



Algorithm Engineering: Milestone 2

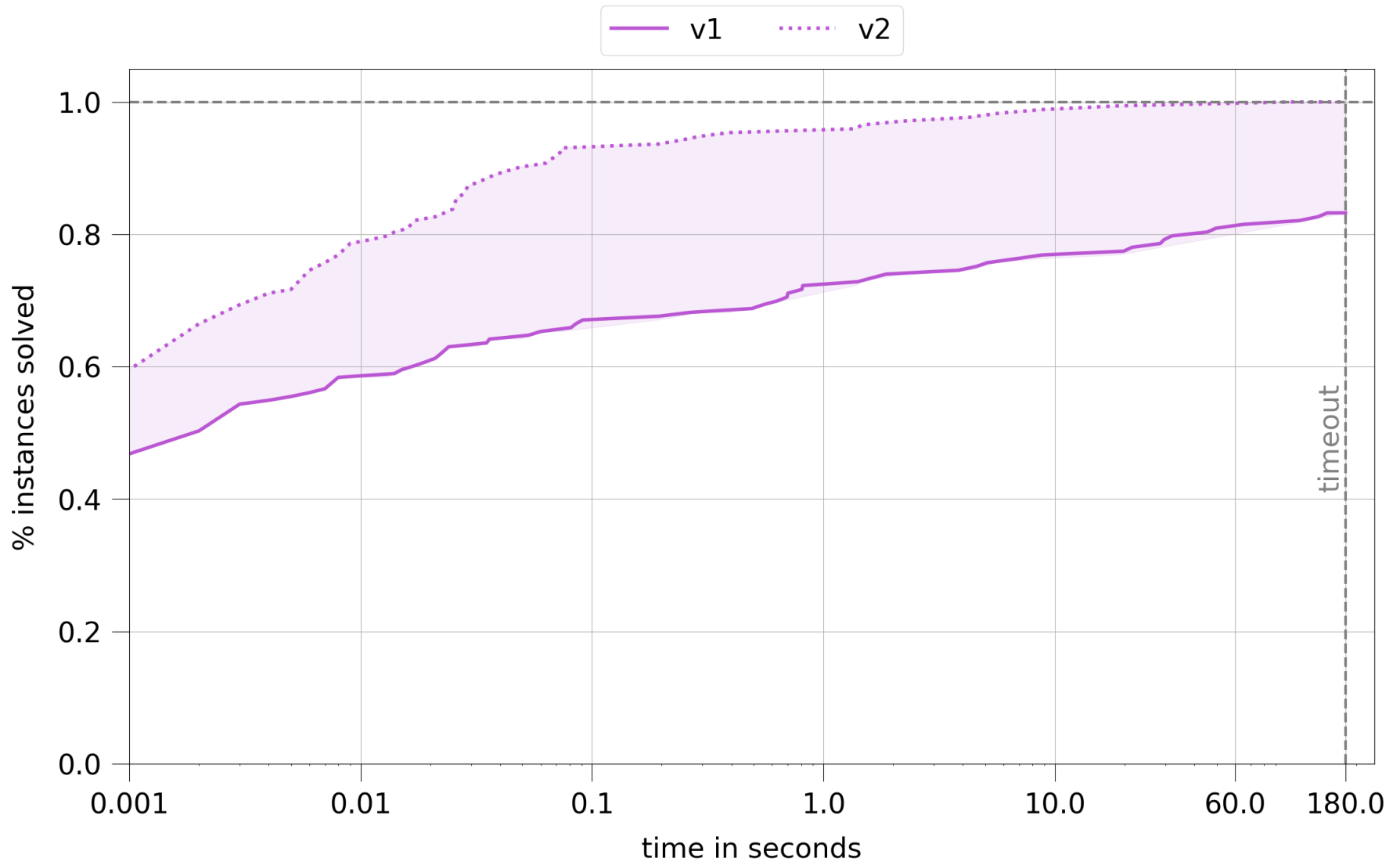
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& Lennart Uhrmacher



Introduction

- Previous state:
 - Base search tree algorithm
 - Finding cycles to branch on with DFS
 - Preprocessing with Tarjan's Algorithm
- Current state:
 - Reduction rules in preprocessing
 - Forbidden nodes
 - Flowers
 - Finding cycles with BFS

Introduction



Introduction

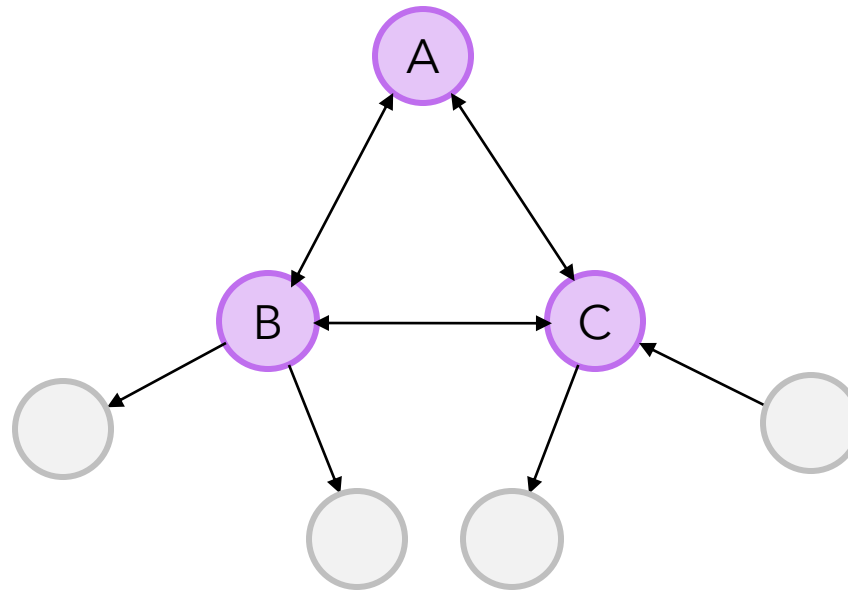
- What will we do today?
 - Go through each new feature
 - Take a closer look on how it works
 - Analyze the importance of the feature

Reduction rules

- Decomposition into cyclic components already used
- Reduction rules:
 - Remove self edges
 - Chain rule
 - Remove trivial vertices
- Advanced triangle reduction rule
- Only used as preprocessing

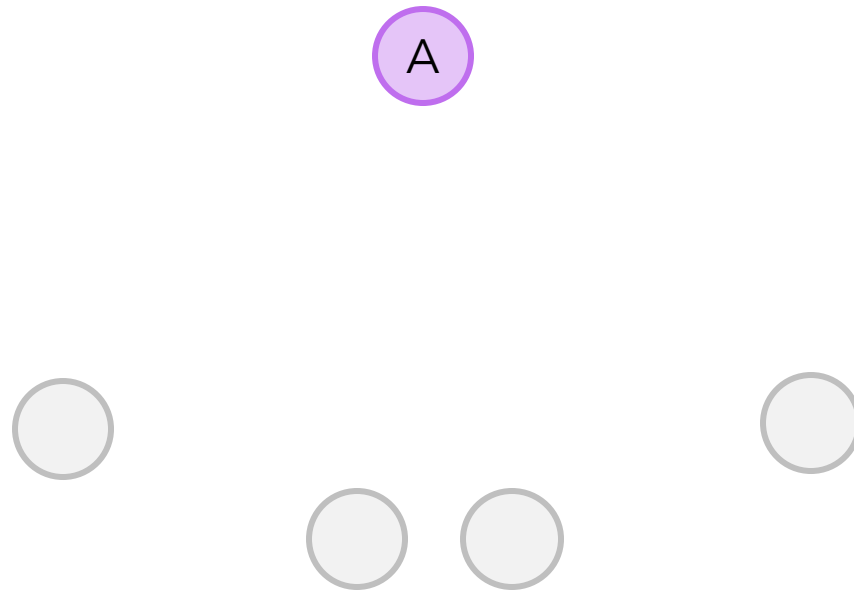
Reduction rules

- Advanced Triangle reduction rule:

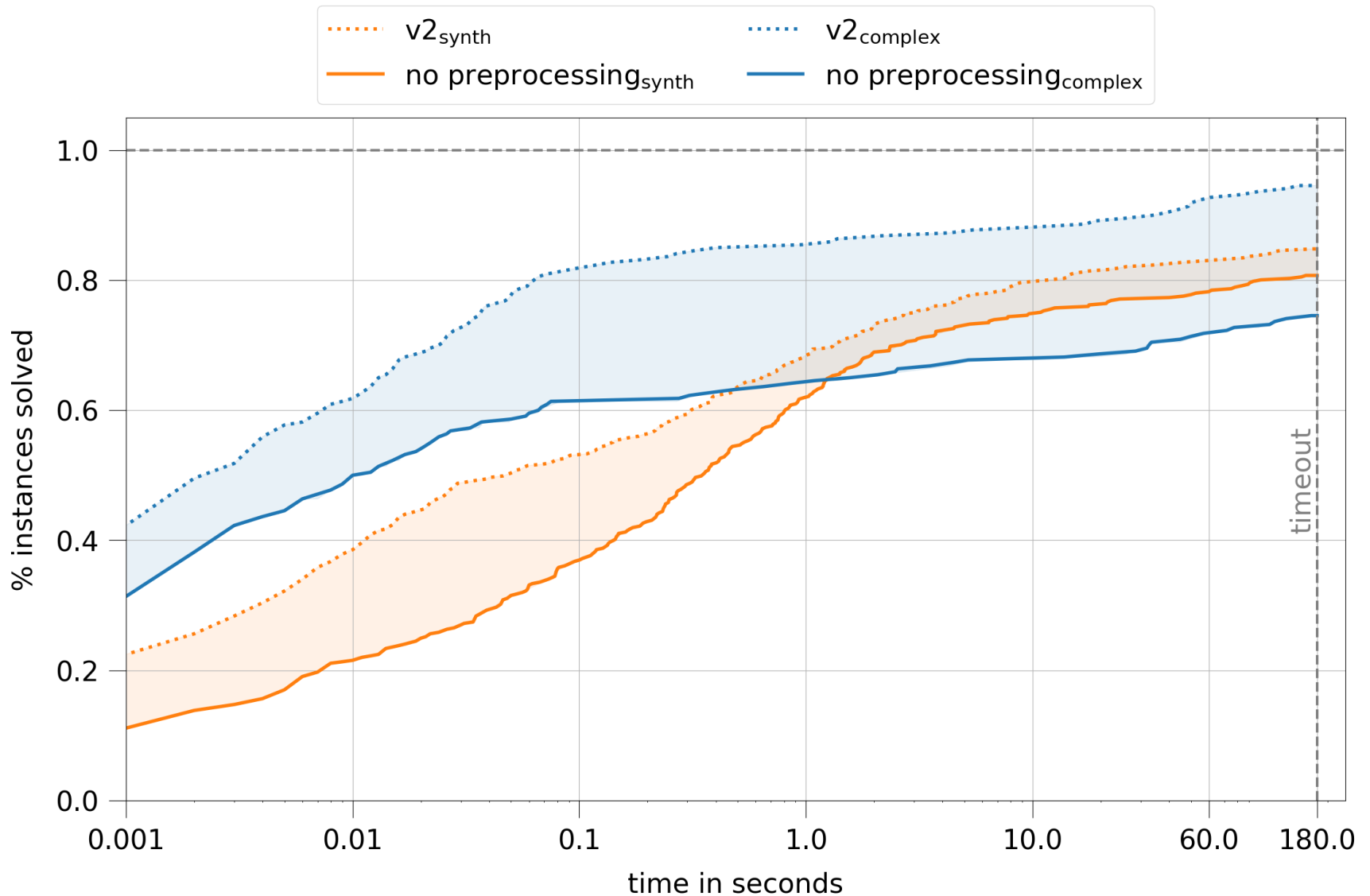


Reduction rules

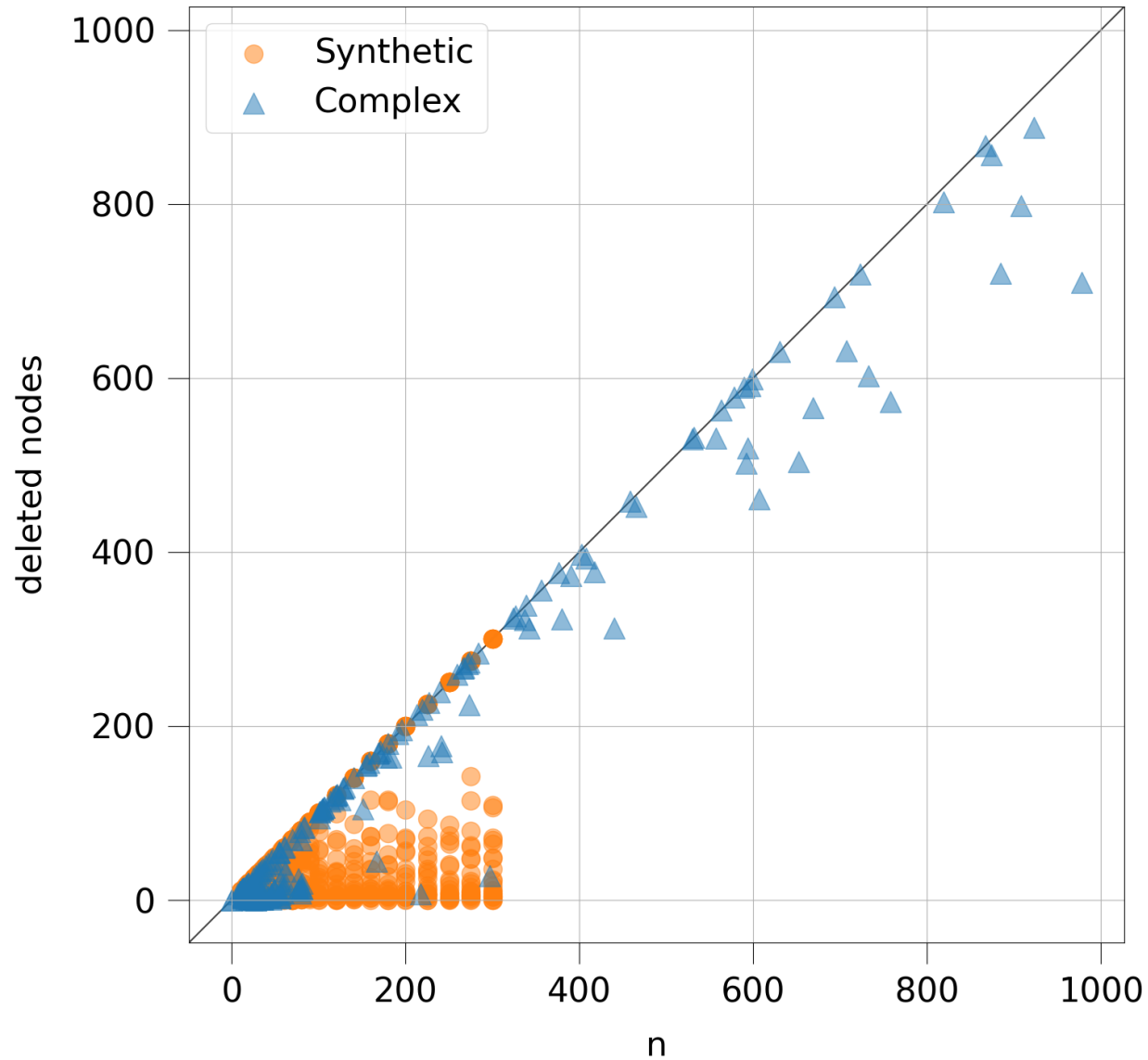
- Advanced Triangle reduction rule:



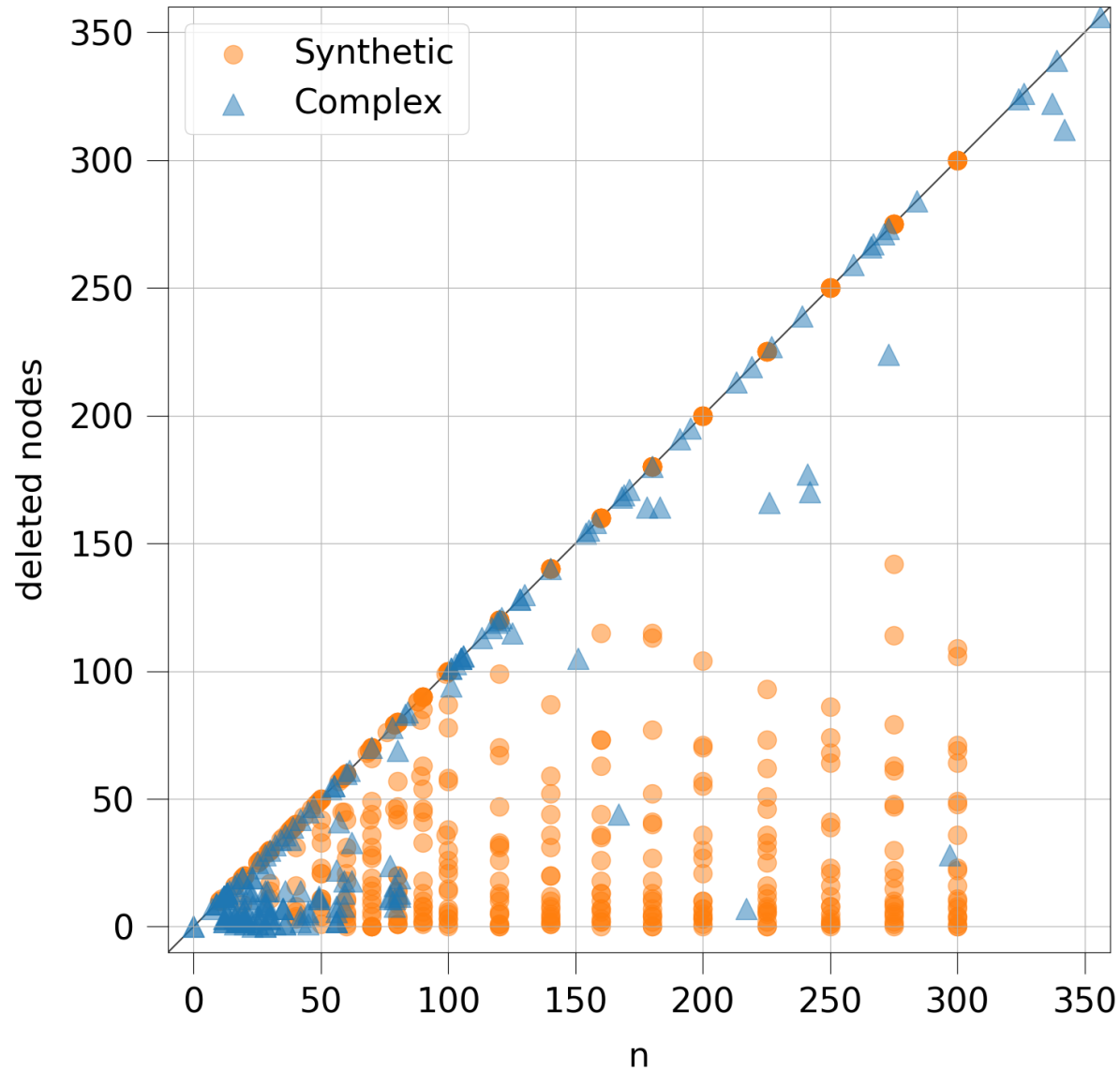
Reduction rules



Reduction rules



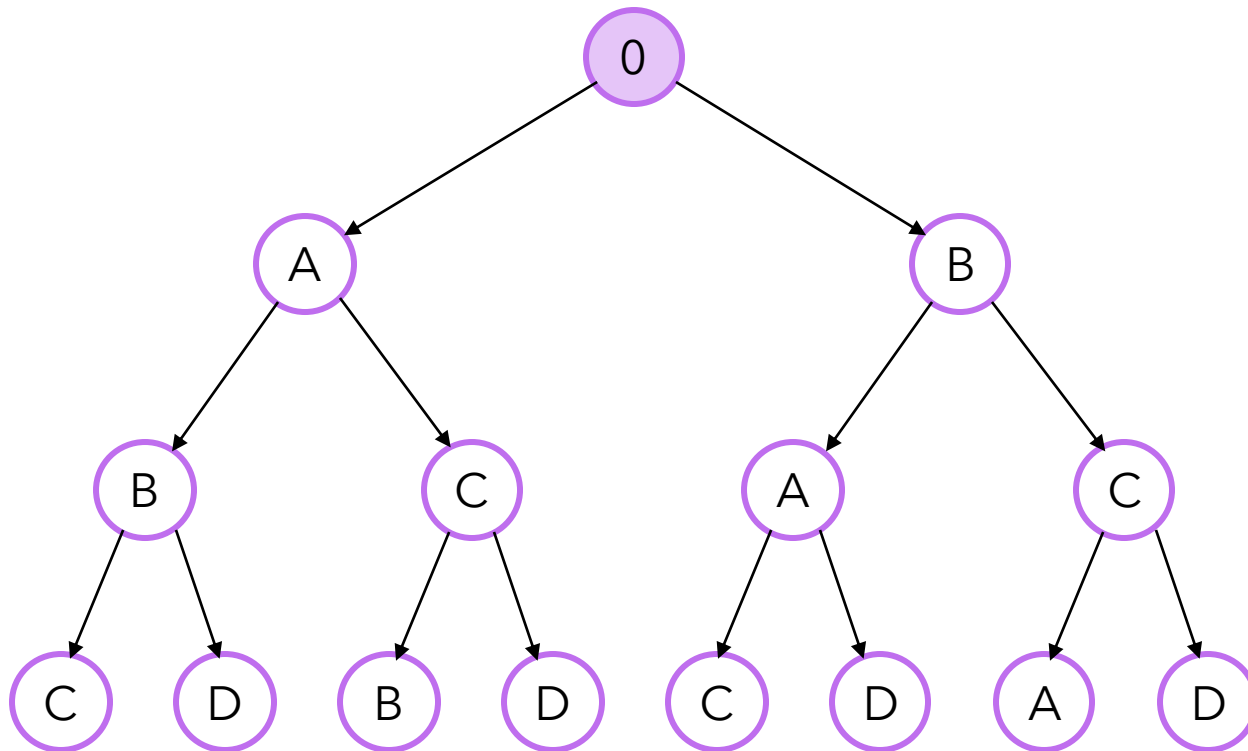
Reduction rules



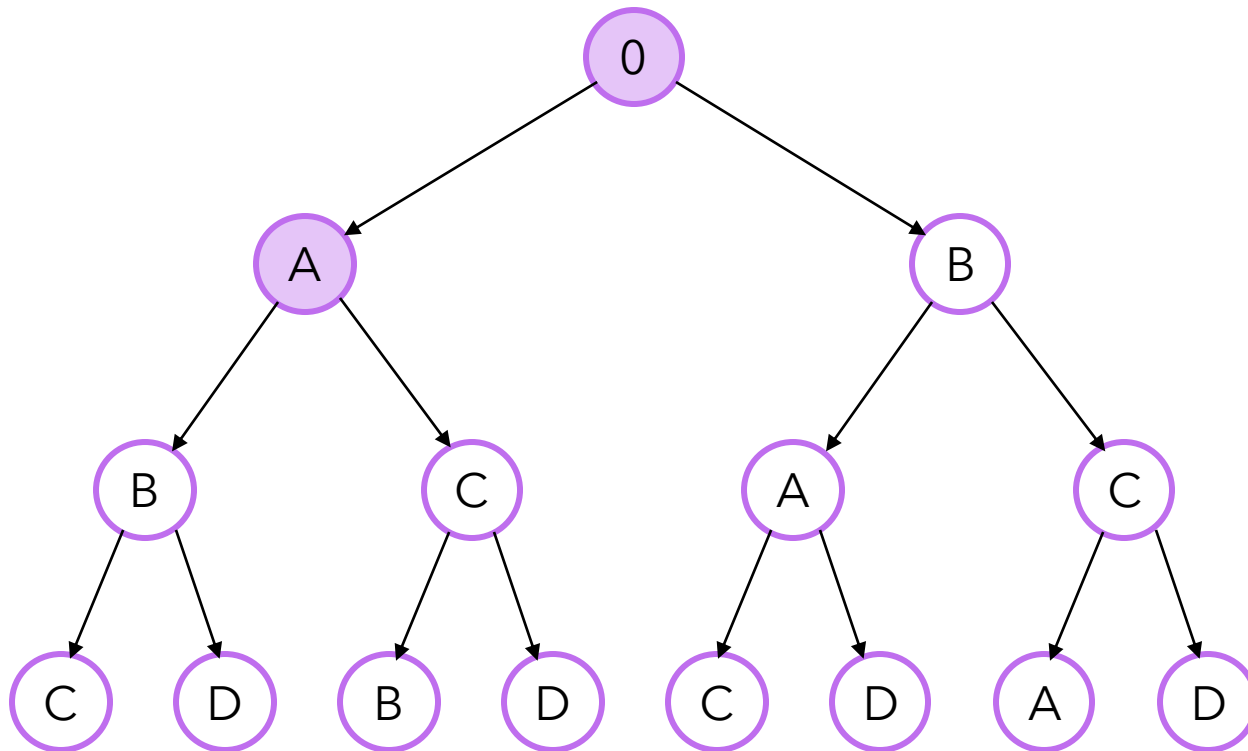
Forbidden nodes

- Forbidden node algorithm:
 - When a node is deleted, it is labeled as 'forbidden'
 - In further branching, 'forbidden' nodes can be skipped

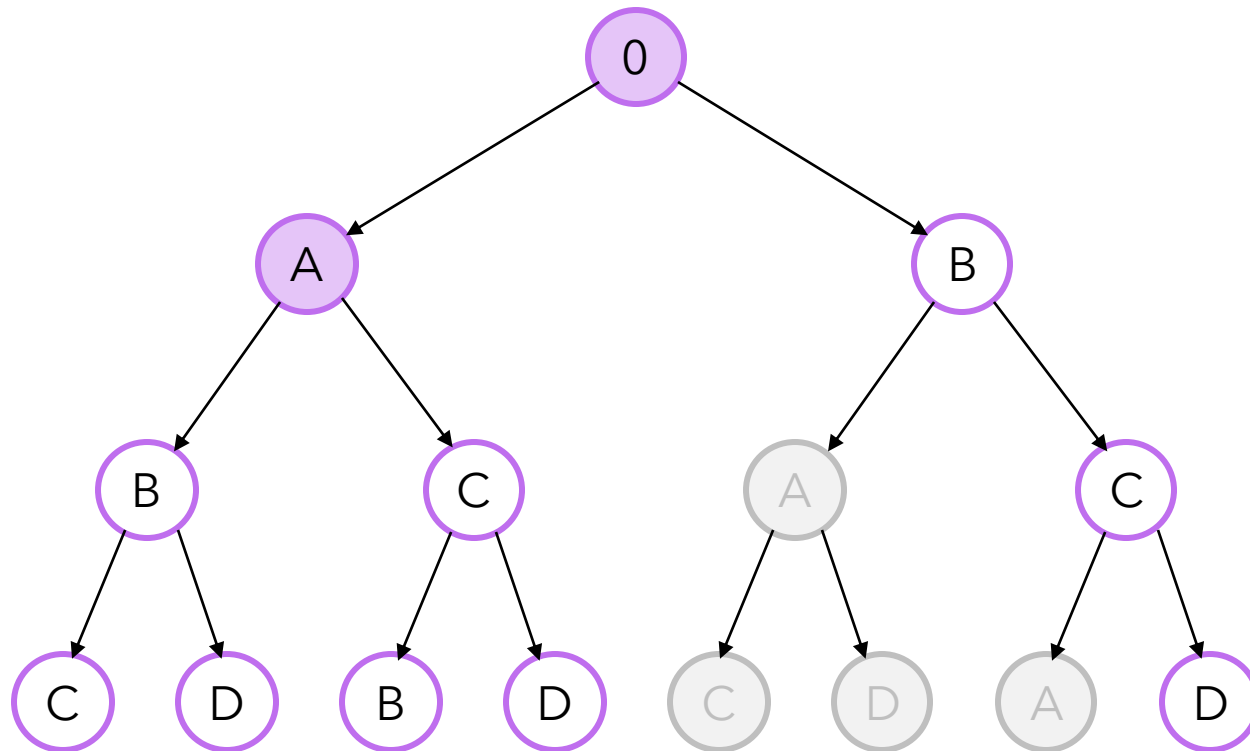
Forbidden nodes

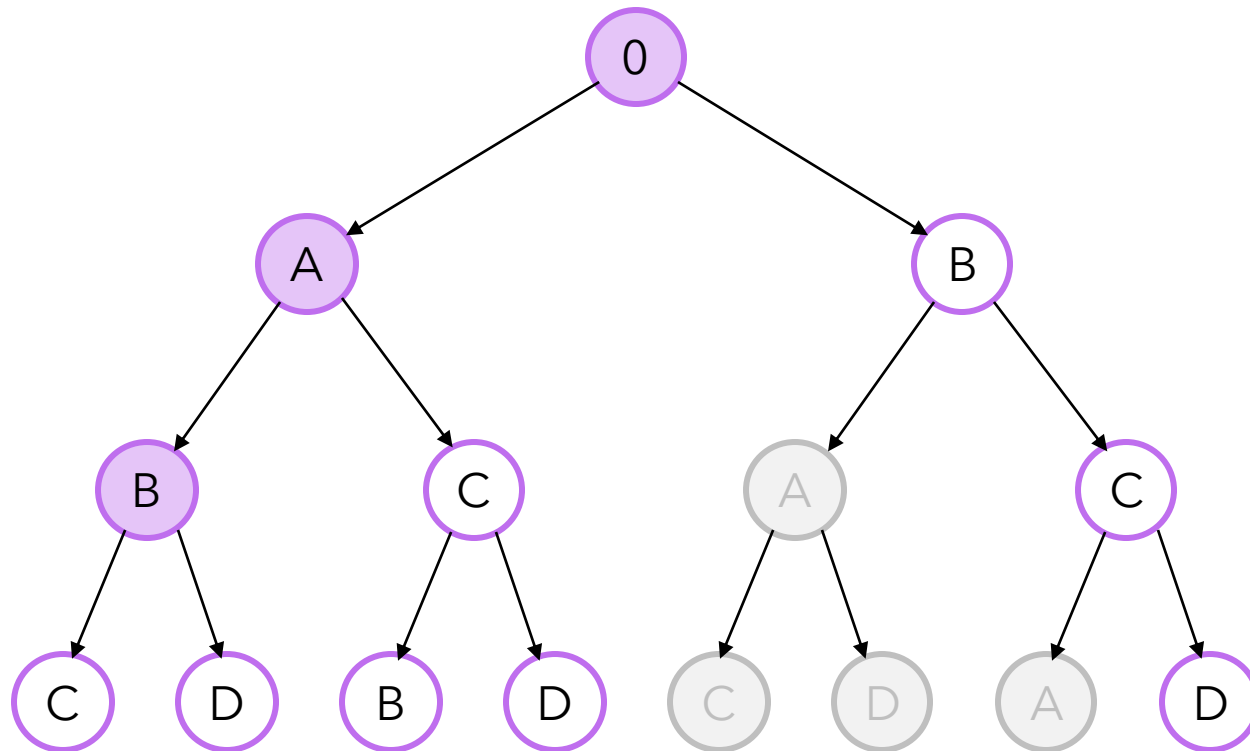


Forbidden nodes

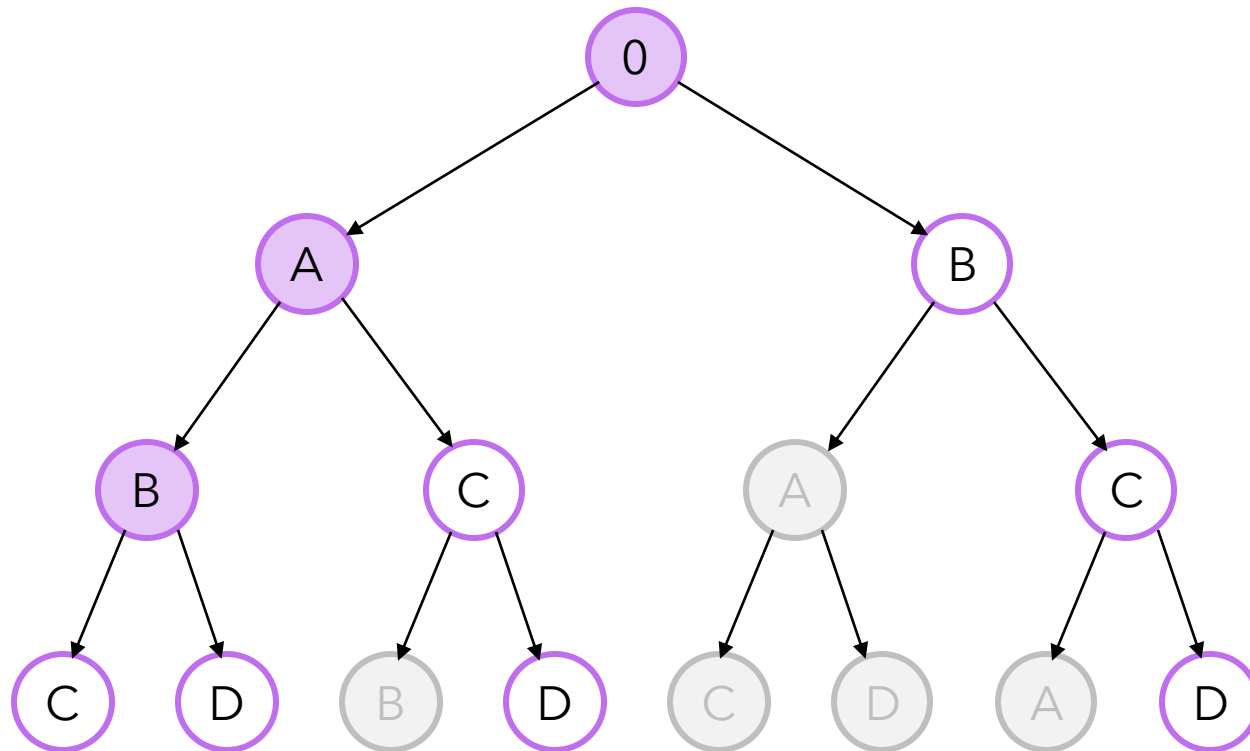


Forbidden nodes

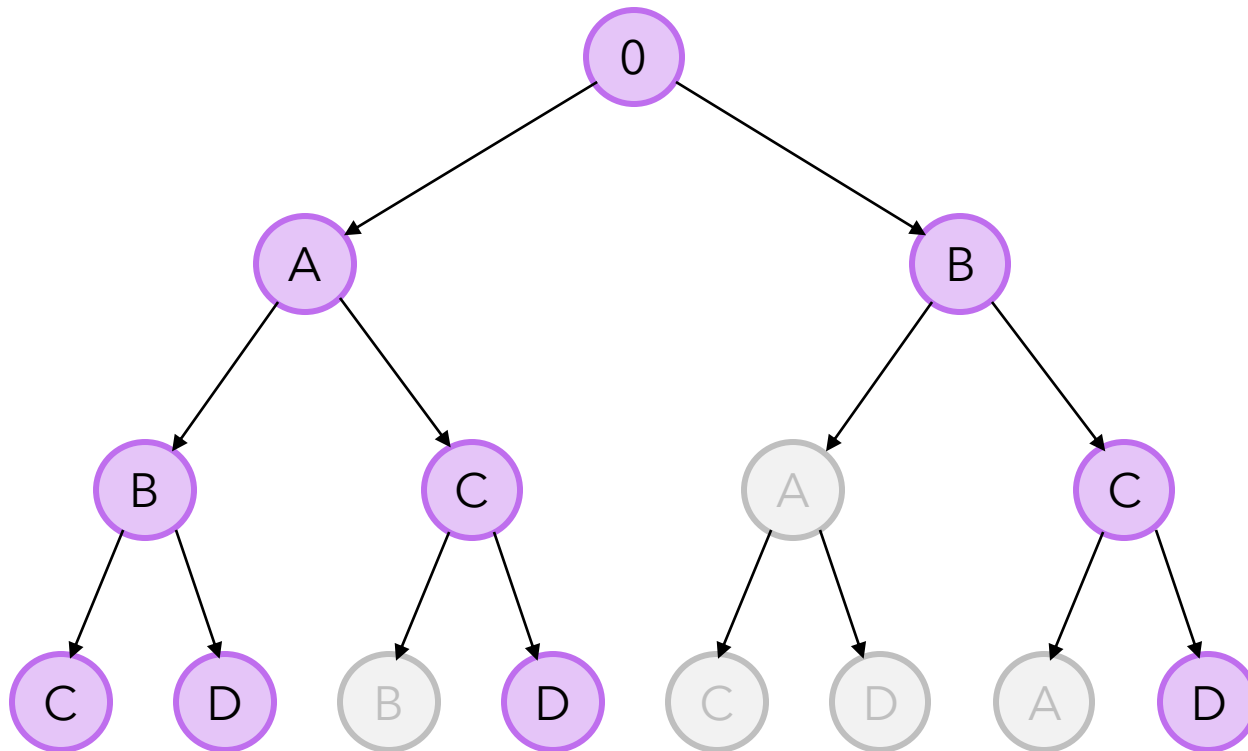




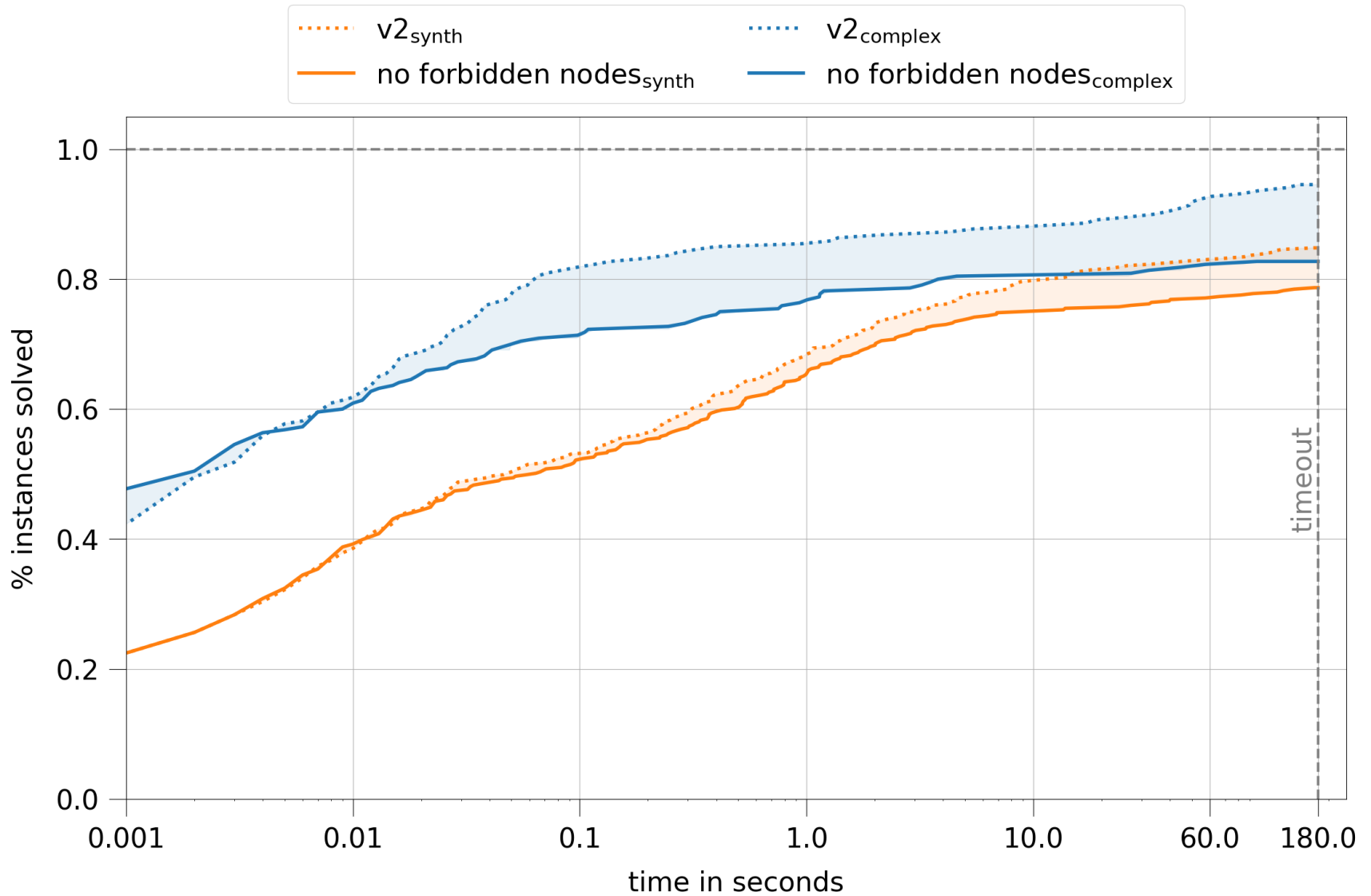
Forbidden nodes



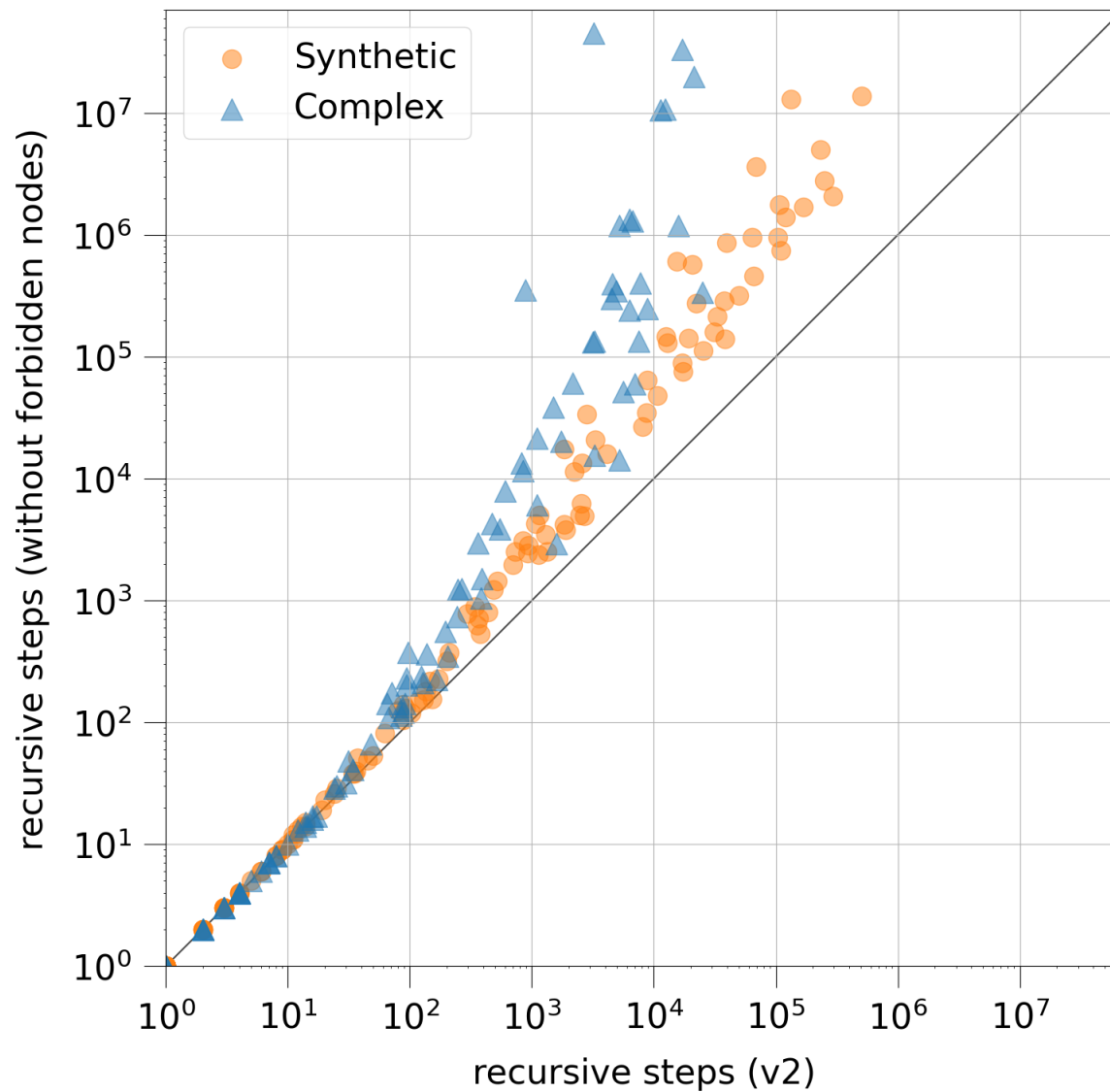
Forbidden nodes



Forbidden nodes



Forbidden nodes



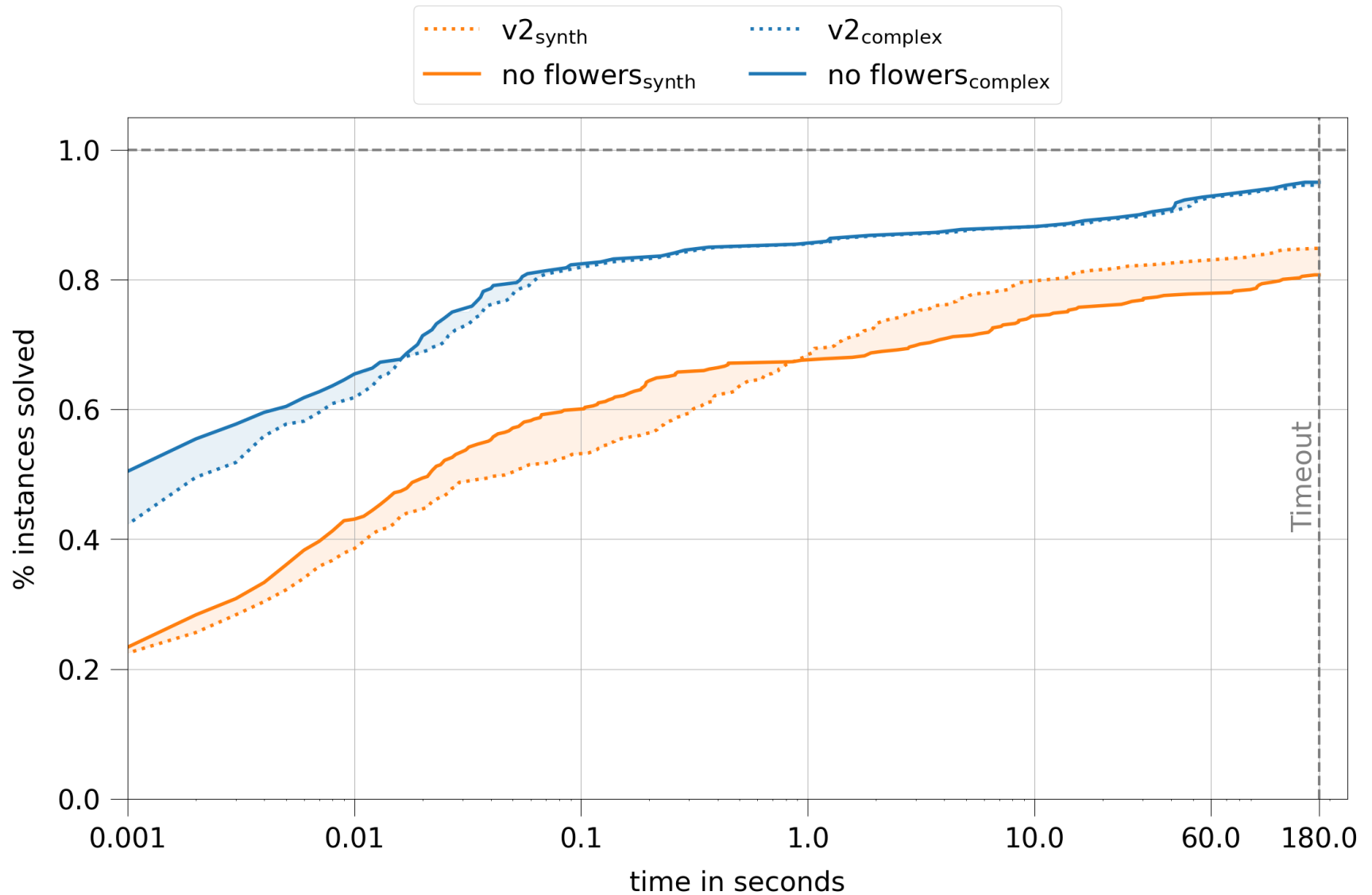
Flowers

- General idea: $\text{petal}(V)$ = number of disjoint cycles connected in V
- Approach from lecture:
 1. Create graph
 2. Calculate Max-Flow/Min-Cut with Ford-Fulkerson
- Optimization: Don't create a new graph for every node

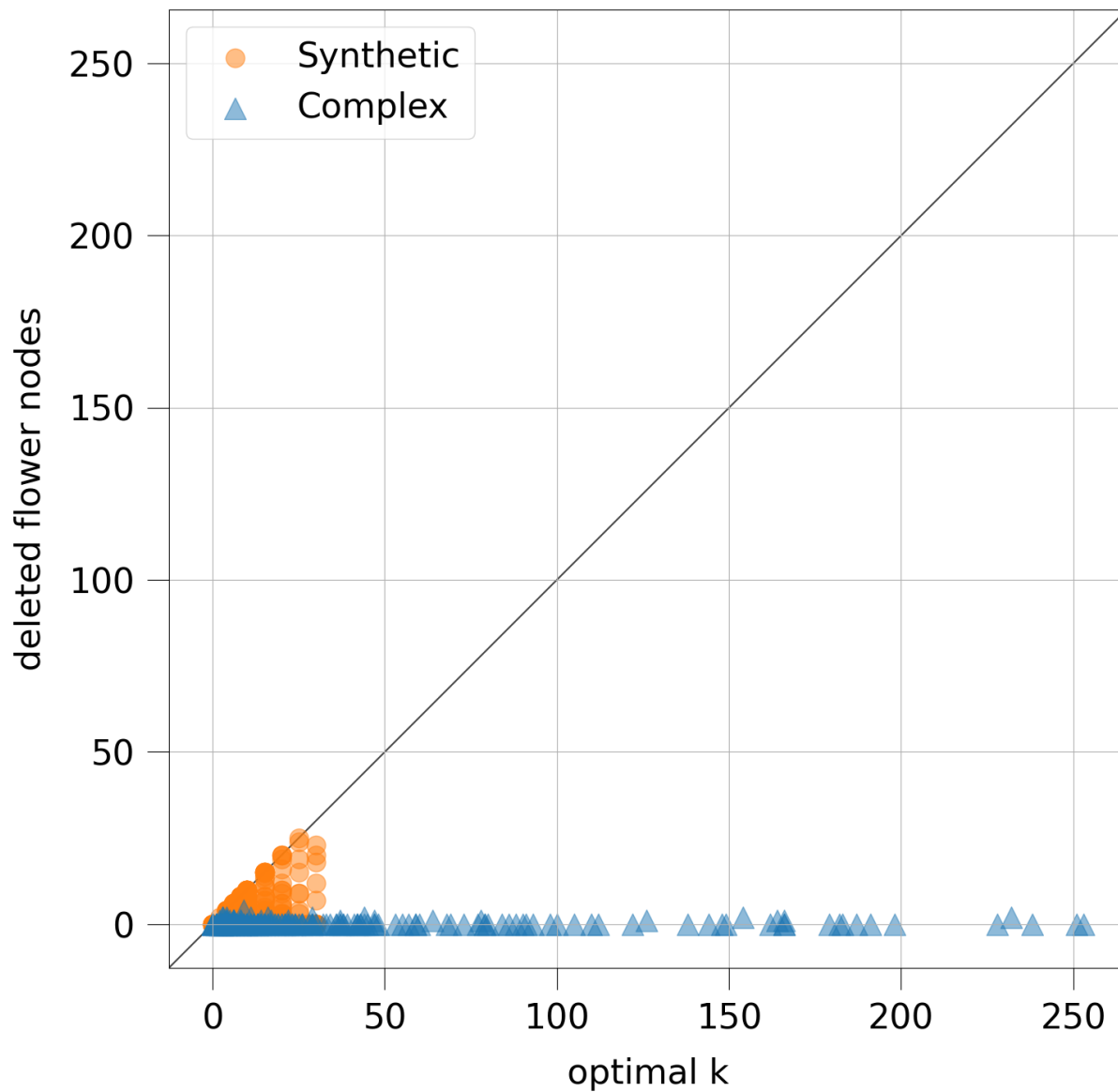
Flowers

- Petal rule: when $\text{petal}(V) > k$, delete V and increase k by one
- How do we use it:
 1. Calculate petal-values for all nodes
 2. Calculate for current k : are there nodes, that can be deleted with the petal rule? \rightarrow If yes, delete one and recalculate petal-values
 3. When the amount of deleted nodes is higher than our current k , there is no solution for this k
- Optimization: reduce petal-values by 1 instead of recalculating
- When do we use it? For each k before we start branching

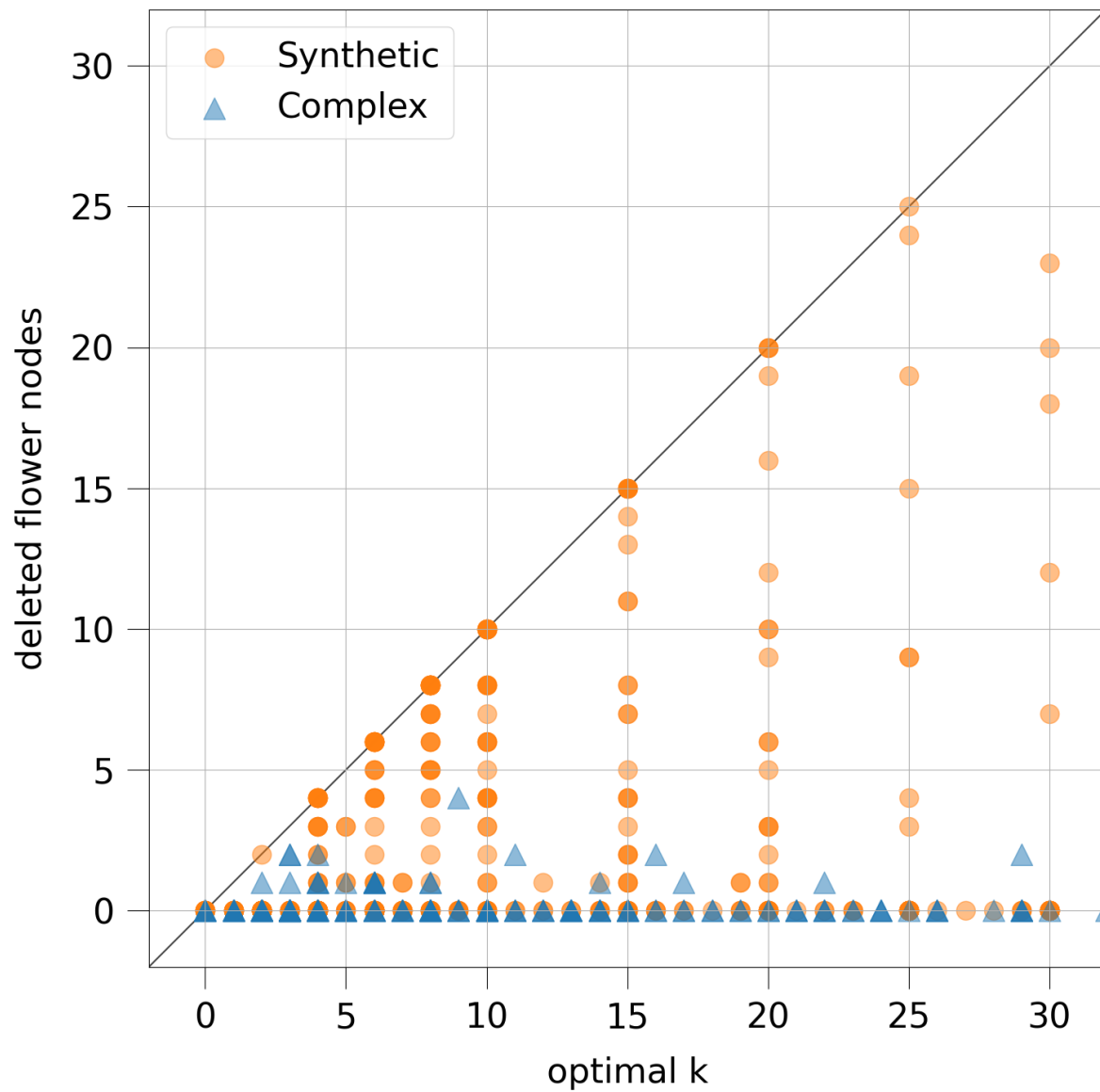
Flowers



Flowers

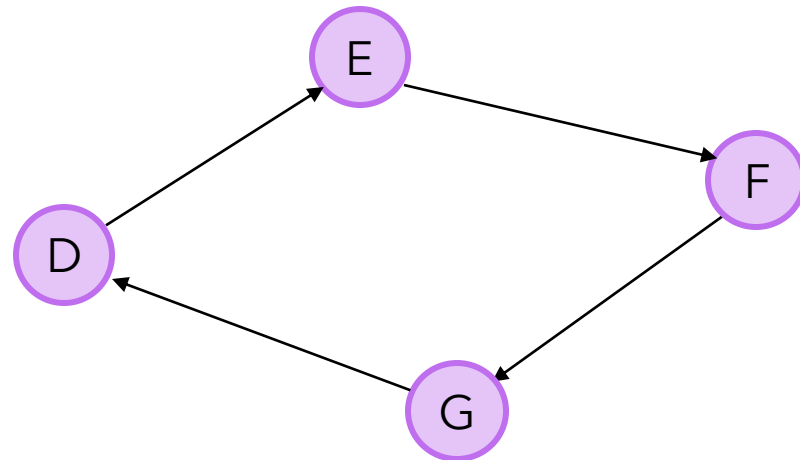
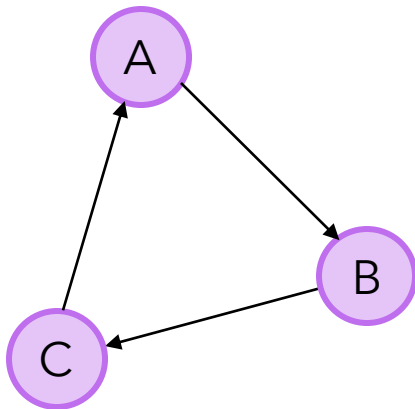


Flowers



Finding cycles

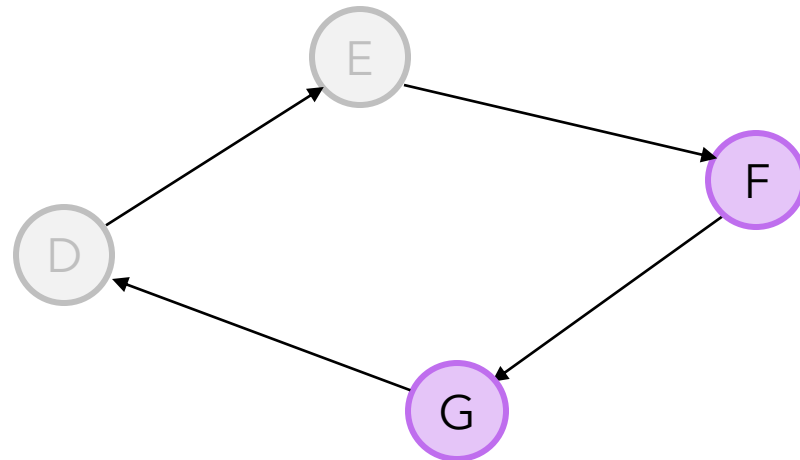
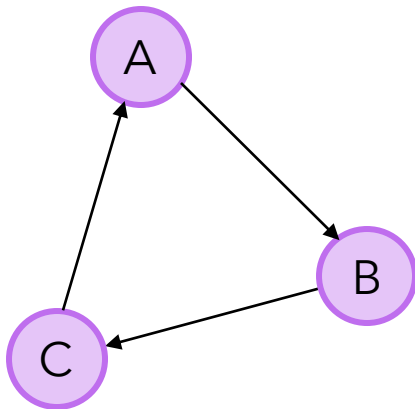
- Forbidden nodes do affect the cycles in the graph:



A-B-C is the shortest cycle, but is it necessarily the best to branch on?

Finding cycles

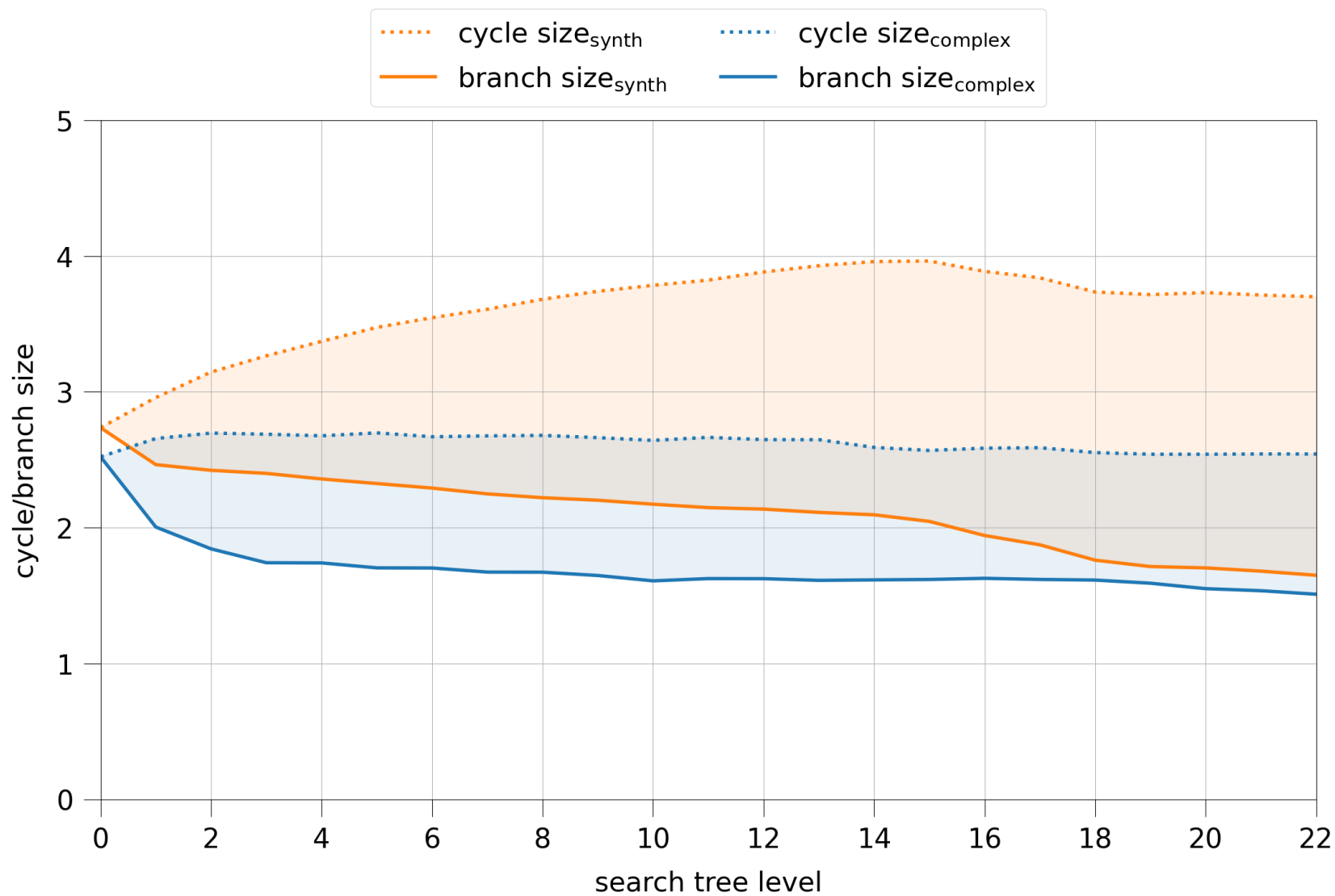
- Forbidden nodes do affect the cycles in the graph:



A-B-C is the shortest cycle, but is it necessarily the best to branch on?

→ No, D-E-F-G branches only 2 instead of 3 times

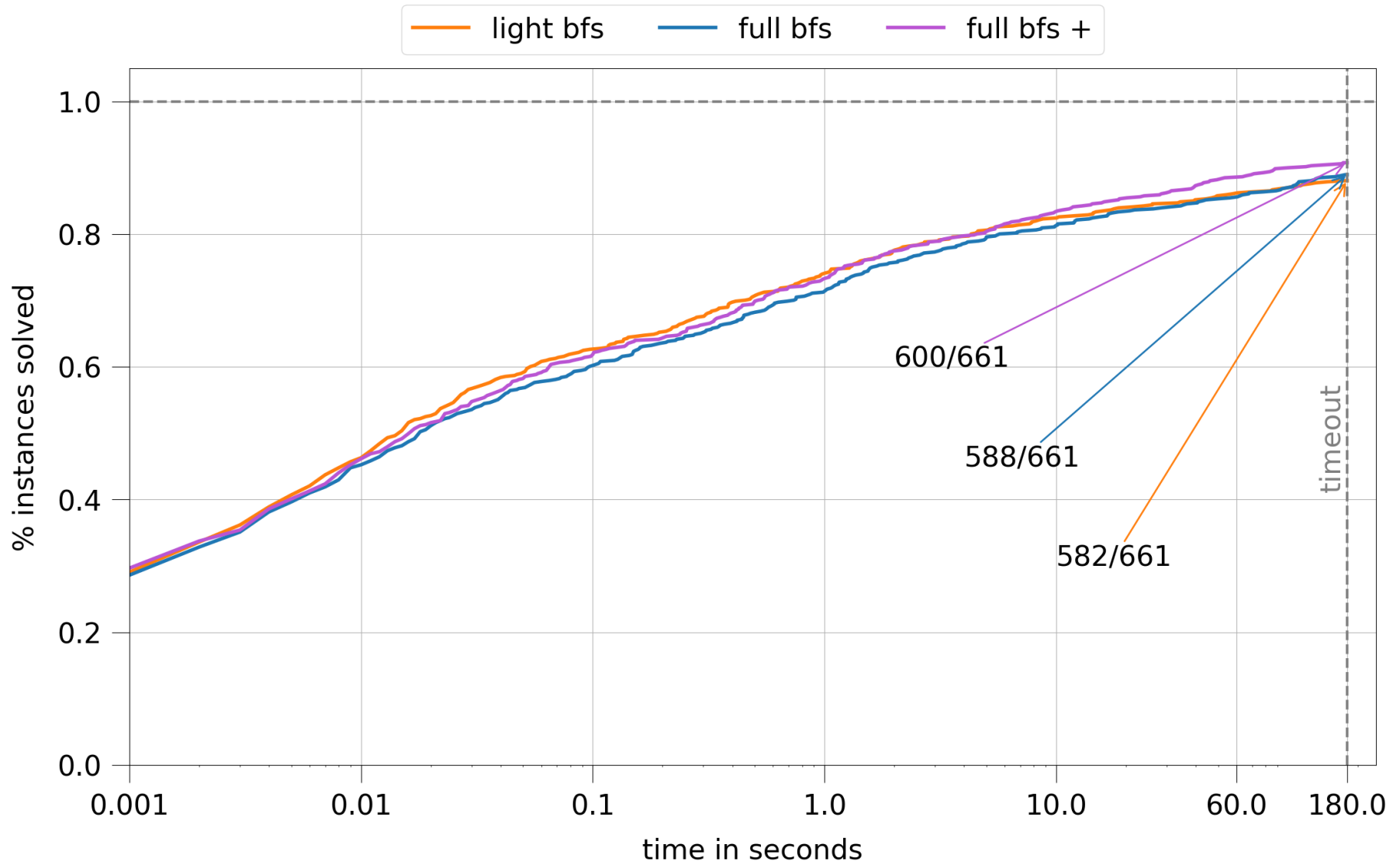
Finding cycles



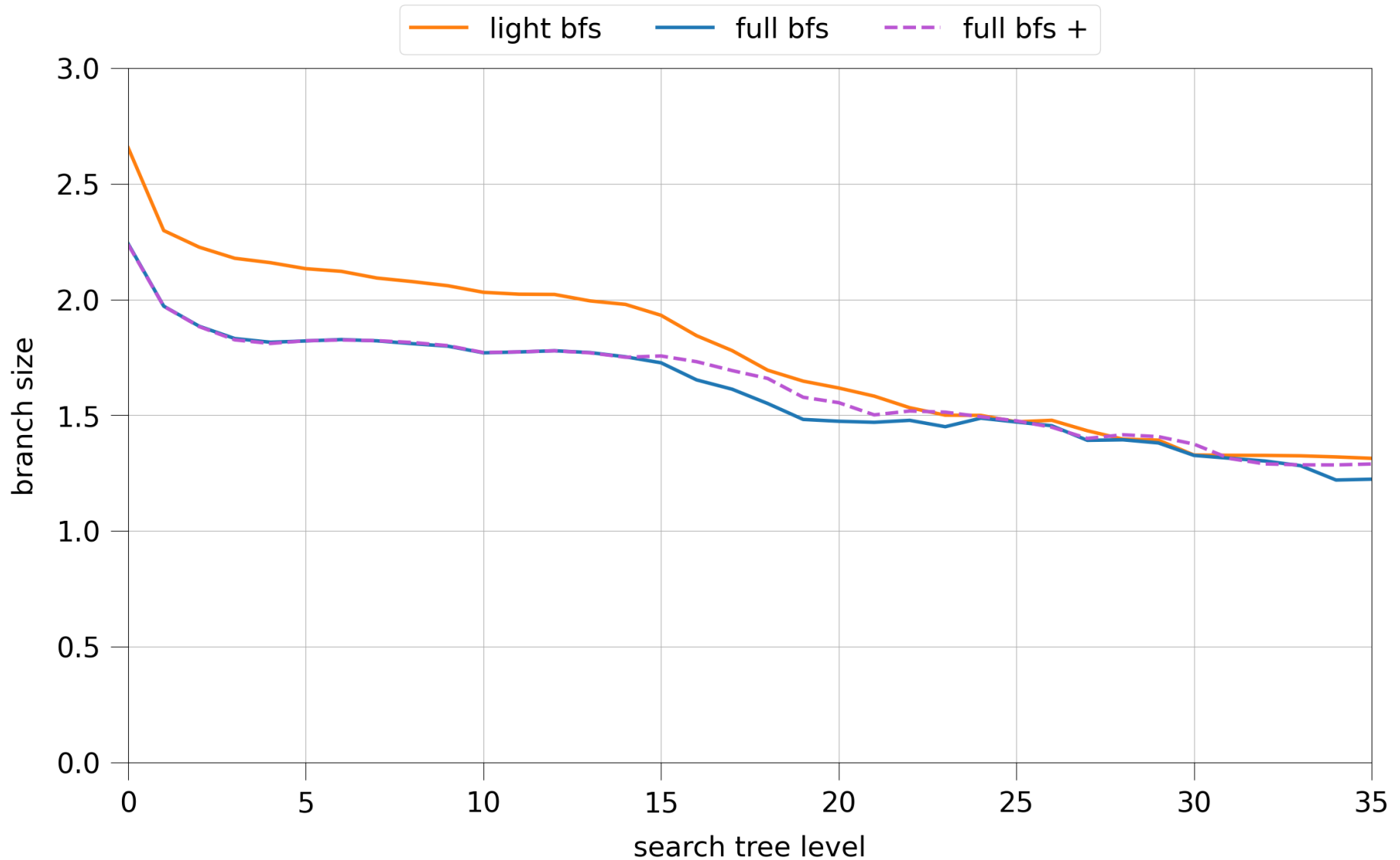
Finding cycles with BFS

- They can be found with breadth-first search
- We tried out three approaches:
 1. Full BFS: in every node of the search tree, run a BFS from every node of the graph and pick the best cycle → always finds the best cycle, high runtime
 2. Full BFS + (with caching): calculate the best cycle for every node only once in the beginning and update it, when their cycle is destroyed → mostly finds the best cycle, medium runtime
 3. Light BFS: In every node of the search tree, run a modified BFS on the graph that returns the first found cycle → not always finds the best cycle, fast runtime
- Light BFS was the best for a long time until we added 1 line of code for the improvement that I just showed you

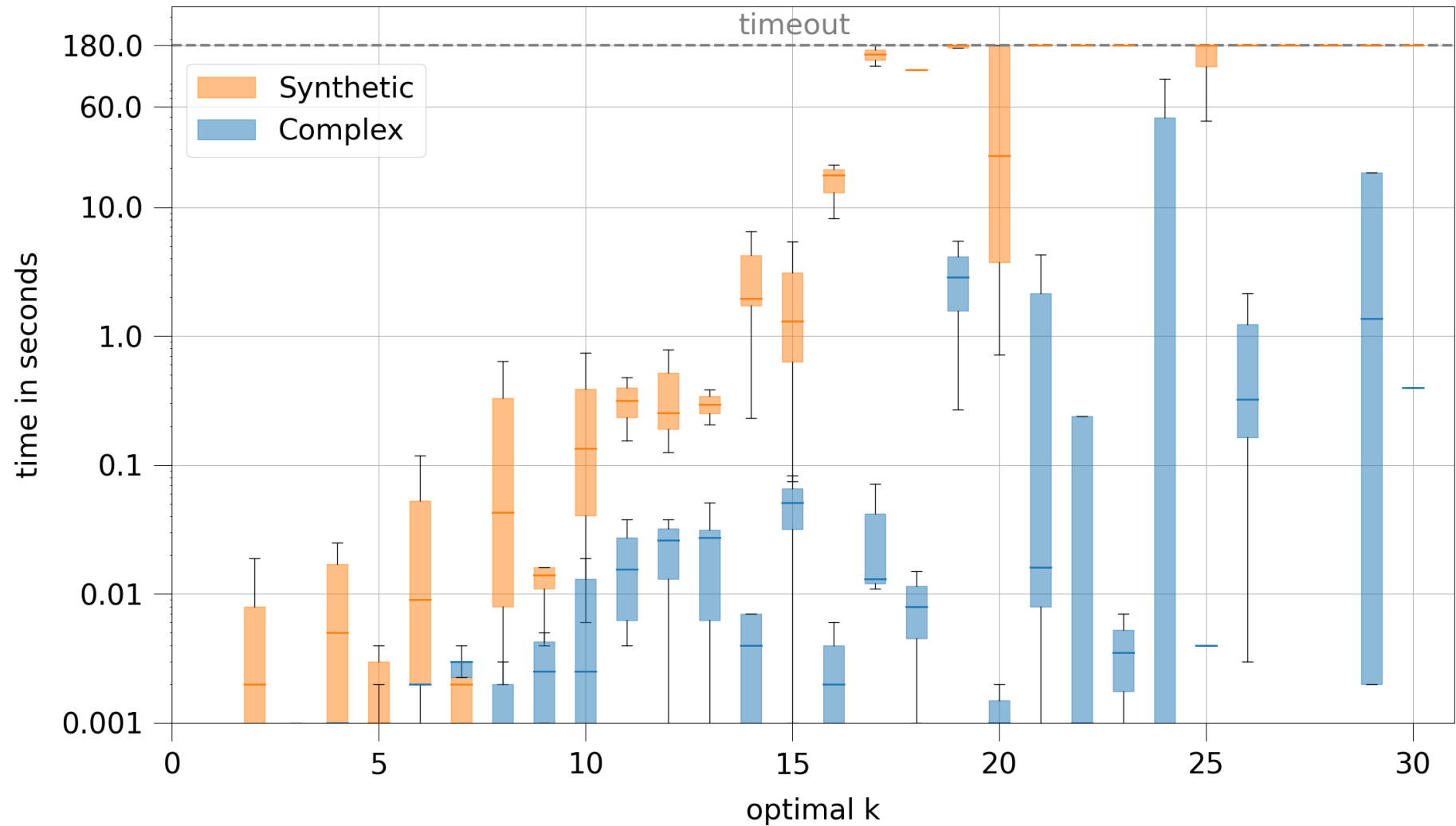
Finding cycles with BFS



Finding cycles with BFS



Summary



Summary & Outlook

- Percentage solved: 88.1% (90.8%)
 - Synthetic: 84.8% (88.7%)
 - Complex: 94.6% (95.0%)
- Goals for next milestone:
 1. Use reductions, cyclic decomposition and flowers in branching
 2. Improve and use the best BFS approach
 3. Use lower bounds (from last lecture)



Do you have questions?