代码库

上海交通大学

2015年10月9日

目录

1	数论		5
	1.1	快速求逆元	5
	1.2	扩展欧几里德算法	5
	1.3	中国剩余定理	5
	1.4	Miller Rabin 素数测试	6
	1.5	Pollard Rho 大数分解	7
	1.6	快速数论变换	8
	1.7	原根	9
	1.8	离散对数	9
	1.9	离散平方根	9
	1.10	佩尔方程求解	9
	1.11	直线下整点个数	9
2	数值		9
	2.1	· 高斯消元	9
	2.2	快速傅立叶变换	9
	2.3	· · · · · · · · · · · · · · · · · · ·	11
	2.4		13
	2.5		13
	2.6		13
	2.7	最小二乘法	13
3	* Ь+□	V/c++6	13
3	数据		
	3.1	1 2003	13
			13
			15
	3.2		18
		—— ——	18
			19
		——	20
		3.2.4 坚固的左偏树	24

	3.3	树上的魔术师	5
		3.3.1 轻重树链剖分 2	5
		3.3.2 Link Cut Tree	27
	3.4	k-d 树 2	27
	⊡ \^		. –
4	图论		27
	4.1	32,2,2,3	27
	4.2	双连通分量	
		4.2.1 点双连通分量 2	_
		4.2.2 边双连通分量 2	
		2-SAT 问题	
	4.4	二分图最大匹配 2	
		4.4.1 Hungary 算法	
		4.4.2 Hopcroft Karp 算法	
	4.5		31
		F 1 1 5 1	31
		4.5.2 扩展 KM 算法	
	4.6	最大流	
	4.7	最小费用最大流 3	
		4.7.1 稀疏图	
		4.7.2 稠密图	
	4.8	一般图最大匹配 3	8
	4.9	70, 3 — 7 3 H 3 H 3	11
	4.10	最小树形图	.2
	4.11	有根树的同构	.2
	4.12	度限制生成树4	.4
	4.13	弦图相关 4	.4
		4.13.1 弦图的判定	.4
		4.13.2 弦图的团数	.4
	4.14	哈密尔顿回路(ORE 性质的图)4	.4
E	字符	<u>串</u> 4	6
5	४ 1 ग	中 模式匹配	_
	5.1	19	-
		# 154	
	E 0	5.1.3 AC 自动机	
	5.2	后缀三姐妹	
		5.2.1 后缀数组	
	- ^	5.2.2 后缀自动机	
	5.3	回文三兄弟	
		5.3.1 Manacher 算法	
	_	5.3.2 回文树	
	54	循环串最小表示 5	12

6	计算	几何																						53
	6.1	二维基	甚础	1.							 										 			53
		6.1.1	봈	类原							 										 			53
		6.1.2	2	1包							 										 			53
		6.1.3	<u>\</u>	έΨ	面多	ই .					 										 			54
		6.1.4	盲	是近	点对	寸.					 										 			54
	6.2	三维基	甚础	┧.							 										 			55
		6.2.1	봈	类原							 										 			55
		6.2.2	2	1包							 										 			55
		6.2.3	生	Ř 轴	旋转	专 .					 										 			55
	6.3	多边形	乡.								 										 			55
		6.3.1	半	刂断	点在	E多	边	形	内	部											 			55
		6.3.2	放	罅	卡克	\bar{\bar{\bar{\bar{\bar{\bar{\bar{					 										 			56
		6.3.3	五	办态	凸É	<u>.</u>					 										 			56
		6.3.4	봈	到	凸É	9的	切切	线			 										 			56
		6.3.5	Ī	1线	5[占包	的	交	点		 										 			56
		6.3.6	7	多	边刑	沙人	的的	最	大	员											 			56
	6.4	员																						56
		6.4.1	5	数							 										 			56
		6.4.2	5	副的	交集	₹ .					 										 			56
		6.4.3	直	是 小、	覆記						 										 			56
		6.4.4	盲	是 小	覆記		Ř .				 										 			56
		6.4.5	半	刂断	圆存	字在	E交	集			 										 			56
		6.4.6		引与																				56
	6.5	三角形	多.								 	•	٠			•			 •		 			56
		6.5.1	Ξ	E角	形的	勺内]心	١			 	•	٠			•			 •		 			56
		6.5.2					_																	56
		6.5.3	Ξ	E角	形的	勺垂	€心	١			 •					•			 •		 			56
7	其他																							56
•		、 某年某	≢ ⊨	其	日長	建	胡	Л.																56
		枚举大																						
		环状最																						57
		搜索.																						
		7.4.1																						57
						, —																		
8	Java																							57
	8.1	基础模	莫板	ί.							 										 			57
9	数学																							58
J		常用数	┟쓷	ά <i>/</i> ,∖:	#																			
	ا . ق	市用奴 9.1.1		·公 文 和																				58
		9.1.1		と心とき波																				59
		9.1.2																						
		J. I. J	T.	шпнн	ΔL	٠.														 				しじ

	9.1.4	五边形数定理	59
	9.1.5	树的计数	59
	9.1.6	牛顿恒等式	60
9.2	平面几	何公式	60
	9.2.1	三角形	60
	9.2.2	四边形	61
	9.2.3	正 n 边形	61
	9.2.4	圆	62
	9.2.5	棱柱	62
	9.2.6	棱锥	62
	9.2.7	棱台	63
	9.2.8	圆柱	63
	9.2.9	圆锥	63
	9.2.10	圆台	63
	9.2.11	球	64
	9.2.12	球台	64
	9.2.13	球扇形	64
9.3	立体几	何公式	64
	9.3.1	球面三角公式	64
	9.3.2	四面体体积公式	65

1 数论

1.1 快速求逆元

```
返回结果:
                                     x^{-1}(mod)
   使用条件: x \in [0, mod) 并且 x \in mod 互质
long long inverse(const long long &x, const long long &mod) {
    if (x == 1) {
        return 1;
    } else {
        return (mod - mod / x) * inverse(mod % x, mod) % mod;
    }
}
1.2 扩展欧几里德算法
    返回结果:
                                 ax + by = gcd(a, b)
    时间复杂度: \mathcal{O}(nlogn)
void solve(const long long &a, const long long &b, long long &x, long long &y) {
    if (b == 0) {
        x = 1;
        y = 0;
    } else {
        solve(b, a \% b, x, y);
        x -= a / b * y;
        std::swap(x, y);
    }
}
1.3 中国剩余定理
    返回结果:
                              x \equiv r_i \pmod{p_i} \ (0 \le i < n)
    使用条件: p_i 无需两两互质
    时间复杂度: \mathcal{O}(nlogn)
bool solve(int n, std::pair<long long, long long> input[],
                  std::pair<long long, long long> &output) {
```

output = std::make_pair(1, 1);

```
for (int i = 0; i < n; ++i) {
        long long number, useless;
        euclid(output.second, input[i].second, number, useless);
        long long divisor = std::__gcd(output.second, input[i].second);
        if ((input[i].first - output.first) % divisor) {
            return false;
        }
        number *= (input[i].first - output.first) / divisor;
        fix(number, input[i].second);
        output.first += output.second * number;
        output.second *= input[i].second / divisor;
        fix(output.first, output.second);
    return true;
}
1.4 Miller Rabin 素数测试
const int BASE[12] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
bool check(const long long &prime, const long long &base) {
    long long number = prime - 1;
    for (; ~number & 1; number >>= 1);
    long long result = power_mod(base, number, prime);
    for (; number != prime - 1 && result != 1 && result != prime - 1; number <<= 1) {
        result = multiply_mod(result, result, prime);
    }
    return result == prime - 1 || (number & 1) == 1;
}
bool miller_rabin(const long long &number) {
    if (number < 2) {
        return false;
    }
    if (number < 4) {
        return true;
    if (~number & 1) {
        return false;
    }
    for (int i = 0; i < 12 && BASE[i] < number; ++i) {</pre>
        if (!check(number, BASE[i])) {
            return false;
```

```
}
    }
    return true;
}
1.5 Pollard Rho 大数分解
    时间复杂度: \mathcal{O}(n^{1/4})
long long pollard_rho(const long long &number, const long long &seed) {
    long long x = rand() \% (number - 1) + 1, y = x;
    for (int head = 1, tail = 2; ; ) {
        x = multiply_mod(x, x, number);
        x = add_mod(x, seed, number);
        if (x == y) {
            return number;
        }
        long long answer = std::_gcd(abs(x - y), number);
        if (answer > 1 && answer < number) {</pre>
            return answer;
        if (++head == tail) {
            y = x;
            tail <<= 1;
        }
    }
}
void factorize(const long long &number, std::vector<long long> &divisor) {
    if (number > 1) {
        if (miller_rabin(number)) {
            divisor.push_back(number);
        } else {
            long long factor = number;
            for (; factor >= number;
                   factor = pollard_rho(number, rand() % (number - 1) + 1));
            factorize(number / factor, divisor);
            factorize(factor, divisor);
        }
    }
}
```

1.6 快速数论变换

```
返回结果:
```

```
c_i = \sum_{0 \le i \le i} a_j \cdot b_{i-j}(mod) \ (0 \le i < n)
    使用说明: magic 是 mod 的原根
    时间复杂度: \mathcal{O}(nlogn)
void solve(long long number[], int length, int type) {
    for (int i = 1, j = 0; i < length - 1; ++i) {
        for (int k = length; j = k >>= 1, ~j & k; );
        if (i < j) {
             std::swap(number[i], number[j]);
        }
    }
    long long unit_p0;
    for (int turn = 0; (1 << turn) < length; ++turn) {</pre>
        int step = 1 << turn, step2 = step << 1;</pre>
        if (type == 1) {
            unit_p0 = power_mod(MAGIC, (MOD - 1) / step2, MOD);
        } else {
            unit_p0 = power_mod(MAGIC, MOD - 1 - (MOD - 1) / step2, MOD);
        for (int i = 0; i < length; i += step2) {</pre>
             long long unit = 1;
             for (int j = 0; j < step; ++j) {
                 long long &number1 = number[i + j + step];
                 long long &number2 = number[i + j];
                 long long delta = unit * number1 % MOD;
                 number1 = (number2 - delta + MOD) % MOD;
                 number2 = (number2 + delta) % MOD;
                 unit = unit * unit_p0 % MOD;
            }
        }
    }
}
void multiply() {
    for (; lowbit(length) != length; ++length);
    solve(number1, length, 1);
    solve(number2, length, 1);
    for (int i = 0; i < length; ++i) {</pre>
        number[i] = number1[i] * number2[i] % MOD;
    }
```

```
solve(number, length, -1);
for (int i = 0; i < length; ++i) {
    answer[i] = number[i] * power_mod(length, MOD - 2, MOD) % MOD;
}</pre>
```

- 1.7 原根
- 1.8 离散对数
- 1.9 离散平方根
- 1.10 佩尔方程求解
- 1.11 直线下整点个数

返回结果:

$$\sum_{0 \le i \le n} \lfloor \frac{a + b \cdot i}{m} \rfloor$$

使用条件: n, m > 0, $a, b \ge 0$ 时间复杂度: $\mathcal{O}(nlogn)$

- 2 数值
- 2.1 高斯消元
- 2.2 快速傅立叶变换

返回结果:

$$c_i = \sum_{0 \le j \le i} a_j \cdot b_{i-j} \ (0 \le i < n)$$

时间复杂度: $\mathcal{O}(nlogn)$

```
void solve(Complex number[], int length, int type) {
    for (int i = 1, j = 0; i < length - 1; ++i) {
        for (int k = length; j = k >>= 1, ~j & k; );
        if (i < j) {
            std::swap(number[i], number[j]);
        }
    }
    Complex unit_p0;
    for (int turn = 0; (1 << turn) < length; ++turn) {</pre>
        int step = 1 << turn, step2 = step << 1;</pre>
        double p0 = PI / step * type;
        sincos(p0, &unit_p0.imag(), &unit_p0.real());
        for (int i = 0; i < length; i += step2) {
            Complex unit = 1;
            for (int j = 0; j < step; ++j) {
                Complex &number1 = number[i + j + step];
                Complex &number2 = number[i + j];
                Complex delta = unit * number1;
                number1 = number2 - delta;
                number2 = number2 + delta;
                unit = unit * unit_p0;
            }
        }
    }
}
void multiply() {
    for (; lowbit(length) != length; ++length);
    solve(number1, length, 1);
    solve(number2, length, 1);
    for (int i = 0; i < length; ++i) {</pre>
        number[i] = number1[i] * number2[i];
    }
    solve(number, length, -1);
    for (int i = 0; i < length; ++i) {</pre>
        answer[i] = (int)(number[i].real() / length + 0.5);
    }
}
```

2.3 单纯形法求解线性规划

返回结果:

```
max\{c_{1\times m}\cdot x_{m\times 1}\mid x_{m\times 1}\geq 0_{m\times 1}, a_{n\times m}\cdot x_{m\times 1}\leq b_{n\times 1}\}
std::vector<double> solve(const std::vector<std::vector<double> > &a,
                            const std::vector<double> &b, const std::vector<double> &c) {
    int n = (int)a.size(), m = (int)a[0].size() + 1;
    std::vector<std::vector<double> > value(n + 2, std::vector<double>(m + 1));
    std::vector<int> index(n + m);
    int r = n, s = m - 1;
    for (int i = 0; i < n + m; ++i) {
        index[i] = i;
    for (int i = 0; i < n; ++i) {
        for (int j = 0; j < m - 1; ++j) {
             value[i][j] = -a[i][j];
        }
        value[i][m - 1] = 1;
        value[i][m] = b[i];
        if (value[r][m] > value[i][m]) {
             r = i;
        }
    }
    for (int j = 0; j < m - 1; ++j) {
        value[n][j] = c[j];
    value[n + 1][m - 1] = -1;
    for (double number; ; ) {
        if (r < n) {
             std::swap(index[s], index[r + m]);
             value[r][s] = 1 / value[r][s];
             for (int j = 0; j <= m; ++j) {
                 if (j != s) {
                     value[r][j] *= -value[r][s];
                 }
             }
             for (int i = 0; i <= n + 1; ++i) {
                 if (i != r) {
                     for (int j = 0; j \le m; ++j) {
                          if (j != s) {
                              value[i][j] += value[r][j] * value[i][s];
                          }
                     }
```

```
}
            }
        }
        r = s = -1;
        for (int j = 0; j < m; ++j) {
             if (s < 0 || index[s] > index[j]) {
                 if (value[n + 1][j] > eps | | value[n + 1][j] > -eps && value[n][j] > eps) {
                     s = j;
                 }
            }
        }
        if (s < 0) {
            break;
        for (int i = 0; i < n; ++i) {
             if (value[i][s] < -eps) {</pre>
                 if (r < 0)
                 \label{eq:continuous} |\ |\ (number = value[r][m] / value[r][s] - value[i][m] / value[i][s]) < -eps
                 | |  number < eps && index[r + m] > index[i + m]) {
                      r = i;
                 }
            }
        if (r < 0) {
            // Solution is unbounded.
            return std::vector<double>();
        }
    }
    if (value[n + 1][m] < -eps) {
        // No solution.
        return std::vector<double>();
    }
    std::vector<double> answer(m - 1);
    for (int i = m; i < n + m; ++i) {</pre>
        if (index[i] < m - 1) {</pre>
             answer[index[i]] = value[i - m][m];
    }
    return answer;
}
```

value[i][s] *= value[r][s];

2.4 自适应辛普森

```
double area(const double &left, const double &right) {
    double mid = (left + right) / 2;
    return (right - left) * (calc(left) + 4 * calc(mid) + calc(right)) / 6;
}
double simpson(const double &left, const double &right,
              const double &eps, const double &area_sum) {
    double mid = (left + right) / 2;
    double area_left = area(left, mid);
    double area_right = area(mid, right);
    double area_total = area_left + area_right;
    if (std::abs(area_total - area_sum) < 15 * eps) {</pre>
       return area_total + (area_total - area_sum) / 15;
    return simpson(left, mid, eps / 2, area_left)
        + simpson(mid, right, eps / 2, area_right);
}
double simpson(const double &left, const double &right, const double &eps) {
    return simpson(left, right, eps, area(left, right));
}
2.5 牛顿迭代法
2.6 多项式方程求解
2.7 最小二乘法
    数据结构
3
3.1 平衡的二叉查找树
3.1.1 Treap
class Node {
public:
    Node *child[2];
    int key;
    int size, priority;
    Node(Node *left, Node *right, int key) : key(key), size(1), priority(rand()) {
        child[0] = left;
```

```
child[1] = right;
    }
    void update() {
        size = child[0]->size + 1 + child[1]->size;
    }
};
Node *null;
void rotate(Node *&x, int dir) {
    Node *y = x->child[dir];
    x->child[dir] = y->child[dir ^ 1];
    y->child[dir ^ 1] = x;
    x->update();
    y->update();
    x = y;
}
void insert(Node *&x, int key) {
    if (x == null) {
        x = new Node(null, null, key);
    } else {
        insert(x->child[key > x->key], key);
        if (x->child[key > x->key]->priority < x->priority) {
            rotate(x, key > x->key);
        }
        x->update();
    }
}
void remove(Node *&x, int key) {
    if (x->key != key) {
        remove(x->child[key > x->key], key);
    } else if (x->child[0] == null && x->child[1] == null) {
        x = null;
        return;
    } else {
        int dir = x->child[0]->priority > x->child[1]->priority;
        rotate(x, dir);
        remove(x->child[dir ^ 1], key);
    }
```

```
x->update();
}
void build() {
    null = new Node(NULL, NULL, 0);
    null->child[0] = null->child[1] = null;
    null->size = 0;
    null->priority = RAND_MAX;
}
3.1.2 Splay
class Node {
public:
    Node *child[2], *father;
    int size;
    int key;
    Node(int key = 0);
    int side() {
        return father->child[1] == this;
    }
    void set(Node *son, int dir) {
        child[dir] = son;
        son->father = this;
    }
    void modify();
    void update() {
        size = child[0]->size + 1 + child[1]->size;
    }
    void release();
};
Node *null, *root;
Node::Node(int key) : size(1), key(key) {
    child[0] = child[1] = father = null;
}
```

```
void Node::modify() {
    if (this == null) {
        return;
    }
}
void rotate(Node *x) {
    int dir = x->side();
    Node *p = x->father;
    p->release();
    x->release();
    p->set(x->child[dir ^ 1], dir);
    p->father->set(x, p->side());
    x->set(p, dir ^ 1);
    if (p == root) {
        root = x;
    }
    p->update();
    x->update();
}
void splay(Node *x, Node *target = null) {
    for (x->release(); x->father != target; ) {
        if (x->father->father == target) {
            rotate(x);
        } else {
            x->side() == x->father->side()
            ? (rotate(x->father), rotate(x))
            : (rotate(x), rotate(x));
        }
    }
    x->update();
}
Node* kth(int size) {
    Node *x = root;
    for (; x->child[0]->size + 1 != size; ) {
        x->release();
        if (x->child[0]->size + 1 > size) {
            x = x->child[0];
        } else {
```

```
size -= x->child[0]->size + 1;
            x = x->child[1];
        }
    }
    return x;
}
void select(int left, int right) {
    splay(kth(right + 2));
    splay(kth(left), root);
}
void insert(int pos, int n, int key[]) {
    select(pos, pos - 1);
    Node *x = root->child[0];
    for (int i = 0; i < n; ++i) {
        Node *now = new Node(key[i]);
        x->set(now, 1);
        x = now;
    }
    splay(x);
}
void solve(int left, int right) {
    select(left, right);
    root->child[0]->child[1]->solve();
    root->child[0]->update();
    root->update();
}
void build() {
    null = new Node();
    null->size = 0;
    root = new Node();
    Node *blank = new Node();
    root->set(blank, 1);
    splay(blank);
}
```

3.2 坚固的数据结构

3.2.1 坚固的线段树

```
class Node {
public:
    Node *left, *right;
    int value;
    Node(Node *left, Node *right, int value) : left(left), right(right), value(value) {}
    Node* modify(int 1, int r, int ql, int qr, int delta);
    int query(int 1, int r, int qx);
};
Node* null;
Node* Node::modify(int 1, int r, int q1, int qr, int delta) {
    if (qr < 1 | | r < q1) {
        return this;
    }
    if (ql <= l && r <= qr) {
        return new Node(this->left, this->right, this->value + delta);
    }
    int mid = 1 + r >> 1;
    return new Node(this->left->modify(1, mid, ql, qr, delta),
                    this->right->modify(mid + 1, r, ql, qr, delta),
                    this->value);
}
int Node::query(int 1, int r, int qx) {
    if (qx < 1 | | r < qx) {
        return 0;
    }
    if (qx \le 1 \&\& r \le qx) {
        return this->value;
    }
    int mid = 1 + r >> 1;
    return this->left->query(1, mid, qx)
         + this->right->query(mid + 1, r, qx)
        + this->value;
}
```

```
void build() {
    null = new Node(NULL, NULL, 0);
    null->left = null->right = null;
}
3.2.2 坚固的平衡树
class Node {
public:
    Node *left, *right;
    int size;
    Node();
    std::pair<Node*, Node*> split(int size);
    Node* update() {
        size = left->size + 1 + right->size;
        return this;
};
bool random(int a, int b) {
    return rand() \% (a + b) < a;
}
Node *null;
Node::Node() : left(null), right(null), size(1) {}
Node* merge(Node *x, Node *y) {
    if (x == null) {
        return y;
    if (y == null) {
        return x;
    }
    if (random(x->size, y->size)) {
        x->right = merge(x->right, y);
        return x->update();
    } else {
        y->left = merge(x, y->left);
        return y->update();
    }
```

```
}
std::pair<Node*, Node*> Node::split(int size) {
    if (this == null) {
        return std::make_pair(null, null);
    }
    if (size <= left->size) {
        std::pair<Node*, Node*> result = left->split(size);
        left = null;
        return std::make_pair(result.first, merge(result.second, this->update()));
    } else {
        std::pair<Node*, Node*> result = right->split(size - left->size - 1);
        right = null;
        return std::make_pair(merge(this->update(), result.first), result.second);
}
void build() {
    null = new Node();
    null->size = 0;
}
3.2.3 坚固的字符串
  1. ext 库中的 rope
    #include <ext/rope>
    using __gnu_cxx::crope;
    using __gnu_cxx::rope;
    crope a, b;
    int main(void) {
        a = b.substr(pos, len);
                                   // [pos, pos + len)
        a = b.substr(pos);
                                    // [pos, pos]
        b.c_str();
                                    // might lead to memory leaks
        b.delete_c_str();
                                    // delete the c str that created before
        a.insert(pos, text);
                                    // insert text before position pos
                                    // erase [pos, pos + len)
        a.erase(pos, len);
    }
```

2. 可持久化平衡树实现的 rope

```
class Rope {
private:
    class Node {
    public:
        Node *left, *right;
        int size;
        char key;
        Node(char key = 0, Node *left = NULL, Node *right = NULL)
               : key(key), left(left), right(right) {
            update();
        }
        void update() {
            size = (left ? left->size : 0) + 1 + (right ? right->size : 0);
        }
        std::string to_string() {
            return (left ? left->to_string() : "") + key
                 + (right ? right->to_string() : "");
        }
    };
    bool random(int a, int b) {
        return rand() \% (a + b) < a;
    }
    Node* merge(Node *x, Node *y) {
        if (!x) {
            return y;
        }
        if (!y) {
            return x;
        if (random(x->size, y->size)) {
            return new Node(x->key, x->left, merge(x->right, y));
        } else {
            return new Node(y->key, merge(x, y->left), y->right);
    }
    std::pair<Node*, Node*> split(Node *x, int size) {
```

```
if (!x) {
            return std::make_pair<Node*, Node*>(NULL, NULL);
        if (size == 0) {
            return std::make_pair<Node*, Node*>(NULL, x);
        if (size > x->size) {
            return std::make_pair<Node*, Node*>(x, NULL);
        }
        if (x->left \&\& size <= x->left->size) {
            std::pair<Node*, Node*> part =
                split(x->left, size);
            return std::make_pair(part.first, new Node(x->key, part.second, x->right));
        } else {
            std::pair<Node*, Node*> part =
                split(x->right, size - (x->left ? x->left->size : 0) - 1);
            return std::make_pair(new Node(x->key, x->left, part.first), part.second);
        }
    }
    Node* build(const std::string &text, int left, int right) {
        if (left > right) {
            return NULL;
        int mid = left + right >> 1;
        return new Node(text[mid],
                        build(text, left, mid - 1),
                        build(text, mid + 1, right));
    }
public:
    Node *root;
    Rope() {
        root = NULL;
    }
    Rope(const std::string &text) {
        root = build(text, 0, (int)text.length() - 1);
    }
    Rope(const Rope &other) {
```

```
root = other.root;
}
Rope& operator = (const Rope &other) {
    if (this == &other) {
        return *this;
    root = other.root;
    return *this;
}
int size() {
    return root ? root->size : 0;
}
void insert(int pos, const std::string &text) {
    if (pos < 0 || pos > size()) {
        throw "Out of range";
    std::pair<Node*, Node*> part = split(root, pos);
    root = merge(merge(part.first, build(text, 0, (int)text.length() - 1)),
                 part.second);
}
void erase(int left, int right) {
    if (left < 0 || left >= size() ||
        right < 1 || right > size()) {
        throw "Out of range";
    }
    if (left >= right) {
        return;
    std::pair<Node*, Node*> part = split(root, left);
    root = merge(part.first, split(part.second, right - left).second);
}
std::string substr(int left, int right) {
    if (left < 0 || left >= size() ||
        right < 1 || right > size()) {
        throw "Out of range";
    if (left >= right) {
```

```
return "";
            }
            return split(split(root, left).second, right - left).first->to_string();
        }
        void copy(int left, int right, int pos) {
             if (left < 0 || left >= size() ||
                right < 1 || right > size() ||
                pos < 0 || pos > size()) {
                throw "Out of range";
            if (left >= right) {
                return;
            std::pair<Node*, Node*> part = split(root, pos);
            root = merge(merge(part.first,
                                split(split(root, left).second, right - left).first),
                         part.second);
        }
    };
3.2.4 坚固的左偏树
class Node {
public:
    Node *left, *right;
    int key, dist;
    Node(int key) : left(NULL), right(NULL), key(key), dist(0) {}
    Node* update() {
        if (!left || (right && left->dist < right->dist)) {
            std::swap(left, right);
        }
        dist = right ? right->dist + 1 : 0;
        return this;
    }
};
Node* merge(Node *x, Node *y) {
    if (!x) {
        return y;
    }
```

```
if (!y) {
        return x;
    }
    if (x->key < y->key) {
        x = new Node(*x);
        x->right = merge(x->right, y);
        return x->update();
    } else {
        y = new Node(*y);
        y->right = merge(x, y->right);
        return y->update();
    }
}
3.3 树上的魔术师
3.3.1 轻重树链剖分
int father[N], height[N], size[N], son[N], top[N], pos[N], data[N];
void build(int root) {
    std::vector<int> queue;
    father[root] = -1;
    height[root] = 0;
    queue.push_back(root);
    for (int head = 0; head < (int)queue.size(); ++head) {</pre>
        int x = queue[head];
        for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
            int y = edge[x][i];
            if (y != father[x]) {
                father[y] = x;
                height[y] = height[x] + 1;
                queue.push_back(y);
            }
        }
    for (int index = n - 1; index >= 0; --index) {
        int x = queue[index];
        size[x] = 1;
        son[x] = -1;
        for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
            int y = edge[x][i];
            if (y != father[x]) {
```

```
size[x] += size[y];
                 if (son[x] == -1 \mid \mid size[son[x]] < size[y]) {
                     son[x] = y;
                }
            }
        }
    }
    std::fill(top, top + n, 0);
    int counter = 0;
    for (int index = 0; index < n; ++index) {</pre>
        int x = queue[index];
        if (top[x] == 0) {
            for (int y = x; y != -1; y = son[y]) {
                top[y] = x;
                pos[y] = ++counter;
                data[counter] = value[y];
            }
        }
    }
    build(1, 1, n);
}
void solve(int x, int y) {
    while (true) {
        if (top[x] == top[y]) {
            if (x == y) {
                 solve(1, 1, n, pos[x], pos[x]);
            } else {
                 if (height[x] < height[y]) {</pre>
                     solve(1, 1, n, pos[x], pos[y]);
                } else {
                     solve(1, 1, n, pos[y], pos[x]);
                }
            }
            break;
        }
        if (height[top[x]] > height[top[y]]) {
            solve(1, 1, n, pos[top[x]], pos[x]);
            x = father[top[x]];
        } else {
            solve(1, 1, n, pos[top[y]], pos[y]);
            y = father[top[y]];
```

```
}
    }
}
3.3.2 Link Cut Tree
3.4 k-d 树
    图论
4.1
    强连通分量
int stamp, comps, top;
int dfn[N], low[N], comp[N], stack[N];
void tarjan(int x) {
    dfn[x] = low[x] = ++stamp;
    stack[top++] = x;
    for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
        int y = edge[x][i];
        if (!dfn[y]) {
            tarjan(y);
            low[x] = std::min(low[x], low[y]);
        } else if (!comp[y]) {
            low[x] = std::min(low[x], dfn[y]);
        }
    }
    if (low[x] == dfn[x]) {
        comps++;
        do {
            int y = stack[--top];
            comp[y] = comps;
        } while (stack[top] != x);
    }
}
void solve() {
    stamp = comps = top = 0;
    std::fill(dfn, dfn + n, 0);
    std::fill(comp, comp + n, 0);
    for (int i = 0; i < n; ++i) {
        if (!dfn[i]) {
            tarjan(i);
```

```
}
    }
}
4.2
    双连通分量
4.2.1 点双连通分量
4.2.2 边双连通分量
4.3 2-SAT 问题
int stamp, comps, top;
int dfn[N], low[N], comp[N], stack[N];
void add(int x, int a, int y, int b) {
    edge[x \ll 1 \mid a].push_back(y \ll 1 \mid b);
}
void tarjan(int x) {
    dfn[x] = low[x] = ++stamp;
    stack[top++] = x;
    for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
        int y = edge[x][i];
        if (!dfn[y]) {
            tarjan(y);
            low[x] = std::min(low[x], low[y]);
        } else if (!comp[y]) {
            low[x] = std::min(low[x], dfn[y]);
        }
    }
    if (low[x] == dfn[x]) {
        comps++;
        do {
            int y = stack[--top];
            comp[y] = comps;
        } while (stack[top] != x);
    }
}
bool solve() {
    int counter = n + n + 1;
    stamp = top = comps = 0;
    std::fill(dfn, dfn + counter, 0);
```

```
std::fill(comp, comp + counter, 0);
    for (int i = 0; i < counter; ++i) {</pre>
        if (!dfn[i]) {
            tarjan(i);
        }
    }
    for (int i = 0; i < n; ++i) {</pre>
        if (comp[i << 1] == comp[i << 1 | 1]) {</pre>
            return false;
        answer[i] = (comp[i << 1 | 1] < comp[i << 1]);</pre>
    }
    return true;
}
4.4 二分图最大匹配
4.4.1 Hungary 算法
    时间复杂度: \mathcal{O}(V \cdot E)
int n, m, stamp;
int match[N], visit[N];
bool dfs(int x) {
    for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
        int y = edge[x][i];
        if (visit[y] != stamp) {
            visit[y] = stamp;
             if (match[y] == -1 || dfs(match[y])) {
                 match[y] = x;
                 return true;
            }
        }
    }
    return false;
}
int solve() {
    std::fill(match, match + m, -1);
    int answer = 0;
    for (int i = 0; i < n; ++i) {</pre>
        stamp++;
        answer += dfs(i);
```

```
}
    return answer;
}
4.4.2 Hopcroft Karp 算法
    时间复杂度: \mathcal{O}(\sqrt{V} \cdot E)
int matchx[N], matchy[N], level[N];
bool dfs(int x) {
    for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
        int y = edge[x][i];
        int w = matchy[y];
        if (w == -1 \mid \mid level[x] + 1 == level[w] && dfs(w)) {
            matchx[x] = y;
            matchy[y] = x;
            return true;
        }
    }
    level[x] = -1;
    return false;
}
int solve() {
    std::fill(matchx, matchx + n, -1);
    std::fill(matchy, matchy + m, -1);
    for (int answer = 0; ; ) {
        std::vector<int> queue;
        for (int i = 0; i < n; ++i) {
             if (matchx[i] == -1) {
                 level[i] = 0;
                 queue.push_back(i);
            } else {
                 level[i] = -1;
            }
        }
        for (int head = 0; head < (int)queue.size(); ++head) {</pre>
             int x = queue[head];
             for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
                 int y = edge[x][i];
                 int w = matchy[y];
                 if (w != -1 \&\& level[w] < 0) {
```

```
level[w] = level[x] + 1;
                    queue.push_back(w);
                }
            }
        }
        int delta = 0;
        for (int i = 0; i < n; ++i) {
            if (matchx[i] == -1 \&\& dfs(i)) {
                delta++;
            }
        if (delta == 0) {
            return answer;
        } else {
            answer += delta;
        }
    }
}
4.5 二分图最大权匹配
4.5.1 KM 算法
    时间复杂度: \mathcal{O}(V^3)
int labelx[N], labely[N], match[N], slack[N];
bool visitx[N], visity[N];
bool dfs(int x) {
    visitx[x] = true;
    for (int y = 0; y < n; ++y) {
        if (visity[y]) {
            continue;
        }
        int delta = labelx[x] + labely[y] - graph[x][y];
        if (delta == 0) {
            visity[y] = true;
            if (match[y] == -1 \mid \mid dfs(match[y])) {
                match[y] = x;
                return true;
            }
        } else {
            slack[y] = std::min(slack[y], delta);
        }
```

```
}
    return false;
}
int solve() {
    for (int i = 0; i < n; ++i) {</pre>
        match[i] = -1;
        labelx[i] = INT_MIN;
        labely[i] = 0;
        for (int j = 0; j < n; ++j) {
            labelx[i] = std::max(labelx[i], graph[i][j]);
        }
    for (int i = 0; i < n; ++i) {
        while (true) {
            std::fill(visitx, visitx + n, 0);
            std::fill(visity, visity + n, 0);
            for (int j = 0; j < n; ++j) {
                slack[j] = INT_MAX;
            }
            if (dfs(i)) {
                break;
            }
            int delta = INT_MAX;
            for (int j = 0; j < n; ++j) {
                if (!visity[j]) {
                    delta = std::min(delta, slack[j]);
                }
            }
            for (int j = 0; j < n; ++j) {
                if (visitx[j]) {
                    labelx[j] -= delta;
                }
                if (visity[j]) {
                    labely[j] += delta;
                } else {
                    slack[j] -= delta;
                }
            }
        }
    }
    int answer = 0;
```

```
for (int i = 0; i < n; ++i) {
        answer += graph[match[i]][i];
    }
    return answer;
}
4.5.2 扩展 KM 算法
4.6 最大流
    时间复杂度: \mathcal{O}(V^2 \cdot E)
struct EdgeList {
    int size;
    int last[N];
    int succ[M], other[M], flow[M];
    void clear(int n) {
        size = 0;
        fill(last, last + n, -1);
    void add(int x, int y, int c) {
        succ[size] = last[x];
        last[x] = size;
        other[size] = y;
        flow[size++] = c;
    }
} e;
int n, source, target;
int dist[N], curr[N];
void add(int x, int y, int c) {
    e.add(x, y, c);
    e.add(y, x, 0);
}
bool relabel() {
    std::vector<int> queue;
    for (int i = 0; i < n; ++i) {
        curr[i] = e.last[i];
        dist[i] = -1;
    queue.push_back(target);
```

dist[target] = 0;

```
for (int head = 0; head < (int)queue.size(); ++head) {</pre>
        int x = queue[head];
        for (int i = e.last[x]; ~i; i = e.succ[i]) {
            int y = e.other[i];
            if (e.flow[i \hat{} 1] && dist[y] == -1) {
                dist[y] = dist[x] + 1;
                queue.push_back(y);
            }
        }
    }
    return ~dist[source];
}
int dfs(int x, int answer) {
    if (x == target) {
        return answer;
    }
    int delta = answer;
    for (int &i = curr[x]; ~i; i = e.succ[i]) {
        int y = e.other[i];
        if (e.flow[i] \&\& dist[x] == dist[y] + 1) {
            int number = dfs(y, std::min(e.flow[i], delta));
            e.flow[i] -= number;
            e.flow[i ^ 1] += number;
            delta -= number;
        if (delta == 0) {
            break;
        }
    }
    return answer - delta;
}
int solve() {
    int answer = 0;
    while (relabel()) {
        answer += dfs(source, INT_MAX));
    return answer;
}
```

4.7 最小费用最大流

4.7.1 稀疏图

```
时间复杂度: \mathcal{O}(V \cdot E^2)
struct EdgeList {
    int size;
    int last[N];
    int succ[M], other[M], flow[M], cost[M];
    void clear(int n) {
        size = 0;
        std::fill(last, last + n, -1);
    }
    void add(int x, int y, int c, int w) {
        succ[size] = last[x];
        last[x] = size;
        other[size] = y;
        flow[size] = c;
        cost[size++] = w;
    }
} e;
int n, source, target;
int prev[N];
void add(int x, int y, int c, int w) {
    e.add(x, y, c, w);
    e.add(y, x, 0, -w);
}
bool augment() {
    static int dist[N], occur[N];
    std::vector<int> queue;
    std::fill(dist, dist + n, INT_MAX);
    std::fill(occur, occur + n, 0);
    dist[source] = 0;
    occur[source] = true;
    queue.push_back(source);
    for (int head = 0; head < (int)queue.size(); ++head) {</pre>
        int x = queue[head];
        for (int i = e.last[x]; ~i; i = e.succ[i]) {
            int y = e.other[i];
            if (e.flow[i] \&\& dist[y] > dist[x] + e.cost[i]) {
```

```
dist[y] = dist[x] + e.cost[i];
                prev[y] = i;
                if (!occur[y]) {
                    occur[y] = true;
                    queue.push_back(y);
                }
            }
        }
        occur[x] = false;
    }
    return dist[target] < INT_MAX;</pre>
}
std::pair<int, int> solve() {
    std::pair<int, int> answer = std::make_pair(0, 0);
    while (augment()) {
        int number = INT_MAX;
        for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
            number = std::min(number, e.flow[prev[i]]);
        }
        answer.first += number;
        for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
            e.flow[prev[i]] -= number;
            e.flow[prev[i] ^ 1] += number;
            answer.second += number * e.cost[prev[i]];
        }
    }
    return answer;
}
4.7.2 稠密图
    使用条件:费用非负
    时间复杂度: \mathcal{O}(V \cdot E^2)
struct EdgeList {
    int size;
    int last[N];
    int succ[M], other[M], flow[M], cost[M];
    void clear(int n) {
        size = 0;
        std::fill(last, last + n, -1);
    }
```

```
void add(int x, int y, int c, int w) {
        succ[size] = last[x];
        last[x] = size;
        other[size] = y;
        flow[size] = c;
        cost[size++] = w;
    }
} e;
int n, source, target, flow, cost;
int slack[N], dist[N];
bool visit[N];
void add(int x, int y, int c, int w) {
    e.add(x, y, c, w);
    e.add(y, x, 0, -w);
}
bool relabel() {
    int delta = INT_MAX;
    for (int i = 0; i < n; ++i) {
        if (!visit[i]) {
            delta = std::min(delta, slack[i]);
        }
        slack[i] = INT_MAX;
    if (delta == INT_MAX) {
        return true;
    }
    for (int i = 0; i < n; ++i) {
        if (visit[i]) {
            dist[i] += delta;
        }
    }
    return false;
}
int dfs(int x, int answer) {
    if (x == target) {
        flow += answer;
        cost += answer * (dist[source] - dist[target]);
        return answer;
```

```
}
    visit[x] = true;
    int delta = answer;
    for (int i = e.last[x]; ~i; i = e.succ[i]) {
        int y = e.other[i];
        if (e.flow[i] > 0 && !visit[y]) {
            if (dist[y] + e.cost[i] == dist[x]) {
                int number = dfs(y, std::min(e.flow[i], delta));
                e.flow[i] -= number;
                e.flow[i ^ 1] += number;
                delta -= number;
                if (delta == 0) {
                    dist[x] = INT_MIN;
                    return answer;
                }
            } else {
                slack[y] = std::min(slack[y], dist[y] + e.cost[i] - dist[x]);
            }
        }
    }
    return answer - delta;
}
std::pair<int, int> solve() {
    flow = cost = 0;
    std::fill(dist, dist + n, 0);
    do {
        do {
            fill(visit, visit + n, 0);
        } while (dfs(source, INT_MAX));
    } while (!relabel());
    return std::make_pair(flow, cost);
}
4.8 一般图最大匹配
    时间复杂度: \mathcal{O}(V^3)
int match[N], belong[N], next[N], mark[N], visit[N];
std::vector<int> queue;
int find(int x) {
    if (belong[x] != x) {
```

```
belong[x] = find(belong[x]);
    }
    return belong[x];
}
void merge(int x, int y) {
    x = find(x);
    y = find(y);
    if (x != y) {
        belong[x] = y;
    }
}
int lca(int x, int y) {
    static int stamp = 0;
    stamp++;
    while (true) {
        if (x != -1) {
            x = find(x);
            if (visit[x] == stamp) {
                return x;
            }
            visit[x] = stamp;
            if (match[x] != -1) {
                x = next[match[x]];
            } else {
                x = -1;
            }
        }
        std::swap(x, y);
    }
}
void group(int a, int p) {
    while (a != p) {
        int b = match[a], c = next[b];
        if (find(c) != p) {
            next[c] = b;
        }
        if (mark[b] == 2) {
            mark[b] = 1;
            queue.push_back(b);
```

```
}
        if (mark[c] == 2) {
            mark[c] = 1;
            queue.push_back(c);
        }
        merge(a, b);
        merge(b, c);
        a = c;
    }
}
void augment(int source) {
    queue.clear();
    for (int i = 0; i < n; ++i) {
        next[i] = visit[i] = -1;
        belong[i] = i;
        mark[i] = 0;
    }
    mark[source] = 1;
    queue.push_back(source);
    for (int head = 0; head < (int)queue.size() && match[source] == -1; ++head) {
        int x = queue[head];
        for (int i = 0; i < (int)edge[x].size(); ++i) {</pre>
            int y = edge[x][i];
            if (match[x] == y \mid \mid find(x) == find(y) \mid \mid mark[y] == 2) {
                 continue;
            }
            if (mark[y] == 1) {
                int r = lca(x, y);
                if (find(x) != r) {
                     next[x] = y;
                }
                 if (find(y) != r) {
                    next[y] = x;
                }
                group(x, r);
                group(y, r);
            } else if (match[y] == -1) {
                next[y] = x;
                for (int u = y; u != -1; ) {
                     int v = next[u];
                     int mv = match[v];
```

```
match[v] = u;
                   match[u] = v;
                   u = mv;
               }
               break;
           } else {
               next[y] = x;
               mark[y] = 2;
               mark[match[y]] = 1;
               queue.push_back(match[y]);
           }
       }
    }
}
int solve() {
    std::fill(match, match + n, -1);
    for (int i = 0; i < n; ++i) {
       if (match[i] == -1) {
           augment(i);
       }
    }
    int answer = 0;
    for (int i = 0; i < n; ++i) {
       answer += (match[i] != -1);
    return answer;
}
4.9 无向图全局最小割
    时间复杂度: \mathcal{O}(V^3)
   注意事项:处理重边时,应该对边权累加
int node[N], dist[N];
bool visit[N];
int solve(int n) {
    int answer = INT_MAX;
    for (int i = 0; i < n; ++i) {
       node[i] = i;
    }
    while (n > 1) {
```

```
int max = 1;
        for (int i = 0; i < n; ++i) {</pre>
            dist[node[i]] = graph[node[0]][node[i]];
            if (dist[node[i]] > dist[node[max]]) {
                max = i;
            }
        }
        int prev = 0;
        memset(visit, 0, sizeof(visit));
        visit[node[0]] = true;
        for (int i = 1; i < n; ++i) {
            if (i == n - 1) {
                answer = std::min(answer, dist[node[max]]);
                for (int k = 0; k < n; ++k) {
                    graph[node[k]][node[prev]] =
                         (graph[node[prev]][node[k]] += graph[node[k]][node[max]]);
                }
                node[max] = node[--n];
            visit[node[max]] = true;
            prev = max;
            max = -1;
            for (int j = 1; j < n; ++j) {
                if (!visit[node[j]]) {
                    dist[node[j]] += graph[node[prev]][node[j]];
                    if (max == -1 || dist[node[max]] < dist[node[j]]) {</pre>
                        max = j;
                    }
                }
            }
        }
    }
    return answer;
}
4.10 最小树形图
4.11 有根树的同构
    时间复杂度: \mathcal{O}(V log V)
```

const unsigned long long MAGIC = 4423;

```
unsigned long long magic[N];
std::pair<unsigned long long, int> hash[N];
void solve(int root) {
    magic[0] = 1;
    for (int i = 1; i <= n; ++i) {
        magic[i] = magic[i - 1] * MAGIC;
    std::vector<int> queue;
    queue.push_back(root);
    for (int head = 0; head < (int)queue.size(); ++head) {</pre>
        int x = queue[head];
        for (int i = 0; i < (int)son[x].size(); ++i) {</pre>
            int y = son[x][i];
            queue.push_back(y);
        }
    }
    for (int index = n - 1; index >= 0; --index) {
        int x = queue[index];
        hash[x] = std::make_pair(0, 0);
        std::vector<std::pair<unsigned long long, int> > value;
        for (int i = 0; i < (int)son[x].size(); ++i) {</pre>
            int y = son[x][i];
            value.push_back(hash[y]);
        std::sort(value.begin(), value.end());
        hash[x].first = hash[x].first * magic[1] + 37;
        hash[x].second++;
        for (int i = 0; i < (int)value.size(); ++i) {</pre>
            hash[x].first = hash[x].first * magic[value[i].second] + value[i].first;
            hash[x].second += value[i].second;
        }
        hash[x].first = hash[x].first * magic[1] + 41;
        hash[x].second++;
    }
}
```

```
4.12 度限制生成树
```

4.13 弦图相关

4.13.1 弦图的判定

4.13.2 弦图的团数

4.14 哈密尔顿回路(ORE 性质的图)

ORE 性质: $\forall x, y \in V \land (x, y) \notin E \quad s.t. \quad deg_x + deg_y \ge n$ 返回结果: 从顶点 1 出发的一个哈密尔顿回路 使用条件: $n \ge 3$ int left[N], right[N], next[N], last[N]; void cover(int x) { left[right[x]] = left[x]; right[left[x]] = right[x]; } int adjacent(int x) { for (int i = right[0]; i <= n; i = right[i]) {</pre> if (graph[x][i]) { return i; } } return 0; } std::vector<int> solve() { for (int i = 1; i <= n; ++i) { left[i] = i - 1;right[i] = i + 1; int head, tail; for (int i = 2; i <= n; ++i) { if (graph[1][i]) { head = 1;tail = i; cover(head); cover(tail);

next[head] = tail;

```
break;
    }
}
while (true) {
    int x;
    while (x = adjacent(head)) {
        next[x] = head;
        head = x;
        cover(head);
    }
    while (x = adjacent(tail)) {
        next[tail] = x;
        tail = x;
        cover(tail);
    }
    if (!graph[head][tail]) {
        for (int i = head, j; i != tail; i = next[i]) {
            if (graph[head][next[i]] && graph[tail][i]) {
                for (j = head; j != i; j = next[j]) {
                    last[next[j]] = j;
                }
                j = next[head];
                next[head] = next[i];
                next[tail] = i;
                tail = j;
                for (j = i; j != head; j = last[j]) {
                    next[j] = last[j];
                }
                break;
            }
        }
    }
    next[tail] = head;
    if (right[0] > n) {
        break;
    }
    for (int i = head; i != tail; i = next[i]) {
        if (adjacent(i)) {
            head = next[i];
            tail = i;
            next[tail] = 0;
            break;
```

```
}
        }
    }
    std::vector<int> answer;
    for (int i = head; ; i = next[i]) {
        if (i == 1) {
            answer.push_back(i);
            for (int j = next[i]; j != i; j = next[j]) {
                answer.push_back(j);
            answer.push_back(i);
            break;
        }
        if (i == tail) {
            break;
        }
    }
    return answer;
}
```

5 字符串

5.1 模式匹配

5.1.1 KMP 算法

```
void build(char *pattern) {
    int length = (int)strlen(pattern + 1);
    fail[0] = -1;
    for (int i = 1, j; i <= length; ++i) {
        for (j = fail[i - 1]; j != -1 && pattern[i] != pattern[j + 1]; j = fail[j]);
        fail[i] = j + 1;
    }
}

void solve(char *text, char *pattern) {
    int length = (int)strlen(text + 1);
    for (int i = 1, j; i <= length; ++i) {
        for (j = match[i - 1]; j != -1 && text[i] != pattern[j + 1]; j = fail[j]);
        match[i] = j + 1;
    }
}</pre>
```

5.1.2 扩展 KMP 算法

5.1.3 AC 自动机

```
class Node {
public:
    Node *child[256], *fail;
    int counter;
    Node() : fail(NULL), counter(0) {
        memset(child, NULL, sizeof(child));
    }
};
void insert(Node *x, char *text) {
    int length = (int)strlen(text);
    for (int i = 0; i < length; ++i) {</pre>
        int token = (int)text[i];
        if (!x->child[token]) {
            x->child[token] = new Node();
        }
        x = x->child[token];
    }
    x->counter++;
}
void build() {
    std::vector<Node*> queue;
    queue.push_back(root->fail = root);
    for (int head = 0; head < (int)queue.size(); ++head) {</pre>
        Node *x = queue[head];
        for (int token = 0; token < 256; ++token) {</pre>
            if (x->child[token]) {
                x->child[token]->fail = (x == root) ? root : x->fail->child[token];
                x->child[token]->counter += x->child[token]->fail->counter;
                queue.push_back(x->child[token]);
            } else {
                x->child[token] = (x == root) ? root : x->fail->child[token];
            }
        }
    }
}
```

5.2 后缀三姐妹

5.2.1 后缀数组

```
int array[N], rank[N], height[N];
int counter[N], new_array[N], new_rank[N][2];
int log2[N], value[N][20];
void build(char *text, int n) {
    memset(counter, 0, sizeof(counter));
    for (int i = 0; i < n; ++i) {
        counter[(int)text[i]]++;
    for (int i = 0; i < 256; ++i) {
        counter[i + 1] += counter[i];
    }
    for (int i = 0; i < n; ++i) {
        rank[i] = counter[(int)text[i]] - 1;
    for (int length = 1; length < n; length <<= 1) {</pre>
        for (int i = 0; i < n; ++i) {
            new_rank[i][0] = rank[i];
            new_rank[i][1] = i + length < n ? rank[i + length] + 1 : 0;</pre>
        memset(counter, 0, sizeof(counter));
        for (int i = 0; i < n; ++i) {
            counter[new rank[i][1]]++;
        for (int i = 0; i < n; ++i) {
            counter[i + 1] += counter[i];
        for (int i = n - 1; i >= 0; --i) {
            new_array[--counter[new_rank[i][1]]] = i;
        }
        memset(counter, 0, sizeof(counter));
        for (int i = 0; i < n; ++i) {
            counter[new_rank[i][0]]++;
        for (int i = 0; i < n; ++i) {
            counter[i + 1] += counter[i];
        for (int i = n - 1; i >= 0; --i) {
            array[--counter[new_rank[new_array[i]][0]]] = new_array[i];
```

```
}
        rank[array[0]] = 0;
        for (int i = 0; i + 1 < n; ++i) {
            rank[array[i + 1]] = rank[array[i]] +
                 (new_rank[array[i]][0] != new_rank[array[i + 1]][0]
              || new_rank[array[i]][1] != new_rank[array[i + 1]][1]);
        }
    }
    for (int i = 0, length = 0; i < n; ++i) {
        if (rank[i]) {
            int j = array[rank[i] - 1];
            while (i + length < n && j + length < n
                    && text[i + length] == text[j + length]) {
                length++;
            height[rank[i]] = length;
            if (length) {
                length--;
            }
        }
    }
    for (int i = 2; i <= n; ++i) {
        log2[i] = log2[i >> 1] + 1;
    for (int i = 1; i < n; ++i) {
        value[i][0] = height[i];
    }
    for (int step = 1; (1 << step) <= n; ++step) {</pre>
        for (int i = 1; i + (1 << step) <= n; ++i) {
            value[i][step] = std::min(value[i][step - 1],
                                       value[i + (1 << step - 1)][step - 1]);</pre>
        }
    }
}
int lcp(int left, int right) {
    if (left > right) {
        std::swap(left, right);
    }
    int step = log2[right - left];
    return std::min(value[left + 1][step], value[right - (1 << step) + 1][step]);</pre>
}
```

5.2.2 后缀自动机

```
class Node {
public:
    Node *child[256], *parent;
    int length;
    Node(int length = 0) : parent(NULL), length(length) {
        memset(child, NULL, sizeof(child));
    }
    Node* extend(Node *start, int token) {
        Node *p = this;
        Node *np = new Node(length + 1);
        for (; p \&\& !p->child[token]; p = p->parent) {
            p->child[token] = np;
        }
        if (!p) {
            np->parent = start;
        } else {
            Node *q = p->child[token];
            if (p->length + 1 == q->length) {
                np->parent = q;
            } else {
                Node *nq = new Node(p->length + 1);
                memcpy(nq->child, q->child, sizeof(q->child));
                nq->parent = q->parent;
                np->parent = q->parent = nq;
                for (; p \&\& p \rightarrow child[token] == q; p = p \rightarrow parent) {
                    p->child[token] = nq;
                }
            }
        }
        return np;
    }
};
5.3
    回文三兄弟
5.3.1 Manacher 算法
void manacher(char *text, int length) {
    palindrome[0] = 1;
```

```
for (int i = 1, j = 0; i < length; ++i) {
        if (j + palindrome[j] <= i) {</pre>
            palindrome[i] = 0;
        } else {
            palindrome[i] = std::min(palindrome[(j << 1) - i], j + palindrome[j] - i);</pre>
        while (i - palindrome[i] >= 0 && i + palindrome[i] < length</pre>
                && text[i - palindrome[i]] == text[i + palindrome[i]]) {
            palindrome[i]++;
        }
        if (i + palindrome[i] > j + palindrome[j]) {
        }
    }
}
5.3.2 回文树
class Node {
public:
    Node *child[256], *fail;
    int length;
    Node(int length) : fail(NULL), length(length) {
        memset(child, NULL, sizeof(child));
    }
};
int size;
int text[N];
Node *odd, *even;
Node* match(Node *now) {
    for (; text[size - now->length - 1] != text[size]; now = now->fail);
    return now;
}
bool extend(Node *&last, int token) {
    text[++size] = token;
    Node *now = last;
    now = match(now);
    if (now->child[token]) {
        last = now->child[token];
```

```
return false;
    }
    last = now->child[token] = new Node(now->length + 2);
    if (now == odd) {
        last->fail = even;
    } else {
        now = match(now->fail);
        last->fail = now->child[token];
    }
    return true;
}
void build() {
    text[size = 0] = -1;
    even = new Node(0), odd = new Node(-1);
    even->fail = odd;
}
5.4 循环串最小表示
int solve(char *text, int length) {
    int i = 0, j = 1, delta = 0;
    while (i < length && j < length && delta < length) {
        char tokeni = text[(i + delta) % length];
        char tokenj = text[(j + delta) % length];
        if (tokeni == tokenj) {
            delta++;
        } else {
            if (tokeni > tokenj) {
                i += delta + 1;
            } else {
                j += delta + 1;
            }
            if (i == j) {
                j++;
            delta = 0;
        }
    }
    return std::min(i, j);
}
```

6 计算几何

6.1 二维基础

6.1.1 点类

6.1.2 凸包

```
std::vector<Point> convex_hull(std::vector<Point> point) {
    std::sort(point.begin(), point.end());
    std::vector<Point> convex;
        std::vector<Point> stack;
        for (int i = 0; i < (int)point.size(); ++i) {</pre>
            while ((int)stack.size() >= 2 &&
                     sgn(det(stack[(int)stack.size() - 2], stack.back(), point[i])) <= 0) {</pre>
                 stack.pop_back();
            }
            stack.push_back(point[i]);
        }
        for (int i = 0; i < (int)stack.size(); ++i) {</pre>
            convex.push_back(stack[i]);
        }
    }
    {
        std::vector<Point> stack;
        for (int i = (int)point.size() - 1; i >= 0; --i) {
            while ((int)stack.size() >= 2 &&
                     sgn(det(stack[(int)stack.size() - 2], stack.back(), point[i])) <= 0) {</pre>
                 stack.pop_back();
            }
            stack.push_back(point[i]);
        }
        for (int i = 1; i < (int)stack.size() - 1; ++i) {</pre>
            convex.push_back(stack[i]);
        }
    }
    return convex;
}
```

6.1.3 半平面交

6.1.4 最近点对

```
bool comparex(const Point &a, const Point &b) {
    return sgn(a.x - b.x) < 0;
}
bool comparey(const Point &a, const Point &b) {
    return sgn(a.y - b.y) < 0;
}
double solve(const std::vector<Point> &point, int left, int right) {
    if (left == right) {
        return INF;
    }
    if (left + 1 == right) {
        return dist(point[left], point[right]);
    int mid = left + right >> 1;
    double result = std::min(solve(left, mid), solve(mid + 1, right));
    std::vector<Point> candidate;
    for (int i = left; i <= right; ++i) {</pre>
        if (std::abs(point[i].x - point[mid].x) <= result) {</pre>
            candidate.push_back(point[i]);
        }
    }
    std::sort(candidate.begin(), candidate.end(), comparey);
    for (int i = 0; i < (int)candidate.size(); ++i) {</pre>
        for (int j = i + 1; j < (int)candidate.size(); ++j) {</pre>
            if (std::abs(candidate[i].y - candidate[j].y) >= result) {
                break;
            }
            result = std::min(result, dist(candidate[i], candidate[j]));
        }
    }
    return result;
}
double solve(std::vector<Point> point) {
    std::sort(point.begin(), point.end(), comparex);
    return solve(point, 0, (int)point.size() - 1);
}
```

```
6.2 三维基础
```

- 6.2.1 点类
- 6.2.2 凸包
- 6.2.3 绕轴旋转
- 6.3 多边形
- 6.3.1 判断点在多边形内部

```
bool point_on_line(const Point &p, const Point &a, const Point &b) {
    return sgn(det(p, a, b)) == 0 && sgn(dot(p, a, b)) <= 0;
}
bool point_in_polygon(const Point &p, const std::vector<Point> &polygon) {
    int counter = 0;
    for (int i = 0; i < (int)polygon.size(); ++i) {</pre>
        Point a = polygon[i], b = polygon[(i + 1) % (int)polygon.size()];
        if (point_on_line(p, a, b)) {
                 Point on the boundary are excluded.
            return false;
        }
        int x = sgn(det(a, p, b));
        int y = sgn(a.y - p.y);
        int z = sgn(b.y - p.y);
        counter += (x > 0 \&\& y \le 0 \&\& z > 0);
        counter -= (x < 0 \&\& z <= 0 \&\& y > 0);
    }
    return counter;
}
```

- 6.3.2 旋转卡壳
- 6.3.3 动态凸包
- 6.3.4 点到凸包的切线
- 6.3.5 直线与凸包的交点
- 6.3.6 凸多边形内的最大圆
- 6.4 圆
- 6.4.1 圆类
- 6.4.2 圆的交集
- 6.4.3 最小覆盖圆
- 6.4.4 最小覆盖球
- 6.4.5 判断圆存在交集
- 6.4.6 圆与多边形的交集
- 6.5 三角形
- 6.5.1 三角形的内心
- 6.5.2 三角形的外心
- 6.5.3 三角形的垂心

7 其他

7.1 某年某月某日是星期几

```
return answer;
}
7.2 枚举大小为 k 的子集
   使用条件: k > 0
void solve(int n, int k) {
   for (int comb = (1 << k) - 1; comb < (1 << n); ) {
       int x = comb \& -comb, y = comb + x;
       comb = (((comb & ~y) / x) >> 1) | y;
   }
}
7.3 环状最长公共子串
    搜索
7.4
7.4.1 Dancing Links X
8
    Java
8.1
   基础模板
import java.io.*;
import java.util.*;
import java.math.*;
public class Main {
   public static void main(String[] args) {
       InputStream inputStream = System.in;
       OutputStream outputStream = System.out;
       InputReader in = new InputReader(inputStream);
       PrintWriter out = new PrintWriter(outputStream);
       Task solver = new Task();
       solver.solve(0, in, out);
       out.close();
   }
}
```

public void solve(int testNumber, InputReader in, PrintWriter out) {

class Task {

```
}
}
class InputReader {
    public BufferedReader reader;
    public StringTokenizer tokenizer;
    public InputReader(InputStream stream) {
        reader = new BufferedReader(new InputStreamReader(stream), 32768);
        tokenizer = null;
    }
    public String next() {
        while (tokenizer == null || !tokenizer.hasMoreTokens()) {
            try {
                tokenizer = new StringTokenizer(reader.readLine());
            } catch (IOException e) {
                throw new RuntimeException(e);
            }
        }
        return tokenizer.nextToken();
    }
    public int nextInt() {
        return Integer.parseInt(next());
    }
    public long nextLong() {
        return Long.parseLong(next());
}
```

9 数学

9.1 常用数学公式

9.1.1 求和公式

1.
$$\sum_{k=1}^{n} (2k-1)^2 = \frac{n(4n^2-1)}{3}$$

2.
$$\sum_{k=1}^{n} k^3 = \left[\frac{n(n+1)}{2}\right]^2$$

3.
$$\sum_{k=1}^{n} (2k-1)^3 = n^2(2n^2-1)$$

4.
$$\sum_{k=1}^{n} k^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$$

5.
$$\sum_{k=1}^{n} k^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12}$$

6.
$$\sum_{k=1}^{n} k(k+1) = \frac{n(n+1)(n+2)}{3}$$

7.
$$\sum_{k=1}^{n} k(k+1)(k+2) = \frac{n(n+1)(n+2)(n+3)}{4}$$

8.
$$\sum_{k=1}^{n} k(k+1)(k+2)(k+3) = \frac{n(n+1)(n+2)(n+3)(n+4)}{5}$$

9.1.2 斐波那契数列

1.
$$fib_0 = 0, fib_1 = 1, fib_n = fib_{n-1} + fib_{n-2}$$

2.
$$fib_{n+2} \cdot fib_n - fib_{n+1}^2 = (-1)^{n+1}$$

3.
$$fib_{-n} = (-1)^{n-1} fib_n$$

4.
$$fib_{n+k} = fib_k \cdot fib_{n+1} + fib_{k-1} \cdot fib_n$$

5.
$$gcd(fib_m, fib_n) = fib_{qcd(m,n)}$$

6.
$$fib_m|fib_n^2 \Leftrightarrow nfib_n|m$$

9.1.3 错排公式

1.
$$D_n = (n-1)(D_{n-2} - D_{n-1})$$

2.
$$D_n = n! \cdot \left(1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + \frac{(-1)^n}{n!}\right)$$

9.1.4 五边形数定理

设 p(n) 是 n 的拆分数,有

$$p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k-1} p\left(n - \frac{k(3k-1)}{2}\right)$$

9.1.5 树的计数

1. 有根树计数: n+1 个结点的有根树的个数为

$$a_{n+1} = \frac{\sum_{j=1}^{n} j \cdot a_j \cdot S_{n,j}}{n}$$

其中,

$$S_{n,j} = \sum_{i=1}^{n/j} a_{n+1-ij} = S_{n-j,j} + a_{n+1-j}$$

2. 无根树计数: 当 n 为奇数时, n 个结点的无根树的个数为

$$a_n - \sum_{i=1}^{n/2} a_i a_{n-i}$$

当 n 为偶数时,n 个结点的无根树的个数为

$$a_n - \sum_{i=1}^{n/2} a_i a_{n-i} + \frac{1}{2} a_{\frac{n}{2}} (a_{\frac{n}{2}} + 1)$$

3. n 个结点的完全图的生成树个数为

$$n^{n-2}$$

4. 矩阵-树定理:图 G 由 n 个结点构成,设 A[G] 为图 G 的邻接矩阵、D[G] 为图 G 的度数矩阵,则图 G 的不同生成树的个数为 C[G] = D[G] - A[G] 的任意一个 n-1 阶主子式的行列式值。

9.1.6 牛顿恒等式

设

$$\prod_{i=1}^{n} (x - x_i) = a_n + a_{n-1}x + \dots + a_1x^{n-1} + a_0x^n$$

$$p_k = \sum_{i=1}^n x_i^k$$

则

$$a_0p_k + a_1p_{k-1} + \dots + a_{k-1}p_1 + ka_k = 0$$

特别地,对于

$$|\mathbf{A} - \lambda \mathbf{E}| = (-1)^n (a_n + a_{n-1}\lambda + \dots + a_1\lambda^{n-1} + a_0\lambda^n)$$

有

$$p_k = Tr(\boldsymbol{A}^k)$$

9.2 平面几何公式

9.2.1 三角形

1. 半周长

$$p=\frac{a+b+c}{2}$$

2. 面积

$$S = \frac{a \cdot H_a}{2} = \frac{ab \cdot sinC}{2} = \sqrt{p(p-a)(p-b)(p-c)}$$

3. 中线

$$M_a = \frac{\sqrt{2(b^2 + c^2) - a^2}}{2} = \frac{\sqrt{b^2 + c^2 + 2bc \cdot cosA}}{2}$$

4. 角平分线

$$T_a = \frac{\sqrt{bc \cdot [(b+c)^2 - a^2]}}{b+c} = \frac{2bc}{b+c} \cos \frac{A}{2}$$

5. 高线

$$H_a = bsinC = csinB = \sqrt{b^2 - (\frac{a^2 + b^2 - c^2}{2a})^2}$$

6. 内切圆半径

$$\begin{split} r &= \frac{S}{p} = \frac{\arcsin\frac{B}{2} \cdot \sin\frac{C}{2}}{\sin\frac{B+C}{2}} = 4R \cdot \sin\frac{A}{2} \sin\frac{B}{2} \sin\frac{C}{2} \\ &= \sqrt{\frac{(p-a)(p-b)(p-c)}{p}} = p \cdot \tan\frac{A}{2} \tan\frac{B}{2} \tan\frac{C}{2} \end{split}$$

7. 外接圆半径

$$R = \frac{abc}{4S} = \frac{a}{2sinA} = \frac{b}{2sinB} = \frac{c}{2sinC}$$

9.2.2 四边形

 D_1, D_2 为对角线,M 对角线中点连线,A 为对角线夹角,p 为半周长

1.
$$a^2 + b^2 + c^2 + d^2 = D_1^2 + D_2^2 + 4M^2$$

2.
$$S = \frac{1}{2}D_1D_2sinA$$

3. 对于圆内接四边形

$$ac + bd = D_1D_2$$

4. 对于圆内接四边形

$$S = \sqrt{(p-a)(p-b)(p-c)(p-d)}$$

9.2.3 正 n 边形

R 为外接圆半径,r 为内切圆半径

1. 中心角

$$A = \frac{2\pi}{n}$$

2. 内角

$$C = \frac{n-2}{n}\pi$$

3. 边长

$$a = 2\sqrt{R^2 - r^2} = 2R \cdot \sin\frac{A}{2} = 2r \cdot \tan\frac{A}{2}$$

4. 面积

$$S = \frac{nar}{2} = nr^2 \cdot tan\frac{A}{2} = \frac{nR^2}{2} \cdot sinA = \frac{na^2}{4 \cdot tan\frac{A}{2}}$$

9.2.4 圆

1. 弧长

$$l = rA$$

2. 弦长

$$a = 2\sqrt{2hr - h^2} = 2r \cdot \sin\frac{A}{2}$$

3. 弓形高

$$h = r - \sqrt{r^2 - \frac{a^2}{4}} = r(1 - \cos\frac{A}{2}) = \frac{1}{2} \cdot \arctan\frac{A}{4}$$

4. 扇形面积

$$S_1 = \frac{rl}{2} = \frac{r^2 A}{2}$$

5. 弓形面积

$$S_2 = \frac{rl - a(r - h)}{2} = \frac{r^2}{2}(A - sinA)$$

9.2.5 棱柱

1. 体积

$$V = Ah$$

A 为底面积,h 为高

2. 侧面积

$$S = lp$$

l 为棱长, p 为直截面周长

3. 全面积

$$T = S + 2A$$

9.2.6 棱锥

1. 体积

$$V = Ah$$

A 为底面积,h 为高

2. 正棱锥侧面积

$$S = lp$$

l 为棱长, p 为直截面周长

3. 正棱锥全面积

$$T = S + 2A$$

9.2.7 棱台

1. 体积

$$V = (A_1 + A_2 + \sqrt{A_1 A_2}) \cdot \frac{h}{3}$$

 A_1, A_2 为上下底面积,h 为高

2. 正棱台侧面积

$$S = \frac{p_1 + p_2}{2}l$$

 p_1, p_2 为上下底面周长,l 为斜高

3. 正棱台全面积

$$T = S + A_1 + A_2$$

9.2.8 圆柱

1. 侧面积

$$S = 2\pi r h$$

2. 全面积

$$T = 2\pi r(h+r)$$

3. 体积

$$V=\pi r^2 h$$

9.2.9 圆锥

1. 母线

$$l = \sqrt{h^2 + r^2}$$

2. 侧面积

$$S = \pi r l$$

3. 全面积

$$T = \pi r(l+r)$$

4. 体积

$$V = \frac{\pi}{3}r^2h$$

9.2.10 圆台

1. 母线

$$l = \sqrt{h^2 + (r_1 - r_2)^2}$$

2. 侧面积

$$S = \pi(r_1 + r_2)l$$

3. 全面积

$$T = \pi r_1(l + r_1) + \pi r_2(l + r_2)$$

4. 体积

$$V = \frac{\pi}{3}(r_1^2 + r_2^2 + r_1 r_2)h$$

9.2.11 球

1. 全面积

$$T = 4\pi r^2$$

2. 体积

$$V = \frac{4}{3}\pi r^3$$

9.2.12 球台

1. 侧面积

$$S = 2\pi r h$$

2. 全面积

$$T = \pi(2rh + r_1^2 + r_2^2)$$

3. 体积

$$V = \frac{\pi h[3(r_1^2 + r_2^2) + h^2]}{6}$$

9.2.13 球扇形

1. 全面积

$$T = \pi r (2h + r_0)$$

h 为球冠高, r_0 为球冠底面半径

2. 体积

$$V = \frac{2}{3}\pi r^2 h$$

9.3 立体几何公式

9.3.1 球面三角公式

设 a,b,c 是边长, A,B,C 是所对的二面角, 有余弦定理

$$cosa = cosb \cdot cosc + sinb \cdot sinc \cdot cosA$$

正弦定理

$$\frac{sinA}{sina} = \frac{sinB}{sinb} = \frac{sinC}{sinc}$$

三角形面积是 $A + B + C - \pi$

9.3.2 四面体体积公式

U, V, W, u, v, w 是四面体的 6 条棱, U, V, W 构成三角形, (U, u), (V, v), (W, w) 互为对棱, 则

$$V = \frac{\sqrt{(s-2a)(s-2b)(s-2c)(s-2d)}}{192uvw}$$

其中

$$\begin{cases} a &= \sqrt{xYZ}, \\ b &= \sqrt{yZX}, \\ c &= \sqrt{zXY}, \\ d &= \sqrt{xyz}, \\ s &= a+b+c+d, \\ X &= (w-U+v)(U+v+w), \\ x &= (U-v+w)(v-w+U), \\ Y &= (U-V+w)(V+w+u), \\ Y &= (V-w+u)(W+w+u), \\ Z &= (V-W+u)(W+u+v), \\ Z &= (W-u+v)(u-v+W) \end{cases}$$