a, b);
    merge(b, c);
    a = c;
}
}

void augment(int source) {
    queue.clear();
    for (int i = 0; i < n; ++i) {
        next[i] = visit[i] = -1;
        belong[i] = i;
        mark[i] = 0;
    }
    mark[source] = 1;
    queue.push_back(source);
    for (int head = 0; head < (int)queue.size() && match[source] == -1; ++head) {
        int x = queue[head];
        for (int i = 0; i < (int)edge[x].size(); ++i) {
            int y = edge[x][i];
            if (match[x] == y || find(x) == find(y) || mark[y] == 2) {
                continue;
            }
            if (mark[y] == 1) {
                int r = lca(x, y);
                if (find(x) != r) {
                    next[x] = y;
                }
                if (find(y) != r) {
                    next[y] = x;
                }
                group(x, r);
                group(y, r);
            } else if (match[y] == -1) {
                next[y] = x;
                for (int u = y; u != -1; ) {
                    int v = next[u];
                    int mv = match[v];
                    match[v] = u;
                    match[u] = v;
                    u = mv;
                }
                break;
            } else {
                next[y] = x;
                mark[y] = 2;
                mark[match[y]] = 1;
                queue.push_back(match[y]);
            }
        }
    }
}

```

```

    }
}

int solve() {
    std::fill(match, match + n, -1);
    for (int i = 0; i < n; ++i) {
        if (match[i] == -1) {
            augment(i);
        }
    }
    int answer = 0;
    for (int i = 0; i < n; ++i) {
        answer += (match[i] != -1);
    }
    return answer;
}

```

## 4.9 无向图全局最小割

时间复杂度： $O(V^3)$

注意事项：处理重边时，应该对边权累加

```

int node[N], dist[N];
bool visit[N];

int solve(int n) {
    int answer = INT_MAX;
    for (int i = 0; i < n; ++i) {
        node[i] = i;
    }
    while (n > 1) {
        int max = 1;
        for (int i = 0; i < n; ++i) {
            dist[node[i]] = graph[node[0]][node[i]];
            if (dist[node[i]] > dist[node[max]]) {
                max = i;
            }
        }
        int prev = 0;
        memset(visit, 0, sizeof(visit));
        visit[node[0]] = true;
        for (int i = 1; i < n; ++i) {
            if (i == n - 1) {
                answer = std::min(answer, dist[node[max]]);
                for (int k = 0; k < n; ++k) {
                    graph[node[k]][node[prev]] = (graph[node[prev]][node[k]] += graph[node[k]][node[prev]]);
                }
                node[max] = node[--n];
            }
            visit[node[max]] = true;
            prev = max;
            max = -1;
            for (int j = 1; j < n; ++j) {
                if (!visit[node[j]]) {
                    dist[node[j]] += graph[node[prev]][node[j]];
                    if (max == -1 || dist[node[max]] < dist[node[j]]) {
                        max = j;
                    }
                }
            }
        }
    }
}

```



```

    }
    }
    }
    }
    }
    return answer;
}

```

#### 4.10 最小树形图

#### 4.11 有根树的同构

时间复杂度： $O(V \log V)$

```

const unsigned long long MAGIC = 4423;

unsigned long long magic[N];
std::pair<unsigned long long, int> hash[N];

void solve(int root) {
    magic[0] = 1;
    for (int i = 1; i <= n; ++i) {
        magic[i] = magic[i - 1] * MAGIC;
    }
    std::vector<int> queue;
    queue.push_back(root);
    for (int head = 0; head < (int)queue.size(); ++head) {
        int x = queue[head];
        for (int i = 0; i < (int)son[x].size(); ++i) {
            int y = son[x][i];
            queue.push_back(y);
        }
    }
    for (int index = n - 1; index >= 0; --index) {
        int x = queue[index];
        hash[x] = std::make_pair(0, 0);

        std::vector<std::pair<unsigned long long, int> > value;
        for (int i = 0; i < (int)son[x].size(); ++i) {
            int y = son[x][i];
            value.push_back(hash[y]);
        }
        std::sort(value.begin(), value.end());

        hash[x].first = hash[x].first * magic[1] + 37;
        hash[x].second++;
        for (int i = 0; i < (int)value.size(); ++i) {
            hash[x].first = hash[x].first * magic[value[i].second] + value[i].first;
            hash[x].second += value[i].second;
        }
        hash[x].first = hash[x].first * magic[1] + 41;
        hash[x].second++;
    }
}

```

## 4.12 度限制生成树

## 4.13 弦图相关

### 4.13.1 弦图的判定

### 4.13.2 弦图的团数

## 4.14 哈密尔顿回路 ( ORE 性质的图 )

ORE 性质 :

$$\forall x, y \in V \wedge (x, y) \notin E \text{ s.t. } \deg_x + \deg_y \geq n$$

返回结果 : 从顶点 1 出发的一个哈密尔顿回路

使用条件 :  $n \geq 3$

```
int left[N], right[N], next[N], last[N];

void cover(int x) {
    left[right[x]] = left[x];
    right[left[x]] = right[x];
}

int adjacent(int x) {
    for (int i = right[0]; i <= n; i = right[i]) {
        if (graph[x][i]) {
            return i;
        }
    }
    return 0;
}

std::vector<int> solve() {
    for (int i = 1; i <= n; ++i) {
        left[i] = i - 1;
        right[i] = i + 1;
    }
    int head, tail;
    for (int i = 2; i <= n; ++i) {
        if (graph[1][i]) {
            head = 1;
            tail = i;
            cover(head);
            cover(tail);
            next[head] = tail;
            break;
        }
    }
    while (true) {
        int x;
        while (x = adjacent(head)) {
            next[x] = head;
            head = x;
            cover(head);
        }
        while (x = adjacent(tail)) {
            next[tail] = x;
            tail = x;
            cover(tail);
        }
    }
}
```

```

    if (!graph[head][tail]) {
        for (int i = head, j; i != tail; i = next[i]) {
            if (graph[head][next[i]] && graph[tail][i]) {
                for (j = head; j != i; j = next[j]) {
                    last[next[j]] = j;
                }
                j = next[head];
                next[head] = next[i];
                next[tail] = i;
                tail = j;
                for (j = i; j != head; j = last[j]) {
                    next[j] = last[j];
                }
                break;
            }
        }
    }
    next[tail] = head;
    if (right[0] > n) {
        break;
    }
    for (int i = head; i != tail; i = next[i]) {
        if (adjacent(i)) {
            head = next[i];
            tail = i;
            next[tail] = 0;
            break;
        }
    }
}
std::vector<int> answer;
for (int i = head; ; i = next[i]) {
    if (i == 1) {
        answer.push_back(i);
        for (int j = next[i]; j != i; j = next[j]) {
            answer.push_back(j);
        }
        answer.push_back(i);
        break;
    }
    if (i == tail) {
        break;
    }
}
return answer;
}

```

## 5 字符串

### 5.1 模式匹配

#### 5.1.1 KMP 算法

```

void build(char *pattern) {
    int length = (int)strlen(pattern + 1);
    fail[0] = -1;
    for (int i = 1, j; i <= length; ++i) {

```

```

        for (j = fail[i - 1]; j != -1 && pattern[i] != pattern[j + 1]; j = fail[j]);
        fail[i] = j + 1;
    }
}

void solve(char *text, char *pattern) {
    int length = (int)strlen(text + 1);
    for (int i = 1, j; i <= length; ++i) {
        for (j = match[i - 1]; j != -1 && text[i] != pattern[j + 1]; j = fail[j]);
        match[i] = j + 1;
    }
}

```

## 5.1.2 扩展 KMP 算法

## 5.1.3 AC 自动机

## 5.2 后缀三姐妹

### 5.2.1 后缀数组

### 5.2.2 后缀自动机

```

class Node {
public:
    Node *child[256], *parent;
    int length;

    Node(int length = 0) : parent(NULL), length(length) {
        memset(child, NULL, sizeof(child));
    }

    Node* extend(Node *start, int token) {
        Node *p = this;
        Node *np = new Node(length + 1);
        for (; p && !p->child[token]; p = p->parent) {
            p->child[token] = np;
        }
        if (!p) {
            np->parent = start;
        } else {
            Node *q = p->child[token];
            if (p->length + 1 == q->length) {
                np->parent = q;
            } else {
                Node *nq = new Node(p->length + 1);
                memcpy(nq->child, q->child, sizeof(q->child));
                nq->parent = q->parent;
                np->parent = q->parent = nq;
                for (; p && p->child[token] == q; p = p->parent) {
                    p->child[token] = nq;
                }
            }
        }
        return np;
    }
};

```

## 5.3 回文三兄弟

### 5.3.1 Manacher 算法

```
void manacher(char *text, int length) {
    palindrome[0] = 1;
    for (int i = 1, j = 0; i < length; ++i) {
        if (j + palindrome[j] <= i) {
            palindrome[i] = 0;
        } else {
            palindrome[i] = std::min(palindrome[(j << 1) - i], j + palindrome[j] - i);
        }
        while (i - palindrome[i] >= 0 && i + palindrome[i] < length
            && text[i - palindrome[i]] == text[i + palindrome[i]]) {
            palindrome[i]++;
        }
        if (i + palindrome[i] > j + palindrome[j]) {
            j = i;
        }
    }
}
```

### 5.3.2 回文树

```
class Node {
public:
    Node *child[256], *fail;
    int length;

    Node(int length) : fail(NULL), length(length) {
        memset(child, NULL, sizeof(child));
    }
};

int size;
int text[N];
Node *odd, *even;

Node* match(Node *now) {
    for (; text[size - now->length - 1] != text[size]; now = now->fail);
    return now;
}

bool extend(Node *&last, int token) {
    text[++size] = token;
    Node *now = last;
    now = match(now);
    if (now->child[token]) {
        last = now->child[token];
        return false;
    }
    last = now->child[token] = new Node(now->length + 2);
    if (now == odd) {
        last->fail = even;
    } else {
        now = match(now->fail);
        last->fail = now->child[token];
    }
    return true;
}
```

```

}

void build() {
    text[size = 0] = -1;
    even = new Node(0), odd = new Node(-1);
    even->fail = odd;
}

```

## 5.4 循环串最小表示

```

int solve(char *text, int length) {
    int i = 0, j = 1, delta = 0;
    while (i < length && j < length && delta < length) {
        char tokeni = text[(i + delta) % length];
        char tokenj = text[(j + delta) % length];
        if (tokeni == tokenj) {
            delta++;
        } else {
            if (tokeni > tokenj) {
                i += delta + 1;
            } else {
                j += delta + 1;
            }
            if (i == j) {
                j++;
            }
            delta = 0;
        }
    }
    return std::min(i, j);
}

```

## 6 计算几何

### 6.1 二维基础

#### 6.1.1 点类

#### 6.1.2 凸包

```

std::vector<Point> convex_hull(std::vector<Point> point) {
    if ((int)point.size() < 3) {
        return point;
    }
    std::sort(point.begin(), point.end());
    std::vector<Point> convex;
    {
        std::vector<Point> stack;
        for (int i = 0; i < (int)point.size(); ++i) {
            while ((int)stack.size() >= 2 &&
                sgn(det(stack[(int)stack.size() - 2], stack.back(), point[i])) <= 0) {
                stack.pop_back();
            }
            stack.push_back(point[i]);
        }
        for (int i = 0; i < (int)stack.size(); ++i) {
            convex.push_back(stack[i]);
        }
    }
}

```

```

    }
    {
        std::vector<Point> stack;
        for (int i = (int)point.size() - 1; i >= 0; --i) {
            while ((int)stack.size() >= 2 &&
                sgn(det(stack[(int)stack.size() - 2], stack.back(), point[i])) <= 0) {
                stack.pop_back();
            }
            stack.push_back(point[i]);
        }
        for (int i = 1; i < (int)stack.size() - 1; ++i) {
            convex.push_back(stack[i]);
        }
    }
    return convex;
}

```

### 6.1.3 半平面交

## 6.2 三维基础

### 6.2.1 点类

### 6.2.2 凸包

### 6.2.3 绕轴旋转

## 6.3 多边形

### 6.3.1 判断点在多边形内部

```

bool point_on_line(const Point &p, const Point &a, const Point &b) {
    return sgn(det(p, a, b)) == 0 && sgn(dot(p, a, b)) <= 0;
}

bool point_in_polygon(const Point &p, const std::vector<Point> &polygon) {
    int counter = 0;
    for (int i = 0; i < (int)polygon.size(); ++i) {
        Point a = polygon[i], b = polygon[(i + 1) % (int)polygon.size()];
        if (point_on_line(p, a, b)) {
            // Point on the boundary are excluded.
            return false;
        }
        int x = sgn(det(a, p, b));
        int y = sgn(a.y - p.y);
        int z = sgn(b.y - p.y);
        counter += (x > 0 && y <= 0 && z > 0);
        counter -= (x < 0 && z <= 0 && y > 0);
    }
    return counter;
}

```

- 6.3.2 旋转卡壳
- 6.3.3 动态凸包
- 6.3.4 点到凸包的切线
- 6.3.5 直线与凸包的交点
- 6.3.6 凸多边形的交集
- 6.3.7 凸多边形内的最大圆

## 6.4 圆

- 6.4.1 圆类
- 6.4.2 圆的交集
- 6.4.3 最小覆盖圆
- 6.4.4 最小覆盖球
- 6.4.5 判断圆存在交集
- 6.4.6 圆与多边形的交集

## 6.5 三角形

- 6.5.1 三角形的内心
- 6.5.2 三角形的外心
- 6.5.3 三角形的垂心

## 6.6 黑暗科技

- 6.6.1 平面图形的转动惯量
- 6.6.2 平面区域处理
- 6.6.3 Voronoi 图

# 7 其他

## 7.1 某年某月某日是星期几

```
int solve(int year, int month, int day) {
    int answer;
    if (month == 1 || month == 2) {
        month += 12;
        year--;
    }
    if ((year < 1752) || (year == 1752 && month < 9) || (year == 1752 && month == 9 && day < 3)) {
        answer = (day + 2 * month + 3 * (month + 1) / 5 + year + year / 4 + 5) % 7;
    } else {
        answer = (day + 2 * month + 3 * (month + 1) / 5 + year + year / 4 - year / 100 + year / 400)
    }
    return answer;
}
```



## 7.2 动态规划

## 7.3 搜索

### 7.3.1 Dancing Links X

# 8 Java

## 8.1 基础模板

```
import java.io.*;
import java.util.*;
import java.math.*;

public class Main {
    public static void main(String[] args) {
        InputStream inputStream = System.in;
        OutputStream outputStream = System.out;
        InputReader in = new InputReader(inputStream);
        PrintWriter out = new PrintWriter(outputStream);
        Task solver = new Task();
        solver.solve(0, in, out);
        out.close();
    }
}

class Task {
    public void solve(int testNumber, InputReader in, PrintWriter out) {

    }
}

class InputReader {
    public BufferedReader reader;
    public StringTokenizer tokenizer;

    public InputReader(InputStream stream) {
        reader = new BufferedReader(new InputStreamReader(stream), 32768);
        tokenizer = null;
    }

    public String next() {
        while (tokenizer == null || !tokenizer.hasMoreTokens()) {
            try {
                tokenizer = new StringTokenizer(reader.readLine());
            } catch (IOException e) {
                throw new RuntimeException(e);
            }
        }
        return tokenizer.nextToken();
    }

    public int nextInt() {
        return Integer.parseInt(next());
    }

    public long nextLong() {
        return Long.parseLong(next());
    }
}
```

}  
}

## 8.2 BigInteger

## 8.3 BigDecimal

# 9 数学

## 9.1 常用积分表

## 9.2 常用数学公式

### 9.2.1 求和公式

1.  $\sum_{k=1}^n (2k-1)^2 = \frac{n(4n^2-1)}{3}$
2.  $\sum_{k=1}^n k^3 = [\frac{n(n+1)}{2}]^2$
3.  $\sum_{k=1}^n (2k-1)^3 = n^2(2n^2-1)$
4.  $\sum_{k=1}^n k^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$
5.  $\sum_{k=1}^n k^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12}$
6.  $\sum_{k=1}^n k(k+1) = \frac{n(n+1)(n+2)}{3}$
7.  $\sum_{k=1}^n k(k+1)(k+2) = \frac{n(n+1)(n+2)(n+3)}{4}$
8.  $\sum_{k=1}^n k(k+1)(k+2)(k+3) = \frac{n(n+1)(n+2)(n+3)(n+4)}{5}$

### 9.2.2 斐波那契数列

1.  $fib_0 = 0, fib_1 = 1, fib_n = fib_{n-1} + fib_{n-2}$
2.  $fib_{n+2} \cdot fib_n - fib_{n+1}^2 = (-1)^{n+1}$
3.  $fib_{-n} = (-1)^{n-1} fib_n$
4.  $fib_{n+k} = fib_k \cdot fib_{n+1} + fib_{k-1} \cdot fib_n$
5.  $gcd(fib_m, fib_n) = fib_{gcd(m,n)}$
6.  $fib_m | fib_n^2 \Leftrightarrow n fib_n | m$

### 9.2.3 错排公式

1.  $D_n = (n-1)(D_{n-2} - D_{n-1})$
2.  $D_n = n! \cdot (1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + \frac{(-1)^n}{n!})$

## 9.3 平面几何公式

### 9.3.1 三角形

1. 半周长

$$p = \frac{a+b+c}{2}$$

2. 面积

$$S = \frac{a \cdot H_a}{2} = \frac{ab \cdot \sin C}{2} = \sqrt{p(p-a)(p-b)(p-c)}$$

3. 中线

$$M_a = \frac{\sqrt{2(b^2 + c^2) - a^2}}{2} = \frac{\sqrt{b^2 + c^2 + 2bc \cdot \cos A}}{2}$$

4. 角平分线

$$T_a = \frac{\sqrt{bc \cdot [(b+c)^2 - a^2]}}{b+c} = \frac{2bc}{b+c} \cos \frac{A}{2}$$

5. 高线

$$H_a = b \sin C = c \sin B = \sqrt{b^2 - \left(\frac{a^2 + b^2 - c^2}{2a}\right)^2}$$

6. 内切圆半径

$$\begin{aligned} r &= \frac{S}{p} = \frac{\arcsin \frac{B}{2} \cdot \sin \frac{C}{2}}{\sin \frac{B+C}{2}} = 4R \cdot \sin \frac{A}{2} \sin \frac{B}{2} \sin \frac{C}{2} \\ &= \sqrt{\frac{(p-a)(p-b)(p-c)}{p}} = p \cdot \tan \frac{A}{2} \tan \frac{B}{2} \tan \frac{C}{2} \end{aligned}$$

7. 外接圆半径

$$R = \frac{abc}{4S} = \frac{a}{2\sin A} = \frac{b}{2\sin B} = \frac{c}{2\sin C}$$

### 9.3.2 四边形

$D_1, D_2$  为对角线,  $M$  为对角线中点连线,  $A$  为对角线夹角,  $p$  为半周长

$$1. a^2 + b^2 + c^2 + d^2 = D_1^2 + D_2^2 + 4M^2$$

$$2. S = \frac{1}{2} D_1 D_2 \sin A$$

3. 对于圆内接四边形

$$ac + bd = D_1 D_2$$

4. 对于圆内接四边形

$$S = \sqrt{(p-a)(p-b)(p-c)(p-d)}$$

### 9.3.3 正 $n$ 边形

$R$  为外接圆半径,  $r$  为内切圆半径

1. 中心角

$$A = \frac{2\pi}{n}$$

2. 内角

$$C = \frac{n-2}{n} \pi$$

3. 边长

$$a = 2\sqrt{R^2 - r^2} = 2R \cdot \sin \frac{A}{2} = 2r \cdot \tan \frac{A}{2}$$

4. 面积

$$S = \frac{nar}{2} = nr^2 \cdot \tan \frac{A}{2} = \frac{nR^2}{2} \cdot \sin A = \frac{na^2}{4 \cdot \tan \frac{A}{2}}$$

### 9.3.4 圆

#### 1. 弧长

$$l = rA$$

#### 2. 弦长

$$a = 2\sqrt{2hr - h^2} = 2r \cdot \sin \frac{A}{2}$$

#### 3. 弓形高

$$h = r - \sqrt{r^2 - \frac{a^2}{4}} = r(1 - \cos \frac{A}{2}) = \frac{1}{2} \cdot \arctan \frac{A}{4}$$

#### 4. 扇形面积

$$S_1 = \frac{rl}{2} = \frac{r^2 A}{2}$$

#### 5. 弓形面积

$$S_2 = \frac{rl - a(r - h)}{2} = \frac{r^2}{2}(A - \sin A)$$

### 9.3.5 棱柱

#### 1. 体积

$$V = Ah$$

$A$  为底面积,  $h$  为高

#### 2. 侧面积

$$S = lp$$

$l$  为棱长,  $p$  为直截面周长

#### 3. 全面积

$$T = S + 2A$$

### 9.3.6 棱锥

#### 1. 体积

$$V = Ah$$

$A$  为底面积,  $h$  为高

#### 2. 正棱锥侧面积

$$S = lp$$

$l$  为棱长,  $p$  为直截面周长

#### 3. 正棱锥全面积

$$T = S + 2A$$

### 9.3.7 棱台

#### 1. 体积

$$V = (A_1 + A_2 + \sqrt{A_1 A_2}) \cdot \frac{h}{3}$$

$A_1, A_2$  为上下底面积,  $h$  为高

#### 2. 正棱台侧面积

$$S = \frac{p_1 + p_2}{2} l$$

$p_1, p_2$  为上下底面周长,  $l$  为斜高

#### 3. 正棱台全面积

$$T = S + A_1 + A_2$$

### 9.3.8 圆柱

#### 1. 侧面积

$$S = 2\pi rh$$

#### 2. 全面积

$$T = 2\pi r(h + r)$$

#### 3. 体积

$$V = \pi r^2 h$$

### 9.3.9 圆锥

#### 1. 母线

$$l = \sqrt{h^2 + r^2}$$

#### 2. 侧面积

$$S = \pi rl$$

#### 3. 全面积

$$T = \pi r(l + r)$$

#### 4. 体积

$$V = \frac{\pi}{3} r^2 h$$

### 9.3.10 圆台

#### 1. 母线

$$l = \sqrt{h^2 + (r_1 - r_2)^2}$$

#### 2. 侧面积

$$S = \pi(r_1 + r_2)l$$

#### 3. 全面积

$$T = \pi r_1(l + r_1) + \pi r_2(l + r_2)$$

#### 4. 体积

$$V = \frac{\pi}{3} (r_1^2 + r_2^2 + r_1 r_2) h$$

### 9.3.11 球

#### 1. 全面积

$$T = 4\pi r^2$$

#### 2. 体积

$$V = \frac{4}{3} \pi r^3$$

### 9.3.12 球台

#### 1. 侧面积

$$S = 2\pi rh$$

#### 2. 全面积

$$T = \pi(2rh + r_1^2 + r_2^2)$$

#### 3. 体积

$$V = \frac{\pi h [3(r_1^2 + r_2^2) + h^2]}{6}$$

### 9.3.13 球扇形

#### 1. 全面积

$$T = \pi r(2h + r_0)$$

$h$  为球冠高,  $r_0$  为球冠底面半径

#### 2. 体积

$$V = \frac{2}{3}\pi r^2 h$$

## 9.4 常用数表

### 9.4.1 梅森数