Design & Analysis of Algorithms

Assignment - 1



BIRLA INSTITUTE OF TECHNOLOGY AND SCIENCE, PILANI

Sr No	Name	ID No	Contribution
1	Bhaskara Hemanth Karthikeya Ganti	2022A7PS0053H	Bron-Kerbosch Degeneracy Code and Project Report
2	Mohammad Jambughodawala	2022А7РЅ0009Н	Tomita Worst-Case Clique Code and Arboricity Clique Generation Code
3	Harsh Vikram Jajodia	2022A7PS0171H	Arboricity Clique Generation Code and Tomita Worst-Case Clique Code
4	Vaishnav Devaguptapu	2022A7PS0085H	Project Website and Bron- Kerbosch Degeneracy Code
5	Anshul Gopal	2022A7PS0009Н	Arboricity Clique Generation Code, Project Report and Readme file

Algorithms

Paper 1: The worst-case time complexity for generating all maximal cliques and computational experiments

The algorithm follows a recursive depth-first search strategy. The main steps are:

- 1. Initialization: The algorithm starts with an empty set Q, representing the current clique being expanded.
- 2. Expansion: At each step, a vertex q is added to Q, and a new subgraph is considered consisting of vertices adjacent to q.
- 3. Pruning: Two pruning techniques are applied:
 - Finished Set (FINI): Tracks vertices that have already been processed to avoid duplicate cliques.
 - Candidate Set (CAND): Keeps track of vertices that can still be added to Q.
- 4. Recursion: The process continues recursively, generating larger complete subgraphs until maximal cliques are found.
- 5. Output Optimization: Instead of storing all maximal cliques explicitly, a tree-like output format is used, reducing space complexity.

The key distinction of this algorithm is its pruning mechanism, which significantly reduces redundant computations compared to previous methods.

The worst-case complexity of the algorithm is derived based on the number of maximal cliques in an n-vertex graph. According to Moon and Moser's theorem, the maximum number of maximal cliques in a graph is $O(3^{(n/3)})$. In the context of space complexity, the algorithm primarily uses recursion and tree-based output, reducing the need to store all cliques explicitly. The space complexity is O(n) for the recursive stack in the worst case.

Paper 2: Listing All Maximal Cliques in Sparse Graphs in Near-optimal Time

This algorithm modifies the classic Bron-Kerbosch algorithm by incorporating:

- Degeneracy Ordering: The graph's vertices are processed in order of degeneracy, ensuring that each recursive call has at most d candidates.
- 2. Pivoting Strategy: A pivot is selected to minimize recursive calls, following the method of Tomita et al.
- 3. Efficient Pruning: The algorithm exploits the degeneracy ordering to restrict search space, reducing redundant computations.
- 4. Graph Data Structure: Uses an adjacency list representation for efficient neighborhood queries.

The algorithm follows these steps:

- Compute a degeneracy ordering of the vertices.
- Process each vertex v in this order, using a modified Bron-Kerbosch approach to find maximal cliques containing v.
- Use pivoting to minimize recursive calls and reduce search space.

The algorithm achieves a worst-case time complexity of O(dn3d/3). This is derived as follows:

- The maximum number of maximal cliques in a d-degenerate graph is bounded by (n-d)3d/3.
- The algorithm ensures that recursive calls are limited by the degeneracy ordering, leading to O(dn3d/3) runtime.
- The space complexity remains O(n+m) due to the adjacency list representation and recursive stack depth being at most O(n).

This complexity is nearly optimal, as it matches the worst-case output size up to a constant factor.

Paper 3: Arboricity & Subgraph Listing Algorithms

1. Initialization:

The graph vertices are numbered in ascending order of their degrees.

Two helper arrays:

 $S[y] \rightarrow Tracks$ neighbors of vertices not in the clique.

 $T[y] \rightarrow Tracks$ neighbors of vertices in the clique.

The algorithm starts with the first vertex and recursively explores cliques.

2. Expansion and Pruning:

At each step, the algorithm adds vertex iii to the current clique CCC.

Pruning techniques:

- S → Tracks candidates that can still be added.
- T→ Tracks finished vertices to avoid duplicates.

If S(y)=0, the vertex is removed, reducing unnecessary expansions.

3. Maximality Test:

- Before printing, the algorithm verifies clique maximality by checking if adding any neighboring vertex forms a larger clique.
- If no valid expansion exists, the clique is confirmed as maximal and printed.
- This step prevents redundant or non-maximal cliques from being printed.

4. Lexicographic Ordering:

- Ensures only the lexicographically largest cliques are considered.
- Vertices are sorted and validated. Smaller vertices cause the clique to be flagged as invalid, avoiding redundant expansions.

Complexity Analysis

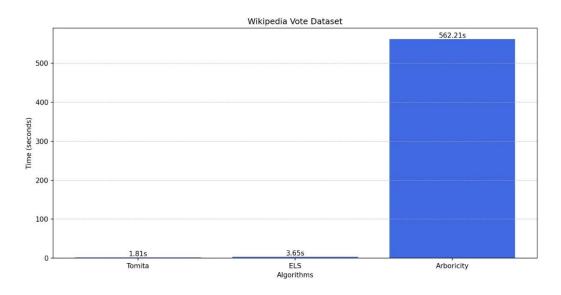
- Time Complexity: 0(3^{n/3}) (worst-case)
- Space Complexity: O(n) for the recursion stack, plus O(n^2) for the S and T sets.

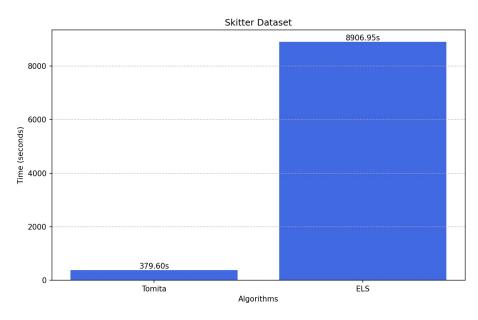
Results:

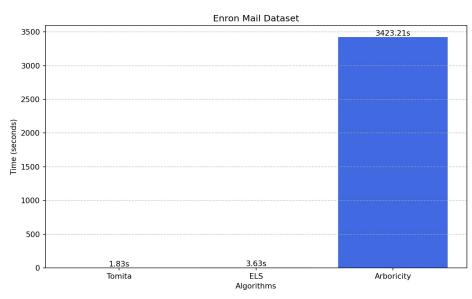
	Enron email network	Wikipedia vote network	Autonomous systems by Skitter
Number of nodes	36692	7115	1696415
Number of edges	183831	103689	11095298
Largest size of the clique in each dataset	20	17	67
Total number of maximal cliques in each dataset	226859	459002	37322355

Execution Time of algorithms on each dataset			
Algorithm	Enron email network	Wikipedia vote network	Autonomous systems by Skitter
Arboricity	3423.21 seconds	562.213 seconds	Ran for 8 hours got
ELS	1.74138 seconds	1.19985 seconds	8906.95 seconds
Tomita	1.83627 seconds	1.81468 seconds	379.6 seconds

Execution time of Algorithms on each dataset:







Distribution of different size cliques:

Enron email network		
Size	Cliques	
2	14070	
3	7077	
4	13319	
5	18143	
6	22715	
7	25896	
8	24766	
9	22884	
10	21393	
11	17833	
12	15181	
13	11487	
14	7417	
15	3157	
16	1178	
17	286	
18	41	
19	10	
20	6	

Size Cliques 2 8655 3 13718 4 27292 5 48416 6 68872 7 83266 8 76732 9 54456 10 35470 11 21736 12 11640 13 5449 14 2329 15 740 16 208 17 23		
3 13718 4 27292 5 48416 6 68872 7 83266 8 76732 9 54456 10 35470 11 21736 12 11640 13 5449 14 2329 15 740 16 208	Size	Cliques
4 27292 5 48416 6 68872 7 83266 8 76732 9 54456 10 35470 11 21736 12 11640 13 5449 14 2329 15 740 16 208	2	8655
5 48416 6 68872 7 83266 8 76732 9 54456 10 35470 11 21736 12 11640 13 5449 14 2329 15 740 16 208	3	13718
6 68872 7 83266 8 76732 9 54456 10 35470 11 21736 12 11640 13 5449 14 2329 15 740 16 208	4	27292
7 83266 8 76732 9 54456 10 35470 11 21736 12 11640 13 5449 14 2329 15 740 16 208	5	48416
8 76732 9 54456 10 35470 11 21736 12 11640 13 5449 14 2329 15 740 16 208	6	68872
9 54456 10 35470 11 21736 12 11640 13 5449 14 2329 15 740 16 208	7	83266
10 35470 11 21736 12 11640 13 5449 14 2329 15 740 16 208	8	76732
11 21736 12 11640 13 5449 14 2329 15 740 16 208	9	54456
12 11640 13 5449 14 2329 15 740 16 208	10	35470
13 5449 14 2329 15 740 16 208	11	21736
14 2329 15 740 16 208	12	11640
15 740 16 208	13	5449
16 208	14	2329
	15	740
17 23	16	208
17 23	17	23

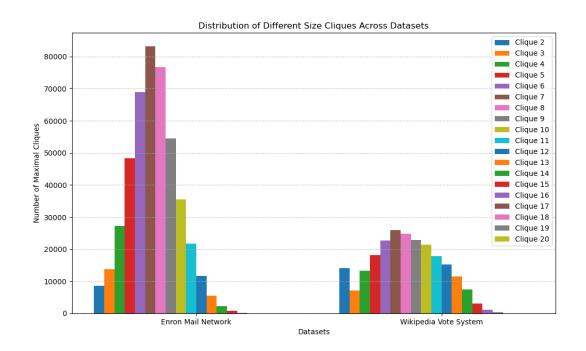
Wikipedia vote network

Size	Cliques	
2	2319807	
3	3171609	
4	1823321	
5	939336	
6	684873	
7	598284	
8	588889	
9	608937	
10	665661	
11	728098	
12	798073	
13	877282	
14	945194	
15	980831	
16	939987	
17	839330	
18	729601	
19	639413	
20	600192	
21	611976	
22	640890	
23	673924	
24	706753	
25	753633	
26	818353	
27	892719	
28	955212	
29	999860	
30	1034106	
31	1055653	
32	1017560	
33	946717	
34	878552	
35 36	809485 744634	
37	663650	
38	583922 520239	
40		
41	474301	
42	420796	
	367879	
43	321829	
44	275995	
45	222461	
	158352	
47	99522 62437	
49	39822	
50 51	30011 25637	
52	17707	
53	9514	
54	3737	
55	2042	
56	1080	
57	546	
58	449	
59	447	
60	405	
	283	
61		
	242	
63	146 84	
65	49	
66	22	
01	4	

Autonomous systems by Skitter

Histogram Representation:

Enron Email Network & Wikipedia Vote System:



Autonomous systems by Skitter:

