## Math 239 Lecture 31

## Graham Cooper

July 22nd, 2015

## **Matchings**

<u>Definition:</u> An <u>Alternating path</u> P with respect to a matching M is a path where consecutive edges alternate between being in M and not in M. An <u>augmenting path</u> is an alternating path that starts and ends with unsaturated vertices.

If there is an augmenting path, we can "switch" edges between the path and M to get a larger matching.

<u>Theorem</u> If there is an augmenting path with respect to M, then M is not a maximum matching. (By switching along the path, we saturate 2 more vertices, and get a matching that has one more edge.)

<u>Theorem</u> If there is no augmenting path with respect to M, then M is a maximum matching

## Vertex Cover

<u>Definition</u> A <u>vertex cover</u> C of a graph G is a set of vertices such that each edge of G has at least one end in C.

General Question: Given a graph, what is the smallest size of a vertex cover?

**Theorem:** If M is any matching and C is any cover, then  $|M| \leq |C|$ .

**PRoof:** For each edge uv in M, at least one of u or v is in C. For different edges in M, different vertices are in C since matchings use distinct vertices So  $|M| \leq |C|$ 

<u>Corollary:</u> If a mathing M and acover C satisfy |M| = |C| then M is a max matching and C is a min cover