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## 12.2 AdjacencyLists: A Graph as a Collection of Lists

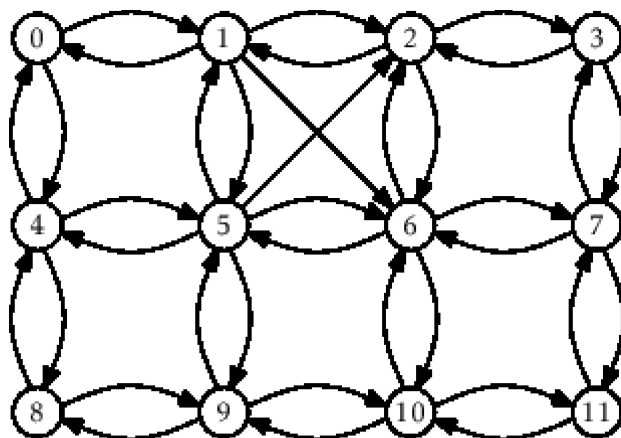
Adjacency list representations of graphs take a more vertex-centric approach. There are many possible implementations of adjacency lists. In this section, we present a simple one. At the end of the section, we discuss different possibilities. In an adjacency list representation, the graph  $G = (V, E)$  is represented as an array, `adj`, of lists. The list `adj[i]` contains a list of all the vertices adjacent to vertex `i`. That is, it contains every index `j` such that  $(i, j) \in E$ .

```

int n;
List<Integer>[] adj;
AdjacencyLists(int n0) {
    n = n0;
    adj = (List<Integer>[])new List[n];
    for (int i = 0; i < n; i++)
        adj[i] = new ArrayStack<Integer>(Integer.class);
}

```

(An example is shown in Figure 12.3.) In this particular implementation, we represent each list in `adj` as an `ArrayStack`, because we would like constant time access by position. Other options are also possible. Specifically, we could have implemented `adj` as a `DLList`.



0	1	2	3	4	5	6	7	8	9	10	11	
1	0	1	2	0	1	5	6	4	8	9	10	

4	2	3	7	5	2	2	3	9	5	6	7	
	6	6		8	6	7	11		10	11		
	5				9	10						
					4							

**Figure 12.3:** A graph and its adjacency lists

The `addEdge(i, j)` operation just appends the value `j` to the list `adj[i]`:

```
void addEdge(int i, int j) {
    adj[i].add(j);
}
```

This takes constant time.

The `removeEdge(i, j)` operation searches through the list `adj[i]` until it finds `j` and then removes it:

```
void removeEdge(int i, int j) {
    Iterator<Integer> it = adj[i].iterator();
    while (it.hasNext()) {
        if (it.next() == j) {
            it.remove();
            return;
        }
    }
}
```

This takes  $O(\text{deg}(i))$  time, where  $\text{deg}(i)$  (the degree of `i`) counts the number of edges in  $E$  that have `i` as their source.

The `hasEdge(i, j)` operation is similar; it searches through the list `adj[i]` until it finds `j` (and returns true), or reaches the end of the list (and returns false):

```
boolean hasEdge(int i, int j) {
    return adj[i].contains(j);
}
```

This also takes  $O(\text{deg}(i))$  time.

The `outEdges(i)` operation is very simple; it returns the list `adj[i]` :

```
List<Integer> outEdges(int i) {
    return adj[i];
}
```



