

CS601: Software Development for Scientific Computing

Autumn 2022

Week7: Motifs – Sparse Matrices (contd.),
Fourier Transforms

Last week..

- Matrix Multiplication
 - ijk variants and recursive matmul
- Efficiency considerations
 - Storage (e.g. cache-oblivious data storage using Z-ordering)
 - Communication cost (data movement cost)
 - Special hardware (FMA, Vector units)
- Motif: Sparse Matrices
 - Triangular Matmul (as an e.g. that exploits structure to accelerate computation)
 - Storage scheme for sparse matrices (e.g. CSR)
 - Banded matrices ($y = y + Ax$ with banded matrix and optimized storage)

$y=y+Ax$ with *Separable* Matrices

Refer to (Section 1 only):


<https://www.math.uci.edu/~chenlong/MathPKU/FMMsimple.pdf>

Matrix Algebra and Efficient Computation

- Pic source: the Parallel Computing Laboratory at U.C. Berkeley: A Research Agenda Based on the Berkeley View (2008)

<i>Motif</i>	Embed	Desktop	Games	DB	ML	HPC	Medicine	Music	Speech	CBIR	Browser		Motif	Desktop	Games	DB	ML	HPC	Medicine	Music	Speech	CBIR	Browser	
1 Finite State Mach.													9 N-Body											
2 Combinational													10 MapReduce											
3 Graph Traversal													11 Backtrack/B&B											
4 Structured Grid													12 Graphical Models											
5 Dense Matrix													13 Unstructured Grid											
6 Sparse Matrix													<i>Temperature Chart of Need</i>				DB = database							
7 Spectral (FFT)												<i>Hot</i>	<i>Warm</i>	<i>Med</i>	<i>Cool</i>	ML = machine learning								
8 Dynamic Prog																HPC = High Perf. Comp.								

Figure 4. Temperature Chart of the 13 Motifs. It shows their importance to each of the original six application areas and then how important each one is to the five compelling applications of Section 3.1. More details on the motifs can be found in (Asanovic, Bodik et al. 2006).

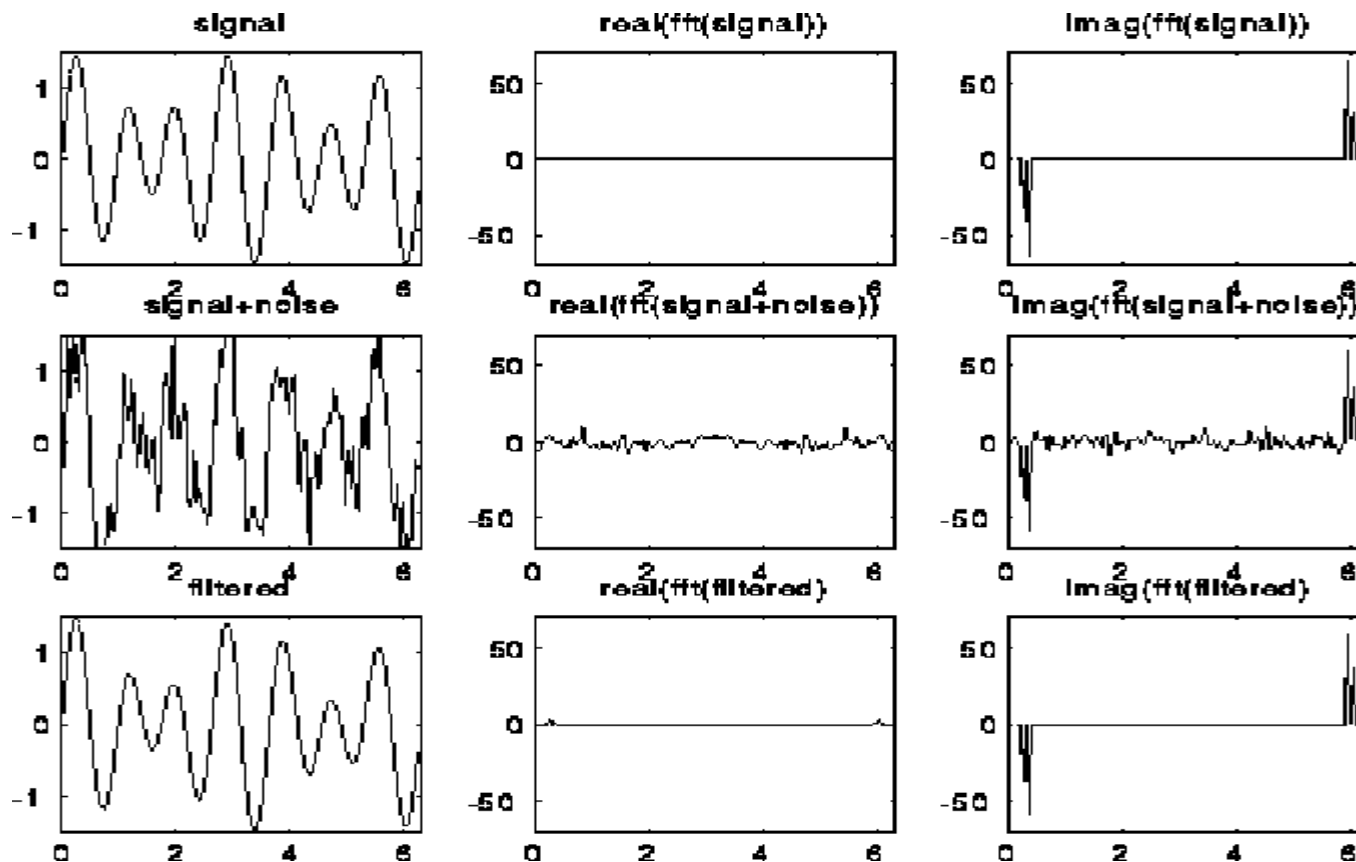
⇒ Seen earlier
 Next..

Faster $y=Ax$: Discrete Fourier Transforms (DFT)

- Very widely used
 - Image compression (jpeg)
 - Signal processing
 - Solving Poisson's Equation
- Represent A with F , a *Fourier Matrix* that has the following (remarkable) properties:
 - F^{-1} is easy to compute
 - Multiplications by F and F^{-1} is fast. (need to do $Fx=y$ and $x= F^{-1} y$ quickly)
- F has complex numbers in its entries.
 - Every entry is a power of a single number w such that $w^n=1$
 - Any entry of a Fourier matrix can be written using $f_{ij} = w^{ij}$ (row and col indices start from 0)

Using the 1D FFT for filtering

- Signal = $\sin(7t) + .5 \sin(5t)$ at 128 points
- Noise = random number bounded by .75
- Filter by zeroing out FFT components $< .25$



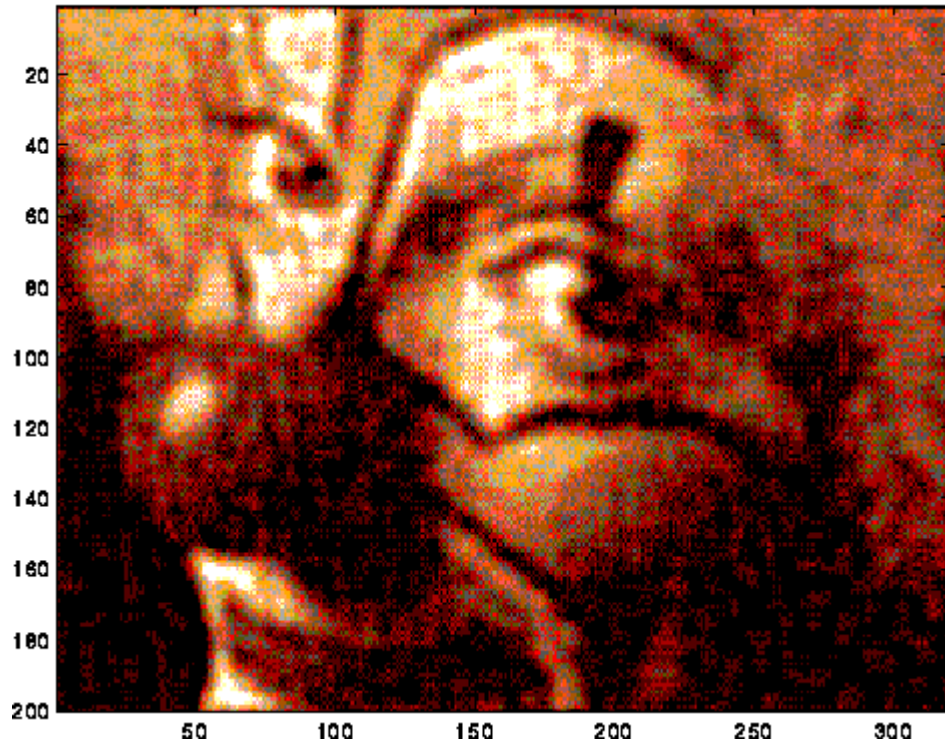
Using the 2D FFT for image compression

- Image = 200x320 matrix of values
- Compress by keeping largest 2.5% of FFT components
- Similar idea used by jpeg

Original Image



Keep only largest 2.5% of entries of 2DFFT



Examples: Fourier Matrix

- $$4 \times 4: F_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & w & w^2 & w^3 \\ 1 & w^2 & w^4 & w^6 \\ 1 & w^3 & w^6 & w^9 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & i & i^2 & i^3 \\ 1 & i^2 & i^4 & i^6 \\ 1 & i^3 & i^6 & i^9 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & w & w^2 & w^3 \\ 1 & w^2 & 1 & w^2 \\ 1 & w^3 & w^2 & w^1 \end{bmatrix}, i = \sqrt{-1}$$

– Here, $w = i$ (also denoted as $w_4 = i$). $w^4 = 1 \Rightarrow i$ is a root.

- $$8 \times 8: F_8 = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & w & w^2 & w^3 & w^4 & w^5 & w^6 & w^7 \\ 1 & w^2 & w^4 & w^6 & w^8 & w^{10} & w^{12} & w^{14} \\ 1 & w^3 & w^6 & w^9 & w^{12} & w^{15} & w^{18} & w^{21} \\ 1 & w^4 & w^8 & w^{12} & w^{16} & w^{20} & w^{24} & w^{28} \\ 1 & w^5 & w^{10} & w^{15} & w^{20} & w^{25} & w^{30} & w^{35} \\ 1 & w^6 & w^{12} & w^{18} & w^{24} & w^{30} & w^{36} & w^{42} \\ 1 & w^7 & w^{14} & w^{21} & w^{28} & w^{35} & w^{42} & w^{49} \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & w & w^2 & w^3 & w^4 & w^5 & w^6 & w^7 \\ 1 & w^2 & w^4 & w^6 & 1 & w^2 & w^4 & w^6 \\ 1 & w^3 & w^6 & w & w^4 & w^7 & w^2 & w^5 \\ 1 & w^4 & 1 & w^4 & 1 & w^4 & 1 & w^4 \\ 1 & w^5 & w^2 & w^7 & w^4 & w^1 & w^6 & w^3 \\ 1 & w^6 & w^4 & w^2 & 1 & w^6 & w^4 & w^2 \\ 1 & w^7 & w^6 & w^5 & w^4 & w^3 & w^2 & w^1 \end{bmatrix}$$

Here, $w = \frac{1+i}{\sqrt{2}}$
(= sqrt of i)

Example: Faster $y=Fx$

Column: 1 2 3 4 5 6 7 8

$$\begin{bmatrix}
 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
 1 & w & w^2 & w^3 & w^4 & w^5 & w^6 & w^7 \\
 1 & w^2 & w^4 & w^6 & 1 & w^2 & w^4 & w^6 \\
 1 & w^3 & w^6 & w & w^4 & w^7 & w^2 & w^5 \\
 1 & w^4 & 1 & w^4 & 1 & w^4 & 1 & w^4 \\
 1 & w^5 & w^2 & w^7 & w^4 & w^1 & w^6 & w^3 \\
 1 & w^6 & w^4 & w^2 & 1 & w^6 & w^4 & w^2 \\
 1 & w^7 & w^6 & w^5 & w^4 & w^3 & w^2 & w^1
 \end{bmatrix}
 =
 \begin{bmatrix}
 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
 1 & w^2 & w^4 & w^6 & w & w^3 & w^5 & w^7 \\
 1 & w^4 & 1 & w^4 & w^2 & w^6 & w^2 & w^6 \\
 1 & w^6 & w^4 & w^2 & w^3 & w & w^7 & w^5 \\
 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 \\
 1 & w^2 & w^4 & w^6 & -w & -w^3 & -w^5 & -w^7 \\
 1 & w^4 & 1 & w^4 & -w^2 & -w^6 & -w^2 & -w^6 \\
 1 & w^6 & w^4 & w^2 & -w^3 & -w & -w^7 & -w^5
 \end{bmatrix}$$

\uparrow
 (Writing columns 1,3,5,7 first and then columns 2,4,6,8)

Example: Faster $y=Fx$

$$\bullet \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & w & w^2 & w^3 & w^4 & w^5 & w^6 & w^7 \\ 1 & w^2 & w^4 & w^6 & 1 & w^2 & w^4 & w^6 \\ 1 & w^3 & w^6 & w & w^4 & w^7 & w^2 & w^5 \\ 1 & w^4 & 1 & w^4 & 1 & w^4 & 1 & w^4 \\ 1 & w^5 & w^2 & w^7 & w^4 & w^1 & w^6 & w^3 \\ 1 & w^6 & w^4 & w^2 & 1 & w^6 & w^4 & w^2 \\ 1 & w^7 & w^6 & w^5 & w^4 & w^3 & w^2 & w^1 \end{bmatrix} = \left[\begin{array}{cccc|cccc} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & \omega^2 & \omega^4 & \omega^6 & \omega & \omega^3 & \omega^5 & \omega^7 \\ 1 & \omega^4 & 1 & \omega^4 & \omega^2 & \omega^6 & \omega^2 & \omega^6 \\ 1 & \omega^6 & \omega^4 & \omega^2 & \omega^3 & \omega & \omega^7 & \omega^5 \\ \hline 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 \\ 1 & \omega^2 & \omega^4 & \omega^6 & -\omega & -\omega^3 & -\omega^5 & -\omega^7 \\ 1 & \omega^4 & 1 & \omega^4 & -\omega^2 & -\omega^6 & -\omega^2 & -\omega^6 \\ 1 & \omega^6 & \omega^4 & \omega^2 & -\omega^3 & -\omega & -\omega^7 & -\omega^5 \end{array} \right]$$



(Partitioning into 4 matrix blocks of size 4x4.)

$$\text{Recall: } F_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & w & w^2 & w^3 \\ 1 & w^2 & w^4 & w^6 \\ 1 & w^3 & w^6 & w^9 \end{bmatrix}, \text{ where } w = i = w_4$$

$$\text{Note: in } F_8, w = \frac{1+i}{\sqrt{2}} = w_8$$

$$\text{therefore, } w_8^2 = w_4$$

Example: Faster $y=Fx$

$$\bullet \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & w & w^2 & w^3 & w^4 & w^5 & w^6 & w^7 \\ 1 & w^2 & w^4 & w^6 & 1 & w^2 & w^4 & w^6 \\ 1 & w^3 & w^6 & w & w^4 & w^7 & w^2 & w^5 \\ 1 & w^4 & 1 & w^4 & 1 & w^4 & 1 & w^4 \\ 1 & w^5 & w^2 & w^7 & w^4 & w^1 & w^6 & w^3 \\ 1 & w^6 & w^4 & w^2 & 1 & w^6 & w^4 & w^2 \\ 1 & w^7 & w^6 & w^5 & w^4 & w^3 & w^2 & w^1 \end{bmatrix} =$$

$$\begin{array}{c|c} F_4 & \Omega_4 F_4 \\ \hline F_4 & -\Omega_4 F_4 \end{array}$$

\uparrow
 (because $w^2 = w_4$)

$$\Omega_4 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & w & 0 & 0 \\ 0 & 0 & w^2 & 0 \\ 0 & 0 & 0 & w^3 \end{bmatrix} \quad (\text{note: } w = \frac{1+i}{\sqrt{2}} = w_8)$$

$$\bullet \text{ So, } F_8 = \begin{bmatrix} F_4 & \Omega_4 F_4 \\ F_4 & -\Omega_4 F_4 \end{bmatrix}$$

FFT

We can obtain 8-point DFT from 4-point DFT. But how do we obtain the result of $y = F_8 x$, from $y_{\text{top}} = F_4 x_{\text{odd}}$ and $y_{\text{bottom}} = F_4 x_{\text{even}}$?

$$\begin{array}{c} \mathbf{F_8} \end{array} \begin{bmatrix} F_4 & \Omega_4 F_4 \\ \hline F_4 & -\Omega_4 F_4 \end{bmatrix} \begin{array}{c} \mathbf{x} \\ \begin{bmatrix} x1 \\ x3 \\ x5 \\ x7 \\ \hline x2 \\ x4 \\ x6 \\ x8 \end{bmatrix} \end{array} = \begin{bmatrix} I_4 & \Omega_4 \\ \hline I_4 & -\Omega_4 \end{bmatrix} \begin{array}{c} \begin{bmatrix} F_4 \\ \hline F_4 \end{bmatrix} \begin{bmatrix} x1 \\ x3 \\ x5 \\ x7 \\ \hline x2 \\ x4 \\ x6 \\ x8 \end{bmatrix} \end{array} = \begin{array}{c} \mathbf{y} \\ \begin{bmatrix} y1 \\ y3 \\ y5 \\ y7 \\ \hline y2 \\ y4 \\ y6 \\ y8 \end{bmatrix} \end{array}$$

Note: can be done with 4 multiplications

$$y1 \text{ to } y4 = y_{\text{top}} + \Omega_4 * y_{\text{bottom}}$$

$$y_{\text{top}} = F_4 x_{\text{odd}}$$

$$y_{\text{bottom}} = F_4 x_{\text{even}}$$

$$y5 \text{ to } y8 = y_{\text{top}} - \Omega_4 * y_{\text{bottom}}$$

(x_{odd} = elements at odd numbered indices of vector x)

(x_{even} = elements at even numbered indices of vector x)

Divide-and-Conquer FFT (D&C FFT)

FFT(v, ω , m) ... assume m is a power of 2

if m = 1 return v[0]

else

v_{even} = FFT(v[0:2:m-2], ω^2 , m/2)

v_{odd} = FFT(v[1:2:m-1], ω^2 , m/2)

ω -vec = [ω^0 , ω^1 , ... $\omega^{(m/2-1)}$]

precomputed



return [v_{even} + (ω -vec .* v_{odd}),

v_{even} - (ω -vec .* v_{odd})]

° Matlab notation: “.*” means component-wise multiply.

Cost: $T(m) = 2T(m/2) + O(m) = O(m \log m)$ operations.

Popularized/published by Cooley-Tuckey in 1965.

FFT

- Refer to Lecture 20 (Spring 2018) at <https://inst.eecs.berkeley.edu/~cs267/archives.html>
- Section 1.4, Matrix Computations, 4th Ed, Golub and Van Loan
- Section 3.5, Linear Algebra and Its Applications, 4th Ed, Gilbert Strang