Graph II MST, Shortest Path

Graph Terminology

- Node (vertex)
- Edge (arc)
- Directed graph, undirected graph
- Degree, in-degree, out-degree
- Subgraph
- Simple path
- Cycle
- Directed acyclic graph
- Weighted graph

Graph representation Adjacency Matrix

- Assume N nodes in graph
- Use Matrix A[0...N-1][0...N-1]
 - □ if vertex i and vertex j are adjacent in graph,A[i][j] = 1,
 - \Box otherwise A[i][j] = 0
 - □ if vertex i has a loop, A[i][i] = 1
 - □ if vertex i has no loop, A[i][i] = 0

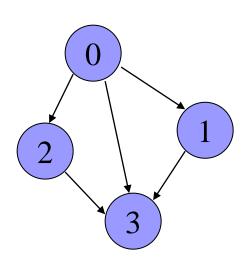
Graph representation

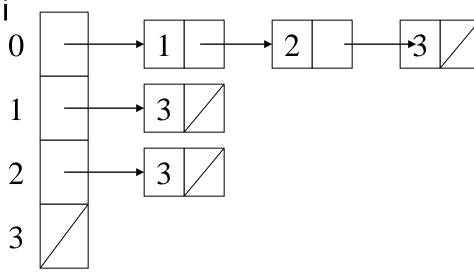
Adjacency List

An array of list

the ith element of the array is a list of vertices that

connect to vertex i





vertex 0 connect to vertex 1, 2 and 3 vertex 1 connects to 3 vertex 2 connects to 3

Graph Traversal

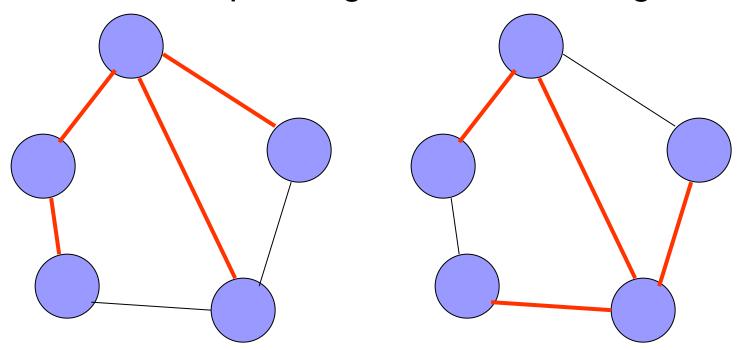
- From one vertex, list out all vertices that can be reached in graph G
- Set of nodes to expand
- Each node has a flag to indicate visited or not
- Depth First Traversal
- Breadth First Traversal

Spanning Tree

- Connected subgraph that includes all vertices of the original connected graph
- Subgraph is a tree
 - □ If original graph has n vertices, the spanning tree has n vertices and n-1 edges.
 - No circle in this subgraph

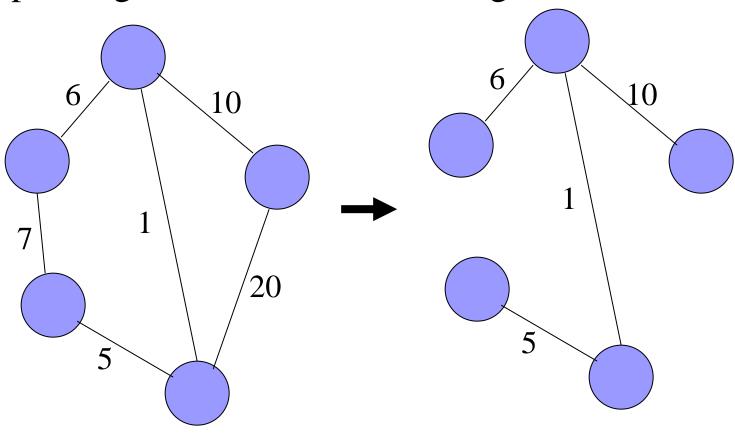
Spanning Tree

- Minimum number of edges to keep it connected
- If N vertices, spanning tree has N-1 edges



Minimum Spanning Tree (MST)

Spanning tree with minimum weight



v.

Prim's Algorithm For Finding MST

- All nodes are unselected, mark node v selected
- 2. For each node of graph,

{

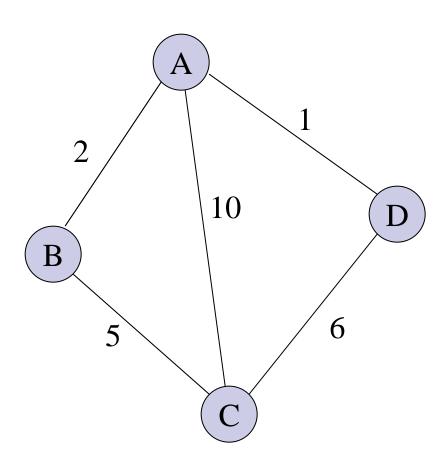
Find the edge with minimum weight that connects an unselected node with a selected node

Mark this unselected node as selected

}

Example

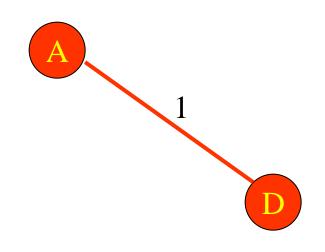
Find the MST of the following graph



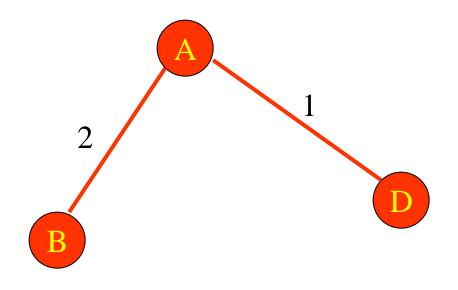
Step 1: mark vertex A as selected



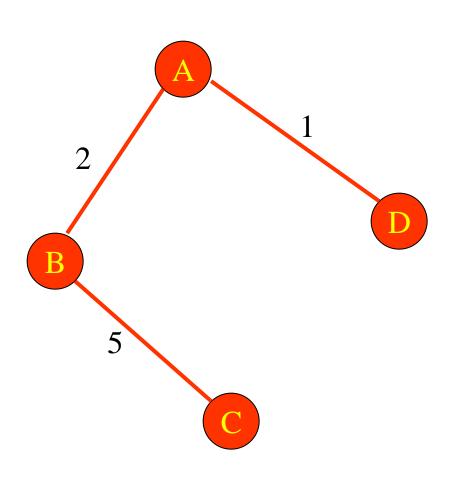
Step 2: find the minimum weighted edge connected to vertex A, and mark the other vertex on this edge as selected.



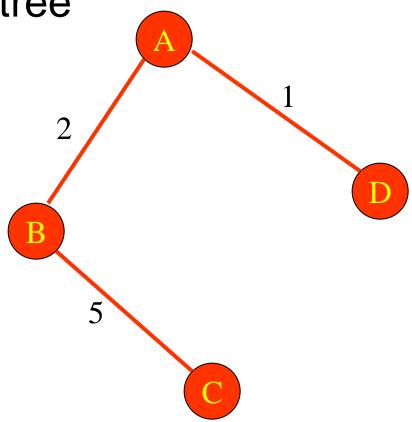
Step 3: find the minimum weighted edge connected to vertices set { A, D }, and mark the other vertex on this edge as selected.



Step 4: find the minimum weighted edge connected to vertices set { A, D, B}, and mark the other vertex on this edge as selected.



Step 5: All vertex are marked as selected, So we find the minimum spanning tree



Pseudo code for Prim's Alg.

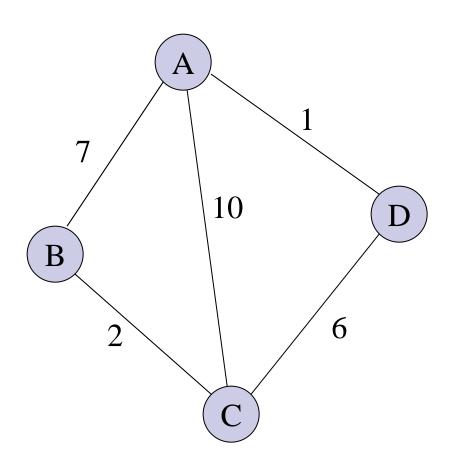
- Minimum-Spanning-Tree-by-Prim(G, weight-function, source)
 - ☐ for each vertex u in graph G
 - set key of u to ∞
 - set parent of u to nil
 - set key of source vertex to zero
 - enqueue all vertices to Q
 - □ while Q is not empty
 - extract vertex u from Q // u is the vertex with the lowest key that is in Q
 - for each adjacent vertex v of u do
 - \Box if (v is still in **Q**) and (weight-function(u, v) < key of v) then
 - set u to be parent of v // in minimum-spanning-tree
 - update v's key to equal weight-function(u, v)

Kruskal's Algorithm For Finding MST

- All edges are unselected
- 2. Sort all edges and store them in set S
- 3. For each edge in S
 {
 If adding the edge to MST does not form a circle, add this edge to MST;
 Delete this edge from S;
 If |edges in MST|=|nodes in graph|-1, exit;
 }

Example

Find the MST of the following graph



Step 1: Initialize S and empty MST

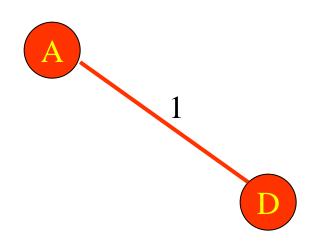


 $\left(\mathbf{B}\right)$



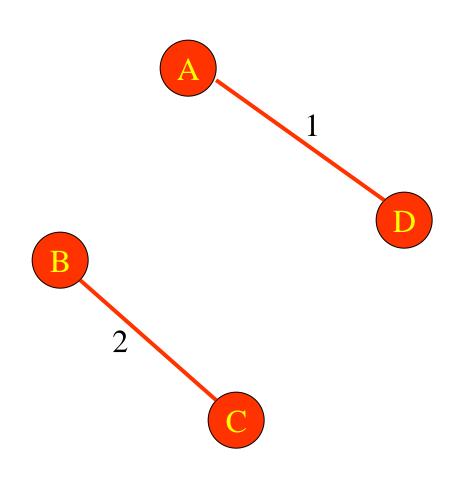


Step 2: Add AD to MST, and delete it edge from S.



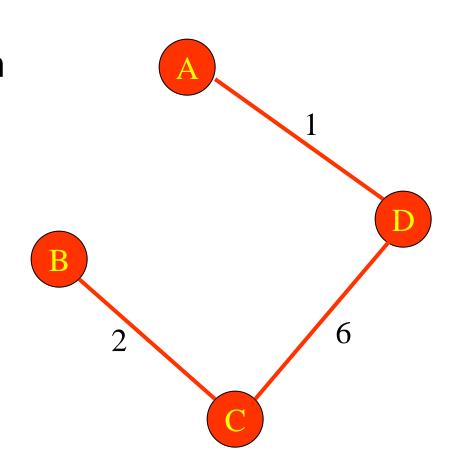
$$S=\{(B,C=)2, (C,D)=6, (A,B)=7, (A,C)=10 \}$$

Step 3: Add BC to MST, and delete it from S.



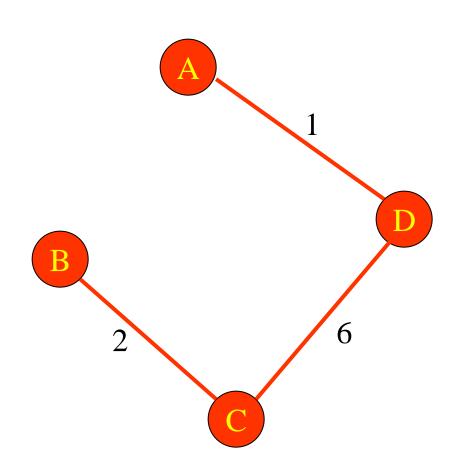
$$S=\{(C,D)=6, (A,B)=7, (A,C)=10 \}$$

Step 4: Add CD to MST, and delete it from S

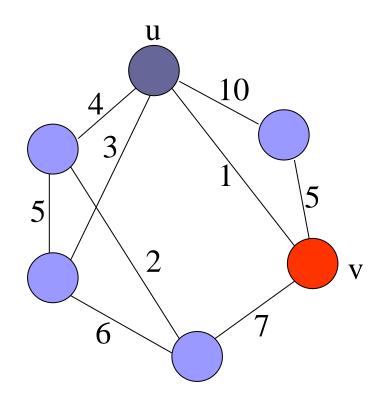


$$S=\{(A,B)=7, (A,C)=10\}$$

Step 5: Satisfy the exiting condition, So we find the minimum spanning tree



Shortest Path Problem

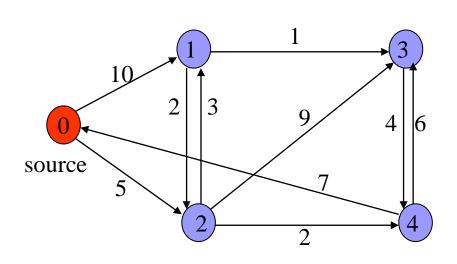


■ Weight: cost, distance, travel time, hop ...

Single Source Shortest Path Problem

- Single source shortest path problem
 - ☐ Find the shortest path to all other nodes
- Dijkstra's shortest path algorithm
 - □ Find shortest path greedily by updating distance to all other nodes
 - Not applicable for negative weights

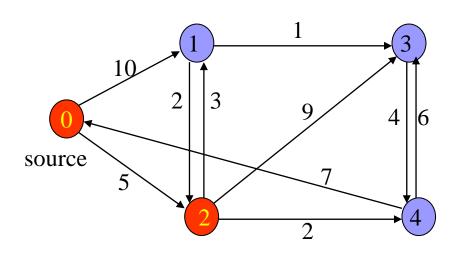
- Greedy Algorithm
- Assume all weight of edge >0



| node | from node V ₀ to other nodes | | | |
|----------------|---|--|--|--|
| V ₁ | 10 (V ₀) | | | |
| V_2 | 5 | | | |
| v 2 | (V ₀) | | | |
| V_3 | ∞ (V ₀) | | | |
| V_4 | ∞ (\/ \ | | | |
| best | (V ₀) | | | |
| best | | | | |

step 1: find the shortest path to node 0

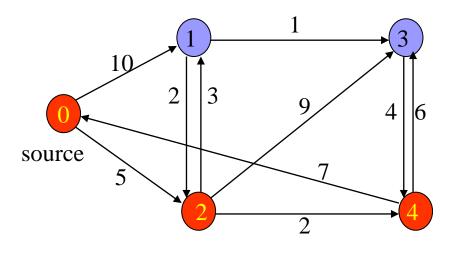
node 2 is selected



| node | from node V ₀ to other nodes | | | |
|----------------|---|--|--|--|
| V ₁ | 10 (V ₀) | | | |
| V ₂ | 5 (V ₀) | | | |
| V ₃ | ∞ (V ₀) | | | |
| V ₄ | ∞ (V ₀) | | | |
| best | V ₂ | | | |

step 2: recalculate the path to all other nodes

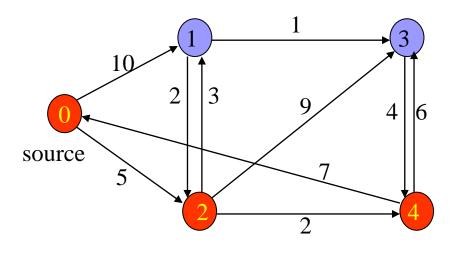
☐ find the shortest path to node 0. Node 4 is



| node | from node V ₀ to other nodes | | | |
|-------|---|-------------------|--|--|
| \/ | 10 | 8 | | |
| V_1 | (V_0) | (V ₂) | | |
| \/ | 5 | | | |
| V_2 | (V_0) | | | |
| \/ | ∞ | 14 | | |
| V_3 | (V_0) | (V ₂) | | |
| V_4 | ∞ | 7 | | |
| | (V_0) | (V ₂) | | |
| best | V_2 | | | |

step 2: recalculate the path to all other nodes

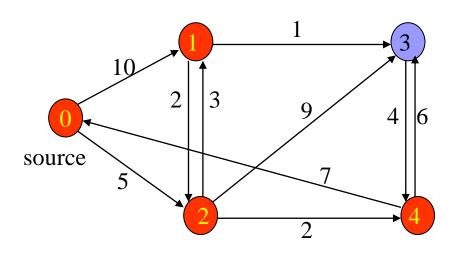
☐ find the shortest path to node 0. Node 4 is



| node | from node V ₀ to other nodes | | | |
|----------------|---|-------------------|--|--|
| \/ | 10 | 8 | | |
| V_1 | (V ₀) | (V ₂) | | |
| \/ | 5 | | | |
| V_2 | (V ₀) | | | |
| \/ | ∞ | 14 | | |
| V_3 | (V_0) | (V_2) | | |
| V ₄ | ∞ | 7 | | |
| | (V ₀) | (V_2) | | |
| best | V_2 | V_4 | | |

step 3: recalculate the path to all other nodes

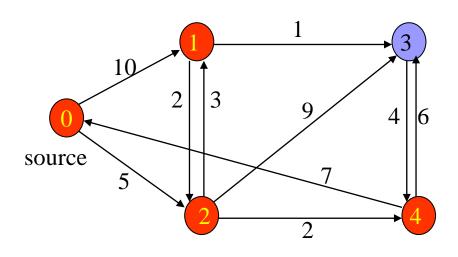
☐ find the shortest path to node 0. node 1 is



| node | from node V ₀ to other nodes | | | |
|-------|---|---------|-------------------|--|
| \/ | 10 | 8 | 8 | |
| V_1 | (V_0) | (V_2) | (V ₂) | |
| \/ | 5 | | | |
| V_2 | (V_0) | | | |
| \/ | ∞ | 14 | 13 | |
| V_3 | (V_0) | (V_2) | (V ₄) | |
| V_4 | 8 | 7 | | |
| | (V_0) | (V_2) | | |
| best | V_2 | V_4 | | |

step 3: recalculate the path to all other nodes

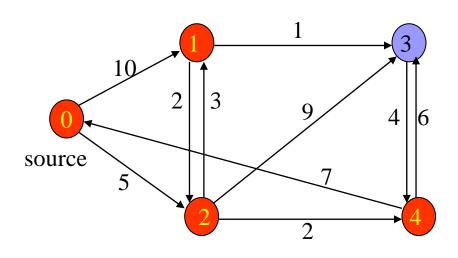
☐ find the shortest path to node 0. node 1 is



| node | from node V ₀ to other nodes | | | |
|-------|---|---------|-------------------|--|
| \/ | 10 | 8 | 8 | |
| V_1 | (V_0) | (V_2) | (V_2) | |
| \/ | 5 | | | |
| V_2 | (V_0) | | | |
| \/ | 8 | 14 | 13 | |
| V_3 | (V_0) | (V_2) | (V ₄) | |
| V_4 | 8 | 7 | | |
| | (V_0) | (V_2) | | |
| best | V_2 | V_4 | V ₁ | |

step 3: recalculate the path to all other nodes

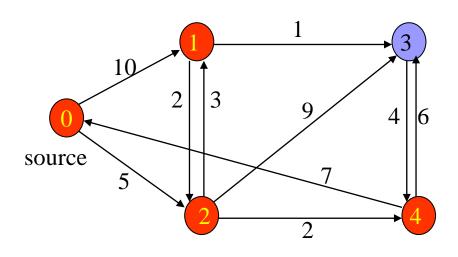
☐ find the shortest path to node 0. node 1 is



| node | from node V ₀ to other nodes | | | |
|-------|---|---------|-------------------|-------------------|
| \/ | 10 | 8 | 8 | |
| V_1 | (V_0) | (V_2) | (V ₂) | |
| \/ | 5 | | | |
| V_2 | (V_0) | | | |
| \/ | ∞ | 14 | 13 | 9 |
| V_3 | (V_0) | (V_2) | (V ₄) | (V ₁) |
| V_4 | 8 | 7 | | |
| | (V_0) | (V_2) | | |
| best | V_2 | V_4 | V ₁ | |

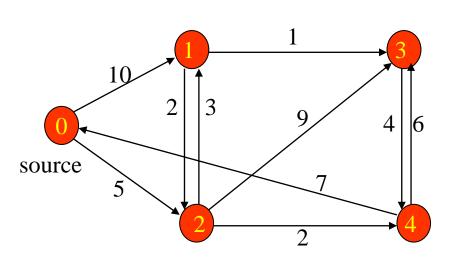
step 3: recalculate the path to all other nodes

☐ find the shortest path to node 0. node 1 is



| node | from node V ₀ to other nodes | | | |
|----------------|---|---------|-------------------|-------------------|
| \/ | 10 | 8 | 8 | |
| V_1 | (V_0) | (V_2) | (V ₂) | |
| \/ | 5 | | | |
| V_2 | (V_0) | | | |
| \/ | 8 | 14 | 13 | 9 |
| V_3 | (V_0) | (V_2) | (V ₄) | (V ₁) |
| V ₄ | 8 | 7 | | |
| | (V_0) | (V_2) | | |
| best | V_2 | V_4 | V ₁ | V_3 |

Now we get all shortest paths to each node



| node | from node V_0 to other nodes | | | |
|----------------|--------------------------------|---------|-------------------|-------------------|
| \/ | 10 | 8 | 8 | |
| V_1 | (V_0) | (V_2) | (V_2) | |
| \/ | 5 | | | |
| V_2 | (V_0) | | | |
| \/ | ∞ | 14 | 13 | 9 |
| V_3 | (V_0) | (V_2) | (V ₄) | (V ₁) |
| V ₄ | ∞ | 7 | | |
| | (V ₀) | (V_2) | | |
| best | V_2 | V_4 | V ₁ | V_3 |



Dijkstra Algorithm

```
Mark source node selected
Initialize all distances to Infinite, source node distance
  to 0.
Make source node the current node.
While (there is unselected node)
      Expand on current node
      Update distance for neighbors of current node
      Find an unselected node with smallest distance,
      and make it current node and mark this node
  selected
```

Pseudo-code For Dijkstra's Algorithm

```
function Dijkstra(G, w, s)
    for each vertex v in V[G] // Initializations
        d[v] := infinity
        previous[v] := undefined
    d[s] := 0
    S := empty set
    Q := V[G]
    while Q is not an empty set // The algorithm itself
        u := Extract_Min(Q)
        S := S union \{u\}
        for each edge (u,v) outgoing from u
            if d[u] + w(u,v) < d[v] // Relax (u,v)
                d[v] := d[u] + w(u,v)
                previous[v] := u
```