

# CS 241 – Week 6 Tutorial

## Regular Languages: DFAs and Regular Expressions

Spring 2015

### Summary

- DFA problems
- Regular Expressions

## 1 Regular Languages Review

An alphabet (denoted  $\Sigma$ ) is a finite set of symbols.

- $\{a, b, c\}$
- $\{b\}$
- $\{to, be, or, not\}$
- $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$

A word (over an alphabet  $\Sigma$ ) is finite sequence of symbols from  $\Sigma$ .

### Example

- bac, aba, c given that  $\Sigma = \{a, b, c\}$
- $\varepsilon, b, bb, bbb$  given that  $\Sigma = \{b\}$
- to be or not to be, not to be (one word formed from the alphabet)  
 $\Sigma = \{to, be, or, not\}$
- DEADBEEF, FACE given that  $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F\}$

A language is a set of words. A Regular Language  $R$  is a sets of words where either:

- $R$  is the empty language
- $R$  contains a single word
- $R$  is the union of two regular languages
- $R$  is the concatenation of two regular languages
- $R = L^* = \cup_{i=0}^{\infty} L^i$  where  $L$  is a regular language,  $L^0 = \{\varepsilon\}$  and for  $i > 0, L^i = L \cdot L^{i-1}$

## Deterministic Finite Automaton (DFA)

A Deterministic Finite Automaton (DFA) is a 5-tuple  $(\Sigma, Q, q_0, \mathcal{A}, \delta)$  where:

$\Sigma$  – the input alphabet

$Q$  – finite set of states

$q_0 \in Q$  – a starting state in the set of states

$\mathcal{A} \subseteq Q$  – set of accepting states

$\delta : Q \times \Sigma \rightarrow Q$  – the transition function

## Regular Expressions

Regular expressions are a means of expressing regular languages using combinations of symbols and specialized operations:

Concatenation  $(ab)$  - a matching word has  $a$  followed by  $b$

Alternation  $(a|b)$  - a matching word has  $a$  or  $b$  but not both

Repetition  $(a^*)$  - a matching word has 0 or more occurrences of  $a$

Furthermore, we can group expressions into subexpressions using parenthesis. For example,  $a(a|b)^*$  matches an  $a$  followed by 0 or more  $a$ 's and  $b$ 's. Note that this is all essentially just shorthand for the rather verbose set notation for regular languages.

## 2 DFA Problems

Draw DFA diagrams for the following languages:

1. The language of strings over  $\Sigma = \{a, b, c\}$  that contain only one  $a$  and an even number of  $c$ 's (no restriction on number of  $b$ 's).
2. The language of strings over  $\Sigma = \{0, 1\}$  that end in 1011.
3. The language of strings over  $\Sigma = \{0, 1, 2, 3\}$  which are integers whose digit sum is 3. Leading zeros are permitted.
4. The language of strings over  $\Sigma = \{a, b, c\}$  that end in  $cab$  and contain an even number of  $a$ 's (no restriction on the number of  $b$ 's or  $c$ 's).

## 3 Regular Expression Problems

Build the following languages using combinations of finite languages with regular operations (set notation):

1. Construct the language of binary strings whose second letter is a '0' and whose 5th is a '1'.
2. Construct the language of binary strings that contain the substring "110101".

Provide a regular expression for each of the following languages:

1.  $\Sigma = \{x, y\}, L = \{xx, xy, yx, yy\}$
2.  $\Sigma = \{G, C, A, T\}, L = \text{all strings containing GACAT}$

3. Convert your solutions to the two regular language problems above into regular expressions.
4. Strings over the alphabet  $\Sigma = \{a, b, +, -, *, /\}$  representing valid arithmetic expressions with no parentheses. All operators should be binary (thus  $a + -b$  is not valid) and multiplication must be written explicitly (thus  $a * b$  is valid but  $ab$  is not).
5.  $\Sigma = \{0, 1, 2\}, L = \{x \in \Sigma^* \mid x \text{ contains an even number of 0's and at least one 1.}\}$