# Lec18 CS241

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### **Grammar Rules**

### Expressions

Address of

 $\frac{E:int}{\&E:int*}$ If E has type int, then &E has type int \*

Dereference:

 $\frac{E:int*}{*E:int}$ IF E has type int \* then \*E has type int

Allocation:

 $\frac{E : int}{newint[E] : int *}$ 

Mult:

 $\frac{E_{1}:intE_{2}:int}{E_{1}*E_{2}:int}$ Similarly for /, %

Addition:

 $\tfrac{E_1:int\cap E_2:int}{E_1+E_2:int}$ 

 $\frac{E_1{:}int{*}{\cap}E_2{:}int}{E_1{+}E_2{:}int{*}}$ 

 $\tfrac{E_1:int\cap E_2:int*}{E_1+E_2:int*}$ 

Substraction:

 $E_1:int\cap E_2:int$  $E_1+E_2:int$ 

 $\frac{E_1:int*\cap E_2:int}{E_1+E_2:int*}$ 

 $E_1:int*\cap E_2:int*$  $E_1+E_2:int$ 

Functions:

 $<\!f,\!(\tau_1,\!\tau_2...\tau_n)\!>\in\!procsE_1\!:\!\tau_1...E_n\!:\!\tau_n$  $F(E_1,...E_n):int$ 

### **Additional Type Constrains**

```
Loops, if statements while(T) \{S\} if (T) \{S_1\} else \{S_2\}
```

Test T hould be boolean, not int or int\*. But there is not boolean type in WLP4

The grammar ensures that T will be a boolean tests (Expr comparison Expr), so as long as the comparison is well-typed, the if/while will be correct

Any expression that has a type is well-typed.  $F:\pi$ 

```
\frac{E:\tau}{well-typed(E)}
```

#### Tests:

$$\frac{E_1:\tau \cap E_2:\tau}{well-typed(E_1==E_2)}$$
similar with !=

$$\begin{array}{l} \frac{E_1:\tau\cap E_2:\tau}{well-typed(E_1< E_2)}\\ \text{Similar with }>,>=,<= \end{array}$$

#### Assignment:

$$\frac{E_1:\tau \cap E_2:\tau}{well-typed(E_1=E_2)}$$

#### Print:

$$\frac{E{:}int}{well-typed(println(E))}$$

#### Deallocation:

$$\frac{E{:}int{*}}{well{-}typed(delete[]E)}$$

#### IF:

$$\frac{well-typed(T)\cap well-typed(S_1)\cap well-typed(S_2)}{well-typed(if(T)(\{S_1\}else\{S_2\})}$$

#### Loops:

```
\frac{well\_typed(T) \cap well\_typed(S)}{well\_typed(while(T)\{S\})}
```

#### Sequences:

 $\overline{well-typed(\epsilon)}$ 

```
\frac{well-typed(S_1)\cap well-typed(S_2)}{well-typed(S_1,S_2)}
Dcls:
\frac{well-typed(intid=NUM)}{well-typed(int*id=NULL)}
Procedures:
\frac{well-typed(dcls)\cap well-typed(stmts)\cap E:int}{well-typed(INT,ID,(params)\{dcls,stmts,return,E;\})}
WAIN:
dcl_2=INT,ID\cap well-typed(dcls)\cap well-typed(stmts)\cap E:int
```

 $well-typed(INT,WAIN(dcl,dcl_2)\{dcls,stmts,return,E;\})$ 

## **Code Generation**

Source  $\to$  Lexical Analysis  $\overset{tokens}{\to}$  parsing  $\to$  semantic analysys  $\overset{parsetree and symbol table}{\to}$  Code Generation  $\to$  assembly

- By now, the source proram is guaranteed to be free of compile-time errors.
- We must now output equivalent MIPS code

There are  $\infty$  many equivalent MIPS programs, which do we choose?

- correct (essential)
- easiest (for cs241)
- shortest
- fastest to run
- fastest to compile

### 0.1 Example:

int wain(int a, int b) { return a;}

#### Convention:

- parameters of wain are in \$1 and \$2
- output will be in \$3

### MIPS:

add \$3, \$1, \$0 ir \$31

int wain (int a, int b) {return b;}

MIPS:

add \$3, \$2, \$0

jr \$31

Parse trees for these programs are the same.

How do we tell one from the other?

Determine which of the two declared ID's the returned ID matches

How do we do that? Use the symbol table.

Add a field to the symbol table to indicate where each symbol is stored

| Name | TYpe | Location |
|------|------|----------|
| a    | int  | \$1      |
| b    | int  | \$2      |

Then when traversing for code gen, look up ID in the symbol table

 $a \rightarrow \$1$ 

 $b \rightarrow \$2$ 

Local Declerations?

- cant all be in regs, probably not enough For simplicity:

• put all local vars on the stack, including params of wain.

int wain(int a, int b){return a;}

sw \$1, -4 (\$30)

sw \$2, -8(\$30)

lis \$4 .word 8 sub \$30, \$30, \$4 lw \$3, 4(\$30) add \$30, \$30, \$4 jr \$31