```
id
         ::=
                                      Values
v
         ::=
                \mathbf{b}
                                         Boolean
                                        Number
                \mathbf{n}
                                         String
                \mathbf{str}
                undefined
                                         Undefined
                null
                                        null
                get(s, id)
                                Μ
                                         GetStore
                                      Store
s
         ::=
                                        PutStore
                put(s, id, v)
                                      VariableDeclaration
vd
         ::=
                id
                                        Declaration
                id = e
                                        Definition
                vd, vd'
                                        Multiple
                                      Expression
e
         ::=
                v
                                         Value
                                        Deref
                id
                                        Ref
                id = e
                                      Statement
m
         ::=
                                         Expression
                e
                \epsilon
                                         Skip
                m; m'
                                         Seq
                                         VarDeclaration
                \mathbf{var} \ vd
Tv
                                      Value Type
         ::=
                                        Number
                number
                                        Boolean
                boolean
                string
                                        String
                                        \quad \text{undefined} \quad
                \quad undefined \\
                null
                                        null
T
         ::=
                                      Expression Type
                                         ValueType
                Tv
                \operatorname{ref}\langle\operatorname{Tv}\rangle
                                        Location Type
Γ
         ::=
                                      Context
                                         EmptyContext
                Ø
                \{\Gamma\}
                                         \\Singleton Context
                \Gamma \cup \Gamma'
                                         UnionContext
                id:T
                                         ContextItem
```

R

::=

Relations

		= ! = ∈ ∉		Equals NotEquals Contains NotContains
Eq	::=	$Eq\ R\ Eq'$ C Γ X T id $func$		Equations
C	::= 	$egin{array}{l} arnothing \ C \cup C' \ C \cap C' \ \{Eq\} \ C \end{array}$		Constraint EmptyConstraint UnionConstraint IntersectConstraint ConstraintEquation BracketedConstraint
X	::= 			Variable EmptyVariable SingletonVariable UnionVariable IntersectVariable FreeVars
func	::=	$dom(\Gamma)$ $FV(T)$		Function
terminals	::= 	$ ightarrow \Gamma$		
formula	::=	judgement		
Jop	::= 	$< m, s > \rightarrow < m', s' >$ $\Gamma(id) = T$ $\Gamma(id) = T$	M M	

X func terminals

$< m, s > \overline{\rightarrow < m', s' >}$

 $\frac{\Gamma(id) = T}{\Gamma(id) = T}$

 $\Gamma \vdash e : T \mid_{X} C, \Gamma'$

$$\frac{\langle v; e, s \rangle \rightarrow \langle e, s \rangle}{\langle e_1; e_2, s \rangle \rightarrow \langle e'_1, s' \rangle} \text{ SEQ2}$$

$$\frac{\langle e_1; e_2, s \rangle \rightarrow \langle e'_1; e_2, s' \rangle}{\langle id = e, s \rangle \rightarrow \langle id = e', s' \rangle} \text{ ASSIGN1}$$

$$\frac{\langle id = v, s \rangle \rightarrow \langle v, put(s, id, v) \rangle}{\langle id = v, s \rangle \rightarrow \langle v, put(s, id, v) \rangle} \text{ ASSIGN2}$$

$$\frac{\langle id = v, s \rangle \rightarrow \langle v, put(s, id, v) \rangle}{\langle id, s \rangle \rightarrow \langle get(s, id), s \rangle} \text{ DEREF}$$

$$\frac{\langle var id = e, s \rangle \rightarrow \langle id = e, s \rangle}{\langle var id, s \rangle \rightarrow \langle undefined, s \rangle} \text{ VAR2}$$

$$\frac{\langle var vd, vd', s \rangle \rightarrow \langle var vd; var vd', s \rangle}{\langle var vd, vd', s \rangle \rightarrow \langle var vd; var vd', s \rangle} \text{ VAR3}$$

Μ

 $\overline{\Gamma \vdash \mathbf{n} : \mathbf{number} \ |_{\varnothing} \ \varnothing, \Gamma} \quad \text{V_NUM}$

```
\overline{\Gamma \vdash \mathbf{b} : \mathbf{boolean} \ |_{\varnothing} \ \varnothing, \Gamma} \quad \text{V\_BOOL}
                                                                                     \overline{\Gamma \vdash \mathbf{str} : \mathbf{string} \ |_{\varnothing} \ \varnothing, \Gamma} \quad \text{V\_STRING}
                                                             \overline{\Gamma \vdash \mathbf{undefined} : \mathbf{undefined} \mid_{\varnothing} \varnothing, \Gamma} \quad \text{V\_UNDEFINED}
                                                                                          \overline{\Gamma \vdash \mathbf{null} : \mathbf{null} \ |_{\varnothing} \ \varnothing, \Gamma} \quad \text{V\_NULL}
                                                                                               \frac{\Gamma(id) = T}{\Gamma \vdash id : T \mid_{\varnothing} \varnothing, \Gamma} \quad \text{IdType}
                                                                  \frac{id \notin dom(\Gamma)}{\Gamma \vdash id: T\mid_{\{T\}} \varnothing, \Gamma \cup \{id:T\}} \quad \text{IdTypeUndef}
                                                                                  \Gamma \vdash e : Tv_1 \mid_{X_1} C_1, \Gamma_1
                                                                                  \Gamma_1 \vdash id : T_2 \mid_{X_2} C_2, \Gamma_2
                                           \frac{X_1 \cap X_2 = \varnothing}{\Gamma \vdash id = e : T_2 \mid_{X_1 \cup X_2} C_1 \cup C_2 \cup \{T_2 = Tv_1\}, \Gamma_2}
\Gamma \vdash m \mid_X C, \Gamma'
                                                                                            \frac{}{\Gamma \vdash \epsilon \mid_{\varnothing} \varnothing, \Gamma} SkipTypable
                                                                     \overline{\Gamma \vdash \mathbf{var} \; id \; |_{\{T\}} \; \varnothing, \Gamma \cup \{id : T\}} \quad \mathsf{DECTYPABLE}
                                                                                   \frac{\Gamma \vdash e : T \mid_{X} C, \Gamma'}{\Gamma \vdash e \mid_{X} C, \Gamma'} \quad \text{ExpTypable}
                                                                                      \Gamma \vdash m_1 \mid_{X_1} C_1, \Gamma_1
                                                                     \Gamma_{1} \vdash m_{2} \mid_{X_{2}} C_{2}, \Gamma_{2}
X_{1} \cap X_{2} = \varnothing
\Gamma \vdash m_{1}; m_{2} \mid_{X_{1} \cup X_{2}} C_{1} \cup C_{2}, \Gamma_{2}
                                                                                                                                                                 SEQTYPABLE
                                                                             \Gamma \vdash \mathbf{var} \ id \mid_{X_1} C, \Gamma_1
                                                                     \Gamma_1 \vdash id = e : T \mid_{X_2} C, \Gamma_2
\frac{X_1 \cap X_2 = \varnothing}{\Gamma \vdash m_1; m_2 \mid_{X_1 \cup X_2} C_1 \cup C_2, \Gamma_2}
                                                                                                                                                                   DefTypable
```

 $\Gamma \vdash \mathbf{var} \ vd \mid_{X_1} C_1, \Gamma_1$ $\Gamma_1 \vdash \mathbf{var} \ vd' \mid_{X_1} C_2, \Gamma_2$

 $\frac{X_1 \cap X_2 = \varnothing}{\Gamma \vdash \mathbf{var} \ vd, \ vd' \mid_{X_1 \cup X_2} C_1 \cup C_2, \Gamma_2}$

Definition rules: 22 good 0 bad Definition rule clauses: 39 good 0 bad MULTIDECTYPABLE