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```
/* Supondo que cada vertice u, o seu
* positivo e 2*u, e negativo e 2*i+1
 * resposta[i]=0, significa que o positivo de i e resposta
 * resposta[i]=1, significa que o negativo de i e resposta
 * chamar Sat(n) , n e o numero de vertices do grafo
 * contando com os negativos .. na maioria dos problemas
 * chamar 2*n;
 * testado em :http://codeforces.com/contest/781/problem/D
int resposta[N];
vi graph[N], rev[N];
int us[N];
stack<int> pilha;
void dfs1(int u)
    us[u] = 1;
    for (int v : graph[u])
       if (!us[v]) dfs1(v);
    pilha.push(u);
void dfs2(int u, int color)
    us[u] = color;
    for (int v : rev[u])
        if (!us[v]) dfs2(v, color);
int Sat(int n)
    for (int i = 0; i < n; i++)</pre>
       if (!us[i]) dfs1(i);
    int color = 1;
    memset(us, 0, sizeof(us));
    while (!pilha.empty()) {
        int topo = pilha.top();
        pilha.pop();
        if (!us[topo]) dfs2(topo, color++);
    for (int i = 0; i < n; i += 2) {</pre>
        if (us[i] == us[i + 1]) return 0;
resposta[i / 2] = (us[i] < us[i + 1]);</pre>
    return 1:
inline void add(int u, int v)
    graph[u].pb(v);
    rev[v].pb(u);
```

2.2 Kosaraju

```
//Retorna os componentes fortemente conectados
//Se o usados[i]=usados[j], temos que i e j
//pertencem ao mesmo componente, col-1= numero
//de componentes fortemente conectados do grafo
class kosaraju {
private:
  vi usados;
  vvi graph;
  vvi trans;
  vi pilha;
  kosaraju(int N)
    graph.resize(N);
    trans.resize(N):
  void AddEdge(int u, int v)
    graph[u].pb(v);
    trans[v].pb(u);
  void dfs(int u, int pass, int color)
    usados[u] = color;
    vi vizinhos;
    if (pass == 1)
      vizinhos = graph[u];
    else
      vizinhos = trans[u];
    for (int j = 0; j < vizinhos.size(); j++) {
  int v = vizinhos[j];</pre>
      if (usados[v] == 0) {
        dfs(v, pass, color);
```

```
pilha.pb(u);
}
int SSC(int n) {
    pilha.clear();
    usados.assign(n, 0);
    for (int i = 0; i < n; i++) {
        if (!usados[i]) dfs(i, 1, 1);
    }
    usados.assign(n, 0);
    int color = 1;
    for (int i = n - 1; i >= 0; i--) {
        if (usados[pilha[i]] == 0) {
            dfs(pilha[i], 2, color);
            color++;
        }
    return color - 1;
}
```

2.3 Tree Isomorphism

```
vvi children, subtreeLabels, tree, L;
vi pred, map;
int n;
bool compare(int a, int b) { return subtreeLabels[a] < subtreeLabels[b]; }</pre>
bool equals(int a, int b) { return subtreeLabels[a] == subtreeLabels[b]; }
void generateMapping(int r1, int r2)
  map.resize(n);
  map[r1] = r2 - n;
  sort(children[r1].begin(), children[r1].end(), compare);
  sort(children[r2].begin(), children[r2].end(), compare);
  for (int i = 0; i < (int)children[r1].size(); i++) {</pre>
   int u = children[r1][i];
int v = children[r2][i];
    generateMapping(u, v);
vi findCenter(int offset = 0)
  int cnt = n;
  vi a;
  vi deg(n);
  for (int i = 0; i < n; i++) {
    deg[i] = tree[i + offset].size();
if (deg[i] <= 1) {</pre>
     a.push_back(i + offset);
      --cnt;
  while (cnt > 0) {
    for (int i = 0; i < (int)a.size(); i++) {</pre>
      int u = a[i];
      for (int j = 0; j < (int)tree[u].size(); j++) {</pre>
        int v = tree[u][j];
        if (--deg[v - offset] == 1) {
          na.push_back(v);
    a = na:
  return a:
int dfs (int u, int p = -1, int depth = 0)
  L[depth].push_back(u);
  for (int i = 0; i < (int)tree[u].size(); i++) {</pre>
    int v = tree[u][i];
    if (v == p) continue;
    pred[v] = u;
    children[u].push_back(v);
    h = max(h, dfs(v, u, depth + 1));
  return h + 1:
bool rootedTreeIsomorphism(int r1, int r2)
```

```
pred.assign(2 * n, -1);
  children.assign(2 * n, vi());
  int h2 = dfs(r2);
  if (h1 != h2) return false;
  int h = h1 - 1;
  vi label(2 * n);
  subtreeLabels.assign(2 * n, vi());
  for (int i = h - 1; i >= 0; i--) {
  for (int j = 0; j < (int)L[i + 1].size(); j++) {
    int v = L[i + 1][j];
}</pre>
      subtreeLabels[pred[v]].push_back(label[v]);
    for (int j = 0; j < (int)L[i].size(); j++) {</pre>
      int v = L[i][j];
      sort(subtreeLabels[v].begin(), subtreeLabels[v].end());
    sort(L[i].begin(), L[i].end(), compare);
    for (int j = 0, cnt = 0; j < (int)L[i].size(); j++) {
  if (j && !equals(L[i][j], L[i][j - 1])) ++cnt;</pre>
      label[L[i][j]] = cnt;
  if (!equals(r1, r2)) return false;
  generateMapping(r1, r2);
  return true;
bool treeIsomorphism()
  vi c1 = findCenter();
  vi c2 = findCenter(n);
  if (c1.size() == c2.size()) {
    if (rootedTreeIsomorphism(c1[0], c2[0]))
      return true:
    else if (c1.size() > 1)
      return rootedTreeIsomorphism(c1[1], c2[0]);
  return false;
int main()
  n = 5;
  vvi t1(n);
  t1[0].pb(1); t1[1].pb(0);
  t1[1].pb(2); t1[2].pb(1);
  t1[1].pb(3); t1[3].pb(1);
  t1[0].pb(4); t1[4].pb(0);
  vvi t2(n);
  t2[0].pb(1); t2[1].pb(0);
  t2[0].pb(4); t2[4].pb(0);
  t2[4].pb(3); t2[3].pb(4);
  t2[4].pb(2); t2[2].pb(4);
  tree.assign(2 * n, vi());
  for (int u = 0; u < n; u++)
    for (int i = 0; i < t1[u].size(); i++) {</pre>
      int v = t1[u][i];
      tree[u].push_back(v);
    for (int i = 0; i < t2[u].size(); i++) {
  int v = t2[u][i];</pre>
      tree[u + n].push_back(v + n);
  bool res = treeIsomorphism();
  cout << res << endl;
    for (int i = 0; i < n; i++)
      cout << map[i] << endl;</pre>
```

L.assign(n, vi());

2.4 LCA

```
//antes de usar as queries de lca, e etc..
//certifique-se de chamar a dfs da arvore e
//process()
const int N = 100000;
const int M = 22;
int P[N][M];
```

```
int big[N][M], low[N][M], level[N];
vii graph[N];
int n;
void dfs(int u, int last, int 1)
  level[u] = 1;
  P[u][0] = last;
  for (ii v : graph[u])
  if (v.first != last) {
     big[v.first][0] = low[v.first][0] = v.second;
dfs(v.first, u, 1 + 1);
void process()
  for (int j = 1; j < M; j++)
    for (int i = 1; i <= n; i++) {
      P[i][j] = P[P[i][j-1]][j-1];

big[i][j] = max(big[i][j-1], big[P[i][j-1]][j-1]);
      low[i][j] = min(low[i][j-1], low[P[i][j-1]][j-1]);
int lca(int u, int v)
  if (level[u] < level[v]) swap(u, v);</pre>
  for (int i = M - 1; i >= 0; i--)
   if (level[u] - (1 << i) >= level[v]) u = P[u][i];
  if (u == v) return u;
  for (int i = M - 1; i >= 0; i--) {
    if (P[u][i] != P[v][i]) u = P[u][i], v = P[v][i];
  return P[u][0];
int maximum(int u, int v, int x)
  int resp = 0;
  for (int i = M - 1; i >= 0; i--)
    if (level[u] - (1 << i) >= level[x]) {
      resp = max(resp, big[u][i]);
      u = P[u][i];
  for (int i = M - 1; i >= 0; i--)
    if (level[v] - (1 << i) >= level[x]) {
      resp = max(resp, big[v][i]);
      v = P[v][i];
  return resp;
int minimum(int u, int v, int x)
  int resp = INF:
  for (int i = M - 1; i >= 0; i--)
   if (level[u] - (1 << i) >= level[x]) {
     resp = min(resp, low[u][i]);
      u = P[u][i];
  for (int i = M - 1; i >= 0; i--)
    if (level[v] - (1 << i) >= level[x]) {
      resp = min(resp, low[v][i]);
      v = P[v][i];
  return resp;
```

2.5 Bridges and Articulation Points

```
class ponte {
  private:
    vvi graph;
    vi usados;
    vi e_articulacao;
    vi dfs_low;
    vi dfs_prof;
    vector
public:
    ponte(int N)
    {
        graph.clear();
        graph.resize(N);
        usados.assign(N, 0);
```

```
dfs_low.assign(N, 0);
  dfs_prof.assign(N, 0);
  e_articulacao.assign(N, 0);
  tempo = 0;
void AddEdge(int u, int v)
  graph[u].pb(v);
  graph[v].pb(u);
void dfs(int u, int pai)
  usados[u] = 1;
  int nf = 0;
  dfs_low[u] = dfs_prof[u] = tempo++;
for (int v : graph[u]) {
    if (!usados[v]) {
      dfs(v, u);
      nf++;
      if (dfs_low[v] >= dfs_prof[u] and pai != -1) e_articulacao[u] = true;
      if (pai == -1 and nf > 1) e_articulacao[u] = true;
      if (dfs_low[v] > dfs_prof[u]) pontes.pb(mp(u, v));
      dfs_low[u] = min(dfs_low[u], dfs_low[v]);
    else if (v != pai)
  dfs_low[u] = min(dfs_low[u], dfs_prof[v]);
void olha_as_pontes()
  for (int i = 0; i < graph.size(); i++)</pre>
    if (!usados[i]) dfs(i, -1);
  if (pontes.size() == 0)
    cout << " Que merda! nao tem ponte!" << endl;</pre>
    for (ii i : pontes) cout << i.first << " " << i.second << endl;
void olha_as_art()
  for (int i = 0; i < graph.size(); i++)</pre>
  if (!usados[i]) dfs(i, -1);
for (int i = 0; i < e_articulacao.size(); i++)
   if (e_articulacao[i]) cout << " OIAAA A PONTE " << i << endl;</pre>
```

2.6 Eulerian Tour

```
multiset<int> graph[N];
stack<int> path;
// -> It suffices to call dfs1 just
// one time leaving from node 0.
// -> To calculate the path,
// call the dfs from the odd degree node.
// -> 0(n * log(n))
void dfs1(int u)
{
  while (graph[u].size())
  {
    int v = *graph[u].begin();
    graph[u].erase(graph[u].begin());
    graph[v].erase(graph[v].find(u));
    dfs1(v);
  }
  path.push(u);
```

2.7 Floyd Warshall

3 Strings

3.1 Aho Corasick

```
//N= tamanho da trie, M tamanho do alfabeto
int to[N][M], Link[N], fim[N];
int idx = 1:
void add_str(string &s)
  int v = 0;
  for (int i = 0; i < s.size(); i++) {
  if (!to[v][s[i]]) to[v][s[i]] = idx++;</pre>
  fim[v] = 1;
void process()
  gueue<int> fila:
  fila.push(0):
  while (!fila.empty()) {
   int cur = fila.front();
    fila.pop();
    int 1 = Link[cur];
    fim[cur] |= fim[1];
    for (int i = 0; i < 200; i++) {
      if (to[cur][i]) {
         if (cur != 0)
           Link[to[cur][i]] = to[1][i];
         else
           Link[to[cur][i]] = 0;
         fila.push(to[cur][i]);
       else {
         to[cur][i] = to[1][i];
int resolve(string &s)
  int v = 0, r = 0;
for (int i = 0; i < s.size(); i++) {
  v = to[v][s[i]];</pre>
    if (fim[v]) r++, v = 0;
  return r:
```

3.2 KMP

```
ret++;
}
return ret;
```

3.3 Suffix Array

```
* O(nlog^2(n)) para o sufix array
 * O(logn) para o LCP(i,j)
 * LCP de i para j;
struct SA {
  const int L;
  string s;
  vvi P:
  vector<pair< ii,int> > M;
  SA(const string \&s) : L(s.size()), s(s), P(1, vi(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] =s[i]-'a';
for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {</pre>
       P.pb(vi(L, 0));
       for (int i = 0; i < L; i++)
        M[i] = mp(mp(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
       sort(M.begin(), M.end());
       for (int i = 0; i < L; i++)
          P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i; 
  vi GetSA() {
    vi v=P.back();
    vi ret(v.size());
    for (int i=0; i < v.size(); i++) {</pre>
      ret[v[i]]=i;
    return ret;
  int LCP(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
   if (P[k][i] == P[k][j]) {
        i += 1 << k;
i += 1 << k;
         len += 1 << k:
    return len;
  vi GetLCP(vi &sa)
    vi lcp(sa.size()-1);
    for (int i=0; i < sa.size() -1; i++) {
    lcp[i]=LCP(sa[i], sa[i+1]);</pre>
    return lcp;
};
```

3.4 Suffix Array 2

```
/************************
 Suffix Array, Builing works in O(NlogN).
 Also LCP array is calculated in O(NlogN).
 This code counts number of different substrings in the string.
 Based on problem I from here: http://codeforces.ru/gym/100133
const int MAXN = 205000;
const int ALPH = 256;
const int MAXLOG = 20;
int n;
char s[MAXN];
int p[MAXN]; // suffix array itself
int pcur[MAXN];
int c[MAXN] [MAXLOG];
int num[MAXN];
int classesNum;
int lcp[MAXN];
```

```
void buildSuffixArray() {
 n++;
  for (int i = 0; i < n; i++)
    num[s[i]]++;
  for (int i = 1; i < ALPH; i++)
    num[i] += num[i - 1];
  for (int i = 0; i < n; i++) {
    p[num[s[i]] - 1] = i;
    num[s[i]]--;
  c[p[0]][0] = 1;
classesNum = 1;
  for (int i = 1; i < n; i++) {
    if (s[p[i]] != s[p[i - 1]])
    c[p[i]][0] = classesNum;
  for (int i = 1; i++) {
    int half = (1 << (i - 1));</pre>
    for (int j = 0; j < n; j++) {
      pcur[j] = p[j] - half;
if (pcur[j] < 0)</pre>
         pcur[j] += n;
    for (int j = 1; j <= classesNum; j++)</pre>
      num[j] = 0;
     for (int j = 0; j < n; j++)
    num[c[pcur[j]][i - 1]]++;
for (int j = 2; j <= classesNum; j++)</pre>
      num[j] += num[j - 1];
    for (int j = n - 1; j >= 0; j--) {
   p[num[c[pcur[j]][i - 1]] - 1] = pcur[j];
   num[c[pcur[j]][i - 1]]--;
    c[p[0]][i] = 1;
    classesNum = 1;
     for (int j = 1; j < n; j++) {
      int p1 = (p[j] + half) % n, p2 = (p[j - 1] + half) % n;
if (c[p[j]][i - 1] != c[p[j - 1]][i - 1] || c[p1][i - 1] != c[p2][i - 1])
         classesNum++;
      c[p[j]][i] = classesNum;
    if ((1 << i) >= n)
      break:
  for (int i = 0; i < n; i++)
   p[i] = p[i + 1];
  n--;
int getLcp(int a, int b) {
  int res = 0;
  for (int i = MAXLOG - 1; i >= 0; i--) {
    int curlen = (1 << i);</pre>
    if (curlen > n)
      continue:
    if (c[a][i] == c[b][i]) {
      res += curlen;
      a += curlen;
      b += curlen;
  return res;
void calcLcpArray() {
  for (int i = 0; i < n - 1; i++)
    lcp[i] = getLcp(p[i], p[i + 1]);
int main() {
 assert(freopen("substr.in", "r", stdin));
  assert(freopen("substr.out", "w", stdout));
  gets(s);
  n = strlen(s);
```

buildSuffixArray();

```
// Now p from 0 to n - 1 contains suffix array of original string
/*for (int i = 0; i < n; i++) {
    printf("%d ", p[i] + 1);
}*/
    calcLcpArray();
long long ans = 0;
for (int i = 0; i < n; i++)
    ans += n - p[i];
for (int i = 1; i < n; i++)
    ans -= lcp[i - 1];
cout << ans << endl;
return 0;</pre>
```

3.5 Suffix Array Dilsunguinha

```
struct SuffixArray{
  const string& s;
  int n;
  vector<int> order, rank, lcp;
  vector<int> count, x, y;
  SuffixArray(const string& s) : s(s), n(s.size()), order(n), rank(n),
                                        count(n + 1), x(n), y(n), lcp(n) {
    build();
    buildLCP();
  void build() {
    //sort suffiixes by the first character
    for(int i = 0; i < n; i++) order[i] = i;
sort(order.begin(), order.end(), [&](int a, int b) {return s[a] < s[b];});</pre>
    rank[order[0]] = 0;
    for(int i = 1; i < n; i++) {
  rank[order[i]] = rank[order[i - 1]];</pre>
      if(s[order[i]] != s[order[i - 1]]) rank[order[i]]++;
     //sort suffixex by the the first 2*p characters, for p in 1, 2, 4, 8,...
    for(int p = 1; p < n, rank[order[n - 1]] < n - 1; p <<= 1) {
   for(int i = 0; i < n; i++) {</pre>
         x[i] = rank[i];
         y[i] = i + p < n ? rank[i + p] + 1 : 0;
       radixPass(v);
       radixPass(x):
       rank[order[0]] = 0;
      for(int i = 1; i < n; i++) {
  rank[order[i]] = rank[order[i - 1]];</pre>
         if(x[order[i]] != x[order[i - 1]] or y[order[i]] != y[order[i - 1]]) rank[order[i]]++;;
  //Stable counting sort
  void radixPass(vector<int>& key) {
    fill(count.begin(), count.end(), 0);
    for(auto index : order) count[key[index]]++;
for(int i = 1; i <= n; i++) count[i] += count[i - 1];
for(int i = n - 1; i >= 0; i--) lcp[--count[key[order[i]]]] = order[i];
    order.swap(lcp);
  //Kasai's algorithm to build the LCP array from order, rank and s
  //For i \ge 1, lcp[i] refers to the suffixes starting at order[i] and order[i - 1]
  void buildLCP() {
    lcp[0] = 0;
    int k = 0;
    for (int i = 0; i < n; i++) {</pre>
      if(rank[i] == n - 1){
         k = 0;
         int j = order[rank[i] + 1];
         while (i + k < n \text{ and } j + k < n \text{ and } s[i + k] == s[j + k]) k++;
         lcp[rank[j]] = k;
         if(k) k--;
```

```
int main() {
  ios::sync_with_stdio(false);
  string s;
  cin >> s;
  SuffixArray sa(s);
  for(int i = 0; i < s.size(); i++) cout << sa.order[i] << '\n';
}</pre>
```

3.6 Manacher Algorithm

```
Manacher's algorithm for finding all subpalindromes in the string.
 Based on problem L from here: http://codeforces.ru/gym/100133
*******************************
const int MAXN = 105000;
string s;
int n;
int odd[MAXN], even[MAXN];
long long ans;
  assert(freopen("palindrome.in","r",stdin));
 assert(freopen("palindrome.out", "w", stdout));
  getline(cin, s);
 n = (int) s.length();
  // Odd case
  1 = r = -1;
  for (int i = 0; i < n; i++) {
   int cur = 1;
     cur = min(r - i + 1, odd[1 + r - i]);
   while (i + cur < n \&\& i - cur >= 0 \&\& s[i - cur] == s[i + cur])
     cur++;
   odd[i] = cur;
   if (i + cur - 1 > r) {
    1 = i - cur + 1:
     r = i + cur - 1:
  // Even case
  1 = r = -1;
  for (int i = 0; i < n; i++) {</pre>
   int cur = 0;
     cur = min(r - i + 1, even[1 + r - i + 1]);
   while (i + cur < n \&\& i - 1 - cur >= 0 \&\& s[i - 1 - cur] == s[i + cur])
    cur++;
   even[i] = cur:
   if (i + cur - 1 > r) {
    1 = i - cur:
     r = i + cur - 1:
  for (int i = 0; i < n; i++) {
   if (odd[i] > 1) {
     ans += odd[i] - 1;
   if (even[i])
     ans += even[i];
 cout << ans << endl:
 return 0;
```

3.7 Splitting String

```
/* String s to be splitted and the delimiter used to split it. */ vector<string> splitstr(string s, string delimiter)
```

```
vector<string> result;
string str = s, token;
size_t pos=0;
while((pos=str.find(delimiter)) != std::string::npos)
{
   token = str.substr(0, pos);
   result.push_back(token);
   str.erase(0, pos + delimiter.length());
}
result.push_back(str);
return result;
```

4 Numerical Algorithms

4.1 Fast Fourier Transform

```
// FFT - The Iterative Version
// Running Time:
     O(n*log n)
// How To Use:
// fft(a,1)
// fft(b,1)
// mul(a,b)
// fft(a,-1)
// INPUT:
    - fft method:
     * The vector representing the polynomial
       * 1 to normal transform
       * -1 to inverse transform
// - mul method:
      * The two polynomials to be multiplyed
// OUTPUT:
// - fft method: Transforms the vector sent.
// - mul method: The result is kept in the first vector.
// NOTES:
// - You can either use the struct defined of define dificil as complex<double>
// * Codeforces Round #296 (Div. 1) D. Fuzzy Search
struct dificil {
  double real;
  double im;
  dificil()
   real=0.0:
   im=0.0;
  dificil(double real, double im):real(real),im(im){}
  dificil operator+(const dificil &o)const {
   return dificil(o.real+real, im+o.im);
  dificil operator/(double v) const {
   return dificil(real/v, im/v);
  dificil operator*(const dificil &o)const {
   return dificil(real*o.real-im*o.im, real*o.im+im*o.real);
  dificil operator-(const dificil &o) const {
    return dificil (real-o.real, im-o.im);
dificil tmp[MAXN*2];
int coco, maiorpot2[MAXN];
void fft(vector<dificil> &A, int s)
  int n = A.size(), p = 0;
  while (n>1) {
   p++;
   n >>= 1;
```

```
\mathbf{n} = (1 << \mathbf{p});
  vector<dificil> a=A;
  for(int i = 0; i < n; ++i) {
    int rev = 0;
    for (int j = 0; j < p; ++j) {
      rev <<= 1;
      rev |= ( (i >> j) & 1 );
    A[i] = a[rev];
  dificil w, wn;
  for(int i = 1; i <= p; ++i) {
    int M = 1 << i;</pre>
    int K = M \gg 1;
     wn = dificil(cos(s*2.0*pi/(double)M), sin(s*2.0*pi/(double)M);
    for (int j = 0; j < n; j += M) {
       w = dificil(1.0, 0.0);
      for(int 1 = j; 1 < K + j; ++1){
        dificil t = w;
        t = t * A[1 + K];
        dificil u = A[1];
        A[1] = A[1] + t;
        u = u-t;
A[1 + K] = u;
        w = wn *w;
  if(s==-1){
    for (int i = 0; i < n; ++i)
      A[i] = A[i] / (double) n;
void mul(vector<dificil> &a, vector<dificil> &b)
  for(int i=0;i<a.size();i++)</pre>
    a[i]=a[i]*b[i];
```

4.2 Fast Fourier Transform 2

```
// FFT - The Recursive Version
// Running Time:
    O(n*log n)
// How To Use:
    fft(&a[0], tam, 0)
// fft(&b[0], tam, 0)
    mul(a,b)
    fft(&a[0], tam, 1)
// INPUT:
    - fft method:
      * The vector representing the polynomial
       * 0 to normal transform
       * 1 to inverse transform
// - mul method:
       * The two polynomials to be multiplyed
// OUTPUT:
// - fft method: Transforms the vector sent.
   - mul method: The result is kept in the first vector.
// NOTES:
// - Tam has to be a power of 2.
// - You can either use the struct defined of define dificil as complex<double>
// SOLVED:
// * Codeforces Round #296 (Div. 1) D. Fuzzy Search
dificil tmp[MAXN*2];
int coco, maiorpot2[MAXN];
void fft(dificil *v, int N, bool inv)
  if(N<=1) return;</pre>
  dificil *vodd = v;
```

```
dificil \starveven = v+N/2;
  for(int i=0; i<N; i++) tmp[i] = v[i];</pre>
  for(int i=0; i<N; i+=2)</pre>
    veven[coco] = tmp[i];
    vodd[coco] = tmp[i+1];
  fft(&vodd[0], N/2, inv);
  fft(&veven[0], N/2, inv);
  dificil w(1);
  double angucomleite = 2.0*PI/(double)N;
  if(inv) angucomleite = -angucomleite;
  dificil wn(cos(angucomleite), sin(angucomleite));
  for(int i=0; i<N/2; i++)
    tmp[i] = veven[i]+w*vodd[i];
    tmp[i+N/2] = veven[i]-w*vodd[i];
    if(inv)
      tmp[i] /= 2:
     tmp[i+N/2] /= 2;
  for(int i=0; i<N; i++) v[i] = tmp[i];</pre>
void mul(vector<dificil> &a, vector<dificil> &b)
  for(int i=0; i<a.size(); i++)</pre>
    a[i] = a[i] *b[i];
void precomp()
  int pot=0:
  for (int i=1; i < MAXN; i++)</pre>
   if((1<<pot)<i) pot++;</pre>
    maiorpot2[i] = (1<<pot);</pre>
```

4.3 Simpson Algorithm

```
const int NPASSOS = 100000;
const int W=1000000;
//W= tamanho do intervalo que eu estou integrando
double integral1()
{
   double h = W / (NPASSOS);
   double a = 0;
   double b = W;
   double s = f(a) + f(b);
   for (double i = 1; i <= NPASSOS; i += 2) s += f(a + i * h) * 4.0;
   for (double i = 2; i <= (NPASSOS - 1); i += 2) s += f(a + i * h) * 2.0;
   return s * h / 3.0;
}</pre>
```

4.4 Matrix Exponentiation

```
//matmul multiplica m1 por m2
//matpow exponencia a matrix m1 por p
//mul vet multiplica a matrix m1 pelo vetor vet
vvi matmul(vvi &m1, vvi &m2)
{
    vvi ans;
    ans.resize(m1.size(), vi(m2.size(), 0));
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
        for (int k = 0; k < n; k++) {
            ans[i][j] += m1[i][k] * m2[k][j];
            ans[i][j] %= MOD;
    }
}
return ans;
}</pre>
```

```
vvi matpow(vvi &ml, ll p)
{
    vvi ans;
    ans.resize(ml.size(), vi(ml.size(), 0));
    for (int i = 0; i < n; i++) ans[i][i] = 1;
    while (p) {
        if (p & 1) ans = matmul(ans, ml);
        ml = matmul(ml, ml);
        p >>= 1;
    }
    return ans;
}
// VETOR TEM N LINHAS E A MATRIZ E QUADRADA
vi mulvet(vvi &ml, vi &vet)
{
    vi ans;
    ans.resize(vet.size(), 0);
    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++) {
            ans[i] += (ml[i][j] * vet[j]);
            ans[i] *= MOD;
    }
    return ans;
}</pre>
```

5 Mathematics

5.1 Big Number

```
void zero_esq(string &resp)
  string retorno = resp;
   reverse(retorno.begin(), retorno.end());
  int i = resp.size() - 1;
  while (retorno[i] == '0' and i > 0) {
    retorno.erase(i);
    i--:
  reverse(retorno.begin(), retorno.end());
  resp = retorno:
string sum_big(string a, string b)
  reverse(a.begin(), a.end());
   reverse(b.begin(), b.end());
  if (a.size() <= b.size()) {
    int carry = 0;
for (int i = 0; i < a.size(); i++) {
   int x = b[i] - '0' + a[i] - '0' + carry;
   resp.push_back((char)(x % 10 + '0'));</pre>
       carry = x / 10;
    for (int i = a.size(); i < b.size(); i++) {</pre>
      int x = b[i] - '0' + carry;
       resp.push_back((char)(x % 10 + '0'));
       carry = x / 10;
    if (carry > 0) resp.push_back((char)(carry + '0'));
    int carry = 0;
    for (int i = 0; i < b.size(); i++) {
  int x = a[i] - '0' + b[i] - '0' + carry;
  resp.push_back((char)(x % 10 + '0'));</pre>
       carry = x / 10;
     for (int i = b.size(); i < a.size(); i++) {</pre>
       int x = a[i] - '0' + carry;
       resp.push_back((char)(x % 10 + '0'));
       carry = x / 10;
    if (carry > 0) resp.push_back((char)(carry + '0'));
  reverse(resp.begin(), resp.end());
   zero_esq(resp);
  return resp;
string mul_big(string a, string b)
  string resp;
  resp.push_back('0');
  string temp;
  int carry = 0;
  reverse(a.begin(), a.end());
```

```
reverse(b.begin(), b.end());
for (int i = 0; i < a.size(); i++) {
    temp.clear();
    for (int k = 0; k < i; k++) temp.push_back('0');
    int x = a[i] - '0';
    for (int j = 0; j < b.size(); j++) {
        int y = b[j] - '0';
        int novo = (x * y * carry);
        temp.push_back((novo % 10) + '0');
        carry = novo / 10;
    }
    if (carry > 0) temp.push_back(carry + '0');
    reverse(temp.begin(), temp.end());
    carry = 0;
    resp = sum_big(temp, resp);
}
zero_esq(resp);
return resp;
```

5.2 Big Number 2

```
Structure implementing long arithmetic in C++
  Analogue to BigInteger in Java.
  Tested on many problems.
  TODO: list some problems
struct BigInt {
  vector <int> num;
  static const int base = 1000 * 1000 * 1000;
  static const int baseDigits = 9:
  string leadingZerosModifier;
  * CONSTRUCTONS & SETTERS
  ****************
  void setLeadingZerosModifier() {
   leadingZerosModifier = "%0xd";
   leadingZerosModifier[2] = '0' + baseDigits;
  void set(int value) {
   num.clear();
   if (value == 0)
     num.push_back(0);
   while (value) {
     num.push_back(value % base);
     value /= base;
  void set(long long value) {
   num.clear();
   if (value == 0)
     num.push_back(0);
   while (value) {
     num.push_back(value % base);
     value /= base;
 void set(string &value) {
   num.clear();
   for (int i = (int)value.length() - 1; i >= 0; i -= baseDigits) {
     int add = 0;
     for (int j = max(0, i - baseDigits + 1); j <= i; j++)
  add = add * 10 + value[j] - '0';</pre>
     num.push_back(add);
  void operator = (int value) {
  void operator = (long long value) {
   set (value);
  void operator = (string &value) {
   set (value);
```

```
BigInt() {
  setLeadingZerosModifier();
  set (0);
BigInt(int value) {
  setLeadingZerosModifier();
  set (value);
BigInt(string value) {
  setLeadingZerosModifier();
  set (value);
* SIZE METHODS
//returns size of vector
int size() {
  return (int) num.size();
//returns length of the number
int digitNum() {
  int result = 0;
  for (int i = 0; i < (int)num.size() - 1; i++)</pre>
    result += baseDigits;
  int lastNum = num.back();
  while (lastNum) {
    lastNum /= 10;
  return result;
* T/O
void read() {
  string s;
  cin >> s;
  num.clear();
  for (int i = (int)s.length() - 1; i >= 0; i -= baseDigits) {
    int add = 0;
    for (int j = max(0, i - baseDigits + 1); j <= i; j++)
      add = add * 10 + s[j] - '0';
    num.push_back(add);
void print() {
  printf("%d", num.back());
  for (int i = (int) num.size() - 2; i >= 0; i--)
   printf (leadingZerosModifier.c_str(), num[i]);
void println() {
  print();
  printf("\n");
string toString() {
  string result = "";
 for (int i = 0; i < (int) num.size(); i++) {
  int cur = num[i];
  for (int j = 1; j <= baseDigits; j++) {
    if (cur == 0 && i == (int) num.size() - 1)</pre>
        break;
      result.append(1, (char) '0' + cur % 10);
      cur /= 10;
  reverse(result.begin(), result.end());
  return result;
* ADDITION
void sumThis(BigInt number) {
  int carry = 0;
  for (int i = 0; i < max((int)num.size(), number.size()) || carry; i++) {</pre>
   if (i == num.size())
      num.push_back(0);
    if (i >= number.size())
      carry = num[i] + carry;
    else
```

```
carry = num[i] + carry + number.num[i];
    num[i] = carry % base;
    carry /= base;
void sumThis(int number) {
  int carry = number;
  for (int i = 0; i < (int)num.size() || carry; i++) {</pre>
    if (i == num.size())
     num.push_back(0);
   carry = num[i] + carry;
num[i] = carry % base;
    carry /= base;
BigInt sum(BigInt number) {
  BigInt result = *this;
  result.sumThis(number);
  return result;
BigInt sum(int number) {
 BigInt result = *this:
  result.sumThis(number):
 return result:
void operator += (BigInt number) {
 sumThis(number);
void operator += (int number) {
  sumThis(number);
BigInt operator + (BigInt number) {
 return sum (number);
BigInt operator + (int number) {
 return sum (number);
* SUBTRACTION
void subThis(BigInt number) {
  int carry = 0;
  for (int i = 0; i < (int)number.size() || carry; i++) {</pre>
    if (i < (int)number.size())</pre>
     num[i] -= carry + number.num[i];
    else
     num[i] -= carry:
    if (num[i] < 0) {
     carry = 1:
     num[i] += base;
    else
      carry = 0;
  while (num.size() > 1 && num.back() == 0)
    num.pop_back();
void subThis(int number) {
 int carry = -number;
for (int i = 0; carry > 0; i++) {
   num[i] -= carry;
    if (num[i] < 0) {
     carry = 1;
      num[i] += base;
    else
     carry = 0;
  while (num.size() > 1 && num.back() == 0)
    num.pop_back();
BigInt sub(BigInt number) {
 BigInt result = *this:
  result.subThis(number);
  return result:
BigInt sub(int number) {
 BigInt result = *this;
  result.subThis(number);
```

return result;

```
void operator -= (BigInt number) {
    subThis(number);
 void operator -= (int number) {
    subThis(number);
BigInt operator - (BigInt number) {
    return sub(number);
BigInt operator - (int number) {
    return sub (number);
 * MULTIPLICATION
BigInt mult(BigInt number) {
    BigInt product;
     product.num.resize(num.size() + number.size());
    for (int i = 0; i < (int) num.size(); i++) {</pre>
         fr (int i = 0; i \ (int); i 
                         0) + carry;
             product.num[i + j] = int (cur % base);
carry = int (cur / base);
     while (product.size() > 1 && product.num.back() == 0)
         product.num.pop_back();
    return product;
void multThis(BigInt number) {
    *this = mult(number);
void multThis(int number) {
   int carry = 0;
for (int i = 0; i < (int) num.size() || carry; ++i) {</pre>
         if (i == num.size())
             num.push_back (0);
         long long cur = carry + num[i] * 111 * number;
num[i] = int (cur % base);
         carry = int (cur / base);
    while (num.size() > 1 && num.back() == 0)
         num.pop_back();
BigInt mult(int number) {
    BigInt result = *this;
    result.multThis(number);
    return result:
 void operator *= (BigInt number) {
    multThis(number);
void operator *= (int number) {
    multThis(number);
BigInt operator * (BigInt number) {
    return mult (number);
BigInt operator * (int number) {
    return mult(number);
 void multThisByPowerOfTen(int power) {
    int baseNums = power / baseDigits;
    int curLen = (int)num.size();
    num.resize(curLen + baseNums);
    for (int i = num.size() - 1; i >= baseNums; i--) {
  num[i] = num[i - baseNums];
    for (int i = baseNums - 1; i >= 0; i--)
       num[i] = 0;
     power %= baseDigits;
    int multBy = (int)pow(10.0, power);
    multThis (multBy);
```

```
void divThis(int number) {
  for (int i= (int) num.size() - 1; i >= 0; i--) {
   long long cur = num[i] + carry * 111 * base;
    num[i] = int (cur / number);
    carry = int (cur % number);
  while (num.size() > 1 && num.back() == 0)
    num.pop_back();
BigInt div(int number) {
  BigInt result = *this;
  result.divThis(number);
  return result;
void operator /= (int number) {
  divThis(number);
BigInt operator / (int number) {
 return div(number);
void divThisByPowerOfTen(int power) {
 int baseNums = power / baseDigits;
  int curLen = (int)num.size();
 for (int i = 0; i < (int)num.size() - baseNums; i++) {
  num[i] = num[i + baseNums];</pre>
  for (int i = 1; i <= baseNums; i++)</pre>
   num.pop_back();
  power %= baseDigits;
  int divBy = (int)pow(10.0, power);
  divThis(divBy);
* MODIILO
void modThis(int number) {
  int carry = 0;
  for (int i= (int) num.size() - 1; i >= 0; i--) {
    long long cur = num[i] + carry * 111 * base;
    num[i] = int (cur / number);
    carry = int (cur % number);
  set (carry);
BigInt mod(int number) {
 BigInt result = *this:
  result.modThis(number);
  return result:
void operator %= (int number) {
BigInt operator % (int number) {
  return mod(number);
* COMPARTSON
//Returns: -1 - this number is less than argument, 0 - equal, 1 - this number is greater
int compareTo(BigInt number) {
  if ((int)num.size() < number.size())</pre>
   return -1;
  if ((int)num.size() > number.size())
  for (int i = (int) num.size() - 1; i >= 0; i--) {
   if (num[i] > number.num[i])
   return 1;
if (num[i] < number.num[i])</pre>
     return -1:
  return 0:
//Returns: -1 - this number is less than argument, 0 - equal, 1 - this number is greater
int compareTo(int number) {
  if (num.size() > 1 || num[0] > number)
    return 1;
```

* DIVISION

```
if (num[0] < number)</pre>
     return -1;
    return 0;
  bool operator < (BigInt number) {</pre>
    return compareTo(number) == -1;
  bool operator < (int number) {</pre>
    return compareTo(number) == -1;
  bool operator <= (BigInt number) {</pre>
    return compareTo(number) != 1;
  bool operator <= (int number)</pre>
    return compareTo(number) != 1;
  bool operator == (BigInt number) {
    return compareTo(number) == 0;
  bool operator == (int number) {
    return compareTo(number) == 0;
  bool operator > (BigInt number) {
    return compareTo(number) == 1;
  bool operator > (int number) {
    return compareTo(number) == 1;
  bool operator >= (BigInt number) {
    return compareTo(number) != -1;
  bool operator >= (int number) {
    return compareTo(number) != 1;
  bool operator != (BigInt number) {
    return compareTo(number) != 0;
 bool operator != (int number)
    return compareTo(number) != 0;
};
```

5.3 Chinese Remainder

```
11 mulmod(11 a, 11 b, 11 m)
  while (b > 0) {
   if (b % 2 != 0) ret = (ret + a) % m;
    a = (a + a) % m;
    b >>= 1;
  return ret;
ll expmod(ll a, ll e, ll m)
  11 ret = 1:
  while (e > 0) {
    if (e % 2 != 0) ret = mulmod(ret, a, m);
    a = mulmod(a, a, m);
  return ret;
11 invmul(11 a, 11 m) { return expmod(a, m - 2, m); }
11 chinese (vector<11> r, vector<11> m)
  int sz = m.size();
 11 M = 1;
for (int i = 0; i < sz; i++) {</pre>
    M *= m[i];
  ll ret = 0;
  for (int i = 0; i < sz; i++) {
```

```
ret += mulmod(mulmod(M / m[i], r[i], M), invmul(M / m[i], M), M);
ret = ret % M;
}
return ret;
```

5.4 Chinese Remainder 2

```
// Chinese remainder theorem (special case): find z such that // z % m1 = r1, z
// % m2 = r2. Here, z is unique modulo M = lcm(m1, m2). // Return (z, M). On
// failure, M = -1;
ii chinese_remainder_theorem(int m1, int r1, int m2, int r2)
  int s, t;
 int g = extended_euclid(m1, m2, s, t);
 if (r1 % g != r2 % g) return mp(0, -1);
 return mp(mod(s * r2 * m1 + t * r1 * m2, m1 * m2) / g, m1 * m2 / g);
// Chinese remainder theorem: find z such that // z % m[i] =
// r[i] for all i
   .Note that the solution is unique modulo M = lcm_i (m[i]).
     Return(z, M)
     .On // failure, M = -1. Note that we do not require the a[i] s
// to be relatively prime.
ii chinese_remainder_theorem(const vi &m, const vi &r)
  ii ret = make_pair(r[0], m[0]);
 for (int i = 1; i < m.size(); i++) {
  ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);</pre>
   if (ret.second == -1) break;
 return ret:
```

5.5 Matrix Exponentiation

```
//matmul multiplica m1 por m2
//matpow exponencia a matrix m1 por p
//mul vet multiplica a matrix m1 pelo vetor vet
vvi matmul(vvi &m1, vvi &m2)
  ans.resize(m1.size(), vi(m2.size(), 0));
 for (int i = 0; i < n; i++)

for (int j = 0; j < n; j++)

for (int k = 0; k < n; k++) {

   ans[i][j] += ml[i][k] * m2[k][j];
         ans[i][j] %= MOD;
  return ans;
vvi matpow(vvi &m1, ll p)
  ans.resize(m1.size(), vi(m1.size(), 0));
  for (int i = 0; i < n; i++) ans[i][i] = 1;</pre>
  while (p) {
    if (p & 1) ans = matmul(ans, m1);
    m1 = matmul(m1, m1);
    p >>= 1;
  return ans;
.
// VETOR TEM N LINHAS E A MATRIZ E QUADRADA
vi mulvet(vvi &m1, vi &vet)
  vi ans;
  ans.resize(vet.size(), 0);
  for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)</pre>
      ans[i] += (m1[i][j] * vet[j]);
      ans[i] %= MOD;
  return ans:
```

5.6 Pascal Triangle

```
//Fazer combinacao de N escolhe M
//por meio do triangulo de pascal
//complexidade: O(m*n)
unsigned long long comb[61][61];
for (int i = 0; i < 61; i++) {
    comb[i][i] = 1;
    comb[i][0] = 1;
}
for (int i = 2; i < 61; i++)
    for (int j = 1; j < i; j++)
    c o mb [i][j] = comb[i - 1][j] + comb[i - 1][j - 1];</pre>
```

5.7 Euler's Totient Function

```
//retorna quantos elementos coprimos
//a N e menores que n existem
int phi (int n)
{
  int result = n;
  for (int i = 2; i * i <= n; ++i)
    if (n % i == 0) {
      while (n % i == 0) n /= i;
      result -= result / i;
    }
  if (n > 1) result -= result / n;
  return result;
```

5.8 Pollard Rho

```
11 u;
11 t:
const int tamteste=5;
11 abss(11 v) { return v>=0 ? v : -v;}
11 randerson()
  ld pseudo=(ld)rand()/(ld)RAND_MAX;
 return (11) (round((1d) range*pseudo))+1LL;
11 mulmod(11 a, 11 b, 11 mod)
  11 ret=0;
  while (b>0)
    if(b%2!=0) ret=(ret+a)%mod;
    a=(a+a) %mod;
    b=b/2LL;
  return ret:
11 expmod(11 a, 11 e, 11 mod)
  11 ret=1;
  while (e>0)
    if(e%2!=0) ret=mulmod(ret,a,mod);
    a=mulmod(a,a,mod);
    e=e/2LL;
  return ret:
bool jeova(ll a, ll n)
  11 \times = expmod(a, u, n);
  ll last=x;
  for(int i=0;i<t;i++)</pre>
    x=mulmod(x,x,n);
    if(x==1 and last!=1 and last!=(n-1)) return true;
  if(x==1) return false;
 return true:
bool isprime(11 n)
  u=n-1;
  t=0;
  while (u%2==0)
```

```
u/=2LL;
  if(n==2) return true;
  if(n==3) return true;
  if(n%2==0) return false;
  if(n<2) return false;</pre>
  for(int i=0;i<tamteste;i++)</pre>
   11 v = randerson()%(n-2)+1;
//cout<<"jeova "<<v<<" "<<n<<end1;</pre>
    if(jeova(v,n)) return false;
  return true:
11 gcd(ll a, ll b) { return !b ? a : gcd(b,a%b);}
ll calc(ll x, ll n, ll c)
  return (mulmod(x,x,n)+c)%n;
ll pollard(ll n)
  11 d=1;
  11 i=1;
 11 k=1:
  11 x=2:
 11 y=x;
  11 c;
    c=randerson()%n;
  } while (c==0 or (c+2) %n==0);
  while (d!=n)
    if(i==k)
        k \star = 2LL;
        y=x;
i=0;
    x=calc(x,n,c);
    i++;
    d=gcd(abss(y-x),n);
    if(d!=1) return d;
vector<ll> getdiv(ll n)
  vector<ll> ret:
  if (n==1) return ret:
  if(isprime(n))
    ret.pb(n):
    return ret:
  11 d = pollard(n);
  ret=getdiv(d);
  vector<ll> ret2=getdiv(n/d);
  for(int i=0;i<ret2.size();i++) ret.pb(ret2[i]);</pre>
```

5.9 Sieve of Eratosthenes

```
//esse crivo gera MAXN primos
const int MAX = 1e6;
int primes(MAX);
void gen_primes()
{
  int i, j;
  for (i = 2; i*i <= MAX; i++)
    if (primes[i])
    for (j = i; j * i < MAX; j++) primes[i * j] = 0;
}</pre>
```

5.10 Extended Euclidean Algorithm

```
//returns g = gcd(a, b);
//finds x,y such that d= ax+by;
```

```
int extended_euclid(int a, int b, int &x, int &y)
{
  int xx = y = 0;
  int yy = x = 1;
  while (b) {
    int q = a / b;
    int t = b;
    b = a % b;
    a = t;
    t = xx;
    xx = x - q * xx;
    x = t;
    t = yy;
    yy = y - q * yy;
    y = t;
}
return a;
```

5.11 Multiplicative Inverse

```
//computes b such that ab = 1 (mod n), returns - 1 on failure
int mod_inverse(int a, int n)
{
   int x, y;
   int g = extended_euclid(a, n, x, y);
   if (g > 1) return -1;
   return (x+n)%n;
}
```

5.12 Multiplicative Inverse 2

```
//inverso multiplicativo de A % MOD
//certifique de MOD estar definido antes bonito!
//complexidade: O(log(a))
ll mul_inv(ll a)
{
    ll pin0 = MOD, pin = MOD, t, q;
    ll x0 = 0, x1 = 1;
    if (pin = 1) return 1;
    while (a > 1) {
        q = a / pin;
        t = pin, pin = a % pin, a = t;
        t = x0, x0 = x1 - q * x0, x1 = t;
    }
    if (x1 < 0) x1 += pin0;
    return x1;
}
```

6 Combinatorial Optimization

6.1 Dinic

```
//grafo bipartido O(Esqrt(v))
//Para recuperar a resposta, e so colocar um bool
//de false na aresta de retorno e fazer uma bfs/dfs
//andando pelos vertices de capacidade =0 e arestas
//que nao sao de retorno
struct Edge {
  int v, rev;
  Edge(int v_, int cap_, int rev_) : v(v_), rev(rev_), cap(cap_) {}
struct MaxFlow {
  vector<vector<Edge> > g;
  vector<int> level;
   queue<int> q;
  int flow, n;
  \begin{array}{l} \texttt{MaxFlow(int } n\_) \ : \ g(n\_) \ , \ level(n\_) \ , \ n(n\_) \ \{\} \\ \textbf{void } \ addEdge(\textbf{int } u, \ \textbf{int } v, \ \textbf{int } cap) \end{array}
     if (u == v) return;
     Edge e(v, cap, int(g[v].size()));
     Edge r(u, 0, int(g[u].size()));
```

```
g[u].push_back(e);
    g[v].push_back(r);
  bool buildLevelGraph(int src, int sink)
    fill(level.begin(), level.end(), -1);
    while (not q.empty()) q.pop();
    level[src] = 0;
     q.push(src);
    while (not q.empty()) {
      int u = q.front();
       g.pop();
      for (auto e = g[u].begin(); e != g[u].end(); ++e) {
  if (not e->cap or level[e->v] != -1) continue;
  level[e->v] = level[u] + 1;
         if (e->v == sink) return true;
         q.push(e->v);
    return false;
  int blockingFlow(int u, int sink, int f)
    if (u == sink or not f) return f;
    int fu = f:
    for (auto e = g[u].begin(); e != g[u].end(); ++e) {
   if (not e->cap or level[e->v] != level[u] + 1) continue;
      int mincap = blockingFlow(e->v, sink, min(fu, e->cap));
      if (mincap) {
         g[e->v][e->rev].cap += mincap;
          e->cap -= mincap;
         fu -= mincap;
    if (f == fu) level[u] = -1;
    return f - fu;
  int maxFlow(int src, int sink)
    while (buildLevelGraph(src, sink))
      flow += blockingFlow(src, sink, numeric_limits<int>::max());
    return flow;
};
```

6.2 Hopcroft-Karp Bipartite Matching

```
\star Matching maximo de grafo bipartido de peso 1 nas arestas
\star supondo que o grafo bipartido seja enumerado de 0-n-1
* chamamos maxMatch(n)
class MaxMatch {
  vi graph[N];
  int match[N], us[N];
  void addEdge(int u, int v) { graph[u].pb(v); }
  int dfs(int u)
    if (us[u]) return 0;
    us[u] = 1;
    for (int v : graph[u]) {
   if (match[v] == -1 or (dfs(match[v]))) {
        match[v] = u;
        return 1;
    return 0;
  int maxMatch(int n)
    memset(match, -1, sizeof(match));
    int ret = 0;
    for (int i = 0; i < n; i++) {</pre>
      memset(us, 0, sizeof(us));
      ret += dfs(i);
    return ret:
1:
```

6.3 Max Bipartite Matching 2

```
// This code performs maximum bipartite matching.
// Running time: O(|E| |V|) -- often much faster in practice
     INPUT: w[i][j] = edge between row node i and column node j
     OUTPUT: mr[i] = assignment for row node i, -1 if unassigned
             mc[j] = assignment for column node j, -1 if unassigned
              function returns number of matches made
#include <vector>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool FindMatch(int i, const VVI &w, VI &mr, VI &mc, VI &seen) {
  for (int j = 0; j < w[i].size(); j++) {
  if (w[i][j] && !seen[j]) {</pre>
      seen[j] = true;
if (mc[j] < 0 || FindMatch(mc[j], w, mr, mc, seen)) {</pre>
       mr[i] = j;
mc[i] = i;
        return true;
  return false;
int BipartiteMatching(const VVI &w, VI &mr, VI &mc) {
  mr = VI(w.size(), -1);
  mc = VI(w[0].size(), -1);
  int ct = 0;
  for (int i = 0; i < w.size(); i++) {</pre>
    VI seen(w[0].size());
    if (FindMatch(i, w, mr, mc, seen)) ct++;
 return ct:
```

6.4 Maximum Matching in General Graphs (Blossom)

```
GETS:
V->number of vertices
E{\operatorname{\operatorname{->}}} number \ of \ edges
pair of vertices as edges (vertices are 1..V)
output of edmonds() is the maximum matching
match[i] is matched pair of i (-1 if there isn't a matched pair)
Code for the SEAGRP problem at CodeChef.
SEAGRP's limits are: 1 <= V, E <= 100.
The problem asked if there is a perfect matching.
#include <bits/stdc++.h>
using namespace std;
struct struct_edge { int v; struct_edge* n; };
typedef struct_edge* edge;
struct_edge pool[M*M*2];
edge adj[M];
int V,E,match[M],qh,qt,q[M],father[M],base[M];
bool inq[M],inb[M],ed[M][M];
void clean()
 memset(ed, false, sizeof(ed));
 topindex=0;
  for (int i = 0; i < M; i++)
    adj[i] = NULL;
void add_edge(int u,int v)
  edge top = &pool[topindex++];
 top->v=v,top->n=adj[u],adj[u]=top;
top = &pool[topindex++];
  top->v=u,top->n=adj[v],adj[v]=top;
int LCA(int root, int u, int v)
```

```
static bool inp[M];
  memset(inp,0,sizeof(inp));
    inp[u=base[u]]=true;
    if (u==root) break;
    u=father[match[u]];
  while (1)
    if (inp[v=base[v]]) return v;
    else v=father[match[v]];
void mark blossom(int lca,int u)
  while (base[u]!=lca)
    int v=match[u];
    inb[base[u]]=inb[base[v]]=true;
    u=father[v];
    if (base[u]!=lca) father[u]=v;
void blossom_contraction(int s,int u,int v)
  int lca=LCA(s,u,v);
  memset(inb,0,sizeof(inb));
  mark blossom(lca.u):
  mark_blossom(lca, v);
  if (base[u]!=lca)
    father[u]=v;
  if (base[v]!=lca)
    father[v]=u;
  for (int u=0; u < V; u++)
    if (inb[base[u]])
      base[u]=lca;
     if (!inq[u])
        inq[q[++qt]=u]=true;
int find_augmenting_path(int s)
  memset(inq,0,sizeof(inq));
  memset (father, -1, sizeof (father));
  for (int i=0;i<V;i++) base[i]=i;</pre>
  inq[q[qh=qt=0]=s]=true;
    int u=q[qh++];
    for (edge e=adj[u];e!=NULL;e=e->n)
      if (base[u]!=base[v]&&match[u]!=v)
        if ((v==s)||(match[v]!=-1 && father[match[v]]!=-1))
          blossom_contraction(s,u,v);
        else if (father[v]==-1)
          father[v]=u;
          if (match[v] ==-1)
            return v;
          else if (!inq[match[v]])
            inq[q[++qt]=match[v]]=true;
  return -1:
int augment path(int s,int t)
  int u=t, v, w;
  while (u!=-1)
    v=father[u];
    w=match[v];
    match[v]=u;
    match[u]=v;
    u=w;
  return t!=-1:
int edmonds()
  int matchc=0;
  memset (match, -1, sizeof (match));
  for (int u=0;u<V;u++)</pre>
    if (match[u]==-1)
      matchc+=augment_path(u, find_augmenting_path(u));
  return matchc;
```

```
int main()
  int u, v, t;
  cin >> t;
  while (t--)
    cin >> V >> E;
    clean();
    while (E--)
      cin >> u >> v;
      if (!ed[u-1][v-1])
         add_edge(u-1,v-1);
         ed[u-1][v-1]=ed[v-1][u-1]=true;
    //cout << "UE\n";
//cout << V << " " << edmonds() << endl;
    //for (int i=0;i<V;i++)
    // if (i<match[i])
// cout<<i+1<<" "<<match[i]+1<<endl;
    //cout << endl;
    if(2*edmonds() == V) cout << "YES\n";</pre>
    else cout << "NO\n";
  return 0:
```

6.5 Min Cost Matching

```
// Min cost bipartite matching via shortest augmenting paths
// This is an O(n^3) implementation of a shortest augmenting path
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
// second.
    cost[i][j] = cost for pairing left node i with right node j
    Lmate[i] = index of right node that left node i pairs with
    Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
#include <algorithm>
#include <cmath>
#include <cstdio>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate)
  int n = int(cost.size());
  // construct dual feasible solution
 VD u(n);
  VD v(n);
  for (int i = 0; i < n; i++) {
   u[i] = cost[i][0];
   for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
  for (int j = 0; j < n; j++) {
   v[j] = cost[0][j] - u[0];
   for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
  // construct primal solution satisfying complementary slackness
  Lmate = VI(n, -1);
  Rmate = VI(n, -1);
 Lmate[i] = j;
Rmate[j] = i;
```

```
mated++;
      break;
VD dist(n);
VI dad(n);
VI seen(n);
// repeat until primal solution is feasible
while (mated < n) {
  // find an unmatched left node
  int s = 0:
  while (Lmate[s] != -1) s++;
  // initialize Dijkstra
  fill(dad.begin(), dad.end(), -1);
  fill(seen.begin(), seen.end(), 0);
  for (int k = 0; k < n; k++) dist[k] = cost[s][k] - u[s] - v[k];
  int i = 0:
  while (true) {
    // find closest
      = -1:
    for (int k = 0; k < n; k++) {
      if (seen[k]) continue;
      if (j == -1 || dist[k] < dist[j]) j = k;</pre>
    seen[j] = 1;
     // termination condition
    if (Rmate[j] == -1) break;
    // relax neighbors
    const int i = Rmate[j];
    for (int k = 0; k < n; k++) {
   if (seen[k]) continue;</pre>
      const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
      if (dist[k] > new_dist) {
        dist[k] = new_dist;
        dad[k] = j;
  // update dual variables
  for (int k = 0; k < n; k++) {
   if (k == j || !seen[k]) continue;</pre>
    const int i = Rmate[k];
    v[k] += dist[k] - dist[j];
    u[i] -= dist[k] - dist[j];
  u[s] += dist[j];
  // augment along path
  while (dad[j] >= 0) {
  const int d = dad[j];
    Rmate[j] = Rmate[d];
    Lmate[Rmate[j]] = j;
    j = d;
  Rmate[j] = s;
  Lmate[s] = j;
  mated++:
double value = 0;
for (int i = 0; i < n; i++) value += cost[i][Lmate[i]];</pre>
return value;
```

6.6 Min Cost Max Flow

```
// Implementation of min cost max flow algorithm using adjacency
// matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge costs to 0.
//
// Running time, O(|V|^2) cost per augmentation
// max flow: O(|V|^3) augmentations
// min cost max flow: O(|V|^4 * MAX_EDGE_COST) augmentations
//
// INPUT:
- graph, constructed using AddEdge()
- source
```

```
- sink
// OUTPUT:
       - (maximum flow value, minimum cost value)
       - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <iostream>
#include <vector>
using namespace std;
typedef vector<VI> VVI;
typedef long long LL;
typedef vector<LL> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const LL INF = numeric_limits<LL>::max() / 4;
struct MinCostMaxFlow {
  int N;
  VVL cap, flow, cost;
  VI found;
  VL dist, pi, width;
  VPII dad:
  MinCostMaxFlow(int N): N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
      found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, LL cap, LL cost)
    this->cap[from][to] = cap;
    this->cost[from][to] = cost;
  void Relax(int s, int k, LL cap, LL cost, int dir)
    LL val = dist[s] + pi[s] - pi[k] + cost;
if (cap && val < dist[k]) {
     dist[k] = val;
dad[k] = make_pair(s, dir);
      width[k] = min(cap, width[s]);
  LL Dijkstra(int s, int t)
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
      int best = -1:
      found[s] = true;
      for (int k = 0; k < N; k++) {
        if (found[k]) continue;
        Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
        Relax(s, k, flow[k][s], -cost[k][s], -1);
        if (best == -1 || dist[k] < dist[best]) best = k;</pre>
    for (int k = 0; k < N; k++) pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<LL, LL> GetMaxFlow(int s, int t)
    LL totflow = 0, totcost = 0;
    while (LL amt = Dijkstra(s, t)) {
      totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
        if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
           totcost += amt * cost[dad[x].first][x];
          flow[x][dad[x].first] -= amt;
          totcost -= amt * cost[x][dad[x].first];
    return make_pair(totflow, totcost);
};
```

6.7 Edmonds Karp

```
struct Edge {
 int at, where;
  11 cap:
  void init(int _at, ll _cap, int _where)
    at = _at, cap = _cap, where = _where;
}:
struct dad {
 int at, up, down;
dad() { at = -1; }
  dad(int _at, int _up, int _down) { at = _at, up = _up, down = _down; }
class MaxFlow {
private:
  vector<vector<Edge> > g;
  11 mf, f;
  int s, t;
  vector<dad> p;
 public:
  void augment(int v, ll minEdge)
    if (v == s) {
     f = minEdge;
     return;
    else if (p[v].at != -1) {
     augment(p[v].at, min(minEdge, g[p[v].at][p[v].up].cap));
      g[p[v].at][p[v].up].cap -= f;
      g[v][p[v].down].cap += f;
  void init(int N)
    for (int i = 0; i < q.size(); i++) g[i].clear();</pre>
    mf = 0, f = 0;
    g.resize(N);
  void addEdge(int u, int v, 11 cap)
    A.init(v, cap, g[v].size());
    B.init(u, 0, g[u].size());
    g[u].pb(A);
    g[v].pb(B);
  int maxFlow(int source, int sink)
    s = source;
    t = sink;
    mf = 0;
    while (true) {
     f = 0;
      vector<int> dist(g.size(), INF);
      dist[s] = 0;
      queue<int> q;
      q.push(s);
      p.clear();
      p.resize(g.size());
      while (!q.empty())
       int u = q.front();
        q.pop();
        if (u == t) break;
        for (int i = 0; i < g[u].size(); i++) {
          Edge prox = g[u][i];
          if (dist[prox.at] == INF and prox.cap > 0) {
            dist[prox.at] = dist[u] + 1;
            q.push(prox.at);
            dad paizao(u, i, prox.where);
            p[prox.at] = paizao;
      augment(t, INF);
      if (f == 0) break;
     mf += f:
    return mf;
1:
```

7 Dynamic Programming

7.1 Convex Hull Trick

```
/∗ Esse convex hull trick e para achar a reta minima!
 * Para maximizar a reta dada , basta trocar o '>' para
 * para '<' na funcao query;
 * Nao chamar query com B ou A vazios! Atualizar dp para
 * depois fazer a query =)
 * ATENCAO COM O DOUBLE!! ESTA EM LONG LONG :)
vi A[N], B[N];
int pont[N];
bool odomeioehlixo(int r1, int r2, int r3, int j)
  (B[j][r1] - B[j][r2]) * (A[j][r3] - A[j][r1]);
void add(ll a, ll b, int j)
  B[j].pb(b);
  A[j].pb(a);
  while (B[j].size() >= 3 and
    odomeioehlixo(B[j].size() - 3, B[j].size() - 2, B[j].size() - 1, j)) {
B[j].erase(B[j].end() - 2);
    A[j].erase(A[j].end() - 2);
11 query(ll x, int j)
 if (pont[j] >= B[j].size()) pont[j] = B[j].size() - 1;
while (pont[j] < B[j].size() - 1 and</pre>
         (A[j][pont[j]] + 1] * x + B[j][pont[j]] + 1] > A[j][pont[j]] * x + B[j][pont[j]]))
    pont[j]++;
 return A[j][pont[j]] * x + B[j][pont[j]];
 * http://www.spoj.com/problems/APIO10A/
 * http://www.spoj.com/problems/ACQUIRE/
```

7.2 Convex Hull Trick 2

```
* Given a set of pairs (m, b) specifying lines of the form y = m * x + b, process
 * set of x-coordinate queries each asking to find the minimum y-value when any
* of
* the given lines are evaluated at the specified x. To instead have the queries
 * optimize for maximum y-value, set the QUERY_MAX flag to true.
 * The following implementation is a fully dynamic variant of the convex hull
 * optimization technique, using a self-balancing binary search tree (std::set)
 \star support the ability to call add_line() and get_best() in any desired order.
 * Explanation: http://wcipeg.com/wiki/Convex_hull_trick#Fully_dynamic_variant
 * Time Complexity: O(n log n) on the total number of calls made to add_line(),
+ for
 * any length n sequence of arbitrarily interlaced add_line() and get_min()
* calls.
 * Each individual call to add_line() is O(log n) amortized and each individual
 * call to qet_best() is O(log n), where n is the number of lines added so far.
 * Space Complexity: O(n) auxiliary on the number of calls made to add_line().
#include <limits> // std::numeric_limits
#include <set>
class hull_optimizer {
  struct line {
    long long m, b, val;
    double xlo;
    bool is_query;
    bool query_max;
    line (long long m, long long b, long long val, bool is query, bool query max)
      this->m = m;
      this->b = b;
```

```
this->val = val;
      this->xlo = -std::numeric_limits<double>::max();
      this->is_query = is_query;
      this->query_max = query_max;
    bool parallel(const line &l) const { return m == 1.m; }
    double intersect(const line &1) const
      if (parallel(1)) return std::numeric_limits<double>::max();
      return (double) (1.b - b) / (m - 1.m);
    bool operator<(const line &1) const
      if (1.is_query) return query_max ? (xlo < 1.val) : (1.val < xlo);</pre>
      return m < 1.m;
  std::set<line> hull;
  bool _query_max;
  typedef std::set<line>::iterator hulliter;
  bool has_prev(hulliter it) const { return it != hull.begin(); }
  bool has_next(hulliter it) const
    return (it != hull.end()) && (++it != hull.end());
  bool irrelevant (hulliter it) const
    if (!has_prev(it) || !has_next(it)) return false;
    hulliter prev = it, next = it;
    --prev;
    ++next;
    return _query_max ? prev->intersect(*next) <= prev->intersect(*it)
                      : next->intersect(*prev) <= next->intersect(*it);
  hulliter update left border(hulliter it)
    if ((_query_max && !has_prev(it)) || (!_query_max && !has_next(it)))
     return it;
    hulliter it2 = it.
    double val = it->intersect(_query_max ? *--it2 : *++it2);
    line l(*it);
    hull.erase(it++);
    return hull.insert(it, 1);
  hull_optimizer(bool query_max = false) { this->_query_max = query_max; }
  void add_line(long long m, long long b)
   line 1(m, b, 0, false, _query_max);
hulliter it = hull.lower_bound(1);
    if (it != hull.end() && it->parallel(1)) {
      if ((_query_max && it->b < b) || (!_query_max && b < it->b))
        hull.erase(it++);
      else
        return;
    it = hull.insert(it, 1);
    if (irrelevant(it)) {
      hull.erase(it);
      return:
    while (has_prev(it) && irrelevant(--it)) hull.erase(it++);
    while (has_next(it) && irrelevant(++it)) hull.erase(it--);
    it = update left border(it);
    if (has_prev(it)) update_left_border(--it);
    if (has_next(++it)) update_left_border(++it);
  long long get_best(long long x) const
    line q(0, 0, x, true, _query_max);
    hulliter it = hull.lower_bound(q);
    if (_query_max) --it;
   return it->m * x + it->b;
1:
/*** Example Usage ***/
#include <cassert>
```

```
hull_optimizer h;
h.add_line(3, 0);
h.add_line(0, 6);
h.add_line(1, 2);
h.add_line(2, 1);
assert(h.get_best(0) == 0);
assert(h.get_best(2) == 4);
assert(h.get_best(1) == 3);
assert(h.get_best(3) == 5);
return 0;
```

7.3 Divide and Conquer

```
//Um exemplo de Divide and conquer:
int MOD = 1e9 + 7;
const int N = 1010;
int dp[N][N], cost[N][N], v[N], pref[N], n, m;
void compDP(int j, int L, int R, int b, int e)
{
    if (L > R) return;
    int mid = (L + R) / 2;
    int idx = -1;
    for (int i = b; i <= min(mid, e); i++)
        if (dp[mid][j] > dp[i][j - 1] + cost[i + 1][mid]) {
        idx = i;
        dp[mid][j] = dp[i][j - 1] + cost[i + 1][mid];
    }
    compDP(j, L, mid - 1, b, idx);
    compDP(j, mid + 1, R, idx, e);
}
//chamada!
for (int i = 1; i <= n; i++) dp[i][0] = cost[1][i];
for (int i = 1; i <= m; i++) compDP(i, 1, n, 1, n);</pre>
```

7.4 Longest Increasing Subsequence

```
//asw -> vetor com resposta!!
//asw.size() o tamanho da maior lis
void lis( const vector< int > & v, vector< int > & asw )
  vector<int> pd(v.size(),0), pd_index(v.size()), pred(v.size());
  int maxi = 0, x=0, j=0, ind=0;
  for(int i=0;i<v.size();i++)</pre>
    j=lower_bound(pd.begin(),pd.begin()+maxi,x) -pd.begin();
   pd[j] = x;
    pd_index[j] = i;
  if(j==maxi)
    maxi++:
    ind = i;
  if(pred[i] == j) pd_index[j-1] = -1;
  int pos=maxi-1,k=v[ind];
  asw.resize( maxi );
  while ( pos >= 0 )
   asw[pos--] = k;
   ind = pred[ind];
   k = v[ind];
```

8 Geometry

8.1 Convex Hull Monotone Chain

```
typedef struct sPoint {
   int x, y;
   sPoint(int _x, int _y)
   {
        x = _x;
        y = _y;
   }
}
```

```
} point;
bool comp(point a, point b)
        if (a.x == b.x) return a.y < b.y;</pre>
        return a.x < b.x;
int cross(point a, point b, point c) // AB x BC
        a.x -= b.x;
        a.y -= b.y;
        b.x -= c.x;
        b.y -= c.y;
        return a.x * b.y - a.y * b.x;
bool isCw(point a, point b, point c) // Clockwise
        return cross(a, b, c) < 0;
// >= if you want to put collinear points on the convex hull
bool isCcw(point a, point b, point c) // Counter Clockwise
        return cross(a, b, c) > 0;
vector<point> convexHull(vector<point> p
        vector<point> u, 1; // Upper and Lower hulls
        sort(p.begin(), p.end(), comp);
        for (unsigned int i = 0; i < p.size(); i++) {</pre>
                while (1.size() > 1 && !isCcw(l[1.size() - 1], l[1.size() - 2], p[i]))
                        1.pop_back();
                l.push_back(p[i]);
        for (int i = p.size() - 1; i >= 0; i--) {
                while (u.size() > 1 && !isCcw(u[u.size() - 1], u[u.size() - 2], p[i]))
                       u.pop back();
               u.push_back(p[i]);
        u.pop_back();
        1.pop_back();
        l.insert(l.end(), u.begin(), u.end());
        return 1;
```

8.2 Minimum Enclosing Circle

```
//6.5- Minimum Enclosing Circle
const double eps = 1e-6;
#define CIRCLE circ
#define PT Ponto
#define MP 101
#define eps 1e-9
#define x first
#define y second
typedef double cood;
typedef int num;
typedef int point;
double resp;
cood x[MP], y[MP], ar, ax, ay;
int p[MP];
typedef pair<double, double> ponto;
typedef pair<double, double> Ponto;
double dista(ponto a, ponto b)
  return sqrt((a.first - b.first) * (a.first - b.first) +
              (a.second - b.second) * (a.second - b.second));
bool in(ponto a, pair<double, ponto> c)
  if (dista(a, c.second) - eps < c.first) return true;</pre>
bool same (point a, point b)
  return (fabs(x[a] - x[b]) < eps && fabs(y[a] - y[b]) < eps);</pre>
bool lexLess(point a, point b)
  if (fabs(x[a] - x[b]) < eps) return y[a] < y[b];</pre>
  return x[a] < x[b];</pre>
```

```
inline cood dist(cood xx, cood yy, point a)
  return sqrt((xx - x[a]) \star (xx - x[a]) + (yy - y[a]) \star (yy - y[a]));
inline cood cP(point a, point b, point c)
  return (x[a] - x[b]) * (y[c] - y[b]) - (x[c] - x[b]) * (y[a] - y[b]);
void findCircle(point a, point b, point c, cood& cx, cood& cy)
  cx = 0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
        (y[b] - y[c]) - 0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
            (y[a] - y[b]),
  cy = 0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
            (x[a] - x[b])
        0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
            (x[b] - x[c]);
  cx /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
  cy /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
void spanCircle2(int k, point p0, point p1, cood& cx, cood& cy, cood& r)
  cx = 0.5 * (x[p0] + x[p1]);
  cy = 0.5 * (y[p0] + y[p1]);
  r = dist(cx, cy, p0);
  for (int i = 0; i < k; i++)
   if (dist(cx, cy, p[i]) > r) {
  findCircle(p0, p1, p[i], cx, cy);
      r = dist(cx, cy, p[i]);
void spanCircle1(int k, point p0, cood& cx, cood& cy, cood& r)
 cx = 0.5 * (x[p0] + x[p[0]]);

cy = 0.5 * (y[p0] + y[p[0]]);
  r = dist(cx, cy, p0);
for (int i = 0; i < k; i++)
    if (dist(cx, cy, p[i]) > r) spanCircle2(i, p0, p[i], cx, cy, r);
void spanCircle(int n, cood& cx, cood& cy, cood& r)
  // Bem importante, retirar repetidos
  sort(p, p + 1, lexLess);
  n = unique(p, p + n) - p;
  random\_shuffle(p, p + n);
  if (n > 1) {
    cx = 0.5 * (x[p[0]] + x[p[1]]);
    cy = 0.5 * (y[p[0]] + y[p[1]]);

r = dist(cx, cy, p[1]);

for (int i = 2; i < n; i++)
      if (dist(cx, cy, p[i]) > r) spanCircle1(i, p[i], cx, cy, r);
  else {
    cx = x[0];
    cy = y[0];
    r = 0.0;
void solve(vector<pair<double, double> >& v)
  int N = v.size();
  for (int i = 0; i < N; i++) {
    x[i] = v[i].first;
    y[i] = v[i].second;</pre>
    p[i] = i;
  spanCircle(N, ax, ay, ar);
```

8.3 Minimum Enclosing Circle 2

```
const double eps = 1e-6;
#define CIRCLE circ
#define PT Ponto
#define MP 101
#define mps 1e-9
#define x first
#define y second
typedef double cood;
typedef int num;
```

```
typedef int point;
double resp;
cood x[MP], y[MP], ar, ax, ay;
int p[MP];
typedef pair<double, double> ponto;
typedef pair<double, double> Ponto;
double dista(ponto a, ponto b)
  return sqrt((a.first - b.first) * (a.first - b.first) +
               (a.second - b.second) * (a.second - b.second));
bool in(ponto a, pair<double, ponto> c)
  if (dista(a, c.second) - eps < c.first) return true;</pre>
  return false:
bool same (point a, point b)
  return (fabs(x[a] - x[b]) < eps && fabs(y[a] - y[b]) < eps);</pre>
bool lexLess(point a, point b)
  if (fabs(x[a] - x[b]) < eps) return y[a] < y[b];
  return x[a] < x[b];</pre>
inline cood dist(cood xx, cood yy, point a)
  return sqrt((xx - x[a]) * (xx - x[a]) + (yy - y[a]) * (yy - y[a]));
inline cood cP(point a, point b, point c)
  return (x[a] - x[b]) * (y[c] - y[b]) - (x[c] - x[b]) * (y[a] - y[b]);
void findCircle(point a, point b, point c, cood& cx, cood& cy)
  cx = 0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (v[b] - v[c])
       0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
           (v[a] - v[b]),
  cy = 0.5 * (x[b] * x[b] + y[b] * y[b] - x[c] * x[c] - y[c] * y[c]) *
           (x[a] - x[b])
       0.5 * (x[a] * x[a] + y[a] * y[a] - x[b] * x[b] - y[b] * y[b]) *
           (x[b] - x[c]);
 cx /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);

cy /= (x[a] - x[b]) * (y[b] - y[c]) - (x[b] - x[c]) * (y[a] - y[b]);
void spanCircle2(int k, point p0, point p1, cood& cx, cood& cy, cood& r)
  cx = 0.5 * (x[p0] + x[p1]);
  cy = 0.5 * (y[p0] + y[p1]);
  r = dist(cx, cy, p0);
for (int i = 0; i < k; i++)
    if (dist(cx, cy, p[i]) > r) {
      findCircle(p0, p1, p[i], cx, cy);
      r = dist(cx, cy, p[i]);
void spanCircle1(int k, point p0, cood& cx, cood& cy, cood& r)
  cx = 0.5 * (x[p0] + x[p[0]]);
 cy = 0.5 * (y[p0] + y[p[0]]);
  r = dist(cx, cy, p0);
for (int i = 0; i < k; i++)</pre>
    if (dist(cx, cy, p[i]) > r) spanCircle2(i, p0, p[i], cx, cy, r);
void spanCircle(int n, cood& cx, cood& cy, cood& r)
  // Bem importante, retirar repetidos
  sort(p, p + 1, lexLess);
  n = unique(p, p + n) - p;
  random_shuffle(p, p + n);
  if (n > 1) {
    cx = 0.5 * (x[p[0]] + x[p[1]]);
    cy = 0.5 * (y[p[0]] + y[p[1]]);
    r = dist(cx, cy, p[1]);
for (int i = 2; i < n; i++)
      if (dist(cx, cy, p[i]) > r) spanCircle1(i, p[i], cx, cy, r);
  else {
    cx = x[0];
    cy = y[0];
    r = 0.0;
void solve(vector<pair<double, double> >& v)
```

```
{
  int N = v.size();
  for (int i = 0; i < N; i++) {
    x[i] = v[i].first;
    y[i] = v[i].second;
    p[i] = i;
}
spanCircle(N, ax, ay, ar);</pre>
```

8.4 Fast Geometry in Cpp

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100:
double EPS = 1e-12;
struct PT {
  double x, y;
  PT (double x, double y) : x(x), y(y) {}
  PT (const PT &p) : x(p.x), y(p.y)
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y);
                               const { return PT(x*c, y*c );
  PT operator * (double c)
 PT operator / (double c)
                               const { return PT(x/c, y/c ); ]
double dot(PT p, PT q)
                           { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q)
                           { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {</pre>
 os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p) { return PT(-p.y,p.x); }
                       { return PT(p.y,-p.x); }
PT RotateCW90 (PT p)
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
 return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
  double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                          double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;</pre>
bool LinesCollinear(PT a, PT b, PT c, PT d) {
 return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
```

```
line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
       dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
     \textbf{if} \ ( \texttt{dot} \, ( \texttt{c-a}, \ \texttt{c-b} ) \ > \ 0 \ \&\& \ \texttt{dot} \, ( \texttt{d-a}, \ \texttt{d-b} ) \ > \ 0 \ \&\& \ \texttt{dot} \, ( \texttt{c-b}, \ \texttt{d-b} ) \ > \ 0 ) 
      return false;
    return true:
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true:
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
\ensuremath{//} strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// \ {\it integer \ arithmetic \ by \ taking \ care \ of \ the \ division \ appropriately}
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0:
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) %p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||
   p[j].y <= q.y && q.y < p[i].y) &&</pre>
       q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
       c = !c;
  return c;
// determine if point is on the boundary of a polygon \,
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    \label{eq:final_problem} \textbf{if} \ (\texttt{dist2}(\texttt{ProjectPointSegment}(\texttt{p[i]},\ \texttt{p[(i+1)\$p.size()]},\ \texttt{q)},\ \texttt{q}) \ < \ \texttt{EPS)}
      return true:
    return false:
// compute intersection of line through points a and b with
// circle centered at c with radius r >
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
  b = b-a;
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R || d+min(r, R) < max(r, R)) return ret;
double x = (d*d-R*R+r*r)/(2*d);</pre>
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret:
```

// determine if line segment from a to b intersects with

```
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
 return fabs (ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p)
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale:
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p)
  for (int i = 0; i < p.size(); i++)
    for (int k = i+1; k < p.size(); k++) {</pre>
     int j = (i+1) % p.size();
int l = (k+1) % p.size();
      if (i == 1 \mid \mid j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "</pre>
       << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
       << ProjectPointSegment (PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
  // expected: 1 0 1
  cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
       << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 0 0 1
  cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
       << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
       << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
  // expected: 1 1 1 0
  cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "</pre>
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
  // expected: (1,2)
  cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
  // expected: (1.1)
  cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
  vector<PT> v:
  v.push_back(PT(0,0));
  v.push_back(PT(5,0));
  v.push_back(PT(5,5));
  v.push_back(PT(0,5));
  // expected: 1 1 1 0 0
```

```
cerr << PointInPolygon(v, PT(2,2)) << " "
     << PointInPolygon(v, PT(2,0)) << " "
     << PointInPolygon(v, PT(0,2)) << " "
     << PointInPolygon(v, PT(5,2)) << " "
     << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "
     << PointOnPolygon(v, PT(2,0)) << " "
     << PointOnPolygon(v, PT(0,2)) << " "
     << PointOnPolygon(v, PT(5,2)) << " "
     << PointOnPolygon(v, PT(2,5)) << endl;
// expected: (1,6)
              (5,4) (4,5)
              blank line
              (4,5) (5,4)
              (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5, 4.5), 10, sqrt(2.0)/2.0);

for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
 u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0); \\  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl; 
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;</pre>
return 0;
```

8.5 Point Inside Polygon O(lg n)

```
* Solution for UVa 11072 - Points
* On this problem you must calculate the convex hull on the
* first set of points.
\star And for each point of the second set, answer if the point
* is inside or outside the convex hull.
typedef struct sPoint {
  11 x, y;
  sPoint (11 _x, 11 _y) : x(_x), y(_y) {}
  bool operator<(const sPoint& other) const
    if(x == other.x) return y < other.y;</pre>
    return x < other.x;</pre>
} point;
vector<point> vp, ch;
ll cross(point a, point b, point c) // AB x BC
  a.x = b.x; a.v = b.v;
  b.x -= c.x; b.y -= c.y;
  return a.x*b.y - a.y*b.x;
vector<point> convexhull()
  sort(vp.begin(), vp.end());
  vector<point> 1, u;
  for(int i = 0; i < vp.size(); i++)</pre>
    while (1.size() > 1 && cross(1[1.size()-2], 1[1.size()-1], vp[i]) <= 0)
      1.pop_back();
    1.pb(vp[i]);
```

```
for(int i = vp.size()-1; i >= 0; i--)
    while (u.size() > 1 && cross(u[u.size()-2], u[u.size()-1], vp[i]) <= 0)
     u.pop_back();
    u.pb(vp[i]);
  1.pop_back(); u.pop_back();
  1.insert(1.end(), u.begin(), u.end());
  return 1;
11 area(point a, point b, point c)
{ return llabs(cross(a, b, c)); }
bool insideTriangle(point a, point b, point c, point p)
  return area(a, b, c) == (area(a, b, p) +
      area(a, c, p) +
      area(b, c, p));
bool isInside(point p)
  if(ch.size() < 3) return false;</pre>
  int i = 2, j = ch.size()-1;
  while(i < j)
    int mid = (i+j)/2;
    11 c = cross(ch[0], ch[mid], p);
    if(c > 0) i = mid+1;
    else j = mid;
  return insideTriangle(ch[0], ch[i], ch[i-1], p);
int main()
  int n;
  while (true)
    ch.clear();
    vp.clear();
    cin >> n;
    if(not cin) break;
    while (n--)
      point p;
      cin >> p.x >> p.y;
      vp.pb(p);
    ch = convexhull();
    cin >> n:
    while (n--)
     point p;
      cin >> p.x >> p.y;
      if(isInside(p)) cout << "inside\n";</pre>
      else cout << "outside\n";</pre>
  return 0;
```

9 Data Structures

9.1 Disjoint Set Union

```
const int N=500010;
int p[N], Rank[N];
void init()
{
   memset(Rank, 0, sizeof(Rank));
   for(int i=0;i<N;i++) p[i]=i;
}
int findset(int i)
{
   if(p[i]==i) return i;
   return p[i]=findset(p[i]);</pre>
```

```
}
bool same(int i, int j)
{
  return (findset(i) == findset(j));
}
void unionSet(int i, int j)
{
   if (!same(i, j)) {
     int x = findset(i), y=findset(j);
   if (Rank[x] > Rank[y])
     p[y] = x;
   else {
     p[x] = y;
     if (Rank[x] == Rank[y]) Rank[y]++;
   }
}
```

9.2 Persistent Segment Tree

```
//PRINTAR O NUMERO DE ELEMENTOS DISTINTOS
//EM UM INTERVALO DO ARRAY
const int N = 30010;
int tr[100 \star N], L[100 \star N], R[100 \star N], root[100 \star N];
int v[N], mapa[100 * N];
int cont = 1:
void build(int node, int b, int e)
  if (b == e) {
    tr[node] = 0;
  else {
    L[node] = cont++;
    R[node] = cont++;
    build(L[node], b, (b + e) / 2);
    build(R[node], (b + e) / 2 + 1, e);
    tr[node] = tr[L[node]] + tr[R[node]];
int update(int node, int b, int e, int i, int val)
  int idx = cont++:
  tr[idx] = tr[node] + val;
  L[idx] = L[node];
  R[idx] = R[node];
  if (b == e) return idx;
  int mid = (b + e) / 2;
  if (i <= mid)</pre>
    L[idx] = update(L[node], b, mid, i, val);
    R[idx] = update(R[node], mid + 1, e, i, val);
  return idx;
int query(int nodeL, int nodeR, int b, int e, int i, int j)
  if (b > j \text{ or } i > e) return 0;
  if (i <= b and j >= e) {
    int p1 = tr[nodeR];
    int p2 = tr[nodeL];
    return p1 - p2;
  int mid = (b + e) / 2;
  return query(L[nodeL], L[nodeR], b, mid, i, j) +
         query(R[nodeL], R[nodeR], mid + 1, e, i, j);
int main()
  int n:
  sc(n):
 memset (mapa, -1, sizeof (mapa));
for (int i = 0; i < n; i++) sc(v[i]);</pre>
  build(1, 0, n - 1);
  for (int i = 0; i < n; i++) {
    if (mapa[v[i]] == -1) {
      root[i + 1] = update(root[i], 0, n - 1, i, 1);
      mapa[v[i]] = i;
      root[i + 1] = update(root[i], 0, n - 1, mapa[v[i]], -1);
      mapa[v[i]] = i;
      root[i + 1] = update(root[i + 1], 0, n - 1, i, 1);
  int q;
  sc(q);
  for (int i = 0; i < q; i++) {
    int 1, r;
```

```
sc2(1, r);
int resp = query(root[1 - 1], root[r], 0, n - 1, 1 - 1, r - 1);
pri(resp);
}
return 0;
```

9.3 Segment Tree

```
//compilar em C++11, essa segment tree
//computa qual e o k's elemento compreendido
//no intervalo entre i, j
//presentes no array
vi tr[5 * N]:
void build(int node, int b, int e)
  if (b == e)
    tr[node].pb(v[b]);
  else {
    build(2 * node, b, (b + e) / 2);
    build(2 * node + 1, (b + e) / 2 + 1, e);
    merget(tr[2 * node], tr[2 * node + 1], tr[node]);
merge(tr[2 * node].begin(), tr[2 * node].end(), tr[2 * node + 1].begin(),
           tr[2 * node + 1].end(), back_inserter(tr[node]));
int query(int node, int b, int e, int i, int j, int k)
  if (i > e \text{ or } b > j) return 0;
  if (i \le b \text{ and } j \ge e) {
    int resp =
         upper_bound(tr[node].begin(), tr[node].end(), k) - tr[node].begin();
    return tr[node].size() - resp;
  return query(2 * node, b, (b + e) / 2, i, j, k) +
query(2 * node + 1, (b + e) / 2 + 1, e, i, j, k);
```

9.4 Sparse Table

9.5 Cartesian Tree

```
if(t) return t->subtr;
                 return 0;
        int key(Node *t) {
                 if(t) return t->val;
                 return 0;
        int sum(Node *t) {
                 if(t) return t->sum;
                 return 0;
        int pref(Node *t) {
                 if(t) return t->pref;
                 return -INF:
        int suf(Node *t) {
                 if(t) return t->suf;
                 return -INF;
        int maximo(Node *t) {
                 if(t) return t->maximo;
                 return -INF:
        void upd(Node* &t){
                if(t){
                          if(!(t->1)){
                                  t->pref= max(t->val, t->val + pref(t->r));
                          else
                                  t->pref= max( pref(t->1), max( sum(t->1) + t->val, sum(t->1) + t->val
                                         + pref(t->r)));
                          if(!(t->r)){
                                  t->suf= max(t->val, t->val + suf(t->1));
                          else{
                                  t->suf= max( suf(t->r), max( sum(t->r) + t->val, sum(t->r) + t->val +
                                         suf(t->1)));
                          t\rightarrow maximo = max(suf(t\rightarrow 1) + t\rightarrow val, suf(t\rightarrow 1) + t\rightarrow val + pref(t\rightarrow r));
                          t->maximo = max(t->maximo , pref(t->r) + t->val);
                         t->maximo = max(t->maximo, max( maximo(t->1), maximo(t->r)));
t->maximo= max(t->maximo, t->val);
                          t\rightarrow sum = sum(t\rightarrow r) + sum(t\rightarrow l) + t\rightarrow val;
                          t->subtr=cnt(t->1) + cnt(t->r) +1;
// junta todos menores que val e todos maiores ou iguais a val
        Node* merge(Node* L, Node *R) {
                if(!L) return R;
                 if(!R) return L;
                if(L->prior > R->prior) {
    L->r = merge(L->r, R);
                          upd(L):
                          return L:
                 R->1 = merge(L, R->1);
                 upd(R);
                 return R;
// separa t em todos menores que val , todos maiores ou igual a val
        pair<Node*, Node*> split(Node* t, int val, int add) {
                          return mp(nullptr, nullptr);
                 int cur_key= add+ cnt(t->1);
                 if(cur key < val){</pre>
                          auto ret= split(t->r, val, cur_key+1);
                          t->r= ret.first;
                          upd(t);
                          return mp(t, ret.second);
                 auto ret= split(t->1, val , add);
                 t->1 = ret second;
                 upd(t);
                 return mp(ret.first, t);
        int querymax(Node *&t, int i, int j) {
                 auto trl= split(t, j+1, 0);
                 auto tr2= split(tr1.first, i, 0);
                 int prefi= pref(tr2.second->r);
                 int sufi= suf(tr2.second->1);
                 int val= key(tr2.second);
                 int r=maximo(tr2.second);
                 auto x= merge(tr2.first, tr2.second);
                 t= merge(x, tr1.second);
```

return r;

```
void insert(Node* &t, int x, int y) {
                Node *aux= new Node(y);
                auto tr= split(t, x,0);
                auto traux=merge(tr.first,aux);
                t=merge(traux, tr.second);
        void replace(Node *&t, int x, int y) {
                Node *aux= new Node(y);
                erase(t, x);
                auto tr=split(t, x, 0);
                t=merge(tr.first,aux);
                //db(pref(t)):
                //db(suf(t));
                t=merge(t, tr.second);
                        db(pref(t));
                         db(suf(t));
        void erase(Node * &t, int x) {
                auto tr=split(t,x+1,0);
                auto tr2=split(tr.first, x,0);
                t= merge(tr2.first, tr.second);
int main()
        int n;
        sc(n);
        for(int i=0;i<n;i++) {</pre>
                int x;
                T.insert(T.root, i, x);
        int q;
        sc(q);
        while (q--) {
                //db(T.cnt(T.root));
                char op;
                cin>>op;
if(op=='I'){
                        int x, y;
                         sc2(x, y);
                         T.insert(T.root, x, y);
                else if(op=='Q'){
                        int 1, r;
                         sc2(1,r);
                         1--, r--;
                         pri(T.querymax(T.root, 1,r));
                else if(op=='R'){
                        int x, y;
                         sc2(x,y);
                         T.replace(T.root, x, y);
                else{
                         int x;
                         sc(x);
                         T.erase(T.root, x);
        return 0:
```

9.6 Cartesian Tree 2

```
int bigrand() { return (rand() <<16) ^rand();}
char r[500001];
struct Node{
    int prior , subtr, sujo;
    int val,add;
    Node *1, *r;
    Node () {}
    Node (int c) : add(0), val(c), prior(bigrand()), l(NULL), r(NULL), subtr(1) {}
};
struct Treap{
    Node *root;
    Treap() : root(NULL) {};
    int cnt(Node *t) {
        if(t) return t->subtr;
    }
}
```

```
return 0;
void upd(Node* &t) {
        if(t){
                 if(t->sujo){
                         swap(t->1, t->r);
                         t->sujo=0;
                         if(t->1) {
                                 t->1->sujo^=1;
                         if(t->r){
                                 t->r->sujo^=1;
                 t->val+=t->add;
                 if(t->1){
                         t->1->add+=t->add;
                 if(t->r){
                         t->r->add+=t->add;
                 t->add=0;
                t->subtr= cnt(t->1) + cnt(t->r) + 1;
Node* merge(Node *L, Node *R) {
        upd(R);
        upd(L);
        if(!L) return R;
        if(!R) return L;
        if(L-> prior > R->prior) {
                L \rightarrow r = merge(L \rightarrow r, R);
                 upd(R);
                 return L;
        R->1 = merge(L,R->1);
        upd(R);
        upd(L);
        return R;
//<. >= val
pair<Node*, Node*> split(Node *t, int val, int add) {
        if(!t) {
                return mp(nullptr, nullptr);
        upd(t);
        int cur_key= add + cnt(t->1);
if(cur_key < val){</pre>
                 auto ret= split(t->r, val , cur_key+1);
                 t->r= ret.first;
                 upd(t);
                 return mp(t, ret.second);
        auto ret= split( t->1, val , add);
        t \rightarrow 1 = ret.second;
        upd(t);
        return mp(ret.first, t);
Node* inverte(Node* &t, int i, int j, int val) {
        if(i>j) return t;
        auto tr1= split(t, j+1, 0);
        auto tr2= split(tr1.first, i, 0);
        if(tr2.second) {
                tr2.second->sujo^=1;
                tr2.second->add+=val;
        auto x=merge(tr2.first,tr2.second);
        x=merge(x,tr1.second);
        return x;
void att(Node* &t, int 1 , int r, int i, int j) {
        t = inverte(t,r+1,i-1,-1);
        t=inverte(t,1,j,1);
void imprime(Node* &t, int add) {
        if(t){
                 upd(t);
                int cur_key= add + cnt(t->1);
                imprime(t->1, add);
                imprime(t->r, cur_key+1);
                 int aux=t->val+t->add;
                 aux%=26;
                 aux+=26;
                 aux%=26;
                 r[cur_key] = aux+'a';
```

```
void poe(Node* &t, string &s){
                 for(int i=0;i<s.size();i++){</pre>
                         Node *aux = new Node(s[i]-'a');
                          auto tr= split(t, i, 0);
                          auto traux= merge(tr.first, aux);
                          t= merge(traux, tr.second);
int main()
        BUFF:
        int X;
        cin>>X;
        while (X--) {
                 Treap T:
                 string s;
                 int op;
                 cin>>s>>op;
                 T.poe(T.root, s);
                 //T.imprime(T.root,0);
                 //for(int i=0;i<s.size();i++) {
                         cout << r[i];
                 //cout<<endl:
                 //assert (T.root!=NULL);
                 while(op--){
                         int 1, r, i, j;
                         cin>>l>>r>>i>>j;
                         1--,r--,i--,j--;
T.att(T.root,l,r,i,j);
                 T.imprime(T.root,0);
                 for(int i=0;i<s.size();i++) cout<<r[i];</pre>
                 cout << end1;
        return 0;
```

10 Miscellaneous

10.1 Hashing

```
//certificar que gethash() foi chamado
//antes de getHash(i,j);
struct Hashing {
 const string &s;
 int n, idx;
vector<ll> hashes,M,B;
  Hashing(const string &s) : s(s), hashes(s.size()){
    M={1000000409, 2000003273, 2000003281, 2000003293};
    B={31, 53, 61, 41};
    srand(time(NULL));
    idx=rand()%4;
    getHash();
  void otherprime(){
    idx=(idx+1)%4;
  ll eleva(ll a, ll b)
    if (b == 0)
     return 1;
    else if (b == 1)
     return a:
    11 \times = eleva(a, b / 2);
    if (b % 2 == 0)
      return (x * x) % M[idx];
      return (a * ((x * x) % M[idx])) % M[idx];
  /*hash da string de 0 ate i*/
  void getHash()
    int n = s.size();
    11 hp = 0;
for (int i = 0; i < s.size(); i++) {
  hp = int_mod(hp * B[idx] + s[i]);</pre>
      hashes[i] = hp;
  /*Hash da string compreendida entre i e j*/
```

```
11 getHash(int i, int j)
{
    if (i == 0) return hashes[j];
    11 h1 = hashes[j];
    11 h2 = (hashes[i - 1] * eleva(B[idx], j - i+1)) % M[idx];
    11 ret = (h1 - h2) % M[idx] + M[idx];
    return ret % M[idx];
};
```

10.2 Invertion Count

```
//conta o numero de inversoes de um array
//x e o tamanho do array, v e o array que quero contar
ll inversoes = 0;
void merge_sort(vi &v, int x)
  if (x == 1) return;
  int tam_esq = (x + 1) / 2, tam_dir = x / 2;
  int esq[tam_esq], dir[tam_dir];
  for (int i = 0; i < tam_esq; i++) esq[i] = v[i];
for (int i = 0; i < tam_dir; i++) dir[i] = v[i + tam_esq];</pre>
  merge_sort(esq, tam_esq);
  merge_sort(dir, tam_dir);
  int i_esq = 0, i_dir = 0, i = 0;
  while (i_esq < tam_esq or i_dir < tam_dir) {
  if (i_esq == tam_esq) {</pre>
       while (i_dir != tam_dir) {
        v[i] = dir[i_dir];
         i_dir++, i++;
    else if (i_dir == tam_dir) {
       while (i_esq != tam_esq) {
         v[i] = esq[i_esq];
         i_esq++, i++;
         inversoes += i_dir;
    else {
      if (esq[i_esq] <= dir[i_dir]) {
  v[i] = esq[i_esq];</pre>
         i++, i esq++;
         inversoes += i_dir;
         v[i] = dir[i_dir];
         i++, i_dir++;
```

10.3 Distinct Elements in ranges

```
const int MOD = 1e9 + 7;
const int N = 1e6 + 10;
int bit[N], v[N], id[N], r[N];
bool compare(int x, int y) { return query[x] < query[y]; }</pre>
void add(int idx, int val)
  while (idx < N) {
    bit[idx] += val;
    idx += idx & -idx:
int sum(int idx)
  while (idx > 0) {
    ret += bit[idx];
    idx -= idx & -idx;
  return ret:
int main()
  memset(bit, 0, sizeof(bit));
  memset (mapa, 0, sizeof (mapa));
  int n;
  sc(n);
```

```
for (int i = 1; i <= n; i++) sc(v[i]);</pre>
int q;
for (int i = 0; i < q; i++) {
  sc2(query[i].second, query[i].first);
sort(id, id + q, compare);
sort(query, query + q);
int j = 1;
for (int i = 0; i < q; i++) {
  int L = query[i].second;
int R = query[i].first;
  while (j \le R) {
   if (mapa[v[j]] > 0) {
  add(mapa[v[j]], -1);
       mapa[v[j]] = j;
       add(mapa[v[j]], 1);
       mapa[v[j]] = j;
       add(mapa[v[j]], 1);
    j++;
  r[id[i]] = sum(R);
  if (L > 1) r[id[i]] -= sum(L - 1);
for (int i = 0; i < q; i++) pri(r[i]);</pre>
return 0:
```

10.4 Maximum Rectangular Area in Histogram

```
* Complexidade : O(N)
ll solve(vi &h)
 int n = h.size();
 11 \text{ resp} = 0;
  stack<int> pilha;
    i = 0;
  while (i < n) {
    if (pilha.empty() or h[pilha.top()] <= h[i]) {</pre>
     pilha.push(i++);
    else (
     int aux = pilha.top();
      pilha.pop();
      resp =
         max(resp, (ll)h[aux] * ((pilha.empty()) ? i : i - pilha.top()-1));
  while (!pilha.empty()) {
    int aux = pilha.top();
    resp = max(resp, (11)h[aux] * ((pilha.empty()) ? n : n - pilha.top()-1));
```

10.5 Multiplying Two LL mod n

10.6 Josephus Problem

```
/* Josephus Problem - It returns the position to be, in order to not die. O(n) \star / \star With k=2, for instance, the game begins with 2 being killed and then n+2, n+4, ... \star / \star 11 josephus(l1 n, l1 k) { if(n=1) return 1; else return (josephus(n-1, k)+k-1)%n+1; }
```

10.7 Josephus Problem 2

10.8 Dynamic MST

```
* Code for URI 1887
* It gives an tree and a bunch of queries to add
* edges from a to b with cost c.
const int MOD = 1e9 + 9;
struct ed{
        int u, v, w, t;
        ed(int _u, int _v, int _w, int _t) { u=_u, v=_v, w=_w, t=_t;}
         ed(){};
        bool operator < ( const ed &a) const {</pre>
                return w<a.w;
const int N=50010;
int p[N],id[N];
void init(int n)
        for (int i=1;i<=n;i++) p[i]=i;</pre>
int findSet(int i)
        if(p[i]==i) return i;
        return p[i]=findSet(p[i]);
bool unionSet(int i, int j)
        int x=findSet(i),y=findSet(j);
        if(x==y) return false;
        return true;
void reduction(int 1,int r,int &n,vector<ed> &graph,int &res)
        vector<ed> g;
        init(n);
        sort(graph.begin(), graph.end());
        for(int i=0;i<graph.size();i++)</pre>
                if(graph[i].t<=r and (graph[i].t>=l or unionSet(graph[i].u,graph[i].v))){
                         g.pb(graph[i]);
        graph=g;
void contraction(int 1,int r,int &n,vector<ed> &graph,int &res)
        vector<ed> g;
        init(n):
        sort(graph.begin(),graph.end());
        for(int i=0;i<(int)graph.size();i++)</pre>
                if(graph[i].t>=1) unionSet(graph[i].u,graph[i].v);
        for (int i=0; i < (int) graph.size(); i++) {</pre>
                if(graph[i].t<l and unionSet(graph[i].u,graph[i].v)){</pre>
                         g.pb(graph[i]);
                         res+=graph[i].w;
        init(n);
        for (int i=0;i<g.size();i++) {</pre>
                unionSet(q[i].u,q[i].v);
        int tot=0;
```

```
for (int i=1; i<=n; i++) id[i]=0;</pre>
        for (int i=1; i<=n; i++) {</pre>
                 int f=findSet(i);
                 if(!id[f]) id[f]=++tot;
                 id[i]=id[f];
        for(int i=0;i<graph.size();i++){</pre>
                 graph[i].u=id[graph[i].u],graph[i].v=id[graph[i].v];
void solve(int l,int r,int n,vector<ed> graph,int res)
        reduction(l,r,n,graph,res);
         contraction(1,r,n,graph,res);
                 sort(graph.begin(),graph.end());
                 for(int i=0;i<(int)graph.size();i++)</pre>
                          if(unionSet(graph[i].u,graph[i].v)){
                                  res+=graph[i].w;
                          pri(res);
                 return:
        int mid=(1+r)/2;
        solve(l,mid,n,graph,res);
        solve(mid+1, r, n, graph, res);
int main()
        int T;
        while (T--)
                 int n,m,q;
                 sc3(n,m,q);
                 vector<ed> graph;
for(int i=1;i<=m;i++)</pre>
                          int u, v, w;
                          sc3(u, v, w);
                          int t=0;
                          graph.pb(ed(u,v,w,t));
                 for(int i=1;i<=q;i++)
                          int u, v, w;
                          sc3(u, v, w);
                          int t=i;
                          graph.pb(ed(u,v,w,t));
                 solve(1,q,n,graph,0);
        return 0:
```

10.9 Ordered Static Set (Examples)

```
#include <bits/stdc++.h>
using namespace std;
#define sc(a) scanf("%d", &a)
#define sc2(a,b) scanf("%d%d", &a, &b) #define sc3(a,b,c) scanf("%d%d%d", &a, &b, &c)
#define pri(x) printf("%d\n", x)
#define mp make_pair
#define pb push_back
#define BUFF ios::sync_with_stdio(false);
#define imprime(v) for(int X=0; X<v.size(); X++) printf("%d ", v[X]); printf("\n");
#define endl "\n"
const int INF= 0x3f3f3f3f3f;
const long double pi= acos(-1);
typedef long long int 11;
typedef long double ld;
typedef pair<int, int> ii;
typedef vector<int> vi;
typedef vector< vector< int > > vvi;
const int MOD=1e9+7;
///USANDO ORDERED STATIC SET PRA ESTRUTURA
#include <ext/pb_ds/assoc_container.hpp> // Common file
#include <ext/pb_ds/tree_policy.hpp> // Including tree_order_statistics_node_update
using namespace __gnu_pbds;
```

```
typedef struct cu {
         int a;
         int b;
         bool operator < (const struct cu &other) const {</pre>
                  if(a != other.a) return a < other.a;</pre>
                  return b < other.b;</pre>
         bool operator == (const struct cu &other) const {
                  return(a == other.a and b == other.b);
cuzao:
bool cmp(const cuzao &a, const cuzao &b) {
         return true;
typedef tree<
cuzao,
         null_type,
         less<cuzao>,
         rb_tree_tag,
         tree_order_statistics_node_update>
         ordered_set;
int main()
         ordered_set os;
         cuzao asd;
         asd.a = 1;
         asd.b = 2;
         os.insert(asd);
         asd.a = 4;
         os.insert(asd);
         cout << (os.find(asd) == end(os)) << endl; //0
         cout<<os order_of_key(asd)<<endl;//1</pre>
         asd.a = 1;
         cout<<os.order_of_key(asd)<<endl;//0
cout<<os.find_by_order(0)->a<<" "<<os.find_by_order(0)->b<<endl;//1 2
cout<<os.find_by_order(1)->a<<" "<<os.find_by_order(1)->b<<endl;//4 2</pre>
         return 0;
//USANDO ORDERED STATIC SET PRA CONTAINER DO STL MESMO
#include <bits/stdc++.h>
using namespace std;
#define sc(a) scanf("%d", &a)
```

```
#define sc2(a,b) scanf("%d%d", &a, &b)
#define sc3(a,b,c) scanf("%d%d%d", &a, &b, &c)
#define pri(x) printf("%d\n", x)
#define mp make_pair
#define pb push_back
#define BUFF ios::sync_with_stdio(false);
#define imprime(v) for(int X=0; X<v.size(); X++) printf("%d ", v[X]); printf("\n");
#define endl "\n"
const int INF= 0x3f3f3f3f3f;
const long double pi= acos(-1);
typedef long long int 11;
typedef long double ld;
typedef pair<int, int> ii;
typedef vector<int> vi;
typedef vector< vector< int > > vvi;
const int MOD=1e9+7;
const 11 LINF=0x3f3f3f3f3f3f3f3f3f3f;
#include <ext/pb_ds/assoc_container.hpp> // Common file
#include <ext/pb_ds/tree_policy.hpp> // Including tree_order_statistics_node_update
using namespace __gnu_pbds;
typedef tree<
         null_type,
         less<int>.
         rb_tree_tag,
tree_order_statistics_node_update>
         ordered_set;//n multi
int main()
         ordered_set os;
         os.insert(1);
         os.insert(10);
         os.insert(1);
         os.insert(15);
         cout<<(os.find(10) == end(os))<<endl;//0 mesma coisa q !count</pre>
         cout<<os.order_of_key(10)<<endl;//1 qual o indice do valor 10, se n tem o indice, pega o</pre>
                proximo
         cout<<os.order_of_key(2)<<endl;//l
cout<<+os.upper_bound(2)<<endl;//locut<+os.ipper_bound(2)<<endl;//locut<+os.find_by_order(0)<<endl;//l
cout<<+os.find_by_order(2)<<endl;//ls
         return 0;
```