

3.1 Rewrite the productions for each of the following nonterminals as right regular grammars: Identifier, Float

Identifier \rightarrow Letter { Letter | Digit }^{*}

Letter \rightarrow a | b | ... | z | A | B | ... | Z

Digit \rightarrow 0 | 1 | 2 | ... | 9

RGrammar

Identifier \rightarrow a { 0 | 1 | 2 | ... | 9 }^{*} identifier | b { 0 | 1 | 2 | ... | 9 }^{*} identifier | ... | z { 0 | 1 | 2 | ... | 9 }^{*} identifier | A { 0 | 1 | 2 | ... | 9 }^{*} identifier | B { 0 | 1 | 2 | ... | 9 }^{*} identifier | ... | Z { 0 | 1 | 2 | ... | 9 }^{*} identifier | a B 2

\rightarrow a identifier

\rightarrow a B identifier

\rightarrow a B 0 2

Integer \rightarrow 0 Int | ... | 9 Int |

\rightarrow 0 9 Int

\rightarrow 0 9 6 Int

identifier \rightarrow letter IdentRest

identifer \rightarrow letter IdentRest | digit identrest

Float \rightarrow Integer . Integer

Rq Float \rightarrow 0 Float | ... | 9 Float | . Float | 0 | ... | 9

4 2 . 3 2 \rightarrow 4 Float

\rightarrow 4 2 Float

\rightarrow 4 2 . Float

\rightarrow 4 2 . 3 Float

\rightarrow 4 2 . 3 2

Float \rightarrow digit IntRest

Intrest \rightarrow . digit DecRest | digit IntRest

DecRest \rightarrow | digit DecRest

3.2 Rewrite the productions for each of the following nonterminals as left regular grammars: identifier, Float

identifier \rightarrow identifier (a | ... | z | A | ... | Z) { 0 | 1 | 2 | ... | 9 }^{*} | λ

0 | 1 | 2 | ... | 9

a b 2 \rightarrow identifier b 2

\rightarrow a b 2

Float \rightarrow Float 0 | ... | Float 9 | Float . | λ

0 | 1 | 2 | ... | 9

Float \rightarrow IntPart . digit | Float digit

IntPart \rightarrow digit | IntPart digit

Identifier \rightarrow letter | Identifier letter | Identifier digit

3 2 . 4 \rightarrow Float 4

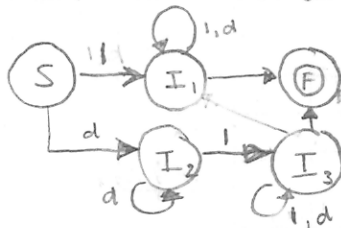
\rightarrow Float . 4

\rightarrow Float 2 . 4

\rightarrow Float 3 2 . 4

\rightarrow 3 2 . 4

3. Draw a DFSA for identifiers that contain only letters and digits, where the identifier must have at least one letter, but it need not be the first character



(S, Kurt123\$) \vdash (I1, urt123\$) \vdash (I1, rt123\$) \vdash (I1, t123\$) \vdash (I1, 123\$) \vdash (I1, 23\$) \vdash (I1, 3\$) \vdash (I1, \$) \vdash (F,)

(S, 12a2\$) \vdash (I2, 2a2\$) \vdash (I2, a2\$) \vdash (I3, 2\$) \vdash (I3, \$) \vdash (F,)

4. Try to define the language $\{a^n b^n\}$ using a regular grammar. Discuss why this might not be possible.

$n=0$

$L \rightarrow a L | b L | a b L$

A regular grammar cannot compensate for all forms of this expression.

10. SciNo Rgrammar

$$\text{SciNo} \rightarrow [+|-]\{0|1\cdots|9\}^* \cdot \{0|1\cdots|9\}^* e \{0|1\cdots|9\}^* \text{SciNo} | \lambda$$

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Float \rightarrow digit Float | .digit IntRest | ExponentIntRest \rightarrow digit IntRest | ExponentExponent \rightarrow E + Expo | E - Expo | E ExpoExpo \rightarrow digit Expo | digit
$$[0-9]^* \cdot ? [0-9]^* E \{+|- \} ? [0-9]^+$$


16. $S \rightarrow |a|(T)$ $\text{First}(S) = \{a(\}$
 $T \rightarrow T, S | S$ $\text{First}(T) = \{a(\}$
 $S' \rightarrow S \$$

