| 1. Singl | le Choice | |
|---|---|--|
| 1) Whic | ch one of the following is NOT TRUE for database language? () | |
| A. | Relational Algebra is a procedural database language. | |
| B. | SQL is a declarative database language. | |
| C. SQL is both a declarative database language and a pure database language | | |
| D. | SQL is both data definition language and data manipulation language. | |
| 2) Whic | ch one of the following is NOT TRUE for relational model? () | |
| A. | In a relation, a foreign key can be a subset of the primary key. | |
| B. In a relation, the primary key can be a subset of a candidate key. | | |
| C. | In a relation, a super key must contain a candidate key. | |
| D. | In a relation, a candidate key can include multiple attributes. | |
| 3) Give | n relation schema R ₁ (A,B,C,D) and R ₂ (B,C), the schema of the relation as the | |
| res | ult of $R_1 \div R_2$ is (). | |
| A. | $(A, R_1.B, R_1.C, R_2.B, R_2.C)$ | |
| | (A) | |
| | (A, B, C) | |
| | (A, D) | |
| | ch one of the following is NOT TRUE for SQL? () | |
| | Views may be defined in terms of other views. | |
| | Result of where clause predicate is treated as false if it evaluates to unknown. | |
| | Set operations automatically eliminate duplicates. | |
| | 'DELETE TABLE r' deletes not only all tuples of r, but also the schema for r. | |
| 5) "Uni | que ((mike, Null), (mike, Null))" is evaluated to () | |
| A. | TRUE | |
| | FALSE | |
| C. | UNKNOWN. | |
| D. | NULL | |
| 6) In "C | CREATE TABLE" statement, unique $(\underline{A_1, A_2,, A_m})$ states that the attributes | |
| | $A_2, \ldots A_m$ form a (). | |
| | Primary key | |
| | Foreign key | |
| | Candidate key | |
| | Relation schema | |
| , | is a one-to-many relationship set from entity set E_1 to E_2 , Which one of the wing is TRUE? (| |
| | If R has any descriptive attribute it can be moved to E_1 . | |
| | • | |
| В. | The primary key of R is the primary key of E_2 . | |
| C. | | |
| | E ₂ must totally participate in R e following statements about weak entity set, () is incorrect. | |
| | Weak entity set is an entity set that does not have a super key | |
| | | |
| В. | The existence of a weak entity set depends on the existence of an identifying | |

| | | entity set |
|---|--------|---|
| | C. | Weak entity set is an entity set that does not have a foreign key |
| D. Weak entity set must relate to the identifying entity set via a total, m | | |
| | | one relationship set |
| 9)] | If and | d only if (), K is a super key of R. |
| | | K→R |
| | B. | R→K |
| | C. | $K \rightarrow (K-R)$ |
| | D. | (R-K)→K |
| 10) | A de | ecomposition of R into R1 and R2 is lossless join if (). |
| | A. | $R_1 \cap R_2 \rightarrow R_1$ is in F^+ |
| | B. | $R_1 \cap R_2 = \Phi$ |
| | C. | $R_1 \cap R_2 \neq \Phi$ |
| | D. | $R_1-R_2 \rightarrow R_1$ |
| | | |
| | | the following blanks |
| 1) 1 | | base systems provide an abstract view of the data, which is achieved through 3 |
| | | l of abstraction: physical level, logical level, and (). |
| | | s are two relations. Suppose a tuple occurs 3 times in r and 5 times in s, then |
| 1 | it oc | curs () times in the execution result of the SQL statement "r |
| | | sect all s". |
| | | ion r has 100 tuples, among these tuples, only 2 have null values on attribute |
| 1 | A, th | ne result of the SQL statement "select count(*), count(A) from r" is |
| . (| | |
| | | grant statement is used to confer authorization, and the () |
| | | ment is used to reclaim authorization. |
| 5) | Colle | ection of operations that form a single logical unit of work in database system |
| | is ca | lled (). |
| | | |
| | | ver the following questions |
| 1) | Brief | ly describe what is referencing constraint. |
| 2) | Brief | ly describe the following concepts about keys in a relational model: Super |
| | | Candidate Key, Primary Key. |
| - | ксу, | Candidate Key, 11mary Key. |
| 3) | In a 1 | bank database, if a loan can be borrowed by more than one customer, should |
| | | ivide relation schema (customer id, loan number, amount), into two schemas |
| | | omer id, loan number) and (loan number, amount)? Describe why. |
| | • | |
| 4) | Brief | ly describe the ACID properties of transactions. |

4、Relational Algebra

Give the result of the following relational algebra expressions.

| R | | | | | |
|----|---|---|--|--|--|
| A | В | C | | | |
| a1 | 6 | 7 | | | |
| a2 | 2 | 3 | | | |
| a1 | 2 | 3 | | | |
| a4 | 4 | 5 | | | |
| a2 | 6 | 7 | | | |
| a3 | 7 | 9 | | | |

| | 3 |
|---|---|
| В | C |
| 6 | 7 |
| 2 | 3 |

- 1) $R1 = \Pi_{B, C, 200} (R)$
- 2) $R2 = \sigma_{B>2} (R)$
- 3) $R_3 = R \bowtie S$
- 4) $R_4 = {}_A\mathbf{Q}\operatorname{sum}(B)$

5. Compose SQL

Consider the following relations (the primary keys are underlined):

Students(<u>snum:integer</u>, <u>sname:string</u>, <u>major:string</u>, <u>level:string</u>, <u>age:integer</u>)

Faculty(fid:integer, fname:string, deptid:integer)

Class(cname:string, meets at:string, room:string, fid:integer)

Enrolled(snum:integer, cname:string)

The meaning of these relations is straightforward; for example, *Enrolled* has one record per student-class pair such that the student is enrolled in the class. Note that *snum* and *cname* in *Enrolled* should correspond with *snum* in *Students* and *cname* in *Class* respectively.

- 1) Write a SQL statement to create relation *Class*. Declare a primary key and foreign keys (if any) on this relation.
- 2) Write a SQL statement to insert into the database the fact that the 22 year-old senior CS student 'Kobe Bryant', with snum 111, is enrolled in class CS411. (hint: Both Students and Enrolled tables need to be updated.)
- 3) Write a SQL statement to delete all the classes taught by "Joe Smith".
- 4) Create a view BusyFaculty that records the ids and names of faculties who teach more than 3 classes.
- 5) Find the names of all students that enroll in a class where students meet in room R128 (i.e., Class.room = R128) or a class in which five or more than five students enroll.
- 6) Find the names of all students who are enrolled in two classes that meet at the same time

6, E/R Diagram

The club *Travel-Often-And-A-Lot* organizes shorter and longer tours for its members. Help them to make a model of their mini world.

Travel-Often-And-A-Lot has members. Each member is represented by her/his full name, address, and birth date.

Some members belong to the board of *Travel-Often-And-A-Lot*. Some members are organizers (of tours). Organizers must be stored with their cell phone number so that they can be reached anytime. Organizers organize tours. Sometimes a tour is organized by several organizers.

Each tour is denoted by a name, e.g. "Museums of Paris, 2004" or "Iceland, 2005". Tours can take place multiple times. "Museums of Paris, 2004", for instances, takes place twice: May 22_{nd} to May 29_{th}, 2004 and June 5_{th} to June 12_{th}, 2004. The cost of a tour depends on the date, e.g. "Museums of Paris, 2004" was cheaper in May than in June. Each *travel* – such as "Museums of Paris, 2004" at June 5_{th} to June 12_{th}, 2004 – is lead by one organizer members participate in travels.

Travel-Often-And-A-Lot wants to keep track of the payments made by its members. A payment can e.g. be the annual club fee, a donation, *etc.* but also the payment for a travel. Mind the subtle distinction between *tour* and travel.

- 1) Create an E-R model that fulfill above requirements.
- 2) Translate the E-R model into relation schemas.

7. Normalization and Schema Design

Consider a relation with schema $R=\{A, B, C, D, E, F\}$ and $F=\{AB \rightarrow CD; A \rightarrow D; D \rightarrow AE; E \rightarrow F\}$ holds on R.

- 1) Give all candidate keys of this relation, motivate. (3 points)
- 2) Indicate all extraneous attributes in F, motivate. (3 points)
- 3) Is this relation in 3NF? If it is not, decompose it into relations in 3NF. (6 points)