

1. Construct an equivalent noncontracting grammar G_L with a nonrecursive start symbol. Give a regular expression.

$$G: S \rightarrow aS | bS | B \\ B \rightarrow bb | C | \lambda \\ C \rightarrow cC | \lambda$$

$$\rightarrow \begin{aligned} S' &\rightarrow S \\ S &\rightarrow aS | bS | B \\ B &\rightarrow bb | C | \lambda \\ C &\rightarrow cC | \lambda \end{aligned}$$

Iteration	NULL	PREV
0	$\{B\}$	
1	$\{B, C\}$	$\{B\}$
2	$\{S, B, C\}$	$\{B, C\}$
3	$\{S', S, B, C\}$	$\{S, B, C\}$
4	$\{S', S, B, C\}$	$\{S', S, B, C\}$

$$S' \rightarrow S | \lambda$$

$$S \rightarrow aS | bS | B | a | b$$

$$B \rightarrow bb | C$$

$$C \rightarrow cC | c$$

$$(a \cup b)^* (bb \cup \lambda) c^*$$

$(bb \cup \lambda)$ covers one application of bb , or null, followed by 0 or more applications of C .

7. Construct an equivalent grammar G_c that does not contain chain rules. Give a regular expression.

$G: S \rightarrow AS A$	$chain(S) = \{S\}$	New	Prev	$Chain(S)$
$A \rightarrow aA bB C$	$chain(A) = \{A\}$	$\{S\}$	\emptyset	$\{S, A\}$ chain S, A
$B \rightarrow bB b$	$chain(B) = \{B\}$	$\{A\}$	$\{S, A\}$	$\{S, A, B, C\}$ chain B, C
$C \rightarrow cC B$	$chain(C) = \{C\}$	$\{B\}$	$\{S, A, B, C\}$	$\{S, A, B, C\}$ no chain, B
		$\{C\}$	$\{S, A, B, C\}$	$\{S, A, B, C\}$ chain C, B

$$G_c: S \rightarrow AS | aA | bB | cC | bB | b \\ A \rightarrow aA | bB | b | cC | bB \\ B \rightarrow bB | b \\ C \rightarrow cC | bB | b$$

14. Construct an equivalent grammar without using useless symbols. Trace the generation of sets TERM and REACH used to construct G_T and G_U

$$G: \begin{aligned} S &\rightarrow AA | CD | bB \\ A &\rightarrow aA | a \\ B &\rightarrow bB | bC \\ C &\rightarrow cB \\ D &\rightarrow dD | d \end{aligned}$$

Term	Prev
$\{D, A\}$	\emptyset
$\{S, D, A\}$	$\{D, A\}$
	$\{S, D, A\}$

$$\begin{aligned} S &\rightarrow AA \\ A &\rightarrow aA | a \\ D &\rightarrow dD | d \end{aligned}$$

Reach	Prev	New
$\{S\}$	\emptyset	$\{S\}$
$\{S, A\}$	$\{S\}$	$\{S, A\}$
$\{S, A\}$	$\{S, A\}$	$\{A\}$

$$\begin{aligned} S &\rightarrow AA \\ A &\rightarrow aA | a \end{aligned}$$

$$23. \begin{aligned} S &\rightarrow A | ABa | AbA \\ A &\rightarrow Aa | \lambda \\ B &\rightarrow Bb | BC \\ C &\rightarrow CB | CA | bB \end{aligned}$$

Remove λ rules

NULL	PREV
$\{A\}$	\emptyset
$\{S, A\}$	$\{A\}$
$\{S, A\}$	$\{S, A\}$

\rightarrow

$$\begin{aligned} S &\rightarrow A | ABa | AbA | Ba | bA | Ab | \lambda \\ A &\rightarrow Aa | a \\ B &\rightarrow Bb | BC \\ C &\rightarrow CB | CA | bB \end{aligned}$$

23 (cont.)

Eliminate chain rules

$$S \rightarrow A | ABa | AbA | Ba | bA | Ab | \lambda$$

$$A \rightarrow Aa | a$$

$$B \rightarrow Bb | BC$$

$$C \rightarrow CB | CA | bB$$

Chain(s)

$$\{S\}$$

$$\{C\}$$

$$\{S, A\}$$

PREV

$$\emptyset$$

$$\{S\}$$

$$\{S, A\}$$

NEW

$$\{S\}$$

$$\{A\}$$

Thus, add rules from A to S

$$S \rightarrow A | ABa | AbA | Ba | bA | Ab | \lambda | Aa | a$$

$$A \rightarrow Aa | a$$

$$B \rightarrow Bb | BC$$

$$C \rightarrow CB | CA | bB$$

Useless symbols

Term Prev

$$\{A\}$$

$$\{S, A\}$$

$$\{S, A\}$$

$$S \rightarrow Aa | a | AbA | bA | Ab | b | \lambda$$

$$A \rightarrow Aa | a$$

Reach Prev New

$$\{S\}$$

$$\emptyset$$

$$\{S, A\}$$

$$\{S\}$$

$$\{S\}$$

$$\{S, A\}$$

$$\{S, A\}$$

$$\{A\}$$

Chomsky transformation

$$S \rightarrow AA' | a | AT_1 | B'A | AB' | b | \lambda$$

$$A \rightarrow AA' | a$$

$$T_1 \rightarrow B'A$$

$$A' \rightarrow a$$

$$B' \rightarrow b$$

26. Use the CYK algorithm to produce the upper diagonal matrix of the Chomsky grammar:

$$S \rightarrow AX | AY | a$$

$$X \rightarrow AX | a$$

$$Y \rightarrow BY | a$$

$$A \rightarrow a$$

$$B \rightarrow b$$

	1	2	3	4	5
1	$\{S, X, Y, A\}$	\emptyset	$\{S\}$	\emptyset	\emptyset
2		$\{B\}$	$\{Y\}$	\emptyset	\emptyset
3			$\{S, X, Y, A\}$	$\{S, X\}$	$\{S, X\}$
4				$\{A, Y, X, S\}$	$\{S, X\}$
5					$\{A, Y, X, S\}$