1. Define an EBNF and obstract syntax for adding record structures to Clibe. The EBNE for a structure reference grounduse the "dot" notation discussed in Chapter 5. The concrete and AS of Declaration, Expression, and Assignment should be modified

Struct employee Type } int id; char name[26]; int age; Float Salary; char debt; Struct employee Type employee; dot > employee 'dot' age = 45; dot: Structure of field of a of Structure

dot Str fld n I not (elem fld str) = error "field not in" 44 Str ++ "."

1 otherwise = structure & fld = n }

Here my tree would produce

- Assignment > Variable Ref

-> Variable

-> my Record

-> Declaration -> Variable Decl

- int

-> Declaration, etc.

type field = a data structure =>[field] } Record -> E Peccod Decl 3 Assignment -> Identifier = Expression \ Identifier = record

Expression -> VarRefl value 1 Bin | Malke Declaration -> Varioble Deci | Array Deci Record Decl Var = Record Record Decl -> Variable v; Type t my REC

Record Dell Record Del name; string age; int

Here I altered Assignment to compensate for a different kind of expression. Identifier & [Record Deci] }. I'd have to modify acid a keyword ' Pecord' to distinguish from a normal expression.

2. Expand the type system to determine when a record reference is will-typed. Make sum to consider born on occurrence on the left hand side and occurrency as expressions.

public Static void V (Program p) } V(p. decpart); V (p. body, typing (p. decpart)); }

Record my Rec }

Type Rule Ky: all records must have unique names. Type Rule Kz; All record's declared variables must have unique names Type Rule kz; field assagnmint int one; } error, duplicate field harms variable's types according to be type map. values must match truit declared

Record myker > error, duplicare record declarations String Medicy;

```
Record my Rec &
   int one = 1; No coercion defined by the type system
  int two = 2.0;
Record myrec }
   int one;
                 - error, this assignment produces an error
myrec. one = 1.5; as the type system does not allow type
                   coercion
```

Final

- 3. I will design an algorithm that determines structural equivolence
 - 1. A record is valid if there is no other record with the same name.
 - 2. Two records are agriculant when the following are true:
 - (a) The length of Record Decl are Egnal.
 - (b) The types of RI meson RZ based on the sorted enumeration of [int, floot, char, bool] being 0.-4. (RI types [0,2,0,0] & RZ types [1,2,0,0])
 - (C) The values of each field in RI must moten the values of each field in RZ

```
This is unneccessary as it would be
Begin: RI; RZ
    If (RI. get Name() == RZ. get Name())
                                                      cangue at compile time becouse
                                                      of the type system.
        error " Records cannot have the same name."
    If (RI. Record Decl. length () == RZ. Record Decl. length())
         for (int 1=0; ( & RI. Record Real, ungen ()-1; itt)
            for (int 1=0; j L RZ. Record Real length ()-1) 1++)
               +1 = RI.get Typel);
                t2 = RZ. get Typel);
                check ( t1 = t2);
       END: Equivalent
    END: Not Equivalent
```

4. Expand type Rule 6.2 for Declarations so that it defines the requirement that the type of each Variable be taken from a small set of available engles.

```
public Static void V (Declarations d) }
  for (inti=0, ikdisizel)-1; int)
    for (int j= i+1; j Ld. size()-1; j++) {
       Declaration di = d. get(i); Type type = diget Type();
Declaration di = d.get(i); Type type2=diget Type();
       if ( type == Type. FLOAT | Type. INT | Type. CHAR | Type. BOOL)
         if (type2 " "...)
          chick (! (di.v. equals (d).v)),
               "duplicate declaration: " + div);
      S.O.S. " Type not supported";
```

5. Consider how Type Ruse 6.2 must be expanded to include the Record type The type rule for duplicate names would be type rules Kland K2 concerted onto the existing rule definition.

++ All records must have unique names

- If a variable is not a reference to a Record, is many nove the same name, as it is considered a different element.
- 6. Add a ped () statement to Clite, where put () takes an expression as an argument and writes the value of the expression to the standard out (stolout)

Statement -> ; | Block | Assignment | If Statement | While Statement | Put

Statement = Skipl Block | Assignment | If Statement | while statement | Put
Put = Expression

Following the definitions of Type Reale 6.4

A Statement is valid with respect to the program's type map if it extistles
the following constraints

{1..53++6. A put statement is valid if all of its expressions are valid.

public static void V (Statement 5, TypeMaptin)

If (Sinstance of Put)

Put e = (Expression) e;

V (e, Typemaptim);

void put (Expression e) {

V(e, typemaphu)

SOS (e);

sample
int main () {

char c)

For (c='A'; c+='Z'; c++)

purchar (c);

return 0;

7. Type correion between record field definitions would work following the Transformer for expressions.

Recod myrec {
int x = 5;
float y = 4.0;
}

This would widen my int and save information loss.

Recording Recretample, Period example.

a = example.x + example.y example.x

If a record field has the type char, int, it may be widered to Float " Char, it may be widered to Int

type error

char a = exm.x + 'c' > the explicit declaration of char a makes this

Type correct a type error

float a = exm.y + 'c' > The unary ops cZi and iZt will allow this to be

type correct

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8. It would appear that Java has a sort of unuseful representation of Infinity under Double and Float types

S.O.S (Double . POSITIVE_INFINITY);
". NEGATIVE_INFINITY); which Print

> Infinity

o. Integer division rounds toward D. The quotient produced for operands in and d that are integers after binary numeric promotion is an interesting quince quinose magnitude is as large as possible while satisfying Id. 9/2/11/

if either operand is D, the result is NaN.