

### 1. 欧几里得算法:

Euclid(m, n)

while n≠0 do

    r=m mod n;

    m=n;

    n=r;

return m;

### 2. 连续整数检测算法:

Function(m, n)

t=min(m, n);

while t≥0 do

    if m mod t == 0 and n mod t == 0

        return t;

    --t;

### 3. 素数筛选法:

Function(n)

for p=2 to n do

    A[p]=p;

for p=2 to sqrt(n) do

    if A[p]≠0

        j=p\*p;

        while j<n do

            A[j]=0;

            j=j+p;

i=0;

for p=2 to n do

    if A[p]≠0

        L[i]=A[p];

        ++i;

Return L;

### 4. 矩阵乘积:

Function(A[0..n-1, 0..n-1], B[0..n-1, 0..n-1])

for i=0 to n-1 do

    for j=0 to n-1 do

        C[i][j]=0;

        for k=0 to n-1 do

            C[i, j]=C[i, j]+A[i, k]\*B[k, j];

return C;

### 5. 汉诺塔 (A→B: A→C, C→B):

void hanoi(int n, int A, int B, int C)

{

    if(n>0)

    {

        hanoi(n-1, A, C, B);

        move(A, B);

        hanoi(n-1, C, A, B);

    }

}

### 6. 选择排序:

(选择当前序列最小值与当前序列首值交换)

SelectionSort(A[0..n-1])

for i=0 to n-1 do

    min=i;

    for j=i+1 to n-1 do

        if A[j]<A[min]

            min=j;

    swap(A[i], A[min]);

### 7. 冒泡排序:

BubbleSort(A[0..n-1])

for i=0 to n-2 do

    for j=i to n-2-i do

        if A[j+1]<A[j]

            swap(A[j+1], A[j])

### 8. 字符串蛮力匹配:

(逐字符匹配, 失败则从下一字符重新开始)

Function(T[0..n-1], P[0..m-1])

for i=0 to n-m do

    j=0;

    while j<m and P[j]=T[i+j] do

        ++j;

    if j==m return i;

return 0;

### 9. 蛮力平面距离最近两点:

Function(p)

d=∞

for i=0 to n-1 do

    for j=i+1 to n do

        d=min(d, sqrt((x<sub>i</sub>-x<sub>j</sub>)<sup>2</sup>+(y<sub>i</sub>-y<sub>j</sub>)<sup>2</sup>));

return d;

### 10. 深度优先搜索遍历:

DFS(G)

count=0;

for each vertex v in V do

    if v is marked with 0

        dfs(v)

dfs(v)//递归访问和 v 相连接未访问顶点

++count; mark v with count;

for each vertex w in V adjacent to v do

if w is marked with 0

dfs(w);

#### 11. 广度优先搜索遍历:

BFS(G)

count=0;

for each vertex v in V do

if v is marked with 0

bfs(v)

bfs(v)

++count;mark v with count and

init a queue with v;

while the queue is not empty do

for each vertex w in V adjacent

to the front vertex do

if w is marked with 0

++count;mark w with count;

add w to the queue;

remove the front vertex from the queue;

#### 12. 插入排序:

(往前找小于的值, 插在其后)

InsertionSort(A[0..n-1])

for i=0 to n-1 do

v=A[i];

j=i-1;

while j>=0 and A[j]>v do

A[j+1]=A[j];

--j;

A[j+1]=v;

#### 13. 拓扑排序:

执行一次 DFS 遍历;

并记住顶点变成死端(即退出遍历栈)的顺序

将该次序反序即得拓扑排序一个解

#### 14. 生成排列算法:

将第一个排列初始化为带左方向标志 12..n;

while 存在一个可移动元素 do

求最大移动元素 k;

把 k 和它箭头指向元素互换;

调转所有大于 k 的元素的方向;

将新排列添加到列表中;

#### 15. 反射格雷码:

BRG(n)

if n=1 表 L 包含位串 0 和 1;

else 调用 BRG(n-1) 生成长度 n-1 位串列表 L1;

把表 L1 倒序后复制给表 L2;

把 0 加到表 L1 中的每个位串前面;

把 1 加到表 L2 中的每个位串后面;

把表 L2 添加到表 L1 后面得到表 L;

Return L;

#### 16. 折半查找:

BinarySearch(A[0..n-1], k)

l=0, r=n-1;

while l<=r do

m=(l+r)/2;

if k==A[m] return m;

else if k<A[m] r=m-1;

else l=m+1;

return -1;

#### 17. 俄式乘法:

n 为偶数:  $n*m=n/2*2m$ ;

n 为奇数:  $n*m=(n-1)/2*2m+m$ ;

n	m	
50	65	
25	130	130
12	260	
6	520	
3	1040	1040
1	2080	2080
		3250

#### 18. 约瑟夫斯问题:

对 n 向左做一次循环移位;

#### 19. 三重查找:

l = 0; r = n - 1;

while l<=r

lmid = l+(r - l) / 3; //一二段分割点

rmid = r-(r - l) / 3; //二三段分割点

if k == A[lmid]

return lmid;

else if k == A[rmid]

return rmid;

else if k < A[lmid] //K 在第一段

r = lmid - 1;

else if k < A[rmid] //K 在第二段

l = lmid + 1, r = rmid - 1;

else //K 在第三段

l = rmid + 1;

return -1;

#### 20. Lomuto 划分:

LomutoPartition(A[l..r])

```

p=A[l];s=l;
for i=l+1 to r do
    if A[i]<p
        ++s;swap(A[s],A[i]);
swap(A[l],A[s]);
return s;
21. 快速选择:
Quickselect(A[l..r],k)
s= LomutoPartition(A[l..r]);
if s=l+k-1 return A[s];
else if s<l+k-1 Quickselect(A[l..s-1],k)
else Quickselect(A[s+1..r],l+k-1-s)
22. 快速排序:
Quicksort(A[l..r])
if l<r
    s= LomutoPartition(A[l..r]);
    Quicksort(A[l..s-1]);
    Quicksort(A[s+1..r]);
23. 合并排序:
Mergesort(A[0..n-1])
if n>1
    copy A[0..n/2-1] to B[0..n/2-1];
    copy A[n/2..n-1] to C[0..n/2-1];
    Mergesort(B[0..n/2-1]);
    Mergesort(C[0..n/2-1]);
    Merge(B,C,A);
24. Merge(B[0..p-1],C[0..q-1],A[0..p+q-1])
i=0;j=0;k=0;
while i<p and j<q do
    if B[i]≤C[j]
        A[k]=B[i];++i;
    else A[k]=C[j];++j;
    ++k;
if i==p
    copy C[j..q-1] to A[k..p+q-1];
else copy B[i..p-1] to A[k..p+q-1];
25. 二叉树遍历:
peroder(BTNode *p)
if p!=null
    printf(p->data);
    perorder(p->lchild);
    perorder(p->rchild);
26. 高斯消去法:
Function(A[1..n,1..n],b[1..n])

```

```

for i=1 to n do A[i,n+1]=b[i];
for i=1 to n-1 do
    for j=i+1 to n do
        for k=n+1 downto i do
            A[j,k]=A[j,k]-
                A[i,k]*A[j,i]/A[i,i];
27. 构造堆:
Function(H[1..n])
for i=n/2 downto 1 do
    k=i;v=H[k];heap=false;
    while not heap and 2*k≤n do
        j=2*k;
        if j<n
            if H[j]<H[j+1]
                ++j;
        if v≥H[j] heap=true;
        else H[k]=H[j];k=j;
H[k]=v;
28. 霍纳法则:
Horner(P[0..n],x)
p=P[n];
for i=n-1 downto 0 do
    p=x*p+P[i];
return p;
29. 从左至右二进制幂:
LRBE(a,b(n))
Product=a;
for i=1 to I do
    product=product*product;
    if bi=1 product=product*a;
return product;
30. 从右至左二进制幂:
RLBE(a,b(n))
term=a;
if bn=1 product=a;
else product=1;
for i=1 to I do
    term=term*term;
    if bi=1 product=product*term;
return product;
31. 比较计数排序:
Function(A[0..n-1])
for i=0 to n-1 do Count[i]=0;
for i=0 to n-2 do

```

```

    for j=i+1 to n-1 do
        if A[i]<A[j]
            ++Count[j];
        else ++Count[i];
    for i=0 to n-1 do
        S[Count[i]]=A[i];
return S;

```

### 32. 分布计数排序:

```

Function(A[0..n-1], l, u)
for j=0 to u-1 do D[j]=0
for i=0 to n-1 do ++D[A[i]-l];
for j=1 to u-1 do D[j]=D[j-1]+D[j];
for i=n-1 downto 0 do
    j=A[i]-l;
    S[D[j]-1]=A[i];
    --D[j];
return S;

```

### 33. 填充移动表:

```

ShiftTable(P[0..m-1])
for i=0 to size-1 do Table[i]=m;
for j=0 to m-2 do Table[P[i]]=m-1-j;
return Table;

```

### 34. Horspool 字符串匹配算法:

```

HorspoolMatching(P[0..m-1], T[0..n-1])
ShiftTable(P[0..m-1]); //生成移动表
i=m-1;
while i≤n-1 do
    k=0;
    while k≤m-1 and P[m-1-k]=T[i-k] do
        ++k;
    if k==m return i-m+1;
    else i=i+Table[T[i]];
return -1;

```

### 35. 币值最大化问题:

```

CoinRow(C[1..n])
F[0]=0; F[1]=C[1];
for i=2 to n do
    F[i]=max(C[i]+F[i-2], F[i-1]);
return F[n];

```

### 36. 找零问题:

```

ChangeMaking(D[1..m], n)
F[0]=0;
for i=1 to n do
    temp=∞; j=1;

```

```

while j≤m and i≥D[j] do
    temp=min(F[i-D[j]], temp);
    ++j;
F[i]=temp+1;
return F[n];

```

### 37. 硬币收集问题:

```

RobotCoinCollection(C[1..n, 1..m])
F[1, 1]=C[1, 1];
for j=2 to m do
    F[1, j]=F[1, j-1]+C[1, j];
for i=2 to n do
    F[i, 1]=F[i-1, 1]+C[i, 1];
    for j=2 to m do
        F[i, j]=max(F[i-1, j], F[i, j-1])+C[i, j];
return F[n, m];

```

### 38. 背包记忆化:

```

Function(i, j)
if F[i, j]<0
    if j<Weights[i]
        value=Function(i-1, j);
    else
        value=max(Function(i-1, j),
            Values[i]+Function(i-1, j-
                Weights[i]));
    F[i, j]=value;
return F[i, j];

```

### 39. 最优二叉查找树:

```

OptimalBST(P[1..n])
for i=1 to n do
    C[i, i-1]=0;
    C[i, i]=P[i];
    R[i, i]=i;
C[n+1, n]=0;
for d=1 to n-1 do
    for i=1 to n-d do
        j=i+d;
        minval=∞
        for k=i to j do
            if C[i, k-1]+C[k+1, j]<minval;
                minval=C[i, k-1]+C[k+1, j];
                kmin=k;
        R[i, j]=kmin;
        sum=P[i];
        for s=i+1 to j do

```

```

        sum=sum+P[s];
        C[i, j]=minval+sum;
return C[1, n], R;
40. Warshall 算法:
Warshall(A[1..n, 1..n])
R(0)=A;
for k=1 to n do
    for i=1 to n do
        for j=1 to n do
            R(k)[i, j]=R(k-1)[i, j] or
            R(k-1)[i, k] and R(k-1)[k, j];
return R(n)
41. Floyd 算法:
Floyd[W[1..n, 1..n]]
D=W;
for k=1 to n do
    for i=1 to n do
        for j=1 to n do
            D[i, j]=min{D[i, j], D[i, k]+D[k, j]};
return D;
42. 最小生成树 Prim 算法:
Prim(G)
V1={v0}; E1=空;
for i=1 to |V|-1 do
    在所有的边(v, u)中, 求权重最小的边
e*=(v*, u*);
    使得 v 在 Vi中, 而 u 在 V-Vi中;
    Vi=Vi ∪ {u*}; Ei=Ei ∪ {e*};
return Ei;
43. 最小生成树 Kruskal 算法:
Kruskal(G)
按照边权重非递减顺序对集合 E 排序
Ei=空; ecounter=0; //初始化边顶点集合及其规模
k=0; //初始化已处理边数量
while ecounter<|V|-1 do
    ++k;
    if Ei ∪ {eik} 无回路
        Ei=Ei ∪ {eik}; ++ecounter;
return Ei;
44. 最短路径 Dijkstra 算法:
Dijkstra(G, s)
Init(Q) //顶点优先队列初始化为空
for V 中每一个顶点
    dv=∞; pv=null;

```

```

    Insert(Q, v, dv) //初始化优先队列顶点优先级
ds=0; Decrease(Q, s, ds) //将 s 的优先级更新为 ds
Vi=空;
for i=0 to |V|-1 do
    u*=DeleteMin(Q) //删除优先级最小的元素
    Vi=Vi ∪ {u*};
    for V-Vi 中每一个和 u*相邻的顶点 u do
        if dus+w(u*, u)<du
            du=dus+w(u*, u); pu=u*;
            Decrease(Q, u, du);
45. 平分法求方程 x3+x-1=0 的根:
do
    mid=(a+b)/2;
    t3=f(mid);
    t1=f(a); t2=f(b);
    if t1*t3>0 a = mid;
    else b=mid;
while fabs(t3)>1e-2
return t3;
46. 试位法求方程 x3+x-1=0 的根:
do
    x=(a*f(b)-b*f(a))/(f(b)-f(a));
    y=f(x);
    if y*f(a) > 0 a = x; fa = y;
    else b = x; fb = y;
while fabs(y)>eps;
return x;
47. 牛顿法求方程 x3+x-1=0 的根:
f1 为原方程, f2 为其导数
do
    x=x0-f1(x0)/f2(x0);
    x0=x;
while fabs(f1(x))>eps;
return x;

```