

APPENDIX

The appendices are organized as follows:

- In Appendix A, we provide a detailed exposition of quaternions and their role in representing rotations, along with several fundamental definitions, notations, and lemmas essential for subsequent analysis.
- In Appendix B, we characterize the recovery accuracy of the solution of (Q-QP) with respect to the ground truth \tilde{q} .
- In Appendix C, we establish the tightness of semidefinite relaxation, which ensures that the rank-one solution of the convex problem (Q-SDP) yields a global optimum for the nonconvex problem (Q-QP).

A. NOTATION AND PRELIMINARIES

A.1. Basic Notations

The fields of real numbers, quaternion numbers, and unit quaternion numbers are denoted by \mathbb{R} , \mathbb{Q} , and \mathbb{U} , respectively. Throughout this paper, scalars, vectors, matrices, and quaternions are denoted by lowercase letters (e.g., x), boldface lowercase letters (e.g., \mathbf{x}), boldface capital letters (e.g., X), and lowercase letters with tilde (e.g., \tilde{x}), respectively. The special orthogonal group $SO(3)$ is the set of three-dimensional rotations that is formally defined by $SO(3) := \{R \in \mathbb{R}^{3 \times 3} \mid R^\top R = I_3, \det(R) = 1\}$. We denote by $\mathbb{S}^{n \times n}$ the set of all symmetric matrices. $\mathbb{1}$ is a vector of all-ones. The notation $\|\cdot\|_2$ denotes the ℓ_2 -norm of vectors or the spectral norm of matrices.

A.2. Quaternion and rotation

A quaternion number $\tilde{q} \in \mathbb{Q}$, proposed by Hamilton, has the form $\tilde{q} = q_0 + q_1\mathbf{i} + q_2\mathbf{j} + q_3\mathbf{k}$, where $q_0, q_1, q_2, q_3 \in \mathbb{R}$ and $\mathbf{i}, \mathbf{j}, \mathbf{k}$ are three imaginary units. We may also write $\tilde{q} = [q_0, q_1, q_2, q_3] = [q_0, \mathbf{q}] \in \mathbb{R}^4$ as the vector representation where $\mathbf{q} = [q_1, q_2, q_3] \in \mathbb{R}^3$ for convenience. We note that we also regard the above representation as a column vector and its transpose $[q_0, \mathbf{q}]^\top$ a row vector. The sum of \tilde{p} and \tilde{q} is defined as $\tilde{p} + \tilde{q} = [p_0 + q_0, \mathbf{p} + \mathbf{q}]$. The product of \tilde{p} and \tilde{q} is defined by

$$\tilde{p}\tilde{q} = [p_0q_0 - \mathbf{p} \cdot \mathbf{q}, p_0\mathbf{q} + q_0\mathbf{p} + \mathbf{p} \times \mathbf{q}],$$

where $\mathbf{p} \cdot \mathbf{q}$ is the dot product, and $\mathbf{p} \times \mathbf{q}$ is the cross product of \mathbf{p} and \mathbf{q} . Thus, in general, $\tilde{p}\tilde{q} \neq \tilde{q}\tilde{p}$, and we have $\tilde{p}\tilde{q} = \tilde{q}\tilde{p}$ if and only if $\mathbf{p} \times \mathbf{q} = \mathbf{0}$, i.e., either $\mathbf{p} = \mathbf{0}$ or $\mathbf{q} = \mathbf{0}$, or $\mathbf{p} = \alpha\mathbf{q}$ for several real number α . The multiplication of quaternions is associative and distributive over vector addition, but is not commutative.

The conjugate of \tilde{q} is the quaternion $\tilde{q}^* = q_0 - q_1\mathbf{i} - q_2\mathbf{j} - q_3\mathbf{k}$. Then, $(\tilde{p}\tilde{q})^* = \tilde{q}^*\tilde{p}^*$ for any $\tilde{p}, \tilde{q} \in \mathbb{Q}$. The magnitude of \tilde{q} is defined by $|\tilde{q}| = \sqrt{\tilde{q}\tilde{q}^*} = \sqrt{\tilde{q}^*\tilde{q}}$. And \tilde{q} is invertible if and only if $|\tilde{q}|$ is positive. In this case, we have $\tilde{q}^{-1} = \tilde{q}^*/|\tilde{q}|$.

The set of all unit quaternions is $\mathbb{U} := \{\tilde{q} \in \mathbb{R}^4 \mid |\tilde{q}| = 1\}$, which can be regarded as a unit sphere in \mathbb{R}^4 . Equivalently, a unit quaternion has the following form: $\tilde{q} = [\cos(\theta/2), \sin(\theta/2)\mathbf{n}]$, where $\mathbf{n} = (n_x, n_y, n_z)$ is a unit vector and θ is an angle. Let a vector $\mathbf{t}_1 \in \mathbb{R}^3$ rotates θ radians around axis \mathbf{n} to reach $\mathbf{t}_2 \in \mathbb{R}^3$. This process can be represented by a quaternion as

$$[0, \mathbf{t}_2] = \tilde{q}[0, \mathbf{t}_1]\tilde{q}^*.$$

Using rotation matrix in $SO(3)$, we also have $\mathbf{t}_2 = R\mathbf{t}_1$, where

$$R = \cos(\theta)I_3 + (1 - \cos(\theta))\mathbf{n}\mathbf{n}^\top + \sin(\theta)\mathbf{n}^\wedge, \quad \text{and} \quad \mathbf{n}^\wedge = \begin{pmatrix} 0 & -n_z & n_y \\ n_z & 0 & -n_x \\ -n_y & n_x & 0 \end{pmatrix}.$$

The relationship between rotation matrix and unit quaternion is given in the next lemma.

Lemma A.1. [12] Given a unit quaternion $\tilde{q} = [q_0, \mathbf{q}] = [q_0, q_1, q_2, q_3] \in \mathbb{U}$ and a vector $\mathbf{t} \in \mathbb{R}^3$. Then $[0, R\mathbf{t}] = \tilde{q}[0, \mathbf{t}]\tilde{q}^*$, where the rotation matrix $R \in SO(3)$ satisfies

$$R = \mathbf{q}\mathbf{q}^\top + q_0^2I + 2q_0\mathbf{q}^\wedge + (\mathbf{q}^\wedge)^2.$$

If the rotation matrix R is compound motion of two rotations, i.e., $R = R_2R_1$, then the corresponding quaternion \tilde{q} can be formulated as $\tilde{q} = \tilde{q}_2\tilde{q}_1$.

A.3. Quaternion matrix

The collections of real, complex and quaternion $m \times n$ matrices are denoted by $\mathbb{R}^{m \times n}$, $\mathbb{C}^{m \times n}$, $\mathbb{Q}^{m \times n}$, respectively.

A quaternion matrix [23] $\tilde{A} = (\tilde{a}_{ij}) \in \mathbb{Q}^{m \times n}$ can be denoted as

$$\tilde{A} = A_0 + A_1 \mathbf{i} + A_2 \mathbf{j} + A_3 \mathbf{k},$$

where $A_0, A_1, A_2, A_3 \in \mathbb{R}^{m \times n}$. The conjugate transpose of \tilde{A} is $\tilde{A}^* = (a_{ji}^*)$. If $\tilde{A} = \tilde{A}^*$, then \tilde{A} is a Hermitian matrix. We denote $\mathbb{H}^{n \times n} = \{\tilde{A} \in \mathbb{Q}^{n \times n} \mid \tilde{A} = \tilde{A}^*\}$ be the set of Hermitian matrices.

For $\tilde{A}, \tilde{B} \in \mathbb{Q}^{m \times n}$, their inner product is defined as

$$\langle \tilde{A}, \tilde{B} \rangle = \text{tr}(\tilde{A}^* \tilde{B}),$$

where $\text{tr}(\tilde{A}^* \tilde{B})$ denotes the trace of matrix $\tilde{A}^* \tilde{B}$. The Frobenius norm of A is given by

$$\|\tilde{A}\|_F = \sqrt{\langle \tilde{A}, \tilde{A} \rangle} = \sqrt{\text{tr}(\tilde{A}^* \tilde{A})} = \sqrt{\sum_{i=1}^m \sum_{j=1}^n |\tilde{a}_{ij}|^2}.$$

In addition, we present some useful properties of quaternion matrices which generalize those in the complex case. It is important to note that, in general, $\text{tr}(\tilde{A}\tilde{B}) \neq \text{tr}(\tilde{B}\tilde{A})$ due to the non-commutativity of quaternion multiplication. This is a key difference from the complex matrix case, and as a result, some properties require careful consideration.

Theorem A.1. Suppose $\tilde{A}, \tilde{B} \in \mathbb{Q}^{n \times n}$ are two any quaternion matrix, then the following statements hold:

- (i) $(\tilde{A}\tilde{B})^* = \tilde{B}^* \tilde{A}^*$.
- (ii) $\text{tr}(\tilde{A}^*) = \text{tr}(\tilde{A})^*$.
- (iii) $\text{tr}(\tilde{A}^* \tilde{B}) = \text{tr}(\tilde{B}^* \tilde{A})^*$, or $\langle \tilde{A}, \tilde{B} \rangle = \langle \tilde{B}, \tilde{A} \rangle^*$.
- (iv) If $\tilde{A} \in \mathbb{H}^{n \times n}$, then $\tilde{x}^* \tilde{A} \tilde{x}$ is a real number for any $\tilde{x} \in \mathbb{Q}^n$.
- (v) $\text{Re}(\text{tr}(\tilde{A}\tilde{B})) = \text{Re}(\text{tr}(\tilde{B}\tilde{A}))$.
- (vi) If $\tilde{A} \in \mathbb{H}^{n \times n}$ and $\tilde{x} \in \mathbb{Q}^n$, then $\text{tr}(\tilde{x}\tilde{x}^* \tilde{A}) + \text{tr}(\tilde{A}\tilde{x}\tilde{x}^*) = 2\tilde{x}^* \tilde{A} \tilde{x}$.

Proof. We omit the proofs of (i)–(v) which are not hard to verify. Here we provide a proof of (vi), a result that will be frequently employed in the subsequent analysis. From (ii) and (iii), we have

$$\text{tr}(\tilde{x}\tilde{x}^* \tilde{A}) = \text{tr}(\tilde{A}\tilde{x}\tilde{x}^*)^*,$$

which implies that $\text{tr}(\tilde{x}\tilde{x}^* \tilde{A}) + \text{tr}(\tilde{A}\tilde{x}\tilde{x}^*) = 2\text{Re}(\text{tr}(\tilde{x}\tilde{x}^* \tilde{A}))$. According to (iv) and (v), we have

$$\text{Re}(\text{tr}(\tilde{x}\tilde{x}^* \tilde{A})) = \text{Re}(\tilde{x}^* \tilde{A} \tilde{x}) = \tilde{x}^* \tilde{A} \tilde{x},$$

which completes the proof. \square

Lemma A.2. ([27] Corollary 6.2) Suppose $\tilde{X} \in \mathbb{H}^{n \times n}$, then there exists a unitary matrix $\tilde{U} \in \mathbb{Q}^{n \times n}$ (that is, $\tilde{U}\tilde{U}^* = \tilde{U}^*\tilde{U} = I$) such that

$$\tilde{U}^* \tilde{X} \tilde{U} = \text{diag}(\lambda_1, \lambda_2, \dots, \lambda_n), \quad (5)$$

where $\lambda_1, \lambda_2, \dots, \lambda_n$ are real eigenvalues of the quaternion matrix \tilde{X} . Furthermore, if \tilde{X} is a positive definite matrix, then $\lambda_1, \lambda_2, \dots, \lambda_n$ are all positive numbers.

For any Hermitian matrix $\tilde{X} \in \mathbb{H}^{n \times n}$, define its spectral norm by $\|\tilde{X}\|_2 = \max_i |\lambda_i|$. We denote the cones of quaternion Hermitian positive semidefinite and positive definite matrices by $\mathbb{H}_+^{n \times n} = \{\tilde{X} \in \mathbb{H}^{n \times n} \mid \tilde{X} \succeq 0\}$ and $\mathbb{H}_{++}^{n \times n} = \{\tilde{X} \in \mathbb{H}^{n \times n} \mid \tilde{X} \succ 0\}$, respectively.

Lemma A.3. Let the dual cone of $\mathbb{H}_+^{n \times n}$ be

$$(\mathbb{H}_+^{n \times n})^* = \{Y \in \mathbb{H}^{n \times n} \mid \text{tr}(\tilde{X}\tilde{Y}) + \text{tr}(\tilde{Y}\tilde{X}) \geq 0, \forall \tilde{X} \in \mathbb{H}_+^{n \times n}\}.$$

Then we have $\mathbb{H}_+^{n \times n} = (\mathbb{H}_+^{n \times n})^*$.

Proof. Given any $\tilde{Y} \in \mathbb{H}_+^{n \times n}$, by Lemma A.2, we have $\tilde{Y} = \sum_{i=1}^n \lambda_i \tilde{u}_i \tilde{u}_i^*$. Then

$$\text{tr}(\tilde{X}\tilde{Y}) + \text{tr}(\tilde{Y}\tilde{X}) = \sum_{i=1}^n \lambda_i \left(\text{tr}(\tilde{X}\tilde{u}_i \tilde{u}_i^*) + \text{tr}(\tilde{u}_i \tilde{u}_i^* \tilde{X}) \right) = 2 \sum_{i=1}^n \lambda_i \tilde{u}_i^* \tilde{X} \tilde{u}_i,$$

where the last equality follows from Theorem A.1. Since $\lambda_i \geq 0$, $\tilde{u}_i^* \tilde{X} \tilde{u}_i \geq 0$, we have $\tilde{Y} \in (\mathbb{H}_+^{n \times n})^*$, which implies $\mathbb{H}_+^{n \times n} \subseteq (\mathbb{H}_+^{n \times n})^*$.

Next we prove $(\mathbb{H}_+^{n \times n})^* \subseteq \mathbb{H}_+^{n \times n}$. Suppose $\tilde{Y} \notin \mathbb{H}_+^{n \times n}$. Then there exists a $\tilde{q} \in \mathbb{Q}^n$ such that $\tilde{q}^* \tilde{Y} \tilde{q} < 0$. Let $\tilde{X} = \tilde{q} \tilde{q}^* \in \mathbb{H}_+^{n \times n}$, and we have

$$\begin{aligned} \text{tr}(\tilde{X}\tilde{Y}) + \text{tr}(\tilde{Y}\tilde{X}) &= \text{tr}(\tilde{q} \tilde{q}^* \tilde{Y}) + \text{tr}(\tilde{Y} \tilde{q} \tilde{q}^*) \\ &= 2 \tilde{q}^* \tilde{Y} \tilde{q} < 0. \end{aligned}$$

Hence, $\tilde{Y} \notin (\mathbb{H}_+^{n \times n})^*$ which implies $(\mathbb{H}_+^{n \times n})^* \subseteq \mathbb{H}_+^{n \times n}$. □

A.4. Distance metrics for rotations

Assume that R_1, R_2 are two rotation matrices, and \tilde{q}_1, \tilde{q}_2 are corresponding quaternion representations. Define the rotation angle θ_{ij} of two different rotations as

$$\theta = \text{dist}_a(R_1, R_2) = \|\text{Log}(R_1^\top R_2)\|_2 = \arccos \frac{\text{tr}(R_1^\top R_2) - 1}{2},$$

where $\text{Log}(R)$ denotes the logarithm map for $SO(3)$ and return its axis-angle representation \mathbf{n} . It can be shown that this distance is a geodesic distance [11], i.e., it is the length of the minimum path between R_1 and R_2 on the manifold $SO(3)$. Then the chordal distance and the quaternion distance between two rotations can be interpreted as

$$\text{dist}_c(R_1, R_2) = \|R_1 - R_2\|_F = 2\sqrt{2} \sin(\theta/2),$$

$$\text{dist}_q(R_1, R_2) = \|\tilde{q}_1 - \tilde{q}_2\|_F = 2 \sin(\theta/4).$$

B. PROOFS FOR SECTION 4.1

B.1. High-dimensional probability for quaternion random matrix

Definition B.1. A quaternion random variable $\tilde{q} \in \mathbb{Q}$ is said to follow a standard quaternion normal distribution if its components q_0, q_1, q_2, q_3 are independent and identically distributed (i.i.d.) with $q_i \sim \mathcal{N}(0, \frac{1}{4})$. Under this distribution, \tilde{q} satisfies $\mathbb{E}(\tilde{q}) = 0$ and $\mathbb{E}|\tilde{q}|^2 = 1$, and we denote $\tilde{q} \sim \mathcal{QN}(0, 1)$.

Prior to the proof of Theorem 4.1, we present two lemmas as follows.

Lemma B.1. Let $\tilde{A} \in \mathbb{H}^{n \times n}$ be a quaternion Hermitian random matrix with i.i.d. off-diagonal entries following a standard quaternion normal distribution $\mathcal{QN}(0, 1)$ and zeros on the diagonal. Then we have

$$\mathbb{P}\{\|\tilde{A}\|_2 > (2 + c_1)\sqrt{n}\} \leq e^{-c_1^2 n/2}, \quad (6)$$

where $c_1 > 0$ is an arbitrary constant.

Proof. The tail bounds for real and complex cases were established in [28, 20, 29] using Slepian's comparison theorem and Gaussian concentration. However, the non-commutativity of quaternion multiplication prevents trivial extension of these results. We divide the proof into three steps.

Step I: We first bound the largest eigenvalue of a quaternion Hermitian random matrix \tilde{A} . Let

$$\begin{aligned} \sigma_{\max}(\tilde{A}) &= \max_{\tilde{q} \in \mathbb{Q}^n: \|\tilde{q}\|_2=1} \tilde{q}^* \tilde{A} \tilde{q} \\ &= \max_{\tilde{q} \in \mathbb{Q}^n: \|\tilde{q}\|_2=1} \frac{1}{2} \langle \tilde{A}, \tilde{q} \tilde{q}^* \rangle + \frac{1}{2} \langle \tilde{q} \tilde{q}^*, \tilde{A} \rangle \end{aligned}$$

denote the largest eigenvalue of \tilde{A} . For any $\tilde{p}, \tilde{q} \in \mathbb{Q}^n$ which satisfy $\|\tilde{p}\|_2 = \|\tilde{q}\|_2 = 1$, the real valued Gaussian process $X_{\tilde{q}} = \tilde{q}^* \tilde{A} \tilde{q}$ satisfies

$$\begin{aligned} \mathbb{E}(X_{\tilde{q}} - X_{\tilde{p}})^2 &= \frac{1}{4} \mathbb{E} \left(\langle \tilde{A}, \tilde{q} \tilde{q}^* - \tilde{p} \tilde{p}^* \rangle + \langle \tilde{q} \tilde{q}^* - \tilde{p} \tilde{p}^*, \tilde{A} \rangle \right)^2 \\ &= \frac{1}{4} \mathbb{E} \left(\sum_{i,j} \tilde{a}_{ij}^* (\tilde{q}_i \tilde{q}_j^* - \tilde{p}_i \tilde{p}_j^*) + \sum_{i,j} (\tilde{q}_j \tilde{q}_i^* - \tilde{p}_j \tilde{p}_i^*) \tilde{a}_{ij} \right)^2. \end{aligned}$$

Note that, for any $1 \leq i, j \leq n$, \tilde{A} is a Hermitian matrix and the last two terms are quaternion conjugates of each other. Then we have

$$\begin{aligned}\mathbb{E}(X_{\tilde{q}} - X_{\tilde{p}})^2 &= \frac{1}{4} \mathbb{E} \left(4 \sum_{i < j} \text{Re}(\tilde{a}_{ij}^* (\tilde{q}_i \tilde{q}_j^* - \tilde{p}_i \tilde{p}_j^*)) \right)^2 \\ &= 4 \sum_{i < j} \mathbb{E} (\text{Re}(\tilde{a}_{ij}^* (\tilde{q}_i \tilde{q}_j^* - \tilde{p}_i \tilde{p}_j^*)))^2 \\ &= \sum_{i < j} |\tilde{q}_i \tilde{q}_j^* - \tilde{p}_i \tilde{p}_j^*|^2.\end{aligned}$$

The second equality follows from that all a_{ij} , $i < j$ are independent and $\mathbb{E}(a_{ij}) = 0$. The last equality follows from the fact that, for any $\tilde{x} \sim \mathcal{QN}(0, 1)$ and $\tilde{y} \in \mathbb{Q}$,

$$\begin{aligned}\mathbb{E}_{\tilde{x}} (\text{Re}(\tilde{x}^* \tilde{y}))^2 &= \mathbb{E}_{\tilde{x}} (x_0 y_0 + x_1 y_1 + x_2 y_2 + x_3 y_3)^2 \\ &= \mathbb{E}_{\tilde{x}} (x_0^2 y_0^2 + x_1^2 y_1^2 + x_2^2 y_2^2 + x_3^2 y_3^2) \\ &= \frac{1}{4} |\tilde{y}|^2.\end{aligned}$$

Further, since for any $\tilde{x}, \tilde{y} \in \mathbb{Q}$,

$$|\tilde{x} \tilde{y}|^2 = |\tilde{x}|^2 |\tilde{y}|^2, \quad (7a)$$

$$\text{Re}(\tilde{x} \tilde{y}) = \text{Re}(\tilde{y} \tilde{x}), \quad (7b)$$

$$\text{Re}(\tilde{x}) + \text{Re}(\tilde{y}) = \text{Re}(\tilde{x} + \tilde{y}), \quad (7c)$$

we obtain

$$\begin{aligned}\sum_{i < j} |\tilde{q}_i \tilde{q}_j^* - \tilde{p}_i \tilde{p}_j^*|^2 &= \frac{1}{2} \sum_{i, j} |\tilde{q}_i \tilde{q}_j^* - \tilde{p}_i \tilde{p}_j^*|^2 \\ &= \frac{1}{2} \sum_{i, j} (|\tilde{q}_i|^2 |\tilde{q}_j|^2 + |\tilde{p}_i|^2 |\tilde{p}_j|^2 - 2 \text{Re}(\tilde{q}_j \tilde{q}_i^* \tilde{p}_i \tilde{p}_j^*)) \quad (\text{using (7a)}) \\ &= \frac{1}{2} \left(\sum_i |\tilde{q}_i|^2 \sum_j |\tilde{q}_j|^2 + \sum_i |\tilde{p}_i|^2 \sum_j |\tilde{p}_j|^2 - 2 \sum_{i, j} \text{Re}(\tilde{q}_j \tilde{q}_i^* \tilde{p}_i \tilde{p}_j^*) \right) \\ &= 1 - \text{Re} \left(\sum_{i, j} \tilde{q}_i^* \tilde{p}_i \tilde{p}_j^* \tilde{q}_j \right) \quad (\text{using (7b) and (7c)}) \\ &= 1 - \text{Re} \left\langle \sum_i (\tilde{p}_i^* \tilde{q}_i), \sum_j (\tilde{p}_j^* \tilde{q}_j) \right\rangle \\ &= 1 - |\tilde{p}^* \tilde{q}|^2.\end{aligned}$$

Due to $\|\tilde{p}\|_2 = \|\tilde{q}\|_2 = 1$, we have $|\tilde{p}^* \tilde{q}| \leq 1$ and

$$\begin{aligned}1 - |\tilde{p}^* \tilde{q}|^2 &\leq 2(1 - |\tilde{p}^* \tilde{q}|) \\ &\leq 2(1 - \text{Re}(\tilde{p}^* \tilde{q})) \\ &= \|\tilde{p} - \tilde{q}\|_2^2.\end{aligned}$$

Therefore, we obtain $\mathbb{E}(X_{\tilde{q}} - X_{\tilde{p}})^2 \leq \|\tilde{p} - \tilde{q}\|_2^2$.

Step II: Define $Y_{\tilde{q}} = \langle g_0, p_0 \rangle + \langle g_1, p_1 \rangle + \langle g_2, p_2 \rangle + \langle g_3, p_3 \rangle$, where $g = [g_0, g_1, g_2, g_3]$ is the canonical Gaussian random vector in \mathbb{R}^{4n} , i.e., g_k are i.i.d. and $g_k \sim \mathcal{N}(0, I_n)$. Regarding both \tilde{q}, \tilde{p} as real vectors, then we have

$$\begin{aligned}\mathbb{E}_g (Y_{\tilde{q}} - Y_{\tilde{p}})^2 &= \mathbb{E}_g \langle g, \tilde{q} - \tilde{p} \rangle^2 \\ &= \mathbb{E}_g (\tilde{q} - \tilde{p})^\top \mathbb{E}(g g^\top) (\tilde{q} - \tilde{p}) \\ &= \|\tilde{p} - \tilde{q}\|_2^2.\end{aligned}$$

Step III: According to Step I and II, we get $\mathbb{E}(X_{\tilde{q}} - X_{\tilde{p}})^2 \leq \mathbb{E}_g (Y_{\tilde{q}} - Y_{\tilde{p}})^2$. Using Slepian's comparison theorem [30], we have

$$\mathbb{E}(\sigma_{\max}(\tilde{A})) = \mathbb{E} \sup_{\|\tilde{q}\|_2=1} X_{\tilde{q}} \leq \mathbb{E} \sup_{\|\tilde{q}\|_2=1} Y_{\tilde{q}} = \mathbb{E} \|g\|_2 \leq \sqrt{\mathbb{E} \|g\|_2^2} = 2\sqrt{n}.$$

Since $\sigma_{\max}(\cdot)$ is a real valued Lipschitz function with Lipschitz constant 1, the concentration in Gauss space [31] implies that

$$\mathbb{P}\{\sigma_{\max}(\tilde{A}) - \mathbb{E}(\sigma_{\max}(\tilde{A})) \geq t\} \leq e^{-t^2/2}.$$

Let $t = c_1\sqrt{n}$, we have $\mathbb{P}\{\|\tilde{A}\|_2 > (2 + c_1)\sqrt{n}\} \leq e^{-c_1^2 n/2}$. \square

Lemma B.2. *Let $\tilde{A} \in \mathbb{H}^{n \times n}$ be a quaternion Hermitian random matrix with i.i.d. off-diagonal entries following a standard quaternion normal distribution $\mathcal{QN}(0, 1)$ and zeros on the diagonal. Then we have*

$$\mathbb{P}\{\|\tilde{A}\mathbf{1}\|_\infty > c_2\sqrt{n \log n}\} \leq (1 + c_2^2)n^{3-2c_2^2}, \quad (8)$$

where $c_2 > 0$ is an arbitrary constant.

Proof. Since $(\tilde{A}\mathbf{1})_i = \sum_j \tilde{a}_{ij}$, and all \tilde{a}_{ij} , $i < j$ are independent, we have $(\tilde{A}\mathbf{1})_i$ also follows quaternion normal distribution, and

$$\mathbb{E}(\tilde{A}\mathbf{1})_i = 0, \quad \mathbb{E}|(\tilde{A}\mathbf{1})_i|^2 = \sum_{j=1}^n \mathbb{E}|\tilde{a}_{ij}|^2 = n - 1.$$

Although $(\tilde{A}\mathbf{1})_i$ are not independent with each other, we can find an upper bound based on subadditivity in probability theory, i.e.,

$$\mathbb{P}\{\|\tilde{A}\mathbf{1}\|_\infty > t\} = \mathbb{P}\left\{\bigcup_i |(\tilde{A}\mathbf{1})_i| > t\right\} \leq \sum_i \mathbb{P}\{|(\tilde{A}\mathbf{1})_i| > t\}. \quad (9)$$

For any quaternion random variable $\tilde{q} \sim \mathcal{QN}(0, 1)$, $Y = 4|\tilde{q}|^2 = (2\tilde{q}_0)^2 + (2\tilde{q}_1)^2 + (2\tilde{q}_2)^2 + (2\tilde{q}_3)^2$ follows Chi-square distribution, i.e., $Y \sim \chi^2(4)$. The probability density function of Y is $f_Y(y) = \frac{1}{4}ye^{-y/2}$. Let $R = \sqrt{Y}$, and then random variable R follows Rayleigh distribution whose probability density function is $f_R(r) = 2rf_Y(r^2) = \frac{r^3}{2}e^{-r^2/2}$. The tail of random variable R satisfies

$$\mathbb{P}\{2|\tilde{q}| > t\} = \int_t^{+\infty} f_R(r)dr = \left(1 + \frac{t^2}{2}\right)e^{-t^2/2}. \quad (10)$$

Combining (9) and (10), we have

$$\begin{aligned} \mathbb{P}\{\|\tilde{A}\mathbf{1}\|_\infty > t\} &\leq \sum_i \mathbb{P}\left\{\left|\frac{1}{\sqrt{n-1}}(\tilde{A}\mathbf{1})_i\right| > \frac{t}{\sqrt{n}}\right\} \\ &\leq \left(1 + \frac{2t^2}{n}\right)ne^{-2t^2/n}. \end{aligned}$$

Let $t = c_2\sqrt{n \log n}$, and we get

$$\mathbb{P}\{\|\tilde{A}\mathbf{1}\|_\infty > c_2\sqrt{n \log n}\} \leq (1 + 2c_2^2 \log n)n^{1-2c_2^2} \leq (1 + c_2^2)n^{3-2c_2^2},$$

where the last inequality holds for any $n \geq 2$. \square

B.2. Proof of Theorem 4.1

Proof. Here we prove that the quaternion random matrix \tilde{A} is \tilde{q} -discordant with high probability. Since

$$\begin{aligned} \|\text{diag}(\tilde{q}^*)\tilde{A}\text{diag}(\tilde{q})\|_2 &= \|\text{diag}(\tilde{q})\|_2^2 \|\tilde{A}\|_2 = \|\tilde{A}\|_2, \\ \|\text{diag}(\tilde{q}^*)\tilde{A}\text{diag}(\tilde{q})\tilde{\mathbf{1}}\|_\infty &= \|\text{diag}(\tilde{q}^*)\tilde{A}\tilde{q}\|_\infty = \|\tilde{A}\tilde{q}\|_\infty, \end{aligned}$$

we have that the quaternion matrix \tilde{A} is \tilde{q} -discordant if and only if $\text{diag}(\tilde{q}^*)\tilde{A}\text{diag}(\tilde{q})$ is $\tilde{\mathbf{1}}$ -discordant.

Furthermore, denote $\tilde{B} = \text{diag}(\tilde{q}^*)\tilde{A}\text{diag}(\tilde{q})$ and we have $\tilde{b}_{ij} = \tilde{q}_i^* \tilde{a}_{ij} \tilde{q}_j$. Since \tilde{a}_{ij} is a standard quaternion normal random variable which has i.i.d. real Gaussian entries, \tilde{b}_{ij} also follows a normal distribution. The expectation $\mathbb{E}(\tilde{b}_{ij}) = \tilde{q}_i^* \mathbb{E}(\tilde{a}_{ij}) \tilde{q}_j = 0$, and the variance $\text{Var}(\tilde{b}_{ij}) = \mathbb{E}(\|\tilde{b}_{ij}\|^2) = 1$. Therefore, \tilde{B} has the same distribution as \tilde{A} , and we may without loss of generality assume $\tilde{q} = \tilde{\mathbf{1}}$ in the proof.

Combining Lemmas B.1 and B.2, we have \tilde{A} is $\mathbf{1}$ -discordant with probability at least $1 - e^{-c_1^2 n/2} - (1 + c_2^2)n^{3-2c_2^2}$. The proof is completed. \square

B.3. Two lemmas

For any feasible point $\tilde{x} \in \mathbb{U}^n$ and $\tilde{z} \in \mathbb{U}$, since the function value $\tilde{x}^* \tilde{C} \tilde{x} = \tilde{z}^* \tilde{x}^* \tilde{C} \tilde{x} \tilde{z}$, the available data are insufficient to distinguish \tilde{x} from $\tilde{x} \tilde{z}$. We denote the equivalence class of \tilde{x} by

$$[\tilde{x}] = \{\tilde{x} \tilde{z} : \tilde{z} \in \mathbb{U}\}.$$

First, we prove that if the function value of \tilde{x} in (Q-QP) is greater than that of the exact rotation \tilde{q} , then the distance between $[\tilde{x}]$ and \tilde{q} under 2-norm can be bounded. In a more general setting, the bound also holds for any $\|\tilde{x}\|_2^2 = n$.

Lemma B.3. *If \tilde{W} is \tilde{q} -discordant with constant $c_1, c_2 > 0$, and $\tilde{x} \in \mathbb{Q}^n$ satisfies $\|\tilde{x}\|_2^2 = n$ and $\tilde{q}^* \tilde{C} \tilde{q} \leq \tilde{x}^* \tilde{C} \tilde{x}$, then we have*

$$\text{dist}(\tilde{q}, [\tilde{x}]) = \min_{\tilde{z} \in \mathbb{U}} \|\tilde{x} \tilde{z} - \tilde{q}\|_2 \leq 4(2 + c_1)\sigma. \quad (11)$$

Proof. First, we reformulated the distance $\text{dist}^2(\tilde{q}, [\tilde{x}])$ as

$$\min_{\tilde{z} \in \mathbb{U}} \|\tilde{x} \tilde{z} - \tilde{q}\|_2^2 = 2 \left(n - \min_{\tilde{z} \in \mathbb{U}} \text{Re}(\tilde{z}^* \tilde{x}^* \tilde{q}) \right) = 2(n - |\tilde{x}^* \tilde{q}|). \quad (12)$$

Combining $\tilde{q}^* \tilde{C} \tilde{q} \leq \tilde{x}^* \tilde{C} \tilde{x}$ with $C = \tilde{q} \tilde{q}^* + \sigma \tilde{W}$, we obtain

$$n^2 - |\tilde{x}^* \tilde{q}|^2 \leq \sigma \left(\tilde{x}^* \tilde{W} \tilde{x} - \tilde{q}^* \tilde{W} \tilde{q} \right).$$

Since $n + |\tilde{x}^* \tilde{q}| \geq n$, we have

$$n - |\tilde{x}^* \tilde{q}| \leq \sigma n^{-1} \left(\tilde{x}^* \tilde{W} \tilde{x} - \tilde{q}^* \tilde{W} \tilde{q} \right). \quad (13)$$

By using the definition of \tilde{q} -discordant and Theorem A.1 (vi), we have

$$\begin{aligned} \tilde{x}^* \tilde{W} \tilde{x} - \tilde{q}^* \tilde{W} \tilde{q} &= \text{Re} \left((\tilde{x} - \tilde{q})^* \tilde{W} (\tilde{x} + \tilde{q}) \right) \\ &\leq \|\tilde{x} - \tilde{q}\|_2 \|\tilde{W}\|_2 \|\tilde{x} + \tilde{q}\|_2 \\ &\leq 2(2 + c_1)n \|\tilde{x} - \tilde{q}\|_2, \end{aligned}$$

where the first equality follows from $\text{Re}(\tilde{x}^* \tilde{W} \tilde{q}) = \text{Re}(\tilde{q}^* \tilde{W} \tilde{x})$. Assume \tilde{x} and \tilde{q} are optimally aligned, i.e., $|\tilde{x}^* \tilde{q}| = \tilde{x}^* \tilde{q}$. Together with (12) and (13), we obtain $\|\tilde{x} - \tilde{q}\|_2^2 \leq 4(2 + c_1)\sigma \|\tilde{x} - \tilde{q}\|_2$, which implies $\|\tilde{x} - \tilde{q}\|_2 \leq 4(2 + c_1)\sigma$. \square

In the next lemma, we establish an upper bound on $\|\tilde{x} - \tilde{q}\|_\infty$ when they are optimally aligned, i.e., $|\tilde{x}^* \tilde{q}| = \tilde{x}^* \tilde{q}$.

Lemma B.4. *If \tilde{W} is \tilde{q} -discordant with constant $c_1, c_2 > 0$, and $\tilde{x} \in \mathbb{U}^n$ is a global optimizer of (Q-QP) which satisfies $|\tilde{x}^* \tilde{q}| = \tilde{x}^* \tilde{q}$, then*

$$\|\tilde{x} - \tilde{q}\|_\infty \leq 2\sigma n^{-1} \left(\|\tilde{W} \tilde{x}\|_\infty + 8(2 + c_1)^2 \sigma \right). \quad (14)$$

Proof. Denote $e_i \in \mathbb{R}^n$ be the i th vector of the canonical basis, whose i th entry is 1 whereas all other entries are 0. Let $\tilde{x} = \tilde{x} + (\tilde{q}_i - \tilde{x}_i)e_i$ whose i th entry is \tilde{q}_i and the others are the same as \tilde{x} . Then $\tilde{x} \in \mathbb{U}^n$ is also a feasible point of problem (Q-QP). Since \tilde{x} is optimal, we have

$$\begin{aligned} \tilde{x}^* \tilde{C} \tilde{x} &\geq \tilde{x}^* \tilde{C} \tilde{x} \\ &= \tilde{x}^* \tilde{C} \tilde{x} + (\tilde{q}_i - \tilde{x}_i)^* \tilde{C}_{ii} (\tilde{q}_i - \tilde{x}_i) + 2 \text{Re}((\tilde{q}_i - \tilde{x}_i)^* (\tilde{C} \tilde{x})_i). \end{aligned} \quad (15)$$

Since $\tilde{C}_{ii} = 1$ and $\tilde{C} = \tilde{q} \tilde{q}^* + \sigma \tilde{W}$, we have

$$\text{Re}((\tilde{q}_i - \tilde{x}_i)^* (\tilde{C} \tilde{x})_i) = \text{Re} \left((\tilde{q}_i - \tilde{x}_i)^* \left(|\tilde{q}^* \tilde{x}| \tilde{q}_i + \sigma (\tilde{W} \tilde{x})_i \right) \right) \leq -\|\tilde{q}_i - \tilde{x}_i\|^2 \leq 0.$$

Then we obtain

$$\begin{aligned} |\tilde{q}^* \tilde{x}| |\tilde{x}_i - \tilde{q}_i|^2 &= 2|\tilde{q}^* \tilde{x}| (1 - \text{Re}(\tilde{x}_i^* \tilde{q}_i)) \\ &= 2|\tilde{q}^* \tilde{x}| \text{Re}((\tilde{q}_i - \tilde{x}_i)^* \tilde{q}_i) \\ &\leq 2\sigma \text{Re} \left((\tilde{x}_i - \tilde{q}_i)^* (\tilde{W} \tilde{x})_i \right) \\ &\leq 2\sigma |\tilde{x}_i - \tilde{q}_i| |(\tilde{W} \tilde{x})_i|. \end{aligned} \quad (16)$$

This holds for all i , hence

$$|\tilde{q}^* \tilde{x}| \|\tilde{x} - \tilde{q}\|_\infty \leq 2\sigma \|\tilde{W} \tilde{x}\|_\infty.$$

According to Lemma B.3, we have $|\hat{q}^* \tilde{x}| \geq n - 8(2 + c_1)^2 \sigma^2$ which implies

$$n \|\tilde{x} - \tilde{q}\|_\infty \leq 2\sigma \|\tilde{W} \tilde{x}\|_\infty + 8(2 + c_1)^2 \sigma^2 \|\tilde{x} - \tilde{q}\|_\infty,$$

or equivalently, by using $\|\tilde{x} - \tilde{q}\|_\infty \leq 2$,

$$\|\tilde{x} - \tilde{q}\|_\infty \leq 2\sigma n^{-1} \left(\|\tilde{W} \tilde{x}\|_\infty + 8(2 + c_1)^2 \sigma \right).$$

□

B.4. Proof of Theorem 4.2

Proof. First, we establish an upper bound of $\|\tilde{W} \tilde{x}\|_\infty$. When \tilde{W} is \tilde{q} -discordant with constant $c_1, c_2 > 0$, we have

$$\begin{aligned} \|\tilde{W} \tilde{x}\|_\infty &\leq \|\tilde{W}(\tilde{x} - \tilde{q})\|_\infty + \|\tilde{W} \tilde{q}\|_\infty \\ &\leq \|\tilde{W}\|_2 \|\tilde{x} - \tilde{q}\|_2 + \|\tilde{W} \tilde{q}\|_\infty \\ &\leq 4(2 + c_1)^2 \sigma \sqrt{n} + c_2 \sqrt{n \log n}, \end{aligned} \tag{17}$$

where the last inequality follows from Definition 4.1. Combining Theorem 4.1, Lemma B.3 and B.4, we obtain

$$\begin{aligned} \|\tilde{x} - \tilde{q}\|_\infty &\leq 2\sigma n^{-1} \left(4(2 + c_1)^2 \sigma \sqrt{n} + c_2 \sqrt{n \log n} + 8(2 + c_1)^2 \sigma \right) \\ &\leq 2\sigma n^{-1/2} \left(4(2 + c_1)^2 \sigma + c_2 \sqrt{\log n} + 8(2 + c_1)^2 \sigma n^{-1/2} \right) \\ &\leq 2\sigma n^{-1/2} \left(4(\sqrt{2} + 1)(2 + c_1)^2 \sigma + c_2 \sqrt{\log n} \right) \end{aligned}$$

with probability at least $1 - e^{-c_1^2 n/2} - (1 + c_2^2) n^{3-2c_2^2}$, where the last inequality holds for any $n \geq 2$. Let $c_1 = 1, c_2 = 2$, and we obtain the conclusion. □

C. PROOFS FOR SECTION 4.2

C.1. Derivation of the dual problem (Q-SDP)

First, recall that (Q-SDP) is

$$\max_{\tilde{X} \in \mathbb{H}^{n \times n}} \text{tr}(\tilde{C} \tilde{X}) + \text{tr}(\tilde{X} \tilde{C}), \text{ s. t. } \text{diag}(\tilde{X}) = \mathbb{1}, \tilde{X} \succeq 0.$$

The constraint $\text{diag}(\tilde{X}) = \mathbb{1}$ can be reformulated as n linear equality constraints, i.e., $\langle A_i, \tilde{X} \rangle = 1$, where $(A_i)_{ii} = 1$ and the other entries are all zeros. Since the dual cone of $\mathbb{H}_+^{n \times n}$ is itself (Lemma A.3), the Lagrangian function is

$$\begin{aligned} \mathcal{L}(\tilde{X}, \tilde{Y}, \mathbf{y}) &= -\text{tr}(\tilde{C} \tilde{X}) - \text{tr}(\tilde{X} \tilde{C}) - \text{tr}(\tilde{Y} \tilde{X}) - \text{tr}(\tilde{X} \tilde{Y}) + \sum_{i=1}^n y_i \left(2 - \text{tr}(A_i \tilde{X}) - \text{tr}(\tilde{X} A_i) \right) \\ &= 2 \sum_{i=1}^n y_i - \text{tr} \left((\tilde{C} + \tilde{Y} + \sum_{i=1}^n y_i A_i) \tilde{X} \right) - \text{tr} \left(\tilde{X} (\tilde{C} + \tilde{Y} + \sum_{i=1}^n y_i A_i) \right). \end{aligned}$$

Then the dual problem $\max_{\mathbf{y} \in \mathbb{R}^n, \tilde{Y} \in \mathbb{H}_+^{n \times n}} \min_{\tilde{X} \in \mathbb{H}^{n \times n}} \mathcal{L}(\tilde{X}, \tilde{Y}, \mathbf{y})$ of (Q-SDP) is

$$\max_{\mathbf{y} \in \mathbb{R}^n, \tilde{Y} \in \mathbb{H}_+^{n \times n}} 2\mathbf{y}^\top \mathbb{1} \quad \text{s.t. } \tilde{Y} + \tilde{C} + \sum_{i=1}^n y_i A_i = 0.$$

Equivalently, since $\text{tr}(\tilde{Y} + \tilde{C}) = \mathbf{y}^\top \mathbb{1}$, the dual problem can be reformulated as

$$\min_{\tilde{Y} \in \mathbb{H}^{n \times n}} \text{tr}(\tilde{Y} + \tilde{C}) \quad \text{s.t. } \tilde{Y} + \tilde{C} \text{ is real diagonal and } \tilde{Y} \succeq 0. \tag{Q-SDP}$$

C.2. Proof of Theorem 4.3

Proof. The proof follows the KKT conditions for a convex optimization problem ([32] Chapter 5). Conditions (i)–(iv) are the primal and dual feasibility, respectively. Condition (v) is complementary slackness. In addition, since the identity matrix I is strictly feasible for the primal problem, the Slater condition holds. Therefore, the KKT conditions are necessary and sufficient for global optimality.

Next, we give an explanation about Condition (v). When $\tilde{X}, \tilde{Y} \in \mathbb{H}^{n \times n}$ are the global optimizers, the strong duality holds and the optimal duality gap is zero, i.e., $\text{tr}(\tilde{C}\tilde{X}) + \text{tr}(\tilde{X}\tilde{C}) = 2 \text{tr}(\tilde{Y} + \tilde{C})$. Note that $\tilde{Y} + \tilde{C}$ is real diagonal and $\text{diag}(\tilde{X}) = \mathbb{1}$, and we have

$$\begin{aligned} \text{tr}(\tilde{C}\tilde{X}) + \text{tr}(\tilde{X}\tilde{C}) &= \text{tr}\left((\tilde{Y} + \tilde{C})\tilde{X}\right) + \text{tr}\left(\tilde{X}(\tilde{Y} + \tilde{C})\right) \\ &= \text{tr}(\tilde{C}\tilde{X}) + \text{tr}(\tilde{X}\tilde{C}) + \text{tr}(\tilde{Y}\tilde{X}) + \text{tr}(\tilde{X}\tilde{Y}), \end{aligned}$$

which implies $\text{tr}(\tilde{Y}\tilde{X}) + \text{tr}(\tilde{X}\tilde{Y}) = 0$. From $\tilde{X}, \tilde{Y} \succeq 0$, we can decompose $\tilde{X} = \tilde{U}^* \tilde{U}$ and $\tilde{Y} = \tilde{V}^* \tilde{V}$. Then we obtain

$$\begin{aligned} 0 &= \text{tr}(\tilde{Y}\tilde{X}) + \text{tr}(\tilde{X}\tilde{Y}) = \text{tr}(\tilde{V}^* \tilde{V} \tilde{U}^* \tilde{U}) + \text{tr}(\tilde{U}^* \tilde{U} \tilde{V}^* \tilde{V}) \\ &= 2 \text{Re}(\text{tr}(\tilde{V}^* \tilde{V} \tilde{U}^* \tilde{U})) \\ &= 2 \text{Re}(\text{tr}(\tilde{V} \tilde{U}^* \tilde{U} \tilde{V}^*)) \\ &= 2 \|\tilde{U} \tilde{V}^*\|_F^2, \end{aligned}$$

where the third equality follows from Theorem A.1 (v). By using $\tilde{U} \tilde{V}^* = 0$, we derive $\tilde{X} \tilde{Y} = \tilde{Y} \tilde{X} = 0$ and vice versa.

The last statement follows from the results in [33]. Assume $\text{rank}(\tilde{Y}) = n - 1$, Theorems 9 and 10 in [33] implies that the primal solution \tilde{X} is unique. Since \tilde{X} is nonzero and $\tilde{X} \tilde{Y} = \tilde{Y} \tilde{X} = 0$, the rank of \tilde{X} must be one. \square

C.3. Riemannian submanifold and Dual certificate

In this subsection, we regard the set of unit quaternions \mathbb{U} as a 3-dimensional sphere, which is a submanifold embedded in \mathbb{Q} . Then the quaternion quadratic programming (Q-QP) is an optimization problem with Riemannian submanifold constraints $\tilde{x} \in \mathbb{U}^n$, whose global optimizer should satisfy the first-order necessary optimality conditions over manifolds. We will then utilize these optimality conditions to identify the optimal dual variable \tilde{Y} that satisfies Theorem 4.3, thereby further proving the tightness of the semidefinite relaxation.

The tangent space [34] of \mathbb{U} at each \tilde{y} and the projector are given respectively by

$$\begin{aligned} \mathcal{T}_{\tilde{y}} \mathbb{U} &= \{\tilde{a} \in \mathbb{Q} \mid \text{Re}(\tilde{y}^* \tilde{a}) = 0\}, \\ \text{Proj}_{\mathcal{T}_{\tilde{y}} \mathbb{U}} : \mathbb{Q} &\rightarrow \mathcal{T}_{\tilde{y}} \mathbb{U} \quad \text{such that} \quad \text{Proj}_{\mathcal{T}_{\tilde{y}} \mathbb{U}}(\tilde{a}) = \tilde{a} - \text{Re}(\tilde{y}^* \tilde{a}) \tilde{y}. \end{aligned}$$

Suppose $\mathcal{M}_1, \mathcal{M}_2$ are two manifolds. Then it holds that $\mathcal{T}_{(u,v)}(\mathcal{M}_1 \times \mathcal{M}_2) = \mathcal{T}_u \mathcal{M}_1 \times \mathcal{T}_v \mathcal{M}_2$ (Proposition 3.20 in [34]). Hence the projector over the tangent space $\mathcal{T}_{\tilde{x}} \mathbb{U}^n$ at each $\tilde{x} = [\tilde{x}_1, \dots, \tilde{x}_n]$ is given by

$$\text{Proj}_{\mathcal{T}_{\tilde{x}} \mathbb{U}^n}(\tilde{v}) = \tilde{v} - \text{Re}(\text{ddiag}(\tilde{v} \tilde{x}^*)) \tilde{x},$$

where $\text{ddiag} : \mathbb{Q}^{n \times n} \rightarrow \mathbb{Q}^{n \times n}$ sets all off-diagonal entries of a matrix to zero. Reformulating the problem (Q-QP) as $\min_{\tilde{x} \in \mathbb{U}^n} f(\tilde{x}) := -\tilde{x}^* \tilde{C} \tilde{x}$, the gradient of f is given by $\nabla f(\tilde{x}) = -2 \tilde{C} \tilde{x}$. Therefore, the Riemannian gradient of f is

$$\begin{aligned} \text{grad } f(\tilde{x}) &= \text{Proj}_{\mathcal{T}_{\tilde{x}} \mathbb{U}^n}(\nabla f(\tilde{x})) \\ &= 2 \left(\text{Re}(\text{ddiag}(\tilde{C} \tilde{x} \tilde{x}^*)) - \tilde{C} \right) \tilde{x}. \end{aligned} \tag{18}$$

When \tilde{x} constitutes a global optimizer of (Q-QP), the first-order necessary optimality condition is satisfied: $\text{grad } f(\tilde{x}) = 0$. Furthermore, the second-order necessary optimality condition requires the Riemannian Hessian of f at \tilde{x} to be positive semidefinite over the tangent space $\mathcal{T}_{\tilde{x}} \mathbb{U}^n$, i.e.,

$$\langle \tilde{v}, \text{Hess } f(\tilde{x})[\tilde{v}] \rangle = 2 \left\langle \tilde{v}, \left(\text{Re}(\text{ddiag}(\tilde{C} \tilde{x} \tilde{x}^*)) - \tilde{C} \right) \tilde{v} \right\rangle \geq 0, \quad \forall \tilde{v} \in \mathcal{T}_{\tilde{x}} \mathbb{U}^n. \tag{19}$$

This observation motivates the variables

$$\tilde{X} = \tilde{x} \tilde{x}^*, \quad \text{and} \quad \tilde{Y} = \text{Re}(\text{ddiag}(\tilde{C} \tilde{x} \tilde{x}^*)) - \tilde{C}. \tag{20}$$

It is not hard to verify that $\text{diag}(\tilde{X}) = \mathbb{1}$, $\tilde{X} \succeq 0$ and $\tilde{Y} + \tilde{C}$ is real diagonal. In addition, (18) shows that $\tilde{Y} \tilde{X} = 0$. Hence, the conditions (i)–(iii) and (v) in Theorem 4.3 are all satisfied. In addition, (18) and (19) also indicate $\tilde{Y} \succeq 0$ (condition (iv)). Therefore, when \tilde{x} constitutes a global optimizer of (Q-QP), \tilde{X} and \tilde{Y} defined in (20) are the global optimizers of (Q-SDP) and (Q-DSDP), respectively. The following lemma demonstrates that \tilde{Y} must necessarily admit the form specified in (20), with no alternative representations possible for $\tilde{X} = \tilde{x} \tilde{x}^*$.

Lemma C.1. *A feasible \tilde{X} with any rank is optimal for (Q-SDP) if and only if $\tilde{Y} = \text{Re}(\text{ddiag}(\tilde{C} \tilde{X})) - \tilde{C}$ is positive semidefinite.*

Proof. We first prove the necessary condition. If \tilde{X} is optimal, there exists $\tilde{Y}_0 \succeq 0$ satisfying $\tilde{Y}_0 \tilde{X} = 0$ and $\tilde{Y}_0 + \tilde{C} = \tilde{D}$ where \tilde{D} is diagonal. Then $(\tilde{D} - \tilde{C})\tilde{X} = \tilde{Y}_0 \tilde{X} = 0$, which implies $\tilde{D} = \text{ddiag}(\tilde{C}\tilde{X})$. Thus, $\tilde{Y}_0 = \text{ddiag}(\tilde{C}\tilde{X}) - \tilde{C}$.

For sufficiency, the conditions (i)–(iv) in Theorem 4.3 are immediately satisfied. Since $\tilde{x}_{ii} = 1$, it follows that

$$\text{tr}(\tilde{Y}\tilde{X}) = 0, \quad \text{and} \quad \text{tr}(\tilde{X}\tilde{Y}) = \text{tr}(\text{Re}(\text{ddiag}(\tilde{C}\tilde{X})) - \tilde{X}\tilde{C}) = 0.$$

Then (v) in Theorem 4.3 holds due to $\tilde{X}, \tilde{Y} \succeq 0$. □

C.4. Proof of Theorem 4.4

Proof. According to the first-order necessary optimality conditions $\text{grad } f(\tilde{x}) = 0$, i.e., $\tilde{Y}\tilde{x} = 0$, and \tilde{x} is nonzero, we have $\text{rank}(\tilde{Y}) \leq n-1$. If for arbitrary $\tilde{v} \in \mathbb{Q}^n$ such that $\tilde{v} \perp \tilde{x}$ and $\tilde{v} \neq 0$, it holds that $\tilde{v}^* \tilde{Y} \tilde{v} > 0$, then the proof is completed. We divide the proof into three steps.

Step I: (The properties of $\text{ddiag}(\tilde{C}\tilde{x}\tilde{x}^*)$) Since $\tilde{Y}\tilde{x} = 0$, we have $\text{Re}((\tilde{C}\tilde{x})_i \tilde{x}_i^*) \tilde{x}_i = (\tilde{C}\tilde{x})_i$. Multiplying both sides by \tilde{x}_i^* , we conclude $(\tilde{C}\tilde{x})_i \tilde{x}_i^*$ is real. In addition, the second-order necessary optimality condition $\langle \tilde{v}, \tilde{Y}\tilde{v} \rangle \geq 0$ holds for all $\tilde{v} \in \mathcal{T}_{\tilde{x}}\mathbb{U}^n$. In particular, we choose $\tilde{v} = [0, \dots, 0, \tilde{x}_i \mathbf{k}, 0, \dots, 0] \in \mathcal{T}_{\tilde{x}}\mathbb{U}^n$ in which the i -th entry is non-zero and \mathbf{k} is an imaginary units. Then we have

$$\langle \tilde{v}, \tilde{Y}\tilde{v} \rangle = (\tilde{C}\tilde{x})_i \tilde{x}_i^* - 1 \geq 0,$$

where the equality follows from \tilde{Y}_{ii} is real and $\tilde{v}^* \tilde{v} = 1$. Thus,

$$(\tilde{C}\tilde{x})_i \tilde{x}_i^* = |(\tilde{C}\tilde{x})_i \tilde{x}_i^*| = |(\tilde{C}\tilde{x})_i| \geq 1. \quad (21)$$

Step II: (The lower bound of $\tilde{v}^* \tilde{Y} \tilde{v}$) Without loss of generality, we assume $\tilde{q}^* \tilde{x} = |\tilde{q}^* \tilde{x}|$. For any $\tilde{v} \in \mathcal{T}_{\tilde{x}}\mathbb{U}^n$, we have

$$\begin{aligned} \tilde{v}^* \tilde{Y} \tilde{v} &= \tilde{v}^* \text{ddiag}(\tilde{C}\tilde{x}\tilde{x}^*) \tilde{v} - \tilde{v}^* \tilde{C} \tilde{v} \\ &= \sum_{i=1}^n |\tilde{v}_i|^2 |(\tilde{C}\tilde{x})_i| - \tilde{v}^* \tilde{C} \tilde{v} \\ &= \sum_{i=1}^n |\tilde{v}_i|^2 |(\tilde{q}\tilde{q}^* \tilde{x})_i + \sigma(\tilde{W}\tilde{x})_i| - |\tilde{v}^* \tilde{q}|^2 - \sigma \tilde{v}^* \tilde{W} \tilde{v} \\ &\geq \sum_{i=1}^n |\tilde{v}_i|^2 |\tilde{q}_i(\tilde{q}^* \tilde{x}) + \sigma(\tilde{W}\tilde{x})_i| - |\tilde{v}^* (\tilde{q} - \tilde{x}) + \tilde{v}^* \tilde{x}|^2 - \sigma \|\tilde{W}\|_2 \|\tilde{v}\|_2^2 \\ &\geq |\tilde{q}^* \tilde{x}| \|\tilde{v}\|_2^2 - \sigma \sum_{i=1}^n |\tilde{v}_i|^2 |(\tilde{W}\tilde{x})_i| - |\tilde{v}^* (\tilde{q} - \tilde{x})|^2 - \sigma \|\tilde{W}\|_2 \|\tilde{v}\|_2^2 \\ &\geq \|\tilde{v}\|_2^2 \left(|\tilde{q}^* \tilde{x}| - \|\tilde{q} - \tilde{x}\|_2^2 - \sigma \|\tilde{W}\|_2 - \sigma \frac{1}{\|\tilde{v}\|_2^2} \sum_{i=1}^n |\tilde{v}_i|^2 |(\tilde{W}\tilde{x})_i| \right) \\ &\geq \|\tilde{v}\|_2^2 \left(|\tilde{q}^* \tilde{x}| - \|\tilde{q} - \tilde{x}\|_2^2 - \sigma \|\tilde{W}\|_2 - \sigma \|\tilde{W}\tilde{x}\|_\infty \right), \end{aligned}$$

where the second equality follows from (21), and the last inequality follows from applying the weighted average. Furthermore, by (13) and Theorem 4.2, we obtain

$$|\tilde{q}^* \tilde{x}| - \|\tilde{q} - \tilde{x}\|_2^2 - \sigma \|\tilde{W}\|_2 - \sigma \|\tilde{W}\tilde{x}\|_\infty \geq n - 216\sigma^2 - 3\sigma\sqrt{n} - \sigma \|\tilde{W}\tilde{x}\|_\infty. \quad (22)$$

Step III: (The upper bound of $\|\tilde{W}\tilde{x}\|_\infty$) By (17) and let $c_1 = 1$, $c_2 = 2$, we have

$$\|\tilde{W}\tilde{x}\|_\infty \leq 36\sigma\sqrt{n} + 2\sqrt{n \log n}.$$

Then for any $\tilde{v} \in \mathcal{T}_{\tilde{x}}\mathbb{U}^n$,

$$\tilde{v}^* \tilde{Y} \tilde{v} \geq \sqrt{n} \|\tilde{v}\|_2^2 f(n), \quad (23)$$

where $f(n) := \sqrt{n} - 216\sigma^2 n^{-1/2} - 3\sigma - 36\sigma^2 - 2\sigma\sqrt{\log n}$. When $\sigma \leq \frac{1}{16} n^{1/4}$ and $n \geq 2$, we obtain $f(n) > 0$, which implies the conclusion. □

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