

19. Convert the grammar

$$\begin{aligned}
 G: S &\rightarrow aAB \mid ABC \mid a \\
 A &\rightarrow aA \mid a \\
 B &\rightarrow bBc \mid b \\
 C &\rightarrow abc
 \end{aligned}$$

to chomsky normal form.  $G$  already satisfies the conditions on the start symbol  $S$ ,  $\lambda$ -rules, useless symbols, and chain rules.

removing terminals from the RHS structure

$$\begin{array}{llll}
 S \rightarrow A'AB'B' & S \rightarrow A'T_1 & S \rightarrow A'T_1 \mid AT_3 \mid a & B \rightarrow B'BC'C \\
 A' \rightarrow a & \rightarrow T_1 \rightarrow AT_2 & T_1 \rightarrow AT_2 & \rightarrow C' \rightarrow C \rightarrow \\
 B' \rightarrow b & \text{for } T_2 \rightarrow B'B & T_2 \rightarrow B'B & \text{for } C' \rightarrow C \\
 & (aAbB) & A' \rightarrow a & (ABC) \quad T_3 \rightarrow BC \quad (bBcC) \\
 & & B' \rightarrow b &
 \end{array}$$

$$\begin{array}{llll}
 B \rightarrow B'T_4 & B \rightarrow B'T_4 \mid b & S \rightarrow A'T_1 \mid AT_3 \mid a & C \rightarrow A'T_6 \\
 T_4 \rightarrow BT_5 & \rightarrow T_4 \rightarrow BT_5 & T_1 \rightarrow AT_2 & T_6 \rightarrow B'C' \\
 T_5 \rightarrow C'C & T_5 \rightarrow C'C & T_2 \rightarrow B'B & \\
 & C' \rightarrow C & T_3 \rightarrow BC & \\
 & & A \rightarrow A'A \mid a & \text{add } C \text{ and terminals} \\
 & & B \rightarrow B'T_4 \mid b & \\
 & & T_4 \rightarrow BT_5 & \\
 & & T_5 \rightarrow C'C & \\
 & & C \rightarrow A'T_6 & \\
 & & T_6 \rightarrow B'C' & \\
 & & A' \rightarrow a & \\
 & & B' \rightarrow b & \\
 & & C' \rightarrow c &
 \end{array}$$

27. Let  $G$  be the grammar

$$\begin{aligned}
 G: S &\rightarrow A \mid B \\
 A &\rightarrow aaB \mid Aab \mid Aba \\
 B &\rightarrow bB \mid Bb \mid aba
 \end{aligned}$$

- Give a regular expression for  $L(G)$
- Construct a grammar  $G'$  that contains no left-recursive rules and is equivalent to  $G$ .

$$a. \{ (ab \cup ba)^* aa \{ b^* aba \}^* \cup \{ b^* aba \} \}$$

$$\begin{aligned}
 b. S &\rightarrow A \mid B \\
 A &\rightarrow \underbrace{Aa}_{u_1} b \mid \underbrace{Aba}_{u_2} \mid \underbrace{aaB}_{v_1} \\
 A &\rightarrow aaBZ \mid aaB \\
 Z &\rightarrow abZ \mid baZ \mid ab \mid ba
 \end{aligned}$$

$$\begin{aligned}
 B &\rightarrow \underbrace{Bb}_{u_1} \mid \underbrace{bB}_{v_1} \mid \underbrace{aba}_{v_2} \\
 B &\rightarrow bBZ_1 \mid abaZ_1 \mid bB \mid aba \\
 Z_1 &\rightarrow bZ_1 \mid b
 \end{aligned}$$

$abaZ$  represents a derived string where  $aba$  terminates the sequence of nonterminal  $B$ 's

30. construct a Greibach normal form grammar equivalent to

$S \rightarrow aAb|a$  remove indirect left recursion

$A \rightarrow SS|b$

$A \rightarrow aAbs|b$

$S \rightarrow aAb \rightarrow a(SS)b \rightarrow a((aAb)S)b \rightarrow a(a(b)b(aAb))b \rightarrow aabba bbb$

$S = 1, A = 2 \quad S \rightarrow aAb|a \rightarrow A \rightarrow aABS|b$   
 $A \rightarrow aAbs|b \quad B \rightarrow b$

$\Rightarrow S \rightarrow aAB|a$   
 $A \rightarrow aABS|b$   
 $B \rightarrow b$

33. convert the following grammar in CNF to GNF. Process the variables according to the order  $S, A, B$

$S \rightarrow BA|AB|A$

$A \rightarrow BB|AA|a$  Remove left recursion

$B \rightarrow AA|b \quad A \rightarrow v_1 Z | v_2 Z | v_1 v_2 \rightarrow A \rightarrow BBZ|aZ|BB|a$   
 $Z \rightarrow AZ|Z \quad Z \rightarrow AZ|Z$

$S = 1, A = 2, B = 3$

$B \rightarrow AA \rightarrow BBZA \rightarrow bBZA$

$B \rightarrow AA \rightarrow aZA$

$B \rightarrow AA \rightarrow BBZA \rightarrow bBA$

$B \rightarrow AA \rightarrow aA$

Revised rules

$S \rightarrow BA|AB|A$

$A \rightarrow BBZ|aZ|BB|a$

$B \rightarrow bBZA|aZA|bBA|aA$

$Z \rightarrow AZ|A$

For  $S$  rules

$S \rightarrow BA \rightarrow bBZAA \quad S \rightarrow bBZAA$

$S \rightarrow BA \rightarrow aZAA \quad S \rightarrow aZAA$

$S \rightarrow BA \rightarrow bBAA \quad S \rightarrow bBAA$

$S \rightarrow BA \rightarrow aAA \quad S \rightarrow aAA$

$S \rightarrow AB \rightarrow BBZ \rightarrow bBZABZ \quad S \rightarrow bBZABZ$

$S \rightarrow AB \rightarrow BBZ \rightarrow aZABZ \quad S \rightarrow aZABZ$

$S \rightarrow AB \rightarrow BBZ \rightarrow bBABZ \quad S \rightarrow bBABZ$

$S \rightarrow AB \rightarrow BBZ \rightarrow aABZ \quad S \rightarrow aABZ$

for  $A$  rules

$A \rightarrow BBZ \rightarrow bBZABZ \quad A \rightarrow bBZABZ$

$A \rightarrow BBZ \rightarrow aZABZ \quad A \rightarrow aZABZ$

$A \rightarrow BBZ \rightarrow bBABZ \quad A \rightarrow bBABZ$

$A \rightarrow BBZ \rightarrow aABZ \quad A \rightarrow aABZ$

$A \rightarrow BB \rightarrow bBZA|B \quad A \rightarrow bBZA|B$

$A \rightarrow BB \rightarrow aZAB \quad A \rightarrow aZAB$

$A \rightarrow BB \rightarrow bBAB \quad A \rightarrow bBAB$

$A \rightarrow BB \rightarrow aAB \quad A \rightarrow aAB$

$Z \rightarrow AZ \rightarrow BBZZ \rightarrow bBZABZZ \quad Z \rightarrow bBZABZZ$

$Z \rightarrow AZ \rightarrow BBZZ \rightarrow aZABZZ \quad Z \rightarrow aZABZZ$

$Z \rightarrow AZ \rightarrow BBZZ \rightarrow bBABZZ \quad Z \rightarrow bBABZZ$

$Z \rightarrow AZ \rightarrow BBZZ \rightarrow aABZZ \quad Z \rightarrow aABZZ$

$S \rightarrow bBZAA|aZAA|bBAA|aAA|bBZABZ|aZABZ|bBABZ|aABZ|A$

$A \rightarrow bBZABZ|aZABZ|bBABZ|aABZ|bBZA|aZAB|bBAB|aAB$

$B \rightarrow bBZA|aZA|bBA|aA$

$Z \rightarrow bBZABZZ|aZABZZ|bBABZZ|aABZZ$