CPMpy, a numpy-based CP modeling environment

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y @TiasGuns

CP 2021 tutorial



CP: Constraint Programming

"Solving constrained optimisation problems"

Vehicle Routing

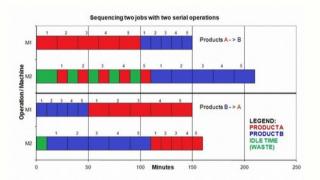


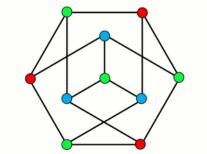
Scheduling

Configuration



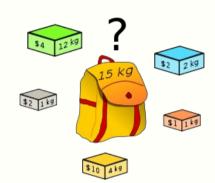
Graph problems & mathematical puzzles





CP Modeling

Knapsack:



Model =

- Variables, with a domain
- Constraints over variables
- Optionally: an objective

- gr, bl, og, ye, gy :: {0,1}
- -12*gr + 2*bl + 1*og + 4*ye + 1*gy <= 15
- maximize(4*gr + 2*bl + 1*og + 10*ye + 2*gy)

Model.solve()

Solver interfaces (cpmpy.solvers)

Expression transformations (cpmpy.transformations)

* » CPMpy: Constraint Programming and Modeling in Python

C Edit on GitHub

https://cpmpy.readthedocs.io/

CPMpy: Constraint Programming and Modeling in **Python**

CPMpy is a Constraint Programming and Modeling library in Python, based on numpy, with direct solver access.

Constraint Programming is a methodology for solving combinatorial optimisation problems like assignment problems or covering, packing and scheduling problems. Problems that require searching over discrete decision variables.

CPMpy allows to model search problems in a high-level manner, by defining decision variables and constraints and an objective over them (similar to MiniZinc and Essence'). You can freely use numpy functions and indexing while doing so. This model is then automatically translated to state-of-theart solver like or-tools, which then compute the optimal answer.

Source code and bug reports at https://github.com/CPMpy/cpmpy

Getting started:

- · Installation instructions
- · Getting started with Constraint Programming and CPMpy
- Quickstart sudoku notebook
- More examples

User Documentation:

- Setting solver parameters and hyperparameter search
- · Obtaining multiple solutions
- · UnSAT core extraction with assumption variables
- · How to debug · Behind the scenes: CPMpy's pipeline

API documentation:

- Expressions (cpmpy.expressions)
- Model (cpmpy.Model)
- Solver interfaces (cpmpy.solvers)

· Expression transformations (cpmpy.transformations)

CP Modeling

Knapsack:

Model =

- Variables, with a domain
- Constraints over variables
- Optionally: an objective

Model.solve()

```
model = Model()
gr,bl,og,ye,gy = boolvar(shape=5)
model += (12*gr + 2*bl + 1*og + 4*ye + 1*gy <= 15)
model.maximize(4*gr + 2*bl + 1*og + 10*ye + 2*gy)
model.solve()</pre>
```

```
print(gr.value(), bl.value(), og.value(), ye.value(), gy.value())
0 1 1 1 1
```

Why CPMpy?

Holy Grail: user specifies, solver solves [Freuder, 1997]

I think we reached it... MiniZinc, Essence'

Why CPMpy?

Holy Grail: user specifies, solver solves [Freuder, 1997]

I think we reached it... MiniZinc, Essence'

"Beyond NP" → CP as an oracle

Existing trend in AI, often with SAT or SMT or MIP solvers.

Growing need for more high-level solvers :: CP

Why CPMpy?

The more practical side of the story:

```
model = Model()
gr,bl,og,ye,gy = boolvar(shape=5)
model += (12*gr + 2*bl + 1*og + 4*ye + 1*gy <= 15)
model.maximize(4*gr + 2*bl + 1*og + 10*ye + 2*gy)
model.solve()</pre>
```

You have your CP model, but...

- Where does the data come from? (python code)
- How can you graphically visualize the result? (python code)
- How do you compare different formulations (python code)
- How can you test what solver params work best (python code)
- How do you make it solve for predicted cost vectors (python code)

What is CPMpy?

CPMpy is a:

- CP modeling library in Python
- based on numpy
- with direct solver access

CPMpy quick start

> pip3 install cpmpy

Will also install ortools (CP-SAT is default solver)

```
> python3
from cpmpy import *
a = boolvar(); b = boolvar()
Model( a & ~b ).solve()
    returns 'True'
```

Python/numpy how?

- All variables are numpy tensors
- Operator overloading
- Array indexing

```
v = boolvar()
print(v)
```

BV5

```
v = boolvar()
print(v)
```

BV5

```
v = boolvar(shape=5)
print(v)
```

[BV6 BV7 BV8 BV9 BV10]

```
v = boolvar()
print(v)
```

BV5

```
v = boolvar(shape=5)
print(v)
```

[BV6 BV7 BV8 BV9 BV10]

```
v = intvar(1,9, shape=5)
print(v)
```

[IV7 IV8 IV9 IV10 IV11]

```
v = boolvar()
print(v)
BV5
v = boolvar(shape=5)
print(v)
[BV6 BV7 BV8 BV9 BV10]
v = intvar(1,9, shape=5)
print(v)
[IV7 IV8 IV9 IV10 IV11]
m = intvar(1,9, shape=(3,3))
print(m)
[[IV12 IV13 IV14]
 [IV15 IV16 IV17]
 [IV18 IV19 IV20]]
```

```
v = boolvar()
print(v)
BV5
v = boolvar(shape=5)
print(v)
[BV6 BV7 BV8 BV9 BV10]
v = intvar(1,9, shape=5)
print(v)
[IV7 IV8 IV9 IV10 IV11]
m = intvar(1,9, shape=(3,3))
print(m)
[[IV12 IV13 IV14]
 [IV15 IV16 IV17]
 [IV18 IV19 IV20]]
t = boolvar(shape=(2,3,4))
print(t)
[[[BV11 BV12 BV13 BV14]
  [BV15 BV16 BV17 BV18]
  [BV19 BV20 BV21 BV22]]
 [[BV23 BV24 BV25 BV26]
  [BV27 BV28 BV29 BV30]
  [BV31 BV32 BV33 BV34]]]
```

```
v = boolvar()
print(v)
BV5
v = boolvar(shape=5)
print(v)
[BV6 BV7 BV8 BV9 BV10]
v = intvar(1,9, shape=5)
print(v)
[IV7 IV8 IV9 IV10 IV11]
m = intvar(1,9, shape=(3,3))
print(m)
[[IV12 IV13 IV14]
 [IV15 IV16 IV17]
 [IV18 IV19 IV20]]
t = boolvar(shape=(2,3,4))
print(t)
[[[BV11 BV12 BV13 BV14]
  [BV15 BV16 BV17 BV18]
  [BV19 BV20 BV21 BV22]]
 [[BV23 BV24 BV25 BV26]
  [BV27 BV28 BV29 BV30]
  [BV31 BV32 BV33 BV34]]]
```

```
puzzle start = np.array([
    [0,3,6],
    [2,4,8],
    [1,7,5]
(dim,dim2) = puzzle start.shape
assert (dim == dim2), "puzzle needs square shape"
n = dim*dim2 - 1 # e.g. an 8-puzzle
# State of puzzle at every step
K = 20
x = intvar(0,n, shape=(K,dim,dim), name="x")
print(x)
[[[x[0,0,0] \ x[0,0,1] \ x[0,0,2]]]
  [x[0,1,0] \ x[0,1,1] \ x[0,1,2]]
  [x[0,2,0] \ x[0,2,1] \ x[0,2,2]]]
 [[x[1.0.0] x[1.0.1] x[1.0.2]]
```

Operator overloading

```
x,y,z = intvar(1,9, shape=3)
print(x + y)
(IV33) + (IV34)
print(x * y)
(IV33) * (IV34)
print(abs(x - y))
abs([(IV33) + (-(IV34))])
a = intvar(1,9, shape=5, name="a")
print( sum(a) )
sum([a[0], a[1], a[2], a[3], a[4]])
c = (abs(sum(a) - (x+y)) == z)
print( c )
(abs([sum([a[0], a[1], a[2], a[3], a[4], -((IV33) + (IV34))]))) == (IV35)
```

Operator overloading

```
x,y,z = intvar(1,9, shape=3)
print(x + y)
                                                                type(c)
(IV33) + (IV34)
                                                                cpmpy.expressions.core.Comparison
                                                                c.name
print(x * y)
                                                                '=='
(IV33) * (IV34)
                                                                c.args[1]
print(abs(x - y))
                                                                IV35
abs([(IV33) + (-(IV34))])
                                                               type(c.args[0])
a = intvar(1,9, shape=5, name="a")
                                                                cpmpy.expressions.core.Operator
print( sum(a) )
sum([a[0], a[1], a[2], a[3], a[4]])
                                                                c.args[0].name
                                                                'abs'
c = (abs(sum(a) - (x+y)) == z)
print( c )
(abs([sum([a[0], a[1], a[2], a[3], a[4], -((IV33) + (IV34))]))) == (IV35)
```

Array indexing

```
x = boolvar(shape=(4,4), name="x")
print(x)
[[x[0,0] x[0,1] x[0,2] x[0,3]]
 [x[1,0] x[1,1] x[1,2] x[1,3]]
 [x[2,0] \ x[2,1] \ x[2,2] \ x[2,3]]
 [x[3,0] x[3,1] x[3,2] x[3,3]]
print(x[0,:])
[x[0,0] \ x[0,1] \ x[0,2] \ x[0,3]]
print(x[:,0])
[x[0,0] \ x[1,0] \ x[2,0] \ x[3,0]]
print(x[:,1:-1])
[[x[0,1] x[0,2]]
 [x[1,1] x[1,2]]
 [x[2,1] x[2,2]]
 [x[3,1] x[3,2]]
```

Python's indexing

Array indexing

```
x = boolvar(shape=4, name="x")
print(x)
[x[0] x[1] x[2] x[3]]
sel = np.array([True, False, True, False])
print(x[sel])
[x[0] x[2]]
print(x[np.arange(4) % 2 == 0])
[x[0] x[2]]
```

Numpy's indexing

Vectorized operations

```
x = intvar(1,9, shape=3, name="x")
y = intvar(1,9, shape=3, name="y")
print(x + y)

[(x[0]) + (y[0]) (x[1]) + (y[1]) (x[2]) + (y[2])]

print(x == [1,2,3])

[x[0] == 1 x[1] == 2 x[2] == 3]
```

Numpy's operator overloading

```
print(x == 1)
[x[0] == 1 x[1] == 1 x[2] == 1]
```

Numpy's broadcasting

Python/numpy how?

- Every variable is a numpy tensor
- Operator overloading
- Array indexing

Let's see it in action...

Sudoku

```
import numpy as np
from cpmpy import *
e = 0 # value for empty cells
qiven = np.array([
   [e, e, e, 2, e, 5, e, e, e],
    [e, 9, e, e, e, e, 7, 3, e],
    [e, e, 2, e, e, 9, e, 6, e],
    [2, e, e, e, e, e, 4, e, 9],
    [e, e, e, e, 7, e, e, e, e],
    [6, e, 9, e, e, e, e, e, 1],
    [e, 8, e, 4, e, e, 1, e, e],
    [e, 6, 3, e, e, e, e, 8, e],
    [e, e, e, 6, e, 8, e, e, e]])
# Variables
puzzle = intvar(1,9, shape=given.shape, name="puzzle")
```

Classic sudoku

```
model = Model()
n = given.shape[0]
# Constraints on rows and columns
for i in range(n):
    model += AllDifferent([puzzle[i,j] for j in range(n)])
    model += AllDifferent([puzzle[j,i] for j in range(n)])
# Constraints on blocks
for i in range(0,9, 3):
    for j in range(0,9, 3):
        model += AllDifferent([puzzle[r,c]
                               for r in range(i,i+3)
                               for c in range(j,j+3)])
# Constraints on values (cells that are not empty)
for r in range(n):
    for c in range(n):
        if given[r,c] != e:
            model += puzzle[r,c] == given[r,c]
model.solve()
```

CPMpy sudoku

```
model = Model()
# Constraints on rows and columns
model += [AllDifferent(row) for row in puzzle]
model += [AllDifferent(col) for col in puzzle.T]
# Constraints on blocks
for i in range(0,9, 3):
    for j in range(0,9, 3):
        model += AllDifferent(puzzle[i:i+3, j:j+3])
# Constraints on values (cells that are not empty)
model += [puzzle[given!=e] == given[given!=e]]
model.solve()
```

Job shop scheduling

```
jobs data = cpm array([ # (job, machine) = duration
    [3,2,2], # job 0
    [2,1,4], # job 1
    [0,4,3], # job 2 (duration 0 = not used)
max dur = sum(jobs data.flat)
n jobs, n machines = jobs data.shape
all jobs = range(n jobs)
all machines = range(n machines)
# Variables
start time = intvar(0, max dur, shape=(n machines,n jobs), name="start")
end time = intvar(0, max dur, shape=(n machines,n jobs), name="stop")
```

Classic jobshop

```
model = Model()
# end = start + dur
for j in all jobs:
    for m in all machines:
        model += (end time[m,j] == start time[m,j] + jobs data[j,m])
# Precedence constraint per job
for j in all jobs:
    for m1, m2 in combinations (all machines, 2): # [0,1,2] \rightarrow [(0,1),(0,2),(1,2)]
        model += (end time[m1,j] <= start time[m2,j])</pre>
# No overlap constraint: one starts before other one ends
for m in all machines:
    for j1,j2 in combinations(all jobs, 2):
        model += (start time[m,j1] >= end time[m,j2]) | 
                  (start time[m,j2] >= end time[m,j1])
# Objective: makespan
makespan = Maximum([end_time[m,j] for m in all_machines for j in all jobs])
model.minimize(makespan)
model.solve()
```

Classic jobshop CPMpy jobshop

```
model = Model()
model = Model()
# end = start + dur
                                           # end = start + dur
for j in all jobs:
                                           model += (end time == start time + jobs data.T)
    for m in all machines:
        model += (end time[m,j] == start
# Precedence constraint per job
                                           # Precedence constraint per job
                                           for m1,m2 in combinations(all machines,2):
for j in all jobs:
    for m1,m2 in combinations(all machine
                                               model += (end time[m1,:] <= start time[m2,:])</pre>
        model += (end time[m1, i] <= start</pre>
# No overlap constraint: one starts befor # No overlap constraint: one starts before other one ends
for m in all machines:
                                           for j1,j2 in combinations(all jobs, 2):
    for j1,j2 in combinations(all jobs, 2
                                               model += (start time[:,j1] >= end time[:,j2]) | \setminus
        model += (start time[m,j1] >= end
                                                        (start time[:,j2] >= end time[:,j1])
                 (start time[m,j2] >= end
# Objective: makespan
                                           # Objective: makespan
makespan = Maximum([end time[m,j] for m i
                                          makespan = max(end time)
model.minimize(makespan)
                                           model.minimize(makespan)
model.solve()
                                           model.solve()
```

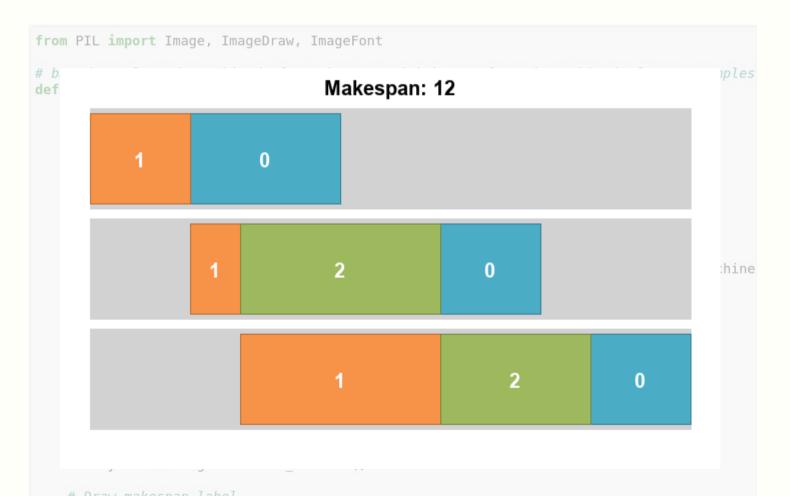
Let's inspect the result...

```
print("Makespan:", makespan.value())
print("Schedule:")
grid = -8*np.ones((n machines, makespan.value()), dtype=int)
for j in all jobs:
    for m in all machines:
         grid[m, start time[m, j].value():end time[m, j].value()] = j
print(grid)
Makespan: 12
Schedule:
[[ 1 1 0 0 0 -8 -8 -8 -8 -8 -8 -8 -8]
[-8 -8 1 2 2 2 2 0 0 -8 -8 -8]
[-8 -8 -8 1 1 1 1 2 2 2 0 0]]
```

Let's visualize the result...

```
from PIL import Image, ImageDraw, ImageFont
# based on Alexander Schiendorfer's https://github.com/Alexander-Schiendorfer/cp-examples
def visualize scheduling(start, end):
    nMachines, nJobs = start.shape
    makespan = max(end.value())
   # Draw solution
   # Define start location image & unit sizes
    start x, start y = 30, 40
    pixel unit = 50
    pixel task height = 100
    vert pad = 10
    imwidth, imheight = makespan * pixel unit + 2 * start x, start y + start x + nMachine
   # Create new Image object
    img = Image.new("RGB", (imwidth, imheight), (255, 255,255))
   # Create rectangle image
    img1 = ImageDraw.Draw(img)
   # Get a font
   try:
        myFont = ImageFont.truetype("arialbd.ttf", 20)
    except:
        myFont = ImageFont.load default()
    # Dunie makeanan lahal
```

Let's visualize the result...



N-puzzle (planning as SAT)

```
# '0' is empty spot
puzzle start = np.array([
    [3,7,5],
    [1,6,4],
    [8,2,0]]) # 19 steps
puzzle end = np.array([
    [1,2,3],
    [4,5,6],
    [7,8,0]])
def n puzzle(puzzle start, puzzle end, K):
    print("Max steps:", K)
   m = Model()
    (dim,dim2) = puzzle start.shape
    assert (dim == dim2), "puzzle needs square shape"
    n = dim*dim2 - 1 # e.g. an 8-puzzle
    # State of puzzle at every step
    x = intvar(0, n, shape=(K, dim, dim), name="x")
```

```
# Start state constraint
m += (x[0] == puzzle start)
# Fnd state constraint
m += (x[-1] == puzzle end)
# define neighbors = allowed moves for the '0'
def neigh(i,j):
    # same, left, right, down, up, if within bounds
    for (rr, cc) in [(0,0),(-1,0),(1,0),(0,-1),(0,1)]:
        if 0 <= i+rr and i+rr < dim and 0 <= j+cc and j+cc < dim:</pre>
            yield (i+rr,j+cc)
# Transition: define next (t) based on prev (t-1) + invariants
for t in range(1, K):
    # Invariant: in each step, all cells are different
    m += AllDifferent(x[t])
    # Invariant: only the '0' position can move
    m += ((x[t-1] == x[t]) | (x[t-1] == 0) | (x[t] == 0))
    # for each position, determine reachability of the '0' position
    for i in range(dim):
        for j in range(dim):
            m \leftarrow (x[t,i,j] == 0).implies(any(x[t-1,r,c] == 0 for r,c in neigh(i,j)))
return (m,x)
```

N-puzzle (planning as SAT)

```
(m,x) = n \text{ puzzle(puzzle start, puzzle end, 10)}
m.solve()
print(m.status())
Max steps: 10
ExitStatus.UNSATISFIABLE (0.010442393000000001 seconds)
(m,x) = n \text{ puzzle}(\text{puzzle start, puzzle end, } 100)
m.solve()
print(m.status())
Max steps: 100
ExitStatus.OPTIMAL (2.20881433 seconds)
```

N-puzzle (planning as SAT)

```
K0 = 5
step = 4
(m,x) = n \text{ puzzle}(\text{puzzle start, puzzle end, K0})
while not m.solve():
    print(m.status())
    K0 = K0 + step
    (m,x) = n \text{ puzzle}(\text{puzzle start, puzzle end, K0})
print(m.status())
Max steps: 5
ExitStatus.UNSATISFIABLE (0.0022607060000000003 seconds)
Max steps: 9
ExitStatus.UNSATISFIABLE (0.008280570000000001 seconds)
Max steps: 13
ExitStatus.UNSATISFIABLE (0.012872023000000002 seconds)
Max steps: 17
ExitStatus.UNSATISFIABLE (0.036358970000000004 seconds)
Max steps: 21
ExitStatus.OPTIMAL (0.113619127 seconds)
```

Other ways to make it faster?

https://github.com/CPMpy/cpmpy/blob/master/examples/advanced/hyperparameter_search.py

```
(m,x) = n puzzle(puzzle start, puzzle end, 20)
from cpmpy.solvers import CPM ortools, param combinations
all params = {'cp model probing level': [0,1,2,3],
              'linearization level': [0,1],
              'symmetry level': [0,1,2]}
configs = [] # (runtime, param)
for params in param combinations(all params):
    s = CPM ortools(m)
    print("Running", params, end='\r')
    s.solve(**params)
    configs.append( (s.status().runtime, params) )
best = sorted(configs)[0]
print("\nFastest in", round(best[0],4), "seconds, config:", best[1])
Max steps: 20
```

Fastest in 0.078 seconds, config: {'cp model probing level': 0, 'linearization level': 1, 'symmetry level': 0}

Running {'cp model probing level': 3, 'linearization level': 1, 'symmetry level': 2}

Other ways to make it faster?

https://github.com/CPMpy/cpmpy/blob/master/examples/advanced/hyperparameter_search.py

```
(m,x) = n \text{ puzzle}(\text{puzzle start, puzzle end, 20})
     from cpmpy.solvers import CPM ortools, param combinations
     all params = {'cp model probing level': [0,1,2,3],
                   'linearization level': [0,1],
                   'symmetry level': [0,1,2]}
     configs = [] # (runtime, param)
                                                            (m,x) = n \text{ puzzle(puzzle start, puzzle end, 100)}
(m,x) = n \text{ puzzle(puzzle start, puzzle end, 100)}
                                                            m = CPM ortools(m)
m.solve()
                                                            m.solve(**best[1])
print(m.status())
                                                            print(m.status())
Max steps: 100
                                                            Max steps: 100
ExitStatus.OPTIMAL (2.20881433 seconds)
                                                            ExitStatus_OPTIMAL (0.41233498100000004 seconds)
     Max steps: 20
     Running {'cp model probing level': 3, 'linearization level': 1, 'symmetry level': 2}
     Fastest in 0.078 seconds, config: {'cp model probing level': 0, 'linearization level': 1, 'symmetry level': 0}
```

CPMpy tutorial

End part 1:

CPMpy is a:

CP modeling library in Python

based on numpy

with direct solver access

CPMpy, a numpy-based CP modeling environment

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CP 2021 tutorial



CPMpy tutorial

Part 1: | CPMpy is a:

CP modeling library in Python

based on numpy

with direct solver access
Part 2: CP as an oracle, repeated solving

Why CPMpy?

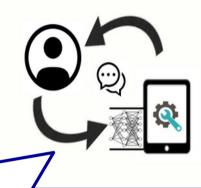
"CP as an oracle"



Consolidator grant (2021-2025)

"Conversational Human-aware Technology for Optimisation"





Towards co-creation of constrained optimisation solutions

- Solver that learns from user and environment
- Towards conversational: explanations and stateful interaction

https://people.cs.kuleuven.be/~tias.guns

@TiasGuns



Why CPMpy?

"CP as an oracle"

- 1. Explanation methods
- 2. Integrated solving and learning

=> requires repeated solving

Multiple solutions

MiniSearch-style:

Returns True (sol. found) or False (no solution)

```
x = intvar(0,3, shape=2)
m = Model(x[0] > x[1])
while m.solve():
    print(x.value())
    m += ~all(x == x.value()) # block solution
[3 0]
[3 2]
                                                      Adds constraint
[2 0]
                                                      to model
[1 0]
                                                      (even if already
                                                       solved before)
```

Diverse solutions

```
# a diversity measure, hamming distance
def hamm(x, y):
    return sum(x != y)
x = intvar(0,3, shape=2)
m = Model(x[0] > x[1])
store = []
while m.solve():
    print(len(store), ":", x.value())
    m += ~all(x == x.value()) # block solution
    store.append(x.value())
    # maximize number of elements that are different
    m.maximize(sum(hamm(x, sol) for sol in store))
```

```
0: [3 0]
1: [2 1]
2: [1 0]
3: [3 2]
4: [2 0]
5: [3 1]

Can change obj. function (even if already solved before)
```

Lazy vs Eager

```
x = intvar(0.30, shape=30)
x = intvar(0.30, shape=30)
                                                    m = Model([x[i-1] < x[i] for i in range(1, len(x))])
m = Model([x[i-1] < x[i] for i in range(1, len(x))])
t0 = time.time()
                                                    t0 = time.time()
                                                    m = CPM ortools(m)
while m.solve():
                                                    while m.solve():
   print(".",end="")
                                                        print(".",end="")
   m += ~all(x == x.value()) # block solution
                                                        m += ~all(x == x.value()) # block solution
print("time:", time.time()-t0)
                                                    print("time:", time.time()-t0)
          ....time: 1.4444539546966
                                                             .....time: 0.61836171150207
```

```
Model() = Lazy
Stores CPMpy expressions
```

SolverInterface() = Eager

Posts constraints to solver

1. Explanation methods

- Multiple and diverse solutions
- Minimal Unsatisfiable Subset (MUS)
- MUS/MSS enumeration (Marco algorithm)
- Optimal Unsatisfiable Subset (OUS)
- Explaining SAT problems step-wise
- Counter-factually explaining OPT problems

```
x = intvar(0.3. shape=4. name="x")
# circular 'bigger then', UNSAT
mus cons = [
    x[0] > x[1],
    x[1] > x[2],
    x[2] > x[0],
    x[3] > x[0],
    (x[3] > x[1]).implies(x[3] > x[2]) & ((x[3] == 3) | (x[1] == x[2]))
i = 0 # we wil dynamically shrink mus vars
while i < len(mus cons):
    # add all other remaining constraints
    assum cons = mus cons[:i] + mus cons[i+1:]
    if Model(assum cons).solve():
        # with all but 'i' it is SAT, so 'i' belongs to the MUS
        print("\tSAT so in MUS:", mus cons[i])
        i += 1
    else:
        # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS:", mus cons[i])
        # overwrite current 'i' and continue
        mus cons = assum cons
        SAT so in MUS: (x[0]) > (x[1])
        SAT so in MUS: (x[1]) > (x[2])
        SAT so in MUS: (x[2]) > (x[0])
        UNSAT so not in MUS: (x[3]) > (x[0])
```

UNSAT so not in MUS: (((x[3]) > (x[1])) -> ((x[3]) > (x[2]))) and ((x[3] == 3) or ((x[1]) == (x[2])))

```
ind = BoolVar(shape=len(mus cons), name="ind")
                                                      for i,bv in enumerate(ind):
                                                          assum model += [bv.implies(mus cons[i])]
x = intvar(0,3, shape=4, name="x")
                                                      # to map indicator variable back to soft constraints
# circular 'bigger then', UNSAT
                                                      indmap = dict((v,i) for (i,v) in enumerate(ind))
mus cons = [
    x[0] > x[1],
                                                      assum solver = CPM ortools(assum model)
    x[1] > x[2],
                                                      assert (not assum solver.solve(assumptions=ind)), "Model must be UNSAT"
    x[2] > x[0],
                                                      # unsat core is an unsatisfiable subset
    x[3] >
                                                     mus vars = assum solver.get core()
           x[1].implies(x[3] > x[2]) & ((x[3] == 3
                                                      print("UNSAT core of size", ten(mus vars))
                                                      # now we shrink the unsatisfiable subset further
                                                      i = 0 # we wil dynamically shrink mus vars
i = 0 # we wil dynamically shrink mus vars
                                                      while i < len(mus vars):</pre>
while i < len(mus cons):</pre>
    # add all other remaining constraints
                                                          # add all other remaining constraints
    assum cons = mus cons[:i] + mus cons[i+1:]
                                                          assum vars = mus vars[:i] + mus vars[i+1:]
    if Model(assum cons).solve():
                                                          if assum solver.solve(assumptions=assum vars):
        # with all but 'i' it is SAT, so 'i' belongs
                                                             # with all but 'i' it is SAI, so 'i' belongs to the MUS
        print("\tSAT so in MUS:", mus cons[i])
                                                              print("\tSAT so in MUS:", mus cons[i])
        i += 1
                                                              i += 1
    else:
                                                          else:
        # still UNSAT, 'i' does not belong to the MU
                                                              # still UNSAT, 'i' does not belong to the MUS
        print("\tUNSAT so not in MUS:", mus cons[i])
                                                              print("\tUNSAT so not in MUS:", mus cons[i])
        # overwrite current 'i' and continue
                                                              # overwrite current 'i' and continue
        mus cons = assum cons
                                                              mus cons = testcons
        SAT so in MUS: (x[0]) > (x[1])
                                                      UNSAT core of size 3
        SAT so in MUS: (x[1]) > (x[2])
                                                              SAT so in MUS: (x[0]) > (x[1])
        SAT so in MUS: (x[2]) > (x[0])
                                                              SAT so in MUS: (x[1]) > (x[2])
        UNSAT so not in MUS: (x[3]) > (x[0])
                                                              SAT so in MUS: (x[2]) > (x[0])
        UNSAT so not in MUS: (((x[3]) > (x[1])) \rightarrow (
```

assum model = Model()

make assumption indicators, add reified constraints

MUS/MSS enumeration (Marco algorithm)

```
from marco musmss enumeration import SubsetSolver, MapSolver
def do marco(model):
    sub solver = SubsetSolver(model.constraints)
   map solver = MapSolver(len(model.constraints))
   while True:
        seed = map solver.next seed()
        if seed is None:
            # all MUS/MSS enumerated
            return
        if sub solver.check subset(seed):
            MSS = sub solver.grow(seed)
            yield ("MSS", [model.constraints[i] for i in MSS])
            map solver.block down(MSS)
        else:
            seed = sub solver.seed from core()
            MUS = sub solver.shrink(seed)
            yield ("MUS", [model.constraints[i] for i in MUS])
            map solver.block up(MUS)
```

MUS/MSS enumeration (Marco algorithm)

```
for kind, exprs in do marco(m):
    print(kind,":")
    for e in sorted(exprs):
        print("\t", e)
MUS:
          (x[2]) > (x[0])
          (x[1]) > (x[2])
          (x[0]) > (x[1])
MSS :
          (x[1]) > (x[2])
          (x[0]) > (x[1])
          (((x[3]) > (x[1])) \rightarrow ((x[3]) > (x[2]))) and ((x[3] == 3)) or ((x[1]) == (x[2])))
          (x[3]) > (x[0])
MSS:
          (x[0]) > (x[1])
          (((x[3]) > (x[1])) \rightarrow ((x[3]) > (x[2]))) and ((x[3] == 3) \text{ or } ((x[1]) == (x[2])))
          (x[3]) > (x[0])
          (x[2]) > (x[0])
MSS:
          (((x[3]) > (x[1])) \rightarrow ((x[3]) > (x[2]))) and ((x[3] == 3)) or ((x[1]) == (x[2])))
          (x[3]) > (x[0])
          (x[2]) > (x[0])
          (x[1]) > (x[2])
```

1. Explanation methods

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Opt. US: implicit hitting set algorithm

```
assum solver = CPM ortools(assum model)
assert (not assum solver.solve(assumptions=ind)), "Model must be UNSAT"
hitset solver = CPM ortools(Model(
                    minimize=sum(weights*ind)))
while(True):
    hitset solver.solve()
   # Get hitting set
    hs = ind[ind.value() == 1]
    if not assum solver.solve(assumptions=hs):
        print("Found Optimal US, total weight:", sum(weights[ind.value() == 1]))
        for i in (ind.value() == 1).nonzero()[0]:
            print("\t", mus cons[i], "w=",weights[i])
        break
    # hs is satisfiable subset, hit one from complement
    C = ind[ind.value() == 0]
    hitset solver += (sum(C) >= 1)
```

Explaining SAT problems step-wise

```
b = boolvar(3, name="b")
m = Model(
    b[1].implies(b[0] | b[2]),
    b[0] | b[1],
    -b[0],
m.solve()
from ocus explanations import explain ocus
r = explain ocus(m.constraints, verbose=True)
Solution intersection: {b[1], ~b[0], b[2]}
Constraint(s): [~b[0]]
  and fact(s): []
           ==> \sim b[0] (cost: 6)
Constraint(s): [(b[0]) or (b[1])
  and fact(s): [~b[0]]
           ==> b[1] (cost: 7)
Constraint(s): [(b[1]) \rightarrow ((b[0]) \text{ or } (b[2]))]
  and fact(s): [ -b[0], b[1] ]
           ==> b[2] (cost: 8)
```

1. Explanation methods

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```
[tias@ishaladvanced python3 counterfactual explain.py
Solution to the following knapsack problem
Values = [45 48 65 68 68 10 84 22 37 88]
Weights = [71 89 89 13 59 66 40 88 47 89]
Capacity: 325
is: [0 3 4 6 8 9]
Resulting in an objective value of 390
Capacity used: 319
====== User Ouerv =======
I would like to change the following to the knapsack you provided:
Leave out items 3.4.6
Put in items 2,7
How should the values corresponding to these items change to make this assignment optimal?
====== Solving the master problem =======
Iteration 1, candidate costs: [45 48 65 68 68 10 84 22 37 88]
  Is foil-based solution now optimal? 212 >=? 390
Iteration 2, candidate costs: [45 48 91 68 0 10 0 22 37 88]
  Is foil-based solution now optimal? 238 >=? 329
Iteration 3, candidate costs: [ 45 48 65 68 68 10 84 200 37 88]
  Is foil-based solution now optimal? 390 >=? 508
Iteration 4, candidate costs: [45 48 65 0 68 10 34 82 37 88]
  Is foil-based solution now optimal? 272 >=? 309
Iteration 5, candidate costs: [45 48 65 0 65 10 0 45 37 88]
  Is foil-based solution now optimal? 235 >=? 263
Iteration 6, candidate costs: [ 45 48 152 31 68 10 84 76 37 88]
  Is foil-based solution now optimal? 353 >=? 431
Iteration 7, candidate costs: [45 48 74 0 37 10 37 45 37 88]
  Is foil-based solution now optimal? 244 >=? 273
Iteration 8, candidate costs: [45 48 74 29 66 10 8 74 37 88]
  Is foil-based solution now optimal? 273 >=? 302
Iteration 9, candidate costs: [45 48 65 15 37 10 22 60 37 88]
  Is foil-based solution now optimal? 250 >=? 250
Values [45 48 65 15 37 10 22 60 37 88] results in an optimal solution satisfying the user query
Optimal knapsack satisfying user query: [2 7 8 9], value 250
  diff: [ 0 0 0 53 31 0 62 -38 0 0]
```

A. Korikov & C. Beck, Counterfactual Explanations via Inverse Constraint

Reimplementation of

Programming, CP2021

```
# cutting plane algorithm
def inverse optimize(d orig, weights, capacity, x d, foil idx):
   Master problem: iteratively find better values for the 'd orig' vector
    (Korikov, A., & Beck, J. C., Counterfactual Explanations via Inverse Constraint Programming (CP2021))
   master model, d, x = make master problem(d orig, weights, capacity, x d, foil idx)
    sub model, x = 0 = make sub problem(d orig, weights, capacity)
   i = 1
   while master model.solve() is not False:
        d star = d.value() # master solution
        if verbose:
            print(f"Iteration {i}, candidate costs: {d star}")
       # solve subproblem
        sub model.maximize(sum(x 0 * d star))
        sub model.solve()
        if verbose:
            print(f" Is foil-based solution now optimal? {sum(d star * x d)} >=? {sum(d star * x 0.value())}")
        if sum(d star * x d) >= sum(d star * x 0.value()):
            return d star # is optimal
        else:
            # add cutting plane to master
            master model += [sum(d * x) >= sum(d * x 0.value())]
        i += 1
    raise ValueError("Master model is UNSAT!")
```

FROM examples/advanced/counterfactual explain.py

2. Integrated machine learning and solving

"CP as an oracle"

- Visual Sudoku (perception + solving)
- Preference learning in vehicle routing
- Decision-focussed learning (predict + optimize)

Why CPMpy?

"CP as an oracle"



Consolidator grant (2021-2025)

"Conversational Human-aware Technology for Optimisation"





CPMpy tutorial

Part 2: CP as an oracle, repeated solving

- Blocking constraints/solutions
- UNSAT core extraction
- Multiple solvers
- Implicit hitting set algorithms
- Cutting plane algorithms

CPMpy tutorial

Part 1: a CP modeling library in Python, based on numpy with direct solver access

Part 2: CP as an oracle, repeated solving

Part 3: peak behind the scenes

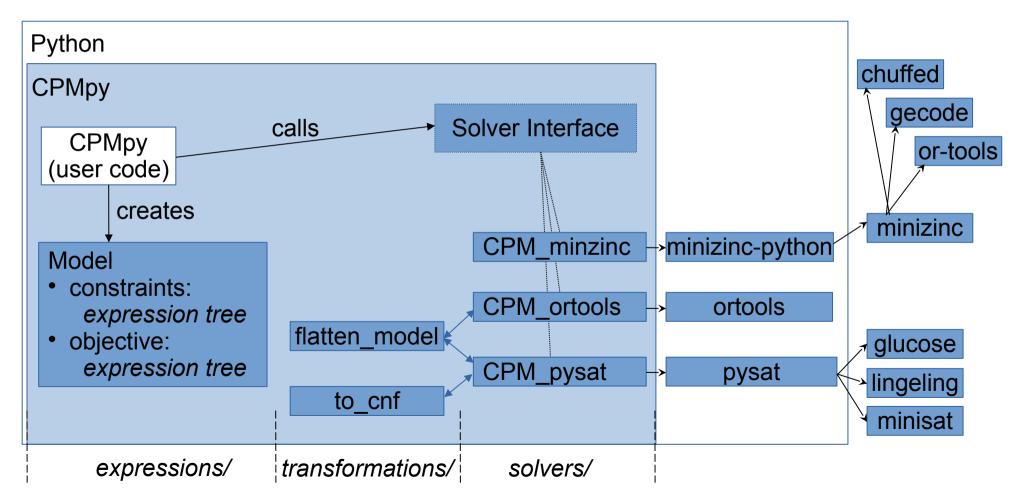
Design

Design principle:

Aim to be a thin layer on top of solver API

Central concept: CPMpy expression

Toolchain

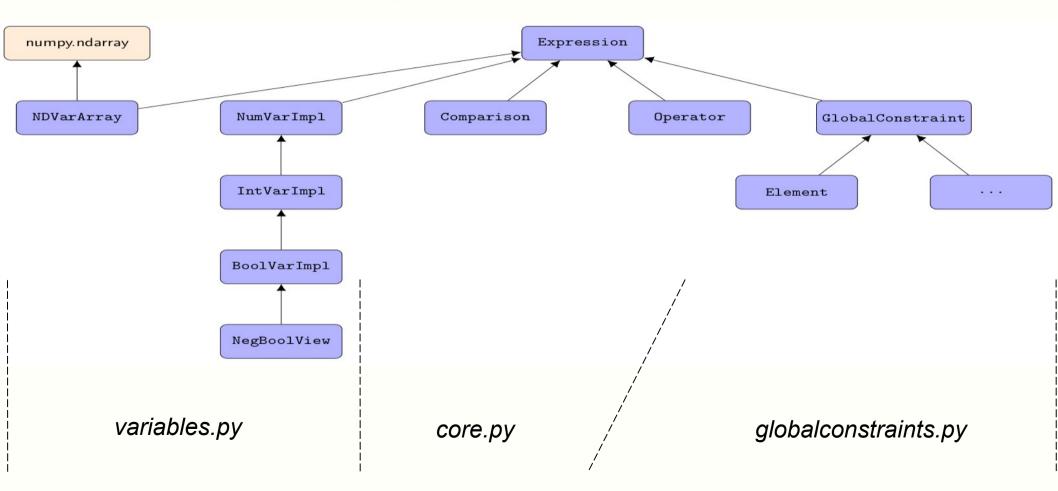


File structure

```
cpmpy/
  init___.py
 model.py
 expressions/
                 Generic container for
  transformations/
                   expressions
  solvers/
                 All kinds of expression objects
docs/
examples/
                 Common methods to rewrite
tests/
                   expressions
```

Classes that translata

Class Diagram of expressions/



Transformations

Key concept: 'Flat Normal Form'

Like Negated Normal Form and Conjunctive Normal Form

but for CP (basically: limited-nesting negated normal form)

In line with what other languages call 'flattening', but as a normal form

Flat Normal Form

```
- Boolean operators: and([Var]), or([Var]), xor([Var]) (CPMpy class 'Operator', is bool())
                                                           (CPMpv class 'Operator', is bool())
    - Boolean impliciation: Var -> Var
                                                           (CPMpy class 'Comparison')
    - Boolean equality: Var == Var
                        Var == Constant
                                                           (CPMpv class 'Comparison')
                                                           (CPMpy class 'GlobalConstraint', is bool())
    - Global constraint (Boolean): global([Var]*)
Comparison constraints: (up to one nesting on one side)
                                                           (CPMpy class 'Comparison')
    - Numeric equality: Numexpr == Var
                         Numexpr == Constant
                                                           (CPMpy class 'Comparison')
   - Numeric disequality: Numexpr != Var
                                                           (CPMpy class 'Comparison')
                          Numexpr != Constant
                                                           (CPMpy class 'Comparison')
                                                           (CPMpy class 'Comparison')
    - Numeric inequality (>=,>,<,<=): Numexpr >=< Var
   Numexpr:
        - Operator (non-Boolean) with all args Var/constant (examples: +,*,/,mod,wsum)
                                                           (CPMpv class 'Operator', not is bool())
       - Global constraint (non-Boolean) (examples: Max, Min, Element)
                                                           (CPMpv class 'GlobalConstraint', not is bool()))
   wsum: wsum([Const],[Var]) represents sum([Const]*[Var]) # TODO: not implemented yet
Reify/imply constraint: (up to two nestings on one side)
    - Reification (double implication): Boolexpr == Var
                                                           (CPMpy class 'Comparison')
    - Implication: Boolexpr -> Var
                                                           (CPMpy class 'Operator', is bool())
                  Var -> Boolexpr
                                                           (CPMpy class 'Operator', is bool())
   Boolexpr:

    Boolean operators: and([Var]), or([Var]), xor([Var]) (CPMpy class 'Operator', is bool())

       - Boolean equality: Var == Var
                                                               (CPMpy class 'Comparison')

    Global constraint (Boolean): global([Var]*)

                                                               (CPMpy class 'GlobalConstraint', is bool())
       - Comparison constraint (see above)
                                                               (CPMpy class 'Comparison')
   Reification of a comparison is the most complex case as it can allow up to 3 levels of nesting in total, e.g.:
       - (wsum([1,2,3],[IV1,IV2,IV3]) > 5) == BV
       - (IV1 == IV2) == BV
       - (BV1 == BV2) == BV3
```

The three families of possible constraints are:

Base constraints: (no nesting)

- Boolean variable

Solvers

We only interface to Python APIs (unfortunately, no Common CP solver API : (

Key principle: solver can implement any subset of expressions!

Solvers can also choose to:

- Support assumptions or not
- Be incremental or not
- Expose own solver parameters

Currently:

- ortools
- pysat
- minizinc

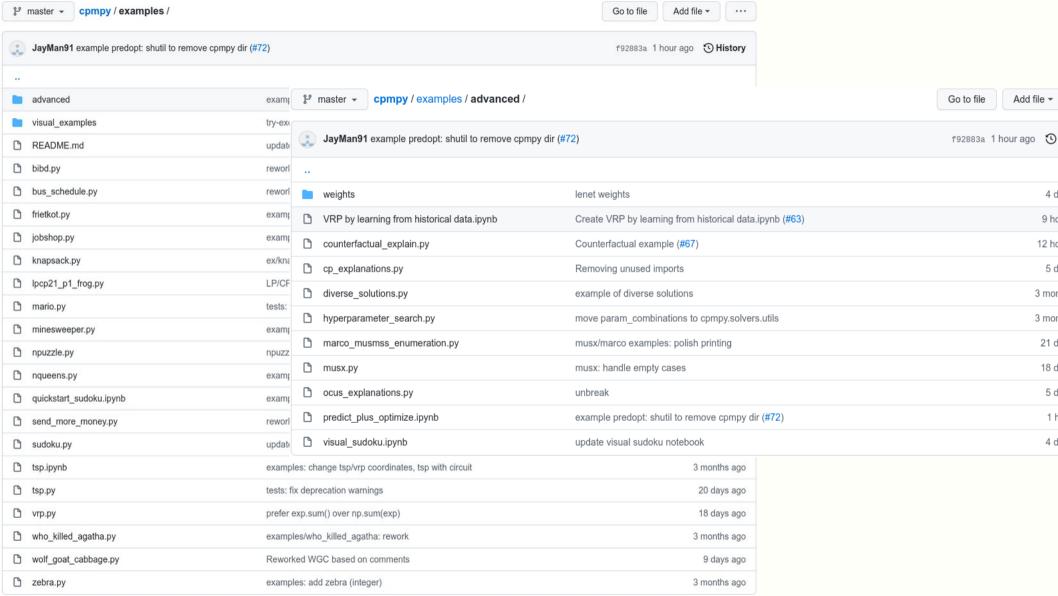
Near future: gurobi, ExactSolver Wishlist: Z3, Mistral2, Geas

CPMpy tutorial

Part 1: a CP modeling library in Python, based on numpy with direct solver access

Part 2: CP as an oracle, repeated solving

Part 3: peak behind the scenes (keep it light)



CPMpy tutorial

If you are a:

- CP user: give CPMpy a try, its fun
- Solver developer:

we can make your solver easier to use, if you make a Python API (contact us)

Report issues and new examples on https://github.com/CPMpy/cpmpy

Future Work

- More examples
- More advanced examples (paper reimplementations)
- More solvers (esp. incremental solvers!)

 towards Conversational, Human-Aware Technology for CP

Solver interfaces (cpmpy.solvers)

Expression transformations (cpmpy.transformations)

* » CPMpy: Constraint Programming and Modeling in Python

C Edit on GitHub

https://cpmpy.readthedocs.io/

CPMpy: Constraint Programming and Modeling in **Python**

CPMpy is a Constraint Programming and Modeling library in Python, based on numpy, with direct solver access.

Constraint Programming is a methodology for solving combinatorial optimisation problems like assignment problems or covering, packing and scheduling problems. Problems that require searching over discrete decision variables.

CPMpy allows to model search problems in a high-level manner, by defining decision variables and constraints and an objective over them (similar to MiniZinc and Essence'). You can freely use numpy functions and indexing while doing so. This model is then automatically translated to state-of-theart solver like or-tools, which then compute the optimal answer.

Source code and bug reports at https://github.com/CPMpy/cpmpy

Getting started:

- · Installation instructions
- · Getting started with Constraint Programming and CPMpy
- · Quickstart sudoku notebook
- More examples

User Documentation:

- Setting solver parameters and hyperparameter search
- · Obtaining multiple solutions
- · UnSAT core extraction with assumption variables
- · How to debug · Behind the scenes: CPMpy's pipeline

API documentation:

- Expressions (cpmpy.expressions)
- Model (cpmpy.Model)
- Solver interfaces (cpmpy.solvers)

· Expression transformations (cpmpy.transformations)