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## Cycles of Functions on Finite Types

We assume some notion of finite types which allows a proof of the pigeonhole principle:

**Proposition 1.** *If  $X$  is a finite type with  $N : \mathbb{N}$  elements, then every for every list  $x_0, \dots, x_N$  of  $N+1$  elements from  $X$ , there are  $i \neq j$  with  $x_i = x_j$ .*

Using this, we can show that by iterating functions on finite types, we must at some point enter a cycle, and we can even give some more details on the cycle length and when they occur.

**Definition 1.** *We define iteration for functions  $f : X \rightarrow X$  and a given starting value  $x : X$  by:*

$$f^0 x := x \quad f^{(n+1)} x := f(f^n x).$$

**Lemma 1.** *Let  $X$  be a type with  $N > 0$  elements. Then for every  $f : X \rightarrow X$  and  $x : X$ , there are  $k < N$  and  $1 < c \leq N$  such that  $f^{c+k} x = f^k x$ , meaning  $f^k x$  is a cycle for  $f$  with length at most length  $c$ .*

*Proof.* Consider the list of  $N+1$  values  $f^0 x, f^1 x, \dots, f^N x$  of type  $X$ . Since  $X$  has  $N$  elements, the pigeonhole principle tells us that at least two elements of the above list must coincide. Let this be  $f^n x$  and  $f^m x$  where  $n < m \leq N$  without loss of generality. We then have

$$f^n x = f^m x = f^{(m-n)+n} x$$

and therefore  $k := n$  and  $c := m - n > 1$  as desired.  $\square$

Note that both  $k$  and  $c$  depend on the specific  $f$  and  $x$  we choose at the start. The final result we want to show will state that there is a kind of *global cycle length*, which holds for all functions and for all starting values, and only depends on the number of elements of the finite type.

**Fact 1.** *Given  $n, m : \mathbb{N}$  we have  $f^{(n+m)} x = f^n(f^m x)$ .*

**Lemma 2.** *Given a function  $f : X \rightarrow X$  and  $k, m : \mathbb{N}$  with  $f^{k+m} x = f^m x$ , if  $m \leq n$  then for all  $d : \mathbb{N}$  we have  $f^{k \cdot d + n} x = f^n x$ .*

*Proof.* We proceed by induction on  $d$ . The case  $d = 0$  is trivial, and we additionally show the statement for  $d = 1$ . By assumption we have  $n - m \geq 0$  and therefore:

$$\begin{aligned} f^{k+n}x &= f^{(n-m)+(k+m)}x = f^{(n-m)}(f^{(k+m)}x) \\ &= f^{(n-m)}(f^m x) = f^{(n-m)+m}x = f^n x \end{aligned}$$

showing the claim. For the inductive step, we have:

$$f^{k \cdot (d+1) + n}x = f^{k + (k \cdot d + n)}x = f^k(f^{k \cdot d + n}x) = f^k(f^n x) = f^{k+n}x = f^n x$$

where we made use of the induction hypothesis and the  $d = 1$  case.  $\square$

**Theorem 1.** *Let  $X$  be finite with  $N > 0$  elements, and  $a : \mathbb{N}$  be divisible by all numbers  $1, \dots, N$ . Then for every function  $f : X \rightarrow X$  and  $x : X$  we have  $f^{a+(N-1)}x = f^{(N-1)}x$ .*

*Proof.* Let  $f : X \rightarrow X$  and  $x : X$  be given. Then by Lemma 1 there are  $k, c : \mathbb{N}$  with  $1 < c \leq N$  and  $f^{c+k}x = f^k x$ . By assumption,  $a$  is divisible by  $c$ , so there is  $d : \mathbb{N}$  with  $a = c \cdot d$ . Since  $k \leq N - 1$  we can make use Lemma 2 to conclude that:

$$f^{a+(N-1)}x = f^{c \cdot d + (N-1)}x = f^{(N-1)}x$$

$\square$

One number which satisfies the property of  $a$  in Theorem 1 is  $N!$ . The smallest number with this property is the least common multiple of the numbers  $1, \dots, N$ . The theorem can be read as saying that not matter the starting value  $x$  of function  $f : X \rightarrow X$ ; after  $(N - 1)$  applications of  $f$  to  $x$ , we get to a cycle and it has at most length  $a$ .