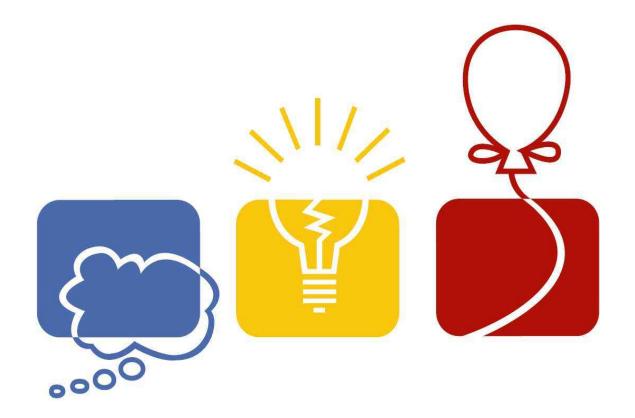
Orzjh 的 XCPC 算法模板

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1 动态规划

1.1 LIS & LCS & LCIS

```
/*
   LIS O(nlogn)实现
   第一问求最长单调不上升子序列
   第二问求最长单调上升子序列
   input: 389 207 155 300 299 170 158 65
   output: 6 2
   f[i]代表当最长单调子序列长度为i时最优的末尾元素
   len为最长单调子序列的长度
   以最长不上升子序列为例
   ①当a[i] <= f[len] 说明可以接在后面
   ②当a[i] > f[len] 在f中找到第一个小于a[i]的数 并替换为a[i]
13
14
   int LIS() {
      while(cin >> a[++n]); n--;
      for(int i = 1; i <= n; i++) cout << a[i] << " "; cout << endl;</pre>
      printf("%d\n%d", LIS_Solve1(), LIS_Solve2());
      //problem1
20
      int len = 1; f[1] = a[1];
      for(int i = 2; i <= n; i++) {
          if(a[i] <= f[len]) f[++len] = a[i];</pre>
          else f[upper_bound(f + 1, f + 1 + len, a[i], greater<int>() ) - f] = a[i];
24
25
      return len;
26
      //problem2
27
      int len = 1; f[1] = a[1];
      for(int i = 2; i <= n; i++) {
          if(a[i] > f[len]) f[++len] = a[i];
          else f[lower_bound(f + 1, f + 1 + len, a[i]) - f] = a[i];
32
      return len;
33
   int LCS() {
      scanf("%s%s", S + 1, T + 1);
37
       int sLen = strlen(S + 1), tLen = strlen(T + 1);
38
      for(int i = 1; i <= sLen; i++) {</pre>
39
         for(int j = 1; j <= tLen; j++) {</pre>
40
             if(S[i] == T[j]) dp[i][j] = dp[i - 1][j - 1] + 1;
41
             else dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
          }
      // 记录路径
      int p = 0, px = sLen, py = tLen;
      while(px != 0 && py != 0) {
          if(dp[px][py] == dp[px][py - 1]) py--;
         else if(dp[px][py] == dp[px - 1][py]) px--;
          else if(dp[px][py] == dp[px - 1][py - 1] + 1) {
50
             ans[++p] = S[px];
51
             px--, py--;
52
53
54
      for(int i = p; i >= 1; i--) printf("%c", ans[i]);
   }
   void LCIS() { // O(nm)
      a[0] = b[0] = -1;
59
      scanf("%d", &n); for(int i = 1; i <= n; i++) scanf("%d", &a[i]);
60
      scanf("%d", &m); for(int i = 1; i <= m; i++) scanf("%d", &b[i]);</pre>
```

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```
62
      for(int i = 1; i <= n; i++) {
63
          int curmax = dp[i - 1][0] + 1, curpos = 0;
          for(int j = 1; j <= m; j++) {
             if(a[i] == b[j]) {
                 if(curmax > dp[i][j]) dp[i][j] = curmax, pre[i][j] = curpos;
             } else dp[i][j] = dp[i - 1][j], pre[i][j] = pre[i - 1][j];
             if(b[j] < a[i] \&\& curmax < dp[i - 1][j] + 1) curmax = dp[i - 1][j] + 1, curpos = j;
          }
73
      int ans = 0, pos = 0;
74
      for(int i = 1; i \le m; i++) if(ans < dp[n][i]) ans = dp[n][i], pos = i;
75
      printf("%d\n", ans);
76
      while(pos) vec.push_back(b[pos]), pos = pre[n][pos];
      reverse(vec.begin(), vec.end());
      for(auto i : vec) printf("%d ", i);
80
   }
```

1.2 背包问题

```
N种物品 容量为V的背包 第i种物品的代价为w[i] 价值为v[i] 数量为m[i]
   void Prework() {
      dp[0] = 0;
      for(int i = 1; i <= V; i++) {</pre>
          //dp[i] = INF; //要求必须装满背包
         dp[i] = 0; //不一定要装满背包
10
   }
11
   void ZeroOnePack() { //01背包
      Prework();
      for(int i = 1; i <= N; i++) {</pre>
          for(int j = V; j >= w[i]; j--) {
             dp[j] = max(dp[j], dp[j - w[i]] + v[i]);
          }
20
   void CompletePack() { //完全背包
22
      Prework();
23
      for(int i = 1; i <= N; i++) {
24
          for(int j = w[i]; j <= V; j++) {</pre>
             dp[j] = max(dp[j], dp[j - w[i]] + v[i]);
          }
      }
28
   }
   void MultiplePack() { //多重背包 利用二进制优化
      Prework();
      int cnt = 0;
      for(int i = 1; i <= N; i++) {
35
          int cur = 1; //将cur个相同物品合并
36
          int cur_w, cur_v;
         while(m[i] >= cur) {
             cur_w = w[i] * cur;
             cur_v = v[i] * cur;
             for(int j = V; j >= cur_w; j--) {
```

```
dp[j] = max(dp[j], dp[j - cur_w] + cur_v);
42
             }
43
             m[i] -= cur;
             cur <<= 1;
          }
          cur_w = w[i] * m[i];
          cur_v = v[i] * m[i];
49
          for(int j = V; j >= cur_w; j--) {
              dp[j] = max(dp[j], dp[j - cur_w] + cur_v);
51
52
53
54
55
56
    void MixedPack() { //混合背包
       for(int i = 1; i <= n; i++) {
          if(m[i] == 0) {
59
              //代表无限个 采用完全背包策略
60
              for(int j = w[i]; j <= V; j++) {</pre>
                 dp[j] = max(dp[j], dp[j - w[i]] + v[i]);
             }
          } else if(m[i] == 1) {
              //01背包策略
             for(int j = V; j >= w[i]; j--) {
66
                 dp[j] = max(dp[j], dp[j - w[i]] + v[i]);
67
          } else if(m[i] > 1) {
             //多重背包策略
             int cur = 1; //将cur个相同物品合并
             int cur_w, cur_v;
             while(m[i] >= cur) {
                 cur w = w[i] * cur;
                 cur_v = v[i] * cur;
                 for(int j = V; j >= cur_w; j--) {
                    dp[j] = max(dp[j], dp[j - cur_w] + cur_v);
                 m[i] -= cur;
79
                 cur <<= 1;
80
             }
81
             cur_w = w[i] * m[i];
             cur_v = v[i] * m[i];
             for(int j = V; j >= cur_w; j--) {
                 dp[j] = max(dp[j], dp[j - cur_w] + cur_v);
86
             }
          }
       }
89
    void TwoDimensionPack() {
92
93
       二维费用背包问题
94
       设第二维容量为T 第i个物品的代价为g[i]
95
       01背包下变量j和k采用逆序循环
       完全背包下采用顺序循环
       多重背包时拆分物品
       */
100
101
       for(int i = 1; i <= n; i++) { //二维01背包
102
          for(int j = V; j >= w[i]; j--) {
              for(int k = T; k >= g[i]; k--) {
104
                 dp[j][k] = max(dp[j][k], dp[j - w[i]][k - g[i]]);
105
106
```

```
}
107
108
110
111
    void DivideGroupPack() { //分组背包
112
113
       这些物品被划分为若干组,每组中的物品互相冲突,最多选一件。
114
       求解将哪些物品装入背包可使这些物品的费用总和不超过背包容量,且价值总和最大。
       设第i件物品的类型为type[i],共有k个类型。
       每组物品有若干种策略:是选择本组的某一件,还是一件都不选。
118
       设dp[i][j]表示前i组物品花费代价j所获得的最大价值
119
       dp[i][j] = max(dp[i - 1][j], dp[i - 1][j - w[cur]] + v[cur]);
120
       cur代表属于组别i的当前物品
121
       可用滚动数组优化 优化时注意三层循环的顺序不能更改
124
       for(int i = 1; i <= N; i++) Items[ type[i] ].push_back(i); //进行分类
       for(int i = 1; i <= K; i++) {
          int num = Items[i].size();
         for(int j = V; j >= 0; j--) {
             for(int k = 0; k < num; k++) {
                //对k的循环要在对j的循环里面 这样能保证每个组别最多选一个物品
131
                int cur = Items[i][k];
132
                if(w[cur] > j) continue;
133
                dp[j] = max(dp[j], dp[j - w[cur]] + v[cur]);
134
             }
          }
139
140
    void PackOnTree(int u) { //树上背包 利用上下界优化
141
      Size[u] = 1;
       for(int i = 0; i < G[u].size(); i++) {</pre>
143
          int v = G[u][i];
144
         PackOnTree(v);
145
         for(int j = min(m + 1, Size[u] + Size[v]); j >= 0; j--) {
146
             for(int k = max(1, j - Size[u]); k < j && k <= Size[v]; k++) {
147
                f[u][j] = max(f[u][j], f[u][j - k] + f[v][k]);
          Size[u] += Size[v];
152
153
    void RollbackPack() { //回退背包
       // select
       for(int i = son[u]; i >= 1; i--) {
157
          for(int j = sze[u]; j \ge sze[v]; j--) f[i][j] = (f[i][j] + f[i - 1][j - sze[v]]) % MOD;
158
159
160
       // delete
161
      for(int i = 1; i <= son[u]; i++) {</pre>
          for(int j = sze[v]; j <= sze[u]; j++) f[i][j] = (f[i][j] - f[i - 1][j - sze[v]] + MOD) % MOD;</pre>
164
   }
165
```

1.3 斜率优化

其一般形式为 $b_i = min_{j < i} \{y_j - k_i x_j\}$ 或 $b_i = max_{j < i} \{y_j - k_i x_j\}$

```
int head, tail, Q[MAXN];
```

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```
int X(int i) {}
   int Y(int i) {}
   void Init() {}
   int Find(int k) {
      int 1 = head, r = tail;
      while(1 < r) {
          int mid = 1 + r + 1 >> 1;
          int i = Q[mid - 1], j = Q[mid];
          double slope = 1.0 * (Y(j) - Y(i)) / (X(j) - X(i));
11
          // 求min 维护下凸包
          if(slope > 1.0 * k) r = mid - 1;
13
         else 1 = mid;
14
          // 求max 维护上凸包
15
          if(slope < 1.0 * k) r = mid - 1;
         else 1 = mid;
      return Q[1];
19
   }
20
   bool Check(int i, int j, int k) {
      double slope_ij = 1.0 * ( Y(j) - Y(i) ) / ( X(j) - X(i) );
      double slope_jk = 1.0 * (Y(j) - Y(k)) / (X(j) - X(k));
      return slope_jk > slope_ij; // 求min 维护下凸包
      return slope_jk > slope_ij; // 求max 维护上凸包
26
   }
27
28
   void Solve() {
      Init();
      head = 1, tail = 0, Q[++tail] = 0;
      for(int i = 1; i <= n; i++) {
          int j = Find(K[i]); // K[i]为斜率
33
          用i更新f[i];
         while(head < tail && !Check(Q[tail - 1], Q[tail], i) ) --tail;</pre>
         Q[++tail] = i;
      return 0;
38
```

1.4 悬线法

```
悬线法的用途: 针对求给定矩阵中满足某条件的极大矩阵,比如"面积最大的长方形、正方形""周长最长的矩形"等等。适合障碍点较密
      集的情况。
   设给定一个n*m的矩形,其中存在障碍点若干。
   维护三个二维数组, left, right, up数组。
   left[i][j]: 代表从(i, j)能到达的最左位置 初始left[i][j] = j;
   right[i][j]: 代表从(i, j)能到达的最右位置 初始right[i][j] = j;
   up[i][j]: 代表从(i, j)向上扩展最长长度 初始up[i][j] = 1;
   void Solve() {
10
     for(int i = 1; i <= n; i++) for(int j = 2; j <= m; j++) left[i][j] = right[i][j] = j, up[i][j] = 1;
11
     for(int i = 1; i <= n; i++) {
        for(int j = 2; j <= m; j++) if(map[i][j]和map[i][j - 1]不是障碍点) left[i][j] = left[i][j - 1];
        for(int j = m - 1; j >= 1; j--) if(map[i][j]和map[i][j + 1]不是障碍点) right[i][j] = right[i][j + 1];
16
17
     int ans = 0;
18
     for(int i = 1; i <= n; i++) {
        for(int j = 1; j <= m; j++) {
           if(map[i][j]和map[i - 1][j]都不是限制点) {
              up[i][j] = up[i - 1][j] + 1;
```

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1.5 最大子段和

```
1 // 在数列的一维方向找到一个连续的子数列,使该子数列的和最大 dp[i] = max(dp[i - 1] + a[i], a[i]);
```

1.6 SOS DP

```
for(int i = 0; i < n; i++) {
    for(int j = 0; j < (1 << n); j++) {
        if((j >> i) & 1) f[j] += f[j ^ (1 << i)];
    }
}</pre>
```

2 数据结构

2.1 并查集

```
struct DSU { // 按秩合并 可撤销
       struct Node {
          int u, v, val;
      int f[MAXN], dep[MAXN], top = 0;
      Node s[MAXN];
      void Init(int n) { for(int i = 1; i <= n; i++) f[i] = i, dep[i] = 0; }</pre>
      int Find(int u) { return u == f[u] ? u : Find(f[u]); }
11
12
       void Union(int u, int v) { // 按秩合并
13
          u = Find(u), v = Find(v);
14
          if(u == v) return ;
          if(dep[u] > dep[v]) swap(u, v);
          int val = (dep[u] == dep[v]);
          s[++top] = (Node)\{u, v, val\};
          f[u] = v, dep[v] += val;
19
      }
20
      bool Same(int u, int v) { return Find(u) == Find(v); }
      void Undo(int cur) {
24
          while(top > cur) {
25
             int u = s[top].u, v = s[top].v, val = s[top].val; // dep[u] <= dep[v]
26
             f[u] = u, dep[v] -= val, top--;
27
          }
28
   }dsu;
```

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2.2 可持久化并查集

```
//#pragma GCC optimize(2)
   #include<bits/stdc++.h>
   using namespace std;
   const int MAXN = 200005;
   const int MOD = 1e9 + 7;
   struct Node {
       int fa, dep, ls, rs;
   }tree[MAXN * 40];
   int cnt, n, m, root[MAXN];
11
12
   int Build(int 1, int r) {
13
       int dir = ++cnt;
14
       if(l == r) { tree[dir].fa = 1; return dir; }
15
       int mid = 1 + r >> 1;
       tree[dir].ls = Build(l, mid);
17
       tree[dir].rs = Build(mid + 1, r);
18
       return dir;
19
   }
20
21
   int ModifyDep(int rt, int l, int r, int loc, int val) {
22
       int dir = ++cnt;
       tree[dir] = tree[rt];
       if(l == r) { tree[dir].dep += val; return dir; }
25
       int mid = (1 + r) >> 1;
26
       if(loc <= mid) tree[dir].ls = ModifyDep(tree[dir].ls, 1, mid, loc, val);</pre>
27
       if(loc > mid) tree[dir].rs = ModifyDep(tree[dir].rs, mid + 1, r, loc, val);
28
       return dir;
30
31
   int ModifyFa(int rt, int l, int r, int loc, int val) {
32
       int dir = ++cnt;
33
       tree[dir] = tree[rt];
34
       if(l == r) { tree[dir].fa = val; return dir; }
35
       int mid = (1 + r) >> 1;
       if(loc <= mid) tree[dir].ls = ModifyFa(tree[dir].ls, l, mid, loc, val);</pre>
       if(loc > mid) tree[dir].rs = ModifyFa(tree[dir].rs, mid + 1, r, loc, val);
38
       return dir;
39
   }
40
41
   int QueryNode(int rt, int 1, int r, int loc) {
       if(1 == r) return rt;
       int mid = (1 + r) >> 1;
       if(loc <= mid) return QueryNode(tree[rt].ls, l, mid, loc);</pre>
45
       if(loc > mid) return QueryNode(tree[rt].rs, mid + 1, r, loc);
46
   }
47
48
   int Find(int rt, int u) { // 返回值为祖先节点在主席树上对应节点编号
       int cur = QueryNode(rt, 1, n, u);
51
       if(u == tree[cur].fa) return cur;
       return Find(rt, tree[cur].fa);
52
   }
53
   void Union(int p, int u, int v) { // 按秩合并
       u = Find(root[p], u), v = Find(root[p], v);
56
57
       if(u == v) return ;
       if(tree[u].dep > tree[v].dep) swap(u, v);
58
       root[p] = ModifyFa(root[p], 1, n, tree[u].fa, tree[v].fa);
59
       if(tree[u].dep == tree[v].dep) root[p] = ModifyDep(root[p], 1, n, tree[v].fa, 1);
60
61
  |signed main()
```

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2.3 K-D Tree

```
//#pragma GCC optimize(2)
   #include<bits/stdc++.h>
   #define x1 x123456789
   #define y1 y123456789
   using namespace std;
   const int MAXN = 5e5 + 5;
   const double alpha = 0.725;
   struct node {
10
      int x, y, v;
11
   }a[MAXN];
12
13
   int n, cnt;
14
   int rt, L[MAXN], R[MAXN], D[MAXN], U[MAXN], d[MAXN], sze[MAXN], sum[MAXN], ls[MAXN], rs[MAXN];
   int g[MAXN], t;
   bool cmp1(int x, int y) { return a[x].x < a[y].x; }</pre>
   bool cmp2(int x, int y) { return a[x].y < a[y].y; }</pre>
19
20
   void maintain(int p) {
21
       sze[p] = sze[ls[p]] + sze[rs[p]] + 1;
       sum[p] = sum[ ls[p] ] + sum[ rs[p] ] + a[p].v;
       L[p] = R[p] = a[p].x;
      D[p] = U[p] = a[p].y;
25
       if(ls[p]) {
26
          L[p] = min(L[p], L[ls[p]]), R[p] = max(R[p], R[ls[p]]);
27
          D[p] = min(D[p], D[ls[p]]), U[p] = max(U[p], U[ls[p]]);
       if(rs[p]) {
30
          L[p] = min(L[p], L[rs[p]]), R[p] = max(R[p], R[rs[p]]);
31
          D[p] = min(D[p], D[rs[p]]), U[p] = max(U[p], U[rs[p]]);
32
33
   }
34
35
   int build(int 1, int r) {
      if(1 > r) return 0;
       int mid = 1 + r \gg 1;
      double ave1 = 0, ave2 = 0, var1 = 0, var2 = 0;
      for(int i = 1; i <= r; i++) ave1 += a[ g[i] ].x, ave2 += a[ g[i] ].y;</pre>
       ave1 /= (r - l + 1), ave2 /= (r - l + 1);
       for(int i = 1; i <= r; i++) {
          var1 += (a[ g[i] ].x - ave1) * (a[ g[i] ].x - ave1);
          var2 += (a[ g[i] ].y - ave2) * (a[ g[i] ].y - ave2);
45
       if(var1 > var2) {
46
          nth_element(g + l, g + mid, g + r + 1, cmp1);
47
          d[g[mid]] = 1;
       } else {
```

```
nth_element(g + 1, g + mid, g + r + 1, cmp2);
50
           d[g[mid]] = 2;
51
        ls[ g[mid] ] = build(l, mid - 1);
        rs[g[mid]] = build(mid + 1, r);
54
        maintain(g[mid]);
55
        return g[mid];
56
    }
57
    void print(int p) {
59
60
        if(!p) return ;
        print(ls[p]);
61
        g[++t] = p;
62
        print(rs[p]);
63
64
    }
    void rebuild(int &p) {
        t = 0;
67
        print(p);
68
        p = build(1, t);
69
    }
70
    bool bad(int p) {
72
        return alpha * sze[p] <= (double) max(sze[ ls[p] ], sze[ rs[p] ]);</pre>
73
74
75
    void insert(int &p, int k) {
76
77
        if(!p) {
           p = k;
           maintain(p);
           return ;
81
        if(d[p] == 1) {
82
            if(a[k].x <= a[p].x) insert(ls[p], k);</pre>
           else insert(rs[p], k);
        } else {
            if(a[k].y <= a[p].y) insert(ls[p], k);</pre>
           else insert(rs[p], k);
87
88
        maintain(p);
89
        if( bad(p) ) rebuild(p);
90
    }
91
    int query(int p, int xl, int yl, int xr, int yr) {
        if(!p || xr < L[p] || xl > R[p] || yr < D[p] || yl > U[p]) return 0;
94
        if(xl \leftarrow L[p] \&\& R[p] \leftarrow xr \&\& yl \leftarrow D[p] \&\& U[p] \leftarrow yr) return sum[p];
95
        int res = 0;
        if(x1 \le a[p].x \& a[p].x \le xr \& y1 \le a[p].y \& a[p].y \le yr) res += a[p].y;
        return query(ls[p], xl, yl, xr, yr) + query(rs[p], xl, yl, xr, yr) + res;
100
    signed main()
101
102
        int lastans = 0, op, x1, y1, x2, y2, A;
103
        scanf("%d", &n); cnt = 0;
104
        while( scanf("%d", &op) ) {
            if(op == 1) {
               cin >> x1 >> y1 >> A;
107
               x1 ^= lastans;
               y1 ^= lastans;
109
               A ^= lastans;
               a[++cnt] = (node)\{x1, y1, A\};
               insert(rt, cnt);
112
113
           if(op == 2) {
114
```

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```
cin >> x1 >> y1 >> x2 >> y2;
115
              x1 ^= lastans;
116
              y1 ^= lastans;
              x2 ^= lastans;
              y2 ^= lastans;
119
              printf("%d\n", lastans = query(rt, x1, y1, x2, y2) );
120
121
           if(op == 3) break;
122
        return 0;
124
125
```

2.4 珂朵莉树

```
#pragma GCC optimize(2)
   #include<bits/stdc++.h>
    #define ll long long
3
4
   using namespace std;
5
    const int MAXN = 1e5 + 5;
6
    namespace ODT {
       struct node {
          11 l, r;
10
          mutable 11 v;
11
          node(11 1, 11 r, 11 v) : 1(1), r(r), v(v) {}
12
          bool operator < (const node& a) const { return 1 < a.1; }</pre>
13
       };
       set<node> tree;
16
       int n, q, sum;
17
18
       set<node>::iterator split(ll pos) {
19
          auto it = tree.lower_bound( node(pos, 0, 0) );
20
           if(it != tree.end() && it->l == pos) return it;
21
          it--;
          11 1 = it \rightarrow l, r = it \rightarrow r, v = it \rightarrow v;
          tree.erase(it);
24
          tree.insert( node(l, pos - 1, v) );
25
          return tree.insert( node(pos, r, v) ).first;
26
       void assign(ll l, ll r, ll v) {
29
           int tot = 0, len = 0;
30
           auto end = split(r + 1), begin = split(l);
31
32
          for(auto it = begin; it != end; it++) {
33
              len += (it->r - it->l + 1);
              tot += it -> v * (it -> r - it -> l + 1);
           }
37
          tree.erase(begin, end);
          tree.insert( node(l, r, v) );
39
40
          if(v == 1) sum += (len - tot);
          else sum -= tot;
43
44
       // codeforces 915e
45
46
       void solve() {
           cin >> n >> q;
47
          tree.insert( node(1, n, 1) );
           sum = n;
          while(q--) {
50
```

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```
int 1, r, k; cin >> 1 >> r >> k;
51
              assign(l, r, k == 1 ? 0 : 1);
52
              cout << sum << "\n";</pre>
           }
55
56
   using namespace ODT;
57
   signed main()
       ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
61
       ODT::solve();
62
       return 0;
63
```

2.5 树状数组

```
int lowbit(int x) { return x & (-x); }
   void Modify(int x, int k) {
3
      while(x <= n) c[x] += k, x += lowbit(x);
5
   int Query(int x) {
      int res = 0;
      while(x > 0) res += c[x], x -= lowbit(x);
      return res;
10
   }
11
12
   // 树状数组二分 O(logn)求全局第k小
   int Kth(int k) {
      int cnt = 0, res = 0;
      for(int i = log2(n); i >= 0; --i) {
16
          res += (1 << i);
17
          if(res >= n || cnt + c[res] >= k) res -= (1 << i);
18
          else cnt += c[res];
19
      return res + 1;
```

2.6 ST 表

```
void Init() {
    for(int k = 1; (1 << k) <= n; k++) {
        for(int i = 1; i <= n; i++) {
            st[i][k] = max(st[i][k - 1], st[i + (1 << (k - 1) )][k - 1]);
        }
    }
}

int Query(int l, int r) {
    int t = floor(log2(r - l + 1));
    return max(st[l][t], st[r - (1 << t) + 1][t]);
}</pre>
```

2.7 单调队列

```
void QueryMin() {
    memset(Index, 0, sizeof(Index)); memset(Value, 0, sizeof(Value));
    int head = 1, tail = 0;
    for(int i = 1; i <= n; i++) {</pre>
```

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```
while(head <= tail && Index[head] < i - k + 1) head++;</pre>
          while(head <= tail && Value[tail] > a[i]) tail--;
6
          Index[++tail] = i;
          Value[tail] = a[i];
       }
   }
10
11
12
   void QueryMax() {
       memset(Index, 0, sizeof(Index)); memset(Value, 0, sizeof(Value));
       int head = 1, tail = 0;
14
       for(int i = 1; i <= n; i++) {
15
          while(head <= tail && Index[head] < i - k + 1) head++;</pre>
16
          while(head <= tail && Value[tail] < a[i]) tail--;</pre>
17
          Index[++tail] = i;
18
          Value[tail] = a[i];
19
       }
   }
22
   void Init() {
       scanf("%d%d", &n, &k);
       for(int i = 1; i <= n; i++) scanf("%d", &a[i]);</pre>
       QueryMin(); QueryMax();
26
```

2.8 单调栈

2.9 块状数据结构

2.9.1 分块

```
#include<bits/stdc++.h>
   using namespace std;
   const int MAXN = (int)50005;
   int a[MAXN], tag[MAXN], ValueAllOne[MAXN], id[MAXN], sum[MAXN], n, S;
   vector<int> b[MAXN];
   pair<int, int> block[MAXN];
   void Init() {
      S = sqrt(n);
      memset(id, -1, sizeof(id));
10
      for(int i = 1; i \le n; i++) id[i] = (i - 1) / S + 1;
12
       int j = 1;
13
       for(int i = id[1]; i <= id[n]; i++) {</pre>
          // 维护区间和
15
          block[i].first = j;
16
          for(; id[j] == i; j++) sum[i] += a[j];
          block[i].second = j - 1;
```

19

```
// 维护区间排序
20
          for(; id[j] == i; j++) b[i].push_back(a[j]);
21
          sort(b[i].begin(), b[i].end());
       }
23
   }
24
25
   void PushDownSqrt(int cid) {
26
       int 1 = block[cid].first, r = block[cid].second;
       if(ValueAllOne[cid] == 1 || tag[cid] == 0) return ;
       bool flag = true;
30
       for(int i = 1; i <= r; i++) {
31
          for(int j = 1; j <= tag[cid]; j++) {</pre>
32
              if(a[i] == 1) break;
33
              a[i] = floor(sqrt(1.0 * a[i]));
          }
          if(a[i] > 1) flag = false;
36
37
       tag[cid] = 0;
       if(flag) ValueAllOne[cid] = 1;
       sum[cid] = 0;
       for(int i = 1; i <= r; i++) sum[cid] += a[i];</pre>
43
44
   void ModifySqrt(int 1, int r) { // 区间元素都开根号
45
       int lid = id[1], rid = id[r];
46
       if(lid == rid) {
47
          SqrtPushDown(lid);
          for(int i = 1; i <= r; i++) {
              int cur = floor(sqrt(1.0 * a[i]));
50
              sum[lid] -= a[i], sum[lid] += cur;
51
              a[i] = cur;
       } else {
54
          PushDown(lid); PushDown(rid);
          for(int i = 1; id[i] == lid; i++) {
56
              int cur = floor(sqrt(1.0 * a[i]));
57
              sum[lid] -= a[i], sum[lid] += cur;
58
              a[i] = cur;
59
          for(int i = r; id[i] == rid; i--) {
              int cur = floor(sqrt(1.0 * a[i]));
              sum[rid] -= a[i], sum[rid] += cur;
63
              a[i] = cur;
65
          for(int i = lid + 1; i < rid; i++) ++tag[i];</pre>
66
       }
   int QuerySum(int 1, int r) { // 区间求和
70
       int lid = id[l], rid = id[r];
71
       if(lid == rid) {
72
          PushDown(lid);
          int res = 0;
          for(int i = 1; i <= r; i++) res += a[i];
          return res;
76
       } else {
          int res = 0;
          PushDown(lid); PushDown(rid);
          for(int i = 1; id[i] == lid; i++) res += a[i];
          for(int i = r; id[i] == rid; i--) res += a[i];
          for(int i = lid + 1; i < rid; i++) PushDown(i), res += sum[i];</pre>
82
          return res;
83
```

```
}
84
    }
85
    void ModifyAdd(int 1, int r, int c) { // 区间元素都加上c
       int lid = id[l], rid = id[r];
       if(lid == rid) {
89
           b[lid].clear();
           for(int i = 1; i <= r; i++) {</pre>
91
              a[i] += c;
              b[lid].push_back(a[i]);
           for(int i = 1 - 1; id[i] == lid; i--) b[lid].push back(a[i]);
95
           for(int i = r + 1; id[i] == lid; i++) b[lid].push_back(a[i]);
96
           sort(b[lid].begin(), b[lid].end());
97
       } else {
98
           b[lid].clear();
           for(int i = 1; id[i] == lid; i++) {
              a[i] += c;
101
              b[lid].push_back(a[i]);
102
103
           for(int i = 1 - 1; id[i] == lid; i--) b[lid].push_back(a[i]);
           sort(b[lid].begin(), b[lid].end());
           b[rid].clear();
107
           for(int i = r; id[i] == rid; i--) {
108
              a[i] += c;
109
              b[rid].push_back(a[i]);
110
111
           for(int i = r + 1; id[i] == rid; i++) b[rid].push_back(a[i]);
           sort(b[rid].begin(), b[rid].end());
           for(int i = lid + 1; i < rid; i++) tag[i] += c;</pre>
115
       }
116
    }
117
    int QuerySumOfSmaller(int 1, int r, int c) { // 区间查询小于c的数字个数
119
       int lid = id[l], rid = id[r];
120
       if(lid == rid) {
121
           int res = 0;
122
           for(int i = 1; i <= r; i++) {
123
              res += (a[i] + tag[lid] < c);
124
           }
           return res;
       } else {
           int res = 0;
           for(int i = 1; id[i] == lid; i++) {
129
              res += (a[i] + tag[lid] < c);
130
           for(int i = r; id[i] == rid; i--) {
              res += (a[i] + tag[rid] < c);
134
           for(int i = lid + 1; i < rid; i++) {</pre>
135
              res += lower_bound(b[i].begin(), b[i].end(), c - tag[i]) - b[i].begin();
136
137
           return res;
       }
    }
140
    int QueryPre(int 1, int r, int c) { // 区间内查询c的前驱(比其小的最大元素)
142
       int lid = id[1], rid = id[r];
143
       if(lid == rid) {
144
           int res = -INF;
           for(int i = 1; i <= r; i++) {
146
               if(a[i] + tag[lid] < c) res = max(res, a[i] + tag[lid]);</pre>
147
148
```

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```
return res;
149
        } else {
150
           int res = -INF;
           for(int i = 1; id[i] == lid; i++) {
               if(a[i] + tag[lid] < c) res = max(res, a[i] + tag[lid]);</pre>
153
           for(int i = r; id[i] == rid; i--) {
155
               if(a[i] + tag[rid] < c) res = max(res, a[i] + tag[rid]);</pre>
           for(int i = lid + 1; i < rid; i++) {</pre>
158
               int cur = lower_bound(b[i].begin(), b[i].end(), c - tag[i]) - b[i].begin() - 1;
159
               if(cur >= 0) res = max(res, b[i][cur] + tag[i]);
160
161
           return res;
162
163
164
    signed main()
166
167
        Init();
168
169
        //Solve();
        return 0;
170
```

2.9.2 分块求区间众数

```
//#pragma GCC optimize(2)
   #include<bits/stdc++.h>
   #define int long long
   #define pir make pair
   #define pii pair<int, int>
   #define fi first
   #define se second
   using namespace std;
   const int MAXN = 40005;
   const int MAXS = 205;
10
   const int MOD = 1e9 + 7;
11
   int n, m, s, a[MAXN], val[MAXN], id[MAXN], 1[MAXS], r[MAXS], c[MAXN], vis[MAXN], sum[MAXS][MAXN], ans[MAXS][MAXS
       ];
   map<int, int> Map;
14
   void Init() {
16
       int cnt = 0;
       for(auto &i : Map) i.se = ++cnt, val[i.se] = i.fi;
       for(int i = 1; i <= n; i++) a[i] = Map[ a[i] ];</pre>
       s = sqrt(n);
21
       for(int i = 1; i <= n; i++) id[i] = (i - 1) / s + 1, r[ id[i] ] = i;
22
       for(int i = n; i >= 1; i--) l[ id[i] ] = i;
23
24
       for(int i = 1; i <= n; i++) sum[ id[i] ][ a[i] ]++;</pre>
25
       for(int i = 1; i <= id[n]; i++) for(int j = 1; j <= n; j++) sum[i][j] += sum[i - 1][j];</pre>
       for(int i = 1; i <= id[n]; i++) {
28
          int res = 0;
29
          for(int j = i; j >= 1; j--) {
30
              for(int k = l[j]; k <= r[j]; k++) {
                 c[ a[k] ]++;
                 if(c[ a[k] ] > c[res] || (c[ a[k] ] == c[res] && a[k] < res) ) res = a[k];</pre>
              ans[j][i] = res;
          for(int j = 1; j <= n; j++) c[j] = 0;
```

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```
}
38
   }
39
   int Query(int L, int R) {
       int lid = id[L], rid = id[R], res = 0;
42
       if(rid - lid + 1 <= 2) {
43
           res = 0;
          for(int i = L; i <= R; i++) {</pre>
              c[ a[i] ]++;
              if(c[ a[i] ] > c[res] || (c[ a[i] ] == c[res] && a[i] < res) ) res = a[i];</pre>
          for(int i = L; i <= R; i++) c[a[i]] = 0;
49
       } else {
50
          res = ans[lid + 1][rid - 1];
51
           c[res] += sum[rid - 1][res] - sum[lid][res];
52
          vis[res] = 1;
          for(int i = L; i <= r[lid]; i++) c[ a[i] ]++;</pre>
          for(int i = l[rid]; i <= R; i++) c[ a[i] ]++;</pre>
          for(int i = L; i <= r[lid]; i++) {</pre>
              if(!vis[ a[i] ]) {
                  c[ a[i] ] += sum[rid - 1][ a[i] ] - sum[lid][ a[i] ];
                 vis[ a[i] ] = 1;
                 if(c[ a[i] ] > c[res] || (c[ a[i] ] == c[res] && a[i] < res) ) res = a[i];</pre>
62
              }
63
64
          for(int i = l[rid]; i <= R; i++) {</pre>
              if(!vis[ a[i] ]) {
                  c[ a[i] ] += sum[rid - 1][ a[i] ] - sum[lid][ a[i] ];
                 vis[ a[i] ] = 1;
                  if(c[ a[i] ] > c[res] || (c[ a[i] ] == c[res] && a[i] < res) ) res = a[i];</pre>
              }
          }
          for(int i = L; i <= r[lid]; i++) c[ a[i] ] = vis[ a[i] ] = 0;</pre>
73
          for(int i = l[rid]; i <= R; i++) c[ a[i] ] = vis[ a[i] ] = 0;</pre>
          c[ ans[lid + 1][rid - 1] ] = vis[ ans[lid + 1][rid - 1] ] = 0;
75
76
       return res;
77
78
79
   signed main()
       ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
82
       cin >> n >> m;
       for(int i = 1; i <= n; i++) cin >> a[i], Map[ a[i] ] = 1;
       Init();
       int x = 0;
       while(m--) {
          int 1, r; cin >> 1 >> r;
89
          l = (1 + x - 1) \% n + 1, r = (r + x - 1) \% n + 1;
90
          if(1 > r) swap(1, r);
91
          x = val[Query(1, r)];
92
          cout << x << "\n";
       return 0;
95
   }
```

2.9.3 块状链表

```
//#pragma GCC optimize(2)
#include<bits/stdc++.h>
```

```
using namespace std;
   const int MAXN = 2010, MAXM = 2010;
4
   struct Node {
       char ch[MAXM + 5];
      int sze, lnode, rnode;
      void push_back(char c) { ch[sze++] = c; }
10
   p[MAXN + 5];
   int pos1 = 0, pos2 = 0; //光标在块中/块内位置
12
13
   int id[MAXN + 5], idx = 0; //可分配的编号池
14
   void Add(int u, int v) { //将节点v插到节点u的右边
15
      p[v].rnode = p[u].rnode;
16
       p[p[v].rnode].lnode = v;
17
       p[u].rnode = v;
      p[v].lnode = u;
   }
20
21
   void Del(int u) { //删除节点u
22
       p[ p[u].lnode ].rnode = p[u].rnode;
       p[ p[u].rnode ].lnode = p[u].lnode;
       p[u].lnode = p[u].rnode = p[u].sze = 0;
       id[++idx] = u;
26
27
28
   void Merge() {
29
      for(int k = p[0].rnode; k; k = p[k].rnode) {
30
          while(p[k].rnode && p[k].sze + p[p[k].rnode ].sze < MAXM) {
             int rnode = p[k].rnode;
             if(pos1 == rnode) pos1 = k, pos2 += p[k].sze; // 与下一条语句顺序不能调换
             for(int i = 0; i < p[rnode].sze; i++) p[k].push_back( p[rnode].ch[i] );</pre>
             Del(rnode);
35
36
          }
   void Move(int k) { //移动到第k个字符后面
40
      pos1 = p[0].rnode;
41
      while (k > p[pos1].sze) k -= p[pos1].sze, pos1 = p[pos1].rnode;
42
      pos2 = k - 1;
43
   }
44
   void Insert(string S, int n) { //在光标后面插入字符串S, 长度为n
       if(pos2 + 1 != p[pos1].sze) { //分裂
47
          int u = id[idx--];
          for(int i = pos2 + 1; i < p[pos1].sze; i++) p[u].push_back( p[pos1].ch[i] );</pre>
          p[pos1].sze = pos2 + 1;
          Add(pos1, u);
       int cur = pos1, i = 0;
53
      while(i < n) {</pre>
54
          int u = id[idx--];
55
          for(; i < n && p[u].sze < MAXN; i++) p[u].push_back(S[i]);</pre>
56
          Add(cur, u);
57
          cur = u;
      Merge();
60
   }
61
   void Delete(int n) { //删除光标后的n个字符
       if(pos2 + 1 + n <= p[pos1].sze) {
          for(int i = pos2 + 1, j = pos2 + 1 + n; j < p[pos1].sze; i++, j++) {
             p[pos1].ch[i] = p[pos1].ch[j];
66
67
```

```
p[pos1].sze -= n;
68
        } else {
69
           n -= (p[pos1].sze - pos2 - 1);
           p[pos1].sze = pos2 + 1;
           while(p[pos1].rnode && n >= p[ p[pos1].rnode ].sze) {
               n -= p[ p[pos1].rnode ].sze;
               Del(p[pos1].rnode);
           int u = p[pos1].rnode;
           for(int i = 0, j = n; j < p[u].sze; i++, j++) {</pre>
               p[u].ch[i] = p[u].ch[j];
79
           p[u].sze -= n;
80
81
       Merge();
82
    }
83
    void Get(int n) { //获取光标后n个字母
85
       if(pos2 + 1 + n \le p[pos1].sze) {
           for(int i = pos2 + 1; i <= pos2 + n; i++) cout << p[pos1].ch[i];</pre>
        } else {
           n -= (p[pos1].sze - pos2 - 1);
           for(int i = pos2 + 1; i < p[pos1].sze; i++) cout << p[pos1].ch[i];</pre>
           int cur = pos1;
           while(p[cur].rnode && n >= p[ p[cur].rnode ].sze) {
92
              n -= p[ p[cur].rnode ].sze;
93
              for(int i = 0; i < p[ p[cur].rnode ].sze; i++) cout << p[ p[cur].rnode ].ch[i];</pre>
94
               cur = p[cur].rnode;
95
           int u = p[cur].rnode;
           for(int i = 0; i < n; i++) cout << p[u].ch[i];</pre>
99
       cout << "\n";
100
    }
101
    void Prev() { //光标前移
        if(pos2) pos2--;
104
        else pos1 = p[pos1].lnode, pos2 = p[pos1].sze - 1;
105
106
107
    void Next() { //光标后移
108
        if(pos2 != p[pos1].sze - 1) pos2++;
109
        else pos1 = p[pos1].rnode, pos2 = 0;
    }
112
    void Solve() {
113
       for(int i = 1; i <= MAXN; i++) id[++idx] = i;</pre>
114
       Insert("\n", 1); Move(1); //预防越界
        int t; cin >> t;
        while(t--) {
           string op, S;
118
           int n;
119
           cin >> op;
120
           if(op == "Insert") {
121
              cin >> n;
              int cur = 0; S = "";
              while(cur < n) {</pre>
                  char ch = getchar();
                  if(32 <= ch && ch <= 126) S.push back(ch), cur++;
               }
              Insert(S, n);
           else if(op == "Move") cin >> n, Move(n + 1);
130
           else if(op == "Delete") cin >> n, Delete(n);
131
           else if(op == "Get") cin >> n, Get(n);
132
```

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2.9.4 块状数组

```
#include<bits/stdc++.h>
   using namespace std;
   const int MAXN = (int)50005;
   int a[MAXN], tag[MAXN], ValueAllOne[MAXN], id[MAXN], sum[MAXN], n, S;
   vector<int> b[MAXN];
   pair<int, int> block[MAXN];
   void Init() {
       S = sqrt(n);
       memset(id, -1, sizeof(id));
       for(int i = 1; i \le n; i++) id[i] = (i - 1) / S + 1;
11
       int j = 1;
13
       for(int i = id[1]; i <= id[n]; i++) {
14
          // 维护区间和
15
          block[i].first = j;
          for(; id[j] == i; j++) sum[i] += a[j];
          block[i].second = j - 1;
          // 维护区间排序
20
          for(; id[j] == i; j++) b[i].push_back(a[j]);
21
          sort(b[i].begin(), b[i].end());
22
       }
   void PushDownSqrt(int cid) {
26
       int 1 = block[cid].first, r = block[cid].second;
27
       if(ValueAllOne[cid] == 1 || tag[cid] == 0) return ;
28
29
       bool flag = true;
       for(int i = 1; i <= r; i++) {
          for(int j = 1; j <= tag[cid]; j++) {</pre>
              if(a[i] == 1) break;
             a[i] = floor(sqrt(1.0 * a[i]));
          if(a[i] > 1) flag = false;
       tag[cid] = 0;
39
       if(flag) ValueAllOne[cid] = 1;
40
       sum[cid] = 0;
41
       for(int i = 1; i <= r; i++) sum[cid] += a[i];</pre>
42
   }
43
   void ModifySqrt(int 1, int r) { // 区间元素开根号
       int lid = id[l], rid = id[r];
46
       if(lid == rid) {
          SqrtPushDown(lid);
          for(int i = 1; i <= r; i++) {</pre>
             int cur = floor(sqrt(1.0 * a[i]));
              sum[lid] -= a[i], sum[lid] += cur;
              a[i] = cur;
52
53
       } else {
54
          PushDown(lid); PushDown(rid);
```

```
for(int i = 1; id[i] == lid; i++) {
56
              int cur = floor(sqrt(1.0 * a[i]));
57
              sum[lid] -= a[i], sum[lid] += cur;
              a[i] = cur;
60
           for(int i = r; id[i] == rid; i--) {
61
              int cur = floor(sqrt(1.0 * a[i]));
              sum[rid] -= a[i], sum[rid] += cur;
              a[i] = cur;
65
66
           for(int i = lid + 1; i < rid; i++) ++tag[i];</pre>
67
    }
68
69
    int QuerySum(int 1, int r) { // 区间求和
70
       int lid = id[l], rid = id[r];
71
       if(lid == rid) {
           PushDown(lid);
73
           int res = 0;
74
           for(int i = 1; i <= r; i++) res += a[i];
75
76
           return res;
77
       } else {
           int res = 0;
           PushDown(lid); PushDown(rid);
           for(int i = 1; id[i] == lid; i++) res += a[i];
80
           for(int i = r; id[i] == rid; i--) res += a[i];
81
           for(int i = lid + 1; i < rid; i++) PushDown(i), res += sum[i];</pre>
82
83
           return res;
       }
    }
    void ModifyAdd(int 1, int r, int c) { // 区间元素都+c
87
       int lid = id[l], rid = id[r];
88
       if(lid == rid) {
89
           b[lid].clear();
           for(int i = 1; i <= r; i++) {
              a[i] += c;
              b[lid].push back(a[i]);
93
94
           for(int i = l - 1; id[i] == lid; i--) b[lid].push back(a[i]);
95
           for(int i = r + 1; id[i] == lid; i++) b[lid].push_back(a[i]);
96
           sort(b[lid].begin(), b[lid].end());
97
        } else {
           b[lid].clear();
           for(int i = 1; id[i] == lid; i++) {
100
              a[i] += c;
101
              b[lid].push_back(a[i]);
102
           for(int i = 1 - 1; id[i] == lid; i--) b[lid].push_back(a[i]);
           sort(b[lid].begin(), b[lid].end());
105
106
           b[rid].clear();
107
           for(int i = r; id[i] == rid; i--) {
108
              a[i] += c;
109
              b[rid].push_back(a[i]);
           for(int i = r + 1; id[i] == rid; i++) b[rid].push_back(a[i]);
           sort(b[rid].begin(), b[rid].end());
113
114
           for(int i = lid + 1; i < rid; i++) tag[i] += c;</pre>
115
       }
116
117
    int QuerySumOfSmaller(int 1, int r, int c) { // 区间查询小于c的数字个数
119
       int lid = id[l], rid = id[r];
120
```

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```
if(lid == rid) {
121
           int res = 0;
122
           for(int i = 1; i <= r; i++) {
               res += (a[i] + tag[lid] < c);
125
           return res;
126
        } else {
127
           int res = 0;
           for(int i = 1; id[i] == lid; i++) {
               res += (a[i] + tag[lid] < c);
130
131
           for(int i = r; id[i] == rid; i--) {
132
               res += (a[i] + tag[rid] < c);
133
134
           for(int i = lid + 1; i < rid; i++) {</pre>
135
              res += lower_bound(b[i].begin(), b[i].end(), c - tag[i]) - b[i].begin();
           return res;
138
        }
139
    }
140
141
    int QueryPre(int 1, int r, int c) { // 区间内查询c的前驱(比其小的最大元素)
142
        int lid = id[1], rid = id[r];
        if(lid == rid) {
           int res = -INF;
145
           for(int i = 1; i <= r; i++) {
146
               if(a[i] + tag[lid] < c) res = max(res, a[i] + tag[lid]);</pre>
147
           }
148
149
           return res;
        } else {
           int res = -INF;
           for(int i = 1; id[i] == lid; i++) {
152
               if(a[i] + tag[lid] < c) res = max(res, a[i] + tag[lid]);</pre>
153
           for(int i = r; id[i] == rid; i--) {
               if(a[i] + tag[rid] < c) res = max(res, a[i] + tag[rid]);</pre>
           for(int i = lid + 1; i < rid; i++) {</pre>
158
               int cur = lower_bound(b[i].begin(), b[i].end(), c - tag[i]) - b[i].begin() - 1;
159
               if(cur >= 0) res = max(res, b[i][cur] + tag[i]);
160
161
162
           return res;
    }
165
    signed main()
166
167
       Init();
168
        //Solve();
169
        return 0;
    }
171
```

2.9.5 莫队

```
#include<bits/stdc++.h>
//#pragma GCC optimize(2)

using namespace std;

const int MAXN = 50000 + 5;
int BLOCK_SIZE;

struct Query {
  int 1, r, id;
```

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```
bool operator < (const Query &a) const {</pre>
11
          if(1 / BLOCK_SIZE != a.1 / BLOCK_SIZE) return 1 < a.1;</pre>
12
          return (1 / BLOCK_SIZE) & 1 ? r < a.r : r > a.r;
   }q[MAXN];
   int n, m, col[MAXN], cnt[MAXN];
   long long sum = 0;
   pair<long long, long long> ans[MAXN];
   void Update(int c, int num) {
21
       sum -= 1ll * cnt[c] * (cnt[c] - 1) / 2;
22
       cnt[c] += num;
23
       sum += 111 * cnt[c] * (cnt[c] - 1) / 2;
24
25
   signed main()
28
       ios::sync_with_stdio(false); cin.tie(0);
29
       cin >> n >> m;
30
       BLOCK_SIZE = (int)ceil( sqrt(1.0 * n) );
       for(int i = 1; i <= n; i++) cin >> col[i];
       for(int i = 1; i <= m; i++) cin >> q[i].l >> q[i].r, q[i].id = i;
       sort(q + 1, q + m + 1);
       for(int i = 1, l = 1, r = 0; i <= m; i++) {
35
          if(q[i].l == q[i].r) {
36
              ans[ q[i].id ] = make_pair(011, 111);
37
              continue;
          while(l > q[i].l) Update(col[--1], 1);
          while(r < q[i].r) Update(col[++r], 1);</pre>
          while(1 < q[i].1) Update(col[1++], -1);</pre>
          while(r > q[i].r) Update(col[r--], -1);
          long long tot = 111 * (r - 1 + 1) * (r - 1) / 2;
          ans[ q[i].id ] = make_pair(sum, tot);
       for(int i = 1; i <= m; i++) {
          if(ans[i].first) {
48
              long long g = __gcd(ans[i].first, ans[i].second);
49
              ans[i].first /= g, ans[i].second /= g;
50
          } else ans[i].second = 1;
51
          cout << ans[i].first << "/" << ans[i].second << "\n";</pre>
       return 0;
   }
```

2.10 线段树

```
/* 区间加法 区间乘法 区间求和 */
   #include<bits/stdc++.h>
   #define 11 long long
   using namespace std;
   const int MAXN = (int)1e5 + 5;
   int n, m, MOD;
   11 mul[MAXN << 2], add[MAXN << 2], sum[MAXN << 2], a[MAXN];</pre>
10
   inline int ls(int x) { return x << 1; }</pre>
11
   inline int rs(int x) { return x << 1 | 1; }</pre>
12
   void PushUp(int p) {
       sum[p] = (sum[ls(p)] + sum[rs(p)]) % MOD;
  |}
16
```

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```
17
   void PushDown(int p, int l, int r) {
18
       int mid = 1 + r \gg 1;
       mul[ls(p)] = mul[ls(p)] * mul[p] % MOD;
       mul[rs(p)] = mul[rs(p)] * mul[p] % MOD;
21
       add[ls(p)] = (add[ls(p)] * mul[p] % MOD + add[p]) % MOD;
       add[rs(p)] = (add[rs(p)] * mul[p] % MOD + add[p]) % MOD;
23
       sum[ls(p)] = (sum[ls(p)] * mul[p] % MOD + add[p] * (mid - 1 + 1) % MOD) % MOD;
       sum[rs(p)] = (sum[rs(p)] * mul[p] % MOD + add[p] * (r - mid) % MOD) % MOD;
       add[p] = 0, mul[p] = 1;
27
28
   void Build(int 1, int r, int p) {
29
       add[p] = 0, mul[p] = 1; // significant
30
31
       if(1 == r) {
          sum[p] = a[1];
          return ;
34
       int mid = 1 + r \gg 1;
35
       Build(l, mid, ls(p));
36
       Build(mid + 1, r, rs(p));
       PushUp(p);
   void ModifyAdd(int nl, int nr, int l, int r, int p, ll k) {
41
       if(nl <= 1 && nr >= r) {
42
          add[p] = (add[p] + k) % MOD;
43
          sum[p] = (sum[p] + 111 * (r - 1 + 1) * k % MOD) % MOD;
44
          return ;
       PushDown(p, 1, r);
       int mid = 1 + r \gg 1;
48
       if(nl <= mid) ModifyAdd(nl, nr, l, mid, ls(p), k);</pre>
49
       if(nr > mid) ModifyAdd(nl, nr, mid + 1, r, rs(p), k);
50
       PushUp(p);
51
52
   void ModifyMul(int nl, int nr, int l, int r, int p, ll k) {
54
       if(nl <= 1 && nr >= r) {
55
          mul[p] = mul[p] * k % MOD;
56
          add[p] = add[p] * k % MOD;
57
          sum[p] = sum[p] * k % MOD;
          return ;
       PushDown(p, 1, r);
61
       int mid = 1 + r \gg 1;
62
       if(nl <= mid) ModifyMul(nl, nr, l, mid, ls(p), k);</pre>
       if(nr > mid) ModifyMul(nl, nr, mid + 1, r, rs(p), k);
       PushUp(p);
67
   11 Query(int nl, int nr, int l, int r, int p) {
68
       if(nl <= 1 && nr >= r) return sum[p];
69
       PushDown(p, 1, r);
70
       int mid = 1 + r >> 1, res = 0;
71
       if(nl <= mid) res = (res + Query(nl, nr, l, mid, ls(p)) ) % MOD;</pre>
       if(nr > mid) res = (res + Query(nl, nr, mid + 1, r, rs(p)) ) % MOD;
       return res;
74
   }
75
   signed main()
       scanf("%d%d%d", &n, &m, &MOD);
       for(int i = 1; i <= n; i++) scanf("%1ld", &a[i]);</pre>
80
       Build(1, n, 1);
81
```

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```
while(m--) {
82
          int op, 1, r; 11 k;
83
          scanf("%d%d%d", &op, &l, &r);
          if(op == 1) {
              scanf("%11d", &k);
86
              ModifyMul(l, r, 1, n, 1, k);
          } else if(op == 2) {
              scanf("%11d", &k);
              ModifyAdd(l, r, 1, n, 1, k);
          } else {
              printf("%lld\n", Query(l, r, 1, n, 1) );
92
93
94
95
       return 0;
96
   }
```

2.10.1 线段树合并分裂

```
// 对每个点开一棵线段树然后能维护的东西,线段树合并都能维护
   #pragma GCC optimize(2)
   #include<bits/stdc++.h>
   #define int long long
   #define pir make_pair
   #define pii pair<int, int>
   #define fi first
   #define se second
   using namespace std;
   const int MAXN = 1e5 + 5;
   const int MOD = 1e9 + 7;
11
   int n, m, Op[MAXN];
   set<int> Set;
14
   struct Node {
16
       int ltype, rtype, ans, sze;
17
       friend Node operator + (Node a, Node b) {
18
19
          if(!a.sze) return b;
          if(!b.sze) return a;
          Node res;
          res.ans = a.ans + b.ans + (a.rtype ^ b.ltype);
22
          res.ltype = a.ltype;
23
          res.rtype = b.rtype;
24
          res.sze = a.sze + b.sze;
25
          return res;
   }tree[MAXN << 2];</pre>
29
   struct Segment_Tree {
30
      int root[MAXN * 40], son[MAXN * 40][3];
31
      Node sum[MAXN * 40];
32
      int pool[MAXN * 40], delcnt = 0, cnt = 0;
33
       void PushUp(int p) { sum[p] = sum[ son[p][0] ] + sum[ son[p][1] ]; }
35
36
      int NewNode() { return delcnt ? pool[delcnt--] : ++cnt; }
37
      void DelNode(int p) {
          pool[++delcnt] = p;
          sum[p] = \{0, 0, 0, 0\};
          son[p][0] = son[p][1] = 0;
42
43
44
      void Insert(int& p, int l, int r, int loc) {
```

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```
if(!p) p = NewNode();
46
           if(1 == r) {
47
              sum[p].ltype = sum[p].rtype = (loc & 1);
              sum[p].ans = 0, sum[p].sze = 1;
              return ;
50
           int mid = 1 + r \gg 1;
52
           if(loc <= mid) Insert(son[p][0], 1, mid, loc);</pre>
           else Insert(son[p][1], mid + 1, r, loc);
           PushUp(p);
55
56
57
       int Merge(int u, int v, int l = 1, int r = n) {
58
           if(!u || !v) return u + v;
59
           if(1 == r) {
60
              sum[u] = sum[u] + sum[v];
              DelNode(v);
              return u;
63
           }
           int mid = 1 + r \gg 1;
65
           son[u][0] = Merge(son[u][0], son[v][0], 1, mid);
           son[u][1] = Merge(son[u][1], son[v][1], mid + 1, r);
           DelNode(v);
           PushUp(u);
           return u;
70
       }
71
72
       void Split(int u, int& v, int k, int flag) { //把u节点分裂, 得到新的放到v里面, 分裂前k个数的节点
73
           if(!k) return ;
74
           v = NewNode();
           if(k >= sum[ son[u][flag] ].sze) {
              Split(son[u][flag ^ 1], son[v][flag ^ 1], k - sum[ son[u][flag] ].sze, flag);
77
              swap(son[u][flag], son[v][flag]);
           } else Split(son[u][flag], son[v][flag], k, flag);
79
           PushUp(u), PushUp(v);
80
81
    }S;
82
83
    void Modify(int loc, int l, int r, int p) {
84
       if(1 == r) {
85
           tree[p] = S.sum[ S.root[loc] ];
86
           if(Op[loc]) swap(tree[p].ltype, tree[p].rtype);
87
           return ;
       }
       int mid = 1 + r \gg 1;
90
       if(loc <= mid) Modify(loc, 1, mid, p << 1);</pre>
91
       else Modify(loc, mid + 1, r, p << 1 | 1);
92
       tree[p] = tree[p << 1] + tree[p << 1 | 1];
93
    }
94
    Node Query(int nl, int nr, int l, int r, int p) {
96
       if(nl <= 1 && nr >= r) return tree[p];
97
       int mid = 1 + r \gg 1;
98
       if(nl > mid) return Query(nl, nr, mid + 1, r, p << 1 | 1);</pre>
99
       else if(nr <= mid) return Query(nl, nr, l, mid, p << 1);</pre>
100
       else return Query(nl, nr, l, mid, p \ll 1) + Query(nl, nr, mid + 1, r, p \ll 1 | 1);
101
    }
102
103
    set<int>::iterator Split(int x) {
104
       auto pos = Set.lower_bound(x), tmp = pos;
105
       if(*pos == x) return pos;
106
       S.Split(S.root[*pos], S.root[x], *tmp - x, Op[*pos] ^ 1);
108
       Modify(*pos, 1, n, 1);
109
       Op[x] = Op[*pos];
110
```

Orzjh-XCPC 第 26 页

```
Modify(x, 1, n, 1);
111
        Set.insert(x);
112
        return Set.lower_bound(x);
113
    }
114
115
    signed main()
116
117
        ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
118
        cin >> n >> m; Set.insert(n + 1);
        for(int i = 1; i <= n; i++) {
120
           int x; cin >> x;
121
           Set.insert(i);
122
           S.Insert(S.root[i], 1, n, x);
123
           Modify(i, 1, n, 1);
124
125
        }
        while(m--) {
           int op, l, r; cin >> op >> l >> r;
           auto pl = Split(l), pr = Split(r + 1);
128
           if(op == 3) {
129
               cout << Query(l, r, 1, n, 1).ans + 1 << "\n";</pre>
130
           } else {
131
               for(auto i = ++pl; i != pr; i++) {
                   S.Merge(S.root[1], S.root[*i]);
134
                  S.root[*i] = 0;
135
                  Modify(*i, 1, n, 1);
136
137
               Set.erase(pl, pr);
138
               Op[1] = op - 1;
               Modify(1, 1, n, 1);
           }
142
        return 0;
143
144
    }
```

2.10.2 扫描线

```
// 求n个矩形的面积并
   #include<bits/stdc++.h>
   #define int long long
   using namespace std;
   const int MAXN = (int)1e5 + 5;
   struct ScanLine {
       int x, lowy, highy, io; //io记录 入边/出边
       ScanLine(){}
       ScanLine(int x, int y1, int y2, int io) : x(x), lowy(y1), highy(y2), io(io){}
10
       bool operator < (const ScanLine &a) const { return x < a.x; }</pre>
11
   }line[MAXN << 1];</pre>
12
13
   int n, ans = 0, tot, cnt, yy[MAXN << 1];</pre>
14
   int length[MAXN << 3], tag[MAXN << 3];</pre>
15
   int ls(int x) { return x << 1; }</pre>
   int rs(int x) { return x << 1 | 1; }</pre>
18
19
   void PushUp(int p, int 1, int r) {
20
       if(tag[p]) length[p] = yy[r] - yy[l];
^{21}
       else if(l + 1 == r) length[p] = 0;
       else length[p] = length[ls(p)] + length[rs(p)];
   }
24
25
   void Build(int 1, int r, int p) { //注意叶子节点的结构 (用区间表示 如[1,2] [2,3])
       if(1 >= r) return ;
```

Orzjh-XCPC 第 27 页

```
tag[p] = length[p] = 0;
28
       if(1 + 1 == r) return;
29
       int mid = (1 + r) >> 1;
       Build(l, mid, ls(p));
       Build(mid, r, rs(p));
32
   }
33
34
   void Update(int nl, int nr, int l, int r, int p, int io) {
35
      if(nl > r || nr < 1) return ;
       if(nl <= l && nr >= r) {
37
          tag[p] += io;
38
          PushUp(p, 1, r);
39
          return ;
40
41
      if(1 + 1 == r) return;
42
      int mid = (1 + r) >> 1;
       if(nl <= mid) Update(nl, nr, l, mid, ls(p), io);</pre>
       if(nr >= mid) Update(nl, nr, mid, r, rs(p), io);
45
      PushUp(p, 1, r);
46
   }
47
   signed main()
49
50
       scanf("%11d", &n);
       cnt = 0;
52
       for(int i = 1; i <= n; i++) {
53
          int x1, y1, x2, y2;
54
          scanf("%lld%lld%lld", &x1, &y1, &x2, &y2); //(x1, y1)为左下角坐标, (x2, y2)为右上角坐标
55
          line[++cnt] = ScanLine(x1, y1, y2, 1);
          yy[cnt] = y1;
          line[++cnt] = ScanLine(x2, y1, y2, -1);
          yy[cnt] = y2;
59
60
61
      sort(yy + 1, yy + 1 + cnt);
       sort(line + 1, line + 1 + cnt);
       tot = unique(yy + 1, yy + 1 + cnt) - (yy + 1); // 离散化, tot记录去重后共多少y值
       ans = 0;
64
       Build(1, tot, 1);
65
       for(int i = 1; i <= cnt; i++) {
66
          ans += length[1] * (line[i].x - line[i - 1].x);
67
          int yl = lower_bound(yy + 1, yy + 1 + tot, line[i].lowy) - yy;
68
          int yr = lower_bound(yy + 1, yy + 1 + tot, line[i].highy) - yy;
          Update(yl, yr, 1, tot, 1, line[i].io);
      printf("%11d", ans);
72
      return 0;
73
   }
74
```

2.10.3 主席树

Orzjh-XCPC 第 28 页

```
int dir = ++cnt;
15
       if(1 == r) {
16
17
          tree[dir].val = a[l];
          return dir;
19
       int mid = (1 + r) >> 1;
20
       tree[dir].ls = Build(l, mid);
21
       tree[dir].rs = Build(mid + 1, r);
22
       return dir;
24
25
   int Modify(int p, int l, int r, int loc, int val) {
26
       int dir = ++cnt;
27
       tree[dir] = tree[p];
28
29
       if(1 == r) {
          tree[dir].val = val;
30
          return dir;
32
       int mid = (1 + r) >> 1;
33
       if(loc <= mid) tree[dir].ls = Modify(tree[dir].ls, 1, mid, loc, val);</pre>
34
       if(loc > mid) tree[dir].rs = Modify(tree[dir].rs, mid + 1, r, loc, val);
35
       return dir;
36
   }
38
   int Query(int p, int l, int r, int loc) {
39
       if(l == r) return tree[p].val;
40
       int mid = (1 + r) >> 1;
41
       if(loc <= mid) return Query(tree[p].ls, 1, mid, loc);</pre>
42
       if(loc > mid) return Query(tree[p].rs, mid + 1, r, loc);
43
   }
   signed main()
46
47
       scanf("%11d%11d", &n, &m);
48
       for(int i = 1; i <= n; i++) scanf("%1ld", &a[i]);</pre>
49
       cnt = 0;
       root[0] = Build(1, n);
51
       for(int i = 1; i <= m; i++) {
52
          int v, op, loc, val;
53
          scanf("%11d%11d%11d", &v, &op, &loc);
54
          if(op == 1) {
55
              scanf("%11d", &val);
              root[i] = Modify(root[v], 1, n, loc, val);
              printf("%lld\n", Query(root[v], 1, n, loc));
59
              root[i] = root[v];
60
           }
61
       }
62
       return 0;
63
```

2.10.4 主席树求静态区间第 k 小

Orzjh-XCPC 第 29 页

```
void Init() { // 离散化
12
       sort(b + 1, b + 1 + tot);
13
       int cnt = 0;
       a[0].val = -0x3f3f3f3f; // significant
       for(int i = 1; i <= n; i++) {
16
          if(a[i].val != a[i - 1].val) ++cnt;
17
          a[i].hval = cnt;
18
19
          Hash[cnt] = a[i].val;
       sort(a + 1, a + 1 + n, cmp2);
21
22
23
    int Build(int 1, int r) {
24
       int dir = ++cnt;
25
26
       tree[dir].val = 0;
       if(1 == r) return dir;
27
       int mid = l + r \gg 1;
       tree[dir].ls = Build(l, mid);
29
       tree[dir].rs = Build(mid + 1, r);
30
       return dir;
31
   }
32
   int Modify(int p, int l, int r, int loc) {
       int dir = ++cnt;
       tree[dir] = tree[p];
36
       tree[dir].val++;
37
       if(1 == r) return dir;
38
       int mid = l + r \gg 1;
39
       if(loc <= mid) tree[dir].ls = Modify(tree[dir].ls, 1, mid, loc);</pre>
       if(loc > mid) tree[dir].rs = Modify(tree[dir].rs, mid + 1, r, loc);
       return dir;
   }
43
44
    int Query(int dl, int dr, int l, int r, int k) {
45
       if(1 == r) return 1;
46
       int mid = l + r \gg 1;
       int x = tree[ tree[dr].ls ].val - tree[ tree[dl].ls ].val;
       if(x >= k) return Query(tree[d1].ls, tree[dr].ls, l, mid, k);
49
       else return Query(tree[dl].rs, tree[dr].rs, mid + 1, r, k - x);
50
   }
51
52
   signed main()
    {
       cin >> n >> m;
       for(int i = 1; i <= n; i++) cin >> a[i], b[++tot] = a[i];
56
       for(int i = 1; i <= m; i++) {</pre>
57
          char ch = getchar();
           if(ch == 'Q') cin >> q[i].l >> q[i].r >> q[i].k;
          else cin >> q[i].x >> q[i].y, b[++tot] = q[i].y;
       Init();
62
63
       return 0;
64
```

2.10.5 李超线段树

```
#include<bits/stdc++.h>
#define int long long

#define fi first

#define se second

#define pir make_pair

#define reg register

#define pdi pair<double, int>
```

```
using namespace std;
   const int MAXN = 100005;
   const int INF = 0x7ffffffff;
   const int MOD1 = 39989;
   const int MOD2 = (int)1e9;
12
13
   struct Tree {
14
15
       int l, r, id;
   }tree[MAXN << 2];</pre>
   struct Seg {
17
       double k, b;
18
       double f(int x) { return 1.0 * k * x + b; }
19
       Seg() {}
20
       Seg(int x0, int y0, int x1, int y1) {
21
          if(x0 == x1) k = 0, b = max(y0, y1);
22
          else {
              k = 1.0 * (y0 - y1) / (x0 - x1);
              b = y0 - 1.0 * k * x0;
25
           }
26
27
   }s[MAXN];
   int n;
   inline int ls(int x) { return x << 1; }</pre>
   inline int rs(int x) { return x << 1 | 1; }
32
   pdi Max(pdi u, pdi v) {
33
       if(u.fi > v.fi) return u;
34
       else if(u.fi < v.fi) return v;</pre>
35
       else {
          if(u.se < v.se) return u;</pre>
          else return v;
       }
39
   }
40
41
   void Build(int 1, int r, int p) {
       tree[p].l = 1;
       tree[p].r = r;
       tree[p].id = 0;
45
       if(1 == r) return ;
46
       int mid = 1 + r \gg 1;
47
       Build(l, mid, ls(p));
48
       Build(mid + 1, r, rs(p));
49
   }
   void Modify(int u, int nl, int nr, int p) {
52
       int l = tree[p].l, r = tree[p].r, v = tree[p].id, mid = l + r >> 1;
53
       double fu = s[u].f(mid), fv = s[v].f(mid);
       if(n1 <= 1 && nr >= r) {
           if(1 == r) {
              if(fu > fv) tree[p].id = u;
              return ;
58
59
          if(s[u].k > s[v].k) {
60
              if(fu > fv) tree[p].id = u, Modify(v, nl, nr, ls(p));
61
              else Modify(u, nl, nr, rs(p));
           } else if(s[u].k < s[v].k) {</pre>
              if(fu > fv) tree[p].id = u, Modify(v, nl, nr, rs(p));
              else Modify(u, nl, nr, ls(p));
           } else {
              if(s[u].b > s[v].b) tree[p].id = u;
          }
          return ;
       if(nl <= mid) Modify(u, nl, nr, ls(p));</pre>
71
       if(nr > mid) Modify(u, nl, nr, rs(p));
72
```

Orzjh-XCPC 第 31 页

```
73
74
    pdi Query(int loc, int p) {
       int l = tree[p].l, r = tree[p].r, id = tree[p].id;
        if(l == r) return pir(s[id].f(loc), id);
77
       int mid = 1 + r \gg 1;
78
        pdi u = pir(s[id].f(loc), id), v;
79
       if(loc <= mid) v = Query(loc, ls(p));</pre>
80
       else v = Query(loc, rs(p));
        return Max(u, v);
82
83
84
    void Modify(int x0, int y0, int x1, int y1, int id) {
85
       s[id] = Seg(x0, y0, x1, y1);
86
87
       Modify(id, x0, x1, 1);
    signed main()
90
91
        scanf("%11d", &n);
92
       Build(1, MOD1, 1);
93
        int last = 0, cnt = 0;
94
        for(int i = 1; i <= n; i++) {
           int op, k, x0, y0, x1, y1;
           scanf("%11d", &op);
97
           if(op == 0) {
98
              scanf("%11d", &k);
99
              k = (k + last - 1 + MOD1) % MOD1 + 1;
100
              last = Query(k, 1).se;
101
              printf("%lld\n", last);
              scanf("%11d%11d%11d", &x0, &y0, &x1, &y1);
104
              x0 = (x0 + last - 1 + MOD1) \% MOD1 + 1;
105
              x1 = (x1 + last - 1 + MOD1) \% MOD1 + 1;
106
              y0 = (y0 + last - 1 + MOD2) \% MOD2 + 1;
107
              y1 = (y1 + last - 1 + MOD2) \% MOD2 + 1;
108
              if(x0 > x1) swap(x0, x1), swap(y0, y1);
109
              Modify(x0, y0, x1, y1, ++cnt);
110
111
112
       return 0;
113
```

2.10.6 可持久化线段树(标记永久化)

```
#include<bits/stdc++.h>
   #define int long long
   using namespace std;
   const int MAXN = 100000 + 5;
   const int MOD = 998244353;
   struct node {
10
       int s, e, p;
   } a[MAXN];
11
12
   struct treenode {
13
       int ls, rs, sum = 0, cnt = 0;
14
   } tree[MAXN * 50];
15
16
   int n, m, tot = 0, root[MAXN];
17
   map<int, int> Map;
18
```

```
int modify(int p, int nl, int nr, int v, int l = 1, int r = n) {
20
       int dir = ++tot;
21
       tree[dir] = tree[p];
       if(nl <= l && nr >= r) {
          tree[dir].sum += v;
          tree[dir].cnt += 1;
25
          return dir;
26
27
       int mid = 1 + r \gg 1;
       if(nl <= mid) tree[dir].ls = modify(tree[dir].ls, nl, nr, v, l, mid);</pre>
       if(nr > mid) tree[dir].rs = modify(tree[dir].rs, nl, nr, v, mid + 1, r);
30
       return dir;
31
   }
32
33
   pair<int, int> query(int p, int loc, int csum = 0, int ccnt = 0, int l = 1, int r = n) {
34
       if(1 == r) return {csum + tree[p].sum, ccnt + tree[p].cnt};
35
       int mid = 1 + r \gg 1;
       if(loc <= mid) return query(tree[p].ls, loc, csum + tree[p].sum, ccnt + tree[p].cnt, l, mid);</pre>
37
       else return query(tree[p].rs, loc, csum + tree[p].sum, ccnt + tree[p].cnt, mid + 1, r);
38
   }
39
40
41
   signed main() {
       ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
       cin >> m >> n;
       for(int i = 1; i <= m; i++) cin >> a[i].s >> a[i].e >> a[i].p, Map[ a[i].p ] = 1;
44
       sort(a + 1, a + 1 + m, [&](const node &a, const node &b) { return a.p < b.p; });</pre>
45
       for(int i = 1; i <= m; i++) root[i] = modify(root[i - 1], a[i].s, a[i].e, a[i].p);</pre>
46
       int pre = 1;
47
       for(int i = 1; i <= n; i++) {
          int x, a, b, c, k; cin >> x >> a >> b >> c;
          k = 1 + (a * pre % c + b) % c;
          int 1 = 1, r = m;
51
          while(1 < r) {
52
              int mid = 1 + r - 1 >> 1;
              if(query(root[mid], x).second < k) l = mid + 1;</pre>
55
              else r = mid;
56
          pre = query(root[1], x).first;
57
          cout << pre << "\n";
58
59
       return 0;
60
   }
```

2.11 平衡树

1 84185 1 89851

2.11.1 普通平衡树

您需要写一种数据结构,来维护一些数,其中需要提供以下操作:

```
1. 插入 x 数
2. 删除 x 数 (若有多个相同的数,因只删除一个)
3. 查询 x 数的排名 (排名定义为比当前数小的数的个数 +1)
4. 查询排名为 x 的数
5. 求 x 的前驱 (前驱定义为小于 x,且最大的数)
6. 求 x 的后继 (后继定义为大于 x,且最小的数)

input:
10
1 106465
4 1
1 317721
1 460929
1 644985
```

```
6 81968
1 492737
5 493598
output:
106465
84185
492737
```

2.11.2 文艺平衡树

题目描述

这是一道模板题。

您需要写一种数据结构,来维护一个序列,其中需要提供以下操作:

翻转一个区间,例如原有序序列是 5 4 3 2 1,翻转区间是 [2,4] 的话,结果是 5 2 3 4 1。

输入格式

第一行为 n,m ,表示初始序列有 n 个数,这个序列依次是 1,2,...n-1,n,m 表示翻转操作次数。接下来 m 行每行两个数 l,r

输出格式

输出一行 n 个数字,表示原始序列经过 m 次变换后的结果。

2.11.3 Treap 普通平衡树

```
#include<bits/stdc++.h>
   using namespace std;
   const int SIZE = (int)1e5 + 5;
   const int INF = 0x3f3f3f3f3f;
   struct Treap {
       int 1, r; // 左右子节点
9
       int val, dat; // 关键码、权值
10
       int cnt, size; // 副本数、子树大小
12
   }a[SIZE];
   int tot = 0, root, n;
13
   int New(int val) { // 创建新节点
15
       a[++tot].val = val;
16
       a[tot].dat = rand();
       a[tot].cnt = a[tot].size = 1;
18
       return tot;
19
   }
20
21
   void Update(int p) { // 更新子树大小
22
       a[p].size = a[a[p].l].size + a[a[p].r].size + a[p].cnt;
23
24
   }
   void Build() { // 建树
26
       New(-INF), New(INF);
27
       root = 1, a[1].r = 2;
28
       Update(root);
29
30
31
32
   void zig(int &p) { // 右旋
       int q = a[p].1;
33
       a[p].1 = a[q].r, a[q].r = p, p = q;
34
       Update(a[p].r), Update(p);
35
36
  |void zag(int &p) { // 左旋
```

```
int q = a[p].r;
39
       a[p].r = a[q].l, a[q].l = p, p = q;
40
41
       Update(a[p].1), Update(p);
    }
43
    void Insert(int &p, int val) { // 插入val数
44
       if(p == 0) {
45
           p = New(val);
          return ;
       if(val == a[p].val) {
49
           a[p].cnt++, Update(p);
50
           return ;
51
52
53
       if(val < a[p].val) {
           Insert(a[p].l, val);
           if( a[p].dat < a[ a[p].l ].dat ) zig(p); // 不满足大根堆性质 右旋
56
       } else {
57
           Insert(a[p].r, val);
           if( a[p].dat < a[ a[p].r ].dat ) zag(p); // 不满足大根堆性质 左旋
       Update(p);
63
    void Remove(int &p, int val) { // 删除val数
64
       if(p == 0) return ;
65
       if(val == a[p].val) {
66
           if(a[p].cnt > 1) {
67
              a[p].cnt--, Update(p);
              return ;
70
           if(a[p].1 || a[p].r) { // 不是叶子节点
              if(a[p].r == 0 || a[ a[p].l ].dat > a[ a[p].r ].dat )
                  zig(p), Remove(a[p].r, val); // 右旋
              else
                 zag(p), Remove(a[p].l, val); // 左旋
76
              Update(p);
77
           } else p = 0;
78
          return ;
79
       }
80
       val < a[p].val ? Remove(a[p].1, val) : Remove(a[p].r, val);</pre>
       Update(p);
83
    }
    int GetRankByVal(int p, int val) { // 查询val数的排名
       if(p == 0) return 0;
       if(val == a[p].val) return a[ a[p].l ].size + 1;
       if(val < a[p].val) return GetRankByVal(a[p].1, val);</pre>
89
       else return GetRankByVal(a[p].r, val) + a[ a[p].l ].size + a[p].cnt;
90
    }
91
92
    int GetValByRank(int p, int rank) { // 查询排名为rank的数
       if(p == 0) return 0;
       if(a[ a[p].l ].size >= rank) return GetValByRank(a[p].l, rank);
       else if(a[ a[p].l ].size + a[p].cnt >= rank) return a[p].val;
96
       else return GetValByRank(a[p].r, rank - a[ a[p].l ].size - a[p].cnt);
97
    }
98
    int GetPre(int val) { // 求val的前驱
100
       int ans = 1; // a[ans].val == -INF;
101
       int p = root;
102
       while(p) {
103
```

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```
if(val == a[p].val) {
104
              if(a[p].1 > 0) {
105
                  p = a[p].1;
                  while(a[p].r > 0) p = a[p].r; // 左子树中一直向右走
                  ans = p;
108
109
              break;
110
111
           if(a[p].val < val && a[p].val > a[ans].val) ans = p; // 更新答案
113
           val < a[p].val ? p = a[p].l : p = a[p].r;
114
115
       return a[ans].val;
116
117
118
    int GetNext(int val) { // 求val的后继
119
       int ans = 2; // a[ans].val = INF;
       int p = root;
121
       while(p) {
           if(val == a[p].val) {
123
              if(a[p].r > 0) {
                  p = a[p].r;
                  while(a[p].1 > 0) p = a[p].1;
                  ans = p;
128
              break;
129
130
           if(a[p].val > val && a[p].val < a[ans].val) ans = p;</pre>
131
           val < a[p].val ? p = a[p].l : p = a[p].r;
       return a[ans].val;
135
    }
136
    signed main()
138
       Build();
140
       srand((unsigned)time(NULL));
141
       scanf("%d", &n);
142
       while(n--) {
143
           int opt, x;
144
           scanf("%d%d", &opt, &x);
           switch(opt) {
              case 1 : Insert(root, x); break;
              case 2 : Remove(root, x); break;
              case 3 : printf("%d\n", GetRankByVal(root, x) - 1 ); break; // 存在-INF
              case 4 : printf("%d\n", GetValByRank(root, x + 1) ); break; // 存在-INF
150
              case 5 : printf("%d\n", GetPre(x) ); break;
              case 6 : printf("%d\n", GetNext(x) ); break;
           }
154
       return 0;
155
156
```

2.11.4 Splay 普通平衡树

```
#include < bits / stdc++.h>

using namespace std;

const int SIZE = (int)1e5 + 5;
const int INF = 0x3f3f3f3f3f;

int son[SIZE][2], fa[SIZE]; // 左右子节点(son[0]代表左儿子, son[1]代表右儿子)、父亲节点
```

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```
int val[SIZE]; // 权值
   int cnt[SIZE], size[SIZE]; // 副本数、子树大小
   int tot = 0, root = 0; // 节点个数、根节点
13
   void Update(int p) { // 更新子树大小
14
      size[p] = size[son[p][0]] + size[son[p][1]] + cnt[p];
15
16
   bool CheckRson(int p) { // 判断节点p是不是右儿子
18
19
      return p == son[ fa[p] ][1];
20
21
   void Clear(int p) { // 销毁节点p
22
      son[p][0] = son[p][1] = fa[p] = val[p] = cnt[p] = size[p] = 0;
23
24
   void Rotate(int p) { //旋转p(根据p的儿子类型判断左旋还是右旋)
26
      int f = fa[p], gf = fa[f]; // f->father gf->grandfather
27
      bool isRson = CheckRson(p);
      son[f][isRson] = son[p][isRson ^ 1];
      if(son[p][isRson ^ 1]) fa[ son[p][isRson ^ 1] ] = f;
      son[p][isRson ^ 1] = f;
      fa[f] = p, fa[p] = gf;
      if(gf) son[gf][ f == son[gf][1] ] = p;
33
      Update(p), Update(f);
34
35
36
   void Splay(int p, int target = 0) { // 节点p旋转到节点target target=0表示根节点
37
      for(int f = fa[p]; (f = fa[p]) != target, f; Rotate(p)) {
          if(fa[f] != target) Rotate( CheckRson(p) == CheckRson(f) ? f : p );
40
      if(!target) root = p;
41
   }
42
   int GetPre() { // 求x的前驱(将x插入 查询x左子树中最右边的节点 删除x)
      int cur = son[root][0];
      if(!cur) return cur;
46
      while(son[cur][1]) cur = son[cur][1];
47
      Splay(cur);
48
      return cur;
49
   }
50
   int GetNext() { // 求x的前驱(将x插入 查询x右子树中最左边的节点 删除x)
      int cur = son[root][1];
53
      if(!cur) return cur;
54
      while(son[cur][0]) cur = son[cur][0];
      Splay(cur);
      return cur;
   int GetRankByVal(int k) { // 查询k数的排名
60
      int res = 0, cur = root;
61
      while(1) {
62
         if(k < val[cur]) cur = son[cur][0];</pre>
             res += size[ son[cur][0] ];
             if(k == val[cur]) {
                Splay(cur);
                return res + 1;
            res += cnt[cur];
             cur = son[cur][1];
          }
72
      }
73
```

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```
74
75
    int GetValByRank(int k) { // 查询排名为k的数
       int cur = root;
       while(1) {
78
           if(son[cur][0] && k <= size[ son[cur][0] ]) cur = son[cur][0];</pre>
79
           else {
80
               k -= size[ son[cur][0] ] + cnt[cur];
               if(k <= 0) {
                  Splay(cur);
                  return val[cur];
85
              cur = son[cur][1];
86
           }
87
88
        }
89
    }
    void Insert(int k) { // 插入k数
91
       if(!root) {
92
           val[++tot] = k, ++cnt[tot];
93
           root = tot, Update(root);
94
           return ;
        int cur = root, f = 0;
       while(1) {
98
           if(val[cur] == k) {
99
              ++cnt[cur];
100
              Update(cur), Update(f);
101
              Splay(cur);
               break;
           f = cur;
105
           cur = son[cur][ val[cur] < k ];</pre>
106
           if(cur == 0) {
107
              val[++tot] = k, ++cnt[tot];
               fa[tot] = f, son[f][val[f] < k] = tot;
               Update(tot), Update(f);
110
               Splay(tot);
111
               break;
112
           }
113
114
    }
115
    void Remove(int k) { // 删除k数
       GetRankByVal(k);
118
        if(cnt[root] > 1) {
119
           --cnt[root];
120
           Update(root);
           return ;
        if(son[root][0] == 0 && son[root][1] == 0) {
124
           Clear(root);
125
           root = 0;
126
        } else if(son[root][0] == 0) {
127
           int cur = root;
           root = son[root][1];
           fa[root] = 0;
           Clear(cur);
131
        } else if(son[root][1] == 0) {
132
           int cur = root;
133
           root = son[root][0];
134
           fa[root] = 0;
           Clear(cur);
136
        } else {
137
           int cur = root, pre = GetPre();
138
```

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```
fa[ son[cur][1] ] = pre;
139
           son[pre][1] = son[cur][1];
140
          Clear(cur);
          Update(root);
143
    }
144
    signed main()
146
       scanf("%d", &n);
148
       while(n--) {
149
           int opt, x;
150
           scanf("%d%d", &opt, &x);
151
           switch(opt) {
152
              case 1 : Insert(x); break;
              case 2 : Remove(x); break;
              case 3 : printf("%d\n", GetRankByVal(x) ); break; // 存在-INF
              case 4 : printf("%d\n", GetValByRank(x) ); break; // 存在-INF
              case 5 : Insert(x), printf("%d\n", val[GetPre()] ), Remove(x); break;
              case 6 : Insert(x), printf("%d\n", val[GetNext()] ), Remove(x); break;
           }
       }
161
       return 0;
```

2.11.5 Splay 文艺平衡树

```
#include<bits/stdc++.h>
   using namespace std;
   const int SIZE = (int)1e5 + 5;
   const int INF = 0x3f3f3f3f;
   int son[SIZE][2], fa[SIZE]; // 左右子节点(son[0]代表左儿子, son[1]代表右儿子)、父亲节点
   int val[SIZE]; // 权值
   int cnt[SIZE], size[SIZE]; // 副本数、子树大小
   bool tag[SIZE]; // 标记
   int tot = 0, root = 0; // 节点个数、根节点
   int n, m, a[SIZE];
11
   void Update(int p) { // 更新子树大小
      size[p] = size[son[p][0]] + size[son[p][1]] + cnt[p];
15
   }
16
   bool CheckRson(int p) { // 判断节点p是不是右儿子
17
      return p == son[ fa[p] ][1];
19
   void Clear(int p) { // 销毁节点p
21
      son[p][0] = son[p][1] = fa[p] = val[p] = cnt[p] = size[p] = 0;
22
23
24
   void Rotate(int p) { //旋转p(根据p的儿子类型判断左旋还是右旋)
      int f = fa[p], gf = fa[f]; // f->father gf->grandfather
      bool isRson = CheckRson(p);
      son[f][isRson] = son[p][isRson ^ 1];
      if(son[p][isRson ^ 1]) fa[ son[p][isRson ^ 1] ] = f;
      son[p][isRson ^ 1] = f;
      fa[f] = p, fa[p] = gf;
      if(gf) son[gf][ f == son[gf][1] ] = p;
      Update(p), Update(f);
34
35
   void Splay(int p, int target = 0) { // 节点p旋转到节点target的儿子下面 target=0表示根节点
      for(int f; (f = fa[p]) != target; Rotate(p)) {
```

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```
if(fa[f] != target) Rotate( CheckRson(p) == CheckRson(f) ? f : p );
38
39
       if(!target) root = p;
    }
42
    void PushDown(int p) { // 下传反转标记
43
       if(p && tag[p]) {
44
           int ls = son[p][0];
45
           int rs = son[p][1];
           tag[ls] ^= 1;
           tag[rs] ^= 1;
48
           swap(son[p][0], son[p][1]);
49
           tag[p] = 0;
50
       }
51
    }
52
    int FindNodeByRank(int k) { // 查询值为k的节点
       int cur = root;
55
       while(1) {
56
           PushDown(cur);
57
           if(son[cur][0] && k <= size[ son[cur][0] ]) cur = son[cur][0];</pre>
           else {
               k \rightarrow size[son[cur][0]] + 1;
              if(!k) return cur;
               else cur = son[cur][1];
62
           }
63
       }
64
65
    }
    int Build(int l, int r, int f) { // 建树
       if(1 > r) return 0;
       int mid = 1 + r \gg 1;
69
       int cur = ++tot;
70
       fa[cur] = f;
71
       cnt[cur] = 1;
       tag[cur] = 0;
       val[cur] = a[mid];
       son[cur][0] = Build(1, mid - 1, cur);
75
       son[cur][1] = Build(mid + 1, r, cur);
76
       Update(cur);
77
       return cur;
78
    }
79
    void Traverse(int p) { // 中序遍历
       if(!p) return ;
82
       PushDown(p);
83
       Traverse(son[p][0]);
       if(abs(val[p]) < INF) printf("%d ", val[p]);</pre>
       Traverse(son[p][1]);
88
    void Reverse(int 1, int r) { // 反转区间[1,r]
89
       int x = FindNodeByRank(1 - 1 + 1); // 存在-INF
90
       int y = FindNodeByRank(r + 1 + 1);
91
       Splay(x, 0);
       Splay(y, x);
       PushDown(root);
       tag[ son[ son[root][1] ][0] ] ^= 1;
95
    }
96
    signed main()
       scanf("%d%d", &n, &m);
100
       a[1] = -INF, a[n + 2] = INF;
101
       for(int i = 1; i <= n; i++) a[i + 1] = i;
102
```

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```
root = Build(1, n + 2, 0);
while(m--) {
    int l, r;
    scanf("%d%d", &l, &r);
    Reverse(l, r);
}

Traverse(root);
return 0;
}
```

2.11.6 FHQ-Treap

```
#include<bits/stdc++.h>
   using namespace std;
   const int MaxSize = (int)1e5 + 5;
   const int INF = 0x3f3f3f3f;
   // T为pair等类型时需要重载运算符
   template <typename T, int MAXN>
   class Treap {
9
10
   private:
      struct Node {
11
          T val;
12
          int ls, rs, sze, priority;
13
      } tree[MaxSize];
14
      int seed, tot, root;
      int Top, Stack[MaxSize];
      int rand() { return seed = (int)(seed * 10483111 % 0x7ffffffff); }
18
19
      void pushup(int p) {
20
          if(p) tree[p].sze = tree[ tree[p].ls ].sze + tree[ tree[p].rs ].sze + 1;
21
      int create(T val) {
24
          int p = Top ? Stack[Top--] : ++tot;
25
          tree[p].val = val;
26
          tree[p].sze = 1;
27
          tree[p].ls = tree[p].rs = 0;
          tree[p].priority = rand();
          return p;
31
32
      // 将根为 p 的子树分裂成 x,y 两部分, x子树中全部小于等于val, y子树中全部大于val
33
      void split(int p, T val, int &x, int &y) {
34
          if(!p) return void(x = y = 0);
          if(val >= tree[p].val) {
             x = p;
37
             split(tree[p].rs, val, tree[p].rs, y);
38
          } else {
39
             y = p;
40
             split(tree[p].ls, val, x, tree[p].ls);
41
          }
          pushup(p);
      }
45
      // 将根为 p 的子树分裂成 x,y 两部分, x子树大小为sze
46
      void split(int p, int sze, int &x, int &y) {
          if(!p) return void(x = y = 0);
          if(tree[ tree[p].ls ].sze + 1 <= sze) {</pre>
             x = p;
             split(tree[p].rs, sze - (tree[ tree[p].ls ].sze + 1), tree[p].rs, y);
          } else {
```

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```
53
              split(tree[p].ls, sze, x, tree[p].ls);
54
           pushup(p);
        }
57
        */
        int merge(int x, int y) {
59
           if(!x \mid | !y) return x + y;
           if(tree[x].priority > tree[y].priority) {
               tree[x].rs = merge(tree[x].rs, y);
               pushup(x);
63
              return x;
64
           } else {
65
              tree[y].ls = merge(x, tree[y].ls);
66
              pushup(y);
67
              return y;
           }
       }
70
    public:
71
       Treap() { seed = (int)(MAXN * 56546311 % 0x7fffffff); }
72
       void insert(T val) {
           int x, y;
           split(root, val - 1, x, y);
           root = merge( merge(x, create(val) ), y );
77
78
79
       void remove(T val) {
80
           int x, y, z;
           split(root, val, x, z);
           split(x, val - 1, x, y);
           if(y) {
              Stack[++Top] = y;
              y = merge(tree[y].ls, tree[y].rs);
           root = merge(merge(x, y), z);
90
       int rank(T val) {
91
           int x, y, res;
92
           split(root, val - 1, x, y);
93
           res = tree[x].sze + 1;
           root = merge(x, y);
           return res;
97
       int val(int rank) {
           int p = root;
           while(1) {
               if(tree[ tree[p].ls ].sze + 1 == rank) break;
              if(tree[ tree[p].ls ].sze + 1 > rank) p = tree[p].ls;
103
               else rank -= (tree[ tree[p].ls ].sze + 1), p = tree[p].rs;
104
105
           return tree[p].val;
106
107
       T prev(T val) {
           int x, y, p; T res;
110
           split(root, val - 1, x, y);
111
           p = x;
112
           while(tree[p].rs) p = tree[p].rs;
           res = tree[p].val;
           root = merge(x, y);
115
           return res;
116
        }
117
```

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```
118
       T next(T val) {
119
           int x, y, p; T res;
           split(root, val, x, y);
           p = y;
122
           while(tree[p].ls) p = tree[p].ls;
           res = tree[p].val;
124
           root = merge(x, y);
125
           return res;
127
128
       bool find(T val) {
129
           int x, y, z;
130
           split(root, val, x, z);
131
           split(x, val - 1, x, y);
132
           bool res = (tree[y].sze > 0);
           root = merge(merge(x, y), z);
           return res;
135
        }
136
137
       int size() { return tree[root].sze; }
    }; Treap<int, 100005> fhqTreap;
139
    signed main()
142
       int q, opt, x; scanf("%d", &q);
143
       while(q--) {
144
           scanf("%d%d", &opt, &x);
145
           switch(opt) {
              case 1 : fhqTreap.insert(x); break;
              case 2 : fhqTreap.remove(x); break;
              case 3 : printf("%d\n", fhqTreap.rank(x) ); break;
              case 4 : printf("%d\n", fhqTreap.val(x) ); break;
              case 5 : printf("%d\n", fhqTreap.prev(x) ); break;
              case 6 : printf("%d\n", fhqTreap.next(x) ); break;
153
154
        return 0;
155
156
```

2.12 树套树

2.12.1 树状数组套主席树求动态区间第 k 小

```
#pragma GCC optimize(2)
   #include<bits/stdc++.h>
   using namespace std;
   const int MAXN = 100000 + 5;
   const int INF = 0x7fffffff;
   struct Query {
       int 1, r, k, x, y;
   }q[MAXN];
10
   map<int, int> Map;
11
   vector<int> lRoot, rRoot;
12
   int n, m, a[MAXN], cnt = 0, num[MAXN << 1], root[MAXN];</pre>
13
   int tot, sum[MAXN << 7], ls[MAXN << 7], rs[MAXN << 7];</pre>
14
   void Modify(int &rt, int 1, int r, int loc, int val) {
16
       if(!rt) rt = ++tot;
17
       sum[rt] += val;
18
       if(1 == r) return ;
19
       int mid = l + r \gg 1;
```

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```
if(loc <= mid) Modify(ls[rt], 1, mid, loc, val);</pre>
21
       else Modify(rs[rt], mid + 1, r, loc, val);
22
   }
23
   int QueryKth(int 1, int r, int k) {
25
       if(l == r) return l;
26
       int mid = 1 + r >> 1, res = 0;
       for(auto i : rRoot) res += sum[ ls[i] ];
       for(auto i : lRoot) res -= sum[ ls[i] ];
       if(res >= k) {
          for(auto &i : lRoot) i = ls[i];
31
          for(auto &i : rRoot) i = ls[i];
32
          return QueryKth(l, mid, k);
33
       } else {
34
          for(auto &i : lRoot) i = rs[i];
35
          for(auto &i : rRoot) i = rs[i];
          return QueryKth(mid + 1, r, k - res);
       }
38
   }
39
40
   void Modify(int pos, int val) {
41
       int loc = Map[ a[pos] ];
       for( ; pos <= n; pos += (pos & (-pos))) Modify(root[pos], 1, cnt, loc, val);</pre>
45
   int Query(int 1, int r, int k) {
46
       lRoot.clear(), rRoot.clear();
47
       for(int pos = 1 - 1; pos; pos -= (pos & (-pos))) lRoot.push_back(root[pos]);
48
       for(int pos = r; pos; pos -= (pos & (-pos))) rRoot.push_back(root[pos]);
       return QueryKth(1, cnt, k);
   }
51
52
   signed main()
53
54
       ios::sync_with_stdio(false); cin.tie(0);
       cin >> n >> m;
       for(int i = 1; i <= n; i++) cin >> a[i], Map[ a[i] ] = 1;
       for(int i = 1; i <= m; i++) {
58
          char ch; cin >> ch;
59
          if(ch == 'Q') cin >> q[i].l >> q[i].r >> q[i].k;
60
          else cin >> q[i].x >> q[i].y, Map[q[i].y] = 1;
61
       for(auto &i : Map) i.second = ++cnt, num[cnt] = i.first;
       for(int i = 1; i <= n; i++) Modify(i, 1);</pre>
       for(int i = 1; i <= m; i++) {
65
          if(q[i].1) cout << num[ Query(q[i].1, q[i].r, q[i].k) ] << "\n";</pre>
          else {
             Modify(q[i].x, -1);
             a[q[i].x] = q[i].y;
             Modify(q[i].x, 1);
          }
71
72
       return 0;
73
```

3 字符串

3.1 字符串双哈希

```
#include<bits/stdc++.h>
using namespace std;

const int MAXN = 6e5 + 10;
const int INF = 0x3f3f3f3f;
```

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```
namespace TwoHash {
       #define fi first
       #define se second
       #define pii pair<int, int>
10
11
       pii Pow[MAXN], Inv[MAXN];
12
       pii h[MAXN];
13
       int n;
       string S;
15
16
       const int BASE1 = 131;
17
       const int BASE2 = 129;
18
       const int MOD1 = 1e9 + 7;
19
       const int MOD2 = 998244353;
20
21
       int qpow(int a, int p, int MOD) {
          a %= MOD, p %= MOD;
23
          int res = 1;
          while(p) {
25
              if(p & 1) res = 1ll * res * a % MOD;
26
              a = 111 * a * a % MOD;
              p >>= 1;
          return res;
30
31
32
       pii Hash(int c, int p) { return {111 * c * Pow[p].fi % MOD1, 111 * c * Pow[p].se % MOD2}; }
33
       pii Add(pii x, pii y) {
          pii res = \{0, 0\};
          res.fi = 111 * (x.fi + y.fi) % MOD1;
37
          res.se = 111 * (x.se + y.se) % MOD2;
38
          return res;
39
40
41
       pii Sub(pii x, pii y) {
42
          pii res = \{0, 0\};
43
          res.fi = 111 * (x.fi - y.fi + MOD1) % MOD1;
44
          res.se = 111 * (x.se - y.se + MOD2) % MOD2;
45
          return res;
46
       }
47
       pii Mul(pii x, pii y) {
          pii res = \{0, 0\};
50
          res.fi = 111 * x.fi * y.fi % MOD1;
51
          res.se = 111 * x.se * y.se % MOD2;
52
          return res;
       pii Div(pii x, pii y) {
56
          pii res = \{0, 0\};
57
          res.fi = 111 * x.fi * qpow(y.fi, MOD1 - 2, MOD1) % MOD1;
58
          res.se = 111 * x.se * qpow(y.se, MOD1 - 2, MOD1) % MOD1;
59
          return res;
       pii HashVal(int 1, int r) { return Mul( Sub(h[r], h[1 - 1]), Inv[1 - 1]); }
63
64
       void Init() {
65
          Pow[0] = Inv[0] = \{1, 1\}; Inv[1] = \{qpow(BASE1, MOD1 - 2, MOD1), qpow(BASE2, MOD2 - 2, MOD2)\};
          for(int i = 1; i <= MAXN - 5; i++) Pow[i] = Mul(Pow[i - 1], {BASE1, BASE2});</pre>
          for(int i = 2; i <= MAXN - 5; i++) Inv[i] = Mul(Inv[i - 1], Inv[1]);</pre>
68
       }
69
```

70

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```
void Build(string T) {
71
          n = T.length(); S = " " + T;
72
          h[0] = \{0, 0\};
          for(int i = 1; i <= n; i++) h[i] = Add(h[i - 1], Hash(S[i], i));
75
   }
76
   signed main()
       TwoHash::Init();
80
       set<pii> Set;
       int T; cin >> T;
82
       while(T--) {
83
          string S; cin >> S;
84
          TwoHash::Build(S);
85
          Set.insert( TwoHash::h[S.length()] );
       cout << Set.size();</pre>
       return 0;
89
   }
```

3.2 Trie

```
int Hash(char ch) { return ch - 'a' + 1; }

void Insert(string S) {
   int u = 0;
   for(auto i : S) {
      if(!trie[u][Hash(i)]) trie[u][Hash(i)] = ++cnt;
      u = trie[u][Hash(i)];
   }
   sum[u]++;
}
```

3.3 KMP

```
#include<bits/stdc++.h>
   using namespace std;
   const int MAXN = 1e6 + 5;
   char S1[MAXN], S2[MAXN];
   int 11, 12, pmt[MAXN], p; //pmt数组(部分匹配表) 向右偏移一位为next数组 并让next[0] = -1
   int main()
11
      cin >> S1 + 1 >> S2 + 1;
      11 = strlen(S1 + 1), 12 = strlen(S2 + 1);
13
      p = pmt[0] = 0;
      for(int i = 2; i <= 12; i++) {
          while(p && S2[p + 1] != S2[i]) p = pmt[p];
          if(S2[p + 1] == S2[i]) p++;
          pmt[i] = p;
18
19
      p = 0;
20
      for(int i = 1; i <= l1; i++) {
21
          while(p && S2[p + 1] != S1[i]) p = pmt[p];
22
          if(S2[p + 1] == S1[i]) p++;
          if(p == 12) {
             printf("%d\n", i - p + 1);
             p = pmt[p];
```

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3.4 exKMP

```
#include<bits/stdc++.h>
   using namespace std;
   const int MAXN = (int)2e7 + 5;
   char S[MAXN], T[MAXN];
   int z[MAXN], lcp[MAXN];
   void Z_Function(char* T) { // z[i] = lcp(s[i ... n-1], s)
       int n = strlen(T);
       z[0] = n;
10
       int 1 = 0, r = 0;
11
12
       for(int i = 1; i < n; i++) {</pre>
          if(i <= r) z[i] = min(z[i - 1], r - i + 1);</pre>
14
          while(i + z[i] < n && T[z[i]] == T[i + z[i]]) z[i]++;
15
          if(i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
16
17
18
       //for(int i = 0; i < n; i++) cout << z[i] << " "; cout << endl;
19
   }
20
21
   void exKMP(char* S, char* T) {
22
       int sLen = strlen(S);
23
       int tLen = strlen(T);
24
       Z_Function(T);
       int p = 0;
       while(S[p] == T[p] \&\& p < min(sLen, tLen)) p++;
       lcp[0] = p;
29
30
       int 1 = 0, r = 0;
31
       for(int i = 1; i < sLen; i++) {</pre>
32
          if(i <= r) lcp[i] = min(z[i - 1], r - i + 1);
          while(i + lcp[i] < sLen && lcp[i] < tLen && S[i + lcp[i]] == T[lcp[i]]) lcp[i]++;
          if(i + lcp[i] - 1 > r) l = i, r = i + lcp[i] - 1;
35
36
37
       //for(int i = 0; i < sLen; i++) cout << lcp[i] << " "; cout << endl;
38
   signed main()
41
42
   {
       scanf("%s%s", S, T);
43
       exKMP(S, T);
44
       return 0;
45
   input:
   aaaabaa
50
   aaaaa
   z function : {5 4 3 2 1}
   lcp function : {4 3 2 1 0 2 1}
   */
```

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3.5 Manacher

```
Manacher算法:
   先在两个字符串中插入某个字符('$') 避免分别处理奇回文和偶回文的情况
   设置两个指针maxR(前i个字符能回文扩展到的最右端) pos(前i个字符中哪个字符能回文扩展到最右端)
   每次扫描到第i个字符时
   ②如果i<maxR,更新f[i] ②暴力拓展maxR(maxR从起点到终点且不会往回退) ③maxR增大时更新maxR和pos
   #include<bits/stdc++.h>
   using namespace std;
11
   const int MAXN = (int)1.1e7 + 5;
   char S[MAXN << 1], T[MAXN << 1];</pre>
   int f[MAXN << 1], n;</pre>
   void Manacher() {
17
      T[0] = '#'; T[1] = '$';
18
      for(int i = 1; i \le n; i++) T[i * 2] = S[i], T[i * 2 + 1] = '$';
19
      n = n * 2 + 1;
20
      for(int i = 0; i <= n; i++) S[i] = T[i];</pre>
21
      int maxR = 0, pos = 0;
      for(int i = 1; i <= n; i++) {
         if(i < maxR) f[i] = min(f[pos * 2 - i], maxR - i);</pre>
         while(i - f[i] - 1 > 0 && i + f[i] + 1 <= n && S[i + f[i] + 1] == S[i - f[i] - 1]) f[i]++;
         if(i + f[i] > maxR) maxR = i + f[i], pos = i;
27
   int main()
31
32
      scanf("%s", S + 1); n = strlen(S + 1);
33
      Manacher();
34
      int ans = 0;
      for(int i = 1; i <= n; i++) ans = max(ans, f[i]);</pre>
      printf("%d", ans);
      return 0;
38
   }
```

3.6 最小最大表示法

```
S的最小表示: 与S循环同构的所有字符串中字典序最小的字符串 (最大表示同理)
   int getMin(string S) {
      int n = S.length(), i = 0, j = 1, k = 0;
      while(i < n && j < n && k < n) \{
         if(S[(i + k) % n] == S[(j + k) % n]) k++;
         else {
            if(S[(i + k) % n] > S[(j + k) % n]) i = i + k + 1;
            else j = j + k + 1;
10
            if(i == j) i++;
            k = 0;
15
      return min(i, j);
16
   }
17
   int getMax(string S) {
      int n = S.length(), i = 0, j = 1, k = 0;
```

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```
while(i < n && j < n && k < n) {
    if(S[(i + k) % n] == S[(j + k) % n]) k++;
    else {
        if(S[(i + k) % n] < S[(j + k) % n]) i = i + k + 1;
        else j = j + k + 1;

        if(i == j) i++;
        k = 0;
    }
}
return min(i, j);
}</pre>
```

3.7 AC 自动机

```
#include<bits/stdc++.h>
   using namespace std;
   const int MAXN = (int)1e6 + 5;
   char S[MAXN], T[MAXN];
   int n, cnt, trie[MAXN][30], fail[MAXN * 30], num[MAXN * 30];
   queue<int> Q;
   inline int Hash(char x) { return x - 'a' + 1; }
10
   void Trie(char* S) {
11
       int cur = 1;
12
       int len = strlen(S);
       for(int i = 0; i < len; i++) {</pre>
           int x = Hash(S[i]);
15
           if(!trie[cur][x]) trie[cur][x] = ++cnt;
16
          cur = trie[cur][x];
17
18
       num[cur]++;
19
   }
20
21
   void GetFail() {
       for(int i = 1; i <= 26; i++) trie[0][i] = 1;</pre>
23
       Q.push(1);
       fail[1] = 0;
25
       while(!Q.empty()) {
          int u = Q.front();
          Q.pop();
28
          int faFail = fail[u];
29
          for(int i = 1; i <= 26; i++) {
30
              int v = trie[u][i];
31
              if(v) fail[v] = trie[faFail][i], Q.push(v);
32
              else trie[u][i] = trie[faFail][i];
33
           }
35
       }
   }
36
   int main()
38
39
       scanf("%d", &n);
       cnt = 1;
       for(int i = 1; i <= n; i++) {
          scanf("%s", S);
43
          Trie(S);
44
45
       GetFail();
       scanf("%s", T);
       int cur = 1, ans = 0, len = strlen(T);
49
```

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```
for(int i = 0; i < len; i++) {
          cur = trie[cur][Hash(T[i])];
          for(int t = cur; t && ~num[t]; t = fail[t]) ans += num[t], num[t] = -1;
}

printf("%d\n", ans);
return 0;
}</pre>
```

3.8 后缀数组

```
/* 注意常数 */
   //#pragma GCC optimize(2)
   #include<bits/stdc++.h>
   using namespace std;
   const int MAXN = 1e6 + 10;
   char S[MAXN];
   int n, m, sa[MAXN], rk[MAXN], oldrk[MAXN << 1], id[MAXN], px[MAXN], cnt[MAXN];</pre>
10
   bool cmp(int x, int y, int w) {
11
       return oldrk[x] == oldrk[y] && oldrk[x + w] == oldrk[y + w];
12
   }
13
   void SuffixArray(char *s) {
15
       // getSA
       n = strlen(s + 1);
       int w, p, i, m = 300, k;
       for(i = 1; i <= n; ++i) ++cnt[ rk[i] = s[i] ];</pre>
       for(i = 1; i <= m; ++i) cnt[i] += cnt[i - 1];</pre>
       for(i = n; i >= 1; --i) sa[ cnt[ rk[i] ]-- ] = i;
22
       for(w = 1; ; w <<= 1, m = p) {
23
          for(p = 0, i = n; i > n - w; --i) id[++p] = i;
24
          for(i = 1; i <= n; ++i) if(sa[i] > w) id[++p] = sa[i] - w;
          for(i = 0; i <= m; ++i) cnt[i] = 0;
          for(i = 1; i <= n; ++i) ++cnt[ px[i] = rk[ id[i] ] ];</pre>
          for(i = 1; i <= m; ++i) cnt[i] += cnt[i - 1];</pre>
          for(i = n; i >= 1; --i) sa[ cnt[ px[i] ]-- ] = id[i];
          for(i = 1; i <= n; ++i) oldrk[i] = rk[i];</pre>
30
          for(p = 0, i = 1; i \le n; ++i) rk[sa[i]] = cmp(sa[i], sa[i - 1], w) ? p : ++p;
          if(p == n) {
              for(int i = 1; i <= n; ++i) sa[ rk[i] ] = i;</pre>
34
              break;
35
          }
36
37
       //for(int i = 1; i <= n; ++i) cout << sa[i] << " "; cout << "\n";
       //getHeight
       for(i = 1, k = 0; i <= n; ++i) {
41
          if(rk[i] == 0) continue;
42
          if(k) --k;
43
          while(S[i + k] == S[sa[rk[i] - 1] + k]) ++k;
          height[ rk[i] ] = k;
48
   signed main()
49
50
       ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
51
       cin >> S + 1;
       SuffixArray(S);
       return 0;
```

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5 | }

3.9 后缀自动机

```
#include<bits/stdc++.h>
   #pragma GCC optimize("-Ofast")
   using namespace std;
   const int MAXN = (int)4e5 + 5;
   struct state {
       int len, link, size;
       int next[30];
   }st[MAXN << 1];</pre>
10
   int sz, last, n, p, q;
11
   long long f[MAXN];
   char s[MAXN];
13
   void init() {
15
       for(int i = 0; i < sz; i++) {</pre>
16
          for(int j = 0; j < 26; j++) st[i].next[j] = 0;</pre>
           st[i].len = st[i].link = st[i].size = 0;
18
19
       sz = 0;
20
21
       st[0].len = 0;
22
       st[0].link = -1;
23
       sz++;
       last = 0;
   }
26
27
   void extend(int c) {
28
       int cur = sz++;
       st[cur].size = 1;
       st[cur].len = st[last].len + 1;
       int p = last;
       while(p != -1 && !st[p].next[c]) {
33
          st[p].next[c] = cur;
34
          p = st[p].link;
35
36
       if(p == -1) st[cur].link = 0;
37
          int q = st[p].next[c];
39
          if(st[p].len + 1 == st[q].len) st[cur].link = q;
40
41
              int clone = sz++;
              st[clone].len = st[p].len + 1;
              for(int i = 0; i < 26; i++) st[clone].next[i] = st[q].next[i];</pre>
              st[clone].link = st[q].link;
              while(p != -1 && st[p].next[c] == q) {
46
                  st[p].next[c] = clone;
47
                  p = st[p].link;
48
49
              st[q].link = st[cur].link = clone;
           }
52
       last = cur;
53
   }
```

4 图论

4.1 DFS 序

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4.2 LCA

```
int n, m, dep[MAXN], vis[MAXN], lca[MAXN][25];
   vector<int> G[MAXN];
   void DFS(int u, int ftr) {
       dep[u] = dep[ftr] + 1;
       vis[u] = 1;
       lca[u][0] = ftr;
       for(int i = 1; i <= 20; i++) {
           if(dep[u] < (1 << i)) break;</pre>
          lca[u][i] = lca[lca[u][i - 1]][i - 1];
       for(int i = 0; i < G[u].size(); i++) {</pre>
12
          int v = G[u][i];
13
           if(vis[v]) continue;
14
          DFS(v, u);
15
16
   }
17
   int LCA(int x, int y) {
19
       int u = x, v = y;
20
       if(dep[u] > dep[v]) swap(u, v); //dep[u] <= dep[v]</pre>
       for(int i = 20; i >= 0; i--) {
22
          if((1 << i) & (dep[v] - dep[u])) v = lca[v][i];</pre>
       for(int i = 20; i >= 0; i--) {
25
          if(lca[u][i] != lca[v][i]) {
26
              u = lca[u][i];
27
              v = lca[v][i];
28
29
       return u == v ? u : lca[u][0];
```

4.3 树的重心

```
int size[MAXN], weight[MAXN], centroid[3];

void GetCentroid(int cur, int fa) {
    size[cur] = 1;
    weight[cur] = 0;
    for(int i = head[cur]; i != -1; i = e[i].nxt) {
        if(e[i].to == fa) continue;
        GetCentroid(e[i].to, cur);
        size[cur] += size[ e[i].to ];
        weight[cur] = max(weight[cur], size[ e[i].to ]);
}
```

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```
weight[cur] = max(weight[cur], n - size[cur]);
if(weight[cur] <= n / 2) centroid[ ++centroid[0] ] = cur;
}</pre>
```

4.4 树的直径

```
int n, d1[MAXN], d2[MAXN], d = 0;
vector<int> G[MAXN];

void DFS(int u, int fa) {
    d1[u] = d2[u] = 0;
    for(auto v : G[u]) {
        if(v == fa) continue;
        DFS(v, u);
        int cur = d1[v] + 1;
        if(cur > d1[u]) d2[u] = d1[u], d1[u] = cur;
        else if(cur > d2[u]) d2[u] = cur;
    }
    d = max(d, d1[u] + d2[u]);
}
```

4.5 最小生成树

```
int Prim() {
       memset(dist, 0x3f, sizeof(dist));
2
       int res = 0;
3
       for(int i = 0; i < n; i++) {</pre>
          int t = -1;
          for(int j = 1; j <= n; j++) {</pre>
              if(!st[j] && (t == -1 || dist[t] > dist[j])) t = j;
          if(i && dist[t] == INF) return INF; // 无解
          if(i) res += dist[t];
          st[t] = true;
          for(int j = 1; j <= n; j++) dist[j] = min(dist[j], G[t][j]);</pre>
14
15
       return res;
16
   }
17
18
   int Kruskal() {
       for(int i = 1; i <= n; i++) f[i] = i;
20
       sort(e + 1, e + 1 + m, cmp); // 接边权排序
       ans = 0;
22
       for(int i = 1; i <= m; i++) {</pre>
23
          int fu = find(e[i].u), fv = find(e[i].v);
          if(fu == fv) continue;
          if(fu > fv) swap(fu, fv);
          f[fv] = fu;
27
          ans += e[i].w;
28
          e[i].flag = 1;
29
          G[ e[i].u ].push_back( make_pair(e[i].v, e[i].w) );
30
          G[ e[i].v ].push_back( make_pair(e[i].u, e[i].w) );
31
       return ans;
33
   }
```

4.6 严格次小生成树

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```
#include<bits/stdc++.h>
   #define int long long
   using namespace std;
   const int MAXN = (int)1e5 + 5;
   const int MAXM = (int)3e5 + 5;
   struct Node {
      int u, v, w, flag;
   }e[MAXM];
   int lca[MAXN][30], maxx1[MAXN][30], maxx2[MAXN][30], dep[MAXN], n, m, f[MAXN], ans, vis[MAXN];
   vector< pair<int, int> > G[MAXN];
   bool cmp(Node &a, Node &b) { return a.w < b.w; }</pre>
13
   int find(int x) { return ( x == f[x] ? x : (f[x] = find(f[x]) ) ); }
14
   void Kruskal() {
       for(int i = 1; i <= n; i++) f[i] = i;</pre>
       sort(e + 1, e + 1 + m, cmp);
18
19
       ans = 0;
20
       for(int i = 1; i <= m; i++) {
21
          int fu = find(e[i].u), fv = find(e[i].v);
22
          if(fu == fv) continue;
          if(fu > fv) swap(fu, fv);
          f[fv] = fu;
          ans += e[i].w;
26
          e[i].flag = 1;
          G[e[i].u].push\_back(make\_pair(e[i].v, e[i].w)), G[e[i].v].push\_back(make\_pair(e[i].u, e[i].w));
       }
   void DFS(int u, int ftr) {
       // maxx1维护最大值 maxx2维护次大值
33
       dep[u] = dep[ftr] + 1;
34
       vis[u] = 1;
35
       for(int i = 1; i <= 20; i++) {
          lca[u][i] = lca[ lca[u][i - 1] ][i - 1];
          \max x1[u][i] = \max(\max x1[u][i - 1], \max x1[lca[u][i - 1]][i - 1]);
          if(maxx1[u][i - 1] == maxx1[ lca[u][i - 1] ][i - 1])
             \max 2[u][i] = \max(\max 2[u][i - 1], \max 2[lca[u][i - 1]][i - 1]);
             \max 2[u][i] = \max(\min(\max x1[u][i-1], \max x1[lca[u][i-1]][i-1]), \max(\max x2[u][i-1], \max x2[u][i-1])
                   lca[u][i - 1] ][i - 1] );
43
44
       for(auto i : G[u]) {
45
          int v = i.first, w = i.second;
46
          if(vis[v]) continue;
47
          lca[v][0] = u; maxx1[v][0] = w; maxx2[v][0] = 0;
48
          DFS(v, u);
   }
51
   int LCA(int u, int v) {
       if(dep[u] < dep[v]) swap(u, v);</pre>
       for(int i = 20; i >= 0; i--) {
          if((1 << i) & (dep[u] - dep[v])) u = lca[u][i];</pre>
       for(int i = 20; i >= 0; i--) {
58
          if(lca[u][i] != lca[v][i]) {
59
             u = lca[u][i];
60
             v = lca[v][i];
       return u == v ? u : lca[u][0];
64
```

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```
65
66
   int Work(int u, int l, int w) {
       int max1 = 0, max2 = 0;
       for(int i = 20; i >= 0; i--) {
69
          if((1 << i) & (dep[u] - dep[l])) {</pre>
70
              max1 = max(max1, maxx1[u][i]);
71
              max2 = max(max2, maxx2[u][i]);
72
              u = lca[u][i];
74
75
       if(w - max1 == 0) return w - max2;
76
       return w - max1;
77
78
79
   signed main()
       scanf("%11d%11d", &n, &m);
82
       for(int i = 1; i <= m; i++) {
83
          int u, v, w;
84
          scanf("%11d%11d%11d", &u, &v, &w);
          e[i].u = u, e[i].v = v, e[i].w = w, e[i].flag = 0;
       Kruskal();
       DFS(1, 0);
89
90
       int d = 1e14;
91
       for(int i = 1; i <= m; i++) {
92
          if(e[i].flag) continue;
          int u = e[i].u, v = e[i].v, w = e[i].w;
          int l = LCA(u, v);
          d = min(d, Work(u, l, w));
96
          d = min(d, Work(v, l, w));
97
       printf("%11d", ans + d);
       return 0;
```

4.7 Dijkstra

```
#include<bits/stdc++.h>
   using namespace std;
   const int INF = 0x3f3f3f3f;
   const int MAXN = (int)1e5 + 5;
   const int MAXM = (int)2e5 + 5;
6
   struct Node {
       int u, w;
       bool operator < (const Node &a) const { return w > a.w; }
9
10
   };
11
   struct Edge {
       int v, w, nxt;
13
   int n, m, s, cnt = 0, head[MAXN], vis[MAXN], dist[MAXN];
14
   priority_queue<Node> Q;
15
16
   void Add(int u, int v, int w) {
18
       e[++cnt].v = v;
       e[cnt].w = w;
       e[cnt].nxt = head[u];
20
       head[u] = cnt;
21
22
  void Dijkstra(int s) {
```

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```
for(int i = 1; i <= n; i++) vis[i] = 0, dist[i] = INF;
25
       dist[s] = 0;
26
27
       Q.push((Node){s, 0});
       while(!Q.empty()) {
          Node now = Q.top();
29
          int u = now.u, w = now.w;
30
          Q.pop();
31
          if(vis[u]) continue;
          vis[u] = 1;
          for(int i = head[u]; i != -1; i = e[i].nxt) {
34
              if(dist[e[i].v] > dist[u] + e[i].w) {
35
                  dist[e[i].v] = dist[u] + e[i].w;
36
                 Q.push((Node){e[i].v, dist[e[i].v]});
37
              }
38
           }
39
40
       }
   }
42
   int main(){
43
       scanf("%d%d%d", &n, &m, &s);
44
       memset(head, -1, sizeof(head));
45
       for(int i = 1; i <= m; i++) {</pre>
           int u, v, w;
          scanf("%d%d%d", &u, &v, &w);
          Add(u, v, w);
49
50
       Dijkstra(s);
51
       for(int i = 1; i <= n; i++) printf("%d ", dist[i]);</pre>
52
       return 0;
   }
```

4.8 SPFA

```
int n; // 总点数
   int h[N], w[N], e[N], ne[N], idx; // 邻接表存边
   int dist[N]; // 存储每个点到1号点的最短距离
   bool st[N]; // 存储每个点是否在队列中
   // 求1号点到n号点的最短距离,如果无法到达返回-1
   int SPFA() {
      memset(dist, 0x3f, sizeof dist);
      queue<int> q;
10
      q.push(1);
      dist[1] = 0, st[1] = true;
12
      while(!q.empty()) {
         auto t = q.front(); q.pop();
13
         st[t] = false;
14
         for(int i = h[t]; i != -1; i = ne[i]) {
15
            int j = e[i];
16
            if(dist[j] > dist[t] + w[i]) {
               dist[j] = dist[t] + w[i];
               if(!st[j]) q.push(j), st[j] = true;
            }
20
         }
      if(dist[n] >= 0x3f3f3f3f) return -1;
      return dist[n];
```

4.9 SPFA 负环

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```
int h[N], w[N], e[N], ne[N], idx; // 邻接表存储所有边
   int dist[N], cnt[N]; // dist[x]存储1号点到x的最短距离, cnt[x]存储1到x的最短路中经过的点数
  bool st[N]; // 存储每个点是否在队列中
  // 如果存在负环,则返回true,否则返回false。
  bool SPFA() {
     // 不需要初始化dist数组
     // 原理:如果某条最短路径上有n个点 (除了自己),那么加上自己之后一共有n+1个点,由抽屉原理一定有两个点相同,所以存在环。
9
     queue<int> q;
11
     for (int i = 1; i \leftarrow n; i++) q.push(i), st[i] = true;
12
13
     while(!q.empty()) {
14
        auto t = q.front(); q.pop();
15
        st[t] = false;
16
        for(int i = h[t]; i != -1; i = ne[i]) {
17
           int j = e[i];
           if(dist[j] > dist[t] + w[i]) {
19
              dist[j] = dist[t] + w[i];
20
              cnt[j] = cnt[t] + 1;
21
              if(cnt[j] >= n) return true; // 如果从1号点到x的最短路中包含至少n个点(不包括自己),则说明存在环
22
              if(!st[j]) q.push(j), st[j] = true;
           }
        }
26
     return false;
27
```

4.10 Floyd

```
void Solve1() {
       for(int i = 1; i <= n; i++) {
          for(int j = 1; j <= n; j++) {</pre>
3
             if(i == j) dp[i][j] = 0;
              else dp[i][j] = INF;
5
          }
6
       for(int k = 1; k <= n; k++) {</pre>
          for(int i = 1; i <= n; i++) {
              for(int j = 1; j <= n; j++) {</pre>
10
                 dp[i][j] = min(dp[i][j], dp[i][k] + dp[k][j]);
11
              }
12
13
          }
14
15
16
17
    已知一个有向图中任意两点之间是否有连边,要求判断任意两点是否连通。
18
   Floyd实现传递闭包 bitset优化
19
   */
20
   void Solve2()
21
   {
       for(int k = 1; k <= n; k++) {
23
          for(int i = 1; i <= n; i++) {
24
             if(G[i][k]) G[i] |= G[k];
25
          }
       for(int k = 1; k <= n; k++) {
          for(int i = 1; i <= n; i++) {
             for(int j = 1; j <= n; j++) {</pre>
30
                 G[i][j] = G[i][k] & G[k][j];
31
32
          }
       }
```

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35 | }

4.11 差分约束

```
#include<bits/stdc++.h>
   using namespace std;
   const int MAXN = 5005;
   const int INF = 0x3f3f3f3f;
   struct Node { int v, w; };
   vector<Node> e[MAXN];
   int head[MAXN], vis[MAXN], cnt[MAXN], dis[MAXN], n, m, ecnt;
10
   queue<int> Q;
11
   bool SPFA(int s) {
13
       for(int i = 0; i \leftarrow n; i++) dis[i] = INF, vis[i] = cnt[i] = 0;
14
       dis[s] = 0, vis[s] = 1; Q.push(s);
15
       while(!Q.empty()) {
16
          int u = Q.front(); Q.pop();
17
          vis[u] = 0;
          for(auto i : e[u]) {
              int v = i.v, w = i.w;
              if(dis[v] > dis[u] + w) {
                 dis[v] = dis[u] + w;
                 cnt[v] = cnt[u] + 1;
                 if(cnt[v] >= n + 1) return 0;
                 if(!vis[v]) Q.push(v), vis[v] = 1;
              }
26
           }
27
28
       return 1;
29
30
31
   signed main()
       scanf("%d%d", &n, &m);
34
       ecnt = 0;
35
       //x_u - x_v <= w
36
       for(int i = 1; i <= m; i++) {
           int u, v, w; scanf("%d%d%d", &u, &v, &w);
          e[u].push_back((Node){v, w});
40
       for(int i = 1; i <= n; i++) Add(0, i, 0);</pre>
41
       if(!SPFA(0)) printf("NO\n");
42
       else for(int i = 1; i <= n; i++) printf("%d ", dis[i]);</pre>
43
       return 0;
44
   }
```

4.12 欧拉路径

```
#pragma GCC optimize(2)
#include<bits/stdc++.h>
using namespace std;

const int MAXN = (int)1e5 + 5;
const int MAXM = (int)2e5 + 5;

stack<int> S;
vector<int> G[MAXN];
int del[MAXN], in[MAXN], out[MAXN], n, m;
```

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```
11
   void DFS(int u) {
12
       for(int i = del[u]; i < G[u].size(); i = del[u]) del[u] = i + 1, DFS(G[u][i]);</pre>
   }
15
16
   signed main()
17
       ios::sync_with_stdio(false); cin.tie(0);
       cin >> n >> m;
       for(int i = 1; i <= m; i++) {
           int u, v; cin >> u >> v;
22
          G[u].push_back(v); out[u]++; in[v]++;
23
24
       for(int i = 1; i <= n; i++) sort(G[i].begin(), G[i].end());</pre>
25
       int s = 1, cnt0 = 0, cnt1 = 0, flag = 1;
       for(int i = 1; i <= n; i++) {
           if(in[i] != out[i]) flag = 0;
           if(out[i] == in[i] + 1) s = i, cnt1++;
           if(in[i] == out[i] + 1) cnt0++;
30
       if(!flag && !(cnt0 == cnt1 && cnt0 == 1) ) return cout << "No", 0;</pre>
       DFS(s);
       while(!S.empty()) cout << S.top() << " ", S.pop();</pre>
       return 0;
35
   }
36
```

4.13 染色法判别二分图

```
//0(n + m)
   int n; // n表示点数
   int h[N], e[M], ne[M], idx; // 邻接表存储图
   int color[N]; // 表示每个点的颜色, -1表示未染色, 0表示白色, 1表示黑色
   // 参数: u表示当前节点, c表示当前点的颜色
   bool dfs(int u, int c) {
      color[u] = c;
      for (int i = h[u]; i != -1; i = ne[i]) {
         int j = e[i];
         if (color[j] == -1) {
10
            if (!dfs(j, !c)) return false;
         } else if (color[j] == c) return false;
      return true;
14
15
   bool check() {
16
      memset(color, -1, sizeof color);
17
      bool flag = true;
18
      for (int i = 1; i <= n; i ++ ) {
19
         if (color[i] == -1) if (!dfs(i, 0)) {
            flag = false;
            break;
         }
      return flag;
   }
```

4.14 匈牙利算法

```
bool Match(int x) {
    for(int i = 1; i <= m; i++) {
        if(G[x][i] && !vis[i]) {
            vis[i] = 1;
        }</pre>
```

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```
if(!a[i] || Match(a[i])) {
                 a[i] = x;
6
                 return true;
              }
          }
10
       return false;
11
   }
12
   int main() {
14
       scanf("%d%d%d", &n, &m, &e);
15
       for(int i = 1; i <= e; i++) {
16
          int u, v; scanf("%d%d", &u, &v);
17
          G[u][v] = 1;
18
19
       }
       int ans = 0;
       for(int i = 1; i <= n; i++) {
          memset(vis, 0, sizeof(vis));
22
          if(Match(i)) ans++;
23
24
       printf("%d", ans);
25
26
       return 0;
   }
```

4.15 Tarjan

```
const int N;
   int dfn[N], low[N], s[N], vis[N], color[N], top = 0, sum = 0, dep = 0;
   int n, m;
   vector<int> G[N];
   // 求强连通分量 复杂度O(E+V)
   void Tarjan(int u) {
       dfn[u] = low[u] = ++dep;
       vis[u] = 1;
10
       s[++top] = u;
11
       for(int i = 0; i < G[u].size(); i++) {</pre>
13
          int v = G[u][i];
14
          if(!dfn[v]) Tarjan(v), low[u] = min(low[u], low[v]);
15
          else if(vis[v]) low[u] = min(low[u], low[v]);
16
17
       if(dfn[u] == low[u]) {
19
20
          color[u] = ++sum;
          vis[u] = 0;
21
          while(s[top] != u) {
22
             color[ s[top] ] = sum;
23
24
             vis[s[top]] = 0;
             top--;
          }
          top--;
27
       }
28
   }
29
30
   // 求割点
   vector<int> cut; // 存储所有割点
   void Tarjan(int u, bool root = true) {
       int tot = 0;
34
       low[u] = dfn[u] = ++dep;
35
       for(auto v : G[u]) {
36
          if(!dfn[v]) {
37
             Tarjan(v, false);
```

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```
low[u] = min(low[u], low[v]);
39
             tot += (low[v] >= dfn[u]); // 统计满足low[v] >= dfn[u]的子节点数目
40
41
          } else low[u] = min(low[u], dfn[v]);
       if (tot > root) // 如果是根, tot需要大于1; 否则只需大于0
43
          cut.push_back(u);
44
   }
45
46
   // 求割桥
   vector<pair<int, int>> bridges; // 存割桥
49
   void Tarjan(int u) {
       low[u] = dfn[u] = ++dep;
50
       for(auto v : G[u]) {
51
          if(!dfn[v]) {
52
             fa[v] = u; // 记录父节点
             Tarjan(v);
             low[u] = min(low[u], low[v]);
             if(low[v] > dfn[u]) bridges.emplace_back(u, v);
56
          } else if (fa[u] != v) // 排除父节点
57
             low[u] = min(low[u], dfn[v]);
59
       }
60
   void Solve() {
      for(int i = 1; i <= n; i++) if(!dfn[i]) {</pre>
63
          Tarjan(i);
64
          Tarjan(i, true);
65
66
       }
   }
```

4.16 Dinic 最大流

```
#include<bits/stdc++.h>
   #define int long long
   using namespace std;
   const int MAXN = 205;
   const int MAXM = 5005;
   struct Edge {
9
       int v, w, nxt;
   }e[MAXM << 1];
11
12
   int n, m, s, t, cnt = 1, dep[MAXN], head[MAXN];
13
14
   void Add(int u, int v, int w) {
15
       e[++cnt].v = v;
16
       e[cnt].w = w;
       e[cnt].nxt = head[u];
18
       head[u] = cnt;
19
   }
20
   bool BFS() {
       for(int i = 1; i <= n + 1; i++) dep[i] = 0;</pre>
       dep[s] = 1;
       queue<int> Q;
       Q.push(s);
26
       while(!Q.empty()) {
27
          int u = Q.front(); Q.pop();
28
          for(int i = head[u]; i != -1; i = e[i].nxt) {
              int v = e[i].v;
              if(dep[v] == 0 \&\& e[i].w > 0) {
                 dep[v] = dep[u] + 1;
```

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```
Q.push(v);
33
             }
34
          }
35
       return (bool)dep[t];
37
   }
38
39
   int DFS(int u, int in) {
40
       if(u == t) return in;
       int out = 0;
       for(int i = head[u]; i != -1 && in > 0; i = e[i].nxt) {
43
          int v = e[i].v;
44
          if(dep[v] == dep[u] + 1 && e[i].w > 0) {
45
              int res = DFS(v, min(e[i].w, in));
46
              e[i].w -= res;
47
             e[i ^ 1].w += res; // 反向边(残量网络)
              in -= res;
              out += res;
50
          }
51
52
       if(out == 0) dep[u] = 0;
53
       return out;
54
   signed main()
57
58
       memset(head, -1, sizeof(head));
59
       scanf("%11d%11d%11d%11d", &n, &m, &s, &t);
60
       for(int i = 1; i <= m; i++) {
61
          int u, v, w; scanf("%11d%11d%11d", &u, &v, &w);
          Add(u, v, w); Add(v, u, 0);
64
       int ans = 0;
65
       while(BFS()) { ans += DFS(s, 1e18); }
       printf("%11d", ans);
67
       return 0;
```

4.17 KM

```
#include<bits/stdc++.h>
   #define int long long
   using namespace std;
   const int MAXN = 505;
   const int INF = (int)1e16;
   int n, m;
   int G[MAXN][MAXN];
   int lmatch[MAXN], rmatch[MAXN];
   int pre[MAXN];
   int lexpect[MAXN], rexpect[MAXN];
   int lvis[MAXN], rvis[MAXN];
   int slack[MAXN];
15
   queue<int> Q;
16
   void aug(int v) {
18
       int temp;
       while(v) {
20
          temp = lmatch[ pre[v] ];
21
          lmatch[ pre[v] ] = v;
22
          rmatch[v] = pre[v];
          v = temp;
```

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```
}
25
    }
26
27
    void BFS(int s) {
       for(int i = 1; i <= n; i++) lvis[i] = rvis[i] = 0, slack[i] = INF;</pre>
29
30
       while(!Q.empty()) Q.pop();
31
       Q.push(s);
32
       while(1) {
34
           while(!Q.empty()) {
35
               int u = Q.front(); Q.pop();
36
               lvis[u] = 1;
37
              for(int v = 1; v <= n; v++) {</pre>
38
                  if(!rvis[v]) {
                      int gap = lexpect[u] + rexpect[v] - G[u][v];
                      if(slack[v] > gap) {
                         slack[v] = gap;
42
                         pre[v] = u;
43
                         if(slack[v] == 0) {
                             rvis[v] = 0;
                             if(!rmatch[v]) { aug(v); return ; }
                             else Q.push(rmatch[v]);
                         }
                      }
49
                  }
50
              }
51
           }
52
           int d = INF;
           for(int i = 1; i <= n; i++)
               if(!rvis[i]) d = min(d, slack[i]);
           for(int i = 1; i <= n; i++) {</pre>
              if(lvis[i]) lexpect[i] -= d;
               if(rvis[i]) rexpect[i] += d;
               else slack[i] -= d;
62
63
64
           for(int i = 1; i <= n; i++) {</pre>
              if(!rvis[i]) {
                  if(slack[i] == 0) {
                      rvis[i] = 1;
                      if(!rmatch[i]) { aug(i); return ; }
69
                      else Q.push(rmatch[i]);
70
                  }
71
              }
72
           }
       }
75
76
    int KM() {
77
       for(int i = 1; i <= n; i++) lmatch[i] = rmatch[i] = lexpect[i] = rexpect[i] = 0;</pre>
78
79
       for(int i = 1; i <= n; i++) {</pre>
           lexpect[i] = G[i][1];
           for(int j = 2; j <= n; j++) lexpect[i] = max(lexpect[i], G[i][j]);</pre>
82
       for(int i = 1; i <= n; i++) BFS(i);</pre>
       int res = 0;
       for(int i = 1; i <= n; i++)</pre>
88
           if(rmatch[i]) res += G[rmatch[i]][i];
89
```

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```
90
        return res;
91
    }
    signed main()
94
95
        scanf("%11d%11d", &n, &m);
96
        for(int i = 1; i <= n; i++) for(int j = 1; j <= n; j++) G[i][j] = -INF;</pre>
97
        for(int i = 1; i <= m; i++) {</pre>
            int u, v, w; scanf("%11d%11d%11d", &u, &v, &w);
           G[u][v] = w;
100
101
        printf("%lld\n", KM());
102
        for(int i = 1; i <= n; i++) printf("%1ld ", rmatch[i]);</pre>
103
        return 0;
104
    }
```

4.18 轻重链剖分

```
#include<bits/stdc++.h>
   #define int long long
2
   using namespace std;
   const int MAXN = (int)1e5 + 5;
   struct Node {
       int v, nxt;
   }e[MAXN << 1];
11
   int sum[MAXN << 2], tag[MAXN << 2];</pre>
   int n, m, r, MOD, cnt, e_cnt, v[MAXN];
13
   int head[MAXN], vis[MAXN], fa[MAXN], sze[MAXN], son[MAXN], dep[MAXN];
14
   int idx[MAXN], a[MAXN], top[MAXN];
15
   int ls(int x) { return x << 1; }</pre>
   int rs(int x) { return x << 1 | 1; }</pre>
18
19
   void F(int 1, int r, int p, int k) {
20
       tag[p] += k; tag[p] %= MOD;
21
       sum[p] += (r - l + 1) * k; sum[p] %= MOD;
22
24
25
   void PushUp(int p) {
       sum[p] = sum[ls(p)] + sum[rs(p)];
26
       sum[p] %= MOD;
27
28
   void PushDown(int 1, int r, int p) {
       int mid = (1 + r) >> 1;
       F(1, mid, ls(p), tag[p]);
32
       F(mid + 1, r, rs(p), tag[p]);
33
       PushUp(p);
       tag[p] = 0;
35
   }
   void Build(int 1, int r, int p) {
       if(1 == r) {
39
          sum[p] = a[1];
40
          tag[p] = 0;
41
          return;
42
       int mid = (1 + r) >> 1;
       Build(l, mid, ls(p));
```

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```
Build(mid + 1, r, rs(p));
46
       PushUp(p);
47
    }
48
    void SegmentTreeModify(int nl, int nr, int l, int r, int p, int k) {
50
        if(nl <= 1 && nr >= r) {
51
           tag[p] += k; tag[p] %= MOD;
52
           sum[p] += (r - 1 + 1) * k; sum[p] %= MOD;
           return ;
55
       PushDown(l, r, p);
56
        int mid = (1 + r) >> 1;
57
       if(nl <= mid) SegmentTreeModify(nl, nr, l, mid, ls(p), k);</pre>
58
        if(nr > mid) SegmentTreeModify(nl, nr, mid + 1, r, rs(p), k);
59
       PushUp(p);
60
61
       return ;
    }
63
    int SegmentTreeQuery(int nl, int nr, int l, int r, int p) {
64
        if(n1 <= 1 && nr >= r) {
65
           return sum[p];
66
        PushDown(l, r, p);
        int res = 0;
        int mid = (1 + r) >> 1;
70
        if(nl <= mid) res += SegmentTreeQuery(nl, nr, l, mid, ls(p)), res %= MOD;</pre>
71
        if(nr > mid) res += SegmentTreeQuery(nl, nr, mid + 1, r, rs(p)), res %= MOD;
72
       return res;
73
    }
74
75
    void Add(int u, int v) {
       e[++e cnt].v = v;
77
        e[e_cnt].nxt = head[u];
78
       head[u] = e_cnt;
79
80
81
    void DFS1(int u, int ftr) {
        fa[u] = ftr;
83
        sze[u] = 1;
84
        dep[u] = dep[ftr] + 1;
85
       vis[u] = 1;
86
        int maxsize = -1;
87
        for(int i = head[u]; i != -1; i = e[i].nxt) {
           int v = e[i].v;
           if(vis[v]) continue;
90
           DFS1(v, u);
91
           sze[u] += sze[v];
92
           if(sze[v] > maxsize) {
               son[u] = v;
               maxsize = sze[v];
           }
96
        }
97
    }
98
99
    void DFS2(int u, int top_u) {
100
       idx[u] = ++cnt;
        a[cnt] = v[u];
       top[u] = top_u;
103
        if(son[u] == 0) return ;
104
       DFS2(son[u], top_u);
105
        for(int i = head[u]; i != -1; i = e[i].nxt) {
106
           int v = e[i].v;
           if(v == fa[u] || v == son[u]) continue;
108
           DFS2(v, v);
109
110
```

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```
}
111
112
113
    void Modify(int u, int v, int w) {
       while(top[u] != top[v]) {
114
           if(dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
115
           SegmentTreeModify(idx[top[u]], idx[u], 1, n, 1, w);
116
           u = fa[top[u]];
117
118
        if(idx[u] > idx[v]) swap(u, v);
        SegmentTreeModify(idx[u], idx[v], 1, n, 1, w);
120
121
122
    int Query(int u, int v) {
123
        int res = 0;
124
       while(top[u] != top[v]) {
125
           if(dep[top[u]] < dep[top[v]]) swap(u, v);</pre>
           res += SegmentTreeQuery(idx[top[u]], idx[u], 1, n, 1);
           res %= MOD;
128
           u = fa[top[u]];
129
130
       if(idx[u] > idx[v]) swap(u, v);
131
       res += SegmentTreeQuery(idx[u], idx[v], 1, n, 1);
        return res % MOD;
135
    signed main()
136
137
        scanf("%11d%11d%11d%11d", &n, &m, &r, &MOD);
138
        cnt = 0, e_cnt = 0; memset(head, -1, sizeof(head));
139
        for(int i = 1; i <= n; i++) scanf("%11d", &v[i]);</pre>
        for(int i = 1; i < n; i++) {
           int u, v; scanf("%11d%11d", &u, &v);
142
           Add(u, v); Add(v, u);
143
144
       DFS1(r, 0); DFS2(r, r); Build(1, n, 1);
145
       while(m--) {
           int opt, x, y, z; scanf("%11d", &opt);
           if(opt == 1) {
148
               scanf("%11d%11d%11d", &x, &y, &z);
149
              Modify(x, y, z);
150
           } else if(opt == 2) {
151
               scanf("%11d%11d", &x, &y);
152
               printf("%lld\n", Query(x, y));
           } else if(opt == 3) {
               scanf("%11d%11d", &x, &z);
155
               SegmentTreeModify(idx[x], idx[x] + sze[x] - 1, 1, n, 1, z);
156
           } else {
               scanf("%11d", &x);
               printf("%lld\n", SegmentTreeQuery(idx[x], idx[x] + sze[x] - 1, 1, n, 1) % MOD);
           }
161
       return 0;
162
    }
163
164
165
    input:
    5 5 2 24
    7 3 7 8 0
168
    1 2
169
    1 5
170
    3 1
171
    4 1
    3 4 2
    3 2
        2
174
    4 5
175
```

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4.19 树上启发式合并

```
// 轻重链剖分
   void DFS(int u, int fa) {
       dfn[u] = ++tot, node[tot] = u, sze[u] = 1;
       for(auto v : G[u]) {
          if(v == fa) continue;
          DFS(v, u);
          if(sze[v] > sze[ son[u] ]) son[u] = v;
          sze[u] += sze[v];
10
11
   void add(int pos) { }
12
   void del(int pos) { }
   int getAns() { }
15
   void DSU(int u, int fa, bool st) {
16
       for(auto v : G[u]) if(v != fa && v != son[u]) DSU(v, u, 0);
17
       if(son[u]) DSU(son[u], u, 1);
       for(auto v : G[u]) {
          if(v == fa || v == son[u]) continue;
          for(int i = dfn[v]; i < dfn[v] + sze[v]; i++) add(node[i]);</pre>
       add(u);
23
       ans[u] = getAns();
24
       if(!st) for(int i = dfn[u]; i < dfn[u] + sze[u]; i++) del(node[i]);</pre>
25
   }
26
   void Solve() {
       DFS(1, 0);
       DSU(1, 0, 0);
30
   }
```

4.20 点分治

```
#include<bits/stdc++.h>
   #define pir make_pair
   #define pii pair<int, int>
   #define fi first
   #define se second
   using namespace std;
   const int MAXN = 1e4 + 5;
   const int MAXV = 1e7 + 5;
   const int MOD = 1e9 + 7;
   int n, m, ans[MAXN], val[MAXN], sze[MAXN], vis[MAXN], centroid;
11
   int Map[MAXN], Cur[MAXN], MapVal[MAXV];
12
   vector<pii> G[MAXN];
14
   void GetCentroid(int u, int fa, int n) {
15
      sze[u] = 1;
16
       int maxx = 0;
       for(auto i : G[u]) {
          int v = i.fi, w = i.se;
          if(v == fa || vis[v]) continue;
```

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```
GetCentroid(v, u, n);
21
          if(centroid != -1) return ;
22
          maxx = max(maxx, sze[v]);
          sze[u] += sze[v];
       maxx = max(maxx, n - sze[u]);
26
       if(maxx <= n / 2) centroid = u, sze[fa] = n - sze[u];</pre>
27
28
   void Calc(int u, int fa, int len) {
30
       if(len > 1e7) return ;
       Cur[ ++Cur[0] ] = len;
32
       for(auto i : G[u]) {
33
           int v = i.fi, w = i.se;
34
           if(vis[v] || v == fa) continue;
35
          Calc(v, u, len + w);
   }
38
39
   void Calc(int u) {
40
       Map[Map[0] = 1] = 0, MapVal[0] = 1;
       for(auto i : G[u]) {
           int v = i.fi, w = i.se;
           if(vis[v]) continue;
          Cur[0] = 0; Calc(v, u, w);
45
          for(int j = 1; j <= Cur[0]; j++)</pre>
46
              for(int k = 1; k <= m; k++)</pre>
47
                 if(val[k] - Cur[j] >= 0)
                     ans[k] |= MapVal[ val[k] - Cur[j] ];
          for(int j = 1; j <= Cur[0]; j++) MapVal[ Cur[j] ] = 1, Map[ ++Map[0] ] = Cur[j];</pre>
       for(int i = 1; i <= Map[0]; i++) MapVal[ Map[i] ] = 0;</pre>
52
   }
53
   void Solve(int u) {
       vis[u] = 1; Calc(u);
       for(auto i : G[u]) {
           int v = i.fi, w = i.se;
58
           if(vis[v]) continue;
59
           centroid = -1; GetCentroid(v, 0, sze[v]);
60
          Solve(centroid);
61
   signed main()
65
66
       ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
       cin >> n >> m;
       for(int i = 1; i < n; i++) {
           int u, v, w; cin >> u >> v >> w;
          G[u].push_back( pir(v, w) ), G[v].push_back( pir(u, w) );
71
72
       for(int i = 1; i <= m; i++) cin >> val[i];
73
       centroid = -1;
74
       GetCentroid(1, 0, n);
       Solve(centroid);
       for(int i = 1; i <= m; i++) cout << (ans[i] ? "AYE" : "NAY") << "\n";</pre>
       return 0;
78
   }
```

4.21 虚树

```
//#pragma GCC optimize(2)
#include<bits/stdc++.h>
```

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```
#define int long long
   #define pir make_pair
   #define pii pair<int, int>
   #define fi first
   #define se second
   using namespace std;
   const int MAXN = 1e6 + 5;
   const int INF = 0x3f3f3f3f;
10
   int n, m, cnt, dfn[MAXN], dep[MAXN], top, s[MAXN], vis[MAXN], lca[MAXN][25];
12
13
   vector<pii> VG[MAXN], vec;
   vector<int> G[MAXN];
14
   unordered_map<int, int> Map;
15
16
   void DFS(int u, int fa) {
17
       dfn[u] = ++cnt; dep[u] = dep[fa] + 1;
       lca[u][0] = fa;
       for(int i = 1; i <= 20; i++) lca[u][i] = lca[ lca[u][i - 1] ][i - 1];
20
       for(auto i : G[u]) if(i != fa) DFS(i, u);
   }
22
   int LCA(int u, int v) {
       if(dep[u] > dep[v]) swap(u, v);
       for(int i = 20; i >= 0; i--) if((1 << i) & (dep[v] - dep[u])) v = lca[v][i];</pre>
       for(int i = 20; i >= 0; i--) if(lca[u][i] != lca[v][i]) u = lca[u][i], v = lca[v][i];
27
       return u == v ? u : lca[u][0];
28
29
   int Dist(int u, int v) {
31
   void Add(int u, int v) {
35
       int w = Dist(u, v);
       VG[u].push_back( pir(v, w) );
       VG[v].push_back( pir(u, w) );
       if(!Map[u]) Map[u] = ++cnt;
       if(!Map[v]) Map[v] = ++cnt;
40
41
42
   void VirtualTreeInsert(int u) {
43
       if(!top) return void(s[++top] = u);
44
       int 1 = LCA(s[top], u);
       while(top > 1 && dep[l] < dep[ s[top - 1] ]) Add(s[top - 1], s[top]), --top;</pre>
       if(dep[1] < dep[ s[top] ]) Add(1, s[top]), --top;</pre>
       if(!top || s[top] != 1) s[++top] = 1;
       if(s[top] != u) s[++top] = u;
49
   void BuildVirtualTree(vector<pii> vec, int k) {
       sort(vec.begin(), vec.end());
53
       for(auto i : Map) VG[i.fi].clear();
       top = cnt = 0; Map.clear();
55
       s[++top] = 1, Map[1] = ++cnt;
56
       for(auto i : vec) VirtualTreeInsert(i.se);
57
       for(int i = 1; i < top; i++) Add(s[i], s[i + 1]);</pre>
   }
60
   signed main()
61
62
       ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
       cin >> n;
       for(int i = 1; i < n; i++) {
          int u, v; cin >> u >> v;
66
          G[u].push_back(v), G[v].push_back(u);
67
```

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```
DFS(1, 0); cnt = 0;
69
       cin >> m;
       while(m--) {
          int k; cin >> k; vec.clear();
72
          for(int i = 1; i <= k; i++) {
73
              int cur; cin >> cur;
              vec.push_back( pir(dfn[cur], cur) );
          BuildVirtualTree(vec, k);
78
       return 0;
79
   }
80
```

4.22 树哈希

```
// 复杂度O(nlogn), n为节点数量
   #include<bits/stdc++.h>
   using namespace std;
   const int MAXN = 1e6 + 5;
   vector<int> g[MAXN];
   unordered_map<int, int> num;
   map<vector<int>, int> Hash;
   int hash_cnt;
9
10
   int dfs(int u, int f) {
^{11}
      vector<int> vec;
      for(auto &v : g[u]) {
          if(v == f) continue;
          vec.push_back( dfs(v, u) );
16
      sort(vec.begin(), vec.end() );
      if(Hash[vec] == 0) Hash[vec] = ++hash_cnt;
      return Hash[vec];
19
   }
20
21
   void init() {
22
      hash\_cnt = 0;
23
      Hash.clear();
24
   }
   signed main() {
      //ios::sync_with_stdio(false); cin.tie(0); cout.tie(0);
      init();
      int n; cin >> n;
      for(int j = 1; j <= n; j++) {</pre>
          int f; cin >> f;
          if(f == 0) continue;
33
          g[f].push_back(j); g[j].push_back(f);
34
35
36
      树哈希返回对应子树的哈希值,
      如需要比较两棵树的哈希值,可以通过求重心的方式固定根节点;
       若重心有两个,分别固定根节点求哈希值即可。
      int h = dfs(centroid, 0);
      if(num[h] == 0) num[h] = i;
      cout << num[h] << "\n";</pre>
      for(int j = 0; j <= n; j++) g[j].clear();</pre>
46
   }
47
```

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5 数学

5.1 快速幂

```
// 快速乘 适用正数
   int qmul(int a, int b, int MOD) {
       int ans = 0;
       while(b > 0) {
          if(b \& 1) ans = (ans + a) % MOD;
          b >>= 1;
          a = (a << 1) \% MOD;
       return ans;
10
   // 快速幂
12
   int qpow(int a, int p, int MOD) {
13
       int ans = 1;
14
       while(p > 0) {
15
          if(p & 1) ans = ans * a % MOD;
          p >>= 1;
          a = a * a % MOD;
19
       return ans;
20
   }
21
22
   // 矩阵快速幂
   struct Martix {
       long long c[MAXN][MAXN];
25
   }A, E;
26
27
   Martix operator * (const Martix &a, const Martix &b) {
28
       Martix Ans;
29
       for(int i = 1; i \leftarrow n; i++) for(int j = 1; j \leftarrow n; j++) Ans.c[i][j] = 0;
       for(int k = 1; k <= n; k++) {</pre>
          for(int i = 1; i <= n; i++) {
              for(int j = 1; j \le n; j++) Ans.c[i][j] = (Ans.c[i][j] + a.c[i][k] * b.c[k][j] % MOD) % MOD;
33
35
       return Ans;
   Martix MartixQuickPow(Martix A, int k) {
39
       Martix ans, now;
40
       for(int i = 1; i <= n; i++) {
41
          for(int j = 1; j <= n; j++) ans.c[i][j] = 0;</pre>
42
          ans.c[i][i] = 1;
       for(int i = 1; i <= n; i++) for(int j = 1; j <= n; j++) now.c[i][j] = A.c[i][j];
       while(k > 0) {
46
          if(1 \& k) ans = ans * now;
47
          now = now * now;
          k = k \gg 1;
       return ans;
51
```

5.2 欧拉函数

```
int phi(int n) {
   int res = n;
   for(int i = 2; i * i <= n; i++) {
      if(n % i == 0) res = res / i * (i - 1);
      while(n % i == 0) n /= i;
}</pre>
```

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```
if(n > 1) res = res / n * (n - 1);
7
      return res;
   }
10
   int phi[MAXN];
11
   void init(int n) {
12
      // 复杂度 O(nloglogn)
13
      for(int i = 1; i <= n; i++) phi[i] = i; // 除1外没有数的欧拉函数是本身, 所以如果phi[i] = i则说明未被筛到
      for(int i = 2; i <= n; i++) {
15
          if(phi[i] == i) { // 未被筛到
16
             for(int j = i; j <= n; j += i) phi[j] = phi[j] / i * (i - 1); // 所有含有该因子的数都进行一次操作
17
          }
18
      }
19
      // 复杂度 O(n)
      phi[1] = 1;
      for(int i = 2; i <= n; i++) {
23
          if(!isnp[i]) primes.push_back(i), phi[i] = i - 1;
          for(int p : primes) {
25
             if(p * i > n) break;
             isnp[p * i] = 1;
             if(i % p == 0) {
                phi[p * i] = phi[i] * p;
                break;
30
             } else {
31
                phi[p * i] = phi[p] * phi[i];
32
             }
          }
      }
   }
37
   int EulerPow(int a, string b, int MOD) {
      int phiMOD = phi(MOD), power = 0, flag = 0;
      for(auto i : b) {
          power = power * 10 + i - '0';
          if(power >= phiMOD) flag = 1;
          power %= phiMOD;
43
44
      if(flag) power += phiMOD;
45
      return qpow(a, power, MOD);
46
   }
```

5.3 线性筛

```
#include<bits/stdc++.h>
2
   using namespace std;
   const int MAXN = 2e7 + 5;
   int n, m, isprime[MAXN], prime[MAXN], cnt;
   void Init(int n = 2e7) {
      cnt = 0;
       isprime[1] = 1;
11
       for(int i = 2; i <= n; i++) {
          if(!isprime[i]) prime[++cnt] = i;
13
          for(int j = 1; j <= cnt && i * prime[j] <= n; j++) {</pre>
             isprime[i * prime[j]] = 1;
15
             if(i % prime[j] == 0) break;
       //for(int i = 1; i <= cnt; i++) cout << prime[i] << " "; cout << endl;
19
```

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5.4 整除分块

```
for(int l = 1, r; l <= n; l = r + 1) r = n / (n / l); // [l, r]为当前整除分块 区间内每个 n / i 相同
```

5.5 组合数处理

```
int inv(int x, int MOD) { // 求逆元
       return qpow(x, MOD - 2, MOD);
   }
3
   void Init() {
       fact[0] = 1;
       for(int i = 1; i <= n; i++) fact[i] = fact[i - 1] * i % MOD;</pre>
       inv[n] = inv(fact[n], MOD);
9
       for(int i = n - 1; i >= 0; i--) inv[i] = inv[i + 1] * (i + 1) % MOD;
10
   }
11
12
   int C(int n, int m) {
13
       if(m > n) return 0;
       return fact[n] * inv[m] % MOD * inv[n - m] % MOD;
15
   }
16
17
   int Lucas(int n, int m) {
       if(m == 0) return 1;
19
       return C(n % MOD, m % MOD) * Lucas(n / MOD, m / MOD) % MOD;
   }
```

6 其他算法

6.1 二分 & 三分

```
const double eps = 1e-10;
   double 1, r, a[MAXN];
   int n;
   double F(double x) {}
   void Solve() {
      // 单调递增序列a中查找 >=x 的数中最小的一个
      while(1 < r) {
         int mid = l + r \gg 1;
         if(a[mid] >= x) r = mid;
10
         else l = mid + 1;
11
      }
      // 单调递增序列a中查找 <=x 的数中最大的一个
      while(l < r) {
         int mid = 1 + r + 1 >> 1;
         if(a[mid] <= x) 1 = mid;</pre>
         else r = mid - 1;
```

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```
20
21
   void Solve() {
      // 浮点数三分 求单峰函数的极大值
      while(r - 1 > eps) {
         double lmid = 1.0 * 1 + 1.0 * (r - 1) / 3.0;
         double rmid = 1.0 * r - 1.0 * (r - 1) / 3.0;
         if(F(lmid) > F(rmid)) r = rmid;
         else 1 = lmid;
30
      // 整数三分 注意特判边界等各种情况 较毒瘤
31
      while(1 < r) {</pre>
32
         int lmid = 1 + (r - 1) / 3;
33
         int rmid = r - (r - 1) / 3;
34
         if(F(lmid) > F(rmid)) r = rmid;
         else l = lmid;
37
      }
   }
```

6.2 前缀和 & 差分

```
// 二维前缀和 & 二维差分
   sum[i][j] = a[i][j] + sum[i - 1][j] + sum[i][j - 1] - sum[i - 1][j - 1];
   d[i][j] = a[i][j] - a[i - 1][j] - a[i][j - 1] + a[i - 1][j - 1];
   int QuerySum() {
      // 以(x1, y1)为左上角,(x2, y2)为右下角的子矩阵的和为:
      return sum[x2][y2] - sum[x1 - 1][y2] - sum[x2][y1 - 1] + sum[x1 - 1][y1 - 1];
   void ModifyD() {
      // 给以(x1, y1)为左上角,(x2, y2)为右下角的子矩阵中的所有元素加上c:
10
      d[x1][y1] += c, d[x2 + 1][y1] -= c, d[x1][y2 + 1] -= c, d[x2 + 1][y2 + 1] += c;
11
   }
12
   树上差分
15
   注意对差分数组进行前缀和操作,得到原数组时顺序是 自叶子节点开始 到 根结束!
16
17
   void SolveNode(int u, int v, int t) { //u--v之间的点全部加上c
18
      val[u] += t, val[v] += t;
19
      int lca = LCA(u, v);
      val[lca] -= t, val[fa[lca]] -= t;
21
22
23
   void SolveEdge(int u, int v, int t) { //u--v之间的边全部加上c
24
      val[u] += t, val[v] += t;
25
      val[LCA(u, v)] -= 2 * t;
26
   }
```

6.3 离散化

```
// a[i] 为初始数组,下标范围为 [1, n]
// len 为离散化后数组的有效长度
std::sort(a + 1, a + 1 + n);
len = std::unique(a + 1, a + n + 1) - a - 1;
// 离散化整个数组的同时求出离散化后本质不同数的个数。
std::lower_bound(a + 1, a + len + 1, x) - a; // 查询 x 离散化后对应的编号
```

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```
void QuickSort(int l, int r) { //快排
       int mid = a[(1 + r) / 2];
       int i = 1, j = r;
       do {
          while(a[i] < mid) i++;</pre>
          while(a[j] > mid) j--;
          if(i <= j) swap(a[i], a[j]), i++, j--;</pre>
       } while(i <= j);</pre>
       if(j > 1) QuickSort(l, j);
9
       if(i < r) QuickSort(i, r);</pre>
10
   }
12
   void MergeSort(int 1, int r) { //归并排序
13
       if (1 >= r) return;
14
       int mid = (1 + r) >> 1;
15
       MergeSort(1, mid);
       MergeSort(mid + 1, r);
       int k = 0, i = 1, j = mid + 1;
18
       while(i <= mid && j <= r) {
19
           if(a[i] <= a[j]) tmp[k++] = a[i++];</pre>
20
          else tmp[k++] = a[j++];
21
22
       while(i <= mid) tmp[k++] = a[i++];
       while(j \le r) tmp[k++] = a[j++];
       for(i = 1, j = 0; i <= r; i++, j++) a[i] = tmp[j];
25
   }
```

6.5 高精度运算

```
// C = A + B, A >= 0, B >= 0
    vector<int> add(vector<int> &A, vector<int> &B) {
       if (A.size() < B.size()) return add(B, A);</pre>
       vector<int> C;
       int t = 0;
       for (int i = 0; i < A.size(); i ++ ) {</pre>
          t += A[i];
           if (i < B.size()) t += B[i];</pre>
          C.push_back(t % 10);
          t /= 10;
10
11
       if (t) C.push_back(t);
12
       return C;
13
   }
15
   // C = A - B, Th3yA >= B, A >= 0, B >= 0
16
    vector<int> sub(vector<int> &A, vector<int> &B) {
17
       vector<int> C;
       for (int i = 0, t = 0; i < A.size(); i ++ ) {</pre>
19
           t = A[i] - t;
           if (i < B.size()) t -= B[i];</pre>
          C.push_back((t + 10) % 10);
^{22}
           if (t < 0) t = 1;
23
          else t = 0;
24
25
       while (C.size() > 1 && C.back() == 0) C.pop_back();
26
       return C;
28
   // C = A * b, A >= 0, b >= 0
30
   vector<int> mul(vector<int> &A, int b) {
31
       vector<int> C;
       int t = 0;
       for (int i = 0; i < A.size() || t; i ++ ) {
```

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```
if (i < A.size()) t += A[i] * b;</pre>
35
          C.push_back(t % 10);
36
37
          t /= 10;
       while (C.size() > 1 && C.back() == 0) C.pop_back();
39
       return C;
40
   }
41
42
   // A / b = C ... r, A >= 0, b > 0
   vector<int> div(vector<int> &A, int b, int &r) {
44
45
       vector<int> C;
       r = 0;
46
       for (int i = A.size() - 1; i >= 0; i -- ) {
47
          r = r * 10 + A[i];
48
          C.push_back(r / b);
49
          r %= b;
51
       }
       reverse(C.begin(), C.end());
52
       while (C.size() > 1 && C.back() == 0) C.pop_back();
53
       return C;
54
   }
55
```

6.6 CDQ 分治

```
#pragma GCC optimize(2)
   #define IOS ios::sync_with_stdio(false); cin.tie(0);
   #include<bits/stdc++.h>
   #define int long long
   using namespace std;
   const int MAXN = (int)1e5 + 5;
   struct Node {
       int val, del, ans;
10
   }a[MAXN];
11
   int n, m, ans, c[MAXN], pos[MAXN];
13
14
   bool cmp1(const Node &a, const Node &b) { return a.val < b.val; }</pre>
15
   bool cmp2(const Node &a, const Node &b) { return a.del < b.del; }</pre>
16
   void Modify(int x, int k) {
       while(x <= n) {
19
          c[x] += k;
20
21
          x += (x & (-x));
22
   }
23
24
   int Query(int x) {
       int res = 0;
       while(x > 0) {
27
          res += c[x];
28
          x -= (x & (-x));
29
30
       return res;
32
33
   void Solve(int 1, int r) {
34
35
       if(l == r) return ;
36
       int mid = 1 + r \gg 1;
       Solve(1, mid), Solve(mid + 1, r);
37
       int i = 1, j = mid + 1;
       while(i <= mid) {</pre>
40
```

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```
while(a[i].val > a[j].val && j <= r) Modify(a[j].del, 1), ++j;</pre>
41
          a[i].ans += Query(m + 1) - Query(a[i].del), ++i;
42
       i = 1, j = mid + 1;
       while(i <= mid) {</pre>
          while(a[i].val > a[j].val && j <= r) Modify(a[j].del, -1), ++j;</pre>
          ++i;
       i = mid, j = r;
50
       while(j > mid) {
51
          while(a[j].val < a[i].val && i >= 1) Modify(a[i].del, 1), --i;
52
          a[j].ans += Query(m + 1) - Query(a[j].del), --j;
53
54
       i = mid, j = r;
55
       while(j > mid) {
          while(a[j].val < a[i].val && i >= 1) Modify(a[i].del, -1), --i;
           --j;
58
       }
59
       sort(a + 1, a + r + 1, cmp1);
60
61
   signed main()
       IOS
65
       cin >> n >> m;
66
       for(int i = 1; i <= n; i++) cin >> a[i].val, pos[ a[i].val ] = i;
67
       for(int i = 1; i <= m; i++) {
          int cur; cin >> cur;
          a[ pos[cur] ].del = i;
       for(int i = 1; i <= n; i++) if(!a[i].del) a[i].del = m + 1;</pre>
72
       for(int i = 1; i <= n; i++) {</pre>
73
          ans += Query(n) - Query(a[i].val);
          Modify(a[i].val, 1);
       for(int i = 1; i <= n; i++) Modify(a[i].val, -1);</pre>
       Solve(1, n);
78
       sort(a + 1, a + n + 1, cmp2);
79
       for(int i = 1; i <= m; i++) {
80
          cout << ans << "\n";</pre>
81
          ans -= a[i].ans;
82
       return 0;
   }
```

7 杂项

7.1 STL

```
// 万能(误)算法头文件(部分)
  #include <algorithm>
   using namespace std;
   int main() {
     iterator begin, end; // 指代某种数据结构首尾迭代器
     T i, x, a, b;
     sort(begin, end, <cmp>); // 排序函数, 默认从小到大
     // 遇到需要特殊排序的需要编写cmp函数 or 重载内部运算符
     next_permutation(begin, end); // 下一个排列
10
     prev_permutation(begin, end); // 前一个排列
11
12
     set_union(begin(a), end(a), begin(b), end(b), begin(c));
13
     // 取两个有序序列a、b的并集,存放到c中
```

```
set_intersection(begin(a), end(a), begin(b), end(b), begin(c));
     // 取两个有序序列a、b的交集,存放到c中
16
     set_difference(begin(a), end(a), begin(b), end(b), begin(c));
17
     // 取两个有序序列a、b的差集,存放到c中
     unique(begin, end); // 有序数据去重
19
     merge(begin(a), end(a), begin(b), end(b), begin(c), cmp);
20
     // 合并两个有序序列a、b, 存放到c中, cmp可定义新序列排列方式
21
22
     lower_bound(begin, end, x); // 返回x的前驱迭代器
     // 在普通的升序序列中, x的前驱指的是第一个大于等于x的值
     upper_bound(begin, end, x); // 返回x的后继迭代器
     // 在普通的升序序列中, x的后继指的是第一个大于x的值
26
     // 上述两个函数时间复杂度为O(log2(n)), 内部实现是二分
27
     // 如果找不到这样的值,会返回end
28
29
     find(begin, end, x); // O(n) 查找x
     binary_search(begin, end, x) // 二分查找x, 返回bool
     min(a, b); max(a, b); // 返回a、b中的最小/最大值
32
     fill(begin, end, x); // 往容器的[begin, end)内填充x
33
     swap(a, b); // 交换a、b的值
34
     return 0;
   // 动态数组(vector)、双向链表(list)
  #include <vector>
  #include <list>
  using namespace std;
42
  int main() {
     Ti;
     unsigned int n, x;
     bool flag;
46
     iterator it;
47
     // 动态数组部分
     // 注意vector的空间需要预留两倍大小
     vector<T> v;
     v.push_back(i); // 往数组尾添加一个元素i
52
     v[x]; // 访问第x - 1个元素
53
     v.begin(); // 返回头元素的迭代器
54
     v.end(); // 返回末尾迭代器 (尾元素的下一个)
55
     n = v.size(); // 数组中元素数量
     v.pop_back(); // 删除最后一个元素
     v.erase(it); // 删除某个的元素
     v.insert(x, i); // 在x位置插入元素i
     // erase、insert时间复杂度为O(n)
     v.clear(); // 清空数组, 不释放空间
     flag = v.empty(); // 判断数组是否为空 (真值)
     // 链表部分
     list<T> li;
     li.push_front(i); // 在链头添加一个元素i
66
     li.push_back(i); // 在链尾添加一个元素i
67
     li.pop_front(i); // 删除链表头元素
     li.pop_back(i); // 删除链表尾元素
     li.erase(it); // 删除某个的元素
     li.insert(x, i); // 在x位置插入元素i O(n)
     li.begin(); // 返回头元素的迭代器
72
     li.end(); // 返回末尾迭代器 (尾元素的下一个)
73
     n = li.size(); // 链表中元素数量
     li.remove(i); // 删除链表中所有值为i的元素
     li.unique(); // 移除所有连续相同元素, 留下一个
     li.reverse(); // 反转链表
     li.clear(); // 情况链表, 不释放空间
78
```

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```
return 0;
   }
81
82
   // 普通队列、双端队列、优先队列
   #include <queue> // 队列头文件
84
   #include <deque> // 双端队列头文件
85
   using namespace std;
   int main() {
      Ti;
      unsigned int n, x;
90
      bool flag;
91
      // 普通队列部分,注意queue没有迭代器
92
      queue<T> q, tmp_q; // 定义普通队列
93
      q.push(i); // 队尾插入元素i
94
      q.pop(); // 弹出队首元素
      i = q.front(); // 访问队首元素
      i = q.back(); // 访问队尾元素
      n = q, size(); // 队内元素数量
      flag = q.empty(); // 判断队列是否为空(真值)
      q.swap(tmp_q); // 交换两个队列元素
100
101
      // 优先队列部分,注意其没有迭代器
      priority_queue<T> pq; // 定义优先队列
      pq.push(i); // 队尾插入元素i
104
      pq.pop(); // 弹出队首元素
105
      i = pq.top(); // 访问队首元素
106
      n = q, size(); // 队内元素数量
107
      flag = q.empty(); // 判断队列是否为空(真值)
108
       q.swap(tmp_q); // 交换两个队列元素
      // 注意优先队列内部是使用<运算符,默认大根堆
      // 可以采用重载运算符或加入运算符类自定义排列方式
111
      // 例: priority_queue<T, vector<T>, greater<T> > 小根堆
112
113
          struct node {
114
115
             int x, y;
116
         bool operator < (node a, node b) {</pre>
117
             // 这里注意是<右边的元素会放在前面
118
            if(a.x != b.x) return a.x < b.x;
119
            else return a.y < b.y;</pre>
120
          }
121
         priority_queue<node>
124
      // 双端队列部分
125
      // 注意deque用到了map来映射,时间复杂度上常数略大
126
      deque<T> dq; // 定义双端队列
127
      // 可以称为vector、list、queue的结合体
       // 用法类似,这里只给代码不做注释
      dq.push back(i);
130
      dq.push_front(i);
131
      dq.front();
132
      dq.back();
133
      dq.pop_front();
134
      dq.pop_back();
      dq.begin();
      dq.end();
137
      dq[x];
138
      n = dq.size();
139
      flag = dq.empty();
140
      dq.insert(x, i);
141
      return 0;
143
  | }
144
```

```
145
    // 栈
146
147
   #include <stack>
   using namespace std;
    int main() {
149
      Τi;
150
       unsigned int n;
151
      bool flag;
152
       stack<T> st; // 注意stack没有迭代器
154
       st.push(i); // 往栈顶加入一个元素
155
       st.pop(); // 弹出栈顶元素
156
       i = st.top(); // 获得栈顶元素的值
157
      flag = st.empty(); // 判断是否为空(真值)
158
       n = st.size(); // 获得栈内元素个数
159
160
161
      return 0;
   }
162
163
   // pair (成组) 、set (有序元素序列)
164
   #include <set>
165
   #include <pair>
    using namespace std;
   int main() {
168
      Ti;
169
      T1 t1;
170
      T2 t2;
171
       iterator it;
172
       unsigned int n;
      bool flag;
       // pair是将两种元素组成一对
176
       pair<T1, T2> p;
177
       p = make_pair(t1, t2); // 将t1、t2构造成一对
       // pair支持比较,遵循字典序
179
       p.first; // 访问第一个元素, 这里是t1
       p.second; // 访问第二个元素, 这里是t2
181
182
       // set内部是RB-tree维护
183
       set<int> st; // 注意, set内元素不重复
184
       st.insert(i); // 往set内插入一个元素i
185
          // 时间复杂度O(log2(n)) 这里会返回一个<pair>迭代器
          // first指向插入元素后所在的迭代器
          // second指向是否插入成功(真值)
       st.begin(); // 返回首迭代器
189
       st.end(); // 范围尾迭代器
190
       st.erase(it); st.erase(i);
191
          // 删除某个元素
192
       st.equal_range(i); // 返回几何中与i相等的上下限两个迭代器
       flag = st.empty();
      n = st.size();
195
      st.clear();
196
       // set内置了lower bound和upeer bound函数
197
       // 用法和algorithm的一样
198
199
       // 可重复元素set
      multiset<int> mst;
       // 用法与set大致相同
202
      // 唯一不同只在删除函数上
203
      mst.erase(i); // 会删除所有值为i的元素
204
205
      return 0;
206
207
   // 如果需要给set自定义排序顺序
209
```

```
struct CMP {
210
       bool operator() (const int& a, const int& b) const {
211
212
          return a > b; // 返回真值则代表左边的值优先级高
    };
214
    multiset<int, CMP> mst;
215
216
    // map (映射)
217
    #include <map>
    using namespace std;
220
    int main() {
       T1 t1;
221
       T2 t2;
222
223
       // map将两种元素做映射, 一种指向另一种
224
225
       // 内部也是RB-tree维护
       map<T1, T2> mp;
       mp[t1] = t2; // 直接让t1对应到t2
227
       mp[t1]; // 访问t1对应的内容, 时间复杂度O(log2(n))
228
       // 如果t1没有指向任何内容,则会返回T2类型的初始值
229
230
       return 0;
231
    // 一些C++的功能/特性
234
    #include <bits/stdc++.h> // 标准库头文件
235
    using namespace std;
236
    int main() {
237
238
        _int128 a; // 128位整数,最大值大概10^38次方
       // C++11以上可用,无法用标准方法读入
241
       cin.tie(0); cout.tie(0);
242
       ios::sync_with_stdio(false);
243
       // 关闭cin、cout同步流,此举后不可混用scanf/printf
244
       auto x; // 自动变量, 可以是任意属性
       // 举个例子
247
       std::set<int> st;
248
       std::for(auto i:st); // C++版for each
249
       // 其中i是auto变量,也可改成set<int>::iterator
250
251
       // 所有的STL容器push/insert操作都可替换为emplcace
       // 速度上减小常数 (不用临时变量)
       // 例:
254
       int i;
255
       std::set<int> st;
256
       st.emplace(i);
257
       std::vector<int> vc;
       vc.emplace_back(i);
260
       return 0;
261
    }
262
263
    // 强大的pb_ds库
264
    #include <bits/stdc++.h>
265
    #include <bits/extc++.h> // 扩展库头文件
267
    * 这里如果没有bits/extc++.h的话需要
268
    * ext/pb_ds/tree_policy.hpp
269
    * ext/pb_ds/assoc_container.hpp
270
     * ext/pb_ds/priority_queue_policy.hpp
    * ext/pb_ds/trie_policy.hpp
    * ext/rope
273
274
```

```
275
   using namespace __gnu_pbds;
276
   using namespace __gnu_cxx;
277
   int main() {
279
      // 哈希表部分,用法与map一样,效率在C++11以下效率高
280
      // 注意, 这部分在namespace gnu pbds下
281
      cc_hash_table<string, int> mp1; // 拉链法
282
      gp_hash_table<string, int> mp2; // 查探法(快一些)
284
      // 优先队列部分,比STL中高级
285
      priority queue<int, std::greater<int>, TAG> pq;
286
287
         第一个参数是数据类型
288
       * 第二个是排序方式
289
       * 第三个是堆的类型
       * 其中堆的类型有下面几种
       * pairing_heap_tag
292
       * thin_heap_tag
293
       * binomial_heap_tag
       * c_binomial_heap_tag
       * binary_heap_tag
       * 其中pairing_heap_tag最快
         并且这个东西是带默认参数的,只需要定义一个int
299
      // 比STL中的优先队列多了join和迭代器
300
      // 例子:
301
302
      priority_queue<int> pq1, pq2;
303
      pq1.join(pq2); // 会将pq2合并到pq1上
      pq1.begin(); pq1.end(); // 可遍历
      // 红黑树 (平衡树) 部分, 与set相似, 但更快
306
      tree <
307
308
         int,
         null_type,
309
          std::less<>,
         rb_tree_tag,
311
         tree order statistics node update
312
      > t, tre;
313
314
       * int 关键字类型
315
       * null_type 无映射 (低版本g++为null_mapped_type)
       * less<int> 从小到大排序
       * rb_tree_tag 红黑树 (splay_tree_tag splay)
       * tree_order_statistics_node_update 结点更新
319
       */
320
      int i, k;
321
      t.insert(i); // 插入
322
      t.erase(i); // 删除
      t.order_of_key(i);
      // 询问这个tree中有多少个比i小的元素
325
      t.find_by_order(k);
326
      // 找第k + 1小的元素的迭代器,如果order太大会返回end()
327
      t.join(tre); // tre合并到t上
328
      t.split(i, tre); // 小于等于i的保留, 其余的属于tre
329
      // 基本操作有size()/empty()/begin()/end()等
      // 同样内置lower_bound/upper_bound
332
      // 可持久化平衡树部分
333
      // 注意, 这部分在namespace __gun_cxx下
334
      rope<char> str;
335
      // 待我学习完后再更新
      return 0;
338
   }
339
```

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7.2 优化

```
// 关闭iostream同步流
   std::ios::sync_with_stdio(false); std::cin.tie(0);
   // 如果编译开启了 C++11 或更高版本, 建议使用 std::cin.tie(nullptr);
   // 注意,此后不可和scanf/printf混用
   // 普通快读快写
   inline void read_int(int &X) {
      X = 0; int w = 0; char ch = 0;
      while(!isdigit(ch)) w |= ch=='-', ch = getchar();
      while( isdigit(ch)) X = (X << 3) + (X << 1) + (ch-48), ch = getchar();
      X = w ? -X : X;
11
   }
12
13
   inline void write_int(int x) {
14
      static int sta[65];
       int top = 0;
      do {
17
          sta[top++] = x % 10, x /= 10;
18
       } while(x);
19
      while(top) putchar(sta[--top] + 48);
20
   }
21
22
   // fread快读
   namespace fastIO {
   #define BUF SIZE 100000
25
       //fread -> read
26
      bool IOerror = 0;
27
       inline char nc() {
          static char buf[BUF_SIZE], *p1 = buf + BUF_SIZE, *pend = buf + BUF_SIZE;
          if (p1 == pend) {
30
             p1 = buf;
31
             pend = buf + fread(buf, 1, BUF_SIZE, stdin);
32
             if (pend == p1) {
33
                 IOerror = 1;
34
                 return -1;
             }
          }
          return *p1++;
39
      inline bool blank(char ch) {
40
          return ch == ' ' || ch == '\n' || ch == '\r' || ch == '\t';
       inline void read(int &x) {
43
          char ch;
          while (blank(ch = nc()));
45
          if (IOerror) return;
46
          for (x = ch - '0'; (ch = nc()) >= '0' && ch <= '9'; x = x * 10 + ch - '0');
47
48
   #undef BUF_SIZE
   };
   using namespace fastIO;
   // 手动扩栈
53
   #pragma comment(linker, "/STACK:1024000000,1024000000")
54
   // 02 03优化
   #pragma GCC optimize(2)
   #pragma GCC optimize(3)
```

7.3 对拍

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```
2 :loop
3 随机数据生成.exe
4 暴力.exe
5 正解.exe
6 fc 暴力.out 正解.out
7 if not errorlevel 1 goto loop
8 pause
9 :end
```

7.4 Java 大整数类

```
// Java大整数
   import java.util.*;
   import java.math.*;
   public class BigInt {
      static Scanner in = new Scanner(System.in); // 定义输入对象
6
      public static void main(String[] args) {
         BigInteger bigInt_1 = new BigInteger("100");
         BigInteger bigInt_2 = BigInteger.valueOf(123);
         //两种定义方式,建议使用第一种
10
         bigInt_1.add(bigInt_2); // 加法
         bigInt_1.subtract(bigInt_2); // 减法
         bigInt_1.multiply(bigInt_2); // 乘法
13
         bigInt_1.divide(bigInt_2);// 除法, 向下取整
         bigInt_1.divideAndRemainder(bigInt_2);
            // 返回一个BigInteger[], 包含商和余数
16
         bigInt_1.remainder(bigInt_2); // 取余数,与this同符号
         bigInt_1.mod(bigInt_2); // 取模,模数只能为正整数
         bigInt_1.pow(10); // a^b
19
         bigInt_1.gcd(bigInt_2); // 最大公约数
20
         bigInt_1.compareTo(bigInt_2);
21
            // 比较大小, <0表示this小, 0表示相等, >0表示this大
22
         bigInt_1.equals(bigInt_2); // 真值为相等
         bigInt_1.negate(); // -a
         bigInt_1.abs(); // 绝对值
         bigInt_1.min(bigInt_2); // 最小值
         bigInt 1.max(bigInt 2); // 最大值
27
         BigInteger a = in.nextBigInteger(); // 读入
28
         System.out.println(bigInt_1); // 输出
29
   }
31
32
   // Java大实数
33
   import java.util.*;
34
   import java.math.*;
35
36
   public class Big {
37
      static Scanner in = new Scanner(System.in); // 定义输入对象
      public static void main(String[] args) {
39
         BigDecimal bigDec 1 = new BigDecimal("123.1");
40
         BigDecimal bigDec_2 = BigDecimal.valueOf(233.213);
41
         // 唯有除法与BigInteger不同,无限小数需要规定保留位数
         int scale = 3; // 保留位数
         bigDec_1.divide(bigDec_2, scale, RoundingMode.HALF_UP);
            第三个参数为保留方式, 有以下几种:
            CEILING 正无穷大方向取整
47
            FLOOR 负无穷大方向取整
49
            DOWN 向 Ø 的方向取整
            UP 正数向正无穷大取整,负数向负无穷大取整
            HALF_UP 常用的4舍5入
            HALF_DOWN 6,7,8,9 向上取整
            HALF_EVEN 小数位是5时,判断整数部分是奇数就进位,6,7,8,9 进位
```

 ${\rm Orzjh\text{-}XCPC}$

```
54 */
55 }
56 }
```