

Math 239 - Lecture 5

Graham Cooper

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Generating Series

Recal: Set S , weight w :

$$\Phi_s(x) = \sum_{\delta \in S} x^{w(\delta)} = \sum_{k \geq 0} a_k x^k$$

a_x = number of things in S of weight k

Example: How many ways can we throw 2 ∞ sided dice to get a sum of k ?

$(a, b) \in N \times N$

Let $S = N \times N$ Define $w(a, b) = a + b$.

Coeff of x^k in $\Phi_s(x)$ is $k-1$

$$\begin{aligned} \text{So } \Phi_s(x) &= x^2 + 2x^3 + 3x^4 + 4x^5 + \dots = \sum_{k \geq 2} (k-1)x^k \\ &= \frac{x^2}{(1-x)^2} \end{aligned}$$

Notes on Generating Series:

1. General Steps (given a counting problem):
 - (a) Define a set of objects S
 - (b) Define Weight function w related to our problem
 - (c) Find Generating series of S with respect to w , $\Phi_s(x)$
 - (d) The answer is in some coeff of $\Phi_s(x)$
2. X is a literal, we do not put values into x . The coeff of x^k keeps track of answers to counting problems.
3. Later problems involve finding generating series first then possibly finding the coefficient.

Q: How many binary strings have no 3 consecutive 1's?

A: The answer is the coeff of x^n in $\frac{1+x+x^2}{1-x-x^2-x^3}$

Formal Power Series

Definition: Let (a_0, a_1, a_2, \dots) be a sequence of numbers. The formal power series associated with $\{a_n\}_{n \geq 0}$ is:

$$\begin{aligned} A(x) &= a_0 + a_1x + a_2x^2 + \dots \\ &= \sum_{k \geq 0} a_k x^k \end{aligned}$$

We say a_k is the coeff of x^k , denoted $[x^k]A(x)$

Example: $A(x) = 1 + 3x + 5x^2$.

Then $[x^2]A(x) = 5$

Let $A(x) = \sum_{k \geq 0} a_k x^k$ and $B(x) = \sum_{k \geq 0} b_k x^k$

$A(x) = B(x)$ if and only if $[x^k]A(x) = [x^k]B(x)$ for all $k \geq 0$

Two Operations:

1. Addition: $A(x) + B(x) = \sum_{k \geq 0} (a_k + b_k)x^k$

2. Multiplication:

Example: $(1 + 2x + 3x^2)(1 - 3x + 5x^2)$

Coeff of x^2 comes from

$1 \cdot 5x^2, 2x(-3x), 3x^2 \cdot 1$

$\implies 5x^2 - 6x^2 + 3x^2 = 2x^2$

In Genera:

$$\begin{aligned} A(x)B(x) &= \left(\sum_{i \geq 0} a_i x^i\right) \left(\sum_{j \geq 0} b_j x^j\right) \\ &= \sum_{n \geq 0} (a_0 b_n + a_1 b_{n-1} + \dots + a_n b_0) \\ &= \sum_{n \geq 0} \left(\sum_{i=0}^n a_i b_{n-i}\right) \\ &= \sum_{i \geq 0} \sum_{j \geq 0} a_i b_j x^{i+j} \end{aligned}$$

Tool:

$$[x^n]x^k A(x)$$

$$A(x) = 1 + 2x + 3x^2$$

$$[x^5]x^4 A(x) = [x^5](x^4 + 2x^5 + 3x^6) = 2 = [x^1]A(x)$$

$$[x^n]x^k A(x) = [x^{n-k} A(x)]_{n \geq k}$$

$$\text{Example: Let } A(x) = \sum_{i \geq 0} x^i \quad B(x) = \sum_{j \geq 0} (j+1)x^j$$

$$A(x) = 1 + x + x^2 + x^3 \dots$$

$$B(x) = 1 + 2x + 3x^2 + 4x^3 + \dots$$

$$\begin{aligned} [x^n]A(x)B(x) &= 1 \cdot (n+1) + 1 \cdot n + 1 \cdot (n-1) + \dots + 1 \cdot 1 \\ &= (n+1) + n + \dots + 1 \\ &= \frac{(n+1)(n+2)}{2} \end{aligned}$$

$$\begin{aligned} [x^n]A(x)B(x) &= \sum_{i=0}^n n[x^i]A(x)[x^{n-i}]B(x) \\ &= \sum_{i=0}^n 1 \cdot (n-i+1) \\ &= \frac{(n+1)(n+2)}{2} \end{aligned}$$