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# 1 Graph Algorithms

### 1.1 2 SAT

8.2

8.3

```
//RODAR O COMPONENTE FORTEMENTE CONECTADO
//RECUPERAR NA ORDEM DE descarga DA PILHA
//CONFORME O EXEMPLO A SEGUIR
const int N = 510;
vi graph[N], rev[N];
int us[N];
stack<int> pilha;
int resposta[N];
void dfs1(int u)
  us[u] = 1;
for (int v : graph[u])
  if (!us[v]) dfs1(v);
  pilha.push(u);
void dfs2(int u, int color)
   us[u] = color;
  for (int v : rev[u])
  if (!us[v]) dfs2(v, color);
int Sat(int n)
  for (int i = 0; i < n; i++)
    if (!us[i]) dfs1(i);
   int color = 1;
  vi r;
   while (!pilha.empty()) {
    int topo = pilha.top();
    r.pb(topo);
     pilha.pop();
     if (!us[topo]) dfs2(topo, color++);
  for (int i = 0; i < n; i += 2) {
  if (us[i] == us[i + 1]) return 0;</pre>
   memset(resposta, -1, sizeof(resposta));
  for (int i = r.size() - 1; i >= 0; i--) {
  int vert = r[i] / 2;
     int ok = r[i] % 2;
    if (resposta[vert] == -1) resposta[vert] = !ok;
inline void add(int u, int v)
   graph[u].pb(v);
  rev[v].pb(u);
inline int pos(int u) { return 2 * u; }
inline int neg(int u) { return 2 * u + 1; }
```

### 1.2 Kosaraju

//Retorna os componentes fortemente conectados
//Se o usados[i]=usados[j], temos que i e j
//pertencem ao mesmo componente, col-1= numero

```
//de componentes fortemente conectados do grafo
class kosaraju {
private:
  vi usados;
  vvi graph;
  vvi trans;
  vi pilha;
 public:
  kosaraju(int N)
    graph.resize(N):
    trans.resize(N);
  void AddEdge (int u. int v)
    graph[u].pb(v);
    trans[v].pb(u);
  void dfs(int u, int pass, int color)
    usados[u] = color;
    vi vizinhos;
    if (pass == 1)
      vizinhos = graph[u];
    else
      vizinhos = trans[u];
    for (int j = 0; j < vizinhos.size(); j++) {
  int v = vizinhos[j];</pre>
      if (usados[v] == 0) {
        dfs(v, pass, color);
    pilha.pb(u);
  int SSC(int n)
    pilha.clear();
    usados.assign(n, 0);
    for (int i = 0; i < n; i++) {
  if (!usados[i]) dfs(i, 1, 1);</pre>
    usados.assign(n, 0);
    int color = 1;
for (int i = n - 1; i >= 0; i--) {
      if (usados[pilha[i]] == 0) {
        dfs(pilha[i], 2, color);
        color++;
    return color - 1:
};
```

# 1.3 Tree Isomorphism

```
//Seque no main um exemplo de utilização do isomorfismo de arvore!
vvichildren, subtreeLabels, tree, L;
vipred, map;
int n:
boolcompare(int a, int b) {
 returnsubtreeLabels[a] <subtreeLabels[b];
bool equals (int a, int b) {
  return subtreeLabels [a] == subtreeLabels [b];
voidgenerateMapping(int r1, int r2) {
 map.resize(n);
 map[r1] = r2 - n;
  sort(children[r1].begin(), children[r1].end(), compare);
  sort(children[r2].begin(), children[r2].end(), compare);
  for (int i = 0; i < (int) children[r1].size(); i++) {
   int u = children[r1][i];
    int v = children[r2][i];
   generateMapping(u, v);
vifindCenter(int offset = 0) {
  int cnt = n;
  vi deg(n);
  for (int i = 0; i < n; i++) {
  deg[i] = tree[i + offset].size();
  if (deg[i] <= 1) {</pre>
     a.push_back(i + offset);
```

--cnt; while (cnt > 0) { for (int i = 0; i < (int) a.size(); i++) { int u = a[i]; for (int j = 0; j < (int) tree[u].size(); j++) { int v = tree[u][j];  $if (--deg[v - offset] == 1) {$ na.push\_back(v); --cnt; a = na;returna; int dfs (int u, int p = -1, int depth = 0) { L[depth].push\_back(u); for (int i = 0; i < (int) tree[u].size(); i++) { int v = tree[u][i]; if (v == p) continue; pred[v] = u: children[u].push\_back(v); h = max(h, dfs(v, u, depth + 1)); return h + 1; boolrootedTreeIsomorphism(int r1, int r2) { L.assign(n, vi()); pred.assign(2  $\star$  n , -1); children.assign(2  $\star$  n, vi()); int h1 = dfs(r1); int h2 = dfs(r2); if (h1 != h2) returnfalse; int h = h1 - 1;vilabel(2 \* n); viliabel(z \* ii);
subtreeLabels.assign(2 \* n, vi());
for (int i = h - 1; i >= 0; i --) {
 for (int j = 0; j < (int) L[i + 1].size(); j++) {
 int v = L[i + 1][j];
}</pre> subtreeLabels[pred[v]].push\_back(label[v]); for (int j = 0; j < (int) L[i].size(); j++) { int v = L[i][j]; sort(subtreeLabels[v].begin(), subtreeLabels[v].end()); sort(L[i].begin(), L[i].end(), compare); for (int j = 0, cnt = 0; j < (int)  $L[i].size(); j++) = \{if (j && !equals(L[i][j], L[i][j-1])\}$ ++cnt; label[L[i][j]] = cnt; if (!equals(r1, r2)) returnfalse; generateMapping(r1, r2); returntrue; booltreeIsomorphism() vi c1 = findCenter(); vi c2 = findCenter(n); if (c1.size() == c2.size()) {
 if (rootedTreeIsomorphism(c1[0], c2[0])) return true; else if (c1.size() > 1)returnrootedTreeIsomorphism(c1[1], c2[0]); return false; int main() { n = 5; vvi t1(n); t1[0].push\_back(1); t1[1].push\_back(0); t1[1].push\_back(2); t1[2].push\_back(1); t1[1].push\_back(3); t1[3].push\_back(1); t1[0].push\_back(4); t1[4].push\_back(0); vvi t2(n); t2[0].push\_back(1); t2[1].push\_back(0);

t2[0].push\_back(4);

```
t2[4].push_back(0);
t2[4].push_back(3);
t2[3].push_back(4);
t2[4].push_back(2);
t2[2].push_back(4);
tree.assign(2 * n , vi());
for (int u = 0; u < n; u++) {
    for (int i = 0; i < tl[u].size(); i++) {
        int v = tl[u][i];
        tree[u].push_back(v);
    }
    for (int i = 0; i < t2[u].size(); i++) {
        int v = t2[u][i];
        tree[u + n].push_back(v + n);
    }
}
bool res = treeIsomorphism();
cout << res << endl;
if (res)
    for (int i = 0; i < n; i++)
    cout << map[i] << endl;</pre>
```

#### 1.4 LCA

```
//antes de usar as queries de lca, e etc..
//certifique-se de chamar a dfs da arvore e
//process()
const int N = 100000;
const int M = 22;
int P[N][M];
int big[N][M], low[N][M], level[N];
vii graph[N];
void dfs(int u, int last, int 1)
  level[u] = 1;
  P[u][0] = last;
  for (ii v : graph[u])
  if (v.first != last) {
    big[v.first][0] = low[v.first][0] = v.second;
      dfs(v.first, u, 1 + 1);
void process()
 int lca(int u, int v)
  if (level[u] < level[v]) swap(u, v);</pre>
  for (int i = M - 1; i >= 0; i--)
   if (level[u] - (1 << i) >= level[v]) u = P[u][i];
  if (u == v) return u;
  for (int i = M - 1; i >= 0; i--) {
     \mbox{if } (P[u][i] \ != P[v][i]) \ u = P[u][i], \ v = P[v][i]; \\
  return P[u][0];
int maximum(int u, int v, int x)
  int resp = 0;
  for (int i = M - 1; i >= 0; i--)
if (level[u] - (1 << i) >= level[x]) {
      resp = max(resp, big[u][i]);
      u = P[u][i];
  for (int i = M - 1; i >= 0; i--)
    if (level[v] - (1 << i) >= level[x]) {
      resp = max(resp, big[v][i]);
      v = P[v][i];
  return resp;
int minimum(int u, int v, int x)
  int i = M - 1; i >= 0; i--)
if (level[u] - (1 << i) >= level[x]) {
      resp = min(resp, low[u][i]);
      u = P[u][i];
  for (int i = M - 1; i >= 0; i--)
```

```
if (level[v] - (1 << i) >= level[x]) {
    resp = min(resp, low[v][i]);
    v = P[v][i];
}
return resp;
```

### 1.5 Bridges and Articulation Points

```
class ponte {
private:
  vvi graph;
  vi usados;
  vi e_articulacao;
  vi dfs_low;
  vi dfs_prof;
  vector<ii> pontes;
  int tempo;
 public:
  ponte(int N)
    graph.clear():
    graph.resize(N);
    usados.assign(N, 0);
    dfs_low.assign(N, 0);
    dfs_prof.assign(N, 0);
    e_articulacao.assign(N, 0);
    tempo = 0;
  void AddEdge(int u, int v)
    graph[u].pb(v);
    graph[v].pb(u);
  void dfs(int u, int pai)
    usados[u] = 1:
    int nf = 0;
    dfs_low[u] = dfs_prof[u] = tempo++;
    for (int v : graph[u]) {
      if (!usados[v]) {
        dfs(v, u);
        if (dfs_low[v] >= dfs_prof[u] and pai != -1) e_articulacao[u] = true;
        if (pai == -1 and nf > 1) e_articulacao[u] = true;
        if (dfs_low[v] > dfs_prof[u]) pontes.pb(mp(u, v));
        dfs_low[u] = min(dfs_low[u], dfs_low[v]);
      else if (v != pai)
        dfs_low[u] = min(dfs_low[u], dfs_prof[v]);
  void olha_as_pontes()
    for (int i = 0; i < graph.size(); i++)</pre>
      if (!usados[i]) dfs(i, -1);
    if (pontes.size() == 0)
      cout << " Que merda! nao tem ponte!" << endl;
      for (ii i : pontes) cout << i.first << " " << i.second << endl;</pre>
  void olha_as_art()
    for (int i = 0; i < graph.size(); i++)</pre>
    if (!usados[i]) dfs(i, -1);
for (int i = 0; i < e_articulacao.size(); i++)</pre>
      if (e_articulacao[i]) cout << " OIAAA A PONTE " << i << endl;</pre>
```

### 1.6 Eulerian Tour

```
multiset<int> graph[N];
stack<int> path;
// -> It suffices to call dfs1 just
// one time leaving from node 0.
// -> To calculate the path,
// call the dfs from the odd degree node.
// -> O(n * log(n))
void dfs1(int u)
{
```

```
while(graph[u].size())
{
  int v = *graph[u].begin();
  graph[u].erase(graph[u].begin());
  graph[v].erase(graph[v].find(u));
  dfs1(v);
}
path.push(u);
```

## 1.7 Floyd Warshall

# 2 Strings

#### 2.1 Aho Corasick

```
//N= tamanho da trie, M tamanho do alfabeto
int to[N][M], Link[N], fim[N];
int idx = 1;
void add_str(string &s)
  int v = 0;
  for (int i = 0; i < s.size(); i++) {</pre>
    if (!to[v][s[i]]) to[v][s[i]] = idx++;
    v = to[v][s[i]];
  fim[v] = 1;
void process()
  queue<int> fila;
  fila.push(0);
  while (!fila.empty()) {
  int cur = fila.front();
    fila.pop():
    int 1 = Link[cur];
    fim[cur] |= fim[l];
for (int i = 0; i < 200; i++) {
      if (to[cur][i]) {
         if (cur != 0)
           Link[to[cur][i]] = to[1][i];
           Link[to[cur][i]] = 0;
         fila.push(to[cur][i]);
       else {
         to[cur][i] = to[1][i];
int resolve(string &s)
  int v = 0, r = 0;
for (int i = 0; i < s.size(); i++) {
  v = to[v][s[i]];</pre>
    if (fim[v]) r++, v = 0;
  return r;
```

#### 2.2 KMP

```
int b[100000];
int sizet, sizep;
void kmpPreprocess(string &text, string &pattern)
  int i = 0, j = -1;
 b[0] = -1;
  while (i < sizep) {
   while (j >= 0 and pattern[i] != pattern[j]) j = b[j];
    i++, j++;
   b[i] = j;
void kmpSearch(string &text, string &pattern)
  kmpPreprocess(text, pattern);
  int i = 0, j = 0;
  while (i < sizet)
    while (j >= 0 and text[i] != pattern[j]) j = b[j];
   if (j == sizep) {
     cout << "Olha a substring do texto " << i - j << endl;</pre>
      j = b[j];
```

# 2.3 Suffix Array

```
* O(nlog^2(n)) para o sufix array
 * O(logn) para o LCP(i,j)
 * LCP de i para j;
struct SA {
  const int L;
  string s:
  vvi P:
  vector<pair< ii,int> > M;
  SA(const string &s) : L(s.size()), s(s), P(1, vi(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = s[i]-'a';
for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
       P.pb(vi(L, 0));
       for (int i = 0; i < L; i++)
    M[i] = mp(mp(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
sort(M.begin(), M.end());</pre>
      for (int i = 0; i < L; i++)
   P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i;
  vi GetSA() {
    vi v=P.back();
     vi ret(v.size());
     for (int i=0; i < v.size(); i++) {</pre>
      ret[v[i]]=i;
     return ret;
  int LCP(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {</pre>
      if (P[k][i] == P[k][j]) {
         i += 1 << k;
          j += 1 << k;
          len += 1 << k;
     return len;
   vi GetLCP(vi &sa)
     vi lcp(sa.size()-1);
     for(int i=0; i < sa.size()-1; i++) {
       lcp[i]=LCP(sa[i],sa[i+1]);
     return lcp;
```

### 2.4 Suffix Array 2

```
Suffix Array. Builing works in O(NlogN).
  Also LCP array is calculated in O(NlogN)
  This code counts number of different substrings in the string.
  Based on problem I from here: http://codeforces.ru/gym/100133
****************************
#include <iostream>
#include <fstream>
#include <cmath>
#include <algorithm>
#include <vector>
#include <set>
#include <map>
#include <stack>
#include <queue>
#include <cstdlib>
#include <cstdio>
#include <string>
#include <cstring>
#include <cassert>
#include <utility>
#include <iomanip>
using namespace std;
const int MAXN = 205000;
const int ALPH = 256;
const int MAXLOG = 20;
int n;
char s[MAXN];
int p[MAXN]; // suffix array itself
int pcur[MAXN];
int c[MAXN][MAXLOG];
int num[MAXN];
int classesNum:
int lcp[MAXN];
void buildSuffixArray() {
 n++;
  for (int i = 0; i < n; i++)
    num[s[i]]++;
  for (int i = 1; i < ALPH; i++)</pre>
   num[i] += num[i - 1];
  for (int i = 0; i < n; i++) {
  p[num[s[i]] - 1] = i;</pre>
    num[s[i]]--;
  c[p[0]][0] = 1;
classesNum = 1;
  for (int i = 1; i < n; i++) {
   if (s[p[i]] != s[p[i - 1]])</pre>
      classesNum++;
    c[p[i]][0] = classesNum;
  for (int i = 1; i++) {
    int half = (1 << (i - 1));</pre>
    for (int j = 0; j < n; j++) {
  pcur[j] = p[j] - half;
  if (pcur[j] < 0)</pre>
        pcur[j] += n;
    for (int j = 1; j <= classesNum; j++)</pre>
      num[j] = 0;
    for (int j = 0; j < n; j++)
      num[c[pcur[j]][i - 1]]++;
    for (int j = 2; j <= classesNum; j++)</pre>
      num[j] += num[j - 1];
    for (int j = n - 1; j >= 0; j--) {
  p[num[c[pcur[j]][i - 1]] - 1] = pcur[j];
  num[c[pcur[j]][i - 1]]--;
    c[p[0]][i] = 1;
    classesNum = 1;
```

```
for (int j = 1; j < n; j++) {
      int p1 = (p[j] + half) % n, p2 = (p[j - 1] + half) % n;
if (c[p[j]][i - 1] != c[p[j - 1]][i - 1] || c[p1][i - 1] != c[p2][i - 1])
       c[p[j]][i] = classesNum;
    if ((1 << i) >= n)
      break;
  for (int i = 0; i < n; i++)
    p[i] = p[i + 1];
  n--:
int getLcp(int a, int b) {
  for (int i = MAXLOG - 1; i >= 0; i--) {
    int curlen = (1 << i);</pre>
    if (curlen > n)
      continue:
    if (c[a][i] == c[b][i]) {
      res += curlen;
      a += curlen:
      b += curlen;
  return res:
void calcLcpArray() {
 for (int i = 0; i < n - 1; i++)
    lcp[i] = getLcp(p[i], p[i + 1]);
int main() {
 assert(freopen("substr.in","r",stdin));
assert(freopen("substr.out","w",stdout));
  gets(s);
 n = strlen(s);
  buildSuffixArray();
  // Now p from 0 to n - 1 contains suffix array of original string
  /*for (int i = 0; i < n; i++) {
  printf("%d ", p[i] + 1);
}*/
    calcLcpArray();
  long long ans = 0;
for (int i = 0; i < n; i++)
ans += n - p[i];</pre>
  for (int i = 1; i < n; i++)
    ans -= lcp[i - 1];
  cout << ans << endl;
  return 0;
```

# 2.5 Suffix Array Disulguinha

```
void build() {
     //sort suffiixes by the first character
    for(int i = 0; i < n; i++) order[i] = i;</pre>
    sort(order.begin(), order.end(), [&](int a, int b){return s[a] < s[b];});</pre>
     rank[order[0]] = 0;
    for(int i = 1; i < n; i++) {
  rank[order[i]] = rank[order[i - 1]];</pre>
      if(s[order[i]] != s[order[i - 1]]) rank[order[i]]++;
    //sort suffixex by the the first 2*p characters, for p in 1, 2, 4, 8,...
for(int p = 1; p < n, rank[order[n - 1]] < n - 1; p <<= 1) {
    for(int i = 0; i < n; i++) {
         x[i] = rank[i];
         y[i] = i + p < n ? rank[i + p] + 1 : 0;
       radixPass(y);
       radixPass(x);
       rank[order[0]] = 0;
      for(int i = 1; i < n; i++) {
  rank[order[i]] = rank[order[i - 1]];</pre>
         if(x[order[i]] != x[order[i - 1]] or y[order[i]] != y[order[i - 1]]) rank[order[i]]++;;
  //Stable counting sort
  void radixPass(vector<int>& key) {
    fill(count.begin(), count.end(), 0);
     for(auto index : order) count[key[index]]++;
    for (int i = 1; i <= n; i++) count[i] += count[i - 1];</pre>
    for (int i = n - 1; i \ge 0; i--) lcp[--count[key[order[i]]]] = order[i];
    order.swap(lcp);
  //Kasai's algorithm to build the LCP array from order, rank and s
  //For i \ge 1, lcp[i] refers to the suffixes starting at order[i] and order[i - 1] void buildLCP() {
    lcp[0] = 0;
    int k = 0;
    for (int i = 0; i < n; i++) {
      if(rank[i] == n - 1){
         k = 0;
       }else{
         int j = order[rank[i] + 1];
         while (i + k < n \text{ and } j + k < n \text{ and } s[i + k] == s[j + k]) k++;
         lcp[rank[j]] = k;
         if(k) k--;
1:
int main(){
 ios::sync_with_stdio(false);
  string s;
  SuffixArray sa(s);
  for(int i = 0; i < s.size(); i++) cout << sa.order[i] << '\n';</pre>
```

# 2.6 Manacher Algorithm

```
Manacher's algorithm for finding all subpalindromes in the string.
Based on problem L from here: http://codeforces.ru/gym/100133

const int MAXN = 105000;

string s;
int n;
int odd[MAXN], even[MAXN];
int l, r;
long long ans;
int main() {
   assert(freopen("palindrome.in","r",stdin));
   assert(freopen("palindrome.out","w",stdout));

getline(cin, s);
   n = (int) s.length();
```

```
// Odd case
  = \mathbf{r} = -1;
for (int i = 0; i < n; i++) {
  int cur = 1;
    cur = min(r - i + 1, odd[1 + r - i]);
  while (i + cur < n && i - cur >= 0 && s[i - cur] == s[i + cur])
   cur++;
  odd[i] = cur;
 if (i + cur - 1 > r) {
   1 = i - cur + 1;
   r = i + cur - 1;
// Even case
1 = r = -1;
for (int i = 0; i < n; i++) {
 int cur = 0;
    cur = min(r - i + 1, even[1 + r - i + 1]);
  while (i + cur < n \&\& i - 1 - cur >= 0 \&\& s[i - 1 - cur] == s[i + cur])
   cur++;
  even[i] = cur;
  if (i + cur - 1 > r) {
   1 = i - cur;
   r = i + cur - 1:
for (int i = 0; i < n; i++) {
 if (odd[i] > 1) {
   ans += odd[i] - 1;
 if (even[i])
   ans += even[i];
cout << ans << endl;
return 0:
```

# 2.7 Splitting String

```
/* String s to be splitted and the delimiter used to split it.*/
vector<string> splitstr(string s, string delimiter)
{
    vector<string> result;
    string str = s, token;
    size_t pos=0;
    while((pos=str.find(delimiter)) != std::string::npos)
    {
        token = str.substr(0, pos);
        result.push_back(token);
        str.erase(0, pos+delimiter.length());
    }
    result.push_back(str);
    return result;
}
```

# 3 Numerical Algorithms

#### 3.1 Fast Fourier Transform

```
// FFT - The Iterative Version
//
Running Time:
// O(n*log n)
//
How To Use:
// fft(a,1)
// fft(b,1)
// mul(a,b)
// fft(a,-1)
//
// INPUT:
// - fft method:
// * The vector representing the polynomial
```

```
* 1 to normal transform
        * -1 to inverse transform
    - mul method:
      * The two polynomials to be multiplyed
// - fft method: Transforms the vector sent.
// - mul method: The result is kept in the first vector.
// - You can either use the struct defined of define dificil as complex<double>
// SOLVED:
// * Codeforces Round #296 (Div. 1) D. Fuzzy Search
struct dificil {
        double real;
        double im;
        dificil() {
                real=0.0;
                im=0.0;
        dificil(double real, double im):real(real),im(im){}
        dificil operator+(const dificil &o)const
                return dificil (o.real+real, im+o.im);
        dificil operator/(double v) const {
                return dificil (real/v, im/v);
        dificil operator * (const dificil &o) const {
                return dificil(real*o.real-im*o.im, real*o.im*im*o.real);
        dificil operator-(const dificil &o) const {
                return dificil (real-o.real, im-o.im);
};
dificil tmp[MAXN*2];
int coco, maiorpot2[MAXN];
void fft(vector<dificil> &A, int s)
        int n = A.size(), p = 0;
        while (n>1) {
                n >>= 1:
        \mathbf{n} = (1 << \mathbf{p});
        vector<dificil> a=A;
        for (int i = 0; i < n; ++i) {
                int rev = 0;
                for (int j = 0; j < p; ++ j) {
                         rev |= ( (i >> j) & 1 );
                A[i] = a[rev];
        dificil w, wn;
        for(int i = 1; i <= p; ++i) {</pre>
                int M = 1 << i;
                int K = M >> 1;
                wn = dificil(cos(s*2.0*pi/(double)M), sin(s*2.0*pi/(double)M));
                for (int j = 0; j < n; j += M) {
                         w = dificil(1.0, 0.0);
                         for (int 1 = j; 1 < K + j; ++1) {
                                 dificil t = w;
                                 t = t *A[1 + K];
                                 dificil u = A[1];
                                 A[1] = A[1] + t;
                                 u = u-t;
A[1 + K] = u;
                                 w = wn*w:
        if(s==-1) {
                for(int i = 0; i < n; ++i)
                        A[i] = A[i] / (double) n;
```

#### 3.2 Fast Fourier Transform 2

```
// FFT - The Recursive Version
// Running Time:
     O(n*log n)
// How To Use:
   fft(&a[0], tam, 0)
// fft(&b[0], tam, 0)
// mul(a,b)
    fft(&a[0], tam, 1)
// INPUT:
// - fft method:
      * The vector representing the polynomial
       * 0 to normal transform
      * 1 to inverse transform
   - mul method:
      * The two polynomials to be multiplyed
    - fft method: Transforms the vector sent.
// - mul method: The result is kept in the first vector.
// NOTES:
// - Tam has to be a power of 2.
// - You can either use the struct defined of define dificil as complex<double>
// SOLVED:
// * Codeforces Round #296 (Div. 1) D. Fuzzy Search
dificil tmp[MAXN*2];
int coco, maiorpot2[MAXN];
void fft(dificil *v, int N, bool inv)
        if(N<=1) return;</pre>
        dificil *vodd = v;
        dificil *veven = v+N/2;
        for(int i=0; i<N; i++) tmp[i] = v[i];</pre>
        for(int i=0; i<N; i+=2)</pre>
                veven[coco] = tmp[i];
                vodd[coco] = tmp[i+1];
                coco++;
        fft(&vodd[0], N/2, inv);
        fft(&veven[0], N/2, inv);
        double angucomleite = 2.0*PI/(double)N;
        if(inv) angucomleite = -angucomleite;
        dificil wn(cos(angucomleite), sin(angucomleite));
        for (int i=0; i<N/2; i++)
                tmp[i] = veven[i]+w*vodd[i];
                tmp[i+N/2] = veven[i]-w*vodd[i];
                w \neq wn:
                if(inv)
                        tmp[i] /= 2;
                        tmp[i+N/2] /= 2;
        for(int i=0; i<N; i++) v[i] = tmp[i];</pre>
void mul(vector<dificil> &a, vector<dificil> &b)
  for(int i=0; i<a.size(); i++)</pre>
    a[i] = a[i] * b[i];
```

```
void precomp()
{
   int pot=0;
   for(int i=1;i<MAXN;i++)
   {
      if((1<<\pot)<i) pot++;
      maiorpot2[i] = (1<<\pot);
   }
}</pre>
```

### 3.3 Simpson Algorithm

```
const int NPASSOS = 100000;
const int W=1000000;
//W= tamanho do intervalo que eu estou integrando
double integral1()
{
    double h = W / (NPASSOS);
    double a = 0;
    double b = W;
    double s = f(a) + f(b);
    for (double i = 1; i <= NPASSOS; i += 2) s += f(a + i * h) * 4.0;
    for (double i = 2; i <= (NPASSOS - 1); i += 2) s += f(a + i * h) * 2.0;
    return s * h / 3.0;
}</pre>
```

# 3.4 Matrix Exponentiation

```
//matmul multiplica m1 por m2
//matpow exponencia a matrix ml por p
//mul vet multiplica a matrix ml pelo vetor vet
vvi matmul(vvi &m1, vvi &m2)
  ans.resize(m1.size(), vi(m2.size(), 0));
 for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
for (int k = 0; k < n; k++) {
    ans[i][j] += ml[i][k] * m2[k][j];</pre>
         ans[i][j] %= MOD;
  return ans;
vvi matpow(vvi &m1, 11 p)
  ans.resize(m1.size(), vi(m1.size(), 0));
  for (int i = 0; i < n; i++) ans[i][i] = 1;</pre>
  while (p) {
   if (p & 1) ans = matmul(ans, m1);
    m1 = matmul(m1, m1);
    p >>= 1;
  return ans;
// VETOR TEM N LINHAS E A MATRIZ E QUADRADA
vi mulvet(vvi &m1, vi &vet)
  ans.resize(vet.size(), 0);
  for (int i = 0; i < n; i++)
    for (int j = 0; j < n; j++) {</pre>
      ans[i] += (m1[i][j] * vet[j]);
      ans[i] %= MOD;
  return ans;
```

### 4 Mathematics

# 4.1 Big Number

```
void zero_esq(string &resp)
  string retorno = resp;
  reverse(retorno.begin(), retorno.end());
  int i = resp.size() - 1;
  while (retorno[i] == '0' and i > 0) {
    retorno.erase(i);
  reverse(retorno.begin(), retorno.end());
  resp = retorno;
string sum_big(string a, string b)
  string resp:
  reverse(a.begin(), a.end());
  reverse(b.begin(), b.end());
  if (a.size() <= b.size()) {
    int carry = 0;
    for (int i = 0; i < a.size(); i++) {
  int x = b[i] - '0' + a[i] - '0' + carry;</pre>
       resp.push_back((char)(x % 10 + '0'));
       carry = x / 10;
    for (int i = a.size(); i < b.size(); i++) {
  int x = b[i] - '0' + carry;
  resp.push_back((char)(x % 10 + '0'));</pre>
       carry = x / 10;
    if (carry > 0) resp.push_back((char)(carry + '0'));
  else {
    for (int i = 0; i < b.size(); i++) {
  int x = a[i] - '0' + b[i] - '0' + carry;
  resp.push_back((char)(x % 10 + '0'));</pre>
       carry = x / 10;
    for (int i = b.size(); i < a.size(); i++) {
  int x = a[i] - '0' + carry;
  resp.push_back((char)(x % 10 + '0'));</pre>
       carry = x / 10;
    if (carry > 0) resp.push_back((char)(carry + '0'));
  reverse(resp.begin(), resp.end());
  zero_esq(resp);
  return resp;
string mul_big(string a, string b)
  string resp;
  resp.push_back('0');
  string temp;
  int carry = 0;
  reverse(a.begin(), a.end());
  reverse(b.begin(), b.end());
  for (int i = 0; i < a.size(); i++) {
    temp.clear();
    for (int k = 0; k < i; k++) temp.push_back('0');
int x = a[i] - '0';</pre>
    for (int j = 0; j < b.size(); j++) {
  int y = b[j] - '0';</pre>
       int novo = (x * y + carry);
       temp.push_back((novo % 10) + '0');
       carry = novo / 10;
    if (carry > 0) temp.push_back(carry + '0');
    reverse(temp.begin(), temp.end());
    carry = 0;
    resp = sum_big(temp, resp);
  zero esq(resp);
  return resp;
```

# 4.2 Big Number 2

```
/*
Structure implementing long arithmetic in C++
Analogue to BigInteger in Java.
Tested on many problems.
TODO: list some problems

**Struct BigInt {
```

```
static const int base = 1000 * 1000 * 1000;
static const int baseDigits = 9;
string leadingZerosModifier;
* CONSTRUCTONS & SETTERS
void setLeadingZerosModifier() +
 leadingZerosModifier = "%0xd";
 leadingZerosModifier[2] = '0' + baseDigits;
void set(int value) {
 num.clear();
 if (value == 0)
    num.push_back(0);
 while (value) {
   num.push_back(value % base);
   value /= base;
void set(long long value) {
 num.clear();
 if (value == 0)
   num.push_back(0);
 while (value) {
   num.push_back(value % base);
   value /= base;
void set(string &value) {
  num.clear();
 for (int i = (int) value.length() - 1; i \ge 0; i = baseDigits) {
   int add = 0;
   for (int j = max(0, i - baseDigits + 1); j <= i; j++)
  add = add * 10 + value[j] - '0';</pre>
   num.push_back(add);
void operator = (int value) {
 set (value);
void operator = (long long value) {
 set (value);
void operator = (string &value) {
 set (value);
BigInt() {
 setLeadingZerosModifier();
 set (0);
BigInt(int value) {
 setLeadingZerosModifier();
 set (value);
BigInt(string value) {
 setLeadingZerosModifier();
 set (value);
* SIZE METHODS
//returns size of vector
int size() {
 return (int) num.size();
//returns length of the number
int digitNum() {
 int result = 0;
 for (int i = 0; i < (int) num.size() - 1; i++)</pre>
   result += baseDigits;
 int lastNum = num.back();
 while (lastNum) {
   result++;
   lastNum /= 10;
 return result;
```

vector <int> num;

\* I/O void read() { string s; cin >> s; num.clear(); for (int  $i = (int)s.length() - 1; i >= 0; i -= baseDigits) {$ int add = 0; for (int j = max(0, i - baseDigits + 1); j <= i; j++)
 add = add \* 10 + s[j] - '0';</pre> num.push\_back(add); void print() { printf("%d", num.back()); for (int i = (int) num.size() - 2; i >= 0; i--) printf (leadingZerosModifier.c\_str(), num[i]); void println() { print();
printf("\n"); string toString() { string result = ""; for (int i = 0; i < (int) num.size(); i++) {</pre> int cur = num[i]; for (int j = 1; j <= baseDigits; j++) {</pre> if (cur == 0 && i == (int) num.size() - 1) break; result.append(1, (char) '0' + cur % 10); cur /= 10; reverse(result.begin(), result.end()); return result; \* ADDITION void sumThis(BigInt number) { int carry = 0;
for (int i = 0; i < max((int)num.size(), number.size()) || carry; i++) {</pre> if (i == num.size()) num.push\_back(0); if (i >= number.size()) carry = num[i] + carry; else carry = num[i] + carry + number.num[i]; num[i] = carry % base; carry /= base; void sumThis(int number) { int carry = number;
for (int i = 0; i < (int) num.size() || carry; i++) {</pre> if (i == num.size()) num.push\_back(0); carry = num[i] + carry; num[i] = carry % base; carry /= base; BigInt sum(BigInt number) { BigInt result = \*this; result.sumThis(number); return result; BigInt sum(int number) { BigInt result = \*this; result.sumThis(number); return result: void operator += (BigInt number) { sumThis(number); void operator += (int number) { sumThis(number);

```
BigInt operator + (BigInt number) {
 return sum(number);
BigInt operator + (int number) {
 return sum(number);
* SUBTRACTION
void subThis(BigInt number) {
 int carry = 0;
for (int i = 0; i < (int) number.size() || carry; i++) {</pre>
   if (i < (int)number.size())</pre>
     num[i] -= carry + number.num[i];
     num[i] -= carry;
   if (num[i] < 0) {
     carry = 1;
     num[i] += base;
   else
     carry = 0;
 while (num.size() > 1 && num.back() == 0)
   num.pop_back();
void subThis(int number) {
 int carry = -number;
 for (int i = 0; carry > 0; i++) {
   num[i] -= carry;
   if (num[i] < 0) {
     carry = 1;
     num[i] += base;
   else
     carry = 0;
 while (num.size() > 1 && num.back() == 0)
   num.pop_back();
BigInt sub(BigInt number) {
 BigInt result = *this;
 result.subThis(number);
 return result;
BigInt sub(int number) {
 BigInt result = +this:
 result.subThis(number):
 return result:
void operator -= (BigInt number) {
 subThis(number);
void operator -= (int number) {
 subThis(number);
BigInt operator - (BigInt number) {
 return sub (number);
BigInt operator - (int number) {
 return sub (number);
* MULTIPLICATION
BigInt mult(BigInt number) {
 BigInt product;
 product.num.resize(num.size() + number.size());
 while (product.size() > 1 && product.num.back() == 0)
   product.num.pop_back();
 return product;
```

```
void multThis(BigInt number) {
  *this = mult(number);
void multThis(int number) {
  int carry = 0;
  for (int i = 0; i < (int)num.size() || carry; ++i) {</pre>
   if (i == num.size())
      num.push_back (0);
   long long cur = carry + num[i] * 111 * number;
   num[i] = int (cur % base);
   carry = int (cur / base);
  while (num.size() > 1 && num.back() == 0)
   num.pop_back();
BigInt mult(int number) {
  BigInt result = *this;
  result.multThis(number);
  return result:
void operator *= (BigInt number) {
 multThis(number):
void operator *= (int number) {
 multThis(number):
BigInt operator * (BigInt number) {
  return mult(number);
BigInt operator * (int number) {
 return mult(number);
void multThisByPowerOfTen(int power) {
 int baseNums = power / baseDigits;
  int curLen = (int)num.size();
  num.resize(curLen + baseNums);
  for (int i = num.size() - 1; i >= baseNums; i--) {
   num[i] = num[i - baseNums];
  for (int i = baseNums - 1; i >= 0; i--)
   num[i] = 0;
  power %= baseDigits;
  int multBy = (int)pow(10.0, power);
 multThis(multBy);
* DIVISION
void divThis(int number) {
 int carry = 0;
for (int i= (int) num.size() - 1; i >= 0; i--) {
   long long cur = num[i] + carry * 111 * base;
    num[i] = int (cur / number);
    carry = int (cur % number);
  while (num.size() > 1 && num.back() == 0)
   num.pop_back();
BigInt div(int number)
 BigInt result = *this;
  result.divThis(number);
  return result;
void operator /= (int number) {
 divThis(number);
BigInt operator / (int number) {
  return div(number);
void divThisByPowerOfTen(int power) {
 int baseNums = power / baseDigits;
  int curLen = (int)num.size();
  for (int i = 0; i < (int)num.size() - baseNums; i++) {</pre>
   num[i] = num[i + baseNums];
  for (int i = 1; i <= baseNums; i++)</pre>
   num.pop_back();
```

```
power %= baseDigits;
  int divBy = (int)pow(10.0, power);
  divThis(divBy);
* MODULO
void modThis(int number) {
  int carry = 0;
  for (int i= (int) num.size() - 1; i >= 0; i--) {
   long long cur = num[i] + carry * 111 * base;
num[i] = int (cur / number);
    carry = int (cur % number);
  set (carry);
BigInt mod(int number) {
  BigInt result = *this;
  result.modThis(number);
  return result;
void operator %= (int number) {
 modThis(number):
BigInt operator % (int number) {
 return mod(number);
* COMPARISON
//Returns: -1 - this number is less than argument, 0 - equal, 1 - this number is greater
int compareTo(BigInt number) {
  if ((int)num.size() < number.size())</pre>
   return -1;
  if ((int)num.size() > number.size())
   return 1;
 for (int i = (int)num.size() - 1; i >= 0; i--) {
   if (num[i] > number.num[i])
      return 1;
    if (num[i] < number.num[i])</pre>
      return -1;
  return 0;
//Returns: -1 - this number is less than argument, 0 - equal, 1 - this number is greater
int compareTo(int number) {
 if (num.size() > 1 || num[0] > number)
   return 1:
  if (num[0] < number)</pre>
   return -1;
  return 0:
bool operator < (BigInt number) {</pre>
  return compareTo(number) == -1;
bool operator < (int number) {
 return compareTo(number) == -1;
bool operator <= (BigInt number) {</pre>
 return compareTo(number) != 1;
bool operator <= (int number) {</pre>
 return compareTo(number) != 1;
bool operator == (BigInt number) {
 return compareTo(number) == 0;
bool operator == (int number) {
 return compareTo(number) == 0;
bool operator > (BigInt number) {
 return compareTo(number) == 1;
bool operator > (int number) {
  return compareTo(number) == 1;
```

```
bool operator >= (BigInt number) {
   return compareTo(number) != -1;
}
bool operator >= (int number) {
   return compareTo(number) != 1;
}
bool operator != (BigInt number) {
   return compareTo(number) != 0;
}
bool operator != (int number) {
   return compareTo(number) != 0;
}
}
```

#### 4.3 Chinese Remainder

```
11 mulmod(11 a, 11 b, 11 m)
  11 \text{ ret} = 0;
  while (b > 0) {
   if (b % 2 != 0) ret = (ret + a) % m;
    a = (a + a) % m;
   b >>= 1:
  return ret:
11 expmod(11 a, 11 e, 11 m)
  11 ret = 1:
  while (e > 0) {
   if (e % 2 != 0) ret = mulmod(ret, a, m);
   a = mulmod(a, a, m);
   e >>= 1;
  return ret;
11 invmul(11 a, 11 m) { return expmod(a, m - 2, m); }
11 chinese(vector<11> r, vector<11> m)
  int sz = m.size();
  11 M = 1:
  for (int i = 0; i < sz; i++) {
   M *= m[i];
  ll ret = 0;
  for (int i = 0; i < sz; i++) {
   ret += mulmod(mulmod(M / m[i], r[i], M), invmul(M / m[i], M), M);
  return ret:
```

### 4.4 Chinese Remainder 2

```
// Chinese remainder theorem (special case): find z such that // z m1 = r1, z
// % m2 = r2. Here, z is unique modulo M = lcm(m1, m2). // Return (z, M). On
// failure, M = -1;
ii chinese remainder theorem (int m1, int r1, int m2, int r2)
 int g = extended_euclid(m1, m2, s, t);
  if (r1 % g != r2 % g) return mp(0, -1);
 return mp(mod(s * r2 * m1 + t * r1 * m2, m1 * m2) / g, m1 * m2 / g);
// Chinese remainder theorem: find z such that // z % m[i] =
// r[i] for all i
     . Note that the solution is unique modulo M = lcm_i (m[i]).
     Return(z, M)
      .On \// failure, M = -1. Note that we do not require the a[i] \/ s
// to be relatively prime.
ii chinese_remainder_theorem(const vi &m, const vi &r)
  ii ret = make_pair(r[0], m[0]);
  for (int i = 1; i < m.size(); i++) {</pre>
    ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
    if (ret.second == -1) break;
```

```
return ret;
```

### 4.5 Matrix Exponentiation

```
//matmul multiplica m1 por m2
//matpow exponencia a matrix m1 por p
//mul vet multiplica a matrix ml pelo vetor vet
vvi matmul(vvi &m1, vvi &m2)
  ans.resize(m1.size(), vi(m2.size(), 0));
  for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
for (int k = 0; k < n; k++)
        ans[i][j] += m1[i][k] * m2[k][j];
        ans[i][j] %= MOD;
  return ans;
vvi matpow(vvi &m1, 11 p)
  ans.resize(m1.size(), vi(m1.size(), 0));
  for (int i = 0; i < n; i++) ans[i][i] = 1;</pre>
  while (p) {
    if (p & 1) ans = matmul(ans, m1);
    m1 = matmul(m1, m1);
    p >>= 1;
  return ans;
// VETOR TEM N LINHAS E A MATRIZ E QUADRADA
vi mulvet(vvi &m1, vi &vet)
  vi ans;
  ans.resize(vet.size(), 0);
  for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++) {
   ans[i] += (m1[i][j] * vet[j]);</pre>
      ans[i] %= MOD;
  return ans:
```

### 4.6 Pascal Triangle

```
//Fazer combinacao de N escolhe M
//por meio do triangulo de pascal
//complexidade: O(m*n)
unsigned long long comb[61][61];
for (int i = 0; i < 61; i++) {
    comb[i][i] = 1;
    comb[i][0] = 1;
}
for (int i = 2; i < 61; i++)
    for (int j = 1; j < i; j++)
        c om b[i][j] = comb[i - 1][j] + comb[i - 1][j - 1];</pre>
```

#### 4.7 Euler's Totient Function

#### 4.8 Pollard Rho

```
11 u;
11 t;
const int tamteste=5;
11 abss(11 v) { return v>=0 ? v : -v;}
  ld pseudo=(ld)rand()/(ld)RAND_MAX;
  return (11) (round((1d) range*pseudo))+1LL;
11 mulmod(11 a, 11 b, 11 mod)
  11 ret=0;
  while (b>0)
    if(b%2!=0) ret=(ret+a)%mod;
    a=(a+a)%mod;
    b=b/2LL;
  return ret;
ll expmod(ll a, ll e, ll mod)
  11 ret=1:
  while (e>0)
    if(e%2!=0) ret=mulmod(ret,a,mod);
    a=mulmod(a,a,mod);
  return ret;
bool jeova(ll a, ll n)
  11 x = expmod(a,u,n);
  ll last=x;
  for(int i=0;i<t;i++)</pre>
    x=mulmod(x,x,n);
    if(x==1 and last!=1 and last!=(n-1)) return true;
    last=x:
  if(x==1) return false;
 return true;
bool isprime(ll n)
  u=n-1;
  while (u%2==0)
    t++;
    u/=2LL:
  if (n==2) return true;
  if(n==3) return true;
  if (n%2==0) return false;
  if(n<2) return false;</pre>
  for(int i=0;i<tamteste;i++)</pre>
    11 v = randerson()%(n-2)+1;
    //cout<<"jeova "<<v<<" "<<n<<endl;
    if(jeova(v,n)) return false;
  return true;
11 gcd(11 a, 11 b) { return !b ? a : gcd(b,a%b);}
11 calc(11 x, 11 n, 11 c)
  return (mulmod(x,x,n)+c)%n;
ll pollard(ll n)
  11 d=1;
  11 i=1;
  11 k=1;
  11 x=2;
11 y=x;
  11 c;
  do
    c=randerson()%n;
  }while(c==0 or (c+2)%n==0);
  while (d!=n)
    if(i==k)
```

```
k \star = 2LL;
        v=x:
        i=0;
    x=calc(x,n,c);
    d=gcd(abss(y-x),n);
    if(d!=1) return d;
vector<ll> getdiv(ll n)
  vector<ll> ret:
  if (n==1) return ret:
  if(isprime(n))
    ret.pb(n);
   return ret;
  11 d = pollard(n);
  ret=getdiv(d);
  vector<11> ret2=getdiv(n/d);
  for(int i=0;i<ret2.size();i++) ret.pb(ret2[i]);</pre>
  return ret;
```

#### 4.9 Sieve of Eratosthenes

```
//esse crivo gera MAXN primos
const int MAX = 1e6;
int primes(MAX);
void gen_primes()
{
  int i, j;
  for (i = 2; i*i <= MAX; i++)
    if (primes[i])
    for (j = i; j * i < MAX; j++) primes[i * j] = 0;
}</pre>
```

### 4.10 Extended Euclidean Algorithm

```
//returns g = gcd(a, b);
//finds x, y such that d= ax+by;
int extended_euclid(int a, int b, int &x, int &y)
{
  int xx = y = 0;
  int yy = x = 1;
  while (b) {
    int q = a / b;
    int t = b;
    b = a % b;
    a = t;
    t = xx;
    xx = x - q * xx;
    x = t;
    t = yy;
    yy = y - q * yy;
    y = t;
}
return a;
}
```

# 4.11 Multiplicative Inverse

```
//computes b such that ab = 1(mod n), returns - 1 on failure
int mod_inverse(int a, int n)
{
  int x, y;
  int g = extended_euclid(a, n, x, y);
  if (g > 1) return -1;
  return (x+n)%n;
}
}
```

# 4.12 Multiplicative Inverse 2

```
//inverso multiplicativo de A % MOD
//certifique de MOD estar definido antes bonito!
//complexidade: O(log(a))
11 mul_inv(11 a)
{
    11 pin0 = MOD, pin = MOD, t, q;
    11 x0 = 0, x1 = 1;
    if (pin == 1) return 1;
    while (a > 1) {
        q = a / pin;
        t = pin, pin = a % pin, a = t;
        t = x0, x0 = x1 - q * x0, x1 = t;
    }
    if (x1 < 0) x1 += pin0;
    return x1;
}</pre>
```

### 4.13 Linear Diophantine Equation

# 5 Combinatorial Optimization

#### 5.1 Dinic

```
//grafo bipartido O(Esqrt(v))
//Para recuperar a resposta, e so colocar um bool
//de false na aresta de retorno e fazer uma bfs/dfs
//andando pelos vertices de capacidade =0 e arestas
//que nao sao de retorno
struct Edge (
 int v, rev;
  int cap:
 Edge(int v_, int cap_, int rev_) : v(v_), rev(rev_), cap(cap_) {}
1:
struct MaxFlow {
 vector<vector<Edge> > g;
  vector<int> level;
  queue<int> q;
  int flow, n;
  \label{eq:maxFlow} \texttt{MaxFlow(int } n\_) \; : \; g(n\_) \, , \; \mbox{level(} n\_) \, , \; \mbox{n(} n\_) \; \ \{ \} 
  void addEdge(int u, int v, int cap)
    if (u == v) return;
    Edge e(v, cap, int(g[v].size()));
    Edge r(u, 0, int(g[u].size()));
    q[u].push back(e);
    g[v].push_back(r);
  bool buildLevelGraph (int src, int sink)
    fill(level.begin(), level.end(), -1);
    while (not q.empty()) q.pop();
    level[src] = 0;
    q.push(src);
    while (not q.empty()) {
      int u = q.front();
       for (auto e = g[u].begin(); e != g[u].end(); ++e) {
        if (not e->cap or level[e->v] != -1) continue;
level[e->v] = level[u] + 1;
         if (e->v == sink) return true;
        q.push(e->v);
    return false;
  int blockingFlow(int u, int sink, int f)
    if (u == sink or not f) return f;
    int fu = f;
    for (auto e = g[u].begin(); e != g[u].end(); ++e) {
   if (not e->cap or level[e->v] != level[u] + 1) continue;
       int mincap = blockingFlow(e->v, sink, min(fu, e->cap));
      if (mincap) {
         g[e->v][e->rev].cap += mincap;
         e->cap -= mincap;
```

```
fu -= mincap;
}
if (f == fu) level[u] = -1;
return f - fu;
}
int maxFlow(int src, int sink)
{
  flow = 0;
  while (buildLevelGraph(src, sink))
    flow += blockingFlow(src, sink, numeric_limits<int>::max());
  return flow;
};
```

# 5.2 Hopcroft-Karp Bipartite Matching

```
/* O(v^3)
* Matching maximo de grafo bipartido de peso 1 nas arestas
* supondo que o grafo bipartido seja enumerado de 0-n-1
class MaxMatch {
  vi graph[N];
  int match[N], us[N];
 public:
  MaxFlow(){};
  void addEdge(int u, int v) { graph[u].pb(v); }
  int dfs(int u)
   if (us[u]) return 0;
    us[u] = 1;
    for (int v : graph[u]) {
     if (match[v] == -1 or (dfs(match[v]))) {
       match[v] = u;
        return 1;
    return 0;
  int maxMatch (int n)
   memset(match, -1, sizeof(match));
   int ret = 0;
for (int i = 0; i < n; i++) {</pre>
     memset(us, 0, sizeof(us));
     ret += dfs(i);
    return ret;
```

# 5.3 Max Bipartite Matching 2

```
// This code performs maximum bipartite matching.
// Running time: O(|E| |V|) -- often much faster in practice
     INPUT: \ w[i][j] = edge \ between \ row \ node \ i \ and \ column \ node \ j
     OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
             mc[i] = assignment for column node i, -1 if unassigned
             function returns number of matches made
#include <vector>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FindMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
  for (int j = 0; j < w[i].size(); j++) {</pre>
    if (w[i][j] && !seen[j]) {
      seen[j] = true;
      if (mc[j] < 0 \mid \mid FindMatch(mc[j], w, mr, mc, seen)) {
        mr[i] = j;
mc[j] = i;
        return true;
  return false;
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
 mr = VI(w.size(), -1);
```

```
mc = VI(w[0].size(), -1);
int ct = 0;
for (int i = 0; i < w.size(); i++) {
    VI seen(w[0].size());
    if (FindMatch(i, w, mr, mc, seen)) ct++;
}
return ct;</pre>
```

# 5.4 Min Cost Matching

```
// Min cost bipartite matching via shortest augmenting paths
// This is an O(n^3) implementation of a shortest augmenting path
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
     cost[i][j] = cost for pairing left node i with right node j
    Lmate[i] = index of right node that left node i pairs with
    Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
#include <algorithm>
#include <cmath>
#include <cstdio>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate)
  int n = int(cost.size());
  // construct dual feasible solution
  VD u(n);
  VD v(n);
  for (int i = 0; i < n; i++) {
    u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
  for (int j = 0; j < n; j++) {
  v[j] = cost[0][j] - u[0];</pre>
    for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);
  // construct primal solution satisfying complementary slackness
  Lmate = VI(n, -1);
  Rmate = VI(n, -1);
 int mated = 0;
int mated = 0;
for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++) {
      if (Rmate[j] != -1) continue;
}</pre>
      if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {</pre>
        Lmate[i] = j;
        Rmate[j] = i;
        mated++;
        break:
  VD dist(n);
  VI dad(n);
  // repeat until primal solution is feasible
  while (mated < n) {
    // find an unmatched left node
    int s = 0:
    while (Lmate[s] != -1) s++;
    // initialize Dijkstra
    fill(dad.begin(), dad.end(), -1);
    fill(seen.begin(), seen.end(), 0);
    for (int k = 0; k < n; k++) dist[k] = cost[s][k] - u[s] - v[k];
    int j = 0;
    while (true) {
```

```
// find closest
     j = -1;
     for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
       if (j == -1 || dist[k] < dist[j]) j = k;</pre>
     seen[j] = 1;
     // termination condition
    if (Rmate[j] == -1) break;
     // relax neighbors
    const int i = Rmate[j];
for (int k = 0; k < n; k++) {</pre>
      if (seen[k]) continue;
const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
       if (dist[k] > new_dist) {
         dist[k] = new_dist;
         dad[k] = j;
   // update dual variables
  for (int k = 0; k < n; k++) {
  if (k == j || !seen[k]) continue;</pre>
    const int i = Rmate[k];
    v[k] += dist[k] - dist[j];
u[i] -= dist[k] - dist[j];
  u[s] += dist[j];
   // augment along path
  while (dad[j] >= 0) {
    const int d = dad[j];
    Rmate[j] = Rmate[d];
    Lmate[Rmate[j]] = j;
    j = d;
  Rmate[j] = s;
Lmate[s] = j;
  mated++:
double value = 0;
for (int i = 0; i < n; i++) value += cost[i][Lmate[i]];</pre>
```

#### 5.5 Min Cost Max Flow

```
int flow[N][N];
vector<pair<int, int> > g[N];
int n, m, k;
inline int ent(int a) { return a * 2; }
inline int out(int a) { return a * 2 + 1; }
inline void addEdge(int a, int b, int custo, int fluxo)
  flow[a][b] += fluxo;
  g[a].push_back(make_pair(b, custo));
  g[b].push_back(make_pair(a, -custo));
int src = N - 1, tgt = N - 2;
int dis[N], pai[N];
inline int dij()
  memset (dis, INF, sizeof dis);
  memset (pai, -1, sizeof pai);
  priority_queue<pair<int, int> > q;
  dis[src] = 0;
  q.push(make_pair(0, src));
  while (!q.empty()) {
    pair<int, int> foo = q.top();
    int x = foo.second, cost = -foo.first;
    if (dis[x] != cost) continue;
for (int i = 0; i < g[x].size(); ++i) {
   int y = g[x][i].first, w = g[x][i].second;
   if (flow[x][y] <= 0) continue;</pre>
       if (dis[y] > dis[x] + w) {
         dis[y] = dis[x] + w;
```

```
pai[v] = x;
        q.push(make_pair(-dis[y], y));
  return dis[tgt] != INF;
int minCost()
  int maxFlow = 0;
  int minC = 0;
  while (dij()) {
   int u = tgt;
int minFlow = INF;
    while (pai[u] != -1) {
     minFlow = min(minFlow, flow[pai[u]][u]);
      u = pai[u];
    maxFlow += minFlow;
    minC += minFlow * dis[tgt];
    while (pai[u] != -1) {
     flow[pai[u]][u] -= minFlow;
      flow[u][pai[u]] += minFlow;
     u = pai[u];
  if (\max Flow != n * k) \min C = -1;
 return minC:
inline void init()
  memset(flow, 0, sizeof flow);
  for (int i = 0; i < N; ++i) {
    g[i].clear();
```

#### 5.6 Min Cost Max Flow 2

```
// Implementation of min cost max flow algorithm using adjacency
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
      max flow:
                          O(|V|^3) augmentations
       min cost max flow: O(|V|^4 * MAX\_EDGE\_COST) augmentations
// INPUT:
       - graph, constructed using AddEdge()
      - source
      - sink
// OUTPUT:
      - (maximum flow value, minimum cost value)
      - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <iostream>
#include <vector>
using namespace std;
typedef vector<VI> VVI;
typedef long long LL;
typedef vector<LL> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const LL INF = numeric_limits<LL>::max() / 4;
struct MinCostMaxFlow {
  int N;
  VVL cap, flow, cost;
  VI found;
  VL dist, pi, width;
  VPII dad:
  \label{eq:minCostMaxFlow(int N): N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),} \\
      found(N), dist(N), pi(N), width(N), dad(N){}
  void AddEdge (int from, int to, LL cap, LL cost)
   this->cap[from][to] = cap;
```

```
this->cost[from][to] = cost;
  void Relax(int s, int k, LL cap, LL cost, int dir)
    LL \ val = dist[s] + pi[s] - pi[k] + cost;
    if (cap && val < dist[k]) {
       dist[k] = val;
       dad[k] = make_pair(s, dir);
       width[k] = min(cap, width[s]);
  LL Dijkstra(int s, int t)
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
      int best = -1;
found[s] = true;
       for (int k = 0; k < N; k++) {
        if (found[k]) continue;
         Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
Relax(s, k, flow[k][s], -cost[k][s], -1);
if (best == -1 || dist[k] < dist[best]) best = k;</pre>
       s = best;
    for (int k = 0; k < N; k++) pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<LL, LL> GetMaxFlow(int s, int t)
    LL totflow = 0, totcost = 0;
while (LL amt = Dijkstra(s, t)) {
       totflow += amt;
       for (int x = t; x != s; x = dad[x].first) {
   if (dad[x].second == 1) {
           flow[dad[x].first][x] += amt;
            totcost += amt * cost[dad[x].first][x];
         else {
            flow[x][dad[x].first] -= amt;
            totcost -= amt * cost[x][dad[x].first];
    return make_pair(totflow, totcost);
1:
```

### 5.7 Edmonds Karp

```
struct Edge {
 int at, where;
  void init(int _at, ll _cap, int _where)
    at = _at, cap = _cap, where = _where;
struct dad {
 int at, up, down;
dad() { at = -1; }
 dad(int _at, int _up, int _down) { at = _at, up = _up, down = _down; }
class MaxFlow {
private:
  vector<vector<Edge> > g;
  11 mf, f;
  vector<dad> p;
 public:
  void augment(int v, ll minEdge)
    if (v == s) {
      f = minEdge:
      return;
    else if (p[v].at != -1) {
```

```
augment(p[v].at, min(minEdge, g[p[v].at][p[v].up].cap));
      g[p[v].at][p[v].up].cap -= f;
      g[v][p[v].down].cap += f;
  void init(int N)
    for (int i = 0; i < g.size(); i++) g[i].clear();</pre>
    mf = 0, f = 0;
    g.resize(N);
  void addEdge(int u, int v, 11 cap)
    A.init(v, cap, g[v].size());
    Edge B;
    B.init(u, 0, q[u].size());
    g[u].pb(A);
    g[v].pb(B);
  int maxFlow(int source, int sink)
    s = source:
    t = sink;
    mf = 0;
    while (true)
      f = 0:
      vector<int> dist(g.size(), INF);
      dist[s] = 0:
      queue<int> q;
      q.push(s);
      p.clear();
      p.resize(g.size());
      while (!q.empty())
        int u = q.front();
        q.pop();
        if (u == t) break;
        for (int i = 0; i < g[u].size(); i++) {
          Edge prox = g[u][i];
          if (dist[prox.at] == INF and prox.cap > 0) {
   dist[prox.at] = dist[u] + 1;
            q.push(prox.at);
            dad paizao(u, i, prox.where);
            p[prox.at] = paizao;
      augment(t, INF);
      if (f == 0) break;
      mf += f;
    return mf;
};
```

# 6 Dynamic Programming

# 6.1 Convex Hull Trick

```
/* Esse convex hull trick e para achar a reta minima!
* Para maximizar a reta dada , basta trocar o '>' para
 * para '<' na funcao query;
 * Nao chamar query com B ou A vazios! Atualizar dp para
* depois fazer a query =)

* ATENCAO COM O DOUBLE!! ESTA EM LONG LONG :)
vi A[N], B[N];
int pont[N]:
bool odomeioehlixo(int r1, int r2, int r3, int j)
  return (B[j][r1] - B[j][r3]) * (A[j][r2] - A[j][r1]) <
         (B[j][r1] - B[j][r2]) * (A[j][r3] - A[j][r1]);
void add(ll a, ll b, int j)
 B[j].pb(b);
  A[j].pb(a);
  while (B[j].size() >= 3 and
    odomeioehlixo(B[j].size() - 3, B[j].size() - 2, B[j].size() - 1, j)) {
B[j].erase(B[j].end() - 2);
    A[j].erase(A[j].end() - 2);
```

#### 6.2 Convex Hull Trick 2

```
* Given a set of pairs (m, b) specifying lines of the form y = m*x + b, process
 \star set of x-coordinate queries each asking to find the minimum y-value when any
* the given lines are evaluated at the specified x. To instead have the queries
 \star optimize for maximum y-value, set the QUERY_MAX flag to true.
 * The following implementation is a fully dynamic variant of the convex hull
 * optimization technique, using a self-balancing binary search tree (std::set)
 * support the ability to call add_line() and get_best() in any desired order.
 * Explanation: http://wcipeg.com/wiki/Convex_hull_trick#Fully_dynamic_variant
 * Time Complexity: O(n log n) on the total number of calls made to add_line(),
 * any length n sequence of arbitrarily interlaced add_line() and get_min()
 * calls.
 * Each individual call to add_line() is O(log n) amortized and each individual
 * call to get_best() is O(log n), where n is the number of lines added so far.
 * Space Complexity: O(n) auxiliary on the number of calls made to add_line().
#include <limits> // std::numeric_limits
#include <set>
class hull_optimizer {
  struct line {
   long long m, b, val;
   double xlo;
   bool is_query;
   bool query_max;
    line(long long m, long long b, long long val, bool is_query, bool query_max)
      this->m = m;
     this->b = b;
     this->val = val;
      this->xlo = -std::numeric limits<double>::max();
     this->is query = is query;
     this->query_max = query_max;
   bool parallel(const line &1) const { return m == 1.m; }
   double intersect(const line &1) const
      if (parallel(l)) return std::numeric_limits<double>::max();
     return (double) (1.b - b) / (m - 1.m);
   bool operator<(const line &1) const
      if (1.is_query) return query_max ? (xlo < 1.val) : (1.val < xlo);</pre>
     return m < 1.m;
  std::set<line> hull;
  bool _query_max;
  typedef std::set<line>::iterator hulliter;
  bool has_prev(hulliter it) const { return it != hull.begin(); }
  bool has_next(hulliter it) const
   return (it != hull.end()) && (++it != hull.end());
  bool irrelevant (hulliter it) const.
   if (!has prev(it) || !has next(it)) return false;
   hulliter prev = it, next = it;
    --prev;
```

```
return _query_max ? prev->intersect(*next) <= prev->intersect(*it)
                       : next->intersect(*prev) <= next->intersect(*it);
  hulliter update_left_border(hulliter it)
    if ((_query_max && !has_prev(it)) || (!_query_max && !has_next(it)))
     return it;
    hulliter it2 = it;
    double val = it->intersect(_query_max ? *--it2 : *++it2);
    line 1(*it);
    1.xlo = val;
    hull.erase(it++):
    return hull.insert(it, 1);
  hull_optimizer(bool query_max = false) { this->_query_max = query_max; }
  void add_line(long long m, long long b)
    line 1(m, b, 0, false, _query_max);
    hulliter it = hull.lower_bound(1);
if (it != hull.end() && it->parallel(1)) {
     if ((_query_max && it->b < b) || (!_query_max && b < it->b))
        hull.erase(it++);
      else
        return:
    it = hull.insert(it, 1);
    if (irrelevant(it)) {
      hull.erase(it);
      return:
    while (has_prev(it) && irrelevant(--it)) hull.erase(it++);
    while (has_next(it) && irrelevant(++it)) hull.erase(it--);
    it = update_left_border(it);
    if (has_prev(it)) update_left_border(--it);
    if (has_next(++it)) update_left_border(++it);
  long long get_best (long long x) const
    line q(0, 0, x, true, _query_max);
    hulliter it = hull.lower bound(q);
    if ( guery max) --it;
    return it->m * x + it->b;
};
/*** Example Usage ***/
#include <cassert>
int main()
 hull_optimizer h;
 h.add_line(3, 0);
 h.add_line(0, 6);
  h.add_line(1, 2);
 h.add_line(2, 1);
  assert(h.get_best(0) == 0);
  assert (h.get_best(2) == 4);
  assert(h.get_best(1) == 3);
  assert(h.get_best(3) == 5);
  return 0;
```

# 6.3 Divide and Conquer

```
//Um exemplo de Divide and conquer:
int MOD = 1e9 + 7;
const int N = 1010;
int dp[N][N], cost[N][N], v[N], pref[N], n, m;
void compDP(int j, int L, int R, int b, int e)
{
    if (L > R) return;
    int mid = (L + R) / 2;
    int idx = -1;
    for (int i = b; i <= min(mid, e); i++)
        if (dp[mid][j] > dp[i][j - 1] + cost[i + 1][mid]) {
        idx = i;
        dp[mid][j] = dp[i][j - 1] + cost[i + 1][mid];
    }
    compDP(j, L, mid - 1, b, idx);
    compDP(j, mid + 1, R, idx, e);
}
//chamada!
```

### 6.4 Longest Increasing Subsequence

```
//asw -> vetor com resposta!!
//asw.size() o tamanho da maior lis
void lis( const vector< int > & v, vector< int > & asw )
 vector<int> pd(v.size(),0), pd_index(v.size()), pred(v.size());
 int maxi = 0, x=0, j=0, ind=0;
 for(int i=0;i<v.size();i++)</pre>
   x = v[i];
   j=lower_bound(pd.begin(),pd.begin()+maxi,x) -pd.begin();
   pd[j] = x;
   pd index[j] = i;
  if(j==maxi)
    ind = i;
  if(pred[i] == j) pd_index[j-1] = -1;
 int pos=maxi-1, k=v[ind];
 asw.resize( maxi );
 while ( pos >= 0 )
   asw[pos--] = k;
ind = pred[ind];
   k = v[ind];
```

# 7 Geometry

### 7.1 Convex Hull Monotone Chain

```
typedef struct sPoint {
        int x, y;
         sPoint(int _x, int _y)
                 y = y;
| point:
bool comp (point a, point b)
        if (a.x == b.x) return a.v < b.v;</pre>
        return a.x < b.x;
int cross(point a, point b, point c) // AB x BC
        a.x = b.x;
        a.y -= b.y;
        b.x -= c.x;
        b.y -= c.y;
        return a.x * b.y - a.y * b.x;
bool isCw(point a, point b, point c) // Clockwise
        return cross(a, b, c) < 0;
// >= if you want to put collinear points on the convex hull
bool isCcw(point a, point b, point c) // Counter Clockwise
        return cross(a, b, c) > 0;
vector<point> convexHull(vector<point> p)
        vector<point> u, 1; // Upper and Lower hulls
        sort(p.begin(), p.end(), comp);
        for (unsigned int i = 0; i < p.size(); i++) {
     while (l.size() > 1 && !isCcw(l[l.size() - 1], l[l.size() - 2], p[i]))
                         1.erase(1.begin() + (1.size() - 1));
                 l.push_back(p[i]);
```

### 7.2 Minimum Enclosing Circle

```
//6.5- Minimum Enclosing Circle
const double eps = 1e-6;
#define CIRCLE circ
#define PT Ponto
#define MP 101
#define eps 1e-9
#define x first
#define v second
typedef double cood;
typedef int num;
typedef int point;
double resp;
cood x[MP], y[MP], ar, ax, ay;
int p[MP];
typedef pair<double, double> ponto;
typedef pair<double, double> Ponto;
double dista(ponto a, ponto b)
  return sqrt((a.first - b.first) * (a.first - b.first) +
               (a.second - b.second) * (a.second - b.second));
bool in(ponto a, pair<double, ponto> c)
  if (dista(a, c.second) - eps < c.first) return true;</pre>
  return false:
bool same (point a, point b)
  return (fabs(x[a] - x[b]) < eps && fabs(y[a] - y[b]) < eps);
bool lexLess(point a, point b)
  if (fabs(x[a] - x[b]) < eps) return y[a] < y[b];
  return x[a] < x[b];</pre>
inline cood dist(cood xx, cood yy, point a)
  return sqrt((xx - x[a]) * (xx - x[a]) + (yy - y[a]) * (yy - y[a]));
inline cood cP(point a, point b, point c)
  return (x[a] - x[b]) * (y[c] - y[b]) - (x[c] - x[b]) * (y[a] - y[b]);
void findCircle(point a, point b, point c, cood& cx, cood& cy)
  cx = 0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (y[b] - y[c])
       0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
           (y[a] - y[b]),
  cy = 0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
           (x[a] - x[b])
       0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (x[b] - x[c]);
 cx /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);

cy /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
void spanCircle2(int k, point p0, point p1, cood& cx, cood& cy, cood& r)
  cx = 0.5 * (x[p0] + x[p1]);
  cy = 0.5 * (y[p0] + y[p1]);
  r = dist(cx, cy, p0);
for (int i = 0; i < k; i++)
   if (dist(cx, cy, p[i]) > r) {
      findCircle(p0, p1, p[i], cx, cy);
      r = dist(cx, cy, p[i]);
```

```
void spanCircle1(int k, point p0, cood& cx, cood& cy, cood& r)
 cx = 0.5 * (x[p0] + x[p[0]]);
 cy = 0.5 * (y[p0] + y[p[0]]);
  r = dist(cx, cy, p0);
  for (int i = 0; i < k; i++)
    if (dist(cx, cy, p[i]) > r) spanCircle2(i, p0, p[i], cx, cy, r);
void spanCircle(int n, cood& cx, cood& cy, cood& r)
  // Bem importante, retirar repetidos
  sort(p, p + 1, lexLess);
 n = unique(p, p + n) - p;
random_shuffle(p, p + n);
  if (n > 1) {
   cx = 0.5 * (x[p[0]] + x[p[1]]);
    cy = 0.5 * (y[p[0]] + y[p[1]]);
    r = dist(cx, cy, p[1]);
   for (int i = 2; i < n; i++)
     if (dist(cx, cy, p[i]) > r) spanCircle1(i, p[i], cx, cy, r);
  else {
   cx = x[0];
   cy = y[0];
r = 0.0;
void solve(vector<pair<double, double> >& v)
 int N = v.size();
for (int i = 0; i < N; i++) {</pre>
   x[i] = v[i].first;
   y[i] = v[i].second;
   p[i] = i;
  spanCircle(N, ax, ay, ar);
```

### 7.3 Minimum Enclosing Circle 2

```
const double eps = 1e-6;
#define CIRCLE circ
#define PT Ponto
#define MP 101
#define eps 1e-9
#define x first
#define y second
typedef double cood;
typedef int num;
typedef int point;
double resp;
cood x[MP], y[MP], ar, ax, ay;
int p[MP]:
typedef pair<double, double> ponto;
typedef pair<double, double> Ponto;
double dista(ponto a, ponto b)
  return sqrt((a.first - b.first) * (a.first - b.first) +
              (a.second - b.second) * (a.second - b.second));
bool in(ponto a, pair<double, ponto> c)
  if (dista(a, c.second) - eps < c.first) return true;</pre>
  return false:
bool same (point a, point b)
  return (fabs(x[a] - x[b]) < eps && fabs(y[a] - y[b]) < eps);
bool lexLess(point a, point b)
  if (fabs(x[a] - x[b]) < eps) return y[a] < y[b];</pre>
  return x[a] < x[b];</pre>
inline cood dist(cood xx, cood yy, point a)
  return sqrt((xx - x[a]) * (xx - x[a]) + (yy - y[a]) * (yy - y[a]));
inline cood cP(point a, point b, point c)
  return (x[a] - x[b]) * (y[c] - y[b]) - (x[c] - x[b]) * (y[a] - y[b]);
```

```
void findCircle(point a, point b, point c, cood& cx, cood& cy)
  cx = 0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (y[b] - y[c])
       0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) 
           (y[a] - y[b]),
  cy = 0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
           (x[a] - x[b]) -
       0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (x[b] - x[c]);
  cx /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
 cy /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
void spanCircle2(int k, point p0, point p1, cood& cx, cood& cy, cood& r)
 cx = 0.5 * (x[p0] + x[p1]);
 cy = 0.5 * (y[p0] + y[p1]);
  r = dist(cx, cy, p0);
for (int i = 0; i < k; i++)
   if (dist(cx, cy, p[i]) > r)
      findCircle(p0, p1, p[i], cx, cy);
      r = dist(cx, cy, p[i]);
void spanCircle1(int k, point p0, cood& cx, cood& cy, cood& r)
 cx = 0.5 * (x[p0] + x[p[0]]);
 cy = 0.5 * (y[p0] + y[p[0]]);
 r = dist(cx, cy, p0);
for (int i = 0; i < k; i++)
   if (dist(cx, cy, p[i]) > r) spanCircle2(i, p0, p[i], cx, cy, r);
void spanCircle(int n, cood& cx, cood& cy, cood& r)
  // Bem importante, retirar repetidos
  sort(p, p + 1, lexLess);
 n = unique(p, p + n) - p;
random_shuffle(p, p + n);
  if (n > 1) {
    cx = 0.5 * (x[p[0]] + x[p[1]]);
    cy = 0.5 * (y[p[0]] + y[p[1]]);
    r = dist(cx, cy, p[1]);
for (int i = 2; i < n; i++)
      if (dist(cx, cy, p[i]) > r) spanCircle1(i, p[i], cx, cy, r);
  else {
    cx = x[0];
    cy = y[0];
    r = 0.0;
void solve(vector<pair<double, double> >& v)
  int N = v.size();
  for (int i = 0; i < N; i++) {
   x[i] = v[i].first;
    y[i] = v[i].second;
    p[i] = i;
  spanCircle(N, ax, ay, ar);
```

# 7.4 Fast Geometry in Cpp

```
#include <iostream>
#include <vector>
#include <vector>
#include <cmath>
#include <cassert>

using namespace std;

double INF = le100;
double EPS = le-12;

struct PT {
    double x, y;
    PT() {}
    PT(double x, double y) : x(x), y(y) {}
    PT(const PT &p) : x(p.x), y(p.y) {}
    PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
    PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
```

```
PT operator * (double c)
                               const { return PT(x*c, y*c ); ]
  PT operator / (double c)
                               const { return PT(x/c, y/c ); }
double dot (PT p, PT q)
                            { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q)
                           { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
  os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x);
PT RotateCW90(PT p)
                       { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
 return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                          double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
  line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
    if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
     return false:
    return true:
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true:
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert (dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
```

// strictly exterior points, and 0 or 1 for the remaining points.

```
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
 // tests for checking point on polygon boundary
bool PointInPolygon (const vector <PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
    if ((p[i].y \le q.y \&\& q.y < p[j].y ||
     p[j].y \le q.y && q.y < p[i].y) &&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
  return c:
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
     return true;
    return false:
// compute intersection of line through points a and b with
// circle centered at c with radius r >
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret:
 b = b-a:
  a = a-c:
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R | | d+min(r, R) < max(r, R)) return ret;</pre>
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if(y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret:
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
 // the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
 return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++)
    for (int k = i+1; k < p.size(); k++) {</pre>
     int j = (i+1) % p.size();
int l = (k+1) % p.size();
      if (i == 1 \mid | j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false:
```

```
return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  << ProjectPointSegment (PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
  // expected: 1 0 1
  cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
       << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
       << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 1 1 1 0
  cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
  // expected: (1,2)
  cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
  // expected: (1,1)
  cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
  vector<PT> v;
  v.push_back(PT(0,0));
  v.push_back(PT(5,0));
  v.push_back(PT(5,5));
  v.push_back(PT(0,5));
  // expected: 1 1 1 0 0
  cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
       << PointInPolygon(v, PT(2,0)) << " "
       << PointInPolygon(v, PT(0,2)) << " "
       << PointInPolygon(v, PT(5,2)) << " "
       << PointInPolygon(v, PT(2,5)) << endl;
  cerr << PointOnPolygon(v, PT(2,2)) << " "
       << PointOnPolygon(v, PT(2,0)) << " "
       << PointOnPolygon(v, PT(0,2)) << " "
       << PointOnPolygon(v, PT(5,2)) << " "
       << PointOnPolygon(v, PT(2,5)) << endl;
  // expected: (1,6)
                (5,4) (4,5)
                blank line
                (4,5) (5,4)
                blank line
                (4,5) (5,4)
  vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
  u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
  u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5); ", cerr << end; for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << end]; u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0); for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << end];
  u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
  // area should be 5.0
  // centroid should be (1.1666666, 1.166666)
  PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
  vector<PT> p(pa, pa+4);
```

```
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;
return 0;</pre>
```

### 8 Data Structures

### 8.1 Disjoint Set Union

```
const int N=500010;
int p[N],rank[N];
void init()
  memset(rank, 0, sizeof(rank));
  for(int i=0;i<N;i++) p[i]=i;</pre>
int findset(int i)
  if(p[i]==i) return i;
  return p[i]=findset(p[i]);
bool same(int i, int j)
  return (findset(i) == findset(j));
void unionSet(int i, int j)
  if (!same(i, j)) {
    int x = findset(i), y=findset(j);
  if (rank[x] > rank[y])
    else {
      if (rank[x] == rank[y]) rank[y]++;
```

## 8.2 Persistent Segment Tree

```
//PRINTAR O NUMERO DE ELEMENTOS DISTINTOS
//EM UM INTERVALO DO ARRAY
const int N = 30010;
int tr[100 * N], L[100 * N], R[100 * N], root[100 * N];
int v[N], mapa[100 * N];
int cont = 0;
void build(int node, int b, int e)
  if (b == e) {
    tr[node] = 0;
  else {
    L[node] = cont++;
    R[node] = cont++;
    build(L[node], b, (b + e) / 2);
    build(R[node], (b + e) / 2 + 1, e);
tr[node] = tr[L[node]] + tr[R[node]];
int update(int node, int b, int e, int i, int val)
  int idx = cont++;
  tr[idx] = tr[node] + val;
  L[idx] = L[node];
  R[idx] = R[node];
  if (b == e) return idx;
  int mid = (b + e) / 2;
  if (i <= mid)
    L[idx] = update(L[node], b, mid, i, val);
    R[idx] = update(R[node], mid + 1, e, i, val);
  return idx;
int query(int nodeL, int nodeR, int b, int e, int i, int j)
  if (b > j \text{ or } i > e) return 0;
  if (i <= b and j >= e) {
    int p1 = tr[nodeR];
```

```
int p2 = tr[nodeL];
   return p1 - p2;
 int mid = (b + e) / 2;
 return query(L[nodeL], L[nodeR], b, mid, i, j) +
        query(R[nodeL], R[nodeR], mid + 1, e, i, j);
int main()
 int n;
 sc(n);
 root[i + 1] = update(root[i], 0, n - 1, i, 1);
     mapa[v[i]] = i;
     root[i + 1] = update(root[i], 0, n - 1, mapa[v[i]], -1);
     mapa[v[i]] = i;
     root[i + 1] = update(root[i + 1], 0, n - 1, i, 1);
 int q;
 sc(a);
 for (int i = 0; i < q; i++) {
   int 1. r:
   sc2(1, r):
   int resp = query(root[1 - 1], root[r], 0, n - 1, 1 - 1, r - 1);
   pri(resp):
 return 0;
```

## 8.3 RMQ of Indices

```
//RMO DE INDICE
class RMO {
 private:
  vi A:
  vi M;
 public:
  RMQ(vi &v)
    M.resize(4 * v.size());
    build(1, 0, v.size() - 1);
  void build(int node, int b, int e)
    if (b == e)
      M[node] = b;
      build(2 * node, b, (b + e) / 2);
build(2 * node + 1, (b + e) / 2 + 1, e);
      if (A[M[2 * node]] <= A[M[2 * node + 1]])</pre>
         M[node] = M[2 * node];
         M[node] = M[2 * node + 1];
  int query(int node, int b, int e, int i, int j)
    int p1, p2;
    if (i > e || j < b) return -1;
if (b >= i and e <= j) return M[node];</pre>
    p1 = query(2 * node, b, (b + e) / 2, i, j);
p2 = query(2 * node + 1, (b + e) / 2 + 1, e, i, j);
    if (p1 == -1) return p2;
    if (p2 == -1) return p1;
    if (A[p1] <= A[p2]) return p1;</pre>
    return p2;
  void atualiza(int node, int b, int e, int i, int val)
    if (i > e || i < b) return;</pre>
    if (e == b) {
      A[i] = val;
      atualiza (2 * node, b, (b + e) / 2, i, val);
       atualiza(2 * node + 1, (b + e) / 2 + 1, e, i, val);
       if (A[M[2 * node]] <= A[M[2 * node + 1]])</pre>
         M[node] = M[2 * node];
```

# 8.4 RSQ with Lazy Propagation

```
//RSQ COM LAZY PROPAGATION!
class RSQ {
 private:
  vll A;
  vll M;
  vll lazy;
 public:
  RSQ(vll &v)
     M.resize(v.size() * 4);
     lazy.assign(v.size() * 4, 0);
     build(1, 0, v.size() - 1);
  void build(int node, int b, int e)
    if (b == e) {
       M[node] = A[b];
       return:
    build(2 * node, b, (b + e) / 2);
build(2 * node + 1, (b + e) / 2 + 1, e);
M[node] = M[2 * node] + M[2 * node + 1];
  void atualiza(int node, int b, int e, int i, int j, ll val)
     if (lazy[node] != 0) {
       M[node] += lazy[node];
       if (b != e) {
         11 inter = (e - b + 1);
         11 \ i1 = (b + e) / 2 - b + 1;
         11 i2 = e - (b + e) / 2;
11 un = lazy[node] / inter;
         lazy[2 * node] += un * i1;
lazy[2 * node + 1] += un * i2;
       lazy[node] = 0;
     if (i > e or j < b) return;</pre>
     if (i <= b and j >= e) {
       11 inter = (e - b + 1);
       M[node] += val * inter;
if (b != e) {
         11 \ i1 = (b + e) / 2 - b + 1;
         11 i2 = e - (b + e) / 2;
         lazy[2 * node] += i1 * (11)val;
         lazy[2 * node + 1] += i2 * (11)val;
       return:
     atualiza(2 * node, b, (b + e) / 2, i, j, val);
    atualiza(2 * node + 1, (b + e) / 2 + 1, e, i, j, val);
M[node] = M[2 * node] + M[2 * node + 1];
   il query (int node, int b, int e, int i, int j)
     if (i > e \text{ or } j < b) return 0;
     11 p1, p2;
     if (lazy[node] != 0) {
       M[node] += lazy[node];
if (b != e) {
    ll inter = (e - b + 1);
         11 \ i1 = (b + e) / 2 - b + 1;
         11 \ i2 = e - (b + e) / 2;
         11 un = lazy[node] / inter;
         lazv[2 * node] += un * i1;
         lazv[2 * node + 1] += un * i2;
       lazy[node] = 0;
     if (i <= b and j >= e) return M[node];
    p1 = query(2 * node, b, (b + e) / 2, i, j);
p2 = query(2 * node + 1, (b + e) / 2 + 1, e, i, j);
     return p1 + p2;
};
```

# 8.5 Segment Tree

```
//compilar em C++11, essa segment tree
//computa qual e o k's elemento compreendido
//no intervalo entre i, j
//presentes no array
vi tr[5 * N];
void build(int node, int b, int e)
  if (b == e)
    tr[node].pb(v[b]);
  else {
    build(2 * node, b, (b + e) / 2);
    build(2 * node + 1, (b + e) / 2 + 1, e);
merget(tr[2 * node], tr[2 * node + 1], tr[node]);
    merge(tr[2 * node].begin(), tr[2 * node].end(), tr[2 * node + 1].begin(),
          tr[2 * node + 1].end(), back_inserter(tr[node]));
int query(int node, int b, int e, int i, int j, int k)
  if (i > e \text{ or } b > j) return 0;
  if (i \le b \text{ and } j \ge e) {
    int resp =
        upper_bound(tr[node].begin(), tr[node].end(), k) - tr[node].begin();
    return tr[node].size() - resp;
  return query(2 * node, b, (b + e) / 2, i, j, k) +
query(2 * node + 1, (b + e) / 2 + 1, e, i, j, k);
```

### 8.6 Sparse Table

```
//comutar RMQ , favor inicializar: dp[i][0]=v[0]
//sendo v[0] o vetor do rmq
//chamar o build!
int dp[200100][22];
int n;
int d[200100];
void build()
{
    d[0] = d[1] = 0;
    for (int i = 2; i < n; i++) d[i] = d[i >> 1] + 1;
    for (int j = 1; j < 22; j++) {
        for (int i = 0; i + (1 << (j - 1)) < n; i++) {
            dp[i][j] = min(dp[i][j - 1], dp[i + (1 << (j - 1))][j - 1]);
        }
    }
}
int query(int i, int j)
{
    int k = d[j - i];
    int x = min(dp[i][k], dp[j - (1 << k) + 1][k]);
    return x;
}</pre>
```

# 9 Miscellaneous

# 9.1 Hashing

```
//certificar que gethash() foi chamado
//antes de getHash(i,j);
struct Hashing {
    const string &s;
    int n, idx;
    vector<ll> hashes,M,B;
    Hashing (const string &s) : s(s), hashes(s.size()){
    M={100000409, 2000003273, 2000003281, 2000003293};
    B={31, 53, 61, 41};
    srand(time(NULL));
    idx=rand()%4;
    getHash();
}
void otherprime(){
    idx=(idx+1)%4;
}
```

```
ll int_mod(ll a) { return (a % M[idx] + M[idx]) % M[idx]; }
  ll eleva(ll a, ll b)
    if (b == 0)
     return 1;
    else if (b == 1)
     return a;
    11 x = eleva(a, b / 2);
    if (b % 2 == 0)
     return (x * x) % M[idx];
    else
      return (a * ((x * x) % M[idx])) % M[idx];
  /*hash da string de 0 ate i*/
  void getHash()
    int n = s.size();
    11 hp = 0;
    for (int i = 0; i < s.size(); i++) {</pre>
      hp = int\_mod(hp * B[idx] + s[i]);
      hashes[i] = hp;
  /*Hash da string compreendida entre i e j*/
  11 getHash(int i, int j)
    if (i == 0) return hashes[j];
    11 h1 = hashes[j];
    11 h2 = (hashes[i - 1] * eleva(B[idx], j - i+1)) % M[idx];
11 ret = (h1 - h2) % M[idx] + M[idx];
    return ret % M[idx];
};
```

#### 9.2 Invertion Count

```
//conta o numero de inversoes de um array
//x e o tamanho do array, v e o array que quero contar
11 inversoes = 0:
void merge_sort(vi &v, int x)
  if (x == 1) return;
  int tam_esq = (x + 1) / 2, tam_dir = x / 2;
  int esq[tam_esq], dir[tam_dir];
for (int i = 0; i < tam_esq; i++) esq[i] = v[i];
for (int i = 0; i < tam_esq; i++) dir[i] = v[i + tam_esq];</pre>
  merge_sort(esq, tam_esq);
  merge_sort(dir, tam_dir);
  int i_esq = 0, i_dir = 0, i = 0;
  while (i_esq < tam_esq or i_dir < tam_dir) {</pre>
    if (i_esq == tam_esq) {
       while (i_dir != tam_dir) {
        v[i] = dir[i_dir];
         i_dir++, i++;
    else if (i_dir == tam_dir)
      while (i_esq != tam_esq)
        v[i] = esq[i_esq];
         i_esq++, i++;
         inversoes += i_dir;
    else {
      if (esq[i_esq] <= dir[i_dir]) {</pre>
         v[i] = esq[i_esq];
         i++, i_esq++;
         inversoes += i_dir;
       else {
        v[i] = dir[i_dir];
         i++, i_dir++;
```

# 9.3 Distinct Elements in ranges

```
const int MOD = 1e9 + 7;
const int N = 1e6 + 10;
int bit[N], v[N], id[N], r[N];
ii query[N];
```

```
int mapa[N];
bool compare(int x, int y) { return query[x] < query[y]; }</pre>
void add(int idx, int val)
  while (idx < N) {
    bit[idx] += val;
    idx += idx & -idx;
int sum(int idx)
  int ret = 0;
  while (idx > 0) {
    ret += bit[idx];
    idx -= idx & -idx;
  return ret;
  memset(bit, 0, sizeof(bit));
  memset(mapa, 0, sizeof(mapa));
  int n;
  sc(n);
  for (int i = 1; i <= n; i++) sc(v[i]);</pre>
  int q;
  sc(a):
  for (int i = 0; i < q; i++) {
    sc2(query[i].second, query[i].first);
    id[i] = i;
  sort(id, id + q, compare);
sort(query, query + q);
  for (int i = 0; i < q; i++) {
    int L = query[i].second;
    int R = query[i].first;
    while (j \le R)
      if (mapa[v[j]] > 0) {
        add(mapa[v[j]], -1);
        mapa[v[j]] = j;
        add(mapa[v[j]], 1);
      else {
        mapa[v[j]] = j;
        add(mapa[v[j]], 1);
    r[id[i]] = sum(R);
    if (L > 1) r[id[i]] -= sum(L - 1);
  for (int i = 0; i < q; i++) pri(r[i]);</pre>
 return 0;
```

## 9.4 Maximum Rectangular Area in Histogram

```
/*
    * Complexidade : O(N)
    */

11 solve(vi &h)
{
    int n = h.size();
```

### 9.5 Multiplying Two LL mod n

### 9.6 Josephus Problem

```
/* Josephus Problem - It returns the position to be, in order to not die. 0(n)*/ /* With k=2, for instance, the game begins with 2 being killed and then n+2, n+4, ... */ li josephus(l1 n, l1 k) {
    if (n==1) return 1;
    else return (josephus(n-1, k)+k-1)%n+1;
}
```

## 9.7 Josephus Problem 2