compiler Design Assignment

CS21B2011

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1) which data structure in a compiler is used for managing information about variables and their attributes?

Ans: In a compiler, the data structure is responsible for managing info about variables and their attributes is a symbol table.

The output of a lexical analyses is

(a) A passe tree (b) Intermediate code.

(c) Machine code (d) stream of tokens.

Ans: (d) stream of tokens.

3 No. of tokens in a FORTRAN statement DO 10 I = 1.25 is; (a) 3 (b) 4 (c) 5 (d) None of these.

Ans: There are 3 tokens because it's 1.25 and not 1,25. TANSWER would be 5 if it would have comma 1,25. : (A) 3;

(4) which of the following machine model can be used in a necessary and sufficient sense for lexical analysis?

(1) Deterministic PDA

(2) Finite Automata

(4) Tusing machine (3) Non Deterministic FA Ans: Finite Automata.

- 5) In a compiler, key words are recognised during?

 Ans: Keywords of a language are recognized during lexical, analysis phase.
- 6) The number of Tokens in the c statement:

 printf ("i=1'd, &i=1'. x", i, &i);

 (a) 3 (b) 26 (c) 10 (d) 21

Ans: (0) 10,

: There are 10 tokens in this Statement.

7 Match elements from Group-1 to Group-2

Group-1

P. Regular Expression

Q. push Down Automata

R. Dataflow Analysis

s. Register Allocation

<u>680010-5</u>

1. syntax Analysis

2. code generation

3. Lexical analysis

4. code optimisation.

AMS: P-3: RE is used in Lexical analysis.

Q-1: PDA is used for syntax Analysis.

R-4: performed during code optimisation.

s-2: part of code generation.

18) Type checking is normally done during!

Ans:- type checking is normally done during the syntax Directed Translation (08) Semantic Analysis.

G. Hemanth 1) Arrange given compileration process in correct order: (i) # Linking (ii) Assembly (iii) Compiling (iv) pre-processing

(a) iv→iii→ii→i (b) i→iii→ii→iV

Ans: (a) iv + iii + ii + i

(1) In some languages, an identifier is permitted to be a letter followed by any number of letters and digits. If L and D denote the sets of letters and digits, How do you define identifier?

Ans: L: da-z, A-z & D: 20-9} then Identifier ⇒ L(L+D)*

1 A lexical analyzer uses following patterns to recognize three tokens, T, T2, T3 over la, b, Cf. Ti: a ?(b(c) * a T2:b?(a/c) * b T3: C?(b(a) *C. Note that "x"! means 0,1 occurrence of symbol x. if string bbaacabc is processed by analyzer which is the tokens it outputs?

(a) T, T2 T3 (b) T, T, T3

(c) T2T1 T3 (d) T8 T3

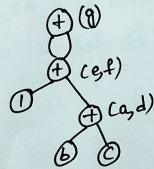
Ans:-(d) T3 T3. Given string: boacabc. Longest matching Prefix is 'bbaac' -> To Remaining Part 'abc' can be made by T3.

(2) consider the following c-code segment, a=b+c; e=a+l; d=b+c; f=d+l; g=e+f; For this, the number of nodes in DAG is?

Ans: Here a and b are same as both have btc. so fand e are also same

: Total nodes = 6 (a, b, c, e, 1,9)

The DAG is,



(13) A top-down parser generals:

(a) Rightmost derivation (b) Rightmost derivation in Reverse.

(c) Leftmost desiration

(d) Leftmost dervation in Reverse.

AMS: (c) Top-down parser is LL parser which means left to right parsing; leftmost derivation. Bottom up parser is LR parser with Left to Right Parsing performing Rightmost derivation in reverse order.

(14) consider the grammer:

s-> as Bld

No. of reduction steps taken by bottom up parser while parsing acad bb is

(a) 5 (b) 6 (c) 7 (d) 8

Ans:- (c) 7,

CS21B2011 G. Hemanth (5) which of the following statements is False?

(a) CFG can be used to specify both lexical & syntax. rules.

(b) Type checking is done before parsing.

(c) both a and b.

(d) None of the above.

Ans: (b) is false.

Type checking is done on parsing tree generated by Parser, so b is wrong.

(16) A grammer is defined as:

A->BC

 $B \rightarrow x/Bx$

c+BID

DAYLEY

Non terminal alphabet of the grammer is:

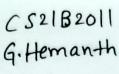
(a) lA,B,C,D,E\$ (b) &B,C,D,E\$ (c) lA,BGD, €, 3,8, ₹\$

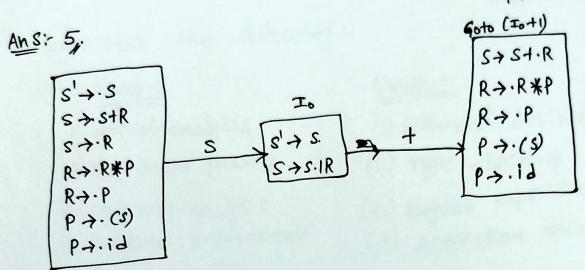
(d) {x, y, z}.

Ansi (a) ¿A, B, C, D, Eg. This set Contains all the symbols from which strings of long can be desired by applying production rules. These are present on Littis of the production rules.

(17) consider the augmented grammer with [+,*,(,),idf as

terminals and if, s'as sastrik Rarpip pacs) lid If Io is the set of two LR(0) items & [s' >s.]. [s->s.+R]& then goto (closusere(I,+4) contains exactly ___ Items.





(B) consider following grammer,

s-asBld B+b

the number of reduction steps taken by a bottom-up

parsser while accepting the string aaadbbb is ____

Ans: Given aaadbbb

(i) s → asB (ii) asB → aasBB

(iii) aasBB→ aaa SBBB (iV) aaa SBBB → aaa dBBB

(V) aaadbbB + aaadbBB (Vi) aaadbBB + aaadbB

(vii) aeadbbB > aaadbbb

:.7 steps.

(9) consider the grammer given here, s>Aa A>BD B>blE D>dlE bet a,b,d and \$ be indexed as, a:3, b:2, d:1,\$:0

compute follow set of B and write the index values of symbols in follow set in descending oxdex.

Ans: Initially, Follow (B) = Fixst(D) = 1d, E3 - 1e3 = 1df

Put & in Prod-2, Follow (B) = Follow (B) U Follow (A) = 1d, a3

: The number = 31 (d,a)

(20) match the following:

GOOD-1

20

P. lexical analysis

a. Top down passing

R. semantic Analysis

s. Runtime Environments

620UP-2

(i) Lettmost desiration

(ii) Type checking.

(iii) Regular Expressions

(iv) Activation secords.

Ansir P-iii, Q-i, R-ii, S-iV

(21) Among simple LR (SLR), canonical LR, and bookahead LR(LALR), which of the following Pairs identify the method that is very easy to implement and the method that is most powerful. (a) SLR, LALR (b) canonical LR, LALR

(c) SLR, canonical LR (d) LALR, canonical LR.

AMS; POWER: LR(0) < SLR(1) < LALR < CLR(1)

complexity:

LR(O) LSLR() LLALR (1) L CLR()

:. SLR, canonical LR > option C.

(22) Match the following:

400UP-1

P. Lexical analysis

Q. Parsing

R. Register Allocation

S. Expression Evaluation 4. Production tree.

680UP-8

1 Graph coloring.

2. DFA minimization.

3. post-oxder traversal.

AMS: P-2; Q-4, R-1; S-3.

@ 23 which of the following is true about shift-Reduce parsing ?

(a) variable prefixes appear only at the buttom of the stack and not inside

(b) viable prefixes appear only at the top of the stack and not inside.

(c) the Stack contains only a set of viable

(d) the stack never contains vaible prefixes

Ans: (c)

The handle is always on top of stack and a Viable prefix will never extend past R-H-S of handle Cietos)

(24) consider grammer defined by S>T*P T->U|T*U P->Q+P|Q Q-id U+id which is true?

(a) + is left * is right associative.

(b) + is right, * is left associative.

(c) Both are right associative.

(d) Both are left associative.

Ans: In T->U/T*U, T->T*U is left recursive. P->Q+P, P->Q+P is right recursive. : * is left, + is right option - (b) //

CS21B2011 9.Hemanth (25) consider the following grammer along with translation S7 S, #TES. Val = S1. Val *T. Val & oules. S> TES. Val = T. Val} T> TIP. RET. Val = T. Val + R. Val & T->R [T.val = R.val] R→id ER. val = id. val f. Here # and % are operators and id is a token that represent a token that represent an integer and id val represents corresponding integer value. Now, And s.val for expression 20#10% 5# 8% 242 Ans: 20# 10% 5#8% 21.2=20*(6:5)*(C8:2)=2) =80/ @consider the following code segment X=u-t; y=x*v; x=y+w; y=t-=; y=x*y; The minimum no of total variables required to convert the above code segment to static assignment Ans: In Static Assignment :. Total = T1 - T5 + (4++, V, W, =) て = い- も T2= T1 *V

27) the purpose of using intermediate cod in

AMS: It increases the chance of reusing the

machine - independent code optimizer in

T3=T2+W

compilers is to -?

other compilers.

T4=t-Z

(28) In the grammer, what is associativity of + and * X > X + X | Y > Z * Y | Z > 19.

AMS: X-X+Y: left Recursive -> ++: left Associative.

Y> Z*Y: Right Recursive > *: right Associative.

(29) A Linker is given object modules for a set of programs that were compiled seperately. What information needed to be included in an object

Ans: Names and locations of all external symbols defined in the object module.

(30) Generation of intermediate code based on an abstract machine model is useful in compilers

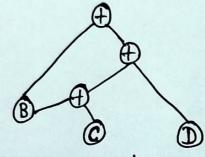
Ans: - It enhances the portability of front - end of

compiler. (31) consider the following SDTS, A > Sb & print 36 s->aA Eprint 13 using SDTS, find output printed by bottom-up parser for input aab

Ans: s > aA printy ->asb print3 +aab print 2

:1324 :: Answer = 231 (32) consider the block given below; a=b+C e=d-bc = a + d a = e + bd = b + c

The nodes and egges in min DAG of above block.



: 6 Nodes & 6 Edges.

(33) which languages necessasily need heap in the

Ans:- Those that allow dynamic data structures.

34) some code optimizations are carried out on the intermediate code because?

Ans:- They enhance the pootability of the compiler to other target processors.

35) which pase of the compiler is responsible for memory allocation & deallocation for variables.

(a) hexical Analysis (b) syntax Analysis

(c) semantic Analysic (d) code generation. Ans: - Semantic Analysis.

(36) Incompatable types work with the

(a) Syntax tree (b) semantic Analyser

(c) code optimiser (d) Lexical Analyser. Ans: (b) semantic Analyses.

CS2182011 G. Hemanth

37) which is NOT intermediate representation of source program? 6) AST

(a) three Address code

(c) CF6

(d) symbol Table.

Ans: (d) symbol.

for names is done in: (38) substitution of values (B) LOOP OPtimization. (A) Local optimization

(d) Strength reduction. (c) constant Folding

Ans:- (c) + constant Folding.

39 which class of statement usually produces no executable code when compiled?

(A) Declaration (B) Assignment statement

(c) I/o Statements (D) Structural Statements.

AMS:- (A) -> Declaration.

consider two modules M, & M2. If M, contains reference will nee to a function defined in M2 the reference will be resolved at:

(A) Edit time (B) Compile time

(c) Link time (D) Load time

Ans:- (c) - link time.

Fach module is compiled seperately and Linked together to make the executable. We can do this Using 9cc.