

# CS406: Compilers

Spring 2022

Week 2: Overview (winding up), Scanners

# Design Considerations

- Compiler and programming language designs influence each other
  - Higher level languages are harder to compile
    - More work to bridge the gap between language and assembly
  - Flexible languages are often harder to compile
    - Dynamic typing (Ruby, Python) makes a language very flexible, but it is hard for a compiler to catch errors (in fact, many simply won't)
  - Influenced by architectures
    - RISC vs. CISC

# Programming Languages and Real-world Usage

- Why are there so many programming languages?
- Why are there new languages?
- What is a good programming language?

# Programming Languages and Real-world Usage

- Why are there so many programming languages?
  - Distinct often conflicting requirements of the application domain

Scientific Computing	Floating-Point Arithmetic, Parallelism Support, Array Manipulation	FORTRAN
Business Applications	No data loss (persistence), Reporting capabilities, Data analysis tools	SQL
Systems Programming	Fine-grained control of system resources, real-time constraints	C/C++

# Programming Languages and Real-world Usage

- Why are there new languages?
  - To fill a technology gap
    - E.g. arrival of Web and Java
    - Java's design closely resembled that of C++

*Training a programmer on a new programming language is a dominant cost*

- Widely-used languages are slow to change
- Easy to start a new language

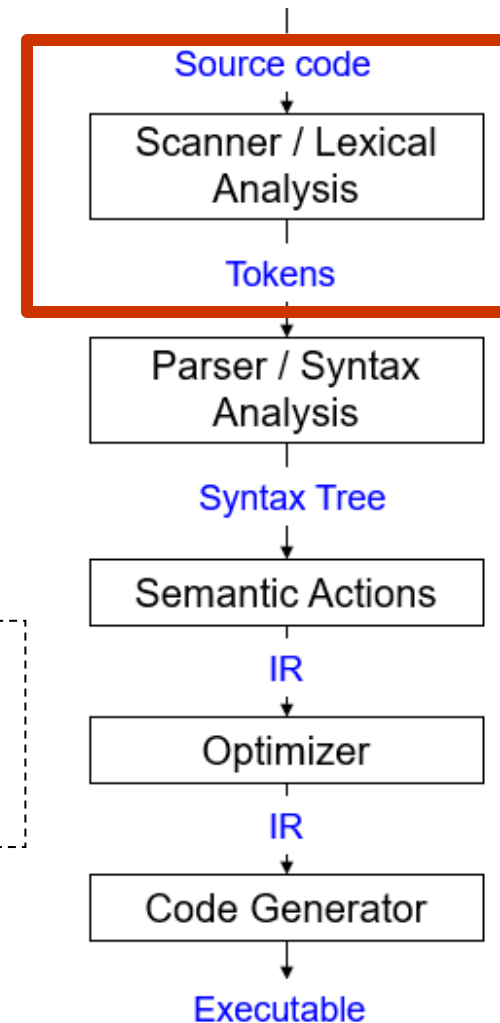
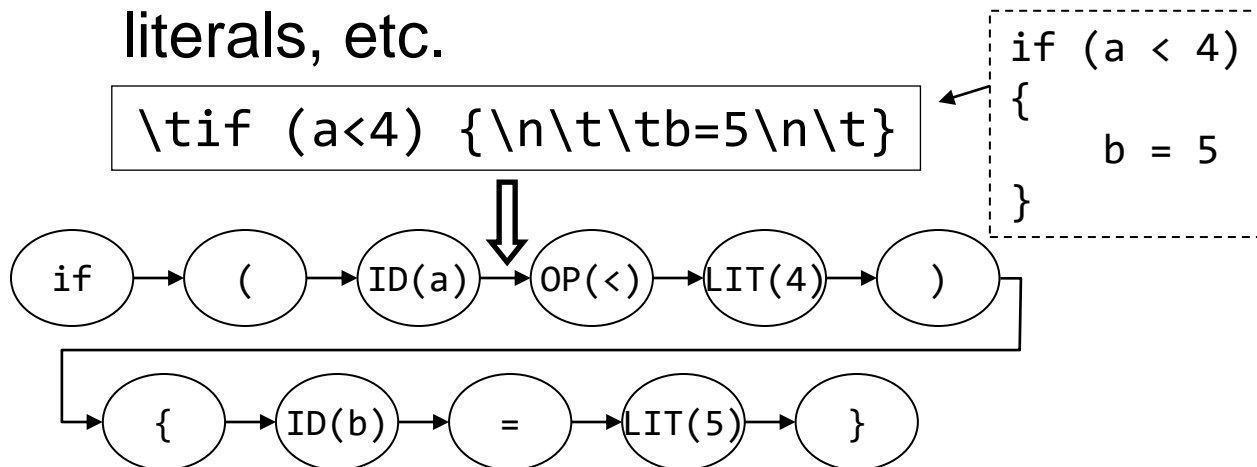
# Programming Languages and Real-world Usage

- What is a good Programming Language?

*No universally accepted argument*

# Scanner - Overview

- Also called lexers / lexical analyzers
- Recall: scanners
  - See program text as a stream of letters
  - break input stream up into a set of tokens: Identifiers, reserved words, literals, etc.



# Scanner - Motivation

- Why have a separate scanner when you can combine this with syntax analyzer (parser)?
  - Simplicity of design
    - E.g. rid parser of handling whitespaces
  - Improve compiler efficiency
    - E.g. sophisticated buffering algorithms for reading input
  - Improve compiler portability
    - E.g. handling ^M character in Linux (CR+LF in Windows)



# Scanner - Tasks

1. Divide the program text into *substrings* or *lexemes*
  - place dividers
2. Identify the *class* of the substring identified
  - Examples: Identifiers, keywords, operators, etc.
    - Identifier – *strings of letters or digits starting with a letter*
    - Integer – *non-empty string of digits*
    - Keyword – *“if”, “else”, “for”* etc.
    - Blankspace - *\t, \n, ‘ ‘*
    - Operator – *(, ), <, =, etc.*
  - *Observation:* substrings follow some pattern

# Categorizing a Substring (English Text)

- What is the English language analogy for *class*?
  - Noun, Verb, Adjective, Article, etc.
  - In an English essay, each of these classes can have a set of strings.
  - Similarly, in a program, each class can have a set of substrings.

# Exercise

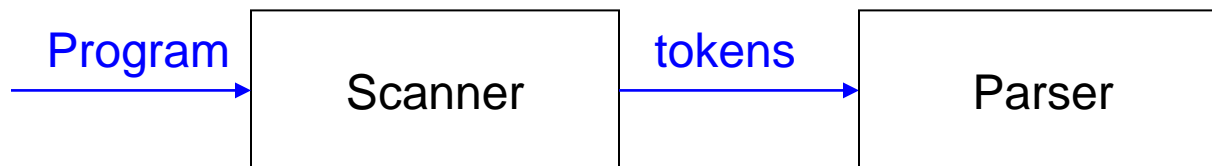

- How many tokens of class *identifier* exist in the code below?

```
for(int i=0;i<10;i++) {  
    printf("hello");  
}
```

# Scanner Output

- A token corresponding to each lexeme
  - Token is a pair: <class, value>

A string / lexeme / substring of program text



E.g. `int x = 0;`

(Keyword, "int"),  
(Identifier, "x"),  
("="),  
(Integer, "0"),  
(";")

# Scanners – interesting examples

- Fortran (white spaces are ignored)

DO 5 I = 1,25 ← DO Loop

DO 5 I = 1.25 ← Assignment statement

- PL/1 (keywords are not reserved)

DECLARE (ARG1, ARG2, . . ., ARGN);

- C++

Nested template: Quad<Square<Box>> b;

Stream input: std::cin >> bx;

# Scanners – interesting examples

- How did we go about recognizing tokens in previous examples?
  - Scan **left-to-right** till a token is identified
  - **One token at a time**: continue scanning the remaining text till the next token is identified...
  - So on...

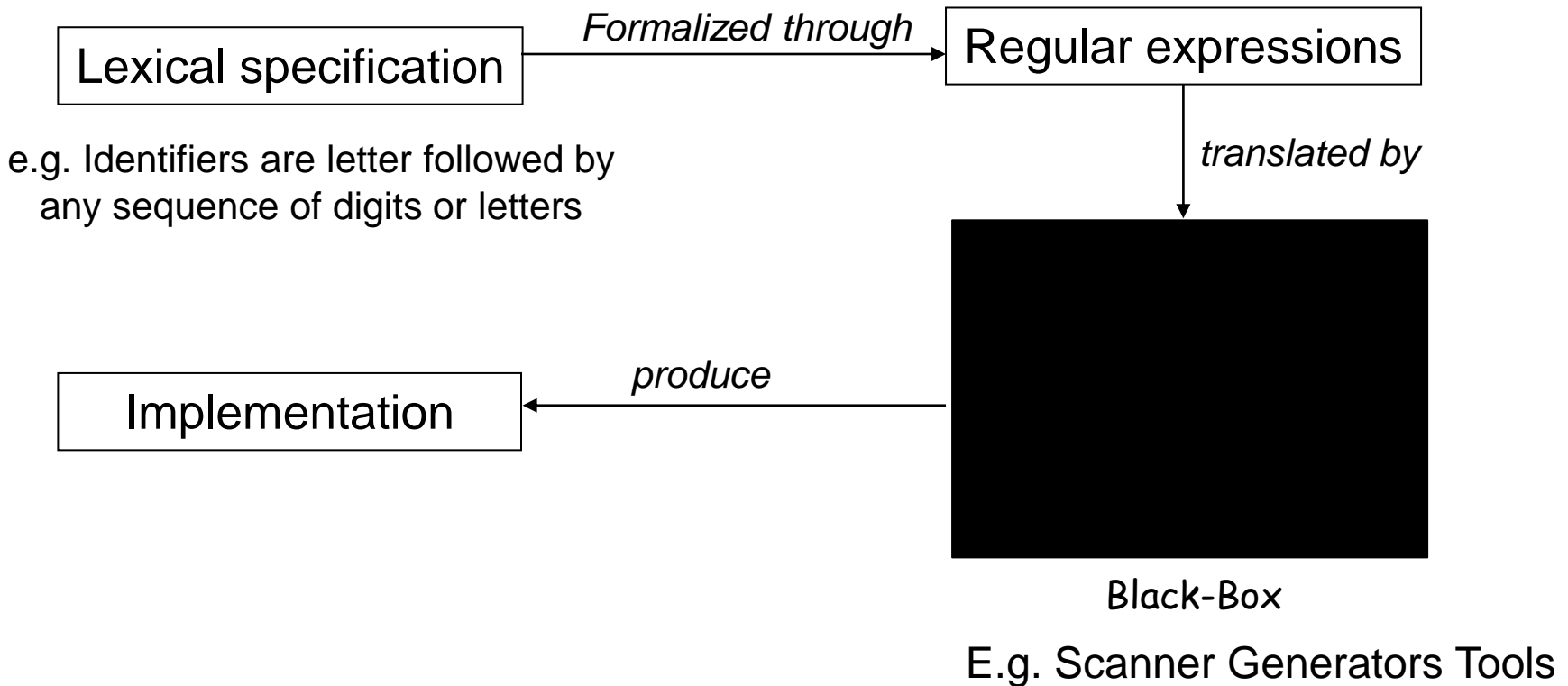
We always need to *look-ahead* to identify tokens

*....but we want to minimize the amount of look-ahead done to simplify scanner implementation*

# Scanners – what do we need to know?

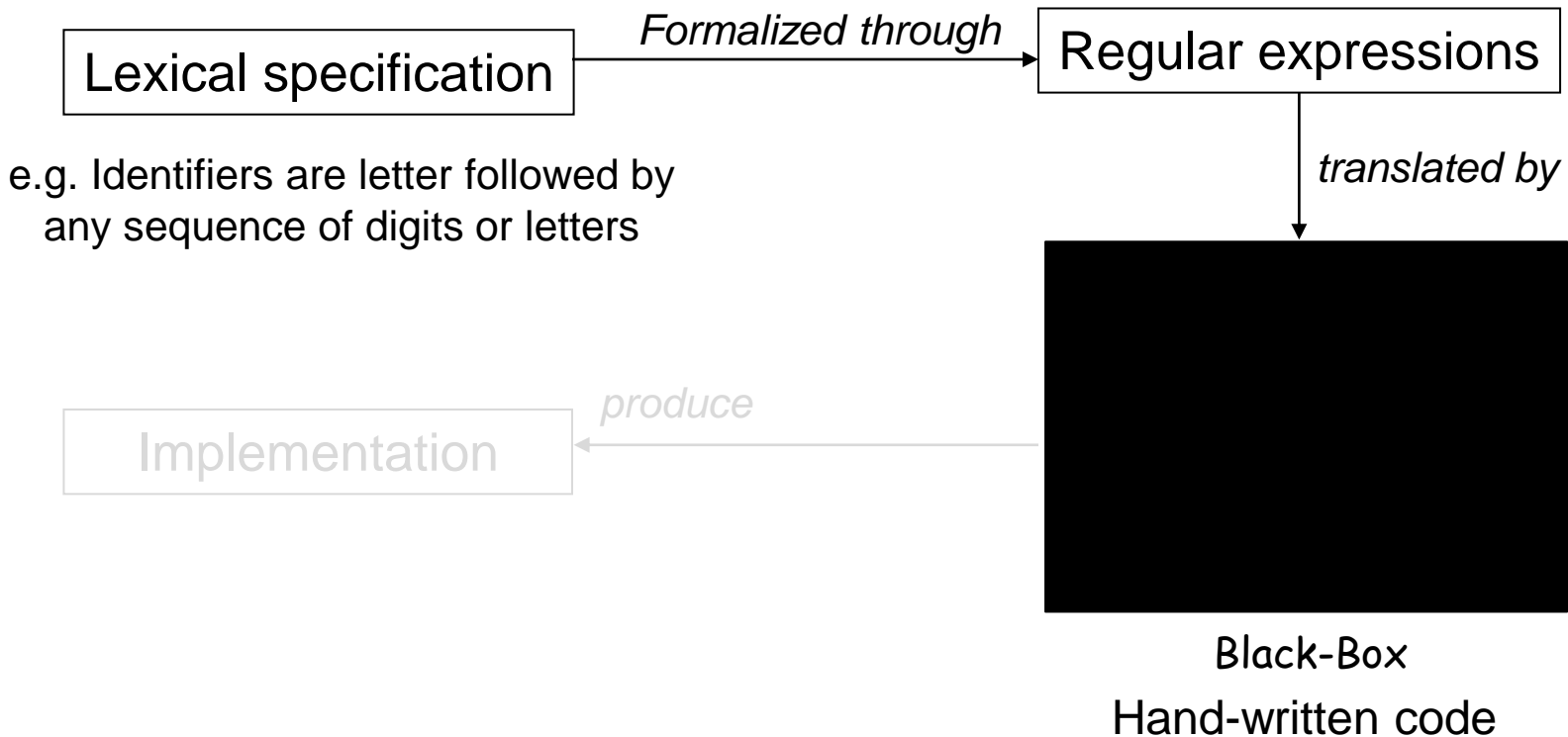
1. How do we define tokens?
  - Regular expressions
2. How do we recognize tokens?
  - build code to find a lexeme that is a prefix and that belongs to one of the classes.
3. How do we write lexers?
  - E.g. use a lexer generator tool such as Flex

# Scanner / Lexical Analyzer - flowchart





# Scanner / Lexical Analyzer - flowchart



# Scanner Generators

- Essentially, tools for converting regular expressions into scanners
  - Lex (Flex) generates C/C++ scanner program
  - ANTLR (ANother Tool for Language Recognition) generates Java program for translating program text (JFlex is a less popular option)
  - Pylexer is a Python-based lexical analyzer (**not a scanner generator**). *It just scans input, matches regexps, and tokenizes. Doesn't produce any program.*

# Exercise

- <https://forms.gle/crJ2cPYKsx3wywNe6>

# Regular Expressions

- Used to define the structure of tokens
  - Regular sets:
    - Formal:** a language that can be defined by regular expressions
    - Informal:** a set of strings defined by regular expressions
- Start with a finite character set or *Vocabulary* (V). Strings are formed using this character set with the following rules:

# Suggested Reading

- Alfred V. Aho, Monica S. Lam, Ravi Sethi and Jeffrey D. Ullman: Compilers: Principles, Techniques, and Tools, 2/E, AddisonWesley 2007
  - Chapter 3 (Sections: 3.1, 3.3, 3.6 to 3.9)
- Fisher and LeBlanc: Crafting a Compiler with C
  - Chapter 3 (Sections 3.1 to 3.4, 3.6, 3.7)