



# Practical Network Defense

*Master's degree in Cybersecurity 2020-21*

## VPN, SSL/TLS and IPSec

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# Today's agenda

- VPN principles
- SSL Tunneling
- VPN device placement
- IPsec



# VPN principles

# Virtual Private Networks

- Definition (NIST SP800-113): A virtual network, built on top of an existing network infrastructure, which can provide a secure communications mechanism for data and other information transferred between two endpoints
- Typically based on the use of encryption, but several possible choices for:
  - How and where to perform the encryption
  - Which parts of communication should be encrypted
- Important subsidiary goal: usability
  - If a solution is too difficult to use, it will not be used → poor usability leads to no security

# Security Goals for a VPN

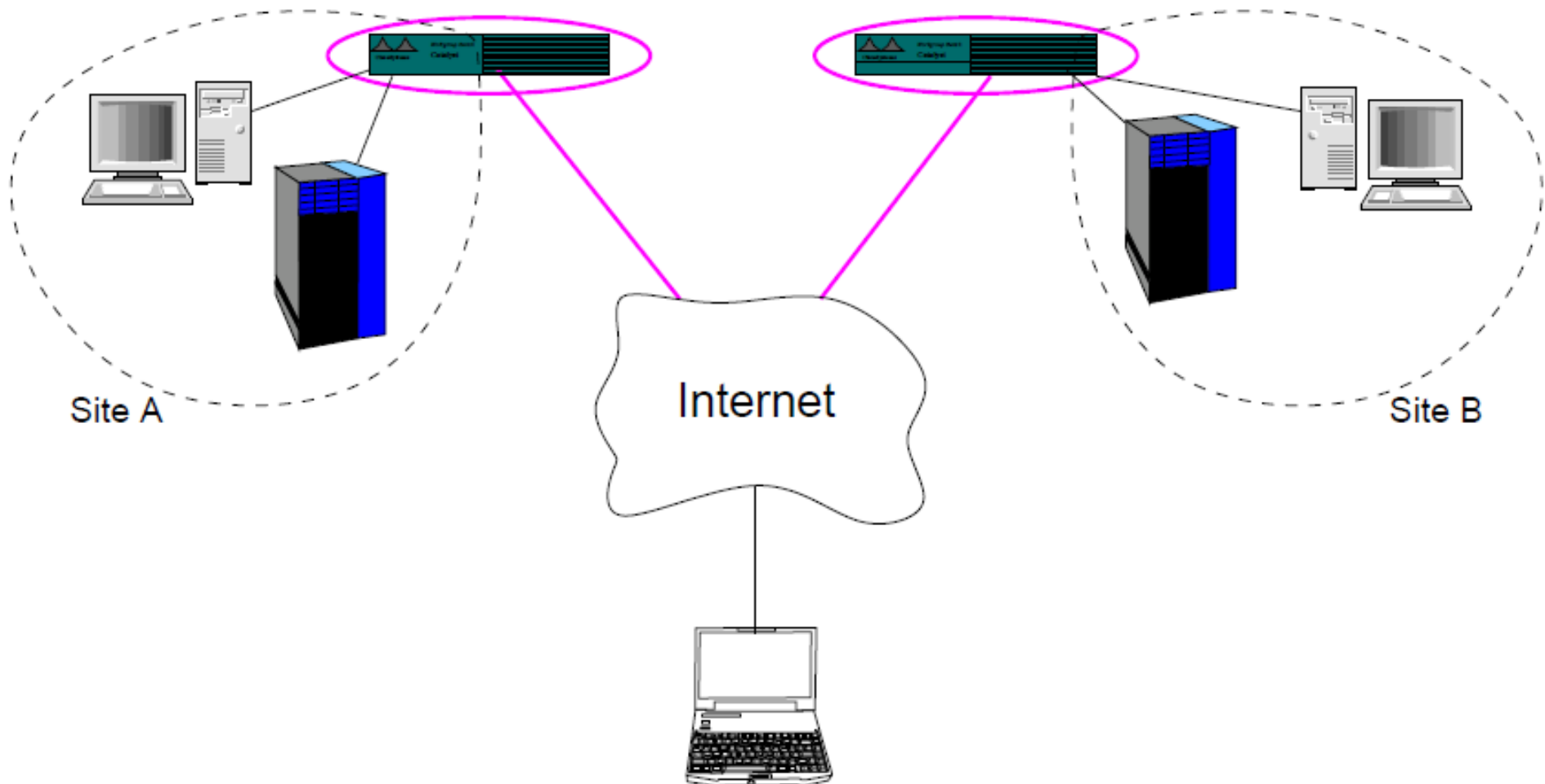
- Traditional
  - Confidentiality of data
  - Integrity of data
  - Peer Authentication
- Extended
  - Replay Protection
  - Access Control
  - Traffic Analysis Protection



# Usability goals

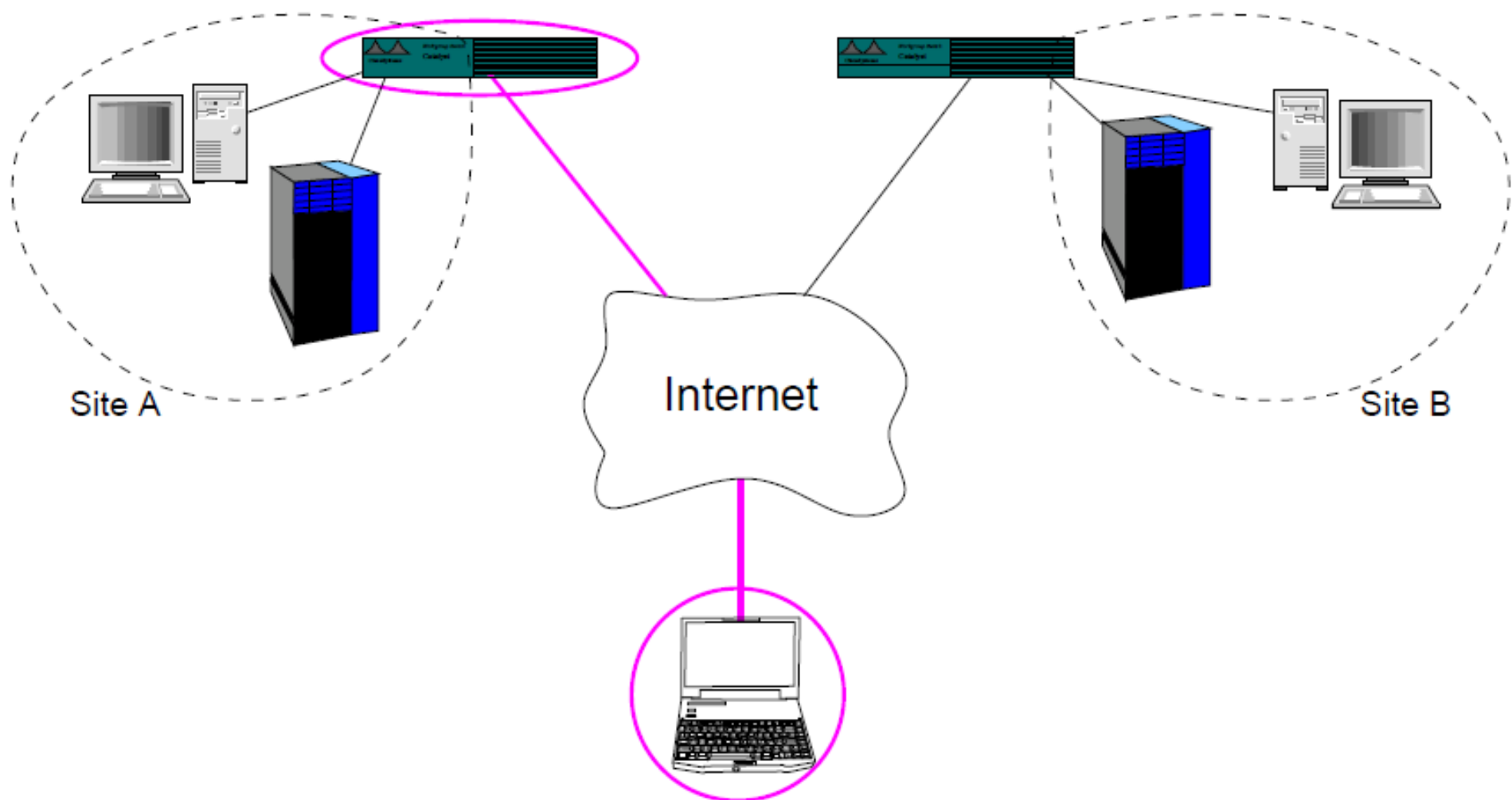
- Transparency
  - VPN should be invisible to users, software, hardware.
- Flexibility
  - VPN can be used between users, applications, hosts, sites.
- Simplicity
  - VPN can be actually used

# Site-to-site security

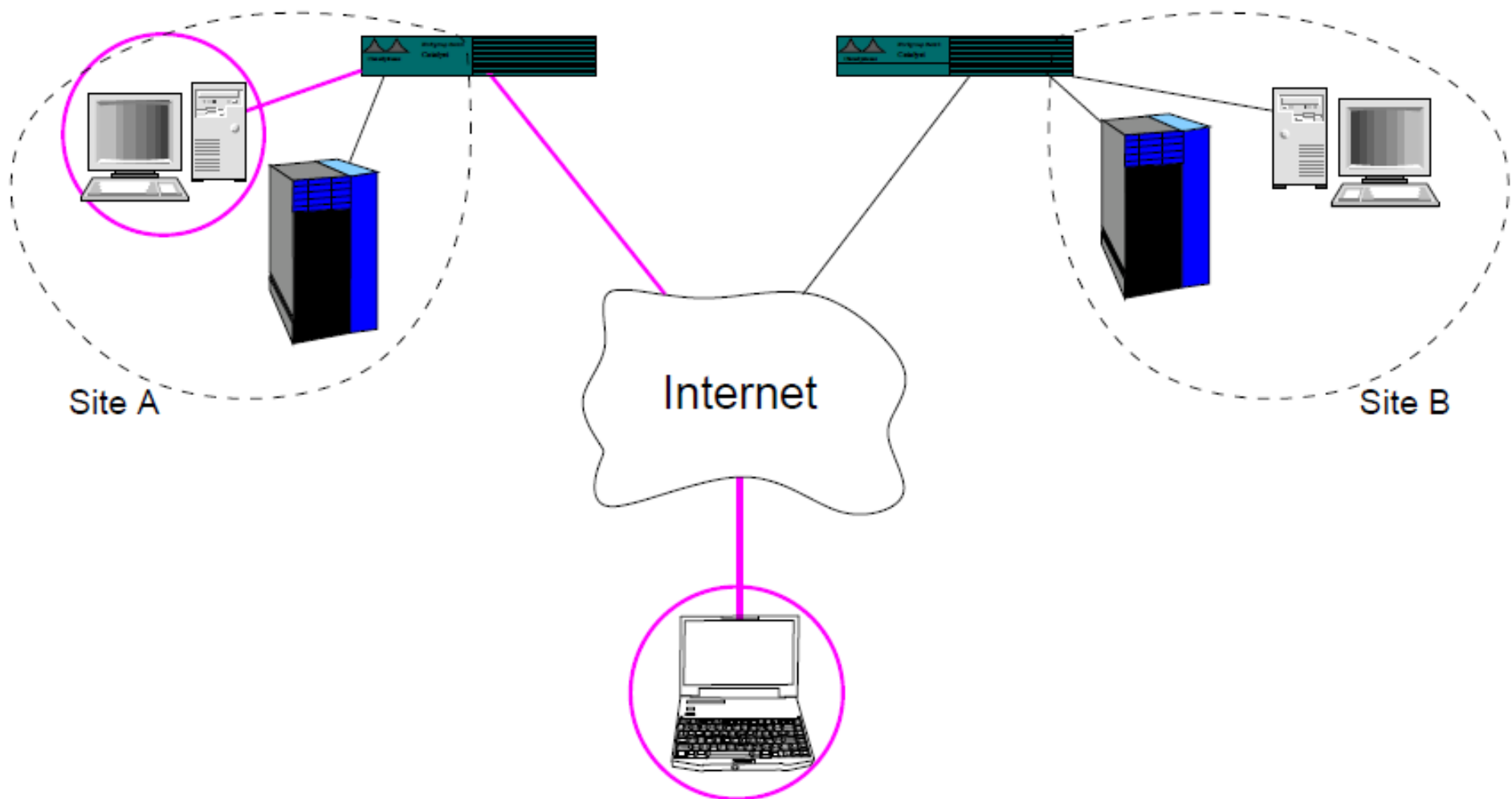




# Host-to-site security

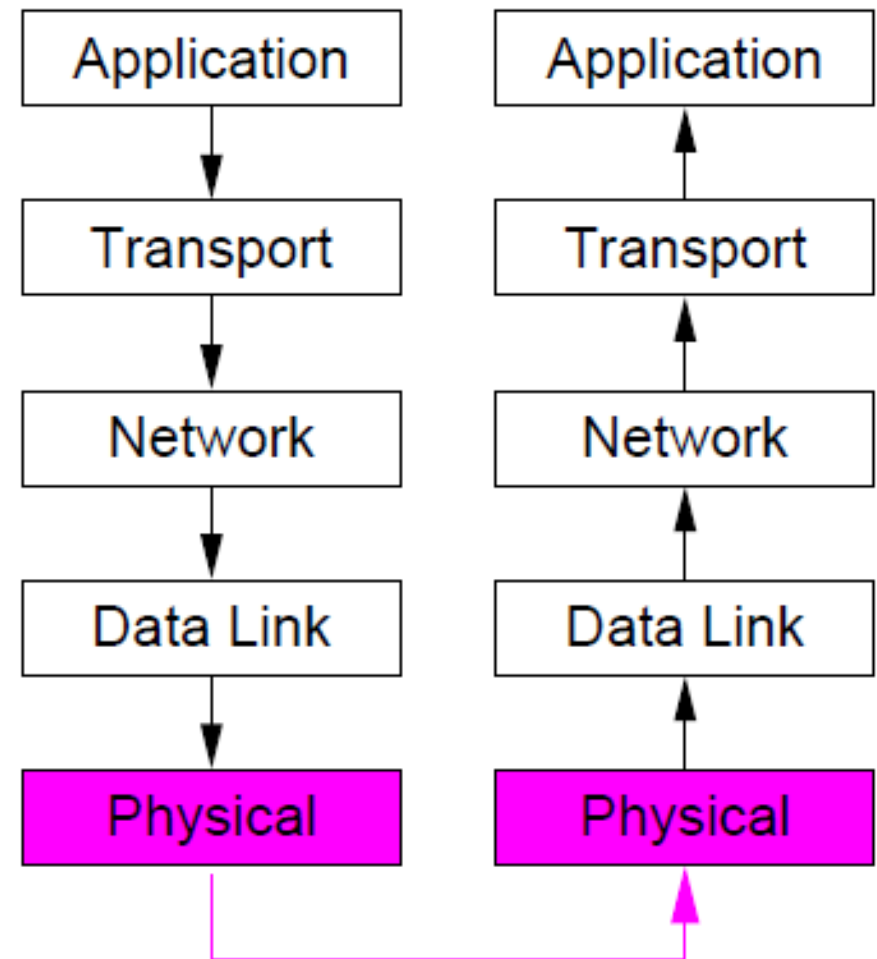


# Host-to-host security

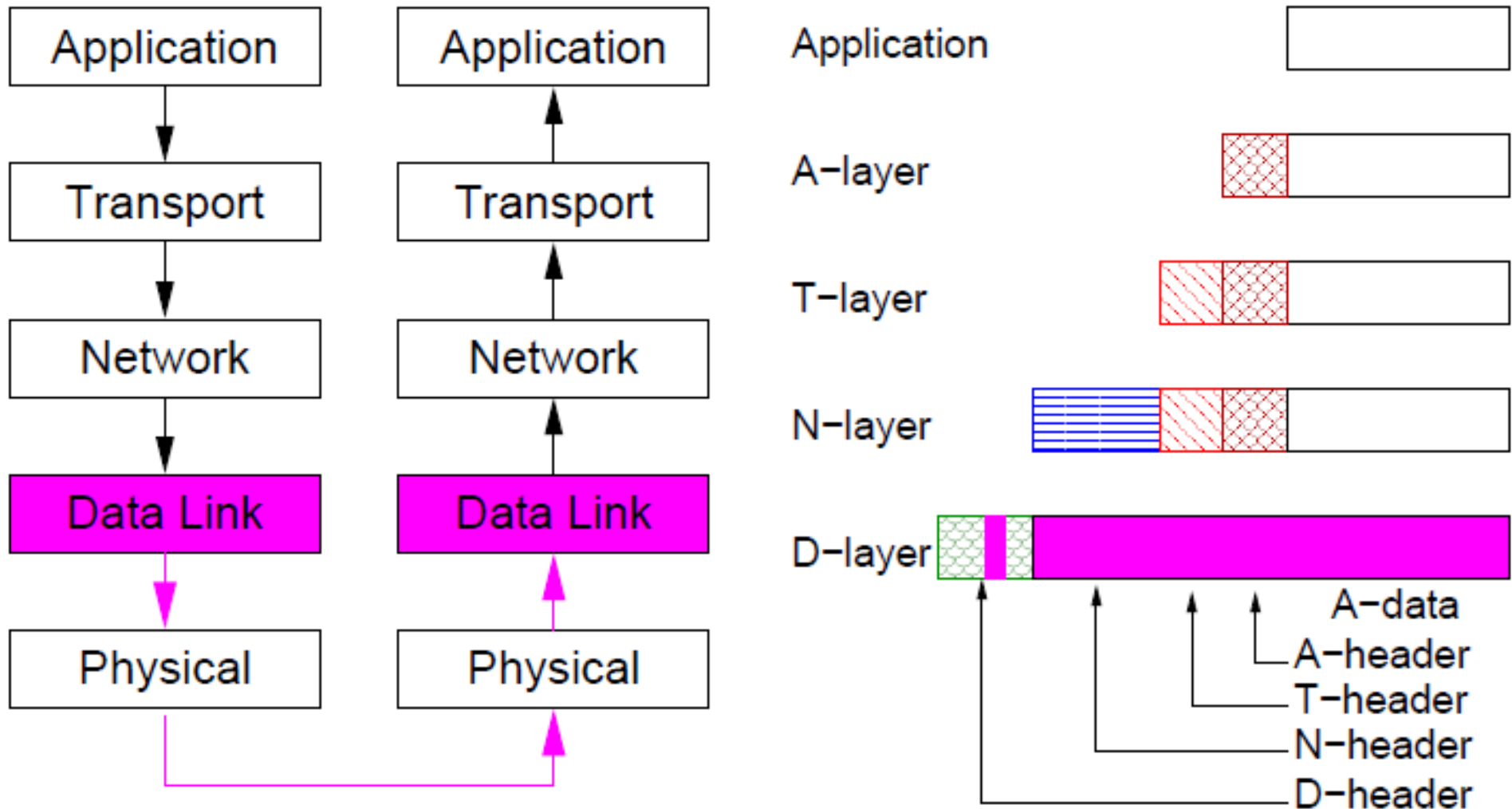


# Physical layer

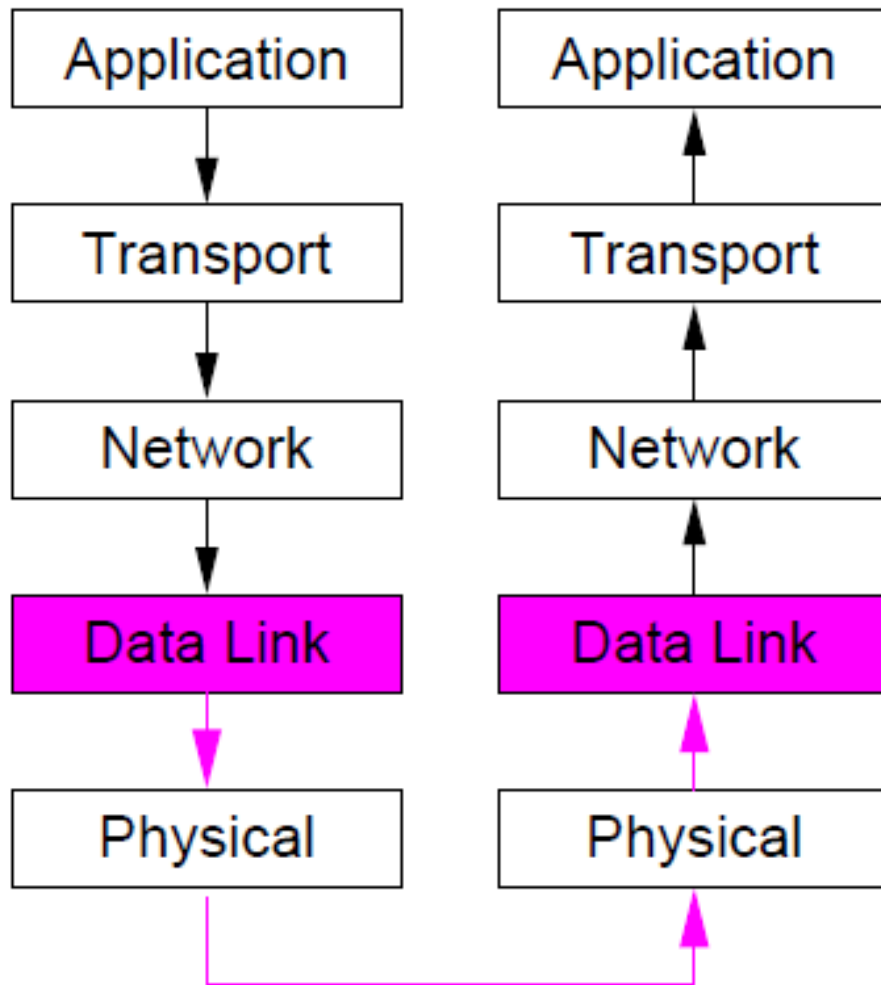
- Confidentiality: on cable
- Integrity: on cable
- Authentication: none
- Replay protection: none
- Traffic analysis protection: on cable
- Access control: physical access
- Transparency: full transparency
- Flexibility: can be hard to add new sites
- Simplicity: excellent!



# Datalink layer: protect a single link

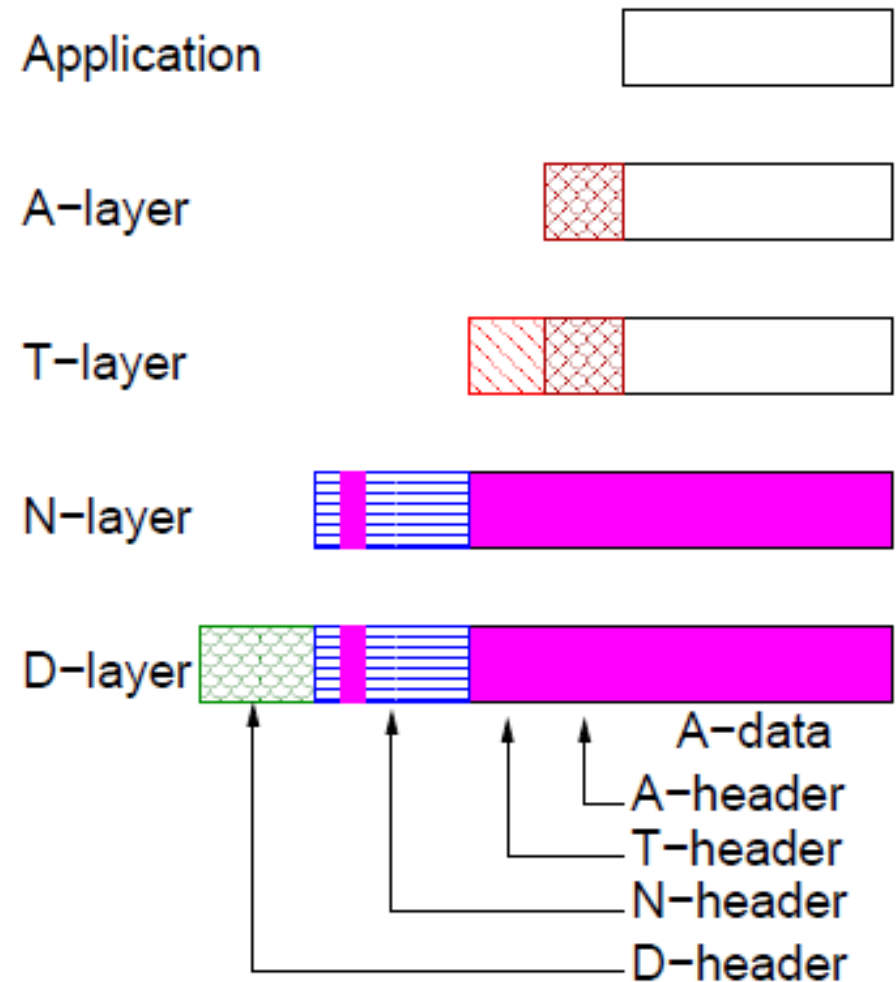
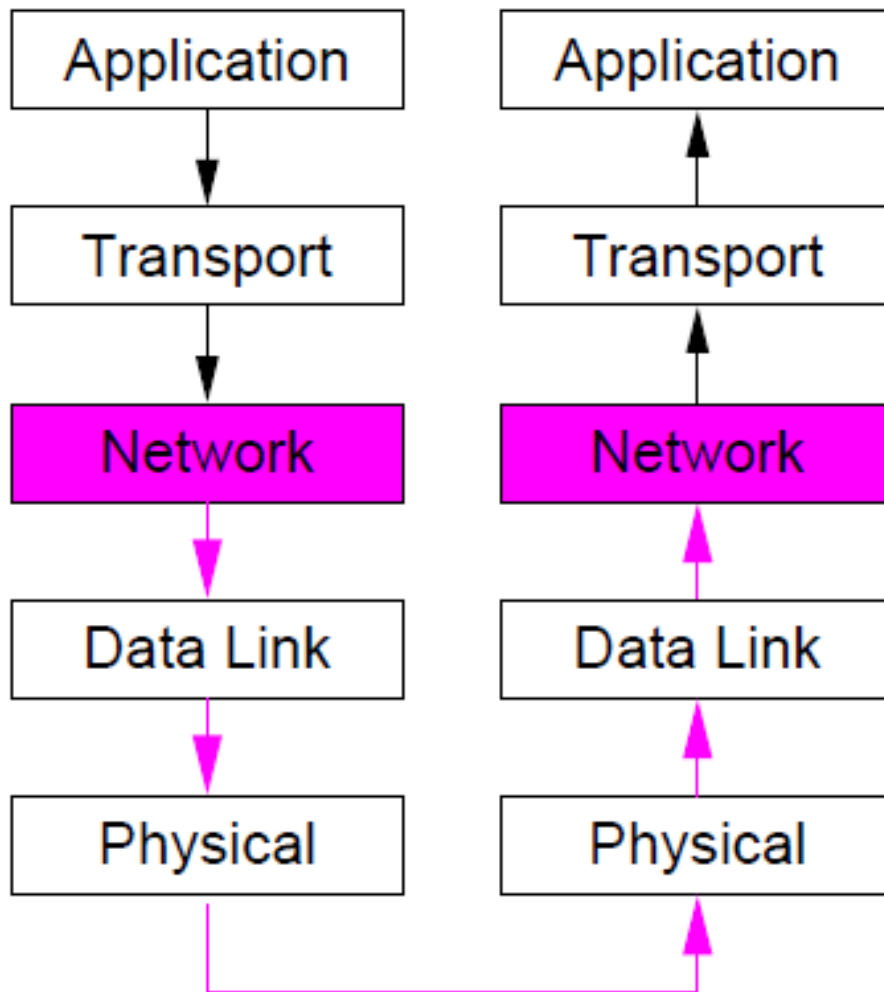


# Datalink layer: protect a single link

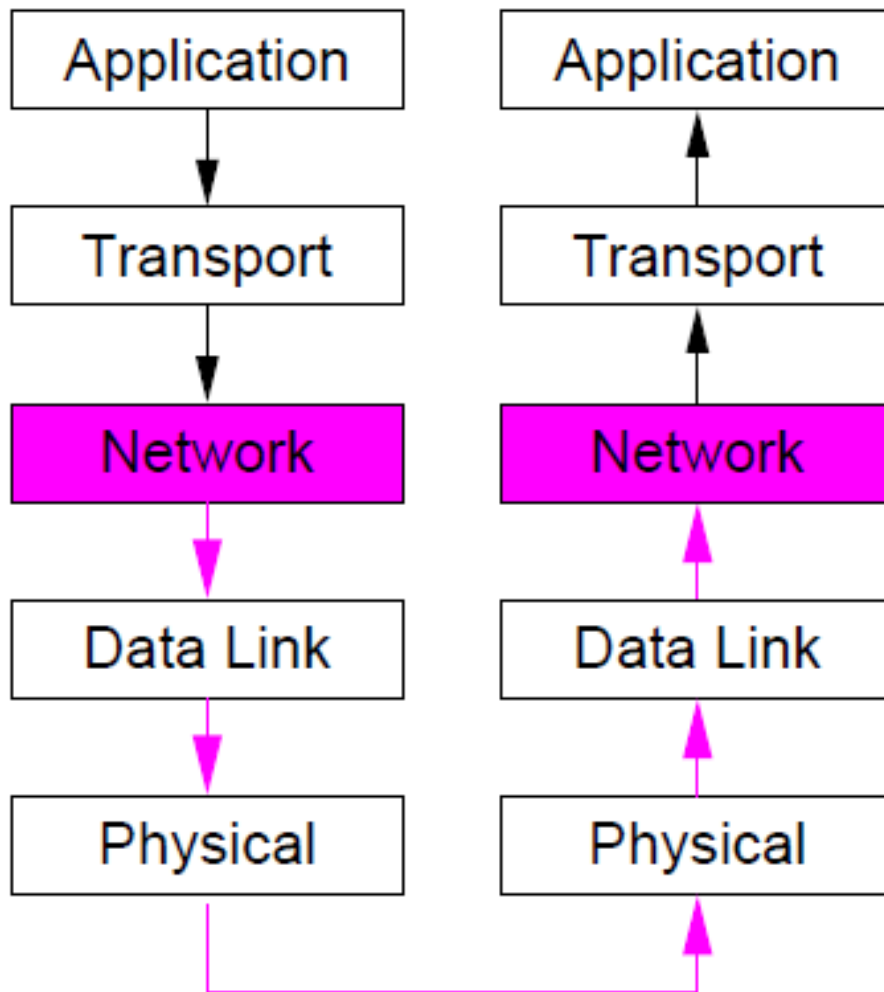


- Confidentiality: on link ("virtual cable")
- Integrity: on link
- Authentication: none
- Replay protection: none
- Traffic analysis protection: on link
- Access control: physical access
- Transparency: full transparency
- Flexibility: can be hard to add new sites
- Simplicity: excellent!

# Network layer: protect end-to-end between systems

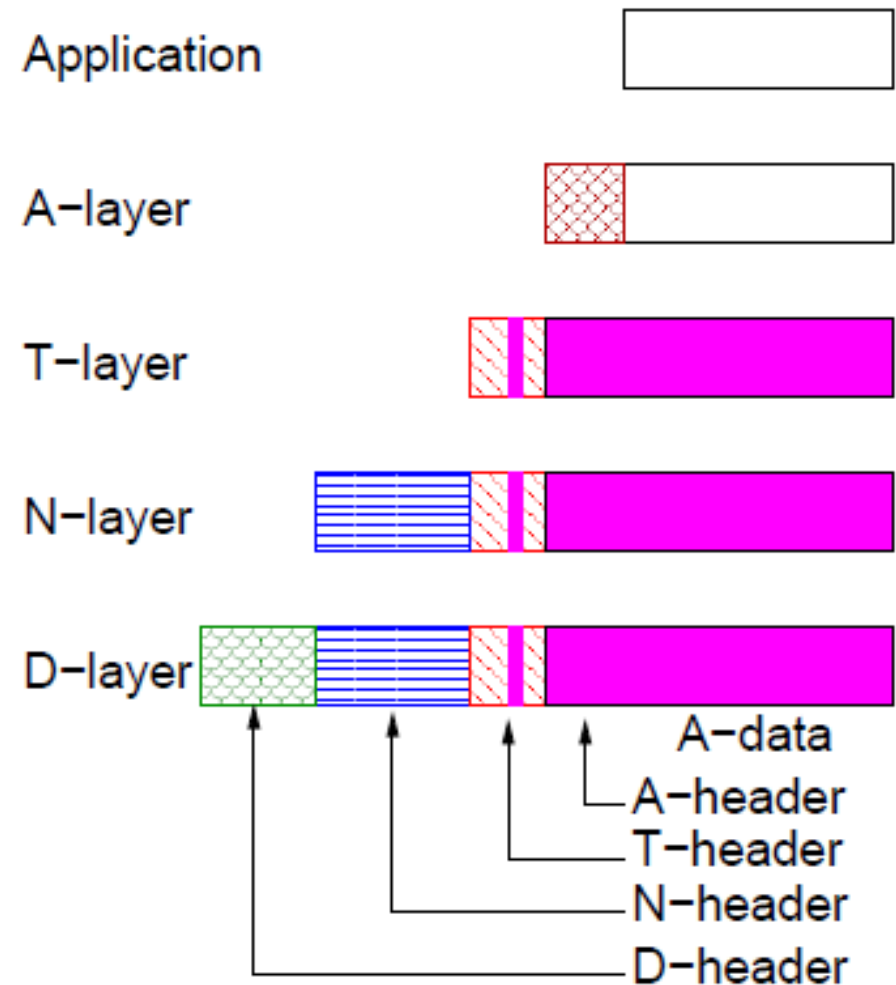
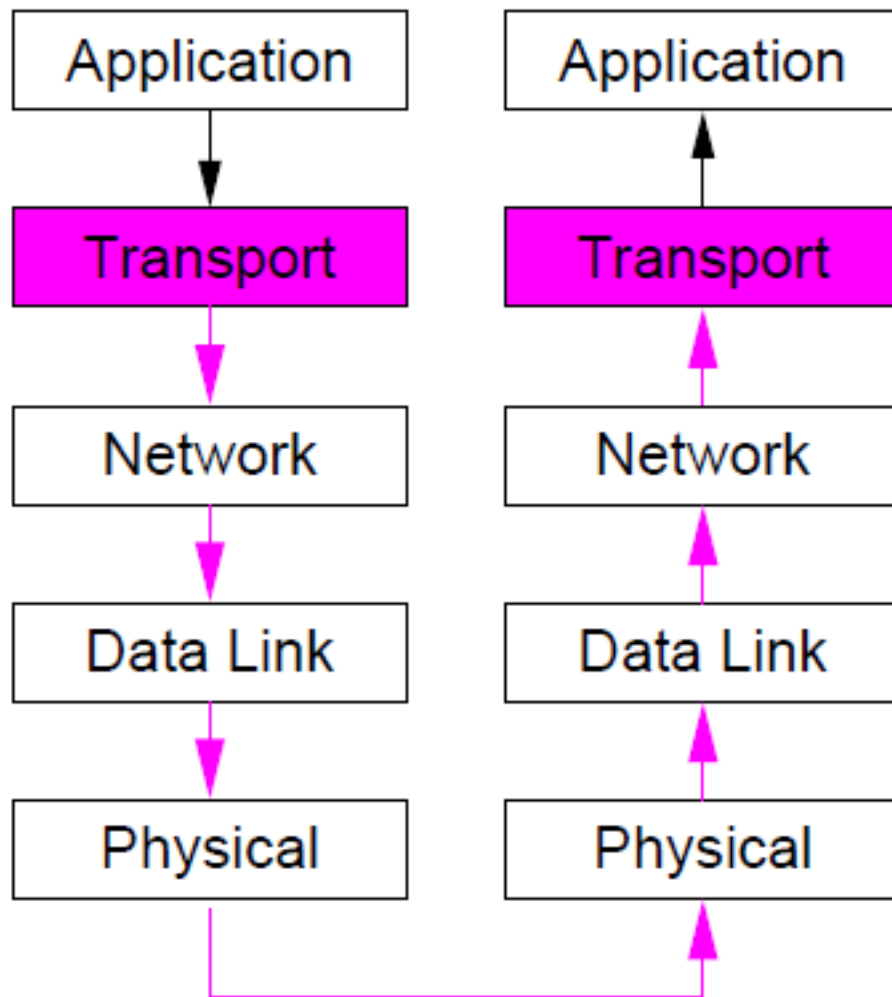


# Network layer: protect end-to-end between systems



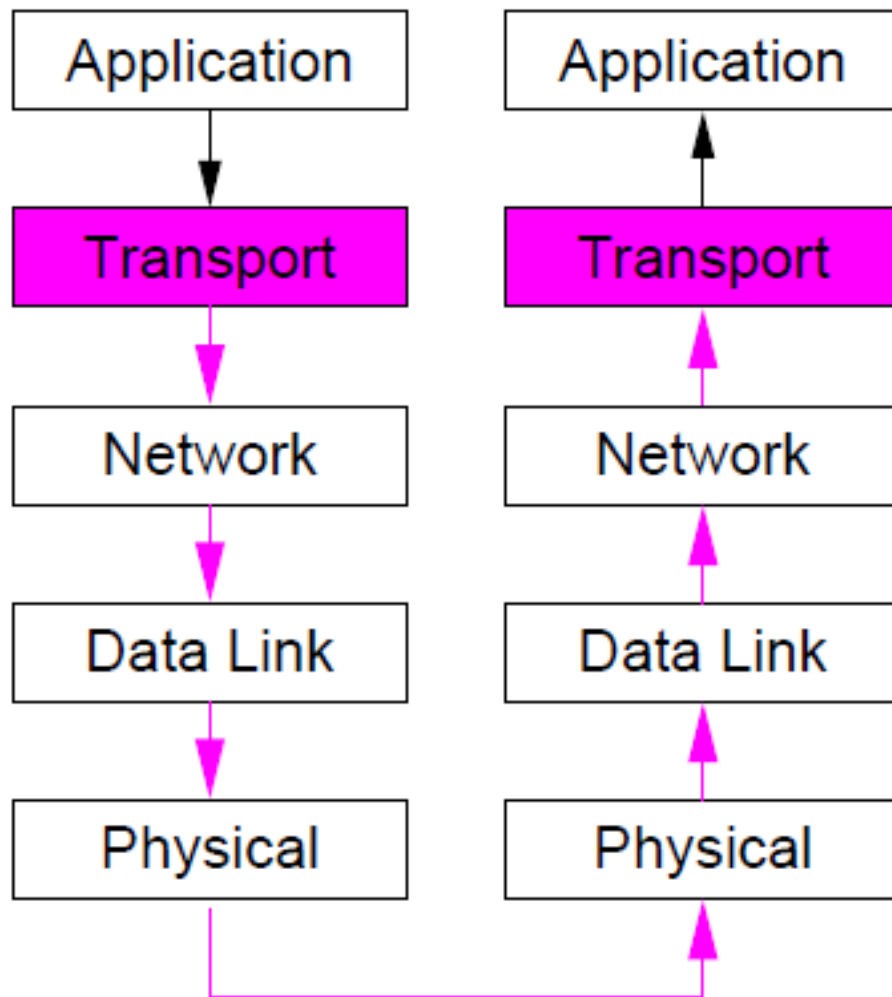
- Confidentiality: between hosts/sites
- Integrity: between hosts/sites
- Authentication: for host or site
- Replay protection: between hosts/sites
- Traffic analysis protection: host/site information exposed
- Access control: to host/site
- Transparency user and SW transparency possible
- Flexibility: may need HW or SW modifications
- Simplicity: good for site-to-site, not good for host-to-site

# Transport layer: Protection end-to-end between processes



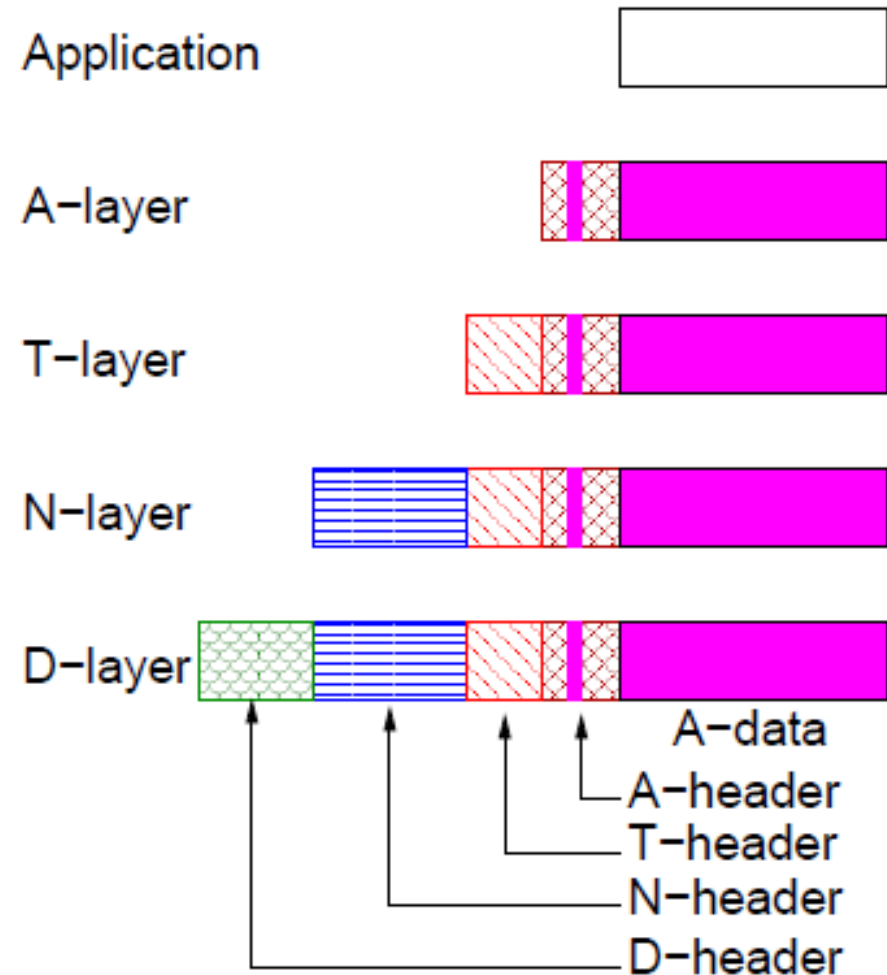
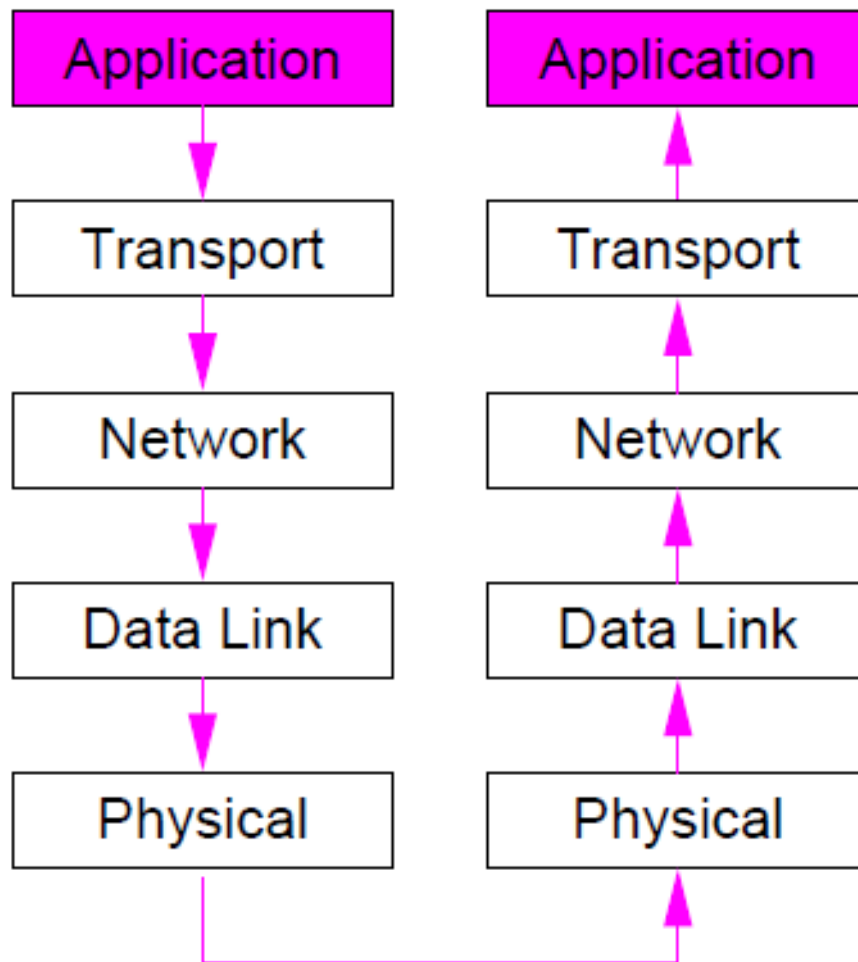


# Transport layer: Protection end-to-end between processes

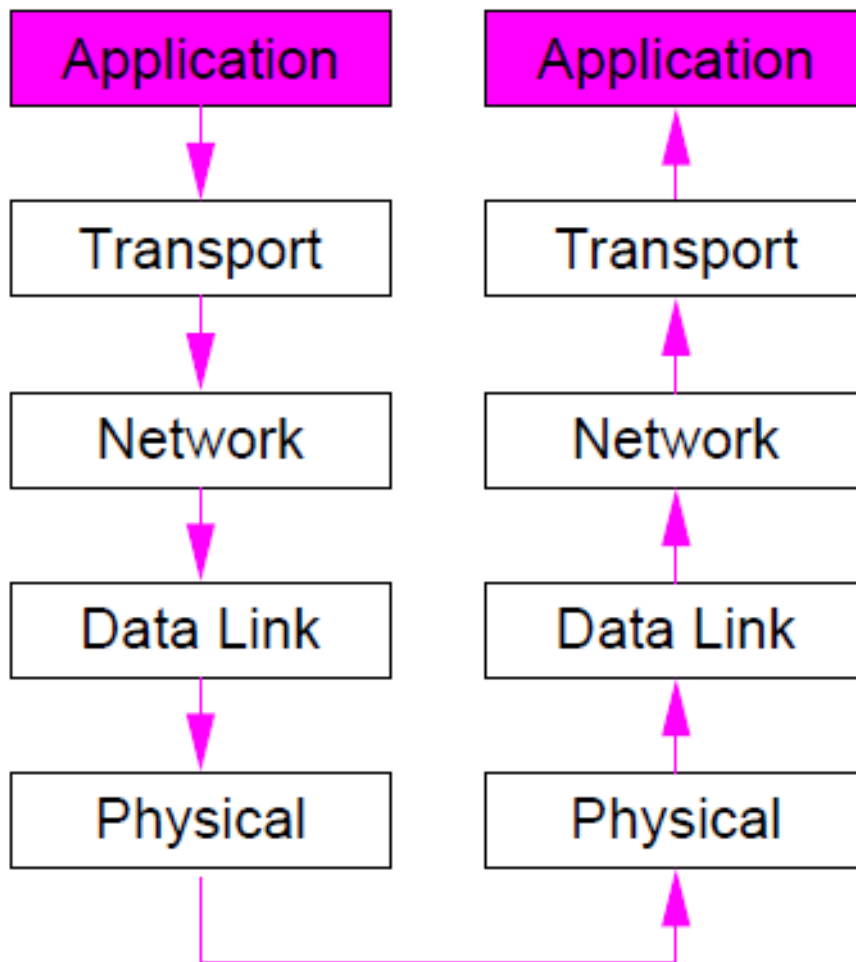


- Confidentiality: between apps/hosts/sites
- Integrity: between apps/hosts/sites
- Authentication: for user, host, site
- Replay protection: between apps/hosts/sites
- Traffic analysis protection: protocol/ host/site info. exposed
- Access control: user/host/site
- Transparency user and SW transparency possible
- Flexibility: HW or SW modifications
- Simplicity: good for site-to-site, not good for host-to-site

# Application layer: Security for a single application



# Application layer: Security for a single application



- Confidentiality: between users/apps
- Integrity: between users/apps
- Authentication: user
- Replay protection: between apps
- Traffic analysis protection: all but data exposed
- Access control: only data access secured
- Transparency: only user transparency
- Flexibility: SW modifications
- Simplicity depends on application

# VPN: then?

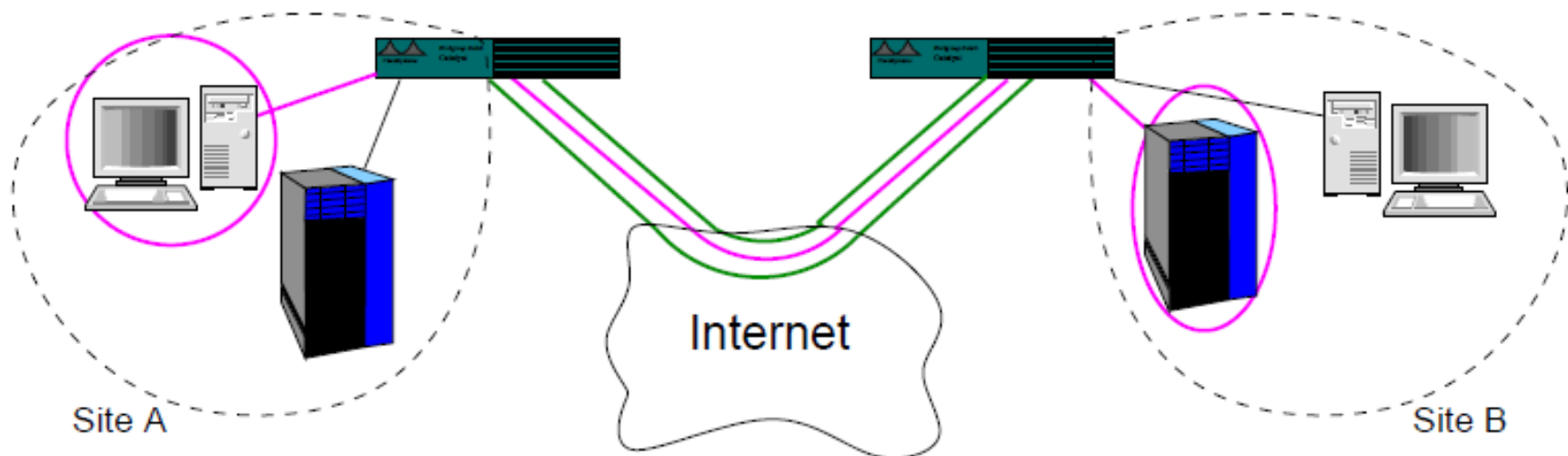
- It looks best to introduce security in the
  - Transport layer
  - Network layer
- These are the most popular choices for VPNs
- Other options:
  - Secure Application layer protocols: only protect a single application, but are often used for specialized purposes, e.g. S/MIME or PGP for secure e-mail
  - Secure Data Link layer protocols: are mostly used with PPP or other modem-based communication. e.g. PPTP, L2TP, LTF



# SSL Tunneling

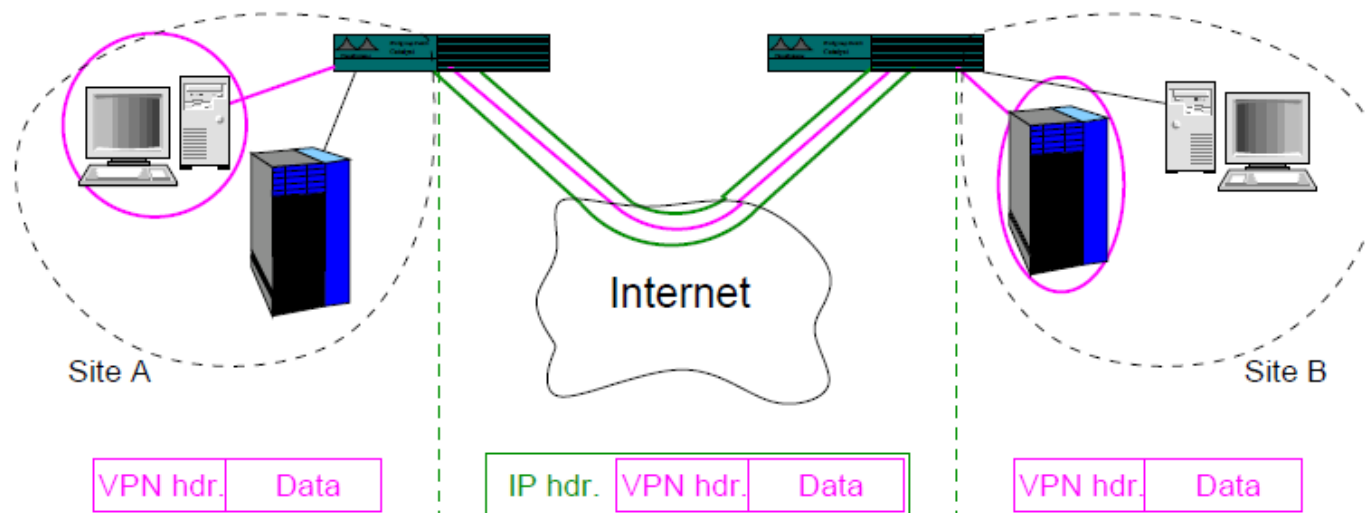
# Tunneling

- Operation of a network connection on top of another network connection
- It allows two hosts or sites to communicate through another network that they do not want to use directly

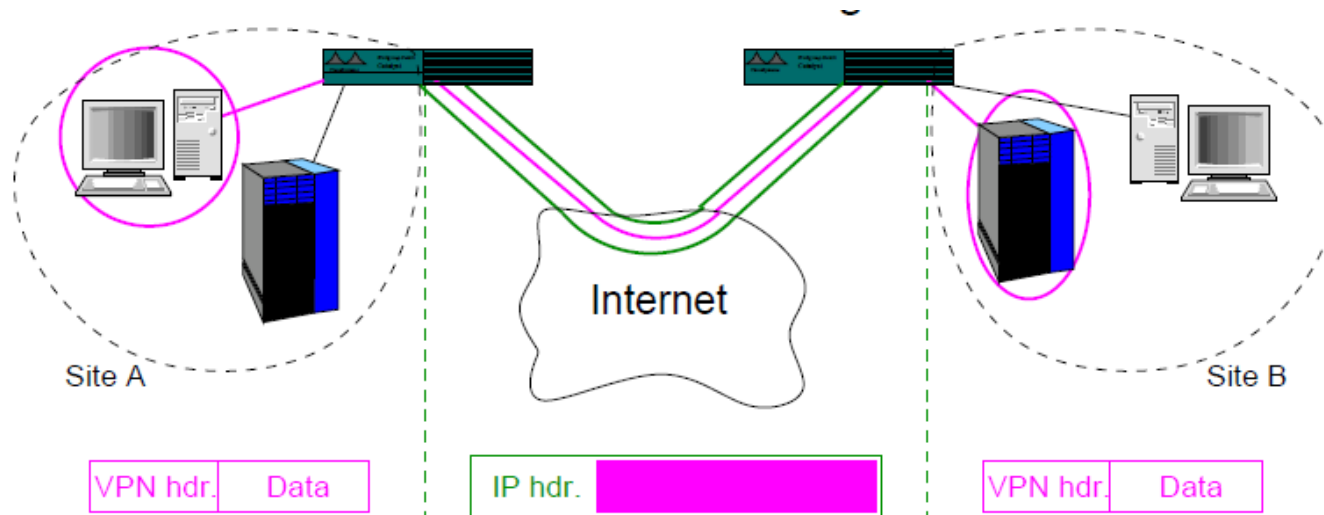


# Site-to-site tunneling

- Enables a PDU to be transported from one site to another without its contents being processed by hosts on the route.
- Idea: Encapsulate the whole PDU in another PDU sent out on the network connecting the two sites.
  - Encapsulation takes place in edge router on src. site.
  - Decapsulation takes place in edge router on dst. site.
- Note that the host-to-host communication does not need to use IP



# Secure tunneling



- Enables a PDU to be transported from one site to another without its contents being seen or changed by hosts on the route.
- Idea: Encrypt the PDU, and then encapsulate it in another PDU sent out on the network connecting the two sites.
  - Encryption can take place in edge router on src. site.
  - Decryption can take place in edge router on dst. site.
- Note: dst. address in IP header is for dst. edge router.



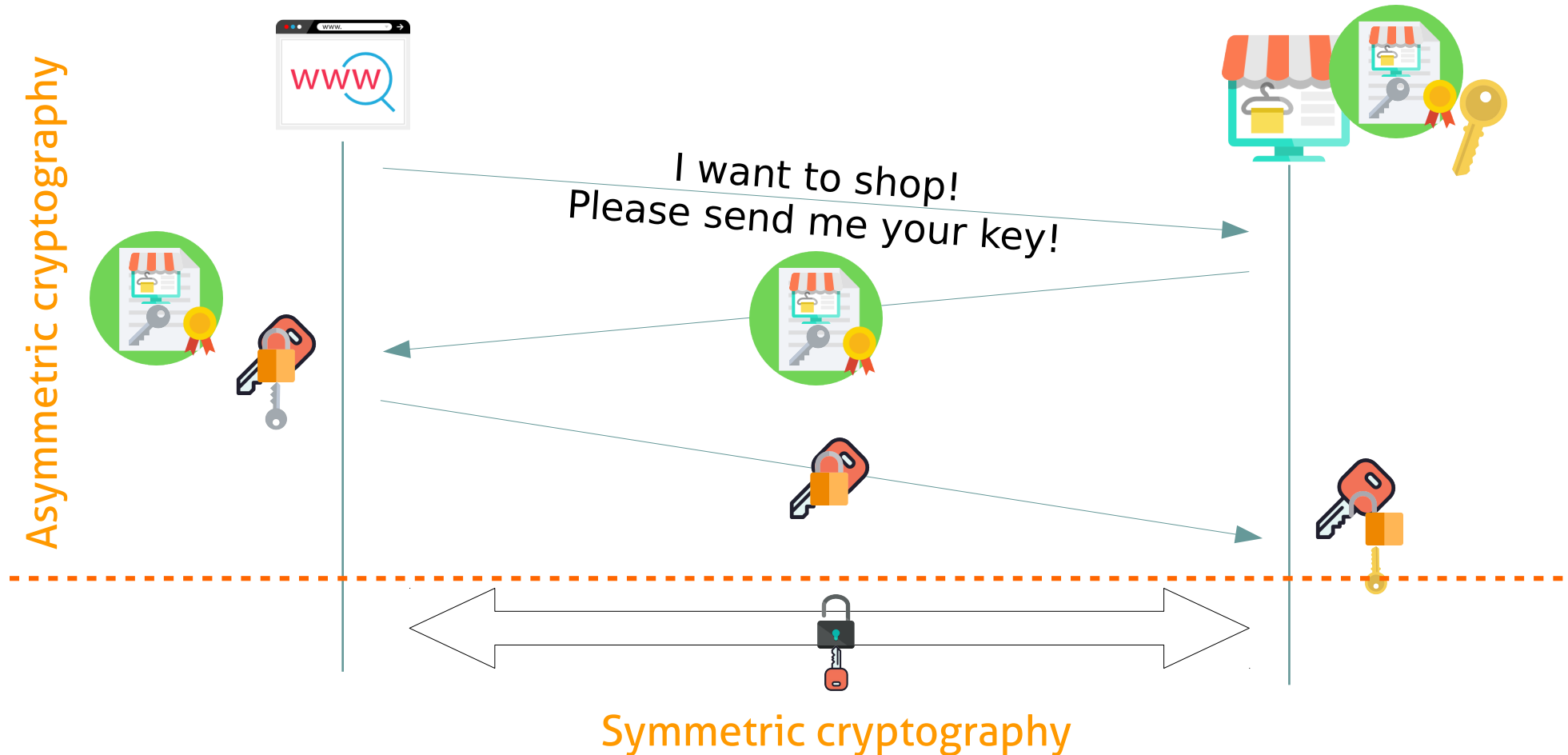
# Tunneling for VPNs

- Tunneling offers the basic method for providing a VPN.
- Where in the network architecture to initiate and terminate the tunnel:
  - Router/firewall?
  - Special box?
  - Host?
  - Application?
- Which layer to do the tunneling in:
  - Transport layer?
  - Network layer?
- Other possibilities (see previous discussion)
- And of course: Is tunneling the only possible technique?

# Secure Socket Layer

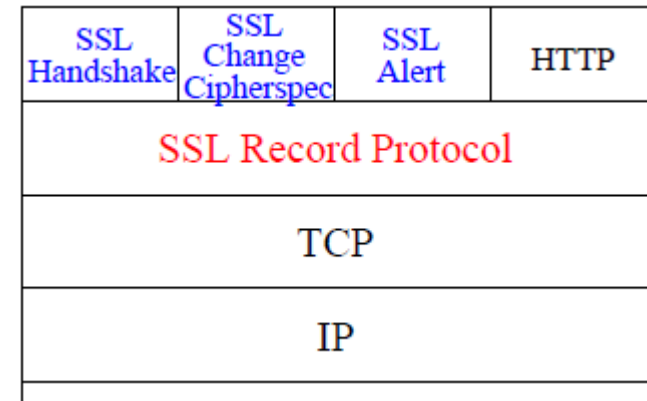
- SSL 3.0 has become TLS standard (RFC 2246) with small changes
- Applies security in the Transport layer.
- Originally designed (by Netscape) to offer security for client-server sessions.
- If implemented on boundary routers (or proxies), can provide a tunnel between two sites – typically LANs.
- Placed on top of TCP, so no need to change TCP/IP stack or OS.
- Provides secure channel (byte stream)
  - Any TCP-based protocol
  - https:// URLs, port 443
  - NNTP, SIP, SMTP...
- Optional server authentication with public key certificates
  - Common on commercial sites

# How HTTPS (HTTP on top of TLS) works



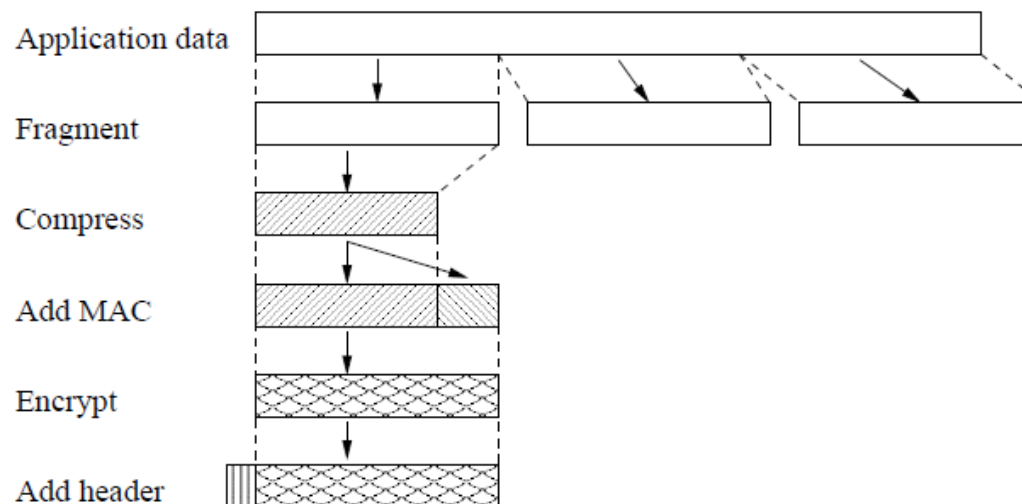
# SSL protocol Architecture

- Adds extra layer between T- and A-layers, and extra elements to A-layer
- Record Protocol: Protocol offering basic encryption and integrity services to applications
- Application Protocols: control operation of the record protocol
  - Handshake: Used to authenticate server (and optionally client) and to agree on encryption keys and algorithms.
  - Change cipher spec: Selects agreed keys and encryption algorithm until further notice.
  - Alert: Transfers information about failures.

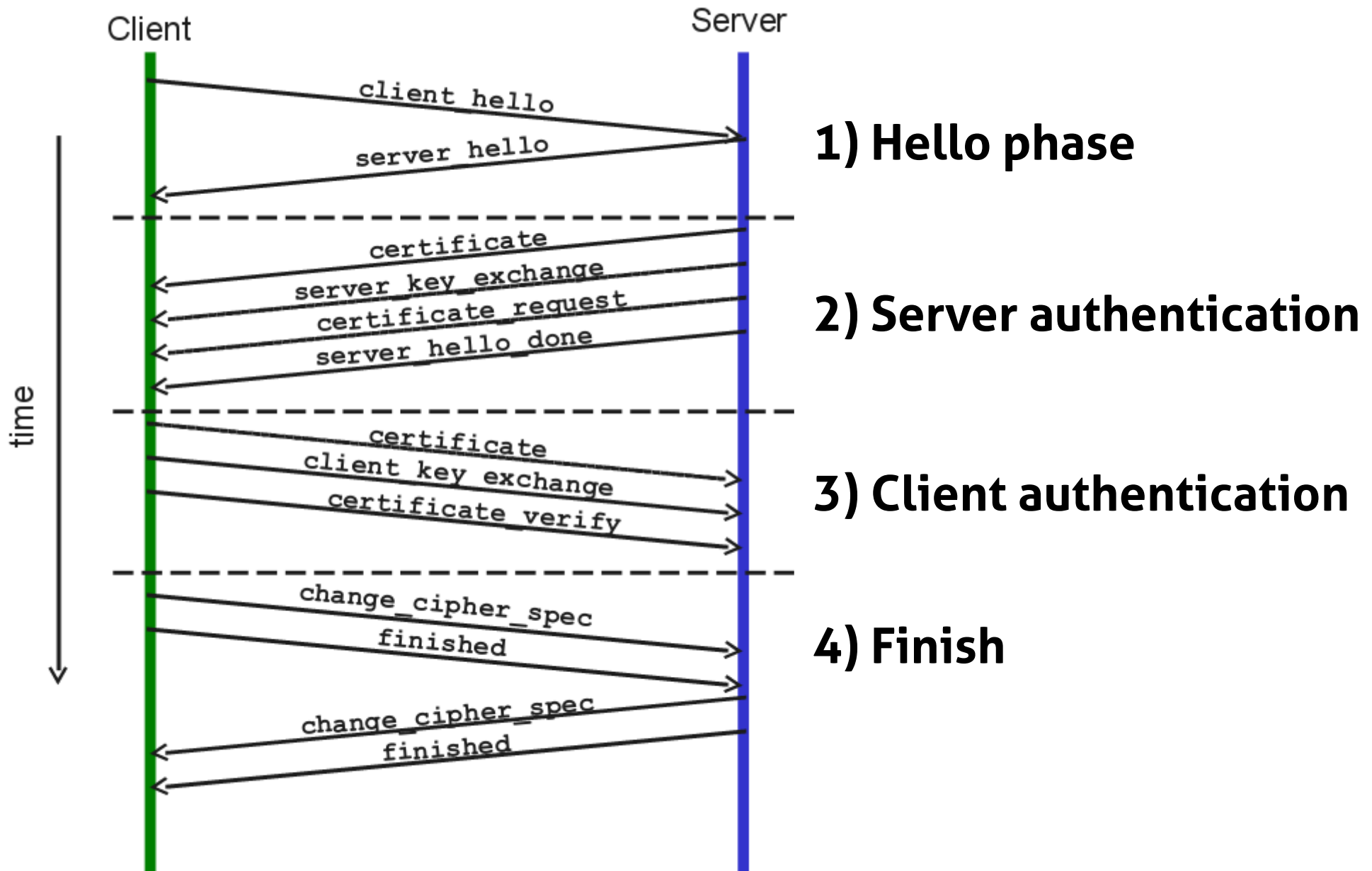


# SSL/TLS Record Protocol

- Offers to apply the following steps to PDUs:
  - 1 Fragmentation into blocks of  $\leq 2^{14}$  bytes.
  - 2 (optional) Lossless compression.
  - 3 Addition of a keyed MAC, using a shared secret MAC key.
  - 4 Encryption, using a shared secret encryption key.
  - 5 Addition of header indicating Application protocol in use.



# SSL Handshake



# SSL/TLS Handshake Protocol

4-phase “Client/Server” protocol to establish parameters of the secure connection (“Client” is the initiator):

- 1) Hello:** Establishment of security capabilities: Client sends list of possibilities, in order of preference. Server selects one, and informs Client of its choice. Parties also exchange random noise for use in key generation.
- 2) Server authentication and key exchange:** Server executes selected key exchange protocol (if needed). Server sends authentication info. (e.g. X.509 cert.) to Client.
- 3) Client authentication and key exchange:** Client executes selected key exchange protocol (mandatory). Client sends authentication info. to Server (optional).
- 4) Finish:** Shared secret key is derived from pre-secrets exch. in 2, 3. Change Cipher Spec. protocol is activated. Summaries of progress of Handshake Protocol are exchanged and checked by both parties.

# SSL/TLS Security Capabilities

- Conventionally expressed by a descriptive string, specifying:
  - Version of SSL/TLS
  - Key exchange algorithm
  - Grade of encryption (previous to TLSv1.1)
  - Encryption algorithm
  - Mode of block encryption (if block cipher used)
  - Cryptographic checksum algorithm
- Example: TLS\_RSA\_WITH\_AES\_128\_CBC\_SHA
  - TLS → (Latest version of) TLS
  - RSA → RSA key exchange
  - WITH → (merely filler...)
  - AES\_128 → 128-bit AES encryption
  - CBC → Cipher Block Chaining
  - SHA → Use HMAC-SHA digest



# Key exchange and authentication

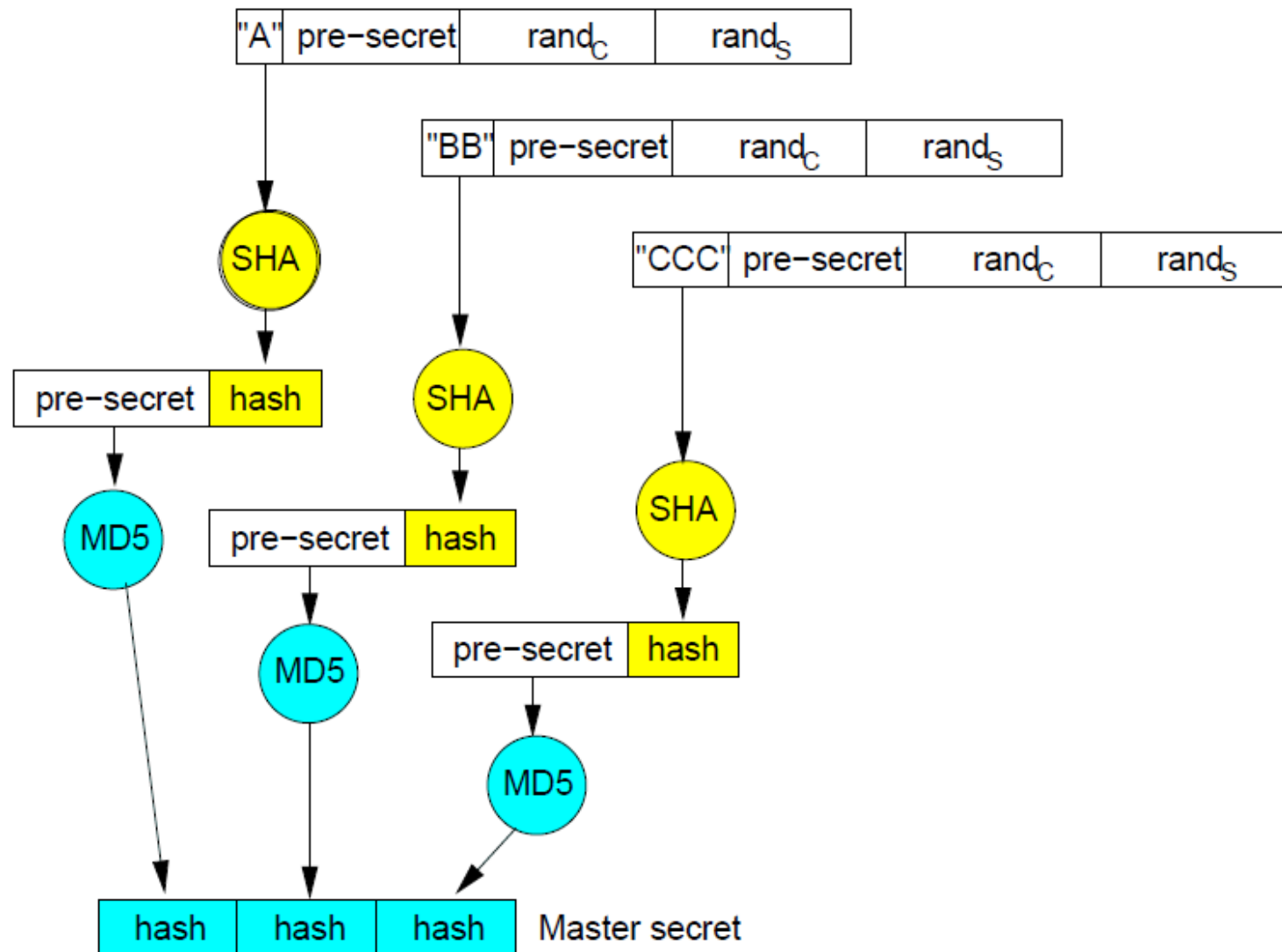
Possible ways of agreeing on secrets in TLS are:

- RSA: RSA key exch. (secret encrypted with recipient's publ. key)
- DHE RSA: Ephemeral Diffie-Hellman with RSA signatures
- DHE DSS: Ephemeral Diffie-Hellman with DSS signatures
- DH DSS: Diffie-Hellman with DSS certificates
- DH RSA: Diffie-Hellman with RSA certificates
- DH anon: Anonymous Diffie-Hellman (no authentication)
- NULL No key exch.

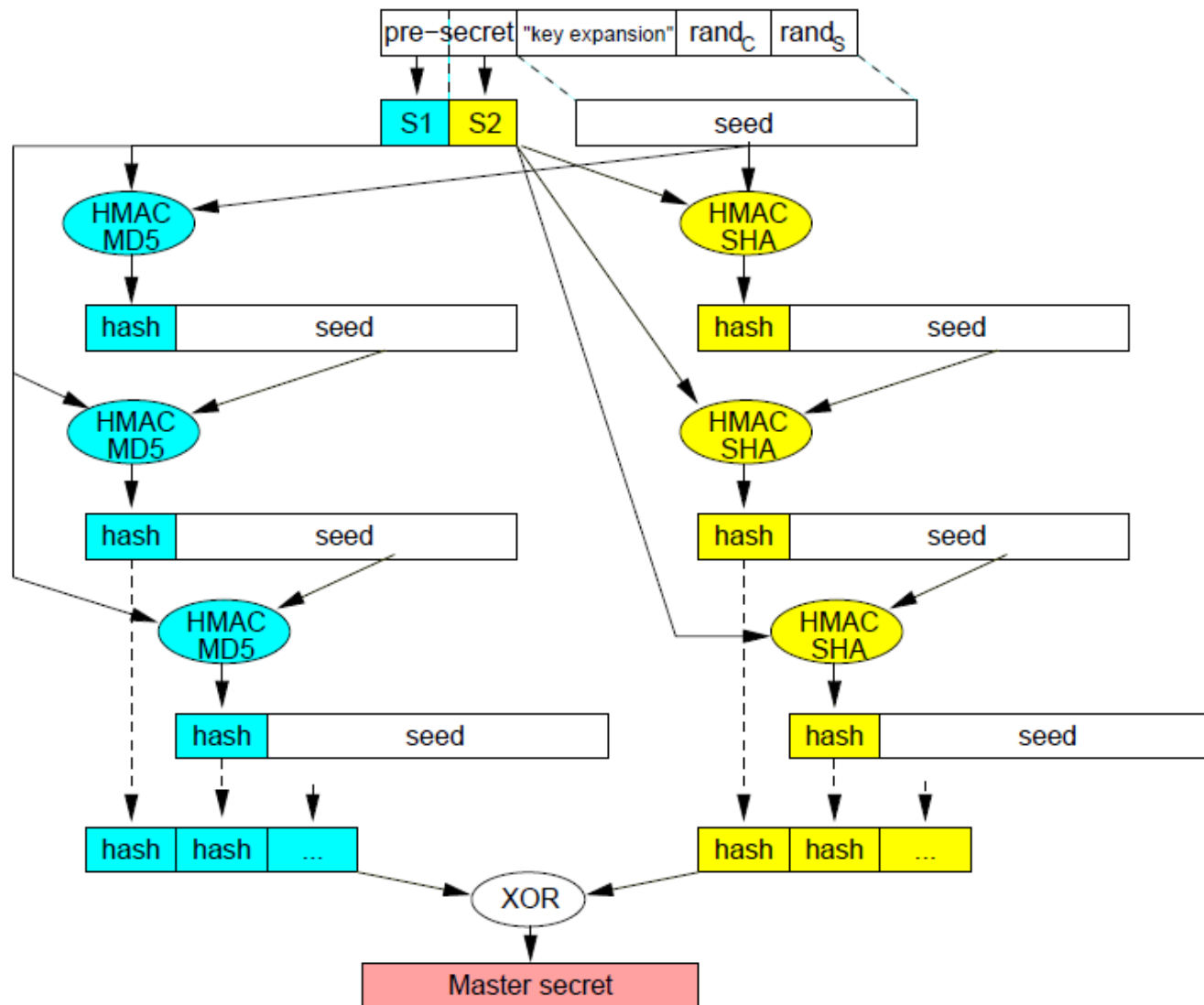
**Variant:** If followed by "EXPORT\_", weak encryption is used. (This option only available prior to TLSv1.1)

- **Note:** "Key exchange" only establishes a pre-secret! From this, a master secret is derived by a pseudo-random function (PRF). Shared secret encryption key is derived by expansion of master secret with another PRF. (In TLS several keys are derived for different purposes.)

# SSL Master Secret



# TLS Master Secret



# SSL/TLS Heartbeat

- It is an extension (RFC 6520) that allows to keep an established session alive
  - That is, as soon as the data exchange between two endpoints terminates, the session will also terminate
- To avoid the re-negotiation of the security parameters for establishing a secure session, we can keep using the same parameters even if there is no exchange of data
- It introduces two messages: **HeartbeatRequest** and **HeartbeatResponse**

# Heartbeat exchange

- When one endpoint sends a HeartbeatRequest message to the other endpoints, the former also starts what is known as the **retransmit timer**
  - During the time interval of the retransmit timer, the sending endpoint will not send another HeartbeatRequest message.
- An SSL/TLS session is considered to have terminated in the absence of a HeartbeatResponse packet within a time interval

# Heartbeat payload

- As a protection against a replay attack, HeartbeatRequest packets include a payload that must be returned without change by the receiver in its HeartbeatResponse packet
- The Heartbeat message is defined as

```
struct {  
    HeartbeatMessageType type;  
    uint16 payload_length;  
    opaque payload[HeartbeatMessage.payload_length];  
    opaque padding[padding_length];  
} HeartbeatMessage;
```

# Heartbleed bug

- Bug in OpenSSL library (4/4/2014)
- The receiver of request did not check that the **size of the payload in the packet** actually equaled the **value given** by the sender to the payload length field in the request packet
  - The attacker sends little data but sets the size to **max**
  - The receiver allcates that amount of memory for the response and copied **max bytes** from the mem location where the request packet was received
  - Then, the actual payload returned could potentially include objects in the memory that **had nothing to do** with the received payload
    - Objects could be private keys, passwords, and such...

# SSL VPN Architecture

- Two primary models:
- 1 SSL Portal VPN
  - Allow remote users to:
    - Connect to VPN gateway from a Web browser
    - Access services from Web site provided on gateway
- 2 SSL Tunnel VPN
  - Allow remote users to:
    - Access network protected by VPN gateway from
    - Web browser allowing active content.
  - More capabilities than portal VPNs, as easier to provide more services.



# SSL VPN functionalities

Most SSL VPNs offer one or more core functionalities:

- Proxying: Intermediate device appears as true server to client. E.g. Web proxy.
- Application Translation: Conversion of information from one protocol to another.
  - e.g. Portal VPN offers translation for applications which are not Web-enabled, so users can use Web browser to access applications with no Web interface.
- Network Extension: Provision of partial or complete network access to remote users, typically via Tunnel VPN.
  - Two variants:
    - Full tunneling: All network traffic goes through tunnel.
    - Split tunneling: Organisation's traffic goes through tunnel, other traffic uses remote user's default gateway.

# SSL VPN Securty Services

Typical services include:

- **Authentication** Via strong authentication methods, such as two-factor authent., X.509 certificates, smartcards, security tokens etc. May be integrated in VPN device or external authent. server.
- **Encryption** and integrity protection: Via the use of the SSL/TLS protocol.
- **Access control**: May be per-user, per-group or per-resource.
- **Endpoint security controls**: Validate the security compliance of clients attempting to use the VPN.
  - e.g. presence of antivirus system, updated patches etc.
- **Intrusion prevention**: Evaluates decrypted data for malicious attacks, malware etc.

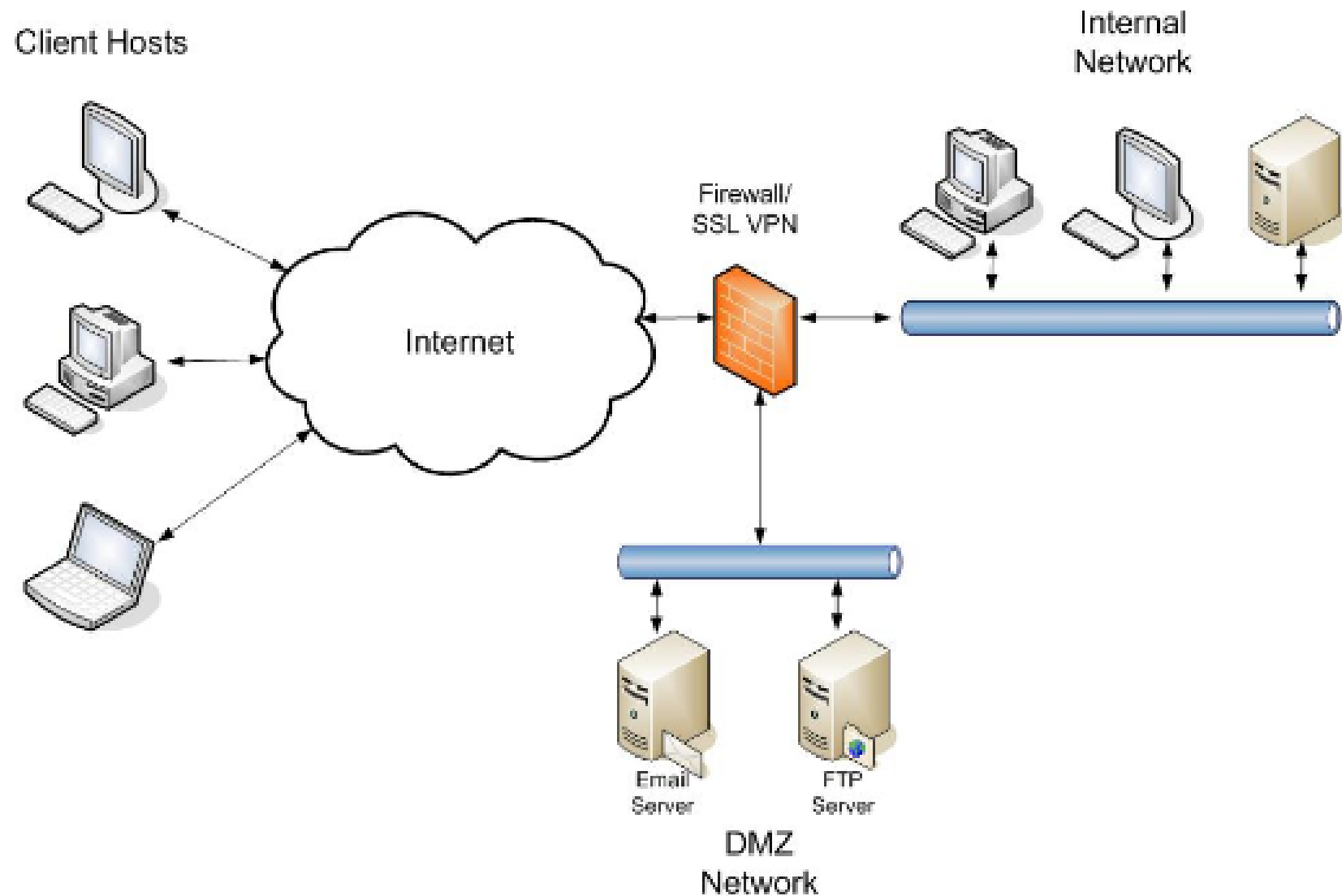


# VPN device placement

# SSL VPN Device placement

- Device placement is a challenge because it affects:
  - Security
  - Functionality
  - Performance
- Main options for placement:
  - VPN functionality in firewall
  - VPN device in internal network
  - Single-interface VPN device in DMZ
  - Dual-interface VPN device in DMZ
- Remember: Cryptographic protection only extends from VPN client systems to the SSL VPN device.

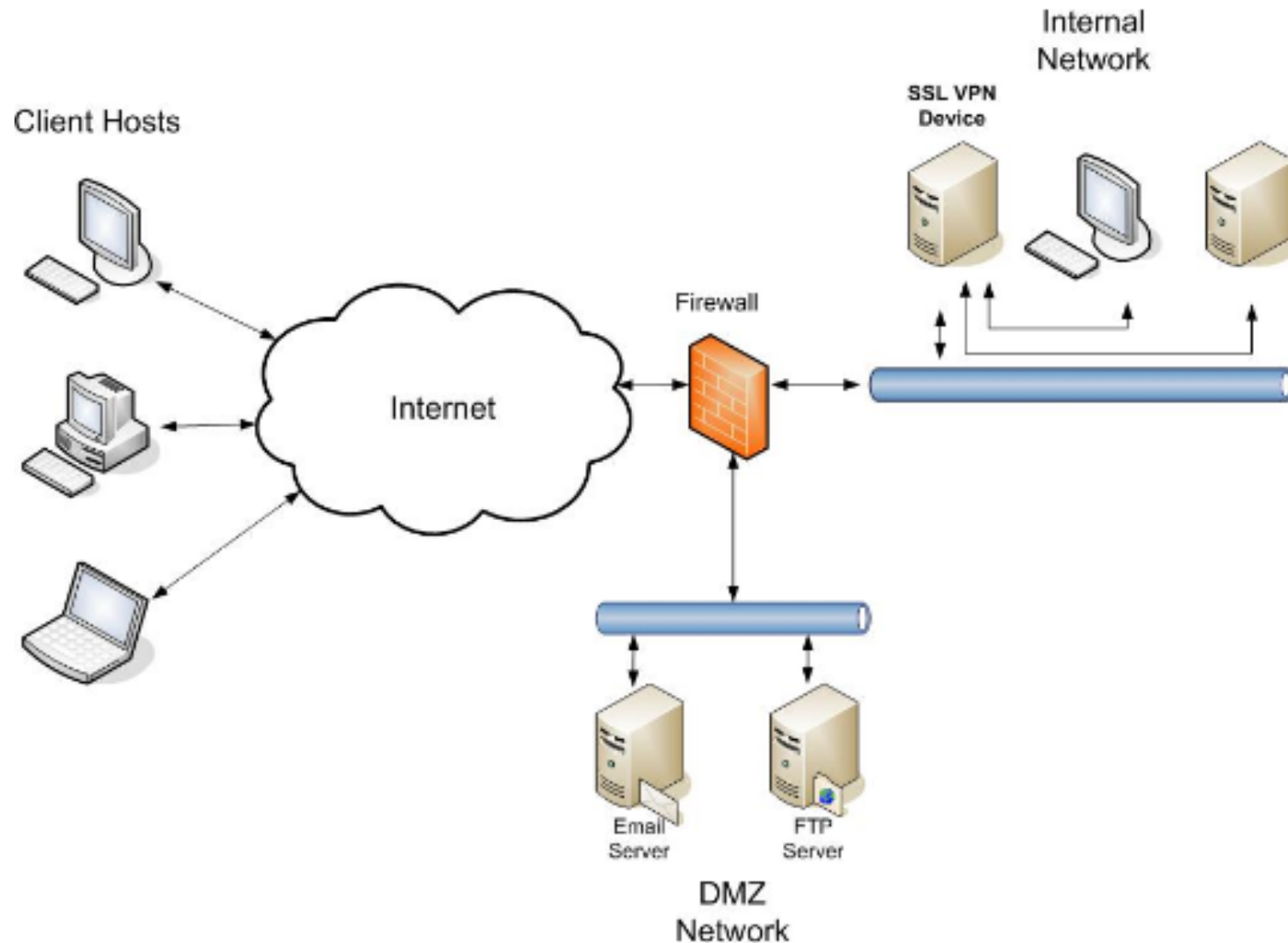
# Firewall with an SSL VPN



# VPN-enabled firewall

- The VPN device communicates directly with internal hosts
- Advantages
  - No holes in FW between external VPN device and internal network.
  - Traffic between device and internal network must go through FW.
  - Simple network administration since only one “box” to administer.
- Disadvantages
  - Limited to VPN functionality offered by FW vendor.
  - FW directly accessible to external users via port 443.
  - Adding VPN functionality to FW can introduce vulnerabilities.
- Note: TCP port 443 (standard) must be open on external FW interface, so clients can initiate connections.

# SSL VPN in internal network

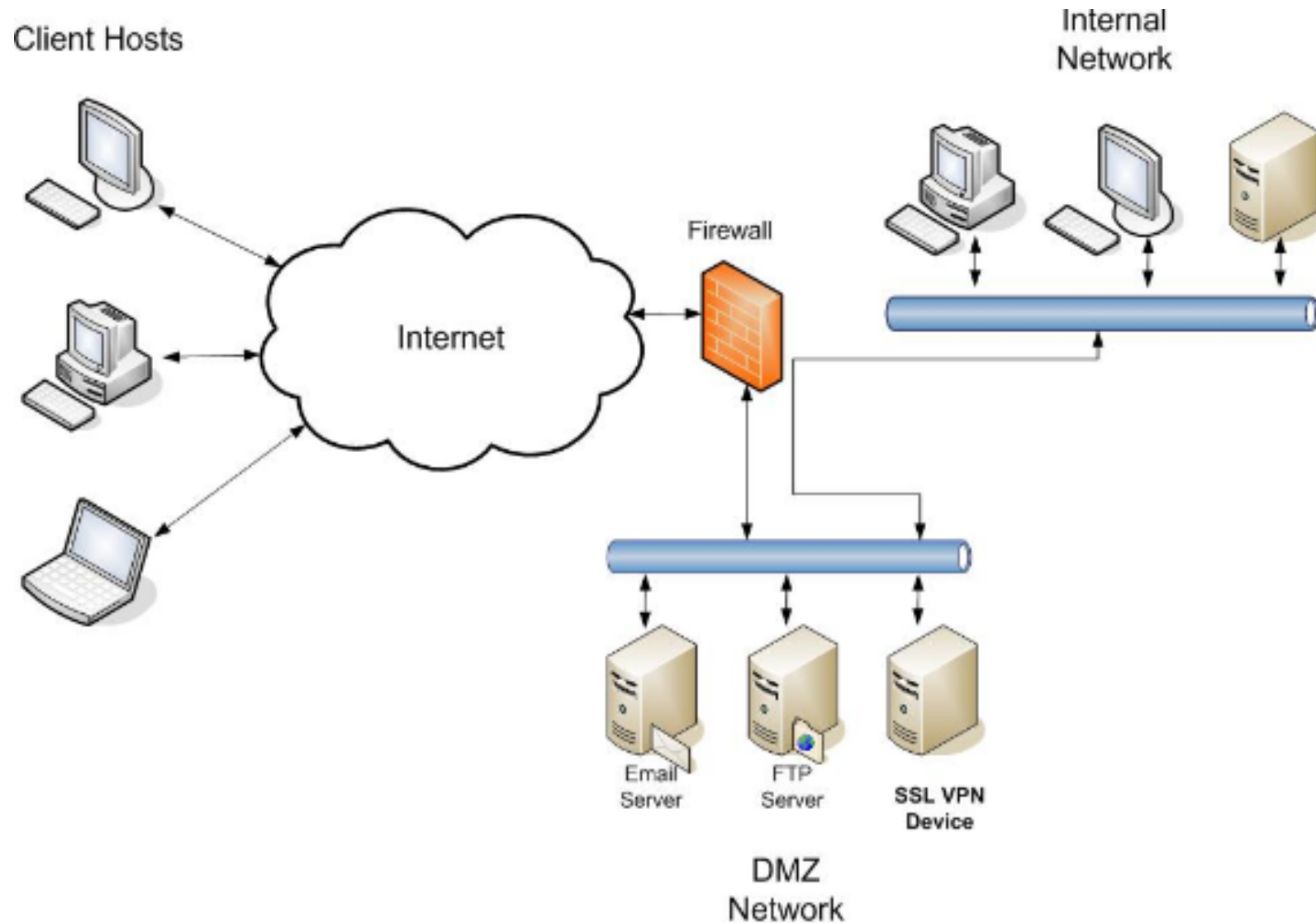


# VPN internal

- Advantages
  - Only single rule for single address to be added to FW.
  - No “holes” needed in FW between VPN device and internal network.
  - VPN traffic is behind FW, so protected from attacks by machines in DMZ.
- Disadvantages
  - VPN traffic passes through FW on tunnel, so it is not analyzed.
  - Unsolicited traffic can be sent into internal network from outside to internal VPN device.
  - Internal network is compromised if VPN device is compromised.
- Note: TCP port 443 (standard) opened on FW for the address of the device.



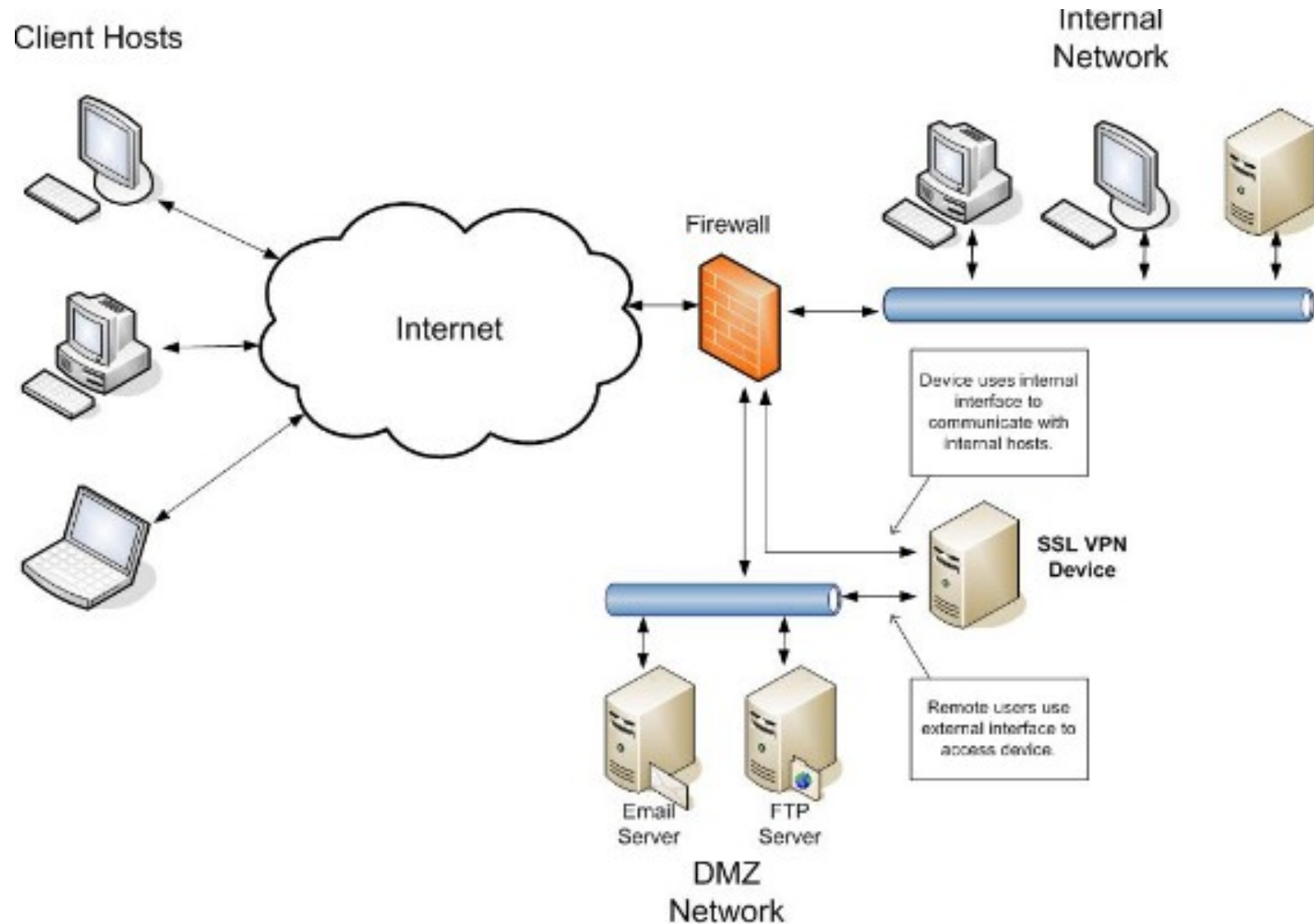
# SSL VPN In DMZ



# DMZ with VPN

- Advantages
  - Internal network protected against compromised VPN device.
  - Traffic between device and internal network must go through FW.
  - IDS in DMZ can analyze traffic destined for internal network.
- Disadvantages
  - Numerous ports open in FW between device and internal hosts.
  - Decrypted traffic from device to internal network must be sent through DMZ.
  - FW bypassed when user traffic is destined for hosts in DMZ.
- Note: TCP port 443 (standard) opened on FW for the address of the device

# Dual interfaces VPN device in DMZ



# VPN with two interfaces in DMZ

- Clients connect to external device interface, internal traffic uses internal interface.
- Advantages
  - All advantages of placing VPN device DMZ.
  - Unencrypted traffic to internal hosts is protected from other hosts in DMZ.
  - Only FW interface connected to device's internal interface needs to permit traffic from VPN device.
- Disadvantages
  - Numerous ports open in FW between device and internal hosts.
  - May introduce additional routing complexity.
  - FW bypassed if split tunneling is not used and user traffic is destined for hosts in DMZ



# IPSec

# IPsec

- A Network Layer protocol suite for providing security over IP.
- Part of IPv6; an add-on for IPv4.
- Can handle all three possible security architectures:

| Feature  | Gateway-to-Gateway | Host-to-Gateway              | Host-to-Host                 |
|--|--------------------|------------------------------|------------------------------|
| Protection between client and local gateway                          | No                 | N/A (client is VPN endpoint) | N/A (client is VPN endpoint) |
| Protection between VPN endpoints                                     | Yes                | Yes                          | Yes                          |
| Protection between remote gateway and remote server (behind gateway) | No                 | No                           | N/A (client is VPN endpoint) |
| Transparency to users  | Yes                | No                           | No                           |
| Transparency to users' systems                                       | Yes                | No                           | No                           |
| Transparency to servers  | Yes                | Yes                          | No                           |

# IPsec services

- Basic functions, provided by separate (sub-)protocols:
  - Authentication Header (AH): Support for data integrity and authentication of IP packets.
  - Encapsulated Security Payload (ESP): Support for encryption and (optionally) authentication.
  - Internet Key Exchange (IKE): Support for key management etc.

| Service                          | AH | ESP (encrypt only) | ESP(encrypt+authent.) |
|----------------------------------|----|--------------------|-----------------------|
| Access Control                   | +  | +                  | +                     |
| Connectionless integrity         | +  |                    | +                     |
| Protection between VPN endpoints | +  |                    | +                     |
| Data origin authentication       | +  |                    | +                     |
| Reject replayed packets          |    | +                  | +                     |
| Payload confidentiality          |    | +                  | +                     |
| Metadata confidentiality         |    | partial            | partial               |
| Traffic flow confidentiality     |    | (*)                | (*)                   |

# IPsec Security Associations

- Think of it as an IPsec connection: all of the parameters needed, like crypto algorithms (AES, SHA1, etc.), modes of operation (CBC, HMAC, etc.), key lengths, traffic to be protected, etc.
- Both sides must agree on the SA for secure communications to work
- For a two-way communication, two SAs must be defined.
- SA parameters must be negotiated (using IKE) between sender and receiver before secure communication can start.
- Each SA is identified by:
  - Security Parameters Index (SPI): 32-bit integer chosen by sender. Enables receiving system to select the required SA.
  - Destination Address: Only unicast IP addresses allowed!
  - Security Protocol Identifier: AH or ESP.



# IPsec modes

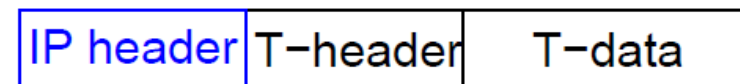
- Transport Mode
  - Provides protection for a T-layer packet embedded as payload in an IP packet.
- Tunnel Mode
  - Provides protection for an IP packet embedded as payload in an IP packet.

|                  | Transport Mode SA  | Tunnel Mode SA  |
|------------------|--|---|
| AH               | Authenticate IP payload and selected parts of IP header and IPv6 extension headers.        | Authenticate entire inner IP packet and selected parts of outer IP header and outer IPv6 extension headers. |
| ESP              | Encrypt IP payload + any IPv6 extension headers after ESP header.                          | Encrypt inner IP packet.  |
| ESP + authentic. | Encrypt IP payload + any IPv6 extension headers after ESP header. Authenticate IP payload. | Encrypt and authenticate inner IP packet.   |

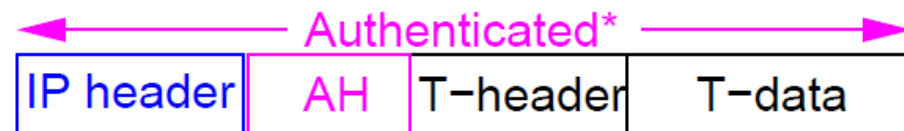
# Authentication with IPv4

- AH header inserted after the outermost IP header – depending on whether Transport or Tunnel mode is used.
  - Do not forget that integrity check (and thus authentication) does not cover any mutable, unpredictable header fields.

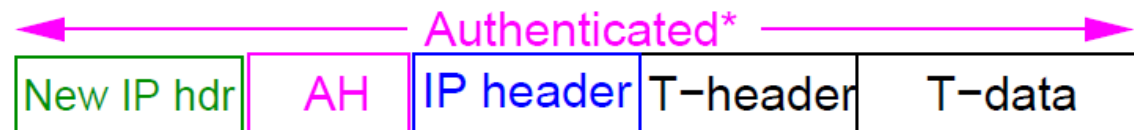
Without IPsec



AH Transport mode



AH Tunnel mode



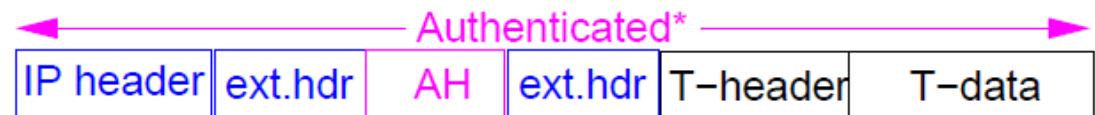
# Authentication with IPv6

- AH header inserted after the outermost IP header – depending on whether Transport or Tunnel mode is used.
  - Do not forget that integrity check (and thus authentication) does not cover any mutable, unpredictable header fields.

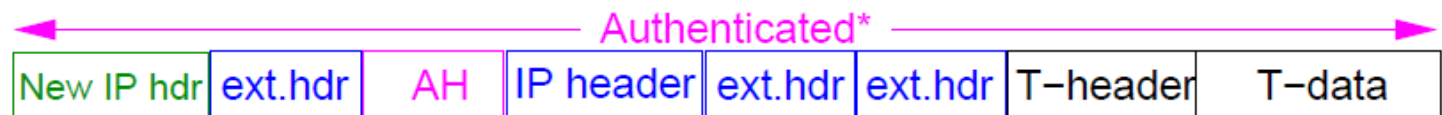
Without IPsec



AH Transport mode



AH Tunnel mode



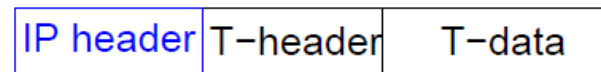
# Authentication Header

- One of the (many possible) IP header fields. Contains:
  - Next Header: Type of following header field.
  - Payload Length: (Length - 2), in 32-bit words, of AH.
  - SPI: Identifies SA in use.
  - Sequence Number: Monotonically increasing packet counter value.
  - Authentication Data (AD): (variable length) HMAC based on MD5 or SHA-1 cryptographic hashing algorithm, or AES-CBC, evaluated over:
    - Immutable or predictable IP header fields. (Other fields assumed zero when MAC is calculated.)
    - Rest of AH header apart from AD field.
    - All embedded payload (from T-layer or embedded IP packet), assumed immutable.
- Immutable fields do not change as the packet traverses the network.
  - Example: Source address.
- Mutable but predictable fields may change, but can be predicted.
  - Example: Destination address.
- Mutable, unpredictable fields include Time-to-live, Header checksum.

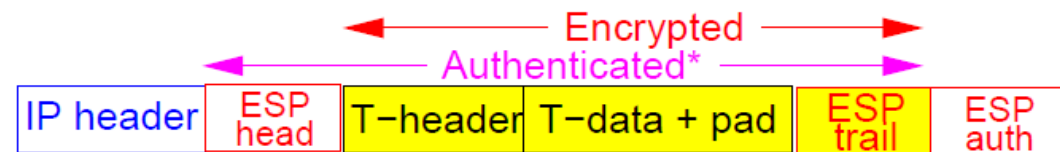
# ESP with IPv4

- ESP header inserted after the outermost IP header – depending on whether Transport or Tunnel mode is used:
- Padding is added to end of T-layer payload to give (a certain amount) of traffic analysis protection.
- ESP trailer and (optional) ESP authentication field added after the end of the padded T-layer payload.
- As usual, integrity check (and thus authentication) does not cover any mutable, unpredictable header fields.

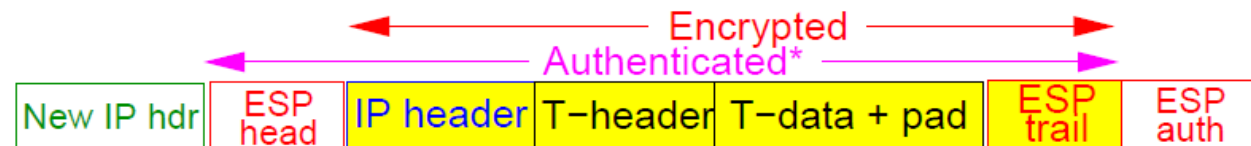
Without IPsec



ESP Transport mode



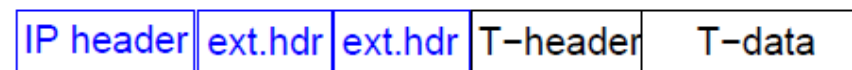
ESP Tunnel mode



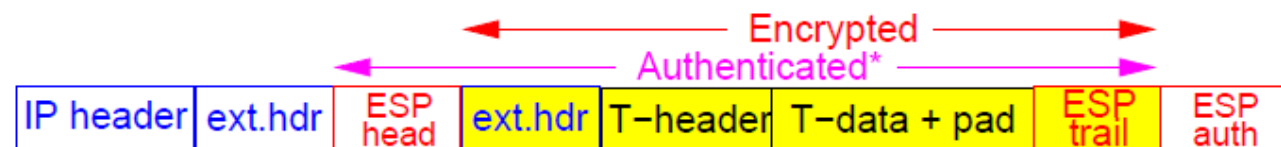
# ESP with IPv6

- ESP header inserted after the outermost IP header – depending on whether Transport or Tunnel mode is used:
- Padding is added to end of T-layer payload to give (a certain amount) of traffic analysis protection.
- ESP trailer and (optional) ESP authentication field added after the end of the padded T-layer payload.
- As usual, integrity check (and thus authentication) does not cover any mutable, unpredictable header fields.

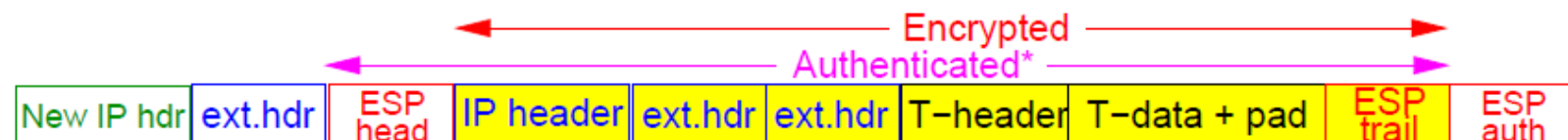
Without IPsec



ESP Transport mode



ESP Tunnel mode



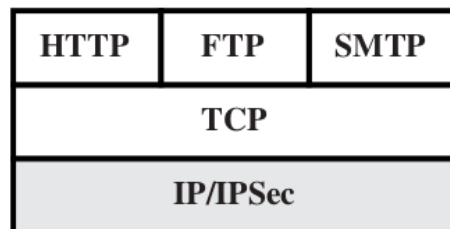
# Encryption + Authentication

A common combination, can be achieved by:

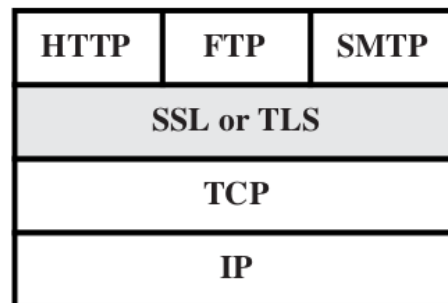
- 1) ESP with Authentication. First apply ESP to data, then add AH field.  
Two subcases:
  - 1) Transport mode: E+A apply to IP payload, but IP header not protected.
  - 2) Tunnel mode: E+A apply to entire inner packet.
- 2) Transport Adjacency. Use bundled SAs, first ESP, then AH.
- 3) Encryption covers original IP payload. Authentication covers ESP + original IP header, including source and destination IP addresses
- 4) Transport-Tunnel bundle. Used to achieve authentication before encryption, for example via inner AH transport SA and outer ESP tunnel SA.
- 5) Authentication covers IP payload + IP immutable header. Encryption is applied to entire authenticated inner packet.

# IPsec vs TLS

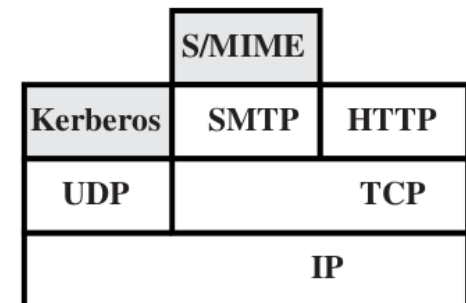
- TLS much more flexible because is in the upper levels
- TLS also provides application end-to-end security, best for web applications → HTTPS
- IPsec has to run in kernel space
- IPsec much more complex and complicated to manage with



(a) Network level



(b) Transport level



(c) Application level



# SSL again

- Crypto is insufficient for Web security
- One issue: linkage between crypto layer and applications
- Trust: what does the server really know about the client?
  - Unless client-side certificates are used, absolutely nothing
  - SSL provides a secure pipe. “Someone” is at the other end; you don’t know who
  - Usually there is no user authentication in SSL, but in the application layer!



# SSL: the Client's Knowledge of the Server

- The client receives the server's certificate
- Does it help?
- A certificate means that someone has attested to the binding of some name to a public key
- Who has done the certification? Is it the right name?
  - Every browser has a list of built-in certificate authorities
  - Hundreds of certificate authorities! Do you trust them all to be honest and competent? Do you even know them all?
- It's all a matter of trust...

# Conclusions on SSL

- The cryptography itself seems correct
  - The human factors are dubious
    - Most users don't know what a certificate is, or how to verify one
- Even when they do know, it's hard to know what it should say in any given situation
- There is no rational basis for deciding whether or not to trust a given CA



# That's all for today

- Questions?
- See you on Thursday (12:40)
- Resources:
  - Chapter 24 textbook
  - "Virtual private networking", Gilbert Held, Wiley ed.
  - [http://www.tcpipguide.com/free/t\\_IPSecurityIPSecProtocols.htm](http://www.tcpipguide.com/free/t_IPSecurityIPSecProtocols.htm)
  - "Guide to IPsec VPNs", NIST800-77
  - "Guide to SSL VPNs", NIST-SP800-113