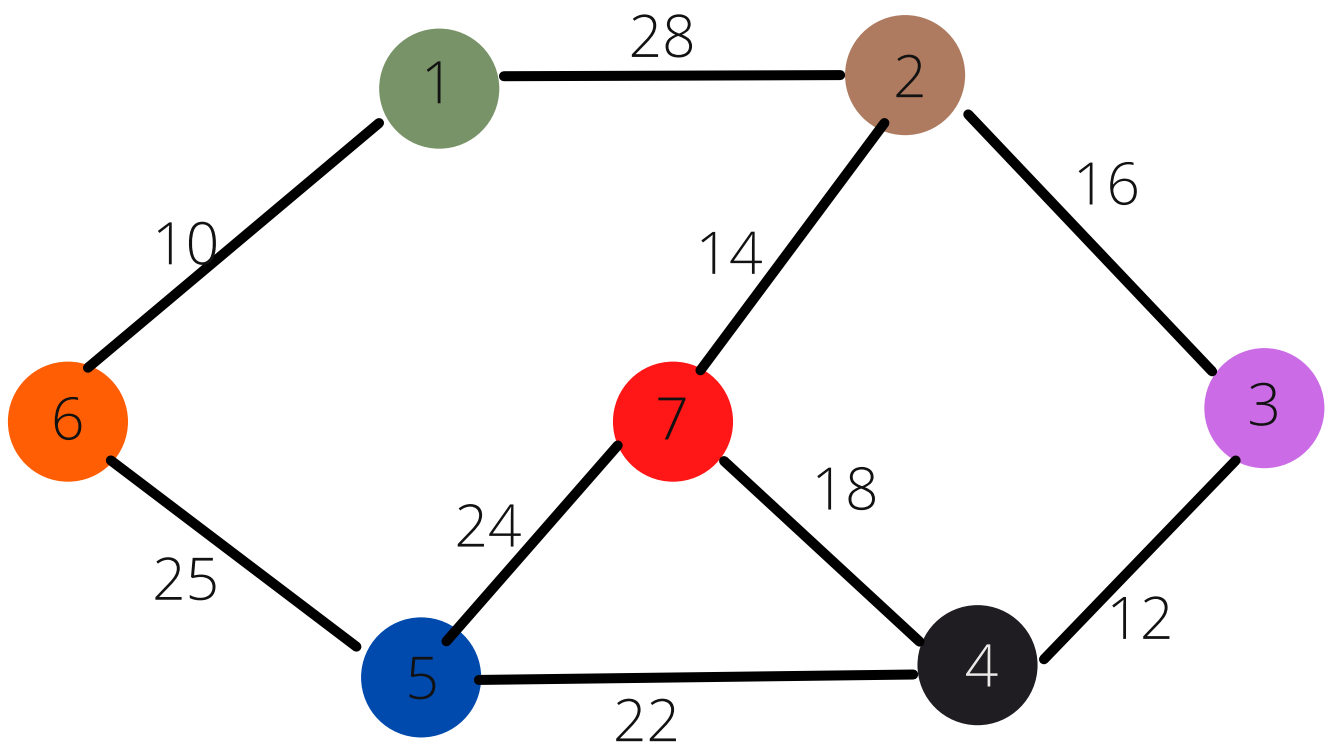


```
$(function){cards();});
$(window).on('resize', function(){cards();});
function cards(){
  var width = $(window).width();
  if(width < 750){
    cardssmallscreen();
  }else{
    cardsbigscreen();
  }
  cardssmallscreen(){
    = $('<div>.card').length;
  }
  cards++;
  length-of-cards;
  cards;
}
```

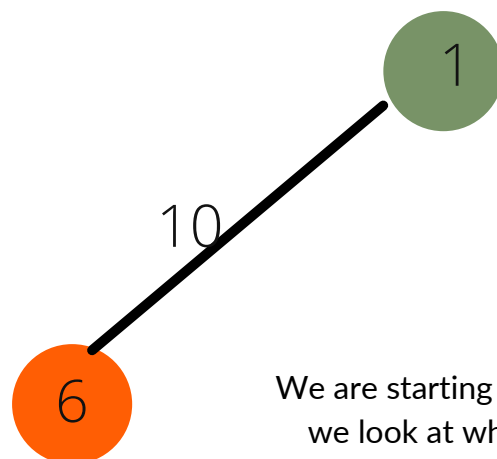
HANDLING MINIMUM SPANNING TREES WITH PRIM'S ALGORITHMS

BY HEATHER MATARUSE

PROCESS



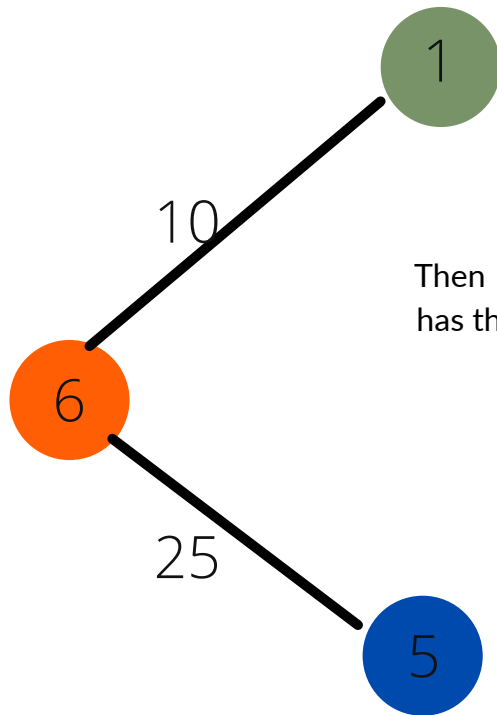
STEP 1



We are starting from 6 -> 1: 10 because when we look at which side has the least weight between 5 and 1 its the 1 which has a weight of 10

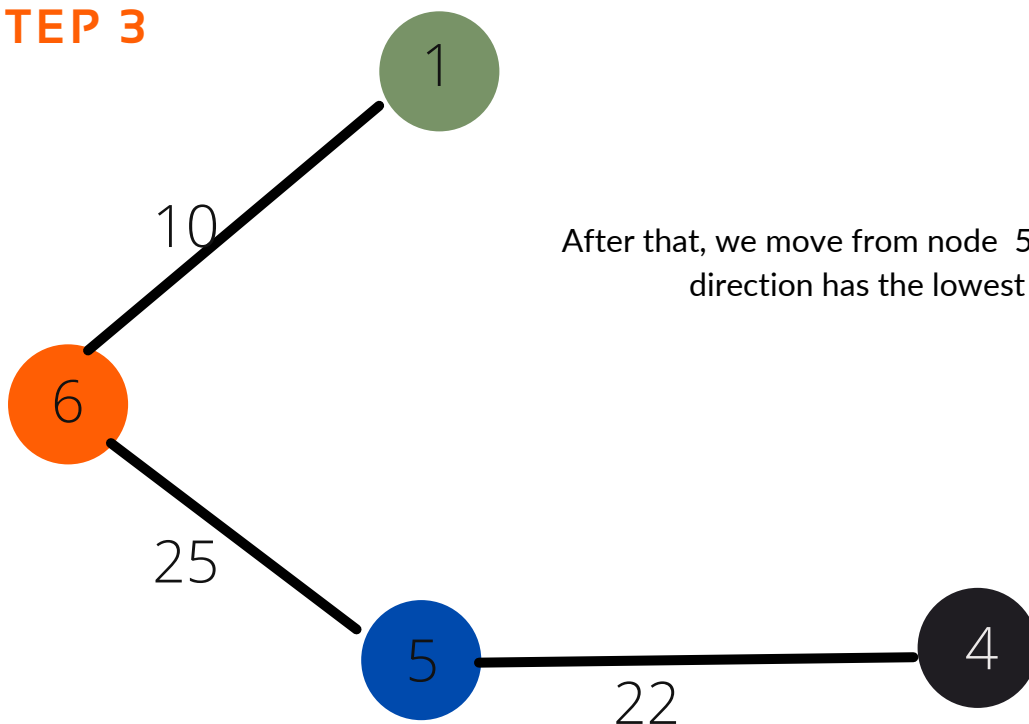
PROCESS

STEP 2



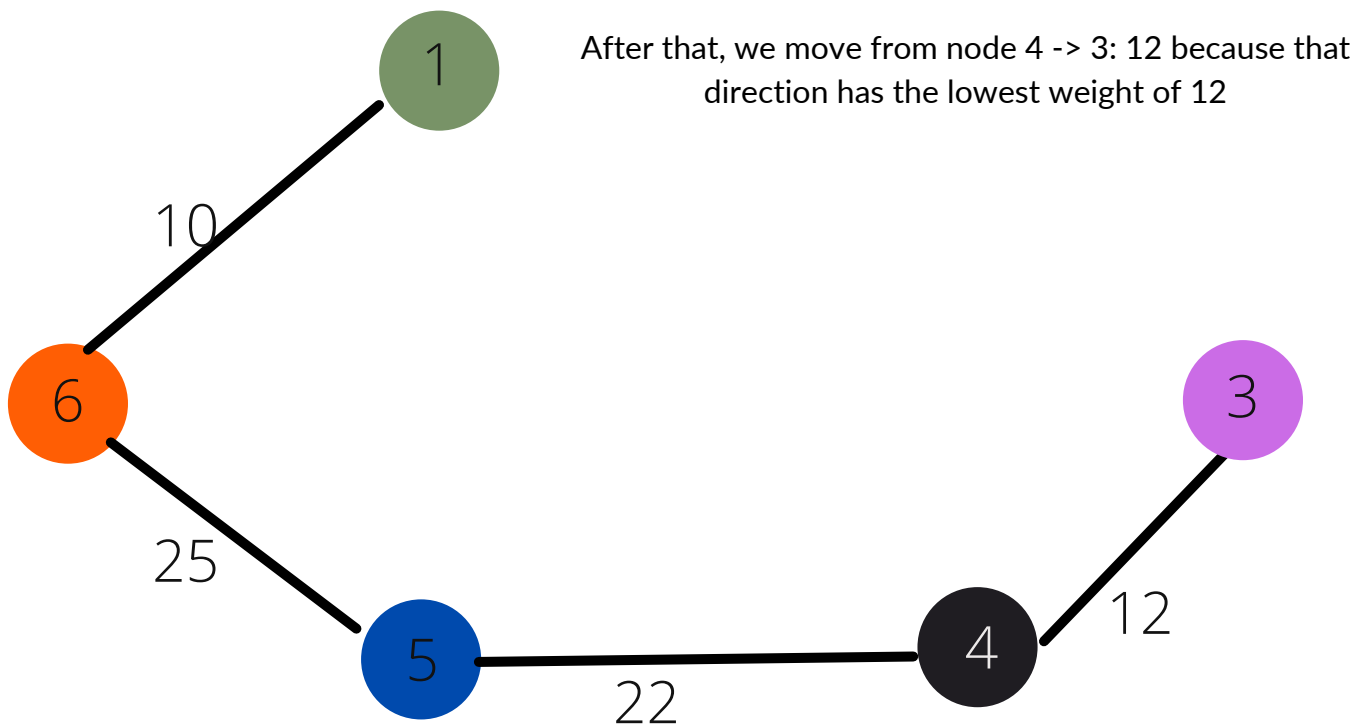
Then from there, we take route 6 -> 5: 25 because node 5 has the weight of 25 and again it is connected to 6 which I have chosen to be the root

STEP 3

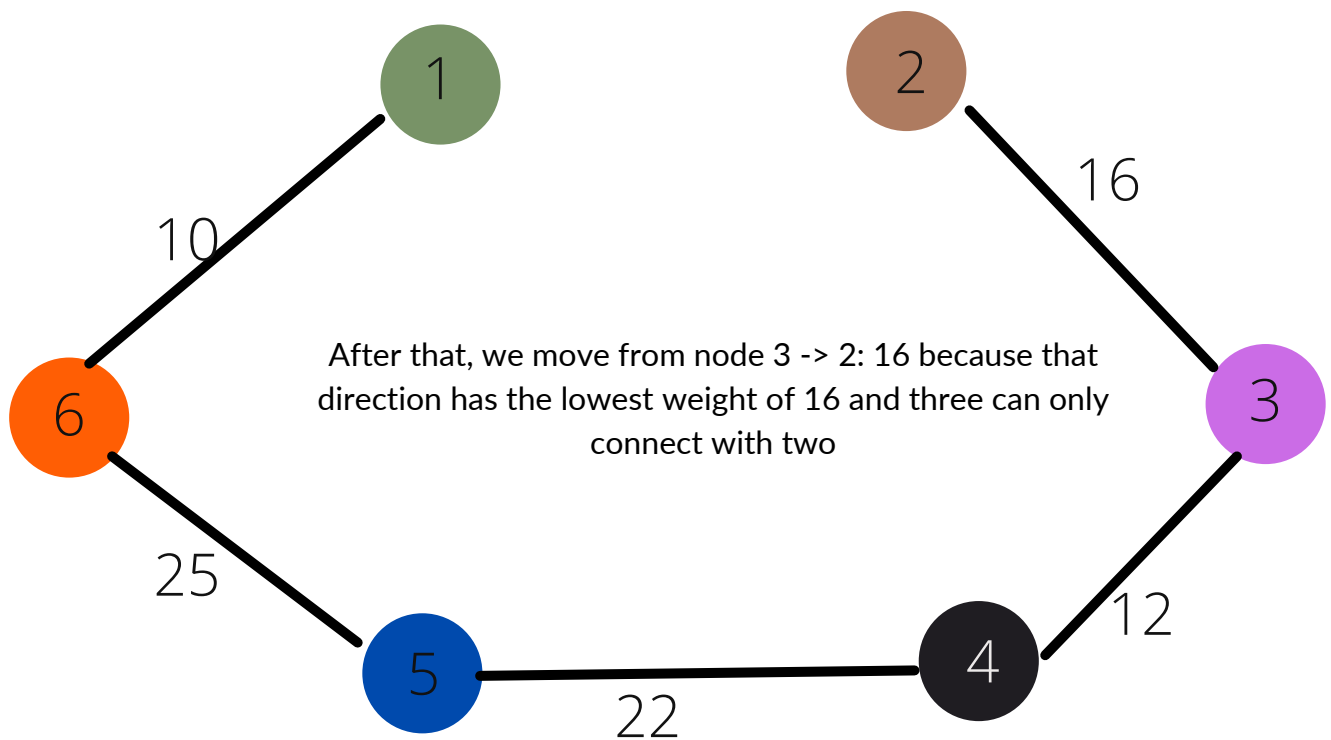


After that, we move from node 5 -> 4: 22 because that direction has the lowest weight of 22

STEP 4



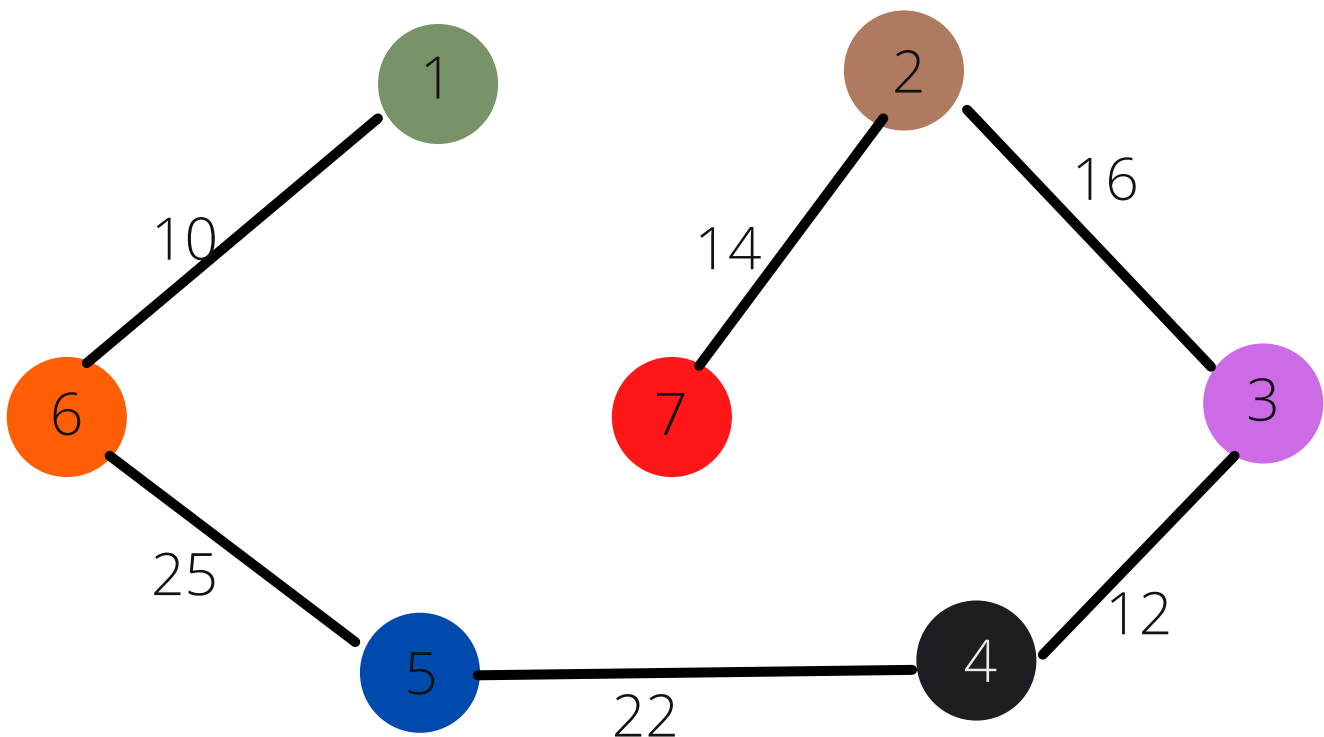
STEP 5



EXPLANATION

STEP 6

After that, we move from node 2 -> 7: 14, because that direction has the lowest weight of 14



Pseudocode

- You build vertex U, which contains the list of visited vertices; The list of vertices that have not been visited is contained in the created vertex V-U.
- By connecting the least weight edge, move vertices from vertex V-U to vertex U one by one.
- We utilize a min-heap to store the vertices not yet included in the MST after traversing all the vertices in the graph using breadth-first search.
- We utilize min-heap as a priority queue to get the minimum weight edge.
- Operations on the min-heap, such as extracting the smallest element and reducing the key value, take $O(\log V)$ time.
- Therefore Cost of Minimum Spanning Tree = Sum of all edge weights = $10 + 25 + 22 + 12 + 16 + 14 = 99$ units

Time Complexity

- The time complexity is $O(V \log V + E \log V) = O(E \log V)$, this is because there is use of a binary heap

Weakness of the Algorithm

- As a new edge is introduced, the list of edges must be searched from the beginning.
- If more than one edge has the same weight, all feasible spanning trees must be identified before the final minimal tree can be obtained.

Code output

```
EDGE : WEIGHT
Node 1 - Node 6 : Weight 10
Node 6 - Node 5 : Weight 25
Node 5 - Node 4 : Weight 22
Node 4 - Node 3 : Weight 12
Node 3 - Node 2 : Weight 16
Node 2 - Node 7 : Weight 14

...Program finished with exit code 0
Press ENTER to exit console.
```

Code in python

```
infinity_variable = 9999999
```

```
no_of_vertices = 7
```

```
matrix = [[0, 28, 0, 0, 0, 10, 0],  
          [28, 0, 16, 0, 0, 0, 14],  
          [0, 16, 0, 12, 0, 0, 0 ],  
          [0, 0, 12, 0, 22, 0, 18],  
          [0, 0, 0, 22, 0, 25, 24],  
          [10, 0, 0, 0, 25, 0, 0],  
          [0, 14, 0, 18, 24, 0, 0]]
```

```
visited_vertices = [0, 0, 0, 0, 0, 0, 0]
```

```
number_of_edges = 0
```

```
visited_vertices[0] = True
```

```
print("EDGE : WEIGHT\n")
```

```
while (number_of_edges < no_of_vertices - 1):
```

```
    mini = infinity_variable
```

```
    a = 0
```

```
    b = 0
```

```
    for k in range(no_of_vertices):
```

```
        if visited_vertices[k]:
```

```
            for i in range(no_of_vertices):
```

```
                if ((not visited_vertices[i]) and matrix[k][i]):
```

```
                    if mini > matrix[k][i]:
```

```
                        mini = matrix[k][i]
```

```
                        a = k
```

```
                        b = i
```

```
    print("Node " + str(a+1) + " -" + " Node " + str(b+1) + " : Weight " + str(matrix[a][b]))
```

```
    visited_vertices[b] = True
```

```
    number_of_edges += 1
```