

DATE:7/6/24

LAB-3

1.You are given a string s, and an array of pairs of indices in the string pairs where pairs[i] = [a, b] indicates 2 indices(0-indexed) of the string.You can swap the characters at any pair of indices in the given pairs any number of times. Return the lexicographically smallest string that s can be changed to after using the swap

CODE:

```
class Graph:
    def __init__(self, V):
        # Number of vertices in the graph
        self.V = V
        # Adjacency list to represent the graph
        self.adj = [[] for i in range(V)]

    def addPair(self, v, w):
        # Adjust indices to be 0-based
        v -= 1
        w -= 1
        # Add a relation (edge) between v and w
        self.adj[v].append(w)
        # Also add the reverse relation (undirected graph)
        self.adj[w].append(v)

    def disjointSets(self):
        # Create a list to track visited vertices
        visited = [False] * self.V
        # List to store groups of connected vertices
        all_sets = []

    def countAll(v, group):
        # Mark the current vertex as visited
        visited[v] = True
```

```
# Add the vertex to the current group (adjust index)
group.append(v + 1)
```

```
# Recursively explore neighbors of the current vertex
for neighbor in self.adj[v]:
    if not visited[neighbor]:
        countAll(neighbor, group)
```

```
for i in range(self.V):
    # If the vertex hasn't been visited,
    # it's the start of a new group
    if not visited[i]:
        # Create a new group
        group = []
        # Find all vertices connected to this one
        countAll(i, group)
        # Add the group to the list of all groups
        all_sets.append(group)
```

```
return all_sets
```

```
# Driver Code
```

```
# Input string
```

```
Str = 'zcxfbe'
```

```
# Number of elements
```

```
N = 6
```

```
# Pairs of indices indicating relationships
```

```
Pairs = [[0, 1], [0, 2], [3, 5]]
```

```
# Create a graph with N vertices
```

```
g = Graph(N)
```

```

# Add relations based on pairs
for i in Pairs:
    g.addPair(i[0], i[1])

# Find groups formed by the relations
groups = g.disjointSets()

key = [] # List to store indices of characters
value = [] # List to store characters

# Sort characters within each group and store the results
for i in range(len(groups)):
    # Temporary list to store characters within a group
    semians = []
    for j in groups[i]:
        # Extract characters from the input string
        semians.append(Str[j])
    # Sort the characters lexicographically
    semians.sort()
    # Sort the indices
    groups[i].sort()

# Add sorted characters to the value list
value.extend(semians)

# Add sorted indices to the key list
key.extend(groups[i])

# Initialize a list to reconstruct the final string
ans = [""] * N

# Reconstruct the final string

```

```
# based on sorted characters and indices
```

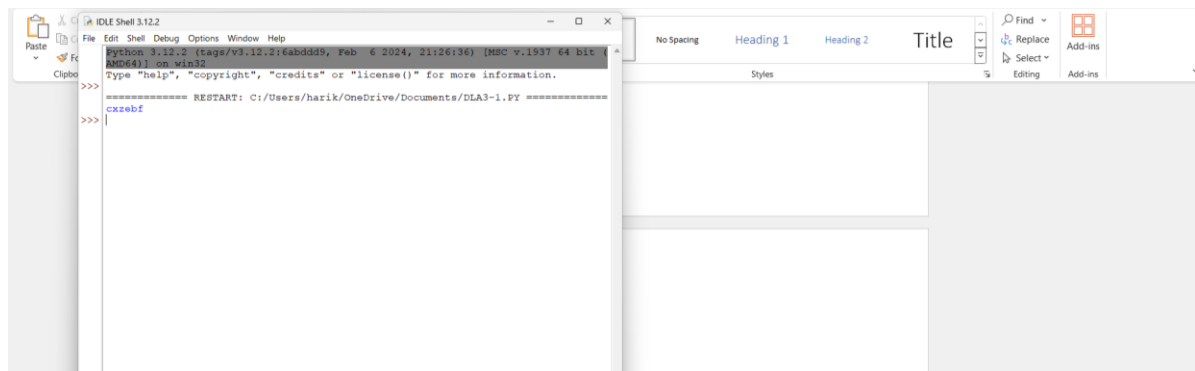
```
for i in range(N):
```

```
ans[key[i]] = value[i]
```

```
# Print the rearranged string
```

```
print("".join(ans))
```

OUTPUT:



2. Given two strings: s1 and s2 with the same size, check if some permutation of string s1 can break some permutation of string s2 or vice-versa. In other words s2 can break s1 or vice-versa. A string x can break string y (both of size n) if $x[i] \geq y[i]$ (in alphabetical order) for all i between 0 and n-1.

Code:

```
def canBreak(s1, s2):
```

```
    # Sort both strings
```

```
    sorted_s1 = sorted(s1)
```

```
    sorted_s2 = sorted(s2)
```

```
    # Check if sorted_s1 can break sorted_s2
```

```
    can_s1_break_s2 = all(c1 >= c2 for c1, c2 in zip(sorted_s1, sorted_s2))
```

```
    # Check if sorted_s2 can break sorted_s1
```

```
    can_s2_break_s1 = all(c2 >= c1 for c1, c2 in zip(sorted_s1, sorted_s2))
```

```
    # Return True if either can break the other
```

```
    return can_s1_break_s2 or can_s2_break_s1
```

Example usage

```
s1 = "abc"
```

```
s2 = "xya"
```

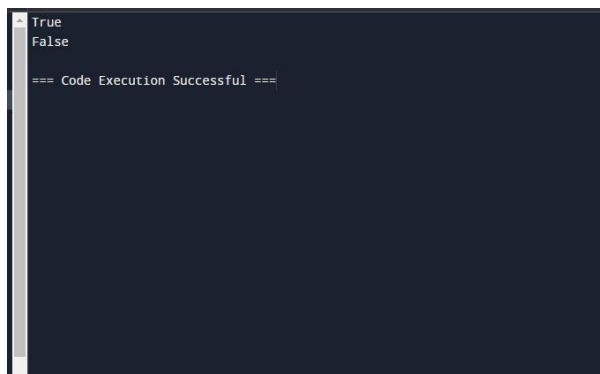
```
print(canBreak(s1, s2)) # Output: True, since "abc" can be permuted to "cba" which can break "axy"
```

```
s1 = "abe"
```

```
s2 = "acd"
```

```
print(canBreak(s1, s2)) # Output: False, neither permutation can break the other
```

OUTPUT:



```
True
False
=== Code Execution Successful ===
```

3. You are given a string s . $s[i]$ is either a lowercase English letter or '?'. For a string t having length m containing only lowercase English letters, we define the function $\text{cost}(i)$ for an index i as the number of characters equal to $t[i]$ that appeared before it, i.e. in the range $[0, i - 1]$. The value of t is the sum of $\text{cost}(i)$ for all indices i .

Code:

```
def minimize_cost_string(s):
```

```
    # Dictionary to keep track of the frequency of each character
```

```
    frequency = {chr(i): 0 for i in range(ord('a'), ord('z') + 1)}
```

```
    result = []
```

```
    for char in s:
```

```
        if char == '?':
```

```
            # Find the character with the minimum frequency
```

```

    min_char = min(frequency, key=frequency.get)
    result.append(min_char)

    # Update the frequency of the chosen character
    frequency[min_char] += 1
else:
    result.append(char)

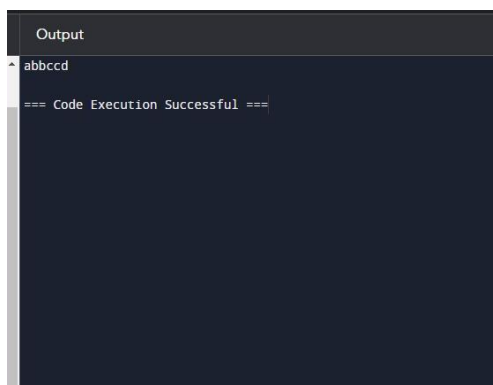
    # Update the frequency of the current character
    frequency[char] += 1

return ''.join(result)

# Example usage
s = "a?b?c?"
print(minimize_cost_string(s)) # Example output: "aabbbc"

OUTPUT:

```



The screenshot shows a dark-themed output window. At the top, the word 'Output' is visible. Below it, the string 'abbccd' is displayed. Underneath that, the text '=== Code Execution Successful ===' is shown. The window has a small upward-pointing arrow icon on the left side.

4. You are given a string *s*. Consider performing the following operation until *s* becomes empty: For every alphabet character from 'a' to 'z', remove the first occurrence of that character in *s* (if it exists).

CODE:

```

def last_state_before_empty(s):
    from string import ascii_lowercase

    while True:
        prev_s = s
        for char in ascii_lowercase:

```

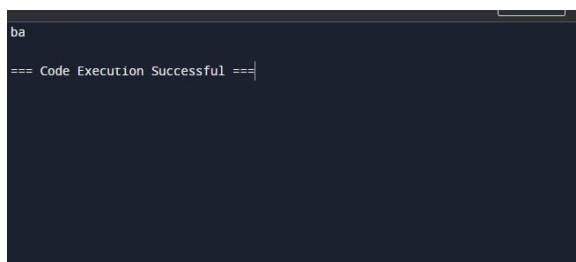
```
s = s.replace(char, "", 1)
if s == "":
    return prev_s
```

Example usage

```
s = "aabcbbca"
```

```
print(last_state_before_empty(s))
```

OUTPUT:



5. Given an integer array nums, find the subarray with the largest sum, and return its sum.

Code:

```
def max_subarray_sum(nums):
```

```
    # Initialize the maximum sum to be the first element
```

```
    max_current = max_global = nums[0]
```

```
    for num in nums[1:]:
```

```
        # Calculate the maximum sum of subarray ending at the current position
```

```
        max_current = max(num, max_current + num)
```

```
    # Update the global maximum sum if the current maximum sum is greater
```

```
    if max_current > max_global:
```

```
        max_global = max_current
```

```
    return max_global
```

Example usage

```
nums = [-2,1,-3,4,-1,2,1,-5,4]
```

```
print(max_subarray_sum(nums))
```

OUTPUT:

```
6
=== Code Execution Successful ===
```

6. You are given an integer array `nums` with no duplicates. A maximum binary tree can be built recursively from `nums` using the following algorithm: Create a root node whose value is the maximum value in `nums`. Recursively build the left subtree on the subarray prefix to the left of the maximum value. Recursively build the right subtree on the subarray suffix to the right of the maximum value. Return the maximum binary tree built from `nums`.

Code:

```
class TreeNode:
```

```
    def __init__(self, val=0, left=None, right=None):
```

```
        self.val = val
```

```
        self.left = left
```

```
        self.right = right
```

```
def constructMaximumBinaryTree(nums):
```

```
    if not nums:
```

```
        return None
```

```
    # Find the index of the maximum value in the current array
```

```
    max_index = nums.index(max(nums))
```

```
    # Create the root node with the maximum value
```

```
    root = TreeNode(nums[max_index])
```

```
    # Recursively build the left and right subtrees
```

```
    root.left = constructMaximumBinaryTree(nums[:max_index])
```

```
    root.right = constructMaximumBinaryTree(nums[max_index + 1:])
```



```
return root
```

```
# Helper function to print the tree in a readable format (preorder traversal)
```

```
def preorderTraversal(root):
```

```
    if not root:
```

```
        return []
```

```
    return [root.val] + preorderTraversal(root.left) + preorderTraversal(root.right)
```

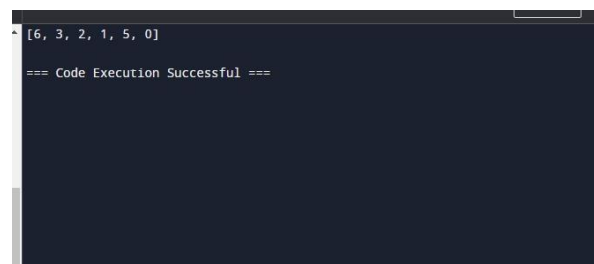
```
# Example usage
```

```
nums = [3, 2, 1, 6, 0, 5]
```

```
tree = constructMaximumBinaryTree(nums)
```

```
print(preorderTraversal(tree))
```

OUTPUT:



```
[6, 3, 2, 1, 5, 0]
=== Code Execution Successful ===
```

7. Given a circular integer array `nums` of length `n`, return the maximum possible sum of a non-empty subarray of `nums`. A circular array means the end of the array connects to the beginning of the array. Formally, the next element of `nums[i]` is `nums[(i + 1) % n]` and the previous element of `nums[i]` is `nums[(i - 1 + n) % n]`. A subarray may only include each element of the fixed buffer `nums` at most once. Formally, for a subarray `nums[i], nums[i + 1], ..., nums[j]`, there does not exist $i \leq k_1, k_2 \leq j$ with $k_1 \% n == k_2 \% n$.

CODE:

```
def maxSubarraySumCircular(nums):
```

```

# Helper function to perform Kadane's algorithm
def kadane(array):
    max_current = max_global = array[0]
    for num in array[1:]:
        max_current = max(num, max_current + num)
        if max_current > max_global:
            max_global = max_current
    return max_global

# Case 1: Max subarray sum for non-circular case
max_kadane = kadane(nums)

# Case 2: Max subarray sum for circular case
total_sum = sum(nums)

# Invert the nums array to find the minimum subarray sum
nums_inverted = [-num for num in nums]

# Find the maximum subarray sum for the inverted array
# This is equivalent to finding the minimum subarray sum for the original array
max_inverted_kadane = kadane(nums_inverted)

max_circular = total_sum + max_inverted_kadane # max_circular considers the wrap-around part

# Handle the edge case where all elements are negative
if max_circular == 0:
    return max_kadane

return max(max_kadane, max_circular)

# Example usage
nums = [5, -3, 5]

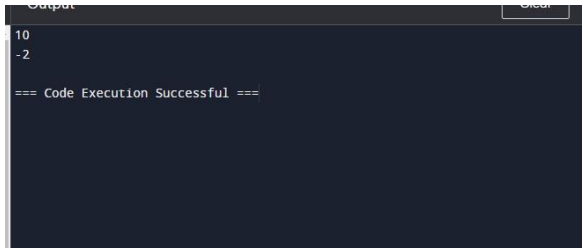
```

```
print(maxSubarraySumCircular(nums))
```

```
nums = [-3, -2, -3]
```

```
print(maxSubarraySumCircular(nums))
```

output:

A screenshot of a code execution window with a dark background. The window has a title bar with 'Output' on the left and a 'Clear' button on the right. The output area shows the number '10' on the first line and '-2' on the second line. Below the output, there is a status message: '=== Code Execution Successful ==='.

```
Output
10
-2
=== Code Execution Successful ===
```

8.You are given an array `nums` consisting of integers. You are also given a 2D array `queries`, where `queries[i] = [posi, xi]`.For query `i`, we first set `nums[posi]` equal to `xi`, then we calculate the answer to query `i` which is the maximum sum of a subsequence of `nums` where no two adjacent elements are selected. Return the sum of the answers to all queries. Since the final answer may be very large, return it modulo $10^9 + 7$. A subsequence is an array that can be derived from another array by deleting some or no elements without changing the order of the remaining elements.

Code:

```
def max_non_adjacent_sum(nums):
    include = 0
    exclude = 0

    for num in nums:
        new_include = exclude + num
        new_exclude = max(include, exclude)
        include = new_include
        exclude = new_exclude

    return max(include, exclude)

def solve(nums, queries):
    MOD = 10**9 + 7
    result = 0
```

```
for pos, val in queries:

    # Update the array
    nums[pos] = val

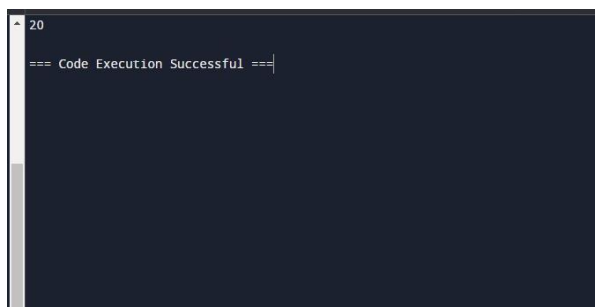
    # Calculate the maximum sum of the non-adjacent subsequence
    max_sum = max_non_adjacent_sum(nums)

    # Add to the result
    result = (result + max_sum) % MOD

return result

# Example usage
nums = [1, 2, 3, 4]
queries = [[0, 2], [3, 5], [1, 1]]
print(solve(nums, queries))
```

OUTPUT:



9. Given an array of points where `points[i] = [xi, yi]` represents a point on the X-Y plane and an integer `k`, return the `k` closest points to the origin `(0, 0)`. The distance between two points on the X-Y plane is the Euclidean distance (i.e., $\sqrt{(x_1 - x_2)^2 + (y_1 - y_2)^2}$). You may return the answer in any order. The answer is guaranteed to be unique (except for the order that it is in).

CODE:

```
import heapq

def kClosest(points, k):
    # Create a min-heap
    heap = []

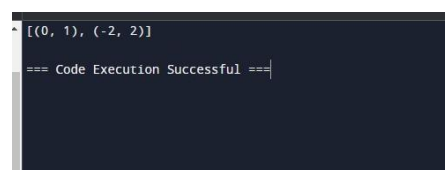
    for (x, y) in points:
        # Calculate the squared distance from the origin
        dist = x * x + y * y
        # Push the distance and the point onto the heap
        heapq.heappush(heap, (dist, (x, y)))

    # Extract the k closest points
    closest_points = [heapq.heappop(heap)[1] for _ in range(k)]

    return closest_points

# Example usage
points = [[1, 3], [-2, 2], [5, 8], [0, 1]]
k = 2
print(kClosest(points, k))
```

OUTPUT:

A screenshot of a code execution environment with a dark background. The first line of output is `[(0, 1), (-2, 2)]`. Below it, a status message reads `=== Code Execution Successful ===` with a cursor at the end.

10. Given two sorted arrays `nums1` and `nums2` of size `m` and `n` respectively, return the median of the two sorted arrays. The overall run time complexity should be $O(\log(m+n))$.

CODE:

```
def findMedianSortedArrays(nums1, nums2):
    # Ensure nums1 is the smaller array
    if len(nums1) > len(nums2):
        nums1, nums2 = nums2, nums1

    m, n = len(nums1), len(nums2)
    imin, imax, half_len = 0, m, (m + n + 1) // 2

    while imin <= imax:
        i = (imin + imax) // 2
        j = half_len - i

        if i < m and nums1[i] < nums2[j - 1]:
            imin = i + 1
        elif i > 0 and nums1[i - 1] > nums2[j]:
            imax = i - 1
        else:
            # i is perfect
            if i == 0: max_of_left = nums2[j - 1]
            elif j == 0: max_of_left = nums1[i - 1]
            else: max_of_left = max(nums1[i - 1], nums2[j - 1])

            if (m + n) % 2 == 1:
                return max_of_left

            if i == m: min_of_right = nums2[j]
            elif j == n: min_of_right = nums1[i]
            else: min_of_right = min(nums1[i], nums2[j])

    return (max_of_left + min_of_right) / 2
```

Example usage

nums1 = [1, 3]

nums2 = [2]

print(findMedianSortedArrays(nums1, nums2))

nums1 = [1, 2]

nums2 = [3, 4]

print(findMedianSortedArrays(nums1, nums2))

OUTPUT:

```
2
2.5
=== Code Execution Successful ===
```