[UFMG] Stenio Garcia (2016-17)

Contents

| 1 | Gran | oh Algorithms 1 |
|---|-------|--------------------------------------|
| | 1.1 | 2-SAT |
| | 1.2 | Kosaraju |
| | 1.3 | Tree Isomorphism |
| | 1.4 | LCA |
| | 1.5 | Bridges and Articulation Points |
| | 1.6 | Eulerian Tour |
| | 1.7 | Floyd Warshall |
| | | |
| 2 | Strin | ags 4 |
| - | 2.1 | Aho-Corasick |
| | 2.2 | KMP |
| | 2.3 | Suffix Array |
| | 2.4 | Suffix Array 2 |
| | 2.4 | Suffix Array Disulguinha |
| | 2.6 | |
| | 2.6 | Manacherś Algorithm |
| _ | | |
| 3 | | nerical Algorithms 6 |
| | 3.1 | Fast Fourier Transform |
| | 3.2 | Fast Fourier Transform 2 |
| | 3.3 | Simpsoné Algorithm |
| | 3.4 | Matrix Exponentiation |
| | | |
| 4 | Matl | hematics 8 |
| | 4.1 | Big Number |
| | 4.2 | Big Number 2 |
| | 4.3 | Chinese Remainder |
| | 4.4 | Chinese Remainder 2 |
| | 4.5 | Matrix Exponentiation |
| | 4.6 | Pascal Triangle |
| | 4.7 | Eulers Totient Function |
| | 4.8 | Pollard Rho |
| | 4.9 | Sieve of Eratosthenes |
| | 4.10 | Extended Euclidean Algorithm |
| | 4.11 | Multiplicative Inverse |
| | 4.12 | Multiplicative Inverse 2 |
| | 4.13 | Linear Diophantine Equation |
| | | |
| 5 | Com | binatorial Optimization 13 |
| 9 | 5.1 | Dinic |
| | 5.2 | Hopcroft-Karp - Bipartite Matching |
| | 5.3 | Max Bipartite Matching 2 |
| | 5.4 | Min Cost Matching |
| | 5.5 | Min Cost Max Flow |
| | 5.6 | Min Cost Max Flow 2 |
| | 5.7 | Edmonds Karp |
| | 3.7 | Edinords Karp |
| c | D | and December 16 |
| 6 | - | amic Programming 16 |
| | 6.1 | Convex Hull Trick |
| | 6.2 | Convex Hull Trick 2 |
| | 6.3 | Divide-and-Conquer |
| | 6.4 | LIS - Longest Increasing Subsequence |
| _ | _ | |
| 7 | Geor | metry 18 |
| | 7.1 | Convex Hull - Monotone Chain |
| | 7.2 | Minimum Enclosing Circle |
| | 7.3 | Minimum Enclosing Circle 2 |
| | 7.4 | Fast Geometry in Cpp |
| | | |
| 8 | Data | Structures 21 |
| | 8.1 | Disjoint Set Union |
| | 8.2 | Persistent Segment Tree |
| | | |

```
      8.3
      RMQ of Indices
      22

      8.4
      RSQ with Lazy-Propagation
      22

      8.5
      Segment Tree
      22

      8.6
      Sparse Table
      23
```

1 Graph Algorithms

1.1 2-SAT

```
//RODAR O COMPONENTE FORTEMENTE CONECTADO
//RECUPERAR NA ORDEM DE descarga DA PILHA
//CONFORME O EXEMPLO A SEGUIR
const int N = 510;
vi graph[N], rev[N];
int us[N];
stack<int> pilha;
int resposta[N];
void dfs1(int u)
  us[u] = 1;
  for (int v : graph[u])
    if (!us[v]) dfs1(v);
  pilha.push(u);
void dfs2(int u, int color)
  us[u] = color;
  for (int v : rev[u])
    if (!us[v]) dfs2(v, color);
  for (int i = 0; i < n; i++)
    if (!us[i]) dfs1(i);
  int color = 1;
  memset(us, 0, sizeof(us));
  while (!pilha.empty()) {
    int topo = pilha.top();
    r.pb(topo);
    pilha.pop();
    if (!us[topo]) dfs2(topo, color++);
  for (int i = 0; i < n; i += 2) {
  if (us[i] == us[i + 1]) return 0;</pre>
   memset (resposta, -1, sizeof (resposta));
  for (int i = r.size() - 1; i >= 0; i--) {
    int vert = r[i] / 2;
    int ok = r[i] % 2;
    if (resposta[vert] == -1) resposta[vert] = !ok;
  return 1:
inline void add(int u, int v)
  graph[u].pb(v);
  rev[v].pb(u);
inline int pos(int u) { return 2 * u; }
inline int neg(int u) { return 2 * u + 1; }
```

1.2 Kosaraju

```
//Retorna os componentes fortemente conectados
//Se o usados[i]=usados[j], temos que i e j
//pertencem ao mesmo componente, col-1= numero
//de componentes fortemente conectados do grafo
class kosaraju {
    private:
    vi usados;
    vvi graph;
    vvi trans;
    vi pilha;

public:
    kosaraju(int N)
    {
        graph.resize(N);
        trans.resize(N);
}
```

```
void AddEdge(int u, int v)
    graph[u].pb(v);
    trans[v].pb(u);
  void dfs(int u, int pass, int color)
    usados[u] = color;
    vi vizinhos;
    if (pass == 1)
      vizinhos = graph[u];
      vizinhos = trans[u];
    for (int j = 0; j < vizinhos.size(); j++) {
  int v = vizinhos[j];</pre>
      if (usados[v] == 0) {
        dfs(v, pass, color);
    pilha.pb(u);
  int SSC(int n)
    pilha.clear();
    usados.assign(n, 0);
    for (int i = 0; i < n; i++) {
  if (!usados[i]) dfs(i, 1, 1);</pre>
    usados.assign(n, 0);
    int color = 1;
    for (int i = n - 1; i >= 0; i--) {
      if (usados[pilha[i]] == 0) {
        dfs(pilha[i], 2, color);
        color++;
    return color - 1;
};
```

1.3 Tree Isomorphism

```
//Seque no main um exemplo de utilizacao do isomorfismo de arvore!
vvichildren, subtreeLabels, tree, L;
vipred, map;
int n;
boolcompare(int a, int b) {
  return subtreeLabels[a] <subtreeLabels[b];
boolequals(int a, int b) {
  return subtreeLabels[a] == subtreeLabels[b];
voidgenerateMapping(intr1, int r2) {
     map.resize(n);
     map[r1] = r2 - n;
     sort (children[r1].begin(), children[r1].end(), compare);
     sort (children[r2].begin(), children[r2].end(), compare);
for (int i = 0; i < (int) children[r1].size(); i++) {</pre>
            int u = children[r1][i];
             int v = children[r2][i];
        generateMapping(u, v);
vifindCenter(int offset = 0) {
     int cnt = n;
     vi a:
     vi deg(n);
      for (int i = 0; i < n; i++) {
            deg[i] = tree[i + offset].size();
             i f (deg[i] <= 1) {
                     a .push_back(i + offset);
     while (cnt > 0) {
            vi na;
             for (int i = 0; i < (int) a.size(); i++) {
                    int u = a[i];
                    for (int j = 0; j < (int) tree[u].size(); j++) {
   int v = tree[u][j];
   if (-deg[v - offset] == 1) {</pre>
                                    n a .push_back(v);
                                      --cnt;
```

```
= na;
    return a;
int dfs (int u, int p = -1, int depth = 0) {
       L [depth].push_back(u);
      for (int i = 0; i < (int) tree[u].size(); i++) {
            int v = tree[u][i];
if (v == p)
            continue;
pred[v] = u;
          c h i l d r e n [u].push_back(v);
    h = max(h, dfs(v, u, depth + 1));
     return h + 1;
boolrootedTreeIsomorphism(int r1, int r2) {
       L .assign(n, vi());
    pred.assign(2 * n, -1);
children.assign(2 * n, vi());
     int h1 = dfs(r1);
int h2 = dfs(r2);
      if (h1!=h2)
    return false;
int h = h1 - 1;
vi label(2 * n);
  s u b t r e e L a b e l s [pred[v]].push_back(label[v]);
             for (int j = 0; j < (int) L[i].size(); j++) {
   int v = L[i][j];</pre>
                    s o r t (subtreeLabels[v].begin(), subtreeLabels[v].end());
            sort (L[i].begin(), L[i].end(), compare);
             for (int j = 0, cnt = 0; j < (int) L[i].size(); j++) { if (j && !equals(L[i][j], L[i][j - 1])) } ++cnt;
                   l a b e l [L[i][j]] = cnt;
       i f (!equals(r1, r2))
          return false;
   generateMapping(r1, r2);
    return true:
booltreeIsomorphism() {
     vi c1 = findCenter();
vi c2 = findCenter(n);
      i f (c1.size() == c2.size()) {
             i f (rootedTreeIsomorphism(c1[0], c2[0]))
                return true;
            else if (c1.size() > 1)
           return rootedTreeIsomorphism(c1[1], c2[0]);
    return false;
int main() {
     n = 5;
vvi t1(n);
     t 1 [0].push_back(1);
      t 1 [1].push_back(0);
      t 1 [1].push_back(2);
      t 1 [2].push_back(1);
      t 1 [1].push_back(3);
      t 1 [3].push back(1);
      t 1 [0].push back(4);
      t 1 [4].push_back(0);
     vvi t2(n);
      t 2 [0].push_back(1);
      t 2 [1].push_back(0);
      t 2 [0].push_back(4);
      t 2 [4].push_back(0);
      t 2 [4].push_back(3);
      t 2 [3].push_back(4);
      t 2 [4].push_back(2);
      t 2 [2] .push_back(4);
     tree.assign(2 * n, vi());
      for (int u = 0; u < n; u++) {
    for (int i = 0; i < t1[u].size(); i++) {
        int v = t1[u][i];
                    tree[u].push_back(v);
              for (int i = 0; i < t2[u].size(); i++) {
                     int v = t2[u][i];
```

t r e e [u + n].push_back(v + n);

```
}
bool res =treeIsomorphism();
cout << res << endl;
if (res)
    for (int i = 0; i < n; i++)
    cout << map[i] << endl;
}</pre>
```

1.4 LCA

```
//antes de usar as queries de lca, e etc..
 //certifique-se de chamar a dfs da arvore e
 //process()
const int N = 100000;
const int M = 22;
int P[N][M];
int big[N][M], low[N][M], level[N];
vii graph[N];
void dfs(int u, int last, int 1)
   level[u] = 1;
   P[u][0] = last;
   for (ii v : graph[u])
      if (v.first != last) {
        big[v.first][0] = low[v.first][0] = v.second;
dfs(v.first, u, 1 + 1);
void process()
  for (int j = 1; j < M; j++)
  for (int i = 1; i <= n; i++) {
    P[i][j] = P[P[i][j-1]][j-1];
    big[i][j] = max(big[i][j-1], big[P[i][j-1]][j-1]);
    low[i][j] = min(low[i][j-1], low[P[i][j-1]][j-1]);</pre>
int lca(int u, int v)
   \textbf{if} \ (\texttt{level}[\texttt{u}] \ \leq \ \texttt{level}[\texttt{v}]) \ \ \mathsf{swap}(\texttt{u}, \ \texttt{v}) \,;
   for (int i = M - 1; i >= 0; i--)
  if (level[u] - (1 << i) >= level[v]) u = P[u][i];
   if (u == v) return u;
for (int i = M - 1; i >= 0; i--) {
   if (P[u][i] != P[v][i]) u = P[u][i], v = P[v][i];
   return P[u][0];
int maximum(int u, int v, int x)
   for (int i = M - 1; i >= 0; i--)
     if (level[u] - (1 << i) >= level[x]) {
        resp = max(resp, big[u][i]);
        u = P[u][i];
   for (int i = M - 1; i >= 0; i--)
  if (level[v] - (1 << i) >= level[x]) {
    resp = max(resp, big[v][i]);
        v = P[v][i];
   return resp:
int minimum(int u, int v, int x)
   for (int i = M - 1; i >= 0; i--)
      if (level[u] - (1 << i) >= level[x]) {
        resp = min(resp, low[u][i]);
        u = P[u][i];
   for (int i = M - 1; i >= 0; i--)
if (level[v] - (1 << i) >= level[x]) {
        resp = min(resp, low[v][i]);
        v = P[v][i];
   return resp;
```

1.5 Bridges and Articulation Points

```
class ponte {
private:
  vvi graph;
  vi usados;
  vi e_articulacao;
  vi dfs_low;
  vi dfs_prof.
  vector<ii>> pontes;
  int tempo;
 public:
  ponte(int N)
    graph.clear():
    graph.resize(N);
    usados.assign(N, 0);
    dfs_low.assign(N, 0);
    dfs_prof.assign(N, 0);
    e_articulacao.assign(N, 0);
    tempo = 0;
  void AddEdge(int u, int v)
    graph[u].pb(v);
    graph[v].pb(u);
  void dfs(int u, int pai)
    usados[u] = 1:
    int nf = 0;
    dfs_low[u] = dfs_prof[u] = tempo++;
    for (int v : graph[u]) {
      if (!usados[v]) {
        dfs(v, u);
        if (dfs_low[v] >= dfs_prof[u] and pai != -1) e_articulacao[u] = true;
        if (pai == -1 and nf > 1) e_articulacao[u] = true;
        if (dfs_low[v] > dfs_prof[u]) pontes.pb(mp(u, v));
        dfs_low[u] = min(dfs_low[u], dfs_low[v]);
      else if (v != pai)
        dfs_low[u] = min(dfs_low[u], dfs_prof[v]);
  void olha as pontes()
    for (int i = 0; i < graph.size(); i++)</pre>
      if (!usados[i]) dfs(i, -1);
    if (pontes.size() == 0)
      cout << " Que merda! nao tem ponte!" << endl;</pre>
    else {
      for (ii i : pontes) cout << i.first << " " << i.second << endl;</pre>
  void olha_as_art()
    for (int i = 0; i < graph.size(); i++)</pre>
    if (!usados[i]) dfs(i, -1);
for (int i = 0; i < e_articulacao.size(); i++)</pre>
      if (e_articulacao[i]) cout << " OIAAA A PONTE " << i << endl;</pre>
```

1.6 Eulerian Tour

```
multiset<int> graph[N];
stack*int> path;
// -> It suffices to call dfs1 just
// one time leaving from node 0.
// -> To calculate the path,
// call the dfs from the odd degree node.
// -> O(n + log(n))
void dfs1(int u)
{
  while(graph[u].size())
  {
    int v = *graph[u].begin();
    graph[u].erase(graph[u].begin());
    graph[v].erase(graph[v].find(u));
    dfs1(v);
  }
  path.push(u);
}
```

1.7 Floyd Warshall

2 Strings

2.1 Aho-Corasick

```
//N= tamanho da trie, M tamanho do alfabeto
int to[N][M], Link[N], fim[N];
int idx = 1;
void add_str(string &s)
   int v = 0;
  for (int i = 0; i < s.size(); i++) {
  if (!to[v][s[i]]) to[v][s[i]] = idx++;</pre>
      v = to[v][s[i]];
   fim[v] = 1;
void process()
   queue<int> fila;
   fila.push(0);
   while (!fila.empty())
     int cur = fila.front();
      fila.pop();
     int 1 = Link[cur];
fim[cur] |= fim[1];
for (int i = 0; i < 200; i++) {</pre>
        if (to[cur][i]) {
  if (cur != 0) {
               \label{eq:link_loss} \operatorname{Link}\left[\operatorname{to}\left[\operatorname{cur}\right]\left[\operatorname{i}\right]\right] \; = \; \operatorname{to}\left[\operatorname{1}\right]\left[\operatorname{i}\right];
            else
              Link[to[cur][i]] = 0;
            fila.push(to[cur][i]);
            to[cur][i] = to[1][i];
int resolve(string &s)
   int v = 0, r = 0;
   for (int i = 0; i < s.size(); i++) {</pre>
      v = to[v][s[i]];
      if (fim[v]) r++, v = 0;
   return r;
```

2.2 KMP

```
int b[100000];
int sizet, sizep;
void kmpPreprocess(string &text, string &pattern) {
  int i = 0, j = -1;
  b[0] = -1;
  while (i < sizep) {</pre>
```

```
while (j >= 0 and pattern[i] != pattern[j]) j = b[j];
    i++, j++;
    b[i] = j;
}

void kmpSearch(string &text, string &pattern)
{
    kmpPreprocess(text, pattern);
    int i = 0, j = 0;
    while (i < sizet) {
        while (j >= 0 and text[i] != pattern[j]) j = b[j];
        i++, j++;
    if (j == sizep) {
        cout << "Otha a substring do texto " << i - j << endl;
        j = b[j];
    }
}</pre>
```

2.3 Suffix Array

```
* O(nlog^2(n)) para o sufix array
 * O(logn) para o LCP(i,j)
 * LCP de i para j;
struct SA (
  const int L:
  string s;
  vvi P:
  vector<pair< ii,int> > M;
  SA(const string &s) : L(s.size()), s(s), P(1, vi(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] =s[i]-a';
for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
      P.pb(vi(L, 0));
for (int i = 0; i < L; i++)
        M[i] = mp(mp(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
       sort(M.begin(), M.end());
       for (int i = 0; i < L; i++)
          P[level][M[i].second] = (i > 0 \&\& M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i; 
  vi GetSA() {
    vi v=P.back();
    vi ret(v.size());
    for (int i=0; i < v.size(); i++) {</pre>
      ret[v[i]]=i;
    return ret;
  int LCP(int i, int j) {
    inf i == j) return L - i;
for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
    if (P[k][i] == P[k][j]) {</pre>
         i += 1 << k;
         i += 1 << k;
         len += 1 << k;
    return len;
  vi GetLCP(vi &sa)
     vi lcp(sa.size()-1);
    for (int i=0;i<sa.size()-1;i++) {</pre>
       lcp[i]=LCP(sa[i],sa[i+1]);
    return lcp;
};
```

2.4 Suffix Array 2

```
/***
Suffix Array. Builing works in O(NlogN).
Also LCP array is calculated in O(NlogN).
This code counts number of different substrings in the string.
Based on problem I from here: http://codeforces.ru/gym/100133
```

```
#include <iostream>
#include <fstream>
#include <cmath>
#include <algorithm>
#include <vector>
#include <set>
#include <map>
#include <stack>
#include <queue>
#include <cstdlib>
#include <cstdio>
#include <string>
#include <cstring>
#include <cassert>
#include <utility>
#include <iomanip>
using namespace std;
const int MAXN = 205000;
const int ALPH = 256;
const int MAXLOG = 20;
int n;
char s[MAXN];
int p[MAXN]; // suffix array itself
int pour[MAXN]:
int c[MAXN] [MAXLOG];
int num[MAXN];
int classesNum;
int lcp[MAXN];
void buildSuffixArray() {
  for (int i = 0; i < n; i++)
    num[s[i]]++;
  for (int i = 1; i < ALPH; i++)
  num[i] += num[i - 1];</pre>
  for (int i = 0; i < n; i++) {
   p[num[s[i]] - 1] = i;</pre>
    num[s[i]]--;
  c[p[0]][0] = 1;
  for (int i = 1; i < n; i++) {
    if (s[p[i]] != s[p[i - 1]])
       classesNum++;
    c[p[i]][0] = classesNum;
  for (int i = 1; ; i++) {
    int half = (1 << (i - 1));</pre>
    for (int j = 0; j < n; j++) {
      pcur[j] = p[j] - half;
if (pcur[j] < 0)</pre>
         pcur[j] += n;
    for (int j = 1; j <= classesNum; j++)</pre>
      num[j] = 0;
    for (int j = 0; j < n; j++)
  num[c[pcur[j]][i - 1]]++;
for (int j = 2; j <= classesNum; j++)</pre>
      num[j] += num[j-1];
    for (int j = n - 1; j >= 0; j--) {
      p[num[c[pcur[j]][i - 1]] - 1] = pcur[j];
       num[c[pcur[j]][i - 1]]--;
    c[p[0]][i] = 1;
    for (int j = 1; j < n; j++) {
   int p1 = (p[j] + half) % n, p2 = (p[j - 1] + half) % n;
   if (c[p[j])[i - 1] != c[p[j - 1]][i - 1] || c[p1][i - 1] != c[p2][i - 1])</pre>
         classesNum++:
       c[p[j]][i] = classesNum;
    if ((1 << i) >= n)
      break;
```

```
for (int i = 0; i < n; i++)
   p[i] = p[i + 1];
int getLcp(int a, int b) {
  for (int i = MAXLOG - 1; i >= 0; i--) {
   int curlen = (1 << i);</pre>
   if (curlen > n)
     continue;
   if (c[a][i] == c[b][i]) {
     res += curlen;
      a += curlen;
     b += curlen;
  return res;
void calcLcpArray() {
 for (int i = 0; i < n - 1; i++)
    lcp[i] = getLcp(p[i], p[i + 1]);
int main() {
  assert(freopen("substr.in", "r", stdin));
  assert (freopen ("substr.out", "w", stdout));
  gets(s):
 n = strlen(s);
 buildSuffixArray();
  // Now p from 0 to n - 1 contains suffix array of original string
  /*for (int i = 0; i < n; i++) {
   printf("%d ", p[i] + 1);
   calcLcpArray();
  long long ans = 0;
  for (int i = 0; i < n; i++)
   ans += n - p[i];
  for (int i = 1; i < n; i++)
   ans -= lcp[i - 1];
  cout << ans << endl;
  return 0:
```

2.5 Suffix Array Disulguinha

```
#include <iostream>
#include <string>
#include <vector>
#include <algorithm>
#include <cstring>
using namespace std;
struct SuffixArray{
  const string& s;
  int n;
  vector<int> order, rank, lcp;
  vector<int> count, x, y;
  SuffixArray(const string& s) : s(s), n(s.size()), order(n), rank(n),
                                    count(n + 1), x(n), y(n), lcp(n) {
    build();
    buildLCP();
  void build() {
    //sort suffiixes by the first character
    for(int i = 0; i < n; i++) order[i] = i;</pre>
    sort(order.begin(), order.end(), [&](int a, int b){return s[a] < s[b];});</pre>
    rank[order[0]] = 0;
    for(int i = 1; i < n; i++) {
  rank[order[i]] = rank[order[i - 1]];</pre>
      if(s[order[i]] != s[order[i - 1]]) rank[order[i]]++;
```

```
//sort suffixex by the the first 2*p characters, for p in 1, 2, 4, 8,...
    for(int p = 1; p < n, rank[order[n - 1]] < n - 1; p <<= 1){</pre>
      for (int i = 0; i < n; i++) {
       x[i] = rank[i];
       y[i] = i + p < n ? rank[i + p] + 1 : 0;
      radixPass(y);
      radixPass(x);
      rank[order[0]] = 0;
      for(int i = 1; i < n; i++) {
  rank[order[i]] = rank[order[i - 1]];</pre>
        if(x[order[i]] != x[order[i - 1]] or y[order[i]] != y[order[i - 1]]) rank[order[i]]++;;
  //Stable counting sort
  void radixPass(vector<int>& key) {
    fill(count.begin(), count.end(), 0);
    for(auto index : order) count[key[index]]++;
    for(int i = 1; i <= n; i++) count[i] += count[i - 1];</pre>
    order.swap(lcp);
  //Kasai's algorithm to build the LCP array from order, rank and s
  //For i \ge 1, lcp[i] refers to the suffixes starting at order[i] and order[i - 1]
  void buildLCP() {
    lcp[0] = 0;
   int k = 0;
for (int i = 0; i < n; i++) {
     if(rank[i] == n - 1){
       int j = order[rank[i] + 1];
        while (i + k < n \text{ and } j + k < n \text{ and } s[i + k] == s[j + k]) k++;
        lcp[rank[j]] = k;
       if(k) k--;
int main(){
  ios::sync_with_stdio(false);
  cin >> s;
  SuffixArray sa(s);
  for(int i = 0; i < s.size(); i++) cout << sa.order[i] << '\n';</pre>
```

2.6 Manacherś Algorithm

```
Manacher's algorithm for finding all subpalindromes in the string.
  Based on problem L from here: http://codeforces.ru/gym/100133
const int MAXN = 105000;
string s;
int n;
int odd[MAXN], even[MAXN];
int 1. r:
long long ans;
int main() {
 assert(freopen("palindrome.in","r",stdin));
 assert (freopen ("palindrome.out", "w", stdout));
  getline(cin, s);
 n = (int) s.length();
  for (int i = 0; i < n; i++) {</pre>
   int cur = 1;
   if (i < r)
     cur = min(r - i + 1, odd[1 + r - i]);
   while (i + cur < n && i - cur >= 0 && s[i - cur] == s[i + cur])
     cur++;
   odd[i] = cur;
   if (i + cur - 1 > r) {
     1 = i - cur + 1;
```

```
r = i + cur - 1;
// Even case
for (int i = 0; i < n; i++) {
  int cur = 0;
  if (i < r)
    cur = min(r - i + 1, even[1 + r - i + 1]);
  while (i + cur < n \&\& i - 1 - cur >= 0 \&\& s[i - 1 - cur] == s[i + cur])
   cur++:
  even[i] = cur;
 if (i + cur - 1 > r) {
   1 = i - cur:
    r = i + cur - 1;
for (int i = 0; i < n; i++) {
 if (odd[i] > 1) {
    ans += odd[i] - 1;
 if (even[i])
    ans += even[i];
cout << ans << endl:
return 0:
```

3 Numerical Algorithms

3.1 Fast Fourier Transform

```
// FFT - The Iterative Version
// Running Time:
    O(n*log n)
// How To Use:
    fft(a,1)
// fft(b,1)
    mul(a,b)
    fft(a,-1)
// INPUT:
   - fft method:
       * The vector representing the polynomial
       * 1 to normal transform
       * -1 to inverse transform
   - mul method:
       * The two polynomials to be multiplyed
// OUTPUT:
// - fft method: Transforms the vector sent.
// - mul method: The result is kept in the first vector.
   - You can either use the struct defined of define dificil as complex<double>
// SOLVED:
// * Codeforces Round #296 (Div. 1) D. Fuzzy Search
struct dificil {
       double real;
       double im;
       dificil() {
               real=0.0;
                im=0.0;
       dificil(double real, double im):real(real),im(im){}
       dificil operator+(const dificil &o)const {
                return dificil(o.real+real, im+o.im);
       dificil operator/(double v) const
                return dificil (real/v, im/v);
       dificil operator * (const dificil &o) const {
```

```
return dificil(real*o.real-im*o.im, real*o.im+im*o.real);
         dificil operator-(const dificil &o) const {
                  return dificil(real-o.real, im-o.im);
};
dificil tmp[MAXN*2];
int coco, maiorpot2[MAXN];
void fft(vector<dificil> &A, int s)
         int n = A.size(), p = 0;
         while (n>1) {
                  n >>= 1;
         \mathbf{n} = (1 << \mathbf{p});
         vector<dificil> a=A:
         for (int i = 0; i < n; ++i) {
                  int rev = 0;
                  for(int j = 0; j < p; ++j) {
                           rev <<= 1;
                           rev |= ( (i >> j) & 1 );
                  A[i] = a[rev];
         dificil w, wn;
         for (int i = 1; i \le p; ++i) {
                  int M = 1 << i;</pre>
                  int K = M >> 1;
                  wn = dificil(\cos(s*2.0*pi/(double)M), \sin(s*2.0*pi/(double)M));
                  for(int j = 0; j < n; j += M) {
    w = dificil(1.0, 0.0);</pre>
                           for(int 1 = j; 1 < K + j; ++1) {
    dificil t = w;</pre>
                                     t = t * A[1 + K];
                                     dificil u = A[1];
                                     A[1] = A[1] + t;
                                    u = u-t;
A[1 + K] = u;
         if(s==-1){
                  for (int i = 0; i < n; ++i)
                           A[i] = A[i]/(double)n;
void mul(vector<dificil> &a, vector<dificil> &b)
         for (int i=0; i < a.size(); i++)</pre>
                  a[i]=a[i]*b[i];
```

3.2 Fast Fourier Transform 2

```
// FFT - The Recursive Version
// Running Time:
// O(n*log n)
//
// How To Use:
    fft(&a[0], tam, 0)
    fft(&b[0], tam, 0)
// mul(a,b)
// fft(&a[0], tam, 1)
//
// INPUT:
// - fft method:
// * The vector representing the polynomial
* to to normal transform
// * 1 to inverse transform
// - mul method:
// * The two polynomials to be multiplyed
```

```
// - fft method: Transforms the vector sent.
// - mul method: The result is kept in the first vector.
// - Tam has to be a power of 2.
// - You can either use the struct defined of define dificil as complex<double>
// SOLVED:
// * Codeforces Round #296 (Div. 1) D. Fuzzy Search
dificil tmp[MAXN*2];
int coco, maiorpot2[MAXN];
void fft(dificil *v, int N, bool inv)
        if(N<=1) return;</pre>
        dificil *vodd = v;
        dificil *veven = v+N/2;
        for(int i=0; i<N; i++) tmp[i] = v[i];</pre>
        for (int i=0; i<N; i+=2)
                veven[coco] = tmp[i];
                vodd[coco] = tmp[i+1];
                coco++;
        fft(&vodd[0], N/2, inv);
        fft(&veven[0], N/2, inv);
        double angucomleite = 2.0*PI/(double)N;
        if(inv) angucomleite = -angucomleite;
        dificil wn(cos(angucomleite), sin(angucomleite));
        for (int i=0; i<N/2; i++)</pre>
                 tmp[i] = veven[i]+w*vodd[i];
                 tmp[i+N/2] = veven[i]-w*vodd[i];
                 w \neq wn;
                if(inv)
                         tmp[i] /= 2;
                         tmp[i+N/2] /= 2;
        for(int i=0; i<N; i++) v[i] = tmp[i];</pre>
void mul(vector<dificil> &a, vector<dificil> &b)
  for(int i=0; i<a.size(); i++)</pre>
    a[i] = a[i] *b[i];
void precomp()
  for (int i=1; i < MAXN; i++)</pre>
    if((1<<pot)<i) pot++;</pre>
    maiorpot2[i] = (1<<pot);</pre>
```

3.3 Simpsonś Algorithm

```
const int NPASSOS = 100000;
const int W=1000000;
//W= tamanho do intervalo que eu estou integrando
double integral1()
{
   double h = W / (NPASSOS);
   double a = 0;
   double b = W;
   double s = f(a) + f(b);
   for (double i = 1; i <= NPASSOS; i += 2) s += f(a + i * h) * 4.0;
   for (double i = 2; i <= (NPASSOS - 1); i += 2) s += f(a + i * h) * 2.0;
   return s * h / 3.0;
}</pre>
```

3.4 Matrix Exponentiation

```
//matmul multiplica m1 por m2
//matpow exponencia a matrix m1 por p
//mul vet multiplica a matrix m1 pelo vetor vet
vvi matmul(vvi &m1, vvi &m2)
  ans.resize(m1.size(), vi(m2.size(), 0));
 for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
for (int k = 0; k < n; k++) {</pre>
        ans[i][j] += m1[i][k] * m2[k][j];
ans[i][j] %= MOD;
  return ans;
vvi matpow(vvi &m1, ll p)
  ans.resize(m1.size(), vi(m1.size(), 0));
  for (int i = 0; i < n; i++) ans[i][i] = 1;
  while (p)
    if (p & 1) ans = matmul(ans, m1);
    m1 = matmul(m1, m1);
    p >>= 1;
  return ans;
// VETOR TEM N LINHAS E A MATRIZ E QUADRADA
vi mulvet(vvi &m1, vi &vet)
  vi ans;
  ans.resize(vet.size(), 0);
  for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++) {
   ans[i] += (m1[i][j] * vet[j]);</pre>
       ans[i] %= MOD;
  return ans:
```

4 Mathematics

4.1 Big Number

```
void zero_esq(string &resp)
  string retorno = resp:
  reverse (retorno.begin(), retorno.end());
  int i = resp.size() - 1;
  while (retorno[i] == '0' and i > 0) {
   retorno.erase(i);
  reverse(retorno.begin(), retorno.end());
string sum_big(string a, string b)
  string resp;
  reverse(a.begin(), a.end());
  reverse(b.begin(), b.end());
  if (a.size() <= b.size()) {
    int carry = 0;
    int carry = 0,
for (int i = 0; i < a.size(); i++) {
  int x = b[i] - '0' + a[i] - '0' + carry;
  resp.push_back((char)(x % 10 + '0'));</pre>
      carry = x / 10;
    for (int i = a.size(); i < b.size(); i++) {</pre>
      int x = b[i] - '0' + carry;
      resp.push_back((char)(x % 10 + '0'));
      carry = x / 10;
    if (carry > 0) resp.push_back((char)(carry + '0'));
  else {
    int carry = 0;
    for (int i = 0; i < b.size(); i++) {
      int x = a[i] - '0' + b[i] - '0' + carry;
```

```
resp.push_back((char)(x % 10 + '0'));
       carry = x / 10;
     for (int i = b.size(); i < a.size(); i++) {</pre>
       int x = a[i] - '0' + carry;
       resp.push_back((char)(x % 10 + '0'));
       carry = x / 10;
    if (carry > 0) resp.push_back((char)(carry + '0'));
  reverse(resp.begin(), resp.end());
  zero_esq(resp);
  return resp;
string mul_big(string a, string b)
  string resp;
  resp.push_back('0');
  string temp;
  int carry = 0;
  reverse(a.begin(), a.end());
  reverse(b.begin(), b.end());
  for (int i = 0; i < a.size(); i++) {
     temp.clear();
    temp.clear();
for (int k = 0; k < i; k++) temp.push_back('0');
int x = a[i] - '0';
for (int j = 0; j < b.size(); j++) {
   int y = b[j] - '0';
   int novo = (x + y + carry);
   temp.push_back(novo % 10) + '0');</pre>
       carry = novo / 10;
    if (carry > 0) temp.push_back(carry + '0');
    reverse(temp.begin(), temp.end());
    carry = 0;
    resp = sum_big(temp, resp);
  zero_esq(resp);
  return resp;
```

4.2 Big Number 2

```
Structure implementing long arithmetic in C++
  Analogue to BigInteger in Java.
  Tested on many problems.
  TODO: list some problems
struct BigInt {
 vector <int> num;
 static const int base = 1000 * 1000 * 1000:
 static const int baseDigits = 9;
 string leadingZerosModifier;
  /****************
  * CONSTRUCTONS & SETTERS
  ****************
  void setLeadingZerosModifier() {
   leadingZerosModifier = "%0xd";
   leadingZerosModifier[2] = '0' + baseDigits;
 void set(int value) {
   num.clear():
   if (value == 0)
     num.push_back(0);
   while (value) {
    num.push_back(value % base);
     value /= base;
  void set(long long value) {
   num.clear();
   if (value == 0)
     num.push back(0):
   while (value) {
    num.push back(value % base);
     value /= base;
```

```
void set(string &value) {
 num.clear();
 for (int i = (int)value.length() - 1; i >= 0; i -= baseDigits) {
   int add = 0;
   for (int j = max(0, i - baseDigits + 1); j <= i; j++)
     add = add * 10 + value[j] - '0';
    num.push_back(add);
void operator = (int value) {
 set (value);
void operator = (long long value) {
 set (value);
void operator = (string &value) {
 set (value);
BigInt() {
 setLeadingZerosModifier();
 set (0);
BigInt(int value) {
 setLeadingZerosModifier();
 set (value):
BigInt(string value) {
 setLeadingZerosModifier();
 set (value);
/*----
* SIZE METHODS
//returns size of vector
int size() {
 return (int) num.size();
//returns length of the number
int digitNum() {
 int result = 0;
 for (int i = 0; i < (int) num.size() - 1; i++)</pre>
   result += baseDigits;
 int lastNum = num.back();
 while (lastNum) {
   result++:
   lastNum /= 10;
 return result:
* I/O
void read() {
 cin >> s;
 num.clear();
 for (int i = (int)s.length() - 1; i >= 0; i -= baseDigits) {
   int add = 0;
   for (int j = max(0, i - baseDigits + 1); j <= i; j++)
add = add * 10 + s[j] - '0';</pre>
   num.push_back(add);
void print() {
 printf("%d", num.back());
for (int i = (int)num.size() - 2; i >= 0; i--)
   printf (leadingZerosModifier.c_str(), num[i]);
void println() {
 print();
 printf("\n");
string toString() {
 string result = "";
 for (int i = 0; i < (int) num.size(); i++) {</pre>
   int cur = num[i];
    for (int j = 1; j <= baseDigits; j++) {</pre>
     if (cur == 0 && i == (int) num.size() - 1)
```

```
break;
      result.append(1, (char) '0' + cur % 10);
      cur /= 10;
  reverse(result.begin(), result.end());
  return result;
/*========
* ADDITION
void sumThis(BigInt number) {
 int carry = 0;
for (int i = 0; i < max((int)num.size(), number.size()) || carry; i++) {</pre>
    if (i == num.size())
      num.push_back(0);
    if (i >= number.size())
      carry = num[i] + carry;
    else
      carry = num[i] + carry + number.num[i];
    num[i] = carry % base;
   carry /= base;
void sumThis(int number) {
 int carry = number;
for (int i = 0; i < (int)num.size() || carry; i++) {</pre>
   if (i == num.size())
     num.push_back(0);
   carry = num[i] + carry;
num[i] = carry % base;
    carry /= base;
BigInt sum(BigInt number) {
 BigInt result = *this;
  result.sumThis(number):
 return result;
BigInt sum(int number) {
 BigInt result = *this;
  result.sumThis(number);
 return result;
void operator += (BigInt number) {
 sumThis(number);
void operator += (int number) {
 sumThis(number);
BigInt operator + (BigInt number) {
 return sum(number);
BigInt operator + (int number) {
 return sum(number);
* SUBTRACTION
void subThis(BigInt number) {
 int carry = 0;
for (int i = 0; i < (int) number.size() || carry; i++) {</pre>
    if (i < (int)number.size())</pre>
      num[i] -= carry + number.num[i];
    else
      num[i] -= carry;
    if (num[i] < 0) {
      num[i] += base;
    else
      carry = 0;
  while (num.size() > 1 && num.back() == 0)
   num.pop_back();
void subThis(int number) {
  int carry = -number;
  for (int i = 0; carry > 0; i++) {
   num[i] -= carry;
```

```
if (num[i] < 0) {
      carry = 1;
      num[i] += base;
    else
      carry = 0;
  while (num.size() > 1 && num.back() == 0)
    num.pop_back();
BigInt sub(BigInt number) {
   BigInt result = *this;
  result.subThis(number);
  return result;
BigInt sub(int number) {
  BigInt result = *this;
  result.subThis(number);
  return result;
void operator -= (BigInt number) {
  subThis(number);
void operator -= (int number) {
  subThis(number):
BigInt operator - (BigInt number) {
  return sub (number);
BigInt operator - (int number) {
  return sub(number);
* MULTIPLICATION
BigInt mult(BigInt number) {
  BigInt product;
  product.num.resize(num.size() + number.size());
  for (int i = 0; i < (int)num.size(); i++) {</pre>
    for (int j = 0, carry = 0; j < (int)number.size() || carry; j++) {
  long long cur = product.num[i + j] + num[i] * 1ll * (j < (int)number.size() ? number.num[j] :</pre>
            0) + carry;
      product.num[i + j] = int (cur % base);
carry = int (cur / base);
  while (product.size() > 1 && product.num.back() == 0)
    product.num.pop_back();
  return product:
void multThis(BigInt number) {
  *this = mult(number);
void multThis(int number) {
  int carry = 0;
  for (int i = 0; i < (int) num.size() || carry; ++i) {</pre>
    if (i == num.size())
      num.push_back (0);
    long long cur = carry + num[i] * 111 * number;
num[i] = int (cur % base);
carry = int (cur / base);
  while (num.size() > 1 && num.back() == 0)
    num.pop_back();
BigInt mult(int number) {
  BigInt result = *this;
  result multThis(number);
  return result;
void operator *= (BigInt number) {
  multThis(number):
void operator *= (int number) {
  multThis(number);
BigInt operator * (BigInt number) {
  return mult(number);
```

```
BigInt operator * (int number) {
 return mult(number);
void multThisByPowerOfTen(int power) {
  int baseNums = power / baseDigits;
  int curLen = (int)num.size();
  num.resize(curLen + baseNums);
  for (int i = num.size() - 1; i >= baseNums; i--) {
    num[i] = num[i - baseNums];
  for (int i = baseNums - 1; i >= 0; i--)
  num[i] = 0;
power %= baseDigits;
  int multBy = (int)pow(10.0, power);
  multThis(multBy);
+ DIVISION
void divThis(int number) {
 int carry = 0;
for (int i= (int)num.size() - 1; i >= 0; i--) {
   long long cur = num[i] + carry * 111 * base,
    num[i] = int (cur / number);
    carry = int (cur % number);
  while (num.size() > 1 && num.back() == 0)
    num.pop_back();
BigInt div(int number) {
  BigInt result = *this;
  result divThis(number);
 return result;
void operator /= (int number) {
  divThis(number):
BigInt operator / (int number) {
 return div(number);
void divThisByPowerOfTen(int power) {
  int baseNums = power / baseDigits;
  int curLen = (int)num.size();
 for (int i = 0; i < (int) num.size() - baseNums; i++) {
  num[i] = num[i + baseNums];</pre>
  for (int i = 1; i <= baseNums; i++)
   num.pop_back();
  power %= baseDigits;
  int divBy = (int)pow(10.0, power);
  divThis(divBy);
* MODULO
void modThis(int number) {
 int carry = 0;
for (int i= (int)num.size() - 1; i >= 0; i--) {
   long long cur = num[i] + carry * 111 * base;
num[i] = int (cur / number);
    carry = int (cur % number);
  set (carry);
BigInt mod(int number) {
  BigInt result = *this;
  result .modThis(number);
  return result;
void operator %= (int number) {
 modThis(number);
BigInt operator % (int number) {
 return mod(number);
* COMPARISON
```

```
//Returns: -1 - this number is less than argument, 0 - equal, 1 - this number is greater
  int compareTo(BigInt number) {
    if ((int)num.size() < number.size())</pre>
     return -1;
    if ((int)num.size() > number.size())
      return 1;
    for (int i = (int) num.size() - 1; i >= 0; i--) {
     if (num[i] > number.num[i])
        return 1;
     if (num[i] < number.num[i])</pre>
        return -1;
    return 0:
  //Returns: -1 - this number is less than argument, 0 - equal, 1 - this number is greater
  int compareTo(int number) {
    if (num.size() > 1 || num[0] > number)
     return 1;
    if (num[0] < number)</pre>
     return -1;
    return 0:
  bool operator < (BigInt number) {
    return compareTo(number) == -1;
  bool operator < (int number) {</pre>
    return compareTo(number) == -1;
  bool operator <= (BigInt number) {</pre>
    return compareTo(number) != 1;
  bool operator <= (int number) {</pre>
    return compareTo(number) != 1;
  bool operator == (BigInt number) {
    return compareTo(number) == 0;
  bool operator == (int number) {
    return compareTo(number) == 0;
  bool operator > (BigInt number) {
    return compareTo(number) == 1;
  bool operator > (int number) {
    return compareTo(number) == 1;
  bool operator >= (BigInt number) {
    return compareTo(number) != -1;
  bool operator >= (int number) {
    return compareTo(number) != 1;
  bool operator != (BigInt number) {
    return compareTo(number) != 0;
  bool operator != (int number) {
    return compareTo(number) != 0;
};
```

4.3 Chinese Remainder

```
11 expmod(11 a, 11 e, 11 m)
  while (e > 0) {
   if (e % 2 != 0) ret = mulmod(ret, a, m);
   a = mulmod(a, a, m);
    e >>= 1;
  return ret;
11 invmul(11 a, 11 m) { return expmod(a, m - 2, m); }
11 chinese (vector<11> r, vector<11> m)
 int sz = m.size():
  11 M = 1;
  for (int i = 0; i < sz; i++) {
   M *= m[i];
  11 ret = 0;
  for (int i = 0; i < sz; i++) {
   ret += mulmod(mulmod(M / m[i], r[i], M), invmul(M / m[i], M), M);
    ret = ret % M;
  return ret:
```

4.4 Chinese Remainder 2

```
// Chinese remainder theorem (special case): find z such that // z % m1 = r1, z
// % m2 = r2. Here, z is unique modulo M = lcm(m1, m2). // Return (z, M). On
// failure, M = -1;
ii chinese_remainder_theorem(int m1, int r1, int m2, int r2)
 int s, t;
 int g = extended_euclid(m1, m2, s, t);
 if (r1 % g != r2 % g) return mp(0, -1);
 return mp (mod(s * r2 * m1 + t * r1 * m2, m1 * m2) / g, m1 * m2 / g);
// Chinese remainder theorem: find z such that // z % m[i] =
//\ r[i]\ for\ all\ i
     .Note that the solution is unique modulo M = lcm i (m[i]).
    Return(z, M)
     .On // failure, M = -1. Note that we do not require the a[i] s
// to be relatively prime.
ii chinese_remainder_theorem(const vi &m, const vi &r)
  ii ret = make_pair(r[0], m[0]);
  for (int i = 1; i < m.size(); i++) {</pre>
   ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
   if (ret.second == -1) break;
  return ret;
```

4.5 Matrix Exponentiation

```
//matmul multiplica m1 por m2
//matpow exponencia a matrix ml por p
//mul vet multiplica a matrix ml pelo vetor vet
vvi matmul(vvi &m1, vvi &m2)
  ans.resize(m1.size(), vi(m2.size(), 0));
  for (int i = 0; i < n; i++)
   for (int j = 0; j < n; j++)
     for (int k = 0; k < n; k++)
        ans[i][j] += m1[i][k] * m2[k][j];
        ans[i][j] %= MOD;
 return ans;
vvi matpow(vvi &m1, ll p)
  ans.resize(m1.size(), vi(m1.size(), 0));
  for (int i = 0; i < n; i++) ans[i][i] = 1;</pre>
 while (p) {
  if (p & 1) ans = matmul(ans, m1);
   m1 = matmul(m1, m1);
   p >>= 1;
  return ans;
```

```
}
// VETOR TEM N LINHAS E A MATRIZ E QUADRADA
vi mulvet(vvi &ml, vi &vet)
{
    vi ans;
    ans.resize(vet.size(), 0);
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++) {
        ans[i] += (ml[i][j] * vet[j]);
        ans[i] %= MOD;
    }
    return ans;</pre>
```

4.6 Pascal Triangle

```
//Fazer combinacao de N escolhe M
//por meio do triangulo de pascal
//Complexidade: O(m*n)
unsigned long long comb[61][61];
for (int i = 0; i < 61; i++) {
    comb[i][0] = 1;
    comb[i][0] = 1;
}
for (int i = 2; i < 61; i++)
    for (int j = 1; j < i; j++)
        c o mb[i][j] = comb[i - 1][j] + comb[i - 1][j - 1];</pre>
```

4.7 Euler's Totient Function

```
//retorna quantos elementos coprimos
//a N e menores que n existem
int phi (int n)
{
  int result = n;
  for (int i = 2; i * i <= n; ++i)
    if (n % i == 0) f /= i;
    result -= result / i;
  }
  if (n > 1) result -= result / n;
  return result;
}
```

4.8 Pollard Rho

```
11 u;
11 t
const int tamteste=5;
11 abss(11 v) { return v>=0 ? v : -v;}
11 randerson()
  ld pseudo=(ld)rand()/(ld)RAND_MAX;
  return (11) (round((1d) range*pseudo))+1LL;
ll mulmod(ll a, ll b, ll mod)
  11 ret=0:
  while (b>0)
    if(b%2!=0) ret=(ret+a)%mod;
    a=(a+a)%mod;
    b=b/2LL;
  return ret;
11 expmod(11 a, 11 e, 11 mod)
  11 ret=1;
  while (e>0)
    if(e%2!=0) ret=mulmod(ret,a,mod);
    a=mulmod(a,a,mod);
    e=e/2LL;
  return ret;
```

```
bool jeova(ll a, ll n)
  11 x = expmod(a, u, n);
  11 last=x;
  for(int i=0;i<t;i++)</pre>
    x=mulmod(x,x,n);
    if(x==1 and last!=1 and last!=(n-1)) return true;
    last=x;
  if(x==1) return false;
  return true;
bool isprime(ll n)
  u=n-1;
  while (u%2==0)
    t++;
    u/=2LL:
  if (n==2) return true;
  if(n==3) return true;
  if(n%2==0) return false:
  if(n<2) return false:
  for(int i=0;i<tamteste;i++)</pre>
    11 v = randerson()%(n-2)+1;
    //cout<<"jeova "<<v<<" "<<n<<endl;
    if(jeova(v,n)) return false;
  return true;
11 gcd(11 a, 11 b) { return !b ? a : gcd(b,a%b);}
11 calc(11 x, 11 n, 11 c)
  return (mulmod(x,x,n)+c)%n;
ll pollard(ll n)
  11 i=1;
  11 k=1;
  11 x=2;
  11 y=x;
  11 c;
  do
    c=randerson()%n;
  }while (c==0 or (c+2) %n==0);
  while (d!=n)
    if(<u>i</u>==k)
        k \star = 2LL;
        y=x;
        i=0;
    x=calc(x,n,c);
    d=\gcd(abss(y-x),n);
    if(d!=1) return d;
vector<ll> getdiv(ll n)
  vector<ll> ret;
  if(n==1) return ret;
  if(isprime(n))
    ret.pb(n);
    return ret;
  11 d = pollard(n);
  ret=getdiv(d);
  vector<ll> ret2=getdiv(n/d);
  for(int i=0;i<ret2.size();i++) ret.pb(ret2[i]);</pre>
 return ret:
```

4.9 Sieve of Eratosthenes

```
//esse crivo gera MAXN primos
const int MAX = le6;
int primes(MAX);
void gen_primes()
{
  int i, j;
  for (i = 2; i*i <= MAX; i++)
    if (primes[i])
    for (j = i; j * i < MAX; j++) primes[i * j] = 0;</pre>
```

4.10 Extended Euclidean Algorithm

```
//returns g = gcd(a, b);
//finds x,y such that d= ax+by;
int extended_euclid(int a, int b, int &x, int &y)
{
   int xx = y = 0;
   int yy = x = 1;
   while (b) {
      int q = a / b;
      int t = b;
      b = a & b;
      a = t;
      t = xx;
      xx = x - q * xx;
      xx = t;
      t = yy;
      yy = y - q * yy;
      y = t;
   }
   return a;
}
```

4.11 Multiplicative Inverse

```
//computes b such that ab = 1(mod n), returns - 1 on failure
int mod_inverse(int a, int n)
{
  int x, y;
  int g = extended_euclid(a, n, x, y);
  if (g > 1) return -1;
  return (x+n)%n;
}
```

4.12 Multiplicative Inverse 2

```
//inverso multiplicativo de A % MOD
//certifique de MOD estar definido antes bonito!
//complexidade: O(log(a))
ll mul_inv(ll a)
{
    ll pin0 = MOD, pin = MOD, t, q;
    ll x0 = 0, x1 = 1;
    if (pin == 1) return 1;
    while (a > 1) {
        q = a / pin;
        t = pin, pin = a % pin, a = t;
        t = x0, x0 = x1 - q * x0, x1 = t;
    }
    if (x1 < 0) x1 += pin0;
    return x1;
}
```

4.13 Linear Diophantine Equation

5 Combinatorial Optimization

5.1 Dinic

```
//grafo bipartido O(Esgrt(v))
//Para recuperar a resposta, e so colocar um bool
//de false na aresta de retorno e fazer uma bfs/dfs
//andando pelos vertices de capacidade =0 e arestas
//que nao sao de retorno
struct Edge {
 int v, rev;
  int cap:
 Edge(int v_, int cap_, int rev_) : v(v_), rev(rev_), cap(cap_) {}
};
struct MaxFlow {
 vector<vector<Edge> > g;
  vector<int> level;
  queue<int> q;
  int flow, n;
  MaxFlow(int n_) : g(n_), level(n_), n(n_) {}
  void addEdge(int u, int v, int cap)
    if (u == v) return;
    Edge e(v, cap, int(g[v].size()));
    Edge r(u, 0, int(g[u].size()));
    g[u].push_back(e);
    g[v].push_back(r);
  bool buildLevelGraph(int src, int sink)
    fill(level.begin(), level.end(), -1);
    while (not q.empty()) q.pop();
    level[src] = 0;
    q.push(src);
    while (not q.empty()) {
      int u = q.front();
      for (auto e = g[u].begin(); e != g[u].end(); ++e) {
        if (not e->cap or level[e->v] != -1) continue;
level[e->v] = level[u] + 1;
        if (e->v == sink) return true;
        q.push(e->v);
    return false;
  int blockingFlow(int u, int sink, int f)
    if (u == sink or not f) return f;
    int fu = f;
    for (auto e = g[u].begin(); e != g[u].end(); ++e) {
     if (not e->cap or level[e->v] != level[u] + 1) continue;
      int mincap = blockingFlow(e->v, sink, min(fu, e->cap));
      if (mincap) {
        g[e->v][e->rev].cap += mincap;
e->cap -= mincap;
        fu -= mincap;
    if (f == fu) level[u] = -1;
    return f - fu;
  int maxFlow(int src, int sink)
    while (buildLevelGraph(src, sink))
      flow += blockingFlow(src, sink, numeric_limits<int>::max());
    return flow:
1:
```

5.2 Hopcroft-Karp - Bipartite Matching

```
/* O(v^3) * Matching maximo de grafo bipartido de peso 1 nas arestas * supondo que o grafo bipartido seja enumerado de 0-n-1
```

```
* chamamos maxMatch(n)
class MaxMatch {
  vi graph[N];
  int match[N], us[N];
  MaxFlow(){};
  void addEdge(int u, int v) { graph[u].pb(v); }
  int dfs(int u)
    if (us[u]) return 0;
    us[u] = 1;
    for (int v : graph[u]) {
   if (match[v] == -1 or (dfs(match[v]))) {
     match[v] = u;
}
         return 1;
    return 0;
  int maxMatch(int n)
    memset(match, -1, sizeof(match));
    int ret = 0;
    for (int i = 0; i < n; i++) {</pre>
      memset(us, 0, sizeof(us));
      ret += dfs(i);
    return ret:
```

5.3 Max Bipartite Matching 2

```
// This code performs maximum bipartite matching.
 // Running time: O(|E| |V|) -- often much faster in practice
      INPUT: w[i][j] = edge between row node i and column node j
     OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
               mc[j] = assignment for column node j, -1 if unassigned
               function returns number of matches made
#include <vector>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
fypedax vectoria vi
bool FindMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
   for (int j = 0; j < w[i].size(); j++) {
      if (w[i][j] && !seen[j]) {</pre>
       seen[j] = true;
       if (mc[j] < 0 || FindMatch(mc[j], w, mr, mc, seen)) {</pre>
         mr[i] = j;
mc[j] = i;
         return true;
  return false:
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
  mr = VI(w.size(), -1);
  mc = VI(w[0].size(), -1);
  int ct = 0;
  for (int i = 0; i < w.size(); i++) {</pre>
     VI seen(w[0].size());
     if (FindMatch(i, w, mr, mc, seen)) ct++;
  return ct;
```

5.4 Min Cost Matching

```
Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
#include <algorithm>
#include <cmath>
#include <cstdio>
#include <vector>
using namespace std;
typedef vector<double> VD:
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate)
  int n = int(cost.size());
  // construct dual feasible solution
  VD u(n);
  VD v(n);
  for (int i = 0; i < n; i++) {
    u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);
  for (int j = 0; j < n; j++) {
  v[j] = cost[0][j] - u[0];</pre>
    for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
  // construct primal solution satisfying complementary slackness
  Rmate = VI(n, -1);
  int mated = 0;
for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++) {
      if (Rmate[j] != -1) continue;
}</pre>
      if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {
        Lmate[i] = j;
Rmate[j] = i;
         mated++;
         break;
  VD dist(n);
  VI dad(n);
  VI seen(n);
  // repeat until primal solution is feasible
  while (mated < n) {
    // find an unmatched left node
    int s = 0;
    while (Lmate[s] != -1) s++;
    // initialize Dijkstra
    fill(dad.begin(), dad.end(), -1);
    fill(seen.begin(), seen.end(), 0);
    for (int k = 0; k < n; k++) dist[k] = cost[s][k] - u[s] - v[k];
    while (true) {
      // find closest
       \frac{1}{1} = -1:
      for (int k = 0; k < n; k++) {
        if (seen[k]) continue;
if (j == -1 || dist[k] < dist[j]) j = k;</pre>
      seen[j] = 1;
       // termination condition
      if (Rmate[j] == -1) break;
       // relax neighbors
       const int i = Rmate[j];
       for (int k = 0; k < n; k++) {
        if (seen[k]) continue;
         const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
        if (dist[k] > new_dist) {
  dist[k] = new_dist;
           dad[k] = j;
     // update dual variables
    for (int k = 0; k < n; k++)
      if (k == j || !seen[k]) continue;
```

```
const int i = Rmate[k];
  v[k] += dist[k] - dist[j];
  u[i] -= dist[k] - dist[j];
}
u[s] += dist[j];

// augment along path
while (dad[j] >= 0) {
  const int d = dad[j];
  Rmate[j] = Rmate[d];
  Lmate[Rmate[j]] = j;
  j = d;
}
Rmate[j] = s;
  Lmate[s] = j;

mated++;
}
double value = 0;
for (int i = 0; i < n; i++) value += cost[i][Lmate[i]];
return value;</pre>
```

5.5 Min Cost Max Flow

```
int flow[N][N];
vector<pair<int, int> > g[N];
int n, m, k;
inline int ent(int a) { return a * 2; }
inline int out(int a) { return a * 2 + 1; }
inline void addEdge(int a, int b, int custo, int fluxo)
  flow[a][b] += fluxo;
  g[a].push_back(make_pair(b, custo));
  g[b].push_back(make_pair(a, -custo));
int src = N - 1, tgt = N - 2;
int dis[N], pai[N];
inline int dij()
  memset(dis, INF, sizeof dis);
  memset(pai, -1, sizeof pai);
  priority_queue<pair<int, int> > q;
  dis[src] = 0;
  q.push(make_pair(0, src));
  while (!q.empty()) {
    pair<int, int> foo = q.top();
    g.pop();
    int x = foo.second, cost = -foo.first;
    if (dis[x] != cost) continue;
for (int i = 0; i < g[x].size(); ++i) {
  int y = g[x][i].first, w = g[x][i].second;</pre>
      if (flow[x][y] <= 0) continue;</pre>
      if (dis[y] > dis[x] + w) {
        dis[y] = dis[x] + w;
        pai[y] = x;
        q.push(make_pair(-dis[y], y));
  return dis[tgt] != INF;
int minCost()
  int maxFlow = 0;
  int minC = 0;
  while (dij()) {
    int u = tgt;
    int minFlow = INF;
    while (pai[u] != -1) {
      minFlow = min(minFlow, flow[pai[u]][u]);
      u = pai[u];
    maxFlow += minFlow;
    minC += minFlow * dis[tgt];
    u = tqt;
    while (pai[u] != -1) {
      flow[pai[u]][u] -= minFlow;
      flow[u][pai[u]] += minFlow;
      u = pai[u];
```

```
}
if (maxFlow != n * k) minC = -1;
return minC;
}
inline void init()
{
  memset(flow, 0, sizeof flow);
  for (int i = 0; i < N; ++i) {
      g[i].clear();
  }
}</pre>
```

5.6 Min Cost Max Flow 2

```
// Implementation of min cost max flow algorithm using adjacency
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
      max flow:
                          O(|V|^3) augmentations
      min cost max flow: O(|V|^4 * MAX\_EDGE\_COST) augmentations
// INPUT:
      - graph, constructed using AddEdge()
      - source
      - sink
// OUTPUT:
      - (maximum flow value, minimum cost value)
       - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <iostream>
#include <vector>
using namespace std:
typedef vector<VI> VVI;
typedef long long LL;
typedef vector<LL> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const LL INF = numeric_limits<LL>::max() / 4;
struct MinCostMaxFlow {
  int N;
  VVL cap, flow, cost;
  VI found:
  VL dist, pi, width;
  VPII dad:
  MinCostMaxFlow(int N): N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
      found(N), dist(N), pi(N), width(N), dad(N){}
  void AddEdge(int from, int to, LL cap, LL cost)
    this->cap[from][to] = cap;
    this->cost[from][to] = cost;
  void Relax(int s, int k, LL cap, LL cost, int dir)
   LL val = dist[s] + pi[s] - pi[k] + cost;
if (cap && val < dist[k]) {
     dist[k] = val;
      dad[k] = make_pair(s, dir);
      width[k] = min(cap, width[s]);
  LL Dijkstra(int s, int t)
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
      int best = -1;
      found[s] = true;
      for (int k = 0; k < N; k++) {
       if (found[k]) continue;
```

```
Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
        Relax(s, k, flow[k][s], -cost[k][s], -1);
        if (best == -1 || dist[k] < dist[best]) best = k;</pre>
    for (int k = 0; k < N; k++) pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<LL, LL> GetMaxFlow(int s, int t)
    LL totflow = 0, totcost = 0;
while (LL amt = Dijkstra(s, t)) {
      totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
          totcost += amt * cost[dad[x].first][x];
          flow[x][dad[x].first] -= amt;
          totcost -= amt * cost[x][dad[x].first];
    return make_pair(totflow, totcost);
};
```

5.7 Edmonds Karp

```
struct Edge {
  int at, where;
  void init(int _at, ll _cap, int _where)
    at = _at, cap = _cap, where = _where;
}:
struct dad {
 int at, up, down;
  dad() \{ at = -1; \}
  dad(int _at, int _up, int _down) { at = _at, up = _up, down = _down; }
private:
  vector<vector<Edge> > g;
  11 mf, f;
  int s, t;
  vector<dad> p;
 public:
  void augment(int v, ll minEdge)
    if (v == s) {
      f = minEdge;
      return;
    else if (p[v].at != -1) {
      augment(p[v].at, min(minEdge, g[p[v].at][p[v].up].cap));
      g[p[v].at][p[v].up].cap -= f;
      g[v][p[v].down].cap += f;
  void init(int N)
   for (int i = 0; i < g.size(); i++) g[i].clear();</pre>
    mf = 0, f = 0;
    g.resize(N);
  void addEdge(int u, int v, 11 cap)
    A.init(v, cap, g[v].size());
    B.init(u, 0, g[u].size());
    g[u].pb(A);
    g[v].pb(B);
  int maxFlow(int source, int sink)
    s = source;
    t = sink;
    mf = 0;
    while (true) {
```

```
f = 0;
      vector<int> dist(g.size(), INF);
      dist[s] = 0;
      queue<int> q;
      q.push(s);
      p.clear();
      p.resize(g.size());
      while (!q.empty())
        int u = q.front();
        q.pop();
        if (u == t) break;
        for (int i = 0; i < g[u].size(); i++) {
          Edge prox = g[u][i];
          if (dist[prox.at] == INF and prox.cap > 0) {
   dist[prox.at] = dist[u] + 1;
             q.push(prox.at);
             dad paizao(u, i, prox.where);
            p[prox.at] = paizao;
      augment(t, INF);
      if (f == 0) break;
     mf += f:
    return mf;
};
```

6 Dynamic Programming

6.1 Convex Hull Trick

```
/\star \ \textit{Esse convex hull trick e para achar a reta minima!}
* Para maximizar a reta dada , basta trocar o '>' para
 * para '<' na funcao query;
 * Nao chamar query com B ou A vazios! Atualizar dp para
* depois fazer a query =)

* ATENCAO COM O DOUBLE!! ESTA EM LONG LONG :)
vi A[N], B[N];
int pont[N];
bool odomeioehlixo(int r1, int r2, int r3, int j)
  return (B[j][r1] - B[j][r3]) * (A[j][r2] - A[j][r1]) <
         (B[j][r1] - B[j][r2]) * (A[j][r3] - A[j][r1]);
void add(ll a, ll b, int j)
 B[j].pb(b);
  A[i].pb(a);
  while (B[j].size() >= 3 and
         odomeioehlixo(B[j].size() - 3, B[j].size() - 2, B[j].size() - 1, j)) {
    B[j].erase(B[j].end() - 2);
    A[j].erase(A[j].end() - 2);
11 query(ll x, int j)
  if (pont[j] >= B[j].size()) pont[j] = B[j].size() - 1;
  while (pont[j] < B[j].size() - 1 and</pre>
         (A[j][pont[j] + 1] * x + B[j][pont[j] + 1] >
          A[j][pont[j]] * x + B[j][pont[j]]))
    pont[i]++;
  return A[j][pont[j]] * x + B[j][pont[j]];
* http://www.spoj.com/problems/APIO10A/
 * http://www.spoj.com/problems/ACQUIRE/
```

6.2 Convex Hull Trick 2

```
/*
 * Given a set of pairs (m, b) specifying lines of the form y = m*x + b, process
 * a
 * set of x-coordinate queries each asking to find the minimum y-value when any
 * of
```

```
\star the given lines are evaluated at the specified x. To instead have the queries
 * optimize for maximum y-value, set the QUERY_MAX flag to true.
 * The following implementation is a fully dynamic variant of the convex hull
 * optimization technique, using a self-balancing binary search tree (std::set)
 * support the ability to call add_line() and get_best() in any desired order.
 * Explanation: http://wcipeg.com/wiki/Convex_hull_trick#Fully_dynamic_variant
 * Time Complexity: O(n log n) on the total number of calls made to add_line(),
 * any length n sequence of arbitrarily interlaced add_line() and get_min()
 \star Each individual call to add_line() is O(log n) amortized and each individual
 * call to get_best() is O(log n), where n is the number of lines added so far.
 * Space Complexity: O(n) auxiliary on the number of calls made to add_line().
#include <limits> // std::numeric_limits
#include <set>
class hull_optimizer {
  struct line {
   long long m, b,
   double xlo;
   bool is query:
   bool query_max;
   line(long long m, long long b, long long val, bool is_query, bool query_max)
      this->m = m:
     this->b = b:
      this->val = val;
     this->xlo = -std::numeric_limits<double>::max();
      this->is_query = is_query;
     this->query_max = query_max;
   bool parallel(const line &l) const { return m == 1.m; }
   double intersect(const line &1) const
      if (parallel(1)) return std::numeric_limits<double>::max();
     return (double) (1.b - b) / (m - 1.m);
   bool operator<(const line &1) const
      if (1.is_query) return query_max ? (xlo < 1.val) : (1.val < xlo);</pre>
     return m < 1.m;
  std::set<line> hull:
  bool _query_max;
  typedef std::set<line>::iterator hulliter;
  bool has prev(hulliter it) const { return it != hull.begin(); }
  bool has next(hulliter it) const
   return (it != hull.end()) && (++it != hull.end());
  bool irrelevant (hulliter it) const
   if (!has_prev(it) || !has_next(it)) return false;
   hulliter prev = it, next = it;
   --prev;
   ++next:
   return _query_max ? prev->intersect(*next) <= prev->intersect(*it)
                     : next->intersect(*prev) <= next->intersect(*it);
  hulliter update left border(hulliter it)
   if ((_query_max && !has_prev(it)) || (!_query_max && !has_next(it)))
     return it;
   hulliter it2 = it;
   double val = it->intersect(_query_max ? *--it2 : *++it2);
   line 1(*it);
    1.xlo = val;
   hull.erase(it++);
   return hull.insert(it, 1);
  hull_optimizer(bool query_max = false) { this->_query_max = query_max; }
  void add_line(long long m, long long b)
   line 1(m, b, 0, false, _query_max);
   hulliter it = hull.lower_bound(1);
   if (it != hull.end() && it->parallel(1)) {
     if ((_query_max && it->b < b) || (!_query_max && b < it->b))
```

hull.erase(it++);

```
else
        return;
    it = hull.insert(it, 1);
    if (irrelevant(it)) {
      hull.erase(it);
      return;
    while (has_prev(it) && irrelevant(--it)) hull.erase(it++);
    while (has_next(it) && irrelevant(++it)) hull.erase(it--);
    it = update_left_border(it);
    if (has_prev(it)) update_left_border(--it);
   if (has_next(++it)) update_left_border(++it);
  long long get_best(long long x) const
    line q(0, 0, x, true, _query_max);
    hulliter it = hull.lower_bound(q);
    if (_query_max) --it;
    return it->m * x + it->b;
};
/*** Example Usage ***/
#include <cassert>
int main()
  hull_optimizer h;
 h.add_line(3, 0);
  h.add_line(0, 6);
  h.add_line(1, 2);
  h.add_line(2, 1);
  assert(h.get_best(0) == 0);
  assert(h.get_best(2) == 4);
  assert(h.get_best(1) == 3);
  assert(h.get_best(3) == 5);
  return 0;
```

6.3 Divide-and-Conquer

```
//Um exemplo de Divide and conquer:
int. MOD = 1e9 + 7:
const int N = 1010;
\quad \textbf{int} \ dp[N][N], \ cost[N][N], \ v[N], \ pref[N], \ n, \ m; \\
void compDP(int j, int L, int R, int b, int e)
  if (L > R) return;
  int mid = (L + R) / 2;
  int idx = -1;
  for (int i = b; i <= min(mid, e); i++)</pre>
    if (dp[mid][j] > dp[i][j - 1] + cost[i + 1][mid]) {
      dp[mid][j] = dp[i][j-1] + cost[i+1][mid];
  compDP(j, L, mid - 1, b, idx);
  compDP(j, mid + 1, R, idx, e);
//chamada!
for(int i=1;i<=n;i++) dp[i][0]=cost[1][i];</pre>
for (int i=1; i <= m; i++) compDP (i, 1, n, 1, n);</pre>
```

6.4 LIS - Longest Increasing Subsequence

```
//asw -> vetor com resposta!!
//asw.size() o tamanho da maior lis
void lis( const vector< int > & v, vector< int > & asw )
{
  vector<int> pd(v.size(),0), pd_index(v.size()), pred(v.size());
  int maxi = 0, x=0, j=0, ind=0;
  for(int i=0;i<v.size();i++)
  {
      x = v[i];
      j=lower_bound(pd.begin(),pd.begin()+maxi,x) -pd.begin();
      pd[j] = x;
      pd_index[j] = i;
    if(j==maxi)
  {
      maxi++;
      ind = i;
    }
}</pre>
```

```
}
if(pred[i] == j) pd_index[j-1] = -1;
}
int pos=maxi-1, k=v[ind];
asw.resize( maxi );
while ( pos >= 0 )
{
    asw[pos--] = k;
    ind = pred[ind];
    k = v[ind];
}
```

7 Geometry

7.1 Convex Hull - Monotone Chain

```
typedef struct sPoint {
        int x, y;
        sPoint (int _x, int _y)
                x = _x;
               y = y;
} point;
bool comp(point a, point b)
        if (a.x == b.x) return a.y < b.y;</pre>
        return a.x < b.x;
int cross(point a, point b, point c) // AB x BC
        a.x = b.x;
        a.v -= b.v;
        b.x -= c.x;
        b.y -= c.y;
        return a.x * b.y - a.y * b.x;
bool isCw(point a, point b, point c) // Clockwise
        return cross(a, b, c) < 0;
// >= if you want to put collinear points on the convex hull
bool isCcw(point a, point b, point c) // Counter Clockwise
        return cross(a, b, c) > 0;
vector<point> convexHull(vector<point> p)
        vector<point> u, 1; // Upper and Lower hulls
        sort(p.begin(), p.end(), comp);
        for (unsigned int i = 0; i < p.size(); i++) {</pre>
                while (1.size() > 1 && !isCow(l[1.size() - 1], l[1.size() - 2], p[i]))
                        1.erase(1.begin() + (1.size() - 1));
                1.push_back(p[i]);
        for (int i = p.size() - 1; i >= 0; i--) {
                while (u.size() > 1 && !isCow(u[u.size() - 1], u[u.size() - 2], p[i]))
                        u.erase(u.begin() + (u.size() - 1));
                u.push_back(p[i]);
        u.erase(u.begin() + (u.size() - 1));
        1.erase(1.begin() + (1.size() - 1));
        1.insert(l.end(), u.begin(), u.end());
        return 1;
```

7.2 Minimum Enclosing Circle

```
//6.5- Minimum Enclosing Circle const double eps = 1e-6; #define CIRCLE circ #define PT Ponto #define MP 101 #define eps 1e-9
```

```
#define x first
#define y second
typedef double cood;
typedef int num;
typedef int point;
double resp;
cood x[MP], y[MP], ar, ax, ay;
int p[MP];
typedef pair<double, double> ponto;
typedef pair<double, double> Ponto;
double dista(ponto a, ponto b)
  return sqrt((a.first - b.first) * (a.first - b.first) +
               (a.second - b.second) * (a.second - b.second));
bool in (ponto a, pair <double, ponto > c)
  if (dista(a, c.second) - eps < c.first) return true;</pre>
  return false;
bool same (point a, point b)
  return (fabs(x[a] - x[b]) < eps && fabs(y[a] - y[b]) < eps);
bool lexLess(point a, point b)
  if (fabs(x[a] - x[b]) < eps) return y[a] < y[b];
 return x[a] < x[b];
inline cood dist(cood xx, cood yy, point a)
  return sqrt((xx - x[a]) * (xx - x[a]) + (yy - y[a]) * (yy - y[a]));
inline cood cP (point a, point b, point c)
  return (x[a] - x[b]) * (y[c] - y[b]) - (x[c] - x[b]) * (y[a] - y[b]);
void findCircle(point a, point b, point c, cood& cx, cood& cy)
  cx = 0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (v[b] - v[c])
       0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
           (y[a] - y[b]),
  cy = 0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
       0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (x[b] - x[c]);
  cx /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
 cy /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
void spanCircle2(int k, point p0, point p1, cood& cx, cood& cy, cood& r)
 cx = 0.5 * (x[p0] + x[p1]);
  cy = 0.5 * (y[p0] + y[p1]);
  r = dist(cx, cy, p0);
for (int i = 0; i < k; i++)
    if (dist(cx, cy, p[i]) > r) {
      findCircle(p0, p1, p[i], cx, cy);
      r = dist(cx, cy, p[i]);
void spanCircle1(int k, point p0, cood& cx, cood& cy, cood& r)
  cx = 0.5 * (x[p0] + x[p[0]]);
  cy = 0.5 * (y[p0] + y[p[0]]);
  r = dist(cx, cy, p0);
for (int i = 0; i < k; i++)
    if (dist(cx, cy, p[i]) > r) spanCircle2(i, p0, p[i], cx, cy, r);
void spanCircle(int n, cood& cx, cood& cy, cood& r)
  // Bem importante, retirar repetidos
  sort(p, p + 1, lexLess);
  n = unique(p, p + n) - p;
  random_shuffle(p, p + n);
  if (n > 1) {
   cx = 0.5 * (x[p[0]] + x[p[1]]);
   cy = 0.5 * (y[p[0]] + y[p[1]]);

r = dist(cx, cy, p[1]);

for (int i = 2; i < n; i++)
      if (dist(cx, cy, p[i]) > r) spanCircle1(i, p[i], cx, cy, r);
    \mathbf{cx} = \mathbf{x}[0];
    cy = y[0];
```

```
r = 0.0;
}

void solve(vector<pair<double, double> >& v)

int N = v.size();
for (int i = 0; i < N; i++) {
    x[i] = v[i].first;
    y[i] = v[i].second;
    p[i] = i;
}
spanCircle(N, ax, ay, ar);
}</pre>
```

7.3 Minimum Enclosing Circle 2

```
const double eps = 1e-6;
#define CIRCLE circ
#define PT Ponto
#define MP 101
#define eps 1e-9
#define x first
#define v second
typedef double cood;
typedef int num;
typedef int point;
double resp;
cood x[MP], y[MP], ar, ax, ay;
int p[MP];
typedef pair<double, double> ponto;
typedef pair<double, double> Ponto;
double dista (ponto a, ponto b)
  return sqrt((a.first - b.first) * (a.first - b.first) +
               (a.second - b.second) * (a.second - b.second));
bool in(ponto a, pair<double, ponto> c)
  if (dista(a, c.second) - eps < c.first) return true;</pre>
  return false:
bool same (point a, point b)
  return (fabs(x[a] - x[b]) < eps && fabs(y[a] - y[b]) < eps);
bool lexLess(point a, point b)
  if (fabs(x[a] - x[b]) < eps) return y[a] < y[b];
  return x[a] < x[b];</pre>
inline cood dist(cood xx, cood yy, point a)
  return sqrt((xx - x[a]) * (xx - x[a]) + (yy - y[a]) * (yy - y[a]));
inline cood cP(point a, point b, point c)
  return (x[a] - x[b]) * (y[c] - y[b]) - (x[c] - x[b]) * (y[a] - y[b]);
void findCircle(point a, point b, point c, cood& cx, cood& cy)
  cx = 0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (y[b] - y[c])
       0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
           (y[a] - y[b]),
  cy = 0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
           (x[a] - x[b])
       0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (x[b] - x[c]);
  cx /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
 cy /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
void spanCircle2(int k, point p0, point p1, cood& cx, cood& cy, cood& r)
  cx = 0.5 * (x[p0] + x[p1]);
  cy = 0.5 * (y[p0] + y[p1]);
  r = dist(cx, cy, p0);
for (int i = 0; i < k; i++)
   if (dist(cx, cy, p[i]) > r) {
  findCircle(p0, p1, p[i], cx, cy);
      r = dist(cx, cy, p[i]);
```

```
void spanCircle1(int k, point p0, cood& cx, cood& cy, cood& r)
 cx = 0.5 * (x[p0] + x[p[0]]);
 cy = 0.5 * (y[p0] + y[p[0]]);
  r = dist(cx, cy, p0);
  for (int i = 0; i < k; i++)
    if (dist(cx, cy, p[i]) > r) spanCircle2(i, p0, p[i], cx, cy, r);
void spanCircle(int n, cood& cx, cood& cy, cood& r)
  // Bem importante, retirar repetidos
  sort(p, p + 1, lexLess);
 n = unique(p, p + n) - p;
random_shuffle(p, p + n);
  if (n > 1) {
    cx = 0.5 * (x[p[0]] + x[p[1]]);
    cy = 0.5 * (y[p[0]] + y[p[1]]);
    r = dist(cx, cy, p[1]);
    for (int i = 2; i < n; i++)
      if (dist(cx, cy, p[i]) > r) spanCircle1(i, p[i], cx, cy, r);
  else {
    cx = x[0];
    cy = y[0];
    r = 0.0;
void solve(vector<pair<double, double> >& v)
 int N = v.size();
for (int i = 0; i < N; i++) {</pre>
   x[i] = v[i].first;
    y[i] = v[i].second;
    p[i] = i;
  spanCircle(N, ax, ay, ar);
```

7.4 Fast Geometry in Cpp

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100:
double EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT (double x, double y) : x(x), y(y) {}
  PT(const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y);
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y);
                                const { return PT(x*c, y*c );
    operator * (double c)
  PT operator / (double c)
                                const { return PT(x/c, y/c ); ]
double dot(PT p, PT q)
                            { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream &os, const PT &p) {
 os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p) { return PT (-p.y,p.x); }
PT RotateCW90(PT p)
                        { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a. PT b. PT c) {
 return a + (b-a) *dot(c-a, b-a) /dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
```

```
double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
     = dot(c-a, b-a)/r;
  if (r < 0) return a;
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z, double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;</pre>
bool LinesCollinear (PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
 // line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    \textbf{if} \ (\texttt{dist2}(\texttt{a, c}) \ \leq \ \texttt{EPS} \ | \ | \ \texttt{dist2}(\texttt{a, d}) \ \leq \ \texttt{EPS} \ | \ |
      dist2(b, c) < EPS \mid \mid dist2(b, d) < EPS) return true;
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
      return false:
    return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
\ensuremath{//} strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
//\ \mbox{(making sure to deal with signs properly)} and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||
   p[j].y <= q.y && q.y < p[i].y) &&</pre>
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
      return true;
    return false:
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
```

```
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
  b = b-a;
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R \mid | d+min(r, R) < max(r, R)) return ret;
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret:
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as // the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
 return fabs (ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale:
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++)
    for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
int l = (k+1) % p.size();
      if (i == 1 \mid \mid j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;
  // expected: (5.2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "</pre>
       << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
       << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
```

```
// expected: 1 0 1
cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
      << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
      << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
      << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
      << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 1 1 1 0
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
// expected: (1,1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push back(PT(5.5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "
      << PointInPolygon(v, PT(2,0)) << " "
      << PointInPolygon(v, PT(0,2)) << " "
      << PointInPolygon(v, PT(5,2)) << " "
      << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
      << PointOnPolygon(v, PT(2,0)) << " "
      << PointOnPolygon(v, PT(0,2)) << " "
      << PointOnPolygon(v, PT(5,2)) << " "
      << PointOnPolygon(v, PT(2,5)) << endl;
// expected: (1,6)
                (5,4) (4,5)
                blank line
                (4,5) (5,4)
                (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5, 4.5), 10, sqrt(2.0)/2.0);

for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;
return 0:
```

8 Data Structures

8.1 Disjoint Set Union

```
const int N=500010;
int p[N],rank[N];
void init()
{
   memset(rank,0,sizeof(rank));
   for(int i=0;i<N;i++) p[i]=i;
}
int findset(int i)</pre>
```

```
{
    if(p[i]==i) return i;
    return p[i]=findset(p[i]);
}
bool same(int i, int j)
{
    return (findset(i) == findset(j));
}
void unionSet(int i, int j)
{
    if (!same(i, j)) {
        int x = findset(i), y=findset(j);
    if (rank[x] > rank[y])
        p[y] = x;
    else {
        p[x] = y;
        if (rank[x] == rank[y]) rank[y]++;
    }
}
```

8.2 Persistent Segment Tree

```
//PRINTAR O NUMERO DE ELEMENTOS DISTINTOS
//EM UM INTERVALO DO ARRAY
const int N = 30010:
int tr[100 * N], L[100 * N], R[100 * N], root[100 * N];
int v[N], mapa[100 * N];
int cont = 0;
void build(int node, int b, int e)
  if (b == e) {
    tr[node] = 0;
  else {
    L[node] = cont++;
    R[node] = cont++;
    build(L[node], b, (b + e) / 2);
build(R[node], (b + e) / 2 + 1, e);
tr[node] = tr[L[node]] + tr[R[node]];
int update(int node, int b, int e, int i, int val)
  int idx = cont++;
  tr[idx] = tr[node] + val;
  L[idx] = L[node];
R[idx] = R[node];
  if (b == e) return idx;
  int mid = (b + e) / 2;
  if (i <= mid)</pre>
    L[idx] = update(L[node], b, mid, i, val);
    R[idx] = update(R[node], mid + 1, e, i, val);
  return idx:
int query (int nodeL, int nodeR, int b, int e, int i, int j)
  if (b > j \text{ or } i > e) \text{ return } 0;
  if (i \le b \text{ and } j \ge e) {
    int p1 = tr[nodeR];
    int p2 = tr[nodeL];
    return p1 - p2;
  int mid = (b + e) / 2;
  return query(L[nodeL], L[nodeR], b, mid, i, j) +
          query(R[nodeL], R[nodeR], mid + 1, e, i, j);
int main()
  int n:
  sc(n):
   memset (mapa, -1, sizeof (mapa));
   for (int i = 0; i < n; i++) sc(v[i]);
  build(1, 0, n - 1);
   for (int i = 0; i < n; i++) {
    if (mapa[v[i]] == -1) {
       root[i + 1] = update(root[i], 0, n - 1, i, 1);
       mapa[v[i]] = i;
      root[i + 1] = update(root[i], 0, n - 1, mapa[v[i]], -1);
mapa[v[i]] = i;
       root[i + 1] = update(root[i + 1], 0, n - 1, i, 1);
  int q;
```

```
sc(q);
for (int i = 0; i < q; i++) {
   int l, r;
   sc2(l, r);
   int resp = query(root[l - 1], root[r], 0, n - 1, l - 1, r - 1);
   pri(resp);
} return 0;</pre>
```

8.3 RMQ of Indices

```
//RMQ DE INDICE
class RMQ {
private:
  vi A:
  vi M:
 public:
  RMQ (vi &v)
     M.resize(4 * v.size());
     build(1, 0, v.size() - 1);
  void build(int node, int b, int e)
     if (b == e)
       M[node] = b;
    else {
      build(2 * node, b, (b + e) / 2);
build(2 * node + 1, (b + e) / 2 + 1, e);
if (A[M[2 * node]] <= A[M[2 * node + 1]])</pre>
         M[node] = M[2 * node];
       else
         M[node] = M[2 * node + 1];
   int query (int node, int b, int e, int i, int j)
    int p1, p2;
    if (i > e \mid \mid j < b) return -1;
     if (b >= i and e <= j) return M[node];</pre>
     p1 = query(2 * node, b, (b + e) / 2, i, j);
     p2 = query(2 * node + 1, (b + e) / 2 + 1, e, i, j);
     if (p1 == -1) return p2;
    if (p2 == -1) return p1;
    if (A[p1] <= A[p2]) return p1;</pre>
     return p2;
   void atualiza(int node, int b, int e, int i, int val)
     if (i > e || i < b) return;</pre>
     if (e == b) {
       A[i] = val;
     else (
       atualiza(2 * node, b, (b + e) / 2, i, val);
atualiza(2 * node + 1, (b + e) / 2 + 1, e, i, val);
if (A[M[2 * node]] <= A[M[2 * node + 1]])
         M[node] = M[2 * node];
       else
         M[node] = M[2 * node + 1];
};
```

8.4 RSQ with Lazy-Propagation

```
//RSQ COM LAZY PROPAGATION!
class RSQ {
private:
    vil A;
    vil M;
    vil lazy;

public:
    RSQ(vil &v) {
        A = v;
        M.resize(v.size() * 4);
        lazy, assign(v.size() * 4, 0);
        build(1, 0, v.size() - 1);
```

```
void build(int node, int b, int e)
    if (b == e) {
      M[node] = A[b];
       return;
    build(2 * node, b, (b + e) / 2);
    build(2 * node + 1, (b + e) / 2 + 1, e);
    M[node] = M[2 * node] + M[2 * node + 1];
  void atualiza(int node, int b, int e, int i, int j, ll val)
    if (lazy[node] != 0) {
      M[node] += lazy[node];
if (b != e) {
        11 inter = (e - b + 1);
         11 \ i1 = (b + e) / 2 - b + 1;
         11 i2 = e - (b + e) / 2;
         11 un = lazy[node] / inter;
         lazy[2 * node] += un * i1;
         lazy[2 * node + 1] += un * i2;
       lazy[node] = 0;
    if (i > e or j < b) return;</pre>
    if (i <= b and j >= e) {
    l1 inter = (e - b + 1);
       M[node] += val * inter;
       if (b != e) {
        11 \ i1 = (b + e) / 2 - b + 1;
         11 i2 = e - (b + e) / 2;
lazy[2 * node] += i1 * (l1) val;
         lazy[2 * node + 1] += i2 * (11)val;
       return;
    atualiza(2 * node, b, (b + e) / 2, i, j, val);
atualiza(2 * node + 1, (b + e) / 2 + 1, e, i, j, val);
    M[node] = M[2 * node] + M[2 * node + 1];
  11 query(int node, int b, int e, int i, int j)
    if (i > e \text{ or } j < b) return 0;
    11 p1, p2;
    if (lazy[node] != 0) {
      M[node] += lazy[node];
       if (b != e) {
        11 inter = (e - b + 1);
         11 \ i1 = (b + e) / 2 - b + 1;
         11 i2 = e - (b + e) / 2;
        11 un = lazy[node] / inter;
         lazy[2 * node] += un * i1;
         lazy[2 * node + 1] += un * i2;
       lazy[node] = 0;
    if (i <= b and j >= e) return M[node];
    p1 = query(2 * node, b, (b + e) / 2, i, j);
p2 = query(2 * node + 1, (b + e) / 2 + 1, e, i, j);
    return p1 + p2;
};
```

8.5 Segment Tree

```
int resp =
    upper_bound(tr[node].begin(), tr[node].end(), k) - tr[node].begin();
    return tr[node].size() - resp;
}
return query(2 * node, b, (b * e) / 2, i, j, k) +
    query(2 * node * 1, (b * e) / 2 * 1, e, i, j, k);
```

8.6 Sparse Table

```
//comutar RMQ , favor inicializar: dp[i][0]=v[0]
//sendo v[0] o vetor do rmq
//chamar o build!
int dp[200100][22];
int n;
```

```
int d[200100];
void build()
{
    d[0] = d[1] = 0;
    for (int i = 2; i < n; i++) d[i] = d[i >> 1] + 1;
    for (int j = 1; j < 22; j++) {
        for (int i = 0; i + (1 << (j - 1)) < n; i++) {
            dp[i][j] = min(dp[i][j - 1], dp[i + (1 << (j - 1))][j - 1]);
        }
    }
}
int query(int i, int j)
{
    int k = d[j - i];
    int x = min(dp[i][k], dp[j - (1 << k) + 1][k]);
    return x;
}</pre>
```