

# Algorithmen und Objektorientierte Programmierung II

Binomialer Heap und Fibonacci Heap

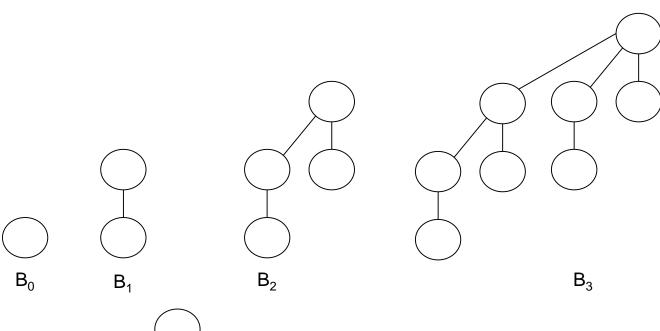
### **Motivation**



Operation	Binärer Heap	Binomialer Heap	Fibonacci-Heap
Make-Heap	O(1)	O(1)	O(1)
Insert	O(log n)	O(log n)	O(1)
Minimum	O(1)	O(log n)	O(1)
Extract-Min	O(log n)	O(log n)	O(log n)*
Decrease-Key/Increase-Key	O(log n)	O(log n)	O(1)*
Delete	O(log n)	O(log n)	O(log n)*
Union	O(n)	O(log n)	O(1)

<sup>(\*)</sup>Amortisierte Kosten (¹)Bei bekannter Position, sonst O(log n)\*





 $B_{k-1}$ 

 $B_{k-1}$ 

 $\boldsymbol{\mathsf{B}_{\mathsf{k}}}$ 



- 1. Ein binomialer Baum ist ein geordneter Baum
- 2. Besteht aus 2k Knoten
- 3. Die Höhe des Baumes ist k
- 4. Wurzel hat Grad k
  - 1. Kinder der Wurzel haben von links nach rechts den Grad k-1, k-2, ... 0
- 5. 4 gilt rekursiv

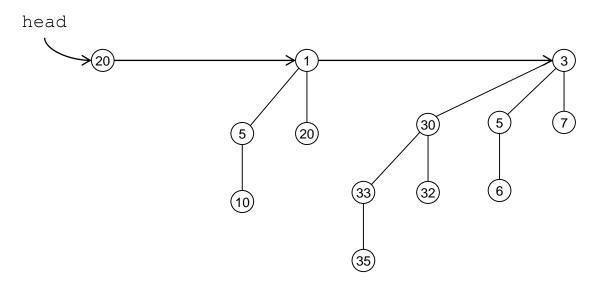
### **Binomialer Heap**



### Ein Binomialer Heap

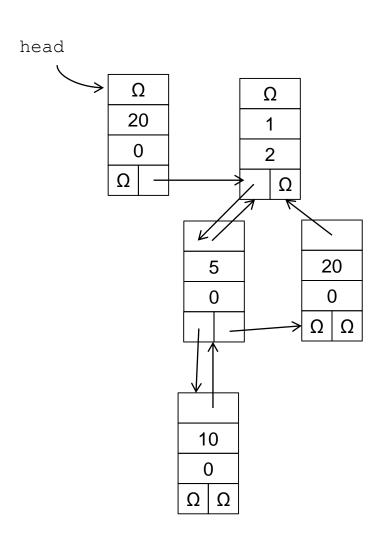
- ist eine Menge von binomialen Bäumen
- · jeder Binomiale Baum erfüllt die Heap-Bedingung
- Für k >= 0 gibt es höchstens einen binomialen Baum, dessen Wurzel den Grad k besitzt

=> Heap mit n Knoten besteht aus maximal  $\lfloor \log n \rfloor + 1$  binomiale Bäume



# **Binomialer Heap**

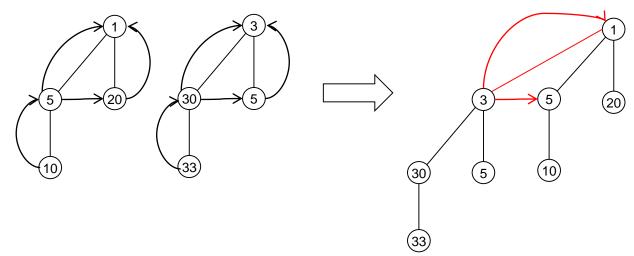




```
class Node<T> {
    T value;
    int degree;
    Node<T> parent;
    Node<T> child;
    Node<T> sibling;
}
```



```
Binomial-Link(y, z) {
    ASSERT(y.degree == z.degree)
    y.parent = z
    y.sibling = z.child
    z.child = y
    z.degree = z.degree + 1
}
```



## **Basisoperationen**



```
Binomial-Union (H_1, H_2) {
     H = Binomial-Heap-Merge(H_1, H_2) //Link roots and sort by increasing degree
     while (|\{B_i\}| > 1, i \epsilon k) {
         Binomial-Link(B_m, B_n) // m == n
     return H
```



```
Binomial-Heap-Insert(H, x) {
     Make-Binomial-Heap (H<sub>1</sub>)
     Insert (H_1, x)
     Binominal-Heap-Union(H, H_1)
Binomial-Heap-ExtractMin(H) {
     Search-Min min and remove
     Make-Binomial-Heap (H<sub>1</sub>)
     for all x in min.children in reverse order
           Insert (H_1, x)
     Binominal-Heap-Union (H, H_1)
     return min
Binomial-Heap-Decrease-Key(H, x, k) {
     ASSERT (k < x.key)
     //Vertausche in Richtung zur Wurzel
Binomial-Heap-Delete(H, x) {
      Binomial-Heap-Decrease-Key(H, x, -\infty)
      Binomial-Heap-ExtractMin(H)
```

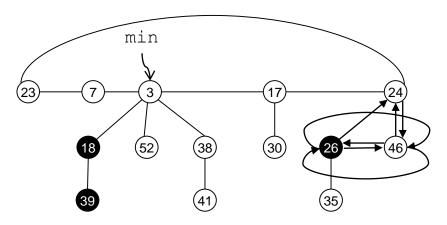
#### Fibonacci Heap



#### Ein Binomialer Heap

- ist eine Menge von ungeordneten Bäumen
- jeder Baum erfüllt die Heap-Bedingung

### => In einem Heap mit n Knoten ist der maximale Grad $D(n) \le \lfloor \log n \rfloor$



```
class Node<T> {
    T value;
    int degree;
    boolean mark;
    Node<T> parent;
    Node<T> child;
    Node<T> leftSibling;
    Node<T> rightSibling;
}
```



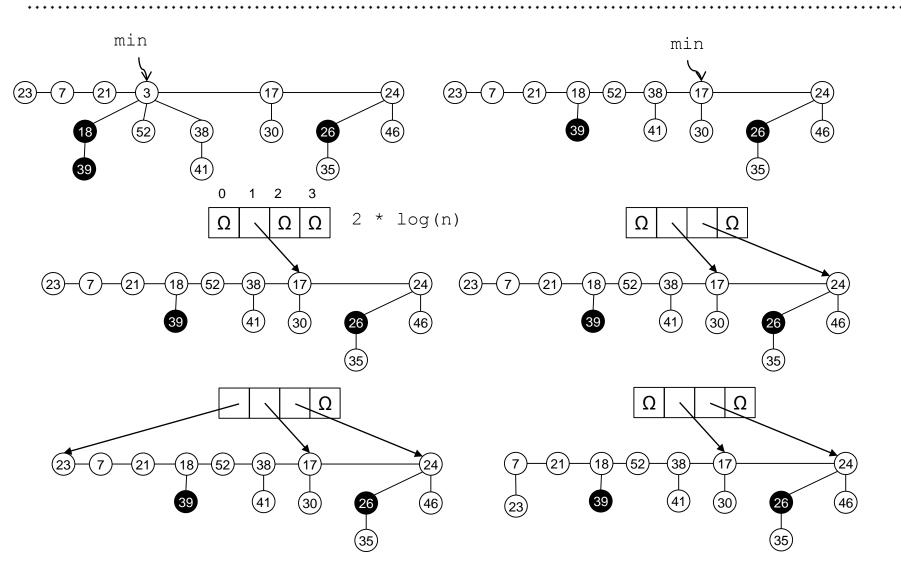
```
FIB-Heap-Insert(H, x) {
      x.degree=0 ... x.mark = false //init x
      Insert(H, x)
      if (H.min == null or x.key < h.min.key)
          H.min = x
     H.n = H.n + 1
FIB-Heap-Union(H<sub>1</sub>, H<sub>2</sub>) {
     FIB-Heap-Make (H)
     H.min = H_1.min //Shallow copy
     Concat(H, H<sub>2</sub>) //link root lists
     if ((H_1.min == null) \mid (H_2.min != null and H_2.min < H_1.min))
          H.min = H_2.min
     H.n = H_1.n + H_2.n
      return H
```



```
FIB-Heap-Extract-Min(H) {
     z = H.min
     if (z != null) {
         for each x of z.children
            FIB-Heap-Insert(H, x)
         Remove (H, z)
         if (z == z.rightSibling)
            H.min = null
         else {
            H.min = z.rightSibling
            FIB-Heap-Consolidate(H)
         H.n = H.n - 1
     return z
```

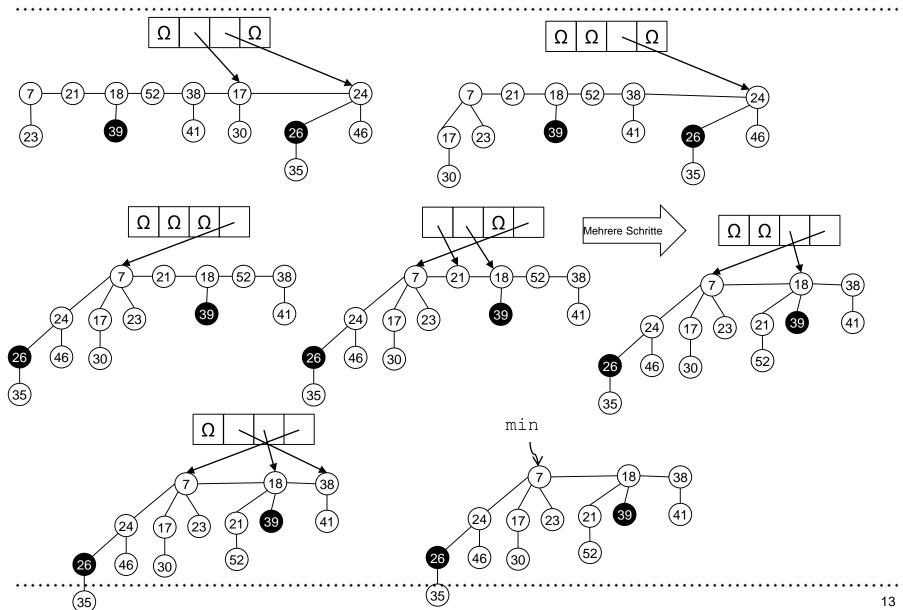
# **Operationen (FIB-Heap-Consolidate 1/2)**





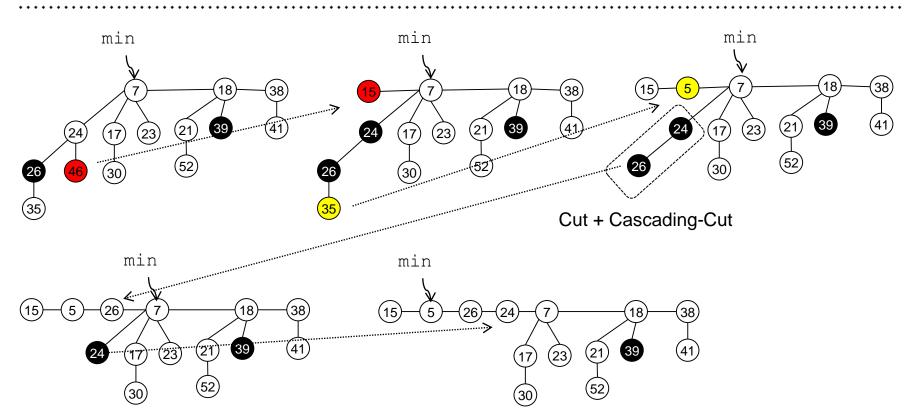
# **Operationen (FIB-Heap-Consolidate 2/2)**





# **Operation (FIB-Heap-Decrease-Key)**





#### **Operation (FIB-Heap-Decrease-Key und FIB-Heap-Delete)**



```
FIB-Cascading-Cut(H, y) {
FIB-Heap-Decrease-Key(H, x, k) {
                                              z = y.parent
     ASSERT (k < x.key)
                                              if (z != null) {
     x.key = k
                                                    if (y.mark == false)
     y = x.parent
                                                         y.mark = true
     if (y != null and x.key < y.key) {
                                                    else {
        FIB-Cut(H, x, y)
                                                         FIB-Cut(H, y, z)
        FIB-Cascading-Cut(H, y)
                                                         FIB-Cascading-Cut(H, z)
     if (x.key < H.min.key)</pre>
        H.min = x
                                           FIB-Heap-Delete(H, x) {
                                              FIB-Heap-Decrease-Key(H, x, -\infty)
FIB-Cut(H, x, y) { //move x to rootList
                                              FIB-Heap-ExtractMin(H)
     Remove x from y.children
     y.degree = y.degree - 1
     x.parent = null
     FIB-Heap-Insert(H, x)
```