

Discrete Mathematics

MATH1064, Lecture 18

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In order
to understand
recursion
you must
first understand
recursion.

Extra exercises for Lecture 18

Section 5.1: Problems 3, 4, 10, 11, 20, 23, 24, 31–34

Section 5.3: Problems 1–4, 7, 23–25

HEY T.Y., MAKE ME SOME
MATH LIMERICKS!

TO BECOME AN INDUCTION PROOF ACE,
ESTABLISH A TERM TO CALL "BASE."
THEN YOU CAN CONFIRM
YOUR CLAIM HOLDS FOR EACH TERM
IF IT'S TRUE FOR THE $N+1$ TH CASE.

IF I TELL YOU THAT I ALWAYS LIE
THEN MY LOGIC IS SURELY AWRY.
YOU CAN SAY WITH CONVICTION
THAT, BY CONTRADICTION,
I'M WRONG, AND PROBABLY HIGH.

OKAY!



A IMPLIES B IS A FACT
SO LET'S SEE HOW THEY INTERACT!
AS LONG AS A'S TRUE,
THEN B MUST BE, TOO!
BUT REVERSING THAT'S NOT QUITE EXACT.

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Another way to think of induction

Remember **modus ponens**?

1. p
2. $p \rightarrow q$
- c. $\therefore q$

Essentially, the **principle of mathematical induction** states that the following argument is valid:

1. $P(0)$
2. $\forall n \in \mathbb{N}, P(n) \rightarrow P(n+1)$
- c. $\therefore \forall n \in \mathbb{N}, P(n)$

Yet another way to think of induction

The principle of mathematical induction (version 2)

Let X be a set of natural numbers with the following properties:

- 1 The number 0 is in X .
- 2 For all $n \in \mathbb{N}$, if n is in X then $n + 1$ is also in X .

Then X is the set of **all** natural numbers, i.e., $X = \mathbb{N}$.

Just think of X as the set of all n for which $P(n)$ is true.

Mathematical induction is not a theorem—it is one of the properties that **defines** the natural numbers.

This is another story for another course!

Another induction example

Question

For which $n \in \mathbb{N}$ is $2^n < n!$?

1. **Exploration**
2. **Conjecture:** $2^n < n!$ for all $n \geq 4$
3. **Proof:** We will use induction (surprise!).

Proposition

For all integers $n \geq 4$, $2^n < n!$.

3. Proof: Let $P(n)$ be the predicate: $2^n < n!$

Basis step: $2^4 = 16$ and $4! = 24$. Since $16 < 24$, $P(4)$ is true.

Inductive step: Suppose $P(n)$ is true. That is, we assume $2^n < n!$. We must prove $P(n+1)$. That is, we must prove $2^{n+1} < (n+1)!$. We have:

$$2^{n+1} = 2 \cdot 2^n < 2 \cdot n!,$$

by the **inductive hypothesis**. So:

$$2^{n+1} < 2 \cdot n! < (n+1) \cdot n! = (n+1)!,$$

using the fact that $n \geq 4$. Therefore $P(n+1)$ is true.

Since $P(4)$ is true and $P(n) \rightarrow P(n+1)$ for $n \geq 4$, it follows by mathematical induction that $P(n)$ is true for all $n \geq 4$. □

Yet another example

Bernoulli's inequality

For all real $x > 0$ and all integers $n \geq 2$, $(1 + x)^n > 1 + nx$.

Proof. Let x be some particular (but arbitrary) positive real.

Basis step: Start with $n = 2$. We need to show $(1 + x)^2 > 1 + 2x$. This follows from $(1 + x)^2 = 1 + 2x + x^2 > 1 + 2x$, since $x > 0$.

Inductive step: Assume that $(1 + x)^n > 1 + nx$ for some $n \geq 2$. We need to show that $(1 + x)^{n+1} > 1 + (n + 1)x$.

We can rewrite the left hand side as

$$(1 + x)^{n+1} = (1 + x)^n(1 + x) > (1 + nx)(1 + x),$$

using the **inductive hypothesis** and the fact that $x > 0$. Then:

$$(1 + x)^{n+1} > (1 + nx)(1 + x) = 1 + (n + 1)x + nx^2 > 1 + (n + 1)x.$$

So: We showed that $(1 + x)^2 > 1 + 2x$, and also showed that

$$(1 + x)^n > 1 + nx \text{ implies } (1 + x)^{n+1} > 1 + (n + 1)x.$$

By mathematical induction, $(1 + x)^n > 1 + nx$ for all $n \geq 2$.



Loading the base case

The principle of strong mathematical induction

Let $P(n)$ be a predicate that is defined for all integers $n \geq a$, and let $b \geq a$. Suppose:

- 1 **Basis step:** $P(a), P(a+1), \dots, P(b-1)$ and $P(b)$ are all true.
- 2 **Inductive step:** For all integers $n \geq b$, if $P(a), P(a+1), \dots, P(n-1)$ and $P(n)$ are all true then $P(n+1)$ is also true.

Then $P(n)$ is true for all integers $n \geq a$.

It can be shown that this is **equivalent** to the ordinary principle of mathematical induction.

Try it!

Recursion

Recursion

Defining a mathematical expression using itself.

Easiest case: define $f : \mathbb{N} \rightarrow \mathbb{R}$ by

- $f(0) = a$ (initial condition)
- $f(n+1) = g(n, f(n))$ where g is known (recurrence relation)

Examples:

Factorial:

Fibonacci sequence:

Squares:

Another example

Remember the Tower of Hanoi?

The sequence $(T_n)_{n \geq 0}$ is defined by:

$$T_0 = 0$$

(initial conditions)

$$T_n = 2T_{n-1} + 1 \text{ for } n \geq 1$$

(recurrence relation)

How to convert this **recursive definition** into an **explicit formula**?

This is difficult in general, and you cannot always do it.

But, if you can **guess the formula**, it is easy to prove!

Just **check that it satisfies the recursive definition**, by substituting in your formula and checking that equality holds.

Tower of Hanoi

The recursive definition:

$$T_0 = 0, \quad T_n = 2T_{n-1} + 1 \text{ for } n \geq 1$$

Compute a few terms:

$$T_0 = 0 \quad T_1 = 1 \quad T_2 = 3 \quad T_3 = 7 \quad T_4 = 15 \quad T_5 = 31$$

Try to guess a formula: $T_n = 2^n - 1$

To prove this, substitute your formula into the recursive definition and check that it works:

- **Initial conditions:**

$$T_0 = 2^0 - 1 = 1 - 1 = 0, \text{ which satisfies the initial conditions}$$

- **Recurrence relation:**

$$T_n = 2^n - 1 = 2 \cdot 2^{n-1} - 1 = 2 \cdot (2^{n-1} - 1) + 1 = 2T_{n-1} + 1, \text{ which satisfies the recurrence relation}$$

Therefore **our explicit formula is correct!**

How can you guess the formula?

Method 1: Expand how each term is constructed

Consider $(a_n)_{n \geq 0}$, with $a_0 = 7$ and $a_n = a_{n-1} + 3$ for $n \geq 1$.

$$a_0 = 7$$

$$a_1 = 7 + 3$$

$$a_2 = 7 + 3 + 3$$

$$a_3 = 7 + 3 + 3 + 3$$

Conjecture: $a_n = 7 + 3n$

Proof:

- $a_0 = 7 + 3 \cdot 0 = 7$, which satisfies the **initial conditions**
- $a_n = 7 + 3n = 7 + 3(n-1) + 3 = a_{n-1} + 3$, which satisfies the **recurrence relation**

Method 2: Compute several terms and look for a pattern

Consider $(a_n)_{n \geq 0}$, with:

$$a_0 = 0, \quad a_1 = 1, \quad a_n = \frac{1}{2}(a_{n-1} + a_{n-2} + 1) + 3(n-1) \text{ for } n \geq 2$$

Some initial terms:

$$a_0 = 0, \quad a_1 = 1, \quad a_2 = 4, \quad a_3 = 9, \quad a_4 = 16, \quad a_5 = 25$$

Conjecture: $a_n = n^2$

Proof:

- $a_0 = 0^2 = 0$ and $a_1 = 1^2 = 1$, satisfying the **initial conditions**
- For the recurrence relation, we start on the right hand side:
$$\begin{aligned} & \frac{1}{2}(a_{n-1} + a_{n-2} + 1) + 3(n-1) = \\ & \frac{1}{2}[(n-1)^2 + (n-2)^2 + 1] + 3(n-1) = \\ & \frac{1}{2}(2n^2 - 6n + 5 + 1) + 3n - 3 = n^2 - 3n + 3 + 3n - 3 = n^2 \\ & = a_n, \text{ which satisfies the } \mathbf{recurrence\ relation} \end{aligned}$$

Other methods of guesswork

Method 3: The Online Encyclopaedia of Integer Sequences (oeis.org)

Method 4: Generating functions

And more: Transform the sequence, examine differences, ...

Guessing can be dangerous. . .

Fermat considered the sequence $(F_n)_{n \geq 0}$, where:

$$F_n = 2^{(2^n)} + 1$$

Some initial terms:

$$F_0 = 3 \quad F_1 = 5 \quad F_2 = 17 \quad F_3 = 257 \quad F_4 = 65537$$

These are **all prime**!

Fermat conjectured that F_n is prime for all n .

$$F_5 = 4294967297 = 641 \times 6700417$$

In fact, it is not known if **any** of these numbers are prime for $n > 4$!

Read about the history of the **Fermat primes** to learn more.