

# Planning

- Representation and reasoning about actions in the situation calculus
- Planning in STRIPS
- Planning as search

Next: Reasoning under uncertainty

Reading: Chaps 13, 14

\*Slides based on those of Hector Levesque and Sheila McIlraith

# Why planning

- So far we studied problem solving by search, and knowledge representation and reasoning
- However, intelligent agents are not simply passive problem solvers or reasoners
- Intelligent agents must act on the world.
- We want them to act in intelligent ways
  - taking purposeful actions,
  - predicting the expected effect of such actions,
  - composing actions together to achieve complex goals

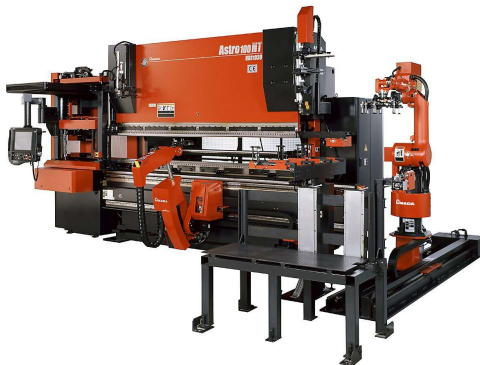
# Activity Planning for the Mars Exploration Rovers (MER)

- Operating MER is a challenging, time-pressured task.
- Each day, the operations team must generate a new plan describing the rover activities for the next day.
- These plans must abide by resource limitations, safety rules, and temporal constraints.
- Automated planning is used in generating the plans.



# Planning in Manufacturing Automation

- Sheet-metal bending machines
- Plan the sequence of bends



# Planning in Bridge Baron

- Bridge Baron is a computer program that plays bridge.
- It won the 1997 world championship of computer bridge.
- It uses AI planning techniques to plan its declarer play.

- Scheduling
  - Supply chain management
  - Hubble Space Telescope scheduler
  - Workflow management
- Air Traffic Control
  - Route aircraft between runways and terminals.
  - Crafts must be kept safely separated.
  - Safe distance depends on craft and mode of transport.
  - Minimize taxi and wait time.

# Other applications

- Character Animation
  - Generate step-by-step character behaviour from high-level spec
- Plan-based intelligent user interfaces
- Web Service Composition
- Genome Rearrangement

- To do planning, we need to reason about what the world will be like after doing a number of actions.
- Only now we want to reason about dynamic environments.
  - `in(robby,room1)`, `lightOn(room1)` are true: will they be true after robby performs the action `turnOffLights`?
  - `in(robby,room1)` is true: what does robby need to do to make `in(robby,room3)` true?
- Reasoning about the effects of actions, and computing what actions can achieve certain effects is at the heart of decision making.



# Planning under Uncertainty

One of the major complexities in planning that we will deal with later is planning under uncertainty.

- Our knowledge of the world will almost certainly be incomplete. We may wish to model this probabilistically.
- Sensing is subject to noise (especially in robots).
- Actions and effectors are also subject to error (uncertainty in their effects).

# Classical Planning

For now we restrict our attention to the deterministic case.

- Complete initial state specifications
- Deterministic effects of actions

What we will cover

- Representing and reasoning about actions in the situation calculus
- Planning in STRIPS
- Planning as search

# Situation Calculus

- First we look at how to represent and reason about dynamic worlds within first-order logic.
- The situation calculus is an important formalism developed for this purpose.
- Building blocks: actions, situations, fluents

- Action functions such as  $pickup(x)$ ,  $stack(x, y)$
- A set of primitive action objects in the domain of individuals.
- Action functions mapping objects to primitive action objects.
- $pickup(x)$  is a function mapping block  $x$  to the primitive action object corresponding to picking up block  $x$
- $stack(x, y)$  is a function mapping blocks  $x$  and  $y$  to the primitive action object corresponding to stacking  $x$  on top of  $y$

# Situations

- Situations denote possible action histories
- A set of situation objects in the domain of individuals.
- A special constant  $S_0$  denoting the initial situation, where no actions have been performed
- A special function  $do(a, s)$  denoting the situation resulting from doing action  $a$  in situation  $s$
- e.g.,  $do(put(a, b), do(put(b, c), S_0))$
- Note that situations are different from states
- e.g., when you flip a switch once and three times, the action histories are different, but the states are the same

- Fluents are predicates or functions whose values may vary from situation to situation
- These are written using predicate or function symbols whose last argument is a situation
- e.g.,  $Holding(r, x, s)$ : robot  $r$  is holding object  $x$  in situation  $s$
- Can have:  
 $\neg Holding(r, x, s) \wedge Holding(r, x, do(pickup(r, x), s))$

# Preconditions and effects

- Actions typically have preconditions: what needs to be true for the action to be performed
  - to do  $pickup(r, x)$  in situation  $s$ , we need to have  $\forall z. \neg Holding(r, z, s) \wedge \neg Heavy(x) \wedge NextTo(r, x, s)$
  - to do  $repair(r, x)$  in situation  $s$ , we need to have  $HasGlue(r, s) \wedge Broken(x, s)$
- Actions typically have effects: the fluents that change as the result of performing the action
  - $Fragile(x) \supset Broken(x, do(drop(r, x), s))$
  - $\neg Broken(x, do(repair(r, x), s))$

# An example of preconditions and effects

Action effects are conditional on their preconditions being true.

To specify action  $pickup(x)$ , we have

$$\begin{aligned} \forall s, x. & \text{ontable}(x, s) \wedge \text{clear}(x, s) \wedge \text{handempty}(s) \\ & \rightarrow \text{holding}(x, \text{do}(\text{pickup}(x), s)) \\ & \wedge \neg \text{handempty}(\text{do}(\text{pickup}(x), s)) \\ & \wedge \neg \text{ontable}(x, \text{do}(\text{pickup}(x), s)) \\ & \wedge \neg \text{clear}(x, \text{do}(\text{pickup}(x), s)) \end{aligned}$$



- There are many ways to generate plans.
- Here we show how to do it by representing actions in the SitCalc and generating a plan via deductive plan synthesis.
- This is not the approach taken by state-of-the-art planners, as we will see later,
- but it is where the field started,
- and is still used for specifying and studying more complex tasks in reasoning about actions and change.

# Plan generation

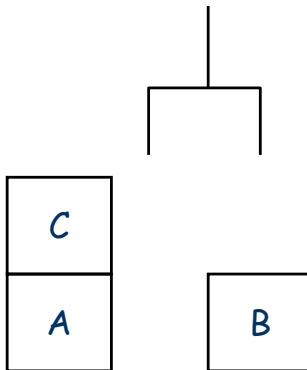
- Let's use Resolution to find a plan.
- The KB contains a description of
  - the initial state of the world
  - the precondition axioms for actions
  - the effect axioms for actions
- The query is the goal condition that we wish to achieve

# Plan generation

1.  $clear(c, S_0)$
2.  $on(c, a, S_0)$
3.  $clear(b, S_0)$
4.  $ontable(a, S_0)$
5.  $ontable(b, S_0)$
6.  $handempty(S_0)$

Query:  $\exists z.holding(b, z)$

7.  $(\neg holding(b, z), ans(z))$



Does there exist a situation in which the robot is holding b?  
If so, what is the name of that situation?

# Plan generation

Convert “pickup” action axiom into clause form:

$$\begin{aligned} \forall s, x. & \text{ontable}(x, s) \wedge \text{clear}(x, s) \wedge \text{handempty}(s) \\ & \rightarrow \text{holding}(x, \text{do}(\text{pickup}(x), s)) \\ & \wedge \neg \text{handempty}(\text{do}(\text{pickup}(x), s)) \\ & \wedge \neg \text{ontable}(x, \text{do}(\text{pickup}(x), s)) \\ & \wedge \neg \text{clear}(x, \text{do}(\text{pickup}(x), s)) \end{aligned}$$

8.  $(\neg \text{ontable}(x, s), \neg \text{clear}(x, s), \neg \text{handempty}(s),$   
 $\text{holding}(x, \text{do}(\text{pickup}(x), s)))$
9.  $(\neg \text{ontable}(x, s), \neg \text{clear}(x, s), \neg \text{handempty}(s), ,$   
 $\neg \text{handempty}(\text{do}(\text{pickup}(x), s)))$
10.  $(\neg \text{ontable}(x, s), \neg \text{clear}(x, s), \neg \text{handempty}(s),$   
 $\neg \text{ontable}(x, \text{do}(\text{pickup}(x), s)))$
11.  $(\neg \text{ontable}(x, s), \neg \text{clear}(x, s), \neg \text{handempty}(s), ,$   
 $\neg \text{clear}(x, \text{do}(\text{pickup}(x), s)))$

# All clauses

1.  $clear(c, S_0)$
2.  $on(c, a, S_0)$
3.  $clear(b, S_0)$
4.  $ontable(a, S_0)$
5.  $ontable(b, S_0)$
6.  $handempty(S_0)$
7.  $(\neg holding(b, z), ans(z))$
8.  $(\neg ontable(x, s), \neg clear(x, s), \neg handempty(s), holding(x, do(pickup(x), s)))$
9.  $(\neg ontable(x, s), \neg clear(x, s), \neg handempty(s), \neg handempty(do(pickup(x), s)))$
10.  $(\neg ontable(x, s), \neg clear(x, s), \neg handempty(s), \neg ontable(x, do(pickup(x), s)))$
11.  $(\neg ontable(x, s), \neg clear(x, s), \neg handempty(s), \neg clear(x, do(pickup(x), s)))$

- 12.  $R[7a,8d] x=b, z=holding(b, do(pickup(x), s))$   
 $(\neg ontable(b, s), \neg clear(b, s), \neg handempty(s), ans(do(pickup(b), s)))$
- 13.  $R[5,12a] s=S_0$   
 $(\neg clear(b, S_0), \neg handempty(S_0), ans(do(pickup(b), S_0)))$
- 14.  $R[3,13a] (\neg handempty(S_0), ans(do(pickup(b), S_0)))$
- 15.  $R[6,14a] ans(do(pickup(b), S_0))$

Thus the plan to achieve “holding(b)” from situation  $S_0$  is simply one action: pickup(b).

# Two common types of reasoning

- Temporal projection:  
Predicting the effects of a given sequence of actions  
*e.g.*,  $on(b, c, do(stack(b, c), do(pickup(b), S_0)))$
- Plan synthesis:  
Computing a sequence of actions that achieve a goal condition  
*e.g.*,  $\exists s. on(b, c, s) \wedge on(c, a, s)$

# The frame problem

Unfortunately, although  $on(c, a, do(pickup(b), S_0))$  should hold, it can't be proved via resolution.

1.  $clear(c, S_0)$
2.  $on(c, a, S_0)$
3.  $clear(b, S_0)$
4.  $ontable(a, S_0)$
5.  $ontable(b, S_0)$
6.  $handempty(S_0)$
7.  $(\neg holding(b, z), ans(z))$
8.  $(\neg ontable(x, s), \neg clear(x, s), \neg handempty(s), holding(x, do(pickup(x), s)))$
9.  $(\neg ontable(x, s), \neg clear(x, s), \neg handempty(s), \neg handempty(do(pickup(x), s)))$
10.  $(\neg ontable(x, s), \neg clear(x, s), \neg handempty(s), \neg ontable(x, do(pickup(x), s)))$
11.  $(\neg ontable(x, s), \neg clear(x, s), \neg handempty(s), \neg clear(x, do(pickup(x), s)))$
12.  $\neg on(c, a, do(pickup(b), S_0))$

Nothing can resolve with 12!



# What's wrong?

- We stated the effects of  $\text{pickup}(b)$ , but did not state that it doesn't affect  $\text{on}(c,a)$ .
- Problem: need to know a vast number of such non-effects.
- Few actions affect the value of a given fluent; most leave it invariant.
- e.g., an object's colour is unaffected by picking things up, opening a door, using the phone, turning on a light, electing a new Prime Minister of Canada, etc.

# The frame problem

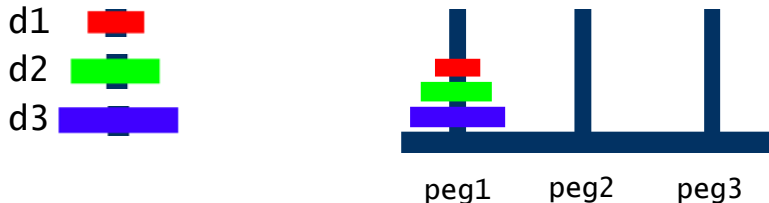
- Finding an effective way of specifying the non-effects of actions, without having to explicitly write them all down is the frame problem.
- Good solutions have been proposed, e.g., successor state axioms.

# Computation problems

- SitCalc is a very powerful representation
- Very good for studying the formal properties of planning and for modelling more complex issues in planning.
- But it is not efficient enough for automated plan generation.
- Next we will study some less expressive representations that support more efficient planning.

# Exercise

将汉诺塔的移动问题形式化为一个规划问题。假设有3个盘子



我们使用6个常量:  $peg1, peg2, peg3, d1, d2, d3$ , 3个谓词:

- ①  $clear(x)$ : 柱 $x$ 上没有盘子, 或者盘 $x$ 上没有其他盘子;
- ②  $on(x, y)$ : 盘 $x$ 在盘 $y$ 上, 或者盘 $x$ 在柱 $y$ 上并且盘 $x$ 不在其他盘子上;
- ③  $smaller(x, y)$ : 盘 $x$ 比盘 $y$ 小, 或者 $x$ 是盘并且 $y$ 是柱子。

假设只有一个动作:  $move(disc, from, to)$ : 将盘 $disc$ 从盘或柱 $from$ 上移到盘或柱 $to$ 上。

试写出 $move$ 动作的公理, 初始知识库, 以及目标。

- Plan generation: Computing a sequence of actions that achieve a goal condition
- Temporal projection: Predicting the effects of a given sequence of actions
- SitCalc: actions, situations, fluents
- The preconditions and effects of actions
- Plan generation in SitCalc via resolution
- The frame problem

# Simplifying the Planning Problem

- Perform Classical Planning. No incomplete or uncertain knowledge.
- Assume complete information about the initial state through the closed-world assumption (CWA).
- Assume action preconditions are restricted to conjunctions of ground atoms.
- Assume action effects are restricted to making ground atoms true or false. No conditional effects, no disjunctive effects, etc.

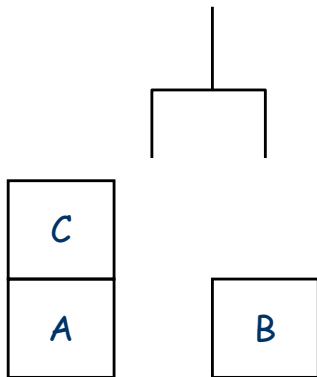
# Closed World Assumption (CWA)

- The knowledge base used to represent a state of the world is a list of positive ground atomic facts. (Like a database.)
- Closed World Assumption (CWA) is the assumption that
  - the constants mentioned in KB are all the domain objects.
  - if a ground atomic fact is not in our list of “known” facts, its negation must be true.
- This gives complete information about the state of the system.

# Querying a Closed World KB

- CWA treats a knowledge base like a database:
  - e.g., if  $employed(John, CIBC)$  is not in the database, we conclude that  $\neg employed(John, CIBC)$  is true.
- Such a KB is called a CW-KB (Closed World Knowledge Base)
- With a CW-KB, we can evaluate the truth or falsity of arbitrarily complex first-order formulas.
  - The CW-KB acts as a model
- This process is very similar to query evaluation in databases.





KB = {handempty  
clear(c), clear(b),  
on(c,a),  
ontable(a), ontable(b)}

Arbitrary queries:

1.  $\text{clear}(c) \wedge \text{clear}(b)?$
2.  $\neg \text{on}(b,c)?$
3.  $\text{on}(a,c) \vee \text{on}(b,c)?$
4.  $\exists X. \text{on}(X,c)?$  ( $D = \{a,b,c\}$ )
5.  $\forall X. \text{ontable}(X)$   
 $\rightarrow X = a \vee X = b?$

# STRIPS representation

- STRIPS (Stanford Research Institute Problem Solver) is an automated planner developed by Fikes and Nilsson in 1971.
- The same name was later used to refer to the formal language of the inputs to this planner.
- Actions are modeled as ways of modifying the world.
- Since the world is represented as a CW-KB, a STRIPS action represents a way of updating the CW-KB.
- An action yields a new KB, describing the new world – the world resulting from executing the action

# SitCalc vs. STRIPS

In the SitCalc, we could

- Have incomplete information (specified by a FOL formula)
- Refer to different situations at the same time in one formula,
  - e.g.,  $on(a, b, s_0) \wedge \neg on(a, b, s_1)$

In STRIPS, we have

- Complete information (specified by a CW-KB)
- Information about the state of the world at different times would have to be captured by two separate databases.

- STRIPS represents an action using 3 lists.
  - A list of action preconditions.
  - A list of action add effects.
  - A list of action delete effects.
- These lists contain variables, so that we can represent a whole class of actions with one specification.
- Each ground instantiation of the variables yields a specific action.

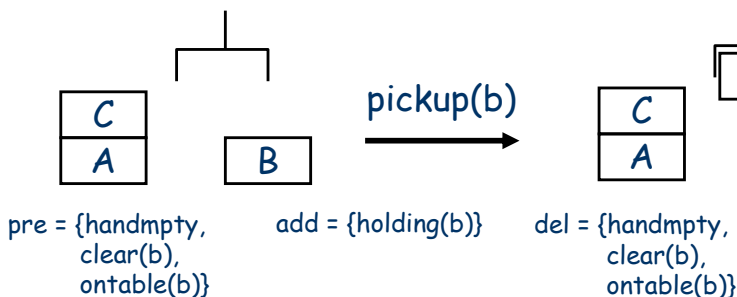
# STRIPS Actions: Example

- `pickup(X)`:
  - Pre:  $\{\text{handempty}, \text{clear}(X), \text{ontable}(X)\}$
  - Adds:  $\{\text{holding}(X)\}$
  - Dels:  $\{\text{handempty}, \text{clear}(X), \text{ontable}(X)\}$
- “`pickup(X)`” is called a STRIPS operator.
- a particular instance (e.g. “`pickup(a)`” ) is called an action.

# Operation of a STRIPS action

- For a particular STRIPS action (ground instance) to be applicable to a state (a CW-KB)
  - Every fact in its precondition list must be true in KB.
  - This amounts to testing membership since we have only atomic facts in the precondition list.
- If the action is applicable, the new state is generated by
  - Removing all facts in Dels from KB, then
  - Adding all facts in Adds to KB.

# Operation of a Strips Action: Example



KB =  $\{\text{handempty}, \text{clear}(c), \text{clear}(b), \text{on}(c,a), \text{ontable}(a), \text{ontable}(b)\}$

KB =  $\{\text{holding}(b), \text{clear}(c), \text{on}(c,a), \text{ontable}(a)\}$

# STRIPS Blocks World Operators (1)

- pickup(X) (from the table)  
Pre: {clear(X), ontable(X), handempty}  
Add: {holding(X)}  
Dels: {clear(X), ontable(X), handempty}
- putdown(X) (on the table)  
Pre: {holding(X)}  
Add: {clear(X), ontable(X), handempty}  
Del: {holding(X)}



## STRIPS Blocks World Operators (2)

- `unstack(X,Y)` (pickup from a stack of blocks)  
Pre:  $\{\text{clear}(X), \text{on}(X,Y), \text{handempty}\}$   
Add:  $\{\text{holding}(X), \text{clear}(Y)\}$   
Del:  $\{\text{clear}(X), \text{on}(X,Y), \text{handempty}\}$
- `stack(X,Y)` (putdown on a block)  
Pre:  $\{\text{holding}(X), \text{clear}(Y)\}$   
Add:  $\{\text{on}(X,Y), \text{handempty}, \text{clear}(X)\}$   
Del:  $\{\text{holding}(X), \text{clear}(Y)\}$

## 8 Puzzle as a planning problem

The constants

- A constant representing each position:  $P_1, \dots, P_9$

P1	P2	P3
P4	P5	P6
P7	P8	P9

- A constant for each tile (and blank):  $B, T_1, \dots, T_8$ .

## 8 Puzzle

The predicates

- $\text{adjacent}(X,Y)$ : position  $X$  is next to position  $Y$ 
  - e.g.,  $\text{adjacent}(P5,P2)$ ,  $\text{adjacent}(P5,P4)$ ,  $\text{adjacent}(P5,P8)$ , ...
- $\text{at}(X,Y)$ : tile  $X$  is at position  $Y$ 
  - e.g.,  $\text{at}(T1,P1)$ ,  $\text{at}(T2,P2)$ ,  $\text{at}(T5,P3)$ , ...

P1	P2	P3
P4	P5	P6
P7	P8	P9

1	2	5
7	8	
6	4	3

## 8 Puzzle

A single operator (creating lots of ground actions)

**slide(T,X,Y)**

Pre: {at(T,X), at(B,Y), adjacent(X,Y)}

Add: {at(B,X), at(T,Y)}

Del: {at(T,X), at(B,Y)}

at(T1,P1), at(T5,P3),  
at(T8,P5), at(B,P6), ...,

at(T1,P1), at(T5,P3),  
at(B,P5), at(T8,P6), ...,

1	2	5
7	8	
6	4	3



slide(T8,P5,P6)

1	2	5
7		8
6	4	3

# STRIPS has no Conditional Effects

- Unlike ADL (Action Description Language) which we'll see later, STRIPS has no conditional effects
- e.g., if we allow conditional effects, we can treat putdown(X) as stack(X,Y) where  $Y = \text{table}$ ,
- Since STRIPS has no conditional effects, we must sometimes utilize extra actions: one for each type of condition.
- We embed the condition in the precondition, and then alter the effects accordingly.

- STRIPS operators are not very expressive
- ADL (Action Description Language) extends the expressivity of STRIPS
- ADL operators add a number of features to STRIPS
  - Their preconditions can be arbitrary formulas
  - They can have conditional and universal effects

# ADL Operators Example

Blocks World Assumptions:

The table has infinite space, so it is always clear.

- If we stack something on the table ( $Y=Table$ ), we cannot delete  $clear(Table)$ ,
- But if  $Y$  is an ordinary block we must delete  $clear(Y)$ .

$move(X,Y,Z)$

Pre:  $on(X,Y) \wedge clear(Z)$

Effs:  $ADD[on(X,Z)]$

$DEL[on(X,Y)]$

$Z \neq table \rightarrow DEL[clear(Z)]$

$Y \neq table \rightarrow ADD[clear(Y)]$

# ADL Operators Example



$\text{move}(c,a,b)$

Pre:  $\text{on}(c,a) \wedge \text{clear}(b)$

Effs:  $\text{ADD}[\text{on}(c,b)]$

$\text{DEL}[\text{on}(c,a)]$

$b \neq \text{table} \rightarrow \text{DEL}[\text{clear}(b)]$

$a \neq \text{table} \rightarrow \text{ADD}[\text{clear}(a)]$

$\text{KB} = \{ \text{clear}(c), \text{clear}(b),$   
 $\text{on}(c,a),$   
 $\text{on}(a,\text{table}),$   
 $\text{on}(b,\text{table}) \}$

$\text{KB} = \{ \text{on}(c,b)$   
 $\text{clear}(c), \text{clear}(a)$   
 $\text{on}(a,\text{table}),$   
 $\text{on}(b,\text{table}) \}$



# ADL Operators with universal effects

clearTable()

Pre:

Effs:  $\forall X.on(X, table) \rightarrow DEL[on(X, table)]$

$KB = \{on(a, table), on(b, table), on(c, table)\} \Rightarrow KB = \{\}$

# Planning as a Search Problem

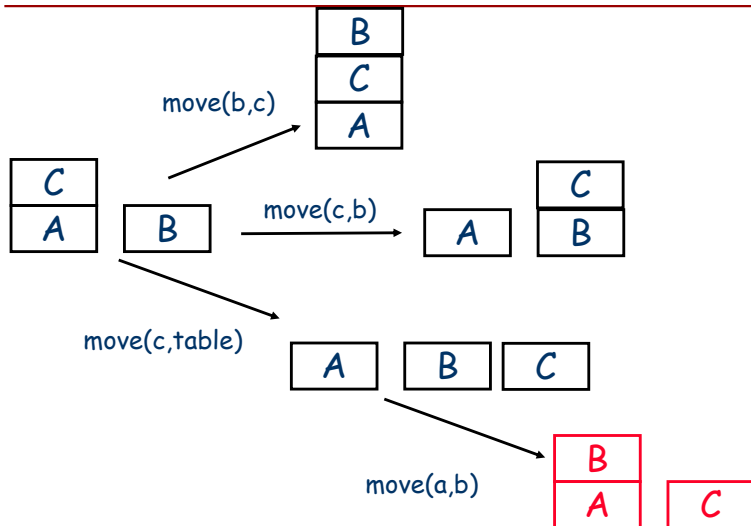
- Given
  - a CW-KB representing the initial state,
  - A set of STRIPS or ADL operators mapping a state to a new state
  - a goal condition (specified as a formula)
- The planning problem is to determine a sequence of actions that when applied to the initial CW-KB yield an updated CW-KB which satisfies the goal.
- This is known as the classical planning task.

# Planning as Search

This is a search problem in which our state representation is a CW-KB.

- The initial CW-KB is the initial state.
- The actions are operators mapping a state (a CW-KB) to a new state (an updated CW-KB).
- It can be checked if the goal is satisfied by any state (CW-KB).
- Typically the goal is a conjunction of ground facts, so we just need to check if all of them are contained in the CW-KB.

# An example



# Problem and solution

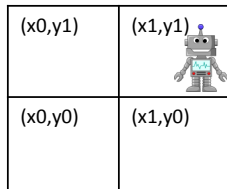
- Problem: Search tree is generally quite large
  - randomly reconfiguring 9 blocks takes thousands of CPU seconds
- But the representation suggests some structure.
  - Each action only affects a small set of facts
- Planning algorithms are designed to take advantage of the fact that the representation makes the “locality” of action changes explicit.
- We will look at one technique: Relaxed Plan heuristics used with heuristic search.
- The heuristics are domain independent. So the technique belongs to domain-independent heuristic search for planning

# Exercise

如图所示，有4个房间。

开始时，机器人在房间 $(x1, y1)$ 。

机器人的目标是访问所有的房间。



我们使用4个常量： $locx0y0, locx0y1, locx1y0, locx1y1$ ，  
3个谓词：

- 1  $at(loc)$ : 机器人在房间 $loc$ ;
- 2  $visited(loc)$ : 机器人访问过房间 $loc$ ;
- 3  $connected(loc1, loc2)$ : 房间 $loc1$ 与房间 $loc2$ 相邻。

假设只有一个动作： $move(from, to)$ : 机器人从房间 $from$ 移动到房间 $to$ ，前提是这两个房间相邻。

试写出 $move$ 动作的STRIPS表示，初始知识库，目标，及一个解

# Planning in STRIPS

- The world is represented as a CW-KB, *i.e.*, a list of ground atoms with CWA
  - the constants mentioned in KB are all the domain objects
  - if a ground atom is not in the list, it is false
- An action is represented using 3 lists: Pre, Add and Del
- ADL extends STRIPS with arbitrary preconditions, conditional and universal effects

# Classical planning

- Given
  - a CW-KB representing the initial state,
  - A set of STRIPS operators mapping a state to a new state
  - a goal condition
- Determine a sequence of actions that transforms the initial CW-KB to a CW-KB satisfying the goal



# Relaxed problem

- We make the assumption that
  - the precondition of each action is a set of positive facts, and
  - the goal is a set of positive facts.
- Recall that we can obtain heuristics by solving a relaxed version of 8-puzzle in which we relax one of the restrictions:
  - move to adjacent field only
  - move to blank field only
- The idea here is similar: consider what happens if we ignore the delete lists of actions.
- This yields a “relaxed problem” that can produce a useful heuristic estimate.

# STRIPS Blocks World Operators – Relaxed

- **pickup(X)**  
Pre: {handempty, ontable(X), clear(X)}  
Add: {holding(X)}  
~~Del: {handempty, ontable(X), clear(X)}~~
- **putdown(X)**  
Pre: {holding(X)}  
Add: {handempty, ontable(X), clear(X)}  
~~Del: {holding(X)}~~
- **unstack(X,Y)**  
Pre: {handempty, clear(X), on(X,Y)}  
Add: {holding(X), clear(Y)}  
~~Del: {handempty, clear(X), on(X,Y)}~~
- **stack(X,Y)**  
Pre: {holding(X), clear(Y)}  
Add: {handempty, clear(X), on(X,Y)}  
~~Del: {holding(X), clear(Y)}~~

**Theorem.** The length of an optimal plan for the relaxed problem is bounded by the length of an optimal plan for the original problem.

Proof:

- Let  $P$  be the original problem, and  $P'$  the relaxed problem.
- We show that  $Sols(P) \subseteq Sols(P')$ .
- Then it follows that  $Minlen(Sols(P')) \leq Minlen(Sols(P))$ .

- Let  $a_0, \dots, a_{n-1}$  be a solution to  $P$ .
- We show that it is also a solution to  $P'$ .
- Let  $s_0$  denote the initial state.
- For  $i < n$ , let  $s_{i+1} = s_i \cup \text{add}(a_i) - \text{del}(a_i)$ .
- Then  $\text{Goal} \subseteq s_n$ , and for  $i < n$ ,  $\text{pre}(a_i) \subseteq s_i$ .
- Let  $s'_0 = s_0$ , and for  $i < n$ , let  $s'_{i+1} = s'_i \cup \text{add}(a_i)$ .
- We show by induction on  $i$  that for  $i \leq n$ ,  $s_i \subseteq s'_i$ .
- Thus  $\text{Goal} \subseteq s_n \subseteq s'_n$ , and for  $i < n$ ,  $\text{pre}(a_i) \subseteq s_i \subseteq s'_i$ .
- So  $a_0, \dots, a_{n-1}$  is also a solution to  $P'$ .

We show by induction on  $i$  that for  $i \leq n$ ,  $s_i \subseteq s'_i$ .

- Basis:  $i = 0$ , obviously.
- Induction: Let  $i < n$ . Assume  $s_i \subseteq s'_i$ . We prove  $s_{i+1} \subseteq s'_{i+1}$ .
- Since  $a_i$  is applicable in  $s_i$ ,  $pre(a_i) \subseteq s_i$ .
- Hence  $pre(a_i) \subseteq s_i \subseteq s'_i$ . So  $a_i$  is applicable in  $s'_i$ .
- So  $s_{i+1} = s_i \cup add(a_i) - del(a_i) \subseteq s'_i \cup add(a_i) = s'_{i+1}$ .

# Computing the heuristic

- So the length of an optimal relaxed plan could serve as an admissible heuristic for  $A^*$ .
- However, it is NP-Hard to compute an optimal relaxed plan
- So how do we compute the heuristic?
- Build a layered structure from a state  $S$  that reaches the goal.
- Count how many actions are required in a relaxed plan.
- Use it as our heuristic estimate of the distance of  $S$  to goal.

# Reachability analysis

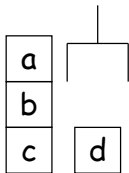
- Start with the initial state  $S_0$ .
- Alternate between state and action layers.
- Find all actions whose preconditions are contained in  $S_0$ .
- These actions comprise the first action layer  $A_0$ .
- The next state layer  $S_1 = S_0 \cup$  all facts added by actions in  $A_0$ .
- Continue this process

# Reachability analysis

- In general,
  - $A_i$ : set of actions not in  $A_{i-1}$  whose preconditions are in  $S_i$
  - $S_i = S_{i-1} \cup$  the add lists of all actions in  $A_i$
- Intuitively
  - the actions at level  $A_i$  are the actions that could be executed at the  $i$ -th step of some plan, and
  - the facts in level  $S_i$  are the facts that could be made true within a plan of length  $i$
- Some of the actions/facts have this property. But not all!
- We continue until
  - the goal  $G$  is contained in the state layer, or
  - the state layer no longer changes (reaches fix point).



# Example

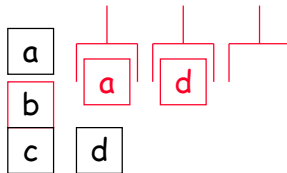


on(a,b),  
on(b,c),  
ontable(c),  
ontable(d),  
clear(a),  
clear(d),  
handempty

$S_0$

**unstack(a,b)**  
**pickup(d)**

$A_0$

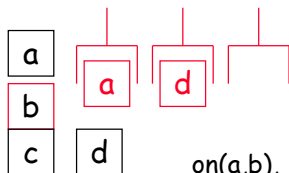


on(a,b),  
on(b,c),  
ontable(c),  
ontable(d),  
clear(a),  
handempty,  
clear(d),  
**holding(a),**  
**clear(b),**  
**holding(d)**

$S_1$

this is not  
a state as  
some of  
these  
facts  
cannot be  
true at the  
same time!

# Example



on(a,b),  
on(b,c),  
ontable(c),  
ontable(d),  
clear(a),  
clear(d),  
handempty,  
holding(a),  
clear(b),  
holding(d)

$S_1$

unstack(a,b)  
pickup(d)

} from  $A_0$

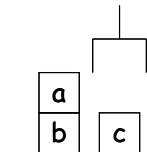
putdown(a),  
putdown(d),  
stack(a,b),  
**stack(a,a),**  
**stack(d,a),**  
**stack(d,b),**  
**stack(d,d),**  
unstack(b,c)

Impossible,  
but we don't  
know because  
we ignore dels.

...

$A_1$

## Reachability



on(a,b),  
ontable(c),  
ontable(b),  
clear(a),  
clear(c),  
handempty

$S_0$

**unstack(a,b)**  
**pickup(c)**

$A_0$

on(a,b),  
ontable(c),  
ontable(b),  
clear(a),  
clear(c),  
handempty,  
**holding(a),**  
**clear(b),**  
**holding(c)**

$S_1$

**stack(c,b)**

...

$A_1$

but  
stack(c,b)  
cannot be  
executed  
after one  
step

to reach  
on(c,b)  
requires 4  
actions

on(c,b),  
...

# Properties of state and action layers

**Proposition.** Let  $a_0, \dots, a_{n-1}$  be a sequence of actions applicable in  $S_0$ . Let  $s_0 = S_0$ , and for  $i < n$ ,  $s_{i+1} = s_i \cup \text{add}(a_i) - \text{del}(a_i)$ . Then for  $i < n$ , there exist  $j, k \leq i$  s.t.  $a_i \in A_k$  and  $s_i \subseteq S_j$ .

We prove by induction on  $i$ :

- Basis:  $i = 0$ , obviously.
- Induction: Let  $i < n$ . Assume there exists  $j \leq i$  s.t.  $s_i \subseteq S_j$ .
- Since  $\text{pre}(a_i) \subseteq s_i$ ,  $\text{pre}(a_i) \subseteq S_j$ .
- Let  $k$  be the least  $u \leq j$  s.t.  $\text{pre}(a_i) \subseteq S_u$ .
- Then  $a_i \in A_k$ .
- So  $\text{add}(a_i) \subseteq S_{k+1} \subseteq S_{j+1}$ .
- Thus  $s_{i+1} \subseteq s_i \cup \text{add}(a_i) \subseteq S_j \cup S_{j+1} = S_{j+1}$ .

**Theorem.** Suppose the state layer stops changing and the goal is not satisfied. Then the original planning problem is not solvable.

Proof:

- Assume that  $a_0, \dots, a_{n-1}$  is a solution to the original problem.
- Then  $Goal \subseteq s_n$ .
- By Proposition, there exists  $m \leq n$  s.t.  $Goal \subseteq s_n \subseteq S_m$ .
- This contradicts the assumption.

- Suppose the goal  $G$  is contained in the state layer
- We want to compute a good relaxed plan
- The idea is to choose a minimum subset of  $A_i$  for each  $i$

# CountActions

CountActions( $G, S_K$ ):

/\* Here  $G$  is contained in  $S_K$ , and we compute the number of actions contained in a relaxed plan achieving  $G$ . \*/

- If  $K = 0$  return 0
- Split  $G$  into  $G_P = G \cap S_{K-1}$  and  $G_N = G - G_P$ 
  - $G_P$  contains the previously achieved (in  $S_{K-1}$ ), and
  - $G_N$  contains the just achieved parts of  $G$  (only in  $S_K$ ).
- Find a **minimal** set of actions  $A$  whose add effects cover  $G_N$ .
  - cannot contain redundant actions,
  - **but may not be the minimum sized set**  
(computing the minimum sized set of actions is the set cover problem and is NP-Hard)
- $\text{NewG} := G_P \cup \text{preconditions of } A$ .
- return  $\text{CountAction}(\text{NewG}, S_{K-1}) + \text{size}(A)$

# The set cover problem

- Given a set of elements  $\{1, 2, \dots, n\}$  (called the universe) and a collection  $S$  of  $m$  sets whose union equals the universe
- Identify the smallest sub-collection of  $S$  whose union equals the universe.
- e.g.,  $U = \{1, 2, 3, 4, 5\}$ ,  $S = \{\{1, 2, 3\}, \{2, 4\}, \{3, 4\}, \{4, 5\}\}$ .  
A smallest cover is  $\{\{1, 2, 3\}, \{4, 5\}\}$ .



# An example

legend: [pre]act[add]

$$S_0 = \{f_1, f_2, f_3\}$$

$$A_0 = \{[f_1]a_1[f_4], [f_2]a_2[f_5]\}$$

$$S_1 = \{f_1, f_2, f_3, f_4, f_5\}$$

$$A_1 = \{[f_2, f_4, f_5]a_3[f_6]\}$$

$$S_2 = \{f_1, f_2, f_3, f_4, f_5, f_6\}$$

**Goal:  $f_6, f_5, f_1$**

**Actions:**

$[f_1]a_1[f_4]$

$[f_2]a_2[f_5]$

$[f_2, f_4, f_5]a_3[f_6]$

# An example

legend: [pre]act[add]

$$S_0 = \{f_1, f_2, f_3\}$$

$$A_0 = \{[f_1]a_1[f_4], [f_2]a_2[f_5]\}$$

$$S_1 = \{f_1, f_2, f_3, f_4, f_5\}$$

$$A_1 = \{[f_2, f_4, f_5]a_3[f_6]\}$$

$$S_2 = \{f_1, f_2, f_3, f_4, f_5, f_6\}$$

$$G = \{f_6, f_5, f_1\}$$

$$G_N = \{f_6\} \text{ (newly achieved)}$$

$$G_p = \{f_5, f_1\} \text{ (achieved before)}$$

# An example

legend: [pre]act[add]

$$S_0 = \{f_1, f_2, f_3\}$$

$$A_0 = \{[f_1]a_1[f_4], [f_2]a_2[f_5]\}$$

$$S_1 = \{f_1, f_2, f_3, f_4, f_5\}$$

$$A_1 = \{[f_2, f_4, f_5]a_3[f_6]\}$$

$$S_2 = \{f_1, f_2, f_3, f_4, f_5, f_6\}$$

$$G = \{f_6, f_5, f_1\}$$

We split  $G$  into  $G_P$  and  $G_N$ :

CountActs( $G, S_2$ )

$G_P = \{f_5, f_1\}$  //already in  $S_1$

$G_N = \{f_6\}$  //New in  $S_2$

$A = \{a_3\}$  //adds all in  $G_N$

//the new goal:  $G_P \cup \text{Pre}(A)$

$G_1 = \{f_5, f_1, f_2, f_4\}$

Return

$1 + \text{CountActs}(G_1, S_1)$

# An example

**Now, we are at level S1**

$$S_0 = \{f_1, f_2, f_3\}$$

$$A_0 = \{[f_1]a_1[f_4], [f_2]a_2[f_5]\}$$

$$S_1 = \{f_1, f_2, f_3, f_4, f_5\}$$

$$A_1 = \{[f_2, f_4, f_5]a_3[f_6]\}$$

$$S_2 = \{f_1, f_2, f_3, f_4, f_5, f_6\}$$

$$G_1 = \{f_5, f_1, f_2, f_4\}$$

**We split  $G_1$  into  $G_P$  and  $G_N$ :**

$$G_N = \{f_5, f_4\}$$

$$G_P = \{f_1, f_2\}$$

**CountActs( $G_1, S_1$ )**

$G_P = \{f_1, f_2\}$  //already in  $S_0$

$G_N = \{f_4, f_5\}$  //New in  $S_1$

$A = \{a_1, a_2\}$  //adds all in  $G_N$

//the new goal:  $G_P \cup \text{Pre}(A)$

$G_2 = \{f_1, f_2\}$

Return

$2 + \text{CountActs}(G_2, S_0)$

Next, we are at level  $S_0$ , simply return 0

So, in total  $\text{CountActs}(G, S_2) = 1 + 2 + 0 = 3$

**Theorem.** Suppose that  $Goal \subseteq S_k$ . For  $i < k$ , let  $A'_{i-1}$  denote the  $A$  obtained by calling  $\text{CountActions}(G_i, S_i)$ . Then  $A'_0, \dots, A'_{k-1}$  is a relaxed plan.

Proof: We prove by induction on  $i \leq k$  that  $A'_0, \dots, A'_{i-1}$  is a relaxed plan for achieving  $G_i$ .

- Basis:  $i = 0$ . We have  $G_0 = G_1 \cap S_0 \cup \text{pre}(A'_0)$ . Since  $\text{Pre}(A_0) \subseteq S_0$ ,  $G_0 \subseteq S_0$ . So the empty plan achieves  $G_0$ .
- Induction: Let  $i < k$ . Assume  $A'_0, \dots, A'_{i-1}$  achieves  $G_i$ .
- We have  $G_P = G_{i+1} \cap S_i$ ,  $G_i = G_P \cup \text{pre}(A'_i)$ , and  $G_N = G_{i+1} - G_P \subseteq \text{add}(A'_i)$ .
- So  $\text{pre}(A'_i) \subseteq G_i$ ,  $G_{i+1} - G_i \subseteq G_{i+1} - G_P \subseteq \text{add}(A'_i)$ .
- Thus  $A'_0, \dots, A'_{i-1}, A'_i$  is a relaxed plan for achieving  $G_{i+1}$ .

However, CountActions does NOT compute the length of an optimal relaxed plan, because.

- The choice of which action set to use to achieve  $G_N$  (“just achieved part of  $G$ ”) is not necessarily optimal
- Even if we picked a true minimum set  $A$  at each stage, we might not obtain a minimum set of actions for the entire plan, since the set  $A$  picked at each stage influences what set can be used at the next stage!

It is NP-Hard to compute an optimal relaxed plan.

- So CountActions cannot be made into an admissible heuristic without making it much harder to compute.
- Empirically, refinements of CountActions performs very well on a number of sample planning domains.

# An exercise: blocks world planning

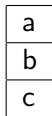
There are a collection of blocks: a block can be on the table, or on top of another block. There are three predicates:

- $clear(x)$ : there is no block on top of block  $x$ ;
- $on(x, y)$ : block  $x$  is on top of block  $y$ ; and
- $onTable(x)$ : block  $x$  is on the table.

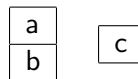
There are three actions:

- $move(x, y, z)$ : move block  $x$  from block  $y$  onto block  $z$ , provided  $x$  is on  $y$ , both  $x$  and  $z$  are clear;
- $moveFromTable(x, y)$ : move block  $x$  from the table onto block  $y$ , provided  $x$  is on the table, both  $x$  and  $y$  are clear; and
- $moveToTable(x, y)$ : move block  $x$  from block  $y$  onto the table, provided  $x$  is on  $y$ , and  $x$  is clear.

The initial state is:



The goal state is:





# An exercise: blocks world planning

- 1 Write the STRIPS representation of the actions, the initial KB, and the goal.
- 2 Use reachability analysis to compute the heuristic value for the initial state. Draw the state and action layers. For each call of CountActions, indicate the values of  $G$ ,  $G_P$ ,  $G_N$ , and  $A$ .