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1. Number Theory**a. Sieve**

```

#define LIM 10000005
#define pb push_back
bool status[LIM]; /// NEED A GLOBAL
ARRAY
vector <int> prime;
void make_primes(int N)
{
    int i, j, sqrtN;
    status[0] = 1, status[1] = 1;
    sqrtN = int( sqrt((double) N )) + 1;
    /// have to check primes up to (sqrt(N))
    for(i = 4; i <= N; i+=2) status[i] =
1;
    for(i = 3; i <= sqrtN; i+= 2){ ///
so, i is a prime, so, discard all the
multiples
        if(status[i] == 0){
            for(j = i*i; j <= N; j+=i+i)
/// j = i * i, because it's the first
number to be colored
                status[j] = 1;
        }
    }
    for(i = 0; i < LIM; i++)
if(status[i] == 0) prime.pb(i);
    ///for(i = 0; i < 10000005; i++) cout
<< v[i] << endl;
}

```

b. Segmented Sieve

```

void sg_sieve(ll m, ll n)
{
    ll i;
    memset(status2,0,sizeof status2);
    for(i = 0; v[i]*v[i] <= n; i++){
        ll tem = (m+v[i]-1)/v[i];
        ll sg = tem*v[i];
        if(sg == v[i]) sg += v[i];
        for(ll st = sg; st <= n; st +=
v[i]){
            status2[st-m+1] = 1;
        }
    }
    for(i = m ; i <= n; i++){
        if(status2[i-m+1] == 0 and i !=
1) cout << i << endl;
    }
    cout << endl;
}

```

c. Bitwise sieve 64bit

```

#define ll long long
#define mx 100000000
#define one (1LL)
vector<ll> prime;
ll status[(mx / 64) + 2];

bool checkprime(ll n) {
    if(status[n / 64] & (one << (n %
64))) return false;
    return true;
}

void setbit(ll n) {
    status[n / 64] = status[n / 64] |
(one << (n % 64));
}

void mkprime(ll n) {
    setbit(0);
    setbit(1);

    prime.push_back(2);
    ll sqrtn = sqrt(n);
    for(ll i = 3; i <= sqrtn; i += 2) {
        if(checkprime(i)) {
            for(ll j = i * i; j <= n; j
+= i + i) {
                setbit(j);
            }
        }
    }

    for(ll i = 3; i <= n; i += 2) {
        if(checkprime(i))
prime.push_back(i);
    }
}

```

d. Prime factorization

```

#define ll long long /// MACRO NEEDED
#define LIM 10005 /// MACRO NEEDED
#define mod 100000007 /// MACRO NEEDED
#define pb push_back /// MACRO NEEDED
#define ff first /// MACRO NEEDED
#define ss second /// MACRO NEEDED

bool status[LIM]; /// NEED A GLOBAL
ARRAY
vector <int> prime;

```

```
vector <pair<int,int> > prime_pow;

void make_prime(int N)
{
    int i, j, sqrtN;
    status[0] = 1, status[1] = 1;
    sqrtN = int( sqrt((double) N )) + 1;
    /// have to check primes up to (sqrt(N))
    for(i = 4; i <= N; i+=2) status[i] = 1;

    for(i = 3; i <= sqrtN; i+= 2){ ///
so, i is a prime, so, discard all the
multiples
        if(status[i] == 0){
            for(j = i*i; j <= N; j+=i+i)
/// j = i * i, because it's the first
number to be colored
                status[j] = 1;
        }
    }
    for(i = 0; i < LIM; i++)
if(status[i] == 0) prime.pb(i);
    ///for(i = 0; i < 10000005; i++) cout
<< v[i] << endl;
}

void factor(int n)
{
    for(int i = 0; i < prime.size() &&
prime[i]*prime[i] <= n; i++){
        int cnt = 0;
        while(n%prime[i] == 0){
            n /= prime[i];
            cnt++;
        }
        if(cnt) prime_pow.pb({prime[i],
cnt});
    }
    if(n != 1) prime_pow.pb({n, 1});
}

make_prime() /// NEED TO CALL THE
FUNCTION FROM MAIN
factor()      /// NEED TO CALL THE
FUNCTION FROM MAIN
```

e. Number of Divisor Sieve

Method(NOD)

```
bool status[LIM+5];
ll nod[LIM+5];
void NOD()
{
    ll i, j, tem, cnt;
```

```
for(i = 0; i <= LIM; i++)
    nod[i] = 1;
for(i = 4; i <= LIM; i+=2){
    status[i] = 1;
    tem = i;
    cnt = 0;
    while(tem%2 == 0){
        tem /= 2;
        cnt++;
    }
    nod[i] *= (cnt+1);
}
for(i = 3; i <= LIM; i+=2){
    if(status[i] == 0){
        for(j = i+i; j <= LIM; j +=
i){
            status[j] = 1;
            tem = j;
            cnt = 0;
            while(tem%i == 0){
                tem /= i;
                cnt++;
            }
            nod[j] *= (cnt+1);
        }
    }
}
nod[0] = 0, nod[1] = 1;
}
```

f. Sum of NOD

Problem : Find the snod of $n = 100000$?
 $\text{nod}(1) + \text{nod}(2) + \text{nod}(3) + \dots + \text{nod}(n) =$??

//code of $O(n)$ solution

```
int snod( int n ) {
    int res = 0;
    for ( int i = 1; i <= n; i++ ) {
        res += n / i;
    }
    return res;
}
```

//code of $O(\sqrt{n})$ solution

```
int snod( int n ) {
    int res = 0;
    int u = sqrt(n);
    for ( int i = 1; i <= u; i++ ) {
        res += ( n / i ) - i; //Step 1
    }
    res *= 2; //Step 2
    res += u; //Step 3
```

```
    return res;
}
```

g. SOD (Sum of Divisor)

Problem : Find the sod of n?

sod(12) = 1 + 2 + 3 + 4 + 6 + 12 = 28;

if $N = p_1^{a_1} \times p_2^{a_2} \times \dots \times p_k^{a_k}$, then

$$SOD(N) = (p_1^0 + p_1^1 + p_1^2 + \dots + p_1^{a_1}) \times (p_2^0 + p_2^1 + p_2^2 + \dots + p_2^{a_2}) \times \dots \times (p_k^0 + p_k^1 + p_k^2 + \dots + p_k^{a_k})$$

//code of O(sqrt(n)) solution

```
int sod(int n){
    int sqrtn = sqrt(n);
    int res = 0;
    for( int i = 1; i <= sqrtn; i++){
        if(n % i == 0){
            res += i; //"i" is a divisor
            res += n / i; //"n/i" is
also a divisor
        }
    }
    if(sqrtn * sqrtn == n) res -= sqrtn;
//same number counted two times on first
loop
    return res;
}
```

//code of sod using prime factorization

```
int sod(int n){
    int res = 1;
    int sqrtn = sqrt(n);
    for(int i = 0; i < prime.size() &&
prime[i] <= sqrtn; i++){
        if(n % prime[i] == 0){
            int tempSum = 1; // Contains
value of (p^0+p^1+...p^a)
            int p = 1;
            while(n % prime[i] == 0){
                n /= prime[i];
                p *= prime[i];
                tempSum += p;
            }
            sqrtn = sqrt(n);
            res *= tempSum;
        }
    }
    if(n != 1){
        res *= (n + 1); // Need to
multiply (p^0+p^1)
    }
    return res;
}
```

h. GCD Extreme

```
#include <bits/stdc++.h>
#define ll long long
#define pb push_back
#define LIM 200001
using namespace std;
ll arr[LIM], ans[LIM];
vector<ll> v[LIM], prime;
void make_prime()
{
    ll i, j;
    for(i = 2; i < LIM; i++){
        ll f = 1;
        for(j = 2; j*j <= i; j++){
            if(i%j == 0) f = 0;
            if(f) prime.pb(i);
        }
    }
}
void ETF()
{
    ll i, j;
    for(i = 0; i < LIM; i++){
        arr[i] = i;
        for(i = 0; i < prime.size(); i++){
            for(j = prime[i]; j < LIM; j +=
prime[i])
                arr[j] =
(arr[j]/prime[i])*(prime[i]-1);
        }
    }
}
void divi()
{
    ll i, j;
    for(i = 2; i < LIM; i++){
        for(j = i+i; j < LIM; j += i){
            v[j].pb(i);
        }
    }
}
void GCD()
{
    ll i, j;
    for(i = 1; i < LIM; i++){
        //cout << i << endl;
        for(j = 0; j < v[i].size();
j++){
            ans[i] +=
v[i][j]*arr[i/v[i][j]];
        }
    }
}
```

```

        //cout << v[i][j] << " " <<
i/v[i][j] << " " << arr[i/v[i][j]] <<
endl;
    }
    ans[i] += arr[i];
}
for(i = 2; i < LIM; i++)
    ans[i] += ans[i-1];
}
int main()
{
    make_prime();
    ETF();
    divi();
    GCD();
    ll n;
    while(cin >> n){
        if(!n) break;
        cout << ans[n]-1 << endl;
    }
    return 0;
}

```

i. LCM Extreme

```

#include <bits/stdc++.h>
#define ull unsigned long long
#define pb push_back
#define LIM 3000005
using namespace std;
ull phi[LIM], ans[LIM];
void divi()
{
    ull i, j;

    for(i = 2; i < LIM; i++){
        for(j = i; j < LIM; j += i){
            ans[j] += j*((i*phi[i])/2);
        }
    }
}
void ETF()
{
    ull i, j;

    for(i = 0; i < LIM; i++)
        phi[i] = i;

    for(i = 2; i < LIM; i++){
        if(phi[i] == i){
            for(j = i; j < LIM; j += i)

```

```

        phi[j] =
        (phi[j]/i)*(i-1);
    }
}

//    for(i = 0; i < LIM; i++){
//        cout << i << ": " << phi[i] <<
endl;
//    }

}

void GCD()
{
    ull i;

    for(i = 1; i < LIM; i++)
        ans[i] += ans[i-1];
}

int main()
{
    //make_prime();
    ETF();
    divi();
    GCD();
    ull t, n, i, cas = 1;
    cin >> t;
    while(t--){
        scanf("%llu", &n);
        printf("Case %llu: %llu\n",
cas++, ans[n]);
    }
    return 0;
}

```

j. Extended GCD

```

void ext_gcd(int a, int b, int *X, int
*Y)
{
    int x, y, r, x1, y1, x2, y2, r1, r2,
q;

    x2 = 1, y2 = 0;
    x1 = 0, y1 = 1;
    for(r2 = a, r1 = b; r1 != 0; r2 =
r1, r1 = r, x2 = x1, x1 = x, y2 = y1, y1
= y){
        q = r2/r1;
        r = r2%r1;
        x = x2-q*x1;
        y = y2-q*y1;

```

```

    }
    *X = x2, *Y = y2;
}

```

k. Linear Diophantine Equation

```

ll gcd(ll a, ll b, ll& x, ll& y) {
    if (b == 0) {
        x = 1;
        y = 0;
        return a;
    }
    ll x1, y1;
    ll d = gcd(b, a % b, x1, y1);
    x = y1;
    y = x1 - y1 * (a / b);
    return d;
}

bool find_any_solution(ll a, ll b, ll c,
ll &x0, ll &y0, ll &g) {
    g = gcd(abs(a), abs(b), x0, y0);
    if (c % g) {
        return false;
    }

    x0 *= c / g;
    y0 *= c / g;
    if (a < 0) x0 = -x0;
    if (b < 0) y0 = -y0;
    return true;
}

void shift_solution(ll & x, ll & y, ll
a, ll b, ll cnt) {
    x += cnt * b;
    y -= cnt * a;
}

ll find_all_solutions(ll a, ll b, ll c,
ll minx, ll maxx, ll miny, ll maxy) {
    ll tem;
    if(a == 0 and b == 0){
        if(c == 0) return
(abs(minx-maxx)+1)*(abs(miny-maxy)+1);
        else return 0;
    }
    else if(a == 0){
        tem = c/b;
        if(c%b == 0 and tem >= miny and
tem <= by) return abs(minx-maxx)+1;
        else return 0;
    }
}

```

```

    else if(b == 0){
        tem = c/a;
        if(c%a == 0 and tem >= minx and
tem <= maxx) return abs(miny-maxy)+1;
        else return 0;
    }
    ll x, y, g;
    if (!find_any_solution(a, b, c, x,
y, g))
        return 0;
    a /= g;
    b /= g;

    ll sign_a = a > 0 ? +1 : -1;
    ll sign_b = b > 0 ? +1 : -1;

    shift_solution(x, y, a, b, (minx -
x) / b);
    if (x < minx)
        shift_solution(x, y, a, b,
sign_b);
    if (x > maxx)
        return 0;
    ll lx1 = x;

    shift_solution(x, y, a, b, (maxx -
x) / b);
    if (x > maxx)
        shift_solution(x, y, a, b,
-sign_b);
    ll rx1 = x;

    shift_solution(x, y, a, b, -(miny -
y) / a);
    if (y < miny)
        shift_solution(x, y, a, b,
-sign_a);
    if (y > maxy)
        return 0;
    ll ly2 = y;

    shift_solution(x, y, a, b, -(maxy -
y) / a);
    if (y > maxy)
        shift_solution(x, y, a, b,
sign_a);
    ll rx2 = y;

    if (ly2 > rx2)
        swap(ly2, rx2);
    ll lx = max(lx1, lx2);
}

```

```

    ll rx = min(rx1, rx2);

    if (lx > rx)
        return 0;
    return (rx - lx) / abs(b) + 1;
}

```

1. Euler Totient Function

$$\begin{aligned}
 \varphi(n) &= \varphi(p_1^{k_1}) \varphi(p_2^{k_2}) \cdots \varphi(p_r^{k_r}) \\
 &= p_1^{k_1-1} (p_1 - 1) p_2^{k_2-1} (p_2 - 1) \cdots p_r^{k_r-1} (p_r - 1) \\
 &= p_1^{k_1} \left(1 - \frac{1}{p_1}\right) p_2^{k_2} \left(1 - \frac{1}{p_2}\right) \cdots p_r^{k_r} \left(1 - \frac{1}{p_r}\right) \\
 &= p_1^{k_1} p_2^{k_2} \cdots p_r^{k_r} \left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_r}\right) \\
 &= n \left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_r}\right).
 \end{aligned}$$

m. Euler totient function using sieve

```

#define MAX 10000000

unsigned long long phi[MAX + 7];

/**
 * It took 0.902 secs to generate up to
 * 1e7.
 */
void generatePhi() {
    phi[1] = 0;
    for (int i = 2; i <= MAX; i++) {
        if(!phi[i]) {
            phi[i] = i-1;
            for(int j = (i << 1); j <=
MAX; j += i) {
                if(!phi[j]) phi[j] = j;
                phi[j] = phi[j] * (i-1) /
i;
            }
        }
    }
}

```

n. Digits of factorial

Base Conversion, $\log_b x = \frac{\log_c x}{\log_c b}$

```

ll digitfac(ll n, ll base){
    double x = 0;
    for (ll i = 1; i <= n; i++){
        x += log10(i) / log10(base); //
Base Conversion

```

```

    }
    ll res = x + 1 + eps; // eps means a
very small number such as 10^(-9)
    return res;
}

```

2. Math

a. BigMod

```

ll bigmod(ll a, ll b)

{
    if(b == 0) return 1;
    ll x = bigmod(a, b/2);
    x = (x*x)%mod;
    if(b%2 == 1) x = (x*a)%mod;
    return x;
}

```

b. Trigonometry functions

```

double PI = acos(-1); //
3.141592653589793238
double angle = 60.0; // angle value in
degree unit and double datatype
cos_val = cos ( angle * PI / 180.0 );
sin_val = sin ( angle * PI / 180.0 );
tan_val = tan ( angle * PI / 180.0 );

cos_inverse_val = acos ( val ) * 180.0 /
PI; //val in double, cos_inverse_val in
degree unit, double data type
sin_inverse_val = asin ( val ) * 180.0 /
PI;
tan_inverse_val = atan ( val ) * 180.0 /
PI;
tan_inverse_y_by_x_val = atan2 (y,x) *
180 / PI;

```

c. Matrix Exponent

```

#define ll long long
#define MAX 100000
#define mat_int long long
using namespace std;

int pw[] = {1, 10, 100, 1000, 10000};

struct matrix{
    int dimR, dimC, MOD = 10000;
    vector <vector <mat_int> > mat;
    matrix(int _dimR, int _dimC){
        dimR = _dimR;
        dimC = _dimC;

```

```

        mat.clear();
        mat.resize(dimR, vector
<mat_int> (dimC, 0));
        if(dimR == dimC){
            for(int i = 0; i < dimR;
i++)
                mat[i][i] = 1;
        }
    }

    matrix operator * (const matrix
&oth)
    {
        int nr = dimR;
        int nc = oth.dimC;
        matrix newMat(nr, nc);
        for(int i = 0; i < nr; i++){
            for(int j = 0; j < nc; j++){
                mat_int sum = 0;
                for(int k = 0; k < dimC;
k++)
                {
                    sum += (mat[i][k] *
oth.mat[k][j]) % MOD;
                    sum %= MOD;
                }
                newMat.mat[i][j] = sum;
            }
        }
        return newMat;
    }

    matrix operator ^ (ll p)
    {
        matrix res(dimR, dimC);
        matrix x = *this;

        while(p){
            if(p % 2 == 1) res = x *
res;

            x = x * x;
            p /= 2;
        }

        return res;
    }

    void print()
    {
        cout << "....." << endl;

```

```

        for(int i = 0; i < dimR; i++){
            for(int j = 0; j < dimC;
j++){
                cout << mat[i][j] << "
";
            }
            cout << endl;
        }
        cout << "....." << endl;
    }
};

```

d. Inclusion-Exclusion

```

ll inclusionExclusion(ll n)
{
    ll ans = 0;
    for(ll i = 1; i < (1LL <<
prime.size()); i++){
        if(__builtin_popcountll(i)%2){
            ll val = 1;
            for(ll pos = 0; pos <
prime.size(); pos++){
                if((i&(1LL << pos))) val
*= prime[pos];
            }
            ans += n/val;
        }
        else{
            ll val = 1;
            for(ll pos = 0; pos <
prime.size(); pos++){
                if((i&(1LL << pos))) val
*= prime[pos];
            }
            ans -= n/val;
        }
    }
    return n - ans;
}

```

3. Counting

a. Derangement

```

ll Derange(ll n)
{
    if(n==0) return 1;
    if(n==1) return 0;
    if(n==2) return 1;

    if(derange[n]!=-1) return
derange[n];

```



```

    return
derange[n]=((n-1)%mod*((Derange(n-1))%mod+
d+(Derange(n-2)%mod))%mod)%mod;
}

```

b. Catalan Number DP Style

```

#define ll long long          ///
MACROS NEEDED
#define LIM 1005             ///
MACROS NEEDED
#define mod 1000000007       ///
MACROS NEEDED

ll catalan[LIM];
void cata()
{
    catalan[0] = 1, catalan[1] = 1;  ///
Two Base Cases
    for(ll i = 2; i <= LIM; i++){
        ll sum = 0;
        for(ll j = 0; j < i; j++){
            sum +=
(catalan[j]*catalan[i-j-1])%mod;  ///
As, catalan[i] =
sum_of(catalan[j]*catalan[i-j-1])
        sum %= mod;  /// Answer with
mod 1e9+7
    }
    catalan[i] = sum;
}
}

```

c. Catalan Numbers Inverse mod

```

#define ll long long  /// MACRO
NEEDED
#define LIM 300005    /// MACRO
NEEDED
#define mod 1000000007 /// MACRO
NEEDED

ll fac[LIM];
ll bigmod(ll a, ll b)
{
    if(b == 0) return 1;
    ll x = bigmod(a, b/2);
    x = (x*x)%mod;
    if(b%2 == 1) x = (x*a)%mod;
    return x;
}
void find_fac()

```

```

{
    ll i;
    for(i = 1, fac[0] = 1; i <
LIM; i++)
        fac[i] = (i*fac[i-1])%mod;
/// factorial generation upto LIM
}
ll ncr(ll n, ll r)
{
    ll up, down, ans;
    up = fac[n];
    down = (fac[r]*fac[n-r])%mod;
    ans = (up*bigmod(down,
mod-2))%mod;  /// modular
multiplicative inverse
    return (ans+mod)%mod;
}
ll cat(ll n)
{
    ll down;
    down = bigmod(n+1,mod-2);
/// modular multiplicative inverse
    return
(ncr(2*n,n)*down+mod)%mod;  ///
catalan[i] = (1/(i+1))*(2*i)Ci
}

/// fac();          NEED TO CALL THE
FUNCTION FROM MAIN
/// val = cal()    NEED TO CALL THE
FUNCTION FROM MAIN

```

d. Occupancy Problems picture

Occupancy Problems	
OI = objects indistinguishable OD = " distinguishable CI = containers indistinguishable CD = " dist	\emptyset = empty containers permitted $\sim \emptyset$ = " " not " K kids in class K-1 of these $\emptyset \mid \emptyset \mid \emptyset \mid \dots \mid$
OI, CI, $\sim \emptyset$ $n > k$ n Snickers Bars, K bags $p(n, K)$	OI, CD, $\sim \emptyset$ n Snickers bars, K kids $n-k$ bars remaining $\binom{k-1+n-k}{k-1} = \binom{n-1}{k-1}$
OI, CI, \emptyset $p(n,1) + p(n,2) + p(n,3) + \dots + p(n,k)$ $p(n+k, K)$	OI, CD, \emptyset n Snickers bars, K kids $\binom{k-1+n}{k-1}$
OD, CI, $\sim \emptyset$ n totally diff candy bars $S(n, K)$	OD, CD, $\sim \emptyset$ n totally diff bars, K kids $K! \cdot S(n, K)$
OD, CI, \emptyset $S(n,1) + S(n,2) + S(n,3) + \dots + S(n,k)$	OI, CD, \emptyset n totally diff candy bar, K kids K^n

e. Pascal Triangle (NCR DP style)

```

ll ncr[1005][1005];
void precal()
{
    ll i, j;
    for(i = 0; i < 1005; i++)
        ncr[i][0] = 1;
    for(i = 1; i < 1005; i++)
        for(j = 1; j < 1005;
            j++)
            ncr[i][j] =
            (ncr[i-1][j] + ncr[i-1][j-1])%mod;
}

```

f. NCR Inverse mod

```

ll fac[LIM];
ll bigmod(ll a, ll b)
{
    if(b == 0) return 1;
    ll x = bigmod(a, b/2);
    x = (x*x)%mod;
    if(b%2 == 1) x = (x*a)%mod;
    return x;
}
void find_fac()
{
    ll i;

```

```

for(i = 1, fac[0] = 1; i < LIM; i++)
    fac[i] = (i*fac[i-1])%mod;

```

```

}
ll ncr(ll n, ll r)
{
    ll up, down, ans;
    up = fac[n];
    down = (fac[r]*fac[n-r])%mod;
    ans = (up*bigmod(down, mod-2))%mod;
    return (ans+mod)%mod;
}

```

// add find_fac() in code

// add ncr() in code

4. Data Structure

a. Segment tree

```

#include <bits/stdc++.h>
using namespace std;
#define MAX 100005
#define left st, (st + en) / 2, nd
+ nd
#define right (st + en) / 2 + 1,
en, nd + nd + 1
int tree[4 * MAX + 5];
int n, arr[MAX + 5];

void build(int st, int en, int nd)
{
    if(st == en)
    {
        tree[nd] = arr[st];
        return;
    }
    build(left); /// left subtree
    build(right); /// right
    subtree
    tree[nd] = tree[nd + nd] +
    tree[nd + nd + 1];
}

int query(int st, int en, int nd,
int L, int R)
{
    if(L <= st && en <= R) return
    tree[nd]; /// if the query segment
    is completely overlapping our tree
    segment/node.
    if(en < L || R < st) return 0;
    return query(left, L, R) +
    query(right, L, R);
}

```

```

}

void update(int st, int en, int
nd, int idx, int v)
{
    if(en < idx || idx < st)
return;
    if(st == en)
    {
        tree[nd] += v;
        return;
    }
    update(left, idx, v);
    update(right, idx, v);
    tree[nd] = tree[nd + nd] +
tree[nd + nd + 1];
}

int main()
{
    /// build(0, n - 1, 1);
    cin >> n;
    for(int i = 0; i < n; i++) cin
>> arr[i];
    build(0, n - 1, 1); /// O(4 *
N)
    int Q;
    cin >> Q;
    while(Q--)
    {
        int com;
        cin >> com;
        if(com == 0) /// update
        {
            int idx, v;
            cin >> idx >> v;
            update(0, n - 1, 1,
idx, v);
        }
        else if(com == 1) ///
query
        {
            int L, R;
            cin >> L >> R;
            cout << query(0, n -
1, 1, L, R) << "\n";
        }
    }
    return 0;
}

```

b. Segment tree with lazy

```

#include <bits/stdc++.h>
#define ll long long
#define LIM 100005
#define left st, (st + ed) / 2, nd
+ nd
#define right (st + ed) / 2 + 1,
ed, nd + nd + 1
using namespace std;

ll arr[LIM], tree[4*LIM],
lazy[4*LIM];

void build(ll st, ll ed, ll nd)
{
    if(st == ed)
    {
        tree[nd] = arr[st];
        return;
    }
    build(left); /// left subtree
    build(right); /// right
subtree

    tree[nd] = tree[nd + nd] +
tree[nd + nd + 1];
}

void update(ll st, ll ed, ll nd,
ll L, ll R, ll val)
{
    if(lazy[nd] != 0){ /// IF ANY
UPDATE IS PENDING
        tree[nd] += (ed - st +
1)*lazy[nd]; /// UPDATE THE TREE
        if(st != ed){
            lazy[nd + nd] +=
lazy[nd]; /// PROPAGATE PENDING
UPDATES TO THE LEFT CHILD
            lazy[nd + nd + 1] +=
lazy[nd]; /// PROPAGATE PENDING
UPDATES TO THE RIGHT CHILD
        }
        lazy[nd] = 0; /// ALL
PENDING ARE CLEAR
    }

    if(ed < L or R < st) return;
    /// INVALID STATE

    if(L <= st and ed <= R){

```

```

        tree[nd] += (ed - st +
1)*val; /// UPDATE THE RANGE WITH
THE INSERTED VALUE 'val'
        if(st != ed){
            lazy[nd + nd] += val;
/// PROPAGATE PENDING UPDATES TO
THE LEFT CHILD
            lazy[nd + nd + 1] +=
val;/// PROPAGATE PENDING UPDATES
TO THE RIGHT CHILD
        }
        return;
    }
    update(left, L, R, val);
/// UPDATE LEFT CHILD
    update(right, L, R, val);
/// UPDATE RIGHT CHILD

    tree[nd] = tree[nd + nd] +
tree[nd + nd + 1];    /// KEEP THE
ANSWERS OF CHILDS IN PARANT NODE
}

ll query(ll st, ll ed, ll nd, ll
L, ll R)
{
    if(ed < L or R < st) return 0;
/// INVALID STATE

    if(lazy[nd] != 0){    /// IF ANY
UPDATE IS PENDING
        tree[nd] += (ed - st +
1)*lazy[nd]; /// UPDATE THE TREE
        if(st != ed){
            lazy[nd + nd] +=
lazy[nd];    /// PROPAGATE PENDING
UPDATES TO THE LEFT CHILD
            lazy[nd + nd + 1] +=
lazy[nd];    /// PROPAGATE PENDING
UPDATES TO THE RIGHT CHILD
        }
        lazy[nd] = 0;    /// ALL
PENDING ARE CLEAR
    }

    if(L <= st and ed <= R){
        return tree[nd];    ///
OVERLAPPED STATE
    }

```

```

        return query(left, L, R) +
query(right, L, R);    /// KEEP THE
ANSWERS OF CHILDS IN PARANT NODE
    }

c. Sparse table
#include <bits/stdc++.h>
#define ll long long
#define LIM 100005
using namespace std;

int arr[LIM], dp[LIM][22],
log_array[LIM];    /// NEVER USE
LOG/SQRT ETC. IN SIZE DECLARATION

///
dp[array_size][log2(array_size)]

/// dp[starting_index][range]

void construct(int n)    /// n =
ARRAY SIZE; O(nlogn)
{
    log_array[1] = 0;
    for(int i = 2; i <= n; i++){
        log_array[i] = 1 +
log_array[i / 2];    ///
GENERATING 2 BASED LOG VALUES 1 TO
n

        for(int i = 0; i < n; i++){
            dp[i][0] = arr[i];
/// SINGLE RANGE RESULTS or BASE
CASE

            int k = log_array[n] + 1;

            for(int j = 1; j <= k; j++){
                for(int i = 0; i + (1 <<
(j - 1)) < n; i++){
                    dp[i][j] = min(dp[i][j
- 1], dp[i + (1 << (j - 1))][j -
1]);
                }
            }
            // dp[0...n-1][1]
            // dp[0...n-1][2]
            // ...
            // dp[0...n-1][K]
        }
    }
}

```

```

int query(int L, int R) /// O(1)
{
    int len = R - L + 1;
    //if(len <= 0) return INT_MAX;
    int lg = log_array[len];
    return min(dp[L][lg], dp[R -
(1 << lg) + 1][lg]); /// RETURN
ANSWER FROM 2 OVERLAPPING RANGES
}

```

d. LCA (Lowest Common Ancestor)

```

#include <bits/stdc++.h>
#define ll long long
#define pb push_back
using namespace std;

int parent[1005][22], depth[1005];
vector <int> edg[1005];

void dfs(int u, int pr) /// O(n)
{
    if(pr != -1)
        depth[u] = depth[pr] + 1;
    parent[u][0] = pr;

    for(int i = 0; i <
edg[u].size(); i++){
        int v = edg[u][i];
        if(v != pr) dfs(v, u);
    }
}

void build(int n) /// O(nlogn)
{
    int lg = log2(n) + 1;

    for(int k = 1; k <= lg; k++){
        for(int i = 1; i <= n;
i++){
            if(parent[i][k-1] !=
-1)
                parent[i][k] =
parent[parent[i][k-1]][k-1];
        }
    }
}

```

```

int findLCA(int n, int u, int v)
/// O(nlogn)
{
    if(depth[u] > depth[v])
swap(u, v);

    int diff = depth[v] -
depth[u];

    int lg = log2(n) + 1;
    for(int i = 0; i <= lg; i++){
        if(diff & (1 << i))
            v = parent[v][i];
    }

    if(u == v) return v;

    for(int i = lg; i >= 0; i--){
        if(parent[u][i] !=
parent[v][i]){
            u = parent[u][i];
            v = parent[v][i];
        }
    }

    return parent[u][0];
}

void init(int n)
{
    int lg = log2(n) + 1;
    for(int i = 0; i <= n; i++){
        edg[i].clear();
        depth[i] = 0;
        for(int j = 0; j <= lg;
j++)
            parent[i][j] = -1;
    }
}

```

e. DSU (Disjoint Set Union)

```

int findP(int u)
{
    if(u == pr[u]) return u;
    return pr[u] = findP(pr[u]);
}

void connect(int u, int v)
{
    u = findP(u);
    v = findP(v);
}

```

```

/// ARTICULATION BRIDGES
#define ll long long    /// MACROS
NEEDED
#define pb push_back    /// MACROS
NEEDED
#define LIM 100005      /// MACROS
NEEDED
ll dis[LIM], vis[LIM], par[LIM],
low[LIM], cnt, n;
vector<ll> ver[LIM];
vector<pair<ll,ll>> bridges;
void dfs(ll u)
{
    dis[u] = low[u] = cnt++;
    for(ll i = 0; i <
ver[u].size(); i++){
        ll v = ver[u][i];
        if(v != par[u]){
            if(vis[v] == 0){ ///
TREE EDGES
                par[v] = u;    ///
SAVING PARTENTS
                vis[v] = 1;
                dfs(v);
                low[u] =
min(low[u], low[v]); /// IF TREE
EDGE, LOW[U] = MIN(LOW[U] OF ALL
OF ITS CHILD)
            }
            else if(vis[v] == 1){
/// BACK EDGES
                low[u] =
min(low[u], dis[v]); /// IF BACK
EDGE, LOW[U] = MIN(MIN OF
DIS[CHILD] AND LOW[U])
            }
            if(dis[u] < low[v])
bridges.pb({min(u,v), max(u,v)});
/// ARTICULATION EDGES

```

```

    }
}
vis[u] = 2;    /// FORWARD
EDGES
}

/// ARTICULATION POINTS
int n, m, tim, dis[LIM], low[LIM],
par[LIM];
vector <int> edg[LIM], artpoint;

void dfs(int u)
{
    dis[u] = low[u] = tim++;
    int child = 0, mx = -inf;
    for(int i = 0; i <
edg[u].size(); i++){
        int v = edg[u][i];
        if(v != par[u]){
            if(dis[v]){    ///
BACK EDGE
                low[u] =
min(low[u], dis[v]); /// IF BACK
EDGE, LOW[U] = MIN(MIN OF
DIS[CHILD] AND LOW[U])
            }
            else{    ///
TREE EDGE
                par[v] = u; ///
SAVING PARENT
                dfs(v);
                low[u] =
min(low[u], low[v]);    /// IF TREE
EDGE, LOW[U] = MIN(LOW[U] OF ALL
OF ITS CHILD)
                if(low[v] >=
dis[u] && u != 1)
                    mx = max(mx,
dis[v]);
                child++;
            }
        }
    }
    if(u != 1 && mx >= dis[u])
artpoint.pb(u);    /// IF THE NODE
IS NOT ROOT
    if(u == 1 && child > 1)
artpoint.pb(u);    /// IF THE

```

```

NODE IS ROOT CHECK IF (# OF CHILD
> 1)
}

```

d. MST (Min/Max Spanning Tree)

```

#define ll long long    /// MACRO
NEEDED
#define pb push_back    /// MACRO
NEEDED

```

```

int n, pr[105], sz[105];
struct abc{
    int w, u, v;
};
vector <abc> edg;

bool cmp(abc x, abc y)
{
    return x.w < y.w;
}

int findP(int u)
{
    if(u == pr[u]) return u;
    return pr[u] = findP(pr[u]);
}

void connect(int u, int v)
{
    u = findP(u);
    v = findP(v);
    if(u != v){
        if(sz[u] < sz[v]){    ///
MARGING SMALLER COMPONENT WITH
LARGER COMPONENT
            pr[u] = v;
            sz[v] += sz[u];
        }
        else{
            pr[v] = u;
            sz[u] += sz[v];
        }
    }
}

```

```

int kruskal(int state)    /// STATE
= 0 (MIN), STATE = 1(MAX)
{
    int i;
    for(i = 0; i <= n; i++){
        pr[i] = i;
    }
}

```

```

        sz[i] = 1;
    }
    sort(edg.begin(), edg.end(),
cmp);
    if(state)
        reverse(edg.begin(),
edg.end());
    int cost = 0;
    for(i = 0; i < edg.size();
i++){
        int parx =
findP(edg[i].u);
        int pary =
findP(edg[i].v);
        if(parx != pary){
            cost += edg[i].w;
            connect(parx, pary);
        }
    }
    return cost;
}

```

e. King, Horse & Adjacent Walk

```
#define ll long long
```

```

ll king_xx[] = {1, -1, 0, 0, 1, 1, -1,
-1};
ll king_yy[] = {0, 0, 1, -1, 1, -1, 1,
-1};

ll horse_xx[] = {1, 2, 2, 1, -1, -2, -2,
-1};
ll horse_yy[] = {2, 1, -1, -2, -2, -1,
1, 2};

ll adj_xx[] = {1, -1, 0, 0};
ll adj_yy[] = {0, 0, 1, -1};

```

6. Strings

a. Trie

```

string s;
int n, trie[LIM][15], cnt = 2;
bool flag[LIM], endpoint[LIM];

void add()
{
    int u = 1;
    for(int i = 0; i < s.size(); i++){
        int edg = s[i] - '0';
        if(!trie[u][edg]){
            trie[u][edg] = cnt++;

```

```

        }
        u = trie[u][edg];
        flag[u] = 1;
    }
    endpoint[u] = 1;
}

bool query()
{
    int u = 1;
    for(int i = 0; i < s.size(); i++){
        int edg = s[i] - '0';
        if(trie[u][edg])
            u = trie[u][edg];
        else
            return 1;
        if(endpoint[u]) return 0;
    }

    return (flag[u] == 0);
}

```

```
void reset(int u)
```

```

{
    for(int i = 0; i < 10; i++){
        if(trie[u][i])
            reset(trie[u][i]);
        trie[u][i] = flag[u] =
endpoint[u] = 0;
    }
}

```

b. KMP

```

#include <bits/stdc++.h>
#define ll long long
#define LIM 1000006
using namespace std;

char temp[LIM];
string text, pat;
int prefix[LIM];

void calcPrefix()
{
    prefix[0] = prefix[1] = 0;

    for(int i = 2; i < pat.size(); i++){
        int p = i - 1;
        while(1){
            if(pat[i] == pat[ prefix[p]
+ 1 ]){

```



```

        prefix[i] = prefix[p] +
1;
        break;
    }
    else if(!p) break;
    else p = prefix[p];
}
}
}

```

```

int kmp() /// returns the # of matched
substrings
{

```

```

    int p = 0, ans = 0;
    for(int j = 0; j < text.size();
j++){
        while(p && pat[p + 1] !=
text[j]){
            p = prefix[p];
        }
        if(pat[p + 1] == text[j]){
            p++;
        }
        if(p + 1 == pat.size()){
            ans++;
        }
    }
    return ans;
}

```

```

int main()
{
    int cas = 1, t;
    scanf("%d", &t);
    while(t--){
        scanf("%s", temp); text = temp;
        scanf("%s", temp); pat = temp;
        pat = "#" + pat;
        calcPrefix();
        printf("Case %d: %d\n", cas++,
kmp());
    }

    return 0;
}

```

c. Hashing

```

#define LIM 1000006

```

```

ll prefHash[2][LIM], basePower[2][LIM];
ll mod[2] = {1000000007, 1000000009};

```

```

ll base[2] = {31, 101};
string str;
char temp[LIM];

void preCal()
{
    prefHash[0][0] = prefHash[1][0] =
str[0] - 'a' + 1;
    basePower[0][0] = basePower[1][0] =
1;

```

```

    for(int j = 0; j < 2; j++){
        for(int i = 1; i < str.size();
i++){
            basePower[j][i] =
(basePower[j][i - 1]*base[j]) %mod[j];
            prefHash[j][i] =
(prefHash[j][i - 1]*base[j] + str[i] -
'a' + 1) %mod[j];
        }
    }
}

```

```

ll evaluateHash(int idx, int L, int R)
{
    if(L == 0) return prefHash[idx][R];
    return (prefHash[idx][R] -
(prefHash[idx][L - 1]*basePower[idx][R -
L + 1]) %mod[idx] + mod[idx]) %mod[idx];
}

```

d. Sub-string palindrome check using Hashing

```

#define ll long long
#define LIM 1000006

```

```

ll prefHash[2][LIM],
prefHashRev[2][LIM], basePower[2][LIM];
ll mod[2] = {1000000007, 1000000009};
ll base[2] = {31, 101};
string str, revStr;
char temp[LIM];

```

```

void preCal()
{
    prefHash[0][0] = prefHash[1][0] =
str[0] - 'a' + 1;
    prefHashRev[0][0] =
prefHashRev[1][0] = revStr[0] - 'a' + 1;
    basePower[0][0] = basePower[1][0] =
1;

```

```

    for(int j = 0; j < 2; j++){
        for(int i = 1; i < str.size();
i++){
            basePower[j][i] =
(basePower[j][i - 1]*base[j]) %mod[j];
            prefHash[j][i] =
(prefHash[j][i - 1]*base[j] + str[i] -
'a' + 1) %mod[j];
            prefHashRev[j][i] =
(prefHashRev[j][i - 1]*base[j] +
revStr[i] - 'a' + 1) %mod[j];
        }
    }

ll getHash(int idx, int L, int R)
{
    if(L == 0) return prefHash[idx][R];
    return (prefHash[idx][R] -
(prefHash[idx][L - 1]*basePower[idx][R -
L + 1]) %mod[idx] + mod[idx]) %mod[idx];
}

ll getHashRev(int idx, int L, int R)
{
    if(L == 0) return
prefHashRev[idx][R];
    return (prefHashRev[idx][R] -
(prefHashRev[idx][L -
1]*basePower[idx][R - L + 1]) %mod[idx]
+ mod[idx]) %mod[idx];
}

bool isPalindrome(int L, int R)
{
    int sz = str.size();
    if(getHash(0, L, R) == getHashRev(0,
sz - R - 1, sz - L - 1) &&
        getHash(1, L, R) == getHashRev(1,
sz - R - 1, sz - L - 1)) return 1;
    return 0;
}

/// preCal() -> CALL FROM MAIN
/// isPalindrome(start, end) -> CALL
FROM MAIN

```

7. Template

```

//start_template
#include <bits/stdc++.h>

```

```

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>

#define ll long long
#define ull unsigned long long
#define pb push_back
#define pii pair<int, int>
#define ff first
#define ss second
#define nl '\n'
#define mod 1000000007
#define inf 1000000009
#define MAXX 10000000000000015
#define LIM 300005
#define eps 1e-9
#define pi acos(-1)

using namespace std;
using namespace __gnu_pbds;
/*
Activate it for pbds set
x.find_by_order(pos) ----> returns the
value at "pos" index in the set 0-based
x.order_of_key(value) ----> returns
the position of "value" in the set
0-based
*/
#define ordered_set tree<pii,
null_type,less<pii>, rb_tree_tag,
tree_order_statistics_node_update>
void FAST_IO() {
ios_base::sync_with_stdio(false);
cin.tie(0); cout.tie(0); }
/*
Partition Numbers:
Formula:  $P(n,k) = P(n-1,k-1) + P(n-k,k)$ 
Base Case:
 $P(n,k) = 0$  if  $n < k$ 
 $P(n,k) = 0$  if  $k = 0$  [ $P(n,0) = 0$ ]
 $P(n,k) = 1$  if  $(k = n)$  [ $P(n,n) = 1$ ]

Catalan Numbers:
Formula:  $C(n) = \text{SumOf}(C(k) * C(n-1-k))$ ,  $n \geq 2$ ,  $0 \leq k < n$ 
Base Case:  $C(0) = C(1) = 1$ 
Formula2:  $C(n) = (1/(n + 1)) * \text{ncr}(2*n, n)$ 

Starling Numbers of 2nd kind:
Formula:  $S(n,k) = S(n-1,k-1) + k * S(n-1,k)$ 

```

Base Case:

```
S(n,k) = 0 if k > n
S(n,k) = 1 if k = 1 [S(n,1) = 1]
S(n,k) = 1 if n = k [S(n,n) = 1]
S(n,k) = 0 if k = 0 [S(n,0) = 0]
```

Moves on Grid:

```
int king_xx[] = {1, -1, 0, 0, 1, 1, -1,
-1};
int king_yy[] = {0, 0, 1, -1, 1, -1, 1,
-1};
```

```
int horse_xx[] = {1, 2, 2, 1, -1, -2,
-2, -1};
int horse_yy[] = {2, 1, -1, -2, -2, -1,
1, 2};
```

```
int adj_xx[] = {1, -1, 0, 0};
int adj_yy[] = {0, 0, 1, -1};
```

```
/// bit manipulations
/// bool checkbit(int mask,int
bit){return mask & (1<<bit);}
/// int setbit(int mask,int bit){ return
mask (1<<bit) ; }
/// int clearbit(int mask,int
bit){return mask & ~(1<<bit);}
/// int togglebit(int mask,int
bit){return mask ^ (1<<bit);}
*/
```

```
int main()
{
    ///MUST READ THE POINTS BELOW BEFORE
    SUBMIT
    FAST_IO();

    return 0;
    ///MUST READ THE POINTS BELOW BEFORE
    SUBMIT
}
/*
1. LOOK SPECIAL CASE N = 1.
2. LOOK FOR OVERFLOW.
3. LOOK FOR OUT OF BOUNDS.
4. ALWAYS TEST WITH HAND-MADE
TEST-CASES BEFORE SUBMIT.
*///end_template
```

8. Mo's ordering

```
struct query {
    ll id, l, r;
    bool operator < (const query &
other) const {
        int block_a = l / block_size;
        int block_b = other.l /
block_size;
        if(block_a == block_b) return r
< other.r;
        return block_a < block_b;
    }
};
ll arr[mx + 5], left, right, res[mx +
5], freq[mx + 5];
query queries[50005];
void add(ll idx) {
    freq[arr[idx]]++;
    if(freq[arr[idx]] == 1)
unique_values++;
}
void remov(ll idx) {
    freq[arr[idx]]--;
    if(freq[arr[idx]] == 0)
unique_values--;
}
//in main function
for(ll i = 0; i < q; i++) {
    while(queries[i].l < left)
add(--left);
    while(queries[i].l > left)
remov(left++);
    while(queries[i].r < right)
remov(right--);
    while(queries[i].r > right)
add(++right);

    res[queries[i].id] =
unique_values;
}
```

9. Maximum Subarray Sum (Kadane's Algo)

//General version

```
vector<int> a(n);
int ans = a[0], sum = 0;
for (int r = 0; r < n; ++r) {
    sum += a[r];
    ans = max(ans, sum);
    sum = max(sum, 0);
}
cout << sum << nl;
```

```
//With index of the subarray of
max-sum(left-right)
vector<int> a(n);
int ans = a[0], ans_l = 0, ans_r = 0;
int sum = 0, minus_pos = -1;
for (int r = 0; r < n; ++r) {
    sum += a[r];
    if (sum > ans) {
        ans = sum;
        ans_l = minus_pos + 1;
        ans_r = r;
    }
    if (sum < 0) {
        sum = 0;
        minus_pos = r;
    }
}
}
```

10. Dynamic Programming

a. How many zeroes

```
ll dp[15][5][5][5], dp2[15][5][5];
ll f2(ll pos, ll issmall, ll num)
{
    /// HOW MANY NUMBERS SMALLER THAN n
    IS POSSIBLE FROM THIS STATE
    (pos,issmall)
    if(pos == s[num].size()) return 1;
    if(dp2[pos][issmall][num] != -1)
return dp2[pos][issmall][num];
    ll lo = 0, hi = s[num][pos] - '0', sum
= 0;
    if(issmall) hi = 9;
    for(; lo <= hi; lo++){
        sum += f2(pos+1,issmall |
(lo<hi), num);
    }
    return dp2[pos][issmall][num] = sum;
}
ll f(ll pos, ll issmall, ll hasstarted,
ll num)
{
    if(pos == s[num].size()) return 0;
    if(dp[pos][issmall][hasstarted][num]
!= -1) return
dp[pos][issmall][hasstarted][num];
    ll lo = 0, hi = s[num][pos] - '0', sum
= 0;
    if(issmall) hi = 9;
    for(; lo <= hi; lo++){
        ll tem = f(pos+1, issmall |
(lo<hi), hasstarted | (lo != 0), num);
```

```
/// ALREADY STARTED AND
RIGHT-NOW WE ARE PUTTING 0 AT CURRENT
POSITION
/// SO WE WILL HAVE TO FIND OUT
IN HOW MANY WAYS THIS ZERO WILL
CONTRIBUTE
    if(hasstarted and lo == 0) tem
+= f2(pos+1,issmall | (lo<hi),num);
    sum += tem;
}
return
dp[pos][issmall][hasstarted][num] = sum;
}
```

11. Searching

a. Ternary Search

```
double ternary_search(double l, double
r) {
    double eps = 1e-9;
    //set the error limit here
    while (r - l > eps) {
        double m1 = l + (r - l) / 3;
        double m2 = r - (r - l) / 3;
        double f1 = f(m1);
        //evaluates the function at m1
        double f2 = f(m2);
        //evaluates the function at m2
        if (f1 < f2)
            l = m1;
        else
            r = m2;
    }
    return f(l);
    //return the maximum of f(x) in [l, r]
}
```