

Math 239 Lecture 23

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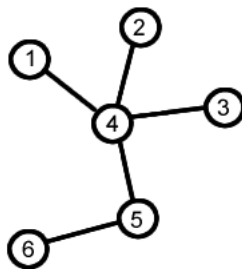
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Items:

- Trees
- Spanning Trees

Trees

Every tree with n vertices has $n-1$ edges



Theorem: Every tree with n vertices has $n-1$ edges.

Proof: By induction on n .

Base Case: When $n = 1$, there is only one tree with 1 vertex and that tree has $n-1 = 0$ edges

Induction Hypothesis: Assume every tree with $n-1$ vertices has $n-2$ edges.

Induction Step: Let T be a tree with n vertices. Let v be a leaf of T and let e be the only edge incident with v . Remove e and v from T to get T' . When we remove e from T there are 2 components, one consists of only v and the other is T' . So T' is connected. Also, T' has no cycles since T has no cycles. So T' is a tree with $n-1$ vertices. By induction hypothesis T' has $n-2$ edges, so T has $n-1$ edges

Q: How many edges does a forest with n vertices and k components have?

A: $n-k$, each component has one fewer edge than vertices. \implies k components mean k fewer edges in total

Theorem: There is a unique path between every pair of vertices in a tree

Proof: Suppose BWOC there are two u,v -paths P_1, P_2 . There is an edge $e=xy$ in one path but not the other. WLOG $e \in P_1$, $P_1 = e, v_1, \dots, v_k, x, y, v_{k+1}, \dots, v$. Then $x, v_k, \dots, u, P_2, v_{k+1}, y$ is an x,y -walk not using e . So there is an x,y -path P_3 , without using e . Then $P_1 + e$ is a cycle in the tree, contradiction.

Theorem: A tree is bipartite

Proof: By induction!

Spanning Trees

Definition: T is a spanning tree of a graph G if T is a subgraph of G that is a tree and uses every vertex in G

