#### 1 ARM Introduction

In this assignment, you'll be writing some short programs directly in machine code. You won't quite write them directly as 1's and 0's—you'll actually type in hexadecimal characters—but you'll directly use the encoding of ARM instructions in binary. This section lays out the instructions you'll need to use for this assignment in detail.

For now, it will suffice to know a few basic formats for data manipulation instructions. Here is the format for manipulating data strictly between registers:

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	0	0	0	O.	рс	od	le	0		R	$l_n$			F	$l_d$		0	0	0	0	0	0	0	0		R	m	

Generally,  $R_n$  and  $R_m$  are the operands of this operation; opcode selects which operator to apply to them, and the result is stored in  $R_d$ . Some mathrelated opcodes useful for this assignment are:

opcode	mnemonic	expression
0000	AND	$R_d := R_n \& R_m$
0001	EOR	$R_d := R_n \hat{R}_m$
0010	SUB	$R_d \coloneqq R_n - R_m$
0011	RSB	$R_d \coloneqq R_m - R_n$
0100	ADD	$R_d \coloneqq R_n + R_m$
1100	ORR	$R_d \coloneqq R_n \mid R_m$
1110	BIC	$R_d \coloneqq R_m \& \sim R_m$

Another form of data manipulation instructions lets us combine values in registers with constants:

31	30	29	28	27	26	25	24	23 2	2 2 1	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	0	0	1	oı	oco	de	0		R	n			R	d		0	0	0	0				In	m	L		

The key difference is the 1 instead of a 0 in bit 25, which makes the instruction operate in immediate mode. Then, Imm is an binary value that is used in place of  $R_m$  above. So, for example, if we used 00001111 as bits 0-7 of of an addition instruction, it would add 15 to the value of the  $R_n$  register and store it in  $R_d$ .

There is another format of instructions which is used for multiply-like instructions:

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	0	0	0	0		op	)	0		R	d		0	0	0	0		R	s		1	0	0	1		R	m	

The multiply operation's opcode is 000, so replacing op above with 000 indicates a multiply. You only need this version of multiply for this assignment,

you can see the others in tabular form in appendix B of the Harris book if you're interested in seeing more.

# 2 Support Code and Example

We've provided support code so that you can simply write machine code instructions as hexadecimal, create a binary file from them, and then execute it with certain starting conditions. Let's consider the following problem:

Assume that r1 contains x and r2 contains y. Write machine code instructions that result in the value x + y being stored in r0.

Our task is to create a binary file containing the (in this case single) instruction that adds the values and stores the result in r0. We can use the specification of add above to get started; we want to take the pattern for data processing instructions, and pick the opcode for ADD (which is 0 1 0 0) and the appropriate registers. Note that the registers are indicated by their unsigned binary representation in 4 bits, so r0 is 0 0 0, r1 is 0 0 0 1, and so on. Here's the first step

31	30	29	28	27	26	25	24	23	22	$^{21}$	20	19	18	17	16	15	14	13	$^{12}$	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	0	0	0	0]	рс	oc	le	0		R	n			R	d		0	0	0	0	0	0	0	0		R	m	
31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
1	1	1	0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1

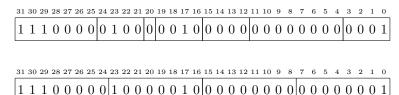
(Here, we picked that the n-labelled source register would hold  $\mathtt{r2}$  (y), and the m-labelled source register would hold  $\mathtt{r1}$  (x). Since addition commutes, we could have picked either order.)

Now we need to take that instruction and turn it into a binary file. Note that a file containing text 1's and 0's is not a binary file [REF cmu binary file link]. We don't want a file containing a sequence of bytes representing ASCII; we want a file containing exactly the four bytes of the instruction above, in the right format and order.

One of the most straightforward ways to do this is with the xxd command. It is useful for a number of things, where one mode of operation is taking files containing hexadecimal text and converting them to the corresponding binary files. So we can take a file that contains purely hexadecimal ASCII text, and produce a binary file. We'll show an example of that command in a moment, but first, we need to think a little bit about endianness and ordering.

All of the instructions we've shown have had the 0th bit all the way to the right. This is a convention for writing out ARM machine code because it reads well left-to-right – first the opcode, then the arguments. However, the actual order of the bytes is the reverse of this. Further, the bytes themselves are interpreted in little-endian order. Let's do an example to process this conversion.

First, we re-organize the breaks to see the instruction in terms of byte chunks, which will help us think through the hexadecimal encoding.



Next we reverse the whole thing to help us see the correct left-to-right order of the whole instruction:

Next, we flip each individual byte, because the bytes are interpreted in little-endian order:

Finally, we convert to hexadecimal:

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30 3	31	
			0	1							0	0							8	2							е	0				

So the hexadecimal encoding of this instruction is 01082e0. We can put that into a file, let's call it add.txt. The xxd command can be used to create a binary from it:

### \$ xxd -p -r add.txt add.bin

And then we can run the file using the provided runner, which lets us give arguments as command-line arguments, puts them into the designated registers, and runs our binary code:

```
$ ./runner add.bin 5 6
f(5, 6) = 11
```

A few extra notes:

• We can use xxd to see our instruction in binary again:

```
$ xxd -b add.bin
00000000: 00000001 00000000 10000010 11100000
```

• An abbreviated process for getting to the appropriate hexadecimal from the conventional binary notation is to convert to hexadecimal, and then simply reverse the order of the four bytes, so the first byte becomes the last, the last the first, and the middle two are swapped:

0 0 0 0	1 0 0 0 0 0 1 0	0 0 0 0 0 0 0	0 0 0 0 0 0 0 1
27 26 25 24	23 22 21 20 19 18 17 16	15 14 13 12 11 10 9 8	7 6 5 4 3 2 1 0
)	82	00	01
_	26 25 24		82 00

82

e0

00

#### 3 Problems

01

You will fill in the provided text files, each with a particular formula using the instructions above. You will hand in your plain-text files with instructions in hexadecimal, using the process described above. In all cases, assume that initially x is stored in r1, y is stored in r2, and the result should appear in r0. You may use any registers from r0 to r6 for your computation.

- 1. In sub2.txt, compute x y y
- 2. In submul.txt, compute (x-y)\*y
- 3. In addsubmul.txt, compute (x y) \* (x + y)
- 4. In quad.txt, compute  $2x^2 + yx 7$

If an operation would produce a number that is too large to represent in 32 bits, you can simply use the part that doesn't overflow (the default for most instructions). For example, the result of multiplying 2,000,000,000,000 by itself requires more than 32 bits, so if that's an input as x in  $\mathtt{quad.txt}$ , it's fine if your program evaluates it to 0 (since the lower 32 bits of the answer will all be 0).

#### 4 README

Along with your code, please also submit a README file which contains your answers to the following questions:

1. Starting from the instruction format in the textbook and this writeup, we have to do two apparent reversing operations: one on the whole instruction, and one on each byte. Please explain why we need these two reverses in your own words.

- 2. Consider the machine code corresponding to r0 = r1 + r2. If we wanted to change the instruction to use some other register instead of r1, how would the machine code change and why?
- 3. Please explain what you did for each statement in quad.txt.

Please limit your answer to each question to no more than 100 words and present them as clearly as possible.

## 5 Handin

Commit and push your four text files to the Github repository that was created for you by 11:59PM on Tuesday, October 10. You can push up to one day late for a 20% penalty. After you push, make sure to check on Github that the files are actually there; we will copy all of the repositories for grading a few minutes after midnight and grade precisely what is there.