



Cohesive Subgraph Detection

Clique Model

Never Stand Still

Faculty of Engineering

Computer Science and Engineering

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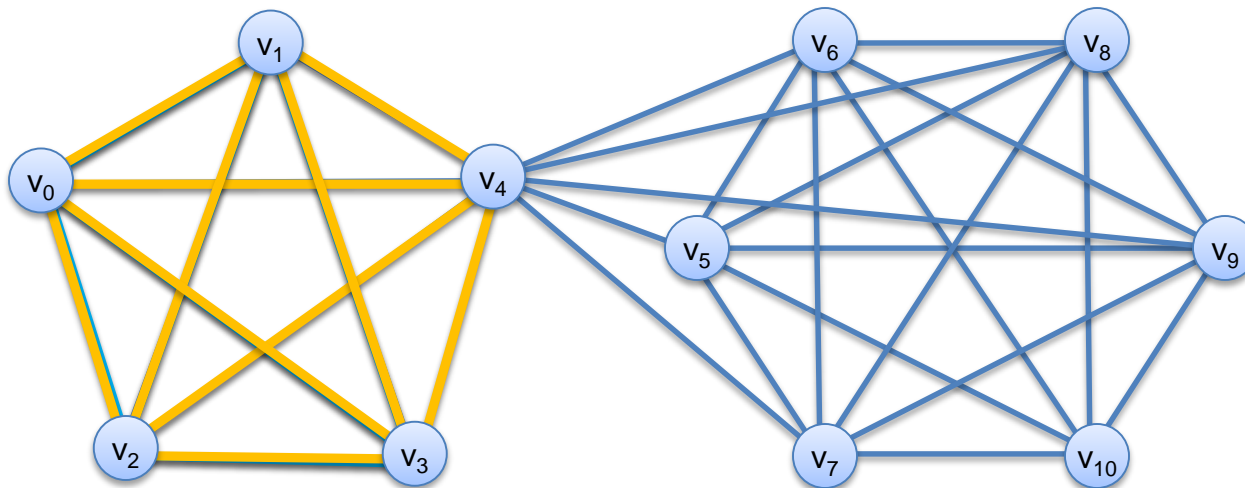
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Outline

- Clique Model
- K-Core vs K-Truss
- K-Edge Connected
- K-Vertex Connected

Clique

- Given a graph G , a clique is a set of nodes such that for any pair of them have an edge
- A clique is called maximal clique if there exist no other bigger cliques that contain it

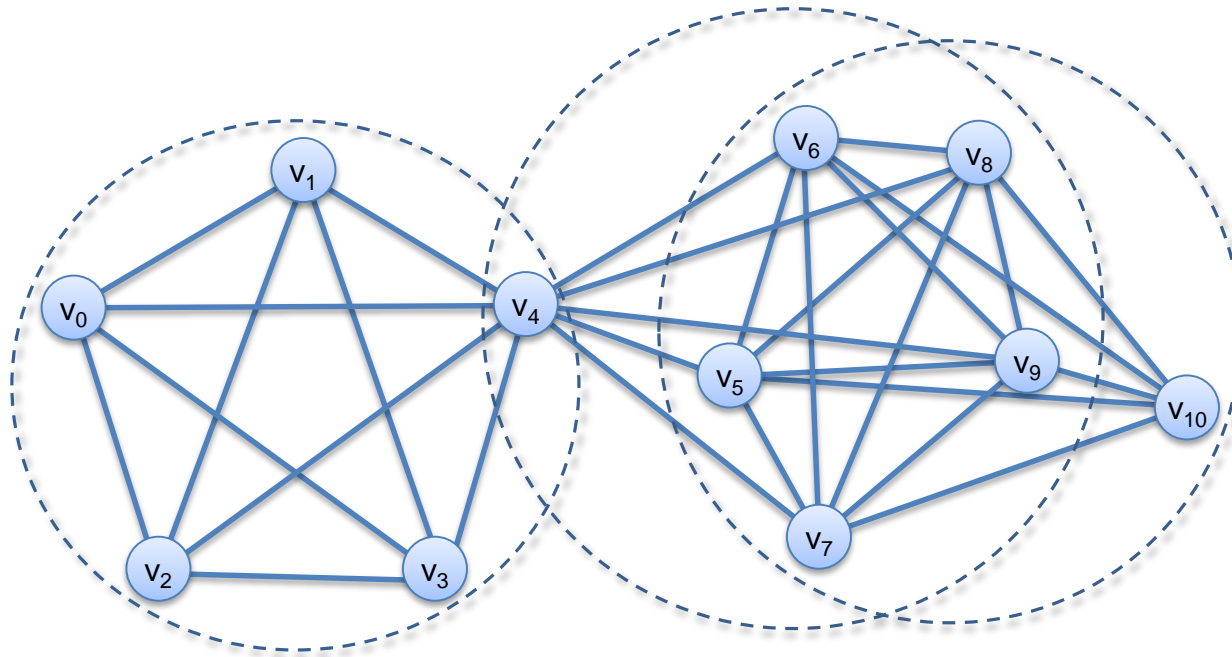


Application

- Community detection
- Gene expression and motif discovery
- Anomaly detection
- Stock market data visualization
- Signal transmission analysis
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Maximal Clique Enumeration

- Given a graph G , find all the maximal cliques in G .
 - NP-Hard Problem



In-Memory MCE

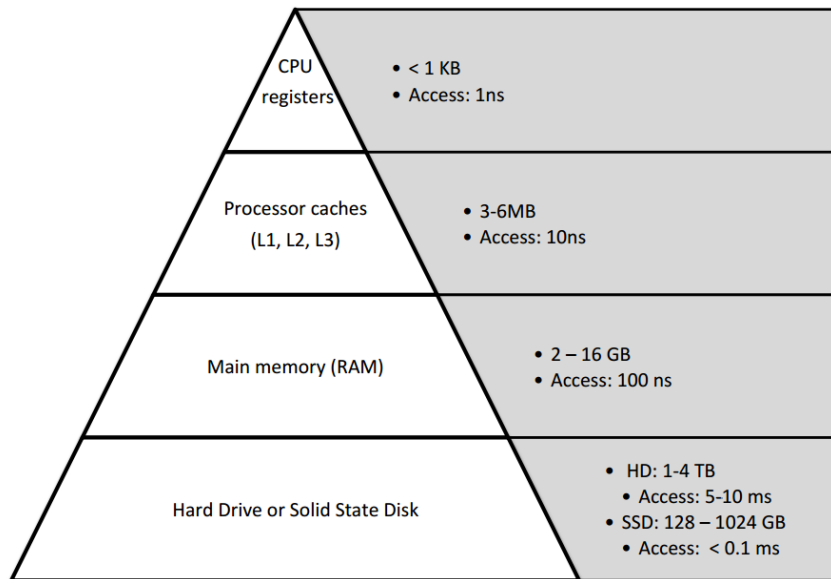
- Bron-Kerbosch Algorithm
 - First Practical In-Memory MCE Algorithm
 - In-Memory means all the input and auxiliary data structure can be loaded in main memory during the computation
 - C. Bron et al., “Algorithm 457: finding all cliques of an undirected graph”, *Communications of the ACM*, **16** (9): 575–577, 1973
 - Based on a recursive backtracking paradigm

I/O Efficient MCE

- Why I/O Efficient?
 - Real graph is massive
 - Facebook contains 1.32 billion nodes and 140 billion edges
 - EU-2015 (sub-domain of web graph) contains 1.07 billion nodes and 91 billion edges
 - Memory is fast but small while disk is slow but large

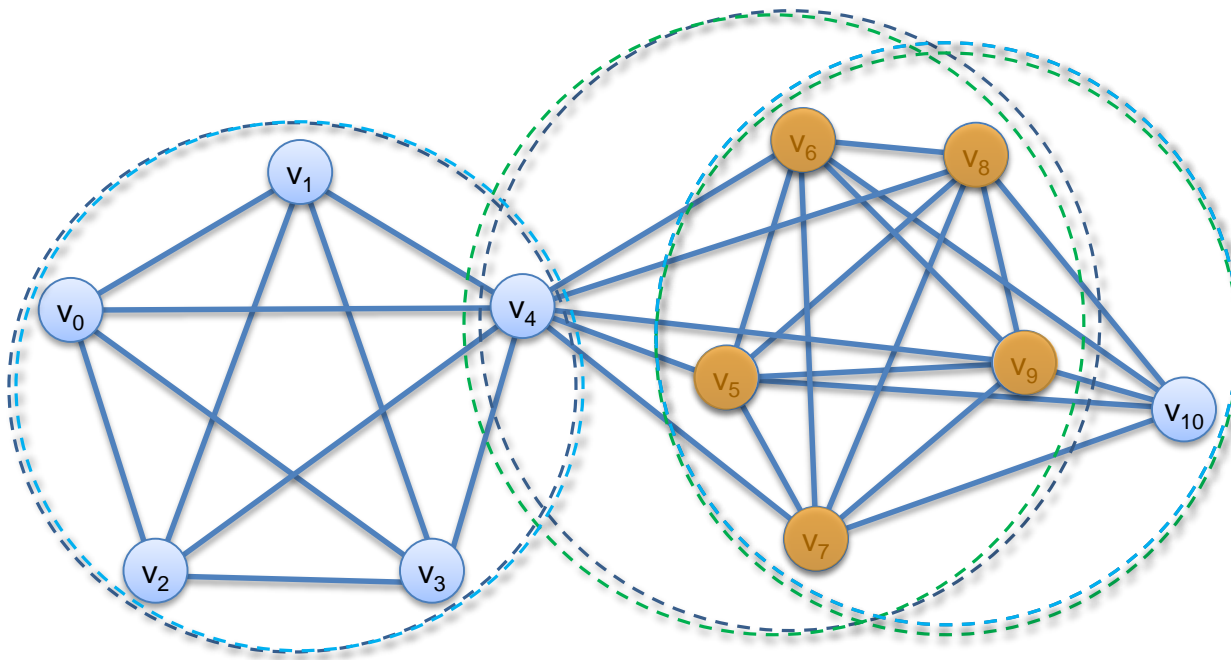
I/O Efficient MCE

- Why I/O Efficient?
 - Memory Hierarchy



Diversified Top-K Clique Search

- Traditional models vs diversified top-k clique



Diversified Top-K Clique Search

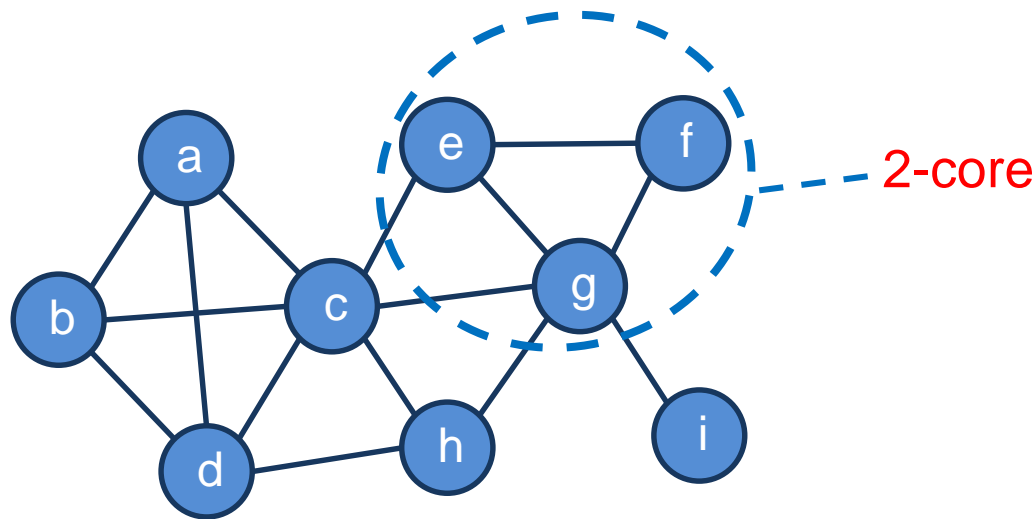
- Our Solution
 - treat it as an online k coverage problem
 - store k maximal cliques in memory
 - update these k candidate maximal cliques while enumerating cliques
 - replace small existing cliques with big new cliques

Diversified Top-K Clique Search

- PNP-Index
 - An naïve implementation for candidate set maintenance needs $O(|\mathcal{A}| * k * |C_{max}|)$
 - With the help of PNP-Index, our algorithm can only take $O(\sum_{C \in \mathcal{A}} |C|)$ time

k -Core

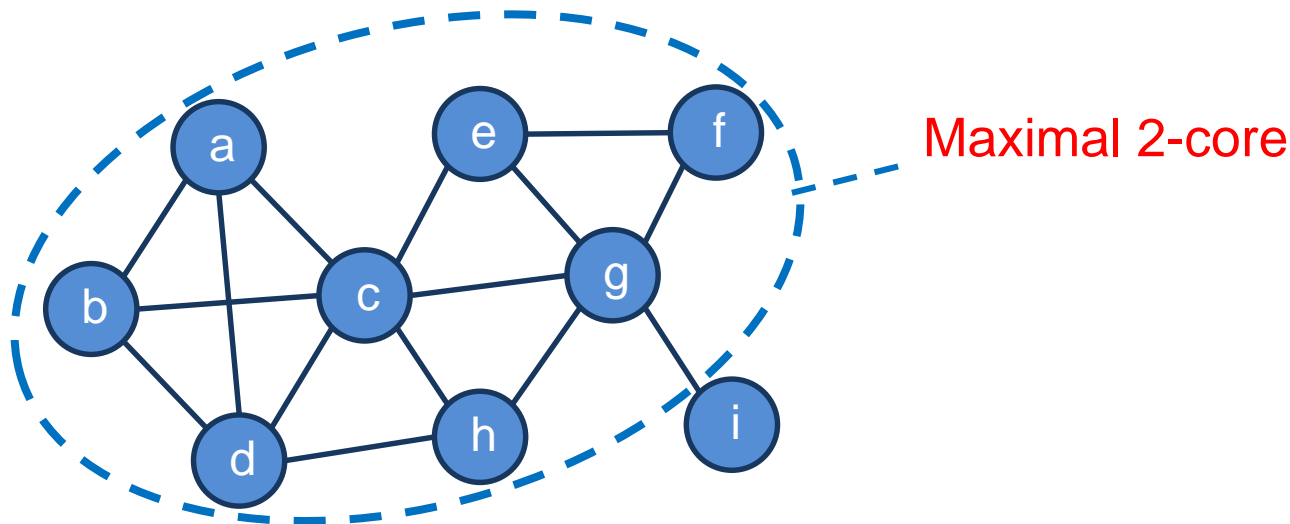
- Given a graph G , the k -core of G is a subgraph where **each node has at least k neighbors** (i.e., k adjacent nodes, or a degree of k).



S. B. Seidman. Network structure and minimum degree. *Social networks*, 5(3):269–287, 1983.

Maximal k -Core

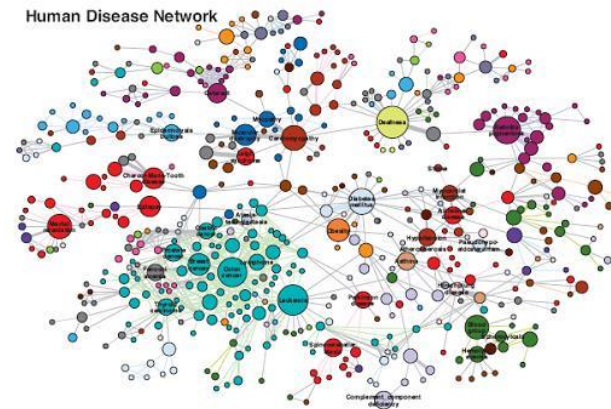
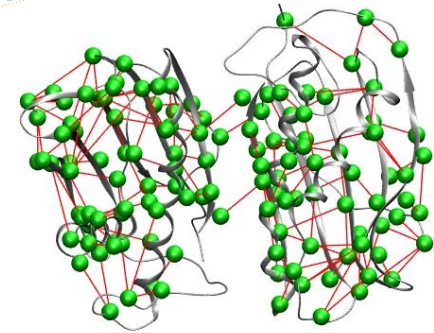
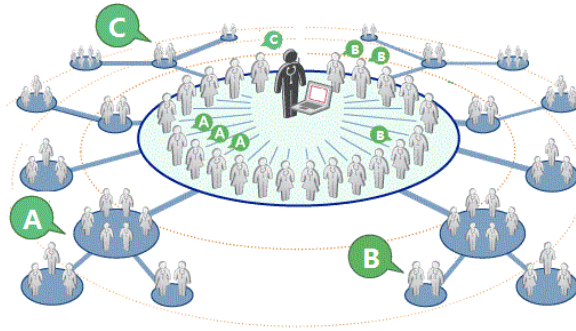
- A k -core C is called maximal if any supergraph of C is not a k -core (i.e., no another k -core which contains C).



S. B. Seidman. Network structure and minimum degree. Social networks, 5(3):269–287, 1983.

Applications

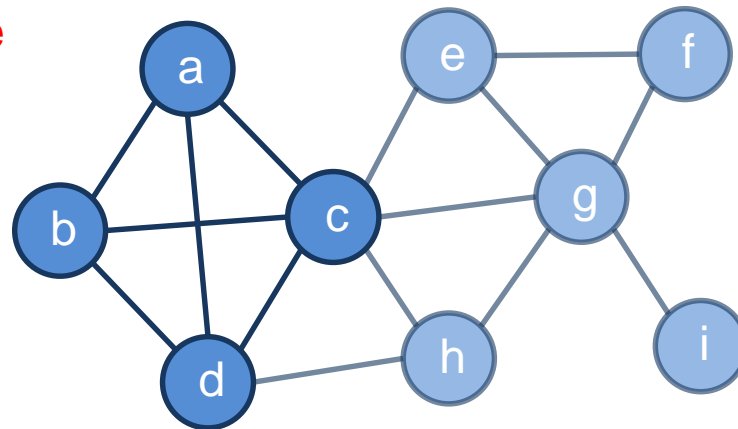
- Community detection
- Social contagion
- User engagement
- Event detection
- Network analysis and visualization
- Influence study
- Graph clustering
- Protein function prediction
- Human Cerebral Cortex
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Compute Maximal k -Core

- Given a graph G , the maximal k -core of G can be computed by recursively deleting every node and its adjacent edges if its degree is less than k .

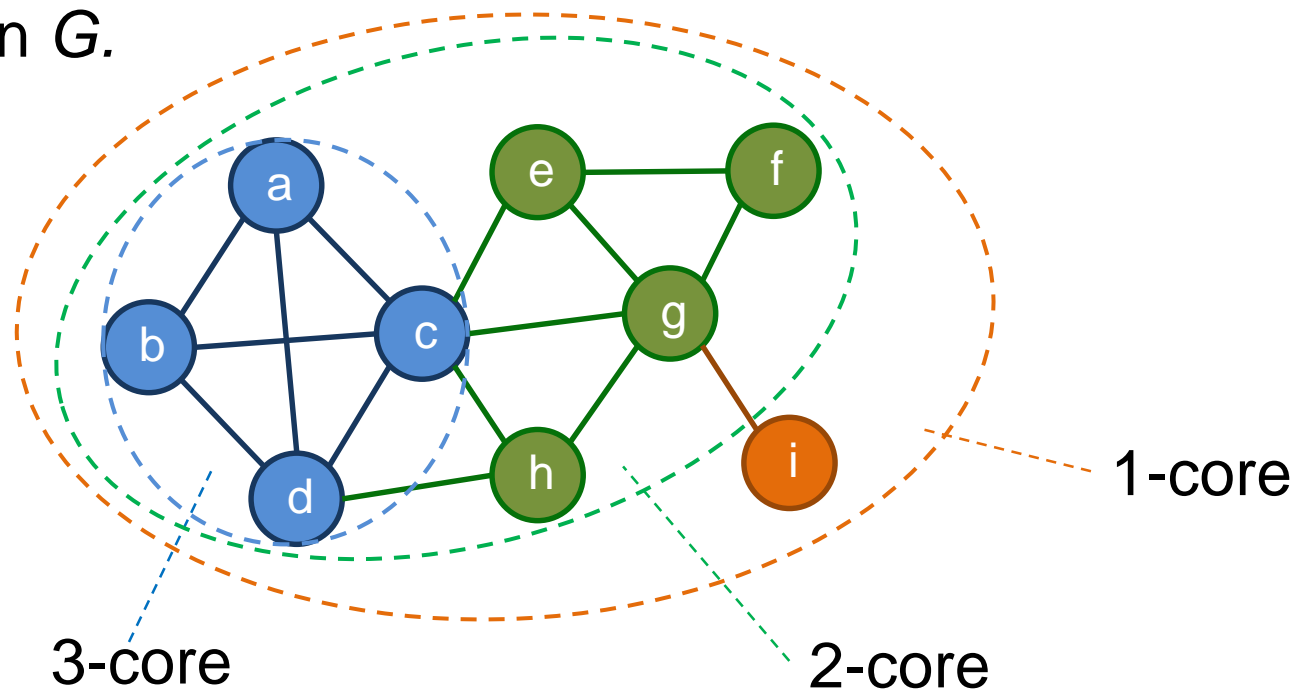
Maximal 3-core



S. B. Seidman. Network structure and minimum degree. Social networks, 5(3):269–287, 1983.

k -Core Decomposition

- **Core number** of a node v : the largest value of k such that there is a k -core containing v .
- Core decomposition: compute the **core number** of each node in G .



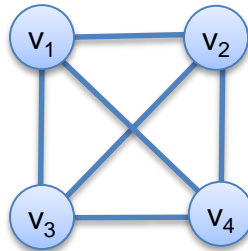
***k*-Truss**

- Given a graph G , the k -truss of G is a subgraph where edge is at least involved in $(k-2)$ triangles.
- k -truss is an enhancement of k -core; each vertex of k -truss has a degree at least $k-1$.

S. B. Seidman. Network structure and minimum degree. Social networks, 5(3):269–287, 1983.

k-edge Connectivity

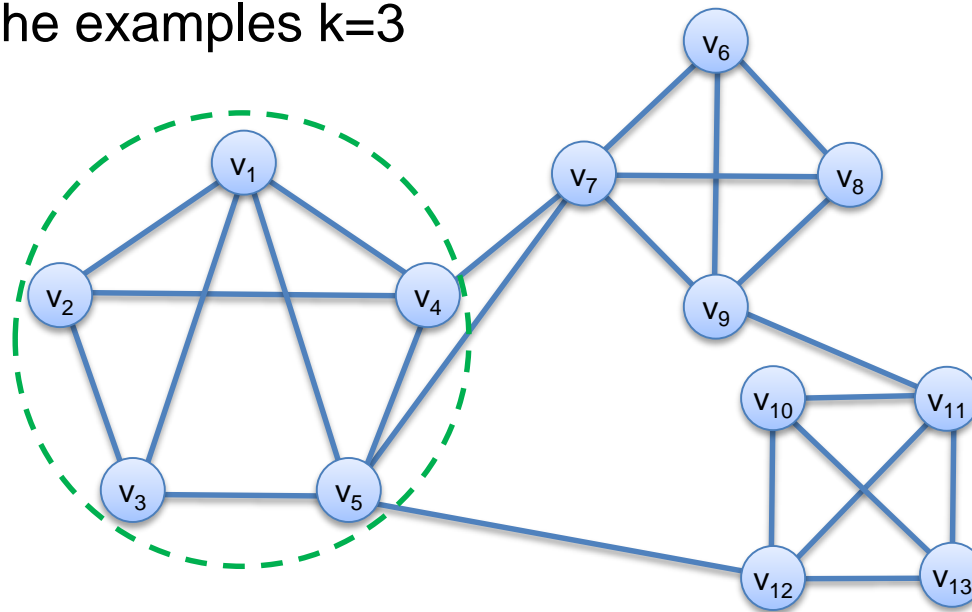
- A graph is k -edge connected if it is still connected after removing any set of $(k-1)$ edges from it



k-edge Connected Component

- A k-edge connected component (k-ECC) of a graph G is a maximal subgraph g of G such that g is k-edge connected

➤ All the examples $k=3$



Application

- Community detection
- Social behaviour mining
- Graph visualization
- Steiner Component Search
- Hierarchy Study in Networks
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k-Vertex Connectivity

- A graph is k -vertex connected if it is still connected after removing any set of $(k-1)$ vertex from it

