

# Part I Syllabus - Fundamental Underlying Layers

Date	Subject	File
Week 1: 9/Jan/2023 11/Jan/2023	Introduction: course logistics and Internet history	M1-L1-Introduction.pptx
	Layered Network Architecture	First part of M1-L2-Network Layer & Physical Resilience.pptx
Week 2: 16/Jan/2023 18/Jan/2023	Physical Layer: Network Resilience	Second part of M1-L2-Network Layer & Physical Resilience.pptx
	Data link layer – Flow control	M1-L3-DLL-Flow Control.pptx
Week 3: 25/Jan/2023	Data link layer – Error control	M1-L4-DLL-Error Control.pptx
Week 4: 30/Jan/2023 01/Feb/2023	Local area network – Introduction	M1-L5-LAN-Introduction.pptx
	Local area network – MAC	M1-L6-LAN-MAC.pptx
Week 5: 06/Feb/2023 08/Feb/2023	Local area network – Ethernet	First part of M1-L7-LAN-Ethernet.pptx
	Local area network – Ethernet Evolutions	Second part of M1-L7-LAN-Ethernet.pptx
Week 6: 13/Feb/2023 15/Feb/2023	Local area network – WLAN	M1-L8-LAN-WLAN.pptx
	Mobile Access Networks	M1-L9-Mobile.pptx
Week 7: 20/Feb/2023 22/Feb/2023	E-learning for Network paradigms	M1-L10-Paradigms.pptx
	Network paradigms	M1-L10-Paradigms.pptx

# Additional Materials

- The related content talked today in [https://eclass.teicrete.gr/modules/document/file.php/TP326/%CE%98%CE%B5%CF%89%CF%81%CE%AF%CE%B1%20\(Lectures\)/Computer\\_Networking\\_A\\_Top-Down\\_Approach.pdf](https://eclass.teicrete.gr/modules/document/file.php/TP326/%CE%98%CE%B5%CF%89%CF%81%CE%AF%CE%B1%20(Lectures)/Computer_Networking_A_Top-Down_Approach.pdf) is as follow:
  - Physical Media: Page 18 - Page 22
  - MAC Protocols: Page 445 - Page 449
  - ALOHA Protocols: Page 450 - Page 453
  - CSMA and CSMA/CD Protocols: Page 453- Page 459
- You can also find other video materials about
  - Physical Layer <https://www.youtube.com/watch?v=rKzDbdGhcdY>
  - CSMA <https://www.youtube.com/watch?v=MAZi6VoekYw>
  - Controlled Access Protocol-Reservation <https://www.youtube.com/watch?v=baaPXiQ44vs>

# What is the problem with the guy?



# *SC2008/CZ3006/CE3005*

# *Computer Network*

## Lecture 8

### Wireless LAN: IEEE 802.11



# Contents

- **WLAN Overview**
  - WLAN Standard
  - WLAN Architecture
  - WLAN Protocol Stack
- **802.11 Physical Layer**
- **802.11 MAC Layer**
  - Hidden and Exposed Terminal Problems
  - CSMA/CA Protocol
  - MAC Management
- **Multi-Access Reservation Protocol**
  - Scheme
  - Throughput Calculation

# WLAN Overview

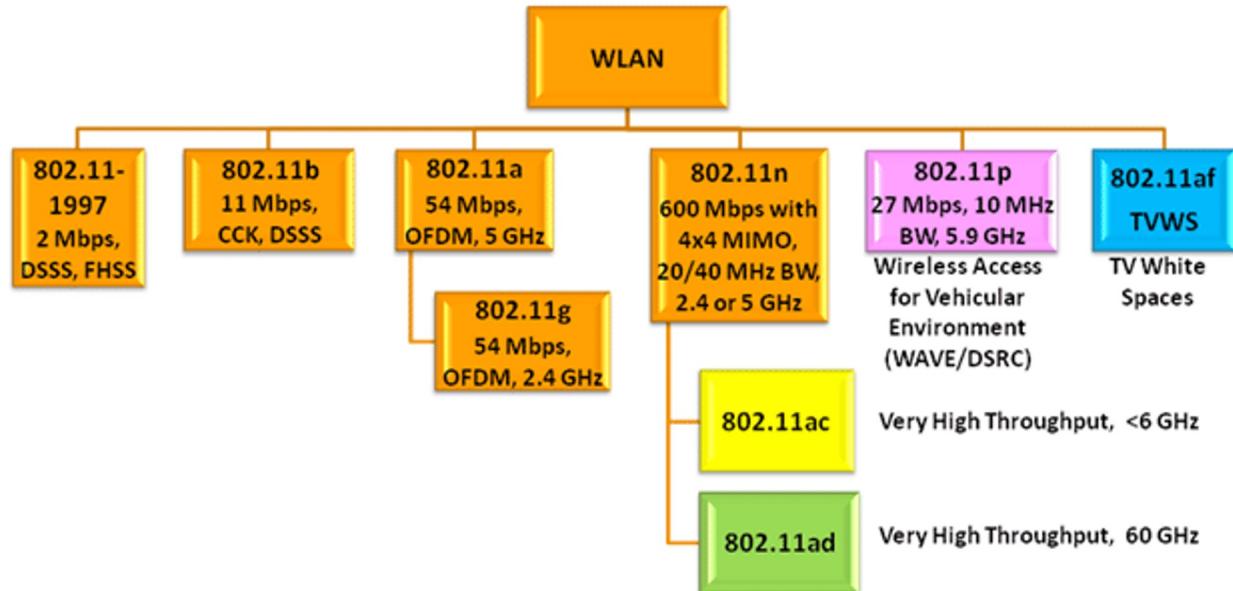
# LAN/WLAN World

- **LANs provide connectivity for interconnecting computing resources at local levels of an organization**
- **Wired LANs**
  - Limitations because of physical, hard-wired infrastructure
- **Wireless LANs**
  - Flexibility
  - Portability
  - Mobility
  - Ease of Installation

# IEEE 802.11 WLAN Standard

not examinable

- In response to lacking standards, IEEE developed the first internationally recognized wireless LAN standard – IEEE 802.11
- IEEE published 802.11 in 1997, after seven years of work
- Most prominent specification for WLANs
- Scope of IEEE 802.11 is limited to Physical and Data Link Layers

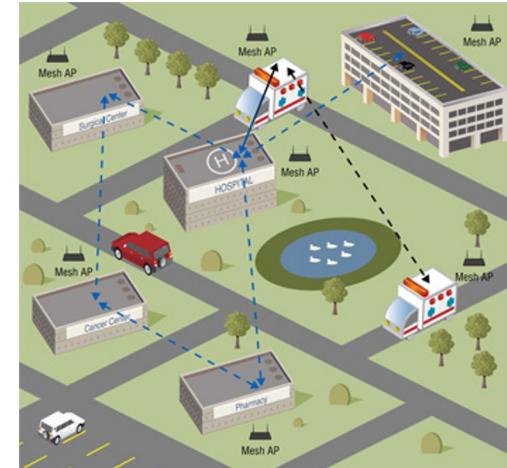


DSRC = Dedicated Short-Range Communications

# Wireless LANs: Characteristics

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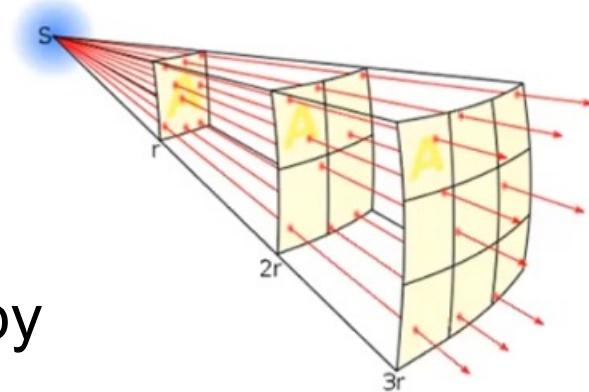
- **Advantages**
  - Flexible deployment
  - Minimal wiring difficulties
  - More robust against disasters (earthquake, etc)
  - Historic buildings, conferences, trade shows,...
- **Disadvantages**
  - Low bandwidth (1-10 Mbit/s) compared to wired networks
  - Proprietary solutions
  - Need to follow wireless spectrum regulations



# Wireless Link Characteristics

not examinable

- Different from wired link ...
    - Decreased signal strength: radio signal attenuates as it propagates through matter (path loss)
    - Interference from other sources: standardized wireless network frequencies (e.g., 2.4 GHz) shared by other devices (e.g., phone); devices interfere as well
    - Multipath propagation: radio signal reflects off objects ground, arriving at destination at slightly different times
- ... make communication over wireless link much more “difficult”



# WLAN Architecture

not examinable

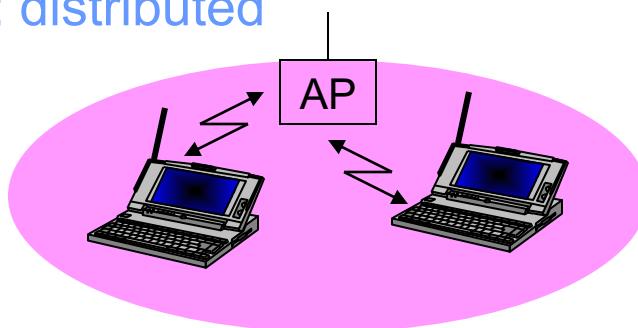
- **Building Modules**

- Station (STA)
  - Mobile node
  - Smartphone, pad, laptop
- Access Point (AP)
  - Stations are connected to access points.

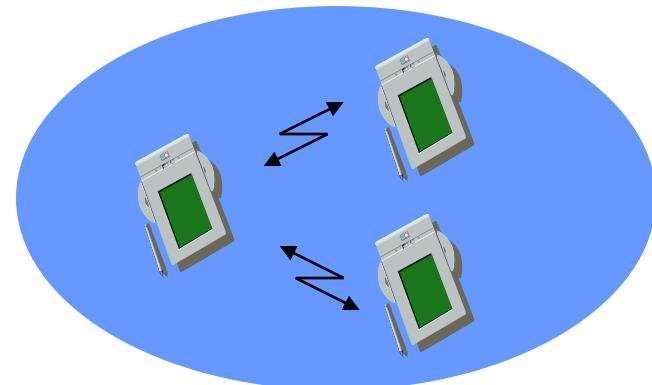


- **Two Architectural Modes**

- Infrastructure: centralized
- Ad Hoc: distributed



Infrastructure



Ad Hoc

# (Extended) Service Set

not examinable

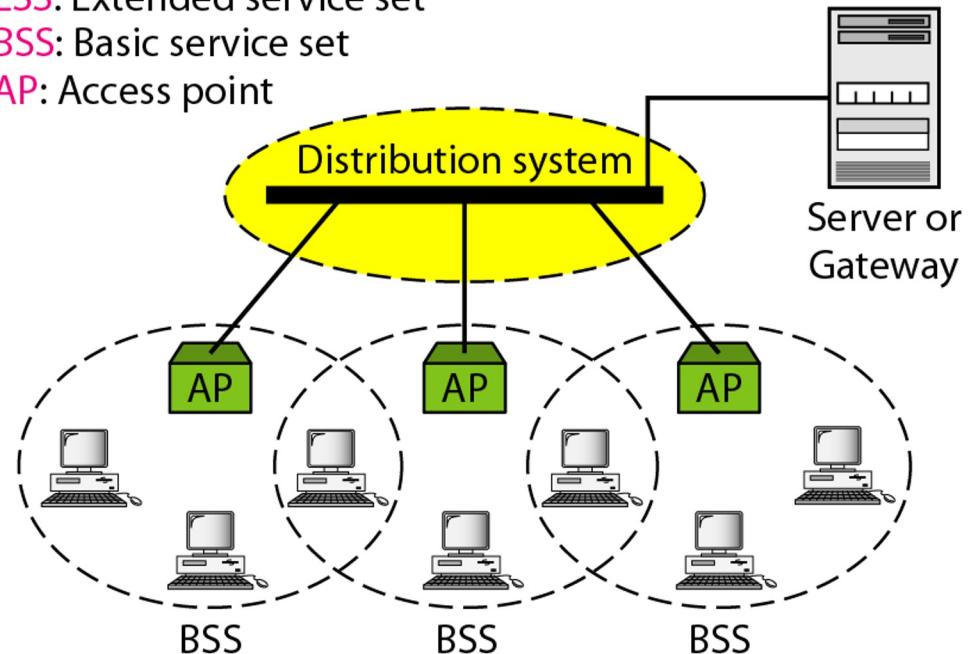
- **Basic Service Set (BSS)**
  - Stations and the AP within the same radio coverage form a BSS.
- **Extended Service Set (ESS)**
  - Several BSSs connected through APs form an ESS.



**ESS:** Extended service set

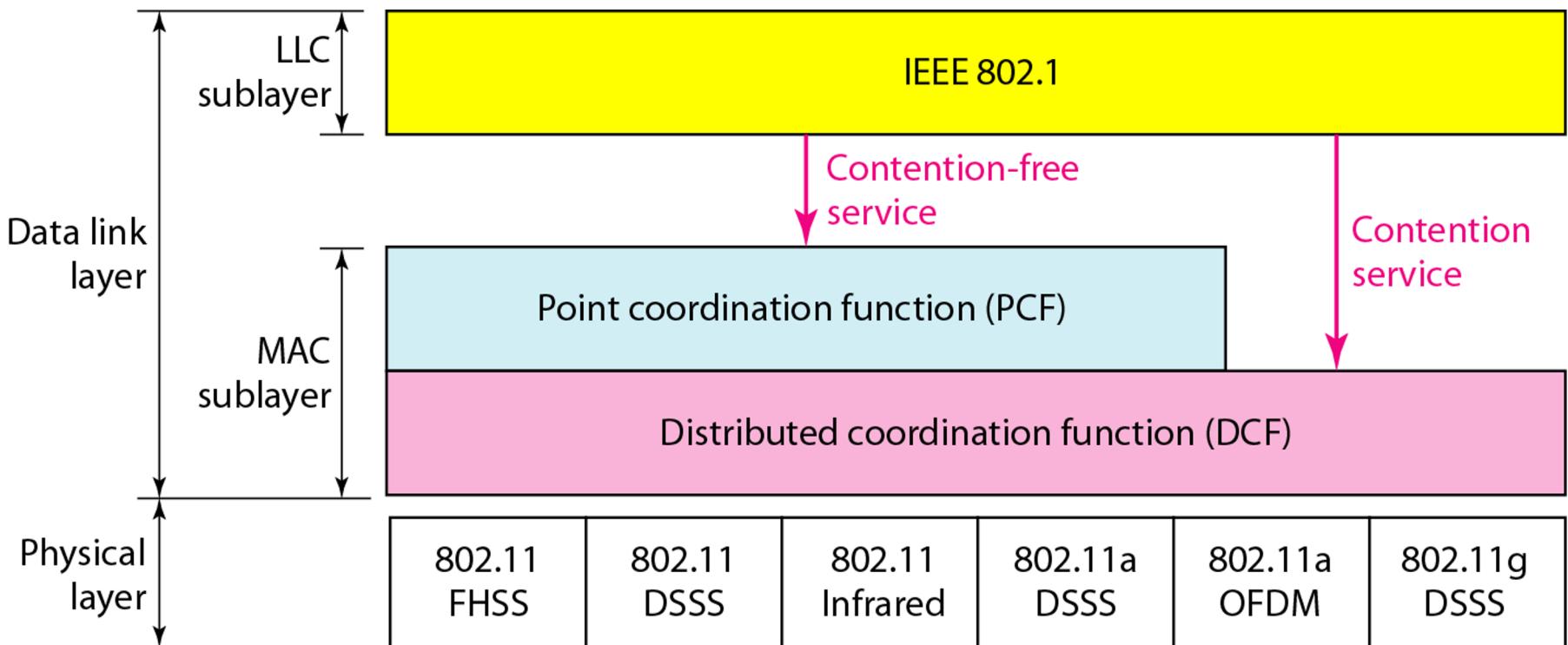
**BSS:** Basic service set

**AP:** Access point



# 802.11 Protocol Stack

not examinable



not examinable

# Wireless Physical Layer

# Radio Spectrum

not examinable

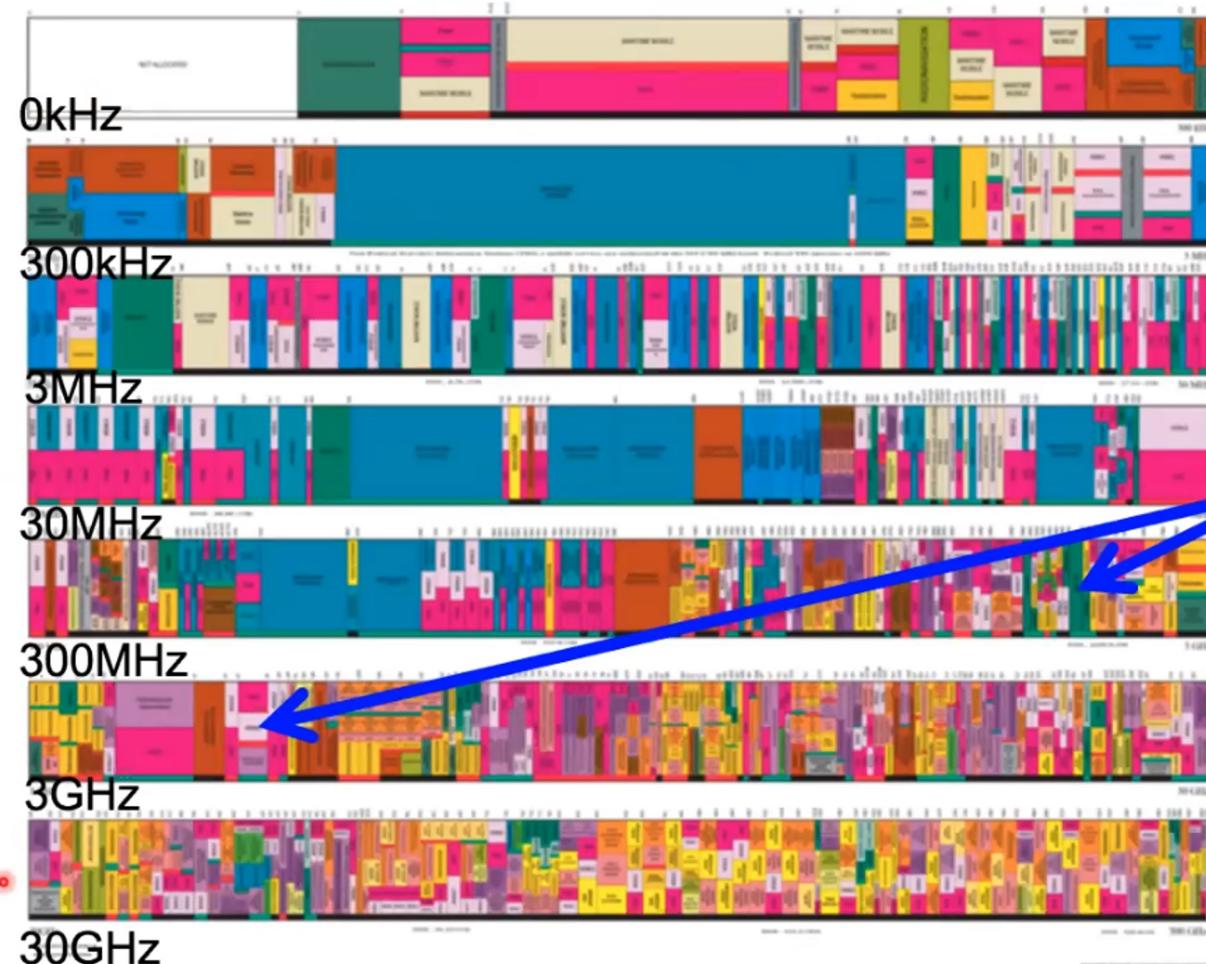
- **Radio Frequency bands are allocated to different applications**
  - The use of most frequency bands needs licenses
  - IEEE 802.11 uses industrial, scientific and medical (ISM) bands that don't require licenses if the radio transmissions follow the national/global regulations

# Sub-THz Radio Spectrum

## **not examinable**

**UNITED  
STATES  
FREQUENCY  
ALLOCATIONS**

## THE RADIO SPECTRUM



**802.11  
2.4GHz  
& 5GHz**



# IEEE 802.11 Physical Layer

not examinable

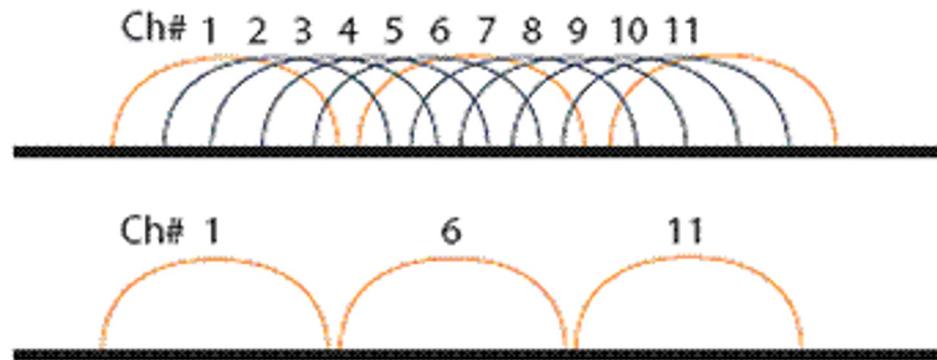
	802.11b	802.11g	802.11a	802.11n	
Frequency Band	2.4GHz	5GHz	2.4GHz	2.4	5
Non-overlapping Channels	3	3	12	3	12
Baseline BW Per Channel	11Mbps	54Mbps	54Mbps	65	65
Max BW Per Channel	11Mbps	54Mbps	54Mbps	130	270
MIMO	1	1	1	4	4
Modulation	DSSS	DSSS/OFDM	OFDM	OFDM	

# IEEE 802.11 Channels, Association

not examinable

- **802.11b:** 2.4GHz-2.485GHz spectrum divided into **11 channels** at different frequencies
  - AP admin chooses frequency for AP
  - Interference possible: channel can be same as that chosen by neighboring AP!

802.11b/g Operating Channels



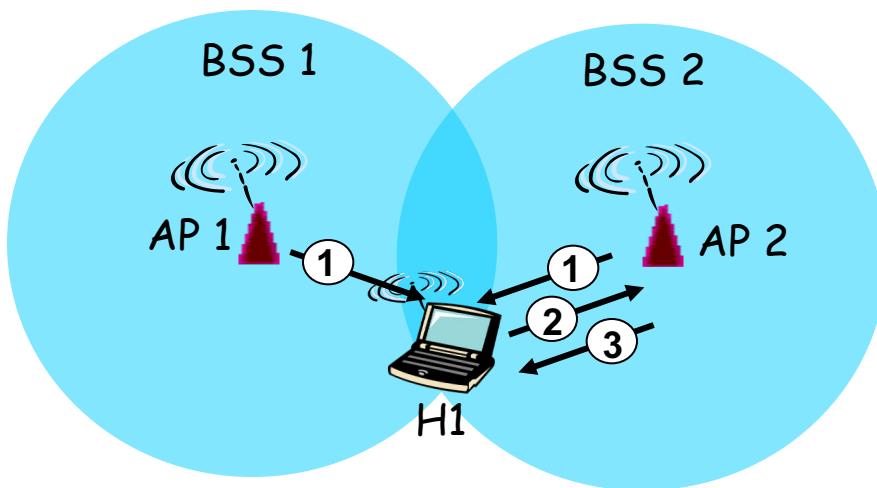
# IEEE 802.11 Channels, Association

not examinable

- **802.11b: 2.4GHz-2.485GHz spectrum divided into 11 channels at different frequencies**
  - AP admin chooses frequency for AP
  - Interference possible: channel can be same as that chosen by neighboring AP!
- **Host: must associate with an AP**
  - Scans channels, listening for beacon frames containing AP's name (SSID) and MAC address
  - Selects AP to associate with
  - May perform authentication (security purpose)
  - Will run DHCP to get IP address in AP's subnet

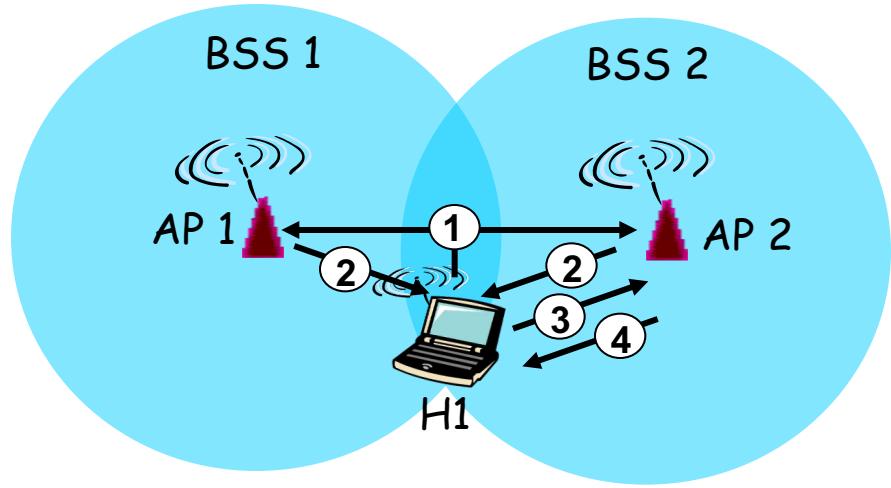
# 802.11 Passive/Active Scanning

not examinable



## Passive Scanning:

- (1) beacon frames sent from APs
- (2) association Request frame sent:  
H1 to selected AP
- (3) association Response frame sent:  
Selected AP to H1



## Active Scanning:

- (1) Probe Request frame broadcast  
from H1
- (2) Probes response frame sent from  
APs
- (3) Association Request frame sent:  
H1 to selected AP
- (4) Association Response frame  
sent: selected AP to H1

**not examinable**

# 802.11 MAC

# 802.11 MAC Sublayer

not examinable

- **New challenges caused by the nature of wireless communications**
  - Broadcast
  - Signal attenuation
  - Pervasive electromagnetic noise
- **Three functional areas**
  - Access control (random access vs controlled access)
  - Reliable data delivery (against noises and collisions)
  - Security (authentication, packet injection, ...)
- **Two additional problems:**
  - Hidden Terminal Problem
  - Exposed Terminal Problem

# Access Control

not examinable

- **Distributed Coordination Function (DCF)**
  - Distributed access protocol
  - Contention-based
  - Makes use of CSMA/CA
  - Suited for ad-hoc network and asynchronous traffic
- **Point Coordination Function (PCF)**
  - Alternative access method on top of DCF
  - Centralized access protocol
  - Contention-free, and works like polling
  - Suited for time-bound services like voice and multimedia

# Reliable Data Delivery

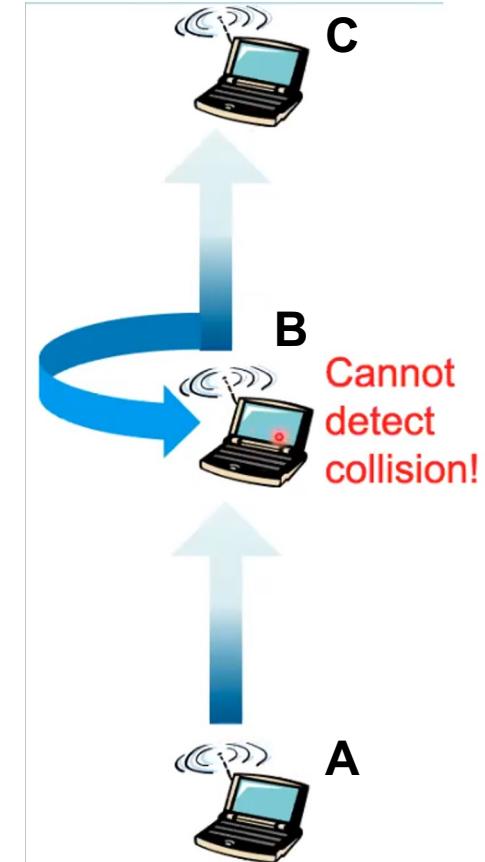
not examinable

- **Loss of frames due to noise, interference and propagation effects**
- **Frame exchange protocol**
  - Sender broadcasts data
  - Receiver responds with acknowledgement (ACK)
  - If sender does not receive ACK, it retransmits frame
- **Four frame exchange for enhanced reliability**
  - Sender issues request-to-send (RTS)
  - Receiver responds with clear-to-send (CTS)
  - Sender transmits data
  - Receiver responds with ACK

# 802.11 Multi-Access

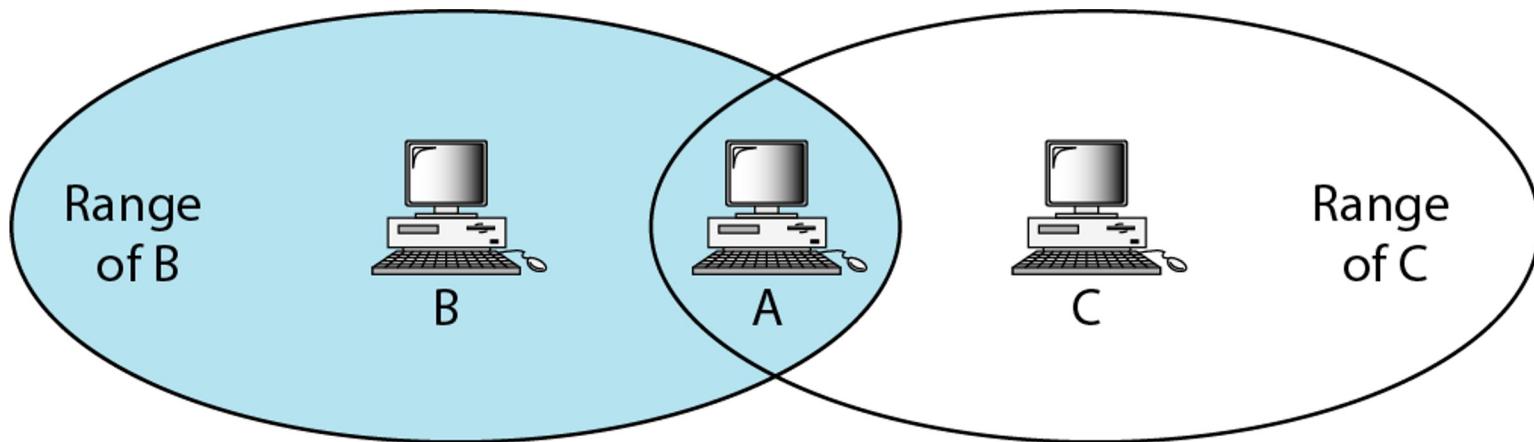
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- **Collision**
  - A receiver hears transmissions from 2<sup>+</sup> nodes at the same time
- **802.11: CSMA - sense before transmitting**
  - Don't collide with ongoing transmission by other node
- **802.11: no collision detection!**
  - Difficult to receive (sense collisions) when transmitting due to weak received signals (fading)
  - Can't sense all carriers & collisions in any case: hidden terminal problem
- **802.11: avoid collisions**
  - CSMA/C(ollision)A(voidance)



# Hidden Terminal Problem

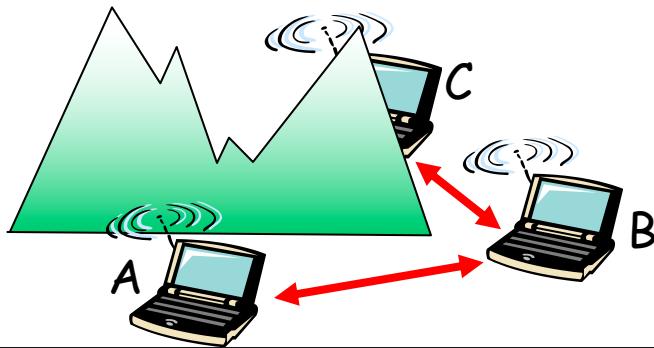
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B and C are hidden from each other with respect to A.

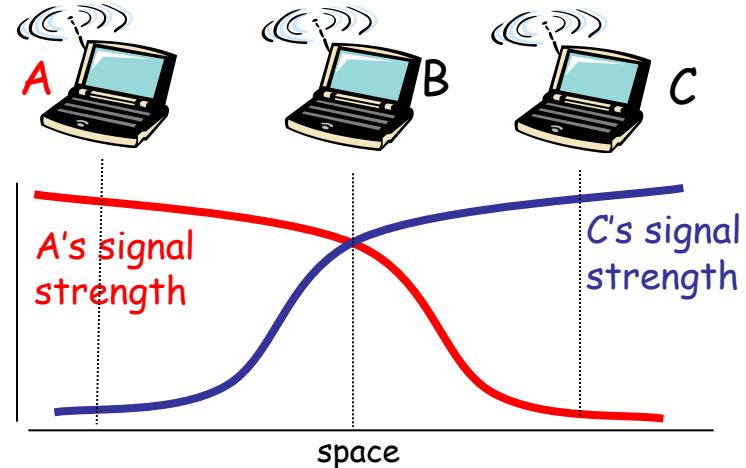
# Hidden Terminals in Wireless LAN

not examinable



## Case 1: Caused by barrier

- B, A hear each other
- B, C hear each other
- A, C can not hear each other
  - A, C unaware of their interference at B
  - A is a hidden terminal to C, vice versa

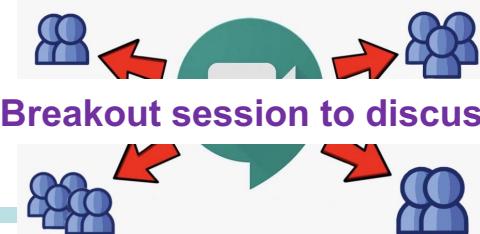


## Case 2: Caused by signal attenuation

- B, A hear each other
- B, C hear each other
- A, C can not hear each other
  - A, C unaware of their interference at B

**Question:** Does Ethernet have hidden terminal problem?

**Answer: No.** Ethernet does not have Case 1 and Case 2 above.



Breakout session to discuss

# Collision Avoidance

not examinable

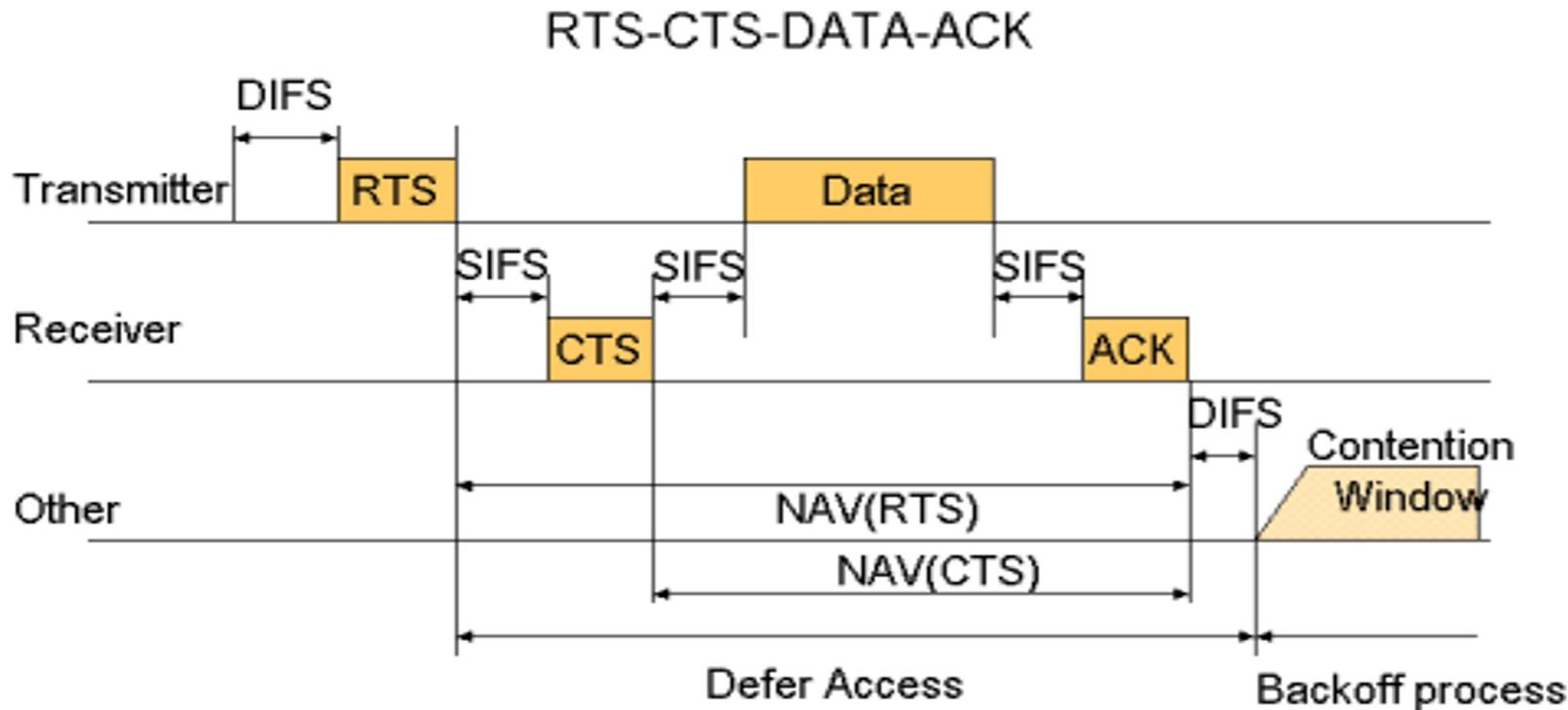
**idea:** Sender to “reserve” channel for a long data frame

- Sender first transmits a *small* **request-to-send (RTS)** packet to receiver using CSMA
  - RTSs may still collide with each other, or an RTS may collide with an ongoing data frame
  - but they’re short
- Receiver broadcasts **clear-to-send (CTS)** in response to RTS
- CTS heard by all nodes
  - Sender transmits data frame
  - Other stations defer transmissions

Avoid data frame collisions completely  
using small reservation packets!

# RTS-CTS-DATA-ACK

not examinable



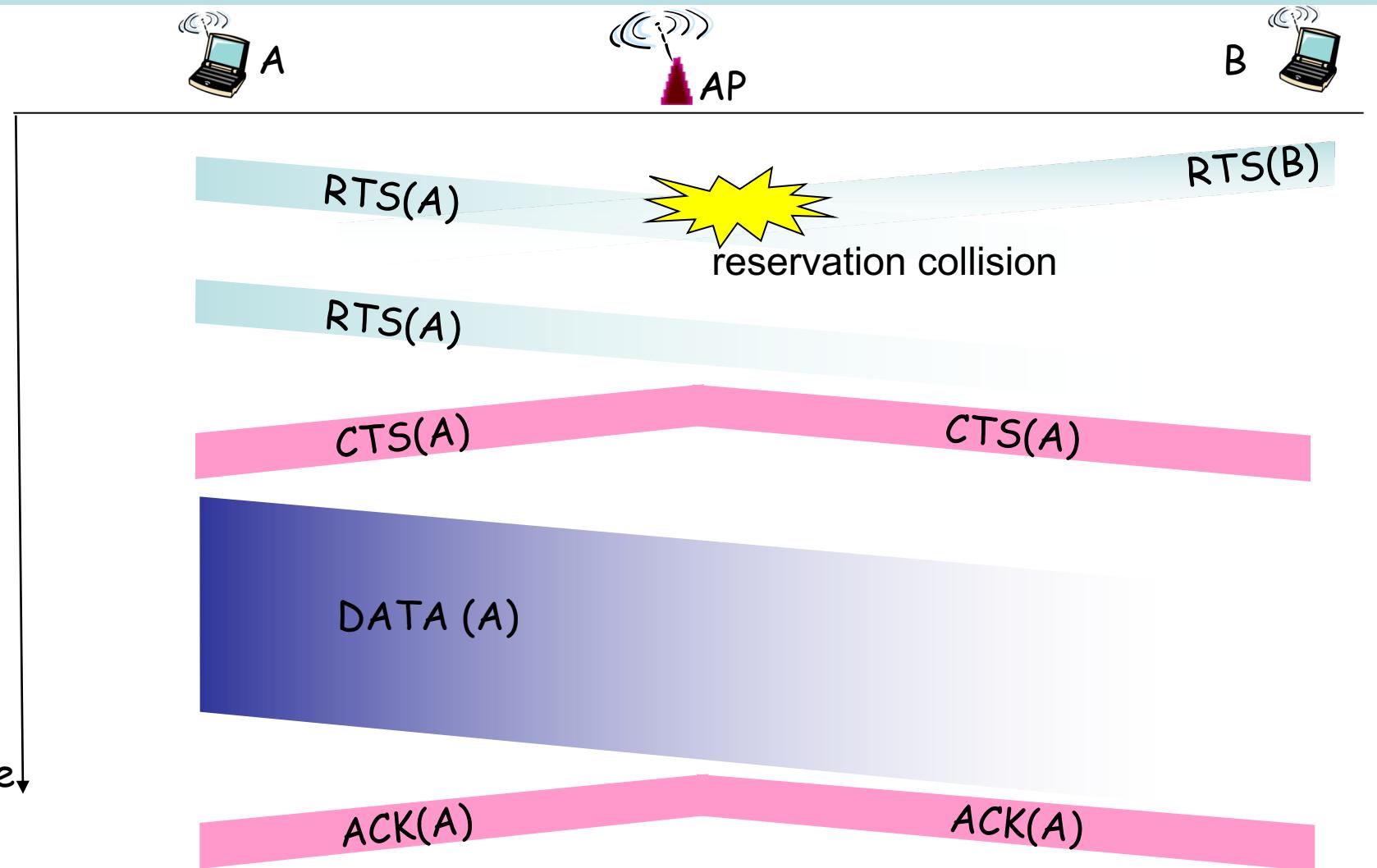
DIFS: Distributed IFS (Inter-frame Space)  
**for carrier sense**

RTS: Request-To-Send  
SIFS: Short IFS

CTS: Clear-To-Send  
ACK: Acknowledgement  
NAV: Network Allocation Vector

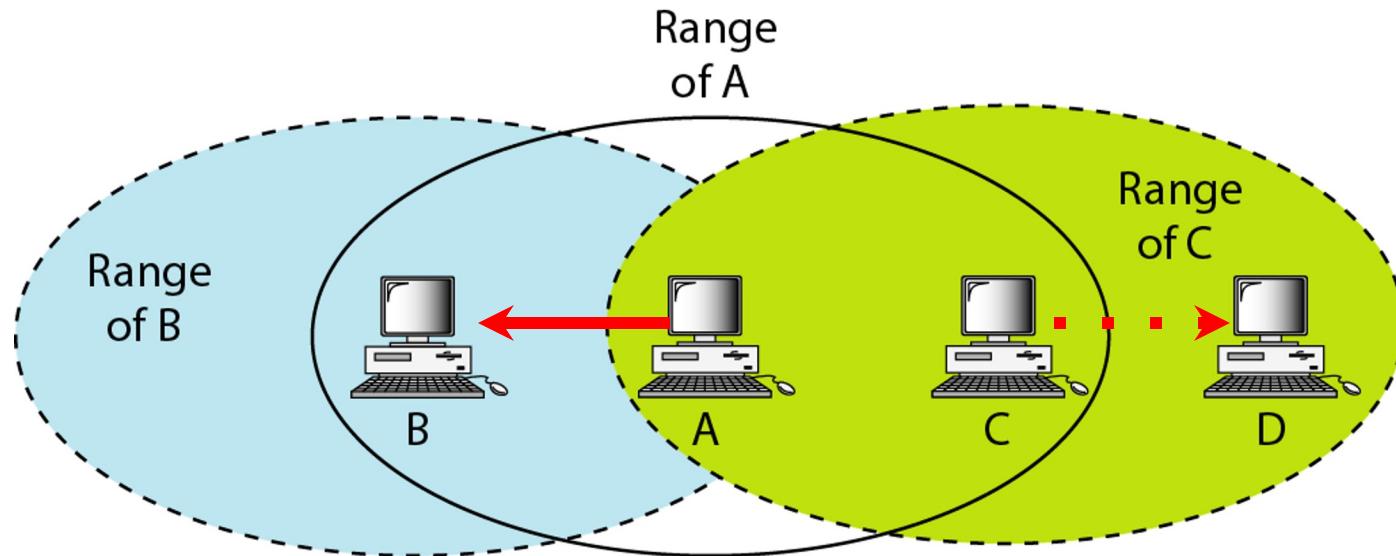
# Handshaking in Hidden Terminal Problem

not examinable



# Exposed Terminal Problem

not examinable

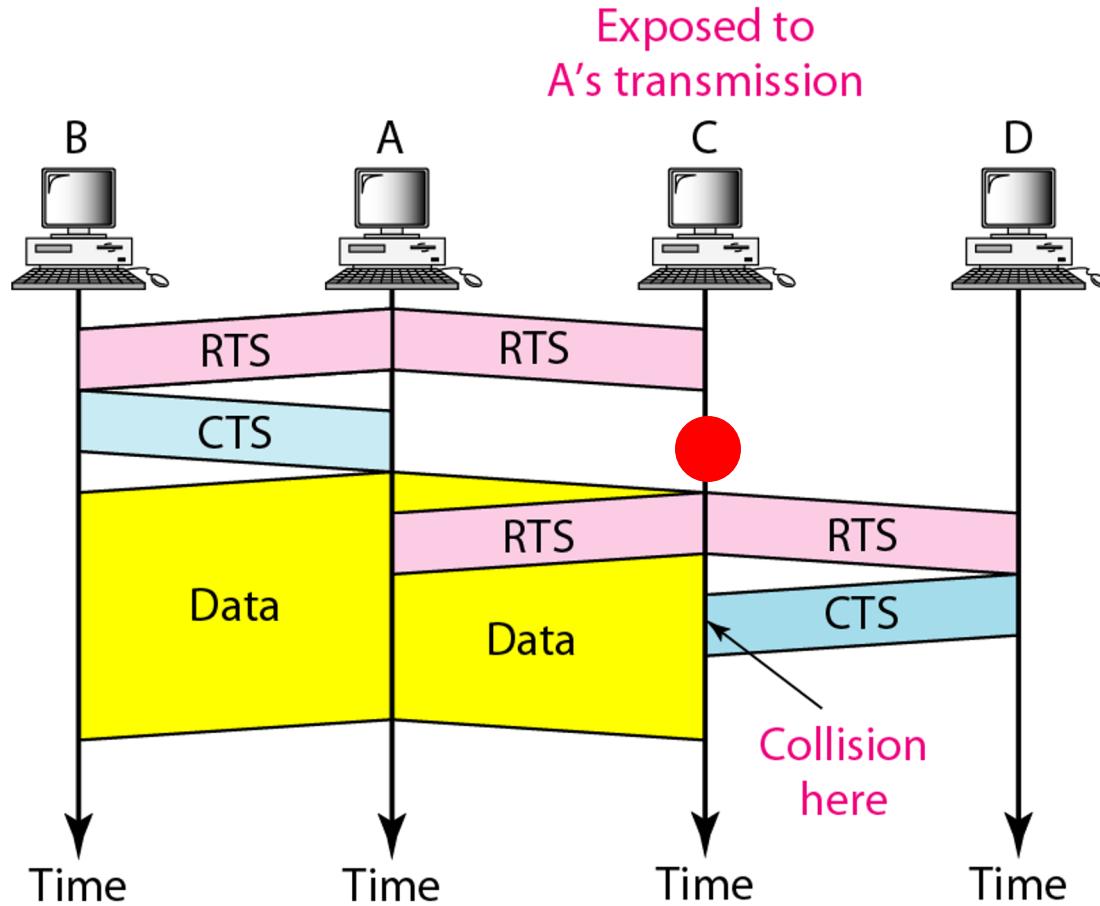


C is exposed to transmission from A to B.

# Handshaking in Exposed Terminal Problem

not examinable

- **RTS-CTS ensures no collision**
- **but doesn't solve the opportunity wasting problem**



# 802.11 Frame

not examinable

2	2	6	6	6	2	6	0 - 2312	4
frame control	duration	address 1	address 2	address 3	seq control	address 4	payload	CRC

Address 1: MAC address of wireless host or AP to receive this frame

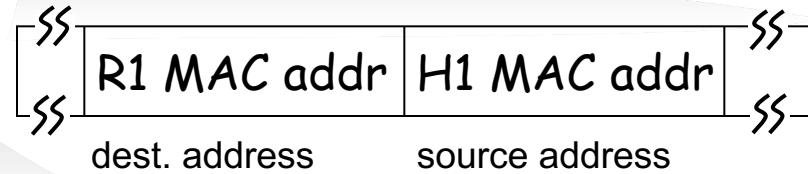
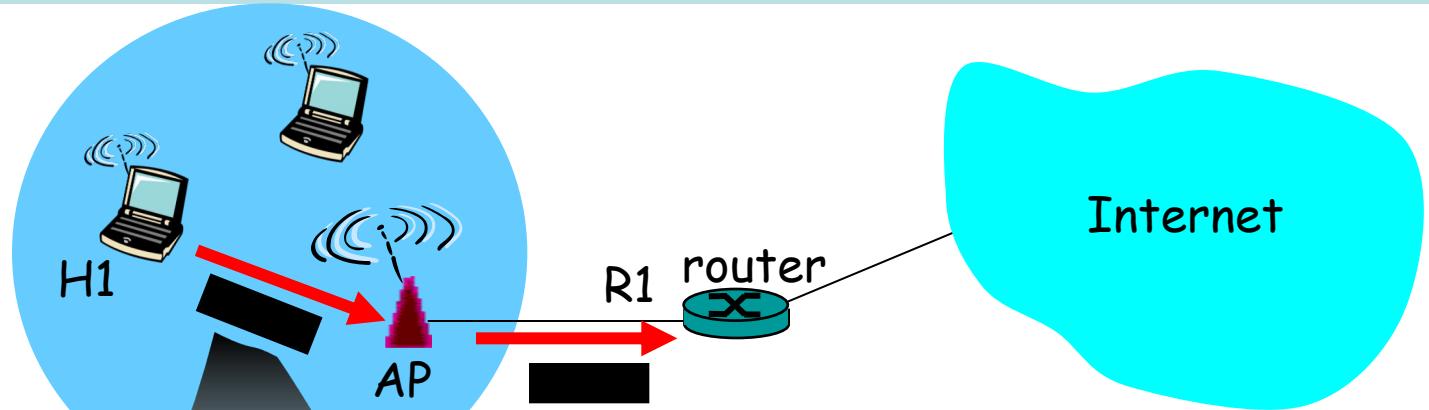
Address 2: MAC address of wireless host or AP transmitting this frame

Address 3: MAC address of router interface to which AP is attached

Address 4: used only in ad hoc mode

# 802.11 Addressing

not examinable



Ethernet frame



802.11 frame

# 802.11 Advanced Capabilities

not examinable

- **Synchronization**
  - finding and staying with a WLAN
  - synchronization functions
- **Power Management**
  - sleeping without missing any messages
  - power management functions
- **Roaming**
  - functions for joining a network
  - changing access points
  - scanning for access points
- **Management information base**

**Examinable**

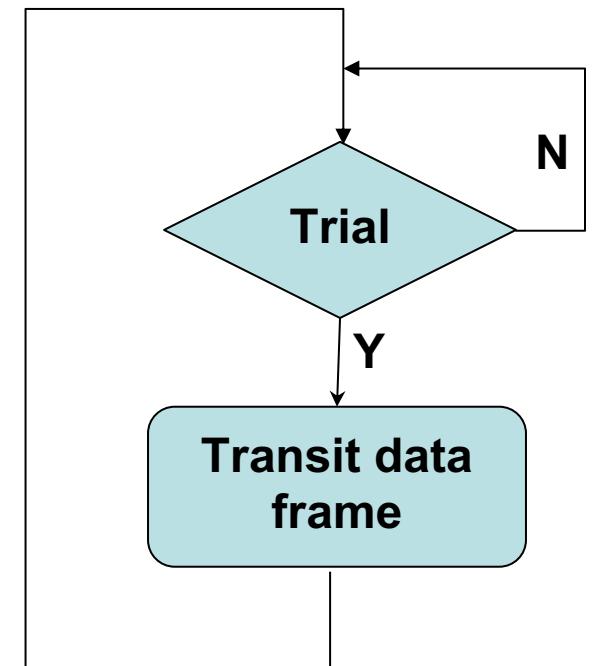
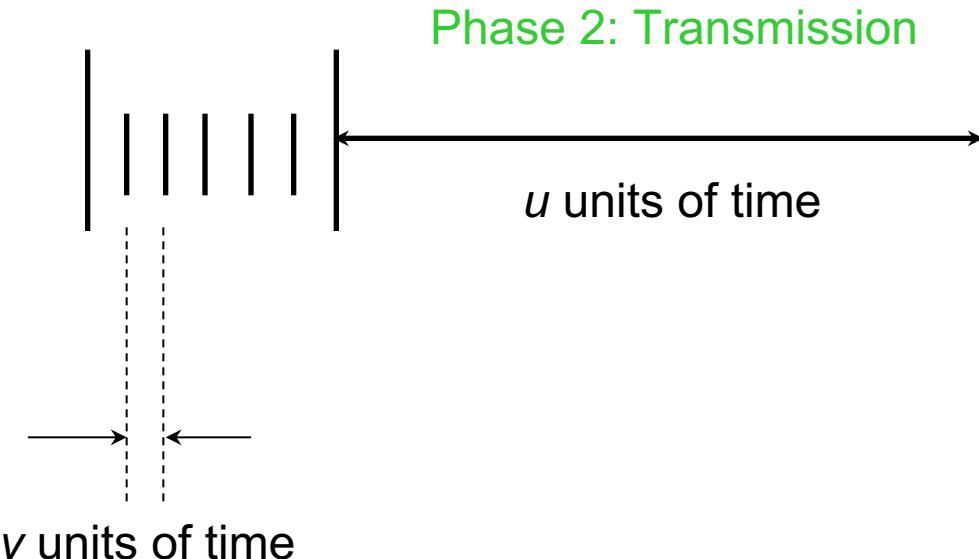
# **Multi-Access Reservation Protocol (MARP)**

# Multi-Access Reservation Protocol

- **Two-Phase Protocol**
  - Phase 1: Channel Reservation
  - Phase 2: Data Transmission

Examinable

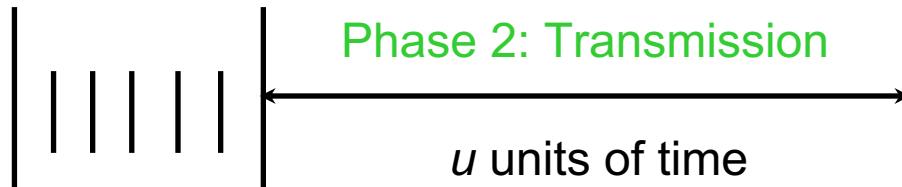
Phase 1: Reservation



# MARP Transmission Window

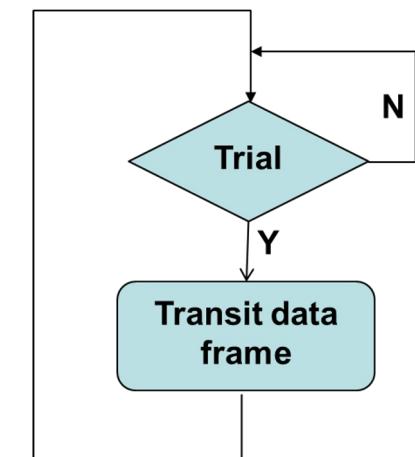
Phase 1: Reservation

Examinable



How many reservation trial frames?

- Assume that the channel utilization in reservation phase:  $S_r$
- Number of reservation trial frames to reserve the channel:  $X$ 
  - $X = 1$  (the first trial succeeds) with probability of  $S_r$
  - $X = k$  (the first  $k-1$  trials fail, the  $k^{\text{th}}$  trial succeeds) with probability of  $S_r(1-S_r)^{k-1}$
  - This is a geometric distribution, so  $E[X] = 1/S_r$
- The average transmission window is  $u + v/S_r$  units of time



# MARP Throughput

Examinable

$$\text{Throughput } S = \frac{\text{Time for message transmission}}{\text{Transmission window}}$$

Case	Message Length	Reservation Phase Length	Throughput
Reservation frame not used for message data bits	$u$	$v/S_r$	$S = \frac{u}{u + v/S_r}$
Reservation frame used for message data bits	$u$	$v/S_r$	$S = \frac{u}{(u - v) + v/S_r}$

# MARP Example

Examinable

Consider an experimental LAN using an MARP for data transmission. The protocol consists of two phases. In phase 1, it adopts some MAC protocol for transmission stations to reserve the channel. In phase 2, when one station reserves the channel, it transmits one frame. The length of reservation frame is 5ms, and the length of the data frame is 1s. No information bit is carried in the reservation frame. If the MAC protocol used in phase 1 has a utilization of 0.5, what is the throughput of the multi-access reservation protocol?

**CRACK Framework:**

**Context:** MARP with no data bits in reservation

**fRamwork:** the throughput of MARP is  $S=1/(1+v/S_r)$

**Apply:**  $v=5\text{ms}$ ,  $S_r=0.5$

**Calculation:**  $S = 1/(1+0.005/0.5)=1/1.01=0.99$

**checK:**  $S \leq 1$

# Local Area Network Summary

MAC Protocols		Transmission Protocol			Throughput/ Utilization	Note									
		Carrier Sensing	Frame Transmission	Collision Detection											
Aloha	Slotted	<ul style="list-style-type: none"> <li>None</li> </ul>	<ul style="list-style-type: none"> <li>Each transmits in a slot immediately with probability <math>p</math></li> </ul>	<ul style="list-style-type: none"> <li>When a collision is detected, the colliding frames are transmitted up to their last bits.</li> </ul>	$S = Np(1 - p)^{(N-1)} = Ge^{-G}$	Number of Stations: $N$ Probability of Attempt: $p$ Attempt Rate: $G = Np$									
	Pure		<ul style="list-style-type: none"> <li>Each transmits immediately with probability <math>p</math></li> </ul>		$S = Np(1 - p)^{2(N-1)} = Ge^{-2G}$										
CSMA	Non-Persistent	<ul style="list-style-type: none"> <li>Must sense channel before transmission</li> </ul>	<ul style="list-style-type: none"> <li>When a busy channel is sensed, a station defers for a random period of time before next sense</li> </ul>	<ul style="list-style-type: none"> <li>When a busy channel is sensed, a station continues to sense until the channel turns idle. Then, with probability <math>p</math>, it transmits, and with probability <math>1 - p</math>, it defers to next time slot.</li> </ul>											
	P-Persistent														
	1-Persistent		<ul style="list-style-type: none"> <li>A special case of P-Persistent where <math>p = 1</math></li> </ul>												
CSMA/CD (Ethernet)		<ul style="list-style-type: none"> <li>Must sense channel before transmission</li> </ul>	<ul style="list-style-type: none"> <li>The same as CSMA</li> </ul>	<ul style="list-style-type: none"> <li>When a collision is detected, transmissions are aborted to reduce the channel wastage.</li> </ul>		Minimum Frame Size <ul style="list-style-type: none"> <li><math>T_{frame} \geq 2\tau</math></li> </ul> Binary Exponential Backoff <ul style="list-style-type: none"> <li>In <math>i</math>-th retransmission, the slot is chosen from a uniformly distributed random variable <math>R_i</math> in the range of <math>[0, 2^K - 1]</math>, where <math>K = \min(i, 10)</math>.</li> </ul>									
CSMA/CA (802.11)		<ul style="list-style-type: none"> <li>Must sense channel before transmission</li> </ul>	<p>Sender:</p> <ul style="list-style-type: none"> <li>If sense channel idle for DIFS, then transmit entire frame (no CD).</li> <li>If sense channel busy, then start random backoff time. Transmits when timer expires.</li> <li>If no ACK, increase random backoff interval</li> </ul> <p>Receiver:</p> <ul style="list-style-type: none"> <li>If frame received OK, return ACK after SIFS</li> </ul>	<ul style="list-style-type: none"> <li>No collision detection due to hidden terminal</li> </ul>	Multi-Access Reservation <ul style="list-style-type: none"> <li>Use random-access with mini-frame (<math>v</math> unit of time) to reserve the channel</li> <li>If reservation successful, transmit <math>u</math> unit of data frame</li> </ul>	<table border="1"> <tr> <td></td><td><math>\frac{u}{u + v/S_r}</math></td><td><math>\frac{u}{(u - v) + \frac{v}{S_r}}</math></td></tr> <tr> <td>Total data length</td><td><math>u</math></td><td><math>u</math></td></tr> <tr> <td>Data bit in mini-frame</td><td>No</td><td>Yes</td></tr> </table>		$\frac{u}{u + v/S_r}$	$\frac{u}{(u - v) + \frac{v}{S_r}}$	Total data length	$u$	$u$	Data bit in mini-frame	No	Yes
	$\frac{u}{u + v/S_r}$	$\frac{u}{(u - v) + \frac{v}{S_r}}$													
Total data length	$u$	$u$													
Data bit in mini-frame	No	Yes													

# Learning Objectives

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- **WLAN Overview**
  - Understand two alternative WLAN architectures
- **802.11 Physical Layer**
  - Understand different transmission schemes
- **802.11 MAC Layer**
  - Understand hidden and exposed terminal problems
  - Understand CSMA/CA protocol
- **Multi-Access Reservation Protocol (MARP)**
  - Understand the scheme of MARP
  - Calculate and maximize throughput for MARP

not examinable

# Reading Material

# Wireless Physical Layer (I)

not examinable

- **Physical layer conforms to OSI (five options)**
  - 1997: **802.11** infrared, FHSS, DHSS
  - 1999: **802.11a** OFDM and **802.11b** HR-DSSS
  - 2001: **802.11g** OFDM
- **802.11 Infrared**
  - Two capacities 1 Mbps or 2 Mbps.
  - Range is 10 to 20 meters and cannot penetrate walls.
  - Does not work outdoors.
- **802.11 FHSS (Frequency Hopping Spread Spectrum)**
  - The main issue is multipath fading.
  - 79 non-overlapping channels, each 1 MHz wide at low end of 2.4 GHz ISM band.
  - Same pseudo-random number generator used by all stations.
  - Dwell time: min. time on channel before hopping (400msec).

# Wireless Physical Layer (II)

not examinable

- **802.11 DSSS (*Direct Sequence Spread Spectrum*)**
  - Spreads signal over entire spectrum using pseudo-random sequence (similar to CDMA see Tanenbaum sec. 2.6.2).
  - Each bit transmitted using an 11 chips Barker sequence, PSK at 1Mbaud.
  - 1 or 2 Mbps.
- **802.11a OFDM (*Orthogonal Frequency Divisional Multiplexing*)**
  - Compatible with European HiperLan2.
  - 54Mbps in wider 5.5 GHz band  transmission range is limited.
  - Uses 52 FDM channels (48 for data; 4 for synchronization).
  - Encoding is complex ( PSM up to 18 Mbps and QAM above this capacity).
  - E.g., at 54Mbps 216 data bits encoded into 288-bit symbols.
  - More difficulty penetrating walls.

# Wireless Physical Layer (III)

not examinable

- **802.11b HR-DSSS (*High Rate Direct Sequence Spread Spectrum*)**
  - 11a and 11b shows a split in the standards committee.
  - 11b approved and hit the market before 11a.
  - Up to 11 Mbps in 2.4 GHz band using 11 million chips/sec.
  - Note in this bandwidth all these protocols have to deal with interference from microwave ovens, cordless phones and garage door openers.
  - Range is 7 times greater than 11a.
  - **11b and 11a are incompatible!!**
- **802.11g OFDM(*Orthogonal Frequency Division Multiplexing*)**
  - An attempt to combine the best of both 802.11a and 802.11b.
  - Supports bandwidths up to 54 Mbps.
  - Uses 2.4 GHz frequency for greater range.
  - Is backward compatible with 802.11b.