

## 1 Infinity of Prime Numbers

### 1.1 Euler Product

We consider the following product

$$\prod_{p \in \mathbb{P}} \frac{1}{1 - \frac{1}{p^s}} = \prod_{p \in \mathbb{P}} \sum_{k=0}^{+\infty} \frac{1}{p^{ks}} \quad (1)$$

As every integer can be represented as

$$n = \prod_{\substack{p_i \in \mathbb{P} \\ \alpha_i \geq 0}} p_i^{\alpha_i} \longrightarrow \frac{1}{n^s} = \prod_{\substack{p_i \in \mathbb{P} \\ \alpha_i \geq 0}} \frac{1}{p_i^{\alpha_i s}} \quad (2)$$

Then

$$\prod_{p \in \mathbb{P}} \sum_{k=0}^{+\infty} \frac{1}{p^{ks}} = \sum_{n=1}^{+\infty} \frac{1}{n^s} \quad (3)$$

### 1.2 Infinity of Prime Numbers

Assume that there are a finite number of prime numbers, denoted as  $\mathbb{P} = \{p_1, p_2, \dots, p_s\}$ , then

$$\sum_{n=1}^N \frac{1}{n} < \sum_{n=1}^{+\infty} \frac{1}{n} = \prod_{p \in \mathbb{P}} \frac{1}{1 - \frac{1}{p^s}} = \prod_{p_i}^s \frac{1}{1 - \frac{1}{p_i^s}} \quad (4)$$

As  $N \rightarrow \infty$ , the harmonic series diverges. Hence the inequality contradicts with the fact that the RHS is a constant. Thus we can conclude that  $s = \infty$ , i.e. prime numbers are infinite many.

## 2 The Greatest Common Divisor Theory

We only give an important lemma of the GCD theory.

### 2.1 Lemma 2.1

$$\text{lcm}(a_1, a_2) \text{gcd}(a_1, a_2) = a_1 a_2 \quad (5)$$

**Proof** It is intuitive that

$$\text{gcd}\left(\frac{a_1}{\text{gcd}(a_1, a_2)}, \frac{a_2}{\text{gcd}(a_1, a_2)}\right) = 1 \rightarrow \text{lcm}\left(\frac{a_1}{\text{gcd}(a_1, a_2)}, \frac{a_2}{\text{gcd}(a_1, a_2)}\right) = \frac{a_1}{\text{gcd}(a_1, a_2)} \cdot \frac{a_2}{\text{gcd}(a_1, a_2)}$$

Then

$$\text{gcd}(a_1, a_2)^2 \text{lcm}\left(\frac{a_1}{\text{gcd}(a_1, a_2)}, \frac{a_2}{\text{gcd}(a_1, a_2)}\right) = a_1 a_2 = \text{lcm}(a_1, a_2) \text{gcd}(a_1, a_2) \quad (6)$$

## 2.2 Fermat's Little Theorem

Fermat's little theorem states that if  $p$  is a prime number, then

$$p \mid a^p - a, \quad a \in \mathbb{Z} \quad (7)$$

$$p \mid a^{p-1} - 1, \quad a \in \mathbb{Z}, \quad \gcd(a, p) = 1 \quad (8)$$

### Proof

**Lemma 2.2** For integer  $1 \leq j \leq p-1$ ,

$$p \mid \binom{p}{j} \quad (9)$$

As  $p$  is a prime number, then for  $1 \leq j \leq p-1$ ,

$$\gcd(p, j) = 1 \rightarrow \gcd(p, j) = \gcd(p, p-j) = 1 \rightarrow \gcd(p, j!(p-j)!) = 1 \quad (10)$$

As the combinatorial number is an integer,

$$\binom{p}{j} \in \mathbb{Z} \rightarrow j!(p-j)! \mid p! \rightarrow j!(p-j)! \mid (p-1)! \quad (11)$$

Then we can conclude that

$$p \mid \frac{p!}{j!(p-j)!} = \binom{p}{j} \quad (12)$$

Then we may use mathematical induction in proving. For  $a = 1$ ,  $p \mid 0$  holds. Assume that for  $a = n$  the theorem holds, then for  $a = n+1$ ,

$$(n+1)^p - (n+1) = \sum_{i=0}^p \binom{p}{i} n^i - n - 1 = n^p - n + \sum_{j=1}^{p-1} \binom{p}{j} n^j \quad (13)$$

Applying the condition and the lemma,

$$p \mid n^p - n, \quad p \mid \binom{p}{j} n^j \rightarrow p \mid (n+1)^p - (n+1)$$

When  $a, p$  are coprime,

$$p \mid a^p - a = a(a^{p-1} - 1) \rightarrow p \mid a^{p-1} - 1 \quad (14)$$

Thus we prove the **Fermat's little theorem**. ■

## 2.3 Lemma 2.2 (Existence of Modular Multiplicative Inverse)

① If  $m \geq 2$ ,  $\gcd(m, a) = 1$ , then integer  $d \leq m-1$  exists, such that  $m \mid a^d - 1$ .

**Proof** For every integer  $1 \leq j \leq m$ , the pseudo-division gives

$$a^j = q_j m + r_j \leftrightarrow a^j \equiv r_j \pmod{m}, \quad r_j \neq 0 \quad (15)$$

As there are  $j - 1$  possible values for  $a^j \pmod m$ , and  $j$  numbers of residual  $r_j$ , then there must exists  $k, b \in [1, m]$ , such that

$$\begin{cases} a^k = q_k m + r_k \\ a^b = q_b m + r_b \end{cases}, \quad r_k = r_b \quad (16)$$

Assume that  $k > b$ , then

$$a^k - a^b = a^b(a^{k-b} - 1) = (q_k - q_b)m \quad (17)$$

As  $\gcd(m, a) = 1$ , it is easy to prove that  $\gcd(m, a^b) = 1$ . Hence,

$$(q_k - q_b)m \mid a^b(a^{k-b} - 1) \rightarrow m \mid (a^{k-b} - 1) \xrightarrow{d=k-b} m \mid a^d - 1 \quad (18)$$

② If  $d_0$  is the least integer in the set of  $d$  in ①, called the **modular exponentiation** of  $a$  to  $m$ , then  $m \mid a^h - 1$  if and only if  $d_0 \mid h$ .

### Proof

**Sufficiency** If  $d_0 \mid h$ , then

$$h = qd_0 \rightarrow a^h - 1 = a^{qd_0} - 1 = (a^{d_0} - 1) \sum_{i=0}^{q-1} a^{id_0} \equiv 0 \pmod m \quad (19)$$

**Necessity** The pseudo-division gives

$$h = qd_0 + r \quad 0 \leq r < d_0 \quad (20)$$

Substitute in we obtain

$$a^h - 1 = a^{qd_0+r} - 1 = a^r(a^{qd_0} - 1) + a^r - 1 \quad (21)$$

If  $a^h - 1 \equiv 0 \pmod m$ , then  $m \mid a^r - 1$ . As  $0 \leq r < d_0$ , and the condition that  $d_0$  is the least integer, the only value  $r$  can take is  $r = 0$ . Thus  $d_0 \mid h$ . ■

## 3 Fundamental Theorem of Arithmetic

### 3.1 Content

#### 3.1.1 Theorem 1

If  $p$  is a prime number, and  $p \mid \prod_{i=1}^k a_i$ , then

$$p \mid a_j \quad 1 \leq j \leq k \quad (22)$$

holds for at least one  $j$ .

### 3.1.2 Theorem 2

Any integer  $a > 1$  can be uniquely represented as

$$a = p_1 p_2 \cdots p_s \quad (23)$$

where  $p_j$   $1 \leq j \leq s$  are all prime numbers.

**Proof** Assume that the prime numbers are arranged in non-decreasing order, i.e.  $p_1 \leq p_2 \leq \cdots \leq p_s$ . If there is another decomposition

$$a = p_1 p_2 \cdots p_s = q_1 q_2 \cdots q_r \quad 1 \leq 2 \leq \cdots \leq q_r \quad (24)$$

As  $q_1 \mid a$ ,  $p_1 \mid a$ , then

$$\begin{cases} q_1 \mid p_1 p_2 \cdots p_s \rightarrow \exists p_i, q_1 \mid p_i \quad 1 \leq i \leq s \rightarrow p_i = q_1 \\ p_1 \mid q_1 q_2 \cdots q_r \rightarrow \exists q_j, p_1 \mid q_j \quad 1 \leq j \leq r \rightarrow p_1 = q_j \end{cases} \longrightarrow p_1 \leq p_i = q_1 \leq q_j = p_1 \quad (25)$$

Hence  $p_1 = q_1$ . Likewise we can derive that  $p_i = q_i$  for  $1 \leq i \leq \min(r, s)$ . Assume that  $r \geq s$ , then

$$q_{s+1} q_{s+2} \cdots q_r = 1 \quad (26)$$

which contradicts with the assumption that  $q_i$  is a prime number, unless  $r = s$ . ■

### 3.2 Corollary 1

If  $p_1, p_2, \dots, p_s$  are all prime integers, and

$$a = \prod_{i=1}^s p_i^{\alpha_i} \quad b = \prod_{i=1}^s p_i^{\beta_i} \quad (27)$$

then

$$\begin{aligned} \gcd(a, b) &= \prod_{i=1}^s p_i^{\delta_i} \quad \delta_i = \min(\alpha_i, \beta_i) \\ \text{lcm}(a, b) &= \prod_{i=1}^s p_i^{\gamma_i} \quad \gamma_i = \max(\alpha_i, \beta_i) \end{aligned} \quad (28)$$

Additionally, the summation of the divisors of integer  $a$ , denoted as  $\sigma(a)$ , can be written as

$$\sigma(a) = \prod_{i=1}^s \sigma(p_i^{\alpha_i}) = \prod_{i=1}^s \sum_{j=0}^{\alpha_i} p_i^j = \prod_{i=1}^s \frac{p_i^{\alpha_i+1} - 1}{p_i - 1} \quad (29)$$

## 4 The Exponent of Prime Factors

For an integer  $n$ , the integer  $\alpha$  for prime number  $p$  such that  $p^\alpha \parallel n!$  can be written as

$$\alpha(p, n) = \sum_{i=1}^{\infty} \left[ \frac{n}{p^i} \right] \quad (30)$$

## 4.1 Derivation

Denote the magnitude of the set  $\{x | x \bmod p^i = 0, 1 \leq x \leq n\}$  as  $c_i$ . Then the magnitude of the set  $\{x | p^i \parallel x, 1 \leq x \leq n\}$  is  $d_i = c_i - c_{i+1}$ . We can conclude that

$$d_i = c_i - c_{i+1} = \left\lfloor \frac{n}{p^i} \right\rfloor - \left\lfloor \frac{n}{p^{i+1}} \right\rfloor \quad (31)$$

Thus

$$\alpha(p, n) = \sum_{i=0}^{\infty} i d_i = \sum_{i=1}^k \left\lfloor \frac{n}{p^i} \right\rfloor = \sum_{i=1}^{\infty} \left\lfloor \frac{n}{p^i} \right\rfloor, \quad p^k \parallel n \quad (32)$$

## 5 Corollary 1

One conclusion can be drawn that

$$n!(m!)^n \mid (mn)! \quad (33)$$

**Proof** Consider an arbitrary prime factor of the integer  $m$ . We need to prove that

$$\alpha(p, n) + n\alpha(p, m) \leq \alpha(p, mn) \Leftrightarrow \sum_{j=1}^{\infty} \left\lfloor \frac{n}{p^j} \right\rfloor + n \sum_{j=1}^{\infty} \left\lfloor \frac{m}{p^j} \right\rfloor \leq \sum_{j=1}^{\infty} \left\lfloor \frac{mn}{p^j} \right\rfloor \quad (34)$$

Consider  $m = p_1^{\alpha_1} p_2^{\alpha_2} \cdots p_s^{\alpha_s}$ , then write  $m = cp^l$ . For  $j \leq l$ , we have

$$\left\lfloor \frac{mn}{p^j} \right\rfloor = cnp^{l-j} = n \left\lfloor \frac{m}{p^j} \right\rfloor \rightarrow \sum_{j=1}^l \left\lfloor \frac{mn}{p^j} \right\rfloor = \sum_{j=1}^l n \left\lfloor \frac{m}{p^j} \right\rfloor \quad (35)$$

For  $j > l$ , we have the pseudo-division  $m = q_j p^j + r_j$ , where  $r_j \in [1, p^j - 1]$ . Thus

$$\begin{aligned} \left\lfloor \frac{mn}{p^j} \right\rfloor &= \left\lfloor nq_j + n \frac{r_j}{p^j} \right\rfloor = nq_j + \left\lfloor \left\{ \frac{m}{p^j} \right\} n \right\rfloor \geq n \left\lfloor \frac{m}{p^j} \right\rfloor + \left\lfloor \frac{n}{p^{j-l}} \right\rfloor \\ &\rightarrow \sum_{j>l} \left\lfloor \frac{mn}{p^j} \right\rfloor \geq n \sum_{j>l} \left\lfloor \frac{m}{p^j} \right\rfloor + \sum_{j-l>0} \left\lfloor \frac{n}{p^{j-l}} \right\rfloor \end{aligned} \quad (36)$$

Sum the two equations up we obtain ■

$$\sum_j \left\lfloor \frac{mn}{p^j} \right\rfloor \geq n \sum_j \left\lfloor \frac{m}{p^j} \right\rfloor + \sum_j \left\lfloor \frac{n}{p^j} \right\rfloor$$

## 6 Euclid Algorithm

### 6.1 Content

If two integers  $u_0, u_1$ ,  $u_1 \nmid u_0$ . Then we can give the following pseudo-division method

$$\begin{aligned} u_0 &= q_0 u_1 + u_2, & 0 < u_2 < |u_1| \\ u_1 &= q_1 u_2 + u_3, & 0 < u_3 < u_2 \\ &\vdots \\ u_{k-1} &= q_{k-1} u_k + u_{k+1}, & 0 < u_k < u_{k-1} \\ u_k &= q_k u_{k+1} \end{aligned} \quad (37)$$

And

$$u_{k+1} = \gcd(u_0, u_1) \quad (38)$$

Or

$$\gcd(u_i, u_j) = \gcd(u_j, u_i \bmod u_i) \quad (39)$$

The equations above indicates that the greatest common divisor of integers  $a_1, a_2, \dots, a_k$ , the coefficients  $x_1, x_2, \dots, x_k$  exists, such that

$$\gcd(a_1, a_2, \dots, a_k) = a_1x_1 + a_2x_2 + \dots + a_kx_k \quad (40)$$

## 6.2 Lemma 1.1

A lemma can be given that

$$\gcd(2^m - 1, 2^n - 1) = 2^{\gcd(m,n)} - 1 \quad (41)$$

**Proof** As

$$m = qn + r \rightarrow 2^m - 1 = 2^{qn+r} - 2^r + 2^r - 1 = 2^r(2^{qn} - 1) + 2^r - 1 \quad (42)$$

Hence

$$\gcd(2^m - 1, 2^n - 1) = \gcd(2^r - 1, 2^n - 1) = \gcd(2^n - 1, 2^{\gcd(m,n)} - 1) = \dots = 2^{\gcd(m,n)} - 1$$

## 7 Extended Euclidean Algorithm

Based on the Euclidean algorithm, the extended one solve the equation

$$ax + by = \gcd(a, b) \quad (43)$$

and calculates the modular multiplicative inverse at the same time.

### 7.1 Implementation

#### 7.1.1 Particular Solution

Assume that there are two equations, where

$$\begin{cases} ax_0 + by_0 = \gcd(a, b) \\ bx_1 + (a \bmod b)y_1 = \gcd(b, a \bmod b) = \gcd(a, b) \end{cases} \quad (44)$$

Then we can conclude that

$$ax_0 + by_0 = bx_1 + (a \bmod b)y_1 = bx_1 + (a - b\lfloor \frac{a}{b} \rfloor)y_1 \quad (45)$$

As  $a, b$  are arbitrary integers, we have

$$b(x_1 - \lfloor \frac{a}{b} \rfloor y_1 - y_0) = a(y_1 - x_0) \rightarrow \begin{cases} x_0 = y_1 \\ y_0 = x_1 - \lfloor \frac{a}{b} \rfloor y_1 \end{cases} \quad (46)$$

Then by recursively repeating the bottom-up recursion that

$$\begin{cases} x_k = y_{k+1} \\ y_k = x_{k+1} - \lfloor \frac{a_k}{b_k} \rfloor y_{k+1} \end{cases} \begin{cases} a_{k+1} = b_k \\ b_{k+1} = a_k \bmod b_k \\ \gcd(a_{k+1}, b_{k+1}) = \gcd(a_k, a_k) \end{cases} \quad (47)$$

The terminal state is the equation  $ax_n + by_n = \gcd(a_n, b_n) = \gcd(a, b)$ , where  $b_n = 0$ . We instantly solve that  $x_n = 1, y_n = 0, a_n = \gcd(a, b), b_n = 0$ . If we assign the terminal value to the bottom variables, then the traceback in every step will calculate the particular solution  $(x_0, y_0)$  for the original equation.

### 7.1.2 Complementary Solution

From the particular solution the complementary solution can be derived. Consider the original equation

$$ax_0 + by_0 = \gcd(a, b) \rightarrow a \left( x_0 - \frac{b}{\gcd(a, b)} t \right) + b \left( y_0 - \frac{a}{\gcd(a, b)} t \right) = \gcd(a, b) \quad (48)$$

Hence, we conclude that the complementary solution for the equation is

$$\begin{cases} x = x_0 - \frac{b}{\gcd(a, b)} t \\ y = y_0 - \frac{a}{\gcd(a, b)} t \end{cases} \quad (49)$$

## 8 Non-zero / Positive Solution to Diphantine Equation