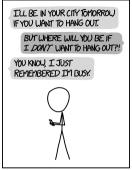
Discrete Mathematics MATH1064, Lecture 2

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WHY I TRY NOT TO BE PEDANTIC ABOUT CONDITIONALS.

Jonathan Spreer Discrete Mathematics 1/16

Propositions (aka Statements)

Definition (Proposition)

A proposition is a sentence that is true or false but not both.

- 1 Is it going to snow tomorrow?
- 4 He is sad.
- 3 Australia has exactly 4 states and 3 territories.
- $a^2 + b^2 = c^2.$
- **5** $a^2 + b^2 = c^2$, a = 3, b = 4, c = 5.
- Go to sleep.

Negation, Conjunction, Disjunction

Let p be a proposition.

Negation: $\neg p$ (read: not p), also written $\sim p$

Conjunction: $p \land q$ (read: p and q) Disjunction: $p \lor q$ (read: p or q)

Defined by the truth tables:

р	q	$p \wedge q$
Т	Т	Т
Т	F	F
F	Т	F
F	F	F

p	q	$p \lor q$
Т	Т	Т
Т	F	Т
F	Т	T
F	F	F

Note: $p \lor q$ is the "inclusive or"

Order of operation

¬ comes first

 \wedge and \vee are equal seconds

So:

$$\neg p \land q = (\neg p) \land q$$

But

$$p \land q \lor r$$

is ambiguous; one needs to use parentheses to say what is meant:

$$(p \land q) \lor r$$
 or $p \land (q \lor r)$

The "exclusive or"

Often written as $p \oplus q$

p	q	either p is true or q is true but not both
Т	Т	F
Т	F	Т
F F	Т	Т
F	F	F

Claim: This can be written as $(p \lor q) \land \neg (p \land q)$

p	q	$p \lor q$	$p \wedge q$	$\mid \neg(p \wedge q) \mid$	$ (p \lor q) \land \neg (p \land q)$
Т	Т				
Т	F				
F	Т				
F	F				

Variables and Compound Propositions

We've been thinking about p and q as propositions with some value (true or false). Now: think of them as variables (also called proposition variables or statement variables).

We use capital letters for compound propositions, which are made up from proposition variables (e.g. p, q, r, ...) and the symbols \neg, \land, \lor with unambiguous parentheses.

Example:
$$P = (p \lor q) \land \neg (p \land q)$$

If there are n variables in P, then the truth table for P has how many rows?

(1)
$$n$$
 (2) $2n$ (3) n^2 (4) 2^n (5) $n!$

Definition (Logical equivalence)

Two compound propositions P and Q are logically equivalent,

$$P \equiv Q$$
,

if they have identical truth values for every possible combination of truth values for their proposition variables.

Example from first lecture:

p	q	$p \wedge q$	$\big \; ((p \; \land \neg q) \lor (p \land q)) \land q$
Т	Т	T	T
Т	F	F	F
F	Т	F	F
F	F	F	F

Hence

$$((p \land \neg q) \lor (p \land q)) \land q \equiv p \land q$$

Which of the following propositions are logically equivalent?

- $Q = \neg (p \lor q)$

The correct answer is: $P \equiv S$ and $Q \equiv R$.

The equivalences are De Morgan's laws:

$$\neg(p \land q) \equiv (\neg p) \lor (\neg q)$$

$$\neg(p\vee q)\equiv(\neg p)\wedge(\neg q)$$

Contradictions and Tautologies

Definition (Contradiction)

A contradiction is a compound proposition which takes the value false for all possible truth values (true/false in all combinations) of its variables.

Example:
$$p \land \neg p$$
, $(q \land p) \land \neg p$
 $p \land (contradiction) \equiv (contradiction)$; $p \lor (contradiction) \equiv p$

Definition (Tautology)

A tautology is a compound proposition which takes the value true for all possible truth values (true/false in all combinations) of its variables.

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Example: p \lor \neg p

p \land (tautology) \equiv p; \quad p \lor (tautology) \equiv (tautology)
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Have a look at the table of logical equivalences on pages 27-28!

The statement

$$(p \land \neg q) \land (\neg p \lor q)$$

is a

- contradiction
- 2 tautology
- neither

The conditional

Let p and q be proposition variables; the conditional from p to q, $p \rightarrow q$, is defined by the following truth table:

р	q	p o q
Т	Т	Т
Т	F	F
F	Т	Т
F	F	Т

p is the hypothesis and q is the conclusion.

Example: If you want to hang out, I will be in your city tomorrow.

p = "You want to hang out", q = "I will be in your city tomorrow"

p	q	$p \rightarrow q$
Т	Т	T
Т	F	F
F	Т	Т
F	F	Т

Example: If you want to hang out, I will be in your city tomorrow.

p = "You want to hang out", q = "I will be in your city tomorrow"

The following propositions are consistent with $p \rightarrow q$:

You want to hang out and I will be in your city tomorrow. $(p \land q)$

You don't want to hang out and I will be in your city tomorrow. $(\neg p \land q)$

You don't want to hang out and I won't be in your city tomorrow. $(\neg p \land \neg q)$

The following proposition is inconsistent with $p \rightarrow q$:

You want to hang out and I won't be in your city tomorrow. $(p \land \neg q)$

Different ways of saying $p \rightarrow q$:

- p implies q
- $oldsymbol{2}$ if p, then q
- q if p
- $\bigcirc p$ only if q
- p is sufficient for q

One way to remember it:

 $p \rightarrow q$ is false if and only if it describes a counterexample; that is, the hypothesis is true but the conclusion is false

In order of operation: third, after \vee and \wedge (which come after \neg)

The conditional

Expressing the conditional in terms of and/or/not

$$p \to q \equiv \neg p \lor q$$

Friday:

Proof by truth table:

p	q	$p \rightarrow q$	$\neg p$	q	$\neg p \lor q$
Т	Т	Т			
Т	F	F			
F	T	T			
F	F	Т			

A final remark on truth tables

What do truth tables do?

- P and Q compound propositions (on n variables). Is $P \equiv Q$?
- Truth tables for P and Q are bit strings t_P and t_Q of length 2^n
- $P \equiv Q$ if and only if $t_P = t_Q$

What about the inverse problem? Given a bitstring t of length 2^n , produce a compound proposition with t as truth table.

Summary

A proposition is a sentence that is true or false but not both.

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Negation: \neg p (read: not p)
Conjunction: p \land q (read: p and q)
Disjunction: p \lor q (read: p or q)
Conditional: p \rightarrow q (read: p implies q)
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Two compound propositions P and Q are logically equivalent, $P \equiv Q$, if they have identical truth values for each possible combination of truth values for their proposition variables.

A contradiction is a compound proposition which takes the value false for all possible truth values of its variables.

A tautology is a compound proposition which takes the value true for all possible truth values of its variables.