# CS601: Software Development for Scientific Computing

Autumn 2022

Week1: Overview

# Let us listen to Jack Dongarra

https://youtu.be/Oe9LRKoE6L0

# Course Takeaways (intended)

- Non-CS majors:
  - 1. Write code and
  - 2. Develop software (not just write standalone code)
    - Numerical software
- CS-Majors:

In addition to the above two:

3. Learn to face mathematical equations and implement them with confidence

## What is this course about?

Software Development

+

Scientific Computing

## This course NOT about...

- Software Engineering
  - Systematic study of Techniques, Methodology, and Tools to build correct software within time and price budget (topics covered in CS305)
    - People, Software life cycle and management etc.
- Scientific Computing
  - Rigorous exploration of numerical methods, their analysis, and theories
  - Programming models (topics covered in CS410)

## Who this course is for?

- You are interested in scientific computing
- You are interested in high-performance computing
- You want to build / add to a large software system

#### GitHub Accounts

https://forms.gle/GTb1vCWj8FpbzaLb8

# Software Development

 Software development is the process of conceiving, specifying, designing, programming, documenting, testing, and bug fixing involved in creating and maintaining applications, frameworks, or other software components.

Software development is a process of writing and maintaining the source code, but in a broader sense, it includes all that is involved between the conception of the desired software through to the final manifestation of the software, ...

- Wikipedia on "Software Development"

# Scientific Computing

- Also called computational science
  - Development of models to understand systems (biological, physical, chemical, engineering, humanities)

Collection of tools, techniques, and theories required to solve on a computer mathematical models of problems in science and engineering

# Why C++ ?

- C/C++/Fortran codes form the majority in scientific computing codes
- Catch a lot of errors early (e.g. at compile-time rather than at run-time)
- Has features for object-oriented software development
- Known to result in codes with better performance

Let us dive into an example....

# **Example - Factorial**

```
• n! = n \times (n-1) \times (n-2) \times ... \times 3 \times 2 \times 1

(n-1)! = (n-1) \times (n-2) \times ... \times 3 \times 2 \times 1

therefore,
```

**Definition1:**  $n! = n \times (n-1)!$  is this definition complete?

plug 0 to n and the equation breaks.

#### **Definition2:**

$$n! = \begin{cases} n \times (n-1)! & \text{when } n>=1 \\ 1 & \text{when } n=0 \end{cases}$$

## Exercise 1

 Does this code implement the definition of factorial correctly?

```
int fact(int n){
   if(n==0)
     return 1;

return n*fact(n-1);
}
```

# Example - Factorial

Definition2: 
$$n! = \begin{cases} n \times (n-1)! & \text{when } n>=1 \\ 1 & \text{when } n=0 \end{cases}$$

is this definition complete?

n! is not defined for negative n

## Solution - Factorial

```
int fact(int n){
   if(n<0)
      return ERROR;
   if(n==0)
      return 1;

return n*fact(n-1);
}</pre>
```

## Exercise 2

In how many flops does the code execute?
 assume 1 flop = 1 step executing any arithmetic operation

```
int fact(int n){
   if(n<0)
       return ERROR;
   if(n==0)
      return 1;

return n*fact(n-1);
}</pre>
```

## Exercise 3

Does the code yield correct results for any n?

```
int fact(int n){
    if(n<0)
        return ERROR;
    if(n==0)
        return 1;

    return n*fact(n-1);
}</pre>
```

#### Who this course is for?

- Anybody who wishes to develop "computational thinking"
  - A skill necessary for everyone, not just computer programmers
  - An approach to problem solving, designing systems, and understanding human behavior that draws on concepts fundamental to computer science.

# Computational Thinking - Examples

- How difficult is the problem to solve? And what is the best way to solve?
- · Modularizing something in anticipation of multiple users
- · Prefetching and caching in anticipation of future use
- Thinking recursively
- Reformulating a seemingly difficult problem into one which we know how to solve by <u>reduction</u>, <u>embedding</u>, <u>transformation</u>, <u>simulation</u>
  - Are approximate solutions accepted?
  - False positives and False negatives allowed? etc.
- Using <u>abstraction</u> and <u>decomposition</u> in tackling large problem

19

# Computational Thinking – 2 As

#### Abstractions

- Our "mental" tools
- Includes: <u>choosing right abstractions</u>, operating at multiple <u>layers</u> of abstractions, and defining <u>relationships</u> among layers

#### Automation

- Our "metal" tools that <u>amplify</u> the power of "mental" tools
- Is mechanizing our abstractions, layers, and relationships
  - Need precise and exact notations / models for the "computer" below ("computer" can be human or machine)

# Computing - 2 As Combined

- Computing is the automation of our abstractions
- Provides us the ability to scale
  - Make infeasible problems feasible
    - E.g. SHA-1 not safe anymore
  - Improve the answer's precision
    - E.g. capture the image of a black-hole

**Summary:** choose the right abstraction and computer

# Recap

- Need to be precise
  - recall: n! = 1 for n=0, not defined for negative n
- Choosing right abstractions
  - recall: use of recursion, correct data type
- Ability to define the complexity
  - recall: flop calculation
- Next?

# Recap

- Need to be precise
  - recall: n! = 1 for n=0, not defined for negative n
- Choosing right abstractions
  - recall: use of recursion, correct data type
- Ability to define the complexity
  - recall: flop calculation
- Choose the right "computer" for mechanizing the abstractions chosen

## General Approach to Computational Problem

- 1. Problem statement: more precise this is, the easier it is to design and implement
- 2. Solution Algorithm: exactly how is the problem going to be solved
- 3. Implementation: breaking the algorithm into manageable pieces and putting it all together to solve the problem using a language of choice.
- **4. Verification:** checking that the implementation solves the original problem.
  - Often most difficult step, because you don't know the correct answer.

#### Scientific Software

- The answer is not a typical yes/no, red/blue/green
- The answer varies continuously. Think of computing the value of pi = 3.141592...
- Uses approximations. Think of discretization
- Employs efficient kernels
  - Kernels are core operations that are executed very frequently
- Should be able to adapt to change.
  - Writing everything from scratch is not an option
- Deals with large-scale problems
  - Lot of input/output data or both
  - Computationally hard

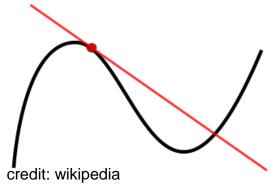
## Exercise

Compute root(s) of:

$$x = \cos x$$
;  $x \in \mathbb{R}$ 

roots, also called zeros, is the value of the argument/input to the function when the function output vanishes i.e. becomes zero

- let y = f(x) $f(x) = \cos(x) - x$
- At  $x = x_n$ , the value of y is  $f(x_n)$ . The coordinates of the point are  $(x_n, f(x_n)) = known$  point.
- From calculus: <u>derivative</u> of a function of single variable at a chosen input value, when it exists, is the <u>slope of</u> <u>the tangent</u> to the graph at that input value.
  - $f'(x_n)$  is the slope of the line that is tangent to f(x) at  $x_n$



From high-school math: point-slope formula for equation of a line

$$y - y_1 = m(x - x_1),$$
  
given the slope m and any known point  $(x_1, y_1)$ 

- Substituting with:
  - $(x_n, f(x_n)) = \text{known point}$
  - $f'(x_n)$  = slope

Equation of the tangent line to graph of f(x) at  $x_n$ :

$$y - f(x_n) = f'(x_n)(x - x_n)$$

- Interested in finding roots i.e. value of x at y=0 i.e. at point  $(x_{np1}, 0)$ .
- Substituting in the equation of the tangent line,

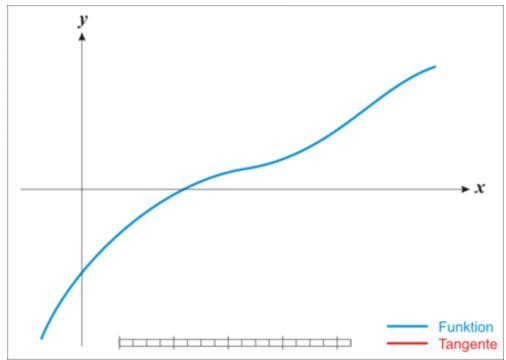
$$y - f(x_n) = f'(x_n)(x - x_n)$$

$$= -f(x_n) = f'(x_n)(x_{np1} - x_n)$$

$$= x_{np1} = x_n - f(x_n) / f'(x_n)$$

Visualizing

(SOURCE: <a href="https://en.wikipedia.org/wiki/Newton's\_method">https://en.wikipedia.org/wiki/Newton's\_method</a>):



The function f is shown in blue and the tangent line is in red. We see that  $x_{n+1}$  is a better approximation than  $x_n$  for the root x of the function f.

$$x_2 = x_1 - f(x_1) I f'(x_1)$$
  
 $x_3 = x_2 - f(x_2) I f'(x_2)$   
 $x_4 = x_3 - f(x_3) I f'(x_3)$ 

# Numerical Analysis

#### Talk to domain experts

- Choosing the initial value of x
- Does the method converge?
- What is an acceptable approximation?
- etc.

# Designing Algorithms and Data Structures

Start with x<sub>1</sub>

$$x_2 = x_1 - f(x_1) I f'(x_1)$$
  
 $x_3 = x_2 - f(x_2) I f'(x_2)$   
 $x_4 = x_3 - f(x_3) I f'(x_3)$ 

. . .

- Repeat for up to maxIterations
- Check for x<sub>n+1</sub> x<sub>n</sub> to be "sufficiently small"
- Choose appropriate data types for x

# Selecting libraries and tools

E.g. use the math library in C++ (cmath)

# Verify the correctness of results

- Compare with 'gold' code / benchmark
- Compare with empirical data

#### Real Numbers R

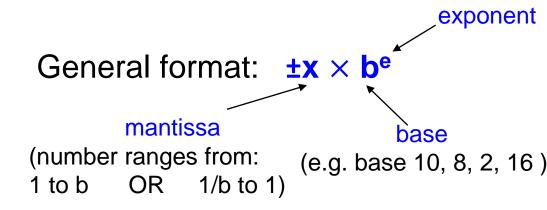
- Most <u>scientific software</u> deal with Real numbers.
   Our toy code dealt with Reals
  - Numerical software is scientific software dealing with Real numbers
- Real numbers include rational numbers (integers and fractions), irrational numbers (pi etc.)
- Used to represent values of <u>continuous quantity</u> such as time, mass, velocity, height, density etc.
  - Infinitely many values possible
  - But computers have limited memory. So, have to use approximations.

## Representing Real Numbers

- Real numbers are stored as floating point numbers
   (floating point system is a scheme to represent real numbers)
- E.g. floating point numbers:

```
-\pi = 3.14159,
```

- $-6.03*10^{23}$
- $-1.60217733*10^{-19}$



# 3-digit Calculator

Suppose base, b=10 and

• 
$$x = \pm d_0 \cdot d_1 d_2 \times 10^e$$
 where 
$$\begin{cases} 1 \le d_0 \le 9, \\ 0 \le d_1 \le 9, \\ 0 \le d_2 \le 9, \\ -9 \le e \le 9 \end{cases}$$

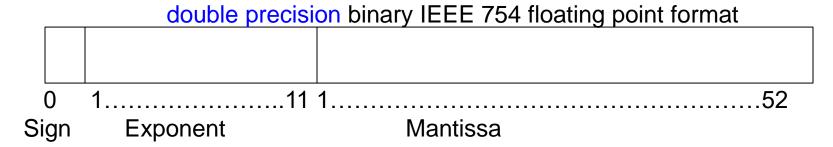
- precision = length of mantissa
  - What is the precision here?
- Exercise: What is the smallest positive number?
- Exercise: What is the largest positive number?
- Exercise: When is this representation not enough?
- Exercise: How many numbers can be represented in this format?

# IEEE 754 Floating Point System

Prescribes single, double, and extended precision formats

| Precision | u                   | Total bits used (sign, exponent, mantissa) |
|-----------|---------------------|--|
| Single    | 6x10 <sup>-8</sup>  | 32 (1, 8, 23)                              |
| Double    | 2x10 <sup>-16</sup> | 64 (1, 11, 52)                             |
| Extended  | 5x10 <sup>-20</sup> | 80 (1, 15, 64)                             |

# IEEE 754 Floating Point Arithmetic



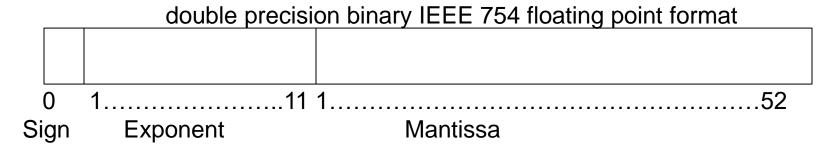
 if exponent bits e<sub>1</sub>-e<sub>11</sub> are not all 1s or 0s, then the normalized number

$$n = \pm (1.m_1 m_2..m_{52})_2 \times 2^{(e_1 e_2..e_{11})_2 - 1023}$$

41

- Machine epsilon is the gap between 1 and the next largest floating point number.  $2^{-52} \approx 10^{-16}$  for double.
- Exercise: What is minimum positive normalized double number?
- Exercise: What is maximum positive normalized double number?

# IEEE 754 Floating Point Arithmetic



if exponent bits e<sub>1</sub>-e<sub>11</sub> are all 0s, then the subnormal number

$$n = \pm (0.m_1 m_2..m_{52})_2 \times 2^{(e_1 e_2..e_{11})_2 - 1022}$$

- if exponent bits  $e_1$ - $e_{11}$  are all 1s, then we can get –inf, NaN, or +inf based on value of  $m_1m_2..m_{52}$ .
  - If any m is non-zero, the number is NaN (not a number)

# IEEE 754 Floating Point Arithmetic

- Order is important
  - Floating point arithmetic is not associative
    - (x+y)+z not the same as x+(y+z)
- Explicit coding of textbook formula may not be the best option to solve
  - $x^2 2px q = 0$  p and q are positive
  - Exercise: find the minimum of the roots.