

# Design & Analysis of Algorithms

## Assignment - 1



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2	Mohammad Jambughodawala	2022A7PS0009H	Tomita Worst-Case Clique Code and Arboricity Clique Generation Code
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4	Vaishnav Devaguptapu	2022A7PS0085H	Project Website and Bron- Kerbosch Degeneracy Code
5	Anshul Gopal	2022A7PS0009H	Readme, Project Report and Arboricity Clique Generation Code

# Algorithms

## Paper 1: The worst-case time complexity for generating all maximal cliques and computational experiments

The algorithm follows a recursive depth-first search strategy. The main steps are:

1. Initialization: The algorithm starts with an empty set  $Q$ , representing the current clique being expanded.
2. Expansion: At each step, a vertex  $q$  is added to  $Q$ , and a new subgraph is considered consisting of vertices adjacent to  $q$ .
3. Pruning: Two pruning techniques are applied:

Finished Set (FINI): Tracks vertices that have already been processed to avoid duplicate cliques.

Candidate Set (CAND): Keeps track of vertices that can still be added to  $Q$ .

4. Recursion: The process continues recursively, generating larger complete subgraphs until maximal cliques are found.
5. Output Optimization: Instead of storing all maximal cliques explicitly, a tree-like output format is used, reducing space complexity.

The key distinction of this algorithm is its pruning mechanism, which significantly reduces redundant computations compared to previous methods.

The worst-case complexity of the algorithm is derived based on the number of maximal cliques in an  $n$ -vertex graph. According to Moon and Moser's theorem, the maximum number of maximal cliques in a graph is  $O(3^{n/3})$ . In the context of space complexity, the algorithm primarily uses recursion and tree-based output, reducing the need to store all cliques explicitly. The space complexity is  $O(n)$  for the recursive stack in the worst case.

## Paper 2: Listing All Maximal Cliques in Sparse Graphs in Near-optimal Time

This algorithm modifies the classic Bron-Kerbosch algorithm by incorporating:

1. **Degeneracy Ordering:** The graph's vertices are processed in order of degeneracy, ensuring that each recursive call has at most  $d$  candidates.
2. **Pivoting Strategy:** A pivot is selected to minimize recursive calls, following the method of Tomita et al.
3. **Efficient Pruning:** The algorithm exploits the degeneracy ordering to restrict search space, reducing redundant computations.
4. **Graph Data Structure:** Uses an adjacency list representation for efficient neighborhood queries.

The algorithm follows these steps:

- Compute a degeneracy ordering of the vertices.
- Process each vertex  $v$  in this order, using a modified Bron-Kerbosch approach to find maximal cliques containing  $v$ .
- Use pivoting to minimize recursive calls and reduce search space.

The algorithm achieves a worst-case time complexity of  $O(dn^{3d/3})$ . This is derived as follows:

- The maximum number of maximal cliques in a  $d$ -degenerate graph is bounded by  $(n-d)^{3d/3}$ .
- The algorithm ensures that recursive calls are limited by the degeneracy ordering, leading to  $O(dn^{3d/3})$  runtime.
- The space complexity remains  $O(n+m)$  due to the adjacency list representation and recursive stack depth being at most  $O(n)$ .

This complexity is nearly optimal, as it matches the worst-case output size up to a constant factor.

## Paper 3: Arboricity & Subgraph Listing Algorithms

Initialization:

The graph vertices are numbered in ascending order of their degrees.

Two helper arrays:

$S[y] \rightarrow$  Tracks neighbors of vertices not in the clique.

$T[y] \rightarrow$  Tracks neighbors of vertices in the clique.

The algorithm starts with the first vertex and recursively explores cliques.

Expansion and Pruning:

At each step, the algorithm adds vertex  $i$  to the current clique  $CCC$ .

Pruning techniques:

- $S \rightarrow$  Tracks candidates that can still be added.
- $T \rightarrow$  Tracks finished vertices to avoid duplicates.

If  $S(y)=0$ , the vertex is removed, reducing unnecessary expansions.

Maximality Test:

- Before printing, the algorithm verifies clique maximality by checking if adding any neighboring vertex forms a larger clique.
- If no valid expansion exists, the clique is confirmed as maximal and printed.
- This step prevents redundant or non-maximal cliques from being printed.

Lexicographic Ordering:

- Ensures only the lexicographically largest cliques are considered.
- Vertices are sorted and validated. Smaller vertices cause the clique to be flagged as invalid, avoiding redundant expansions.

Complexity Analysis

- Time Complexity:  $O(3^{\lfloor n/3 \rfloor})$  (worst-case)
- Space Complexity:  $O(n)$  for the recursion stack, plus  $O(n^2)$  for the  $S$  and  $T$  sets.

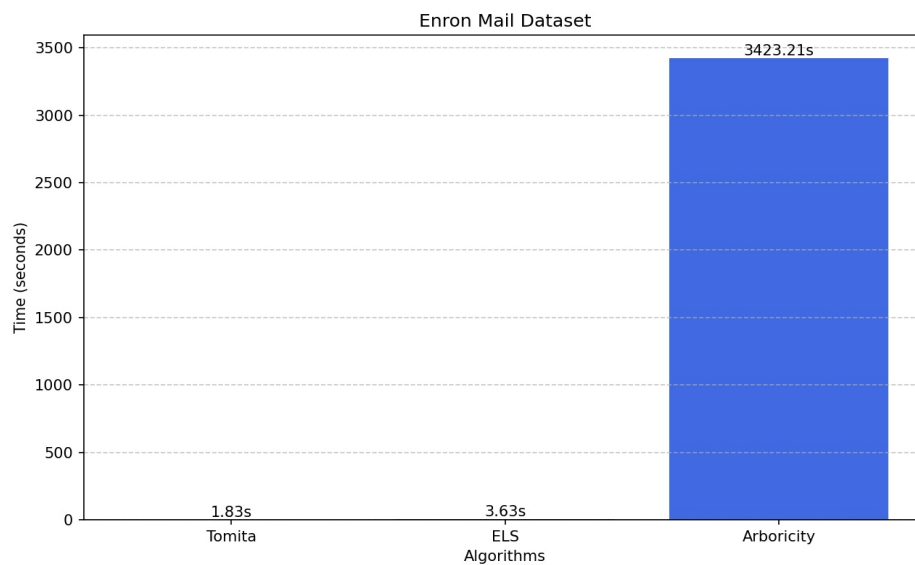
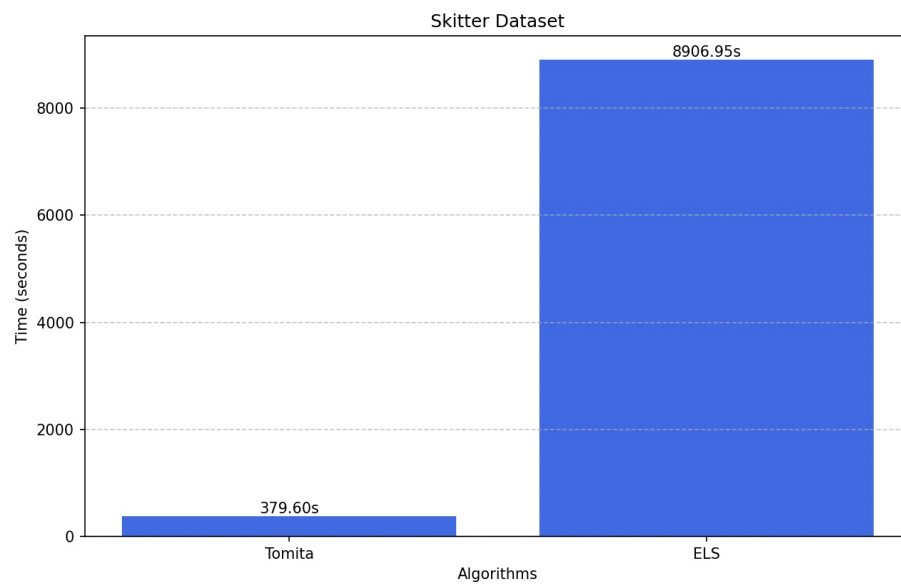
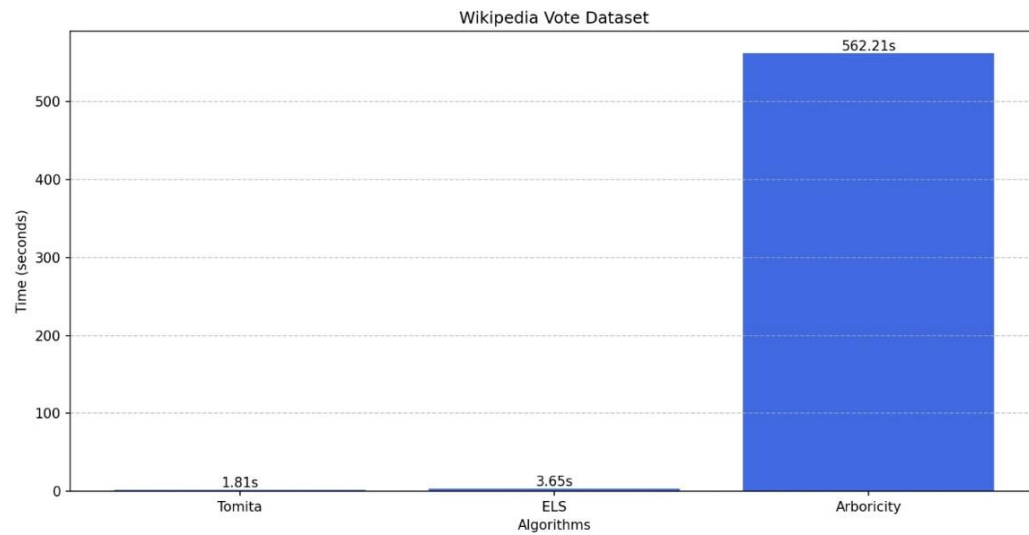
# Results:

	Enron email network	Wikipedia vote network	Autonomous systems by Skitter
Largest size of the clique in each dataset	20	17	67
Total number of maximal cliques in each dataset	226859	459002	37322355

## Execution Time of algorithms on each dataset

Algorithm	Enron email network	Wikipedia vote network	Autonomous systems by Skitter
Arboricity	3423.21 seconds	562.213 seconds	Ran for 8 hours got 146232 cliques
ELS	1.74138 seconds	1.19985 seconds	8906.95 seconds
Tomita	1.83627 seconds	1.81468 seconds	379.6 seconds

## Execution time of Algorithms on each dataset:



## Distribution of different size cliques:

Enron email network		Wikipedia vote network		Autonomous systems by Skitter	
Size	Cliques	Size	Cliques	Size	Cliques
2	14070	2	8655	2	2319807
3	7077	3	13718	3	3171609
4	13319	4	27292	4	1823321
5	18143	5	48416	5	939336
6	22715	6	68872	6	684873
7	25896	7	83266	7	598284
8	24766	8	76732	8	588889
9	22884	9	54456	9	608937
10	21393	10	35470	10	665661
11	17833	11	21736	11	728098
12	15181	12	11640	12	798073
13	11487	13	5449	13	877282
14	7417	14	2329	14	945194
15	3157	15	740	15	980831
16	1178	16	208	16	939987
17	286	17	23	17	839330
18	41			18	729601
19	10			19	639413
20	6			20	600192
				21	611976
				22	640890
				23	673924
				24	706753
				25	753633
				26	818353
				27	892719
				28	955212
				29	999860
				30	1034106
				31	1055653
				32	1017560
				33	946717
				34	878552
				35	809485
				36	744634
				37	663650
				38	583922
				39	520239
				40	474301
				41	420796
				42	367879
				43	321829
				44	275995
				45	222461
				46	158352
				47	99522
				48	62437
				49	39822
				50	30011
				51	25637
				52	17707
				53	9514
				54	3737
				55	2042
				56	1080
				57	546
				58	449
				59	447
				60	405
				61	283
				62	242
				63	146
				64	84
				65	49
				66	22
				67	4





*Thankyou*