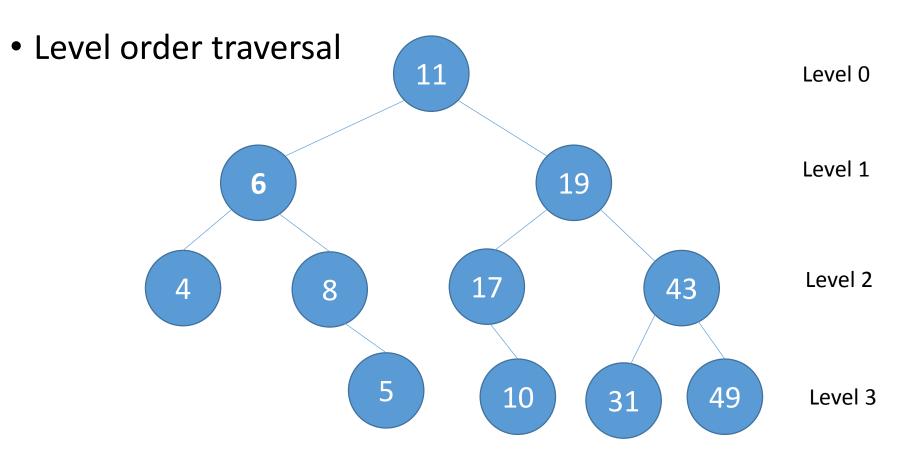
ECE264: Advanced C Programming

Summer 2019

Week 7: Binary Tree Traversal (contd.), Binary Search Trees, Misc. topics (const, variadic functions, macros, bitwise operations, bit fields), Parallel programming using threads

Breadth First Traversal (of a tree)



• 11, 6, 19, 4, 8, 17, 43, 5, 10, 31, 49

Breadth first traversal (of a tree)

```
void LOT(Node * root) {
 Queue q;
 push(&q, root);
 while (IsEmpty(&q) == false) {
  Node* headNode = Dequeue(&q)
  print(headNode->val)
  Enqueue(&q, headNode->leftChild);
  Enqueue(&q headNode->rightChild);
```

Depth first traversal (of a tree) – iterative code

Recall Preorder, Inorder, and Postorder were written as recursive

```
codes
                                       Inorder(Node* n) {
Preorder(Node* n) {
                                       if(n->val == NULL)
if(n->val == NULL)
                                              return;
      return;
                                       Inorder(n->leftChild);
print(n->val)
                                       print(n->val)
Preorder(n->leftChild);
                                       Inorder(n->rightChild);
Preorder(n->rightChild);
               PostOrder(Node* n) {
               if(n->val == NULL)
                      return;
               Postorder(n->leftChild);
               Postorder(n->rightChild);
               print(n->val)
```

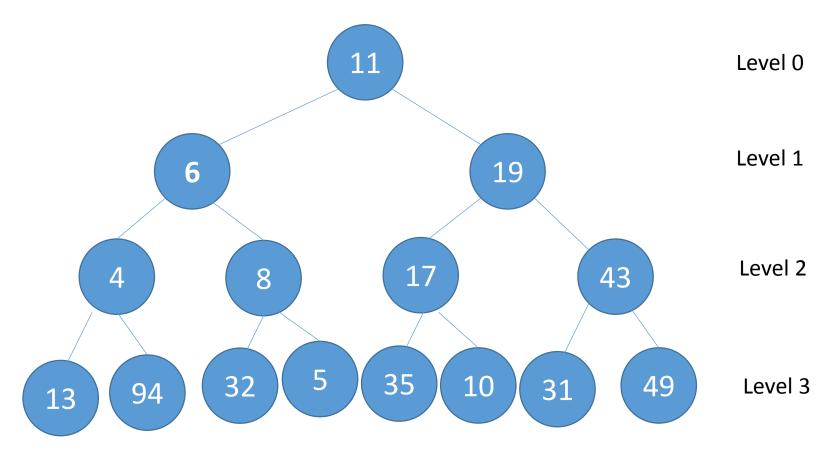
```
void Preorder(Node * root) {
 stack s;
 push(&s, root);
while (IsEmpty(&s) == false) {
  Node* topNode = Pop(&s)
  print(topNode->val)
  Push(&q, topNode->rightChild);
  Push(&q topNode->leftChild);
```

Exercise

What data structure do you need to use for writing an iterative code of Postorder traversal?

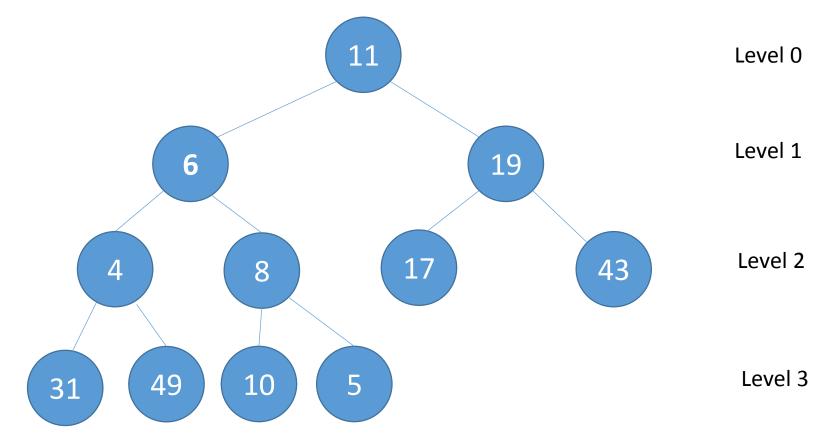
Full Binary Tree

• Every node except leaf has two children



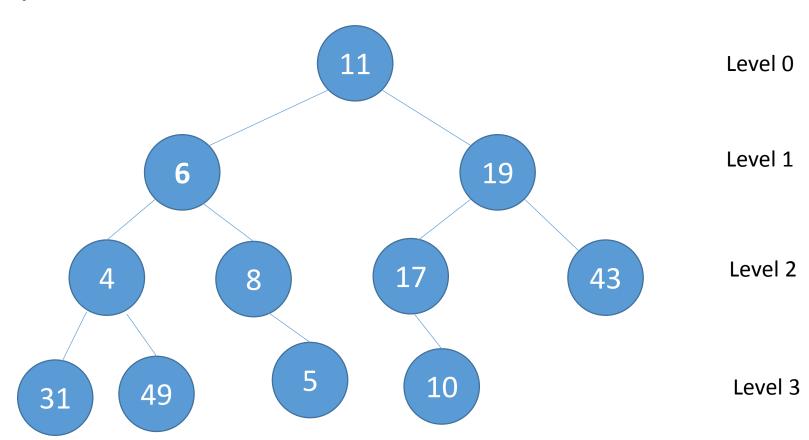
Complete Binary Tree

 Every level except the last is filled and all nodes at the last level are as far left as possible



Exercise

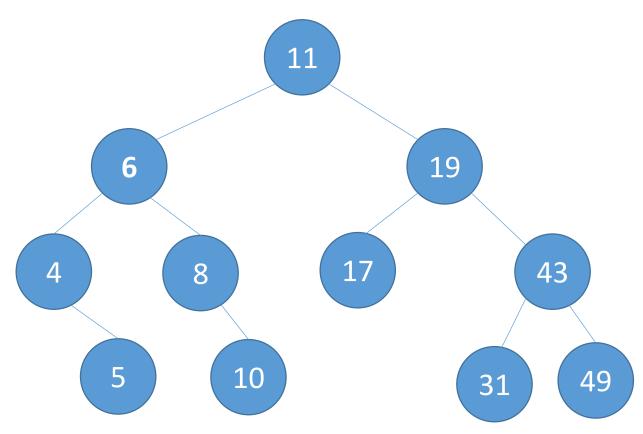
• Complete or Full ?



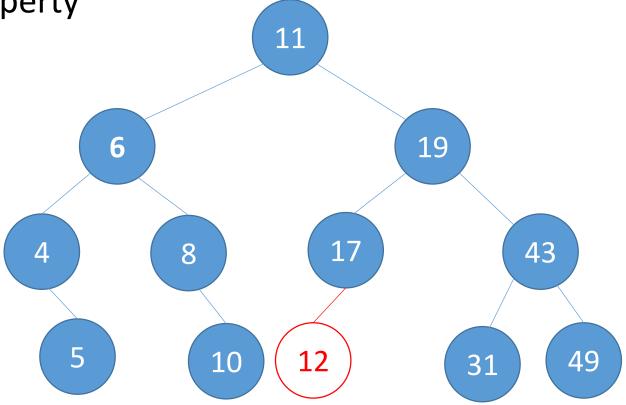
Binary Search Trees (BST)

- For efficient sorting, searching, retrieving
- BST Property:
 - Keys in left subtree are lesser than parent node key
 - Keys in right subtree are greater than parent node key
 - Duplicate keys not allowed

• Example



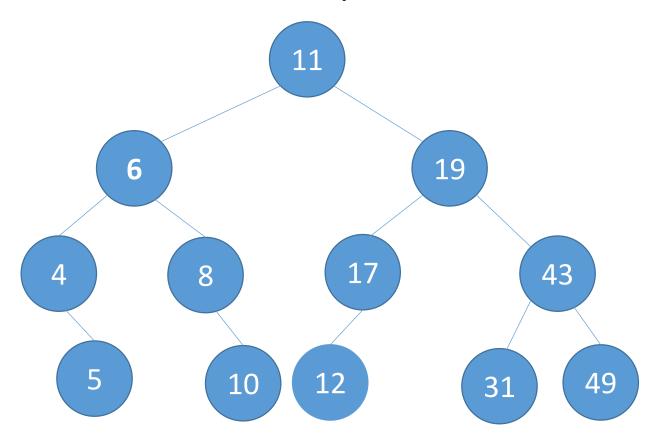
Insertion: inserts element without violating the BST property



Insertion

```
bool add(TreeNode **rootPtr, int key) {
       if (*rootPtr == NULL) {
              *rootPtr = buildNode(key);
3
4
              return true;
5
       } else if ((*rootPtr)->val == key) {
6
              return false;
       } else if ((*rootPtr)->val < key) {</pre>
              return add(&((*rootPtr)->right), key);
9
       } else {
10
              return add(&((*rootPtr)->left), key);
       }
11
12 }
```

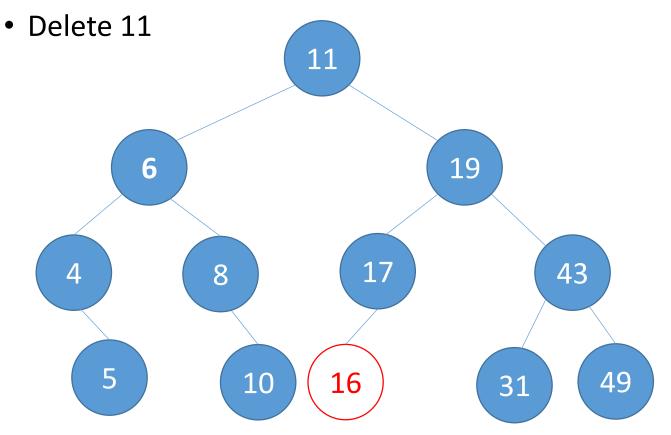
• Search: returns true if key exists. False otherwise.



Search

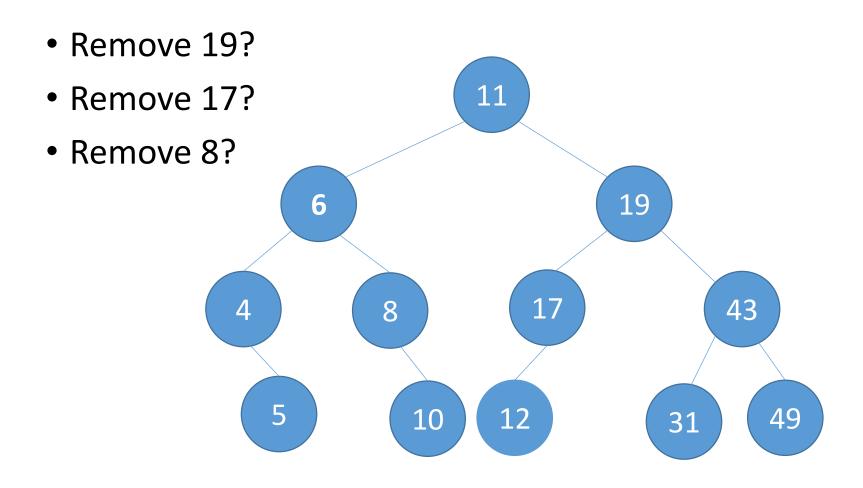
```
bool Contains(Node* n, int key) {
   if(n == NULL)
     return false;
   if(n->val == key)
     return true;
   else if (n->val > key)
     return Contains(n->leftChild, key);
   else
     return Contains(n->rightChild, key);
}
```

Removal: remove without violating BST property



- Removal cases
 - Not in a tree
 - Is a leaf
 - Has one or more children
- Return true if key removed. False otherwise.

Exercise



BST remove node

• Removal code: see bst.c

Applications – parsing of expression trees

- Goal: turn 2 + 3 into 2 3 +
 - We did this using stacks
- We can use binary trees to do the same job
- Binary trees allow us to create a more useful program
 - earlier we never checked if the input was a valid infix expression

We can build a basic compiler!

- Expressions (algebraic notation) are the normal way we are used to seeing them. E.g. 2 + 3
- Fully-parenthesized expressions are simpler versions: every binary operation is enclosed in parenthesis
 - E.g. (2 + (3 * 7))
 - So can ignore order-of-operations (PEMDAS rule)

Fully-parenthesized expression – definition

- Recursive definition
 - A number (floating point in our example)
 - 2. Open parenthesis '(' followed by fully-parenthesized expression followed by an operator ('+', '-', '*', '/') followed by fully-parenthesized expression followed by closed parenthesis ')'

Fully-parenthesized expression – notation

- 1. E -> lit
- 2. $E \rightarrow (E \text{ op } E)$

Expression parsing

Parsing is:

- 1. The process of determining if an expression is a valid fully-parenthesized expression
- 2. Breaking the expression into components
 - Why do we need this step?

We need not worry if a number has single digit, or multiple digits, or how many blank spaces separate two components etc.

Parsing

Rules: 1) E -> lit 2) E -> (E op E)

- Get the next token
- If the next token is a VAL (matches rule 1), return true.
- If the next token is an LPAREN match all of rule 2:
 - We have already seen the LPAREN, so the next thing we expect to see is a fully-parenthesized expression. We can just call this same function recursively to do that! If the recursive call returns true, it means we have found a fully-parenthesized expression
 - The next part of rule 2 is to match an operation, so we grab the next token and see if it is an ADD, SUB, MUL, or DIV. If it is, we continue.
 - Then we call this same function recursively again to find another fullyparenthesized expression.
 - Finally, we grab the next token to see if it is an RPAREN. If it is, we have matched rule 2, and this is a fully-parenthesized expression, so we return true.

Parsing

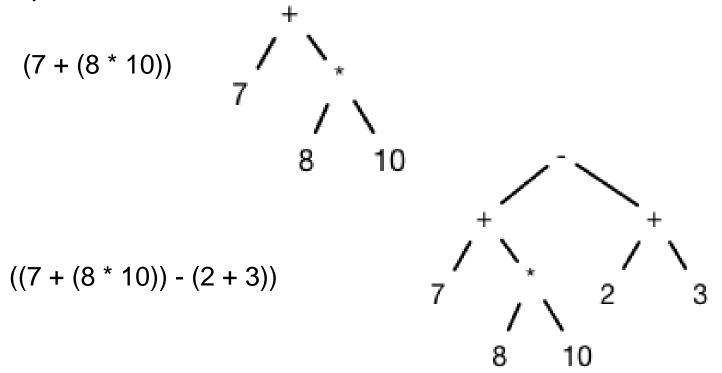
Example of a recursive descent parser

Can check if an expression is fully parenthesized

Can't check if a C program is valid

Expression trees

 Each leaf node is a number, non-leaf (interior) node a binary operation.



Building expression trees

- Can build while parsing a fully parenthesized expression
 Via bottom-up building of the tree
- Create subtrees, make those subtrees left- and right-children of a newly created root.

Modify recursive parser:

- 1. If token == VAL, return a pointer to newly created node containing a number
- 2. Else
 - store pointers to nodes that are left- and right- expression subtrees
 - 2. Create a new node with value = 'OP'

Expression trees

• Example: (7 + (8 * 10))

Exercise

What traversal order needs to be followed for tree deletion?

Type Qualifiers

- const, volatile, restrict
- Examples:

```
const int x=10; //equivalent to: int const x=10;
volatile int y=0; //eq to: int volatile y;
int *restrict c;
```

Const Qualifier

- The type is a constant (cannot be modified).
- const is the keyword

const int x=10; //x is a constant integer (hence, in RO memory). x cannot be modified.

We can also declare a const variable as:

```
int const x=10;
```

Const Properties

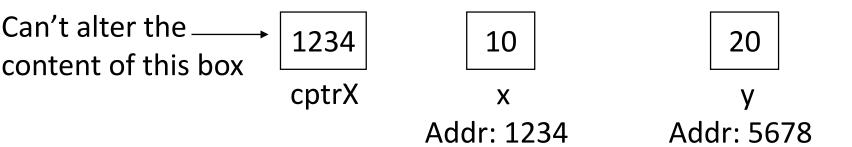
- Needs to be initialized at the time of definition
- Can't modify after definition
- const int x=10;
 x=20; //compiler would throw an error
- int const x=10;
 x=10; //can't even assign the same value
- int const y; //uninitialized const variable y. Useless.
 - 10 Can't alter the content of this box

Const Example1 (error)

```
/*ptrCX is a pointer to a constant integer. So, can't
modify what ptrCX points to.*/
const int* ptrCX;
int const* ptrCX;
int const x=10;
ptrCX = &x;
*ptrCX = 20; //Error
                              Can't alter the content of this box
               10
 1234
                              using ptrCX or x
ptrCX
                X
           Addr: 1234
```

Const Example2 (error)

```
/*cptrX is a constant pointer to an integer. So, can't
point to anything else after initialized.*/
int x=10, y=20;
int *const cptrX=&x;
cptrX = &y; //Error
```



Const Example3 (error)

```
/*cptrXC is a constant pointer to a constant integer. So,
      can't point to anything else after initialized. Also,
      can't modify what cptrXC points to.*/
      const int x=10, y=20;
      const int *const cptrXC=&x;
      int const *const cptrXC2=&x; //equivalent to prev. defn.
      cptrXC = &y; //Error
      *cptrX = 40; //Error
                                             Can't alter the content of
Can't alter the ____
                     1234
                                   10
                                             this box using cptrCX or x
content of this box
                    cptrXC
                                   X
                              Addr: 1234
```

Const Example4 (warning)

```
/*p2x is a pointer to an integer. So, we can use p2x to
alter the contents of the memory location that it points
to. However, the memory location contains read-only data -
cannot be altered. */
const int x=10;
const int *p1x=&x;
int *p2x=&x; //warning
*p2x = 20; //goes through. Might crash depending on memory
location accessed
                                      Can't alter the content of
1234
              1234
                            10
                                      this box using p1x or x.
p2x
                                      Can alter using p2x.
              p1x
                             X
                        Addr: 1234
```

Const Example5 (no warning, no error)

```
/*p1x is a pointer to a constant integer. So, we can't use p1x to alter the content of the memory location that it points to. However, the memory location it points to can be altered (through some other means e.g. using x)*/
```

```
int x=10;
const int *p1x=&x;
```

1234 Can't alter the content of this box using p1x.

p1x x Can alter using x.
Addr: 1234

Const Example6 (warning)

```
/*p1x is a constant pointer to an integer. So, we can use p1x to alter the contents of the memory location that it points to (and we can't let p1x point to something else other than x). However, the memory location contains readonly data - cannot be altered. */
```

```
const int x=10;
int *const p1x=&x;//warning
*p1x = 20; //goes through. Might crash depending on memory
location accessed
```

Can't alter the _____ 1234 = ____ Can't alter the content of this box using x. Can alter using p1x.

Addr: 1234

Const Example 7 (no warning, no error)

```
/*p1x is a pointer to a constant integer. So, we can't use p1x to alter the content of the memory location that it points to. However, the memory location it points to can be altered (through some other means e.g. using x)*/
```

```
int x=10;
const int *const p1x=&x;
```

Can't alter the _____ 1234 10 \leftarrow Can't alter the content of this box using p1x. Can alter using x.

Addr: 1234

const Case Study - strchr

- strchr is a library function that accepts a string and a char and returns a pointer to the first occurrence of the char
 - char* strchr(const char* str, char c)
- Returns a pointer to a char. So, we could modify the content of the memory location that the return value (a pointer) points to!
 - Exercise: is this an error or warning?

- Hint to the compiler indicating that a variable can change in unexpected ways (is volatile)
 - signals the compiler to not do any optimizations with the variable
 - Example:

```
volatile int* x = (volatile int *)0x1234; //x
is a pointer to a memory location with address
0x1234
```

- Hint to the compiler indicating that a variable can change in unexpected ways (is volatile)
 - signals the compiler to not do any optimizations with the variable
 - Example:

```
volatile int* x = (volatile int *)0x1234; //x
is a pointer to a memory location with address
0x1234
```

- Special memory locations in embedded systems programming are assigned certain addresses
 - Control registers, output buffers, input buffers
- For example, memory location at address 0x1234 could be a control register.
 - We can then access this register as we would access an unsigned int:

```
unsigned int *ctrlReg = (unsigned int *) 0x1234;
printf("current val of ctrl reg: %u", ctrlReg);
*ctrlReg=0x00000001; /*setting least significant bit to
indicate that input buffer has some data (we put some
data in input buffer and whoever is interested may
consume it)*/
```

```
unsigned int* ctrlReg = (unsigned int *)0x1234;
while (0 == *ctrlReg) {
//no data in input buffer. do some other work
sample assembly code (when optimizations turned on):
mov ctrlReg, #0x1234
mov a, @ctrlReg
loop:
bz loop
```

```
volatile unsigned int* ctrlReg = (volatile unsigned int
*)0x1234;
while (0 == *ctrlReg) {
//no data in input buffer. do some other work
sample assembly code (when optimizations turned on):
mov ctrlReg, #0x1234
loop:
mov a, @ctrlReg
bz loop
```

restrict Qualifier

- Introduced in C99
- May only be used with pointers
- Tells that the pointer is the only way to access a memory location

```
int * restrict source;
```

Example:

https://en.wikipedia.org/wiki/Restrict