

ECE264: Advanced C Programming

Summer 2019

Week 5: Examples of Recursive Algorithms (Mergesort, Depth-first search, Enums, Unions, Complex Structures, Dynamic data structures (Linked lists, Stacks, Queues))

Merge sort

- Based on the principle of divide-conquer
 1. Divide the array into roughly two equal halves (sub-arrays)
 2. Sort the divided sub-arrays separately
 3. Merge the sorted sub-arrays to produce the larger array
- Always takes the same amount of time to run (best-case or worst-case)

Recursive code skeleton

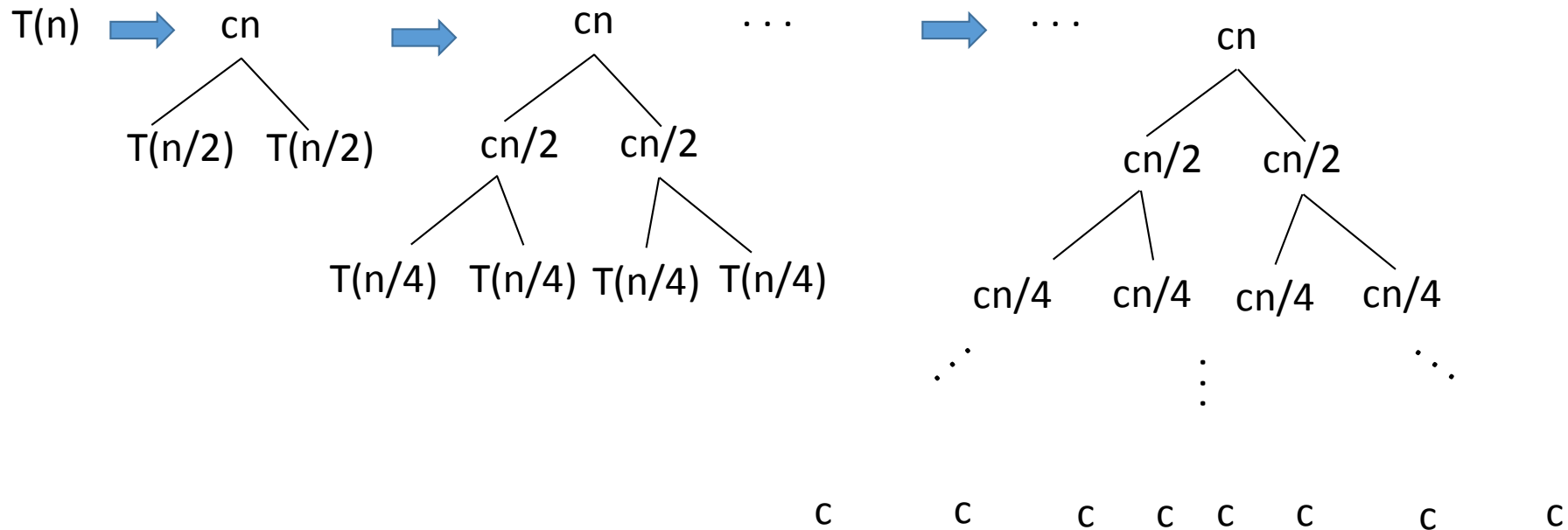
```
void Mergesort(int* arr, int left, int right) {  
    if (left >= right)  
        return;  
    //compute middle index. Left-subarray  
    always >= right sub-array  
    int nels = (right - left + 1) / 2;  
    int nelsLeft = (nels + 1) / 2 ;  
    int mid = left + nelsLeft - 1;  
    //Recursively sort  
    Mergesort(arr, left, mid)  
    Mergesort(arr, mid+1, right);  
    Merge(arr, left, mid, right);  
}
```

Recursive code skeleton

```
void Merge(int* arr, int l, int m, int r) {
1  numL = (m - l + 1); numR = (r - m);
2  //reserve space for 1 element more than numL and numR
   in left- and right-subarrays L and R
4  //copy arr[l] to arr[m] into L[0] to L[numL-1]
5  //copy arr[m+1] to arr[r] into R[0] to R[numR-1]
6  i=0;j=0; //initialize indices to L and R
7  L[numL]=INFINITY; R[numR]=INFINITY;
8  for(p=l; p<=r; p++) {
9      if(L[i] <= R[j]) {
10         arr[p]=L[i]; i++;
11     } else {
12         arr[p]=R[j]; j++;
13     }
14 }
15 }
```

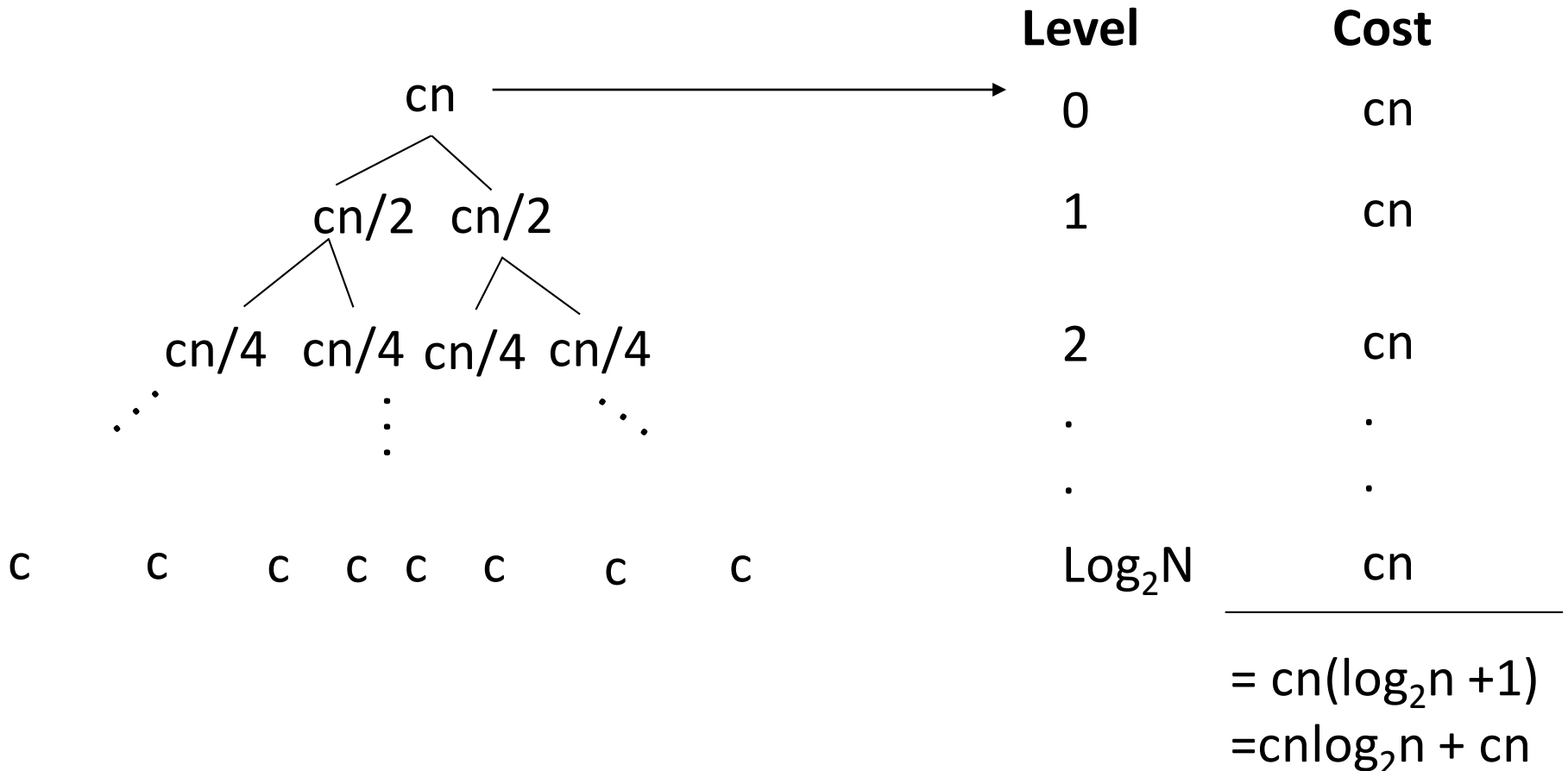
Merge sort analysis

- Assume n is perfect power of two



Merge sort analysis

- Assume n is perfect power of two



Recursion – more examples

- Depth first search
recall dictionary lookup:

“This is an increasingly common occurrence in our political **discourse**.”

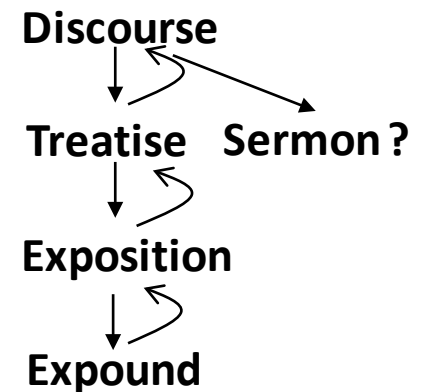
[Washington Post Jun 25, 2019](#)

discourse: a formal discussion of a subject in speech or writing, as a dissertation, **treatise**, sermon, etc.

treatise: a formal and systematic **exposition** in writing of the principles of a subject, generally longer and more detailed than an essay.

exposition: the act of **expounding**, setting forth, or explaining.

expound: To set forth or state in detail



Depth first search

We looked up the *first word we did not know* until we knew a definition completely. Then, we started backing up.

- Recursive procedure!
- One more example in PA?? (next week)

Enumerations (enums)

- Way to create user defined data type
- An alternative to using *magic numbers*
 - *Gives names to numbers – makes it easier to read and maintain programs*

```
typedef enum {<const1>, <const2>, ... <const n>}  
<typename>;
```

Enumerations (enums)

```
enum days{MON,TUE,WED,THU,FRI,SAT,SUN};
```



1. This **defines** a **type** called days

```
typedef enum {MON,TUE,WED,THU,FRI,SAT,SUN} days;
```



2. This **defines** a **type** called days (same as 1)

```
days dayOfWeek;
```



- This **defines** a **variable** called dayOfWeek, whose type is days

Enum - example

```
#include<stdio.h>
typedef enum days{MON,TUE,WED,THU,FRI,SAT,SUN};
int main(){
    days day = MON;
    days later = WED;
    days nineDaysLater = ?; //insert code here.
    printf("Now is %d",day);//prints "Now is 0"
    printf("Later is %d",later);//prints "Later is 2"
}
```

Enum – contd..

```
typedef enum days{MON,TUE,WED,THU,FRI,SAT,SUN};  
Value = 0,    1,    2,    3,    4,    5,    6
```

```
typedef enum days{MON=100,TUE,WED,THU,FRI,SAT,SUN};  
Value = 100,    101, 102, 103, 104, 105, 106
```

```
typedef enum days{MON,TUE,WED,THU=100,FRI,SAT,SUN};  
Value = 0,    1,    2,    100,    101, 102, 103
```

Unions

- Another way to create user defined data type

```
typedef union {  
    int i;  
    double d;  
    char c;  
}mpm;
```

Syntax is very similar to structure/enum definition

Union – accessing members

```
#include<stdio.h>
typedef union{
    int i;
    double d;
    char c;
}mpm;
int main(){
    mpm unionVar;//declaring an object of type mpm
    unionVar.i = 10;//accessing members of unionVar
    printf(“%d”,unionVar.i);//prints “10”
    unionVar.d = 3.14;
    printf(“%f”,unionVar.d);//prints “3.14”
    unionVar.c = ‘A’;
    printf(“%c”,unionVar.c);//prints A
}
```

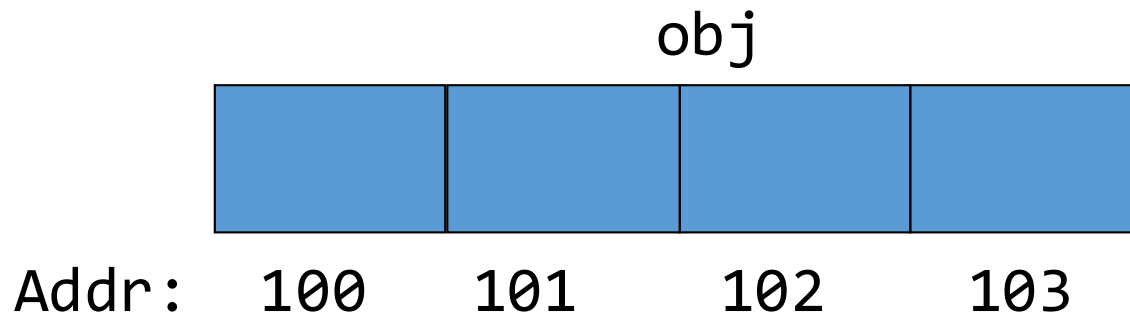
Union – overwriting memory

```
int main(){
    mpm unionVar;
    unionVar.i = 0x12345678;
    printf(“%x”,unionVar.i);//prints “12345678”
    unionVar.d = 3.14;
    printf(“%f”,unionVar.d);//prints “3.14”
    unionVar.c = ‘A’;
    printf(“%c”,unionVar.c);//prints A
    printf(“%x”,unionVar.i);//prints 51eb8541
    printf(“sizeof(mpm):%zu\n”,sizeof(mpm));
    //prints 8
}
```

Union – memory layout

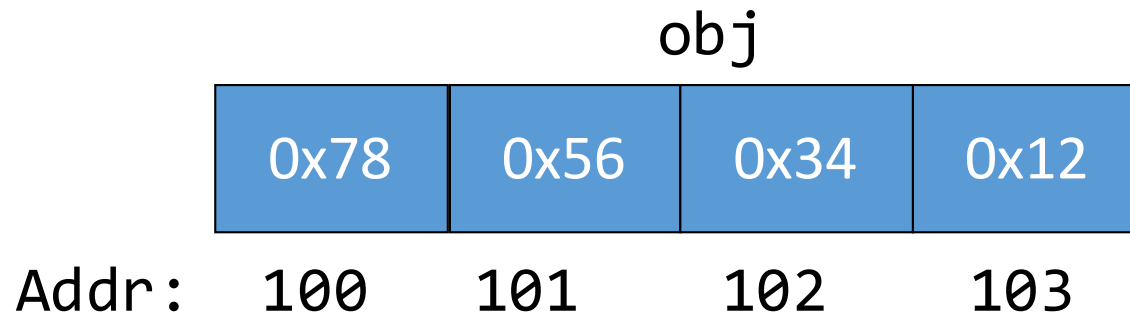
```
typedef union {  
    int i;  
    char c;  
}su;
```

```
su obj; //size of obj is 4 bytes (assuming int  
occupies 4 bytes)
```



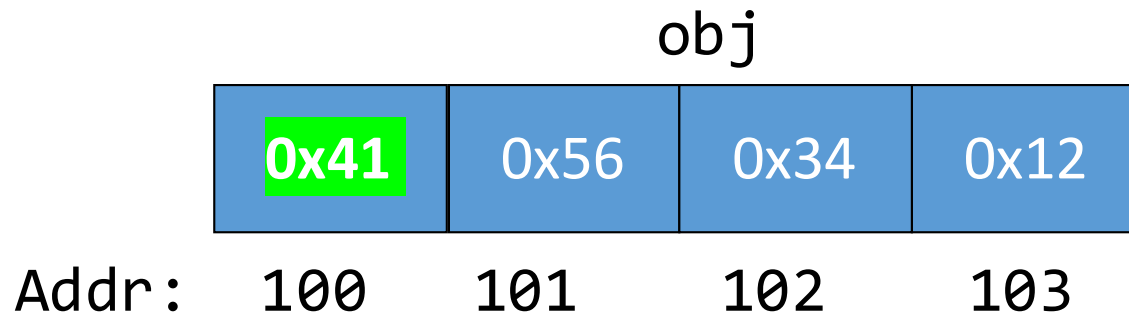
Union – memory layout

```
obj.i = 0x12345678;
```



Union – memory layout

```
obj.c = 'A';
```



Ascii value of 'A' is 65 = 0x41

Complex structures

- Can have struct members within (addr is a member of student, and its type is struct Addr)

```
typedef struct {  
    int ID;  
    char name[MAX_LEN];  
    Address addr;  
}Student;
```

```
typedef struct {  
    char* st;  
    int zip;  
}Addr;
```

Complex structures

- Accessing members is not different from accessing structure members

```
char* str=malloc(3); //allocate memory on heap
str[0]='E'; str[1]='C'; str[2]='E'; //initialize memory
Student stu1; //defines a variable stu1 of type Student
stu1.addr //accesses the addr member
stu1.addr.st //accesses the st member of Addr type
stu.addr.st = str; //assigns address of the heap memory
                  location to st
```

Complex structures

- Careful while assigning (shallow-copy)

```
char* str=malloc(3); //allocate memory on heap
str[0]='E'; str[1]='C'; str[2]='E'; //initialize memory
Student stu1, stu2;
stu1.addr.st = str; //assigns a dynamically
                    allocated char array to st
stu2=stu1 //copies the value of str.NOT what str points to
free(stu1.addr.st) //frees memory allocated to str
stu1.addr.st = NULL; //resets the pointer
```

stu2.addr.st still points to memory released

2D Arrays

- 1D array gives us access to *a row* of data
- 2D array gives us access to *multiple rows* of data
 - A 2D array is basically an *array of arrays*
- Consider a fixed-length 1D array
`int arr1[4];`//defines array of 4 elements; every element is an integer. Reserves contiguous memory to store 4 integers.



Starting addr:

100

104

108

112

2D Arrays (fixed-length)

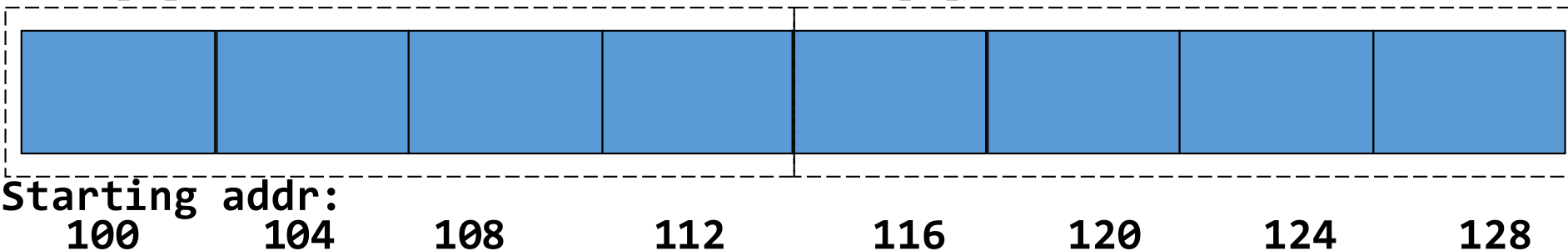
- Consider a fixed-length 2D array (*array of arrays*). Think:
 - array of integers => every element is an `int`
 - array of characters => every element is a `char`
 - array of array => every element is an *array*
- Example:

```
int arr[2][4];
```

//defines array of 2 elements; every element is an array of 4 integers. Therefore, reserves contiguous memory to store 8 integers

arr[0]

arr[1]



2D Arrays (on heap)

- What if we don't know the length of the array upfront?

E.g. A line in a file contains number of people riding a bus every trip.

Multiple trips happen per day and the number can vary depending on the traffic.

Day1 numbers: 10 23 45 44

Day2 numbers: 5 33 38 34 10 4

Day3 numbers: 9 17 10

.....

DayN numbers: 13 15 28 22 26 23 22 21

//we need array arr of N elements; every element is an array of M integers. Both N and M vary with every file input.

2D Arrays (on heap)

1. First, we need to create an array `arr2D` of N elements. So, get the number of lines in the input file.

- But what is the *type* of every element? - array of M elements, where every element is an integer (i.e. every element is an integer array). `int *`
- What is the type of `arr2D`? (array of array of integers)

Think:

type of an integer => `int`

type of array of integers => `int *`

(append a `*` to the type for every occurrence of the term array)

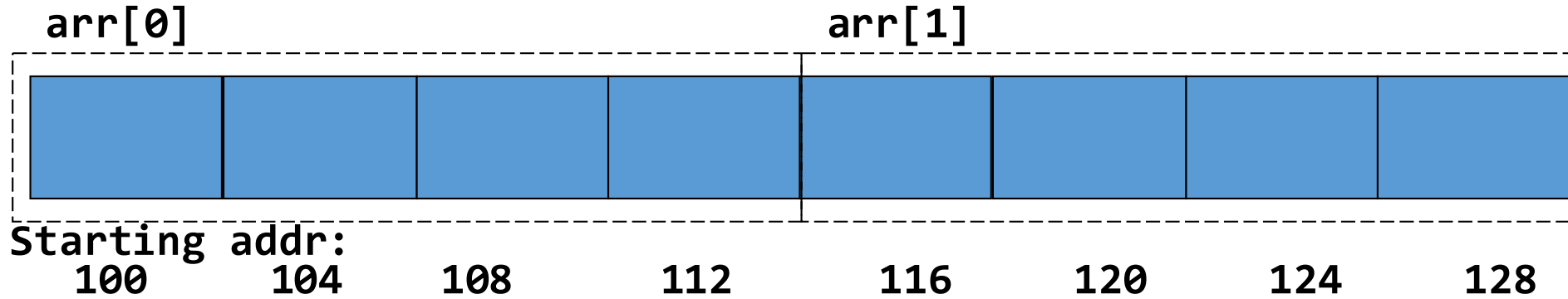
type of array of array of integers => `int **`

2D Arrays (on heap)

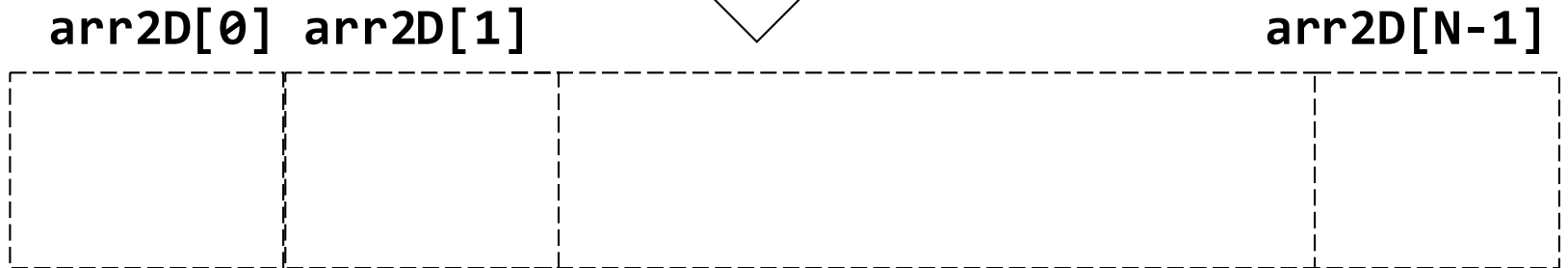
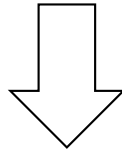
1. First, we need to create an array `arr2D` of N elements. So, get the number of lines in the input file.
 - What is the type of `arr2D`? (`int **`)

```
int N = GetNumberOfLinesFromFile(argv[1]);  
int** arr2D = malloc(sizeof(int *) * N)
```

Recall boxes with dashed lines in `int arr[2][4];`

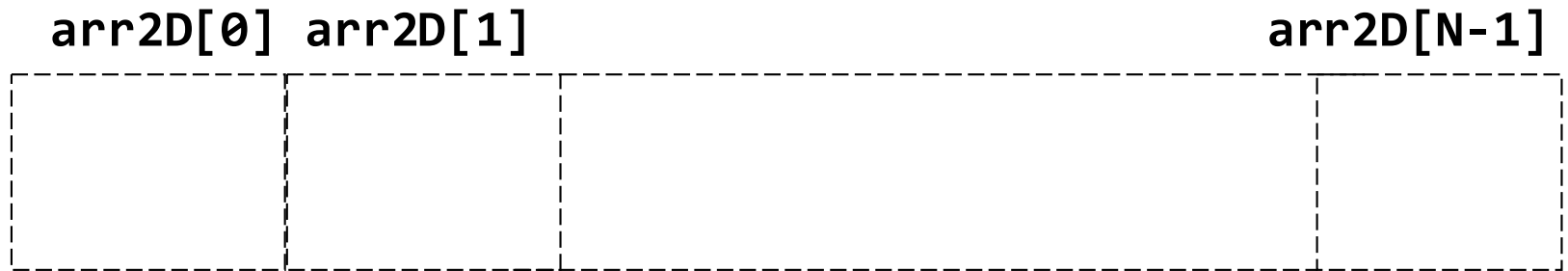


```
int N = GetNumberOfLinesFromFile(argv[1]);
int** arr2D = malloc(sizeof(int *) * N)
```



Starting addr(assuming 64-bit machine/pointer stored in 8 bytes):

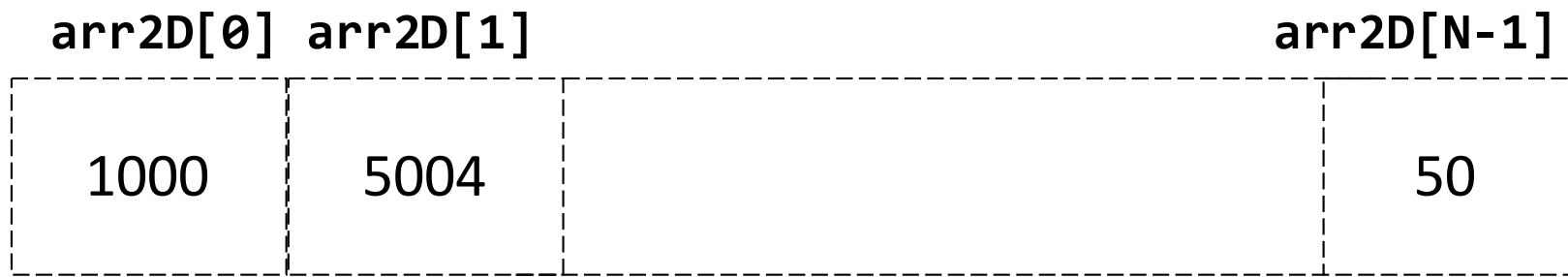
100	108	$100 + (N-1) * 8$
-----	-----	-------------------



Starting addr(assuming 64-bit machine/pointer stored in 8 bytes):
100 108 100+(N-1)*8

2. `arr2D[0]`, `arr2D[1]`, etc. are not initialized. They hold garbage values. How do we initialize them?

```
for(int i=0;i<N;i++) {  
    char* line = ReadLineFromFile(argv[1]);  
    int M = GetNumberOfIntegersPerLine(line);  
    arr2D[i] = malloc(sizeof(int) * M)  
}
```



Starting addr(assuming 64-bit machine/pointer stored in 8 bytes):

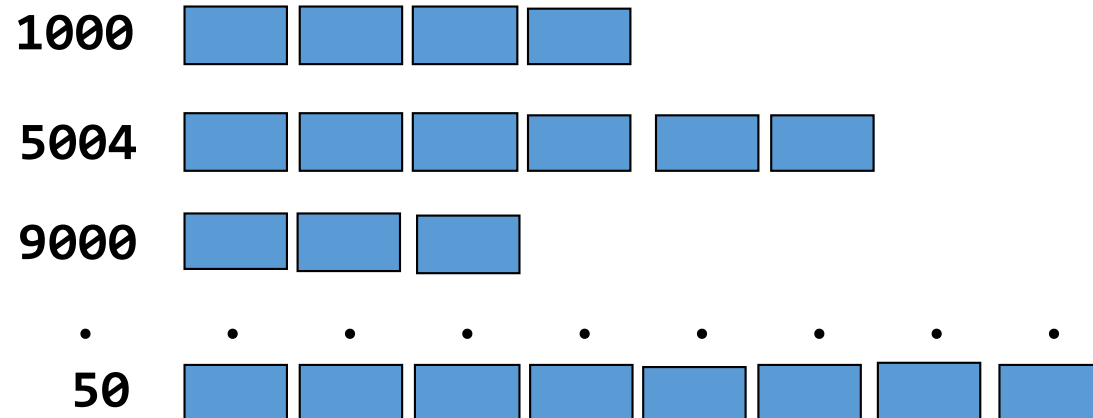
100

108

$100 + (N-1) * 8$

```
for(int i=0;i<N;i++) {
    char* line = ReadLineFromFile(argv[1]);
    int M = GetNumberOfIntegersPerLine(line);
    arr2D[i] = malloc(sizeof(int) * M)
}
```

Starting addr:



2D Arrays (on heap)

2. Now we need to initialize each of the array elements (in the second dimension)

Summary:

Creation: 2-steps

Initializing: 2-steps

Freeing: 2-steps

```
for(int i=0;i<N;i++)  
    free(arr2D[i]); //frees memory at 1000, 5004, etc.  
free(arr2D); //frees memory at 100
```

2D Arrays (trivia)

- Notation used to refer to elements different from cartesian coordinates

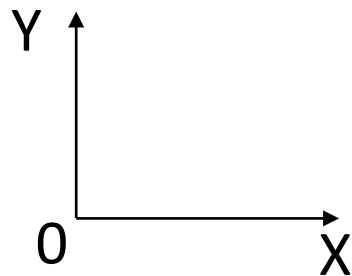
`arr2D[0][0]` accesses 1st row, 1st element

`arr2D[0][1]` accesses 1st row, 2nd element

`arr2D[1][1]` accesses 2nd row, 2nd element

`arr2D[N][M]` accesses $N+1^{\text{th}}$ row, $M+1^{\text{th}}$ element

- Cartesian:



- (M, N) = move M along X axis, N along Y axis
- In 2D array notation (M, N) means move to $M+1^{\text{th}}$ row (along Y axis), to $N+1^{\text{th}}$ column (along X axis)!

- From the previous bus trip data, what if we wanted to:

Day1 numbers: 10 23 45 44

Day2 numbers: 5 33 38 34 10 4

Day3 numbers: 9 17 10

.....

DayN numbers: 13 15 28 22 26 23 22 21

- Drop certain days as we analyzed arr2D?
- Add more days to (read from another file) to arr2D ?

i.e.

modify arr2D as program executes?

Dynamic Data Structures

- We use *dynamic data structures*
 - Allocate more space as we realize that we need to store more data
 - Free up space when we realize that we are storing less data
- Example:
 - Linked Lists, Trees, Stacks, Queues etc.

Linked Lists

- Most basic dynamic data structure
- Create a linked set of *structures* (struct objects)
 - Each structure holds a piece of data, and a *pointer* to the *next* structure in the list, which holds the next piece of data

How can we create such a structure?

Linked Lists

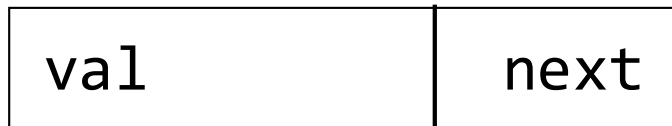
- Use recursive structure definition
 - Pointer from *within* the structure to another structure *of the same type*
 - Example: (structure holds integer data)

```
typedef struct Node {  
    int val;  
    struct Node* next;  
}Node;
```

- Note the 'struct Node' as part of the type for member next

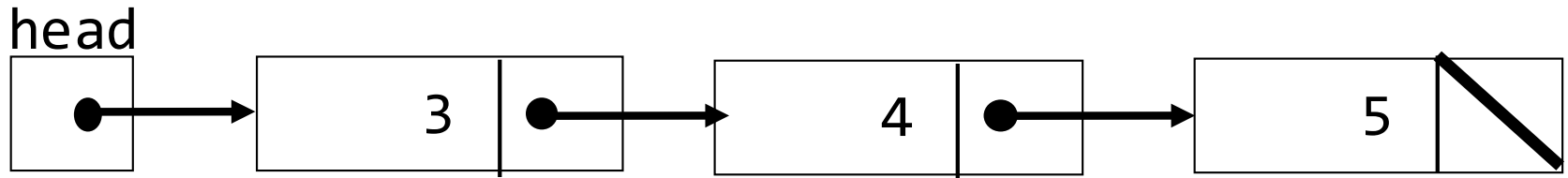
Linked Lists

- Graphical representation of Node



- Creating a list of integers:

```
Node* head = malloc(sizeof(Node));  
head->val=3;  
head->next = malloc(sizeof(Node));  
head->next->val=4;  
head->next->next = malloc(sizeof(Node));  
head->next->next->val=5;  
Head->next->next->next = NULL;
```

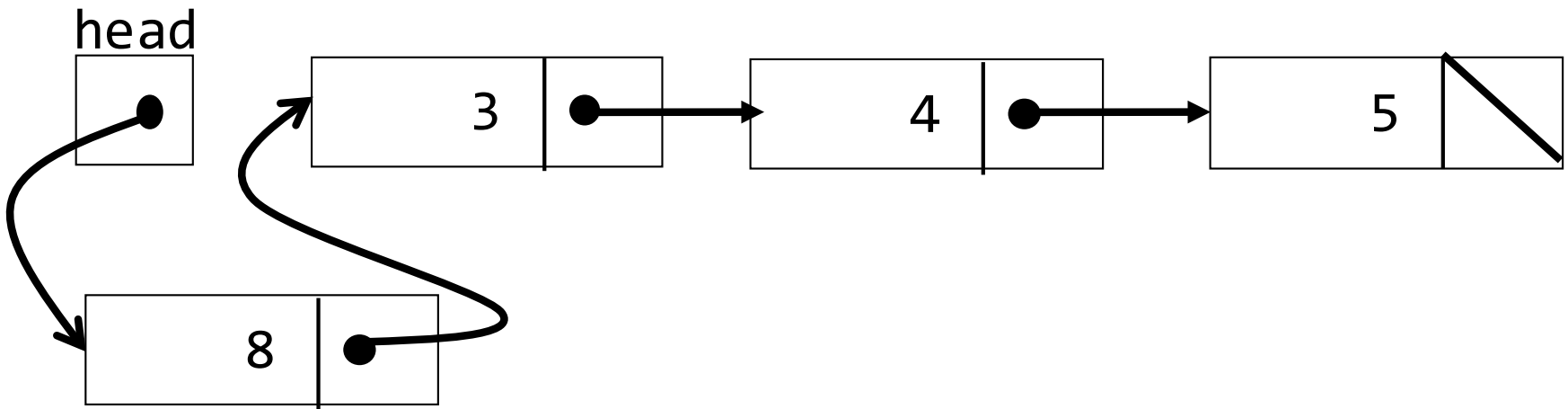


- head is a Node * just like next: use same sized box to represent head
- Using head, we can get to any node in the list
 - Just follow the next pointers
- Next field of last node is NULL . So, represent it with a slash

Linked Lists (updating the list)

- Add a new number

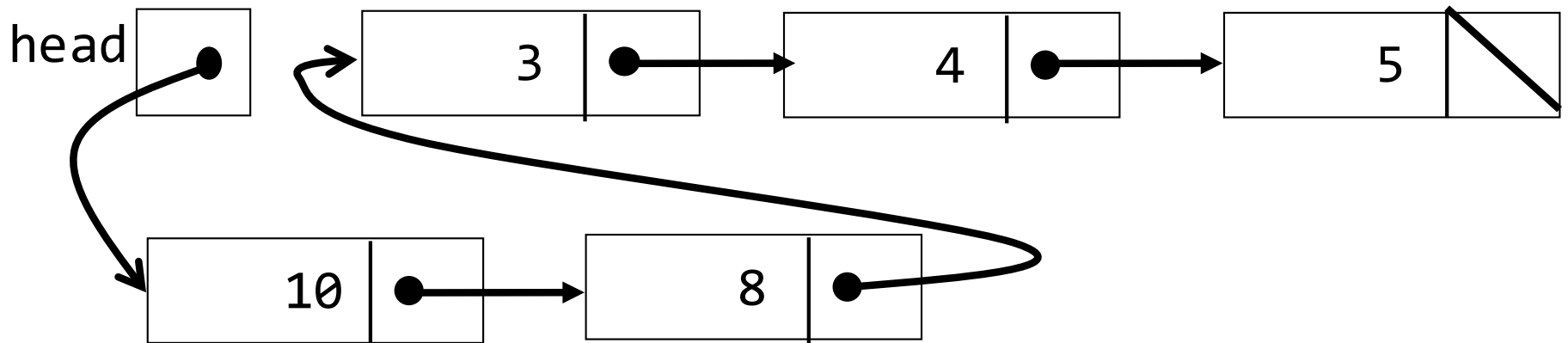
```
Node* newNode = malloc(sizeof(Node));  
newNode->val=8;  
newNode->next = head;  
head = newNode;
```



Linked Lists (updating the list)

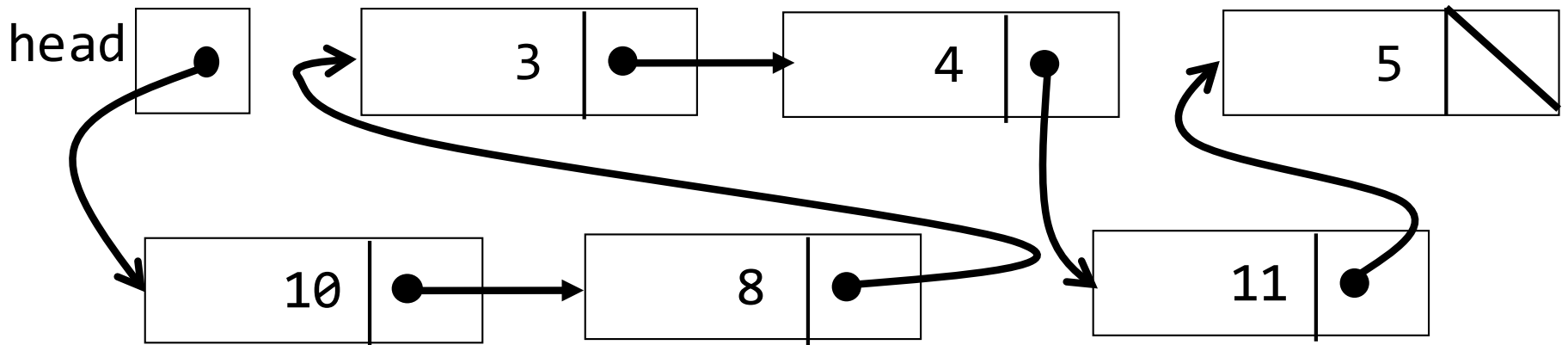
- Add a new number:

```
void insert(Node** loc, int val) {  
    Node* newNode = malloc(sizeof(Node));  
    newNode->val=val;  
    newNode->next = *loc;  
    *loc = newNode;  
}  
insert(&head, 10);
```



Linked Lists (updating the list)

```
Node* cur=head->next->next->next->next;//points to 4  
insert(&cur, 11);
```

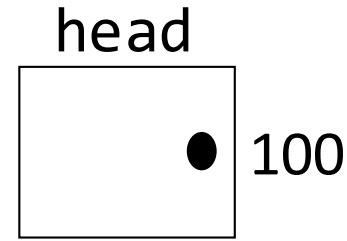


List Creation (Analysis)

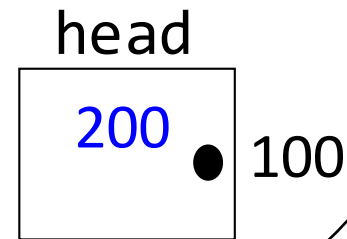
```
Node* head = malloc(sizeof(Node));  
head->val=3;  
head->next = malloc(sizeof(Node));  
head->next->val=4;  
head->next->next = malloc(sizeof(Node));  
head->next->next->val=5;  
Head->next->next->next = NULL;
```

List Creation (Analysis)

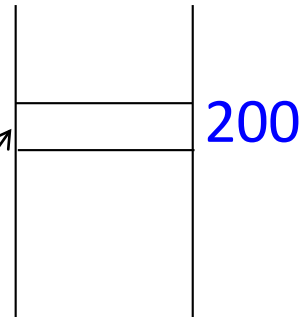
```
Node* head;
```



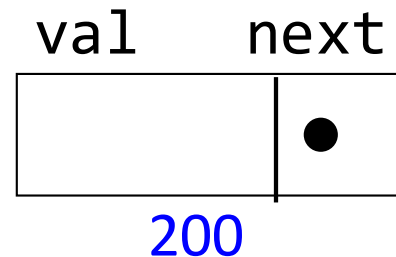
```
head = malloc(sizeof(Node));
```



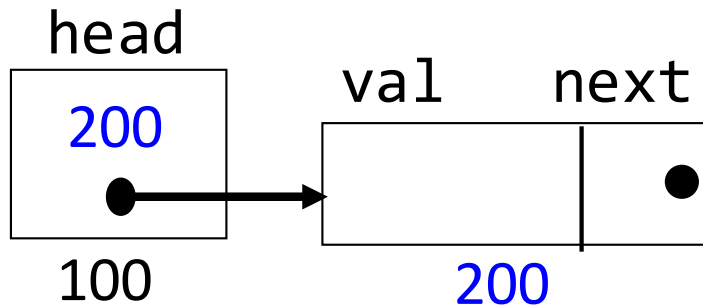
Heap



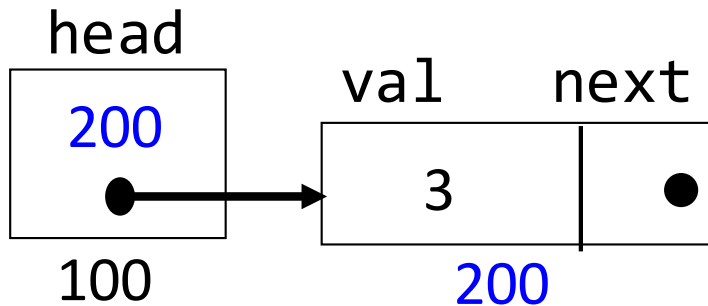
Malloc grabs this box from the heap (address of the box = 200)
That box is:



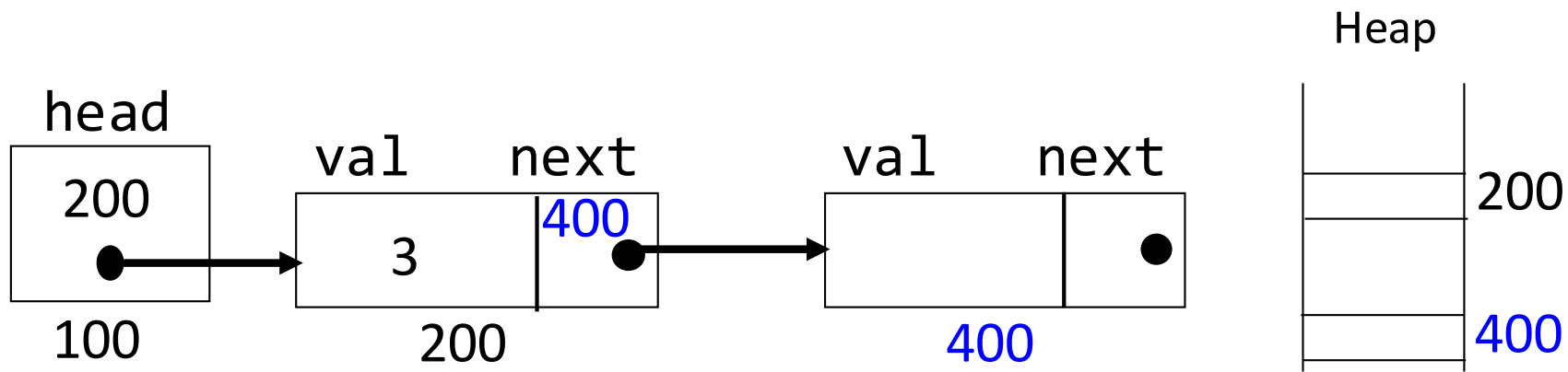
We denote it graphically as:



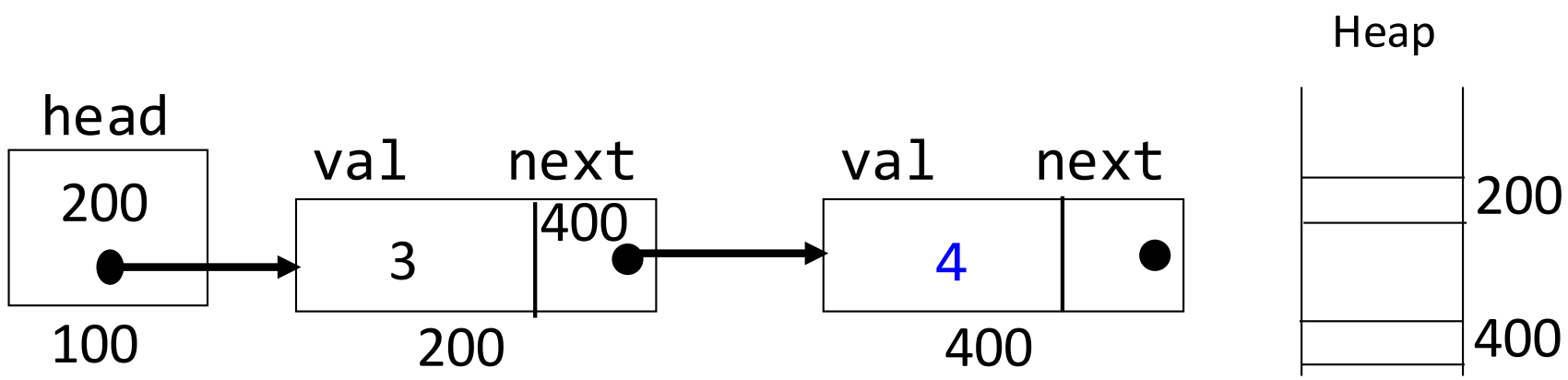
`head->val = 3;` //same as writing `(*head).val`, same as “get data at head and then get val from that data. Assign 3 to that val”. Head is 200. Data at 200 consists of (val, next) pair.



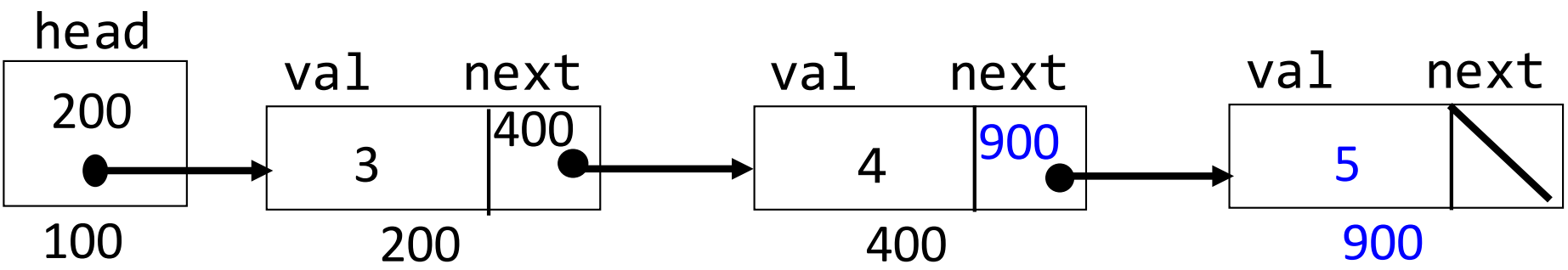
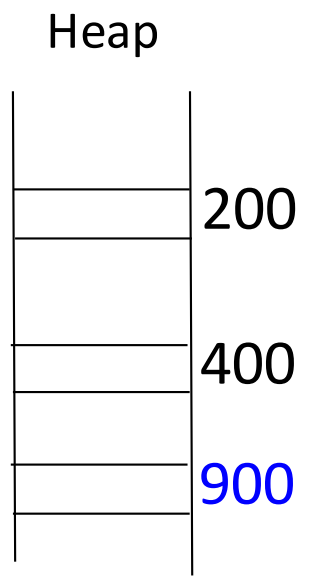
```
head->next=malloc(sizeof(Node)); //malloc returns box 400
```



```
head->next->val=4;
```

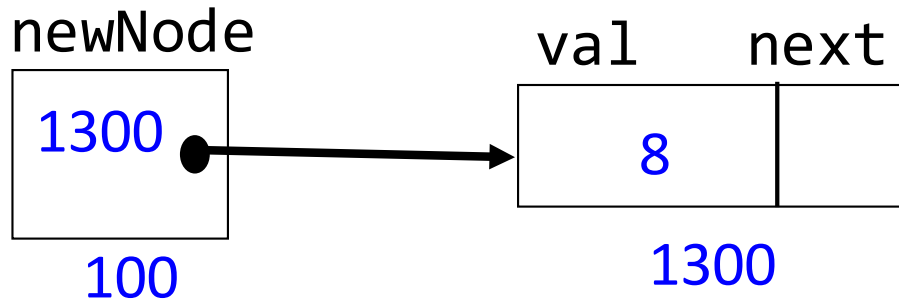


```
head->next->next=malloc(sizeof(Node));  
  
head->next->next->val=5;  
  
head->next->next->next=NULL;
```

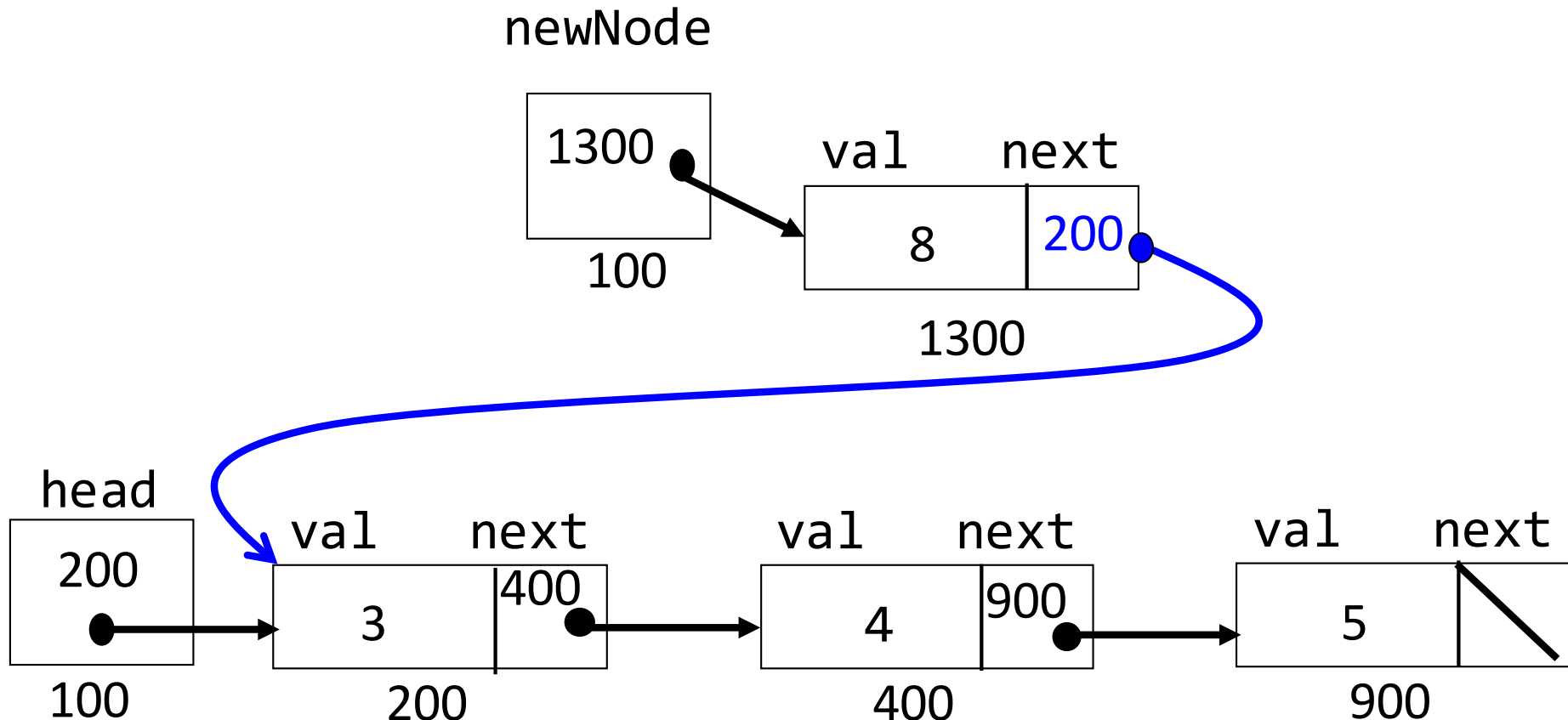


Insert at beginning (Analysis)

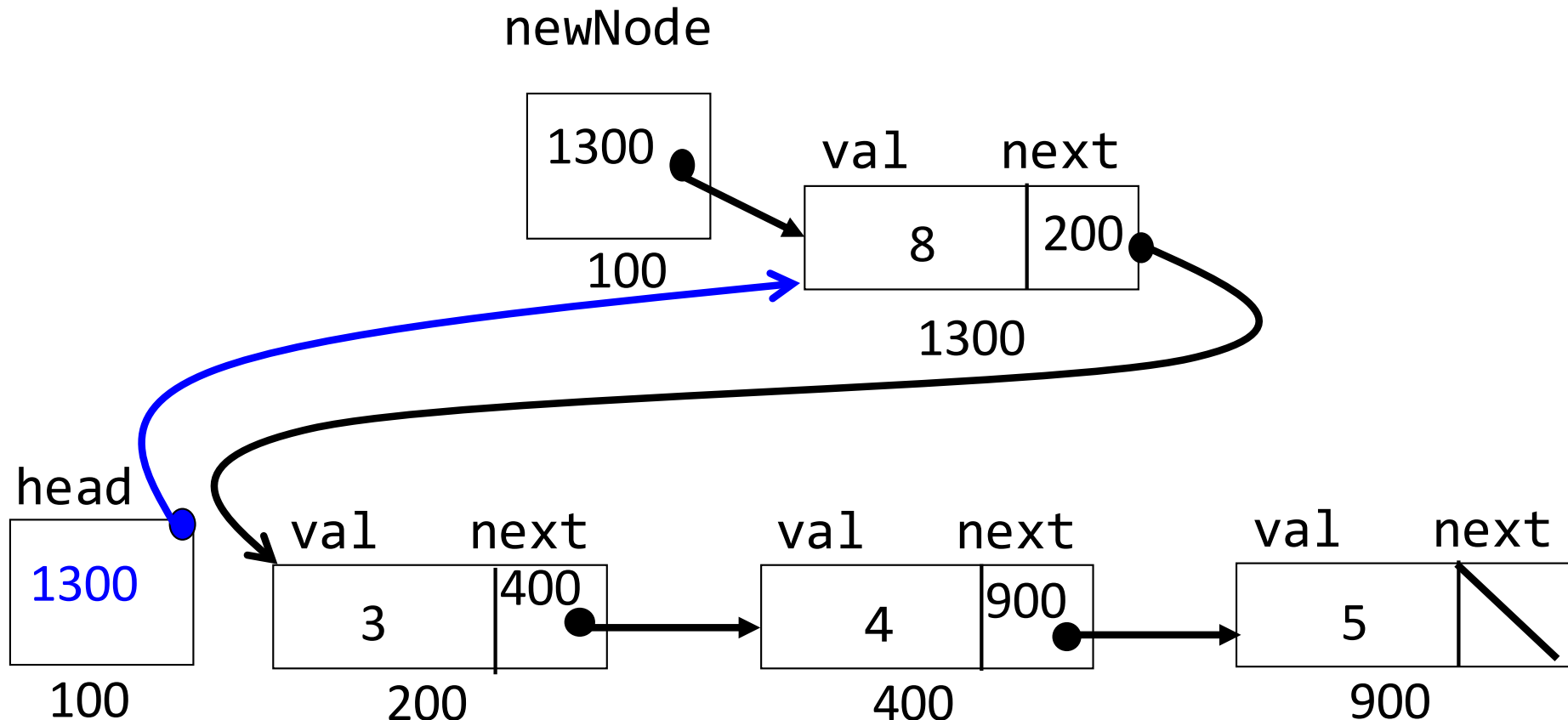
```
Node* newNode = malloc(sizeof(Node));  
newNode->val=8;  
newNode->next = head;  
head = newNode;
```



```
Node* newNode = malloc(sizeof(Node));  
newNode->val=8;  
newNode->next = head; //current value of head is 200  
head = newNode;
```

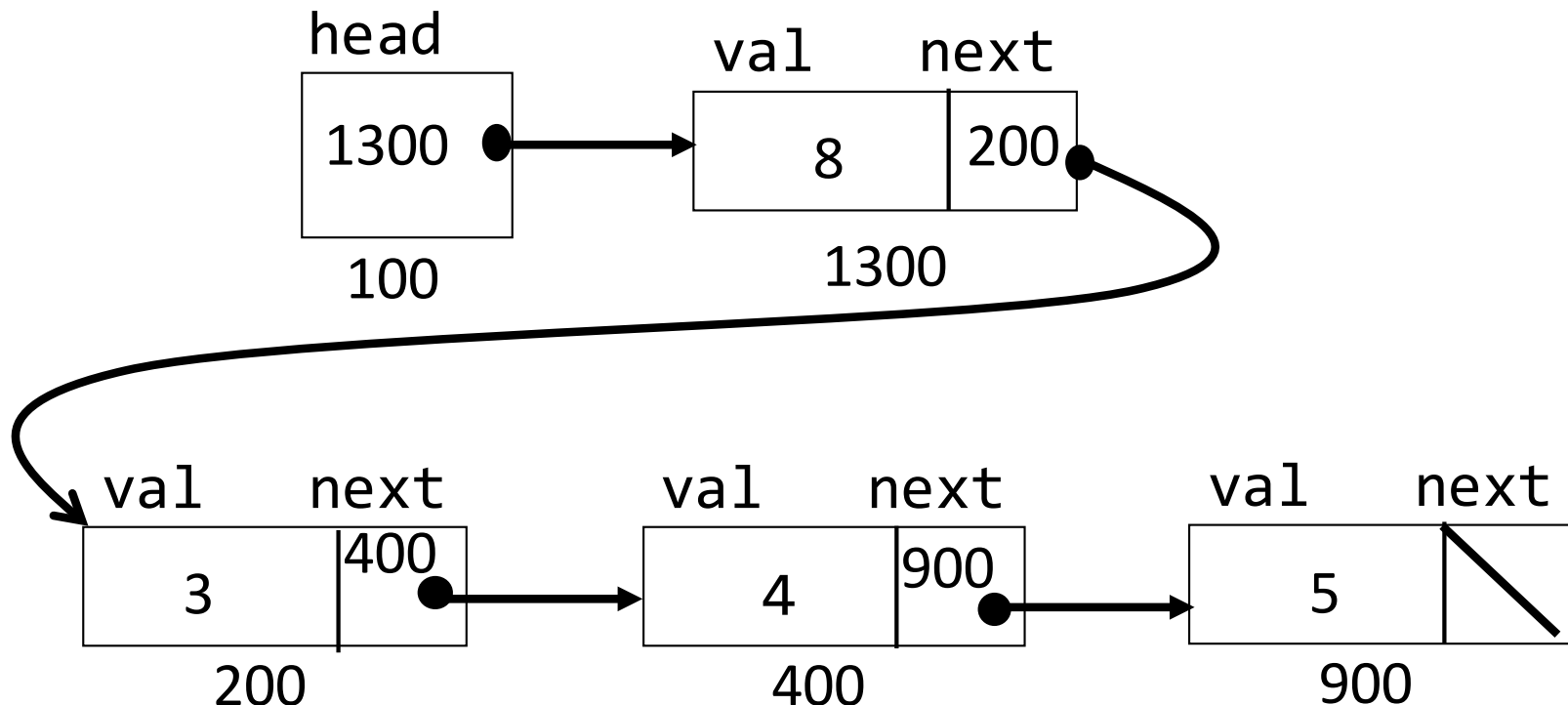


```
Node* newNode = malloc(sizeof(Node));  
newNode->val=8;  
newNode->next = head;  
head = newNode; //head is now 1300
```



The `insert` function:

- we should be able to insert anywhere in the list
- if we expect `insert` to change head, then must pass the address of head



//insert anywhere

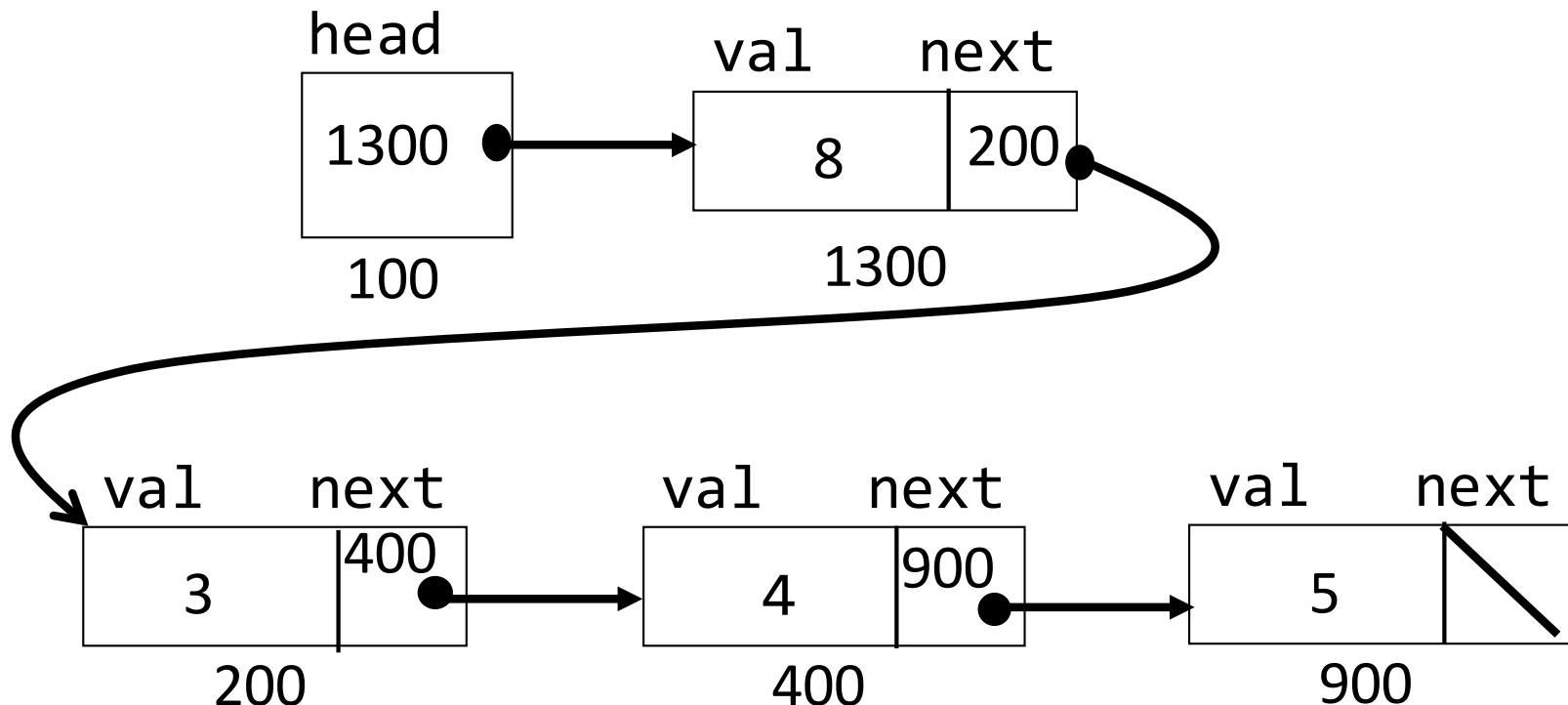
```
void insert(Node** loc, int val) {  
    Node* newNode = malloc(sizeof(Node));  
    newNode->val=val;  
    newNode->next = *loc;  
    *loc = newNode;  
}  
insert(&head, 10);
```

//insert at the beginning

```
Node* newNode = malloc(sizeof(Node));  
newNode->val=8;  
newNode->next = head;  
head = newNode;
```

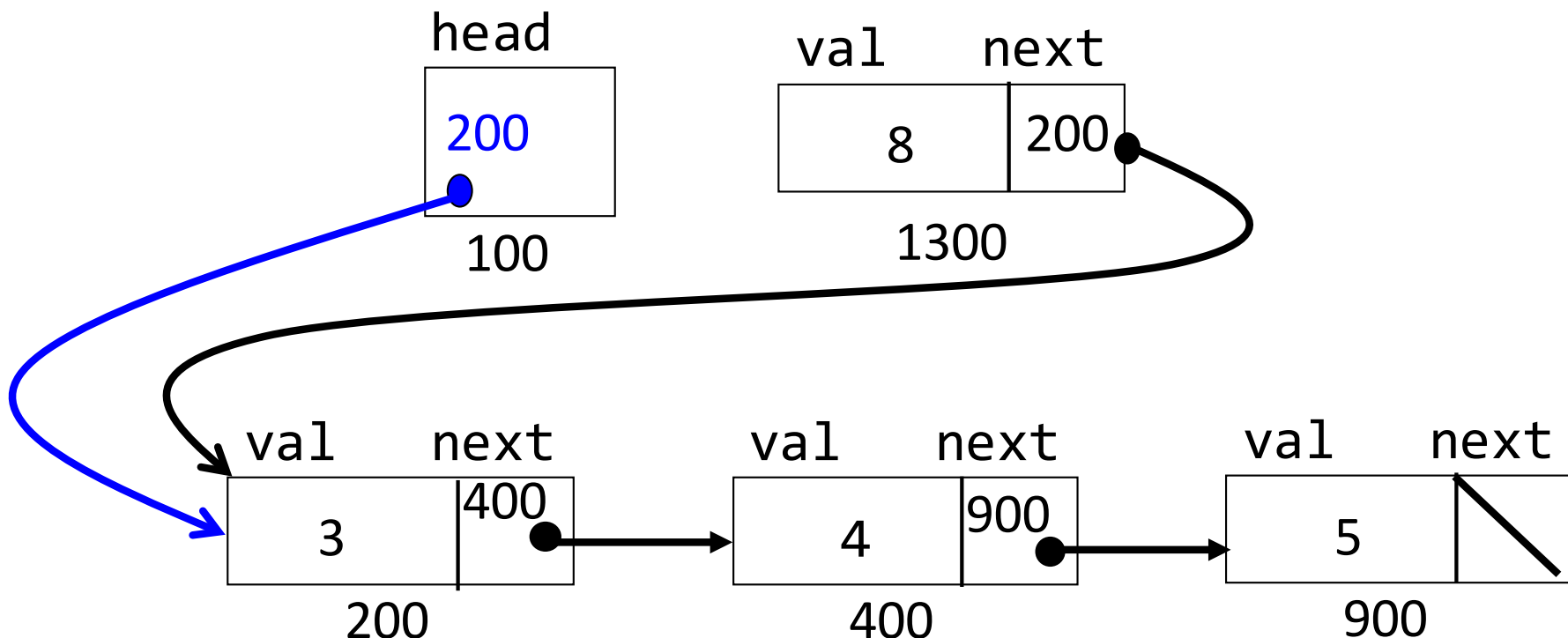
Linked Lists (removing from list)

`head = head->next; //what is the problem here?`



Linked Lists (removing from list)

`head = head->next; //what is the problem here?`

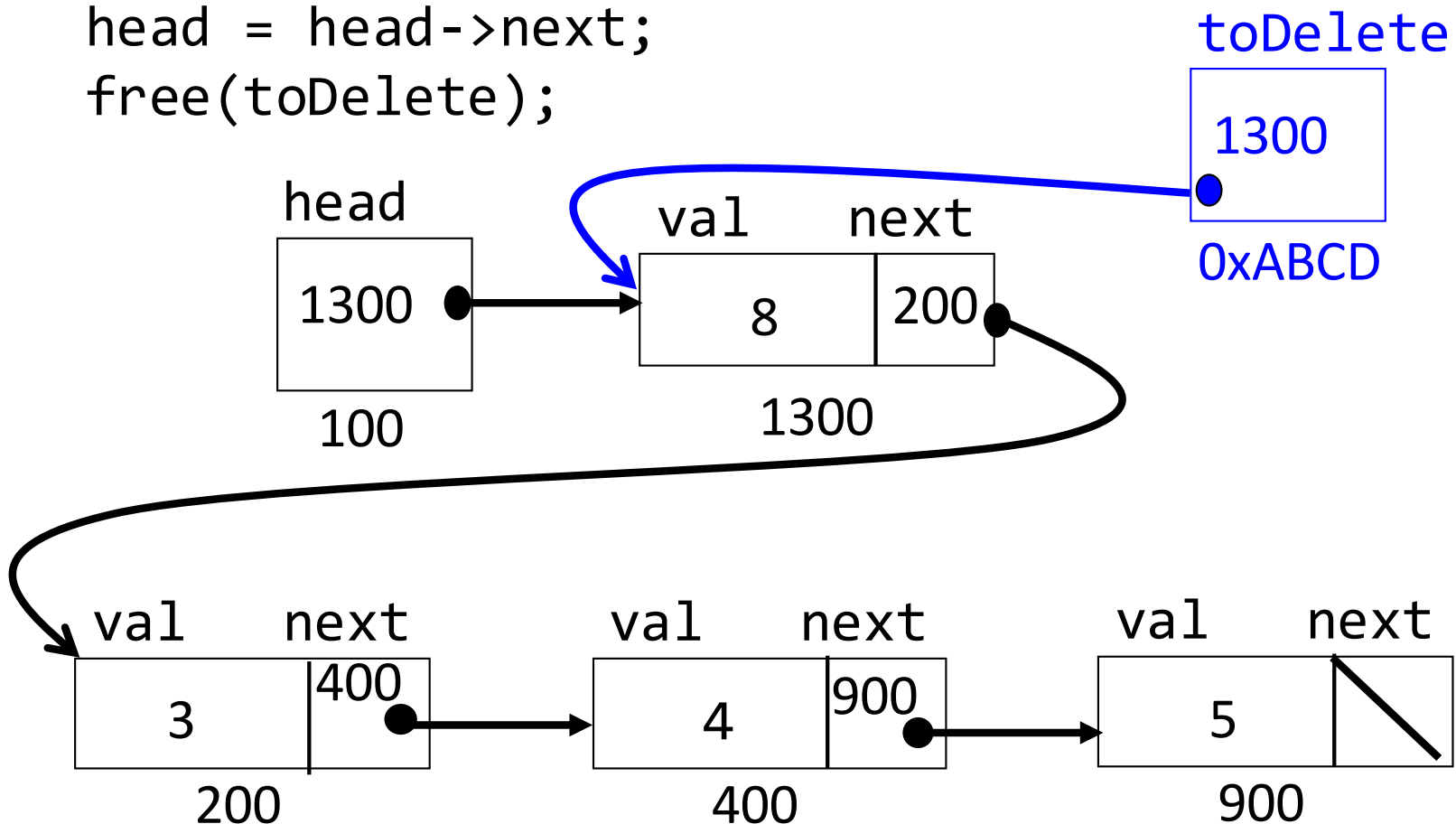


Removing from list (Correct)

```
Node* toDelete = head;
```

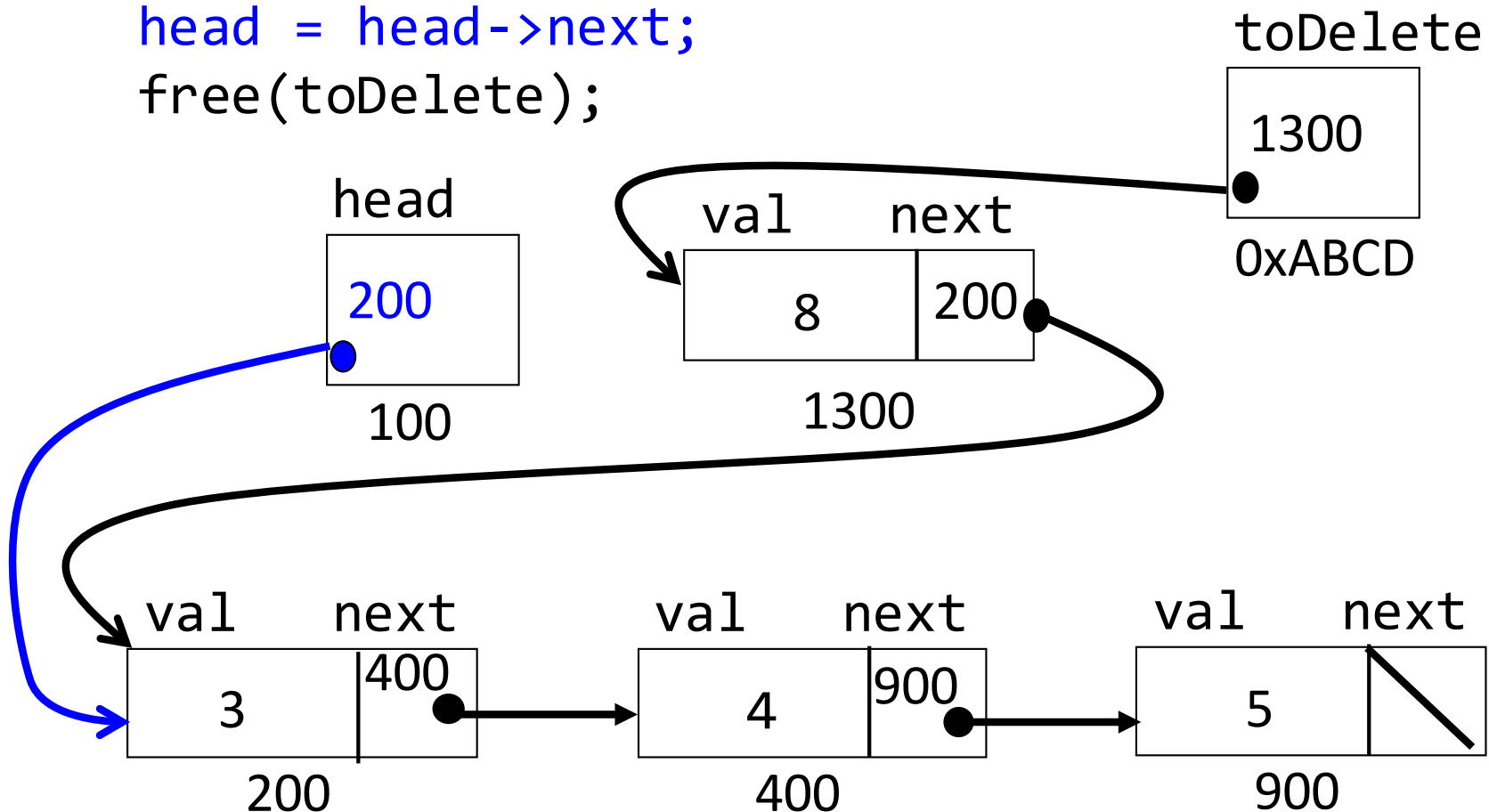
```
head = head->next;
```

```
free(toDelete);
```



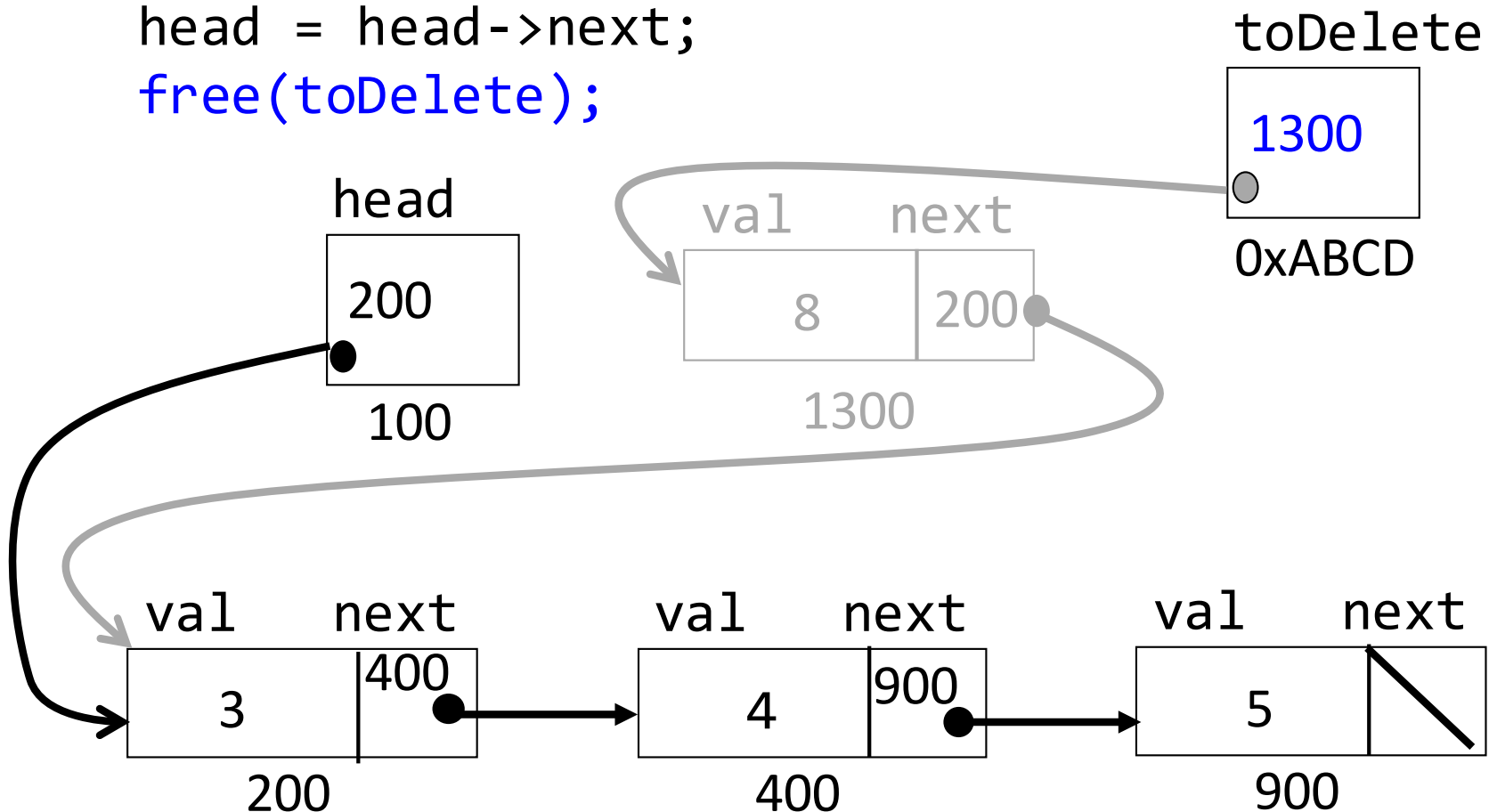
Removing from list (Correct)

```
Node* toDelete = head;  
head = head->next;  
free(toDelete);
```



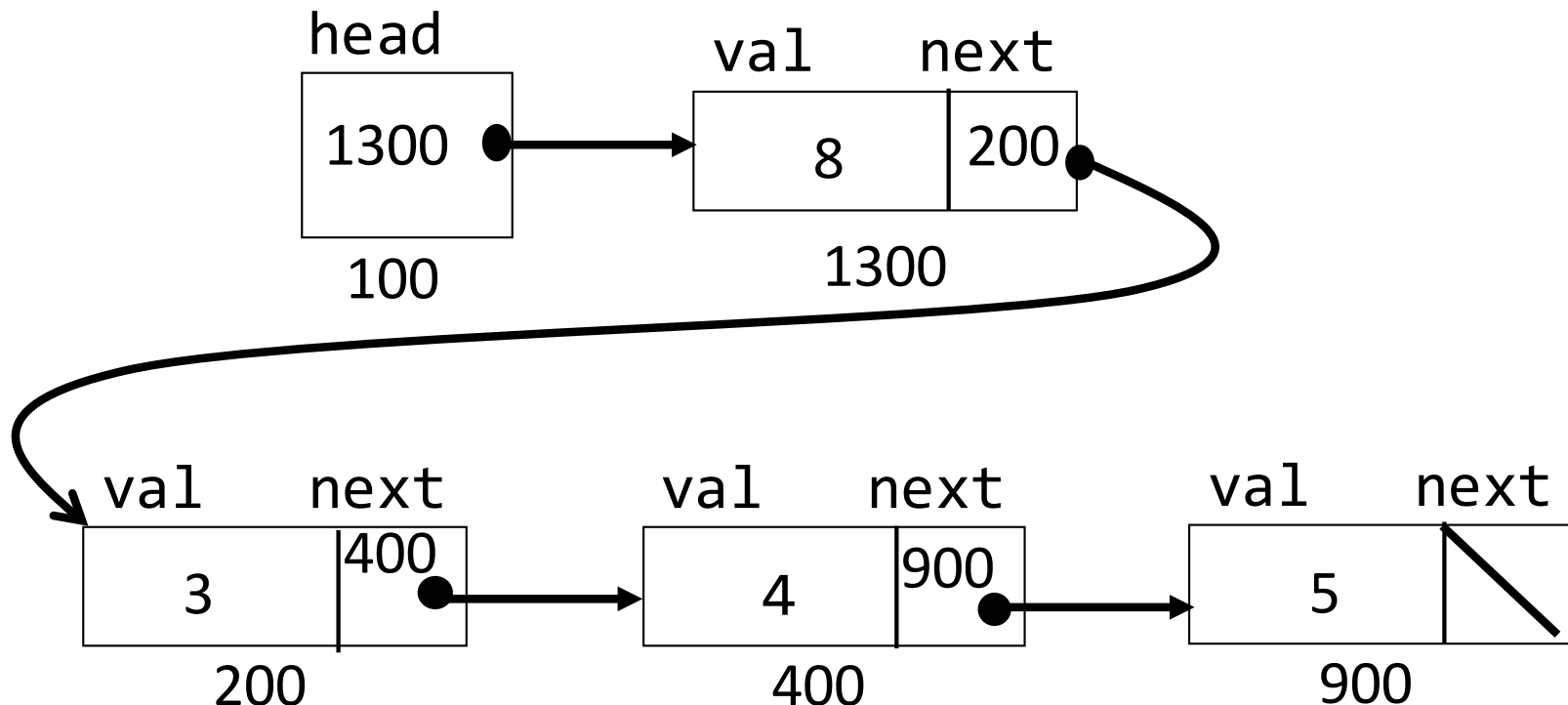
Removing from list (Correct)

```
Node* toDelete = head;  
head = head->next;  
free(toDelete);
```



The remove function:

- we should be able to remove any node in the list
- if we expect remove to change head, then must pass the address of head



//removes the first node (node pointed to by head)

```
Node* toDelete = head;
```

```
head = head->next;
```

```
free(toDelete);
```

//remove from anywhere in the list

```
void remove(Node** loc) {
```

```
    Node* toDelete = *loc;
```

```
    *loc = (*loc)->next;
```

```
    free(toDelete);
```

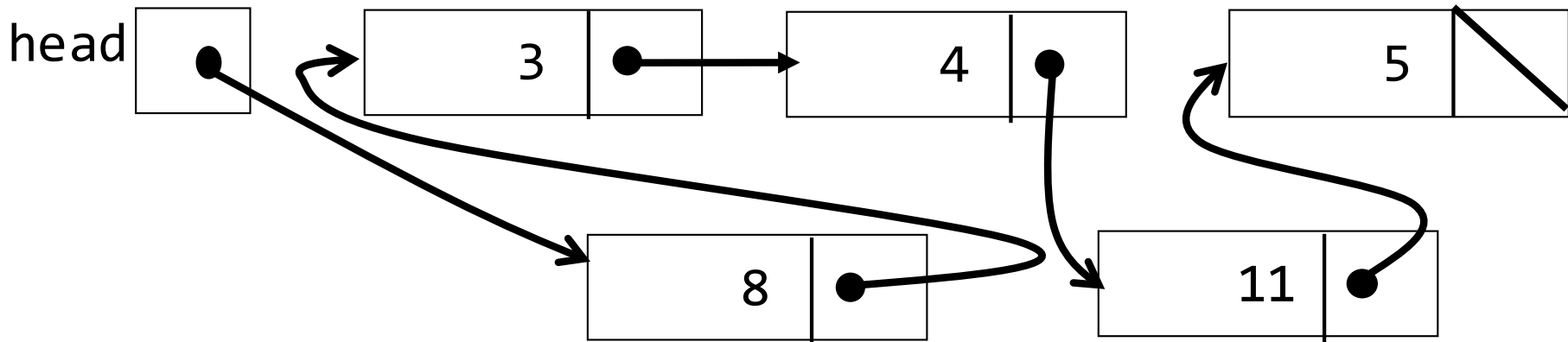
```
}
```

```
remove(&head) //removes the first node (node that head  
points to)
```

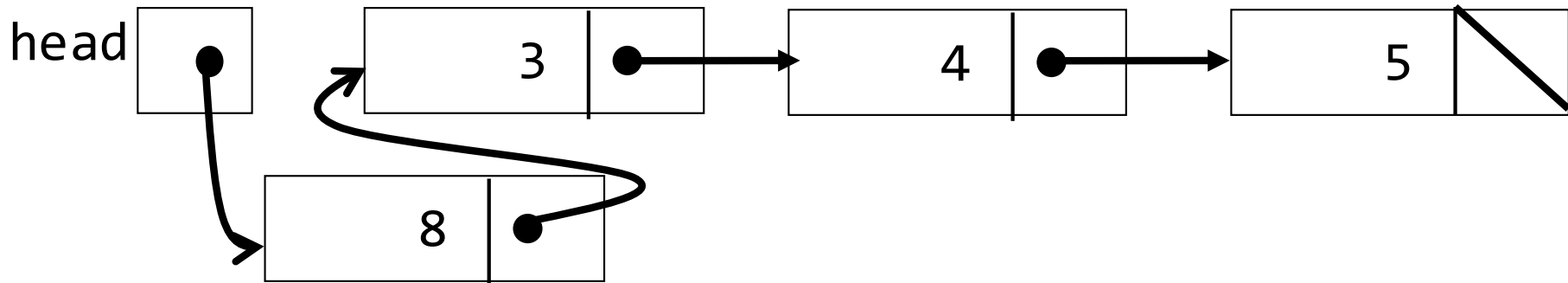
```
remove(&cur) //removes the node that cur points to
```

Linked Lists (removing from list)

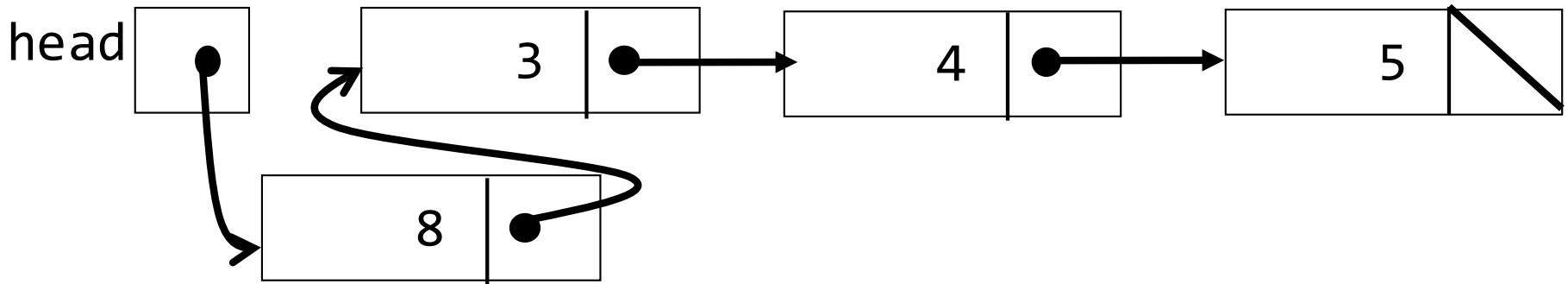
Example:



```
Node* del = head->next->next; //points to 4  
remove(&(del->next)); //removes 11
```



Exercise (removing from list)



```
void remove(Node** loc) {  
    Node* toDelete = *loc;  
    *loc = (*loc)->next;  
    free(toDelete);  
}
```

```
Node* del = head->next->next->next; //points to 5  
remove(&(del->next)); //What happens here?
```

Linked Lists (searching a list)

1. Step through each node in the list
2. Check if val in each node matches key
3. Return true if matches, false if no match found till end of list

```
bool contains(Node* head, int key) {  
    Node* cur = head;  
    while(cur != NULL) {  
        if(cur->val == key) return true;  
        cur = cur->next;  
    }  
    return false;  
}
```

List manipulation

- Given a key, delete that node:
- 1) Find the key. If found return the address of the next pointer that points to that node. If not, return address of last next pointer

```
Node** findEq(Node** loc, int key) {  
    while((*loc) != NULL) {  
        if((*loc)->val == key) return loc;  
        loc = &((*loc)->next);  
    }  
    return loc;  
}
```

List manipulation contd..

2) pass the returned address from findEq to remove.

```
Node** toRemove = findEq(&head, 8);  
if((*toRemove) != NULL)  
    remove(toRemove);
```

Linked List vs. Array

Linked List	Array
Sequential Access	Random Access
Capacity manipulation is easy	Impossible once defined
Always on heap	Stack or heap
E.g. a paper roll	E.g. A book

Stacks

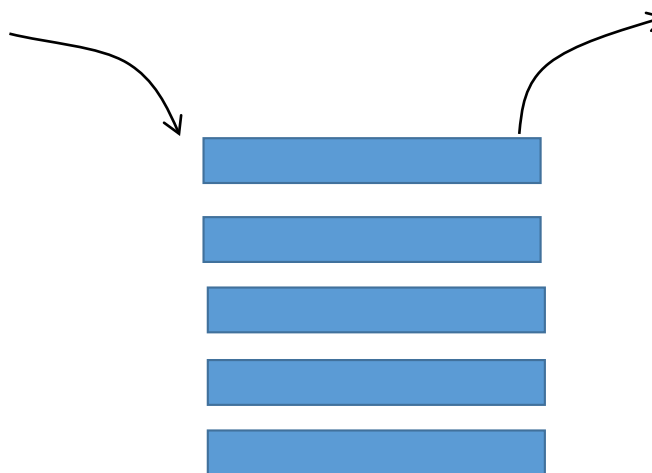
- Sequential data structure
- Objects are inserted and removed according to LIFO (last-in first-out) principle
- Operations allowed (push, pop, top, empty, full)
- Recursive definition:
 - Either empty OR
 - There is top and rest is a stack

Stacks

- Examples:
 - Support for recursion (call stack)
 - Undoing operation in text editors
 - Parsing of arithmetic expressions ($(2+(3+4)*5)/2*((10+6)*8)$)

Insert / Push

Remove / Pop



Stacks – Implementation

- Multiple implementations possible
 - Array-based, Linked-list based
- Permissible operations - push, pop (for data access), isEmpty, isFull (for error check)
 - Further, some implementations define constant time operations for push and pop

Stacks – Implementation

- Array-based
 - Array of MAX elements

```
typedef struct{  
    int vals[MAX];  
    int top;  
}stack;
```

```
stack s={.top=0}; //defines a stack that can hold a  
maximum of MAX elements
```

Stacks – Implementation

- push

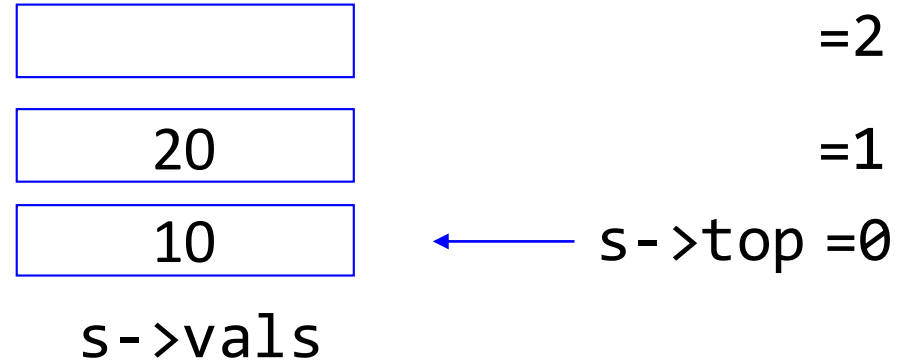
```
int push(stack* s, int val) {  
    if(isFull(s))  
        return ERROR;  
    s->vals[s->top] = val;  
    (s->top)++;  
    return SUCCESS;  
}
```

Stacks – Implementation

```
#define MAX 3
```

```
push(&s, 10)
```

```
push(&s, 20)
```



Stacks – Implementation

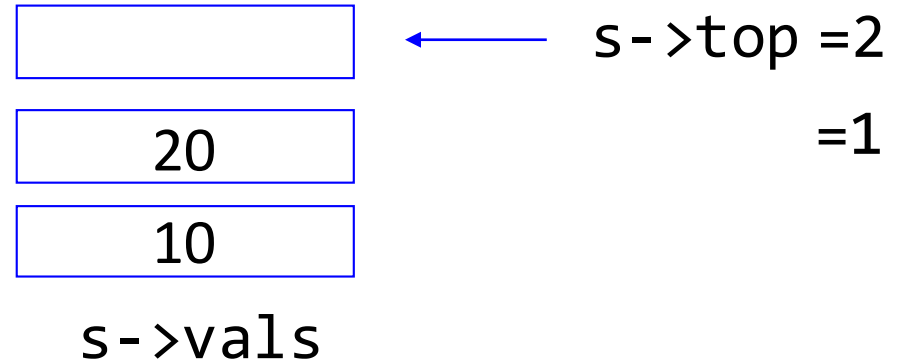
- pop

```
int pop(stack* s) {  
    if(isEmpty(s))  
        return ERROR;  
    (s->top)--;  
    return s->vals[s->top];  
}
```

Stacks – Implementation

```
#define MAX 3
```

```
pop(&s);
```



Stacks – Implementation

- isEmpty and isFull

```
bool isEmpty(stack* s) {  
    if(s->top==0)  
        return true;  
    return false;  
}
```

```
bool isFull(stack* s) {  
    if(s->top==MAX)  
        return true;  
    return false;  
}
```

//if returns true called underflow

//if returns true called overflow

Stacks – Implementation

- Linked-List based

```
typedef struct Node{  
int val;  
struct Node* next;  
}Node;
```

```
typedef struct{  
Node* head;  
}stack;
```

```
stack s={.head=NULL}; //defines a stack that contains  
no elements in the linked list
```

Stacks – Implementation

- **push** (always insert at the beginning of the list. Never implement a loop – e.g. inserting at the end of list)

```
int push(stack* s, int val) {
    Node* newNode = malloc(sizeof(Node));
    if(newNode == NULL)
        return ERROR; //overflow
    newNode->val = val;
    newNode->next = s->head;
    s->head=newNode;
    return SUCCESS;
}
```

Stacks – Implementation

- pop

```
int pop(stack* s) {  
    if(s->head == NULL)  
        return ERROR; //underflow  
    Node* toDelete = s->head;  
    s->head = s->head->next;  
    int retVal = toDelete->val;  
    free(toDelete);  
    return retVal;  
}
```

Stacks – Implementation

- isEmpty and isFull

```
bool isEmpty(stack* s) {  
    if(s->head==NULL)  
        return true;  
    return false;  
}
```

```
bool isFull(stack* s) {  
    ?  
}
```

//if returns true called underflow

//if returns true called overflow

Stack vs. Linked List

Linked List	Stack
Sequential Access	Sequential Access
Capacity manipulation is easy	Depends on implementation
Always on heap	Stack or heap
Can insert and remove from anywhere	Only at the top
E.g. a paper roll	E.g. Stack of books

Application of Stacks (evaluating arithmetic expressions)

- Needed while parsing arithmetic expressions
- Two stages:
 1. Break the expression and translate it
 - infix to postfix
 2. Parse the translated expression
 - Parse postfix expression

Evaluating arithmetic expressions

- Infix expression: $2 + 3$
 - The operator '+' appears in between operand.
 - Intuitive and is most commonly used.
 - Can be ambiguous if the operator precedence is not known.
 - PEMDAS (parenthesis, exponentiation, multiplication, division, addition, subtraction) for precedence
- Postfix expression: $23+$
 - The operator follows operands
 - Free of parenthesis
 - Also called Reverse Polish Notation (RPN)

Evaluating arithmetic expressions

- Stage 1: Translating infix expression to postfix
 1. See an integer print it to output: `print(n)`
 2. See an operator `op`:
 - a. Push `op` onto the stack only if top of the stack has an operator with lower precedence than `op` or top contains `'('`.
 - b. Otherwise, pop from the stack until an operator on top of the stack has lower precedence than `op`. Then push `op`. If you encounter a `'('`, stop and push `op`. Print popped symbols.
 3. See a `'('`, push it on to the stack
 4. See a `') '` pop from the stack until a `'('` seen (including the `'('` symbol). Print popped symbols, discard `'('`.
 5. No more input symbols, pop everything from the stack and print.

Evaluating arithmetic expressions

- Stage 1 Example: $2*(3+4)$

Symbol	Action/Rule	Output	Stack
2	Print(2) / 1	2	Empty
*	Push('*') / 2.a	2	<div> <div>*</div> </div>
(Push('(') / 3	2	<div> <div>(</div> <div>*</div> </div>
3	Print(3) / 1	23	Same
+	Push('+') / 2.b	23	<div> <div>+</div> <div>(</div> <div>*</div> </div>
4	Print(4) / 1	234	Same
)	Pop / 4	234+	<div> <div>*</div> </div>
	Pop / 5	234+*	Empty

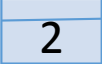


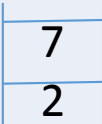
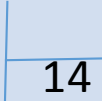
- Postfix expression: $234+*$

Evaluating arithmetic expressions

- Stage 2: Parsing postfix expression
 1. See an integer push onto stack
 2. See an operator op :
 - a. Pop twice ($A = \text{pop}()$, $B = \text{pop}()$). Evaluate $B \ op \ A$. Push result on stack.
 3. When end of expression is reached, pop and print the result

Evaluating arithmetic expressions

- Stage 2 Example: 234+*

Symbol	Action/Rule	Output	Stack
2	Push(2) / 1	None	
3	Push(3) / 1	None	
4	Push(4) / 1	None	
+	4=Pop() 3=Pop() Eval(3+4) Push(7) / 2	None	
*	7=Pop() 2=Pop() -> Eval(2*7) -> Push(14) / 2	None	
	Pop(4) / 3	14	Empty

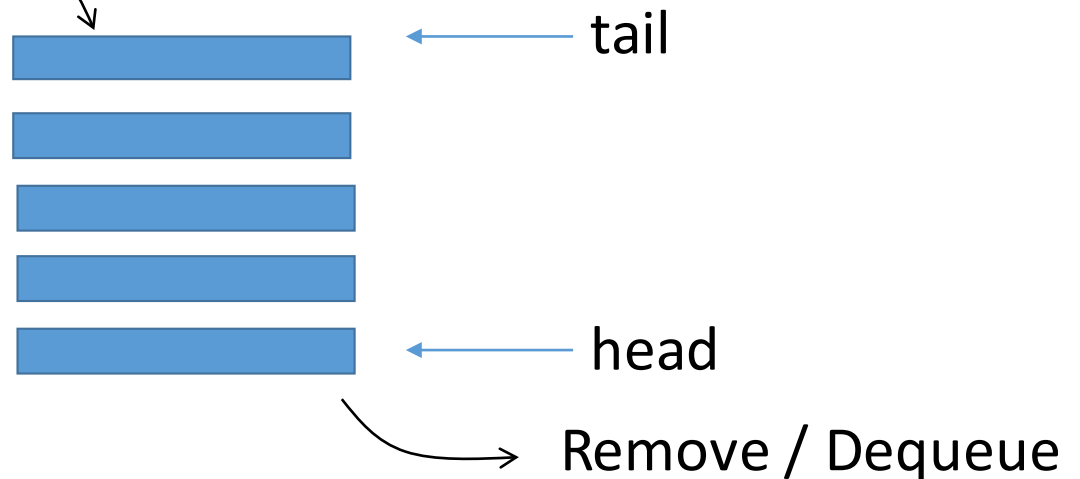
Queues

- Sequential data structure
- Objects are inserted and removed according to FIFO (first-in first-out) principle
- Operations allowed (enqueue, deque, empty, full)

Queues

- Examples:
 - Line in the food court
 - Job scheduling in computer science

Insert / Enqueue



Queues – Implementation

- Multiple implementations possible
 - Array-based (requires resizing), Linked-list based (most commonly used)
- Permissible operations - enqueue, dequeue (for data access), isEmpty, isFull (for error check)

Queues – Implementation

- Linked-List based

```
typedef struct Node{  
int val;  
struct Node* next;  
}Node;
```

```
typedef struct{  
Node* head, *tail;  
}queue;
```

```
queue q={.head=NULL, .tail=NULL}; //defines a queue  
that contains no elements in the linked list
```

Queues – Implementation

- enqueue

```
int enqueue(queue* q, int val) {
    Node* newNode = malloc(sizeof(Node));
    if(newNode == NULL) return ERROR;
    newNode->val = val; //initialize newNode
    newNode->next = NULL;
    if(q->tail == NULL) //initialize head and tail
        q->head = q->tail = newNode;
    else {
        q->tail->next = newNode;
        q->tail = newNode;
    }
    return SUCCESS;
}
```


Queues – Implementation

- dequeue

```
int dequeue(queue* q) {  
    if(q->head == NULL)  
        return ERROR; //empty queue  
    newNode->val = val; //initialize newNode  
    Node* toDelete = q->front;  
    q->front = q->front->next;  
    free(toDelete);  
    if(q->front == NULL)  
        q->head = q->tail = NULL;  
    return SUCCESS;  
}
```

Queues – Implementation

- isEmpty

```
bool isEmpty(queue* s) {  
    if(q->head==NULL)  
        return true;  
    return false;  
}
```