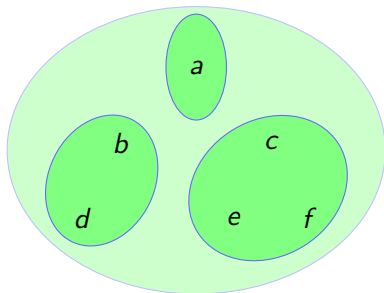


Disjoint Sets



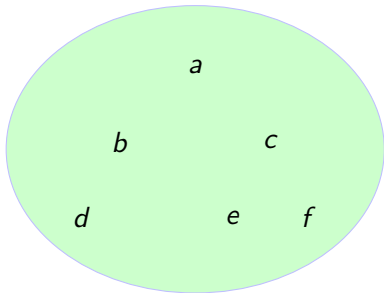
set of sets

$$\{\{a\}, \{b, d\}, \{c, e, f\}\}$$

equivalence relation

	<i>a</i>	<i>b</i>	<i>d</i>	<i>c</i>	<i>e</i>	<i>f</i>
<i>a</i>	1	0	0	0	0	0
<i>b</i>	0	1	1	0	0	0
<i>d</i>	0	1	1	0	0	0
<i>c</i>	0	0	0	1	1	1
<i>e</i>	0	0	0	1	1	1
<i>f</i>	0	0	0	1	1	1

Disjoint Set Operations



make-set(a)

make-set(b)

make-set(c)

make-set(d)

make-set(e)

make-set(f)

find-set(f) == f

union-sets(f, c)

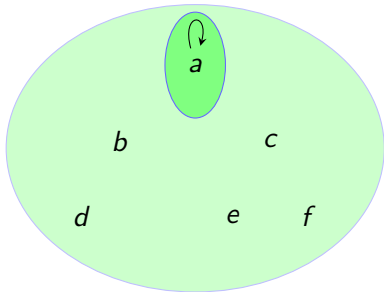
find-set(f) == c

union-sets(d, b)

union-sets(f, e)

find-set(f) == e

Disjoint Set Operations



→ make-set(a)
make-set(b)
make-set(c)
make-set(d)
make-set(e)
make-set(f)

find-set(f) == f

union-sets(f, c)

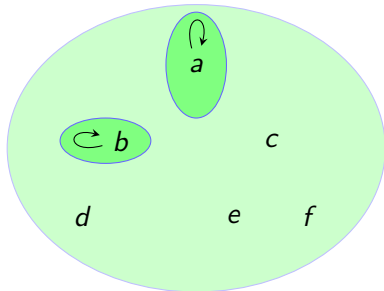
find-set(f) == c

union-sets(d, b)

union-sets(f, e)

find-set(f) == e

Disjoint Set Operations



→ make-set(a)
make-set(b)
make-set(c)
make-set(d)
make-set(e)
make-set(f)

find-set(f) == f

union-sets(f, c)

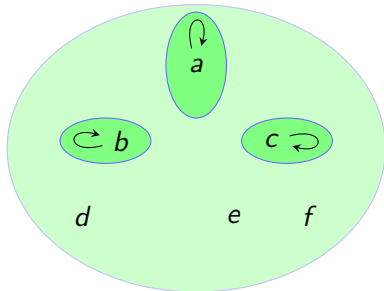
find-set(f) == c

union-sets(d, b)

union-sets(f, e)

find-set(f) == e

Disjoint Set Operations



→ make-set(a)
make-set(b)
make-set(c)
make-set(d)
make-set(e)
make-set(f)

find-set(f) == f

union-sets(f, c)

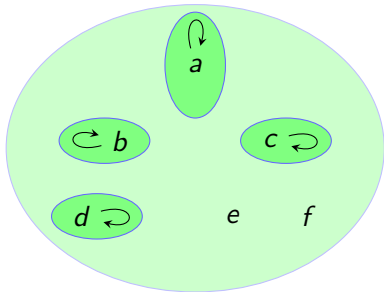
find-set(f) == c

union-sets(d, b)

union-sets(f, e)

find-set(f) == e

Disjoint Set Operations



→ make-set(a)
make-set(b)
make-set(c)
make-set(d)
make-set(e)
make-set(f)

find-set(f) == f

union-sets(f, c)

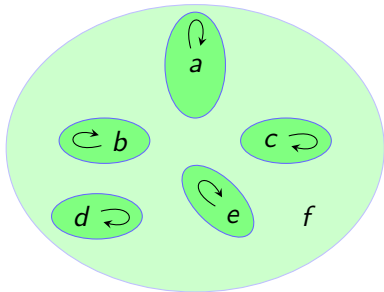
find-set(f) == c

union-sets(d, b)

union-sets(f, e)

find-set(f) == e

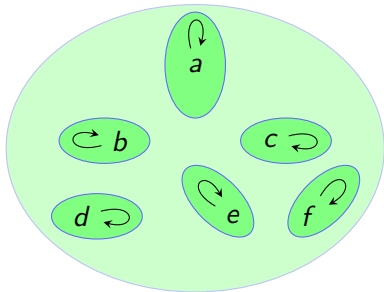
Disjoint Set Operations



make-set(a)
make-set(b)
make-set(c)
make-set(d)
make-set(e)
→ make-set(f)

find-set(f) == f
union-sets(f, c)
find-set(f) == c
union-sets(d, b)
union-sets(f, e)
find-set(f) == e

Disjoint Set Operations



make-set(a)

make-set(b)

make-set(c)

make-set(d)

make-set(e)

make-set(f)



find-set(f) == f

union-sets(f, c)

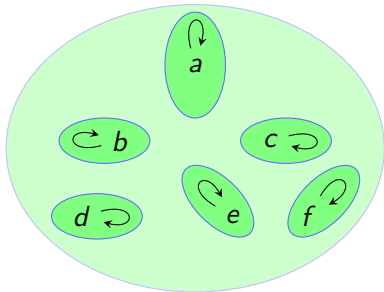
find-set(f) == c

union-sets(d, b)

union-sets(f, e)

find-set(f) == e

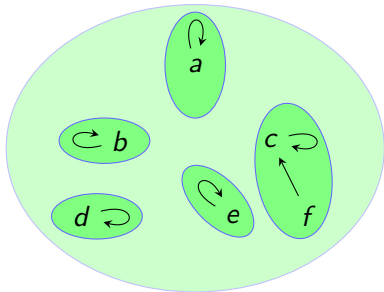
Disjoint Set Operations



make-set(a)
make-set(b)
make-set(c)
make-set(d)
make-set(e)
make-set(f)

→ find-set(f) == f
union-sets(f, c)
find-set(f) == c
union-sets(d, b)
union-sets(f, e)
find-set(f) == e

Disjoint Set Operations



make-set(a)

make-set(b)

make-set(c)

make-set(d)

make-set(e)

make-set(f)

find-set(f) == f

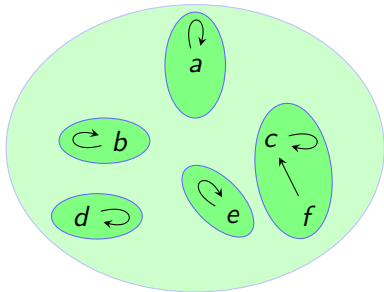
→ union-sets(f, c)
find-set(f) == c

union-sets(d, b)

union-sets(f, e)

find-set(f) == e

Disjoint Set Operations



make-set(a)

make-set(b)

make-set(c)

make-set(d)

make-set(e)

make-set(f)

find-set(f) == f

union-sets(f, c)

find-set(f) == c

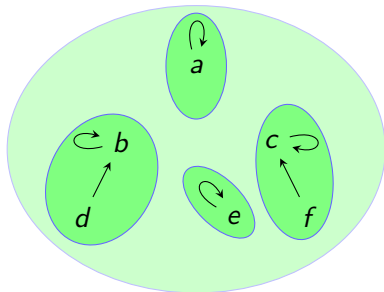


union-sets(d, b)

union-sets(f, e)

find-set(f) == e

Disjoint Set Operations



make-set(a)
make-set(b)
make-set(c)
make-set(d)
make-set(e)
make-set(f)

find-set(f) == f

union-sets(f, c)

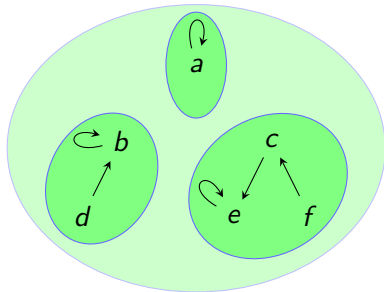
find-set(f) == c

union-sets(d, b)

→ union-sets(f, e)

find-set(f) == e

Disjoint Set Operations



make-set(a)

make-set(b)

make-set(c)

make-set(d)

make-set(e)

make-set(f)

find-set(f) == f

union-sets(f, c)

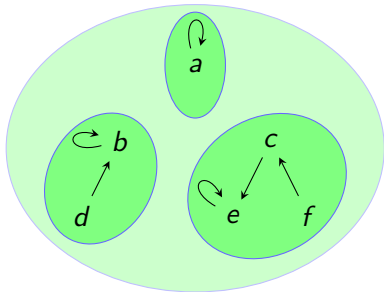
find-set(f) == c

union-sets(d, b)

union-sets(f, e)

→ find-set(f) == e

Disjoint Set Operations



make-set(a)

make-set(b)

make-set(c)

make-set(d)

make-set(e)

make-set(f)

find-set(f) == f

union-sets(f, c)

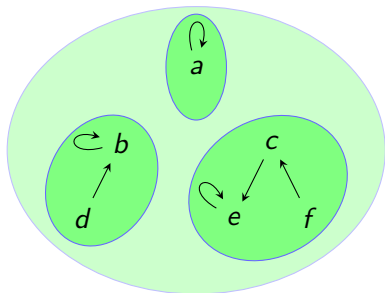
find-set(f) == c

union-sets(d, b)

union-sets(f, e)

→ find-set(f) == e

Disjoint Set Forests



tree = equivalence class

root = representative

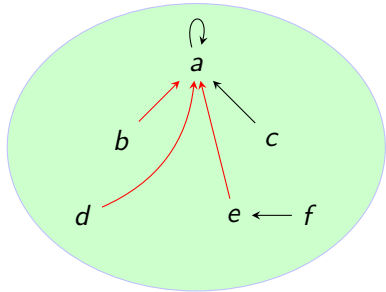
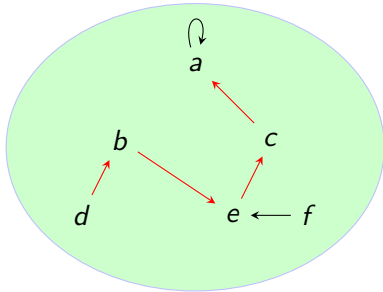
parent array

$$\begin{bmatrix} a & b & c & d & e & f \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ a & b & e & b & e & c \end{bmatrix}$$

parent relation

	<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>e</i>	<i>f</i>
<i>a</i>	1	0	0	0	0	0
<i>b</i>	0	1	0	0	0	0
<i>c</i>	0	0	0	0	1	0
<i>d</i>	0	1	0	0	0	0
<i>e</i>	0	0	0	0	1	0
<i>f</i>	0	0	1	0	0	0

Path Compression



$\text{find-set}(d) == a$

Implementation

make-set(x):

$p[x] = x$

$rank[x] = 0$

find-set(x):

if $x \neq p[x]$:

$p[x] = \text{find-set}(p[x])$

return $p[x]$

union-sets(x, y):

$r = \text{find-set}(x)$

$s = \text{find-set}(y)$

if $rank[r] > rank[s]$:

$p[s] = r$

else:

$p[r] = s$

if $rank[r] == rank[s]$:

$rank[s] += 1$

- union by rank makes smaller tree a subtree of larger tree
- $rank[x]$ is an upper bound on the depth of tree with root x
- time complexity almost linear in the total number of operations